

# RM0090 Reference manual

STM32F405xx, STM32F407xx, STM32F415xx and STM32F417xx advanced ARM-based 32-bit MCUs

#### Introduction

This reference manual targets application developers. It provides complete information on how to use the STM32F405xx, STM32F407xx, STM32F415xx and STM32F417xx microcontroller memory and peripherals. The STM32F405xx, STM32F407xx, STM32F415xx and STM32F417xx will be referred to as STM32F40x and STM32F41x throughout the document, unless otherwise specified.

The STM32F40x and STM32F41x constitute a family of microcontrollers with different memory sizes, packages and peripherals.

For ordering information, mechanical and electrical device characteristics please refer to the STM32F40x and STM32F41x datasheets.

For information on programming, erasing and protection of the internal Flash memory please refer to the *STM32F40x* and *STM32F41x* Flash programming manual.

For information on the ARM Cortex<sup>™</sup>-M4F core, please refer to the *Cortex<sup>™</sup>-M4F Technical Reference Manual*.

#### Related documents

Related documents Available from www.arm.com:

- STM32F40x and STM32F41x datasheets
- STM32F40x and STM32F41x Flash programming manual (PM0081)
- For information on the ARM Cortex<sup>™</sup>-M4 core with FPU, please refer to the STM32F3xx/F4xxx Cortex<sup>™</sup>-M4 programming manual (PM0214).

Table 1. Applicable products

Product family	Part numbers and product categories
Microcontrollers	STM32F405xx, STM32F407xx STM32F415xx, STM32F417xx

October 2012 Doc ID 018909 Rev 2 1/1340

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## 1 Documentation conventions

The Cortex-M4 core with FPU will be referred to as Cortex-M4F throughout this document.

# 1.1 List of abbreviations for registers

The following abbreviations are used in register descriptions:

read/write (rw) Software can read and write to these bits. read-only (r) Software can only read these bits. write-only (w) Software can only write to this bit. Reading the bit returns the reset value. Software can read as well as clear this bit by writing 1. Writing '0' has read/clear (rc\_w1) no effect on the bit value. read/clear (rc\_w0) Software can read as well as clear this bit by writing 0. Writing '1' has no effect on the bit value. read/clear by read Software can read this bit. Reading this bit automatically clears it to '0'. (rc\_r) Writing '0' has no effect on the bit value. read/set (rs) Software can read as well as set this bit. Writing '0' has no effect on the bit value. read-only write Software can read this bit. Writing '0' or '1' triggers an event but has no trigger (rt\_w) effect on the bit value. Software can only toggle this bit by writing '1'. Writing '0' has no effect. toggle (t)

# 1.2 Peripheral availability

Reserved (Res.)

For peripheral availability and number across all STM32F40x and STM32F41x sales types, please refer to the STM32F40x and STM32F41x datasheets.

Reserved bit, must be kept at reset value.

# 2 Memory and bus architecture

# 2.1 System architecture

The main system consists of 32-bit multilayer AHB bus matrix that interconnects:

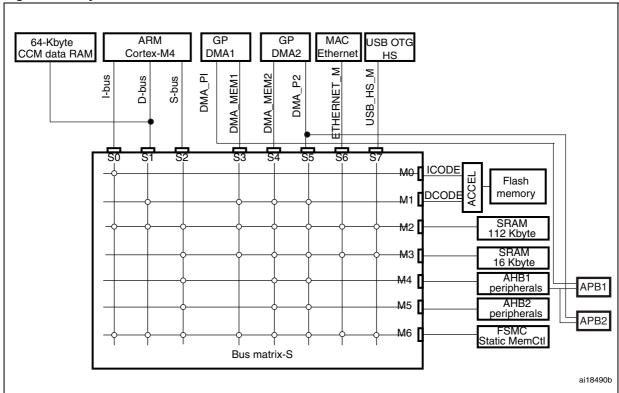
- Eight masters:
  - Cortex<sup>™</sup>-M4F core I-bus, D-bus and S-bus
  - DMA1 memory bus
  - DMA2 memory bus
  - DMA2 peripheral bus
  - Ethernet DMA bus
  - USB OTG HS DMA bus
- Seven slaves:
  - Internal Flash memory ICode bus
  - Internal Flash memory DCode bus
  - Main internal SRAM1 (112 KB)
  - Auxiliary internal SRAM2 (16 KB)
  - AHB1peripherals including AHB to APB bridges and APB peripherals
  - AHB2 peripherals
  - FSMC

The bus matrix provides access from a master to a slave, enabling concurrent access and efficient operation even when several high-speed peripherals work simultaneously. This architecture is shown in *Figure 1*.

Note:

The 64-Kbyte CCM (core coupled memory) data RAM is not part of the bus matrix (see Figure 1: System architecture). It can be accessed only through the CPU.

Figure 1. System architecture



#### 2.1.1 S0: I-bus

This bus connects the Instruction bus of the Cortex<sup>™</sup>-M4F core to the BusMatrix. This bus is used by the core to fetch instructions. The target of this bus is a memory containing code (internal Flash memory/SRAM or external memories through the FSMC).

### 2.1.2 S1: D-bus

This bus connects the databus of the Cortex<sup>™</sup>-M4F and the 64-Kbyte CCM data RAM to the BusMatrix. This bus is used by the core for literal load and debug access. The target of this bus is a memory containing code or data (internal Flash memory or external memories through the FSMC).

#### 2.1.3 S2: S-bus

This bus connects the system bus of the Cortex<sup>TM</sup>-M4F core to a BusMatrix. This bus is used to access data located in a peripheral or in SRAM. Instructions may also be fetch on this bus (less efficient than ICode). The targets of this bus are the 112, and 16 Kbytes

internal SRAMs, the AHB1 peripherals including the APB peripherals, the AHB2 peripherals and the external memories through the FSMC.

### 2.1.4 S3, S4: DMA memory bus

This bus connects the DMA memory bus master interface to the BusMatrix. It is used by the DMA to perform transfer to/from memories. The targets of this bus are data memories: internal SRAMs (112, and 16 Kbytes) and external memories through the FSMC.

## 2.1.5 S5: DMA peripheral bus

This bus connects the DMA peripheral master bus interface to the BusMatrix. This bus is used by the DMA to access AHB peripherals or to perform memory-to-memory transfers. The targets of this bus are the AHB and APB peripherals plus data memories: internal SRAM and external memories through the FSMC.

### 2.1.6 S6: Ethernet DMA bus

This bus connects the Ethernet DMA master interface to the BusMatrix. This bus is used by the Ethernet DMA to load/store data to a memory. The targets of this bus are data memories: internal SRAMs (112, and 16 Kbytes) and external memories through the FSMC.

#### 2.1.7 S7: USB OTG HS DMA bus

This bus connects the USB OTG HS DMA master interface to the BusMatrix. This bus is used by the USB OTG DMA to load/store data to a memory. The targets of this bus are data memories: internal SRAMs (112, and 16 Kbytes) and external memories through the FSMC.

### 2.1.8 BusMatrix

The BusMatrix manages the access arbitration between masters. The arbitration uses a round-robin algorithm.

### 2.1.9 AHB/APB bridges (APB)

The two AHB/APB bridges, APB1 and APB2, provide full synchronous connections between the AHB and the two APB buses, allowing flexible selection of the peripheral frequency.

Refer to the device datasheets for more details on APB1 and APB2 maximum frequencies, and to *Table 2 on page 50* for the address mapping of AHB and APB peripherals.

After each device reset, all peripheral clocks are disabled (except for the SRAM and Flash memory interface). Before using a peripheral you have to enable its clock in the RCC\_AHBxENR or RCC\_APBxENR register.

Note:

When a 16- or an 8-bit access is performed on an APB register, the access is transformed into a 32-bit access: the bridge duplicates the 16- or 8-bit data to feed the 32-bit vector.

# 2.2 Memory organization

Program memory, data memory, registers and I/O ports are organized within the same linear 4 Gbyte address space.

The bytes are coded in memory in little endian format. The lowest numbered byte in a word is considered the word's least significant byte and the highest numbered byte, the word's most significant.

For the detailed mapping of peripheral registers, please refer to the related chapters.

The addressable memory space is divided into 8 main blocks, each of 512 MB.

All the memory areas that are not allocated to on-chip memories and peripherals are considered "Reserved"). Refer to the memory map figure in the product datasheet.

# 2.3 Memory map

See the datasheet corresponding to your device for a comprehensive diagram of the memory map. *Table 2* gives the boundary addresses of the peripherals available in all STM32F4xx devices.

Table 2. STM32F4xx register boundary addresses

Boundary address	Peripheral	Bus	Register map
0xA000 0000 - 0xA000 0FFF	FSMC control register	AHB3	Section 31.6.9: FSMC register map on page 1293 Section 36.8: FMC register map on page 1592
0x5006 0800 - 0X5006 0BFF	RNG		Section 20.4.4: RNG register map on page 532
0x5006 0400 - 0X5006 07FF	HASH		Section 21.4.8: HASH register map on page 551
0x5006 0000 - 0X5006 03FF	CRYP	AHB2	Section 19.6.11: CRYP register map on page 527
0x5005 0000 - 0X5005 03FF	DCMI	711152	Section 12.8.12: DCMI register map on page 296
0x5000 0000 - 0X5003 FFFF	USB OTG FS		Section 29.16.6: OTG_FS register map on page 1031

Table 2. STM32F4xx register boundary addresses (continued)

Boundary address	Peripheral	Bus	Register map
0x4004 0000 - 0x4007 FFFF	USB OTG HS		Section 30.12.6: OTG_HS register map on page 1172
0x4002 9000 - 0x4002 93FF			
0x4002 8C00 - 0x4002 8FFF			
0x4002 8800 - 0x4002 8BFF	ETHERNET MAC		Section 28.8.5: Ethernet register maps on page 942
0x4002 8400 - 0x4002 87FF			
0x4002 8000 - 0x4002 83FF			
0x4002 6400 - 0x4002 67FF	DMA2		Costian 9 F 11: DMA register man on page 10F
0x4002 6000 - 0x4002 63FF	DMA1		Section 8.5.11: DMA register map on page 195
0x4002 4000 - 0x4002 4FFF	BKPSRAM		
0x4002 3C00 - 0x4002 3FFF	Flash interface register	AHB1	See Flash programming manual
0x4002 3800 - 0x4002 3BFF	RCC		Section 5.3.25: RCC register map on page 137
0x4002 3000 - 0x4002 33FF	CRC		Section 3.4.4: CRC register map on page 62
0x4002 2000 - 0x4002 23FF	GPIOI		
0x4002 1C00 - 0x4002 1FFF	GPIOH		
0x4002 1800 - 0x4002 1BFF	GPIOG		
0x4002 1400 - 0x4002 17FF	GPIOF		
0x4002 1000 - 0x4002 13FF	GPIOE		Section 6.4.11: GPIO register map on page 156
0X4002 0C00 - 0x4002 0FFF	GPIOD		
0x4002 0800 - 0x4002 0BFF	GPIOC		
0x4002 0400 - 0x4002 07FF	GPIOB	7	
0x4002 0000 - 0x4002 03FF	GPIOA		
0x4001 4800 - 0x4001 4BFF	TIM11		Section 15.6.11: TIM10/11/13/14 register map on
0x4001 4400 - 0x4001 47FF	TIM10		page 469
0x4001 4000 - 0x4001 43FF	TIM9	APB2	Section 15.5.14: TIM9/12 register map on page 459
0x4001 3C00 - 0x4001 3FFF	EXTI		Section 9.3.7: EXTI register map on page 209
0x4001 3800 - 0x4001 3BFF	SYSCFG		Section 7.2.8: SYSCFG register maps on page 156
0x4001 3000 - 0x4001 33FF	SPI1		Section 25.5.10: SPI register map on page 726
0x4001 2C00 - 0x4001 2FFF	SDIO		Section 26.9.16: SDIO register map on page 782
0x4001 2000 - 0x4001 23FF	ADC1 - ADC2 - ADC3	<u> </u>	Section 10.13.18: ADC register map on page 252
0x4001 1400 - 0x4001 17FF	USART6	APB2	Section 24.6.8: USART register map on page 674
0x4001 1000 - 0x4001 13FF	USART1		Geolion 24.0.0. GGART Tegister Map on page 674
0x4001 0400 - 0x4001 07FF	TIM8		Section 13.4.21: TIM1&TIM8 register map on
0x4001 0000 - 0x4001 03FF	TIM1		page 365

Table 2. STM32F4xx register boundary addresses (continued)

Boundary address	Peripheral	Bus	Register map
0x4000 7400 - 0x4000 77FF	DAC		Section 11.5.15: DAC register map on page 274
0x4000 7000 - 0x4000 73FF	PWR		Section 4.4.3: PWR register map on page 82
0x4000 6800 - 0x4000 6BFF	CAN2		Section 27.9.5: bxCAN register map on page 825
0x4000 6400 - 0x4000 67FF	CAN1		Gection 27.3.3. bxOAN register map on page 623
0x4000 5C00 - 0x4000 5FFF	I2C3		
0x4000 5800 - 0x4000 5BFF	I2C2		Section 23.6.10: I2C register map on page 622
0x4000 5400 - 0x4000 57FF	I2C1		
0x4000 5000 - 0x4000 53FF	UART5		
0x4000 4C00 - 0x4000 4FFF	UART4		Section 24.6.8: USART register map on page 674
0x4000 4800 - 0x4000 4BFF	USART3		Section 24.0.0. USANT Tegister map on page 074
0x4000 4400 - 0x4000 47FF	USART2		
0x4000 4000 - 0x4000 43FF	I2S3ext		
0x4000 3C00 - 0x4000 3FFF	SPI3 / I2S3		Section 25.5.10: SPI register map on page 726
0x4000 3800 - 0x4000 3BFF	SPI2 / I2S2	APB1	Section 25.5.10. SET register map on page 726
0x4000 3400 - 0x4000 37FF	I2S2ext	/	
0x4000 3000 - 0x4000 33FF	IWDG		Section 17.4.5: IWDG register map on page 485
0x4000 2C00 - 0x4000 2FFF	WWDG		Section 18.6.4: WWDG register map on page 492
0x4000 2800 - 0x4000 2BFF	RTC & BKP Registers		Section 22.6.21: RTC register map on page 589
0x4000 2000 - 0x4000 23FF	TIM14		Section 15.6.11: TIM10/11/13/14 register map on
0x4000 1C00 - 0x4000 1FFF	TIM13		page 469
0x4000 1800 - 0x4000 1BFF	TIM12		Section 15.5.14: TIM9/12 register map on page 459
0x4000 1400 - 0x4000 17FF	TIM7		Section 16.4.9: TIM6&TIM7 register map on
0x4000 1000 - 0x4000 13FF	TIM6		page 480
0x4000 0C00 - 0x4000 0FFF	TIM5		
0x4000 0800 - 0x4000 0BFF	TIM4		Section 14.4.21: TIMy register man on page 425
0x4000 0400 - 0x4000 07FF	TIM3		Section 14.4.21: TIMx register map on page 425
0x4000 0000 - 0x4000 03FF	TIM2		

## 2.3.1 Embedded SRAM

The STM32F40x and STM32F41x feature 4 Kbytes of backup SRAM (see *Section 4.1.2: Battery backup domain*) plus 192 Kbytes of system SRAM.

The system SRAM can be accessed as bytes, half-words (16 bits) or full words (32 bits). Read and write operations are performed at CPU speed with 0 wait state. The system SRAM is split up into three blocks:

- 112 Kbyte and 16 Kbyte blocks mapped at address 0x2000 0000 and accessible by all AHB masters.
- A 64 Kbyte block mapped at address 0x1000 0000 and accessible only by the CPU through the D-bus. AHB masters support concurrent SRAM accesses (from the Ethernet or the USB OTG HS): for instance, the Ethernet MAC can read/write from/to the 16 Kbyte SRAM while the CPU is reading/writing from/to the 112 Kbyte SRAM.

The CPU can access the system SRAM through the System Bus or through the I-Code/D-Code buses when boot from SRAM is selected or when physical remap is selected (*Section 7.2.1: SYSCFG memory remap register (SYSCFG\_MEMRMP)* in the SYSCFG controller). To get the max performance on SRAM execution, physical remap should be selected (boot or software selection).

## 2.3.2 Bit banding

The Cortex<sup>™</sup>-M4F memory map includes two bit-band regions. These regions map each word in an alias region of memory to a bit in a bit-band region of memory. Writing to a word in the alias region has the same effect as a read-modify-write operation on the targeted bit in the bit-band region.

In the STM32F4xx devices both the peripheral registers and the SRAM are mapped to a bit-band region, so that single bit-band write and read operations are allowed. The operations are only available for Cortex<sup>™</sup>-M4F accesses, and not from other bus masters (e.g. DMA).

A mapping formula shows how to reference each word in the alias region to a corresponding bit in the bit-band region. The mapping formula is:

```
bit_word_addr = bit_band_base + (byte_offset x 32) + (bit_number x 4)
where:
```

- bit\_word\_addr is the address of the word in the alias memory region that maps to the targeted bit
- bit band base is the starting address of the alias region
- byte\_offset is the number of the byte in the bit-band region that contains the targeted bit
- bit\_number is the bit position (0-7) of the targeted bit

#### Example

The following example shows how to map bit 2 of the byte located at SRAM address 0x20000300 to the alias region:

```
0x22006008 = 0x22000000 + (0x300*32) + (2*4)
```

Writing to address 0x22006008 has the same effect as a read-modify-write operation on bit 2 of the byte at SRAM address 0x20000300.

Reading address 0x22006008 returns the value (0x01 or 0x00) of bit 2 of the byte at SRAM address 0x20000300 (0x01: bit set; 0x00: bit reset).

For more information on bit-banding, please refer to the *Cortex™-M4F programming manual* (see *Related documents on page 1*).

### 2.3.3 Embedded Flash memory

The Flash memory has the following main features:

- Capacity up to 1 Mbyte
- 128 bits wide data read
- Byte, half-word, word and double word write
- Sector and mass erase
- Memory organization

The Flash memory is organized as follows:

- A main memory block divided into 4 sectors of 16 Kbytes, 1 sector of 64 Kbytes, and 7 sectors of 128 Kbytes
- System memory from which the device boots in System memory boot mode
- 512 OTP (one-time programmable) bytes for user data
   The OTP area contains 16 additional bytes used to lock the corresponding OTP data block.
- Option bytes to configure read and write protection, BOR level, watchdog software/hardware and reset when the device is in Standby or Stop mode.

Table 3. Flash module organization

Block	Name	Block base addresses	Size
	Sector 0	0x0800 0000 - 0x0800 3FFF	16 Kbytes
	Sector 1	0x0800 4000 - 0x0800 7FFF	16 Kbytes
	Sector 2	0x0800 8000 - 0x0800 BFFF	16 Kbytes
	Sector 3	0x0800 C000 - 0x0800 FFFF	16 Kbytes
	Sector 4	0x0801 0000 - 0x0801 FFFF	64 Kbytes
Main memory	Sector 5	0x0802 0000 - 0x0803 FFFF	128 Kbytes
	Sector 6	0x0804 0000 - 0x0805 FFFF	128 Kbytes
			·
	Sector 11	0x080E 0000 - 0x080F FFFF	128 Kbytes
System	memory	0x1FFF 0000 - 0x1FFF 77FF	30 Kbytes
0	TP	0x1FFF 7800 - 0x1FFF 7A0F	528 bytes
Option	n bytes	0x1FFF C000 - 0x1FFF C00F	16 bytes

### 2.3.4 Flash memory read interface

### Relation between CPU clock frequency and Flash memory read time

To correctly read data from Flash memory, the number of wait states (LATENCY) must be correctly programmed in the *Flash access control register (FLASH\_ACR) for STM32F40x and STM32F41x* and according to the frequency of the Cortex<sup>TM</sup>-M4F clock and the supply voltage of the device. The correspondence between wait states and core clock frequency is

given in *Table 4*. The number of wait states for each voltage range must be extended to acheive de maximum frequency at low voltage.

Note: When VOS = '0' the maximum value of  $f_{HCLK} = 144$  MHz.

Table 4. Number of wait states according to CPU clock (HCLK) frequency

		HCLK	(MHz)	
Wait states (WS) (LATENCY)	Voltage range 2.7 V - 3.6 V	Voltage range 2.4 V - 2.7 V	Voltage range 2.1 V - 2.4 V	Voltage range 1.8 V - 2.1 V <sup>(1)</sup> Prefetch OFF
0 WS (1 CPU cycle)	0 <hclk≤ 30<="" td=""><td>0 <hclk 24<="" td="" ≤=""><td>0 <hclk 22<="" td="" ≤=""><td>0 &lt; HCLK ≤ 20</td></hclk></td></hclk></td></hclk≤>	0 <hclk 24<="" td="" ≤=""><td>0 <hclk 22<="" td="" ≤=""><td>0 &lt; HCLK ≤ 20</td></hclk></td></hclk>	0 <hclk 22<="" td="" ≤=""><td>0 &lt; HCLK ≤ 20</td></hclk>	0 < HCLK ≤ 20
1 WS (2 CPU cycles)	30 <hclk 60<="" td="" ≤=""><td>24 &lt; HCLK≤ 48</td><td>22 <hclk 44<="" td="" ≤=""><td>20 <hclk 40<="" td="" ≤=""></hclk></td></hclk></td></hclk>	24 < HCLK≤ 48	22 <hclk 44<="" td="" ≤=""><td>20 <hclk 40<="" td="" ≤=""></hclk></td></hclk>	20 <hclk 40<="" td="" ≤=""></hclk>
2 WS (3 CPU cycles)	60 <hclk 90<="" td="" ≤=""><td>48 &lt; HCLK≤ 72</td><td>44 &lt; HCLK≤ 66</td><td>40 &lt; HCLK≤ 60</td></hclk>	48 < HCLK≤ 72	44 < HCLK≤ 66	40 < HCLK≤ 60
3 WS (4 CPU cycles)	90 <hclk 120<="" td="" ≤=""><td>72 &lt; HCLK≤ 96</td><td>66 <hclk 88<="" td="" ≤=""><td>60 &lt; HCLK≤ 80</td></hclk></td></hclk>	72 < HCLK≤ 96	66 <hclk 88<="" td="" ≤=""><td>60 &lt; HCLK≤ 80</td></hclk>	60 < HCLK≤ 80
4 WS (5 CPU cycles)	120 <hclk 150<="" td="" ≤=""><td>96 &lt; HCLK≤ 120</td><td>88 &lt; HCLK≤ 110</td><td>80 &lt; HCLK≤ 100</td></hclk>	96 < HCLK≤ 120	88 < HCLK≤ 110	80 < HCLK≤ 100
5 WS (6 CPU cycles)	150 <hclk 168<="" td="" ≤=""><td>120 <hclk 144<="" td="" ≤=""><td>110 &lt; HCLK≤ 132</td><td>100 &lt; HCLK≤ 120</td></hclk></td></hclk>	120 <hclk 144<="" td="" ≤=""><td>110 &lt; HCLK≤ 132</td><td>100 &lt; HCLK≤ 120</td></hclk>	110 < HCLK≤ 132	100 < HCLK≤ 120
6 WS (7 CPU cycles)		144 <hclk 168<="" td="" ≤=""><td>132 &lt; HCLK≤ 154</td><td>120 &lt; HCLK≤ 140</td></hclk>	132 < HCLK≤ 154	120 < HCLK≤ 140
7 WS (8 CPU cycles)			154 <hclk 168<="" td="" ≤=""><td>140 &lt; HCLK≤ 160</td></hclk>	140 < HCLK≤ 160

<sup>1.</sup> If an inverted reset is applied to PDR\_ON, this value can be lowered to 1.7 V.

After reset, the CPU clock frequency is 16 MHz and 0 wait state (WS) is configured in the FLASH\_ACR register.

It is highly recommended to use the following software sequences to tune the number of wait states needed to access the Flash memory with the CPU frequency.

### Increasing the CPU frequency

- Program the new number of wait states to the LATENCY bits in the FLASH\_ACR register
- Check that the new number of wait states is used to access the Flash memory by reading the FLASH\_ACR register
- Modify the CPU clock source by writing the SW bits in the RCC clock configuration register (RCC\_CFGR)
- If needed, modify the CPU clock prescaler by writing the HPRE bits in RCC\_CFGR
- Check that the new CPU clock source or/and the new CPU clock prescaler value is/are taken into account by reading the clock source status (SWS bits) or/and the AHB prescaler value (HPRE bits), respectively, in the RCC\_CFGR register

### **Decreasing the CPU frequency**

- Modify the CPU clock source by writing the SW bits in the RCC\_CFGR register
- If needed, modify the CPU clock prescaler by writing the HPRE bits in RCC\_CFGR
- Check that the new CPU clock source or/and the new CPU clock prescaler value is/are taken into account by reading the clock source status (SWS bits) or/and the AHB prescaler value (HPRE bits), respectively, in the RCC\_CFGR register
- Program the new number of wait states to the LATENCY bits in FLASH ACR
- Check that the new number of wait states is used to access the Flash memory by reading the FLASH\_ACR register

Note:

A change in CPU clock configuration or wait state (WS) configuration may not be effective straight away. To make sure that the current CPU clock frequency is the one you have configured, you can check the AHB prescaler factor and clock source status values. To make sure that the number of WS you have programmed is effective, you can read the FLASH\_ACR register.

The FLASH\_ACR register is used to enable/disable the acceleration features and control the Flash memory access time according to CPU frequency. The tables below provides the bit map and bit descriptions for this register.

For complete information on Flash memory operations and register configurations, please refer to the STM32F40x and STM32F41x Flash programming manual (PM0059).

### Flash access control register (FLASH\_ACR) for STM32F40x and STM32F41x

The Flash access control register is used to enable/disable the acceleration features and control the Flash memory access time according to CPU frequency.

Address offset: 0x00 Reset value: 0x0000 0000

Access: no wait state, word, half-word and byte access

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Res	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Dagamya	J	DCRST	ICRST	DCEN	ICEN	PRFTEN			Dagam, 40.0	1			LATENCY	,
	Reserved	1	rw	W	rw	rw	rw			Reserved	ı		rw	rw	rw

Bits 31:11 Reserved, must be kept cleared.

Bit 12 DCRST: Data cache reset

0: Data cache is not reset

1: Data cache is reset

This bit can be written only when the D cache is disabled.

Bit 11 ICRST: Instruction cache reset

0: Instruction cache is not reset

1: Instruction cache is reset

This bit can be written only when the I cache is disabled.

Bit 10 DCEN: Data cache enable

0: Data cache is disabled

1: Data cache is enabled

Bit 9 ICEN: Instruction cache enable

0: Instruction cache is disabled

1: Instruction cache is enabled

Bit 8 PRFTEN: Prefetch enable

0: Prefetch is disabled1: Prefetch is enabled

Bits 7:3 Reserved, must be kept cleared.

Bits 2:0 LATENCY: Latency

These bits represent the ratio of the CPU clock period to the Flash memory access time.

000: Zero wait state

001: One wait state

010: Two wait states

011: Three wait states

100: Four wait states

101: Five wait states

110: Six wait states

111: Seven wait states

### 2.3.5 Adaptive real-time memory accelerator (ART Accelerator™)

The ART Accelerator<sup>™</sup> is a memory accelerator which is optimized for STM32 industry-standard Cortex<sup>™</sup>-M4F processors. It balances the inherent performance advantage of the ARM Cortex<sup>™</sup>-M4F over Flash memory technologies, which normally requires the processor to wait for the Flash memory at higher operating frequencies. Thanks to the ART Accelerator<sup>™</sup>, the CPU can operate up to 168 MHz frequency without wait states, thereby increasing the overall system speed and efficiency (see *Table 4*).

To release the processor 210 DMIPS performance at this frequency, the accelerator implements an instruction prefetch queue and branch cache, which enables program execution from Flash memory at up to 168 MHz without wait states.

# 2.4 Boot configuration

Due to its fixed memory map, the code area starts from address 0x0000 0000 (accessed through the ICode/DCode buses) while the data area (SRAM) starts from address 0x2000 0000 (accessed through the system bus). The Cortex<sup>™</sup>-M4F CPU always fetches the reset vector on the ICode bus, which implies to have the boot space available only in the code area (typically, Flash memory). STM32F4xx microcontrollers implement a special mechanism to be able to boot from other memories (like the internal SRAM).

In the STM32F4xx, three different boot modes can be selected through the BOOT[1:0] pins as shown in *Table 5*.

Table 5. Boot modes

Boot mode s	election pins	Boot mode	Aliasing				
BOOT1	воото	Boot mode	Allasiliy				
Х	0	Main Flash memory	Main Flash memory is selected as the boot spa				

	20010000						
Boot mode s	selection pins	Boot mode	Aliasing				
BOOT1	воото	Boot mode	Allasing				
0	1	System memory	System memory is selected as the boot space				
1	1	Embedded SRAM	Embedded SRAM is selected as the boot space				

Table 5. Boot modes

The values on the BOOT pins are latched on the 4th rising edge of SYSCLK after a reset. It is up to the user to set the BOOT1 and BOOT0 pins after reset to select the required boot mode.

BOOT0 is a dedicated pin while BOOT1 is shared with a GPIO pin. Once BOOT1 has been sampled, the corresponding GPIO pin is free and can be used for other purposes.

The BOOT pins are also resampled when the device exits the Standby mode. Consequently, they must be kept in the required Boot mode configuration when the device is in the Standby mode. After this startup delay is over, the CPU fetches the top-of-stack value from address 0x0000 0000, then starts code execution from the boot memory starting from 0x0000 0004.

Note:

When the device boots from SRAM, in the application initialization code, you have to relocate the vector table in SRAM using the NVIC exception table and the offset register.

#### Embedded bootloader

The embedded bootloader mode is used to reprogram the Flash memory using one of the following serial interfaces:

- USART1(PA9/PA10)
- USART3(PB10/11 and PC10/11)
- CAN2(PB5/13)
- USB OTG FS(PA11/12) in Device mode (DFU: device firmware upgrade).

The USART peripherals operate at the internal 16 MHz oscillator (HSI) frequency, while the CAN and USB OTG FS require an external clock (HSE) multiple of 1 MHz (ranging from 4 to 26 MHz).

The embedded bootloader code is located in system memory. It is programmed by ST during production. For additional information, refer to application note AN2606.

#### Physical remap

Once the boot pins are selected, the application software can modify the memory accessible in the code area (in this way the code can be executed through the ICode bus in place of the System bus). This modification is performed by programming the *Section 7.2.1:* SYSCFG memory remap register (SYSCFG\_MEMRMP) in the SYSCFG controller.

The following memories can thus be remapped:

- Main Flash memory
- System memory
- Embedded SRAM1 (112 KB)
- FSMC Bank 1 (NOR/PSRAM 1 and 2)

Table 6. Memory mapping vs. Boot mode/physical remap

Addresses	Boot/Remap in main Flash memory	Boot/Remap in embedded SRAM	Boot/Remap in System memory	Remap in FSMC
0x2001 C000 - 0x2001 FFFF	SRAM2 (16 KB)	SRAM2 (16 KB)	SRAM2 (16 KB)	SRAM2 (16 KB)
0x2000 0000 - 0x2001 BFFF	SRAM1 (112 KB)	SRAM1 (112 KB)	SRAM1 (112 KB)	SRAM1 (112 KB)
0x1FFF 0000 - 0x1FFF 77FF	System memory	System memory	System memory	System memory
0x0810 0000 - 0x0FFF FFFF	Reserved	Reserved	Reserved	Reserved
0x0800 0000 - 0x080F FFFF	Flash memory	Flash memory	Flash memory	Flash memory
0x0400 0000 - 0x07FF FFFF	Reserved	Reserved	Reserved	FSMC Bank1 NOR/PSRAM 2 (Aliased)
0x0000 0000 - 0x03FF FFFF <sup>(1)(2)</sup>	Flash (1 MB) Aliased	SRAM1 (112 KB) Aliased	System memory (30 KB) Aliased	FSMC Bank1 NOR/PSRAM 1 (Aliased)

When the FSMC is remapped at address 0x0000 0000, only the first two regions of Bank 1 memory controller (Bank1 NOR/PSRAM 1 and NOR/PSRAM 2) can be remapped. In remap mode, the CPU can access the external memory via ICode bus instead of System bus which boosts up the performance.

<sup>2.</sup> Even when aliased in the boot memory space, the related memory is still accessible at its original memory space.

RM0090 CRC calculation unit

### 3 CRC calculation unit

This section applies to the whole STM32F4xx family, unless otherwise specified.

### 3.1 CRC introduction

The CRC (cyclic redundancy check) calculation unit is used to get a CRC code from a 32-bit data word and a fixed generator polynomial.

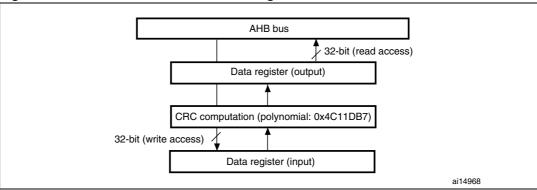
Among other applications, CRC-based techniques are used to verify data transmission or storage integrity. In the scope of the EN/IEC 60335-1 standard, they offer a means of verifying the Flash memory integrity. The CRC calculation unit helps compute a signature of the software during runtime, to be compared with a reference signature generated at link-time and stored at a given memory location.

### 3.2 CRC main features

- Uses CRC-32 (Ethernet) polynomial: 0x4C11DB7
   X<sup>32</sup> + X<sup>26</sup> + X<sup>23</sup> + X<sup>22</sup> + X<sup>16</sup> + X<sup>12</sup> + X<sup>11</sup> + X<sup>11</sup> + X<sup>10</sup> + X<sup>8</sup> + X<sup>7</sup> + X<sup>5</sup> + X<sup>4</sup> + X<sup>2</sup> + X + 1
- Single input/output 32-bit data register
- CRC computation done in 4 AHB clock cycles (HCLK)
- General-purpose 8-bit register (can be used for temporary storage)

The block diagram is shown in Figure 2.

Figure 2. CRC calculation unit block diagram



# 3.3 CRC functional description

The CRC calculation unit mainly consists of a single 32-bit data register, which:

- is used as an input register to enter new data in the CRC calculator (when writing into the register)
- holds the result of the previous CRC calculation (when reading the register)

CRC calculation unit

Each write operation into the data register creates a combination of the previous CRC value and the new one (CRC computation is done on the whole 32-bit data word, and not byte per byte).

The write operation is stalled until the end of the CRC computation, thus allowing back-to-back write accesses or consecutive write and read accesses.

The CRC calculator can be reset to 0xFFFF FFFF with the RESET control bit in the CRC\_CR register. This operation does not affect the contents of the CRC\_IDR register.

# 3.4 CRC registers

The CRC calculation unit contains two data registers and a control register. The peripheral The CRC registers have to be accessed by words (32 bits).

# 3.4.1 Data register (CRC\_DR)

Address offset: 0x00

Reset value: 0xFFFF FFFF

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							DR [3	31:16]							
rw	rw	rw	rw	rw	rw	rw	rw	rw							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							DR [	15:0]							
rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bits 31:0 Data register bits

Used as an input register when writing new data into the CRC calculator. Holds the previous CRC calculation result when it is read.

### 3.4.2 Independent data register (CRC\_IDR)

Address offset: 0x04

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			Rese	ruad							IDR	[7:0]			
			nese	erved				rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:8 Reserved, must be kept at reset value.

RM0090 CRC calculation unit

#### Bits 7:0 General-purpose 8-bit data register bits

Can be used as a temporary storage location for one byte.

This register is not affected by CRC resets generated by the RESET bit in the CRC\_CR register.

# 3.4.3 Control register (CRC\_CR)

Address offset: 0x08

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							_								RESET
							Reserved	i							w

Bits 31:1 Reserved, must be kept at reset value.

#### Bit 0 RESET bit

Resets the CRC calculation unit and sets the data register to 0xFFFF FFFF.

This bit can only be set, it is automatically cleared by hardware.

## 3.4.4 CRC register map

The following table provides the CRC register map and reset values.

Table 7. CRC calculation unit register map and reset values

Offset	Register	31-24	23-16	15-8	7	6	5	4	3	2	1	0
0x00	CRC_DR Reset value						registe					
0x04	CRC_IDR Reset value		Reserved	ı	Independent data register 0x00							
0x08	CRC_CR Reset value				Reserved							

Power control (PWR) RM0090

# 4 Power control (PWR)

This section applies to the whole STM32F4xx family, unless otherwise specified.

# 4.1 Power supplies

The device requires a 1.8-to-3.6 V operating voltage supply (V<sub>DD</sub>). An embedded linear voltage regulator is used to supply the internal 1.2 V digital power.

The real-time clock (RTC), the RTC backup registers, and the backup SRAM (BKP SRAM) can be powered from the  $V_{BAT}$  voltage when the main  $V_{DD}$  supply is powered off.

Note:

Depending on the operating power supply range, some peripheral may be used with limited functionality and performance. For more details refer to section "General operating conditions" in STM32F4xx datasheets.

Power control (PWR)

 $V_{BAT}$ Backup circuitry Power switch (OSC32K,RTC, VBAT = Wakeup logic 1.65 to 3.6V Backup registers, backup RĂM) ΟUT shifter Ю GP I/Os Logic evel Kernel logic (CPU, V<sub>CAP\_1</sub> digital  $2 \times 2.2 \mu F$ & RAM)  $V_{DD}$ V<sub>DD</sub> 1/2/...14/15 Voltage regulator  $15 \times 100 \text{ nF}$  $+ 1 \times 4.7 \mu F$ Flash memory BYPASS\_REG Reset PDR\_ON controller  $V_{DDA}$ V<sub>REF+</sub> Analog: 100 nF 100 nF **ADC** V<sub>REF</sub>-RCs, PLL  $+1 \mu F$  $V_{SSA}$ MS19911V2

Figure 3. Power supply overview

1.  $V_{DDA}$  and  $V_{SSA}$  must be connected to  $V_{DD}$  and  $V_{SS}$ , respectively.

### 4.1.1 Independent A/D converter supply and reference voltage

To improve conversion accuracy, the ADC has an independent power supply which can be separately filtered and shielded from noise on the PCB.

- The ADC voltage supply input is available on a separate V<sub>DDA</sub> pin.
- An isolated supply ground connection is provided on pin V<sub>SSA</sub>.

To ensure a better accuracy of low voltage inputs, the user can connect a separate external reference voltage ADC input on  $V_{REF}$  The voltage on  $V_{REF}$  ranges from 1.8 V to  $V_{DDA}$ .

Power control (PWR) RM0090

### 4.1.2 Battery backup domain

### **Backup domain description**

To retain the content of the RTC backup registers, backup SRAM, and supply the RTC when  $V_{DD}$  is turned off,  $V_{BAT}$  pin can be connected to an optional standby voltage supplied by a battery or by another source.

To allow the RTC to operate even when the main digital supply  $(V_{DD})$  is turned off, the  $V_{BAT}$  pin powers the following blocks:

- The RTC
- The LSE oscillator
- The backup SRAM when the low power backup regulator is enabled
- PC13 to PC15 I/Os, plus PI8 I/O (when available)

The switch to the V<sub>BAT</sub> supply is controlled by the power-down reset embedded in the Reset block.

#### Warning:

During  $t_{RSTTEMPO}$  (temporization at  $V_{DD}$  startup) or after a PDR is detected, the power switch between  $V_{BAT}$  and  $V_{DD}$  remains connected to  $V_{BAT}$ .

During the startup phase, if  $V_{DD}$  is established in less than  $t_{RSTTEMPO}$  (Refer to the datasheet for the value of  $t_{RSTTEMPO}$ ) and  $V_{DD} > V_{BAT} + 0.6$  V, a current may be injected into  $V_{BAT}$  through an internal diode connected between  $V_{DD}$  and the power switch  $(V_{BAT})$ .

If the power supply/battery connected to the  $V_{BAT}$  pin cannot support this current injection, it is strongly recommended to connect an external low-drop diode between this power supply and the  $V_{BAT}$  pin.

If no external battery is used in the application, it is recommended to connect the  $V_{BAT}$  pin to  $V_{DD}$  with a 100 nF external decoupling ceramic capacitor in parallel.

When the backup domain is supplied by  $V_{DD}$  (analog switch connected to  $V_{DD}$ ), the following functions are available:

- PC14 and PC15 can be used as either GPIO or LSE pins
- PC13 can be used as a GPIO or as the RTC\_AF1 pin (refer to *Table 18: RTC\_AF1 pin* for more details about this pin configuration)

Note:

Due to the fact that the switch only sinks a limited amount of current (3 mA), the use of GPIOs PI8 and PC13 to PC15 are restricted: only one I/O at a time can be used as an output, the speed has to be limited to 2 MHz with a maximum load of 30 pF and these I/Os must not be used as a current source (e.g. to drive an LED).

When the backup domain is supplied by  $V_{BAT}$  (analog switch connected to  $V_{BAT}$  because  $V_{DD}$  is not present), the following functions are available:

- PC14 and PC15 can be used as LSE pins only
- PC13 can be used as the RTC\_AF1 pin (refer to Table 18: RTC\_AF1 pin for more details about this pin configuration)
- PI8 can be used as RTC AF2

### **Backup domain access**

After reset, the backup domain (RTC registers, RTC backup register and backup SRAM) is protected against possible unwanted write accesses. To enable access to the backup domain, proceed as follows:

- Access to the RTC and RTC backup registers
- Enable the power interface clock by setting the PWREN bits in the RCC APB1
   peripheral clock enable register (RCC\_APB1ENR)
- 2. Set the DBP bit in the *PWR power control register (PWR\_CR)* to enable access to the backup domain
- 3. Select the RTC clock source: see Section 5.2.8: RTC/AWU clock
- 4. Enable the RTC clock by programming the RTCEN [15] bit in the RCC Backup domain control register (RCC\_BDCR)
- Access to the backup SRAM
- Enable the power interface clock by setting the PWREN bits in the RCC APB1
   peripheral clock enable register (RCC\_APB1ENR)
- 2. Set the DBP bit in the *PWR power control register (PWR\_CR)* to enable access to the backup domain
- 3. Enable the backup SRAM clock by setting BKPSRAMEN bit in the *RCC AHB1* peripheral clock register (*RCC\_AHB1ENR*).

#### RTC and RTC backup registers

The real-time clock (RTC) is an independent BCD timer/counter. The RTC provides a time-of-day clock/calendar, two programmable alarm interrupts, and a periodic programmable wakeup flag with interrupt capability. The RTC contains 20 backup data registers (80 bytes) which are reset when a tamper detection event occurs. For more details refer to Section 22: Real-time clock (RTC).

#### **Backup SRAM**

The backup domain includes 4 Kbytes of backup SRAM accessible only from the CPU, and address in 32-bit, 16-bit or 8-bit mode. Its content is retained even in Standby or  $V_{BAT}$  mode when the low power backup regulator is enabled. It can be considered as an internal EEPROM when  $V_{BAT}$  is always present.

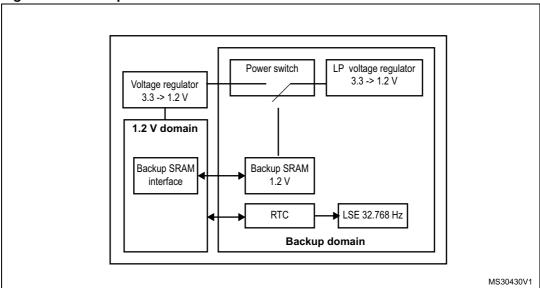
When the backup domain is supplied by  $V_{DD}$  (analog switch connected to  $V_{DD}$ ), the backup SRAM is powered from  $V_{DD}$  which replaces the  $V_{BAT}$  power supply to save battery life.

When the backup domain is supplied by  $V_{BAT}$  (analog switch connected to  $V_{BAT}$  because  $V_{DD}$  is not present), the backup SRAM is powered by a dedicated low power regulator. This regulator can be ON or OFF depending whether the application needs the backup SRAM function in Standby and  $V_{BAT}$  modes or not. The power-down of this regulator is controlled by a dedicated bit, the BRE control bit of the PWR\_CSR register (see Section 4.4.2: PWR power control/status register (PWR\_CSR)).

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The backup SRAM is not mass erased by an tamper event. It is read protected to prevent confidential data, such as cryptographic private key, from being accessed. The backup SRAM can be erased only through the Flash interface when a protection level change from level 1 to level 0 is requested. Refer to the description of Read protection (RDP) in the Flash programming manual.





### 4.1.3 Voltage regulator

An embedded linear voltage regulator supplies all the digital circuitries except for the backup domain and the Standby circuitry. The regulator output voltage is around 1.2 V.

This voltage regulator requires two external capacitors to be connected to two dedicated pins,  $V_{CAP\_1}$  and  $V_{CAP\_2}$  available in all packages. Specific pins must be connected either to  $V_{SS}$  or  $V_{DD}$  to activate or deactivate the voltage regulator. These pins depend on the package.

When activated by software, the voltage regulator is always enabled after Reset. It works in three different modes depending on the application modes.

- In **Run mode**, the regulator supplies full power to the 1.2 V domain (core, memories and digital peripherals). In this mode, the regulator output voltage (around 1.2 V) can be scaled by software to different voltage values (scale 1 or scale 2 configured through the VOS bit of the PWR\_CR register). The voltage scaling allows optimizing the power consumption when the device is clocked below the maximum system frequency (see Section 4.4.1: PWR power control register (PWR\_CR).
- In **Stop mode** the regulator supplies low power to the 1.2 V domain, preserving the content of registers and internal SRAM.
- In **Standby mode**, the regulator is powered down. The content of the registers and SRAM are lost except for the Standby circuitry and the backup domain.

Note: For more details, refer to the voltage regulator section in the STM32F40x and STM32F41x datasheets.

Note:

# 4.2 Power supply supervisor

## 4.2.1 Power-on reset (POR)/power-down reset (PDR)

The device has an integrated POR/PDR circuitry that allows proper operation starting from to 1.8 V.

The device remains in Reset mode when  $V_{DD}/V_{DDA}$  is below a specified threshold,  $V_{POR/PDR}$ , without the need for an external reset circuit. For more details concerning the power on/power-down reset threshold, refer to the electrical characteristics of the datasheet.

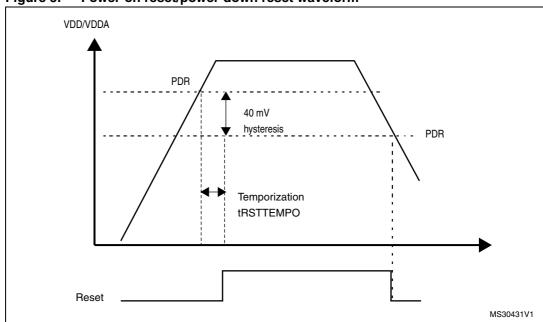


Figure 5. Power-on reset/power-down reset waveform

# 4.2.2 Brownout reset (BOR)

During power on, the Brownout reset (BOR) keeps the device under reset until the supply voltage reaches the specified  $V_{BOR}$  threshold.

 $V_{BOR}$  is configured through device option bytes. By default, BOR is off. 4 programmable  $V_{BOR}$  thresholds can be selected.

- BOR off (V<sub>BOR0</sub>): reset threshold level for 1.8 to 2.10 V voltage range
- BOR Level 1 (V<sub>BOR1</sub>): reset threshold level for 2.10 to 2.40 V voltage range
- BOR Level 2 (V<sub>BOR2</sub>): reset threshold level for 2.40 to 2.70 V voltage range
- BOR Level 3 (V<sub>BOR3</sub>): reset threshold level for 2.70 to 3.60 V voltage range

When the supply voltage  $(V_{DD})$  drops below the selected  $V_{BOR}$  threshold, a device reset is generated.

BOR can be disabled by programming the device option bytes. To disable the BOR function,  $V_{DD}$  must have been higher than  $V_{BOR0}$  to start the device option byte programming sequence. The power-down is then monitored by the PDR (see *Section 4.2.1: Power-on reset (POR)/power-down reset (PDR)*)

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The BOR threshold hysteresis is ~100 mV (between the rising and the falling edge of the supply voltage).

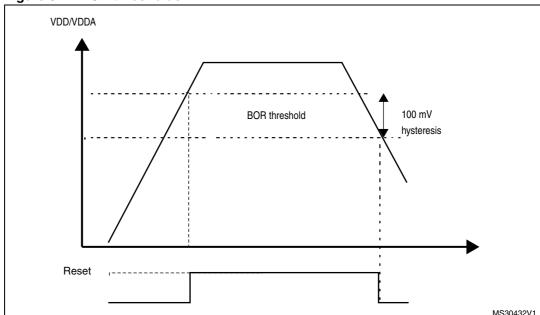


Figure 6. BOR thresholds

### 4.2.3 Programmable voltage detector (PVD)

You can use the PVD to monitor the  $V_{DD}$  power supply by comparing it to a threshold selected by the PLS[2:0] bits in the *PWR power control register (PWR\_CR)*.

The PVD is enabled by setting the PVDE bit.

A PVDO flag is available, in the *PWR power control/status register (PWR\_CSR)*, to indicate if  $V_{DD}$  is higher or lower than the PVD threshold. This event is internally connected to the EXTI line16 and can generate an interrupt if enabled through the EXTI registers. The PVD output interrupt can be generated when  $V_{DD}$  drops below the PVD threshold and/or when  $V_{DD}$  rises above the PVD threshold depending on EXTI line16 rising/falling edge configuration. As an example the service routine could perform emergency shutdown tasks.

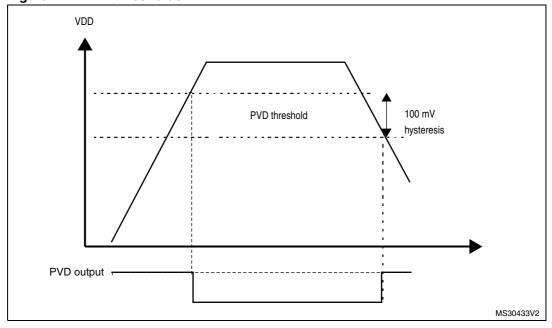


Figure 7. PVD thresholds

# 4.3 Low-power modes

By default, the microcontroller is in Run mode after a system or a power-on reset. In Run mode the CPU is clocked by HCLK and the program code is executed. Several low-power modes are available to save power when the CPU does not need to be kept running, for example when waiting for an external event. It is up to the user to select the mode that gives the best compromise between low-power consumption, short startup time and available wakeup sources.

The devices feature three low-power modes:

- Sleep mode (Cortex<sup>™</sup>-M4F core stopped, peripherals kept running)
- Stop mode (all clocks are stopped)
- Standby mode (1.2 V domain powered off)

In addition, the power consumption in Run mode can be reduce by one of the following means:

- Slowing down the system clocks
- Gating the clocks to the APBx and AHBx peripherals when they are unused.

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Table 8. Low-power mode summary

Mode name	Entry	Wakeup	Effect on 1.2 V domain clocks	Effect on V <sub>DD</sub> domain clocks	Voltage regulator
Sleep	WFI	Any interrupt	CPU CLK OFF		
(Sleep now or Sleep-on-exit)	WFE	Wakeup event	clocks or analog	None	ON
Stop	PDDS and LPDS bits + SLEEPDEEP bit + WFI or WFE	Any EXTI line (configured in the EXTI registers, internal and external lines)	All 1.2 V domain clocks OFF	HSI and HSE oscillators OFF	ON or in low- power mode (depends on PWR power control register (PWR_CR))
Standby	PDDS bit + SLEEPDEEP bit + WFI or WFE	WKUP pin rising edge, RTC alarm (Alarm A or Alarm B), RTC Wakeup event, RTC tamper events, RTC time stamp event, external reset in NRST pin, IWDG reset	All 1.2 V domain clocks OFF	HSI and HSE oscillators OFF	OFF

### 4.3.1 Slowing down system clocks

In Run mode the speed of the system clocks (SYSCLK, HCLK, PCLK1, PCLK2) can be reduced by programming the prescaler registers. These prescalers can also be used to slow down peripherals before entering Sleep mode.

For more details refer to Section 5.3.3: RCC clock configuration register (RCC\_CFGR).

### 4.3.2 Peripheral clock gating

In Run mode, the HCLKx and PCLKx for individual peripherals and memories can be stopped at any time to reduce power consumption.

To further reduce power consumption in Sleep mode the peripheral clocks can be disabled prior to executing the WFI or WFE instructions.

Peripheral clock gating is controlled by the AHB1 peripheral clock enable register (RCC\_AHB1ENR), AHB2 peripheral clock enable register (RCC\_AHB2ENR), AHB3 peripheral clock enable register (RCC\_AHB3ENR) (see RCC APB1 peripheral clock enable register (RCC\_APB1ENR) and RCC APB2 peripheral clock enable register (RCC\_APB2ENR)).

Disabling the peripherals clocks in Sleep mode can be performed automatically by resetting the corresponding bit in RCC AHBxLPENR and RCC APBxLPENR registers.

### 4.3.3 Sleep mode

### **Entering Sleep mode**

The Sleep mode is entered by executing the WFI (Wait For Interrupt) or WFE (Wait for Event) instructions. Two options are available to select the Sleep mode entry mechanism, depending on the SLEEPONEXIT bit in the Cortex<sup>TM</sup>-M4F System Control register:

- Sleep-now: if the SLEEPONEXIT bit is cleared, the MCU enters Sleep mode as soon as WFI or WFE instruction is executed.
- Sleep-on-exit: if the SLEEPONEXIT bit is set, the MCU enters Sleep mode as soon as it exits the lowest priority ISR.

Refer to Table 9 and Table 10 for details on how to enter Sleep mode.

### **Exiting Sleep mode**

If the WFI instruction is used to enter Sleep mode, any peripheral interrupt acknowledged by the nested vectored interrupt controller (NVIC) can wake up the device from Sleep mode.

If the WFE instruction is used to enter Sleep mode, the MCU exits Sleep mode as soon as an event occurs. The wakeup event can be generated either by:

- Enabling an interrupt in the peripheral control register but not in the NVIC, and enabling the SEVONPEND bit in the Cortex™-M4F System Control register. When the MCU resumes from WFE, the peripheral interrupt pending bit and the peripheral NVIC IRQ channel pending bit (in the NVIC interrupt clear pending register) have to be cleared.
- Or configuring an external or internal EXTI line in event mode. When the CPU resumes
  from WFE, it is not necessary to clear the peripheral interrupt pending bit or the NVIC
  IRQ channel pending bit as the pending bit corresponding to the event line is not set.

This mode offers the lowest wakeup time as no time is wasted in interrupt entry/exit.

Refer to Table 9 and Table 10 for more details on how to exit Sleep mode.

Table 9. Sleep-now entry and exit

Sleep-now mode	Description
Mode entry	WFI (Wait for Interrupt) or WFE (Wait for Event) while:  - SLEEPDEEP = 0 and  - SLEEPONEXIT = 0  Refer to the Cortex™-M4F System Control register.
Mode exit	If WFI was used for entry: Interrupt: Refer to <i>Table 32: Vector table</i> If WFE was used for entry Wakeup event: Refer to <i>Section 9.2.3: Wakeup event management</i>
Wakeup latency	None

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Table 10. Sleep-on-exit entry and exit

Sleep-on-exit	Description
Mode entry	WFI (wait for interrupt) while:  - SLEEPDEEP = 0 and  - SLEEPONEXIT = 1  Refer to the Cortex™-M4F System Control register.
Mode exit	Interrupt: refer to Table 32: Vector table.
Wakeup latency	None

# 4.3.4 Stop mode

The Stop mode is based on the Cortex<sup>™</sup>-M4F deepsleep mode combined with peripheral clock gating. The voltage regulator can be configured either in normal or low-power mode. In Stop mode, all clocks in the 1.2 V domain are stopped, the PLLs, the HSI and the HSE RC oscillators are disabled. Internal SRAM and register contents are preserved.

By setting the FPDS bit in the PWR\_CR register, the Flash memory also enters power-down mode when the device enters Stop mode. When the Flash memory is in power-down mode, an additional startup delay is incurred when waking up from Stop mode (see *Table 11: Stop operating modes*).

Table 11. Stop operating modes

Stop mode	LPDS bit	FPDS bit	Wake-up latency
STOP MR ( Main regulator)	0	0	HSI RC startup time
STOP MR-FPD	0	1	HSI RC startup time + Flash wakeup time from Power Down mode
STOP LP	1	0	HSI RC startup time + regulator wakeup time from LP mode
STOP LP-FPD	1	1	HSI RC startup time + Flash wakeup time from Power Down mode + regulator wakeup time from LP mode

# **Entering Stop mode**

Refer to *Table 12* for details on how to enter the Stop mode.

To further reduce power consumption in Stop mode, the internal voltage regulator can be put in low power mode. This is configured by the LPDS bit of the *PWR power control register (PWR\_CR)*.

If Flash memory programming is ongoing, the Stop mode entry is delayed until the memory access is finished.

If an access to the APB domain is ongoing, The Stop mode entry is delayed until the APB access is finished.

In Stop mode, the following features can be selected by programming individual control bits:

- Independent watchdog (IWDG): the IWDG is started by writing to its Key register or by hardware option. Once started it cannot be stopped except by a Reset. See Section 17.3 in Section 17: Independent watchdog (IWDG).
- Real-time clock (RTC): this is configured by the RTCEN bit in the RCC Backup domain control register (RCC\_BDCR)
- Internal RC oscillator (LSI RC): this is configured by the LSION bit in the RCC clock control & status register (RCC\_CSR).
- External 32.768 kHz oscillator (LSE OSC): this is configured by the LSEON bit in the *RCC Backup domain control register (RCC\_BDCR)*.

The ADC or DAC can also consume power during the Stop mode, unless they are disabled before entering it. To disable them, the ADON bit in the ADC\_CR2 register and the ENx bit in the DAC\_CR register must both be written to 0.

### **Exiting Stop mode**

Refer to Table 12 for more details on how to exit Stop mode.

When exiting Stop mode by issuing an interrupt or a wakeup event, the HSI RC oscillator is selected as system clock.

When the voltage regulator operates in low power mode, an additional startup delay is incurred when waking up from Stop mode. By keeping the internal regulator ON during Stop mode, the consumption is higher although the startup time is reduced.

Table 12. Stop mode entry and ex	Table 12.	Stop mode entry and	exit
----------------------------------	-----------	---------------------	------

Stop mode	Description
Mode entry	WFI (Wait for Interrupt) or WFE (Wait for Event) while:  - Set SLEEPDEEP bit in Cortex™-M4F System Control register  - Clear PDDS bit in Power Control register (PWR_CR)  - Select the voltage regulator mode by configuring LPDS bit in PWR_CR.  Note: To enter the Stop mode, all EXTI Line pending bits (in Pending register (EXTI_PR)), the RTC Alarm (Alarm A and Alarm B), RTC wakeup, RTC tamper, and RTC time stamp flags, must be reset. Otherwise, the Stop
	mode entry procedure is ignored and program execution continues.
Mode exit	If WFI was used for entry:  All EXTI lines configured in Interrupt mode (the corresponding EXTI Interrupt vector must be enabled in the NVIC). Refer to <i>Table 32: Vector table on page 199</i> .  If WFE was used for entry:  All EXTI Lines configured in event mode. Refer to <i>Section 9.2.3: Wakeup event management on page 204</i>
Wakeup latency	Table 11: Stop operating modes

# 4.3.5 Standby mode

The Standby mode allows to achieve the lowest power consumption. It is based on the Cortex<sup>™</sup>-M4F deepsleep mode, with the voltage regulator disabled. The 1.2 V domain is consequently powered off. The PLLs, the HSI oscillator and the HSE oscillator are also switched off. SRAM and register contents are lost except for registers in the backup domain

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> (RTC registers, RTC backup register and backup SRAM), and Standby circuitry (see Figure 3).

### **Entering Standby mode**

Refer to *Table 13* for more details on how to enter Standby mode.

In Standby mode, the following features can be selected by programming individual control bits:

- Independent watchdog (IWDG): the IWDG is started by writing to its Key register or by hardware option. Once started it cannot be stopped except by a reset. See Section 17.3 in Section 17: Independent watchdog (IWDG).
- Real-time clock (RTC): this is configured by the RTCEN bit in the backup domain control register (RCC\_BDCR)
- Internal RC oscillator (LSI RC): this is configured by the LSION bit in the Control/status register (RCC\_CSR).
- External 32.768 kHz oscillator (LSE OSC): this is configured by the LSEON bit in the backup domain control register (RCC\_BDCR)

# **Exiting Standby mode**

Table 13.

The microcontroller exits Standby mode when an external Reset (NRST pin), an IWDG Reset, a rising edge on WKUP pin, an RTC alarm, a tamper event, or a time stamp event is detected. All registers are reset after wakeup from Standby except for PWR power control/status register (PWR\_CSR).

After waking up from Standby mode, program execution restarts in the same way as after a Reset (boot pins sampling, vector reset is fetched, etc.). The SBF status flag in the PWR power control/status register (PWR\_CSR) indicates that the MCU was in Standby mode.

Refer to *Table 13* for more details on how to exit Standby mode.

Standby mode entry and exit

Standby mode	Description
	WFI (Wait for Interrupt) or WFE (Wait for Event) while:
	<ul> <li>Set SLEEPDEEP in Cortex<sup>™</sup>-M4F System Control register</li> </ul>

Standby mode	Description
Mode entry	WFI (Wait for Interrupt) or WFE (Wait for Event) while:  - Set SLEEPDEEP in Cortex™-M4F System Control register  - Set PDDS bit in Power Control register (PWR_CR)  - Clear WUF bit in Power Control/Status register (PWR_CSR)  - Clear the RTC flag corresponding to the chosen wakeup source (RTC Alarm A, RTC Alarm B, RTC wakeup, Tamper or Timestamp flags)
Mode exit	WKUP pin rising edge, RTC alarm (Alarm A and Alarm B), RTC wakeup, tamper event, time stamp event, external reset in NRST pin, IWDG reset.
Wakeup latency	Reset phase.

#### I/O states in Standby mode

In Standby mode, all I/O pins are high impedance except for:

- Reset pad (still available)
- RTC\_AF1 pin (PC13) if configured for tamper, time stamp, RTC Alarm out, or RTC clock calibration out
- WKUP pin (PA0), if enabled

# **Debug mode**

By default, the debug connection is lost if the application puts the MCU in Stop or Standby mode while the debug features are used. This is due to the fact that the Cortex<sup>™</sup>-M4F core is no longer clocked.

However, by setting some configuration bits in the DBGMCU\_CR register, the software can be debugged even when using the low-power modes extensively. For more details, refer to *Section 32.16.1: Debug support for low-power modes*.

# 4.3.6 Programming the RTC alternate functions to wake up the device from the Stop and Standby modes

The MCU can be woken up from a low-power mode by an RTC alternate function.

The RTC alternate functions are the RTC alarms (Alarm A and Alarm B), RTC wakeup, RTC tamper event detection and RTC time stamp event detection.

These RTC alternate functions can wake up the system from the Stop and Standby low-power modes.

The system can also wake up from low-power modes without depending on an external interrupt (Auto-wakeup mode), by using the RTC alarm or the RTC wakeup events.

The RTC provides a programmable time base for waking up from the Stop or Standby mode at regular intervals.

For this purpose, two of the three alternate RTC clock sources can be selected by programming the RTCSEL[1:0] bits in the *RCC Backup domain control register* (*RCC\_BDCR*):

- Low-power 32.768 kHz external crystal oscillator (LSE OSC)
   This clock source provides a precise time base with a very low-power consumption (additional consumption of less than 1 μA under typical conditions)
- Low-power internal RC oscillator (LSI RC)
   This clock source has the advantage of saving the cost of the 32.768 kHz crystal. This internal RC oscillator is designed to use minimum power.

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# RTC alternate functions to wake up the device from the Stop mode

- To wake up the device from the Stop mode with an RTC alarm event, it is necessary to:
  - a) Configure the EXTI Line 17 to be sensitive to rising edges (Interrupt or Event modes)
  - b) Enable the RTC Alarm Interrupt in the RTC\_CR register
  - c) Configure the RTC to generate the RTC alarm
- To wake up the device from the Stop mode with an RTC tamper or time stamp event, it is necessary to:
  - a) Configure the EXTI Line 21 to be sensitive to rising edges (Interrupt or Event modes)
  - b) Enable the RTC time stamp Interrupt in the RTC\_CR register or the RTC tamper interrupt in the RTC\_TAFCR register
  - c) Configure the RTC to detect the tamper or time stamp event
- To wake up the device from the Stop mode with an RTC wakeup event, it is necessary to:
  - a) Configure the EXTI Line 22 to be sensitive to rising edges (Interrupt or Event modes)
  - b) Enable the RTC wakeup interrupt in the RTC CR register
  - c) Configure the RTC to generate the RTC Wakeup event

#### RTC alternate functions to wake up the device from the Standby mode

- To wake up the device from the Standby mode with an RTC alarm event, it is necessary to:
  - a) Enable the RTC alarm interrupt in the RTC CR register
  - b) Configure the RTC to generate the RTC alarm
- To wake up the device from the Standby mode with an RTC tamper or time stamp event, it is necessary to:
  - Enable the RTC time stamp interrupt in the RTC\_CR register or the RTC tamper interrupt in the RTC\_TAFCR register
  - b) Configure the RTC to detect the tamper or time stamp event
- To wake up the device from the Standby mode with an RTC wakeup event, it is necessary to:
  - a) Enable the RTC wakeup interrupt in the RTC\_CR register
  - b) Configure the RTC to generate the RTC wakeup event

# Safe RTC alternate function wakeup flag clearing sequence

If the selected RTC alternate function is set before the PWR wakeup flag (WUTF) is cleared, it will not be detected on the next event as detection is made once on the rising edge.

To avoid bouncing on the pins onto which the RTC alternate functions are mapped, and exit correctly from the Stop and Standby modes, it is recommended to follow the sequence below before entering the Standby mode:

- When using RTC alarm to wake up the device from the low-power modes:
  - a) Disable the RTC alarm interrupt (ALRAIE or ALRBIE bits in the RTC\_CR register)
  - b) Clear the RTC alarm (ALRAF/ALRBF) flag
  - c) Clear the PWR Wakeup (WUF) flag
  - d) Enable the RTC alarm interrupt
  - e) Re-enter the low-power mode
- When using RTC wakeup to wake up the device from the low-power modes:
  - a) Disable the RTC Wakeup interrupt (WUTIE bit in the RTC\_CR register)
  - b) Clear the RTC Wakeup (WUTF) flag
  - c) Clear the PWR Wakeup (WUF) flag
  - d) Enable the RTC Wakeup interrupt
  - e) Re-enter the low power mode
- When using RTC tamper to wake up the device from the low-power modes:
  - a) Disable the RTC tamper interrupt (TAMPIE bit in the RTC\_TAFCR register)
  - b) Clear the Tamper (TAMP1F/TSF) flag
  - c) Clear the PWR Wakeup (WUF) flag
  - d) Enable the RTC tamper interrupt
  - e) Re-enter the low-power mode
- When using RTC time stamp to wake up the device from the low-power modes:
  - a) Disable the RTC time stamp interrupt (TSIE bit in RTC\_CR)
  - b) Clear the RTC time stamp (TSF) flag
  - c) Clear the PWR Wakeup (WUF) flag
  - d) Enable the RTC TimeStamp interrupt
  - e) Re-enter the low-power mode

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# 4.4 Power control registers

# 4.4.1 PWR power control register (PWR CR)

Address offset: 0x00

Reset value: 0x0000 4000 (reset by wakeup from Standby mode)

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Res.	vos		Rese	erved		FPDS	DBP		PLS[2:0]		PVDE	CSBF	CWUF	PDDS	LPDS
	rw					rw	rw	rw	rw	rw	rw	rc_w1	rc_w1	rw	rw

Bits 31:15 Reserved, must be kept at reset value.

#### Bit 14 VOS: Regulator voltage scaling output selection

This bit controls the main internal voltage regulator output voltage to achieve a trade-off between performance and power consumption when the device does not operate at the maximum frequency (refer to the STM32F40x and STM32F41x datasheets for more details).

- 0: Scale 2 mode
- 1: Scale 1 mode (default value at reset)

Bits 13:10 Reserved, must be kept at reset value.

#### Bit 9 FPDS: Flash power-down in Stop mode

When set, the Flash memory enters power-down mode when the device enters Stop mode. This allows to achieve a lower consumption in stop mode but a longer restart time.

- 0: Flash memory not in power-down when the device is in Stop mode
- 1: Flash memory in power-down when the device is in Stop mode

#### Bit 8 DBP: Disable backup domain write protection

In reset state, the RCC\_BDCR register, the RTC registers (including the backup registers), and the BRE bit of the PWR\_CSR register, are protected against parasitic write access. This bit must be set to enable write access to these registers.

- 0: Access to RTC and RTC Backup registers and backup SRAM disabled
- 1: Access to RTC and RTC Backup registers and backup SRAM enabled

#### Bits 7:5 PLS[2:0]: PVD level selection

These bits are written by software to select the voltage threshold detected by the Power Voltage Detector

000: 2.0 V

001: 2.1 V

010: 2.3 V

011: 2.5 V

011. Z.5 V

100: 2.6 V

101: 2.7 V 110: 2.8 V

111: 2.9 V

Note: Refer to the electrical characteristics of the datasheet for more details.

Bit 4 PVDE: Power voltage detector enable

This bit is set and cleared by software.

0: PVD disabled

1: PVD enabled

Bit 3 CSBF: Clear standby flag

This bit is always read as 0.

0: No effect

1: Clear the SBF Standby Flag (write).

Bit 2 CWUF: Clear wakeup flag

This bit is always read as 0.

0: No effect

1: Clear the WUF Wakeup Flag after 2 System clock cycles

Bit 1 PDDS: Power-down deepsleep

This bit is set and cleared by software. It works together with the LPDS bit.

0: Enter Stop mode when the CPU enters deepsleep. The regulator status depends on the LPDS bit.

1: Enter Standby mode when the CPU enters deepsleep.

Bit 0 LPDS: Low-power deepsleep

This bit is set and cleared by software. It works together with the PDDS bit.

0: Voltage regulator on during Stop mode

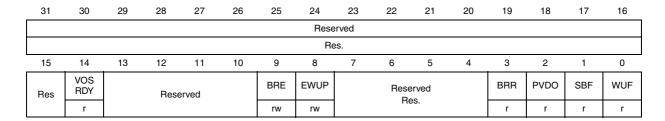
1: Voltage regulator in low-power mode during Stop mode

# 4.4.2 PWR power control/status register (PWR\_CSR)

Address offset: 0x04

Reset value: 0x0000 0000 (not reset by wakeup from Standby mode)

Additional APB cycles are needed to read this register versus a standard APB read.



Bits 31:15 Reserved, must be kept at reset value.

Bit 14 VOSRDY: Regulator voltage scaling output selection ready bit

0: Not ready 1: Ready Power control (PWR) RM0090

#### Bits 13:10 Reserved, must be kept at reset value.

#### Bit 9 BRE: Backup regulator enable

When set, the Backup regulator (used to maintain backup SRAM content in Standby and  $V_{BAT}$  modes) is enabled. If BRE is reset, the backup regulator is switched off. The backup SRAM can still be used but its content will be lost in the Standby and  $V_{BAT}$  modes. Once set, the application must wait that the Backup Regulator Ready flag (BRR) is set to indicate that the data written into the RAM will be maintained in the Standby and  $V_{BAT}$  modes.

- 0: Backup regulator disabled
- 1: Backup regulator enabled

Note: This bit is not reset when the device wakes up from Standby mode, by a system reset, or by a power reset.

#### Bit 8 **EWUP**: Enable WKUP pin

This bit is set and cleared by software.

- 0: WKUP pin is used for general purpose I/O. An event on the WKUP pin does not wakeup the device from Standby mode.
- 1: WKUP pin is used for wakeup from Standby mode and forced in input pull down configuration (rising edge on WKUP pin wakes-up the system from Standby mode).

Note: This bit is reset by a system reset.

#### Bits 7:4 Reserved, must be kept at reset value.

#### Bit 3 BRR: Backup regulator ready

Set by hardware to indicate that the Backup Regulator is ready.

- 0: Backup Regulator not ready
- 1: Backup Regulator ready

Note: This bit is not reset when the device wakes up from Standby mode or by a system reset or power reset.

#### Bit 2 PVDO: PVD output

This bit is set and cleared by hardware. It is valid only if PVD is enabled by the PVDE bit.

- 0:  $V_{DD}$  is higher than the PVD threshold selected with the PLS[2:0] bits.
- 1:  $V_{DD}$  is lower than the PVD threshold selected with the PLS[2:0] bits.

Note: The PVD is stopped by Standby mode. For this reason, this bit is equal to 0 after Standby or reset until the PVDE bit is set.

#### Bit 1 SBF: Standby flag

This bit is set by hardware and cleared only by a POR/PDR (power-on reset/power-down reset) or by setting the CSBF bit in the *PWR power control register (PWR\_CR)* 

- 0: Device has not been in Standby mode
- 1: Device has been in Standby mode

#### Bit 0 WUF: Wakeup flag

This bit is set by hardware and cleared only by a POR/PDR (power-on reset/power-down reset) or by setting the CWUF bit in the *PWR power control register (PWR\_CR)* 

- 0: No wakeup event occurred
- 1: A wakeup event was received from the WKUP pin or from the RTC alarm (Alarm A or Alarm B), RTC Tamper event, RTC TimeStamp event or RTC Wakeup).

Note: An additional wakeup event is detected if the WKUP pin is enabled (by setting the EWUP bit) when the WKUP pin level is already high.

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# 4.4.3 PWR register map

The following table summarizes the PWR registers.

Refer to *Table 2 on page 50* for the register boundary addresses.

Table 14. PWR - register map and reset values for

Offset	Register	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	6	8	7	9	2	4	3	2	1	0
0x000	PWR_CR								Res	serv	ed								NOS	R	ese	rve	d	FPDS	DBP	PL	.S[2	:0]	PVDE	CSBF	CWUF	PDDS	LPDS
	Reset value																		1					0	0	0	0	0	0		0	0	0
0x004	PWR_CSR								Res	serv	ed								VOSRDY	R	ese	rve	d	BRE	EWUP	F	Rese	erve	d	BRR	PVDO	SBF	WUF
	Reset value																		0					0	0					0	0	0	0

# 5 Reset and clock control for (RCC)

# 5.1 Reset

There are three types of reset, defined as system Reset, power Reset and backup domain Reset.

# 5.1.1 System reset

A system reset sets all registers to their reset values except the reset flags in the clock controller CSR register and the registers in the Backup domain (see *Figure 4*).

A system reset is generated when one of the following events occurs:

- 1. A low level on the NRST pin (external reset)
- 2. Window watchdog end of count condition (WWDG reset)
- 3. Independent watchdog end of count condition (IWDG reset)
- 4. A software reset (SW reset) (see Software reset)
- 5. Low-power management reset (see Low-power management reset)

#### Software reset

The reset source can be identified by checking the reset flags in the *RCC clock control & status register (RCC\_CSR)*.

The SYSRESETREQ bit in Cortex<sup>™</sup>-M4F Application Interrupt and Reset Control Register must be set to force a software reset on the device. Refer to the Cortex<sup>™</sup>-M4F technical reference manual for more details.

#### Low-power management reset

There are two ways of generating a low-power management reset:

- Reset generated when entering the Standby mode:

  This type of reset is enabled by resetting the pRST\_STDRY bit in the user of the process.

  The type of reset is enabled by resetting the pRST\_STDRY bit in the user of the process.
  - This type of reset is enabled by resetting the nRST\_STDBY bit in the user option bytes. In this case, whenever a Standby mode entry sequence is successfully executed, the device is reset instead of entering the Standby mode.
- 2. Reset when entering the Stop mode:

This type of reset is enabled by resetting the nRST\_STOP bit in the user option bytes. In this case, whenever a Stop mode entry sequence is successfully executed, the device is reset instead of entering the Stop mode.

For further information on the user option bytes, refer to the STM32F40x and STM32F41x Flash programming manual available from your ST sales office.

#### 5.1.2 Power reset

A power reset is generated when one of the following events occurs:

- 1. Power-on/power-down reset (POR/PDR reset) or brownout (BOR) reset
- 2. When exiting the Standby mode

A power reset sets all registers to their reset values except the Backup domain (see Figure 4)

These sources act on the NRST pin and it is always kept low during the delay phase. The RESET service routine vector is fixed at address 0x0000\_0004 in the memory map.

The system reset signal provided to the device is output on the NRST pin. The pulse generator guarantees a minimum reset pulse duration of 20  $\mu$ s for each internal reset source. In case of an external reset, the reset pulse is generated while the NRST pin is asserted low.

External reset NRST Filter System reset

Pulse generator (min 20 µs) System reset

WWDG reset IWDG reset Power reset
Software reset Low-power management reset

Figure 8. Simplified diagram of the reset circuit

The Backup domain has two specific resets that affect only the Backup domain (see *Figure 4*).

# 5.1.3 Backup domain reset

The backup domain reset sets all RTC registers and the RCC\_BDCR register to their reset values. The BKPSRAM is not affected by this reset. The only way of resetting the BKPSRAM is through the Flash interface by requesting a protection level change from 1 to 0.

A backup domain reset is generated when one of the following events occurs:

- 1. Software reset, triggered by setting the BDRST bit in the *RCC Backup domain control register (RCC\_BDCR)*.
- 2.  $V_{DD}$  or  $V_{BAT}$  power on, if both supplies have previously been powered off.

# 5.2 Clocks

Three different clock sources can be used to drive the system clock (SYSCLK):

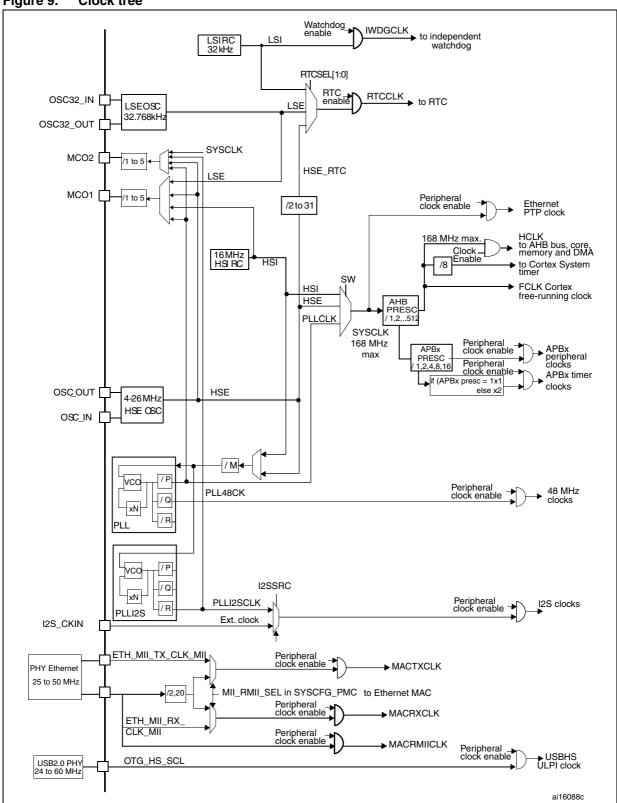
- HSI oscillator clock
- HSE oscillator clock
- Main PLL (PLL) clock

The devices have the two following secondary clock sources:

- 32 kHz low-speed internal RC (LSI RC) which drives the independent watchdog and, optionally, the RTC used for Auto-wakeup from the Stop/Standby mode.
- 32.768 kHz low-speed external crystal (LSE crystal) which optionally drives the RTC clock (RTCCLK)

Each clock source can be switched on or off independently when it is not used, to optimize power consumption.

Figure 9. **Clock tree** 



For full details about the internal and external clock source characteristics, refer to the Electrical characteristics section in the device datasheet.

The clock controller provides a high degree of flexibility to the application in the choice of the external crystal or the oscillator to run the core and peripherals at the highest frequency and, guarantee the appropriate frequency for peripherals that need a specific clock like Ethernet, USB OTG FS and HS, I2S and SDIO.

Several prescalers are used to configure the AHB frequency, the high-speed APB (APB2) and the low-speed APB (APB1) domains. The maximum frequency of the AHB domain is 168 MHz. The maximum allowed frequency of the high-speed APB2 domain is 84 MHz. The maximum allowed frequency of the low-speed APB1 domain is 42 MHz

All peripheral clocks are derived from the system clock (SYSCLK) except for:

- The USB OTG FS clock (48 MHz), the random analog generator (RNG) clock (≤ 48 MHz) and the SDIO clock (≤ 48 MHz) which are coming from a specific output of PLL (PLL48CLK)
- The I2S clock

To achieve high-quality audio performance, the I2S clock can be derived either from a specific PLL (PLLI2S) or from an external clock mapped on the I2S\_CKIN pin. For more information about I2S clock frequency and precision, refer to *Section 25.4.4: Clock generator*.

- The USB OTG HS (60 MHz) clock which is provided from the external PHY
- The Ethernet MAC clocks (TX, RX and RMII) which are provided from the external PHY. For further information on the Ethernet configuration, please refer to Section 28.4.4: MII/RMII selection in the Ethernet peripheral description. When the Ethernet is used, the AHB clock frequency must be at least 25 MHz.

The RCC feeds the external clock of the Cortex System Timer (SysTick) with the AHB clock (HCLK) divided by 8. The SysTick can work either with this clock or with the Cortex clock (HCLK), configurable in the SysTick control and status register.

The timer clock frequencies are automatically set by hardware. There are two cases:

- If the APB prescaler is 1, the timer clock frequencies are set to the same frequency as that of the APB domain to which the timers are connected.
- 2. Otherwise, they are set to twice (x2) the frequency of the APB domain to which the timers are connected.

FCLK acts as Cortex<sup>™</sup>-M4F free-running clock. For more details, refer to the Cortex<sup>™</sup>-M4F technical reference manual.

# 5.2.1 HSE clock

The high speed external clock signal (HSE) can be generated from two possible clock sources:

- HSE external crystal/ceramic resonator
- HSE external user clock

The resonator and the load capacitors have to be placed as close as possible to the oscillator pins in order to minimize output distortion and startup stabilization time. The loading capacitance values must be adjusted according to the selected oscillator.

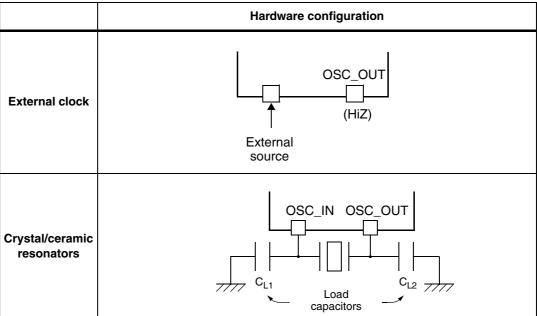


Figure 10. HSE/ LSE clock sources

# **External source (HSE bypass)**

In this mode, an external clock source must be provided. You select this mode by setting the HSEBYP and HSEON bits in the *RCC clock control register (RCC\_CR)*. The external clock signal (square, sinus or triangle) with ~50% duty cycle has to drive the OSC\_IN pin while the OSC\_OUT pin should be left hi-Z. See *Figure 10*.

# External crystal/ceramic resonator (HSE crystal)

The HSE has the advantage of producing a very accurate rate on the main clock.

The associated hardware configuration is shown in *Figure 10*. Refer to the electrical characteristics section of the *datasheet* for more details.

The HSERDY flag in the *RCC clock control register (RCC\_CR)* indicates if the high-speed external oscillator is stable or not. At startup, the clock is not released until this bit is set by hardware. An interrupt can be generated if enabled in the *RCC clock interrupt register (RCC\_CIR)*.

The HSE Crystal can be switched on and off using the HSEON bit in the *RCC clock control register (RCC\_CR)*.

#### 5.2.2 HSI clock

The HSI clock signal is generated from an internal 16 MHz RC oscillator and can be used directly as a system clock, or used as PLL input.

The HSI RC oscillator has the advantage of providing a clock source at low cost (no external components). It also has a faster startup time than the HSE crystal oscillator however, even with calibration the frequency is less accurate than an external crystal oscillator or ceramic resonator.

#### Calibration

RC oscillator frequencies can vary from one chip to another due to manufacturing process variations, this is why each device is factory calibrated by ST for 1% accuracy at  $T_A$ = 25 °C.

After reset, the factory calibration value is loaded in the HSICAL[7:0] bits in the *RCC clock control register (RCC\_CR)*.

If the application is subject to voltage or temperature variations this may affect the RC oscillator speed. You can trim the HSI frequency in the application using the HSITRIM[4:0] bits in the *RCC clock control register* (*RCC\_CR*).

The HSIRDY flag in the *RCC clock control register (RCC\_CR)* indicates if the HSI RC is stable or not. At startup, the HSI RC output clock is not released until this bit is set by hardware.

The HSI RC can be switched on and off using the HSION bit in the *RCC clock control register (RCC CR)*.

The HSI signal can also be used as a backup source (Auxiliary clock) if the HSE crystal oscillator fails. Refer to *Section 5.2.7: Clock security system (CSS) on page 90*.

# 5.2.3 PLL configuration

The STM32F4xx devices feature two PLLs:

- A main PLL (PLL) clocked by the HSE or HSI oscillator and featuring two different output clocks:
  - The first output is used to generate the high speed system clock (up to 168 MHz)
  - The second output is used to generate the clock for the USB OTG FS (48 MHz), the random analog generator (≤48 MHz) and the SDIO (≤ 48 MHz).
- A dedicated PLL (PLLI2S) used to generate an accurate clock to achieve high-quality audio performance on the I2S interface.

Since the main-PLL configuration parameters cannot be changed once PLL is enabled, it is recommended to configure PLL before enabling it (selection of the HSI or HSE oscillator as PLL clock source, and configuration of division factors M, N, P, and Q).

The PLLI2S uses the same input clock as PLL (PLLM[5:0] and PLLSRC bits are common to both PLLs). However, the PLLI2S has dedicated enable/disable and division factors (N and R) configuration bits. Once the PLLI2S is enabled, the configuration parameters cannot be changed.

The two PLLs are disabled by hardware when entering Stop and Standby modes, or when an HSE failure occurs when HSE or PLL (clocked by HSE) are used as system clock. *RCC PLL configuration register (RCC\_PLLCFGR)* and *RCC clock configuration register (RCC\_CFGR)* can be used to configure PLL and PLLI2S, respectively.

# 5.2.4 LSE clock

The LSE crystal is a 32.768 kHz low-speed external (LSE) crystal or ceramic resonator. It has the advantage providing a low-power but highly accurate clock source to the real-time clock peripheral (RTC) for clock/calendar or other timing functions.

The LSE crystal is switched on and off using the LSEON bit in *RCC Backup domain control register (RCC\_BDCR)*.

The LSERDY flag in the *RCC Backup domain control register (RCC\_BDCR)* indicates if the LSE crystal is stable or not. At startup, the LSE crystal output clock signal is not released until this bit is set by hardware. An interrupt can be generated if enabled in the *RCC clock interrupt register (RCC\_CIR)*.

#### **External source (LSE bypass)**

In this mode, an external clock source must be provided. It must have a frequency up to 1 MHz. You select this mode by setting the LSEBYP and LSEON bits in the *RCC Backup domain control register (RCC\_BDCR)*. The external clock signal (square, sinus or triangle) with ~50% duty cycle has to drive the OSC32\_IN pin while the OSC32\_OUT pin should be left Hi-Z. See *Figure 10*.

#### 5.2.5 LSI clock

The LSI RC acts as an low-power clock source that can be kept running in Stop and Standby mode for the independent watchdog (IWDG) and Auto-wakeup unit (AWU). The clock frequency is around 32 kHz. For more details, refer to the electrical characteristics section of the datasheets.

The LSI RC can be switched on and off using the LSION bit in the RCC clock control & status register (RCC\_CSR).

The LSIRDY flag in the *RCC clock control & status register (RCC\_CSR)* indicates if the low-speed internal oscillator is stable or not. At startup, the clock is not released until this bit is set by hardware. An interrupt can be generated if enabled in the *RCC clock interrupt register (RCC\_CIR)*.

# 5.2.6 System clock (SYSCLK) selection

After a system reset, the HSI oscillator is selected as the system clock. When a clock source is used directly or through PLL as the system clock, it is not possible to stop it.

A switch from one clock source to another occurs only if the target clock source is ready (clock stable after startup delay or PLL locked). If a clock source that is not yet ready is selected, the switch occurs when the clock source is ready. Status bits in the *RCC clock control register (RCC\_CR)* indicate which clock(s) is (are) ready and which clock is currently used as the system clock.

# 5.2.7 Clock security system (CSS)

The clock security system can be activated by software. In this case, the clock detector is enabled after the HSE oscillator startup delay, and disabled when this oscillator is stopped.

If a failure is detected on the HSE clock, this oscillator is automatically disabled, a clock failure event is sent to the break inputs of advanced-control timers TIM1 and TIM8, and an interrupt is generated to inform the software about the failure (clock security system interrupt CSSI), allowing the MCU to perform rescue operations. The CSSI is linked to the Cortex<sup>TM</sup>-M4F NMI (non-maskable interrupt) exception vector.

Note:

When the CSS is enabled, if the HSE clock happens to fail, the CSS generates an interrupt, which causes the automatic generation of an NMI. The NMI is executed indefinitely unless the CSS interrupt pending bit is cleared. As a consequence, the application has to clear the CSS interrupt in the NMI ISR by setting the CSSC bit in the Clock interrupt register (RCC\_CIR).

If the HSE oscillator is used directly or indirectly as the system clock (indirectly meaning that it is directly used as PLL input clock, and that PLL clock is the system clock) and a failure is detected, then the system clock switches to the HSI oscillator and the HSE oscillator is disabled.

If the HSE oscillator clock was the clock source of PLL used as the system clock when the failure occurred, PLL is also disabled. In this case, if the PLLI2S was enabled, it is also disabled when the HSE fails.

# 5.2.8 RTC/AWU clock

Once the RTCCLK clock source has been selected, the only possible way of modifying the selection is to reset the power domain.

The RTCCLK clock source can be either the HSE 1 MHz (HSE divided by a programmable prescaler), the LSE or the LSI clock. This is selected by programming the RTCSEL[1:0] bits in the RCC Backup domain control register (RCC\_BDCR) and the RTCPRE[4:0] bits in RCC clock configuration register (RCC\_CFGR). This selection cannot be modified without resetting the Backup domain.

If the LSE is selected as the RTC clock, the RTC will work normally if the backup or the system supply disappears. If the LSI is selected as the AWU clock, the AWU state is not guaranteed if the system supply disappears. If the HSE oscillator divided by a value between 2 and 31 is used as the RTC clock, the RTC state is not guaranteed if the backup or the system supply disappears.

The LSE clock is in the Backup domain, whereas the HSE and LSI clocks are not. As a consequence:

- If LSE is selected as the RTC clock:
  - The RTC continues to work even if the V<sub>DD</sub> supply is switched off, provided the V<sub>BAT</sub> supply is maintained.
- If LSI is selected as the Auto-wakeup unit (AWU) clock:
  - The AWU state is not guaranteed if the V<sub>DD</sub> supply is powered off. Refer to Section 5.2.5: LSI clock on page 90 for more details on LSI calibration.
- If the HSE clock is used as the RTC clock:
  - The RTC state is not guaranteed if the V<sub>DD</sub> supply is powered off or if the internal voltage regulator is powered off (removing power from the 1.2 V domain).

Note:

To read the RTC calendar register when the APB1 clock frequency is less than seven times the RTC clock frequency ( $f_{APB1} < 7xf_{RTCLCK}$ ), the software must read the calendar time and date registers twice. The data are correct if the second read access to RTC\_TR gives the same result than the first one. Otherwise a third read access must be performed.

# 5.2.9 Watchdog clock

If the independent watchdog (IWDG) is started by either hardware option or software access, the LSI oscillator is forced ON and cannot be disabled. After the LSI oscillator temporization, the clock is provided to the IWDG.

# 5.2.10 Clock-out capability

Two microcontroller clock output (MCO) pins are available:

#### MCO1

You can output four different clock sources onto the MCO1 pin (PA8) using the configurable prescaler (from 1 to 5):

- HSI clock
- LSE clock
- HSE clock
- PLL clock

The desired clock source is selected using the MCO1PRE[2:0] and MCO1[1:0] bits in the *RCC clock configuration register (RCC\_CFGR)*.

#### MCO2

You can output four different clock sources onto the MCO2 pin (PC9) using the configurable prescaler (from 1 to 5):

- HSE clock
- PLL clock
- System clock (SYSCLK)
- PLLI2S clock

The desired clock source is selected using the MCO2PRE[2:0] and MCO2 bits in the RCC clock configuration register (RCC\_CFGR).

For the different MCO pins, the corresponding GPIO port has to be programmed in alternate function mode.

The selected clock to output onto MCO must not exceed 100 MHz (the maximum I/O speed).

# 5.2.11 Internal/external clock measurement using TIM5/TIM11

It is possible to indirectly measure the frequencies of all on-board clock source generators by means of the input capture of TIM5 channel4 and TIM11 channel1 as shown in *Figure 11* and *Figure 11*.

#### Internal/external clock measurement using TIM5 channel4

TIM5 has an input multiplexer which allows choosing whether the input capture is triggered by the I/O or by an internal clock. This selection is performed through the TI4\_RMP [1:0] bits in the TIM5 OR register.

The primary purpose of having the LSE connected to the channel4 input capture is to be able to precisely measure the HSI (this requires to have the HSI used as the system clock source). The number of HSI clock counts between consecutive edges of the LSE signal provides a measurement of the internal clock period. Taking advantage of the high precision of LSE crystals (typically a few tens of ppm) we can determine the internal clock frequency with the same resolution, and trim the source to compensate for manufacturing-process and/or temperature- and voltage-related frequency deviations.

The HSI oscillator has dedicated, user-accessible calibration bits for this purpose.

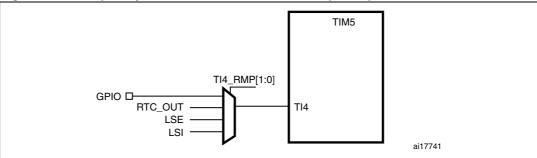
The basic concept consists in providing a relative measurement (e.g. HSI/LSE ratio): the precision is therefore tightly linked to the ratio between the two clock sources. The greater the ratio, the better the measurement.

It is also possible to measure the LSI frequency: this is useful for applications that do not have a crystal. The ultralow-power LSI oscillator has a large manufacturing process deviation: by measuring it versus the HSI clock source, it is possible to determine its frequency with the precision of the HSI. The measured value can be used to have more accurate RTC time base timeouts (when LSI is used as the RTC clock source) and/or an IWDG timeout with an acceptable accuracy.

Use the following procedure to measure the LSI frequency:

- Enable the TIM5 timer and configure channel4 in Input capture mode.
- 2. Set the TI4\_RMP bits in the TIM5\_OR register to 0x01 to connect the LSI clock internally to TIM5 channel4 input capture for calibration purposes.
- 3. Measure the LSI clock frequency using the TIM5 capture/compare 4 event or interrupt.
- 4. Use the measured LSI frequency to update the prescaler of the RTC depending on the desired time base and/or to compute the IWDG timeout.

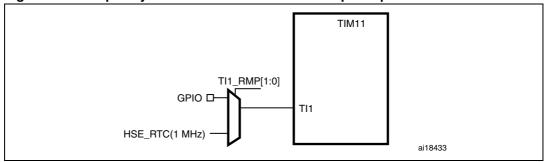
Figure 11. Frequency measurement with TIM5 in Input capture mode



# Internal/external clock measurement using TIM11 channel1

TIM11 has an input multiplexer which allows choosing whether the input capture is triggered by the I/O or by an internal clock. This selection is performed through TI1\_RMP [1:0] bits in the TIM11\_OR register. The HSE\_RTC clock (HSE divided by a programmable prescaler) is connected to channel 1 input capture to have a rough indication of the external crystal frequency. This requires that the HSI is the system clock source. This can be useful for instance to ensure compliance with the IEC 60730/IEC 61335 standards which require to be able to determine harmonic or subharmonic frequencies (–50/+100% deviations).

Figure 12. Frequency measurement with TIM11 in Input capture mode



# 5.3 RCC registers

Refer to *Section 1.1: List of abbreviations for registers* for a list of abbreviations used in register descriptions.

# 5.3.1 RCC clock control register (RCC\_CR)

Address offset: 0x00

Reset value: 0x0000 XX83 where X is undefined.

Access: no wait state, word, half-word and byte access

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	Rese	erved		PLLI2S RDY	PLLI2S ON	PLLRDY	PLLON		Rese	erved		CSS ON	HSE BYP	HSE RDY	HSE ON
				r	rw	r	rw					rw	rw	r	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	HSICAL[7:0] HSITRIM[4:0]										:0]		Res.	HSI RDY	HSION
r	r	r	r	r	r	r	r	rw	rw	rw	rw	rw	nes.	r	rw

Bits 31:28 Reserved, must be kept at reset value.

#### Bit 27 PLLI2SRDY: PLLI2S clock ready flag

Set by hardware to indicate that the PLLI2S is locked.

0: PLLI2S unlocked 1: PLLI2S locked

#### Bit 26 PLLI2SON: PLLI2S enable

Set and cleared by software to enable PLLI2S.

Cleared by hardware when entering Stop or Standby mode.

0: PLLI2S OFF 1: PLLI2S ON

#### Bit 25 PLLRDY: Main PLL (PLL) clock ready flag

Set by hardware to indicate that PLL is locked.

0: PLL unlocked 1: PLL locked

#### Bit 24 PLLON: Main PLL (PLL) enable

Set and cleared by software to enable PLL.

Cleared by hardware when entering Stop or Standby mode. This bit cannot be reset if PLL clock is used as the system clock.

0: PLL OFF 1: PLL ON

#### Bits 23:20 Reserved, must be kept at reset value.

#### Bit 19 CSSON: Clock security system enable

Set and cleared by software to enable the clock security system. When CSSON is set, the clock detector is enabled by hardware when the HSE oscillator is ready, and disabled by hardware if an oscillator failure is detected.

0: Clock security system OFF (Clock detector OFF)

1: Clock security system ON (Clock detector ON if HSE oscillator is stable, OFF if not)

#### Bit 18 HSEBYP: HSE clock bypass

Set and cleared by software to bypass the oscillator with an external clock. The external clock must be enabled with the HSEON bit, to be used by the device.

The HSEBYP bit can be written only if the HSE oscillator is disabled.

0: HSE oscillator not bypassed

1: HSE oscillator bypassed with an external clock

#### Bit 17 HSERDY: HSE clock ready flag

Set by hardware to indicate that the HSE oscillator is stable. After the HSEON bit is cleared, HSERDY goes low after 6 HSE oscillator clock cycles.

0: HSE oscillator not ready

1: HSE oscillator ready

#### Bit 16 HSEON: HSE clock enable

Set and cleared by software.

Cleared by hardware to stop the HSE oscillator when entering Stop or Standby mode. This bit cannot be reset if the HSE oscillator is used directly or indirectly as the system clock.

0: HSE oscillator OFF

1: HSE oscillator ON

# Bits 15:8 HSICAL[7:0]: Internal high-speed clock calibration

These bits are initialized automatically at startup.

## Bits 7:3 HSITRIM[4:0]: Internal high-speed clock trimming

These bits provide an additional user-programmable trimming value that is added to the HSICAL[7:0] bits. It can be programmed to adjust to variations in voltage and temperature that influence the frequency of the internal HSI RC.

Bit 2 Reserved, must be kept at reset value.

#### Bit 1 HSIRDY: Internal high-speed clock ready flag

Set by hardware to indicate that the HSI oscillator is stable. After the HSION bit is cleared, HSIRDY goes low after 6 HSI clock cycles.

0: HSI oscillator not ready

1: HSI oscillator ready

### Bit 0 HSION: Internal high-speed clock enable

Set and cleared by software.

Set by hardware to force the HSI oscillator ON when leaving the Stop or Standby mode or in case of a failure of the HSE oscillator used directly or indirectly as the system clock. This bit cannot be cleared if the HSI is used directly or indirectly as the system clock.

0: HSI oscillator OFF

1: HSI oscillator ON

#### 5.3.2 RCC PLL configuration register (RCC\_PLLCFGR)

Address offset: 0x04

Reset value: 0x2400 3010

Access: no wait state, word, half-word and byte access.

This register is used to configure the PLL clock outputs according to the formulas:

- $f_{(VCO clock)} = f_{(PLL clock input)} \times (PLLN / PLLM)$
- $f_{(PLL \text{ general clock output})} = f_{(VCO \text{ clock})} / PLLP$
- $f_{\text{(USB OTG FS, SDIO, RNG clock output)}} = f_{\text{(VCO clock)}} / PLLQ$

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	Rese	erved		PLLQ3	PLLQ2	PLLQ1	PLLQ0	Reserv	PLLSR C		Rese	erved		PLLP1	PLLP0
				rw	rw	rw	rw	eu	rw			rw	rw		
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Reserv					PLLN					PLLM5	PLLM4	PLLM2	PLLM1	PLLM0	
ed	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	

Bit 31:28 Reserved, must be kept at reset value.

Bits 27:24 PLLQ: Main PLL (PLL) division factor for USB OTG FS, SDIO and random number generator

Set and cleared by software to control the frequency of USB OTG FS clock, the random number generator clock and the SDIO clock. These bits should be written only if PLL is disabled.

Caution: The USB OTG FS requires a 48 MHz clock to work correctly. The SDIO and the random number generator need a frequency lower than or equal to 48 MHz to work

USB OTG FS clock frequency = VCO frequency / PLLQ with 2  $\leq$  PLLQ  $\leq$  15

0000: PLLQ = 0, wrong configuration 0001: PLLQ = 1, wrong configuration

0010: PLLQ = 2 0011: PLLQ = 3 0100: PLLQ = 4

1111: PLLQ = 15

Bit 23 Reserved, must be kept at reset value.

Bit 22 PLLSRC: Main PLL(PLL) and audio PLL (PLLI2S) entry clock source

Set and cleared by software to select PLL and PLLI2S clock source. This bit can be written only when PLL and PLLI2S are disabled.

0: HSI clock selected as PLL and PLLI2S clock entry

1: HSE oscillator clock selected as PLL and PLLI2S clock entry

Bits 21:18 Reserved, must be kept at reset value.

#### Bits 17:16 PLLP: Main PLL (PLL) division factor for main system clock

Set and cleared by software to control the frequency of the general PLL output clock. These bits can be written only if PLL is disabled.

Caution: The software has to set these bits correctly not to exceed 168 MHz on this domain.

```
PLL output clock frequency = VCO frequency / PLLP with PLLP = 2, 4, 6, or 8 00: PLLP = 2
```

00: PLLP = 2 01: PLLP = 4 10: PLLP = 6 11: PLLP = 8

#### Bits 14:6 PLLN: Main PLL (PLL) multiplication factor for VCO

Set and cleared by software to control the multiplication factor of the VCO. These bits can be written only when PLL is disabled. Only half-word and word accesses are allowed to write these bits.

**Caution:** The software has to set these bits correctly to ensure that the VCO output frequency is between 192 and 432 MHz.

```
VCO output frequency = VCO input frequency \times PLLN with 192 \le PLLN \le 432 000000000: PLLN = 0, wrong configuration 00000001: PLLN = 1, wrong configuration ... 011000000: PLLN = 192 ... 110110000: PLLN = 432 110110001: PLLN = 433, wrong configuration ... 111111111: PLLN = 511, wrong configuration
```

Bits 5:0 **PLLM:** Division factor for the main PLL (PLL) and audio PLL (PLLI2S) input clock

Set and cleared by software to divide the PLL and PLLI2S input clock before the VCO.

These bits can be written only when the PLL and PLLI2S are disabled.

**Caution:** The software has to set these bits correctly to ensure that the VCO input frequency ranges from 1 to 2 MHz. It is recommended to select a frequency of 2 MHz to limit PLL jitter.

```
VCO input frequency = PLL input clock frequency / PLLM with 2 \le PLLM \le 63 000000: PLLM = 0, wrong configuration 000001: PLLM = 1, wrong configuration 000010: PLLM = 2 000011: PLLM = 3 000100: PLLM = 4 ... 111110: PLLM = 62 111111: PLLM = 63
```

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# 5.3.3 RCC clock configuration register (RCC\_CFGR)

Address offset: 0x08

Reset value: 0x0000 0000

Access: 0 ≤ wait state ≤ 2, word, half-word and byte access

1 or 2 wait states inserted only if the access occurs during a clock source switch.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
МС	002	МС	:O2 PRE[	2:0]	МС	O1 PRE[	2:0]	I2SSC R	МС	01		R	TCPRE[4	l:0]	
rw		rw	rw	rw	rw	rw	rw	rw	rw		rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
F	PRE2[2:0	0]	Р	PRE1[2:	0]	Poor	arvad		HPRI	E[3:0]		SWS1	SWS0	SW1	SW0
rw	rw	rw	rw	rw	rw	Reserved		rw	rw	rw	rw	r	r	rw	rw

#### Bits 31:30 MCO2[1:0]: Microcontroller clock output 2

Set and cleared by software. Clock source selection may generate glitches on MCO2. It is highly recommended to configure these bits only after reset before enabling the external oscillators and the PLLs.

00: System clock (SYSCLK) selected

01: PLLI2S clock selected

10: HSE oscillator clock selected

11: PLL clock selected

#### Bits 27:29 MCO2PRE: MCO2 prescaler

Set and cleared by software to configure the prescaler of the MCO2. Modification of this prescaler may generate glitches on MCO2. It is highly recommended to change this prescaler only after reset before enabling the external oscillators and the PLLs.

0xx: no division 100: division by 2 101: division by 3 110: division by 4 111: division by 5

#### Bits 24:26 MCO1PRE: MCO1 prescaler

Set and cleared by software to configure the prescaler of the MCO1. Modification of this prescaler may generate glitches on MCO1. It is highly recommended to change this prescaler only after reset before enabling the external oscillators and the PLL.

0xx: no division 100: division by 2 101: division by 3 110: division by 4 111: division by 5

#### Bit 23 I2SSRC: I2S clock selection

Set and cleared by software. This bit allows to select the I2S clock source between the PLLI2S clock and the external clock. It is highly recommended to change this bit only after reset and before enabling the I2S module.

0: PLLI2S clock used as I2S clock source

1: External clock mapped on the I2S\_CKIN pin used as I2S clock source

#### Bits 22:21 MCO1: Microcontroller clock output 1

Set and cleared by software. Clock source selection may generate glitches on MCO1. It is highly recommended to configure these bits only after reset before enabling the external oscillators and PLL.

00: HSI clock selected01: LSE oscillator selected10: HSE oscillator clock selected11: PLL clock selected

#### Bits 20:16 RTCPRE: HSE division factor for RTC clock

Set and cleared by software to divide the HSE clock input clock to generate a 1 MHz clock for RTC

Caution: The software has to set these bits correctly to ensure that the clock supplied to the RTC is 1 MHz. These bits must be configured if needed before selecting the RTC clock source.

00000: no clock 00001: no clock 00010: HSE/2 00011: HSE/3 00100: HSE/4 ... 11110: HSE/30 11111: HSE/31

#### Bits 15:13 PPRE2: APB high-speed prescaler (APB2)

Set and cleared by software to control APB high-speed clock division factor.

Caution: The software has to set these bits correctly not to exceed 84 MHz on this domain. The clocks are divided with the new prescaler factor from 1 to 16 AHB cycles after PPRE2 write.

0xx: AHB clock not divided 100: AHB clock divided by 2 101: AHB clock divided by 4 110: AHB clock divided by 8 111: AHB clock divided by 16

#### Bits 12:10 **PPRE1:** APB Low speed prescaler (APB1)

Set and cleared by software to control APB low-speed clock division factor.

Caution: The software has to set these bits correctly not to exceed 42 MHz on this domain. The clocks are divided with the new prescaler factor from 1 to 16 AHB cycles after PPRE1 write.

0xx: AHB clock not divided 100: AHB clock divided by 2 101: AHB clock divided by 4 110: AHB clock divided by 8 111: AHB clock divided by 16

Bits 9:8 Reserved, must be kept at reset value.

#### Bits 7:4 HPRE: AHB prescaler

Set and cleared by software to control AHB clock division factor.

Caution: The clocks are divided with the new prescaler factor from 1 to 16 AHB cycles after

HPRE write.

Caution: The AHB clock frequency must be at least 25 MHz when the Ethernet is used.

0xxx: system clock not divided 1000: system clock divided by 2 1001: system clock divided by 4 1010: system clock divided by 8 1011: system clock divided by 16 1100: system clock divided by 64 1101: system clock divided by 128 1110: system clock divided by 256 1111: system clock divided by 512

#### Bits 3:2 SWS: System clock switch status

Set and cleared by hardware to indicate which clock source is used as the system clock.

00: HSI oscillator used as the system clock

01: HSE oscillator used as the system clock

10: PLL used as the system clock

11: not applicable

# Bits 1:0 SW: System clock switch

Set and cleared by software to select the system clock source.

Set by hardware to force the HSI selection when leaving the Stop or Standby mode or in case of failure of the HSE oscillator used directly or indirectly as the system clock.

00: HSI oscillator selected as system clock

01: HSE oscillator selected as system clock

10: PLL selected as system clock

11: not allowed

# 5.3.4 RCC clock interrupt register (RCC\_CIR)

Address offset: 0x0C

Reset value: 0x0000 0000

Access: no wait state, word, half-word and byte access

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
			Rese	erved				CSSC	Reser ved	PLLI2S RDYC	PLL RDYC	HSE RDYC	HSI RDYC	LSE RDYC	LSI RDYC
								w	veu	w	W	w	w	W	w
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Page	Reserved		PLL RDYIE	HSE RDYIE	HSI RDYIE	LSE RDYIE	LSI RDYIE	CSSF	Reser	PLLI2S RDYF	PLL RDYF	HSE RDYF	HSI RDYF	LSE RDYF	LSI RDYF
nese	Heserved		rw	rw	rw	rw	rw	r	ved	r	r	r	r	r	r

- Bits 31:24 Reserved, must be kept at reset value.
  - Bit 23 CSSC: Clock security system interrupt clear

This bit is set by software to clear the CSSF flag.

0: No effect

1: Clear CSSF flag

- Bits 22 Reserved, must be kept at reset value.
- Bit 21 PLLI2SRDYC: PLLI2S ready interrupt clear

This bit is set by software to clear the PLLI2SRDYF flag.

0: No effect

1: PLLI2SRDYF cleared

Bit 20 PLLRDYC: Main PLL(PLL) ready interrupt clear

This bit is set by software to clear the PLLRDYF flag.

0: No effect

1: PLLRDYF cleared

Bit 19 HSERDYC: HSE ready interrupt clear

This bit is set by software to clear the HSERDYF flag.

0: No effect

1: HSERDYF cleared

Bit 18 HSIRDYC: HSI ready interrupt clear

This bit is set software to clear the HSIRDYF flag.

0: No effect

1: HSIRDYF cleared

Bit 17 LSERDYC: LSE ready interrupt clear

This bit is set by software to clear the LSERDYF flag.

0: No effect

1: LSERDYF cleared

Bit 16 LSIRDYC: LSI ready interrupt clear

This bit is set by software to clear the LSIRDYF flag.

0: No effect

1: LSIRDYF cleared

- Bits 15:12 Reserved, must be kept at reset value.
  - Bit 13 PLLI2SRDYIE: PLLI2S ready interrupt enable

Set and cleared by software to enable/disable interrupt caused by PLLI2S lock.

0: PLLI2S lock interrupt disabled

1: PLLI2S lock interrupt enabled

Bit 12 PLLRDYIE: Main PLL (PLL) ready interrupt enable

Set and cleared by software to enable/disable interrupt caused by PLL lock.

0: PLL lock interrupt disabled

1: PLL lock interrupt enabled

Bit 11 HSERDYIE: HSE ready interrupt enable

Set and cleared by software to enable/disable interrupt caused by the HSE oscillator stabilization.

0: HSE ready interrupt disabled

1: HSE ready interrupt enabled

#### Bit 10 HSIRDYIE: HSI ready interrupt enable

Set and cleared by software to enable/disable interrupt caused by the HSI oscillator stabilization.

0: HSI ready interrupt disabled

1: HSI ready interrupt enabled

#### Bit 9 LSERDYIE: LSE ready interrupt enable

Set and cleared by software to enable/disable interrupt caused by the LSE oscillator stabilization.

0: LSE ready interrupt disabled

1: LSE ready interrupt enabled

#### Bit 8 LSIRDYIE: LSI ready interrupt enable

Set and cleared by software to enable/disable interrupt caused by LSI oscillator stabilization.

0: LSI ready interrupt disabled

1: LSI ready interrupt enabled

#### Bit 7 CSSF: Clock security system interrupt flag

Set by hardware when a failure is detected in the HSE oscillator.

Cleared by software setting the CSSC bit.

0: No clock security interrupt caused by HSE clock failure

1: Clock security interrupt caused by HSE clock failure

#### Bits 6 Reserved, must be kept at reset value.

#### Bit 5 PLLI2SRDYF: PLLI2S ready interrupt flag

Set by hardware when the PLLI2S locks and PLLI2SRDYDIE is set.

Cleared by software setting the PLLRI2SDYC bit.

0: No clock ready interrupt caused by PLLI2S lock

1: Clock ready interrupt caused by PLLI2S lock

#### Bit 4 PLLRDYF: Main PLL (PLL) ready interrupt flag

Set by hardware when PLL locks and PLLRDYDIE is set.

Cleared by software setting the PLLRDYC bit.

0: No clock ready interrupt caused by PLL lock

1: Clock ready interrupt caused by PLL lock

#### Bit 3 HSERDYF: HSE ready interrupt flag

Set by hardware when External High Speed clock becomes stable and HSERDYDIE is set.

Cleared by software setting the HSERDYC bit.

0: No clock ready interrupt caused by the HSE oscillator

1: Clock ready interrupt caused by the HSE oscillator

#### Bit 2 HSIRDYF: HSI ready interrupt flag

Set by hardware when the Internal High Speed clock becomes stable and HSIRDYDIE is set.

Cleared by software setting the HSIRDYC bit.

0: No clock ready interrupt caused by the HSI oscillator

1: Clock ready interrupt caused by the HSI oscillator

#### Bit 1 LSERDYF: LSE ready interrupt flag

Set by hardware when the External Low Speed clock becomes stable and LSERDYDIE is set.

Cleared by software setting the LSERDYC bit.

0: No clock ready interrupt caused by the LSE oscillator

1: Clock ready interrupt caused by the LSE oscillator

#### Bit 0 LSIRDYF: LSI ready interrupt flag

Set by hardware when the internal low speed clock becomes stable and LSIRDYDIE is set.

Cleared by software setting the LSIRDYC bit.

0: No clock ready interrupt caused by the LSI oscillator

1: Clock ready interrupt caused by the LSI oscillator

# 5.3.5 RCC AHB1 peripheral reset register (RCC\_AHB1RSTR)

Address offset: 0x10

Reset value: 0x0000 0000

Access: no wait state, word, half-word and byte access.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
Res	erved	OTGHS RST		Reserved	l	ETHMAC RST	Rese	erved	DMA2 RST	DMA1 RST					
		rw				rw			rw	rw					
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Reserve		CRCRS T Reserved			ed	GPIOI RST	GPIOH RST	GPIOGG RST	GPIOF RST	GPIOE RST	GPIOD RST	GPIOC RST	GPIOB RST	GPIOA RST
			rw				rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:30 Reserved, must be kept at reset value.

Bit 29 OTGHSRST: USB OTG HS module reset

Set and cleared by software.

0: does not reset the USB OTG HS module

1: resets the USB OTG HS module

Bits 28:26 Reserved, must be kept at reset value.

Bit 25 ETHMACRST: Ethernet MAC reset

Set and cleared by software.

0: does not reset Ethernet MAC

1: resets Ethernet MAC

Bits 24:23 Reserved, must be kept at reset value.

Bit 22 DMA2RST: DMA2 reset

Set and cleared by software.

0: does not reset DMA2

1: resets DMA2

Bit 21 DMA1RST: DMA2 reset

Set and cleared by software.

0: does not reset DMA2

1: resets DMA2

Bits 20:13 Reserved, must be kept at reset value.

Bit 12 CRCRST: CRC reset

Set and cleared by software.

0: does not reset CRC

1: resets CRC

Bits 11:9 Reserved, must be kept at reset value.

#### Bit 8 GPIOIRST: IO port I reset

Set and cleared by software.

0: does not reset IO port I

1: resets IO port I

#### Bit 7 GPIOHRST: IO port H reset

Set and cleared by software.

0: does not reset IO port H

1: resets IO port H

#### Bits 6 GPIOGRST: IO port G reset

Set and cleared by software.

0: does not reset IO port G

1: resets IO port G

#### Bit 5 GPIOFRST: IO port F reset

Set and cleared by software.

0: does not reset IO port F

1: resets IO port F

#### Bit 4 GPIOERST: IO port E reset

Set and cleared by software.

0: does not reset IO port E

1: resets IO port E

#### Bit 3 GPIODRST: IO port D reset

Set and cleared by software.

0: does not reset IO port D

1: resets IO port D

#### Bit 2 GPIOCRST: IO port C reset

Set and cleared by software.

0: does not reset IO port C

1: resets IO port C

## Bit 1 GPIOBRST: IO port B reset

Set and cleared by software.

0: does not reset IO port B

1:resets IO port B

# Bit 0 GPIOARST: IO port A reset

Set and cleared by software.

0: does not reset IO port A

1: resets IO port A

# 5.3.6 RCC AHB2 peripheral reset register (RCC\_AHB2RSTR)

Address offset: 0x14

Reset value: 0x0000 0000

Access: no wait state, word, half-word and byte access

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
Reserved															
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Reserved								RNG RST	HASH RST	CRYP RST	Reserved			DCMI RST
								rw	rw	rw	rw				rw

Bits 31:8 Reserved, must be kept at reset value.

Bit 7 OTGFSRST: USB OTG FS module reset

Set and cleared by software.

0: does not reset the USB OTG FS module

1: resets the USB OTG FS module

Bit 6 RNGRST: Random number generator module reset

Set and cleared by software.

0: does not reset the random number generator module

1: resets the random number generator module

Bit 5 HASHRST: Hash module reset

Set and cleared by software.

0: does not reset the HASH module

1: resets the HASH module

Bit 4 CRYPRST: Cryptographic module reset

Set and cleared by software.

0: does not reset the cryptographic module

1: resets the cryptographic module

Bit 3:1 Reserved, must be kept at reset value.

Bit 0 DCMIRST: Camera interface reset

Set and cleared by software.

0: does not reset the Camera interface

1: resets the Camera interface

# 5.3.7 RCC AHB3 peripheral reset register (RCC\_AHB3RSTR)

Address offset: 0x18

Reset value: 0x0000 0000

Access: no wait state, word, half-word and byte access.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
Reserved															
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Reserved													FSMCRST	
														rw	

Bits 31:1 Reserved, must be kept at reset value.

Bit 0 FSMCRST: Flexible static memory controller module reset

Set and cleared by software.

0: does not reset the FSMC module

1: resets the FSMC module

# 5.3.8 RCC APB1 peripheral reset register (RCC\_APB1RSTR)

Address offset: 0x20

Reset value: 0x0000 0000

Access: no wait state, word, half-word and byte access.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
Rese	erved	DACRST	PWR RST	Reser-	CAN2 RST	CAN1 RST	Reser-	I2C3 RST	I2C2 RST	I2C1 RST	UART5 RST	UART4 RST	UART3 RST	UART2 RST	Reser-
			rw	ved		rw	ved	rw	rw	rw	rw	rw	rw	rw	ved
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
SPI3 RST	SPI2 RST	Rese	Reserved		WDG RST Reserved		TIM14 RST	TIM13 RST	TIM12 RST	TIM7 RST	TIM6 RST	TIM5 RST	TIM4 RST	TIM3 RST	TIM2 RST
rw	rw			rw			rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:30 Reserved, must be kept at reset value.

Bit 29 DACRST: DAC reset

Set and cleared by software.

0: does not reset the DAC interface

1: resets the DAC interface

Bit 28 PWRRST: Power interface reset

Set and cleared by software.

0: does not reset the power interface

1: resets the power interface

Bit 27 Reserved, must be kept at reset value.

#### Bit 26 CAN2RST: CAN2 reset

Set and cleared by software.

0: does not reset CAN2

1: resets CAN2

#### Bit 25 CAN1RST: CAN1 reset

Set and cleared by software.

0: does not reset CAN1

1: resets CAN1

#### Bit 24 Reserved, must be kept at reset value.

#### Bit 23 I2C3RST: I2C3 reset

Set and cleared by software.

0: does not reset I2C3

1: resets I2C3

#### Bit 22 I2C2RST: I2C2 reset

Set and cleared by software.

0: does not reset I2C2

1: resets I2C2

#### Bit 21 I2C1RST: I2C1 reset

Set and cleared by software.

0: does not reset I2C1

1: resets I2C1

#### Bit 20 UART5RST: UART5 reset

Set and cleared by software.

0: does not reset UART5

1: resets UART5

#### Bit 19 UART4RST: USART4 reset

Set and cleared by software.

0: does not reset UART4

1: resets UART4

## Bit 18 USART3RST: USART3 reset

Set and cleared by software.

0: does not reset USART3

1: resets USART3

### Bit 17 USART2RST: USART2 reset

Set and cleared by software.

0: does not reset USART2

1: resets USART2

#### Bit 16 Reserved, must be kept at reset value.

#### Bit 15 SPI3RST: SPI3 reset

Set and cleared by software.

0: does not reset SPI3

1: resets SPI3

#### Bit 14 SPI2RST: SPI2 reset

Set and cleared by software.

0: does not reset SPI2

1: resets SPI2

#### Bits 13:12 Reserved, must be kept at reset value.

#### Bit 11 WWDGRST: Window watchdog reset

Set and cleared by software.

0: does not reset the window watchdog

1: resets the window watchdog

#### Bits 10:9 Reserved, must be kept at reset value.

#### Bit 8 TIM14RST: TIM14 reset

Set and cleared by software.

0: does not reset TIM14

1: resets TIM14

#### Bit 7 TIM13RST: TIM13 reset

Set and cleared by software.

0: does not reset TIM13

1: resets TIM13

#### Bit 6 TIM12RST: TIM12 reset

Set and cleared by software.

0: does not reset TIM12

1: resets TIM12

#### Bit 5 TIM7RST: TIM7 reset

Set and cleared by software.

0: does not reset TIM7

1: resets TIM7

#### Bit 4 TIM6RST: TIM6 reset

Set and cleared by software.

0: does not reset TIM6

1: resets TIM6

# Bit 3 TIM5RST: TIM5 reset

Set and cleared by software.

0: does not reset TIM5

1: resets TIM5

# Bit 2 TIM4RST: TIM4 reset

Set and cleared by software.

0: does not reset TIM4

1: resets TIM4

#### Bit 1 TIM3RST: TIM3 reset

Set and cleared by software.

0: does not reset TIM3

1: resets TIM3

#### Bit 0 TIM2RST: TIM2 reset

Set and cleared by software.

0: does not reset TIM2

1: resets TIM2

## 5.3.9 RCC APB2 peripheral reset register (RCC\_APB2RSTR)

Address offset: 0x24

Reset value: 0x0000 0000

Access: no wait state, word, half-word and byte access.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
						Reserve	ed						TIM11 RST	TIM10 RST	TIM9 RST
													rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Reser-	SYSCF G RST	Reser- ved	SPI1 RST	SDIO RST	Res	erved	ADC RST	Rese	erved	USART6 RST	USART1 RST	Rese	erved	TIM8 RST	TIM1 RST
veu	rw	veu	rw	rw			rw			rw	rw			rw	rw

Bits 31:19 Reserved, must be kept at reset value.

Bit 18 TIM11RST: TIM11 reset

Set and cleared by software.

0: does not reset TIM11

1: resets TIM14

Bit 17 TIM10RST: TIM10 reset

Set and cleared by software.

0: does not reset TIM10

1: resets TIM10

Bit 16 TIM9RST: TIM9 reset

Set and cleared by software.

0: does not reset TIM9

1: resets TIM9

Bit 15 Reserved, must be kept at reset value.

Bit 14 SYSCFGRST: System configuration controller reset

Set and cleared by software.

0: does not reset the System configuration controller

1: resets the System configuration controller

Bit 13 Reserved, must be kept at reset value.

Bit 12 SPI1RST: SPI1 reset

Set and cleared by software.

0: does not reset SPI1

1: resets SPI1

Bit 11 SDIORST: SDIO reset

Set and cleared by software.

0: does not reset the SDIO module

1: resets the SDIO module

Bits 10:9 Reserved, must be kept at reset value.

Bit 8 ADCRST: ADC interface reset (common to all ADCs)

Set and cleared by software.

0: does not reset the ADC interface

1: resets the ADC interface

Bits 7:6 Reserved, must be kept at reset value.

Bit 5 USART6RST: USART6 reset

Set and cleared by software.

0: does not reset USART6

1: resets USART6

Bit 4 USART1RST: USART1 reset

Set and cleared by software.

0: does not reset USART1

1: resets USART1

Bits 3:2 Reserved, must be kept at reset value.

Bit 1 TIM8RST: TIM8 reset

Set and cleared by software.

0: does not reset TIM8

1: resets TIM8

Bit 0 TIM1RST: TIM1 reset

Set and cleared by software.

0: does not reset TIM1

1: resets TIM1

## 5.3.10 RCC AHB1 peripheral clock register (RCC\_AHB1ENR)

Address offset: 0x30

Reset value: 0x0010 0000

Access: no wait state, word, half-word and byte access.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
Reser- ved	OTGHS ULPIEN	OTGHS	ETHMA CPTPE N	ETHMA CRXEN		ETHMA CEN	Rese		DMA2EN	DMA1EN	CCMDATA RAMEN	Res.	BKPSR AMEN		erved
	rw	rw	rw	rw	rw	rw			rw	rw			rw		
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Reserved	d	CRCEN		Reserved	d	GPIOIE N	GPIOH EN	GPIOGE N	GPIOFE N	GPIOEEN	GPIOD EN	GPIOC EN	GPIOB EN	GPIOA EN
			rw				rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31 Reserved, must be kept at reset value.

Bit 30 OTGHSULPIEN: USB OTG HSULPI clock enable

Set and cleared by software.

0: USB OTG HS ULPI clock disabled 1: USB OTG HS ULPI clock enabled

Bit 29 OTGHSEN: USB OTG HS clock enable

Set and cleared by software.
0: USB OTG HS clock disabled
1: USB OTG HS clock enabled

Bit 28 ETHMACPTPEN: Ethernet PTP clock enable

Set and cleared by software.

0: Ethernet PTP clock disabled

1: Ethernet PTP clock enabled

Bit 27 ETHMACRXEN: Ethernet Reception clock enable

Set and cleared by software.

0: Ethernet Reception clock disabled

1: Ethernet Reception clock enabled

Bit 26 ETHMACTXEN: Ethernet Transmission clock enable

Set and cleared by software.

0: Ethernet Transmission clock disabled

1: Ethernet Transmission clock enabled

Bit 25 ETHMACEN: Ethernet MAC clock enable

Set and cleared by software.

0: Ethernet MAC clock disabled

1: Ethernet MAC clock enabled

Bits 24:23 Reserved, must be kept at reset value.

Bit 22 DMA2EN: DMA2 clock enable

Set and cleared by software.

0: DMA2 clock disabled

1: DMA2 clock enabled

#### Bit 21 DMA1EN: DMA1 clock enable

Set and cleared by software.

0: DMA1 clock disabled

1: DMA1 clock enabled

#### Bit 20 **CCMDATARAMEN**: CCM data RAM clock enable

Set and cleared by software.

0: CCM data RAM clock disabled

1: CCM data RAM clock enabled

#### Bits 19 Reserved, must be kept at reset value.

## Bit 18 **BKPSRAMEN:** Backup SRAM interface clock enable

Set and cleared by software.

0: Backup SRAM interface clock disabled

1: Backup SRAM interface clock enabled

### Bits 17:13 Reserved, must be kept at reset value.

#### Bit 12 CRCEN: CRC clock enable

Set and cleared by software.

0: CRC clock disabled

1: CRC clock enabled

#### Bits 11:9 Reserved, must be kept at reset value.

#### Bit 8 GPIOIEN: IO port I clock enable

Set and cleared by software.

0: IO port I clock disabled

1: IO port I clock enabled

### Bit 7 **GPIOHEN:** IO port H clock enable

Set and cleared by software.

0: IO port H clock disabled

1: IO port H clock enabled

## Bit 6 GPIOGEN: IO port G clock enable

Set and cleared by software.

0: IO port G clock disabled

1: IO port G clock enabled

## Bit 5 GPIOFEN: IO port F clock enable

Set and cleared by software.

0: IO port F clock disabled

1: IO port F clock enabled

### Bit 4 GPIOEEN: IO port E clock enable

Set and cleared by software.

0: IO port E clock disabled

1: IO port E clock enabled

## Bit 3 GPIODEN: IO port D clock enable

Set and cleared by software.

0: IO port D clock disabled

1: IO port D clock enabled

Bit 2 GPIOCEN: IO port C clock enable

Set and cleared by software.

0: IO port C clock disabled

1: IO port C clock enabled

Bit 1 GPIOBEN: IO port B clock enable

Set and cleared by software.

0: IO port B clock disabled

1: IO port B clock enabled

Bit 0 GPIOAEN: IO port A clock enable

Set and cleared by software.

0: IO port A clock disabled

1: IO port A clock enabled

## 5.3.11 RCC AHB2 peripheral clock enable register (RCC\_AHB2ENR)

Address offset: 0x34

Reset value: 0x0000 0000

Access: no wait state, word, half-word and byte access.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Res	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			Rese	erved				OTGFS EN	RNG EN	HASH EN	CRYP EN		Reserved	I	DCMI EN
								rw	rw	rw	rw				rw

Bits 31:8 Reserved, must be kept at reset value.

Bit 7 OTGFSEN: USB OTG FS clock enable

Set and cleared by software.

0: USB OTG FS clock disabled

1: USB OTG FS clock enabled

Bit 6 RNGEN: Random number generator clock enable

Set and cleared by software.

0: Random number generator clock disabled

1: Random number generator clock enabled

Bit 5 HASHEN: Hash modules clock enable

Set and cleared by software.

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0: Hash modules clock disabled

1: Hash modules clock enabled

Bit 4 CRYPEN: Cryptographic modules clock enable

Set and cleared by software.

0: cryptographic module clock disabled1: cryptographic module clock enabled

Bit 3:1 Reserved, must be kept at reset value.

Bit 0 **DCMIEN:** Camera interface enable
Set and cleared by software.
0: Camera interface clock disabled

1: Camera interface clock enabled

## 5.3.12 RCC AHB3 peripheral clock enable register (RCC\_AHB3ENR)

Address offset: 0x38

Reset value: 0x0000 0000

Access: no wait state, word, half-word and byte access.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Res	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							Reserved								FSMCEN
							neserved								rw

Bits 31:1 Reserved, must be kept at reset value.

Bit 0 FSMCEN: Flexible static memory controller module clock enable

Set and cleared by software.
0: FSMC module clock disabled
1: FSMC module clock enabled

# 5.3.13 RCC APB1 peripheral clock enable register (RCC\_APB1ENR)

Address offset: 0x40

Reset value: 0x0000 0000

Access: no wait state, word, half-word and byte access.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
Rese	erved	DAC EN	PWR EN	Reser- ved	CAN2 EN	CAN1 EN	Reser- ved	I2C3 EN	I2C2 EN	I2C1 EN	UART5 EN	UART4 EN	USART3 EN	USART2 EN	Reser-
		rw	rw	veu	rw	rw	veu	rw	rw	rw	rw	rw	rw	rw	veu
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
SPI3 EN	SPI2 EN	Rese	erved	WWDG EN	Rese	Reserved		TIM13 EN	TIM12 EN	TIM7 EN	TIM6 EN	TIM5 EN	TIM4 EN	TIM3 EN	TIM2 EN
rw	rw			rw			rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:30 Reserved, must be kept at reset value.

Bit 29 DACEN: DAC interface clock enable

Set and cleared by software.

0: DAC interface clock disabled

1: DAC interface clock enable

Bit 28 PWREN: Power interface clock enable

Set and cleared by software.

0: Power interface clock disabled

1: Power interface clock enable

Bit 27 Reserved, must be kept at reset value.

Bit 26 CAN2EN: CAN 2 clock enable

Set and cleared by software.

0: CAN 2 clock disabled

1: CAN 2 clock enabled

Bit 25 CAN1EN: CAN 1 clock enable

Set and cleared by software.

0: CAN 1 clock disabled

1: CAN 1 clock enabled

Bit 24 Reserved, must be kept at reset value.

Bit 23 I2C3EN: I2C3 clock enable

Set and cleared by software.

0: I2C3 clock disabled

1: I2C3 clock enabled

Bit 22 I2C2EN: I2C2 clock enable

Set and cleared by software.

0: I2C2 clock disabled

1: I2C2 clock enabled

Bit 21 I2C1EN: I2C1 clock enable

Set and cleared by software.

0: I2C1 clock disabled

1: I2C1 clock enabled

Bit 20 UART5EN: UART5 clock enable

Set and cleared by software.

0: UART5 clock disabled

1: UART5 clock enabled

Bit 19 UART4EN: UART4 clock enable

Set and cleared by software.

0: UART4 clock disabled

1: UART4 clock enabled

Bit 18 USART3EN: USART3 clock enable

Set and cleared by software.

0: USART3 clock disabled

1: USART3 clock enabled

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#### Bit 17 USART2EN: USART2 clock enable

Set and cleared by software.

0: USART2 clock disabled

1: USART2 clock enabled

Bit 16 Reserved, must be kept at reset value.

#### Bit 15 SPI3EN: SPI3 clock enable

Set and cleared by software.

0: SPI3 clock disabled

1: SPI3 clock enabled

#### Bit 14 SPI2EN: SPI2 clock enable

Set and cleared by software.

0: SPI2 clock disabled

1: SPI2 clock enabled

#### Bits 13:12 Reserved, must be kept at reset value.

### Bit 11 WWDGEN: Window watchdog clock enable

Set and cleared by software.

0: Window watchdog clock disabled

1: Window watchdog clock enabled

#### Bit 10:9 Reserved, must be kept at reset value.

#### Bit 8 TIM14EN: TIM14 clock enable

Set and cleared by software.

0: TIM14 clock disabled

1: TIM14 clock enabled

#### Bit 7 TIM13EN: TIM13 clock enable

Set and cleared by software.

0: TIM13 clock disabled

1: TIM13 clock enabled

#### Bit 6 TIM12EN: TIM12 clock enable

Set and cleared by software.

0: TIM12 clock disabled

1: TIM12 clock enabled

#### Bit 5 TIM7EN: TIM7 clock enable

Set and cleared by software.

0: TIM7 clock disabled

1: TIM7 clock enabled

#### Bit 4 TIM6EN: TIM6 clock enable

Set and cleared by software.

0: TIM6 clock disabled

1: TIM6 clock enabled

#### Bit 3 TIM5EN: TIM5 clock enable

Set and cleared by software.

0: TIM5 clock disabled

1: TIM5 clock enabled

Bit 2 TIM4EN: TIM4 clock enable

Set and cleared by software.

0: TIM4 clock disabled

1: TIM4 clock enabled

Bit 1 TIM3EN: TIM3 clock enable

Set and cleared by software.

0: TIM3 clock disabled

1: TIM3 clock enabled

Bit 0 TIM2EN: TIM2 clock enable

Set and cleared by software.

0: TIM2 clock disabled

1: TIM2 clock enabled

# 5.3.14 RCC APB2 peripheral clock enable register (RCC\_APB2ENR)

Address offset: 0x44

Reset value: 0x0000 0000

Access: no wait state, word, half-word and byte access.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
						Reserved							TIM11 EN	TIM10 EN	TIM9 EN
													rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Reser- ved	SYSCF G EN	Reser- ved	SPI1 EN	SDIO EN	ADC3 EN	ADC2 EN	ADC1 EN	Rese	erved	USART6 EN	USART1 EN		erved	TIM8 EN	TIM1 EN
ved	rw	Veu	rw	rw	rw	rw	rw			rw	rw	•		rw	rw

Bits 31:19 Reserved, must be kept at reset value.

Bit 18 TIM11EN: TIM11 clock enable

Set and cleared by software.

0: TIM11 clock disabled

1: TIM11 clock enabled

Bit 17 TIM10EN: TIM10 clock enable

Set and cleared by software.

0: TIM10 clock disabled

1: TIM10 clock enabled

Bit 16 TIM9EN: TIM9 clock enable

Set and cleared by software.

0: TIM9 clock disabled

1: TIM9 clock enabled

Bit 15 Reserved, must be kept at reset value.

Bit 14 SYSCFGEN: System configuration controller clock enable

Set and cleared by software.

0: System configuration controller clock disabled

1: System configuration controller clock enabled

Bit 13 Reserved, must be kept at reset value.

Bit 12 SPI1EN: SPI1 clock enable

Set and cleared by software.

0: SPI1 clock disabled

1: SPI1 clock enabled

Bit 11 SDIOEN: SDIO clock enable

Set and cleared by software.

0: SDIO module clock disabled

1: SDIO module clock enabled

Bit 10 ADC3EN: ADC3 clock enable

Set and cleared by software.

0: ADC3 clock disabled

1: ADC3 clock disabled

#### Bit 9 ADC2EN: ADC2 clock enable

Set and cleared by software.

0: ADC2 clock disabled

1: ADC2 clock disabled

#### Bit 8 ADC1EN: ADC1 clock enable

Set and cleared by software.

0: ADC1 clock disabled

1: ADC1 clock disabled

## Bits 7:6 Reserved, must be kept at reset value.

#### Bit 5 USART6EN: USART6 clock enable

Set and cleared by software.

0: USART6 clock disabled

1: USART6 clock enabled

#### Bit 4 USART1EN: USART1 clock enable

Set and cleared by software.

0: USART1 clock disabled

1: USART1 clock enabled

#### Bits 3:2 Reserved, must be kept at reset value.

#### Bit 1 TIM8EN: TIM8 clock enable

Set and cleared by software.

0: TIM8 clock disabled

1: TIM8 clock enabled

#### Bit 0 TIM1EN: TIM1 clock enable

Set and cleared by software.

0: TIM1 clock disabled

1: TIM1 clock enabled

# 5.3.15 RCC APB2 peripheral clock enable register (RCC\_APB2ENR)

Address offset: 0x44

Reset value: 0x0000 0000

Access: no wait state, word, half-word and byte access.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
				Rese	rved					SPI6EN	SPI5EN	Res.	TIM11 EN	TIM10 EN	TIM9 EN
										rw	rw		rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Reser-	SYSCF G EN	SPI4EN	SPI1 EN	SDIO EN	ADC3 EN	ADC2 EN	ADC1 EN	Rese		USART6 EN	USART1 EN		erved	TIM8 EN	TIM1 EN
veu	rw	rw	rw	rw	rw	rw	rw			rw	rw			rw	rw

Bits 31:22 Reserved, must be kept at reset value.

Bit 21 SPI6EN: SPI6 clock enable

Set and cleared by software.

0: SPI6 clock disabled

1: SPI6 clock enabled

Bit 20 SPI5EN: SPI5 clock enable

Set and cleared by software.

0: SPI5 clock disabled

1: SPI5 clock enabled

Bit 19 Reserved, must be kept at reset value.

Bit 18 TIM11EN: TIM11 clock enable

Set and cleared by software.

0: TIM11 clock disabled

1: TIM11 clock enabled

Bit 17 TIM10EN: TIM10 clock enable

Set and cleared by software.

0: TIM10 clock disabled

1: TIM10 clock enabled

Bit 16 TIM9EN: TIM9 clock enable

Set and cleared by software.

0: TIM9 clock disabled

1: TIM9 clock enabled

Bit 15 Reserved, must be kept at reset value.

Bit 14 SYSCFGEN: System configuration controller clock enable

Set and cleared by software.

0: System configuration controller clock disabled

1: System configuration controller clock enabled

Bit 13 SPI4EN: SPI4 clock enable

Set and cleared by software.

0: SPI4 clock disabled

1: SPI4 clock enabled

#### Bit 12 SPI1EN: SPI1 clock enable

Set and cleared by software.

0: SPI1 clock disabled

1: SPI1 clock enabled

#### Bit 11 SDIOEN: SDIO clock enable

Set and cleared by software.

0: SDIO module clock disabled

1: SDIO module clock enabled

#### Bit 10 ADC3EN: ADC3 clock enable

Set and cleared by software.

0: ADC3 clock disabled

1: ADC3 clock disabled

#### Bit 9 ADC2EN: ADC2 clock enable

Set and cleared by software.

0: ADC2 clock disabled

1: ADC2 clock disabled

#### Bit 8 ADC1EN: ADC1 clock enable

Set and cleared by software.

0: ADC1 clock disabled

1: ADC1 clock disabled

#### Bits 7:6 Reserved, must be kept at reset value.

#### Bit 5 USART6EN: USART6 clock enable

Set and cleared by software.

0: USART6 clock disabled

1: USART6 clock enabled

#### Bit 4 USART1EN: USART1 clock enable

Set and cleared by software.

0: USART1 clock disabled

1: USART1 clock enabled

#### Bits 3:2 Reserved, must be kept at reset value.

#### Bit 1 TIM8EN: TIM8 clock enable

Set and cleared by software.

0: TIM8 clock disabled

1: TIM8 clock enabled

#### Bit 0 TIM1EN: TIM1 clock enable

Set and cleared by software.

0: TIM1 clock disabled

1: TIM1 clock enabled

# 5.3.16 RCC AHB1 peripheral clock enable in low power mode register (RCC\_AHB1LPENR)

Address offset: 0x50

Reset value: 0x7E67 91FF

Access: no wait state, word, half-word and byte access.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
Reser- ved	OTGHS ULPILPEN		ETHPTP LPEN	ETHRX LPEN	ETHTX LPEN	ETHMAC LPEN		erved	DMA2 LPEN	DMA1 LPEN	Rese		BKPSRA M LPEN	SRAM2 LPEN	SRAM1 LPEN
	rw	rw	rw	rw	rw	rw			rw	rw			rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
FLITF LPEN	Reser	ved	CRC LPEN				GPIOI LPEN	GPIOH LPEN	GPIOGG LPEN	GPIOF LPEN	GPIOE LPEN	GPIOD LPEN	GPIOC LPEN	GPIOB LPEN	GPIOA LPEN
rw			rw	Reserved			rw	rw	rw	rw	rw	rw	rw	rw	rw

Bit 31 Reserved, must be kept at reset value.

Bit 30 OTGHSULPILPEN: USB OTG HS ULPI clock enable during Sleep mode

Set and cleared by software.

0: USB OTG HS ULPI clock disabled during Sleep mode

1: USB OTG HS ULPI clock enabled during Sleep mode

Bit 29 OTGHSLPEN: USB OTG HS clock enable during Sleep mode

Set and cleared by software.

0: USB OTG HS clock disabled during Sleep mode

1: USB OTG HS clock enabled during Sleep mode

Bit 28 ETHMACPTPLPEN: Ethernet PTP clock enable during Sleep mode

Set and cleared by software.

0: Ethernet PTP clock disabled during Sleep mode

1: Ethernet PTP clock enabled during Sleep mode

Bit 27 ETHMACRXLPEN: Ethernet reception clock enable during Sleep mode

Set and cleared by software.

0: Ethernet reception clock disabled during Sleep mode

1: Ethernet reception clock enabled during Sleep mode

Bit 26 ETHMACTXLPEN: Ethernet transmission clock enable during Sleep mode

Set and cleared by software.

0: Ethernet transmission clock disabled during sleep mode

1: Ethernet transmission clock enabled during sleep mode

Bit 25 ETHMACLPEN: Ethernet MAC clock enable during Sleep mode

Set and cleared by software.

0: Ethernet MAC clock disabled during Sleep mode

1: Ethernet MAC clock enabled during Sleep mode

Bits 24:23 Reserved, must be kept at reset value.

Bit 22 DMA2LPEN: DMA2 clock enable during Sleep mode

Set and cleared by software.

0: DMA2 clock disabled during Sleep mode

1: DMA2 clock enabled during Sleep mode

Bit 21 DMA1LPEN: DMA1 clock enable during Sleep mode

Set and cleared by software.

- 0: DMA1 clock disabled during Sleep mode
- 1: DMA1 clock enabled during Sleep mode
- Bits 20:19 Reserved, must be kept at reset value.
  - Bit 18 BKPSRAMLPEN: Backup SRAM interface clock enable during Sleep mode

Set and cleared by software.

- 0: Backup SRAM interface clock disabled during Sleep mode
- 1: Backup SRAM interface clock enabled during Sleep mode
- Bit 17 SRAM2LPEN: SRAM 2 interface clock enable during Sleep mode

Set and cleared by software.

- 0: SRAM 2 interface clock disabled during Sleep mode
- 1: SRAM 2 interface clock enabled during Sleep mode
- Bit 16 SRAM1LPEN: SRAM 1interface clock enable during Sleep mode

Set and cleared by software.

- 0: SRAM 1 interface clock disabled during Sleep mode
- 1: SRAM 1 interface clock enabled during Sleep mode
- Bit 15 FLITFLPEN: Flash interface clock enable during Sleep mode

Set and cleared by software.

- 0: Flash interface clock disabled during Sleep mode
- 1: Flash interface clock enabled during Sleep mode
- Bits 14:13 Reserved, must be kept at reset value.
  - Bit 12 CRCLPEN: CRC clock enable during Sleep mode

Set and cleared by software.

- 0: CRC clock disabled during Sleep mode
- 1: CRC clock enabled during Sleep mode
- Bits 11:9 Reserved, must be kept at reset value.
  - Bit 8 GPIOILPEN: IO port I clock enable during Sleep mode

Set and cleared by software.

- 0: IO port I clock disabled during Sleep mode
- 1: IO port I clock enabled during Sleep mode
- Bit 7 GPIOHLPEN: IO port H clock enable during Sleep mode

Set and cleared by software.

- 0: IO port H clock disabled during Sleep mode
- 1: IO port H clock enabled during Sleep mode
- Bits 6 GPIOGLPEN: IO port G clock enable during Sleep mode

Set and cleared by software.

- 0: IO port G clock disabled during Sleep mode
- 1: IO port G clock enabled during Sleep mode
- Bit 5 GPIOFLPEN: IO port F clock enable during Sleep mode

Set and cleared by software.

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- 0: IO port F clock disabled during Sleep mode
- 1: IO port F clock enabled during Sleep mode

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Bit 4 GPIOELPEN: IO port E clock enable during Sleep mode

Set and cleared by software.

0: IO port E clock disabled during Sleep mode

1: IO port E clock enabled during Sleep mode

Bit 3 GPIODLPEN: IO port D clock enable during Sleep mode

Set and cleared by software.

0: IO port D clock disabled during Sleep mode

1: IO port D clock enabled during Sleep mode

Bit 2 GPIOCLPEN: IO port C clock enable during Sleep mode

Set and cleared by software.

0: IO port C clock disabled during Sleep mode

1: IO port C clock enabled during Sleep mode

Bit 1 GPIOBLPEN: IO port B clock enable during Sleep mode

Set and cleared by software.

0: IO port B clock disabled during Sleep mode

1: IO port B clock enabled during Sleep mode

Bit 0 GPIOALPEN: IO port A clock enable during sleep mode

Set and cleared by software.

0: IO port A clock disabled during Sleep mode

1: IO port A clock enabled during Sleep mode

# 5.3.17 RCC AHB2 peripheral clock enable in low power mode register (RCC\_AHB2LPENR)

Address offset: 0x54

Reset value: 0x0000 00F1

Access: no wait state, word, half-word and byte access.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Res	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			Res	erved				OTGFS LPEN	RNG LPEN	HASH LPEN	CRYP LPEN		Reserved	i	DCMI LPEN
								rw	rw	rw	rw				rw

Bits 31:8 Reserved, must be kept at reset value.

Bit 7 OTGFSLPEN: USB OTG FS clock enable during Sleep mode

Set and cleared by software.

0: USB OTG FS clock disabled during Sleep mode

1: USB OTG FS clock enabled during Sleep mode

Bit 6 RNGLPEN: Random number generator clock enable during Sleep mode

Set and cleared by software.

0: Random number generator clock disabled during Sleep mode

1: Random number generator clock enabled during Sleep mode

Bit 5 HASHLPEN: Hash modules clock enable during Sleep mode

Set and cleared by software.

0: Hash modules clock disabled during Sleep mode

1: Hash modules clock enabled during Sleep mode

Bit 4 CRYPLPEN: Cryptography modules clock enable during Sleep mode

Set and cleared by software.

0: cryptography modules clock disabled during Sleep mode

1: cryptography modules clock enabled during Sleep mode

Bit 3:1 Reserved, must be kept at reset value.

Bit 0 DCMILPEN: Camera interface enable during Sleep mode

Set and cleared by software.

0: Camera interface clock disabled during Sleep mode

1: Camera interface clock enabled during Sleep mode

# 5.3.18 RCC AHB3 peripheral clock enable in low power mode register (RCC\_AHB3LPENR)

Address offset: 0x58

Reset value: 0x0000 0001

Access: no wait state, word, half-word and byte access.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							Reserved	l							FSMC LPEN
															rw

Bits 31:1Reserved, must be kept at reset value.

FSMCLPEN: Flexible static memory controller module clock enable during Sleep mode

Set and cleared by software.

0: FSMC module clock disabled during Sleep mode

1: FSMC module clock enabled during Sleep mode

# 5.3.19 RCC APB1 peripheral clock enable in low power mode register (RCC\_APB1LPENR)

Address offset: 0x60

Reset value: 0x36FE C9FF

Access: no wait state, word, half-word and byte access.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
Rese	erved	DAC LPEN	PWR LPEN	RESER VED	CAN2 LPEN	CAN1 LPEN	Reser-	I2C3 LPEN	I2C2 LPEN	I2C1 LPEN	UART5 LPEN	UART4 LPEN	USART3 LPEN	USART2 LPEN	Reser- ved
		rw	rw		rw	rw	veu	rw	rw	rw	rw	rw	rw	rw	veu
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
15 SPI3 LPEN	14 SPI2 LPEN	13 Rese		11 WWDG LPEN		9 erved	8 TIM14 LPEN	7 TIM13 LPEN	6 TIM12 LPEN	5 TIM7 LPEN	4 TIM6 LPEN	3 TIM5 LPEN	2 TIM4 LPEN	1 TIM3 LPEN	0 TIM2 LPEN

Bits 31:30 Reserved, must be kept at reset value.

Bit 29 DACLPEN: DAC interface clock enable during Sleep mode

Set and cleared by software.

0: DAC interface clock disabled during Sleep mode

1: DAC interface clock enabled during Sleep mode

Bit 28 **PWRLPEN:** Power interface clock enable during Sleep mode

Set and cleared by software.

0: Power interface clock disabled during Sleep mode

1: Power interface clock enabled during Sleep mode

Bit 27 Reserved, must be kept at reset value.

Bit 26 CAN2LPEN: CAN 2 clock enable during Sleep mode

Set and cleared by software.

0: CAN 2 clock disabled during sleep mode

1: CAN 2 clock enabled during sleep mode

Bit 25 CAN1LPEN: CAN 1 clock enable during Sleep mode

Set and cleared by software.

0: CAN 1 clock disabled during Sleep mode

1: CAN 1 clock enabled during Sleep mode

Bit 24 Reserved, must be kept at reset value.

Bit 23 I2C3LPEN: I2C3 clock enable during Sleep mode

Set and cleared by software.

0: I2C3 clock disabled during Sleep mode

1: I2C3 clock enabled during Sleep mode

Bit 22 I2C2LPEN: I2C2 clock enable during Sleep mode

Set and cleared by software.

0: I2C2 clock disabled during Sleep mode

1: I2C2 clock enabled during Sleep mode

Bit 21 I2C1LPEN: I2C1 clock enable during Sleep mode

Set and cleared by software.

0: I2C1 clock disabled during Sleep mode

1: I2C1 clock enabled during Sleep mode

#### Bit 20 UART5LPEN: UART5 clock enable during Sleep mode

Set and cleared by software.

0: UART5 clock disabled during Sleep mode

1: UART5 clock enabled during Sleep mode

#### Bit 19 UART4LPEN: UART4 clock enable during Sleep mode

Set and cleared by software.

0: UART4 clock disabled during Sleep mode

1: UART4 clock enabled during Sleep mode

#### Bit 18 USART3LPEN: USART3 clock enable during Sleep mode

Set and cleared by software.

0: USART3 clock disabled during Sleep mode

1: USART3 clock enabled during Sleep mode

#### Bit 17 USART2LPEN: USART2 clock enable during Sleep mode

Set and cleared by software.

0: USART2 clock disabled during Sleep mode

1: USART2 clock enabled during Sleep mode

#### Bit 16 Reserved, must be kept at reset value.

#### Bit 15 SPI3LPEN: SPI3 clock enable during Sleep mode

Set and cleared by software.

0: SPI3 clock disabled during Sleep mode

1: SPI3 clock enabled during Sleep mode

#### Bit 14 SPI2LPEN: SPI2 clock enable during Sleep mode

Set and cleared by software.

0: SPI2 clock disabled during Sleep mode

1: SPI2 clock enabled during Sleep mode

## Bits 13:12 Reserved, must be kept at reset value.

## Bit 11 WWDGLPEN: Window watchdog clock enable during Sleep mode

Set and cleared by software.

0: Window watchdog clock disabled during sleep mode

1: Window watchdog clock enabled during sleep mode

### Bits 10:9 Reserved, must be kept at reset value.

### Bit 8 TIM14LPEN: TIM14 clock enable during Sleep mode

Set and cleared by software.

0: TIM14 clock disabled during Sleep mode

1: TIM14 clock enabled during Sleep mode

#### Bit 7 TIM13LPEN: TIM13 clock enable during Sleep mode

Set and cleared by software.

0: TIM13 clock disabled during Sleep mode

1: TIM13 clock enabled during Sleep mode

#### Bit 6 TIM12LPEN: TIM12 clock enable during Sleep mode

Set and cleared by software.

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0: TIM12 clock disabled during Sleep mode

1: TIM12 clock enabled during Sleep mode

Bit 5 TIM7LPEN: TIM7 clock enable during Sleep mode

Set and cleared by software.

0: TIM7 clock disabled during Sleep mode

1: TIM7 clock enabled during Sleep mode

Bit 4 TIM6LPEN: TIM6 clock enable during Sleep mode

Set and cleared by software.

0: TIM6 clock disabled during Sleep mode

1: TIM6 clock enabled during Sleep mode

Bit 3 TIM5LPEN: TIM5 clock enable during Sleep mode

Set and cleared by software.

0: TIM5 clock disabled during Sleep mode

1: TIM5 clock enabled during Sleep mode

Bit 2 TIM4LPEN: TIM4 clock enable during Sleep mode

Set and cleared by software.

0: TIM4 clock disabled during Sleep mode

1: TIM4 clock enabled during Sleep mode

Bit 1 TIM3LPEN: TIM3 clock enable during Sleep mode

Set and cleared by software.

0: TIM3 clock disabled during Sleep mode

1: TIM3 clock enabled during Sleep mode

Bit 0 **TIM2LPEN:** TIM2 clock enable during Sleep mode

Set and cleared by software.

0: TIM2 clock disabled during Sleep mode

1: TIM2 clock enabled during Sleep mode

# 5.3.20 RCC APB2 peripheral clock enabled in low power mode register (RCC\_APB2LPENR)

Address offset: 0x64

Reset value: 0x0007 5F33

Access: no wait state, word, half-word and byte access.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
						Reserve	ed						TIM11 LPEN	TIM10 LPEN	TIM9 LPEN
													rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Reser- ved	SYSC FG LPEN	Reser- ved	SPI1 LPEN	SDIO LPEN	ADC3 LPEN	ADC2 LPEN	ADC1 LPEN	Rese	rved	USART6 LPEN	USART1 LPEN	Rese	erved	TIM8 LPEN	TIM1 LPEN
	rw		rw	rw	rw	rw	rw			rw	rw			rw	rw

Bits 31:19 Reserved, must be kept at reset value.

Bit 18 TIM11LPEN: TIM11 clock enable during Sleep mode

Set and cleared by software.

0: TIM11 clock disabled during Sleep mode

1: TIM11 clock enabled during Sleep mode

Bit 17 TIM10LPEN: TIM10 clock enable during Sleep mode

Set and cleared by software.

0: TIM10 clock disabled during Sleep mode

1: TIM10 clock enabled during Sleep mode

Bit 16 TIM9LPEN: TIM9 clock enable during sleep mode

Set and cleared by software.

0: TIM9 clock disabled during Sleep mode

1: TIM9 clock enabled during Sleep mode

Bit 15 Reserved, must be kept at reset value.

Bit 14 SYSCFGLPEN: System configuration controller clock enable during Sleep mode

Set and cleared by software.

0: System configuration controller clock disabled during Sleep mode

1: System configuration controller clock enabled during Sleep mode

Bits 13 Reserved, must be kept at reset value.

Bit 12 SPI1LPEN: SPI1 clock enable during Sleep mode

Set and cleared by software.

0: SPI1 clock disabled during Sleep mode

1: SPI1 clock enabled during Sleep mode

Bit 11 SDIOLPEN: SDIO clock enable during Sleep mode

Set and cleared by software.

0: SDIO module clock disabled during Sleep mode

1: SDIO module clock enabled during Sleep mode

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Bit 10 ADC3LPEN: ADC 3 clock enable during Sleep mode

Set and cleared by software.

0: ADC 3 clock disabled during Sleep mode

1: ADC 3 clock disabled during Sleep mode

Bit 9 ADC2LPEN: ADC2 clock enable during Sleep mode

Set and cleared by software.

0: ADC2 clock disabled during Sleep mode

1: ADC2 clock disabled during Sleep mode

Bit 8 ADC1LPEN: ADC1 clock enable during Sleep mode

Set and cleared by software.

0: ADC1 clock disabled during Sleep mode

1: ADC1 clock disabled during Sleep mode

Bits 7:6 Reserved, must be kept at reset value.

Bit 5 USART6LPEN: USART6 clock enable during Sleep mode

Set and cleared by software.

0: USART6 clock disabled during Sleep mode

1: USART6 clock enabled during Sleep mode

Bit 4 USART1LPEN: USART1 clock enable during Sleep mode

Set and cleared by software.

0: USART1 clock disabled during Sleep mode

1: USART1 clock enabled during Sleep mode

Bits 3:2 Reserved, must be kept at reset value.

Bit 1 TIM8LPEN: TIM8 clock enable during Sleep mode

Set and cleared by software.

0: TIM8 clock disabled during Sleep mode

1: TIM8 clock enabled during Sleep mode

Bit 0 TIM1LPEN: TIM1 clock enable during Sleep mode

Set and cleared by software.

0: TIM1 clock disabled during Sleep mode

1: TIM1 clock enabled during Sleep mode

## 5.3.21 RCC Backup domain control register (RCC\_BDCR)

Address offset: 0x70

Reset value:  $0x0000\ 0000$ , reset by Backup domain reset. Access:  $0 \le \text{wait state} \le 3$ , word, half-word and byte access

Wait states are inserted in case of successive accesses to this register.

The LSEON, LSEBYP, RTCSEL and RTCEN bits in the *RCC Backup domain control register (RCC\_BDCR)* are in the Backup domain. As a result, after Reset, these bits are write-protected and the DBP bit in the *Power control register (PWR\_CR)* has to be set before these can be modified. Refer to *Section 5.1.2 on page 67* for further information. These bits are only reset after a Backup domain Reset (see *Section 5.1.3: Backup domain reset*). Any internal or external Reset will not have any effect on these bits.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							5 .								BDRST
							Reserved								rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
RTCEN			Reserved	ı		RTCS	EL[1:0]			Reserved	ı		LSEBYP	LSERDY	LSEON
rw			neserved	ı		rw	rw			neserved	ı		rw	r	rw

Bits 31:17 Reserved, must be kept at reset value.

Bit 16 BDRST: Backup domain software reset

Set and cleared by software.

0: Reset not activated

1: Resets the entire Backup domain

Note: The BKPSRAM is not affected by this reset, the only way of resetting the BKPSRAM is through the Flash interface when a protection level change from level 1 to level 0 is requested.

#### Bit 15 RTCEN: RTC clock enable

Set and cleared by software.

0: RTC clock disabled

1: RTC clock enabled

Bits 14:10 Reserved, must be kept at reset value.

#### Bits 9:8 RTCSEL[1:0]: RTC clock source selection

Set by software to select the clock source for the RTC. Once the RTC clock source has been selected, it cannot be changed anymore unless the Backup domain is reset. The BDRST bit can be used to reset them.

00: No clock

01: LSE oscillator clock used as the RTC clock

10: LSI oscillator clock used as the RTC clock

11: HSE oscillator clock divided by a programmable prescaler (selection through the RTCPRE[4:0] bits in the RCC clock configuration register (RCC\_CFGR)) used as the RTC clock

Bits 7:3 Reserved, must be kept at reset value.

### Bit 2 LSEBYP: External low-speed oscillator bypass

Set and cleared by software to bypass oscillator in debug mode. This bit can be written only when the LSE clock is disabled.

0: LSE oscillator not bypassed

1: LSE oscillator bypassed

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#### Bit 1 LSERDY: External low-speed oscillator ready

Set and cleared by hardware to indicate when the external 32 kHz oscillator is stable. After the LSEON bit is cleared, LSERDY goes low after 6 external low-speed oscillator clock cycles.

0: LSE clock not ready 1: LSE clock ready

#### Bit 0 LSEON: External low-speed oscillator enable

Set and cleared by software.

0: LSE clock OFF 1: LSE clock ON

## 5.3.22 RCC clock control & status register (RCC\_CSR)

Address offset: 0x74

Reset value: 0x0E00 0000, reset by system reset, except reset flags by power reset only.

Access:  $0 \le \text{ wait state} \le 3$ , word, half-word and byte access

Wait states are inserted in case of successive accesses to this register.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
LPWR RSTF	WWDG RSTF	IWDG RSTF	SFT RSTF	POR RSTF	PIN RSTF	BORRS TF	RMVF				Res				
rw	rw	rw	rw	rw	rw	rw	rw								
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
						Dana	a							LSIRDY	LSION
	Reserved														rw

#### Bit 31 LPWRRSTF: Low-power reset flag

Set by hardware when a Low-power management reset occurs.

Cleared by writing to the RMVF bit.

0: No Low-power management reset occurred

1: Low-power management reset occurred

For further information on Low-power management reset, refer to *Low-power management* reset

#### Bit 30 WWDGRSTF: Window watchdog reset flag

Set by hardware when a window watchdog reset occurs.

Cleared by writing to the RMVF bit.

0: No window watchdog reset occurred

1: Window watchdog reset occurred

#### Bit 29 IWDGRSTF: Independent watchdog reset flag

Set by hardware when an independent watchdog reset from  $V_{DD}$  domain occurs.

Cleared by writing to the RMVF bit.

0: No watchdog reset occurred

1: Watchdog reset occurred

#### Bit 28 SFTRSTF: Software reset flag

Set by hardware when a software reset occurs.

Cleared by writing to the RMVF bit.

0: No software reset occurred

1: Software reset occurred

#### Bit 27 PORRSTF: POR/PDR reset flag

Set by hardware when a POR/PDR reset occurs.

Cleared by writing to the RMVF bit.

0: No POR/PDR reset occurred

1: POR/PDR reset occurred

#### Bit 26 PINRSTF: PIN reset flag

Set by hardware when a reset from the NRST pin occurs.

Cleared by writing to the RMVF bit.

0: No reset from NRST pin occurred

1: Reset from NRST pin occurred

#### Bit 25 BORRSTF: BOR reset flag

Cleared by software by writing the RMVF bit.

Set by hardware when a POR/PDR or BOR reset occurs.

0: No POR/PDR or BOR reset occurred

1: POR/PDR or BOR reset occurred

#### Bit 24 RMVF: Remove reset flag

Set by software to clear the reset flags.

0: No effect

1: Clear the reset flags

## Bits 23:2 Reserved, must be kept at reset value.

#### Bit 1 LSIRDY: Internal low-speed oscillator ready

Set and cleared by hardware to indicate when the internal RC 40 kHz oscillator is stable. After the LSION bit is cleared, LSIRDY goes low after 3 LSI clock cycles.

0: LSI RC oscillator not ready

1: LSI RC oscillator ready

#### Bit 0 LSION: Internal low-speed oscillator enable

Set and cleared by software.

0: LSI RC oscillator OFF

1: LSI RC oscillator ON

## 5.3.23 RCC spread spectrum clock generation register (RCC\_SSCGR)

Address offset: 0x80

Reset value: 0x0000 0000

Access: no wait state, word, half-word and byte access.

The spread spectrum clock generation is available only for the main PLL.

The RCC\_SSCGR register must be written either before the main PLL is enabled or after

the main PLL disabled.

Note: For full details about PLL spread spectrum clock generation (SSCG) characteristics, refer to

the "Electrical characteristics" section in your device datasheet.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
SSCG EN	SPR EAD SEL	Rese	erved						INCS	STEP					
rw	rw			rw	rw	rw		rw	rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	INCSTEF	)							MODPEF	}					
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31 SSCGEN: Spread spectrum modulation enable

Set and cleared by software.

- 0: Spread spectrum modulation DISABLE. (To write after clearing CR[24]=PLLON bit)
- 1: Spread spectrum modulation ENABLE. (To write before setting CR[24]=PLLON bit)

#### Bit 30 SPREADSEL: Spread Select

Set and cleared by software.

To write before to set CR[24]=PLLON bit.

0: Center spread

1: Down spread

Bit 29:28 Reserved, must be kept at reset value.

#### Bit 27:13 INCSTEP: Incrementation step

Set and cleared by software. To write before setting CR[24]=PLLON bit.

Configuration input for modulation profile amplitude.

### Bit 12:0 MODPER: Modulation period

Set and cleared by software. To write before setting CR[24]=PLLON bit.

Configuration input for modulation profile period.

## 5.3.24 RCC PLLI2S configuration register (RCC\_PLLI2SCFGR)

Address offset: 0x84

Reset value: 0x2000 3000

Access: no wait state, word, half-word and byte access.

This register is used to configure the PLLI2S clock outputs according to the formulas:

f<sub>(VCO clock)</sub> = f<sub>(PLLI2S clock input)</sub> × (PLLI2SN / PLLM)

• f<sub>(PLL I2S clock output)</sub> = f<sub>(VCO clock)</sub> / PLLI2SR

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16		
Reserv	PLLI2S R2	PLLI2S R1	PLLI2S R0					Rese	erved								
eu	rw	rw	rw														
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0		
Reserv	PLLI2SN 8	PLLI2SN 7	PLLI2SN 6	PLLI2SN 5	PLLI2SN 4	PLLI2SN 3	PLLI2SN 2	PLLI2SN 1	PLLI2SN 0	Reserved							
ed	rw																

Bit 31 Reserved, must be kept at reset value.

Bits 30:28 PLLI2SR: PLLI2S division factor for I2S clocks

Set and cleared by software to control the I2S clock frequency. These bits should be written only if the PLLI2S is disabled. The factor must be chosen in accordance with the prescaler values inside the I2S peripherals, to reach 0.3% error when using standard crystals and 0% error with audio crystals. For more information about I2S clock frequency and precision, refer to *Section 25.4.4: Clock generator* in the I2S chapter.

**Caution:** The I2Ss requires a frequency lower than or equal to 192 MHz to work correctly. I2S clock frequency = VCO frequency / PLLR with  $2 \le PLLR \le 7$ 

000: PLLR = 0, wrong configuration 001: PLLR = 1, wrong configuration

010: PLLR = 2

111: PLLR = 7

Bits 27:15 Reserved, must be kept at reset value.

Bits 14:6 PLLI2SN: PLLI2S multiplication factor for VCO

Set and cleared by software to control the multiplication factor of the VCO. These bits can be written only when the PLLI2S is disabled. Only half-word and word accesses are allowed to write these bits.

**Caution:** The software has to set these bits correctly to ensure that the VCO output frequency is between 192 and 432 MHz.

```
VCO output frequency = VCO input frequency × PLLI2SN with 192 ≤ PLLI2SN ≤ 432 000000000: PLLI2SN = 0, wrong configuration 000000001: PLLI2SN = 1, wrong configuration ...
01100000: PLLI2SN = 192 011000001: PLLI2SN = 193 011000010: PLLI2SN = 194 ...
110110000: PLLI2SN = 432 110110000: PLLI2SN = 433, wrong configuration ...
111111111: PLLI2SN = 511, wrong configuration
```

Bits 5:0 Reserved, must be kept at reset value.

## 5.3.25 RCC register map

Table 15 gives the register map and reset values.

Table 15. RCC register map and reset values

Addr. offset	Register name	31	30	29	28	27	56	25	24	23	22	21	20	19	18	17	16	15	14	13	12	1	10	6	8	7	9	2	4	ဗ	2	1	0
0x00	RCC_CR	F	Rese	erve	d	PLL I2SRDY	PLL I2SON	PLL RDY	PLL ON	F	Rese	erve	d	CSSON	HSEBYP	HSERDY	HSEON	HSICAL 7	HSICAL 6	HSICAL 5	HSICAL 4	HSICAL 3	HSICAL 2	HSICAL 1	HSICAL 0	HSITRIM 4	HSITRIM 3	HSITRIM 2	HSITRIM 1	HSITRIM 0	Reserved	HSIRDY	NOISH
0x04	RCC_PLLCF GR	F	Rese	erve	d	PLLQ 3		PLLQ 1	0 DTTd	Reserved	PLLSRC	F	Rese	erve	d	PLLP 1	DLLP 0	Reserved	PLLN 8	PLLN 7	9 NTTA	PLLN 5	PLLN 4	E NTTA	PLLN 2	PLLN 1	PLLN 0	PLLM 5	PLLM 4	PLLM 3	2 MJJA	PLLM 1	PLLM 0
0x08	RCC_CFGR	MCO2 1	MCO2 0	MCO2PRE2	MCO2PRE1	MCO2PRE0	MCO1PRE2	MCO1PRE1	MCO1PRE0	12SSRC	MCO11	MCO1 0	RTCPRE 4	RTCPRE 3	RTCPRE 2	RTCPRE 1	RTCPRE 0	PPRE2 2	PPRE2 1	PPRE20	PPRE12	PPRE1 1	PPRE1 0	Becomed	0000	HPRE 3	HPRE 2	HPRE 1	HPRE 0	SWS 1	0 SMS	SW 1	0 MS
0x0C	RCC_CIR			F	Rese	erve	d			CSSC	Reserved	PLLI2SRDYC	PLLRDYC	HSERDYC	HSIRDYC	LSERDYC	LSIRDYC	poracoo	00000	PLLI2SRDYIE	PLLRDYIE	HSERDYIE	HSIRDYIE	LSERDYIE	LSIRDYIE	CSSF	Reserved	PLLI2SRDYF	PLLRDYF	HSERDYF	HSIRDYF	LSERDYF	LSIRDYF
0x10	RCC_AHB1R STR	Beserved		OTGHSRST		Reserved		ETHMACRST	Beenwed	000	DMA2RST	DMA1RST			F	Rese	erve	d			CRCRST		Reserved		GPIOIRST	GPIOHRST	GPIOGRST	GPIOFRST	GPIOERST	GPIODRST	GPIOCRST	GPIOBRST	GPIOARST
0x14	RCC_AHB2R STR		OTGFSRST RNGRST HSAHRST CRYPRST CRYPRST CRYPRST CRYPRST																														
0x18	RCC_AHB3R STR		Reserved Fig. 12 O C C C C C C C C C C C C C C C C C C																														
0x1C	Reserved															F	Rese	erve	d														
0x20	RCC_APB1R STR	Beserved		DACRST	PWRRST	Reserved	CAN2RST	CAN1RST	Reserved	12C3RST	I2C2RST	12C1RST	UART5RST	UART4RST	UART3RST	UART2RST	Reserved	SPI3RST	SPI2RST	Beconved		WWDGRST	Beserved		TIM14RST	TIM13RST	TIM12RST	TIM7RST	TIMGRST	TIM5RST	TIM4RST	TIM3RST	TIM2RST
0x24	RCC_APB2R STR						Re	serv	/ed						TIM11RST	TIM10RST	TIM9RST	Reserved	SYSCFGRST	Reserved	SPI1RST	SDIORST	Beserved		ADCRST	D0,0000	חמא ושפשרו	USARTERST	USART1RST	Doynosod		TIMBRST	TIM1RST
0x28	Reserved															F	Rese	erve	d														
0x2C	Reserved															F	Rese	erve	d														
0x30	RCC_AHB1E NR		OTGHSULPIEN	OTGHSEN	ETHMACPTPEN	ETHMACRXEN	ETHMACTXEN	ETHMACEN	Beserved		DMA2EN	DMA1EN	CCMDATARAMEN	Reserved	BKPSRAMEN			Reserved			CRCEN		Reserved		GPIOIEN	GPIOHEN	GPIOGEN	GPIOFEN	GPIOEEN	GPIODEN	GPIOCEN	GPIOBEN	GPIOAEN
0x34	RCC_AHB2E NR		OTGFSEN  Reserved  CRYPEN  DCMIEN																														
0x38	RCC_AHB3E NR		Reserved Sun L																														
0x3C	Reserved															F	Rese	erve	d														

Table 15. RCC register map and reset values (continued)

Table	J. NCC		<b>J</b>								_			· (,			iiu		,													1	
Addr. offset	Register name	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	6	8	2	9	2	4	3	7	-	0
0x40	RCC_APB1E NR	Reserved		DACEN	PWREN	Reserved	CANZEN	CAN1EN	Reserved	I2C3EN	IZCZEN	12C1 EN	<b>UART5EN</b>	UART4EN	<b>USART3EN</b>	<b>USART2EN</b>	Reserved	SPI3EN	SPIZEN	Reserved		WWDGEN	Beserved		TIM14EN	TIM13EN	TIM12EN	TIM7EN	TIM6EN	TIMSEN	TIM4EN	TIM3EN	TIM2EN
0x44	RCC_APB2E NR						Re	serve	ed						TIM11EN	TIM10EN	TIM9EN	Reserved	S	Reserved	SPI1EN	SDIOEN	ADC3EN	ADC2EN	ADC1EN	Boong	000000000000000000000000000000000000000	USART6EN	USART1EN	poracional	חפאפט	TIMBEN	TIM1EN
0x48	Reserved															F	Rese	erve	d														
0x4C	Reserved															F	Rese	erve	d														
0x50	RCC_AHB1L PENR	Reserved											GPIOCLPEN	GPIOBLPEN	GPIOALPEN																		
0x54	RCC_AHB2L PENR															DCMILPEN																	
0x58	RCC_AHB3L PENR		Reserved Reserved														FSMCLPEN																
0x5C	Reserved															F	Rese	erve	d														
0x60	RCC_APB1L PENR	Reserved		DACLPEN	PWRLPEN	Reserved	CANZLPEN	CAN1LPEN	Reserved	I2C3LPEN	IZCZLPEN	12C1LPEN	UARTSLPEN	UART4LPEN	<b>USART3LPEN</b>	USART2LPEN	Reserved	SPI3LPEN	SPI2LPEN	Reserved		WWDGLPEN	Reserved	10001	TIM14LPEN	TIM13LPEN	TIM12LPEN	TIM7LPEN	TIM6LPEN	TIMSLPEN	TIM4LPEN	TIM3LPEN	TIM2LPEN
0x64	RCC_APB2L PENR						Re	serve	ed						TIM11LPEN	TIM10LPEN	TIM9LPEN	Reserved	SYSCFGLPEN	Reserved	SPI1LPEN	SDIOLPEN	<b>ADC3LPEN</b>	ADC2LPEN	ADC1LPEN	Beserved	000000	USART6LPEN	USART1LPEN	Postagoad	חפסם ומפסבו	TIM8LPEN	TIM1LPEN
0x68	Reserved															F	Rese	erve	d														
0x6C	Reserved															F	Rese	erve	d														
0x70	RCC_BDCR							Res	erv	ed .							BDRST	RTCEN		Re	serv	ed		RTCSEL 1	RTCSEL 0		Re	ser	ved		LSEBYP	LSERDY	LSEON
0x74	RCC_CSR												NOIST																				
0x78	Reserved																Rese																
0x7C	Reserved															F	Rese	erve	d														
0x80	RCC_SSCGR	SSC SPRE																															
0x84	RCC_PLLI2S CFGR	Reserved													F	Rese	erve	d															

Refer to *Table 2 on page 50* for the register boundary addresses.

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## 6 General-purpose I/Os (GPIO)

This section applies to the whole STM32F4xx family, unless otherwise specified.

## 6.1 GPIO introduction

Each general-purpose I/O port has four 32-bit configuration registers (GPIOx\_MODER, GPIOx\_OTYPER, GPIOx\_OSPEEDR and GPIOx\_PUPDR), two 32-bit data registers (GPIOx\_IDR and GPIOx\_ODR), a 32-bit set/reset register (GPIOx\_BSRR), a 32-bit locking register (GPIOx\_LCKR) and two 32-bit alternate function selection register (GPIOx\_AFRH and GPIOx\_AFRL).

## 6.2 GPIO main features

- Up to 16 I/Os under control
- Output states: push-pull or open drain + pull-up/down
- Output data from output data register (GPIOx\_ODR) or peripheral (alternate function output)
- Speed selection for each I/O
- Input states: floating, pull-up/down, analog
- Input data to input data register (GPIOx\_IDR) or peripheral (alternate function input)
- Bit set and reset register (GPIOx\_BSRR) for bitwise write access to GPIOx\_ODR
- Locking mechanism (GPIOx\_LCKR) provided to freeze the I/O configuration
- Analog function
- Alternate function input/output selection registers (at most 16 AFs per I/O)
- Fast toggle capable of changing every two clock cycles
- Highly flexible pin multiplexing allows the use of I/O pins as GPIOs or as one of several peripheral functions

## 6.3 GPIO functional description

Subject to the specific hardware characteristics of each I/O port listed in the datasheet, each port bit of the general-purpose I/O (GPIO) ports can be individually configured by software in several modes:

- Input floating
- Input pull-up
- Input-pull-down
- Analog
- Output open-drain with pull-up or pull-down capability
- Output push-pull with pull-up or pull-down capability
- Alternate function push-pull with pull-up or pull-down capability
- Alternate function open-drain with pull-up or pull-down capability



Each I/O port bit is freely programmable, however the I/O port registers have to be accessed as 32-bit words, half-words or bytes. The purpose of the GPIOx\_BSRR register is to allow atomic read/modify accesses to any of the GPIO registers. In this way, there is no risk of an IRQ occurring between the read and the modify access.

*Figure 13* shows the basic structure of a 5 V tolerant I/O port bit. *Table 20* gives the possible port bit configurations.

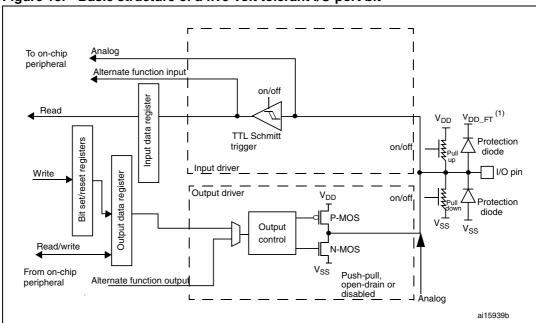


Figure 13. Basic structure of a five-volt tolerant I/O port bit

Table 16. Port bit configuration table<sup>(1)</sup>

MODER(i) [1:0]	OTYPER(i)	OSPEEDR(i) [B:A]		DR(i) :0]	I/O conf	iguration
	0		0	0	GP output	PP
	0		0	1	GP output	PP + PU
	0		1	0	GP output	PP + PD
01	0	SPEED	1	1	Reserved	
01	1	[B:A]	0	0	GP output	OD
	1		0	1	GP output	OD + PU
	1		1	0	GP output	OD + PD
	1		1	1	Reserved (GP ou	itput OD)

<sup>1.</sup>  $V_{DD\_FT}$  is a potential specific to five-volt tolerant I/Os and different from  $V_{DD}$ .

MODER(i) [1:0]	OTYPER(i)		EEDR(i) B:A]		DR(i) :0]	I/O con	figuration
	0			0	0	AF	PP
	0			0	1	AF	PP + PU
	0			1	0	AF	PP + PD
10	0	SP	EED	1	1	Reserved	
10	1	[E	3:A]	0	0	AF	OD
	1			0	1	AF	OD + PU
	1			1	0	AF	OD + PD
	1			1	1	Reserved	
	Х	х	Х	0	0	Input	Floating
00	Х	х	х	0	1	Input	PU
00	Х	х	Х	1	0	Input	PD
	х	х	х	1	1	Reserved (input	floating)
	Х	х	Х	0	0	Input/output	Analog
11	Х	Х	Х	0	1		
	X	Х	Х	1	0	Reserved	
	X	х	х	1	1		

Table 16. Port bit configuration table<sup>(1)</sup> (continued)

## 6.3.1 General-purpose I/O (GPIO)

During and just after reset, the alternate functions are not active and the I/O ports are configured in input floating mode.

The debug pins are in AF pull-up/pull-down after reset:

- PA15: JTDI in pull-up
- PA14: JTCK/SWCLK in pull-down
- PA13: JTMS/SWDAT in pull-up
- PB4: NJTRST in pull-up
- PB3: JTDO in floating state

When the pin is configured as output, the value written to the output data register (GPIOx\_ODR) is output on the I/O pin. It is possible to use the output driver in push-pull mode or open-drain mode (only the N-MOS is activated when 0 is output).

The input data register (GPIOx\_IDR) captures the data present on the I/O pin at every AHB1 clock cycle.

All GPIO pins have weak internal pull-up and pull-down resistors, which can be activated or not depending on the value in the GPIOx\_PUPDR register.

<sup>1.</sup> GP = general-purpose, PP = push-pull, PU = pull-up, PD = pull-down, OD = open-drain, AF = alternate function

## 6.3.2 I/O pin multiplexer and mapping

The microcontroller I/O pins are connected to onboard peripherals/modules through a multiplexer that allows only one peripheral's alternate function (AF) connected to an I/O pin at a time. In this way, there can be no conflict between peripherals sharing the same I/O pin.

Each I/O pin has a multiplexer with sixteen alternate function inputs (AF0 to AF15) that can be configured through the GPIOx\_AFRL (for pin 0 to 7) and GPIOx\_AFRH (for pin 8 to 15) registers:

- After reset all I/Os are connected to the system's alternate function 0 (AF0)
- The peripherals' alternate functions are mapped from AF1 to AF13
- Cortex<sup>™</sup>-M4F EVENTOUT is mapped on AF15

This structure is shown in Figure 14 below.

In addition to this flexible I/O multiplexing architecture, each peripheral has alternate functions mapped onto different I/O pins to optimize the number of peripherals available in smaller packages.

To use an I/O in a given configuration, proceed as follows:

#### 1. System function

Connect the I/O to AF0 and configure it depending on the function used:

- JTAG/SWD, after each device reset these pins are assigned as dedicated pins immediately usable by the debugger host (not controlled by the GPIO controller)
- RTC\_50Hz: this pin should be configured in Input floating mode
- MCO1 and MCO2: these pins have to be configured in alternate function mode.

Note:

You can disable some or all of the JTAG/SWD pins and so release the associated pins for GPIO usage.

For more details please refer to Section 5.2.10: Clock-out capability.

SWJ I/O pin assigned PA13/ PA14/ Available debug ports PA15/ **PB3** / PB4/ JTMS/ JTCK/ JTDI **JTDO NJTRST SWDIO SWCLK** Full SWJ (JTAG-DP + SW-DP) - Reset state Х Χ Χ Χ Χ Full SWJ (JTAG-DP + SW-DP) but without Χ Χ Χ Χ **NJTRST** JTAG-DP Disabled and SW-DP Enabled Х Χ JTAG-DP Disabled and SW-DP Disabled Released

Table 17. Flexible SWJ-DP pin assignment

#### 2. GPIO

Configure the desired I/O as output or input in the GPIOx\_MODER register.

#### 3. Peripheral alternate function

For the ADC and DAC, configure the desired I/O as analog in the GPIOx\_MODER register.

For other peripherals:

- Configure the desired I/O as an alternate function in the GPIOx\_MODER register
- Select the type, pull-up/pull-down and output speed via the GPIOx\_OTYPER, GPIOx\_PUPDR and GPIOx\_OSPEEDER registers, respectively
- Connect the I/O to the desired AFx in the GPIOx\_AFRL or GPIOx\_AFRH register

#### 4. EVENTOUT

Configure the I/O pin used to output the Cortex<sup>™</sup>-M4F EVENTOUT signal by connecting it to AF15

Note: EVENTOUT is not mapped onto the following I/O pins: PC13, PC14, PC15, PH0, PH1 and PI8.

Please refer to the "Alternate function mapping" table in the datasheets for the detailed mapping of the system and peripherals' alternate function I/O pins.

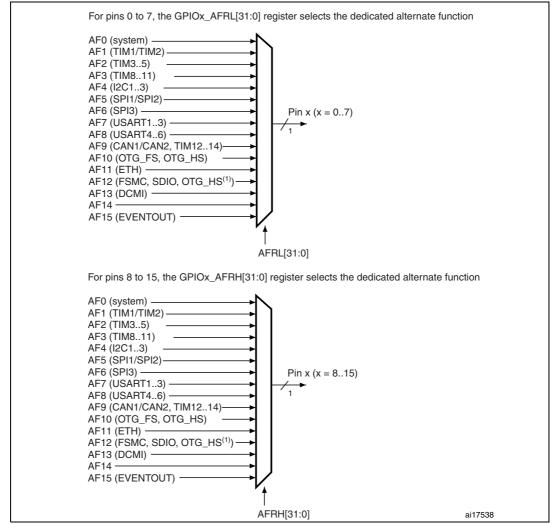


Figure 14. Selecting an alternate function

1. Configured in FS.

## 6.3.3 I/O port control registers

Each of the GPIOs has four 32-bit memory-mapped control registers (GPIOx\_MODER, GPIOx\_OTYPER, GPIOx\_OSPEEDR, GPIOx\_PUPDR) to configure up to 16 I/Os. The GPIOx\_MODER register is used to select the I/O direction (input, output, AF, analog). The GPIOx\_OTYPER and GPIOx\_OSPEEDR registers are used to select the output type (pushpull or open-drain) and speed (the I/O speed pins are directly connected to the corresponding GPIOx\_OSPEEDR register bits whatever the I/O direction). The GPIOx\_PUPDR register is used to select the pull-up/pull-down whatever the I/O direction.

## 6.3.4 I/O port data registers

Each GPIO has two 16-bit memory-mapped data registers: input and output data registers (GPIOx\_IDR and GPIOx\_ODR). GPIOx\_ODR stores the data to be output, it is read/write accessible. The data input through the I/O are stored into the input data register (GPIOx\_IDR), a read-only register.

See Section 6.4.5: GPIO port input data register (GPIOx\_IDR) (x = A..I) and Section 6.4.6: GPIO port output data register (GPIOx\_ODR) (x = A..I) for the register descriptions.

### 6.3.5 I/O data bitwise handling

The bit set reset register (GPIOx\_BSRR) is a 32-bit register which allows the application to set and reset each individual bit in the output data register (GPIOx\_ODR). The bit set reset register has twice the size of GPIOx\_ODR.

To each bit in GPIOx\_ODR, correspond two control bits in GPIOx\_BSRR: BSRR(i) and BSRR(i+SIZE). When written to 1, bit BSRR(i) **sets** the corresponding ODR(i) bit. When written to 1, bit BSRR(i+SIZE) **resets** the ODR(i) corresponding bit.

Writing any bit to 0 in GPIOx\_BSRR does not have any effect on the corresponding bit in GPIOx\_ODR. If there is an attempt to both set and reset a bit in GPIOx\_BSRR, the set action takes priority.

Using the GPIOx\_BSRR register to change the values of individual bits in GPIOx\_ODR is a "one-shot" effect that does not lock the GPIOx\_ODR bits. The GPIOx\_ODR bits can always be accessed directly. The GPIOx\_BSRR register provides a way of performing atomic bitwise handling.

There is no need for the software to disable interrupts when programming the GPIOx\_ODR at bit level: it is possible to modify one or more bits in a single atomic AHB1 write access.

### 6.3.6 GPIO locking mechanism

It is possible to freeze the GPIO control registers by applying a specific write sequence to the GPIOx\_LCKR register. The frozen registers are GPIOx\_MODER, GPIOx\_OTYPER, GPIOx\_OSPEEDR, GPIOx\_PUPDR, GPIOx\_AFRL and GPIOx\_AFRH.

To write the GPIOx\_LCKR register, a specific write / read sequence has to be applied. When the right LOCK sequence is applied to bit 16 in this register, the value of LCKR[15:0] is used to lock the configuration of the I/Os (during the write sequence the LCKR[15:0] value must be the same). When the LOCK sequence has been applied to a port bit, the value of the port bit can no longer be modified until the next reset. Each GPIOx\_LCKR bit freezes the corresponding bit in the control registers (GPIOx\_MODER, GPIOx\_OTYPER, GPIOx\_OSPEEDR, GPIOx\_PUPDR, GPIOx\_AFRL and GPIOx\_AFRH).

The LOCK sequence (refer to Section 6.4.8: GPIO port configuration lock register (GPIOx\_LCKR) (x = A..I)) can only be performed using a word (32-bit long) access to the GPIOx\_LCKR register due to the fact that GPIOx\_LCKR bit 16 has to be set at the same time as the [15:0] bits.

For more details please refer to LCKR register description in Section 6.4.8: GPIO port configuration lock register (GPIOx\_LCKR) (x = A..I).

### 6.3.7 I/O alternate function input/output

Two registers are provided to select one out of the sixteen alternate function inputs/outputs available for each I/O. With these registers, you can connect an alternate function to some other pin as required by your application.

This means that a number of possible peripheral functions are multiplexed on each GPIO using the GPIOx\_AFRL and GPIOx\_AFRH alternate function registers. The application can thus select any one of the possible functions for each I/O. The AF selection signal being

common to the alternate function input and alternate function output, a single channel is selected for the alternate function input/output of one I/O.

To know which functions are multiplexed on each GPIO pin, refer to the datasheets.

Note:

The application is allowed to select one of the possible peripheral functions for each I/O at a time.

### 6.3.8 External interrupt/wakeup lines

All ports have external interrupt capability. To use external interrupt lines, the port must be configured in input mode, refer to Section 9.2: External interrupt/event controller (EXTI) and Section 9.2.3: Wakeup event management.

### 6.3.9 Input configuration

When the I/O port is programmed as Input:

- the output buffer is disabled
- the Schmitt trigger input is activated
- the pull-up and pull-down resistors are activated depending on the value in the GPIOx\_PUPDR register
- The data present on the I/O pin are sampled into the input data register every AHB1 clock cycle
- A read access to the input data register provides the I/O State

Figure 15 shows the input configuration of the I/O port bit.

Read Input data  $V_{DD}$ Bit set/reset registers on/off TTL Schmitt protection trigge diode Write up Output data register I/O pin output drive protection ٧ss VSS Read/write ai15940b

Figure 15. Input floating/pull up/pull down configurations

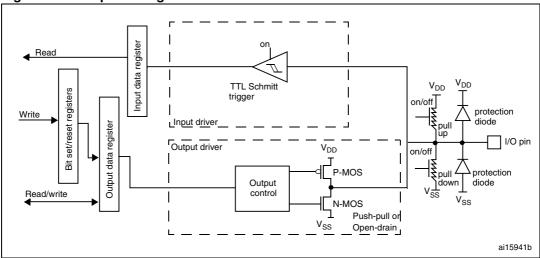
### 6.3.10 Output configuration

When the I/O port is programmed as output:

- The output buffer is enabled:
  - Open drain mode: A "0" in the Output register activates the N-MOS whereas a "1" in the Output register leaves the port in Hi-Z (the P-MOS is never activated)
  - Push-pull mode: A "0" in the Output register activates the N-MOS whereas a "1" in the Output register activates the P-MOS
- The Schmitt trigger input is activated
- The weak pull-up and pull-down resistors are activated or not depending on the value in the GPIOx\_PUPDR register
- The data present on the I/O pin are sampled into the input data register every AHB1 clock cycle
- A read access to the input data register gets the I/O state
- A read access to the output data register gets the last written value

Figure 16 shows the output configuration of the I/O port bit.

Figure 16. Output configuration



## 6.3.11 Alternate function configuration

When the I/O port is programmed as alternate function:

- The output buffer can be configured as open-drain or push-pull
- The output buffer is driven by the signal coming from the peripheral (transmitter enable and data)
- The Schmitt trigger input is activated
- The weak pull-up and pull-down resistors are activated or not depending on the value in the GPIOx\_PUPDR register
- The data present on the I/O pin are sampled into the input data register every AHB1 clock cycle
- A read access to the input data register gets the I/O state

*Figure 17* shows the Alternate function configuration of the I/O port bit.

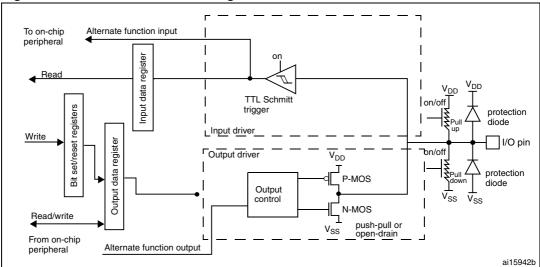


Figure 17. Alternate function configuration

#### 6.3.12 **Analog configuration**

When the I/O port is programmed as analog configuration:

- The output buffer is disabled
- The Schmitt trigger input is deactivated, providing zero consumption for every analog value of the I/O pin. The output of the Schmitt trigger is forced to a constant value (0).
- The weak pull-up and pull-down resistors are disabled
- Read access to the input data register gets the value "0"

In the analog configuration, the I/O pins cannot be 5 Volt tolerant.

Figure 18 shows the high-impedance, analog-input configuration of the I/O port bit.

Figure 18. High impedance-analog configuration Analog To on-chip peripheral Input data register registers TTL Schmitt trigger protection Write diode Bit set/reset Output data register I/O pin protection diode Read/write From on-chip Analog periphera ai15943

Note:

# 6.3.13 Using the OSC32\_IN/OSC32\_OUT pins as GPIO PC14/PC15 port pins

The LSE oscillator pins OSC32\_IN and OSC32\_OUT can be used as general-purpose PC14 and PC15 I/Os, respectively, when the LSE oscillator is off. The PC14 and PC15 I/Os are only configured as LSE oscillator pins OSC32\_IN and OSC32\_OUT when the LSE oscillator is ON. This is done by setting the LSEON bit in the RCC\_BDCR register. The LSE has priority over the GPIO function.

Note:

The PC14/PC15 GPIO functionality is lost when the 1.2 V domain is powered off (by the device entering the standby mode) or when the backup domain is supplied by  $V_{BAT}$  ( $V_{DD}$  no more supplied). In this case the I/Os are set in analog input mode.

### 6.3.14 Using the OSC IN/OSC OUT pins as GPIO PH0/PH1 port pins

The HSE oscillator pins OSC\_IN/OSC\_OUT can be used as general-purpose PH0/PH1 I/Os, respectively, when the HSE oscillator is OFF. (after reset, the HSE oscillator is off). The PH0/PH1 I/Os are only configured as OSC\_IN/OSC\_OUT HSE oscillator pins when the HSE oscillator is ON. This is done by setting the HSEON bit in the RCC\_CR register. The HSE has priority over the GPIO function.

### 6.3.15 Selection of RTC AF1 and RTC AF2 alternate functions

The STM32F4xx feature two GPIO pins RTC\_AF1 and RTC\_AF2 that can be used for the detection of a tamper or time stamp event, or AFO\_ALARM, or AFO\_CALIB RTC outputs.

The RTC\_AF1 (PC13) can be used for the following purposes:

- RTC AFO\_ALARM output: this output can be RTC Alarm A, RTC Alarm B or RTC Wakeup depending on the OSEL[1:0] bits in the RTC\_CR register
- RTC AFO\_CALIB output: this feature is enabled by setting the COE[23] in the RTC\_CR register
- RTC AFI\_TAMPER1: tamper event detection
- RTC AFI\_TIMESTAMP: time stamp event detection

The RTC\_AF2 (PI8) can be used for the following purposes:

- RTC AFI TAMPER1: tamper event detection
- RTC AFI\_TAMPER2: tamper event detection
- RTC AFI\_TIMESTAMP: time stamp event detection

The selection of the corresponding pin is performed through the RTC\_TAFCR register as follows:

- TAMP1INSEL is used to select which pin is used as the AFI\_TAMPER1 tamper input
- TSINSEL is used to select which pin is used as the AFI TIMESTAMP time stamp input
- ALARMOUTTYPE is used to select whether the RTC AFO\_ALARM is output in pushpull or open-drain mode

The output mechanism follows the priority order listed in *Table 18* and *Table 19*.

Table 18. RTC\_AF1 pin<sup>(1)</sup>

Pin configuration and function	AFO_ALARM enabled	AFO_CALIB enabled	Tamper enabled	Time stamp enabled	TAMP1INSEL TAMPER1 pin selection	TSINSEL TIMESTAMP pin selection	ALARMOUTTYPE AFO_ALARM configuration
Alarm out output OD	1	Don't care	Don't care	Don't care	Don't care	Don't care	0
Alarm out output PP	1	Don't care	Don't care	Don't care	Don't care	Don't care	1
Calibration out output PP	0	1	Don't care	Don't care	Don't care	Don't care	Don't care
TAMPER1 input floating	0	0	1	0	0	Don't care	Don't care
TIMESTAMP and TAMPER1 input floating	0	0	1	1	0	0	Don't care
TIMESTAMP input floating	0	0	0	1	Don't care	0	Don't care
Standard GPIO	0	0	0	0	Don't care	Don't care	Don't care

<sup>1.</sup> OD: open drain; PP: push-pull.

Table 19. RTC\_AF2 pin

Pin configuration and function	Tamper enabled	Time stamp enabled	TAMP1INSEL TAMPER1 pin selection	TSINSEL TIMESTAMP pin selection	ALARMOUTTYPE AFO_ALARM configuration
TAMPER1 input floating	1	0	1	Don't care	Don't care
TIMESTAMP and TAMPER1 input floating	1	1	1	1	Don't care
TIMESTAMP input floating	0	1	Don't care	1	Don't care
Standard GPIO	0	0	Don't care	Don't care	Don't care

## 6.4 GPIO registers

This section gives a detailed description of the GPIO registers.

For a summary of register bits, register address offsets and reset values, refer to Table 20.

The GPIO registers can be accessed by byte (8 bits), half-words (16 bits) or words (32 bits).

### 6.4.1 GPIO port mode register (GPIOx\_MODER) (x = A..I)

Address offset: 0x00

Reset values:

0xA800 0000 for port A

• 0x0000 0280 for port B

• 0x0000 0000 for other ports

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
MODER	R15[1:0]	MODER	R14[1:0]	MODE	R13[1:0]	MODER	R12[1:0]	MODER	R11[1:0]	MODER	R10[1:0]	MODE	R9[1:0]	MODE	R8[1:0]
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
MODE	R7[1:0]	MODE	R6[1:0]	MODE	R5[1:0]	MODE	R4[1:0]	MODE	R3[1:0]	MODE	R2[1:0]	MODE	R1[1:0]	MODE	R0[1:0]
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 2y:2y+1 **MODERy[1:0]:** Port x configuration bits (y = 0..15)

These bits are written by software to configure the I/O direction mode.

00: Input (reset state)

01: General purpose output mode

10: Alternate function mode

11: Analog mode

## 6.4.2 GPIO port output type register (GPIOx\_OTYPER) (x = A..I)

Address offset: 0x04

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Res	served							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
OT15	OT14	OT13	OT12	OT11	OT10	ОТ9	OT8	OT7	OT6	OT5	OT4	OT3	OT2	OT1	ОТ0
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:16 Reserved, must be kept at reset value.

Bits 15:0 **OTy[1:0]:** Port x configuration bits (y = 0..15)

These bits are written by software to configure the output type of the I/O port.

0: Output push-pull (reset state)

1: Output open-drain

# 6.4.3 GPIO port output speed register (GPIOx\_OSPEEDR) (x = A..I)

Address offset: 0x08

Reset values:

0x0000 00C0 for port B

0x0000 0000 for other ports

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
OSPEED	DR15[1:0]	OSPEED	)R14[1:0]	OSPEED	R13[1:0]	OSPEED	R12[1:0]	OSPEED	)R11[1:0]	OSPEED	R10[1:0]	OSPEE	DR9[1:0]	OSPEE	DR8[1:0]
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
OSPEE	DR7[1:0]	OSPEE	DR6[1:0]	OSPEE	DR5[1:0]	OSPEE	DR4[1:0]	OSPEE	DR3[1:0]	OSPEE	DR2[1:0]	OSPEE	DR1[1:0]	OSPEE	DR0[1:0]
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 2y:2y+1 **OSPEEDRy[1:0]:** Port x configuration bits (y = 0..15)

These bits are written by software to configure the I/O output speed.

00: 2 MHz Low speed 01: 25 MHz Medium speed 10: 50 MHz Fast speed

11: 100 MHz High speed on 30 pF (80 MHz Output max speed on 15 pF)

# 6.4.4 GPIO port pull-up/pull-down register (GPIOx\_PUPDR) (x = A..I)

Address offset: 0x0C

Reset values:

- 0x6400 0000 for port A0x0000 0100 for port B
- 0x0000 0000 for other ports

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
PUPDF	R15[1:0]	PUPDF	R14[1:0]	PUPDF	R13[1:0]	PUPDF	R12[1:0]	PUPDF	R11[1:0]	PUPDF	R10[1:0]	PUPDI	R9[1:0]	PUPDI	R8[1:0]
rw	rw	rw	rw	rw	rw										
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
PUPDI	R7[1:0]	PUPD	R6[1:0]	PUPD	R5[1:0]	PUPD	R4[1:0]	PUPDI	R3[1:0]	PUPDI	R2[1:0]	PUPDI	R1[1:0]	PUPDI	R0[1:0]
rw	rw	rw	rw	rw	rw										

Bits 2y:2y+1 **PUPDRy[1:0]:** Port x configuration bits (y = 0..15)

These bits are written by software to configure the I/O pull-up or pull-down

00: No pull-up, pull-down

01: Pull-up 10: Pull-down 11: Reserved

## 6.4.5 GPIO port input data register (GPIOx\_IDR) (x = A..I)

Address offset: 0x10

Reset value: 0x0000 XXXX (where X means undefined)

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Res	served							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
IDR15	IDR14	IDR13	IDR12	IDR11	IDR10	IDR9	IDR8	IDR7	IDR6	IDR5	IDR4	IDR3	IDR2	IDR1	IDR0
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:16 Reserved, must be kept at reset value.

Bits 15:0 **IDRy[15:0]**: Port input data (y = 0..15)

These bits are read-only and can be accessed in word mode only. They contain the input value of the corresponding I/O port.

## 6.4.6 GPIO port output data register (GPIOx\_ODR) (x = A..I)

Address offset: 0x14

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	rved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ODR15	ODR14	ODR13	ODR12	ODR11	ODR10	ODR9	ODR8	ODR7	ODR6	ODR5	ODR4	ODR3	ODR2	ODR1	ODR0
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:16 Reserved, must be kept at reset value.

Bits 15:0 **ODRy[15:0]:** Port output data (y = 0..15)

These bits can be read and written by software.

Note: For atomic bit set/reset, the ODR bits can be individually set and reset by writing to the  $GPIOx\_BSRR$  register (x = A..I).

### 6.4.7 GPIO port bit set/reset register (GPIOx\_BSRR) (x = A..I)

Address offset: 0x18

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
BR15	BR14	BR13	BR12	BR11	BR10	BR9	BR8	BR7	BR6	BR5	BR4	BR3	BR2	BR1	BR0
w	w	w	w	w	w	W	W	W	w	W	w	W	W	w	w
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
BS15	BS14	BS13	BS12	BS11	BS10	BS9	BS8	BS7	BS6	BS5	BS4	BS3	BS2	BS1	BS0
w	W	W	W	W	W	W	W	W	W	W	W	W	W	W	w

Bits 31:16 **BRy:** Port x reset bit y (y = 0..15)

These bits are write-only and can be accessed in word, half-word or byte mode. A read to these bits returns the value 0x0000.

- 0: No action on the corresponding ODRx bit
- 1: Resets the corresponding ODRx bit

Note: If both BSx and BRx are set, BSx has priority.

#### Bits 15:0 **BSy:** Port x set bit y (y = 0..15)

These bits are write-only and can be accessed in word, half-word or byte mode. A read to these bits returns the value 0x0000.

- 0: No action on the corresponding ODRx bit
- 1: Sets the corresponding ODRx bit

# 6.4.8 GPIO port configuration lock register (GPIOx\_LCKR) (x = A..I)

This register is used to lock the configuration of the port bits when a correct write sequence is applied to bit 16 (LCKK). The value of bits [15:0] is used to lock the configuration of the GPIO. During the write sequence, the value of LCKR[15:0] must not change. When the LOCK sequence has been applied on a port bit, the value of this port bit can no longer be modified until the next reset.

Note:

A specific write sequence is used to write to the GPIOx\_LCKR register. Only word access (32-bit long) is allowed during this write sequence.

Each lock bit freezes a specific configuration register (control and alternate function registers).

Address offset: 0x1C

Reset value: 0x0000 0000

Access: 32-bit word only, read/write register

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
						Г	Reserved								LCKK
						Г	ieserveu								rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
LCK15	LCK14	LCK13	LCK12	LCK11	LCK10	LCK9	LCK8	LCK7	LCK6	LCK5	LCK4	LCK3	LCK2	LCK1	LCK0
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:17 Reserved, must be kept at reset value.

### Bit 16 LCKK[16]: Lock key

This bit can be read any time. It can only be modified using the lock key write sequence.

0: Port configuration lock key not active

1: Port configuration lock key active. The GPIOx\_LCKR register is locked until an MCU reset occurs.

LOCK key write sequence:

WR LCKR[16] = '1' + LCKR[15:0]

WR LCKR[16] = '0' + LCKR[15:0]

WR LCKR[16] = '1' + LCKR[15:0]

RD LCKR

RD LCKR[16] = '1' (this read operation is optional but it confirms that the lock is active)

Note: During the LOCK key write sequence, the value of LCK[15:0] must not change.

Any error in the lock sequence aborts the lock.

After the first lock sequence on any bit of the port, any read access on the LCKK bit will return '1' until the next CPU reset.

Bits 15:0 **LCKy:** Port x lock bit y (y=0..15)

These bits are read/write but can only be written when the LCKK bit is '0.

0: Port configuration not locked

1: Port configuration locked

## 6.4.9 GPIO alternate function low register (GPIOx\_AFRL) (x = A..I)

Address offset: 0x20 Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	AFRL	.7[3:0]			AFRL	6[3:0]			AFRL	.5[3:0]			AFRL	.4[3:0]	,
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	AFRL	.3[3:0]		AFRL2[3:0]					AFRL	.1[3:0]			AFRL	.0[3:0]	
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:0 **AFRLy:** Alternate function selection for port x bit y (y = 0..7)

These bits are written by software to configure alternate function I/Os

AFRLy selection:

<b>,</b>	
0000: AF0	1000: AF8
0001: AF1	1001: AF9
0010: AF2	1010: AF10
0011: AF3	1011: AF11
0100: AF4	1100: AF12
0101: AF5	1101: AF13
0110: AF6	1110: AF14
0111: AF7	1111: AF15

# 6.4.10 GPIO alternate function high register (GPIOx\_AFRH) (x = A..I)

Address offset: 0x24

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	AFRH	15[3:0]			AFRH <sup>-</sup>	14[3:0]			AFRH	13[3:0]			AFRH	12[3:0]	
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	AFRH	11[3:0]			AFRH <sup>-</sup>	10[3:0]			AFRH	19[3:0]			AFRH	18[3:0]	
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:0 **AFRHy:** Alternate function selection for port x bit y (y = 8..15)

These bits are written by software to configure alternate function I/Os

AFRHy selection:

0000: AF0	1000: AF8
0001: AF1	1001: AF9
0010: AF2	1010: AF10
0011: AF3	1011: AF11
0100: AF4	1100: AF12
0101: AF5	1101: AF13
0110: AF6	1110: AF14
0111: AF7	1111: AF15

## 6.4.11 GPIO register map

The following table gives the GPIO register map and the reset values.

Table 20. GPIO register map and reset values

Offset	Register	30	28 28	27 26	25	23	20	18	17	15	13	1 0	တ ထ	<b>2</b> 9	2 4	e 0	- 0
011001	Hogiotoi																
0x00	GPIOA_MODER	MODER15[1:0]	MODER14[1:0]	.MODER13[1:0]	.MODER12[1:0]	MODER11[1:0]	MODER10[1:0]	MODER9[1:0]	MODER8[1:0]	MODER7[1:0]	MODER6[1:0]	MODER5[1:0]	MODER4[1:0]	. MODER3[1:0]	MODER2[1:0]	MODER1[1:0]	MODER0[1:0]
	Reset value	1 0	1 0	1 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0
0x00	GPIOB_MODER	MODER15[1:0]	MODER14[1:0]	MODER13[1:0]	MODER12[1:0]	MODER11[1:0]	MODER10[1:0]	MODER9[1:0]	MODER8[1:0]	MODER7[1:0]	MODER6[1:0]	MODER5[1:0]	MODER4[1:0]	MODER3[1:0]	MODER2[1:0]	MODER1[1:0]	MODER0[1:0]
	Reset value	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	1 0	1 0	0 0	0 0	0 0
0x00	GPIOx_MODER (where x = CI)	MODER15[1:0]	0 MODER14[1:0] 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0		MODER12[1:0]	MODER11[1:0]	MODER10[1:0]	. MODER9[1:0]	. MODER8[1:0]	. MODER7[1:0]	MODER6[1:0]	MODER5[1:0]	. MODER4[1:0]	MODER3[1:0]	MODER2[1:0]	. MODER1[1:0]	. MODER0[1:0]
	Reset value	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0
0x04	GPIOx_OTYPER (where x = AI)				Rese	erved				OT15	OT13	OT11 OT10	OT9 OT8	OT7 OT6	OT5	OT3 OT2	OT0
	Reset value		0 0 0 0							0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0
0x08	GPIOx_OSPEED ER (where x = AI except B)	OSPEEDR15[1:0]	OSPEEDR14[1:0] OSPEEDR14[1:0] OSPEEDR13[1:0] OOSPEEDR13[1:0] OOSPEEDR13[1:0] OOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOO		OSPEEDR12[1:0]	OSPEEDR11[1:0]	OSPEEDR10[1:0]	OSPEEDR9[1:0]	OSPEEDR8[1:0]	OSPEEDR7[1:0]	OSPEEDR6[1:0]	OSPEEDR5[1:0]	OSPEEDR4[1:0]	OSPEEDR3[1:0]	OSPEEDR2[1:0]	OSPEEDR1[1:0]	OSPEEDR0[1:0]
	Reset value	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0
0x08	GPIOB_OSPEED ER	OSPEEDR15[1:0]	OSPEEDR15[1:0] OSPEEDR15[1:0] OSPEEDR14[1:0] OSPEEDR14[1:0]		OSPEEDR12[1:0]	OSPEEDR11[1:0]	OSPEEDR10[1:0]	OSPEEDR9[1:0]	OSPEEDR8[1:0]	OSPEEDR7[1:0]	OSPEEDR6[1:0]	OSPEEDR5[1:0]	OSPEEDR4[1:0]	OSPEEDR3[1:0]	OSPEEDR2[1:0]	OSPEEDR1[1:0]	OSPEEDR0[1:0]
	Reset value				0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	1 1	0 0	0 0	0 0
0x0C	GPIOA_PUPDR	PUPDR15[1:0]	PUPDR14[1:0]	PUPDR13[1:0]	PUPDR12[1:0]	PUPDR11[1:0]	PUPDR10[1:0]	PUPDR9[1:0]	PUPDR8[1:0]	PUPDR7[1:0]	PUPDR6[1:0]	PUPDR5[1:0]	PUPDR4[1:0]	PUPDR3[1:0]	PUPDR2[1:0]	PUPDR1[1:0]	PUPDR0[1:0]
	Reset value	0 1	1 0	0 1	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0
0x0C	GPIOB_PUPDR	PUPDR15[1:0]	PUPDR14[1:0]	PUPDR13[1:0]	PUPDR12[1:0]	PUPDR11[1:0]	PUPDR10[1:0]	PUPDR9[1:0]	. PUPDR8[1:0]	PUPDR7[1:0]	PUPDR6[1:0]	PUPDR5[1:0]	. PUPDR4[1:0]	PUPDR3[1:0]	. PUPDR2[1:0]	PUPDR1[1:0]	. PUPDR0[1:0]
	Reset value	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0 1	0 0	0 0	0 0	0 0
0x0C	GPIOx_PUPDR (where x = CI)	O PUPDR15[1:0]	O PUPDR14[1:0]	O PUPDR13[1:0]	O PUPDR12[1:0]	O PUPDR11[1:0]	O PUPDR10[1:0]	O PUPDR9[1:0]	PUPDR8[1:0]	O PUPDR7[1:0]	O PUPDR6[1:0]	PUPDR5[1:0]	O PUPDR4[1:0]	PUPDR3[1:0]	O PUPDR2[1:0]	O PUPDR1[1:0]	o PUPDR0[1:0]

Table 20. GPIO register map and reset values (continued)

Offset	Register	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	1	10	6	œ	7	9	2	4	က	7	1	0
0x10	GPIOx_IDR (where x = A <b>I</b> )							F	Rese	erve	d							IDR15	IDR14	IDR13	IDR12	IDR11	IDR10	IDR9	IDR8	IDR7	IDR6	IDR5	IDR4	IDR3	IDR2	IDR1	IDR0
	Reset value																	х	х	х	х	х	Х	х	Х	х	х	х	х	Х	х	х	х
0x14	GPIOx_ODR (where x = A <b>I</b> )							F	Rese	erve	d							ODR15	ODR14	ODR13	ODR12	ODR11	ODR10	ODR9	ODR8	ODR7	ODR6	ODR5	ODR4	ODR3	ODR2	ODR1	ODR0
	Reset value																	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x18	GPIOx_BSRR (where x = A <b>I</b> )	BR15	BR14	BR13	BR12	BR11	BR10	BR9	BR8	BR7	BR6	BR5	BR4	BR3	BR2	BR1	BR0	BS15	BS14	BS13	BS12	BS11	BS10	BS9	BS8	BS7	BS6	BS5	BS4	BS3	BS2	BS1	BS0
	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x1C	GPIOx_LCKR (where x = AI)							Re	serv	/ed							LCKK	LCK15	LCK14	LCK13	LCK12	LCK11	LCK10	LCK9	LCK8	LCK7	LCK6	LCK5	LCK4	LCK3	LCK2	LCK1	LCK0
	Reset value																0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x20	GPIOx_AFRL (where x = AI)	Α	FRL	L7[3:0] AFRL6[3:0] AFRL5[3:0] AF											FRL	.4[3:	:0]	Α	FRL	3[3:	:0]	Al	FRL	.2[3:	0]	Α	FRL	.1[3:	0]	Α	FRL	.0[3:	0]
	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0			
0x24	GPIOx_AFRH (where x = AI)	AF	RH	H15[3:0] AFRH14[3:0] AFRH13[3:0] AFR										RH	12[3	3:0]	AF	RH	11[3	3:0]	AF	RH	10[3	:0]	Al	FRH	19[3	:0]	Al	RH	18[3:	0]	
	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

Refer to *Table 2 on page 50* for the register boundary addresses.

## 7 System configuration controller (SYSCFG)

The system configuration controller is mainly used to remap the memory accessible in the code area, select the Ethernet PHY interface and manage the external interrupt line connection to the GPIOs.

This section applies to the whole STM32F4xx family, unless otherwise specified.

## 7.1 I/O compensation cell

By default the I/O compensation cell is not used. However when the I/O output buffer speed is configured in 50 MHz or 100 MHz mode, it is recommended to use the compensation cell for slew rate control on I/O  $t_{f(IO)out}$ / $t_{r(IO)out}$  commutation to reduce the I/O noise on power supply.

When the compensation cell is enabled, a READY flag is set to indicate that the compensation cell is ready and can be used. The I/O compensation cell can be used only when the supply voltage ranges from 2.4 to 3.6 V.

## 7.2 SYSCFG registers

### 7.2.1 SYSCFG memory remap register (SYSCFG\_MEMRMP)

This register is used for specific configurations on memory remap:

- Two bits are used to configure the type of memory accessible at address 0x0000 0000.
   These bits are used to select the physical remap by software and so, bypass the BOOT pins.
- After reset these bits take the value selected by the BOOT pins. When booting from main Flash memory with BOOT pins set to 10 [(BOOT1,BOOT0) = (1,0)] this register takes the value 0x00.

When the FSMC is remapped at address 0x0000 0000, only the first two regions of Bank 1 memory controller (Bank1 NOR/PSRAM 1 and NOR/PSRAM 2) can be remapped. In remap mode, the CPU can access the external memory via ICode bus instead of System bus which boosts up the performance.

Address offset: 0x00

Reset value: 0x0000 000X (X is the memory mode selected by the BOOT pins

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
						Poor	erved							MEM_	MODE
						Hese	erved							rw	rw

Bits 31:2 Reserved, must be kept at reset value.

Bits 1:0 MEM\_MODE: Memory mapping selection

Set and cleared by software. This bit controls the memory internal mapping at address 0x0000 0000. After reset these bits take the value selected by the Boot pins (except for FSMC).

00: Main Flash memory mapped at 0x0000 0000 01: System Flash memory mapped at 0x0000 0000

10: FSMC Bank1 (NOR/PSRAM 1 and 2) mapped at 0x0000 0000

11: Embedded SRAM (112kB) mapped at 0x0000 0000

### 7.2.2 SYSCFG peripheral mode configuration register (SYSCFG\_PMC)

Address offset: 0x04

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
			Rese	erved				MII_RMII _SEL rw				Reserved	I		
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							Res	erved							

Bits 31:24 Reserved, must be kept at reset value.

Bit 23 MII\_RMII\_SEL: Ethernet PHY interface selection

Set and Cleared by software. These bits control the PHY interface for the Ethernet MAC.

0: MII interface is selected

1: RMII Why interface is selected

Note: This configuration must be done while the MAC is under reset and before enabling the MAC clocks.

Bits 22:0 Reserved, must be kept at reset value.

# 7.2.3 SYSCFG external interrupt configuration register 1 (SYSCFG\_EXTICR1)

Address offset: 0x08
Reset value: 0x0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	EXTI	3[3:0]			EXTI	2[3:0]			EXTI	1[3:0]			EXTI	0[3:0]	
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

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Bits 31:16 Reserved, must be kept at reset value.

#### Bits 15:0 **EXTIx[3:0]**: EXTI x configuration (x = 0 to 3)

These bits are written by software to select the source input for the EXTIx

external interrupt. 0000: PA[x] pin

0001: PB[x] pin

0010: PC[x] pin

0011: PD[x] pin

0100: PE[x] pin

0101: PF[C] pin

0110: PG[x] pin 0111: PH[x] pin

1000: PI[x] pin

# 7.2.4 SYSCFG external interrupt configuration register 2 (SYSCFG\_EXTICR2)

Address offset: 0x0C Reset value: 0x0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	EXTI	7[3:0]			EXTI	6[3:0]			EXTI	5[3:0]			EXT	4[3:0]	
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:16 Reserved, must be kept at reset value.

#### Bits 15:0 **EXTIx[3:0]**: EXTI x configuration (x = 4 to 7)

These bits are written by software to select the source input for the EXTIx external interrupt.

0000: PA[x] pin

0001: PB[x] pin

0010: PC[x] pin

0011: PD[x] pin

0100: PE[x] pin

0101: PF[x] pin

0110: PG[x] pin

0111: PH[x] pin

1000: PI[x] pin

# 7.2.5 SYSCFG external interrupt configuration register 3 (SYSCFG\_EXTICR3)

Address offset: 0x10 Reset value: 0x0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	EXTI1	1[3:0]			EXTI1	0[3:0]			EXTI	9[3:0]			EXTI	8[3:0]	
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:16 Reserved, must be kept at reset value.

Bits 15:0 **EXTIx[3:0]**: EXTI x configuration (x = 8 to 11)

These bits are written by software to select the source input for the EXTIx external interrupt.

0000: PA[x] pin 0001: PB[x] pin 0010: PC[x] pin 0011: PD[x] pin 0100: PE[x] pin 0101: PF[x] pin 0110: PG[x] pin 0111: PH[x] pin 1000: PI[x] pin

# 7.2.6 SYSCFG external interrupt configuration register 4 (SYSCFG\_EXTICR4)

Address offset: 0x14 Reset value: 0x0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	EXTI1	5[3:0]			EXTI1	4[3:0]			EXTI1	3[3:0]			EXTI1	12[3:0]	
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

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Bits 31:16 Reserved, must be kept at reset value.

Bits 15:0 **EXTIx[3:0]**: EXTI x configuration (x = 12 to 15)

These bits are written by software to select the source input for the EXTIx external

interrupt.

0000: PA[x] pin

0001: PB[x] pin

0010: PC[x] pin 0011: PD[x] pin

0100: PE[x] pin

0101: PF[x] pin

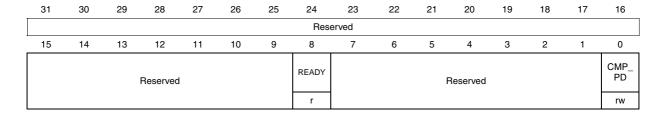
0110: PG[x] pin

0111: PH[x] pin
Note: PI[15:12] are not used.

### 7.2.7 Compensation cell control register (SYSCFG\_CMPCR)

Address offset: 0x20

Reset value: 0x0000 0000



Bits 31:9 Reserved, must be kept at reset value.

Bit 8 READY: Compensation cell ready flag

0: I/O compensation cell not ready

1: O compensation cell ready

Bits 7:2 Reserved, must be kept at reset value.

Bit 0 CMP\_PD: Compensation cell power-down

0: I/O compensation cell power-down mode

1: I/O compensation cell enabled

## 7.2.8 SYSCFG register maps

The following table gives the SYSCFG register map and the reset values.

Table 21. SYSCFG register map and reset values

Offset	Register	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	6	8	7	9	2	4	3	2	_	0
0x00	SYSCFG_MEMRM														F	lese	erve	d														MEM_MODE	
	Reset value																															х	х

Table 21. SYSCFG register map and reset values

Offset	Register	25 23 30 31 26 27 28 29	25	23	22	21	20	19	18	17	16	2	14	13	12	7	10	6	8	7	9	2	4	3	7	_	0
0x04	SYSCFG_PMC	Reserved	i	MII_RMII_SEL												serv											
	Reset value			0																							
0x08	SYSCFG_EXTICR1		Rese	rve	d						ı.			3[3:0	- 1		XTI2	٠.			XTI	٠.	-		XTIC	٠.	- 1
	Reset value										C	)	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x0C	SYSCFG_EXTICR2		Rese	rvor	Ч							EX	(TI7	'[3:0	)]	E	ΧTΙ	3[3:	0]	Е	XTI	5[3:	0]	E	XTI4	[3:0	)]
0,00	Reset value		110001	vec	u						C	)	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x10	SYSCFG_EXTICR3		Rese	n 101	٨						E	EX	TI1	1[3:	[0	ΕX	TI1	0[3	:0]	Е	XTI	9[3:	[[0	E	XTI8	[3:0	)]
0.00	Reset value		nesei	vec	u						C	)	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x14	SYSCFG_EXTICR4		Rese	n 101	٨						E	EX	TI1	5[3:	0]	EX	TI1	4[3	:0]	E	KTI1	3[3	[0]	ΕX	TI1:	2[3:	0]
UX 14	Reset value		nesei	vec	u						C	)	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x20	SYSCFG_CMPCR					F	Rese	rve	d											READY		F	lese	erve	d		CMP_PD
	Reset value																			0							0

Refer to *Table 2 on page 50* for the register boundary addresses.

## 8 DMA controller (DMA)

This section applies to the whole STM32F4xx family, unless otherwise specified.

### 8.1 DMA introduction

Direct memory access (DMA) is used in order to provide high-speed data transfer between peripherals and memory and between memory and memory. Data can be quickly moved by DMA without any CPU action. This keeps CPU resources free for other operations.

The DMA controller combines a powerful dual AHB master bus architecture with independent FIFO to optimize the bandwidth of the system, based on a complex bus matrix architecture.

The two DMA controllers have 16 streams in total (8 for each controller), each dedicated to managing memory access requests from one or more peripherals. Each stream can have up to 8 channels (requests) in total. And each has an arbiter for handling the priority between DMA requests.

### 8.2 DMA main features

The main DMA features are:

- Dual AHB master bus architecture, one dedicated to memory accesses and one dedicated to peripheral accesses
- AHB slave programming interface supporting only 32-bit accesses
- 8 streams for each DMA controller, up to 8 channels (requests) per stream
- Four separate 32 first-in, first-out memory buffers (FIFOs) per stream, that can be used in FIFO mode or direct mode:
  - FIFO mode: with threshold level software selectable between 1/4, 1/2 or 3/4 of the FIFO size
  - Direct mode

Each DMA request immediately initiates a transfer from/to the memory. When it is configured in direct mode (FIFO disabled), to transfer data in memory-to-peripheral mode, the DMA preloads only one data from the memory to the internal

FIFO to ensure an immediate data transfer as soon as a DMA request is triggered by a peripheral.

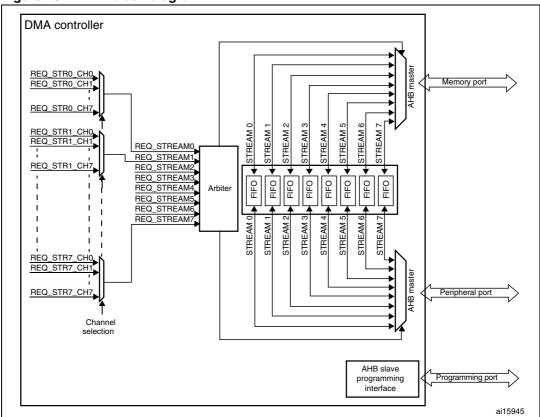
- Each stream can be configured by hardware to be:
  - a regular channel that supports peripheral-to-memory, memory-to-peripheral and memory-to-memory transfers
  - a double buffer channel that also supports double buffering on the memory side
- Each of the 8 streams are connected to dedicated hardware DMA channels (requests)
- Priorities between DMA stream requests are software-programmable (4 levels consisting of very high, high, medium, low) or hardware in case of equality (request 0 has priority over request 1, etc.)
- Each stream also supports software trigger for memory-to-memory transfers (only available for the DMA2 controller)
- Each stream request can be selected among up to 8 possible channel requests. This selection is software-configurable and allows several peripherals to initiate DMA requests
- The number of data items to be transferred can be managed either by the DMA controller or by the peripheral:
  - DMA flow controller: the number of data items to be transferred is softwareprogrammable from 1 to 65535
  - Peripheral flow controller: the number of data items to be transferred is unknown and controlled by the source or the destination peripheral that signals the end of the transfer by hardware
- Independent source and destination transfer width (byte, half-word, word): when the
  data widths of the source and destination are not equal, the DMA automatically
  packs/unpacks the necessary transfers to optimize the bandwidth. This feature is only
  available in FIFO mode
- Incrementing or nonincrementing addressing for source and destination
- Supports incremental burst transfers of 4, 8 or 16 beats. The size of the burst is software-configurable, usually equal to half the FIFO size of the peripheral
- Each stream supports circular buffer management
- 5 event flags (DMA Half Transfer, DMA Transfer complete, DMA Transfer Error, DMA FIFO Error, Direct Mode Error) logically ORed together in a single interrupt request for each stream

## 8.3 DMA functional description

### 8.3.1 General description

Figure 19 shows the block diagram of a DMA.

Figure 19. DMA block diagram



The DMA controller performs direct memory transfer: as an AHB master, it can take the control of the AHB bus matrix to initiate AHB transactions.

It can carry out the following transactions:

- peripheral-to-memory
- memory-to-peripheral
- memory-to-memory

The DMA controller provides two AHB master ports: the *AHB memory port*, intended to be connected to memories and the *AHB peripheral port*, intended to be connected to peripherals. However, to allow memory-to-memory transfers, the *AHB peripheral port* must also have access to the memories.

The AHB slave port is used to program the DMA controller (it supports only 32-bit accesses).

See *Figure 20* for the implementation of the system of two DMA controllers.

DMA controller (DMA) RM0090

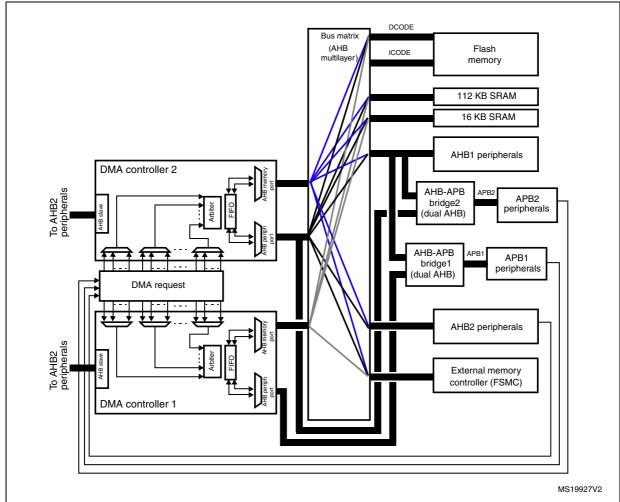


Figure 20. System implementation of the two DMA controllers

The DMA1 controller AHB peripheral port is not connected to the bus matrix like in the case of the DMA2 controller, thus
only DMA2 streams are able to perform memory-to-memory transfers.

### 8.3.2 DMA transactions

A DMA transaction consists of a sequence of a given number of data transfers. The number of data items to be transferred and their width (8-bit, 16-bit or 32-bit) are software-programmable.

Each DMA transfer consists of three operations:

- A loading from the peripheral data register or a location in memory, addressed through the DMA SxPAR or DMA SxM0AR register
- A storage of the data loaded to the peripheral data register or a location in memory addressed through the DMA\_SxPAR or DMA\_SxMOAR register
- A post-decrement of the DMA\_SxNDTR register, which contains the number of transactions that still have to be performed

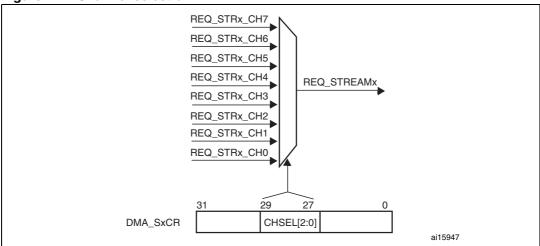
After an event, the peripheral sends a request signal to the DMA controller. The DMA controller serves the request depending on the channel priorities. As soon as the DMA controller accesses the peripheral, an Acknowledge signal is sent to the peripheral by the

DMA controller. The peripheral releases its request as soon as it gets the Acknowledge signal from the DMA controller. Once the request has been deasserted by the peripheral, the DMA controller releases the Acknowledge signal. If there are more requests, the peripheral can initiate the next transaction.

### 8.3.3 Channel selection

Each stream is associated with a DMA request that can be selected out of 8 possible channel requests. The selection is controlled by the CHSEL[2:0] bits in the DMA\_SxCR register.

Figure 21. Channel selection



The 8 requests from the peripherals (TIM, ADC, SPI, I2C, etc.) are independently connected to each channel and their connection depends on the product implementation.

Table 22 and Table 23 give examples of DMA request mappings.

Table 22. DMA1 request mapping

Peripheral requests	Stream 0	Stream 1	Stream 2	Stream 3	Stream 4	Stream 5	Stream 6	Stream 7
Channel 0	SPI3_RX		SPI3_RX	SPI2_RX	SPI2_TX	SPI3_TX		SPI3_TX
Channel 1	I2C1_RX		TIM7_UP		TIM7_UP	I2C1_RX	I2C1_TX	I2C1_TX
Channel 2	TIM4_CH1		I2S3_EXT_ RX	TIM4_CH2	I2S2_EXT_ TX	I2S3_EXT_ TX	TIM4_UP	TIM4_CH3
Channel 3	I2S3_EXT_ RX	TIM2_UP TIM2_CH3	I2C3_RX	I2S2_EXT_ RX	I2C3_TX	TIM2_CH1	TIM2_CH2 TIM2_CH4	TIM2_UP TIM2_CH4
Channel 4	UART5_RX	USART3_RX	UART4_RX	USART3_TX	UART4_TX	USART2_RX	USART2_TX	UART5_TX
Channel 5			TIM3_CH4 TIM3_UP		TIM3_CH1 TIM3_TRIG	TIM3_CH2		TIM3_CH3
Channel 6	TIM5_CH3 TIM5_UP	TIM5_CH4 TIM5_TRIG	TIM5_CH1	TIM5_CH4 TIM5_TRIG	TIM5_CH2		TIM5_UP	
Channel 7		TIM6_UP	I2C2_RX	I2C2_RX	USART3_TX	DAC1	DAC2	I2C2_TX

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Table 23. DMA2 request mapping

Peripheral requests	Stream 0	Stream 1	Stream 2	Stream 3	Stream 4	Stream 5	Stream 6	Stream 7
			TIM8_CH1				TIM1_CH1	
Channel 0	ADC1		TIM8_CH2		ADC1		TIM1_CH2	
			TIM8_CH3				TIM1_CH3	
Channel 1		DCMI	ADC2	ADC2				DCMI
Channel 2	ADC3	ADC3				CRYP_OUT	CRYP_IN	HASH_IN
Channel 3	SPI1_RX		SPI1_RX	SPI1_TX		SPI1_TX		
Channel 4			USART1_RX	SDIO		USART1_RX	SDIO	USART1_TX
Channel 5		USART6_RX	USART6_RX				USART6_TX	USART6_TX
					TIM1_CH4			
Channel 6	TIM1_TRIG	TIM1_CH1	TIM1_CH2	TIM1_CH1	TIM1_TRIG	TIM1_UP	TIM1_CH3	
					TIM1_COM			
								TIM8_CH4
Channel 7		TIM8_UP	TIM8_CH1	TIM8_CH2	TIM8_CH3			TIM8_TRIG
								TIM8_COM

### 8.3.4 Arbiter

An arbiter manages the 8 DMA stream requests based on their priority for each of the two AHB master ports (memory and peripheral ports) and launches the peripheral/memory access sequences.

Priorities are managed in two stages:

- Software: each stream priority can be configured in the DMA\_SxCR register. There are four levels:
  - Very high priority
  - High priority
  - Medium priority
  - Low priority
- Hardware: If two requests have the same software priority level, the stream with the lower number takes priority over the stream with the higher number. For example, Stream 2 takes priority over Stream 4.

### 8.3.5 DMA streams

Each of the 8 DMA controller streams provides a unidirectional transfer link between a source and a destination.

Each stream can be configured to perform:

- Regular type transactions: memory-to-peripherals, peripherals-to-memory or memory-to-memory transfers
- Double-buffer type transactions: double buffer transfers using two memory pointers for the memory (while the DMA is reading/writing from/to a buffer, the application can write/read to/from the other buffer).

The amount of data to be transferred (up to 65535) is programmable and related to the source width of the peripheral that requests the DMA transfer connected to the peripheral AHB port. The register that contains the amount of data items to be transferred is decremented after each transaction.

### 8.3.6 Source, destination and transfer modes

Both source and destination transfers can address peripherals and memories in the entire 4 GB area, at addresses comprised between 0x0000 0000 and 0xFFFF FFFF.

The direction is configured using the DIR[1:0] bits in the DMA\_SxCR register and offers three possibilities: memory-to-peripheral, peripheral-to-memory or memory-to-memory transfers. *Table 24* describes the corresponding source and destination addresses.

Table 24. Oddiec and destination address							
Bits DIR[1:0] of the DMA_SxCR register	Direction	Source address	Destination address				
00	Peripheral-to-memory	DMA_SxPAR	DMA_SxM0AR				
01	Memory-to-peripheral	DMA_SxM0AR	DMA_SxPAR				
10	Memory-to-memory	DMA_SxPAR	DMA_SxM0AR				
11	reserved	-	-				

Table 24. Source and destination address

When the data width (programmed in the PSIZE or MSIZE bits in the DMA\_SxCR register) is a half-word or a word, respectively, the peripheral or memory address written into the DMA\_SxPAR or DMA\_SxM0AR/M1AR registers has to be aligned on a word or half-word address boundary, respectively.

#### Peripheral-to-memory mode

Figure 22 describes this mode.

When this mode is enabled (by setting the bit EN in the DMA\_SxCR register), each time a peripheral request occurs, the stream initiates a transfer from the source to fill the FIFO.

When the threshold level of the FIFO is reached, the contents of the FIFO are drained and stored into the destination.

The transfer stops once the DMA\_SxNDTR register reaches zero, when the peripheral requests the end of transfers (in case of a peripheral flow controller) or when the EN bit in the DMA\_SxCR register is cleared by software.

In direct mode (when the DMDIS value in the DMA\_SxFCR register is '0'), the threshold level of the FIFO is not used: after each single data transfer from the peripheral to the FIFO, the corresponding data are immediately drained and stored into the destination.

The stream has access to the AHB source or destination port only if the arbitration of the corresponding stream is won. This arbitration is performed using the priority defined for each stream using the PL[1:0] bits in the DMA\_SxCR register.

DMA controller (DMA) RM0090

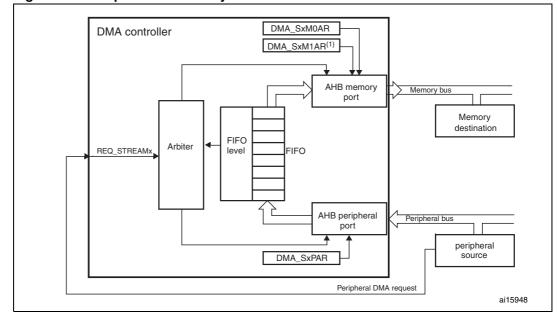


Figure 22. Peripheral-to-memory mode

1. For double-buffer mode.

### Memory-to-peripheral mode

Figure 23 describes this mode.

When this mode is enabled (by setting the EN bit in the DMA\_SxCR register), the stream immediately initiates transfers from the source to entirely fill the FIFO.

Each time a peripheral request occurs, the contents of the FIFO are drained and stored into the destination. When the level of the FIFO is lower than or equal to the predefined threshold level, the FIFO is fully reloaded with data from the memory.

The transfer stops once the DMA\_SxNDTR register reaches zero, when the peripheral requests the end of transfers (in case of a peripheral flow controller) or when the EN bit in the DMA\_SxCR register is cleared by software.

In direct mode (when the DMDIS value in the DMA\_SxFCR register is '0'), the threshold level of the FIFO is not used. Once the stream is enabled, the DMA preloads the first data to transfer into an internal FIFO. As soon as the peripheral requests a data transfer, the DMA transfers the preloaded value into the configured destination. It then reloads again the empty internal FIFO with the next data to be transfer. The preloaded data size corresponds to the value of the PSIZE bitfield in the DMA\_SxCR register.

The stream has access to the AHB source or destination port only if the arbitration of the corresponding stream is won. This arbitration is performed using the priority defined for each stream using the PL[1:0] bits in the DMA\_SxCR register.

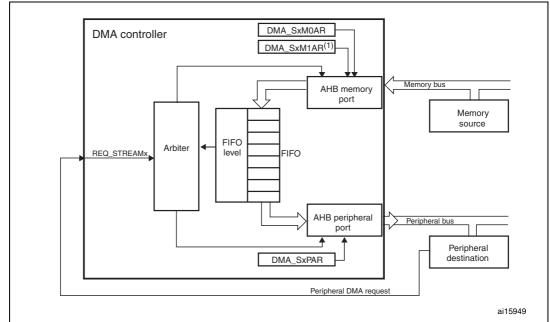


Figure 23. Memory-to-peripheral mode

1. For double-buffer mode.

### Memory-to-memory mode

The DMA channels can also work without being triggered by a request from a peripheral. This is the memory-to-memory mode, described in *Figure 24*.

When the stream is enabled by setting the Enable bit (EN) in the DMA\_SxCR register, the stream immediately starts to fill the FIFO up to the threshold level. When the threshold level is reached, the FIFO contents are drained and stored into the destination.

The transfer stops once the DMA\_SxNDTR register reaches zero or when the EN bit in the DMA\_SxCR register is cleared by software.

The stream has access to the AHB source or destination port only if the arbitration of the corresponding stream is won. This arbitration is performed using the priority defined for each stream using the PL[1:0] bits in the DMA\_SxCR register.

Note: When memory-to-memory mode is used, the Circular and direct modes are not allowed.

Only the DMA2 controller is able to perform memory-to-memory transfers.

DMA controller (DMA) RM0090

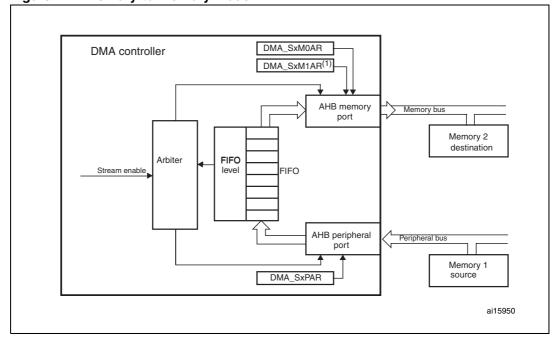


Figure 24. Memory-to-memory mode

1. For double-buffer mode.

### 8.3.7 Pointer incrementation

Peripheral and memory pointers can optionally be automatically post-incremented or kept constant after each transfer depending on the PINC and MINC bits in the DMA\_SxCR register.

Disabling the Increment mode is useful when the peripheral source or destination data are accessed through a single register.

If the Increment mode is enabled, the address of the next transfer will be the address of the previous one incremented by 1 (for bytes), 2 (for half-words) or 4 (for words) depending on the data width programmed in the PSIZE or MSIZE bits in the DMA\_SxCR register.

In order to optimize the packing operation, it is possible to fix the increment offset size for the peripheral address whatever the size of the data transferred on the AHB peripheral port. The PINCOS bit in the DMA\_SxCR register is used to align the increment offset size with the data size on the peripheral AHB port, or on a 32-bit address (the address is then incremented by 4). The PINCOS bit has an impact on the AHB peripheral port only.

If PINCOS bit is set, the address of the next transfer is the address of the previous one incremented by 4 (automatically aligned on a 32-bit address) whatever the PSIZE value. The AHB memory port, however, is not impacted by this operation.

The PINC or the MINC bit needs to be set if the burst transaction is requested on the AHB peripheral port or the AHB memory port, respectively, to satisfy the AMBA protocol (burst is not allowed in the fixed address mode).

#### 8.3.8 Circular mode

The Circular mode is available to handle circular buffers and continuous data flows (e.g. ADC scan mode). This feature can be enabled using the CIRC bit in the DMA\_SxCR register.

When the circular mode is activated, the number of data items to be transferred is automatically reloaded with the initial value programmed during the stream configuration phase, and the DMA requests continue to be served.

Note:

In the circular mode, it is mandatory to respect the following rule in case of a burst mode configured for memory:

DMA  $SxNDTR = Multiple of ((Mburst beat) \times (Msize)/(Psize)), where:$ 

- (Mburst beat) = 4, 8 or 16 (depending on the MBURST bits in the DMA\_SxCR register)
- ((Msize)/(Psize)) = 1, 2, 4, 1/2 or 1/4 (Msize and Psize represent the MSIZE and PSIZE bits in the DMA\_SxCR register. They are byte dependent)
- DMA\_SxNDTR = Number of data items to transfer on the AHB peripheral port

For example: Mburst beat = 8 (INCR8), MSIZE = '00' (byte) and PSIZE = '01' (half-word), in this case:  $DMA\_SxNDTR$  must be a multiple of  $(8 \times 1/2 = 4)$ .

If this formula is not respected, the DMA behavior and data integrity are not guaranteed.

NDTR must also be a multiple of the Peripheral burst size multiplied by the peripheral data size, otherwise this could result in a bad DMA behavior.

#### 8.3.9 Double buffer mode

This mode is available for all the DMA1 and DMA2 streams.

The Double buffer mode is enabled by setting the DBM bit in the DMA\_SxCR register.

A double-buffer stream works as a regular (single buffer) stream with the difference that it has two memory pointers. When the Double buffer mode is enabled, the Circular mode is automatically enabled (CIRC bit in DMA\_SxCR is don't care) and at each end of transaction, the memory pointers are swapped.

In this mode, the DMA controller swaps from one memory target to another at each end of transaction. This allows the software to process one memory area while the second memory area is being filled/used by the DMA transfer. The double-buffer stream can work in both directions (the memory can be either the source or the destination) as described in *Table 25: Source and destination address registers in Double buffer mode (DBM=1)*.

Note:

In Double buffer mode, it is possible to update the base address for the AHB memory port on-the-fly (DMA\_SxM0AR or DMA\_SxM1AR) when the stream is enabled, by respecting the following conditions:

- When the CT bit is '0' in the DMA\_SxCR register, the DMA\_SxM1AR register can be written. Attempting to write to this register while CT = '1' sets an error flag (TEIF) and the stream is automatically disabled.
- When the CT bit is '1' in the DMA\_SxCR register, the DMA\_SxM0AR register can be written. Attempting to write to this register while CT = '0', sets an error flag (TEIF) and the stream is automatically disabled.

To avoid any error condition, it is advised to change the base address as soon as the TCIF flag is asserted because, at this point, the targeted memory must have changed from

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memory 0 to 1 (or from 1 to 0) depending on the value of CT in the DMA\_SxCR register in accordance with one of the two above conditions.

For all the other modes (except the Double buffer mode), the memory address registers are write-protected as soon as the stream is enabled.

Table 25.	Source and destination address registers in Double buffer mode (DBM=1)
-----------	--

Bits DIR[1:0] of the DMA_SxCR register	Direction	Source address	Destination address
00	Peripheral-to-memory	DMA_SxPAR	DMA_SxM0AR / DMA_SxM1AR
01	Memory-to-peripheral	DMA_SxM0AR / DMA_SxM1AR	DMA_SxPAR
10	Not allowed <sup>(1)</sup>		
11	Reserved	-	-

When the Double buffer mode is enabled, the Circular mode is automatically enabled. Since the memoryto-memory mode is not compatible with the Circular mode, when the Double buffer mode is enabled, it is not allowed to configure the memory-to-memory mode.

### 8.3.10 Programmable data width, packing/unpacking, endianess

The number of data items to be transferred has to be programmed into DMA\_SxNDTR (number of data items to transfer bit, NDT) before enabling the stream (except when the flow controller is the peripheral, PFCTRL bit in DMA\_SxCR is set).

When using the internal FIFO, the data widths of the source and destination data are programmable through the PSIZE and MSIZE bits in the DMA\_SxCR register (can be 8-, 16- or 32-bit).

When PSIZE and MSIZE are not equal:

- The data width of the number of data items to transfer, configured in the DMA\_SxNDTR register is equal to the width of the peripheral bus (configured by the PSIZE bits in the DMA\_SxCR register). For instance, in case of peripheral-to-memory, memory-to-peripheral or memory-to-memory transfers and if the PSIZE[1:0] bits are configured for half-word, the number of bytes to be transferred is equal to 2 × NDT.
- The DMA controller only copes with little-endian addressing for both source and destination. This is described in *Table 26: Packing/unpacking & endian behavior (bit PINC = MINC = 1)*.

This packing/unpacking procedure may present a risk of data corruption when the operation is interrupted before the data are completely packed/unpacked. So, to ensure data coherence, the stream may be configured to generate burst transfers: in this case, each group of transfers belonging to a burst are indivisible (refer to *Section 8.3.11: Single and burst transfers*).

In direct mode (DMDIS = 0 in the DMA\_SxFCR register), the packing/unpacking of data is not possible. In this case, it is not allowed to have different source and destination transfer data widths: both are equal and defined by the PSIZE bits in the DMA\_SxCR MSIZE bits are don't care).

Table 26. Packing/unpacking & endian behavior (bit PINC = MINC = 1)

AHB AHB		Number	Memor		Doubbou	Peripheral port address / byte lane	
memor y port width	peripher al port width	of data items to transfer (NDT)	y transfe r numbe r	Memory port address / byte lane	Peripher al transfer number	PINCOS = 1	PINCOS = 0
8	8	4	1 2 3 4	0x0 / B0[7:0] 0x1 / B1[7:0] 0x2 / B2[7:0] 0x3 / B3[7:0]	1 2 3 4	0x0 / B0[7:0] 0x4 / B1[7:0] 0x8 / B2[7:0] 0xC / B3[7:0]	0x0 / B0[7:0] 0x1 / B1[7:0] 0x2 / B2[7:0] 0x3 / B3[7:0]
8	16	2	1 2 3 4	0x0 / B0[7:0] 0x1 / B1[7:0] 0x2 / B2[7:0] 0x3 / B3[7:0]	1 2	0x0 / B1 B0[15:0] 0x4 / B3 B2[15:0]	0x0 / B1 B0[15:0] 0x2 / B3 B2[15:0]
8	32	1	1 2 3 4	0x0 / B0[7:0] 0x1 / B1[7:0] 0x2 / B2[7:0] 0x3 / B3[7:0]	1	0x0 / B3 B2 B1 B0[31:0]	0x0 / B3IB2IB1IB0[31:0]
16	8	4	2	0x0 / B1 B0[15:0] 0x2 / B3 B2[15:0]	1 2 3 4	0x0 / B0[7:0] 0x4 / B1[7:0] 0x8 / B2[7:0] 0xC / B3[7:0]	0x0 / B0[7:0] 0x1 / B1[7:0] 0x2 / B2[7:0] 0x3 / B3[7:0]
16	16	2	1 2	0x0 / B1 B0[15:0] 0x2 / B1 B0[15:0]	1 2	0x0 / B1lB0[15:0] 0x4 / B3lB2[15:0]	0x0 / B1 B0[15:0] 0x2 / B3 B2[15:0]
16	32	1	1 2	0x0 / B1 B0[15:0] 0x2 / B3 B2[15:0]	1	0x0 / B3 B2 B1 B0[31:0]	0x0 / B3lB2lB1lB0[31:0]
32	8	4	1	0x0 / B3IB2IB1IB0[31:0]	1 2 3 4	0x0 / B0[7:0] 0x4 / B1[7:0] 0x8 / B2[7:0] 0xC / B3[7:0]	0x0 / B0[7:0] 0x1 / B1[7:0] 0x2 / B2[7:0] 0x3 / B3[7:0]
32	16	2	1	0x0 /B3IB2IB1IB0[31:0]	1 2	0x0 / B1 B0[15:0] 0x4 / B3 B2[15:0]	0x0 / B1 B0[15:0] 0x2 / B3 B2[15:0]
32	32	1	1	0x0 /B3 B2 B1 B0 [31:0]	1	0x0 /B3 B2 B1 B0 [31:0]	0x0 / B3 B2 B1 B0[31:0]

Note:

Peripheral port may be the source or the destination (it could also be the memory source in the case of memory-to-memory transfer).

PSIZE, MSIZE and NDT[15:0] have to be configured so as to ensure that the last transfer will not be incomplete. This can occur when the data width of the peripheral port (PSIZE bits) is lower than the data width of the memory port (MSIZE bits). This constraint is summarized in *Table 27*.

Table 27. Restriction on NDT versus PSIZE and MSIZE

PSIZE[1:0] of DMA_SxCR	MSIZE[1:0] of DMA_SxCR	NDT[15:0] of DMA_SxNDTR
00 (8-bit)	01 (16-bit)	must be a multiple of 2
00 (8-bit)	10 (32-bit)	must be a multiple of 4
01 (16-bit)	10 (32-bit)	must be a multiple of 2

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### 8.3.11 Single and burst transfers

The DMA controller can generate single transfers or incremental burst transfers of 4, 8 or 16 beats.

The size of the burst is configured by software independently for the two AHB ports by using the MBURST[1:0] and PBURST[1:0] bits in the DMA\_SxCR register.

The burst size indicates the number of beats in the burst, not the number of bytes transferred.

To ensure data coherence, each group of transfers that form a burst are indivisible: AHB transfers are locked and the arbiter of the AHB bus matrix does not degrant the DMA master during the sequence of the burst transfer.

Depending on the single or burst configuration, each DMA request initiates a different number of transfers on the AHB peripheral port:

- When the AHB peripheral port is configured for single transfers, each DMA request generates a data transfer of a byte, half-word or word depending on the PSIZE[1:0] bits in the DMA\_SxCR register
- When the AHB peripheral port is configured for burst transfers, each DMA request generates 4,8 or 16 beats of byte, half word or word transfers depending on the PBURST[1:0] and PSIZE[1:0] bits in the DMA\_SxCR register.

The same as above has to be considered for the AHB memory port considering the MBURST and MSIZE bits.

In direct mode, the stream can only generate single transfers and the MBURST[1:0] and PBURST[1:0] bits are forced by hardware.

The address pointers (DMA\_SxPAR or DMA\_SxM0AR registers) must be chosen so as to ensure that all transfers within a burst block are aligned on the address boundary equal to the size of the transfer.

The burst configuration has to be selected in order to respect the AHB protocol, where bursts must *not* cross the 1 KB address boundary because the minimum address space that can be allocated to a single slave is 1 KB. This means that the 1 KB address boundary should not be crossed by a burst block transfer, otherwise an AHB error would be generated, that is not reported by the DMA registers.

Note:

The Burst mode is allowed only when incremetation is enabled:

- When the PINC bit is at '0', the PBURST bits should also be cleared to '00'
- When the MINC bit is at '0', the MBURST bits should also be cleared to '00'.

### 8.3.12 FIFO

#### **FIFO structure**

The FIFO is used to temporarily store data coming from the source before transmitting them to the destination.

Each stream has an independent 4-word FIFO and the threshold level is software-configurable between 1/4, 1/2, 3/4 or full.

To enable the use of the FIFO threshold level, the direct mode must be disabled by setting the DMDIS bit in the DMA SxFCR register.

The structure of the FIFO differs depending on the source and destination data widths, and is described in *Figure 25: FIFO structure*.

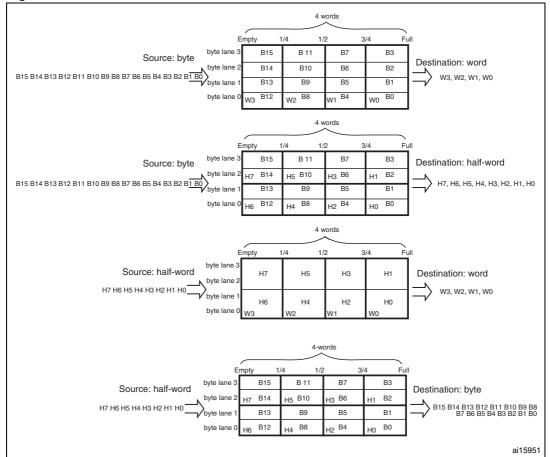


Figure 25. FIFO structure

### FIFO threshold and burst configuration

Caution is required when choosing the FIFO threshold (bits FTH[1:0] of the DMA\_SxFCR register) and the size of the memory burst (MBURST[1:0] of the DMA\_SxCR register): The content pointed by the FIFO threshold must exactly match to an integer number of memory burst transfers. If this is not in the case, a FIFO error (flag FEIFx of the DMA\_HISR or DMA\_LISR register) will be generated when the stream is enabled, then the stream will be automatically disabled. The allowed and forbidden configurations are described in the *Table 28: FIFO threshold configurations*.

Table 28. FIFO threshold configurations

MSIZE	FIFO level	MBURST = INCR4	MBURST = INCR8	MBURST = INCR16			
Byte	1/4	1 burst of 4 beats	forbidden				
	1/2	2 bursts of 4 beats	1 burst of 8 beats	forbidden			
	3/4	3 bursts of 4 beats	forbidden				
	Full	4 bursts of 4 beats	2 bursts of 8 beats	1 burst of 16 beats			

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MSIZE	FIFO level	MBURST = INCR4	MBURST = INCR8	MBURST = INCR16
	1/4	forbidden		
Half-word	1/2	1 burst of 4 beats	forbidden	
Tiali-word	3/4	forbidden		
	Full 2 bursts of 4 beats 1 burst of 8 beats	1 burst of 8 beats	forbidden	
	1/4			lorbidderi
Word	1/2	forbidden	forbidden	
vvora	3/4			
	Full	1 burst of 4 beats		

Table 28. FIFO threshold configurations (continued)

In all cases, the burst size multiplied by the data size must not exceed the FIFO size (data size can be: 1 (byte), 2 (half-word) or 4 (word)).

Incomplete Burst transfer at the end of a DMA transfer may happen if one of the following conditions occurs:

- For the AHB peripheral port configuration: the total number of data items (set in the DMA\_SxNDTR register) is not a multiple of the burst size multiplied by the data size
- For the AHB memory port configuration: the number of remaining data items in the FIFO to be transferred to the memory is not a multiple of the burst size multiplied by the data size

In such cases, the remaining data to be transferred will be managed in single mode by the DMA, even if a burst transaction was requested during the DMA stream configuration.

Note:

When burst transfers are requested on the peripheral AHB port and the FIFO is used (DMDIS = 1 in the DMA\_SxCR register), it is mandatory to respect the following rule to avoid permanent underrun or overrun conditions, depending on the DMA stream direction:

If  $(PBURST \times PSIZE) = FIFO\_SIZE$  (4 words),  $FIFO\_Threshold = 3/4$  is forbidden with PSIZE = 1, 2 or 4 and PBURST = 4, 8 or 16.

This rule ensures that enough FIFO space at a time will be free to serve the request from the peripheral.

### FIFO flush

The FIFO can be flushed when the stream is disabled by resetting the EN bit in the DMA\_SxCR register and when the stream is configured to manage peripheral-to-memory or memory-to-memory transfers: If some data are still present in the FIFO when the stream is disabled, the DMA controller continues transferring the remaining data to the destination (even though stream is effectively disabled). When this flush is completed, the transfer complete status bit (TCIFx) in the DMA\_LISR or DMA\_HISR register is set.

The remaining data counter DMA\_SxNDTR keeps the value in this case to indicate how many data items are currently available in the destination memory.

Note that during the FIFO flush operation, if the number of remaining data items in the FIFO to be transferred to memory (in bytes) is less than the memory data width (for example 2 bytes in FIFO while MSIZE is configured to word), data will be sent with the data width set in the MSIZE bit in the DMA\_SxCR register. This means that memory will be written with an

undesired value. The software may read the DMA\_SxNDTR register to determine the memory area that contains the good data (start address and last address).

If the number of remaining data items in the FIFO is lower than a burst size (if the MBURST bits in DMA\_SxCR register are set to configure the stream to manage burst on the AHB memory port), single transactions will be generated to complete the FIFO flush.

#### **Direct mode**

By default, the FIFO operates in direct mode (DMDIS bit in the DMA\_SxFCR is reset) and the FIFO threshold level is not used. This mode is useful when the system requires an immediate and single transfer to or from the memory after each DMA request.

When the DMA is configured in direct mode (FIFO disabled), to transfer data in memory-toperipheral mode, the DMA preloads one data from the memory to the internal FIFO to ensure an immediate data transfer as soon as a DMA request is triggered by a peripheral.

To avoid saturating the FIFO, it is recommended to configure the corresponding stream with a high priority.

This mode is restricted to transfers where:

- The source and destination transfer widths are equal and both defined by the PSIZE[1:0] bits in DMA\_SxCR (MSIZE[1:0] bits are don't care)
- Burst transfers are not possible (PBURST[1:0] and MBURST[1:0] bits in DMA\_SxCR are don't care)

Direct mode must not be used when implementing memory-to-memory transfers.

### 8.3.13 DMA transfer completion

Different events can generate an end of transfer by setting the TCIFx bit in the DMA\_LISR or DMA\_HISR status register:

- In DMA flow controller mode:
  - The DMA\_SxNDTR counter has reached zero in the memory-to-peripheral mode
  - The stream is disabled before the end of transfer (by clearing the EN bit in the DMA\_SxCR register) and (when transfers are peripheral-to-memory or memoryto-memory) all the remaining data have been flushed from the FIFO into the memory
- In Peripheral flow controller mode:
  - The last external burst or single request has been generated from the peripheral and (when the DMA is operating in peripheral-to-memory mode) the remaining data have been transferred from the FIFO into the memory
  - The stream is disabled by software, and (when the DMA is operating in peripheralto-memory mode) the remaining data have been transferred from the FIFO into the memory

Note:

The transfer completion is dependent on the remaining data in FIFO to be transferred into memory only in the case of peripheral-to-memory mode. This condition is not applicable in memory-to-peripheral mode.

If the stream is configured in noncircular mode, after the end of the transfer (that is when the number of data to be transferred reaches zero), the DMA is stopped (EN bit in DMA\_SxCR register is cleared by Hardware) and no DMA request is served unless the software reprograms the stream and re-enables it (by setting the EN bit in the DMA\_SxCR register).

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## 8.3.14 DMA transfer suspension

At any time, a DMA transfer can be suspended to be restarted later on or to be definitively disabled before the end of the DMA transfer.

There are two cases:

- The stream disables the transfer with no later-on restart from the point where it was stopped. There is no particular action to do, except to clear the EN bit in the DMA\_SxCR register to disable the stream. The stream may take time to be disabled (ongoing transfer is completed first). The transfer complete interrupt flag (TCIF in the DMA\_LISR or DMA\_HISR register) is set in order to indicate the end of transfer. The value of the EN bit in DMA\_SxCR is now '0' to confirm the stream interruption. The DMA\_SxNDTR register contains the number of remaining data items at the moment when the stream was stopped so that the software can determine how many data items have been transferred before the stream was interrupted.
- The stream suspends the transfer before the number of remaining data items to be transferred in the DMA\_SxNDTR register reaches 0. The aim is to restart the transfer later by re-enabling the stream. In order to restart from the point where the transfer was stopped, the software has to read the DMA\_SxNDTR register after disabling the stream by writing the EN bit in DMA\_SxCR register (and then checking that it is at '0') to know the number of data items already collected. Then:
  - The peripheral and/or memory addresses have to be updated in order to adjust the address pointers
  - The SxNDTR register has to be updated with the remaining number of data items to be transferred (the value read when the stream was disabled)
  - The stream may then be re-enabled to restart the transfer from the point it was stopped

Note:

Note that a Transfer complete interrupt flag (TCIF in DMA\_LISR or DMA\_HISR) is set to indicate the end of transfer due to the stream interruption.

## 8.3.15 Flow controller

The entity that controls the number of data to be transferred is known as the flow controller. This flow controller is configured independently for each stream using the PFCTRL bit in the DMA\_SxCR register.

The flow controller can be:

- The DMA controller: in this case, the number of data items to be transferred is programmed by software into the DMA\_SxNDTR register before the DMA stream is enabled.
- The peripheral source or destination: this is the case when the number of data items to be transferred is unknown. The peripheral indicates by hardware to the DMA controller when the last data are being transferred. This feature is only supported for peripherals which are able to signal the end of the transfer, that is:
  - SDIO

When the peripheral flow controller is used for a given stream, the value written into the DMA\_SxNDTR has no effect on the DMA transfer. Actually, whatever the value written, it will

be forced by hardware to 0xFFFF as soon as the stream is enabled, to respect the following schemes:

- Anticipated stream interruption: EN bit in DMA\_SxCR register is reset to 0 by the software to stop the stream before the last data hardware signal (single or burst) is sent by the peripheral. In such a case, the stream is switched off and the FIFO flush is triggered in the case of a peripheral-to-memory DMA transfer. The TCIFx flag of the corresponding stream is set in the status register to indicate the DMA completion. To know the number of data items transferred during the DMA transfer, read the DMA\_SxNDTR register and apply the following formula:
  - Number of data transferred = 0xFFFF DMA SxNDTR
- Normal stream interruption due to the reception of a last data hardware signal: the stream is automatically interrupted when the peripheral requests the last transfer (single or burst) and when this transfer is complete. the TCIFx flag of the corresponding stream is set in the status register to indicate the DMA transfer completion. To know the number of data items transferred, read the DMA\_SxNDTR register and apply the same formula as above.
- The DMA\_SxNDTR register reaches 0: the TCIFx flag of the corresponding stream is set in the status register to indicate the forced DMA transfer completion. The stream is automatically switched off even though the last data hardware signal (single or burst) has not been yet asserted. The already transferred data will not be lost. This means that a maximum of 65535 data items can be managed by the DMA in a single transaction, even in peripheral flow control mode.

Note: When configured in memory-to-memory mode, the DMA is always the flow controller and the PFCTRL bit is forced to 0 by hardware.

The Circular mode is forbidden in the peripheral flow controller mode.

## 8.3.16 Summary of the possible DMA configurations

Table 29 summarizes the different possible DMA configurations.

Table 29. Possible DMA configurations

DMA transfer mode	Source	Destination	Flow controller	Circular mode	Transfer type	Direct mode	Double buffer mode
			DMA	possible	single	possible	possible
Peripheral-to-	AHB	AHB	DIVIA	possible	burst	forbidden	possible
memory	peripheral port	memory port	Peripheral	forbidden	single	possible	forbidden
			i empirerai	lorbidderi	burst	forbidden	IOIDIQUEII
			DMA	possible	single	possible	possible
Memory-to-	AHB	AHB	DIVIA	possible	burst	forbidden	possible
peripheral	memory port	peripheral port	Peripheral	forbidden	single	possible	forbidden
			renpherai	lorbidderi	burst	forbidden	Torbidden
	AHB	AHB	DMA only	forbidden	single	forbidden	forbidden
memory	peripheral port	memory port	DIVIA OHIY	lorbidden	burst	ioibiddeii	IOIDIQUEII

## 8.3.17 Stream configuration procedure

The following sequence should be followed to configure a DMA stream x (where x is the stream number):

- 1. If the stream is enabled, disable it by resetting the EN bit in the DMA\_SxCR register, then read this bit in order to confirm that there is no ongoing stream operation. Writing this bit to 0 is not immediately effective since it is actually written to 0 once all the current transfers have finished. When the EN bit is read as 0, this means that the stream is ready to be configured. It is therefore necessary to wait for the EN bit to be cleared before starting any stream configuration. All the stream dedicated bits set in the status register (DMA\_LISR and DMA\_HISR) from the previous data block DMA transfer should be cleared before the stream can be re-enabled.
- 2. Set the peripheral port register address in the DMA\_SxPAR register. The data will be moved from/ to this address to/ from the peripheral port after the peripheral event.
- 3. Set the memory address in the DMA\_SxMA0R register (and in the DMA\_SxMA1R register in the case of a double buffer mode). The data will be written to or read from this memory after the peripheral event.
- Configure the total number of data items to be transferred in the DMA\_SxNDTR register. After each peripheral event or each beat of the burst, this value is decremented.
- 5. Select the DMA channel (request) using CHSEL[2:0] in the DMA SxCR register.
- 6. If the peripheral is intended to be the flow controller and if it supports this feature, set the PFCTRL bit in the DMA\_SxCR register.
- 7. Configure the stream priority using the PL[1:0] bits in the DMA\_SxCR register.
- 8. Configure the FIFO usage (enable or disable, threshold in transmission and reception)
- Configure the data transfer direction, peripheral and memory incremented/fixed mode, single or burst transactions, peripheral and memory data widths, Circular mode, Double buffer mode and interrupts after half and/or full transfer, and/or errors in the DMA\_SxCR register.
- 10. Activate the stream by setting the EN bit in the DMA\_SxCR register.

As soon as the stream is enabled, it can serve any DMA request from the peripheral connected to the stream.

Once half the data have been transferred on the AHB destination port, the half-transfer flag (HTIF) is set and an interrupt is generated if the half-transfer interrupt enable bit (HTIE) is set. At the end of the transfer, the transfer complete flag (TCIF) is set and an interrupt is generated if the transfer complete interrupt enable bit (TCIE) is set.

Warning:

To switch off a peripheral connected to a DMA stream request, it is mandatory to, first, switch off the DMA stream to which the peripheral is connected, then to wait for EN bit = 0. Only then can the peripheral be safely disabled.

## 8.3.18 Error management

The DMA controller can detect the following errors:

- Transfer error: the transfer error interrupt flag (TEIFx) is set when:
  - A bus error occurs during a DMA read or a write access
  - A write access is requested by software on a memory address register in Double buffer mode whereas the stream is enabled and the current target memory is the one impacted by the write into the memory address register (refer to Section 8.3.9: Double buffer mode)
- FIFO error: the FIFO error interrupt flag (FEIFx) is set if:
  - A FIFO underrun condition is detected
  - A FIFO overrun condition is detected (no detection in memory-to-memory mode because requests and transfers are internally managed by the DMA)
  - The stream is enabled while the FIFO threshold level is not compatible with the size of the memory burst (refer to Table 28: FIFO threshold configurations)
- Direct mode error: the direct mode error interrupt flag (DMEIFx) can only be set in the peripheral-to-memory mode while operating in direct mode and when the MINC bit in the DMA\_SxCR register is cleared. This flag is set when a DMA request occurs while the previous data have not yet been fully transferred into the memory (because the memory bus was not granted). In this case, the flag indicates that 2 data items were be transferred successively to the same destination address, which could be an issue if the destination is not able to manage this situation

In direct mode, the FIFO error flag can also be set under the following conditions:

- In the peripheral-to-memory mode, the FIFO can be saturated (overrun) if the memory bus is not granted for several peripheral requests
- In the memory-to-peripheral mode, an underrun condition may occur if the memory bus has not been granted before a peripheral request occurs

If the TEIFx or the FEIFx flag is set due to incompatibility between burst size and FIFO threshold level, the faulty stream is automatically disabled through a hardware clear of its EN bit in the corresponding stream configuration register (DMA\_SxCR).

If the DMEIFx or the FEIFx flag is set due to an overrun or underrun condition, the faulty stream is not automatically disabled and it is up to the software to disable or not the stream by resetting the EN bit in the DMA\_SxCR register. This is because there is no data loss when this kind of errors occur.

When the stream's error interrupt flag (TEIF, FEIF, DMEIF) in the DMA\_LISR or DMA\_HISR register is set, an interrupt is generated if the corresponding interrupt enable bit (TEIE, FEIE, DMIE) in the DMA\_SxCR or DMA\_SxFCR register is set.

Note:

When a FIFO overrun or underrun condition occurs, the data are not lost because the peripheral request is not acknowledged by the stream until the overrun or underrun condition is cleared. If this acknowledge takes too much time, the peripheral itself may detect an overrun or underrun condition of its internal buffer and data might be lost.

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# 8.4 DMA interrupts

For each DMA stream, an interrupt can be produced on the following events:

- Half-transfer reached
- Transfer complete
- Transfer error
- Fifo error (overrun, underrun or FIFO level error)
- Direct mode error

Separate interrupt enable control bits are available for flexibility as shown in *Table 30*.

Table 30. DMA interrupt requests

Interrupt event	Event flag	Enable control bit
Half-transfer	HTIF	HTIE
Transfer complete	TCIF	TCIE
Transfer error	TEIF	TEIE
FIFO overrun/underrun	FEIF	FEIE
Direct mode error	DMEIF	DMEIE

Note:

Before setting an Enable control bit to '1', the corresponding event flag should be cleared, otherwise an interrupt is immediately generated.

# 8.5 DMA registers

The DMA registers can be accessed by words (32 bits).

## 8.5.1 DMA low interrupt status register (DMA\_LISR)

Address offset: 0x00

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	Rese	rved		TCIF3	HTIF3	TEIF3	DMEIF3	Reserv	FEIF3	TCIF2	HTIF2	TEIF2	DMEIF2	Reserv	FEIF2
r	r	r	r	r	r	r	r	ed	r	r	r	r	r	ed	r
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Rese	rved		TCIF1	HTIF1	TEIF1	DMEIF1	Reserv	FEIF1	TCIF0	HTIF0	TEIF0	DMEIF0	Reserv	FEIF0
r	r	r	r	r	r	r	r	ed	r	r	r	r	r	ed	r

Bits 31:28, 15:12 Reserved, must be kept at reset value.

Bits 27, 21, 11, 5 **TCIFx**: Stream x transfer complete interrupt flag (x = 3..0)

This bit is set by hardware. It is cleared by software writing 1 to the corresponding bit in the DMA\_LIFCR register.

0: No transfer complete event on stream x

1: A transfer complete event occurred on stream x

Bits 26, 20, 10, 4 **HTIFx**: Stream x half transfer interrupt flag (x=3..0)

This bit is set by hardware. It is cleared by software writing 1 to the corresponding bit in the DMA\_LIFCR register.

0: No half transfer event on stream x

1: A half transfer event occurred on stream x

Bits 25, 19, 9, 3 **TEIFx**: Stream x transfer error interrupt flag (x=3..0)

This bit is set by hardware. It is cleared by software writing 1 to the corresponding bit in the DMA\_LIFCR register.

0: No transfer error on stream x

1: A transfer error occurred on stream x

Bits 24, 18, 8, 2 **DMEIFx**: Stream x direct mode error interrupt flag (x=3..0)

This bit is set by hardware. It is cleared by software writing 1 to the corresponding bit in the DMA\_LIFCR register.

0: No Direct Mode Error on stream x

1: A Direct Mode Error occurred on stream x

Bits 23, 17, 7, 1 Reserved, must be kept at reset value.

Bits 22, 16, 6, 0 **FEIFx**: Stream x FIFO error interrupt flag (x=3..0)

This bit is set by hardware. It is cleared by software writing 1 to the corresponding bit in the DMA\_LIFCR register.

0: No FIFO Error event on stream x

1: A FIFO Error event occurred on stream x

# 8.5.2 DMA high interrupt status register (DMA\_HISR)

Address offset: 0x04

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	Rese	rved		TCIF7	HTIF7	TEIF7	DMEIF7	Reserv ed	FEIF7	TCIF6	HTIF6	TEIF6	DMEIF6	Reserv	FEIF6
				r	r	r	r	eu	r	r	r	r	r	ed	r
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	15 14 13 1 Reserved			TCIF5	HTIF5	TEIF5	DMEIF5	Reserv ed	FEIF5	TCIF4	HTIF4	TEIF4	DMEIF4	Reserv	FEIF4
				r	r	r	r	eu	r	r	r	r	r	ed	r

Bits 31:28, 15:12 Reserved, must be kept at reset value.

Bits 27, 21, 11, 5 TCIFx: Stream x transfer complete interrupt flag (x=7..4)

This bit is set by hardware. It is cleared by software writing 1 to the corresponding bit in the DMA\_HIFCR register.

0: No transfer complete event on stream x

1: A transfer complete event occurred on stream x

Bits 26, 20, 10, 4 HTIFx: Stream x half transfer interrupt flag (x=7..4)

This bit is set by hardware. It is cleared by software writing 1 to the corresponding bit in the DMA\_HIFCR register.

0: No half transfer event on stream x

1: A half transfer event occurred on stream x

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Bits 25, 19, 9, 3 **TEIFx**: Stream x transfer error interrupt flag (x=7..4)

This bit is set by hardware. It is cleared by software writing 1 to the corresponding bit in the DMA\_HIFCR register.

0: No transfer error on stream x

1: A transfer error occurred on stream x

Bits 24, 18, 8, 2 **DMEIFx**: Stream x direct mode error interrupt flag (x=7..4)

This bit is set by hardware. It is cleared by software writing 1 to the corresponding bit in the DMA\_HIFCR register.

0: No Direct mode error on stream x

1: A Direct mode error occurred on stream x

Bits 23, 17, 7, 1 Reserved, must be kept at reset value.

Bits 22, 16, 6, 0 FEIFx: Stream x FIFO error interrupt flag (x=7..4)

This bit is set by hardware. It is cleared by software writing 1 to the corresponding bit in the DMA\_HIFCR register.

0: No FIFO error event on stream x

1: A FIFO error event occurred on stream x

# 8.5.3 DMA low interrupt flag clear register (DMA\_LIFCR)

Address offset: 0x08

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	Rese	ryod		CTCIF3	CHTIF3	CTEIF3	CDMEIF3	Reserved	CFEIF3	CTCIF2	CHTIF2	CTEIF2	CDMEIF2	Reserved	CFEIF2
	nese	iveu		rw	rw	rw	rw	neserveu	rw	rw	rw	rw	rw	neserveu	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Poso	ryod		CTCIF1	CHTIF1	CTEIF1	CDMEIF1	Reserved	CFEIF1	CTCIF0	CHTIF0	CTEIF0	CDMEIF0	Reserved	CFEIF0
	Reserved	rw	rw	rw	rw	neserveu	rw	rw	rw	rw	rw	neserveu	rw		

Bits 31:28, 15:12 Reserved, must be kept at reset value.

Bits 27, 21, 11, 5 CTCIFx: Stream x clear transfer complete interrupt flag (x = 3..0)

This bit is set and cleared by software.

0: No effect

1: Clears the corresponding TCIFx flag in the DMA\_LISR register

Bits 26, 20, 10, 4 **CHTIFx**: Stream x clear half transfer interrupt flag (x = 3..0)

This bit is set and cleared by software.

0: No effect

1: Clears the corresponding HTIFx flag in the DMA\_LISR register

Bits 25, 19, 9, 3 **CTEIFx**: Stream x clear transfer error interrupt flag (x = 3..0)

This bit is set and cleared by software.

0: No effect

1: Clears the corresponding TEIFx flag in the DMA\_LISR register

Bits 24, 18, 8, 2 **CDMEIFx**: Stream x clear direct mode error interrupt flag (x = 3..0)

This bit is set and cleared by software.

0: No effect

1: Clears the corresponding DMEIFx flag in the DMA\_LISR register

Bits 23, 17, 7, 1 Reserved, must be kept at reset value.

Bits 22, 16, 6, 0 **CFEIFx**: Stream x clear FIFO error interrupt flag (x = 3..0)

This bit is set and cleared by software.

0: No effect

1: Clears the corresponding CFEIFx flag in the DMA\_LISR register

## 8.5.4 DMA high interrupt flag clear register (DMA\_HIFCR)

Address offset: 0x0C

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	Rese	rvod		CTCIF7	CHTIF7	CTEIF7	CDMEIF7	Reserved	CFEIF7	CTCIF6	CHTIF6	CTEIF6	CDMEIF6	Reserved	CFEIF6
	nese	iveu		rw	rw	rw	rw	neserveu	rw	rw	rw	rw	rw	nesei veu	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Door	nod		CTCIF5	CHTIF5	CTEIF5	CDMEIF5	Reserved	CFEIF5	CTCIF4	CHTIF4	CTEIF4	CDMEIF4	Reserved	CFEIF4
	Reserved	rw	rw	rw	rw	neserveu	rw	rw	rw	rw	rw	neserveu	rw		

Bits 31:28, 15:12 Reserved, must be kept at reset value.

Bits 27, 21, 11, 5 **CTCIFx**: Stream x clear transfer complete interrupt flag (x = 7..4)

This bit is set and cleared by software.

0: No effect

1: Clears the corresponding TCIFx flag in the DMA\_HISR register

Bits 26, 20, 10, 4 **CHTIFx**: Stream x clear half transfer interrupt flag (x = 7..4)

This bit is set and cleared by software.

0: No effect

1: Clears the corresponding HTIFx flag in the DMA\_HISR register

Bits 25, 19, 9, 3 **CTEIFx**: Stream x clear transfer error interrupt flag (x = 7..4)

This bit is set and cleared by software.

0: No effect

1: Clears the corresponding TEIFx flag in the DMA\_HISR register

Bits 24, 18, 8, 2 **CDMEIFx**: Stream x clear direct mode error interrupt flag (x = 7..4)

This bit is set and cleared by software.

0: No effect

1: Clears the corresponding DMEIFx flag in the DMA\_HISR register

Bits 23, 17, 7, 1 Reserved, must be kept at reset value.

Bits 22, 16, 6, 0 **CFEIFx**: Stream x clear FIFO error interrupt flag (x = 7..4)

This bit is set and cleared by software.

0: No effect

1: Clears the corresponding CFEIFx flag in the DMA\_HISR register

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## 8.5.5 DMA stream x configuration register (DMA\_SxCR) (x = 0..7)

This register is used to configure the concerned stream.

Address offset: 0x10 + 0x18 × stream number

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	Reser	ved		(	CHSEL[3:	0]	MBURS	T [1:0]	PBUF	RST[1:0]	Reserv	СТ	DBM or reserved	PL[	1:0]
				rw	rw	rw	rw	rw	rw	rw	eu	rw	rw or r	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
PINCOS	MSIZ	E[1:0]	PSIZI	Ξ[1:0]	MINC	PINC	CIRC	DIR	[1:0]	PFCTRL	TCIE	HTIE	TEIE	DMEIE	EN
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:28 Reserved, must be kept at reset value.

### Bits 27:25 CHSEL[2:0]: Channel selection

These bits are set and cleared by software.

000: channel 0 selected

001: channel 1 selected

010: channel 2 selected

011: channel 3 selected

100: channel 4 selected

101: channel 5 selected

110: channel 6 selected111: channel 7 selected

These bits are protected and can be written only if EN is '0'

### Bits 24:23 MBURST: Memory burst transfer configuration

These bits are set and cleared by software.

00: single transfer

01: INCR4 (incremental burst of 4 beats)

10: INCR8 (incremental burst of 8 beats)

11: INCR16 (incremental burst of 16 beats)

These bits are protected and can be written only if EN is '0'

In direct mode, these bits are forced to 0x0 by hardware as soon as bit EN= '1'.

### Bits 22:21 PBURST[1:0]: Peripheral burst transfer configuration

These bits are set and cleared by software.

00: single transfer

01: INCR4 (incremental burst of 4 beats)

10: INCR8 (incremental burst of 8 beats)

11: INCR16 (incremental burst of 16 beats)

These bits are protected and can be written only if EN is '0'

In direct mode, these bits are forced to 0x0 by hardware.

### Bits 20 Reserved, must be kept at reset value.

### Bits 19 CT: Current target (only in double buffer mode)

This bits is set and cleared by hardware. It can also be written by software.

0: The current target memory is Memory 0 (addressed by the DMA\_SxM0AR pointer)

1: The current target memory is Memory 1 (addressed by the DMA\_SxM1AR pointer)

This bit can be written only if EN is '0' to indicate the target memory area of the first transfer. Once the stream is enabled, this bit operates as a status flag indicating which memory area

is the current target.

#### Bits 18 DBM: Double buffer mode

This bits is set and cleared by software.

0: No buffer switching at the end of transfer

1: Memory target switched at the end of the DMA transfer

This bit is protected and can be written only if EN is '0'.

### Bits 17:16 PL[1:0]: Priority level

These bits are set and cleared by software.

00: Low

01: Medium

10: High

11: Very high

These bits are protected and can be written only if EN is '0'.

#### Bits 15 PINCOS: Peripheral increment offset size

This bit is set and cleared by software

0: The offset size for the peripheral address calculation is linked to the PSIZE

1: The offset size for the peripheral address calculation is fixed to 4 (32-bit alignment).

This bit has no meaning if bit PINC = '0'.

This bit is protected and can be written only if EN = '0'.

This bit is forced low by hardware when the stream is enabled (bit EN = '1') if the direct mode is selected or if PBURST are different from "00".

#### Bits 14:13 MSIZE[1:0]: Memory data size

These bits are set and cleared by software.

00: byte (8-bit)

01: half-word (16-bit)

10: word (32-bit)

11: reserved

These bits are protected and can be written only if EN is '0'.

In direct mode, MSIZE is forced by hardware to the same value as PSIZE as soon as bit EN = '1'.

### Bits 12:11 PSIZE[1:0]: Peripheral data size

These bits are set and cleared by software.

00: Byte (8-bit)

01: Half-word (16-bit)

10: Word (32-bit)

11: reserved

These bits are protected and can be written only if EN is '0'

## Bits 10 MINC: Memory increment mode

This bit is set and cleared by software.

0: Memory address pointer is fixed

1: Memory address pointer is incremented after each data transfer (increment is done according to MSIZE)

This bit is protected and can be written only if EN is '0'.

### Bits 9 PINC: Peripheral increment mode

This bit is set and cleared by software.

0: Peripheral address pointer is fixed

1: Peripheral address pointer is incremented after each data transfer (increment is done according to PSIZE)

This bit is protected and can be written only if EN is '0'.

#### Bits 8 CIRC: Circular mode

This bit is set and cleared by software and can be cleared by hardware.

0: Circular mode disabled

1: Circular mode enabled

When the peripheral is the flow controller (bit PFCTRL=1) and the stream is enabled (bit EN=1), then this bit is automatically forced by hardware to 0.

It is automatically forced by hardware to 1 if the DBM bit is set, as soon as the stream is enabled (bit EN = '1').

### Bits 7:6 DIR[1:0]: Data transfer direction

These bits are set and cleared by software.

00: Peripheral-to-memory

01: Memory-to-peripheral

10: Memory-to-memory

11: reserved

These bits are protected and can be written only if EN is '0'.

### Bits 5 PFCTRL: Peripheral flow controller

This bit is set and cleared by software.

0: The DMA is the flow controller

1: The peripheral is the flow controller

This bit is protected and can be written only if EN is '0'.

When the memory-to-memory mode is selected (bits DIR[1:0]=10), then this bit is automatically forced to 0 by hardware.

### Bits 4 TCIE: Transfer complete interrupt enable

This bit is set and cleared by software.

0: TC interrupt disabled

1: TC interrupt enabled

### Bits 3 HTIE: Half transfer interrupt enable

This bit is set and cleared by software.

0: HT interrupt disabled

1: HT interrupt enabled

### Bits 2 TEIE: Transfer error interrupt enable

This bit is set and cleared by software.

0: TE interrupt disabled

1: TE interrupt enabled

### Bits 1 **DMEIE**: Direct mode error interrupt enable

This bit is set and cleared by software.

0: DME interrupt disabled

1: DME interrupt enabled

Bits 0 EN: Stream enable / flag stream ready when read low

This bit is set and cleared by software.

- 0: Stream disabled
- 1: Stream enabled

This bit may be cleared by hardware:

- on a DMA end of transfer (stream ready to be configured)
- if a transfer error occurs on the AHB master buses
- when the FIFO threshold on memory AHB port is not compatible with the size of the burst

When this bit is read as 0, the software is allowed to program the Configuration and FIFO bits registers. It is forbidden to write these registers when the EN bit is read as 1.

Note: Before setting EN bit to '1' to start a new transfer, the event flags corresponding to the stream in DMA\_LISR or DMA\_HISR register must be cleared.

## 8.5.6 DMA stream x number of data register (DMA\_SxNDTR) (x = 0..7)

Address offset: 0x14 + 0x18 × stream number

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Res	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							NDT	[15:0]							
rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bits 31:16 Reserved, must be kept at reset value.

#### Bits 15:0 NDT[15:0]: Number of data items to transfer

Number of data items to be transferred (0 up to 65535). This register can be written only when the stream is disabled. When the stream is enabled, this register is read-only, indicating the remaining data items to be transmitted. This register decrements after each DMA transfer.

Once the transfer has completed, this register can either stay at zero (when the stream is in normal mode) or be reloaded automatically with the previously programmed value in the following cases:

- when the stream is configured in Circular mode.
- when the stream is enabled again by setting EN bit to '1'

If the value of this register is zero, no transaction can be served even if the stream is enabled.

## 8.5.7 DMA stream x peripheral address register (DMA\_SxPAR) (x = 0..7)

Address offset: 0x18 + 0x18 × stream number

Reset value: 0x0000 0000

3	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
								PAR	[31:16]							
r	w	rw	rw	rw	rw	rw	rw	rw	rw	rw						

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15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							PAR	[15:0]							
rw	rw	rw	rw	rw	rw	rw	rw	rw							

### Bits 31:0 PAR[31:0]: Peripheral address

Base address of the peripheral data register from/to which the data will be read/written.

These bits are write-protected and can be written only when bit EN = '0' in the DMA\_SxCR register.

# 8.5.8 DMA stream x memory 0 address register (DMA\_SxM0AR) (x = 0..7)

Address offset: 0x1C + 0x18 × stream number

Reset value: 0x0000 0000

3	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
								MOA	[31:16]							
r	w	rw	rw	rw	rw	rw	rw	rw	rw	rw						
1	5	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
								MOA	[15:0]							
r	w	rw	rw	rw	rw	rw	rw	rw	rw	rw						

Bits 31:0 MOA[31:0]: Memory 0 address

Base address of Memory area 0 from/to which the data will be read/written.

These bits are write-protected. They can be written only if:

- the stream is disabled (bit EN= '0' in the DMA\_SxCR register) or
- the stream is enabled (EN='1' in DMA\_SxCR register) and bit CT = '1' in the DMA\_SxCR register (in Double buffer mode).

## 8.5.9 DMA stream x memory 1 address register (DMA\_SxM1AR) (x = 0..7)

Address offset: 0x20 + 0x18 × stream number

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	ā.					-	M1A	[31:16]	ā.						
rw	rw	rw	rw	rw	rw	rw	rw	rw							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							M1A	[15:0]							
rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bits 31:0 M1A[31:0]: Memory 1 address (used in case of Double buffer mode)

Base address of Memory area 1 from/to which the data will be read/written.

This register is used only for the Double buffer mode.

These bits are write-protected. They can be written only if:

- the stream is disabled (bit EN= '0' in the DMA\_SxCR register) or
- the stream is enabled (EN='1' in DMA\_SxCR register) and bit CT = '0' in the DMA\_SxCR register.

## 8.5.10 DMA stream x FIFO control register (DMA\_SxFCR) (x = 0..7)

Address offset: 0x24 + 0x24 × stream number

Reset value: 0x0000 0021

	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
								Re	eserved							
_	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
				Por	served				FEIE	Reser		FS[2:0]		DMDIS	FTH	[1:0]
				nes	serveu				rw	ved	r	r	r	rw	rw	rw

Bits 31:8 Reserved, must be kept at reset value.

### Bits 7 FEIE: FIFO error interrupt enable

This bit is set and cleared by software.

0: FE interrupt disabled1: FE interrupt enabled

Bits 6 Reserved, must be kept at reset value.

### Bits 5:3 FS[2:0]: FIFO status

These bits are read-only. 000:  $0 < \text{fifo\_level} < 1/4$  001:  $1/4 \le \text{fifo\_level} < 1/2$  010:  $1/2 \le \text{fifo\_level} < 3/4$  011:  $3/4 \le \text{fifo\_level} < \text{full}$  100: FIFO is empty 101: FIFO is full

These bits are not relevant in the direct mode (DMDIS bit is zero).

### Bits 2 DMDIS: Direct mode disable

others: no meaning

This bit is set and cleared by software. It can be set by hardware.

0: Direct mode enabled1: Direct mode disabled

This bit is protected and can be written only if EN is '0'.

This bit is set by hardware if the memory-to-memory mode is selected (DIR bit in DMA\_SxCR are "10") and the EN bit in the DMA\_SxCR register is '1' because the direct mode is not allowed in the memory-to-memory configuration.

### Bits 1:0 FTH[1:0]: FIFO threshold selection

These bits are set and cleared by software.

00: 1/4 full FIFO 01: 1/2 full FIFO 10: 3/4 full FIFO 11: full FIFO

These bits are not used in the direct mode when the DMIS value is zero.

These bits are protected and can be written only if EN is '1'.

# 8.5.11 DMA register map

Table 31 summarizes the DMA registers.

Table 31. DMA register map and reset values

DMA_HISR   Reserved   DMA_BISR   Reserved   DMA_BISR   Reserved   DMA_SORR   Reserved	Offset	Register	E 8 8 8	<del> </del>		24 2	1 1	21	20	19	18	17	16	15	13	12	10	6	8	7	9	2	4	3	7	-	0
Reservative	0x0000		Reserved		o TEIF3	O DMEIF3 Reserved				O TEIF2	O DMEIF2	Reserved		Rese	rved				O DMEIF1	Reserved		l		O TEIF0	O DMEIF0	Reserved	o FEIF0
Male	0x0004		Reserved	-			-					Reserved		Rese	rved					Reserved	_					Reserved	o FEIF4
DMA_HIFCR   Reset value   DMA_SOCR   DMA_SOCR   DMA_SOCR   Reset value   DMA_SOCR	0x0008		Reserved									Reserved		Rese	rved				CDMEIF1	Reserved				cTEIF0	CDMEIF0	Reserved	O CFEIF0
DMA_SOCR   Reserved   DMA_SOCR   DMA_SOCR   Reserved   DMA_SOCR   DMA_SOCR   DMA_SOCR   Reserved   DMA_SOCR   Reserved   DMA_SOCR	0x000C	DMA_HIFCR	Reserved	CTCIF7 CHTIF7	CTEIF7	CDIMEIF/ leserved	CFEIF7	CTCIF6	CHTIF6	CTEIF6	CDMEIF6	Reserved	CFEIF6	Rese	rved	CTCIF5	CHTIF5	CTEIF5	CDMEIF5	Reserved	CFEIF5	CTCIF4	CHTIF4	CTEIF4	CDMEIF4	Reserved	O CFEIF4
Reset value	0x0010		Reserved	$\vdash$		<u> </u>	+	-				DI [4:0]		PINCOS MSIZE[1:0]						10.10						DMEIE	EN
Name		Reset value		0 0	0					0	0	0	0	0 0	0	0 0	0	0	0	0	0	0	0	0	0	0	0
Name		DMA_S0NDTR					1 1											1	NDT	[15:	:.]	<u> </u>			<u> </u>		_
Name	0x0014	Reset value			He	eserv	ed							0 0	0	0 0	0	0	0	0	0	0	0	0	0	0	0
DMA_SOMOAR  Reset value  O O O O O O O O O O O O O O O O O O O	0x0018	DMA_S0PAR											PA[3	31:0]								ı	ı				
Reset value		Reset value	0 0 0 0	0 0	0	0 0	0	0	0	0	0	0	0	0 0	0	0 0	0	0	0	0	0	0	0	0	0	0	0
DMA_S0M1AR    Paset value   DMA_S1CR   Reset value   DMA_S1CR   DMA_S1CR   DMA_S1CR   Reset value   DMA_S1CR   D	0x001C	DMA_S0M0AR										٨	/IOA[	[31:0]													
Reset value		Reset value	0 0 0 0	0 0	0	0 0	0	0	0	0	0	0	0	0 0	0	0 0	0	0	0	0	0	0	0	0	0	0	0
DMA_S0FCR   Reset value   DMA_S1CR   Reset value   DMA_S1DTR   Reset value   DMA_S1NDTR   Reset value   DMA_S1NDTR   Reset value   DMA_S1NDTR   Reset value   DMA_S1PAR   DMA_S1PAR   DMA_S1PAR   DMA_S1PAR   DMA_S1PAR   DMA_S1PAR   DMA_S1PAR   DMA_S1PAR   DMA_S1DTR   DMA_S1DTR   DMA_S1PAR   DMA_S1PAR   DMA_S1PAR   DMA_S1PAR   DMA_S1DTR   DMA_SDTR   D	0x0020	DMA_S0M1AR										٨	/11A[	[31:0]													
Reset value		Reset value	0 0 0 0	0 0	0	0 0	0	0	0	0	0	0	0	0 0	0	0 0	0	0	0	0	0	0	0			0	0
Reset value	0x0024	DMA_S0FCR						R	lese	rve	d									FEE	served		S[2:	0]	DMDIS	FT [1:	
Reset value		Reset value																		0	Re	1	0			0	1
Ox002C         DMA_S1NDTR         Reserved         Image: color of the c	0x0028	DMA_S1CR	Reserved	CHSEL [9:0]		MBURST[1:]	PBURST[1:0]		AC		DBM	FJ Id	- -	PING			MIN	PINC	CIRC	DID[4:0]		ı —		_	Ŧ	IWQ	N EN
0x002C         Reset value         Reserved         0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0				0 0	0	0 0	0	0	0	0	0	0	0	0 0	0	0 0	0	0	0	0	0	0	0	0	0	0	0
Reset value	0x002C	DMA_S1NDTR			Be	eserv	ed											1	NDT	[15	:.]						
0x0030	5.0020	Reset value												0 0	0	0 0	0	0	0	0	0	0	0	0	0	0	0
Reset value 10101010101010101010101010101010101010	0x0030	DMA_S1PAR											PA[3	31:0]	•	•	•							<u> </u>			
		Reset value	0 0 0 0	0 0	0	0 0	0	0	0	0	0	0	0	0 0	0	0 0	0	0	0	0	0	0	0	0	0	0	0

Table 31. DMA register map and reset values (continued)

Table .	JI. DINA	egister map and reset values (continued)	
Offset	Register	2 3 4 5 6 6 7 8 8 9 9 9 9 9 9 9 9 9 9 9 9 9 9 9 9 9	- 0
0x0034	DMA_S1M0AR	M0A[31:0]	
	Reset value		0 0
0x0038	DMA_S1M1AR	M1A[31:0]	
	Reset value		0 0
0x003C	DMA_S1FCR	Reserved	FTH [1:0]
	Reset value	0 2 1 0 0 0	0 1
0x0040	DMA_S2CR		DMEIE
	Reset value		0 0
0x0044	DMA_S2NDTR	NDT[15:.]	
	Reset value		0 0
0x0048	DMA_S2PAR	PA[31:0]	
	Reset value		0 0
0x004C	DMA_S2M0AR	M0A[31:0]	
	Reset value		0 0
0x0050	DMA_S2M1AR	M1A[31:0]	
	Reset value		0 0
0x0054	DMA_S2FCR	Reserved	FTH [1:0]
	Reset value	0 2 1 0 0 0	0 1
0x0058	DMA_S3CR	CHSEL[2:0]	DMEIE
	Reset value	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	0 0
0x005C	DMA_S3NDTR	Reserved NDT[15:.]	
0,0000	Reset value	0 0 0 0 0 0 0 0 0 0 0 0	0 0
0x0060	DMA_S3PAR	PA[31:0]	
	Reset value		0 0
0x0064	DMA_S3M0AR	M0A[31:0]	
	Reset value		0 0
0x0068	DMA_S3M1AR	M1A[31:0]	
	Reset value		0 0
0x006C	DMA_S3FCR	Reserved By FS[2:0] ON DATE OF THE PART OF	FTH [1:0]
	Reset value	0 2 1 0 0 0	0 1

DMA controller (DMA) RM0090

Table 31. DMA register map and reset values (continued)

011		<del> </del>	<del>i i i</del>		г т		s (co	_	т т				Τ												$\neg$
Offset	Register	30 28 28 28	27 26	25 24	23	22	20	18	17	16	15	13	12	11	10	6	8	7	9	2	4	3	2	-	0
0x0070	DMA_S4CR	Reserved	CHSEL[2:0]	MBUBST[1-0]	[o:-]-o	PBURST[1:0]	ACK	MBQ	PL[1:0]		PINCOS	MSIZE[1:0]	DC17E14-01	r 312E[1.0]		PINC	CIRC	DIR	[1:0]	PFCTRL	TCIE		EIE	DMEIE	EN
	Reset value		0 0	0 0	0	0 0	0 0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x0074	DMA_S4NDTR			Rese	ervec	d											IDT	_	_						
	Reset value											0 0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x0078	DMA_S4PAR		1 - 1 - 1	- 1 -						PA[3				_			_				_				_
	Reset value	0 0 0 0	0 0	0 0	0	0 0	0 0	0	0	0		0 0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x007C	DMA_S4M0AR			0.10		010	1010	١.			31:0		<u> </u>	_		_		_	_		_		_		
	Reset value	0 0 0 0	0 0	0 0	0	0 0	0 0	0	0	0		0 0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x0080	DMA_S4M1AR			0 1 0	I o I	0 0		10			31:0		Ι ο	Ι ο		^	_	۱ ۵	١.	Ι ο			0		_
	Reset value	0 0 0 0	0 0	0 0	0	0 0	0 0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x0084	DMA_S4FCR					F	Reserve	ed										FEIE	Reserved	F	S[2:	_	SIQWQ	FT [1:	
	Reset value																	0	Re	1	0	0	0	0	1
0x0088	DMA_S5CR	Reserved	CHSEL[2:0]	MBURST[1:0]		PBURST[1:0]	ACK	DBM	PL[1:0]		PINCOS	MSIZE[1:0]	DC17E14-01	r 312E[1.0]	MINC	PINC	CIRC	DIB[1:0]	5	PFCTRL	TCIE	HTIE	TEIE	DMEIE	EN
	Reset value		0 0	0 0	0	0 0	0 0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0
	DMA_S5NDTR				<u> </u>				<u> </u>							N	IDT	[15:	.]						
0x008C	Reset value			Rese	ervec	d				-	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0
	DMA_S5PAR								Р	PA[3	1:0]			<u> </u>			<b>I</b>	l	l	1					
0x0090	Reset value	0 0 0 0	0 0	0 0	0	0 0	0 0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0
0,,0004	DMA_S5M0AR		1 1 1	U	11				М	0A[	31:0	]						l		1					
0x0094	Reset value	0 0 0 0	0 0	0 0	0	0 0	0 0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x0098	DMA_S5M1AR		1 1 1						М	1A[	31:0														
0,0030	Reset value	0 0 0 0	0 0	0 0	0	0 0	0 0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x009C	DMA_S5FCR			_		F	Reserve	ed										FEIE	Reserved	F	S[2:	0]	DMDIS	FT [1:	H 0]
	Reset value																	0	Re	1	0	0	0	0	1
0x00A0	DMA_S6CR	Reserved	CHSEL[2:0]	ABUBST[1-0]		PBURST[1:0]	ACK	DBM	PL[1:0]		PINCOS	MSIZE[1:0]	DC17E[4-0]	r 312E[1.0]	MINC	PINC	CIRC	DIRIT-01	5	PFCTRL	TCIE	HTIE	TEIE	DMEIE	EN
	Reset value		0 0	0 0	0	0 0	0 0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0
0.0044	DMA_S6NDTR				11				<u> </u>				<u> </u>	<u> </u>		١	IDT	[15:	.]	1					
0x00A4	Reset value			Rese	ervec	d				-	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0
00040	DMA_S6PAR								P	PA[3	1:0]		1	<u> </u>			<b>I</b>	l	l	<u> </u>	ш				
0x00A8	Reset value	0 0 0 0	0 0	0 0	0	0 0	0 0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x00AC	DMA_S6M0AR		1 1		11				М	0A[	31:0	]													
UXUUAC	Reset value	0 0 0 0	0 0	0 0	0	0 0	0 0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x00B0	DMA_S6M1AR								М	1A[	31:0														
0,0000	Reset value	0 0 0 0	0 0	0 0	0	0 0	0 0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x00B4	DMA_S6FCR					F	Reserve	ed										FEIE	Reserved	F	S[2:	0]	SIGWG	FT [1:	
	Reset value																	0	Res	1	0	0	0	0	1

Table 31. DMA register map and reset values (continued)

Offset	Register	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	6	8	7	9	5	4	3	2	-	0
0x00B8	DMA_S7CR	F	Rese	rve	d		CHSEL[2:0]		MBURST[1:0]		PBURST[1:0]		ACK	CT	DBM	PI [1:0]	[O: -] .	PINCOS	MSIZE[1:0]		PSIZE[1:0]	[0]	MINC	PINC	CIRC	DIR[1:0]		PFCTRL	TCIE	HTIE	TEIE	DMEIE	EN
	Reset value					0	0	0	0	0	0	0	0	0	0	0	0		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x00BC	DMA_S7NDTR							В	lese	rvo	٠.													١	IDT[	15:	.]						
OXOODC	Reset value								1636	ive	J							0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x00C0	DMA_S7PAR															- 1	PA[3	31:0	]														
0.0000	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x00C4	DMA_S7M0AR															N	10A[	[31:	0]														
0.00004	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x00C8	DMA_S7M1AR															N	11A[	[31:	0]					•									
0,0000	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x00CC	DMA_S7FCR											F	Rese	rve	d											FEIE	served	F	S[2:	0]	DMDIS	FT [1:	ΓΗ :0]
	Reset value																									0	Re	1	0	0	0	0	1

Refer to *Table 2 on page 50* for the register boundary addresses.

# 9 Interrupts and events

This Section applies to the whole STM32F4xx family, unless otherwise specified.

# 9.1 Nested vectored interrupt controller (NVIC)

## 9.1.1 NVIC features

The nested vector interrupt controller NVIC includes the following features:

- 82 maskable interrupt channels (not including the 16 interrupt lines of Cortex<sup>™</sup>-M4F)
- 16 programmable priority levels (4 bits of interrupt priority are used)
- low-latency exception and interrupt handling
- power management control
- implementation of system control registers

The NVIC and the processor core interface are closely coupled, which enables low latency interrupt processing and efficient processing of late arriving interrupts.

All interrupts including the core exceptions are managed by the NVIC. For more information on exceptions and NVIC programming see *Chapter 5: Exceptions & Chapter 8: Nested Vectored Interrupt Controller* in the ARM *Cortex<sup>TM</sup>-M4F Technical Reference Manual.* 

## 9.1.2 SysTick calibration value register

The SysTick calibration value is fixed to 18750, which gives a reference time base of 1 ms with the SysTick clock set to 18.75 MHz (HCLK/8, with HCLK set to 150 MHz).

# 9.1.3 Interrupt and exception vectors

See *Table 32*, for the vector table for the STM32F40x and STM32F41x devices.

# 9.2 External interrupt/event controller (EXTI)

The external interrupt/event controller consists of up to 23 edge detectors for generating event/interrupt requests. Each input line can be independently configured to select the type (interrupt or event) and the corresponding trigger event (rising or falling or both). Each line can also masked independently. A pending register maintains the status line of the interrupt requests.

Table 32. Vector table

Position	Priority	Type of priority	Acronym	Description	Address
	-	-	-	Reserved	0x0000 0000
	-3	fixed	Reset	Reset	0x0000 0004

Table 32. Vector table (continued)

Iabit	· 02.	100101	table (continued)	,	
Position	Priority	Type of priority	Acronym	Description	Address
	-2	fixed	NMI	Non maskable interrupt. The RCC Clock Security System (CSS) is linked to the NMI vector.	0x0000 0008
	-1	fixed	HardFault	All class of fault	0x0000 000C
	0	settable	MemManage	Memory management	0x0000 0010
	1	settable	BusFault	Pre-fetch fault, memory access fault	0x0000 0014
	2	settable	UsageFault	Undefined instruction or illegal state	0x0000 0018
	1	1	-	Reserved	0x0000 001C - 0x0000 002B
	3	settable	SVCall	System service call via SWI instruction	0x0000 002C
	4	settable	Debug Monitor	Debug Monitor	0x0000 0030
	-	-	-	Reserved	0x0000 0034
	5	settable	PendSV	Pendable request for system service	0x0000 0038
	6	settable	SysTick	System tick timer	0x0000 003C
0	7	settable	WWDG	Window Watchdog interrupt	0x0000 0040
1	8	settable	PVD	PVD through EXTI line detection interrupt	0x0000 0044
2	9	settable	TAMP_STAMP	Tamper and TimeStamp interrupts through the EXTI line	0x0000 0048
3	10	settable	RTC_WKUP	RTC Wakeup interrupt through the EXTI line	0x0000 004C
4	11	settable	FLASH	Flash global interrupt	0x0000 0050
5	12	settable	RCC	RCC global interrupt	0x0000 0054
6	13	settable	EXTI0	EXTI Line0 interrupt	0x0000 0058
7	14	settable	EXTI1	EXTI Line1 interrupt	0x0000 005C
8	15	settable	EXTI2	EXTI Line2 interrupt	0x0000 0060
9	16	settable	EXTI3	EXTI Line3 interrupt	0x0000 0064
10	17	settable	EXTI4	EXTI Line4 interrupt	0x0000 0068
11	18	settable	DMA1_Stream0	DMA1 Stream0 global interrupt	0x0000 006C
12	19	settable	DMA1_Stream1	DMA1 Stream1 global interrupt	0x0000 0070
13	20	settable	DMA1_Stream2	DMA1 Stream2 global interrupt	0x0000 0074
14	21	settable	DMA1_Stream3	DMA1 Stream3 global interrupt	0x0000 0078
15	22	settable	DMA1_Stream4	DMA1 Stream4 global interrupt	0x0000 007C
			•		

Table 32. Vector table (continued)

			table (continued)		
Position	Priority	Type of priority	Acronym	Description	Address
16	23	settable	DMA1_Stream5	DMA1 Stream5 global interrupt	0x0000 0080
17	24	settable	DMA1_Stream6	DMA1 Stream6 global interrupt	0x0000 0084
18	25	settable	ADC	ADC1, ADC2 and ADC3 global interrupts	0x0000 0088
19	26	settable	CAN1_TX	CAN1 TX interrupts	0x0000 008C
20	27	settable	CAN1_RX0	CAN1 RX0 interrupts	0x0000 0090
21	28	settable	CAN1_RX1	CAN1 RX1 interrupt	0x0000 0094
22	29	settable	CAN1_SCE	CAN1 SCE interrupt	0x0000 0098
23	30	settable	EXTI9_5	EXTI Line[9:5] interrupts	0x0000 009C
24	31	settable	TIM1_BRK_TIM9	TIM1 Break interrupt and TIM9 global interrupt	0x0000 00A0
25	32	settable	TIM1_UP_TIM10	TIM1 Update interrupt and TIM10 global interrupt	0x0000 00A4
26	33	settable	TIM1_TRG_COM_T IM11	TIM1 Trigger and Commutation interrupts and TIM11 global interrupt	0x0000 00A8
27	34	settable	TIM1_CC	TIM1 Capture Compare interrupt	0x0000 00AC
28	35	settable	TIM2	TIM2 global interrupt	0x0000 00B0
29	36	settable	ТІМЗ	TIM3 global interrupt	0x0000 00B4
30	37	settable	TIM4	TIM4 global interrupt	0x0000 00B8
31	38	settable	I2C1_EV	I <sup>2</sup> C1 event interrupt	0x0000 00BC
32	39	settable	I2C1_ER	I <sup>2</sup> C1 error interrupt	0x0000 00C0
33	40	settable	I2C2_EV	I <sup>2</sup> C2 event interrupt	0x0000 00C4
34	41	settable	I2C2_ER	I <sup>2</sup> C2 error interrupt	0x0000 00C8
35	42	settable	SPI1	SPI1 global interrupt	0x0000 00CC
36	43	settable	SPI2	SPI2 global interrupt	0x0000 00D0
37	44	settable	USART1	USART1 global interrupt	0x0000 00D4
38	45	settable	USART2	USART2 global interrupt	0x0000 00D8
39	46	settable	USART3	USART3 global interrupt	0x0000 00DC
40	47	settable	EXTI15_10	EXTI Line[15:10] interrupts	0x0000 00E0
41	48	settable	RTC_Alarm	RTC Alarms (A and B) through EXTI line interrupt	0x0000 00E4
42	49	settable	OTG_FS WKUP	USB On-The-Go FS Wakeup through EXTI line interrupt	0x0000 00E8
43	50	settable	TIM8_BRK_TIM12	TIM8 Break interrupt and TIM12 global interrupt	0x0000 00EC

Table 32. Vector table (continued)

Position	Priority	Type of priority	Acronym	Description	Address
44	51	settable	TIM8_UP_TIM13	TIM8 Update interrupt and TIM13 global interrupt	0x0000 00F0
45	52	settable	TIM8_TRG_COM_ TIM14	TIM8 Trigger and Commutation interrupts and TIM14 global interrupt	0x0000 00F4
46	53	settable	TIM8_CC	TIM8 Capture Compare interrupt	0x0000 00F8
47	54	settable	DMA1_Stream7	DMA1 Stream7 global interrupt	0x0000 00FC
48	55	settable	FSMC	FSMC global interrupt	0x0000 0100
49	56	settable	SDIO	SDIO global interrupt	0x0000 0104
50	57	settable	TIM5	TIM5 global interrupt	0x0000 0108
51	58	settable	SPI3	SPI3 global interrupt	0x0000 010C
52	59	settable	UART4	UART4 global interrupt	0x0000 0110
53	60	settable	UART5	UART5 global interrupt	0x0000 0114
54	61	settable	TIM6_DAC	TIM6 global interrupt, DAC1 and DAC2 underrun error interrupts	0x0000 0118
55	62	settable	TIM7	TIM7 global interrupt	0x0000 011C
56	63	settable	DMA2_Stream0	DMA2 Stream0 global interrupt	0x0000 0120
57	64	settable	DMA2_Stream1	DMA2 Stream1 global interrupt	0x0000 0124
58	65	settable	DMA2_Stream2	DMA2 Stream2 global interrupt	0x0000 0128
59	66	settable	DMA2_Stream3	DMA2 Stream3 global interrupt	0x0000 012C
60	67	settable	DMA2_Stream4	DMA2 Stream4 global interrupt	0x0000 0130
61	68	settable	ETH	Ethernet global interrupt	0x0000 0134
62	69	settable	ETH_WKUP	Ethernet Wakeup through EXTI line interrupt	0x0000 0138
63	70	settable	CAN2_TX	CAN2 TX interrupts	0x0000 013C
64	71	settable	CAN2_RX0	CAN2 RX0 interrupts	0x0000 0140
65	72	settable	CAN2_RX1	CAN2 RX1 interrupt	0x0000 0144
66	73	settable	CAN2_SCE	CAN2 SCE interrupt	0x0000 0148
67	74	settable	OTG_FS	USB On The Go FS global interrupt	0x0000 014C
68	75	settable	DMA2_Stream5	DMA2 Stream5 global interrupt	0x0000 0150
69	76	settable	DMA2_Stream6	DMA2 Stream6 global interrupt	0x0000 0154
70	77	settable	DMA2_Stream7	DMA2 Stream7 global interrupt	0x0000 0158
71	78	settable	USART6	USART6 global interrupt	0x0000 015C
72	79	settable	I2C3_EV	I <sup>2</sup> C3 event interrupt	0x0000 0160

Table 32. Vector table (continued)

Position	Priority	Type of priority	Acronym	Description	Address
73	80	settable	I2C3_ER	I <sup>2</sup> C3 error interrupt	0x0000 0164
74	81	settable	OTG_HS_EP1_OUT	USB On The Go HS End Point 1 Out global interrupt	0x0000 0168
75	82	settable	OTG_HS_EP1_IN	USB On The Go HS End Point 1 In global interrupt	0x0000 016C
76	83	settable	OTG_HS_WKUP	USB On The Go HS Wakeup through EXTI interrupt	0x0000 0170
77	84	settable	OTG_HS	USB On The Go HS global interrupt	0x0000 0174
78	85	settable	DCMI	DCMI global interrupt	0x0000 0178
79	86	settable	CRYP	CRYP crypto global interrupt	0x0000 017C
80	87	settable	HASH_RNG	Hash and Rng global interrupt	0x0000 0180
81	88	settable	FPU	FPU global interrupt	0x0000 0184

## 9.2.1 EXTI main features

The main features of the EXTI controller are the following:

- independent trigger and mask on each interrupt/event line
- dedicated status bit for each interrupt line
- generation of up to 23 software event/interrupt requests
- detection of external signals with a pulse width lower than the APB2 clock period. Refer
  to the electrical characteristics section of the STM32F4xx datasheets for details on this
  parameter.

## 9.2.2 EXTI block diagram

Figure 26 shows the block diagram.

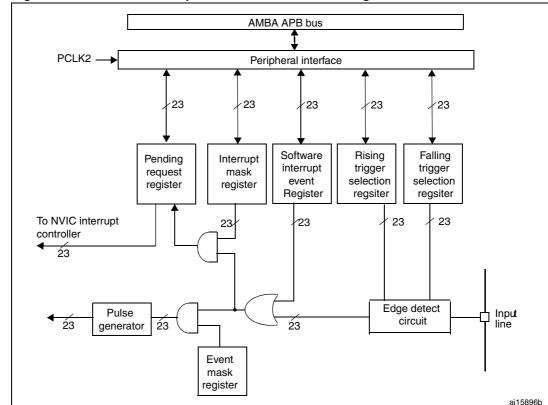


Figure 26. External interrupt/event controller block diagram

## 9.2.3 Wakeup event management

The STM32F4xx are able to handle external or internal events in order to wake up the core (WFE). The wakeup event can be generated either by:

- enabling an interrupt in the peripheral control register but not in the NVIC, and enabling
  the SEVONPEND bit in the Cortex™-M4F System Control register. When the MCU
  resumes from WFE, the peripheral interrupt pending bit and the peripheral NVIC IRQ
  channel pending bit (in the NVIC interrupt clear pending register) have to be cleared.
- or configuring an external or internal EXTI line in event mode. When the CPU resumes from WFE, it is not necessary to clear the peripheral interrupt pending bit or the NVIC IRQ channel pending bit as the pending bit corresponding to the event line is not set.

To use an external line as a wakeup event, refer to Section 9.2.4: Functional description.

## 9.2.4 Functional description

To generate the interrupt, the interrupt line should be configured and enabled. This is done by programming the two trigger registers with the desired edge detection and by enabling the interrupt request by writing a '1' to the corresponding bit in the interrupt mask register. When the selected edge occurs on the external interrupt line, an interrupt request is generated. The pending bit corresponding to the interrupt line is also set. This request is reset by writing a '1' in the pending register.

To generate the event, the event line should be configured and enabled. This is done by programming the two trigger registers with the desired edge detection and by enabling the

event request by writing a '1' to the corresponding bit in the event mask register. When the selected edge occurs on the event line, an event pulse is generated. The pending bit corresponding to the event line is not set.

An interrupt/event request can also be generated by software by writing a '1' in the software interrupt/event register.

### Hardware interrupt selection

To configure the 23 lines as interrupt sources, use the following procedure:

- Configure the mask bits of the 23 interrupt lines (EXTI\_IMR)
- Configure the Trigger selection bits of the interrupt lines (EXTI\_RTSR and EXTI\_FTSR)
- Configure the enable and mask bits that control the NVIC IRQ channel mapped to the
  external interrupt controller (EXTI) so that an interrupt coming from one of the 23 lines
  can be correctly acknowledged.

### Hardware event selection

To configure the 23 lines as event sources, use the following procedure:

- Configure the mask bits of the 23 event lines (EXTI\_EMR)
- Configure the Trigger selection bits of the event lines (EXTI\_RTSR and EXTI\_FTSR)

## Software interrupt/event selection

The 23 lines can be configured as software interrupt/event lines. The following is the procedure to generate a software interrupt.

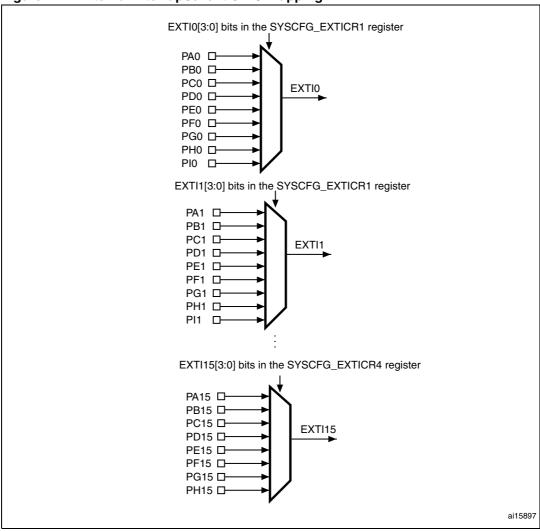
- Configure the mask bits of the 23 interrupt/event lines (EXTI\_IMR, EXTI\_EMR)
- Set the required bit in the software interrupt register (EXTI\_SWIER)

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## 9.2.5 External interrupt/event line mapping

Up to 140 GPIOs are connected to the 16 external interrupt/event lines in the following manner:

Figure 27. External interrupt/event GPIO mapping



The seven other EXTI lines are connected as follows:

- EXTI line 16 is connected to the PVD output
- EXTI line 17 is connected to the RTC Alarm event
- EXTI line 18 is connected to the USB OTG FS Wakeup event
- EXTI line 19 is connected to the Ethernet Wakeup event
- EXTI line 20 is connected to the USB OTG HS (configured in FS) Wakeup event
- EXTI line 21 is connected to the RTC Tamper and TimeStamp events
- EXTI line 22 is connected to the RTC Wakeup event

# 9.3 EXTI registers

Refer to Section 1.1 on page 46 for a list of abbreviations used in register descriptions.

# 9.3.1 Interrupt mask register (EXTI\_IMR)

Address offset: 0x00

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
				Reserve	4				MR22	MR21	MR20	MR19	MR18	MR17	MR16
				neserve	u				rw						
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
MR15	MR14	MR13	MR12	MR11	MR10	MR9	MR8	MR7	MR6	MR5	MR4	MR3	MR2	MR1	MR0
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:23 Reserved, must be kept at reset value.

Bits 22:0 MRx: Interrupt mask on line x

0: Interrupt request from line x is masked1: Interrupt request from line x is not masked

# 9.3.2 Event mask register (EXTI\_EMR)

Address offset: 0x04 Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
				Reserve	4				MR22	MR21	MR20	MR19	MR18	MR17	MR16
				neserve	J				rw						
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
MR15	MR14	MR13	MR12	MR11	MR10	MR9	MR8	MR7	MR6	MR5	MR4	MR3	MR2	MR1	MR0
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:23 Reserved, must be kept at reset value.

Bits 22:0 MRx: Event mask on line x

0: Event request from line x is masked1: Event request from line x is not masked

# 9.3.3 Rising trigger selection register (EXTI\_RTSR)

Address offset: 0x08 Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
				Reserve	d				TR22	TR21	TR20	TR19	TR18	TR17	TR16
				neserve	u				rw						

RM0090 Interrupts and events

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
TR15	TR14	TR13	TR12	TR11	TR10	TR9	TR8	TR7	TR6	TR5	TR4	TR3	TR2	TR1	TR0
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:23 Reserved, must be kept at reset value.

Bits 22:0 TRx: Rising trigger event configuration bit of line x

0: Rising trigger disabled (for Event and Interrupt) for input line 1: Rising trigger enabled (for Event and Interrupt) for input line

Note:

The external wakeup lines are edge triggered, no glitch must be generated on these lines. If a rising edge occurs on the external interrupt line while writing to the EXTI\_RTSR register, the pending bit is be set.

Rising and falling edge triggers can be set for the same interrupt line. In this configuration, both generate a trigger condition.

# 9.3.4 Falling trigger selection register (EXTI\_FTSR)

Address offset: 0x0C Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
				Reserve	4				TR22	TR21	TR20	TR19	TR18	TR17	TR16
				neserve	u				rw						
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
TR15	TR14	TR13	TR12	TR11	TR10	TR9	TR8	TR7	TR6	TR5	TR4	TR3	TR2	TR1	TR0
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:23 Reserved, must be kept at reset value.

Bits 22:0 TRx: Falling trigger event configuration bit of line x

0: Falling trigger disabled (for Event and Interrupt) for input line

1: Falling trigger enabled (for Event and Interrupt) for input line.

Note:

The external wakeup lines are edge triggered, no glitch must be generated on these lines. If a falling edge occurs on the external interrupt line while writing to the EXTI\_FTSR register, the pending bit is not set.

Rising and falling edge triggers can be set for the same interrupt line. In this configuration, both generate a trigger condition.

# 9.3.5 Software interrupt event register (EXTI\_SWIER)

Address offset: 0x10 Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
				December	J				SWIER 22	SWIER 21	SWIER 20	SWIER 19	SWIER 18	SWIER 17	SWIER 16
				Reserved	1				rw						

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
SWIER 15	SWIER 14	SWIER 13	SWIER 12	SWIER 11	SWIER 10	SWIER 9	SWIER 8	SWIER 7	SWIER 6	SWIER 5	SWIER 4	SWIER 3	SWIER 2	SWIER 1	SWIER 0
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:23 Reserved, must be kept at reset value.

### Bits 22:0 SWIERx: Software Interrupt on line x

Writing a 1 to this bit when it is at 0 sets the corresponding pending bit in EXTI\_PR. If the interrupt is enabled on this line on the EXTI\_IMR and EXTI\_EMR, an interrupt request is generated.

This bit is cleared by clearing the corresponding bit in EXTI\_PR (by writing a 1 to the bit).

## 9.3.6 Pending register (EXTI\_PR)

Address offset: 0x14 Reset value: undefined

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
				Reserve	4				PR22	PR21	PR20	PR19	PR18	PR17	PR16
				neserve	J				rc_w1						
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
PR15	PR14	PR13	PR12	PR11	PR10	PR9	PR8	PR7	PR6	PR5	PR4	PR3	PR2	PR1	PR0
rc_w1	rc_w1	rc_w1	rc_w1	rc_w1	rc_w1	rc_w1	rc_w1	rc_w1	rc_w1	rc_w1	rc_w1	rc_w1	rc_w1	rc_w1	rc_w1

Bits 31:23 Reserved, must be kept at reset value.

Bits 22:0 PRx: Pending bit

0: No trigger request occurred

1: selected trigger request occurred

This bit is set when the selected edge event arrives on the external interrupt line. This bit is cleared by writing a 1 to the bit or by changing the sensitivity of the edge detector.

# 9.3.7 EXTI register map

Table 33 gives the EXTI register map and the reset values.

Table 33. External interrupt/event controller register map and reset values

Offset	Register	31 30 28 27 27 26 25 25 24 23	2 2 2 3 4 2 9 6 7 7 2 3 4 7 9 9 6 7 9 7 7 7 8 7 7 7 8 7 7 7 8 7 7 7 8 7
0x00	EXTI_IMR	Reserved	MR[22:0]
	Reset value		0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
0x04	EXTI_EMR	Reserved	MR[22:0]
	Reset value		0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
0x08	EXTI_RTSR	Reserved	TR[22:0]
	Reset value		0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0

Table 33. External interrupt/event controller register map and reset values

					_ •									J				1															
Offset	Register	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	6	8	7	9	5	4	3	2	1	0
0x0C	EXTI_FTSR				Re	serv	ed														TF	R[22	:0]										
	Reset value									-	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x10 EXTI_SWIER Reserved							SWIER[22:0]																										
	Reset value									-	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x14	EXTI_PR		Reserved															PF	R[22	::0]													
	Reset value										0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

Refer to *Table 2 on page 50* for the register boundary addresses.

# 10 Analog-to-digital converter (ADC)

This section applies to the whole STM32F4xx family, unless otherwise specified.

## 10.1 ADC introduction

The 12-bit ADC is a successive approximation analog-to-digital converter. It has up to 19 multiplexed channels allowing it to measure signals from 16 external sources, two internal sources, and the  $V_{BAT}$  channel. The A/D conversion of the channels can be performed in single, continuous, scan or discontinuous mode. The result of the ADC is stored into a left-or right-aligned 16-bit data register.

The analog watchdog feature allows the application to detect if the input voltage goes beyond the user-defined, higher or lower thresholds.

## 10.2 ADC main features

- 12-bit, 10-bit, 8-bit or 6-bit configurable resolution
- Interrupt generation at the end of conversion, end of injected conversion, and in case of analog watchdog or overrun events
- Single and continuous conversion modes
- Scan mode for automatic conversion of channel 0 to channel 'n'
- Data alignment with in-built data coherency
- Channel-wise programmable sampling time
- External trigger option with configurable polarity for both regular and injected conversions
- Discontinuous mode
- Dual/Triple mode (on devices with 2 ADCs or more)
- Configurable DMA data storage in Dual/Triple ADC mode
- Configurable delay between conversions in Dual/Triple interleaved mode
- ADC conversion type (refer to the datasheets)
- ADC supply requirements: 2.4 V to 3.6 V at full speed and down to 1.8 V at slower speed
- ADC input range: V<sub>REF</sub> ≤ V<sub>IN</sub> ≤ V<sub>REF</sub>
- DMA request generation during regular channel conversion

Figure 28 shows the block diagram of the ADC.

Note:  $V_{REF-}$ , if available (depending on package), must be tied to  $V_{SSA}$ .

# 10.3 ADC functional description

Figure 28 shows a single ADC block diagram and Table 34 gives the ADC pin description.

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Figure 28. Single ADC block diagram

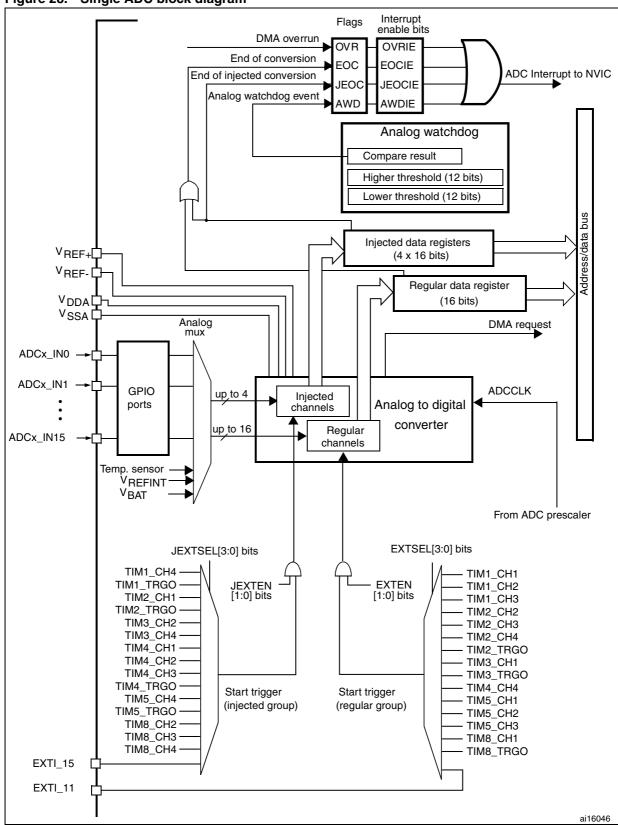


Table 34. ADC pins

Name	Signal type	Remarks
V <sub>REF+</sub>	Input, analog reference positive	The higher/positive reference voltage for the ADC, 1.8 V $\leq$ V <sub>REF+</sub> $\leq$ V <sub>DDA</sub>
$V_{DDA}$	Input, analog supply	Analog power supply equal to $V_{DD}$ and 2.4 $V \le V_{DDA} \le V_{DD}$ (3.6 V) for full speed 1.8 $V \le V_{DDA} \le V_{DD}$ (3.6 V) for reduced speed
V <sub>REF</sub> _	Input, analog reference negative	The lower/negative reference voltage for the ADC, $V_{REF-} = V_{SSA}$
V <sub>SSA</sub>	Input, analog supply ground	Ground for analog power supply equal to V <sub>SS</sub>
ADCx_IN[15:0]	Analog input signals	16 analog input channels

## 10.3.1 ADC on-off control

The ADC is powered on by setting the ADON bit in the ADC\_CR2 register. When the ADON bit is set for the first time, it wakes up the ADC from the Power-down mode.

Conversion starts when either the SWSTART or the JSWSTART bit is set.

You can stop conversion and put the ADC in power down mode by clearing the ADON bit. In this mode the ADC consumes almost no power (only a few  $\mu$ A).

### 10.3.2 ADC clock

The ADC features two clock schemes:

- Clock for the analog circuitry: ADCCLK, common to all ADCs
   This clock is generated from the APB2 clock divided by a programmable prescaler that allows the ADC to work at f<sub>PCLK2</sub>/2, /4, /6 or /8. Refer to the datasheets for the maximum value of ADCCLK.
- Clock for the digital interface (used for registers read/write access)
   This clock is equal to the APB2 clock. The digital interface clock can be enabled/disabled individually for each ADC through the RCC APB2 peripheral clock enable register (RCC\_APB2ENR).

## 10.3.3 Channel selection

There are 16 multiplexed channels. It is possible to organize the conversions in two groups: regular and injected. A group consists of a sequence of conversions that can be done on any channel and in any order. For instance, it is possible to implement the conversion sequence in the following order: ADC\_IN3, ADC\_IN8, ADC\_IN2, ADC\_IN2, ADC\_IN0, ADC\_IN2, ADC\_IN2, ADC\_IN15.

- A regular group is composed of up to 16 conversions. The regular channels and their
  order in the conversion sequence must be selected in the ADC\_SQRx registers. The
  total number of conversions in the regular group must be written in the L[3:0] bits in the
  ADC\_SQR1 register.
- An injected group is composed of up to 4 conversions. The injected channels and their order in the conversion sequence must be selected in the ADC JSQR register.

The total number of conversions in the injected group must be written in the L[1:0] bits in the ADC\_JSQR register.

If the ADC\_SQRx or ADC\_JSQR registers are modified during a conversion, the current conversion is reset and a new start pulse is sent to the ADC to convert the newly chosen group.

## Temperature sensor, V<sub>REFINT</sub> and V<sub>BAT</sub> internal channels

The temperature sensor is connected to channel ADC1\_IN16 and the internal reference voltage V<sub>REFINT</sub> is connected to ADC1\_IN17. These two internal channels can be selected and converted as injected or regular channels.

The V<sub>BAT</sub> channel is connected to channel ADC1\_IN18. It can also be converted as an injected or regular channel.

Note:

The temperature sensor,  $V_{REFINT}$  and the  $V_{BAT}$  channel are available only on the master ADC1 peripheral.

# 10.3.4 Single conversion mode

In Single conversion mode the ADC does one conversion. This mode is started with the CONT bit at 0 by either:

- setting the SWSTART bit in the ADC\_CR2 register (for a regular channel only)
- setting the JSWSTART bit (for an injected channel)
- external trigger (for a regular or injected channel)

Once the conversion of the selected channel is complete:

- If a regular channel was converted:
  - The converted data are stored into the 16-bit ADC\_DR register
  - The EOC (end of conversion) flag is set
  - An interrupt is generated if the EOCIE bit is set
- If an injected channel was converted:
  - The converted data are stored into the 16-bit ADC JDR1 register
  - The JEOC (end of conversion injected) flag is set
  - An interrupt is generated if the JEOCIE bit is set

Then the ADC stops.

### 10.3.5 Continuous conversion mode

In continuous conversion mode, the ADC starts a new conversion as soon as it finishes one. This mode is started with the CONT bit at 1 either by external trigger or by setting the SWSTRT bit in the ADC\_CR2 register (for regular channels only).

After each conversion:

- If a regular group of channels was converted:
  - The last converted data are stored into the 16-bit ADC DR register
  - The EOC (end of conversion) flag is set
  - An interrupt is generated if the EOCIE bit is set

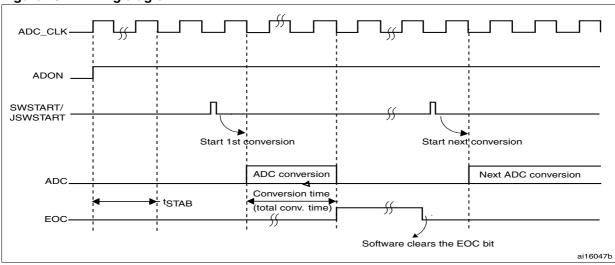
Note:

Injected channels cannot be converted continuously. The only exception is when an injected channel is configured to be converted automatically after regular channels in continuous mode (using JAUTO bit), refer to Auto-injection section).

## 10.3.6 Timing diagram

As shown in *Figure 29*, the ADC needs a stabilization time of t<sub>STAB</sub> before it starts converting accurately. After the start of the ADC conversion and after 15 clock cycles, the EOC flag is set and the 16-bit ADC data register contains the result of the conversion.





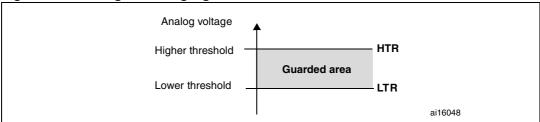
## 10.3.7 Analog watchdog

The AWD analog watchdog status bit is set if the analog voltage converted by the ADC is below a lower threshold or above a higher threshold. These thresholds are programmed in the 12 least significant bits of the ADC\_HTR and ADC\_LTR 16-bit registers. An interrupt can be enabled by using the AWDIE bit in the ADC\_CR1 register.

The threshold value is independent of the alignment selected by the ALIGN bit in the ADC\_CR2 register. The analog voltage is compared to the lower and higher thresholds before alignment.

*Table 35* shows how the ADC\_CR1 register should be configured to enable the analog watchdog on one or more channels.

Figure 30. Analog watchdog's guarded area



Channels guarded by the analog	ADC_CR1 register control bits (x = don't care)										
watchdog	AWDSGL bit	AWDEN bit	JAWDEN bit								
None	x	0	0								
All injected channels	0	0	1								
All regular channels	0	1	0								
All regular and injected channels	0	1	1								
Single <sup>(1)</sup> injected channel	1	0	1								
Single <sup>(1)</sup> regular channel	1	1	0								
Single (1) regular or injected channel	1	1	1								

Table 35. Analog watchdog channel selection

### 10.3.8 Scan mode

This mode is used to scan a group of analog channels.

The Scan mode is selected by setting the SCAN bit in the ADC\_CR1 register. Once this bit has been set, the ADC scans all the channels selected in the ADC\_SQRx registers (for regular channels) or in the ADC\_JSQR register (for injected channels). A single conversion is performed for each channel of the group. After each end of conversion, the next channel in the group is converted automatically. If the CONT bit is set, regular channel conversion does not stop at the last selected channel in the group but continues again from the first selected channel.

If the DMA bit is set, the direct memory access (DMA) controller is used to transfer the data converted from the regular group of channels (stored in the ADC\_DR register) to SRAM after each regular channel conversion.

The EOC bit is set in the ADC\_SR register:

- At the end of each regular group sequence if the EOCS bit is cleared to 0
- At the end of each regular channel conversion if the EOCS bit is set to 1

The data converted from an injected channel are always stored into the ADC\_JDRx registers.

## 10.3.9 Injected channel management

## **Triggered injection**

To use triggered injection, the JAUTO bit must be cleared in the ADC\_CR1 register.

- 1. Start the conversion of a group of regular channels either by external trigger or by setting the SWSTART bit in the ADC\_CR2 register.
- 2. If an external injected trigger occurs or if the JSWSTART bit is set during the conversion of a regular group of channels, the current conversion is reset and the injected channel sequence switches to Scan-once mode.
- 3. Then, the regular conversion of the regular group of channels is resumed from the last interrupted regular conversion.
  - If a regular event occurs during an injected conversion, the injected conversion is not

<sup>1.</sup> Selected by the AWDCH[4:0] bits

interrupted but the regular sequence is executed at the end of the injected sequence. *Figure 31* shows the corresponding timing diagram.

Note:

When using triggered injection, one must ensure that the interval between trigger events is longer than the injection sequence. For instance, if the sequence length is 30 ADC clock cycles (that is two conversions with a sampling time of 3 clock periods), the minimum interval between triggers must be 31 ADC clock cycles.

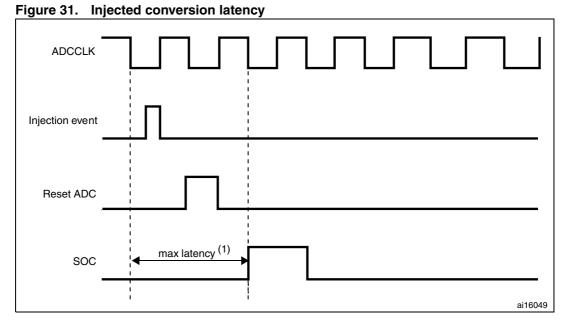
#### **Auto-injection**

If the JAUTO bit is set, then the channels in the injected group are automatically converted after the regular group of channels. This can be used to convert a sequence of up to 20 conversions programmed in the ADC\_SQRx and ADC\_JSQR registers.

In this mode, external trigger on injected channels must be disabled.

If the CONT bit is also set in addition to the JAUTO bit, regular channels followed by injected channels are continuously converted.

Note: It is not possible to use both the auto-injected and discontinuous modes simultaneously.



 The maximum latency value can be found in the electrical characteristics of the STM32F40x and STM32F41x datasheets.

#### 10.3.10 Discontinuous mode

#### Regular group

This mode is enabled by setting the DISCEN bit in the ADC\_CR1 register. It can be used to convert a short sequence of n conversions ( $n \le 8$ ) that is part of the sequence of conversions selected in the ADC\_SQRx registers. The value of n is specified by writing to the DISCNUM[2:0] bits in the ADC\_CR1 register.

When an external trigger occurs, it starts the next n conversions selected in the ADC\_SQRx registers until all the conversions in the sequence are done. The total sequence length is defined by the L[3:0] bits in the ADC\_SQR1 register.

#### Example:

n = 3, channels to be converted = 0, 1, 2, 3, 6, 7, 9, 10

1st trigger: sequence converted 0, 1, 2 2nd trigger: sequence converted 3, 6, 7

3rd trigger: sequence converted 9, 10 and an EOC event generated

4th trigger: sequence converted 0, 1, 2

Note: When a regular group is converted in discontinuous mode, no rollover occurs.

When all subgroups are converted, the next trigger starts the conversion of the first subgroup. In the example above, the 4th trigger reconverts the channels 0, 1 and 2 in the 1st subgroup.

### Injected group

This mode is enabled by setting the JDISCEN bit in the ADC\_CR1 register. It can be used to convert the sequence selected in the ADC\_JSQR register, channel by channel, after an external trigger event.

When an external trigger occurs, it starts the next channel conversions selected in the ADC\_JSQR registers until all the conversions in the sequence are done. The total sequence length is defined by the JL[1:0] bits in the ADC\_JSQR register.

#### Example:

n = 1, channels to be converted = 1, 2, 3

1st trigger: channel 1 converted 2nd trigger: channel 2 converted

3rd trigger: channel 3 converted and EOC and JEOC events generated

4th trigger: channel 1

Note:

When all injected channels are converted, the next trigger starts the conversion of the first injected channel. In the example above, the 4th trigger reconverts the 1st injected channel 1

It is not possible to use both the auto-injected and discontinuous modes simultaneously.

Discontinuous mode must not be set for regular and injected groups at the same time. Discontinuous mode must be enabled only for the conversion of one group.

# 10.4 Data alignment

The ALIGN bit in the ADC\_CR2 register selects the alignment of the data stored after conversion. Data can be right- or left-aligned as shown in *Figure 32* and *Figure 33*.

The converted data value from the injected group of channels is decreased by the userdefined offset written in the ADC\_JOFRx registers so the result can be a negative value. The SEXT bit represents the extended sign value.

For channels in a regular group, no offset is subtracted so only twelve bits are significant.

Figure 32. Right alignment of 12-bit data

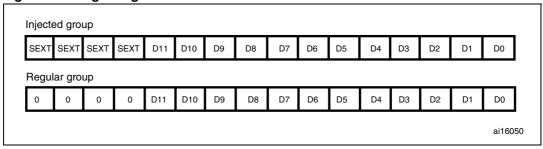
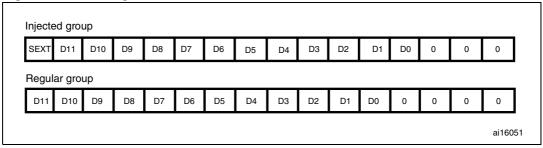
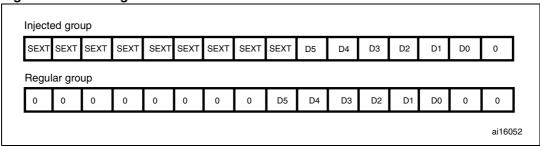


Figure 33. Left alignment of 12-bit data



Special case: when left-aligned, the data are aligned on a half-word basis except when the resolution is set to 6-bit. in that case, the data are aligned on a byte basis as shown in *Figure 34*.

Figure 34. Left alignment of 6-bit data



# 10.5 Channel-wise programmable sampling time

The ADC samples the input voltage for a number of ADCCLK cycles that can be modified using the SMP[2:0] bits in the ADC\_SMPR1 and ADC\_SMPR2 registers. Each channel can be sampled with a different sampling time.

The total conversion time is calculated as follows:

$$T_{conv}$$
 = Sampling time + 12 cycles

#### Example:

With ADCCLK = 30 MHz and sampling time = 3 cycles:

$$T_{conv} = 3 + 12 = 15$$
 cycles = 0.5  $\mu$ s with APB2 at 60 MHz

# 10.6 Conversion on external trigger and trigger polarity

Conversion can be triggered by an external event (e.g. timer capture, EXTI line). If the EXTEN[1:0] control bits (for a regular conversion) or JEXTEN[1:0] bits (for an injected conversion) are different from "0b00", then external events are able to trigger a conversion with the selected polarity. *Table 36* provides the correspondence between the EXTEN[1:0] and JEXTEN[1:0] values and the trigger polarity.

Table 36. Configuring the trigger polarity

Source	EXTEN[1:0] / JEXTEN[1:0]
Trigger detection disabled	00
Detection on the rising edge	01
Detection on the falling edge	10
Detection on both the rising and falling edges	11

Note:

The polarity of the external trigger can be changed on the fly.

The EXTSEL[3:0] and JEXTSEL[3:0] control bits are used to select which out of 16 possible events can trigger conversion for the regular and injected groups.

Table 37 gives the possible external trigger for regular conversion.

Table 37. External trigger for regular channels

Source	Туре	EXTSEL[3:0]
TIM1_CH1 event		0000
TIM1_CH2 event		0001
TIM1_CH3 event		0010
TIM2_CH2 event		0011
TIM2_CH3 event		0100
TIM2_CH4 event		0101
TIM2_TRGO event		0110
TIM3_CH1 event	Internal signal from on-chip timers	0111
TIM3_TRGO event		1000
TIM4_CH4 event		1001
TIM5_CH1 event		1010
TIM5_CH2 event		1011
TIM5_CH3 event		1100
TIM8_CH1 event		1101
TIM8_TRGO event		1110
EXTI line11	External pin	1111

*Table 38* gives the possible external trigger for injected conversion.

Table 38. External trigger for injected channels

Source	Connection type	JEXTSEL[3:0]
TIM1_CH4 event		0000
TIM1_TRGO event		0001
TIM2_CH1 event		0010
TIM2_TRGO event		0011
TIM3_CH2 event		0100
TIM3_CH4 event		0101
TIM4_CH1 event		0110
TIM4_CH2 event	Internal signal from on-chip timers	0111
TIM4_CH3 event		1000
TIM4_TRGO event		1001
TIM5_CH4 event		1010
TIM5_TRGO event		1011
TIM8_CH2 event		1100
TIM8_CH3 event		1101
TIM8_CH4 event		1110
EXTI line15	External pin	1111

Software source trigger events can be generated by setting SWSTART (for regular conversion) or JSWSTART (for injected conversion) in ADC\_CR2.

A regular group conversion can be interrupted by an injected trigger.

Note:

The trigger selection can be changed on the fly. However, when the selection changes, there is a time frame of 1 APB clock cycle during which the trigger detection is disabled. This is to avoid spurious detection during transitions.

### 10.7 Fast conversion mode

It is possible to perform faster conversion by reducing the ADC resolution. The RES bits are used to select the number of bits available in the data register. The minimum conversion time for each resolution is then as follows:

- 12 bits: 3 + 12 = 15 ADCCLK cycles
- 10 bits: 3 + 10 = 13 ADCCLK cycles
- 8 bits: 3 + 8 = 11 ADCCLK cycles
- 6 bits: 3 + 6 = 9 ADCCLK cycles

# 10.8 Data management

### 10.8.1 Using the DMA

Since converted regular channel values are stored into a unique data register, it is useful to use DMA for conversion of more than one regular channel. This avoids the loss of the data already stored in the ADC\_DR register.

When the DMA mode is enabled (DMA bit set to 1 in the ADC\_CR2 register), after each conversion of a regular channel, a DMA request is generated. This allows the transfer of the converted data from the ADC\_DR register to the destination location selected by the software.

Despite this, if data are lost (overrun), the OVR bit in the ADC\_SR register is set and an interrupt is generated (if the OVRIE enable bit is set). DMA transfers are then disabled and DMA requests are no longer accepted. In this case, if a DMA request is made, the regular conversion in progress is aborted and further regular triggers are ignored. It is then necessary to clear the OVR flag and the DMAEN bit in the used DMA stream, and to reinitialize both the DMA and the ADC to have the wanted converted channel data transferred to the right memory location. Only then can the conversion be resumed and the data transfer, enabled again. Injected channel conversions are not impacted by overrun errors.

When OVR = 1 in DMA mode, the DMA requests are blocked after the last valid data have been transferred, which means that all the data transferred to the RAM can be considered as valid.

At the end of the last DMA transfer (number of transfers configured in the DMA controller's DMA\_SxRTR register):

- No new DMA request is issued to the DMA controller if the DDS bit is cleared to 0 in the ADC\_CR2 register (this avoids generating an overrun error). However the DMA bit is not cleared by hardware. It must be written to 0, then to 1 to start a new transfer.
- Requests can continue to be generated if the DDS bit is set to 1. This allows configuring the DMA in double-buffer circular mode.

To recover the ADC from OVR state when the DMA is used, follow the steps below:

- 1. Reinitialize the DMA (adjust destination address and NDTR counter)
- 2. Clear the ADC OVR bit in ADC\_SR register
- 3. Trigger the ADC to start the conversion.

# 10.8.2 Managing a sequence of conversions without using the DMA

If the conversions are slow enough, the conversion sequence can be handled by the software. In this case the EOCS bit must be set in the ADC\_CR2 register for the EOC status bit to be set at the end of each conversion, and not only at the end of the sequence. When EOCS = 1, overrun detection is automatically enabled. Thus, each time a conversion is complete, EOC is set and the ADC\_DR register can be read. The overrun management is the same as when the DMA is used.

To recover the ADC from OVR state when the EOCS is set, follow the steps below:

- Clear the ADC OVR bit in ADC\_SR register
- 2. Trigger the ADC to start the conversion.

#### 10.8.3 Conversions without DMA and without overrun detection

It may be useful to let the ADC convert one or more channels without reading the data each time (if there is an analog watchdog for instance). For that, the DMA must be disabled (DMA=0) and the EOC bit must be set at the end of a sequence only (EOCS=0). In this configuration, overrun detection is disabled.

# 10.9 Multi ADC mode

In devices with two ADCs or more, the Dual (with two ADCs) and Triple (with three ADCs) ADC modes can be used (see *Figure 35*).

In multi ADC mode, the start of conversion is triggered alternately or simultaneously by the ADC1 master to the ADC2 and ADC3 slaves, depending on the mode selected by the MULTI[4:0] bits in the ADC\_CCR register.

Note:

In multi ADC mode, when configuring conversion trigger by an external event, the application must set trigger by the master only and disable trigger by slaves to prevent spurious triggers that would start unwanted slave conversions.

The four possible modes below are implemented:

- Injected simultaneous mode
- Regular simultaneous mode
- Interleaved mode
- Alternate trigger mode

It is also possible to use the previous modes combined in the following ways:

- Injected simultaneous mode + Regular simultaneous mode
- Regular simultaneous mode + Alternate trigger mode

Note:

In multi ADC mode, the converted data can be read on the multi-mode data register (ADC\_CDR). The status bits can be read in the multi-mode status register (ADC\_CSR).

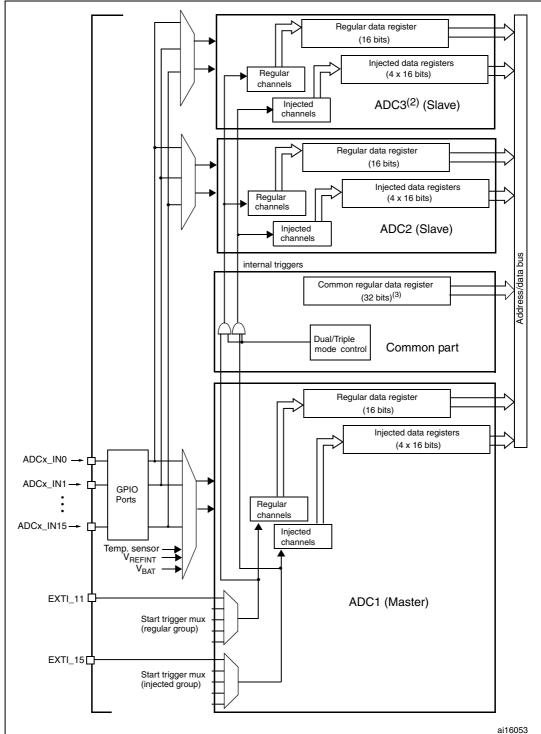


Figure 35. Multi ADC block diagram<sup>(1)</sup>

- 1. Although external triggers are present on ADC2 and ADC3 they are not shown in this diagram.
- 2. In the Dual ADC mode, the ADC3 slave part is not present.
- 3. In Triple ADC mode, the ADC common data register (ADC\_CDR) contains the ADC1, ADC2 and ADC3's regular converted data. All 32 register bits are used according to a selected storage order. In Dual ADC mode, the ADC common data register (ADC\_CDR) contains both the ADC1 and ADC2's regular converted data. All 32 register bits are used.

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DMA requests in Multi ADC mode:

In Multi ADC mode the DMA may be configured to transfer converted data in three different modes. In all cases, the DMA streams to use are those connected to the ADC:

 DMA mode 1: On each DMA request (one data item is available), a half-word representing an ADC-converted data item is transferred.

In Dual ADC mode, ADC1 data are transferred on the first request, ADC2 data are transferred on the second request and so on.

In Triple ADC mode, ADC1 data are transferred on the first request, ADC2 data are transferred on the second request and ADC3 data are transferred on the third request; the sequence is repeated. So the DMA first transfers ADC1 data followed by ADC2 data followed by ADC3 data and so on.

DMA mode 1 is used in regular simultaneous triple mode.

#### **Example:**

Regular simultaneous triple mode: 3 consecutive DMA requests are generated (one for each converted data item)

```
1st request: ADC_CDR[31:0] = ADC1_DR[15:0]
2nd request: ADC_CDR[31:0] = ADC2_DR[15:0]
3rd request: ADC_CDR[31:0] = ADC3_DR[15:0]
4th request: ADC_CDR[31:0] = ADC1_DR[15:0]
```

 DMA mode 2: On each DMA request (two data items are available) two half-words representing two ADC-converted data items are transferred as a word.

In Dual ADC mode, both ADC2 and ADC1 data are transferred on the first request (ADC2 data take the upper half-word and ADC1 data take the lower half-word) and so on.

In Triple ADC mode, three DMA requests are generated. On the first request, both ADC2 and ADC1 data are transferred (ADC2 data take the upper half-word and ADC1 data take the lower half-word). On the second request, both ADC1 and ADC3 data are transferred (ADC1 data take the upper half-word and ADC3 data take the lower half-word). On the third request, both ADC3 and ADC2 data are transferred (ADC3 data take the upper half-word and ADC2 data take the lower half-word) and so on.

DAM mode 2 is used in interleaved mode and in regular simultaneous mode (for Dual ADC mode only).

#### **Example:**

 a) Interleaved dual mode: a DMA request is generated each time 2 data items are available:

```
1st request: ADC_CDR[31:0] = ADC2_DR[15:0] | ADC1_DR[15:0] 2nd request: ADC_CDR[31:0] = ADC2_DR[15:0] | ADC1_DR[15:0]
```

b) Interleaved triple mode: a DMA request is generated each time 2 data items are available

```
1st request: ADC_CDR[31:0] = ADC2_DR[15:0] | ADC1_DR[15:0] 2nd request: ADC_CDR[31:0] = ADC1_DR[15:0] | ADC3_DR[15:0] 3rd request: ADC_CDR[31:0] = ADC3_DR[15:0] | ADC2_DR[15:0] 4th request: ADC_CDR[31:0] = ADC2_DR[15:0] | ADC1_DR[15:0]
```

DMA mode 3: This mode is similar to the DMA mode 2. The only differences are that the on each DMA request (two data items are available) two bytes representing two ADC converted data items are transferred as a half-word. The data transfer order is similar to that of the DMA mode 2.

DMA mode 3 is used in interleaved mode in 6-bit and 8-bit resolutions.

#### **Example:**

 a) Interleaved dual mode: a DMA request is generated each time 2 data items are available

b) Interleaved triple mode: a DMA request is generated each time 2 data items are available

```
1st request: ADC_CDR[15:0] = ADC2_DR[7:0] | ADC1_DR7:0]
2nd request: ADC_CDR[15:0] = ADC1_DR[7:0] | ADC3_DR[15:0]
3rd request: ADC_CDR[15:0] = ADC3_DR[7:0] | ADC2_DR[7:0]
4th request: ADC_CDR[15:0] = ADC2_DR[7:0] | ADC1_DR7:0]
```

**Overrun detection:** If an overrun is detected on one of the concerned ADCs (ADC1 and ADC2 in dual and triple modes, ADC3 in triple mode only), the DMA requests are no longer issued to ensure that all the data transferred to the RAM are valid. It may happen that the EOC bit corresponding to one ADC remains set because the data register of this ADC contains valid data.

### 10.9.1 Injected simultaneous mode

This mode converts an injected group of channels. The external trigger source comes from the injected group multiplexer of ADC1 (selected by the JEXTSEL[3:0] bits in the ADC1\_CR2 register). A simultaneous trigger is provided to ADC2 and ADC3.

Note:

Do not convert the same channel on the two/three ADCs (no overlapping sampling times for the two/three ADCs when converting the same channel).

In simultaneous mode, one must convert sequences with the same length or ensure that the interval between triggers is longer than the longer of the 2 sequences (Dual ADC mode) /3 sequences (Triple ADC mode). Otherwise, the ADC with the shortest sequence may restart while the ADC with the longest sequence is completing the previous conversions.

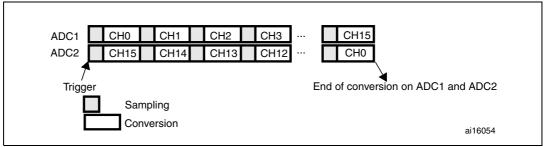
Regular conversions can be performed on one or all ADCs. In that case, they are independent of each other and are interrupted when an injected event occurs. They are resumed at the end of the injected conversion group.

#### **Dual ADC mode**

At the end of conversion event on ADC1 or ADC2:

- The converted data are stored into the ADC JDRx registers of each ADC interface.
- A JEOC interrupt is generated (if enabled on one of the two ADC interfaces) when the ADC1/ADC2's injected channels have all been converted.

Figure 36. Injected simultaneous mode on 4 channels: dual ADC mode

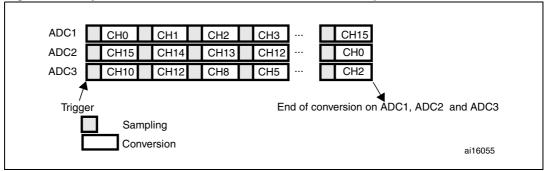


### **Triple ADC mode**

At the end of conversion event on ADC1, ADC2 or ADC3:

- The converted data are stored into the ADC\_JDRx registers of each ADC interface.
- A JEOC interrupt is generated (if enabled on one of the three ADC interfaces) when the ADC1/ADC2/ADC3's injected channels have all been converted.

Figure 37. Injected simultaneous mode on 4 channels: triple ADC mode



### 10.9.2 Regular simultaneous mode

This mode is performed on a regular group of channels. The external trigger source comes from the regular group multiplexer of ADC1 (selected by the EXTSEL[3:0] bits in the ADC1\_CR2 register). A simultaneous trigger is provided to ADC2 and ADC3.

Note:

Do not convert the same channel on the two/three ADCs (no overlapping sampling times for the two/three ADCs when converting the same channel).

In regular simultaneous mode, one must convert sequences with the same length or ensure that the interval between triggers is longer than the long conversion time of the 2 sequences (Dual ADC mode) /3 sequences (Triple ADC mode). Otherwise, the ADC with the shortest sequence may restart while the ADC with the longest sequence is completing the previous conversions.

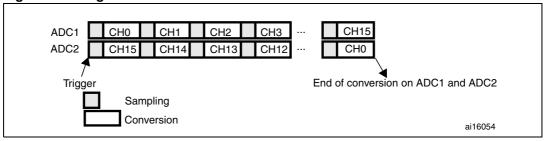
Injected conversions must be disabled.

#### **Dual ADC mode**

At the end of conversion event on ADC1 or ADC2:

- A 32-bit DMA transfer request is generated (if DMA[1:0] bits in the ADC\_CCR register are equal to 0b10). This request transfers the ADC2 converted data stored in the upper half-word of the ADC\_CDR 32-bit register to the SRAM and then the ADC1 converted data stored in the lower half-word of ADC\_CCR to the SRAM.
- An EOC interrupt is generated (if enabled on one of the two ADC interfaces) when the ADC1/ADC2's regular channels have all been converted.

Figure 38. Regular simultaneous mode on 16 channels: dual ADC mode

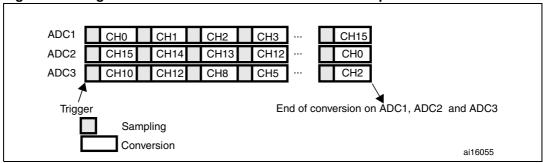


### **Triple ADC mode**

At the end of conversion event on ADC1, ADC2 or ADC3:

- Three 32-bit DMA transfer requests are generated (if DMA[1:0] bits in the ADC\_CCR register are equal to 0b01). Three transfers then take place from the ADC\_CDR 32-bit register to SRAM: first the ADC1 converted data, then the ADC2 converted data and finally the ADC3 converted data. The process is repeated for each new three conversions.
- An EOC interrupt is generated (if enabled on one of the three ADC interfaces) when the ADC1/ADC2/ADC3's regular channels are have all been converted.

Figure 39. Regular simultaneous mode on 16 channels: triple ADC mode



#### 10.9.3 Interleaved mode

This mode can be started only on a regular group (usually one channel). The external trigger source comes from the regular channel multiplexer of ADC1.

#### **Dual ADC mode**

After an external trigger occurs:

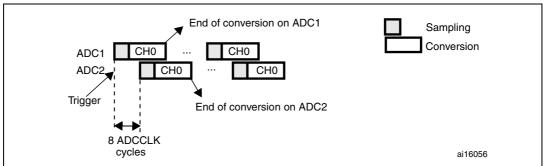
- ADC1 starts immediately
- ADC2 starts after a delay of several-ADC clock cycles

The minimum delay which separates 2 conversions in interleaved mode is configured in the DELAY bits in the ADC\_CCR register. However, an ADC cannot start a conversion if the complementary ADC is still sampling its input (only one ADC can sample the input signal at a given time). In this case, the delay becomes the sampling time + 2 ADC clock cycles. For instance, if DELAY = 5 clock cycles and the sampling takes 15 clock cycles on both ADCs, then 17 clock cycles will separate conversions on ADC1 and ADC2).

If the CONT bit is set on both ADC1 and ADC2, the selected regular channels of both ADCs are continuously converted.

After an EOC interrupt is generated by ADC2 (if enabled through the EOCIE bit) a 32-bit DMA transfer request is generated (if the DMA[1:0] bits in ADC\_CCR are equal to 0b10). This request first transfers the ADC2 converted data stored in the upper half-word of the ADC\_CDR 32-bit register into SRAM, then the ADC1 converted data stored in the register's lower half-word into SRAM.

Figure 40. Interleaved mode on 1 channel in continuous conversion mode: dual ADC mode



#### **Triple ADC mode**

After an external trigger occurs:

- ADC1 starts immediately and
- ADC2 starts after a delay of several ADC clock cycles
- ADC3 starts after a delay of several ADC clock cycles referred to the ADC2 conversion

The minimum delay which separates 2 conversions in interleaved mode is configured in the DELAY bits in the ADC\_CCR register. However, an ADC cannot start a conversion if the complementary ADC is still sampling its input (only one ADC can sample the input signal at a given time). In this case, the delay becomes the sampling time + 2 ADC clock cycles. For instance, if DELAY = 5 clock cycles and the sampling takes 15 clock cycles on the three ADCs, then 17 clock cycles will separate the conversions on ADC1, ADC2 and ADC3).

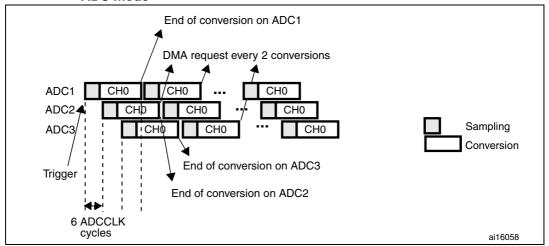
229/1340 Doc ID 018909 Rev 2

If the CONT bit is set on ADC1, ADC2 and ADC3, the selected regular channels of all ADCs are continuously converted.

In this mode a DMA request is generated each time 2 data items are available, (if the DMA[1:0] bits in the ADC\_CCR register are equal to 0b10). The request first transfers the first converted data stored in the lower half-word of the ADC\_CDR 32-bit register to SRAM, then it transfers the second converted data stored in ADC\_CDR's upper half-word to SRAM. The sequence is the following:

- 1st request: ADC\_CDR[31:0] = ADC2\_DR[15:0] | ADC1\_DR[15:0]
- 2nd request: ADC\_CDR[31:0] = ADC1\_DR[15:0] | ADC3\_DR[15:0]
- 3rd request: ADC\_CDR[31:0] = ADC3\_DR[15:0] | ADC2\_DR[15:0]
- 4th request: ADC\_CDR[31:0] = ADC2\_DR[15:0] | ADC1\_DR[15:0], ...

Figure 41. Interleaved mode on 1 channel in continuous conversion mode: triple ADC mode



### 10.9.4 Alternate trigger mode

This mode can be started only on an injected group. The source of external trigger comes from the injected group multiplexer of ADC1.

Note:

Regular conversions can be enabled on one or all ADCs. In this case the regular conversions are independent of each other. A regular conversion is interrupted when the ADC has to perform an injected conversion. It is resumed when the injected conversion is finished.

The time interval between 2 trigger events must be greater than or equal to 1 ADC clock period. The minimum time interval between 2 trigger events that start conversions on the same ADC is the same as in the single ADC mode.

### **Dual ADC mode**

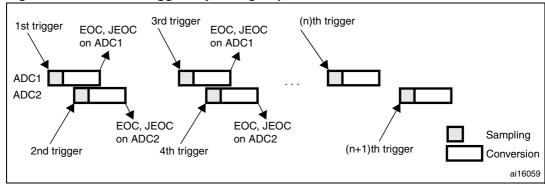
- When the 1st trigger occurs, all injected ADC1 channels in the group are converted
- When the 2nd trigger occurs, all injected ADC2 channels in the group are converted
- and so on

A JEOC interrupt, if enabled, is generated after all injected ADC1 channels in the group have been converted.

A JEOC interrupt, if enabled, is generated after all injected ADC2 channels in the group have been converted.

If another external trigger occurs after all injected channels in the group have been converted then the alternate trigger process restarts by converting the injected ADC1 channels in the group.

Figure 42. Alternate trigger: injected group of each ADC



If the injected discontinuous mode is enabled for both ADC1 and ADC2:

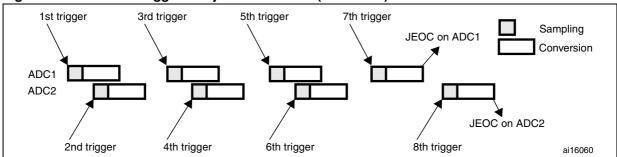
- When the 1st trigger occurs, the first injected ADC1 channel is converted.
- When the 2nd trigger occurs, the first injected ADC2 channel are converted
- and so on

A JEOC interrupt, if enabled, is generated after all injected ADC1 channels in the group have been converted.

A JEOC interrupt, if enabled, is generated after all injected ADC2 channels in the group have been converted.

If another external trigger occurs after all injected channels in the group have been converted then the alternate trigger process restarts.

Figure 43. Alternate trigger: 4 injected channels (each ADC) in discontinuous mode



#### **Triple ADC mode**

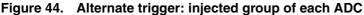
- When the 1st trigger occurs, all injected ADC1 channels in the group are converted.
- When the 2nd trigger occurs, all injected ADC2 channels in the group are converted.
- When the 3rd trigger occurs, all injected ADC3 channels in the group are converted.
- and so on

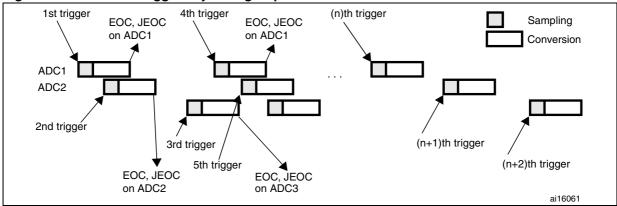
A JEOC interrupt, if enabled, is generated after all injected ADC1 channels in the group have been converted.

A JEOC interrupt, if enabled, is generated after all injected ADC2 channels in the group have been converted.

A JEOC interrupt, if enabled, is generated after all injected ADC3 channels in the group have been converted.

If another external trigger occurs after all injected channels in the group have been converted then the alternate trigger process restarts by converting the injected ADC1 channels in the group.





### 10.9.5 Combined regular/injected simultaneous mode

It is possible to interrupt the simultaneous conversion of a regular group to start the simultaneous conversion of an injected group.

Note:

In combined regular/injected simultaneous mode, one must convert sequences with the same length or ensure that the interval between triggers is longer than the long conversion time of the 2 sequences (Dual ADC mode) /3 sequences (Triple ADC mode). Otherwise, the ADC with the shortest sequence may restart while the ADC with the longest sequence is completing the previous conversions.

### 10.9.6 Combined regular simultaneous + alternate trigger mode

It is possible to interrupt the simultaneous conversion of a regular group to start the alternate trigger conversion of an injected group. *Figure 45* shows the behavior of an alternate trigger interrupting a simultaneous regular conversion.

The injected alternate conversion is immediately started after the injected event. If regular conversion is already running, in order to ensure synchronization after the injected conversion, the regular conversion of all (master/slave) ADCs is stopped and resumed synchronously at the end of the injected conversion.

Note:

In combined regular simultaneous + alternate trigger mode, one must convert sequences with the same length or ensure that the interval between triggers is longer than the long conversion time of the 2 sequences (Dual ADC mode) /3 sequences (Triple ADC mode). Otherwise, the ADC with the shortest sequence may restart while the ADC with the longest sequence is completing the previous conversions.

ADC2 ini

synchro not lost

ai16062

1st trigger ADC1 reg CH<sub>0</sub> CH<sub>1</sub> CH<sub>2</sub> CH2 CH3 СНЗ CH4 CH<sub>0</sub> ADC1 inj СНЗ CH5 CH6 CH6 CH7 CH7 CH8 ADC2 reg

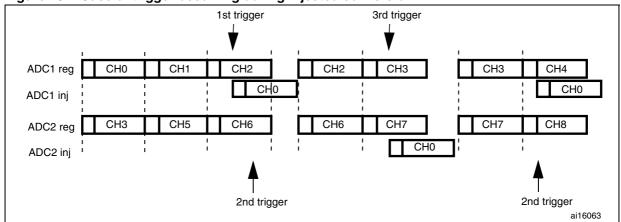
Alternate + regular simultaneous

If a trigger occurs during an injected conversion that has interrupted a regular conversion, it is ignored. Figure 46 shows the behavior in this case (2nd trigger is ignored).

CH0

2nd trigger

Figure 46. Case of trigger occurring during injected conversion



#### 10.10 **Temperature sensor**

The temperature sensor can be used to measure the ambient temperature  $(T_{\mbox{\scriptsize A}})$  of the device.

The temperature sensor is internally connected to the ADC1 IN16 input channel which is used to convert the sensor's output voltage to a digital value.

Figure 47 shows the block diagram of the temperature sensor.

When not in use, the sensor can be put in power down mode.

Note:

The TSVREFE bit must be set to enable the conversion of both internal channels: ADC1\_IN16 (temperature sensor) and ADC1\_IN17 ( $V_{REFINT}$ ).

#### Main features

- Supported temperature range: -40 to 125 °C
- Precision: ±1.5 °C

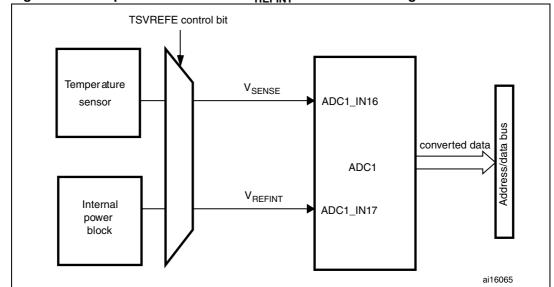


Figure 47. Temperature sensor and V<sub>REFINT</sub> channel block diagram

### Reading the temperature

To use the sensor:

- 1. Select the ADC1\_IN16 input channel
- 2. Select a sampling time greater than the minimum sampling time specified in the datasheet.
- 3. Set the TSVREFE bit in the ADC\_CCR register to wake up the temperature sensor from power down mode
- 4. Start the ADC conversion by setting the SWSTART bit (or by external trigger)
- 5. Read the resulting V<sub>SENSE</sub> data in the ADC data register
- 6. Calculate the temperature using the following formula:

Temperature (in °C) =  $\{(V_{SENSE} - V_{25}) / Avg\_Slope\} + 25$ Where:

- V<sub>25</sub> = V<sub>SENSE</sub> value for 25° C
- Avg\_Slope = average slope of the temperature vs. V<sub>SENSE</sub> curve (given in mV/°C or μV/°C)

Refer to the datasheet's electrical characteristics section for the actual values of  $V_{25}$  and  $Avg\_Slope$ .

Note:

The sensor has a startup time after waking from power down mode before it can output  $V_{SENSE}$  at the correct level. The ADC also has a startup time after power-on, so to minimize the delay, the ADON and TSVREFE bits should be set at the same time.

# 10.11 Battery charge monitoring

The VBATE bit in the ADC\_CCR register is used to switch to the battery voltage. As the  $V_{BAT}$  voltage could be higher than  $V_{DDA}$ , to ensure the correct operation of the ADC, the  $V_{BAT}$  pin is internally connected to a bridge divider.

When the VBATE is set, the bridge is automatically enabled to connect:

- VBAT/2 to the ADC1\_IN18 input channel, on STM32F40xx and STM32F41xx devices
- VBAT/4 to the ADC1\_IN18 input channel, on STM32F42xx and STM32F43xx devices

# 10.12 ADC interrupts

An interrupt can be produced on the end of conversion for regular and injected groups, when the analog watchdog status bit is set and when the overrun status bit is set. Separate interrupt enable bits are available for flexibility.

Two other flags are present in the ADC\_SR register, but there is no interrupt associated with them:

- JSTRT (Start of conversion for channels of an injected group)
- STRT (Start of conversion for channels of a regular group)

Table 39. ADC interrupts

Interrupt event	Event flag	Enable control bit
End of conversion of a regular group	EOC	EOCIE
End of conversion of an injected group	JEOC	JEOCIE
Analog watchdog status bit is set	AWD	AWDIE
Overrun	OVR	OVRIE

# 10.13 ADC registers

Refer to Section 1.1 on page 46 for a list of abbreviations used in register descriptions.

The peripheral registers must be written at word level (32 bits). Read accesses can be done by bytes (8 bits), half-words (16 bits) or words (32 bits).

### 10.13.1 ADC status register (ADC\_SR)

Address offset: 0x00

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Res	served							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
				Dag	erved					OVR	STRT	JSTRT	JEOC	EOC	AWD
				nes	ervea					rc_w0	rc_w0	rc_w0	rc_w0	rc_w0	rc_w0

Bits 31:6 Reserved, must be kept at reset value.

#### Bit 5 OVR: Overrun

This bit is set by hardware when data are lost (either in single mode or in dual/triple mode). It is cleared by software. Overrun detection is enabled only when DMA = 1 or EOCS = 1.

- 0: No overrun occurred
- 1: Overrun has occurred

#### Bit 4 STRT: Regular channel start flag

This bit is set by hardware when regular channel conversion starts. It is cleared by software.

- 0: No regular channel conversion started
- 1: Regular channel conversion has started

#### Bit 3 JSTRT: Injected channel start flag

This bit is set by hardware when injected group conversion starts. It is cleared by software.

- 0: No injected group conversion started
- 1: Injected group conversion has started

#### Bit 2 **JEOC:** Injected channel end of conversion

This bit is set by hardware at the end of the conversion of all injected channels in the group. It is cleared by software.

- 0: Conversion is not complete
- 1: Conversion complete

#### Bit 1 EOC: Regular channel end of conversion

This bit is set by hardware at the end of the conversion of a regular group of channels. It is cleared by software or by reading the ADC\_DR register.

- 0: Conversion not complete (EOCS=0), or sequence of conversions not complete (EOCS=1)
- 1: Conversion complete (EOCS=0), or sequence of conversions complete (EOCS=1)

#### Bit 0 AWD: Analog watchdog flag

This bit is set by hardware when the converted voltage crosses the values programmed in the ADC\_LTR and ADC\_HTR registers. It is cleared by software.

- 0: No analog watchdog event occurred
- 1: Analog watchdog event occurred

### 10.13.2 ADC control register 1 (ADC\_CR1)

Address offset: 0x04

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
		Reserve	od		OVRIE	RE	S	AWDEN	JAWDEN			Poor	mund		
		neserv	eu		rw	rw	rw	rw	rw	Reserved					
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
DIS	SCNUM[	2:0]	JDISCE N	DISC EN	JAUTO	AWDSG L	SCAN	JEOCIE	AWDIE	EOCIE	AWDCH[4:0]				
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:27 Reserved, must be kept at reset value.

#### Bit 26 OVRIE: Overrun interrupt enable

This bit is set and cleared by software to enable/disable the Overrun interrupt.

0: Overrun interrupt disabled

1: Overrun interrupt enabled. An interrupt is generated when the OVR bit is set.

#### Bits 25:24 RES[1:0]: Resolution

These bits are written by software to select the resolution of the conversion.

00: 12-bit (15 ADCCLK cycles) 01: 10-bit (13 ADCCLK cycles) 10: 8-bit (11 ADCCLK cycles) 11: 6-bit (9 ADCCLK cycles)

#### Bit 23 AWDEN: Analog watchdog enable on regular channels

This bit is set and cleared by software.

0: Analog watchdog disabled on regular channels

1: Analog watchdog enabled on regular channels

#### Bit 22 JAWDEN: Analog watchdog enable on injected channels

This bit is set and cleared by software.

0: Analog watchdog disabled on injected channels

1: Analog watchdog enabled on injected channels

#### Bits 21:16 Reserved, must be kept at reset value.

#### Bits 15:13 DISCNUM[2:0]: Discontinuous mode channel count

These bits are written by software to define the number of regular channels to be converted in discontinuous mode, after receiving an external trigger.

000: 1 channel 001: 2 channels ... 111: 8 channels

#### Bit 12 JDISCEN: Discontinuous mode on injected channels

This bit is set and cleared by software to enable/disable discontinuous mode on the injected channels of a group.

0: Discontinuous mode on injected channels disabled

1: Discontinuous mode on injected channels enabled

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#### Bit 11 DISCEN: Discontinuous mode on regular channels

This bit is set and cleared by software to enable/disable Discontinuous mode on regular channels.

- 0: Discontinuous mode on regular channels disabled
- 1: Discontinuous mode on regular channels enabled

#### Bit 10 JAUTO: Automatic injected group conversion

This bit is set and cleared by software to enable/disable automatic injected group conversion after regular group conversion.

- 0: Automatic injected group conversion disabled
- 1: Automatic injected group conversion enabled

#### Bit 9 AWDSGL: Enable the watchdog on a single channel in scan mode

This bit is set and cleared by software to enable/disable the analog watchdog on the channel identified by the AWDCH[4:0] bits.

- 0: Analog watchdog enabled on all channels
- 1: Analog watchdog enabled on a single channel

#### Bit 8 SCAN: Scan mode

This bit is set and cleared by software to enable/disable the Scan mode. In Scan mode, the inputs selected through the ADC\_SQRx or ADC\_JSQRx registers are converted.

- 0: Scan mode disabled
- 1: Scan mode enabled

Note: An EOC interrupt is generated if the EOCIE bit is set:

- At the end of each regular group sequence if the EOCS bit is cleared to 0
- At the end of each regular channel conversion if the EOCS bit is set to 1

Note: A JEOC interrupt is generated only on the end of conversion of the last channel if the JEOCIE bit is set.

#### Bit 7 **JEOCIE:** Interrupt enable for injected channels

This bit is set and cleared by software to enable/disable the end of conversion interrupt for injected channels.

- 0: JEOC interrupt disabled
- 1: JEOC interrupt enabled. An interrupt is generated when the JEOC bit is set.

#### Bit 6 AWDIE: Analog watchdog interrupt enable

This bit is set and cleared by software to enable/disable the analog watchdog interrupt.

- 0: Analog watchdog interrupt disabled
- 1: Analog watchdog interrupt enabled

### Bit 5 EOCIE: Interrupt enable for EOC

This bit is set and cleared by software to enable/disable the end of conversion interrupt.

- 0: EOC interrupt disabled
- 1: EOC interrupt enabled. An interrupt is generated when the EOC bit is set.

#### Bits 4:0 AWDCH[4:0]: Analog watchdog channel select bits

These bits are set and cleared by software. They select the input channel to be guarded by the analog watchdog.

Note: 00000: ADC analog input Channel0

00001: ADC analog input Channel1

• • •

01111: ADC analog input Channel15 10000: ADC analog input Channel16 10001: ADC analog input Channel17 10010: ADC analog input Channel18

Other values reserved

### 10.13.3 ADC control register 2 (ADC\_CR2)

Address offset: 0x08

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
reserve d	SWST ART	EXT	EN		EXTS	EL[3:0]		reserve	JSWST ART	JEXT	ΓΕΝ		JEXTS	EL[3:0]	
u	rw	rw	rw	rw	rw	rw	rw	u	rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
reserved ALIGN EOCS DDS DMA							Decembed					CONT	ADON		
	rese	iveu		rw	rw	rw	rw	Reserved					rw	rw	

Bit 31 Reserved, must be kept at reset value.

#### Bit 30 SWSTART: Start conversion of regular channels

This bit is set by software to start conversion and cleared by hardware as soon as the conversion starts.

0: Reset state

1: Starts conversion of regular channels

Note: This bit can be set only when ADON = 1 otherwise no conversion is launched.

#### Bits 29:28 EXTEN: External trigger enable for regular channels

These bits are set and cleared by software to select the external trigger polarity and enable the trigger of a regular group.

00: Trigger detection disabled

01: Trigger detection on the rising edge

10: Trigger detection on the falling edge

11: Trigger detection on both the rising and falling edges

### Bits 27:24 EXTSEL[3:0]: External event select for regular group

These bits select the external event used to trigger the start of conversion of a regular group:

0000: Timer 1 CC1 event

0001: Timer 1 CC2 event

0010: Timer 1 CC3 event

0011: Timer 2 CC2 event

0100: Timer 2 CC3 event

0101: Timer 2 CC4 event

0110: Timer 2 TRGO event

0111: Timer 3 CC1 event

1000: Timer 3 TRGO event 1001: Timer 4 CC4 event

1010: Timer 5 CC1 event

1011: Timer 5 CC2 event

1100: Timer 5 CC3 event

1101: Timer 8 CC1 event 1110: Timer 8 TRGO event

1111: EXTI line11

Bit 23 Reserved, must be kept at reset value.

#### Bit 22 JSWSTART: Start conversion of injected channels

This bit is set by software and cleared by hardware as soon as the conversion starts.

0: Reset state

1: Starts conversion of injected channels

Note: This bit can be set only when ADON = 1 otherwise no conversion is launched.

#### Bits 21:20 **JEXTEN:** External trigger enable for injected channels

These bits are set and cleared by software to select the external trigger polarity and enable the trigger of an injected group.

00: Trigger detection disabled

01: Trigger detection on the rising edge

10: Trigger detection on the falling edge

11: Trigger detection on both the rising and falling edges

#### Bits 19:16 JEXTSEL[3:0]: External event select for injected group

These bits select the external event used to trigger the start of conversion of an injected group.

0000: Timer 1 CC4 event

0001: Timer 1 TRGO event

0010: Timer 2 CC1 event

0011: Timer 2 TRGO event

0100: Timer 3 CC2 event

0101: Timer 3 CC4 event

0110: Timer 4 CC1 event

0111: Timer 4 CC2 event

1000: Timer 4 CC3 event 1001: Timer 4 TRGO event

1010: Timer 5 CC4 event

1010. Timer 5 004 event

1011: Timer 5 TRGO event 1100: Timer 8 CC2 event

1101: Timer 8 CC3 event

1110: Timer 8 CC3 event

1111: EXTI line15

#### Bits 15:12 Reserved, must be kept at reset value.

#### Bit 11 ALIGN: Data alignment

This bit is set and cleared by software. Refer to Figure 32 and Figure 33.

0: Right alignment

1: Left alignment

#### Bit 10 EOCS: End of conversion selection

This bit is set and cleared by software.

- 0: The EOC bit is set at the end of each sequence of regular conversions. Overrun detection is enabled only if DMA=1.
- 1: The EOC bit is set at the end of each regular conversion. Overrun detection is enabled.

#### Bit 9 DDS: DMA disable selection (for single ADC mode)

This bit is set and cleared by software.

- 0: No new DMA request is issued after the last transfer (as configured in the DMA controller)
- 1: DMA requests are issued as long as data are converted and DMA=1

#### Bit 8 DMA: Direct memory access mode (for single ADC mode)

This bit is set and cleared by software. Refer to the DMA controller chapter for more details.

- 0: DMA mode disabled
- 1: DMA mode enabled

#### Bit 22 JSWSTART: Start conversion of injected channels

This bit is set by software and cleared by hardware as soon as the conversion starts.

0: Reset state

1: Starts conversion of injected channels

Note: This bit can be set only when ADON = 1 otherwise no conversion is launched.

#### Bits 21:20 **JEXTEN:** External trigger enable for injected channels

These bits are set and cleared by software to select the external trigger polarity and enable the trigger of an injected group.

00: Trigger detection disabled

01: Trigger detection on the rising edge

10: Trigger detection on the falling edge

11: Trigger detection on both the rising and falling edges

#### Bits 19:16 JEXTSEL[3:0]: External event select for injected group

These bits select the external event used to trigger the start of conversion of an injected group.

0000: Timer 1 CC4 event

0001: Timer 1 TRGO event

0010: Timer 2 CC1 event

0011: Timer 2 TRGO event

0100: Timer 3 CC2 event

0101: Timer 3 CC4 event

0110: Timer 4 CC1 event

0111: Timer 4 CC2 event

1000: Timer 4 CC3 event

1001: Timer 4 TRGO event 1010: Timer 5 CC4 event

1011: Timer 5 TRGO event

1011: Timer 5 TRGO eveni

1100: Timer 8 CC2 event

1101: Timer 8 CC3 event

1110: Timer 8 CC4 event

1111: EXTI line15

#### Bits 15:12 Reserved, must be kept at reset value.

#### Bit 11 ALIGN: Data alignment

This bit is set and cleared by software. Refer to Figure 32 and Figure 33.

0: Right alignment

1: Left alignment

#### Bit 10 EOCS: End of conversion selection

This bit is set and cleared by software.

- 0: The EOC bit is set at the end of each sequence of regular conversions. Overrun detection is enabled only if DMA=1.
- 1: The EOC bit is set at the end of each regular conversion. Overrun detection is enabled.

#### Bit 9 DDS: DMA disable selection (for single ADC mode)

This bit is set and cleared by software.

- 0: No new DMA request is issued after the last transfer (as configured in the DMA controller)
- 1: DMA requests are issued as long as data are converted and DMA=1

#### Bit 8 DMA: Direct memory access mode (for single ADC mode)

This bit is set and cleared by software. Refer to the DMA controller chapter for more details.

0: DMA mode disabled

1: DMA mode enabled

Bits 7:2 Reserved, must be kept at reset value.

Bit 1 CONT: Continuous conversion

This bit is set and cleared by software. If it is set, conversion takes place continuously until it is cleared.

0: Single conversion mode

1: Continuous conversion mode

Bit 0 ADON: A/D Converter ON / OFF

This bit is set and cleared by software.

Note: 0: Disable ADC conversion and go to power down mode

1: Enable ADC

### 10.13.4 ADC sample time register 1 (ADC\_SMPR1)

Address offset: 0x0C

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	Ь	eserved			5	SMP18[2:0	0]	S	MP17[2:0	0]	S	MP16[2:0	0]	SMP1	5[2:1]
	neserveu				rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
SMP15_0	MP15_0 SMP14[2:0]		S	MP13[2:	0]	S	SMP12[2:0	0]	S	MP11[2:0	0]	S	MP10[2:0	0]	
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31: 27 Reserved, must be kept at reset value.

#### Bits 26:0 SMPx[2:0]: Channel x sampling time selection

These bits are written by software to select the sampling time individually for each channel. During sampling cycles, the channel selection bits must remain unchanged.

Note: 000: 3 cycles 001: 15 cycles 010: 28 cycles 011: 56 cycles 100: 84 cycles 101: 112 cycles 110: 144 cycles 111: 480 cycles

# 10.13.5 ADC sample time register 2 (ADC\_SMPR2)

Address offset: 0x10

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
Boos	ruad		SMP9[2:0	]		SMP8[2:0	)]	,	SMP7[2:0	]	:	SMP6[2:0	]	SMP	5[2:1]
nese	served rw rw rw		rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
SMP 5_0	;	SMP4[2:0	)]	SMP3[2:0]		)]	;	SMP2[2:0	]	;	SMP1[2:0	]		SMP0[2:0	]
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:30 Reserved, must be kept at reset value.

#### Bits 29:0 SMPx[2:0]: Channel x sampling time selection

These bits are written by software to select the sampling time individually for each channel. During sample cycles, the channel selection bits must remain unchanged.

Note: 000: 3 cycles 001: 15 cycles 010: 28 cycles 011: 56 cycles 100: 84 cycles 101: 112 cycles 110: 144 cycles 111: 480 cycles

# 10.13.6 ADC injected channel data offset register x (ADC\_JOFRx)(x=1..4)

Address offset: 0x14-0x20 Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	
					Reserved											
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
	Rese	ruod			JOFFSETx[11:0]											
	nese	erveu	•	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	

Bits 31:12 Reserved, must be kept at reset value.

#### Bits 11:0 JOFFSETx[11:0]: Data offset for injected channel x

These bits are written by software to define the offset to be subtracted from the raw converted data when converting injected channels. The conversion result can be read from in the ADC\_JDRx registers.

# 10.13.7 ADC watchdog higher threshold register (ADC\_HTR)

Address offset: 0x24

Reset value: 0x0000 0FFF

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
					Reserved										
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
HT[11:0]															
	nese	riveu	•	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:12 Reserved, must be kept at reset value.

Bits 11:0 HT[11:0]: Analog watchdog higher threshold

These bits are written by software to define the higher threshold for the analog watchdog.

# 10.13.8 ADC watchdog lower threshold register (ADC\_LTR)

Address offset: 0x28

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Res	served							•
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
LT[11:0]															
rw rw rw rw							rw	rw	rw	rw	rw	rw	rw		

Bits 31:12 Reserved, must be kept at reset value.

Bits 11:0 LT[11:0]: Analog watchdog lower threshold

These bits are written by software to define the lower threshold for the analog watchdog.

### 10.13.9 ADC regular sequence register 1 (ADC\_SQR1)

Address offset: 0x2C

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
			Pos	erved					L[3	3:0]			SQ1	6[4:1]	
			nese	erveu				rw	rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
SQ16_0		;	SQ15[4:0	]				SQ14[4:0]	]				SQ13[4:0	]	
rw	rw	rw	rw	rw	rw	rw	rw				rw	rw	rw	rw	rw

Bits 31:24 Reserved, must be kept at reset value.

Bits 23:20 L[3:0]: Regular channel sequence length

These bits are written by software to define the total number of conversions in the regular channel conversion sequence.

0000: 1 conversion 0001: 2 conversions

...

1111: 16 conversions

Bits 19:15 SQ16[4:0]: 16th conversion in regular sequence

These bits are written by software with the channel number (0..18) assigned as the 16th in the conversion sequence.

Bits 14:10 SQ15[4:0]: 15th conversion in regular sequence

Bits 9:5 SQ14[4:0]: 14th conversion in regular sequence

Bits 4:0 SQ13[4:0]: 13th conversion in regular sequence

### 10.13.10 ADC regular sequence register 2 (ADC\_SQR2)

Address offset: 0x30

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
Rese	SQ12[4:0]								SQ1	SQ10[4:1]					
nese	iveu	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
SQ10_0			SQ9[4:0]					SQ8[4:0]					SQ7[4:0]		
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:30 Reserved, must be kept at reset value.

Bits 29:26 SQ12[4:0]: 12th conversion in regular sequence

These bits are written by software with the channel number (0..18) assigned as the 12th in the sequence to be converted.

Bits 24:20 SQ11[4:0]: 11th conversion in regular sequence

Bits 19:15 SQ10[4:0]: 10th conversion in regular sequence

Bits 14:10 SQ9[4:0]: 9th conversion in regular sequence

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Bits 9:5 **SQ8[4:0]:** 8th conversion in regular sequence

Bits 4:0 SQ7[4:0]: 7th conversion in regular sequence

### 10.13.11 ADC regular sequence register 3 (ADC\_SQR3)

Address offset: 0x34

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
Rese	nuad	SQ6[4:0]					SQ5[4:0] S								
nese	erveu	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
SQ4_0			SQ3[4:0]					SQ2[4:0]					SQ1[4:0]		
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:30 Reserved, must be kept at reset value.

Bits 29:25 SQ6[4:0]: 6th conversion in regular sequence

These bits are written by software with the channel number (0..18) assigned as the 6th in the sequence to be converted.

Bits 24:20 SQ5[4:0]: 5th conversion in regular sequence

Bits 19:15 SQ4[4:0]: 4th conversion in regular sequence

Bits 14:10 SQ3[4:0]: 3rd conversion in regular sequence

Bits 9:5 **SQ2[4:0]:** 2nd conversion in regular sequence

Bits 4:0 SQ1[4:0]: 1st conversion in regular sequence

### 10.13.12 ADC injected sequence register (ADC\_JSQR)

Address offset: 0x38

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
				Pos	erved					JL[	1:0]		JSQ	4[4:1]	
				nes	erveu					rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
JSQ4[0]			JSQ3[4:0	]				JSQ2[4:0]	]				JSQ1[4:0	]	
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:22 Reserved, must be kept at reset value.

Bits 21:20 JL[1:0]: Injected sequence length

These bits are written by software to define the total number of conversions in the injected channel conversion sequence.

00: 1 conversion 01: 2 conversions 10: 3 conversions 11: 4 conversions Bits 19:15 JSQ4[4:0]: 4th conversion in injected sequence (when JL[1:0]=3, see note below)

These bits are written by software with the channel number (0..18) assigned as the 4th in the sequence to be converted.

Bits 14:10 JSQ3[4:0]: 3rd conversion in injected sequence (when JL[1:0]=3, see note below)

Bits 9:5 JSQ2[4:0]: 2nd conversion in injected sequence (when JL[1:0]=3, see note below)

Bits 4:0 JSQ1[4:0]: 1st conversion in injected sequence (when JL[1:0]=3, see note below)

Note:

When JL[1:0]=3 (4 injected conversions in the sequencer), the ADC converts the channels in the following order: JSQ1[4:0], JSQ2[4:0], JSQ3[4:0], and JSQ4[4:0].

When JL=2 (3 injected conversions in the sequencer), the ADC converts the channels in the following order: JSQ2[4:0], JSQ3[4:0], and JSQ4[4:0].

When JL=1 (2 injected conversions in the sequencer), the ADC converts the channels in starting from JSQ3[4:0], and then JSQ4[4:0].

When JL=0 (1 injected conversion in the sequencer), the ADC converts only JSQ4[4:0] channel.

### 10.13.13 ADC injected data register x (ADC\_JDRx) (x= 1..4)

Address offset: 0x3C - 0x48 Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Res	served							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		_	_	_	_		JDAT	A[15:0]	_						
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:16 Reserved, must be kept at reset value.

Bits 15:0 JDATA[15:0]: Injected data

These bits are read-only. They contain the conversion result from injected channel x. The data are left -or right-aligned as shown in *Figure 32* and *Figure 33*.

### 10.13.14 ADC regular data register (ADC\_DR)

Address offset: 0x4C

Reset value: 0x0000 0000

	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
								Res	served							
	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
								DAT	A[15:0]							
Γ	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:16 Reserved, must be kept at reset value.

Bits 15:0 DATA[15:0]: Regular data

These bits are read-only. They contain the conversion result from the regular channels. The data are left- or right-aligned as shown in *Figure 32* and *Figure 33*.

### 10.13.15 ADC Common status register (ADC\_CSR)

Address offset: 0x00 (this offset address is relative to ADC1 base address + 0x300)

Reset value: 0x0000 0000

This register provides an image of the status bits of the different ADCs. Nevertheless it is read-only and does not allow to clear the different status bits. Instead each status bit must be cleared by writing it to 0 in the corresponding ADC\_SR register.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
										OVR3	STRT3	JSTRT3	JEOC 3	EOC3	AWD3
				Res	served							AD	C3		
										r	r	r	r	r	r
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		OVR2	STRT2	JSTRT 2	JEOC2	EOC2	AWD2			OVR1	STRT1	JSTRT1	JEOC 1	EOC1	AWD1
Rese	rved			AD	C2			Res	erved			AD	C1		
		r	r	r	r	r	r			r	r	r	r	r	r

Bits 31:22 Reserved, must be kept at reset value.

Bit 21 **OVR3:** Overrun flag of ADC3

This bit is a copy of the OVR bit in the ADC3\_SR register.

Bit 20 STRT3: Regular channel Start flag of ADC3

This bit is a copy of the STRT bit in the ADC3\_SR register.

Bit 19 JSTRT3: Injected channel Start flag of ADC3

This bit is a copy of the JSTRT bit in the ADC3\_SR register.

Bit 18 **JEOC3:** Injected channel end of conversion of ADC3

This bit is a copy of the JEOC bit in the ADC3\_SR register.

Bit 17 EOC3: End of conversion of ADC3

This bit is a copy of the EOC bit in the ADC3\_SR register.

Bit 16 AWD3: Analog watchdog flag of ADC3

This bit is a copy of the AWD bit in the ADC3\_SR register.

Bits 15:14 Reserved, must be kept at reset value.

Bit 13 OVR2: Overrun flag of ADC2

This bit is a copy of the OVR bit in the ADC2\_SR register.

Bit 12 STRT2: Regular channel Start flag of ADC2

This bit is a copy of the STRT bit in the ADC2\_SR register.

Bit 11 JSTRT2: Injected channel Start flag of ADC2

This bit is a copy of the JSTRT bit in the ADC2\_SR register.

Bit 10 **JEOC2**: Injected channel end of conversion of ADC2

This bit is a copy of the JEOC bit in the ADC2\_SR register.

Bit 9 EOC2: End of conversion of ADC2

This bit is a copy of the EOC bit in the ADC2\_SR register.

Bit 8 AWD2: Analog watchdog flag of ADC2

This bit is a copy of the AWD bit in the ADC2\_SR register.

Bits 7:6 Reserved, must be kept at reset value.

Bit 5 OVR1: Overrun flag of ADC1

This bit is a copy of the OVR bit in the ADC1\_SR register.

Bit 4 STRT1: Regular channel Start flag of ADC1

This bit is a copy of the STRT bit in the ADC1\_SR register.

Bit 3 JSTRT1: Injected channel Start flag of ADC1

This bit is a copy of the JSTRT bit in the ADC1\_SR register.

Bit 2 **JEOC1:** Injected channel end of conversion of ADC1

This bit is a copy of the JEOC bit in the ADC1\_SR register.

Bit 1 EOC1: End of conversion of ADC1

This bit is a copy of the EOC bit in the ADC1\_SR register.

Bit 0 AWD1: Analog watchdog flag of ADC1

This bit is a copy of the AWD bit in the ADC1\_SR register.

### 10.13.16 ADC common control register (ADC\_CCR)

Address offset: 0x04 (this offset address is relative to ADC1 base address + 0x300)

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
			Poo	erved				TSVREFE	VBATE		Rese	ruod		ADO	CPRE
			nes	erveu				rw	rw		nese	erveu		rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
DMA	\[1:0]	DDS	Res.		DELA	Y[3:0]		ь	eserved	MULTI[4:0]					
rw	rw	rw	nes.	rw	rw	rw	rw	]	esei veu	rw rw rw		rw			

Bits 31:24 Reserved, must be kept at reset value.

Bit 23  $\,$  TSVREFE: Temperature sensor and  $\rm V_{REFINT}$  enable

This bit is set and cleared by software to enable/disable the temperature sensor and the  $V_{\mathsf{RFFINT}}$  channel.

0: Temperature sensor and V<sub>REFINT</sub> channel disabled

1: Temperature sensor and  $V_{\mbox{\scriptsize REFINT}}$  channel enabled

Bit 22 VBATE: V<sub>BAT</sub> enable

This bit is set and cleared by software to enable/disable the V<sub>BAT</sub> channel.

0: V<sub>BAT</sub> channel disabled

1: V<sub>BAT</sub> channel enabled

Bits 21:18 Reserved, must be kept at reset value.

Bits 17:16 ADCPRE: ADC prescaler

Set and cleared by software to select the frequency of the clock to the ADC. The clock is common for all the ADCs.

Note: 00: PCLK2 divided by 2

01: PCLK2 divided by 4 10: PCLK2 divided by 6 11: PCLK2 divided by 8 Bits 15:14 DMA: Direct memory access mode for multi ADC mode

This bit-field is set and cleared by software. Refer to the DMA controller section for more details.

00: DMA mode disabled

01: DMA mode 1 enabled (2 / 3 half-words one by one - 1 then 2 then 3)

10: DMA mode 2 enabled (2 / 3 half-words by pairs - 2&1 then 1&3 then 3&2)

11: DMA mode 3 enabled (2 / 3 bytes by pairs - 2&1 then 1&3 then 3&2)

Bit 13 DDS: DMA disable selection (for multi-ADC mode)

This bit is set and cleared by software.

0: No new DMA request is issued after the last transfer (as configured in the DMA controller). DMA bits are not cleared by hardware, however they must have been cleared and set to the wanted mode by software before new DMA requests can be generated.

1: DMA requests are issued as long as data are converted and DMA = 01, 10 or 11.

Bit 12 Reserved, must be kept at reset value.

#### Bit 11:8 **DELAY:** Delay between 2 sampling phases

Set and cleared by software. These bits are used in dual or triple interleaved modes.

0000: 5 \* T<sub>ADCCLK</sub> 0001: 6 \* T<sub>ADCCLK</sub> 0010: 7 \* T<sub>ADCCLK</sub> ... 1111: 20 \* T<sub>ADCCLK</sub>

Bits 7:5 Reserved, must be kept at reset value.

#### Bits 4:0 MULTI[4:0]: Multi ADC mode selection

These bits are written by software to select the operating mode.

- All the ADCs independent:

00000: Independent mode

- 00001 to 01001: Dual mode, ADC1 and ADC2 working together, ADC3 is independent

00001: Combined regular simultaneous + injected simultaneous mode

00010: Combined regular simultaneous + alternate trigger mode

00011: Reserved

00101: Injected simultaneous mode only

00110: Regular simultaneous mode only

00111: interleaved mode only

01001: Alternate trigger mode only

- 10001 to 11001: Triple mode: ADC1, 2 and 3 working together

10001: Combined regular simultaneous + injected simultaneous mode

10010: Combined regular simultaneous + alternate trigger mode

10011: Reserved

10101: Injected simultaneous mode only

10110: Regular simultaneous mode only

10111: interleaved mode only

11001: Alternate trigger mode only

All other combinations are reserved and must not be programmed

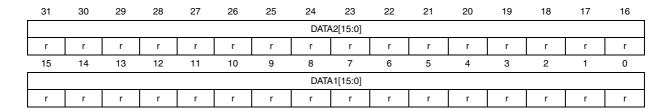
Note: In multi mode, a change of channel configuration generates an abort that can cause a loss of synchronization. It is recommended to disable the multi ADC mode before any configuration change.

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# 10.13.17 ADC common regular data register for dual and triple modes (ADC\_CDR)

Address offset: 0x08 (this offset address is relative to ADC1 base address + 0x300)

Reset value: 0x0000 0000



Bits 31:16 DATA2[15:0]: 2nd data item of a pair of regular conversions

- In dual mode, these bits contain the regular data of ADC2. Refer to *Dual ADC mode*.
- In triple mode, these bits contain alternatively the regular data of ADC2, ADC1 and ADC3.
   Refer to *Triple ADC mode*.

Bits 15:0 DATA1[15:0]: 1st data item of a pair of regular conversions

- In dual mode, these bits contain the regular data of ADC1. Refer to *Dual ADC mode*
- In triple mode, these bits contain alternatively the regular data of ADC1, ADC3 and ADC2.
   Refer to *Triple ADC mode*.

### 10.13.18 ADC register map

The following table summarizes the ADC registers.

Table 40. ADC global register map

Offset	Register
0x000 - 0x04C	ADC1
0x050 - 0x0FC	Reserved
0x100 - 0x14C	ADC2
0x118 - 0x1FC	Reserved
0x200 - 0x24C	ADC3
0x250 - 0x2FC	Reserved
0x300 - 0x308	Common registers

Table 41. ADC register map and reset values for each ADC

lable	41. ADC re		-	_	-			_			_					<del></del>														_
Offset	Register	30	20	28	27	26	25	23	22	21	20	19	18	17	16	15	13	12	11	10	6	8	7	9	2	4	ဗ	7	-	0
0x00	ADC_SR										R	Reser	ved	l											OVR	STRT	JSTRT	- 1	EOC	7,,,
	Reset value																								0	0	0	0	0 (	0
0x04	ADC_CR1	Re	ese	rved	į	OVAIR	RES[1:0]	AWDEN	JAWDEN		В	Reser	ved	l		D NUI	ISC M [2:0]	JDISCEN	DISCEN	JAUTO	AWD SGL	SCAN	JEOCIE	AWDIE	EOCIE	A	WD	CH[	4:0]	
	Reset value					0 (	0 (	0	0							0	0 0	0	0	0	0	0	0	0	0	0	0	0	0 (	0
0x08	ADC_CR2	Re se se d SWSTART		EXTEN[1:0]	EXT				JSWSTART	JEXTEN[1:0]			[3:0	_		R	eserve	d			SQQ	DMA		F	Rese	ervec	i		CONT	200
	Reset value	0	0	0	0	0	0 (	)	0	0	0	0			0				0	0		0							0 (	0
0x0C	ADC_SMPR1																SMPx.													
	Reset value	0 0	0	0	0	0	0 (	0	0	0	0					0		0	0	0	0	0	0	0	0	0	0	0	0 (	0
0x10	ADC_SMPR2	0.10	T ~	Ι.		<u> </u>	<u> </u>	\	Ι.	I o I			÷				SMPx.		_	_		_	_	I 0		_	^	$\overline{}$	0 1	_
	Reset value	0 0	0	0	0	0 (	0 0	0	0	0	0	0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0 11:0	0	0	0 (	0
0x14	ADC_JOFR1							ı	Res	erve	d								_	Δ.	٥							0	0 (	
	Reset value ADC_JOFR2																		0	0	0	0	0	0		0 11:0	0	0	0 (	0
0x18	Reset value							ı	Res	erve	d								0	0	0	0	0	0	0	0	0	0	0 (	0
	ADC_JOFR3																		Ŭ	U	U					11:0			<u> </u>	_
0x1C	Reset value							ı	Res	erve	d								0	0	0	0	0	0	0	0	0	0	0 (	0
	ADC_JOFR4																		Ť	Ŭ	Ů					11:0		Ů	,	-
0x20	Reset value							ı	Res	erve	d								0	0	0	0	0	0	0	0	0	0	0 (	0
	ADC_HTR																			<u> </u>			<u> </u>	HT[·	11:0	]				-
0x24	Reset value								Hes	erve	a								1	1	1	1	1	1	1	1	1	1	1	1
0x28	ADC_LTR								Doo	erve	۸													LT[1	1:0]				- 1	
0,20	Reset value							. '	103	SI VC	u .								0	0	0	0	0	0	0	0	0	0	0 (	0
0x2C	ADC_SQR1			Rese	erved					3:0]							Regula	r cha	anne		que	ence	SC	λx_x	bits	3				
0,120	Reset value							0	0	0	0					0			0	0	0	0	0	0	0	0	0	0	0 (	0
	ADC_SQR2	ved	L								F	Regul	ar	char	nne	l se	quence	SC	)x_x	bits	3									
0x30	Reset value	Reserved Reserved	0	0	0	0 (	0 (	0	0	0	0				0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0 (	0
0.04	ADC_SQR3	.ved	L								F	Regul	ar	char	nne	l se	quence	SC	)x_>	bits	}									
0x34	Reset value	Reser	0	0	0	0 (	0 (	0	0	0	0	0	0	0	0		0 0		0	0	0	0	0	0	0	0	0	0	0 (	0
0x38	ADC_JSQR	_		F	Reser	ved				JL[1		0.1.	<u>. Т</u>	_	<u> </u>		jected		_								_			_
	Reset value									0	0	0	0	0	0	0	0 0	0	0	0	0	0	0		0	0	0	0	0 (	0
0x3C	ADC_JDR1						Re	serve	ed						-	0	0 1 0	0	١٨	_		ATA			_	0	0	0	0 (	
	Reset value ADC_JDR2															U	0 0	U	U	U		)ATA			U	U	U	U	0 (	_
0x40	Reset value						Re	serve	ed						F	0	0 0	0	0	0			•	0	٥	0	0	0	0 (	0
	ADC_JDR3															Ů	0 0	Ŭ	Ů	Ŭ		)ATA			Ů	Ü	Ů	Ů	<u> </u>	-
0x44	Reset value						Re	serve	ed						-	0	0 0	0	0	0			•	0	0	0	0	0	0 (	0
	ADC_JDR4															-	-   -					)ATA		1		-	-		<u> </u>	$\exists$
0x48	Reset value						Re	serve	ed						-	0	0 0	0	0	0			-	0	0	0	0	0	0 (	0
010	ADC_DR						_									1	- 1	1	·					A[15		<u> </u>	1			$\dashv$
0x4C	Reset value						Re	serve	ea						ľ	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0 (	0
	l																						ı			<u> </u>	1		!_	_

Table 42. ADC register map and reset values (common ADC registers)

		_				_																	_										
Offset	Register	31	30	29	88	22	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	9	9	4	ε	2	1	0
	ADC_CSR											OVR	STRT	JSTRT	JEOC	EOC	AWD	Res		OVR	STRT	JSTRT	JEOC	EOC	AWD	Re	ser	OVR	STRT	JSTRT	JEOC	EOC	AWD
0x00	Reset value				۲	Rese	rve	d				0	0	0	0	0	0	ve		0	0	0	0	0	0	VE		0	0	0	0	0	0
														AD	C3							AD	C2							AD	C1		
0x04	ADC_CCR			F	Rese	erved	d			TSVREFE	VBATE	F	Rese	erve	d	ADCPBE[1:0]	מייין ייין	DMA[1:0]		DDS	Re se rv ed	DE	ΞLΑ	Y [3	:0]	Re	serv	/ed		MUI	_TI [	4:0]	
	Reset value									0	0					0	0	0	0	0		0	0	0	0				0	0	0	0	0
0x08	ADC_CDR						Re	gula	ır D/	λTΑ	2[15	:0]											Re	gula	ar D	AΤΑ	1[15	5:0]					
0,000	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

Refer to *Table 2 on page 50* for the register boundary addresses.

# 11 Digital-to-analog converter (DAC)

This section applies to the whole STM32F4xx family, unless otherwise specified.

# 11.1 DAC introduction

The DAC module is a 12-bit, voltage output digital-to-analog converter. The DAC can be configured in 8- or 12-bit mode and may be used in conjunction with the DMA controller. In 12-bit mode, the data could be left- or right-aligned. The DAC has two output channels, each with its own converter. In dual DAC channel mode, conversions could be done independently or simultaneously when both channels are grouped together for synchronous update operations. An input reference pin,  $V_{\text{REF+}}$  (shared with ADC) is available for better resolution.

# 11.2 DAC main features

- Two DAC converters: one output channel each
- Left or right data alignment in 12-bit mode
- Synchronized update capability
- Noise-wave generation
- Triangular-wave generation
- Dual DAC channel for independent or simultaneous conversions
- DMA capability for each channel
- DMA underrun error detection
- External triggers for conversion
- Input voltage reference, V<sub>REF+</sub>

Figure 48 shows the block diagram of a DAC channel and Table 43 gives the pin description.

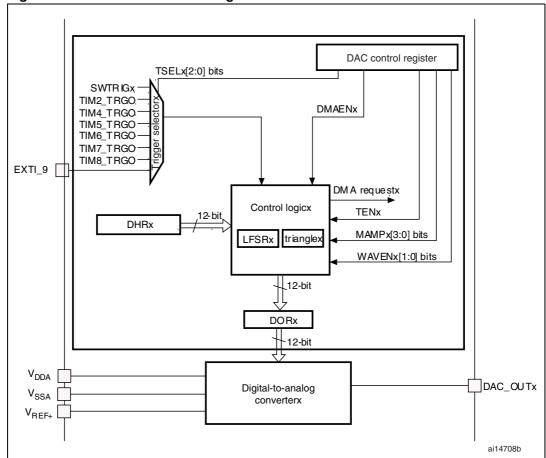


Figure 48. DAC channel block diagram

Table 43. DAC pins

Name	Signal type	Remarks
V <sub>REF+</sub>	Input, analog reference positive	The higher/positive reference voltage for the DAC, 1.8 V $\leq$ V <sub>REF+</sub> $\leq$ V <sub>DDA</sub>
$V_{DDA}$	Input, analog supply	Analog power supply
V <sub>SSA</sub>	Input, analog supply ground	Ground for analog power supply
DAC_OUTx	Analog output signal	DAC channelx analog output

Note:

Once the DAC channelx is enabled, the corresponding GPIO pin (PA4 or PA5) is automatically connected to the analog converter output (DAC\_OUTx). In order to avoid parasitic consumption, the PA4 or PA5 pin should first be configured to analog (AIN).

#### 11.3 **DAC** functional description

#### 11.3.1 **DAC** channel enable

Each DAC channel can be powered on by setting its corresponding ENx bit in the DAC\_CR register. The DAC channel is then enabled after a startup time twakfur

Note:

The ENx bit enables the analog DAC Channelx macrocell only. The DAC Channelx digital interface is enabled even if the ENx bit is reset.

#### **DAC** output buffer enable 11.3.2

The DAC integrates two output buffers that can be used to reduce the output impedance, and to drive external loads directly without having to add an external operational amplifier. Each DAC channel output buffer can be enabled and disabled using the corresponding BOFFx bit in the DAC CR register.

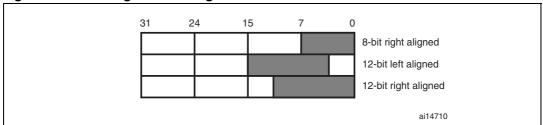
#### 11.3.3 **DAC** data format

Depending on the selected configuration mode, the data have to be written into the specified register as described below:

- Single DAC channelx, there are three possibilities:
  - 8-bit right alignment: the software has to load data into the DAC\_DHR8Rx [7:0] bits (stored into the DHRx[11:4] bits)
  - 12-bit left alignment: the software has to load data into the DAC\_DHR12Lx [15:4] bits (stored into the DHRx[11:0] bits)
  - 12-bit right alignment: the software has to load data into the DAC\_DHR12Rx [11:0] bits (stored into the DHRx[11:0] bits)

Depending on the loaded DAC\_DHRyyyx register, the data written by the user is shifted and stored into the corresponding DHRx (data holding registerx, which are internal nonmemory-mapped registers). The DHRx register is then loaded into the DORx register either automatically, by software trigger or by an external event trigger.

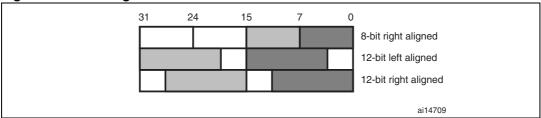
Figure 49. Data registers in single DAC channel mode



- Dual DAC channels, there are three possibilities:
  - 8-bit right alignment: data for DAC channel1 to be loaded into the DAC\_DHR8RD
     [7:0] bits (stored into the DHR1[11:4] bits) and data for DAC channel2 to be loaded into the DAC\_DHR8RD [15:8] bits (stored into the DHR2[11:4] bits)
  - 12-bit left alignment: data for DAC channel1 to be loaded into the DAC\_DHR12LD [15:4] bits (stored into the DHR1[11:0] bits) and data for DAC channel2 to be loaded into the DAC\_DHR12LD [31:20] bits (stored into the DHR2[11:0] bits)
  - 12-bit right alignment: data for DAC channel1 to be loaded into the DAC\_DHR12RD [11:0] bits (stored into the DHR1[11:0] bits) and data for DAC channel2 to be loaded into the DAC\_DHR12LD [27:16] bits (stored into the DHR2[11:0] bits)

Depending on the loaded DAC\_DHRyyyD register, the data written by the user is shifted and stored into DHR1 and DHR2 (data holding registers, which are internal non-memory-mapped registers). The DHR1 and DHR2 registers are then loaded into the DOR1 and DOR2 registers, respectively, either automatically, by software trigger or by an external event trigger.

Figure 50. Data registers in dual DAC channel mode



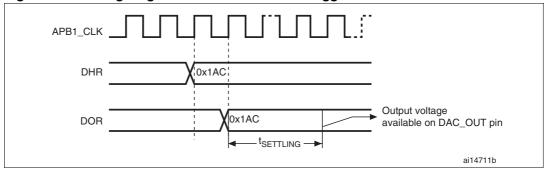
## 11.3.4 DAC conversion

The DAC\_DORx cannot be written directly and any data transfer to the DAC channelx must be performed by loading the DAC\_DHRx register (write to DAC\_DHR8Rx, DAC\_DHR12Lx, DAC\_DHR12Rx, DAC\_DHR8RD, DAC\_DHR12LD or DAC\_DHR12LD).

Data stored in the DAC\_DHRx register are automatically transferred to the DAC\_DORx register after one APB1 clock cycle, if no hardware trigger is selected (TENx bit in DAC\_CR register is reset). However, when a hardware trigger is selected (TENx bit in DAC\_CR register is set) and a trigger occurs, the transfer is performed three APB1 clock cycles later.

When DAC\_DORx is loaded with the DAC\_DHRx contents, the analog output voltage becomes available after a time t<sub>SETTLING</sub> that depends on the power supply voltage and the analog output load.

Figure 51. Timing diagram for conversion with trigger disabled TEN = 0



# 11.3.5 DAC output voltage

Digital inputs are converted to output voltages on a linear conversion between 0 and V<sub>RFF+</sub>.

The analog output voltages on each DAC channel pin are determined by the following equation:

DACoutput = 
$$V_{REF} \times \frac{DOR}{4095}$$

# 11.3.6 DAC trigger selection

If the TENx control bit is set, conversion can then be triggered by an external event (timer counter, external interrupt line). The TSELx[2:0] control bits determine which out of 8 possible events will trigger conversion as shown in *Table 44*.

Table 44. External triggers

Source	Туре	TSEL[2:0]
Timer 6 TRGO event		000
Timer 8 TRGO event		001
Timer 7 TRGO event	Internal signal from on-chip	010
Timer 5 TRGO event	timers	011
Timer 2 TRGO event		100
Timer 4 TRGO event		101
EXTI line9	External pin	110
SWTRIG	Software control bit	111

Each time a DAC interface detects a rising edge on the selected timer TRGO output, or on the selected external interrupt line 9, the last data stored into the DAC\_DHRx register are transferred into the DAC\_DORx register. The DAC\_DORx register is updated three APB1 cycles after the trigger occurs.

If the software trigger is selected, the conversion starts once the SWTRIG bit is set. SWTRIG is reset by hardware once the DAC\_DORx register has been loaded with the DAC\_DHRx register contents.

Note:

TSELx[2:0] bit cannot be changed when the ENx bit is set.

When software trigger is selected, the transfer from the DAC\_DHRx register to the DAC\_DORx register takes only one APB1 clock cycle.

# 11.3.7 DMA request

Each DAC channel has a DMA capability. Two DMA channels are used to service DAC channel DMA requests.

A DAC DMA request is generated when an external trigger (but not a software trigger) occurs while the DMAENx bit is set. The value of the DAC\_DHRx register is then transferred into the DAC\_DORx register.

In dual mode, if both DMAENx bits are set, two DMA requests are generated. If only one DMA request is needed, you should set only the corresponding DMAENx bit. In this way, the application can manage both DAC channels in dual mode by using one DMA request and a unique DMA channel.

#### **DMA** underrun

The DAC DMA request is not queued so that if a second external trigger arrives before the acknowledgement for the first external trigger is received (first request), then no new request is issued and the DMA channelx underrun flag DMAUDRx in the DAC\_SR register is set, reporting the error condition. DMA data transfers are then disabled and no further DMA request is treated. The DAC channelx continues to convert old data.

The software should clear the DMAUDRx flag by writing "1", clear the DMAEN bit of the used DMA stream and re-initialize both DMA and DAC channelx to restart the transfer correctly. The software should modify the DAC trigger conversion frequency or lighten the DMA workload to avoid a new DMA underrun. Finally, the DAC conversion could be resumed by enabling both DMA data transfer and conversion trigger.

For each DAC channelx, an interrupt is also generated if its corresponding DMAUDRIEx bit in the DAC CR register is enabled.

# 11.3.8 Noise generation

In order to generate a variable-amplitude pseudonoise, an LFSR (linear feedback shift register) is available. DAC noise generation is selected by setting WAVEx[1:0] to "01". The preloaded value in LFSR is 0xAAA. This register is updated three APB1 clock cycles after each trigger event, following a specific calculation algorithm.

Figure 52. DAC LFSR register calculation algorithm

The LFSR value, that may be masked partially or totally by means of the MAMPx[3:0] bits in the DAC\_CR register, is added up to the DAC\_DHRx contents without overflow and this value is then stored into the DAC\_DORx register.

If LFSR is 0x0000, a '1 is injected into it (antilock-up mechanism).

It is possible to reset LFSR wave generation by resetting the WAVEx[1:0] bits.

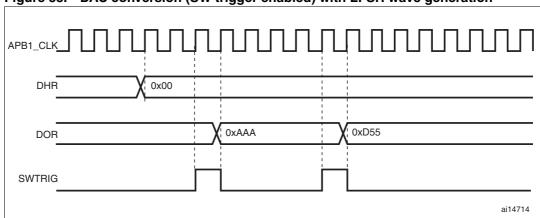


Figure 53. DAC conversion (SW trigger enabled) with LFSR wave generation

Note:

The DAC trigger must be enabled for noise generation by setting the TENx bit in the DAC\_CR register.

# 11.3.9 Triangle-wave generation

It is possible to add a small-amplitude triangular waveform on a DC or slowly varying signal. DAC triangle-wave generation is selected by setting WAVEx[1:0] to "10". The amplitude is configured through the MAMPx[3:0] bits in the DAC\_CR register. An internal triangle counter is incremented three APB1 clock cycles after each trigger event. The value of this counter is then added to the DAC\_DHRx register without overflow and the sum is stored into the DAC\_DORx register. The triangle counter is incremented as long as it is less than the maximum amplitude defined by the MAMPx[3:0] bits. Once the configured amplitude is reached, the counter is decremented down to 0, then incremented again and so on.

It is possible to reset triangle wave generation by resetting the WAVEx[1:0] bits.

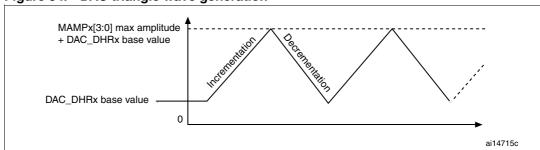


Figure 54. DAC triangle wave generation

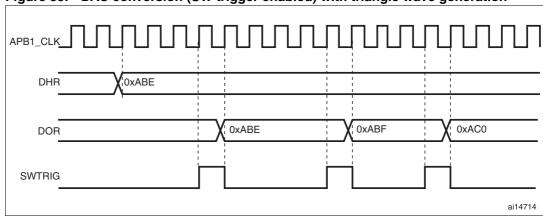


Figure 55. DAC conversion (SW trigger enabled) with triangle wave generation

Note:

The DAC trigger must be enabled for noise generation by setting the TENx bit in the DAC\_CR register.

The MAMPx[3:0] bits must be configured before enabling the DAC, otherwise they cannot be changed.

# 11.4 Dual DAC channel conversion

To efficiently use the bus bandwidth in applications that require the two DAC channels at the same time, three dual registers are implemented: DHR8RD, DHR12RD and DHR12LD. A unique register access is then required to drive both DAC channels at the same time.

Eleven possible conversion modes are possible using the two DAC channels and these dual registers. All the conversion modes can nevertheless be obtained using separate DHRx registers if needed.

All modes are described in the paragraphs below.

# 11.4.1 Independent trigger without wave generation

To configure the DAC in this conversion mode, the following sequence is required:

- Set the two DAC channel trigger enable bits TEN1 and TEN2
- Configure different trigger sources by setting different values in the TSEL1[2:0] and TSEL2[2:0] bits
- Load the dual DAC channel data into the desired DHR register (DAC\_DHR12RD, DAC\_DHR12LD or DAC\_DHR8RD)

When a DAC channel1 trigger arrives, the DHR1 register is transferred into DAC\_DOR1 (three APB1 clock cycles later).

When a DAC channel2 trigger arrives, the DHR2 register is transferred into DAC\_DOR2 (three APB1 clock cycles later).

# 11.4.2 Independent trigger with single LFSR generation

To configure the DAC in this conversion mode, the following sequence is required:

- Set the two DAC channel trigger enable bits TEN1 and TEN2
- Configure different trigger sources by setting different values in the TSEL1[2:0] and TSEL2[2:0] bits
- Configure the two DAC channel WAVEx[1:0] bits as "01" and the same LFSR mask value in the MAMPx[3:0] bits
- Load the dual DAC channel data into the desired DHR register (DHR12RD, DHR12LD or DHR8RD)

When a DAC channel1 trigger arrives, the LFSR1 counter, with the same mask, is added to the DHR1 register and the sum is transferred into DAC\_DOR1 (three APB1 clock cycles later). Then the LFSR1 counter is updated.

When a DAC channel2 trigger arrives, the LFSR2 counter, with the same mask, is added to the DHR2 register and the sum is transferred into DAC\_DOR2 (three APB1 clock cycles later). Then the LFSR2 counter is updated.

# 11.4.3 Independent trigger with different LFSR generation

To configure the DAC in this conversion mode, the following sequence is required:

- Set the two DAC channel trigger enable bits TEN1 and TEN2
- Configure different trigger sources by setting different values in the TSEL1[2:0] and TSEL2[2:0] bits
- Configure the two DAC channel WAVEx[1:0] bits as "01" and set different LFSR masks values in the MAMP1[3:0] and MAMP2[3:0] bits
- Load the dual DAC channel data into the desired DHR register (DAC\_DHR12RD, DAC\_DHR12LD or DAC\_DHR8RD)

When a DAC channel1 trigger arrives, the LFSR1 counter, with the mask configured by MAMP1[3:0], is added to the DHR1 register and the sum is transferred into DAC\_DOR1 (three APB1 clock cycles later). Then the LFSR1 counter is updated.

When a DAC channel2 trigger arrives, the LFSR2 counter, with the mask configured by MAMP2[3:0], is added to the DHR2 register and the sum is transferred into DAC\_DOR2 (three APB1 clock cycles later). Then the LFSR2 counter is updated.

# 11.4.4 Independent trigger with single triangle generation

To configure the DAC in this conversion mode, the following sequence is required:

- Set the two DAC channel trigger enable bits TEN1 and TEN2
- Configure different trigger sources by setting different values in the TSEL1[2:0] and TSEL2[2:0] bits
- Configure the two DAC channel WAVEx[1:0] bits as "1x" and the same maximum amplitude value in the MAMPx[3:0] bits
- Load the dual DAC channel data into the desired DHR register (DAC\_DHR12RD, DAC\_DHR12LD or DAC\_DHR8RD)

When a DAC channel1 trigger arrives, the DAC channel1 triangle counter, with the same triangle amplitude, is added to the DHR1 register and the sum is transferred into

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DAC\_DOR1 (three APB1 clock cycles later). The DAC channel1 triangle counter is then updated.

When a DAC channel2 trigger arrives, the DAC channel2 triangle counter, with the same triangle amplitude, is added to the DHR2 register and the sum is transferred into DAC\_DOR2 (three APB1 clock cycles later). The DAC channel2 triangle counter is then updated.

# 11.4.5 Independent trigger with different triangle generation

To configure the DAC in this conversion mode, the following sequence is required:

- Set the two DAC channel trigger enable bits TEN1 and TEN2
- Configure different trigger sources by setting different values in the TSEL1[2:0] and TSEL2[2:0] bits
- Configure the two DAC channel WAVEx[1:0] bits as "1x" and set different maximum amplitude values in the MAMP1[3:0] and MAMP2[3:0] bits
- Load the dual DAC channel data into the desired DHR register (DAC\_DHR12RD, DAC\_DHR12LD or DAC\_DHR8RD)

When a DAC channel1 trigger arrives, the DAC channel1 triangle counter, with a triangle amplitude configured by MAMP1[3:0], is added to the DHR1 register and the sum is transferred into DAC\_DOR1 (three APB1 clock cycles later). The DAC channel1 triangle counter is then updated.

When a DAC channel2 trigger arrives, the DAC channel2 triangle counter, with a triangle amplitude configured by MAMP2[3:0], is added to the DHR2 register and the sum is transferred into DAC\_DOR2 (three APB1 clock cycles later). The DAC channel2 triangle counter is then updated.

## 11.4.6 Simultaneous software start

To configure the DAC in this conversion mode, the following sequence is required:

 Load the dual DAC channel data to the desired DHR register (DAC\_DHR12RD, DAC\_DHR12LD or DAC\_DHR8RD)

In this configuration, one APB1 clock cycle later, the DHR1 and DHR2 registers are transferred into DAC\_DOR1 and DAC\_DOR2, respectively.

# 11.4.7 Simultaneous trigger without wave generation

To configure the DAC in this conversion mode, the following sequence is required:

- Set the two DAC channel trigger enable bits TEN1 and TEN2
- Configure the same trigger source for both DAC channels by setting the same value in the TSEL1[2:0] and TSEL2[2:0] bits
- Load the dual DAC channel data to the desired DHR register (DAC\_DHR12RD, DAC\_DHR12LD or DAC\_DHR8RD)

When a trigger arrives, the DHR1 and DHR2 registers are transferred into DAC\_DOR1 and DAC\_DOR2, respectively (after three APB1 clock cycles).

# 11.4.8 Simultaneous trigger with single LFSR generation

To configure the DAC in this conversion mode, the following sequence is required:

- Set the two DAC channel trigger enable bits TEN1 and TEN2
- Configure the same trigger source for both DAC channels by setting the same value in the TSEL1[2:0] and TSEL2[2:0] bits
- Configure the two DAC channel WAVEx[1:0] bits as "01" and the same LFSR mask value in the MAMPx[3:0] bits
- Load the dual DAC channel data to the desired DHR register (DHR12RD, DHR12LD or DHR8RD)

When a trigger arrives, the LFSR1 counter, with the same mask, is added to the DHR1 register and the sum is transferred into DAC\_DOR1 (three APB1 clock cycles later). The LFSR1 counter is then updated. At the same time, the LFSR2 counter, with the same mask, is added to the DHR2 register and the sum is transferred into DAC\_DOR2 (three APB1 clock cycles later). The LFSR2 counter is then updated.

# 11.4.9 Simultaneous trigger with different LFSR generation

To configure the DAC in this conversion mode, the following sequence is required:

- Set the two DAC channel trigger enable bits TEN1 and TEN2
- Configure the same trigger source for both DAC channels by setting the same value in the TSEL1[2:0] and TSEL2[2:0] bits
- Configure the two DAC channel WAVEx[1:0] bits as "01" and set different LFSR mask values using the MAMP1[3:0] and MAMP2[3:0] bits
- Load the dual DAC channel data into the desired DHR register (DAC\_DHR12RD, DAC\_DHR12LD or DAC\_DHR8RD)

When a trigger arrives, the LFSR1 counter, with the mask configured by MAMP1[3:0], is added to the DHR1 register and the sum is transferred into DAC\_DOR1 (three APB1 clock cycles later). The LFSR1 counter is then updated.

At the same time, the LFSR2 counter, with the mask configured by MAMP2[3:0], is added to the DHR2 register and the sum is transferred into DAC\_DOR2 (three APB1 clock cycles later). The LFSR2 counter is then updated.

# 11.4.10 Simultaneous trigger with single triangle generation

To configure the DAC in this conversion mode, the following sequence is required:

- Set the two DAC channel trigger enable bits TEN1 and TEN2
- Configure the same trigger source for both DAC channels by setting the same value in the TSEL1[2:0] and TSEL2[2:0] bits
- Configure the two DAC channel WAVEx[1:0] bits as "1x" and the same maximum amplitude value using the MAMPx[3:0] bits
- Load the dual DAC channel data into the desired DHR register (DAC\_DHR12RD, DAC\_DHR12LD or DAC\_DHR8RD)

When a trigger arrives, the DAC channel1 triangle counter, with the same triangle amplitude, is added to the DHR1 register and the sum is transferred into DAC\_DOR1 (three APB1 clock cycles later). The DAC channel1 triangle counter is then updated.

At the same time, the DAC channel2 triangle counter, with the same triangle amplitude, is

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added to the DHR2 register and the sum is transferred into DAC\_DOR2 (three APB1 clock cycles later). The DAC channel2 triangle counter is then updated.

# 11.4.11 Simultaneous trigger with different triangle generation

To configure the DAC in this conversion mode, the following sequence is required:

- Set the two DAC channel trigger enable bits TEN1 and TEN2
- Configure the same trigger source for both DAC channels by setting the same value in the TSEL1[2:0] and TSEL2[2:0] bits
- Configure the two DAC channel WAVEx[1:0] bits as "1x" and set different maximum amplitude values in the MAMP1[3:0] and MAMP2[3:0] bits
- Load the dual DAC channel data into the desired DHR register (DAC\_DHR12RD, DAC\_DHR12LD or DAC\_DHR8RD)

When a trigger arrives, the DAC channel1 triangle counter, with a triangle amplitude configured by MAMP1[3:0], is added to the DHR1 register and the sum is transferred into DAC\_DOR1 (three APB1 clock cycles later). Then the DAC channel1 triangle counter is updated.

At the same time, the DAC channel2 triangle counter, with a triangle amplitude configured by MAMP2[3:0], is added to the DHR2 register and the sum is transferred into DAC\_DOR2 (three APB1 clock cycles later). Then the DAC channel2 triangle counter is updated.

# 11.5 DAC registers

Refer to Section 1.1 on page 46 for a list of abbreviations used in register descriptions.

The peripheral registers have to be accessed by words (32 bits).

# 11.5.1 DAC control register (DAC CR)

Address offset: 0x00

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
Rese	erved	DMAU DRIE2	DMA EN2		MAM	P2[3:0]		WAVE	[2[1:0]	٦	SEL2[2:0	)]	TEN2	BOFF2	EN2
		rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Rese	erved	DMAU DRIE1	DMA EN1		MAM	P1[3:0]		WAVE	[1:0]	٦	SEL1[2:0	)]	TEN1	BOFF1	EN1
		rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:30 Reserved, must be kept at reset value.

Bits 29 **DMAUDRIE2**: DAC channel2 DMA underrun interrupt enable

This bit is set and cleared by software.

0: DAC channel2 DMA underrun interrupt disabled

1: DAC channel2 DMA underrun interrupt enabled

Bit 28 DMAEN2: DAC channel2 DMA enable

This bit is set and cleared by software.

0: DAC channel2 DMA mode disabled

1: DAC channel2 DMA mode enabled

#### Bit 27:24 MAMP2[3:0]: DAC channel2 mask/amplitude selector

These bits are written by software to select mask in wave generation mode or amplitude in triangle generation mode.

0000: Unmask bit0 of LFSR/ triangle amplitude equal to 1

0001: Unmask bits[1:0] of LFSR/ triangle amplitude equal to 3

0010: Unmask bits[2:0] of LFSR/ triangle amplitude equal to 7

0011: Unmask bits[3:0] of LFSR/ triangle amplitude equal to 15

0100: Unmask bits[4:0] of LFSR/ triangle amplitude equal to 31

0101: Unmask bits[5:0] of LFSR/ triangle amplitude equal to 63

0110: Unmask bits[6:0] of LFSR/ triangle amplitude equal to 127

0111: Unmask bits[7:0] of LFSR/ triangle amplitude equal to 255

1000: Unmask bits[8:0] of LFSR/ triangle amplitude equal to 511

1001: Unmask bits[9:0] of LFSR/ triangle amplitude equal to 1023

1010: Unmask bits[10:0] of LFSR/ triangle amplitude equal to 2047

≥ 1011: Unmask bits[11:0] of LFSR/ triangle amplitude equal to 4095

### Bit 23:22 WAVE2[1:0]: DAC channel2 noise/triangle wave generation enable

These bits are set/reset by software.

00: wave generation disabled

01: Noise wave generation enabled

1x: Triangle wave generation enabled

Note: Only used if bit TEN2 = 1 (DAC channel2 trigger enabled)

## Bits 21:19 TSEL2[2:0]: DAC channel2 trigger selection

These bits select the external event used to trigger DAC channel2

000: Timer 6 TRGO event

001: Timer 8 TRGO event

010: Timer 7 TRGO event

011: Timer 5 TRGO event

100: Timer 2 TRGO event

101: Timer 4 TRGO event

110: External line9111: Software trigger

Note: Only used if bit TEN2 = 1 (DAC channel2 trigger enabled).

## Bit 18 TEN2: DAC channel2 trigger enable

This bit is set and cleared by software to enable/disable DAC channel2 trigger

0: DAC channel2 trigger disabled and data written into the DAC\_DHRx register are transferred one APB1 clock cycle later to the DAC\_DOR2 register

1: DAC channel2 trigger enabled and data from the DAC\_DHRx register are transferred three APB1 clock cycles later to the DAC\_DOR2 register

Note: When software trigger is selected, the transfer from the DAC\_DHRx register to the DAC\_DOR2 register takes only one APB1 clock cycle.

## Bit 17 BOFF2: DAC channel2 output buffer disable

This bit is set and cleared by software to enable/disable DAC channel2 output buffer.

0: DAC channel2 output buffer enabled

1: DAC channel2 output buffer disabled

## Bit 16 EN2: DAC channel2 enable

This bit is set and cleared by software to enable/disable DAC channel2.

0: DAC channel2 disabled

1: DAC channel2 enabled

Bits 15:14 Reserved, must be kept at reset value.

### Bit 13 DMAUDRIE1: DAC channel1 DMA Underrun Interrupt enable

This bit is set and cleared by software.

0: DAC channel1 DMA Underrun Interrupt disabled

1: DAC channel1 DMA Underrun Interrupt enabled

#### Bit 12 DMAEN1: DAC channel1 DMA enable

This bit is set and cleared by software.

0: DAC channel1 DMA mode disabled

1: DAC channel1 DMA mode enabled

### Bits 11:8 MAMP1[3:0]: DAC channel1 mask/amplitude selector

These bits are written by software to select mask in wave generation mode or amplitude in triangle generation mode.

0000: Unmask bit0 of LFSR/ triangle amplitude equal to 1

0001: Unmask bits[1:0] of LFSR/ triangle amplitude equal to 3

0010: Unmask bits[2:0] of LFSR/ triangle amplitude equal to 7

0011: Unmask bits[3:0] of LFSR/ triangle amplitude equal to 15

0100: Unmask bits[4:0] of LFSR/ triangle amplitude equal to 31

0101: Unmask bits[5:0] of LFSR/ triangle amplitude equal to 63

0110: Unmask bits[6:0] of LFSR/ triangle amplitude equal to 127

0111: Unmask bits[7:0] of LFSR/ triangle amplitude equal to 255

1000: Unmask bits[8:0] of LFSR/ triangle amplitude equal to 511

1001: Unmask bits[9:0] of LFSR/ triangle amplitude equal to 1023

1010: Unmask bits[10:0] of LFSR/ triangle amplitude equal to 2047

≥ 1011: Unmask bits[11:0] of LFSR/ triangle amplitude equal to 4095

### Bits 7:6 WAVE1[1:0]: DAC channel1 noise/triangle wave generation enable

These bits are set and cleared by software.

00: wave generation disabled

01: Noise wave generation enabled

1x: Triangle wave generation enabled

Note: Only used if bit TEN1 = 1 (DAC channel1 trigger enabled).

## Bits 5:3 TSEL1[2:0]: DAC channel1 trigger selection

These bits select the external event used to trigger DAC channel1.

000: Timer 6 TRGO event

001: Timer 8 TRGO event

010: Timer 7 TRGO event

011: Timer 5 TRGO event

100: Timer 2 TRGO event

101: Timer 4 TRGO event

110: External line9

111: Software trigger

Note: Only used if bit TEN1 = 1 (DAC channel1 trigger enabled).

## Bit 2 TEN1: DAC channel1 trigger enable

This bit is set and cleared by software to enable/disable DAC channel1 trigger.

0: DAC channel1 trigger disabled and data written into the DAC\_DHRx register are transferred one APB1 clock cycle later to the DAC\_DOR1 register

1: DAC channel1 trigger enabled and data from the DAC\_DHRx register are transferred three APB1 clock cycles later to the DAC\_DOR1 register

Note: When software trigger is selected, the transfer from the DAC\_DHRx register to the DAC\_DOR1 register takes only one APB1 clock cycle.



## Bit 1 BOFF1: DAC channel1 output buffer disable

This bit is set and cleared by software to enable/disable DAC channel1 output buffer.

0: DAC channel1 output buffer enabled1: DAC channel1 output buffer disabled

#### Bit 0 EN1: DAC channel1 enable

This bit is set and cleared by software to enable/disable DAC channel1.

0: DAC channel1 disabled 1: DAC channel1 enabled

# 11.5.2 DAC software trigger register (DAC\_SWTRIGR)

Address offset: 0x04 Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							F	Reserved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
						Day	served							SWTRIG2	SWTRIG1
						nes	sei veu							w	w

Bits 31:2 Reserved, must be kept at reset value.

## Bit 1 SWTRIG2: DAC channel2 software trigger

This bit is set and cleared by software to enable/disable the software trigger.

0: Software trigger disabled

1: Software trigger enabled

Note: This bit is cleared by hardware (one APB1 clock cycle later) once the DAC\_DHR2 register value has been loaded into the DAC\_DOR2 register.

### Bit 0 SWTRIG1: DAC channel1 software trigger

This bit is set and cleared by software to enable/disable the software trigger.

0: Software trigger disabled

1: Software trigger enabled

Note: This bit is cleared by hardware (one APB1 clock cycle later) once the DAC\_DHR1 register value has been loaded into the DAC\_DOR1 register.

# 11.5.3 DAC channel1 12-bit right-aligned data holding register (DAC\_DHR12R1)

Address offset: 0x08 Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Res	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Poor	mod							DACC1D	HR[11:0]					
	Rese	erveu		rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:12 Reserved, must be kept at reset value.

## Bit 11:0 DACC1DHR[11:0]: DAC channel1 12-bit right-aligned data

These bits are written by software which specifies 12-bit data for DAC channel1.

# 11.5.4 DAC channel1 12-bit left aligned data holding register (DAC\_DHR12L1)

Address offset: 0x0C

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
					DACC1E	HR[11:0]							Rese	rvod	
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw		nese	riveu	

Bits 31:16 Reserved, must be kept at reset value.

Bit 15:4 DACC1DHR[11:0]: DAC channel1 12-bit left-aligned data

These bits are written by software which specifies 12-bit data for DAC channel1.

Bits 3:0 Reserved, must be kept at reset value.

# 11.5.5 DAC channel1 8-bit right aligned data holding register (DAC\_DHR8R1)

Address offset: 0x10

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Res	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			Pos	erved							DACC1	DHR[7:0]			
			nes	erveu				rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:8 Reserved, must be kept at reset value.

Bits 7:0 DACC1DHR[7:0]: DAC channel1 8-bit right-aligned data

These bits are written by software which specifies 8-bit data for DAC channel1.

# 11.5.6 DAC channel2 12-bit right aligned data holding register (DAC\_DHR12R2)

Address offset: 0x14

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Rese	.m.co.d							DACC2D	HR[11:0]					
	nese	riveu		rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:12 Reserved, must be kept at reset value.

Bits 11:0 DACC2DHR[11:0]: DAC channel2 12-bit right-aligned data

These bits are written by software which specifies 12-bit data for DAC channel2.

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# 11.5.7 DAC channel2 12-bit left aligned data holding register (DAC\_DHR12L2)

Address offset: 0x18

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Res	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
					DACC2E	DHR[11:0]							Rese	rvod	
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw		nese	ei veu	

Bits 31:16 Reserved, must be kept at reset value.

Bits 15:4 DACC2DHR[11:0]: DAC channel2 12-bit left-aligned data

These bits are written by software which specify 12-bit data for DAC channel2.

Bits 3:0 Reserved, must be kept at reset value.

# 11.5.8 DAC channel2 8-bit right-aligned data holding register (DAC\_DHR8R2)

Address offset: 0x1C

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
 15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			Pos	erved							DACC2	DHR[7:0]			
			nes	erveu				rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:8 Reserved, must be kept at reset value.

Bits 7:0 DACC2DHR[7:0]: DAC channel2 8-bit right-aligned data

These bits are written by software which specifies 8-bit data for DAC channel2.

# 11.5.9 Dual DAC 12-bit right-aligned data holding register (DAC\_DHR12RD)

Address offset: 0x20

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	Rese	vrvod							DACC2D	HR[11:0]					
	nese	erveu		rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Reserved							DACC1D	HR[11:0]						
	nese	ei veu		rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

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Bits 31:28 Reserved, must be kept at reset value.

Bits 27:16 DACC2DHR[11:0]: DAC channel2 12-bit right-aligned data

These bits are written by software which specifies 12-bit data for DAC channel2.

Bits 15:12 Reserved, must be kept at reset value.

Bits 11:0 DACC1DHR[11:0]: DAC channel1 12-bit right-aligned data

These bits are written by software which specifies 12-bit data for DAC channel1.

# 11.5.10 DUAL DAC 12-bit left aligned data holding register (DAC\_DHR12LD)

Address offset: 0x24

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
					DACC2E	HR[11:0]							Poor	erved	
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw		nest	erveu	
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	DACC1DHR[11:0]												Pos	arvod	
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	Reserved			

Bits 31:20 DACC2DHR[11:0]: DAC channel2 12-bit left-aligned data

These bits are written by software which specifies 12-bit data for DAC channel2.

Bits 19:16 Reserved, must be kept at reset value.

Bits 15:4 DACC1DHR[11:0]: DAC channel1 12-bit left-aligned data

These bits are written by software which specifies 12-bit data for DAC channel1.

Bits 3:0 Reserved, must be kept at reset value.

# 11.5.11 DUAL DAC 8-bit right aligned data holding register (DAC\_DHR8RD)

Address offset: 0x28

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Res	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	DACC2DHR[7:0]										DACC1	DHR[7:0]			
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:16 Reserved, must be kept at reset value.

Bits 15:8 DACC2DHR[7:0]: DAC channel2 8-bit right-aligned data

These bits are written by software which specifies 8-bit data for DAC channel2.

Bits 7:0 DACC1DHR[7:0]: DAC channel1 8-bit right-aligned data

These bits are written by software which specifies 8-bit data for DAC channel1.

# 11.5.12 DAC channel1 data output register (DAC\_DOR1)

Address offset: 0x2C

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
Reserved															
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Reserved								DACC1D	OR[11:0]					
	nese	erved		r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:12 Reserved, must be kept at reset value.

Bit 11:0 DACC1DOR[11:0]: DAC channel1 data output

These bits are read-only, they contain data output for DAC channel1.

# 11.5.13 DAC channel2 data output register (DAC\_DOR2)

Address offset: 0x30 Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
Reserved															
 15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Rese	nyod						_	DACC2D	OR[11:0]					
	nese	i veu		r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:12 Reserved, must be kept at reset value.

Bit 11:0 DACC2DOR[11:0]: DAC channel2 data output

These bits are read-only, they contain data output for DAC channel2.

# 11.5.14 DAC status register (DAC\_SR)

Address offset: 0x34

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
Posc	erved	DMAUDR2							Reserve	4					
nese	erveu	rc_w1							neserve	u					
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Page	erved	DMAUDR1							Pagarya	4					
nese	erveu	rc_w1						Reserved							

Bits 31:30 Reserved, must be kept at reset value.

Bit 29 DMAUDR2: DAC channel2 DMA underrun flag

This bit is set by hardware and cleared by software (by writing it to 1).

- 0: No DMA underrun error condition occurred for DAC channel2
- 1: DMA underrun error condition occurred for DAC channel2 (the currently selected trigger is driving DAC channel2 conversion at a frequency higher than the DMA service capability rate)
- Bits 28:14 Reserved, must be kept at reset value.
  - Bit 13 DMAUDR1: DAC channel1 DMA underrun flag

This bit is set by hardware and cleared by software (by writing it to 1).

- 0: No DMA underrun error condition occurred for DAC channel1
- 1: DMA underrun error condition occurred for DAC channel1 (the currently selected trigger is driving DAC channel1 conversion at a frequency higher than the DMA service capability rate)

Bits 12:0 Reserved, must be kept at reset value.

# 11.5.15 DAC register map

Table 45 summarizes the DAC registers.

Table 45. DAC register map

Address offset	Register name	31	30	29	28	27	26	25	24	23	1 2	2	19	18	17	16	15	14	13	12	11	10	6	æ	7	9	5	4	က	2	1	0
0x00	DAC_CR	Reserved		<b>DMAUDRIE2</b>	DMAEN2	M	AMP:	2[3:0		WAVE 2[2:0]	T	SEL	2[2:0	TEN2	BOFF2	EN2	Reserved	200	DMAUDRIE1	DMAEN1	MAN	MP1	[3:0		WA\ 1[2:	VE :0]	TSE	EL1[:	2:0]	TEN1	BOFF1	EN1
0x04	DAC_SWT RIGR		Reserved ELM									SWTRIG2	SWTRIG1																			
0x08	DAC_DHR1 2R1		Reserved								D/	AC(	C1D	HR[	11:0	)]																
0x0C	DAC_DHR1 2L1		Reserved DACC1DHR[1							11:0	]				Res	erv	ed															
0x10	DAC_DHR8 R1		Reserved										I	DAC	CC1	DHF	R[7:0	0]														
0x14	DAC_DHR1 2R2									Res	serv	ved												D	AC(	C2D	HR[	11:0	)]			
0x18	DAC_DHR1 2L2							Re	ese	rved										D	ACC	C2DI	HR[	11:0	]				Res	erv	ed	
0x1C	DAC_DHR8 R2											Re	serve	d											DAC	CC2	DHF	R[7:0	0]			
0x20	DAC_DHR1 2RD	R	lese	erve	d				D	ACC2	DH	IR[1	1:0]				F	Rese	rve	t				D	AC(	C1D	HR[	11:0	)]			
0x24	DAC_DHR1 2LD				D	AC	C2DI	HR[1	1:0	)]				Res	erve	d				С	ACC	C1DI	HR[	11:0	]				F	lese	rve	b
0x28	DAC_DHR8 RD	Reserved DACC2DHR[7:0]										R[7:0	•																			
0x2C	DAC_DOR1		Reserved											D	AC(	C1D	OR[	11:0	)]													
0x30	DAC_DOR2									Res	serv	ved												DA	ACC	C2D	OR[	11:0	)]			
0x34	DAC_SR	Reserved	Reserved  Reserved  Reserved  Reserved									Re	serv	/ed																		

Refer to Table 2 on page 50 for the register boundary addresses.

# 12 Digital camera interface (DCMI)

This section applies to the whole STM32F4xx family, unless otherwise specified.

# 12.1 DCMI introduction

The digital camera is a synchronous parallel interface able to receive a high-speed data flow from an external 8-, 10-, 12- or 14-bit CMOS camera module. It supports different data formats: YCbCr4:2:2/RGB565 progressive video and compressed data (JPEG).

This interface is for use with black & white cameras, X24 and X5 cameras, and it is assumed that all pre-processing like resizing is performed in the camera module.

# 12.2 DCMI main features

- 8-, 10-, 12- or 14-bit parallel interface
- Embedded/external line and frame synchronization
- Continuous or snapshot mode
- Crop feature
- Supports the following data formats:
  - 8/10/12/14- bit progressive video: either monochrome or raw bayer
  - YCbCr 4:2:2 progressive video
  - RGB 565 progressive video
  - Compressed data: JPEG

# 12.3 DCMI pins

Table 46 shows the DCMI pins.

Table 46. DCMI pins

Name	Signal type
D[0:13]	Data inputs
HSYNC	Horizontal synchronization input
VSYNC	Vertical synchronization input
PIXCLX	Pixel clock input

# 12.4 DCMI clocks

The digital camera interface uses two clock domains PIXCLK and HCLK. The signals generated with PIXCLK are sampled on the rising edge of HCLK once they are stable. An enable signal is generated in the HCLK domain, to indicate that data coming from the camera are stable and can be sampled. The maximum PIXCLK period must be higher than 2.5 HCLK periods.

#### 12.5 **DCMI functional overview**

The digital camera interface is a synchronous parallel interface that can receive high-speed (up to 54 Mbytes/s) data flows. It consists of up to 14 data lines (D13-D0) and a pixel clock line (PIXCLK). The pixel clock has a programmable polarity, so that data can be captured on either the rising or the falling edge of the pixel clock.

The data are packed into a 32-bit data register (DCMI\_DR) and then transferred through a general-purpose DMA channel. The image buffer is managed by the DMA, not by the camera interface.

The data received from the camera can be organized in lines/frames (raw YUB/RGB/Bayer modes) or can be a sequence of JPEG images. To enable JPEG image reception, the JPEG bit (bit 3 of DCMI CR register) must be set.

The data flow is synchronized either by hardware using the optional HSYNC (horizontal synchronization) and VSYNC (vertical synchronization) signals or by synchronization codes embedded in the data flow.

Figure 56 shows the DCMI block diagram.

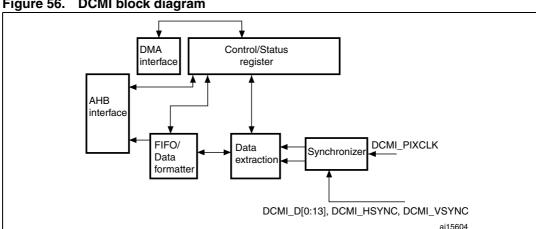
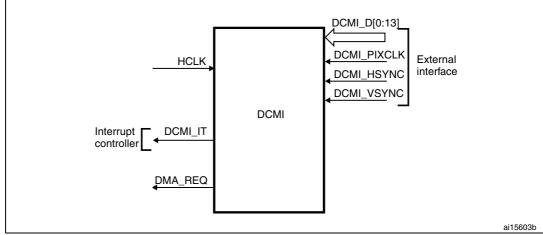


Figure 56. **DCMI block diagram** 

Figure 57. Top-level block diagram



## 12.5.1 DMA interface

The DMA interface is active when the CAPTURE bit in the DCMI\_CR register is set. A DMA request is generated each time the camera interface receives a complete 32-bit data block in its register.

# 12.5.2 DCMI physical interface

The interface is composed of 11/13/15/17 inputs. Only the Slave mode is supported.

The camera interface can capture 8-bit, 10-bit, 12-bit or 14-bit data depending on the EDM[1:0] bits in the DCMI\_CR register. If less than 14 bits are used, the unused input pins must be connected to ground.

Table 47. DCMI signals

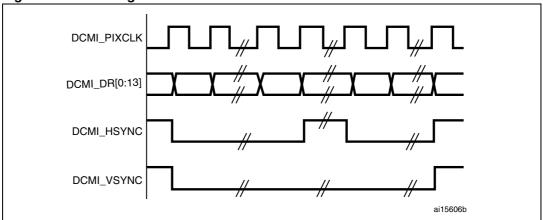
	Signal name	Signal description
8 bits 10 bits 12 bits 14 bits	D[07] D[09] D[011] D[013]	Data
PIXCLK		Pixel clock
HSYNC		Horizontal synchronization / Data valid
VSYNC		Vertical synchronization

The data are synchronous with PIXCLK and change on the rising/falling edge of the pixel clock depending on the polarity.

The HSYNC signal indicates the start/end of a line.

The VSYNC signal indicates the start/end of a frame

Figure 58. DCMI signal waveforms



The capture edge of DCMI\_PIXCLK is the falling edge, the active state of DCMI\_HSYNC and DCMI\_VSYNC is 1.

1. DCMI\_HSYNC and DCMI\_VSYNC can change states at the same time.

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#### 8-bit data

When EDM[1:0] in DCMI\_CR are programmed to "00" the interface captures 8 LSB's at its input (D[0:7]) and stores them as 8-bit data. The D[13:8] inputs are ignored. In this case, to capture a 32-bit word, the camera interface takes four pixel clock cycles.

The first captured data byte is placed in the LSB position in the 32-bit word and the 4<sup>th</sup> captured data byte is placed in the MSB position in the 32-bit word. *Table 48* gives an example of the positioning of captured data bytes in two 32-bit words.

Table 48. Positioning of captured data bytes in 32-bit words (8-bit width)

Byte address	31:24	23:16	15:8	7:0
0	D <sub>n+3</sub> [7:0]	D <sub>n+2</sub> [7:0]	D <sub>n+1</sub> [7:0]	D <sub>n</sub> [7:0]
4	D <sub>n+7</sub> [7:0]	D <sub>n+6</sub> [7:0]	D <sub>n+5</sub> [7:0]	D <sub>n+4</sub> [7:0]

## 10-bit data

When EDM[1:0] in DCMI\_CR are programmed to "01", the camera interface captures 10-bit data at its input D[0..9] and stores them as the 10 least significant bits of a 16-bit word. The remaining most significant bits in the DCMI\_DR register (bits 11 to 15) are cleared to zero. So, in this case, a 32-bit data word is made up every two pixel clock cycles.

The first captured data are placed in the LSB position in the 32-bit word and the 2<sup>nd</sup> captured data are placed in the MSB position in the 32-bit word as shown in *Table 49*.

Table 49. Positioning of captured data bytes in 32-bit words (10-bit width)

Byte address	31:26	25:16	15:10	9:0
0	0	D <sub>n+1</sub> [9:0]	0	D <sub>n</sub> [9:0]
4	0	D <sub>n+3</sub> [9:0]	0	D <sub>n+2</sub> [9:0]

## 12-bit data

When EDM[1:0] in DCMI\_CR are programmed to "10", the camera interface captures the 12-bit data at its input D[0..11] and stores them as the 12 least significant bits of a 16-bit word. The remaining most significant bits are cleared to zero. So, in this case a 32-bit data word is made up every two pixel clock cycles.

The first captured data are placed in the LSB position in the 32-bit word and the 2<sup>nd</sup> captured data are placed in the MSB position in the 32-bit word as shown in *Table 50*.

Table 50. Positioning of captured data bytes in 32-bit words (12-bit width)

Byte address	31:28	27:16	15:12	11:0
0	0	D <sub>n+1</sub> [11:0]	0	D <sub>n</sub> [11:0]
4	0	D <sub>n+3</sub> [11:0]	0	D <sub>n+2</sub> [11:0]

## 14-bit data

When EDM[1:0] in DCMI\_CR are programmed to "11", the camera interface captures the 14-bit data at its input D[0..13] and stores them as the 14 least significant bits of a 16-bit

word. The remaining most significant bits are cleared to zero. So, in this case a 32-bit data word is made up every two pixel clock cycles.

The first captured data are placed in the LSB position in the 32-bit word and the 2<sup>nd</sup> captured data are placed in the MSB position in the 32-bit word as shown in *Table 51*.

Table 51. Positioning of captured data bytes in 32-bit words (14-bit width)

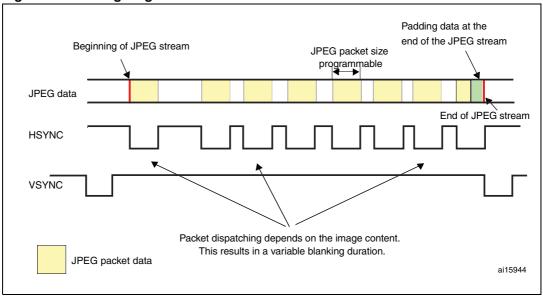
Byte address	31:30	29:16	29:16 15:14	
0	0	D <sub>n+1</sub> [13:0]	0	D <sub>n</sub> [13:0]
4	0	D <sub>n+3</sub> [13:0]	0	D <sub>n+2</sub> [13:0]

# 12.5.3 Synchronization

The digital camera interface supports embedded or hardware (HSYNC & VSYNC) synchronization. When embedded synchronization is used, it is up to the digital camera module to make sure that the 0x00 and 0xFF values are used ONLY for synchronization (not in data). Embedded synchronization codes are supported only for the 8-bit parallel data interface width (that is, in the DCMI\_CR register, the EDM[1:0] bits should be cleared to "00").

For compressed data, the DCMI supports only the hardware synchronization mode. In this case, VSYNC is used as a start/end of the image, and HSYNC is used as a Data Valid signal. *Figure 59* shows the corresponding timing diagram.

Figure 59. Timing diagram



## Hardware synchronization mode

In hardware synchronisation mode, the two synchronization signals (HSYNC/VSYNC) are used.

Depending on the camera module/mode, data may be transmitted during horizontal/vertical synchronisation periods. The HSYNC/VSYNC signals act like blanking signals since all the data received during HSYNC/VSYNC active periods are ignored.

In order to correctly transfer images into the DMA/RAM buffer, data transfer is synchronized with the VSYNC signal. When the hardware synchronisation mode is selected, and capture is enabled (CAPTURE bit set in DCMI\_CR), data transfer is synchronized with the deactivation of the VSYNC signal (next start of frame).

Transfer can then be continuous, with successive frames transferred by DMA to successive buffers or the same/circular buffer. To allow the DMA management of successive frames, a VSIF (Vertical synchronization interrupt flag) is activated at the end of each frame.

## **Embedded data synchronization mode**

In this synchronisation mode, the data flow is synchronised using 32-bit codes embedded in the data flow. These codes use the 0x00/0xFF values that are *not* used in data anymore. There are 4 types of codes, all with a 0xFF0000XY format. The embedded synchronization codes are supported only in 8-bit parallel data width capture (in the DCMI\_CR register, the EDM[1:0] bits should be programmed to "00"). For other data widths, this mode generates unpredictable results and must not be used.

Note:

Camera modules can have 8 such codes (in interleaved mode). For this reason, the interleaved mode is not supported by the camera interface (otherwise, every other half-frame would be discarded).

## Mode 2

Four embedded codes signal the following events

- Frame start (FS)
- Frame end (FE)
- Line start (LS)
- Line end (LE)

The XY values in the 0xFF0000XY format of the four codes are programmable (see Section 12.8.7: DCMI embedded synchronization code register (DCMI\_ESCR)).

A 0xFF value programmed as a "frame end" means that all the unused codes are interpreted as valid frame end codes.

In this mode, once the camera interface has been enabled, the frame capture starts after the first occurrence of the frame end (FE) code followed by a frame start (FS) code.

## Mode 1

An alternative coding is the camera mode 1. This mode is ITU656 compatible.

The codes signal another set of events:

- SAV (active line) line start
- EAV (active line) line end
- SAV (blanking) end of line during interframe blanking period
- EAV (blanking) end of line during interframe blanking period

This mode can be supported by programming the following codes:

- FS ≤ 0xFF
- FE ≤ 0xFF
- LS ≤ SAV (active)
- LE ≤ EAV (active)

An embedded unmask code is also implemented for frame/line start and frame/line end codes. Using it, it is possible to compare only the selected unmasked bits with the programmed code. You can therefore select a bit to compare in the embedded code and detect a frame/line start or frame/line end. This means that there can be different codes for the frame/line start and frame/line end with the unmasked bit position remaining the same.

# Example

FS = 0xA5

Unmask code for FS = 0x10

In this case the frame start code is embedded in the bit 4 of the frame start code.

# 12.5.4 Capture modes

This interface supports two types of capture: snapshot (single frame) and continuous grab.

## **Snapshot mode (single frame)**

In this mode, a single frame is captured (CM = '1' in the DCMI\_CR register). After the CAPTURE bit is set in DCMI\_CR, the interface waits for the detection of a start of frame before sampling the data. The camera interface is automatically disabled (CAPTURE bit cleared in DCMI\_CR) after receiving the first complete frame. An interrupt is generated (IT\_FRAME) if it is enabled.

In case of an overrun, the frame is lost and the CAPTURE bit is cleared.

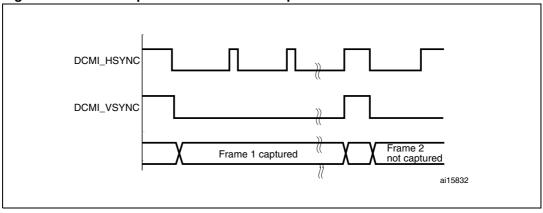


Figure 60. Frame capture waveforms in Snapshot mode

- 1. Here, the active state of DCMI\_HSYNC and DCMI\_VSYNC is 1.
- 2. DCMI\_HSYNC and DCMI\_VSYNC can change states at the same time.

# Continuous grab mode

In this mode (CM bit = '0' in DCMI CR), once the CAPTURE bit has been set in DCMI CR, the grabbing process starts on the next VSYNC or embedded frame start depending on the mode. The process continues until the CAPTURE bit is cleared in DCMI CR. Once the CAPTURE bit has been cleared, the grabbing process continues until the end of the current frame.

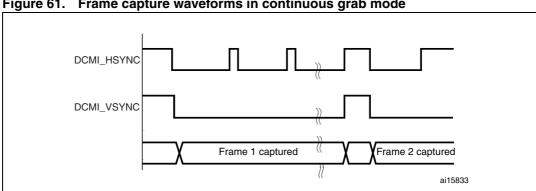


Figure 61. Frame capture waveforms in continuous grab mode

- 1. Here, the active state of DCMI\_HSYNC and DCMI\_VSYNC is 1.
- 2. DCMI\_HSYNC and DCMI\_VSYNC can change states at the same time.

In continuous grab mode, you can configure the FCRC bits in DCMI CR to grab all pictures, every second picture or one out of four pictures to decrease the frame capture rate.

Note:

In the hardware synchronization mode (ESS = '0' in DCMI CR), the IT VSYNC interrupt is generated (if enabled) even when CAPTURE = '0' in DCMI CR so, to reduce the frame capture rate even further, the IT\_VSYNC interrupt can be used to count the number of frames between 2 captures in conjunction with the Snapshot mode. This is not allowed by embedded data synchronization mode.

#### 12.5.5 **Crop feature**

With the crop feature, the camera interface can select a rectangular window from the received image. The start (upper left corner) coordinates and size (horizontal dimension in number of pixel clocks and vertical dimension in number of lines) are specified using two 32bit registers (DCMI\_CWSTRT and DCMI\_CWSIZE). The size of the window is specified in number of pixel clocks (horizontal dimension) and in number of lines (vertical dimension).

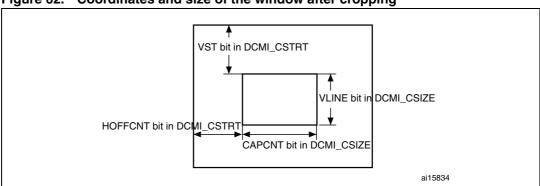


Figure 62. Coordinates and size of the window after cropping

These registers specify the coordinates of the starting point of the capture window as a line number (in the frame, starting from 0) and a number of pixel clocks (on the line, starting from 0), and the size of the window as a line number and a number of pixel clocks. The CAPCNT value can only be a multiple of 4 (two least significant bits are forced to 0) to allow the correct transfer of data through the DMA.

If the VSYNC signal goes active before the number of lines is specified in the DCMI\_CWSIZE register, then the capture stops and an IT\_FRAME interrupt is generated when enabled.

DCMI\_HSYNC

DCMI\_VSYNC

Data not captured in this phase

Data captured in this phase

ai15833

Figure 63. Data capture waveforms

- 1. Here, the active state of DCMI\_HSYNC and DCMI\_VSYNC is 1.
- 2. DCMI\_HSYNC and DCMI\_VSYNC can change states at the same time.

## 12.5.6 JPEG format

To allow JPEG image reception, it is necessary to set the JPEG bit in the DCMI\_CR register. JPEG images are not stored as lines and frames, so the VSYNC signal is used to start the capture while HSYNC serves as a data enable signal. The number of bytes in a line may not be a multiple of 4, you should therefore be careful when handling this case since a DMA request is generated each time a complete 32-bit word has been constructed from the captured data. When an end of frame is detected and the 32-bit word to be transferred has not been completely received, the remaining data are padded with '0s' and a DMA request is generated.

The crop feature and embedded synchronization codes cannot be used in the JPEG format.

## 12.5.7 FIFO

A four-word FIFO is implemented to manage data rate transfers on the AHB. The DCMI features a simple FIFO controller with a read pointer incremented each time the camera interface reads from the AHB, and a write pointer incremented each time the camera interface writes to the FIFO. There is no overrun protection to prevent the data from being overwritten if the AHB interface does not sustain the data transfer rate.

In case of overrun or errors in the synchronization signals, the FIFO is reset and the DCMI interface waits for a new start of frame.

# 12.6 Data format description

# 12.6.1 Data formats

Three types of data are supported:

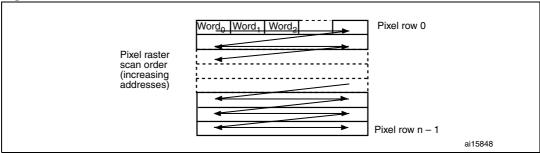
- 8-bit progressive video: either monochrome or raw Bayer format
- YCbCr 4:2:2 progressive video
- RGB565 progressive video. A pixel coded in 16 bits (5 bits for blue, 5 bits for red, 6 bits for green) takes two clock cycles to be transferred.

Compressed data: JPEG

For B&W, YCbCr or RGB data, the maximum input size is  $2048 \times 2048$  pixels. No limit in JPEG compressed mode.

For monochrome, RGB & YCbCr, the frame buffer is stored in raster mode. 32-bit words are used. Only the little endian format is supported.

Figure 64. Pixel raster scan order



# 12.6.2 Monochrome format

Characteristics:

- Raster format
- 8 bits per pixel

Table 52 shows how the data are stored.

Table 52. Data storage in monochrome progressive video format

Byte address	31:24	23:16	15:8	7:0
0	n + 3	n + 2	n + 1	n
4	n + 7	n + 6	n + 5	n + 4

# 12.6.3 RGB format

Characteristics:

- Raster format
- RGB
- Interleaved: one buffer: R, G & B interleaved: BRGBRGBRG, etc.
- Optimized for display output

The RGB planar format is compatible with standard OS frame buffer display formats. Only 16 BPP (bits per pixel): RGB565 (2 pixels per 32-bit word) is supported. The 24 BPP (palletized format) and grayscale formats are not supported. Pixels are stored in a raster scan order, that is from top to bottom for pixel rows, and from left to right within a pixel row. Pixel components are R (red), G (green) and B (blue). All components have the same spatial resolution (4:4:4 format). A frame is stored in a single part, with the components interleaved on a pixel basis.

Table 53 shows how the data are stored.

Table 53. Data storage in RGB progressive video format

Byte address	31:27	26:21	20:16	15:11	10:5	4:0
0	Red n + 1	Green n + 1	Blue n + 1	Red n	Green n	Blue n
4	Red n + 4	Green n + 3	Blue n + 3	Red n + 2	Green n + 2	Blue n + 2

## 12.6.4 YCbCr format

Characteristics:

- Raster format
- YCbCr 4:2:2
- Interleaved: one Buffer: Y, Cb & Cr interleaved: CbYCrYCbYCr, etc.

Pixel components are Y (luminance or "luma"), Cb and Cr (chrominance or "chroma" blue and red). Each component is encoded in 8 bits. Luma and chroma are stored together (interleaved) as shown in *Table 54*.

Table 54. Data storage in YCbCr progressive video format

Byte address	31:24	23:16	15:8	7:0
0	Y n + 1	Cr n	Υn	Cb n
4	Y n + 3	Cr n + 2	Y n + 2	Cb n + 2

# 12.7 DCMI interrupts

Five interrupts are generated. All interrupts are maskable by software. The global interrupt (IT\_DCMI) is the OR of all the individual interrupts. *Table 55* gives the list of all interrupts.

Table 55. DCMI interrupts

Interrupt name	Interrupt event		
IT_LINE	Indicates the end of line		
IT_FRAME	Indicates the end of frame capture		
IT_OVR	indicates the overrun of data reception		
IT_VSYNC	Indicates the synchronization frame		
IT_ERR	Indicates the detection of an error in the embedded synchronization frame detection		
IT_DCMI	Logic OR of the previous interrupts		

# 12.8 DCMI register description

All DCMI registers have to be accessed as 32-bit words, otherwise a bus error occurs.

# 12.8.1 DCMI control register 1 (DCMI\_CR)

Address offset: 0x00

Reset value: 0x0000 0x0000

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 PCKPOL CROP JPEG Reserved ENABL ESS EDM **FCRC** S Reserved rw rw rw rw rw rw rw rw rw rw

Bit 31:15 Reserved, must be kept at reset value.

Bit 14 ENABLE: DCMI enable

0: DCMI disabled 1: DCMI enabled

Note: The DCMI configuration registers should be programmed correctly before enabling this Bit

Bit 13: 12 Reserved, must be kept at reset value.

## 11:10 EDM[1:0]: Extended data mode

00: Interface captures 8-bit data on every pixel clock01: Interface captures 10-bit data on every pixel clock10: Interface captures 12-bit data on every pixel clock

11: Interface captures 14-bit data on every pixel clock

## 9:8 FCRC[1:0]: Frame capture rate control

These bits define the frequency of frame capture. They are meaningful only in Continuous grab mode. They are ignored in snapshot mode.

00: All frames are captured

01: Every alternate frame captured (50% bandwidth reduction)

10: One frame in 4 frames captured (75% bandwidth reduction)

11: reserved

## Bit 7 VSPOL: Vertical synchronization polarity

This bit indicates the level on the VSYNC pin when the data are not valid on the parallel interface.

0: VSYNC active low

1: VSYNC active high

## Bit 6 HSPOL: Horizontal synchronization polarity

This bit indicates the level on the HSYNC pin when the data are not valid on the parallel interface.

0: HSYNC active low

1: HSYNC active high

## Bit 5 PCKPOL: Pixel clock polarity

This bit configures the capture edge of the pixel clock

0: Falling edge active.

1: Rising edge active.

## Bit 4 ESS: Embedded synchronization select

- 0: Hardware synchronization data capture (frame/line start/stop) is synchronized with the HSYNC/VSYNC signals.
- 1: Embedded synchronization data capture is synchronized with synchronization codes embedded in the data flow.

Note: Valid only for 8-bit parallel data. HSPOL/VSPOL are ignored when the ESS bit is set.

This bit is disabled in JPEG mode.

## Bit 3 JPEG: JPEG format

- 0: Uncompressed video format
- 1: This bit is used for JPEG data transfers. The HSYNC signal is used as data enable. The crop and embedded synchronization features (ESS bit) cannot be used in this mode.

#### Bits 2 CROP: Crop feature

- 0: The full image is captured. In this case the total number of bytes in an image frame should be a multiple of 4
- 1: Only the data inside the window specified by the crop register will be captured. If the size of the crop window exceeds the picture size, then only the picture size is captured.

### Bit 1 CM: Capture mode

- 0: Continuous grab mode The received data are transferred into the destination memory through the DMA. The buffer location and mode (linear or circular buffer) is controlled through the system DMA.
- 1: Snapshot mode (single frame) Once activated, the interface waits for the start of frame and then transfers a single frame through the DMA. At the end of the frame, the CAPTURE bit is automatically reset.

## Bit 0 CAPTURE: Capture enable

- 0: Capture disabled.
- 1: Capture enabled.

The camera interface waits for the first start of frame, then a DMA request is generated to transfer the received data into the destination memory.

In snapshot mode, the CAPTURE bit is automatically cleared at the end of the 1st frame received.

In continuous grab mode, if the software clears this bit while a capture is ongoing, the bit will be effectively cleared after the frame end.

Note: The DMA controller and all DCMI configuration registers should be programmed correctly before enabling this bit.

# 12.8.2 DCMI status register (DCMI\_SR)

Address offset: 0x04

Reset value: 0x0000 0x0000

 $31 \ \ 30 \ \ 29 \ \ 28 \ \ 27 \ \ 26 \ \ 25 \ \ 24 \ \ 23 \ \ 22 \ \ 21 \ \ 20 \ \ 19 \ \ 18 \ \ 17 \ \ 16 \ \ 15 \ \ 14 \ \ 13 \ \ 12 \ \ 11 \ \ 10 \ \ 9 \ \ 8 \ \ 7 \ \ 6 \ \ 5 \ \ 4 \ \ 3 \ \ 2 \ \ 1 \ \ 0$ 

Reserved	FNE	VSYNC	HSYNC
	r	r	r

Bit 31:3 Reserved, must be kept at reset value.

Bit 2 FNE: FIFO not empty

This bit gives the status of the FIFO

1: FIFO contains valid data

0: FIFO empty

#### Bit 1 VSYNC

This bit gives the state of the VSYNC pin with the correct programmed polarity. When embedded synchronization codes are used, the meaning of this bit is the following:

0: active frame

1: synchronization between frames

In case of embedded synchronization, this bit is meaningful only if the CAPTURE bit in DCMI\_CR is set.

## Bit 0 HSYNC

This bit gives the state of the HSYNC pin with the correct programmed polarity. When embedded synchronization codes are used, the meaning of this bit is the following:

0: active line

1: synchronization between lines

In case of embedded synchronization, this bit is meaningful only if the CAPTURE bit in DCMI\_CR is set.

## 12.8.3 DCMI raw interrupt status register (DCMI\_RIS)

Address offset: 0x08

Reset value: 0x0000 0x0000

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	LINE_RIS	VSYNC_RIS	ERR_RIS	OVR_RIS	FRAME_RIS
	r	r	r	r	r

DCMI\_RIS gives the raw interrupt status and is accessible in read only. When read, this register returns the status of the corresponding interrupt before masking with the DCMI\_IER register value.

Bit 31:5 Reserved, must be kept at reset value.

### Bit 4 LINE\_RIS: Line raw interrupt status

This bit gets set when the HSYNC signal changes from the inactive state to the active state. It goes high even if the line is not valid.

In the case of embedded synchronization, this bit is set only if the CAPTURE bit in DCMI\_CR is set.

It is cleared by writing a '1' to the LINE\_ISC bit in DCMI\_ICR.

### Bit 3 VSYNC\_RIS: VSYNC raw interrupt status

This bit is set when the VSYNC signal changes from the inactive state to the active state.

In the case of embedded synchronization, this bit is set only if the CAPTURE bit is set in  $\mathsf{DCMI\_CR}$ .

It is cleared by writing a '1' to the VSYNC\_ISC bit in DCMI\_ICR.

### Bit 2 ERR\_RIS: Synchronization error raw interrupt status

0: No synchronization error detected

1: Embedded synchronization characters are not received in the correct order.

This bit is valid only in the embedded synchronization mode. It is cleared by writing a '1' to the ERR\_ISC bit in DCMI\_ICR.

Note: This bit is available only in embedded synchronization mode.

### Bit 1 **OVR\_RIS:** Overrun raw interrupt status

0: No data buffer overrun occurred

1: A data buffer overrun occurred and the data FIFO is corrupted.

This bit is cleared by writing a '1' to the OVR\_ISC bit in DCMI\_ICR.

### Bit 0 FRAME\_RIS: Capture complete raw interrupt status

0: No new capture

1: A frame has been captured.

This bit is set when a frame or window has been captured.

In case of a cropped window, this bit is set at the end of line of the last line in the crop. It is set even if the captured frame is empty (e.g. window cropped outside the frame).

This bit is cleared by writing a '1' to the FRAME\_ISC bit in DCMI\_ICR.

## 12.8.4 DCMI interrupt enable register (DCMI\_IER)

Address offset: 0x0C

Reset value: 0x0000 0x0000

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved

Reserved

Reserved

The DCMI\_IER register is used to enable interrupts. When one of the DCMI\_IER bits is set, the corresponding interrupt is enabled. This register is accessible in both read and write.

Bit 31:5 Reserved, must be kept at reset value.

### Bit 4 LINE\_IE: Line interrupt enable

- 0: No interrupt generation when the line is received
- 1: An Interrupt is generated when a line has been completely received

### Bit 3 VSYNC\_IE: VSYNC interrupt enable

- 0: No interrupt generation
- 1: An interrupt is generated on each VSYNC transition from the inactive to the active state

The active state of the VSYNC signal is defined by the VSPOL bit.

### Bit 2 **ERR\_IE:** Synchronization error interrupt enable

- 0: No interrupt generation
- 1: An interrupt is generated if the embedded synchronization codes are not received in the correct order.

Note: This bit is available only in embedded synchronization mode.

### Bit 1 **OVR\_IE:** Overrun interrupt enable

- 0: No interrupt generation
- 1: An interrupt is generated if the DMA was not able to transfer the last data before new data (32-bit) are received.

### Bit 0 FRAME\_IE: Capture complete interrupt enable

- 0: No interrupt generation
- 1: An interrupt is generated at the end of each received frame/crop window (in crop mode).

## 12.8.5 DCMI masked interrupt status register (DCMI\_MIS)

This DCMI\_MIS register is a read-only register. When read, it returns the current masked status value (depending on the value in DCMI\_IER) of the corresponding interrupt. A bit in this register is set if the corresponding enable bit in DCMI\_IER is set and the corresponding bit in DCMI\_RIS is set.

Address offset: 0x10 Reset value: 0x0000

3	1 3	30 2	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
													Re	eserv	ed /													CINE_MIS	VSYNC_MIS	ERR_MIS	OVR_MIS	FRAME_MIS
																												r	r	r	r	r

Bit 31:5 Reserved, must be kept at reset value.

### Bit 4 LINE\_MIS: Line masked interrupt status

This bit gives the status of the masked line interrupt

0: No interrupt generation when the line is received

1: An Interrupt is generated when a line has been completely received and the LINE\_IE bit is set in DCMI\_IER.

### Bit 3 VSYNC\_MIS: VSYNC masked interrupt status

This bit gives the status of the masked VSYNC interrupt

0: No interrupt is generated on VSYNC transitions

1: An interrupt is generated on each VSYNC transition from the inactive to the active state and the VSYNC\_IE bit is set in DCMI\_IER.

The active state of the VSYNC signal is defined by the VSPOL bit.

## Bit 2 **ERR\_MIS:** Synchronization error masked interrupt status

This bit gives the status of the masked synchronization error interrupt

0: No interrupt is generated on a synchronization error

1: An interrupt is generated if the embedded synchronization codes are not received in the correct order and the ERR\_IE bit in DCMI\_IER is set.

Note: This bit is available only in embedded synchronization mode.

### Bit 1 **OVR\_MIS:** Overrun masked interrupt status

This bit gives the status of the masked overflow interrupt

0: No interrupt is generated on overrun

1: An interrupt is generated if the DMA was not able to transfer the last data before new data (32-bit) are received and the OVR\_IE bit is set in DCMI\_IER.

### Bit 0 FRAME\_MIS: Capture complete masked interrupt status

This bit gives the status of the masked capture complete interrupt

0: No interrupt is generated after a complete capture

1: An interrupt is generated at the end of each received frame/crop window (in crop mode) and the FRAME\_IE bit is set in DCMI\_IER.

## 12.8.6 DCMI interrupt clear register (DCMI\_ICR)

Address offset: 0x14

Reset value: 0x0000 0x0000

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved

Reserved

W W W W W W W

The DCMI\_ICR register is write-only. Writing a '1' into a bit of this register clears the corresponding bit in the DCMI\_RIS and DCMI\_MIS registers. Writing a '0' has no effect.

Bit 15:5 Reserved, must be kept at reset value.

Bit 4 LINE\_ISC: line interrupt status clear
Writing a '1' into this bit clears LINE\_RIS in the DCMI\_RIS register

Bit 3 VSYNC\_ISC: Vertical synch interrupt status clear
Writing a '1' into this bit clears the VSYNC\_RIS bit in DCMI\_RIS

Bit 2 **ERR\_ISC:** Synchronization error interrupt status clear

Writing a '1' into this bit clears the ERR\_RIS bit in DCMI\_RIS

Note: This bit is available only in embedded synchronization mode.

Bit 1 **OVR\_ISC:** Overrun interrupt status clear Writing a '1' into this bit clears the OVR\_RIS bit in DCMI\_RIS

Bits 0 FRAME\_ISC: Capture complete interrupt status clear
Writing a '1' into this bit clears the FRAME\_RIS bit in DCMI\_RIS

## 12.8.7 DCMI embedded synchronization code register (DCMI\_ESCR)

Address offset: 0x18

Reset value: 0x0000 0x0000

 $31 \quad 30 \quad 29 \quad 28 \quad 27 \quad 26 \quad 25 \quad 24 \quad 23 \quad 22 \quad 21 \quad 20 \quad 19 \quad 18 \quad 17 \quad 16 \quad 15 \quad 14 \quad 13 \quad 12 \quad 11 \quad 10 \quad 9 \quad 8 \quad 7 \quad 6 \quad 5 \quad 4 \quad 3 \quad 2 \quad 1 \quad 0$ 

			FE	EC							LE	EC							LS	SC							FS	SC			
rv	/ rw	rw																													

### Bit 31:24 FEC: Frame end delimiter code

This byte specifies the code of the frame end delimiter. The code consists of 4 bytes in the form of 0xFF, 0x00, 0x00, FEC.

If FEC is programmed to 0xFF, all the unused codes (0xFF0000XY) are interpreted as frame end delimiters.

### Bit 23:16 **LEC:** Line end delimiter code

This byte specifies the code of the line end delimiter. The code consists of 4 bytes in the form of 0xFF, 0x00, 0x00, LEC.

### Bit 15:8 LSC: Line start delimiter code

This byte specifies the code of the line start delimiter. The code consists of 4 bytes in the form of 0xFF, 0x00, 0x00, LSC.

## Bit 7:0 FSC: Frame start delimiter code

This byte specifies the code of the frame start delimiter. The code consists of 4 bytes in the form of 0xFF, 0x00, 0x00, FSC.

If FSC is programmed to 0xFF, no frame start delimiter is detected. But, the 1<sup>st</sup> occurrence of LSC after an FEC code will be interpreted as a start of frame delimiter.

## 12.8.8 DCMI embedded synchronization unmask register (DCMI\_ESUR)

Address offset: 0x1C

Reset value: 0x0000 0x0000

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

			FE	U							LE	U							LS	SU							FS	3U			-
rw																															

### Bit 31:24 FEU: Frame end delimiter unmask

This byte specifies the mask to be applied to the code of the frame end delimiter.

- 0: The corresponding bit in the FEC byte in DCMI\_ESCR is masked while comparing the frame end delimiter with the received data.
- 1: The corresponding bit in the FEC byte in DCMI\_ESCR is compared while comparing the frame end delimiter with the received data

### Bit 23:16 LEU: Line end delimiter unmask

This byte specifies the mask to be applied to the code of the line end delimiter.

- 0: The corresponding bit in the LEC byte in DCMI\_ESCR is masked while comparing the line end delimiter with the received data
- 1: The corresponding bit in the LEC byte in DCMI\_ESCR is compared while comparing the line end delimiter with the received data

### Bit 15:8 LSU: Line start delimiter unmask

This byte specifies the mask to be applied to the code of the line start delimiter.

- 0: The corresponding bit in the LSC byte in DCMI\_ESCR is masked while comparing the line start delimiter with the received data
- 1: The corresponding bit in the LSC byte in DCMI\_ESCR is compared while comparing the line start delimiter with the received data

### Bit 7:0 FSU: Frame start delimiter unmask

This byte specifies the mask to be applied to the code of the frame start delimiter.

- 0: The corresponding bit in the FSC byte in DCMI\_ESCR is masked while comparing the frame start delimiter with the received data
- 1: The corresponding bit in the FSC byte in DCMI\_ESCR is compared while comparing the frame start delimiter with the received data

## 12.8.9 DCMI crop window start (DCMI\_CWSTRT)

Address offset: 0x20

Reset value: 0x0000 0x0000

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Bosoniod						VS	ST[12	2:0						Reserv						НО	FFCI	NT[1	3:0]					
Reserved	rw	rw	rw	rw	rw	rw	rw	ed	rw	rw	rw	rw	rw	rw	rw	rw												

Bits 31:29 Reserved, must be kept at reset value.

Bit 28:16 VST[12:0]: Vertical start line count

The image capture starts with this line number. Previous line data are ignored.

 $0x0000 \Rightarrow line 1$   $0x0001 \Rightarrow line 2$  $0x0002 \Rightarrow line 3$ 

Bits 15:14 Reserved, must be kept at reset value.

Bit 13:0 HOFFCNT[13:0]: Horizontal offset count

This value gives the number of pixel clocks to count before starting a capture.

## 12.8.10 DCMI crop window size (DCMI\_CWSIZE)

Address offset: 0x24

Reset value: 0x0000 0x0000

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

ved						٧	/LINE	E13:0	0]						ved						CA	PCN	IT[13	3:0]					
Reser	rw	rw	rw	rw	rw	rw	rw	rw	Reser	rw	rw	rw	rw	rw	rw	rw	rw												

Bits 31:30 Reserved, must be kept at reset value.

Bit 29:16 VLINE[13:0]: Vertical line count

This value gives the number of lines to be captured from the starting point.

0x0000 => 1 line 0x0001 => 2 lines 0x0002 => 3 lines

Bits 15:14 Reserved, must be kept at reset value.

### Bit 13:0 CAPCNT[13:0]: Capture count

This value gives the number of pixel clocks to be captured from the starting point on the same line. It value should corresponds to word-aligned data for different widths of parallel interfaces.

0x0000 => 1 pixel 0x0001 => 2 pixels 0x0002 => 3 pixels

## 12.8.11 DCMI data register (DCMI\_DR)

Address offset: 0x28

Reset value: 0x0000 0x0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	/	6	5	4	3	2	1	U
			Ву	te3							Ву	te2							Ву	te1							Ву	te0			
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:24 Data byte 3

Bit 23:16 Data byte 2

Bits 15:8 Data byte 1

Bit 7:0 Data byte 0

The digital camera Interface packages all the received data in 32-bit format before requesting a DMA transfer. A 4-word deep FIFO is available to leave enough time for DMA transfers and avoid DMA overrun conditions.

## 12.8.12 DCMI register map

Table 56 summarizes the DCMI registers.

Table 56. DCMI register map and reset values

Offset	Register	31	30	29	28	27	9	25	24	23	22	2	i 6	δ 5	2 2	17	16	15	4	13	12	7	우	6	œ	7	9	2	4	က	7	-	0
0x00	DCMI_CR							F	Res	serv	/ed								ENABLE	Reserved		ED	М	FCF	RC	VSPOL	HSPOL	PCKPOL	ESS	JPEG	CROP	CM	CAPTURE
	Reset value																		0	<u> </u>		0	0	0	0	0	0	0	0	0	0	0	0
0x04	DCMI_SR														Re	esei	ved														FNE	VSYNC	HSYNC
	Reset value																														0	0	0
0x08	DCMI_RIS													R	eser	ved	l												LINE_RIS	VSYNC_RIS	ERR_RIS	OVR_RIS	FRAME_RIS
	Reset value																												0	0	0	0	0
0x0C	DCMI_IER													R	eser	ved	l												LINE_IE	VSYNC_IE	ERR_IE	OVR_IE	FRAME_IE
	Reset value																												0	0	0	0	
0x10	DCMI_MIS													R	eser	ved	ı												LINE_MIS	VSYNC_MIS	ERR_MIS	OVR_MIS	FRAME_MIS
	Reset value																												0	0	0	0	0
0x14	DCMI_ICR													R	eser	ved	1												LINE_ISC	VSYNC_ISC	ERR_ISC	OVR_ISC	FRAME_ISC
	Reset value																												0	0	0	0	0

Table 56. DCMI register map and reset values (continued)

Offset	Register	31	30	6	8	72	9	25	7	ည	2	7.	20	19	18	17	16	2	14	13	2	11	10	6	æ	7	9	5	4	က	2	_	0
		(·)	(.)	"	(1	(1	(1	"	"	"	14	(1	"	_	_	_	_	-	_	_	_	_	_			Ė			_				Ĺ
0x18	DCMI_ESCR				FI	ΞC							LE	EC							LS	SC							FS	SC			
	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x1C	DCMI_ESUR				FI	ΞU							LE	ΞU							LS	SU							FS	SU			
	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x20	DCMI_CWSTRT	Re	serv	ved						vs	T[12	2:0						Reserved						ı	HOF	FC	NT[	13:0	0]				
	Reset value				0	0	0	0	0	0	0	0	0	0	0	0	0	Re		0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x24	DCMI_CWSIZE	Reserved							VI	INE	Ξ13:	0]						Reserved							CA	PCN	NT[1	3:0	]				
	Reset value	Be		0	0	0	0	0	0	0	0	0	0	0	0	0	0	Re		0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x28	DCMI_DR				Ву	te3							Ву	te2							Ву	te1							Ву	te0			
	Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

Refer to *Table 2 on page 50* for the register boundary addresses.

# 13 Advanced-control timers (TIM1&TIM8)

This section applies to the whole STM32F4xx family, unless otherwise specified.

## 13.1 TIM1&TIM8 introduction

The advanced-control timers (TIM1&TIM8) consist of a 16-bit auto-reload counter driven by a programmable prescaler.

It may be used for a variety of purposes, including measuring the pulse lengths of input signals (input capture) or generating output waveforms (output compare, PWM, complementary PWM with dead-time insertion).

Pulse lengths and waveform periods can be modulated from a few microseconds to several milliseconds using the timer prescaler and the RCC clock controller prescalers.

The advanced-control (TIM1&TIM8) and general-purpose (TIMx) timers are completely independent, and do not share any resources. They can be synchronized together as described in *Section 13.3.20*.

## 13.2 TIM1&TIM8 main features

TIM1&TIM8 timer features include:

- 16-bit up, down, up/down auto-reload counter.
- 16-bit programmable prescaler allowing dividing (also "on the fly") the counter clock frequency either by any factor between 1 and 65536.
- Up to 4 independent channels for:
  - Input Capture
  - Output Compare
  - PWM generation (Edge and Center-aligned Mode)
  - One-pulse mode output
- Complementary outputs with programmable dead-time
- Synchronization circuit to control the timer with external signals and to interconnect several timers together.
- Repetition counter to update the timer registers only after a given number of cycles of the counter.
- Break input to put the timer's output signals in reset state or in a known state.

- Interrupt/DMA generation on the following events:
  - Update: counter overflow/underflow, counter initialization (by software or internal/external trigger)
  - Trigger event (counter start, stop, initialization or count by internal/external trigger)
  - Input capture
  - Output compare
  - Break input
- Supports incremental (quadrature) encoder and hall-sensor circuitry for positioning purposes
- Trigger input for external clock or cycle-by-cycle current management

CK\_TIM18 from RCC TIMx ETR ITR0 ITR1 ITR2 eset, Enable, Up/Down, Count ITR3 TI1FP1 TI2FP2 Stop, Clear or Up/Dov CC1 DTG[7:0] registers TIMx\_CH1 TIMx\_CH2 TIMx\_CH3 TIMx CH4 TIMx\_BKIN

Figure 65. Advanced-control timer block diagram

## 13.3 TIM1&TIM8 functional description

### 13.3.1 Time-base unit

The main block of the programmable advanced-control timer is a 16-bit counter with its related auto-reload register. The counter can count up, down or both up and down. The counter clock can be divided by a prescaler.

The counter, the auto-reload register and the prescaler register can be written or read by software. This is true even when the counter is running.

The time-base unit includes:

- Counter register (TIMx\_CNT)
- Prescaler register (TIMx\_PSC)
- Auto-reload register (TIMx\_ARR)
- Repetition counter register (TIMx\_RCR)

The auto-reload register is preloaded. Writing to or reading from the auto-reload register accesses the preload register. The content of the preload register are transferred into the shadow register permanently or at each update event (UEV), depending on the auto-reload preload enable bit (ARPE) in TIMx\_CR1 register. The update event is sent when the counter reaches the overflow (or underflow when downcounting) and if the UDIS bit equals 0 in the TIMx\_CR1 register. It can also be generated by software. The generation of the update event is described in detailed for each configuration.

The counter is clocked by the prescaler output CK\_CNT, which is enabled only when the counter enable bit (CEN) in TIMx\_CR1 register is set (refer also to the slave mode controller description to get more details on counter enabling).

Note that the counter starts counting 1 clock cycle after setting the CEN bit in the TIMx\_CR1 register.

### Prescaler description

The prescaler can divide the counter clock frequency by any factor between 1 and 65536. It is based on a 16-bit counter controlled through a 16-bit register (in the TIMx\_PSC register). It can be changed on the fly as this control register is buffered. The new prescaler ratio is taken into account at the next update event.

*Figure 66* and *Figure 67* give some examples of the counter behavior when the prescaler ratio is changed on the fly:

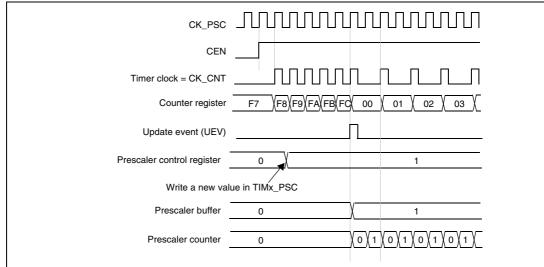
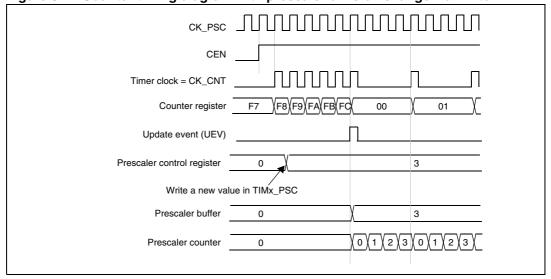


Figure 66. Counter timing diagram with prescaler division change from 1 to 2





## 13.3.2 Counter modes

### **Upcounting mode**

In upcounting mode, the counter counts from 0 to the auto-reload value (content of the TIMx\_ARR register), then restarts from 0 and generates a counter overflow event.

If the repetition counter is used, the update event (UEV) is generated after upcounting is repeated for the number of times programmed in the repetition counter register plus one (TIMx\_RCR+1). Else the update event is generated at each counter overflow.

Setting the UG bit in the TIMx\_EGR register (by software or by using the slave mode controller) also generates an update event.

The UEV event can be disabled by software by setting the UDIS bit in the TIMx\_CR1 register. This is to avoid updating the shadow registers while writing new values in the

preload registers. Then no update event occurs until the UDIS bit has been written to 0. However, the counter restarts from 0, as well as the counter of the prescaler (but the prescale rate does not change). In addition, if the URS bit (update request selection) in TIMx\_CR1 register is set, setting the UG bit generates an update event UEV but without setting the UIF flag (thus no interrupt or DMA request is sent). This is to avoid generating both update and capture interrupts when clearing the counter on the capture event.

When an update event occurs, all the registers are updated and the update flag (UIF bit in TIMx\_SR register) is set (depending on the URS bit):

- The repetition counter is reloaded with the content of TIMx\_RCR register,
- The auto-reload shadow register is updated with the preload value (TIMx\_ARR),
- The buffer of the prescaler is reloaded with the preload value (content of the TIMx\_PSC register).

The following figures show some examples of the counter behavior for different clock frequencies when TIMx\_ARR=0x36.



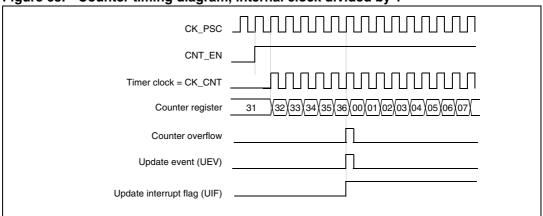
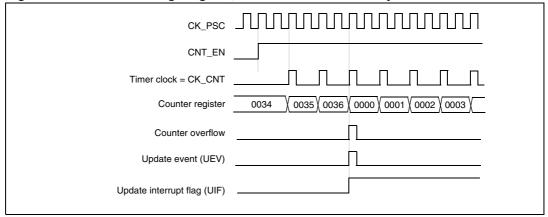
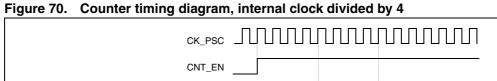


Figure 69. Counter timing diagram, internal clock divided by 2





Timer clock = CK\_CNT 0035 Counter register Counter overflow Update event (UEV) Update interrupt flag (UIF)

Figure 71. Counter timing diagram, internal clock divided by N

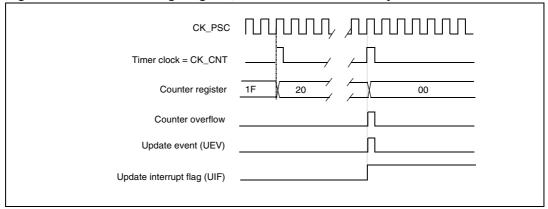
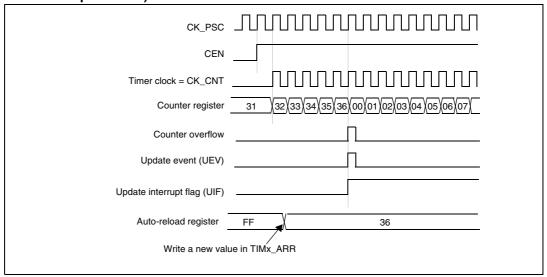


Figure 72. Counter timing diagram, update event when ARPE=0 (TIMx\_ARR not preloaded)



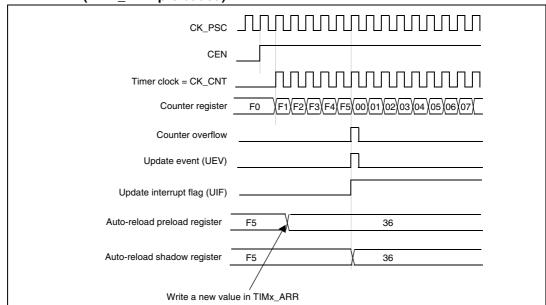


Figure 73. Counter timing diagram, update event when ARPE=1 (TIMx ARR preloaded)

### **Downcounting mode**

In downcounting mode, the counter counts from the auto-reload value (content of the TIMx\_ARR register) down to 0, then restarts from the auto-reload value and generates a counter underflow event.

If the repetition counter is used, the update event (UEV) is generated after downcounting is repeated for the number of times programmed in the repetition counter register plus one (TIMx\_RCR+1). Else the update event is generated at each counter underflow.

Setting the UG bit in the TIMx\_EGR register (by software or by using the slave mode controller) also generates an update event.

The UEV update event can be disabled by software by setting the UDIS bit in TIMx\_CR1 register. This is to avoid updating the shadow registers while writing new values in the preload registers. Then no update event occurs until UDIS bit has been written to 0. However, the counter restarts from the current auto-reload value, whereas the counter of the prescaler restarts from 0 (but the prescale rate doesn't change).

In addition, if the URS bit (update request selection) in TIMx\_CR1 register is set, setting the UG bit generates an update event UEV but without setting the UIF flag (thus no interrupt or DMA request is sent). This is to avoid generating both update and capture interrupts when clearing the counter on the capture event.

When an update event occurs, all the registers are updated and the update flag (UIF bit in TIMx\_SR register) is set (depending on the URS bit):

- The repetition counter is reloaded with the content of TIMx\_RCR register
- The buffer of the prescaler is reloaded with the preload value (content of the TIMx\_PSC register)
- The auto-reload active register is updated with the preload value (content of the TIMx\_ARR register). Note that the auto-reload is updated before the counter is reloaded, so that the next period is the expected one

The following figures show some examples of the counter behavior for different clock frequencies when TIMx\_ARR=0x36.

Figure 74. Counter timing diagram, internal clock divided by 1

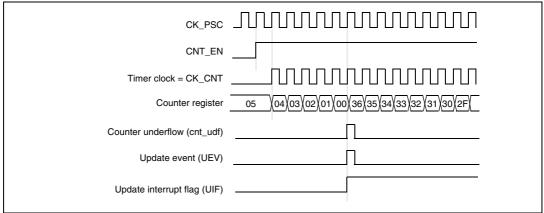


Figure 75. Counter timing diagram, internal clock divided by 2

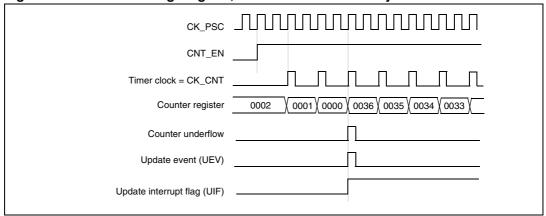
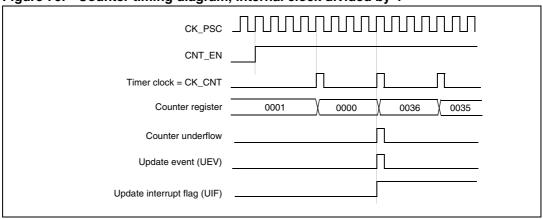


Figure 76. Counter timing diagram, internal clock divided by 4



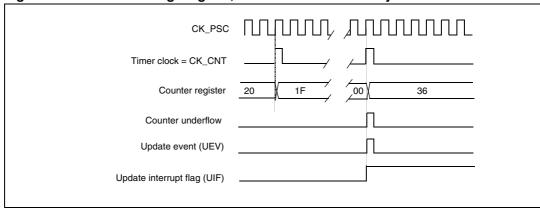
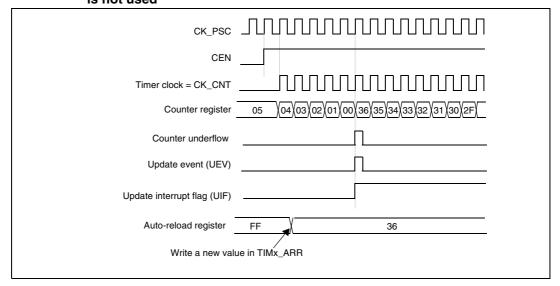


Figure 77. Counter timing diagram, internal clock divided by N

Figure 78. Counter timing diagram, update event when repetition counter is not used



### Center-aligned mode (up/down counting)

In center-aligned mode, the counter counts from 0 to the auto-reload value (content of the  $TIMx\_ARR$  register) – 1, generates a counter overflow event, then counts from the auto-reload value down to 1 and generates a counter underflow event. Then it restarts counting from 0.

Center-aligned mode is active when the CMS bits in TIMx\_CR1 register are not equal to '00'. The Output compare interrupt flag of channels configured in output is set when: the counter counts down (Center aligned mode 1, CMS = "01"), the counter counts up (Center aligned mode 2, CMS = "10") the counter counts up and down (Center aligned mode 3, CMS = "11").

In this mode, the DIR direction bit in the TIMx\_CR1 register cannot be written. It is updated by hardware and gives the current direction of the counter.

The update event can be generated at each counter overflow and at each counter underflow or by setting the UG bit in the TIMx\_EGR register (by software or by using the slave mode controller) also generates an update event. In this case, the counter restarts counting from 0, as well as the counter of the prescaler.

The UEV update event can be disabled by software by setting the UDIS bit in the TIMx\_CR1 register. This is to avoid updating the shadow registers while writing new values in the preload registers. Then no update event occurs until UDIS bit has been written to 0. However, the counter continues counting up and down, based on the current auto-reload value.

In addition, if the URS bit (update request selection) in TIMx\_CR1 register is set, setting the UG bit generates an UEV update event but without setting the UIF flag (thus no interrupt or DMA request is sent). This is to avoid generating both update and capture interrupts when clearing the counter on the capture event.

When an update event occurs, all the registers are updated and the update flag (UIF bit in TIMx\_SR register) is set (depending on the URS bit):

- The repetition counter is reloaded with the content of TIMx\_RCR register
- The buffer of the prescaler is reloaded with the preload value (content of the TIMx\_PSC register)
- The auto-reload active register is updated with the preload value (content of the TIMx\_ARR register). Note that if the update source is a counter overflow, the autoreload is updated before the counter is reloaded, so that the next period is the expected one (the counter is loaded with the new value).

The following figures show some examples of the counter behavior for different clock frequencies.

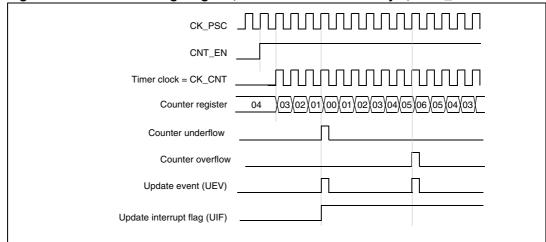


Figure 79. Counter timing diagram, internal clock divided by 1, TIMx\_ARR = 0x6

1. Here, center-aligned mode 1 is used (for more details refer to Section 13.4: TIM1&TIM8 registers on page 339).

Figure 80. Counter timing diagram, internal clock divided by 2

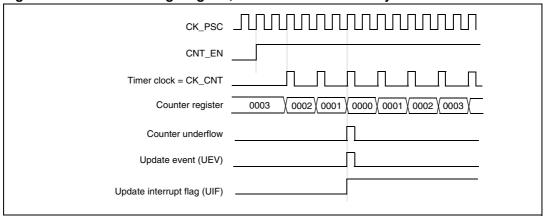
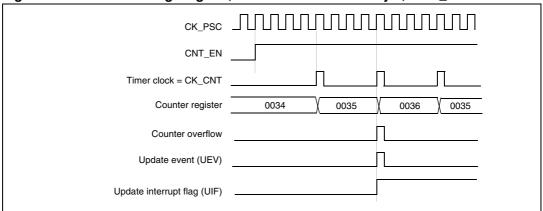
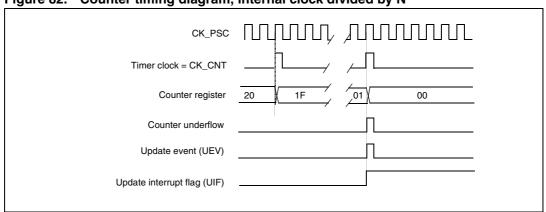


Figure 81. Counter timing diagram, internal clock divided by 4, TIMx\_ARR=0x36



1. Center-aligned mode 2 or 3 is used with an UIF on overflow.

Figure 82. Counter timing diagram, internal clock divided by N



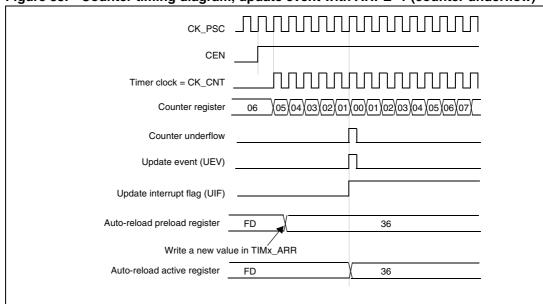
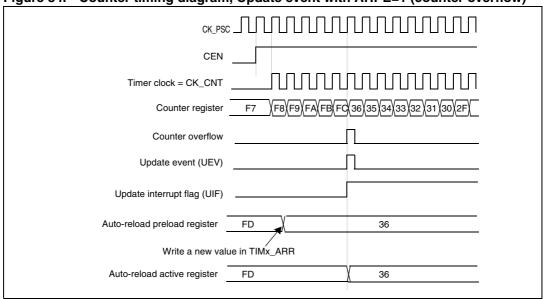


Figure 83. Counter timing diagram, update event with ARPE=1 (counter underflow)





## 13.3.3 Repetition counter

Section 13.3.1: Time-base unit describes how the update event (UEV) is generated with respect to the counter overflows/underflows. It is actually generated only when the repetition counter has reached zero. This can be useful when generating PWM signals.

This means that data are transferred from the preload registers to the shadow registers (TIMx\_ARR auto-reload register, TIMx\_PSC prescaler register, but also TIMx\_CCRx capture/compare registers in compare mode) every N+1 counter overflows or underflows, where N is the value in the TIMx\_RCR repetition counter register.

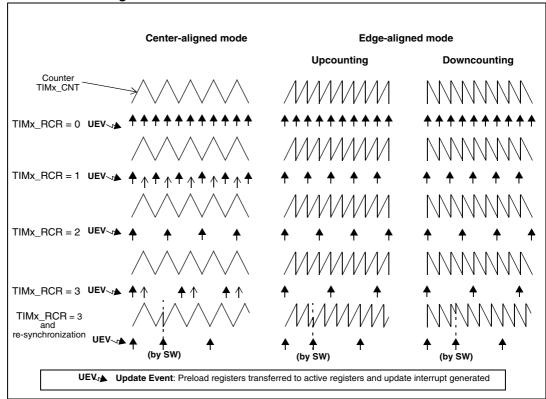
The repetition counter is decremented:

- At each counter overflow in upcounting mode,
- At each counter underflow in downcounting mode,
- At each counter overflow and at each counter underflow in center-aligned mode.
   Although this limits the maximum number of repetition to 128 PWM cycles, it makes it possible to update the duty cycle twice per PWM period. When refreshing compare registers only once per PWM period in center-aligned mode, maximum resolution is 2xT<sub>ck</sub>, due to the symmetry of the pattern.

The repetition counter is an auto-reload type; the repetition rate is maintained as defined by the TIMx\_RCR register value (refer to *Figure 85*). When the update event is generated by software (by setting the UG bit in TIMx\_EGR register) or by hardware through the slave mode controller, it occurs immediately whatever the value of the repetition counter is and the repetition counter is reloaded with the content of the TIMx\_RCR register.

In center-aligned mode, for odd values of RCR, the update event occurs either on the overflow or on the underflow depending on when the RCR register was written and when the counter was started. If the RCR was written before starting the counter, the UEV occurs on the overflow. If the RCR was written after starting the counter, the UEV occurs on the underflow. For example for RCR = 3, the UEV is generated on each 4th overflow or underflow event depending on when RCR was written.

Figure 85. Update rate examples depending on mode and TIMx\_RCR register settings



#### 13.3.4 Clock selection

The counter clock can be provided by the following clock sources:

- Internal clock (CK INT)
- External clock mode1: external input pin
- External clock mode2: external trigger input ETR
- Internal trigger inputs (ITRx): using one timer as prescaler for another timer, for example, you can configure Timer 1 to act as a prescaler for Timer 2. Refer to Using one timer as prescaler for another for more details.

## Internal clock source (CK\_INT)

If the slave mode controller is disabled (SMS=000), then the CEN, DIR (in the TIMx\_CR1 register) and UG bits (in the TIMx\_EGR register) are actual control bits and can be changed only by software (except UG which remains cleared automatically). As soon as the CEN bit is written to 1, the prescaler is clocked by the internal clock CK\_INT.

Figure 86 shows the behavior of the control circuit and the upcounter in normal mode, without prescaler.

Internal clock CEN=CNT\_EN UG Counter clock = CK\_CNT = CK\_PSC Counter register \(\)32\(\)33\(\)34\(\)35\(\)36\(\)00\(\)01\(\)02\(\)03\(\)04\(\)05\(\)06\(\)07\(\)

Figure 86. Control circuit in normal mode, internal clock divided by 1

### External clock source mode 1

This mode is selected when SMS=111 in the TIMx\_SMCR register. The counter can count at each rising or falling edge on a selected input.

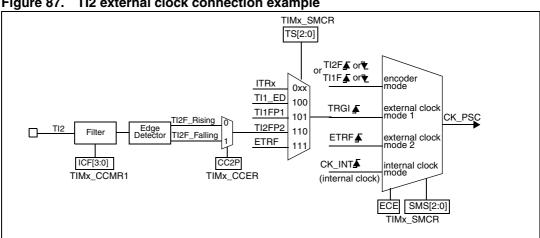


Figure 87. TI2 external clock connection example

For example, to configure the upcounter to count in response to a rising edge on the TI2 input, use the following procedure:

- 1. Configure channel 2 to detect rising edges on the TI2 input by writing CC2S = '01' in the TIMx\_CCMR1 register.
- 2. Configure the input filter duration by writing the IC2F[3:0] bits in the TIMx\_CCMR1 register (if no filter is needed, keep IC2F=0000).
- 3. Select rising edge polarity by writing CC2P=0 and CC2NP=0 in the TIMx\_CCER register.
- 4. Configure the timer in external clock mode 1 by writing SMS=111 in the TIMx\_SMCR register.
- 5. Select TI2 as the trigger input source by writing TS=110 in the TIMx\_SMCR register.
- 6. Enable the counter by writing CEN=1 in the TIMx\_CR1 register.

Note: The capture prescaler is not used for triggering, so you don't need to configure it.

When a rising edge occurs on TI2, the counter counts once and the TIF flag is set.

The delay between the rising edge on TI2 and the actual clock of the counter is due to the resynchronization circuit on TI2 input.

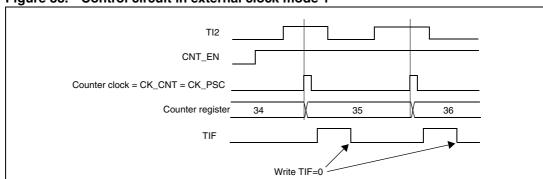


Figure 88. Control circuit in external clock mode 1

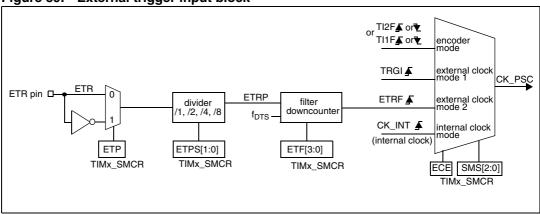
### External clock source mode 2

This mode is selected by writing ECE=1 in the TIMx SMCR register.

The counter can count at each rising or falling edge on the external trigger input ETR.

Figure 89 gives an overview of the external trigger input block.

Figure 89. External trigger input block



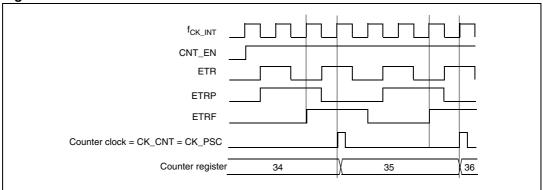
For example, to configure the upcounter to count each 2 rising edges on ETR, use the following procedure:

- As no filter is needed in this example, write ETF[3:0]=0000 in the TIMx\_SMCR register.
- 2. Set the prescaler by writing ETPS[1:0]=01 in the TIMx\_SMCR register
- 3. Select rising edge detection on the ETR pin by writing ETP=0 in the TIMx\_SMCR register
- 4. Enable external clock mode 2 by writing ECE=1 in the TIMx\_SMCR register.
- 5. Enable the counter by writing CEN=1 in the TIMx\_CR1 register.

The counter counts once each 2 ETR rising edges.

The delay between the rising edge on ETR and the actual clock of the counter is due to the resynchronization circuit on the ETRP signal.

Figure 90. Control circuit in external clock mode 2



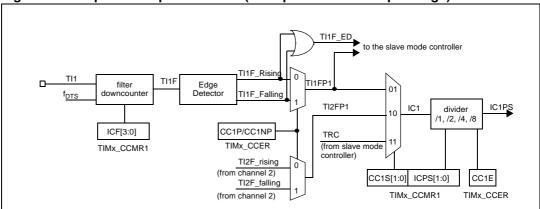
## 13.3.5 Capture/compare channels

Each Capture/Compare channel is built around a capture/compare register (including a shadow register), a input stage for capture (with digital filter, multiplexing and prescaler) and an output stage (with comparator and output control).

Figure 91 to Figure 94 give an overview of one Capture/Compare channel.

The input stage samples the corresponding TIx input to generate a filtered signal TIxF. Then, an edge detector with polarity selection generates a signal (TIxFPx) which can be used as trigger input by the slave mode controller or as the capture command. It is prescaled before the capture register (ICxPS).

Figure 91. Capture/compare channel (example: channel 1 input stage)



The output stage generates an intermediate waveform which is then used for reference: OCxRef (active high). The polarity acts at the end of the chain.

APB Bus MCU-peripheral interface 16-bit) 8 8 high Š write CCR1H Ξ read CCR1H S write\_in\_progress read\_in\_progress write CCR1L Capture/compare preload register read CCR1 output CC1S[1] compare\_transfer capture\_transfer CC1S[0] input CC1S[1] OC1PE mode OC1PE Capture/compare shadow register CC1S[0] UEV TIM1\_CCMR1 from time comparator IC1PS capture base unit) CC1E CNT>CCR1 Counter CNT=CCR1 CC1G TIM1\_EGR

Figure 92. Capture/compare channel 1 main circuit

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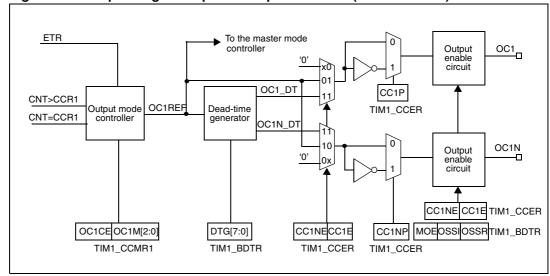
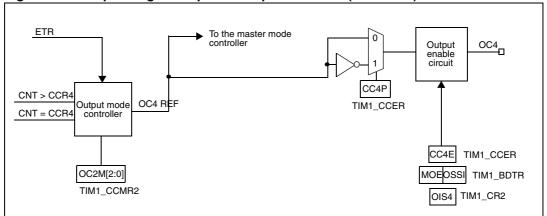


Figure 93. Output stage of capture/compare channel (channel 1 to 3)

Figure 94. Output stage of capture/compare channel (channel 4)



The capture/compare block is made of one preload register and one shadow register. Write and read always access the preload register.

In capture mode, captures are actually done in the shadow register, which is copied into the preload register.

In compare mode, the content of the preload register is copied into the shadow register which is compared to the counter.

## 13.3.6 Input capture mode

In Input capture mode, the Capture/Compare Registers (TIMx\_CCRx) are used to latch the value of the counter after a transition detected by the corresponding ICx signal. When a capture occurs, the corresponding CCXIF flag (TIMx\_SR register) is set and an interrupt or a DMA request can be sent if they are enabled. If a capture occurs while the CCxIF flag was already high, then the over-capture flag CCxOF (TIMx\_SR register) is set. CCxIF can be cleared by software by writing it to '0' or by reading the captured data stored in the TIMx\_CCRx register. CCxOF is cleared when you write it to '0'.

The following example shows how to capture the counter value in TIMx\_CCR1 when TI1 input rises. To do this, use the following procedure:

- Select the active input: TIMx\_CCR1 must be linked to the TI1 input, so write the CC1S bits to 01 in the TIMx\_CCMR1 register. As soon as CC1S becomes different from 00, the channel is configured in input and the TIMx\_CCR1 register becomes read-only.
- Program the input filter duration you need with respect to the signal you connect to the timer (by programming ICxF bits in the TIMx\_CCMRx register if the input is a TIx input). Let's imagine that, when toggling, the input signal is not stable during at must 5 internal clock cycles. We must program a filter duration longer than these 5 clock cycles. We can validate a transition on TI1 when 8 consecutive samples with the new level have been detected (sampled at f<sub>DTS</sub> frequency). Then write IC1F bits to 0011 in the TIMx\_CCMR1 register.
- Select the edge of the active transition on the TI1 channel by writing CC1P and CC1NP bits to 0 in the TIMx\_CCER register (rising edge in this case).
- Program the input prescaler. In our example, we wish the capture to be performed at each valid transition, so the prescaler is disabled (write IC1PS bits to '00' in the TIMx\_CCMR1 register).
- Enable capture from the counter into the capture register by setting the CC1E bit in the TIMx\_CCER register.
- If needed, enable the related interrupt request by setting the CC1IE bit in the TIMx\_DIER register, and/or the DMA request by setting the CC1DE bit in the TIMx\_DIER register.

When an input capture occurs:

- The TIMx\_CCR1 register gets the value of the counter on the active transition.
- CC1IF flag is set (interrupt flag). CC1OF is also set if at least two consecutive captures occurred whereas the flag was not cleared.
- An interrupt is generated depending on the CC1IE bit.
- A DMA request is generated depending on the CC1DE bit.

In order to handle the overcapture, it is recommended to read the data before the overcapture flag. This is to avoid missing an overcapture which could happen after reading the flag and before reading the data.

Note:

IC interrupt and/or DMA requests can be generated by software by setting the corresponding CCxG bit in the TIMx\_EGR register.

## 13.3.7 PWM input mode

This mode is a particular case of input capture mode. The procedure is the same except:

- Two ICx signals are mapped on the same TIx input.
- These 2 ICx signals are active on edges with opposite polarity.
- One of the two TIxFP signals is selected as trigger input and the slave mode controller is configured in reset mode.

For example, you can measure the period (in TIMx\_CCR1 register) and the duty cycle (in TIMx\_CCR2 register) of the PWM applied on TI1 using the following procedure (depending on CK\_INT frequency and prescaler value):

- Select the active input for TIMx\_CCR1: write the CC1S bits to 01 in the TIMx\_CCMR1 register (TI1 selected).
- Select the active polarity for TI1FP1 (used both for capture in TIMx\_CCR1 and counter clear): write the CC1P and CC1NP bits to '0' (active on rising edge).
- Select the active input for TIMx\_CCR2: write the CC2S bits to 10 in the TIMx\_CCMR1 register (TI1 selected).
- Select the active polarity for TI1FP2 (used for capture in TIMx\_CCR2): write the CC2P and CC2NP bits to '1' (active on falling edge).
- Select the valid trigger input: write the TS bits to 101 in the TIMx\_SMCR register (TI1FP1 selected).
- Configure the slave mode controller in reset mode: write the SMS bits to 100 in the TIMx\_SMCR register.
- Enable the captures: write the CC1E and CC2E bits to '1' in the TIMx\_CCER register.

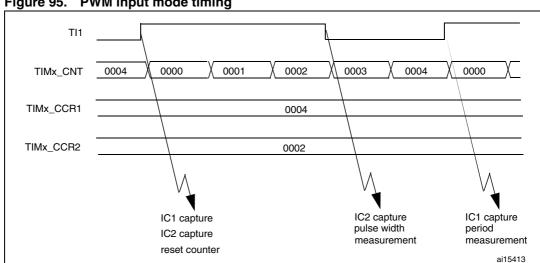


Figure 95. PWM input mode timing

## 13.3.8 Forced output mode

In output mode (CCxS bits = 00 in the TIMx\_CCMRx register), each output compare signal (OCxREF and then OCx/OCxN) can be forced to active or inactive level directly by software, independently of any comparison between the output compare register and the counter.

To force an output compare signal (OCXREF/OCx) to its active level, you just need to write 101 in the OCxM bits in the corresponding TIMx\_CCMRx register. Thus OCXREF is forced high (OCxREF is always active high) and OCx get opposite value to CCxP polarity bit.

For example: CCxP=0 (OCx active high) => OCx is forced to high level.

The OCxREF signal can be forced low by writing the OCxM bits to 100 in the TIMx\_CCMRx register.

Anyway, the comparison between the TIMx\_CCRx shadow register and the counter is still performed and allows the flag to be set. Interrupt and DMA requests can be sent accordingly. This is described in the output compare mode section below.

## 13.3.9 Output compare mode

This function is used to control an output waveform or indicating when a period of time has elapsed.

When a match is found between the capture/compare register and the counter, the output compare function:

- Assigns the corresponding output pin to a programmable value defined by the output compare mode (OCxM bits in the TIMx\_CCMRx register) and the output polarity (CCxP bit in the TIMx\_CCER register). The output pin can keep its level (OCXM=000), be set active (OCxM=001), be set inactive (OCxM=010) or can toggle (OCxM=011) on match.
- Sets a flag in the interrupt status register (CCxIF bit in the TIMx\_SR register).
- Generates an interrupt if the corresponding interrupt mask is set (CCXIE bit in the TIMx\_DIER register).
- Sends a DMA request if the corresponding enable bit is set (CCxDE bit in the TIMx\_DIER register, CCDS bit in the TIMx\_CR2 register for the DMA request selection).

The TIMx\_CCRx registers can be programmed with or without preload registers using the OCxPE bit in the TIMx\_CCMRx register.

In output compare mode, the update event UEV has no effect on OCxREF and OCx output. The timing resolution is one count of the counter. Output compare mode can also be used to output a single pulse (in One Pulse mode).

### Procedure:

- 1. Select the counter clock (internal, external, prescaler).
- 2. Write the desired data in the TIMx ARR and TIMx CCRx registers.
- 3. Set the CCxIE bit if an interrupt request is to be generated.
- 4. Select the output mode. For example:
  - Write OCxM = 011 to toggle OCx output pin when CNT matches CCRx
  - Write OCxPE = 0 to disable preload register
  - Write CCxP = 0 to select active high polarity
  - Write CCxE = 1 to enable the output
- 5. Enable the counter by setting the CEN bit in the TIMx\_CR1 register.

The TIMx\_CCRx register can be updated at any time by software to control the output waveform, provided that the preload register is not enabled (OCxPE='0', else TIMx\_CCRx shadow register is updated only at the next update event UEV). An example is given in *Figure 96*.

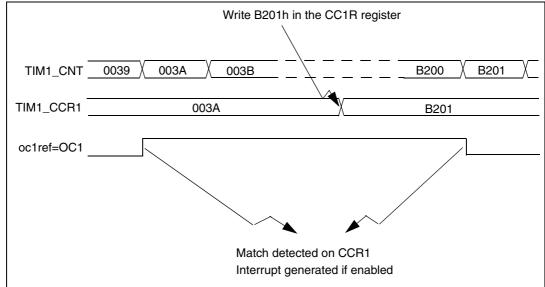


Figure 96. Output compare mode, toggle on OC1.

### 13.3.10 PWM mode

Pulse Width Modulation mode allows you to generate a signal with a frequency determined by the value of the TIMx\_ARR register and a duty cycle determined by the value of the TIMx\_CCRx register.

The PWM mode can be selected independently on each channel (one PWM per OCx output) by writing '110' (PWM mode 1) or '111' (PWM mode 2) in the OCxM bits in the TIMx\_CCMRx register. You must enable the corresponding preload register by setting the OCxPE bit in the TIMx\_CCMRx register, and eventually the auto-reload preload register (in upcounting or center-aligned modes) by setting the ARPE bit in the TIMx\_CR1 register.

As the preload registers are transferred to the shadow registers only when an update event occurs, before starting the counter, you have to initialize all the registers by setting the UG bit in the TIMx EGR register.

OCx polarity is software programmable using the CCxP bit in the TIMx\_CCER register. It can be programmed as active high or active low. OCx output is enabled by a combination of the CCxE, CCxNE, MOE, OSSI and OSSR bits (TIMx\_CCER and TIMx\_BDTR registers). Refer to the TIMx\_CCER register description for more details.

In PWM mode (1 or 2), TIMx\_CNT and TIMx\_CCRx are always compared to determine whether TIMx\_CCRx  $\leq$  TIMx\_CNT or TIMx\_CNT  $\leq$  TIMx\_CCRx (depending on the direction of the counter).

The timer is able to generate PWM in edge-aligned mode or center-aligned mode depending on the CMS bits in the TIMx\_CR1 register.

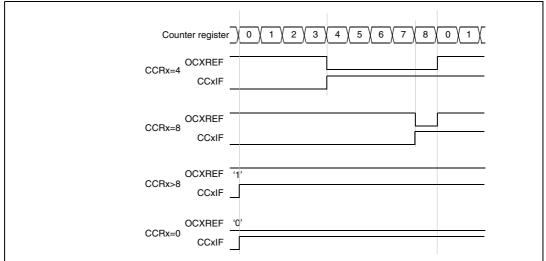
## PWM edge-aligned mode

Upcounting configuration

Upcounting is active when the DIR bit in the TIMx\_CR1 register is low. Refer to Section: Upcounting mode on page 302.

In the following example, we consider PWM mode 1. The reference PWM signal OCxREF is high as long as TIMx\_CNT < TIMx\_CCRx else it becomes low. If the compare value in TIMx\_CCRx is greater than the auto-reload value (in TIMx\_ARR) then OCxREF is held at '1'. If the compare value is 0 then OCxRef is held at '0'. *Figure 97* shows some edge-aligned PWM waveforms in an example where TIMx\_ARR=8.





Downcounting configuration

Downcounting is active when DIR bit in TIMx\_CR1 register is high. Refer to *Section : Downcounting mode on page 305* 

In PWM mode 1, the reference signal OCxRef is low as long as TIMx\_CNT > TIMx\_CCRx else it becomes high. If the compare value in TIMx\_CCRx is greater than the auto-reload value in TIMx\_ARR, then OCxREF is held at '1'. 0% PWM is not possible in this mode.

### PWM center-aligned mode

Center-aligned mode is active when the CMS bits in TIMx\_CR1 register are different from '00' (all the remaining configurations having the same effect on the OCxRef/OCx signals). The compare flag is set when the counter counts up, when it counts down or both when it counts up and down depending on the CMS bits configuration. The direction bit (DIR) in the TIMx\_CR1 register is updated by hardware and must not be changed by software. Refer to Section: Center-aligned mode (up/down counting) on page 307.

*Figure 98* shows some center-aligned PWM waveforms in an example where:

- TIMx ARR=8,
- PWM mode is the PWM mode 1,
- The flag is set when the counter counts down corresponding to the center-aligned mode 1 selected for CMS=01 in TIMx\_CR1 register.

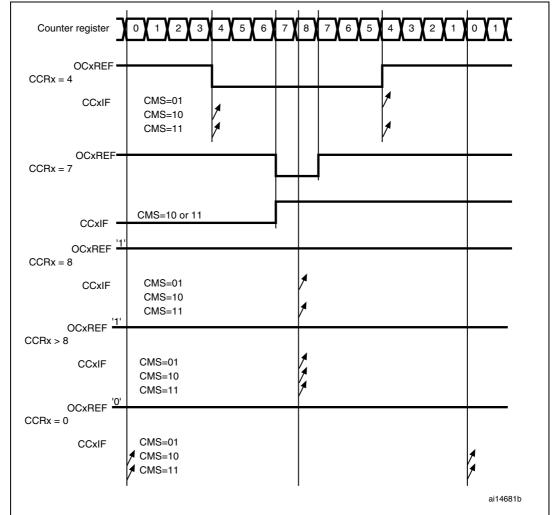


Figure 98. Center-aligned PWM waveforms (ARR=8)

### Hints on using center-aligned mode:

- When starting in center-aligned mode, the current up-down configuration is used. It
  means that the counter counts up or down depending on the value written in the DIR bit
  in the TIMx\_CR1 register. Moreover, the DIR and CMS bits must not be changed at the
  same time by the software.
- Writing to the counter while running in center-aligned mode is not recommended as it can lead to unexpected results. In particular:
  - The direction is not updated if you write a value in the counter that is greater than the auto-reload value (TIMx\_CNT>TIMx\_ARR). For example, if the counter was counting up, it continues to count up.
  - The direction is updated if you write 0 or write the TIMx\_ARR value in the counter but no Update Event UEV is generated.
- The safest way to use center-aligned mode is to generate an update by software (setting the UG bit in the TIMx\_EGR register) just before starting the counter and not to write the counter while it is running.

## 13.3.11 Complementary outputs and dead-time insertion

The advanced-control timers (TIM1&TIM8) can output two complementary signals and manage the switching-off and the switching-on instants of the outputs.

This time is generally known as dead-time and you have to adjust it depending on the devices you have connected to the outputs and their characteristics (intrinsic delays of level-shifters, delays due to power switches...)

You can select the polarity of the outputs (main output OCx or complementary OCxN) independently for each output. This is done by writing to the CCxP and CCxNP bits in the TIMx\_CCER register.

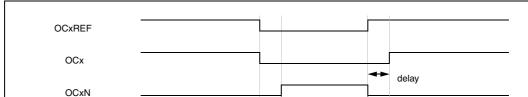
The complementary signals OCx and OCxN are activated by a combination of several control bits: the CCxE and CCxNE bits in the TIMx\_CCER register and the MOE, OISx, OISxN, OSSI and OSSR bits in the TIMx\_BDTR and TIMx\_CR2 registers. Refer to *Table 59:* Output control bits for complementary OCx and OCxN channels with break feature on page 357 for more details. In particular, the dead-time is activated when switching to the IDLE state (MOE falling down to 0).

Dead-time insertion is enabled by setting both CCxE and CCxNE bits, and the MOE bit if the break circuit is present. DTG[7:0] bits of the TIMx\_BDTR register are used to control the dead-time generation for all channels. From a reference waveform OCxREF, it generates 2 outputs OCx and OCxN. If OCx and OCxN are active high:

- The OCx output signal is the same as the reference signal except for the rising edge, which is delayed relative to the reference rising edge.
- The OCxN output signal is the opposite of the reference signal except for the rising edge, which is delayed relative to the reference falling edge.

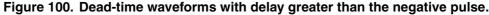
If the delay is greater than the width of the active output (OCx or OCxN) then the corresponding pulse is not generated.

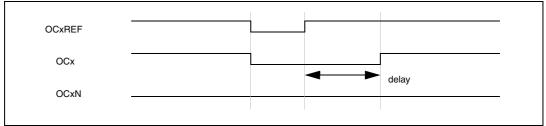
The following figures show the relationships between the output signals of the dead-time generator and the reference signal OCxREF. (we suppose CCxP=0, CCxNP=0, MOE=1, CCxE=1 and CCxNE=1 in these examples)



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Figure 99. Complementary output with dead-time insertion.





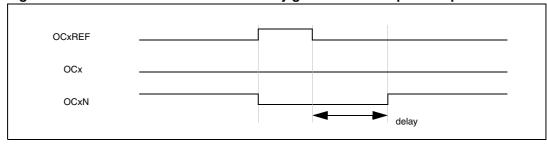


Figure 101. Dead-time waveforms with delay greater than the positive pulse.

The dead-time delay is the same for each of the channels and is programmable with the DTG bits in the TIMx\_BDTR register. Refer to *Section 13.4.18: TIM1&TIM8 break and dead-time register (TIMx\_BDTR) on page 361* for delay calculation.

### Re-directing OCxREF to OCx or OCxN

In output mode (forced, output compare or PWM), OCxREF can be re-directed to the OCx output or to OCxN output by configuring the CCxE and CCxNE bits in the TIMx\_CCER register.

This allows you to send a specific waveform (such as PWM or static active level) on one output while the complementary remains at its inactive level. Other alternative possibilities are to have both outputs at inactive level or both outputs active and complementary with dead-time.

Note:

When only OCxN is enabled (CCxE=0, CCxNE=1), it is not complemented and becomes active as soon as OCxREF is high. For example, if CCxNP=0 then OCxN=OCxRef. On the other hand, when both OCx and OCxN are enabled (CCxE=CCxNE=1) OCx becomes active when OCxREF is high whereas OCxN is complemented and becomes active when OCxREF is low.

## 13.3.12 Using the break function

When using the break function, the output enable signals and inactive levels are modified according to additional control bits (MOE, OSSI and OSSR bits in the TIMx\_BDTR register, OISx and OISxN bits in the TIMx\_CR2 register). In any case, the OCx and OCxN outputs cannot be set both to active level at a given time. Refer to *Table 59: Output control bits for complementary OCx and OCxN channels with break feature on page 357* for more details.

The break source can be either the break input pin or a clock failure event, generated by the Clock Security System (CSS), from the Reset Clock Controller. For further information on the Clock Security System, refer to Section 5.2.7: Clock security system (CSS).

When exiting from reset, the break circuit is disabled and the MOE bit is low. You can enable the break function by setting the BKE bit in the TIMx\_BDTR register. The break input polarity can be selected by configuring the BKP bit in the same register. BKE and BKP can be modified at the same time. When the BKE and BKP bits are written, a delay of 1 APB clock cycle is applied before the writing is effective. Consequently, it is necessary to wait 1 APB clock period to correctly read back the bit after the write operation.

Because MOE falling edge can be asynchronous, a resynchronization circuit has been inserted between the actual signal (acting on the outputs) and the synchronous control bit (accessed in the TIMx\_BDTR register). It results in some delays between the asynchronous and the synchronous signals. In particular, if you write MOE to 1 whereas it was low, you

must insert a delay (dummy instruction) before reading it correctly. This is because you write the asynchronous signal and read the synchronous signal.

When a break occurs (selected level on the break input):

- The MOE bit is cleared asynchronously, putting the outputs in inactive state, idle state
  or in reset state (selected by the OSSI bit). This feature functions even if the MCU
  oscillator is off.
- Each output channel is driven with the level programmed in the OISx bit in the TIMx\_CR2 register as soon as MOE=0. If OSSI=0 then the timer releases the enable output else the enable output remains high.
- When complementary outputs are used:
  - The outputs are first put in reset state inactive state (depending on the polarity).
     This is done asynchronously so that it works even if no clock is provided to the timer.
  - If the timer clock is still present, then the dead-time generator is reactivated in order to drive the outputs with the level programmed in the OISx and OISxN bits after a dead-time. Even in this case, OCx and OCxN cannot be driven to their active level together. Note that because of the resynchronization on MOE, the dead-time duration is a bit longer than usual (around 2 ck\_tim clock cycles).
  - If OSSI=0 then the timer releases the enable outputs else the enable outputs remain or become high as soon as one of the CCxE or CCxNE bits is high.
- The break status flag (BIF bit in the TIMx\_SR register) is set. An interrupt can be generated if the BIE bit in the TIMx\_DIER register is set. A DMA request can be sent if the BDE bit in the TIMx\_DIER register is set.
- If the AOE bit in the TIMx\_BDTR register is set, the MOE bit is automatically set again at the next update event UEV. This can be used to perform a regulation, for instance. Else, MOE remains low until you write it to '1' again. In this case, it can be used for security and you can connect the break input to an alarm from power drivers, thermal sensors or any security components.

Note:

The break inputs is acting on level. Thus, the MOE cannot be set while the break input is active (neither automatically nor by software). In the meantime, the status flag BIF cannot be cleared.

The break can be generated by the BRK input which has a programmable polarity and an enable bit BKE in the TIMx\_BDTR Register.

There are two solutions to generate a break:

- By using the BRK input which has a programmable polarity and an enable bit BKE in the TIMx\_BDTR register
- By software through the BG bit of the TIMx\_EGR register.

In addition to the break input and the output management, a write protection has been implemented inside the break circuit to safeguard the application. It allows you to freeze the configuration of several parameters (dead-time duration, OCx/OCxN polarities and state when disabled, OCxM configurations, break enable and polarity). You can choose from 3 levels of protection selected by the LOCK bits in the TIMx\_BDTR register. Refer to Section 13.4.18: TIM1&TIM8 break and dead-time register (TIMx\_BDTR) on page 361. The LOCK bits can be written only once after an MCU reset.

Figure 102 shows an example of behavior of the outputs in response to a break.



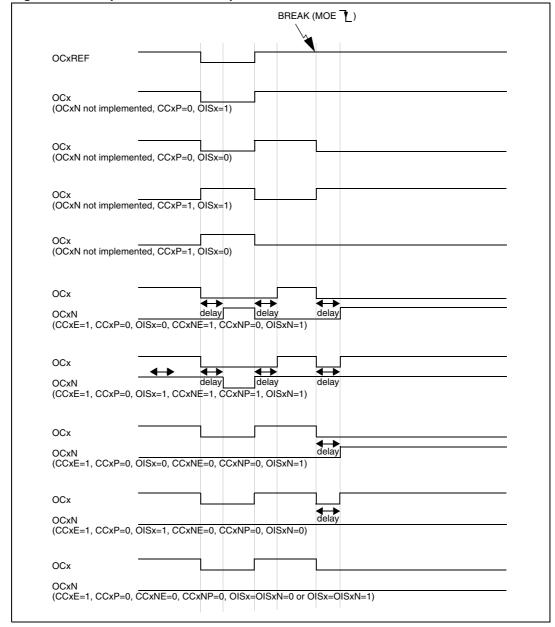


Figure 102. Output behavior in response to a break.

#### 13.3.13 Clearing the OCxREF signal on an external event

The OCxREF signal for a given channel can be driven Low by applying a High level to the ETRF input (OCxCE enable bit of the corresponding TIMx CCMRx register set to '1'). The OCxREF signal remains Low until the next update event, UEV, occurs.

This function can only be used in output compare and PWM modes, and does not work in forced mode.

For example, the ETR signal can be connected to the output of a comparator to be used for current handling. In this case, the ETR must be configured as follow:

- The External Trigger Prescaler should be kept off: bits ETPS[1:0] of the TIMx SMCR register set to '00'.
- 2. The external clock mode 2 must be disabled: bit ECE of the TIMx\_SMCR register set to '0'.
- 3. The External Trigger Polarity (ETP) and the External Trigger Filter (ETF) can be configured according to the user needs.

Figure 103 shows the behavior of the OCxREF signal when the ETRF Input becomes High, for both values of the enable bit OCxCE. In this example, the timer TIMx is programmed in PWM mode.

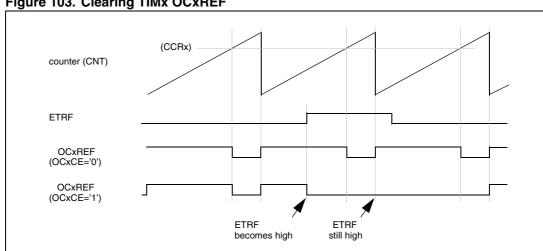


Figure 103. Clearing TIMx OCxREF

Note:

In case of a PWM with a 100% duty cycle (if CCRx>ARR), then OCxREF is enabled again at the next counter overflow.

# 13.3.14 6-step PWM generation

When complementary outputs are used on a channel, preload bits are available on the OCxM, CCxE and CCxNE bits. The preload bits are transferred to the shadow bits at the COM commutation event. Thus you can program in advance the configuration for the next step and change the configuration of all the channels at the same time. COM can be generated by software by setting the COM bit in the TIMx\_EGR register or by hardware (on TRGI rising edge).

A flag is set when the COM event occurs (COMIF bit in the TIMx\_SR register), which can generate an interrupt (if the COMIE bit is set in the TIMx\_DIER register) or a DMA request (if the COMDE bit is set in the TIMx\_DIER register).

*Figure 104* describes the behavior of the OCx and OCxN outputs when a COM event occurs, in 3 different examples of programmed configurations.

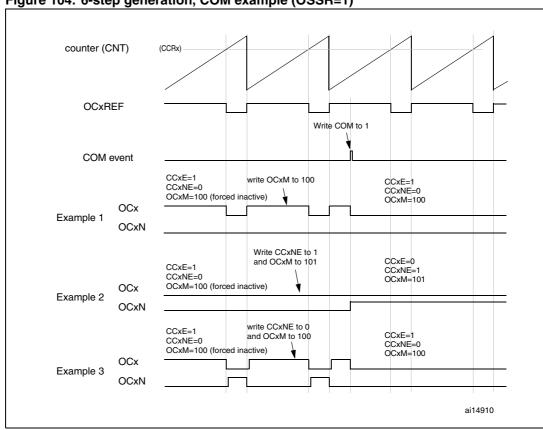


Figure 104. 6-step generation, COM example (OSSR=1)

# 13.3.15 One-pulse mode

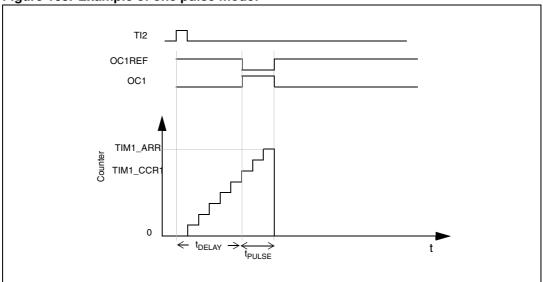
One-pulse mode (OPM) is a particular case of the previous modes. It allows the counter to be started in response to a stimulus and to generate a pulse with a programmable length after a programmable delay.

Starting the counter can be controlled through the slave mode controller. Generating the waveform can be done in output compare mode or PWM mode. You select One-pulse mode by setting the OPM bit in the TIMx\_CR1 register. This makes the counter stop automatically at the next update event UEV.

A pulse can be correctly generated only if the compare value is different from the counter initial value. Before starting (when the timer is waiting for the trigger), the configuration must be:

- In upcounting: CNT < CCRx ≤ ARR (in particular, 0 < CCRx)</li>
- In downcounting: CNT > CCRx

Figure 105. Example of one pulse mode.



For example you may want to generate a positive pulse on OC1 with a length of t<sub>PULSE</sub> and after a delay of t<sub>DELAY</sub> as soon as a positive edge is detected on the TI2 input pin.

Let's use TI2FP2 as trigger 1:

- Map TI2FP2 to TI2 by writing CC2S='01' in the TIMx\_CCMR1 register.
- TI2FP2 must detect a rising edge, write CC2P='0' and CC2NP='0' in the TIMx\_CCER register.
- Configure TI2FP2 as trigger for the slave mode controller (TRGI) by writing TS='110' in the TIMx\_SMCR register.
- TI2FP2 is used to start the counter by writing SMS to '110' in the TIMx\_SMCR register (trigger mode).

The OPM waveform is defined by writing the compare registers (taking into account the clock frequency and the counter prescaler).

- The t<sub>DELAY</sub> is defined by the value written in the TIMx\_CCR1 register.
- The t<sub>PULSE</sub> is defined by the difference between the auto-reload value and the compare value (TIMx\_ARR - TIMx\_CCR1).
- Let's say you want to build a waveform with a transition from '0' to '1' when a compare match occurs and a transition from '1' to '0' when the counter reaches the auto-reload value. To do this you enable PWM mode 2 by writing OC1M=111 in the TIMx\_CCMR1 register. You can optionally enable the preload registers by writing OC1PE='1' in the TIMx\_CCMR1 register and ARPE in the TIMx\_CR1 register. In this case you have to write the compare value in the TIMx\_CCR1 register, the auto-reload value in the TIMx\_ARR register, generate an update by setting the UG bit and wait for external trigger event on TI2. CC1P is written to '0' in this example.

In our example, the DIR and CMS bits in the TIMx\_CR1 register should be low.

You only want 1 pulse (Single mode), so you write '1 in the OPM bit in the TIMx\_CR1 register to stop the counter at the next update event (when the counter rolls over from the auto-reload value back to 0). When OPM bit in the TIMx\_CR1 register is set to '0', so the Repetitive Mode is selected.

#### Particular case: OCx fast enable:

In One-pulse mode, the edge detection on TIx input set the CEN bit which enables the counter. Then the comparison between the counter and the compare value makes the output toggle. But several clock cycles are needed for these operations and it limits the minimum delay  $t_{\text{DFI AY}}$  min we can get.

If you want to output a waveform with the minimum delay, you can set the OCxFE bit in the TIMx\_CCMRx register. Then OCxRef (and OCx) are forced in response to the stimulus, without taking in account the comparison. Its new level is the same as if a compare match had occurred. OCxFE acts only if the channel is configured in PWM1 or PWM2 mode.

# 13.3.16 Encoder interface mode

To select Encoder Interface mode write SMS='001' in the TIMx\_SMCR register if the counter is counting on TI2 edges only, SMS='010' if it is counting on TI1 edges only and SMS='011' if it is counting on both TI1 and TI2 edges.

Select the TI1 and TI2 polarity by programming the CC1P and CC2P bits in the TIMx\_CCER register. When needed, you can program the input filter as well. CC1NP and CC2NP must be kept low.

The two inputs TI1 and TI2 are used to interface to an incremental encoder. Refer to *Table 57*. The counter is clocked by each valid transition on TI1FP1 or TI2FP2 (TI1 and TI2 after input filter and polarity selection, TI1FP1=TI1 if not filtered and not inverted, TI2FP2=TI2 if not filtered and not inverted) assuming that it is enabled (CEN bit in TIMx\_CR1 register written to '1'). The sequence of transitions of the two inputs is evaluated and generates count pulses as well as the direction signal. Depending on the sequence the counter counts up or down, the DIR bit in the TIMx\_CR1 register is modified by hardware accordingly. The DIR bit is calculated at each transition on any input (TI1 or TI2), whatever the counter is counting on TI1 only, TI2 only or both TI1 and TI2.

Encoder interface mode acts simply as an external clock with direction selection. This means that the counter just counts continuously between 0 and the auto-reload value in the TIMx\_ARR register (0 to ARR or ARR down to 0 depending on the direction). So you must

configure TIMx\_ARR before starting. in the same way, the capture, compare, prescaler, repetition counter, trigger output features continue to work as normal. Encoder mode and External clock mode 2 are not compatible and must not be selected together.

In this mode, the counter is modified automatically following the speed and the direction of the incremental encoder and its content, therefore, always represents the encoder's position. The count direction correspond to the rotation direction of the connected sensor. The table summarizes the possible combinations, assuming TI1 and TI2 don't switch at the same time.

	Level on	TI1FP1	signal	TI2FP2	signal
Active edge	opposite signal (TI1FP1 for TI2, TI2FP2 for TI1)	Rising	Falling	Rising	Falling
Counting on	High	Down	Up	No Count	No Count
TI1 only	Low	Up	Down	No Count	No Count
Counting on	High	No Count	No Count	Up	Down
TI2 only	Low	No Count	No Count	Down	Up
Counting on	High	Down	Up	Up	Down
TI1 and TI2	Low	Up	Down	Down	Up

Table 57. Counting direction versus encoder signals

An external incremental encoder can be connected directly to the MCU without external interface logic. However, comparators are normally be used to convert the encoder's differential outputs to digital signals. This greatly increases noise immunity. The third encoder output which indicate the mechanical zero position, may be connected to an external interrupt input and trigger a counter reset.

*Figure 106* gives an example of counter operation, showing count signal generation and direction control. It also shows how input jitter is compensated where both edges are selected. This might occur if the sensor is positioned near to one of the switching points. For this example we assume that the configuration is the following:

- CC1S='01' (TIMx\_CCMR1 register, TI1FP1 mapped on TI1).
- CC2S='01' (TIMx\_CCMR2 register, TI1FP2 mapped on TI2).
- CC1P='0', CC1NP='0', and IC1F = '0000' (TIMx\_CCER register, TI1FP1 non-inverted, TI1FP1=TI1).
- CC2P='0', CC2NP='0', and IC2F = '0000' (TIMx\_CCER register, TI1FP2 non-inverted, TI1FP2= TI2).
- SMS='011' (TIMx\_SMCR register, both inputs are active on both rising and falling edges).
- CEN='1' (TIMx\_CR1 register, Counter enabled).

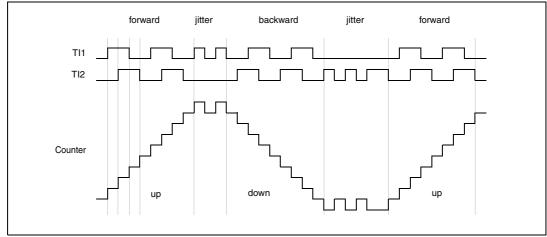


Figure 106. Example of counter operation in encoder interface mode.

Figure 107 gives an example of counter behavior when TI1FP1 polarity is inverted (same configuration as above except CC1P='1').

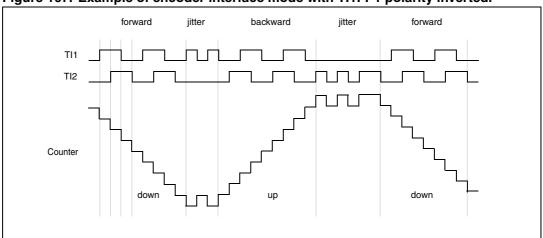


Figure 107. Example of encoder interface mode with TI1FP1 polarity inverted.

The timer, when configured in Encoder Interface mode provides information on the sensor's current position. You can obtain dynamic information (speed, acceleration, deceleration) by measuring the period between two encoder events using a second timer configured in capture mode. The output of the encoder which indicates the mechanical zero can be used for this purpose. Depending on the time between two events, the counter can also be read at regular times. You can do this by latching the counter value into a third input capture register if available (then the capture signal must be periodic and can be generated by another timer). when available, it is also possible to read its value through a DMA request generated by a real-time clock.

# 13.3.17 Timer input XOR function

The TI1S bit in the TIMx\_CR2 register, allows the input filter of channel 1 to be connected to the output of a XOR gate, combining the three input pins TIMx\_CH1, TIMx\_CH2 and TIMx\_CH3.

The XOR output can be used with all the timer input functions such as trigger or input capture. An example of this feature used to interface Hall sensors is given in *Section 13.3.18* below.

# 13.3.18 Interfacing with Hall sensors

This is done using the advanced-control timers (TIM1 or TIM8) to generate PWM signals to drive the motor and another timer TIMx (TIM2, TIM3, TIM4 or TIM5) referred to as "interfacing timer" in *Figure 108*. The "interfacing timer" captures the 3 timer input pins (TIMx\_CH1, TIMx\_CH2, and TIMx\_CH3) connected through a XOR to the TI1 input channel (selected by setting the TI1S bit in the TIMx\_CR2 register).

The slave mode controller is configured in reset mode; the slave input is TI1F\_ED. Thus, each time one of the 3 inputs toggles, the counter restarts counting from 0. This creates a time base triggered by any change on the Hall inputs.

On the "interfacing timer", capture/compare channel 1 is configured in capture mode, capture signal is TRC (see *Figure 91: Capture/compare channel (example: channel 1 input stage) on page 315*). The captured value, which corresponds to the time elapsed between 2 changes on the inputs, gives information about motor speed.

The "interfacing timer" can be used in output mode to generate a pulse which changes the configuration of the channels of the advanced-control timer (TIM1 or TIM8) (by triggering a COM event). The TIM1 timer is used to generate PWM signals to drive the motor. To do this, the interfacing timer channel must be programmed so that a positive pulse is generated after a programmed delay (in output compare or PWM mode). This pulse is sent to the advanced-control timer (TIM1 or TIM8) through the TRGO output.

Example: you want to change the PWM configuration of your advanced-control timer TIM1 after a programmed delay each time a change occurs on the Hall inputs connected to one of the TIMx timers.

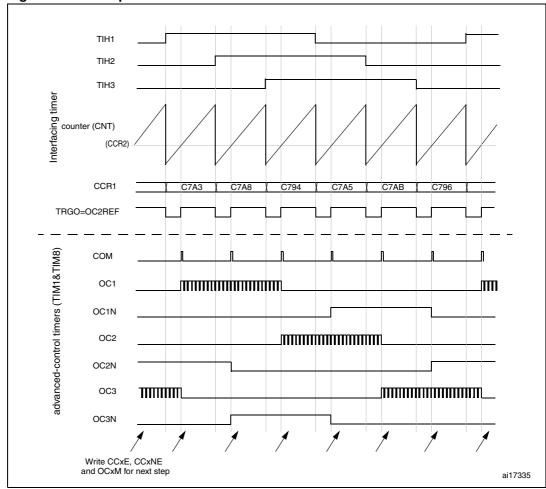
- Configure 3 timer inputs ORed to the TI1 input channel by writing the TI1S bit in the TIMx\_CR2 register to '1',
- Program the time base: write the TIMx\_ARR to the max value (the counter must be cleared by the TI1 change. Set the prescaler to get a maximum counter period longer than the time between 2 changes on the sensors,
- Program channel 1 in capture mode (TRC selected): write the CC1S bits in the TIMx CCMR1 register to '11'. You can also program the digital filter if needed,
- Program channel 2 in PWM 2 mode with the desired delay: write the OC2M bits to '111' and the CC2S bits to '00' in the TIMx\_CCMR1 register,
- Select OC2REF as trigger output on TRGO: write the MMS bits in the TIMx\_CR2 register to '101',

In the advanced-control timer TIM1, the right ITR input must be selected as trigger input, the timer is programmed to generate PWM signals, the capture/compare control signals are preloaded (CCPC=1 in the TIMx\_CR2 register) and the COM event is controlled by the trigger input (CCUS=1 in the TIMx\_CR2 register). The PWM control bits (CCxE, OCxM) are

written after a COM event for the next step (this can be done in an interrupt subroutine generated by the rising edge of OC2REF).

Figure 108 describes this example.

Figure 108. Example of hall sensor interface



# 13.3.19 TIMx and external trigger synchronization

The TIMx timer can be synchronized with an external trigger in several modes: Reset mode, Gated mode and Trigger mode.

#### Slave mode: Reset mode

The counter and its prescaler can be reinitialized in response to an event on a trigger input. Moreover, if the URS bit from the TIMx\_CR1 register is low, an update event UEV is generated. Then all the preloaded registers (TIMx\_ARR, TIMx\_CCRx) are updated.

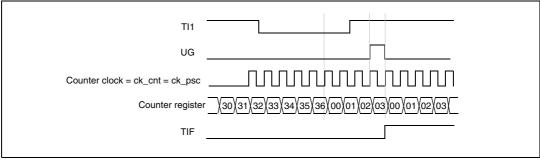
In the following example, the upcounter is cleared in response to a rising edge on TI1 input:

- Configure the channel 1 to detect rising edges on TI1. Configure the input filter duration (in this example, we don't need any filter, so we keep IC1F=0000). The capture prescaler is not used for triggering, so you don't need to configure it. The CC1S bits select the input capture source only, CC1S = 01 in the TIMx\_CCMR1 register. Write CC1P=0 and CC1NP='0' in TIMx\_CCER register to validate the polarity (and detect rising edges only).
- Configure the timer in reset mode by writing SMS=100 in TIMx\_SMCR register. Select TI1 as the input source by writing TS=101 in TIMx\_SMCR register.
- Start the counter by writing CEN=1 in the TIMx\_CR1 register.

The counter starts counting on the internal clock, then behaves normally until TI1 rising edge. When TI1 rises, the counter is cleared and restarts from 0. In the meantime, the trigger flag is set (TIF bit in the TIMx\_SR register) and an interrupt request, or a DMA request can be sent if enabled (depending on the TIE and TDE bits in TIMx\_DIER register).

The following figure shows this behavior when the auto-reload register TIMx\_ARR=0x36. The delay between the rising edge on TI1 and the actual reset of the counter is due to the resynchronization circuit on TI1 input.





#### Slave mode: Gated mode

The counter can be enabled depending on the level of a selected input.

In the following example, the upcounter counts only when TI1 input is low:

- Configure the channel 1 to detect low levels on TI1. Configure the input filter duration
  (in this example, we don't need any filter, so we keep IC1F=0000). The capture
  prescaler is not used for triggering, so you don't need to configure it. The CC1S bits
  select the input capture source only, CC1S=01 in TIMx\_CCMR1 register. Write
  CC1P=1 and CC1NP='0' in TIMx\_CCER register to validate the polarity (and detect low
  level only).
- Configure the timer in gated mode by writing SMS=101 in TIMx\_SMCR register. Select TI1 as the input source by writing TS=101 in TIMx SMCR register.
- Enable the counter by writing CEN=1 in the TIMx\_CR1 register (in gated mode, the counter doesn't start if CEN=0, whatever is the trigger input level).

The counter starts counting on the internal clock as long as TI1 is low and stops as soon as TI1 becomes high. The TIF flag in the TIMx\_SR register is set both when the counter starts or stops.

The delay between the rising edge on TI1 and the actual stop of the counter is due to the resynchronization circuit on TI1 input.

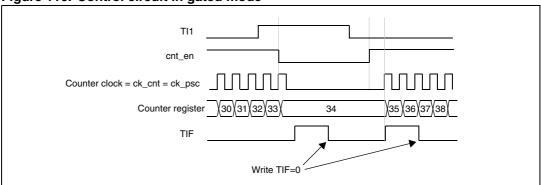


Figure 110. Control circuit in gated mode

# Slave mode: Trigger mode

The counter can start in response to an event on a selected input.

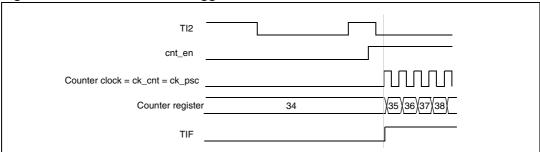
In the following example, the upcounter starts in response to a rising edge on TI2 input:

- Configure the channel 2 to detect rising edges on TI2. Configure the input filter duration (in this example, we don't need any filter, so we keep IC2F=0000). The capture prescaler is not used for triggering, so you don't need to configure it. The CC2S bits are configured to select the input capture source only, CC2S=01 in TIMx\_CCMR1 register. Write CC2P=1 and CC2NP=0 in TIMx\_CCER register to validate the polarity (and detect low level only).
- Configure the timer in trigger mode by writing SMS=110 in TIMx\_SMCR register. Select TI2 as the input source by writing TS=110 in TIMx\_SMCR register.

When a rising edge occurs on TI2, the counter starts counting on the internal clock and the TIF flag is set.

The delay between the rising edge on TI2 and the actual start of the counter is due to the resynchronization circuit on TI2 input.

Figure 111. Control circuit in trigger mode



### Slave mode: external clock mode 2 + trigger mode

The external clock mode 2 can be used in addition to another slave mode (except external clock mode 1 and encoder mode). In this case, the ETR signal is used as external clock input, and another input can be selected as trigger input (in reset mode, gated mode or trigger mode). It is recommended not to select ETR as TRGI through the TS bits of TIMx SMCR register.

In the following example, the upcounter is incremented at each rising edge of the ETR signal as soon as a rising edge of TI1 occurs:

- Configure the external trigger input circuit by programming the TIMx\_SMCR register as follows:
  - ETF = 0000: no filter
  - ETPS=00: prescaler disabled
  - ETP=0: detection of rising edges on ETR and ECE=1 to enable the external clock mode 2.

- 2. Configure the channel 1 as follows, to detect rising edges on TI:
  - IC1F=0000: no filter.
  - The capture prescaler is not used for triggering and does not need to be configured.
  - CC1S=01 in TIMx\_CCMR1 register to select only the input capture source
  - CC1P=0 and CC1NP='0' in TIMx\_CCER register to validate the polarity (and detect rising edge only).
- 3. Configure the timer in trigger mode by writing SMS=110 in TIMx\_SMCR register. Select TI1 as the input source by writing TS=101 in TIMx\_SMCR register.

A rising edge on TI1 enables the counter and sets the TIF flag. The counter then counts on ETR rising edges.

The delay between the rising edge of the ETR signal and the actual reset of the counter is due to the resynchronization circuit on ETRP input.

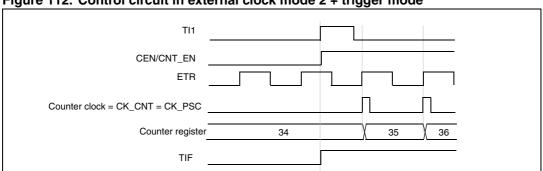


Figure 112. Control circuit in external clock mode 2 + trigger mode

# 13.3.20 Timer synchronization

The TIM timers are linked together internally for timer synchronization or chaining. Refer to *Section 14.3.15: Timer synchronization on page 398* for details.

## 13.3.21 **Debug mode**

When the microcontroller enters debug mode (Cortex<sup>TM</sup>-M4F core halted), the TIMx counter either continues to work normally or stops, depending on DBG\_TIMx\_STOP configuration bit in DBG module. For more details, refer to *Section 32.16.2: Debug support for timers, watchdog, bxCAN and I2C*.

# 13.4 TIM1&TIM8 registers

Refer to Section 1.1 on page 46 for a list of abbreviations used in register descriptions.

The peripheral registers must be written by half-words (16 bits) or words (32 bits). Read accesses can be done by bytes (8 bits), half-word (16 bits) or words (32 bits).

# 13.4.1 TIM1&TIM8 control register 1 (TIMx\_CR1)

Address offset: 0x00 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
		Rese	nuod			CKE	[1:0]	ARPE	CMS	8[1:0]	DIR	ОРМ	URS	UDIS	CEN	
		nese	erveu			rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	

Bits 15:10 Reserved, must be kept at reset value.

### Bits 9:8 CKD[1:0]: Clock division

This bit-field indicates the division ratio between the timer clock (CK\_INT) frequency and the dead-time and sampling clock (t<sub>DTS</sub>)used by the dead-time generators and the digital filters (ETR, Tlx),

00: t<sub>DTS</sub>=t<sub>CK\_INT</sub> 01: t<sub>DTS</sub>=2\*t<sub>CK\_INT</sub> 10: t<sub>DTS</sub>=4\*t<sub>CK\_INT</sub>

11: Reserved, do not program this value

#### Bit 7 ARPE: Auto-reload preload enable

0: TIMx\_ARR register is not buffered1: TIMx\_ARR register is buffered

### Bits 6:5 CMS[1:0]: Center-aligned mode selection

00: Edge-aligned mode. The counter counts up or down depending on the direction bit (DIR).

01: Center-aligned mode 1. The counter counts up and down alternatively. Output compare interrupt flags of channels configured in output (CCxS=00 in TIMx\_CCMRx register) are set only when the counter is counting down.

10: Center-aligned mode 2. The counter counts up and down alternatively. Output compare interrupt flags of channels configured in output (CCxS=00 in TIMx\_CCMRx register) are set only when the counter is counting up.

11: Center-aligned mode 3. The counter counts up and down alternatively. Output compare interrupt flags of channels configured in output (CCxS=00 in TIMx\_CCMRx register) are set both when the counter is counting up or down.

Note: It is not allowed to switch from edge-aligned mode to center-aligned mode as long as the counter is enabled (CEN=1)

#### Bit 4 DIR: Direction

0: Counter used as upcounter

1: Counter used as downcounter

Note: This bit is read only when the timer is configured in Center-aligned mode or Encoder mode.

### Bit 3 OPM: One pulse mode

- 0: Counter is not stopped at update event
- 1: Counter stops counting at the next update event (clearing the bit CEN)

#### Bit 2 URS: Update request source

This bit is set and cleared by software to select the UEV event sources.

0: Any of the following events generate an update interrupt or DMA request if enabled. These events can be:

- Counter overflow/underflow
- Setting the UG bit
- Update generation through the slave mode controller
- 1: Only counter overflow/underflow generates an update interrupt or DMA request if enabled.

#### Bit 1 UDIS: Update disable

This bit is set and cleared by software to enable/disable UEV event generation.

0: UEV enabled. The Update (UEV) event is generated by one of the following events:

- Counter overflow/underflow
- Setting the UG bit
- Update generation through the slave mode controller

Buffered registers are then loaded with their preload values.

1: UEV disabled. The Update event is not generated, shadow registers keep their value (ARR, PSC, CCRx). However the counter and the prescaler are reinitialized if the UG bit is set or if a hardware reset is received from the slave mode controller.

#### Bit 0 CEN: Counter enable

- 0: Counter disabled
- 1: Counter enabled

Note: External clock, gated mode and encoder mode can work only if the CEN bit has been previously set by software. However trigger mode can set the CEN bit automatically by hardware.

# 13.4.2 TIM1&TIM8 control register 2 (TIMx\_CR2)

Address offset: 0x04 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Res.	OIS4	OIS3N	OIS3	OIS2N	OIS2	OIS1N	OIS1	TI1S		MMS[2:0]		CCDS	ccus	Res.	CCPC
nes.	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	nes.	rw

Bit 15 Reserved, must be kept at reset value.

Bit 14 OIS4: Output Idle state 4 (OC4 output)

refer to OIS1 bit

Bit 13 OIS3N: Output Idle state 3 (OC3N output)

refer to OIS1N bit

Bit 12 OIS3: Output Idle state 3 (OC3 output)

refer to OIS1 bit

Bit 11 OIS2N: Output Idle state 2 (OC2N output)

refer to OIS1N bit

Bit 10 OIS2: Output Idle state 2 (OC2 output)

refer to OIS1 bit

Bit 9 OIS1N: Output Idle state 1 (OC1N output)

0: OC1N=0 after a dead-time when MOE=0

1: OC1N=1 after a dead-time when MOE=0

Note: This bit can not be modified as long as LOCK level 1, 2 or 3 has been programmed (LOCK bits in TIMx\_BDTR register).

Bit 8 OIS1: Output Idle state 1 (OC1 output)

0: OC1=0 (after a dead-time if OC1N is implemented) when MOE=0

1: OC1=1 (after a dead-time if OC1N is implemented) when MOE=0

Note: This bit can not be modified as long as LOCK level 1, 2 or 3 has been programmed (LOCK bits in TIMx\_BDTR register).

Bit 7 TI1S: TI1 selection

0: The TIMx\_CH1 pin is connected to TI1 input

1: The TIMx\_CH1, CH2 and CH3 pins are connected to the TI1 input (XOR combination)

#### Bits 6:4 MMS[1:0]: Master mode selection

These bits allow to select the information to be sent in master mode to slave timers for synchronization (TRGO). The combination is as follows:

000: **Reset** - the UG bit from the TIMx\_EGR register is used as trigger output (TRGO). If the reset is generated by the trigger input (slave mode controller configured in reset mode) then the signal on TRGO is delayed compared to the actual reset.

001: **Enable** - the Counter Enable signal CNT\_EN is used as trigger output (TRGO). It is useful to start several timers at the same time or to control a window in which a slave timer is enable. The Counter Enable signal is generated by a logic OR between CEN control bit and the trigger input when configured in gated mode. When the Counter Enable signal is controlled by the trigger input, there is a delay on TRGO, except if the master/slave mode is selected (see the MSM bit description in TIMx\_SMCR register).

010: **Update** - The update event is selected as trigger output (TRGO). For instance a master timer can then be used as a prescaler for a slave timer.

011: **Compare Pulse** - The trigger output send a positive pulse when the CC1IF flag is to be set (even if it was already high), as soon as a capture or a compare match occurred. (TRGO).

100: Compare - OC1REF signal is used as trigger output (TRGO)

101: Compare - OC2REF signal is used as trigger output (TRGO)

110: Compare - OC3REF signal is used as trigger output (TRGO)

111: Compare - OC4REF signal is used as trigger output (TRGO)

Bit 3 CCDS: Capture/compare DMA selection

0: CCx DMA request sent when CCx event occurs

1: CCx DMA requests sent when update event occurs

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- Bit 2 CCUS: Capture/compare control update selection
  - 0: When capture/compare control bits are preloaded (CCPC=1), they are updated by setting the COMG bit only
  - 1: When capture/compare control bits are preloaded (CCPC=1), they are updated by setting the COMG bit or when an rising edge occurs on TRGI

Note: This bit acts only on channels that have a complementary output.

- Bit 1 Reserved, must be kept at reset value.
- Bit 0 CCPC: Capture/compare preloaded control
  - 0: CCxE, CCxNE and OCxM bits are not preloaded
  - 1: CCxE, CCxNE and OCxM bits are preloaded, after having been written, they are updated only when a commutation event (COM) occurs (COMG bit set or rising edge detected on TRGI, depending on the CCUS bit).

Note: This bit acts only on channels that have a complementary output.

# 13.4.3 TIM1&TIM8 slave mode control register (TIMx\_SMCR)

Address offset: 0x08 Reset value: 0x0000

	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Ī	ETP	ECE	ETPS	S[1:0]		ETF[3:0]			MSM		TS[2:0]		Res.		SMS[2:0]	
Ī	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	Res.	rw	rw	rw

### Bit 15 ETP: External trigger polarity

This bit selects whether ETR or ETR is used for trigger operations

0: ETR is non-inverted, active at high level or rising edge.

1: ETR is inverted, active at low level or falling edge.

#### Bit 14 ECE: External clock enable

This bit enables External clock mode 2.

0: External clock mode 2 disabled

1: External clock mode 2 enabled. The counter is clocked by any active edge on the ETRF signal.

Note: 1: Setting the ECE bit has the same effect as selecting external clock mode 1 with TRGI connected to ETRF (SMS=111 and TS=111).

2: It is possible to simultaneously use external clock mode 2 with the following slave modes: reset mode, gated mode and trigger mode. Nevertheless, TRGI must not be connected to ETRF in this case (TS bits must not be 111).

**3:** If external clock mode 1 and external clock mode 2 are enabled at the same time, the external clock input is ETRF.

### Bits 13:12 ETPS[1:0]: External trigger prescaler

External trigger signal ETRP frequency must be at most 1/4 of TIMxCLK frequency. A prescaler can be enabled to reduce ETRP frequency. It is useful when inputting fast external clocks.

00: Prescaler OFF

01: ETRP frequency divided by 2

10: ETRP frequency divided by 4

11: ETRP frequency divided by 8

### Bits 11:8 ETF[3:0]: External trigger filter

This bit-field then defines the frequency used to sample ETRP signal and the length of the digital filter applied to ETRP. The digital filter is made of an event counter in which N events are needed to validate a transition on the output:

```
0000: No filter, sampling is done at f_{DTS} 0001: f_{SAMPLING} = f_{CK\_INT}, N=2 0010: f_{SAMPLING} = f_{CK\_INT}, N=4 0011: f_{SAMPLING} = f_{CK\_INT}, N=8 0100: f_{SAMPLING} = f_{DTS}/2, N=6 0101: f_{SAMPLING} = f_{DTS}/4, N=6 0111: f_{SAMPLING} = f_{DTS}/4, N=6 0111: f_{SAMPLING} = f_{DTS}/4, N=8 1000: f_{SAMPLING} = f_{DTS}/8, N=6 1001: f_{SAMPLING} = f_{DTS}/8, N=8 1010: f_{SAMPLING} = f_{DTS}/16, N=5 1011: f_{SAMPLING} = f_{DTS}/16, N=6 1100: f_{SAMPLING} = f_{DTS}/16, N=8 1101: f_{SAMPLING} = f_{DTS}/16, N=8 1101: f_{SAMPLING} = f_{DTS}/32, N=5 1110: f_{SAMPLING} = f_{DTS}/32, N=6
```

#### Bit 7 MSM: Master/slave mode

1111: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/32, N=8

0: No action

1: The effect of an event on the trigger input (TRGI) is delayed to allow a perfect synchronization between the current timer and its slaves (through TRGO). It is useful if we want to synchronize several timers on a single external event.

### Bits 6:4 TS[2:0]: Trigger selection

This bit-field selects the trigger input to be used to synchronize the counter.

```
000: Internal Trigger 0 (ITR0)
001: Internal Trigger 1 (ITR1)
010: Internal Trigger 2 (ITR2)
011: Internal Trigger 3 (ITR3)
100: TI1 Edge Detector (TI1F_ED)
101: Filtered Timer Input 1 (TI1FP1)
110: Filtered Timer Input 2 (TI2FP2)
111: External Trigger input (ETRF)
```

See *Table 58: TIMx Internal trigger connection on page 345* for more details on ITRx meaning for each Timer.

Note: These bits must be changed only when they are not used (e.g. when SMS=000) to avoid wrong edge detections at the transition.

Bit 3 Reserved, must be kept at reset value.

#### Bits 2:0 SMS: Slave mode selection

When external signals are selected the active edge of the trigger signal (TRGI) is linked to the polarity selected on the external input (see Input Control register and Control Register description.

000: Slave mode disabled - if CEN = '1' then the prescaler is clocked directly by the internal clock.

001: Encoder mode 1 - Counter counts up/down on TI2FP2 edge depending on TI1FP1 level.

010: Encoder mode 2 - Counter counts up/down on TI1FP1 edge depending on TI2FP2 level.

011: Encoder mode 3 - Counter counts up/down on both TI1FP1 and TI2FP2 edges depending on the level of the other input.

100: Reset Mode - Rising edge of the selected trigger input (TRGI) reinitializes the counter and generates an update of the registers.

101: Gated Mode - The counter clock is enabled when the trigger input (TRGI) is high. The counter stops (but is not reset) as soon as the trigger becomes low. Both start and stop of the counter are controlled.

110: Trigger Mode - The counter starts at a rising edge of the trigger TRGI (but it is not reset). Only the start of the counter is controlled.

111: External Clock Mode 1 - Rising edges of the selected trigger (TRGI) clock the counter.

Note: The gated mode must not be used if TI1F\_ED is selected as the trigger input (TS='100'). Indeed, TI1F\_ED outputs 1 pulse for each transition on TI1F, whereas the gated mode checks the level of the trigger signal.

Table 58. TIMx Internal trigger connection

Slave TIM	ITR0 (TS = 000)	ITR1 (TS = 001)	ITR2 (TS = 010)	ITR3 (TS = 011)
TIM1	TIM5	TIM2	TIM3	TIM4
TIM8	TIM1	TIM2	TIM4	TIM5

# 13.4.4 TIM1&TIM8 DMA/interrupt enable register (TIMx\_DIER)

Address offset: 0x0C Reset value: 0x0000

	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Ī	Res.	TDE	COMDE	CC4DE	CC3DE	CC2DE	CC1DE	UDE	BIE	TIE	COMIE	CC4IE	CC3IE	CC2IE	CC1IE	UIE
	nes.	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bit 15 Reserved, must be kept at reset value.

Bit 14 TDE: Trigger DMA request enable

0: Trigger DMA request disabled1: Trigger DMA request enabled

Bit 13 **COMDE**: COM DMA request enable

0: COM DMA request disabled 1: COM DMA request enabled

Bit 12 CC4DE: Capture/Compare 4 DMA request enable

0: CC4 DMA request disabled1: CC4 DMA request enabled

- Bit 11 CC3DE: Capture/Compare 3 DMA request enable
  - 0: CC3 DMA request disabled
  - 1: CC3 DMA request enabled
- Bit 10 CC2DE: Capture/Compare 2 DMA request enable
  - 0: CC2 DMA request disabled
  - 1: CC2 DMA request enabled
- Bit 9 CC1DE: Capture/Compare 1 DMA request enable
  - 0: CC1 DMA request disabled
  - 1: CC1 DMA request enabled
- Bit 8 **UDE**: Update DMA request enable
  - 0: Update DMA request disabled
  - 1: Update DMA request enabled
- Bit 7 BIE: Break interrupt enable
  - 0: Break interrupt disabled
  - 1: Break interrupt enabled
- Bit 6 TIE: Trigger interrupt enable
  - 0: Trigger interrupt disabled
  - 1: Trigger interrupt enabled
- Bit 5 COMIE: COM interrupt enable
  - 0: COM interrupt disabled
  - 1: COM interrupt enabled
- Bit 4 CC4IE: Capture/Compare 4 interrupt enable
  - 0: CC4 interrupt disabled
  - 1: CC4 interrupt enabled
- Bit 3 CC3IE: Capture/Compare 3 interrupt enable
  - 0: CC3 interrupt disabled
  - 1: CC3 interrupt enabled
- Bit 2 CC2IE: Capture/Compare 2 interrupt enable
  - 0: CC2 interrupt disabled
  - 1: CC2 interrupt enabled
- Bit 1 CC1IE: Capture/Compare 1 interrupt enable
  - 0: CC1 interrupt disabled
  - 1: CC1 interrupt enabled
- Bit 0 UIE: Update interrupt enable
  - 0: Update interrupt disabled
  - 1: Update interrupt enabled

# 13.4.5 TIM1&TIM8 status register (TIMx\_SR)

Address offset: 0x10 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Reserved		CC4OF	CC3OF	CC2OF	CC10F	Res.	BIF	TIF	COMIF	CC4IF	CC3IF	CC2IF	CC1IF	UIF	
Reserved	rc_w0	rc_w0	rc_w0	rc_w0	Res.	rc_w0									

Bits 15:13 Reserved, must be kept at reset value.

Bit 12 CC40F: Capture/Compare 4 overcapture flag

refer to CC1OF description

Bit 11 CC3OF: Capture/Compare 3 overcapture flag

refer to CC1OF description

Bit 10 CC2OF: Capture/Compare 2 overcapture flag

refer to CC1OF description

Bit 9 CC10F: Capture/Compare 1 overcapture flag

This flag is set by hardware only when the corresponding channel is configured in input capture mode. It is cleared by software by writing it to '0'.

0: No overcapture has been detected.

1: The counter value has been captured in TIMx\_CCR1 register while CC1IF flag was already set

Bit 8 Reserved, must be kept at reset value.

Bit 7 BIF: Break interrupt flag

This flag is set by hardware as soon as the break input goes active. It can be cleared by software if the break input is not active.

0: No break event occurred.

1: An active level has been detected on the break input.

Bit 6 TIF: Trigger interrupt flag

This flag is set by hardware on trigger event (active edge detected on TRGI input when the slave mode controller is enabled in all modes but gated mode. It is set when the counter starts or stops when gated mode is selected. It is cleared by software.

0: No trigger event occurred.

1: Trigger interrupt pending.

Bit 5 COMIF: COM interrupt flag

This flag is set by hardware on COM event (when Capture/compare Control bits - CCxE, CCxNE, OCxM - have been updated). It is cleared by software.

0: No COM event occurred.

1: COM interrupt pending.

Bit 4 CC4IF: Capture/Compare 4 interrupt flag

refer to CC1IF description

Bit 3 CC3IF: Capture/Compare 3 interrupt flag

refer to CC1IF description

Bit 2 CC2IF: Capture/Compare 2 interrupt flag

refer to CC1IF description

Bit 1 CC1IF: Capture/Compare 1 interrupt flag

#### If channel CC1 is configured as output:

This flag is set by hardware when the counter matches the compare value, with some exception in center-aligned mode (refer to the CMS bits in the TIMx\_CR1 register description). It is cleared by software.

0: No match.

1: The content of the counter TIMx\_CNT matches the content of the TIMx\_CCR1 register. When the contents of TIMx\_CCR1 are greater than the contents of TIMx\_ARR, the CC1IF bit goes high on the counter overflow (in upcounting and up/down-counting modes) or underflow (in downcounting mode)

## If channel CC1 is configured as input:

This bit is set by hardware on a capture. It is cleared by software or by reading the TIMx\_CCR1 register.

0: No input capture occurred

1: The counter value has been captured in TIMx\_CCR1 register (An edge has been detected on IC1 which matches the selected polarity)

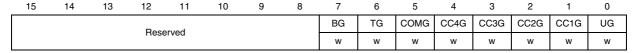
### Bit 0 UIF: Update interrupt flag

This bit is set by hardware on an update event. It is cleared by software.

- 0: No update occurred.
- 1: Update interrupt pending. This bit is set by hardware when the registers are updated:
  - -At overflow or underflow regarding the repetition counter value (update if repetition counter = 0) and if the UDIS=0 in the TIMx\_CR1 register.
- -When CNT is reinitialized by software using the UG bit in TIMx\_EGR register, if URS=0 and UDIS=0 in the TIMx\_CR1 register.
- -When CNT is reinitialized by a trigger event (refer to Section 13.4.3: TIM1&TIM8 slave mode control register (TIMx\_SMCR)), if URS=0 and UDIS=0 in the TIMx\_CR1 register.

# 13.4.6 TIM1&TIM8 event generation register (TIMx\_EGR)

Address offset: 0x14 Reset value: 0x0000



Bits 15:8 Reserved, must be kept at reset value.

## Bit 7 BG: Break generation

This bit is set by software in order to generate an event, it is automatically cleared by hardware.

0: No action

1: A break event is generated. MOE bit is cleared and BIF flag is set. Related interrupt or DMA transfer can occur if enabled.

### Bit 6 TG: Trigger generation

This bit is set by software in order to generate an event, it is automatically cleared by hardware.

0: No action

1: The TIF flag is set in TIMx\_SR register. Related interrupt or DMA transfer can occur if enabled.

#### Bit 5 **COMG**: Capture/Compare control update generation

This bit can be set by software, it is automatically cleared by hardware

0: No action

1: When CCPC bit is set, it allows to update CCxE, CCxNE and OCxM bits

Note: This bit acts only on channels having a complementary output.

### Bit 4 CC4G: Capture/Compare 4 generation

refer to CC1G description

### Bit 3 CC3G: Capture/Compare 3 generation

refer to CC1G description

### Bit 2 CC2G: Capture/Compare 2 generation

refer to CC1G description

### Bit 1 CC1G: Capture/Compare 1 generation

This bit is set by software in order to generate an event, it is automatically cleared by hardware.

0: No action

1: A capture/compare event is generated on channel 1:

### If channel CC1 is configured as output:

CC1IF flag is set, Corresponding interrupt or DMA request is sent if enabled.

### If channel CC1 is configured as input:

The current value of the counter is captured in TIMx\_CCR1 register. The CC1IF flag is set, the corresponding interrupt or DMA request is sent if enabled. The CC1OF flag is set if the CC1IF flag was already high.

## Bit 0 **UG**: Update generation

This bit can be set by software, it is automatically cleared by hardware.

0: No action

1: Reinitialize the counter and generates an update of the registers. Note that the prescaler counter is cleared too (anyway the prescaler ratio is not affected). The counter is cleared if the center-aligned mode is selected or if DIR=0 (upcounting), else it takes the auto-reload value (TIMx\_ARR) if DIR=1 (downcounting).



# 13.4.7 TIM1&TIM8 capture/compare mode register 1 (TIMx\_CCMR1)

Address offset: 0x18 Reset value: 0x0000

The channels can be used in input (capture mode) or in output (compare mode). The direction of a channel is defined by configuring the corresponding CCxS bits. All the other bits of this register have a different function in input and in output mode. For a given bit, OCxx describes its function when the channel is configured in output, ICxx describes its function when the channel is configured in input. So you must take care that the same bit can have a different meaning for the input stage and for the output stage.

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
OC2 CE	(	OC2M[2:0	]	OC2 PE	OC2 FE	CC2S	S[1:0]	OC1 CE	(	OC1M[2:0	)]	OC1 PE	OC1 FE	CC18	S[1:0]
	IC2F	[3:0]		IC2PS	C[1:0]				IC1F	[3:0]		IC1PS	SC[1:0]		
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

## Output compare mode:

Bit 15 OC2CE: Output Compare 2 clear enable

Bits 14:12 OC2M[2:0]: Output Compare 2 mode

Bit 11 OC2PE: Output Compare 2 preload enable

Bit 10 OC2FE: Output Compare 2 fast enable

Bits 9:8 CC2S[1:0]: Capture/Compare 2 selection

This bit-field defines the direction of the channel (input/output) as well as the used input.

00: CC2 channel is configured as output

01: CC2 channel is configured as input, IC2 is mapped on TI2

10: CC2 channel is configured as input, IC2 is mapped on TI1

11: CC2 channel is configured as input, IC2 is mapped on TRC. This mode is working only if

an internal trigger input is selected through the TS bit (TIMx\_SMCR register)

Note: CC2S bits are writable only when the channel is OFF (CC2E = '0' in TIMx\_CCER).

Bit 7 OC1CE: Output Compare 1 clear enable

OC1CE: Output Compare 1 Clear Enable

0: OC1Ref is not affected by the ETRF Input

1: OC1Ref is cleared as soon as a High level is detected on ETRF input

#### Bits 6:4 OC1M: Output Compare 1 mode

These bits define the behavior of the output reference signal OC1REF from which OC1 and OC1N are derived. OC1REF is active high whereas OC1 and OC1N active level depends on CC1P and CC1NP bits.

000: Frozen - The comparison between the output compare register TIMx\_CCR1 and the counter TIMx\_CNT has no effect on the outputs.(this mode is used to generate a timing base).

001: Set channel 1 to active level on match. OC1REF signal is forced high when the counter TIMx\_CNT matches the capture/compare register 1 (TIMx\_CCR1).

010: Set channel 1 to inactive level on match. OC1REF signal is forced low when the counter TIMx\_CNT matches the capture/compare register 1 (TIMx\_CCR1).

011: Toggle - OC1REF toggles when TIMx\_CNT=TIMx\_CCR1.

100: Force inactive level - OC1REF is forced low.

101: Force active level - OC1REF is forced high.

110: PWM mode 1 - In upcounting, channel 1 is active as long as TIMx\_CNT<TIMx\_CCR1 else inactive. In downcounting, channel 1 is inactive (OC1REF='0') as long as TIMx\_CNT>TIMx\_CCR1 else active (OC1REF='1').

111: PWM mode 2 - In upcounting, channel 1 is inactive as long as TIMx\_CNT<TIMx\_CCR1 else active. In downcounting, channel 1 is active as long as TIMx\_CNT>TIMx\_CCR1 else inactive.

Note: 1: These bits can not be modified as long as LOCK level 3 has been programmed (LOCK bits in TIMx\_BDTR register) and CC1S='00' (the channel is configured in output).

2: In PWM mode 1 or 2, the OCREF level changes only when the result of the comparison changes or when the output compare mode switches from "frozen" mode to "PWM" mode.

3: On channels having a complementary output, this bit field is preloaded. If the CCPC bit is set in the TIMx\_CR2 register then the OC1M active bits take the new value from the preloaded bits only when a COM event is generated.

#### Bit 3 OC1PE: Output Compare 1 preload enable

0: Preload register on TIMx\_CCR1 disabled. TIMx\_CCR1 can be written at anytime, the new value is taken in account immediately.

1: Preload register on TIMx\_CCR1 enabled. Read/Write operations access the preload register. TIMx\_CCR1 preload value is loaded in the active register at each update event.

Note: 1: These bits can not be modified as long as LOCK level 3 has been programmed (LOCK bits in TIMx\_BDTR register) and CC1S='00' (the channel is configured in output).

2: The PWM mode can be used without validating the preload register only in one pulse mode (OPM bit set in TIMx\_CR1 register). Else the behavior is not guaranteed.

### Bit 2 OC1FE: Output Compare 1 fast enable

This bit is used to accelerate the effect of an event on the trigger in input on the CC output. 0: CC1 behaves normally depending on counter and CCR1 values even when the trigger is ON. The minimum delay to activate CC1 output when an edge occurs on the trigger input is 5 clock cycles.

1: An active edge on the trigger input acts like a compare match on CC1 output. Then, OC is set to the compare level independently from the result of the comparison. Delay to sample the trigger input and to activate CC1 output is reduced to 3 clock cycles. OCFE acts only if the channel is configured in PWM1 or PWM2 mode.

### Bits 1:0 CC1S: Capture/Compare 1 selection

This bit-field defines the direction of the channel (input/output) as well as the used input.

00: CC1 channel is configured as output

01: CC1 channel is configured as input, IC1 is mapped on TI1

10: CC1 channel is configured as input, IC1 is mapped on TI2

11: CC1 channel is configured as input, IC1 is mapped on TRC. This mode is working only if an internal trigger input is selected through TS bit (TIMx\_SMCR register)

Note: CC1S bits are writable only when the channel is OFF (CC1E = '0' in TIMx\_CCER).

## Input capture mode

```
Bits 15:12 IC2F: Input capture 2 filter
```

### Bits 11:10 IC2PSC[1:0]: Input capture 2 prescaler

### Bits 9:8 CC2S: Capture/Compare 2 selection

This bit-field defines the direction of the channel (input/output) as well as the used input.

00: CC2 channel is configured as output

01: CC2 channel is configured as input, IC2 is mapped on TI2

10: CC2 channel is configured as input, IC2 is mapped on TI1

11: CC2 channel is configured as input, IC2 is mapped on TRC. This mode is working only if an internal trigger input is selected through TS bit (TIMx\_SMCR register)

Note: CC2S bits are writable only when the channel is OFF (CC2E = '0' in TIMx\_CCER).

#### Bits 7:4 IC1F[3:0]: Input capture 1 filter

This bit-field defines the frequency used to sample TI1 input and the length of the digital filter applied to TI1. The digital filter is made of an event counter in which N events are needed to validate a transition on the output:

```
0000: No filter, sampling is done at fDTS
```

0001: f<sub>SAMPLING</sub>=f<sub>CK INT</sub>, N=2

0010: f<sub>SAMPLING</sub>=f<sub>CK INT</sub>, N=4

0011: f<sub>SAMPLING</sub>=f<sub>CK INT</sub>, N=8

0100: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/2, N=6

0101: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/2, N=8

0110: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/4, N=6

0111:  $f_{SAMPLING} = f_{DTS}/4$ , N=8

1000:  $f_{SAMPLING} = f_{DTS}/8$ , N=6

1001: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/8, N=8

1010: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/16, N=5

1011: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/16, N=6

1100: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/16, N=8

1101: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/32, N=5 1110: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/32, N=6

1110. ISAMPLING-1015/02, 11-0

1111:  $f_{SAMPLING} = f_{DTS}/32$ , N=8

### Bits 3:2 IC1PSC: Input capture 1 prescaler

This bit-field defines the ratio of the prescaler acting on CC1 input (IC1).

The prescaler is reset as soon as CC1E='0' (TIMx\_CCER register).

00: no prescaler, capture is done each time an edge is detected on the capture input

01: capture is done once every 2 events

10: capture is done once every 4 events

11: capture is done once every 8 events

Bits 1:0 CC1S: Capture/Compare 1 Selection

This bit-field defines the direction of the channel (input/output) as well as the used input.

00: CC1 channel is configured as output

01: CC1 channel is configured as input, IC1 is mapped on TI1

10: CC1 channel is configured as input, IC1 is mapped on TI2

11: CC1 channel is configured as input, IC1 is mapped on TRC. This mode is working only if an internal trigger input is selected through TS bit (TIMx\_SMCR register)

Note: CC1S bits are writable only when the channel is OFF (CC1E = '0' in TIMX CCER).

# 13.4.8 TIM1&TIM8 capture/compare mode register 2 (TIMx\_CCMR2)

Address offset: 0x1C Reset value: 0x0000

Refer to the above CCMR1 register description.

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
OC4 CE	(	OC4M[2:0	]	OC4 PE	OC4 FE	CC48	OC3 CE.		(	OC3M[2:0	]	OC3 PE	OC3 FE	CC38	S[1:0]
	IC4F	[3:0]		IC4PS	C[1:0]		_		IC3F	[3:0]		IC3PS	SC[1:0]		
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

## **Output compare mode**

Bit 15 **OC4CE:** Output compare 4 clear enable

Bits 14:12 OC4M: Output compare 4 mode

Bit 11 OC4PE: Output compare 4 preload enable

Bit 10 OC4FE: Output compare 4 fast enable

Bits 9:8 CC4S: Capture/Compare 4 selection

This bit-field defines the direction of the channel (input/output) as well as the used input.

00: CC4 channel is configured as output

01: CC4 channel is configured as input, IC4 is mapped on TI4

10: CC4 channel is configured as input, IC4 is mapped on TI3

11: CC4 channel is configured as input, IC4 is mapped on TRC. This mode is working only if an internal trigger input is selected through TS bit (TIMx\_SMCR register)

Note: CC4S bits are writable only when the channel is OFF (CC4E = '0' in TIMx\_CCER).

Bit 7 OC3CE: Output compare 3 clear enable

Bits 6:4 OC3M: Output compare 3 mode

Bit 3 OC3PE: Output compare 3 preload enable

Bit 2 OC3FE: Output compare 3 fast enable

Bits 1:0 CC3S: Capture/Compare 3 selection

This bit-field defines the direction of the channel (input/output) as well as the used input.

00: CC3 channel is configured as output

01: CC3 channel is configured as input, IC3 is mapped on TI3

10: CC3 channel is configured as input, IC3 is mapped on TI4

11: CC3 channel is configured as input, IC3 is mapped on TRC. This mode is working only if an internal trigger input is selected through TS bit (TIMx\_SMCR register)

Note: CC3S bits are writable only when the channel is OFF (CC3E = '0' in TIMx\_CCER).

# Input capture mode

Bits 15:12 IC4F: Input capture 4 filter

Bits 11:10 IC4PSC: Input capture 4 prescaler

Bits 9:8 CC4S: Capture/Compare 4 selection

This bit-field defines the direction of the channel (input/output) as well as the used input.

00: CC4 channel is configured as output

01: CC4 channel is configured as input, IC4 is mapped on TI4

10: CC4 channel is configured as input, IC4 is mapped on TI3

11: CC4 channel is configured as input, IC4 is mapped on TRC. This mode is working only if an internal trigger input is selected through TS bit (TIMx\_SMCR register)

Note: CC4S bits are writable only when the channel is OFF (CC4E = '0' in TIMx\_CCER).

Bits 7:4 IC3F: Input capture 3 filter

Bits 3:2 IC3PSC: Input capture 3 prescaler

Bits 1:0 CC3S: Capture/compare 3 selection

This bit-field defines the direction of the channel (input/output) as well as the used input.

00: CC3 channel is configured as output

01: CC3 channel is configured as input, IC3 is mapped on TI3

10: CC3 channel is configured as input, IC3 is mapped on TI4

11: CC3 channel is configured as input, IC3 is mapped on TRC. This mode is working only if an internal trigger input is selected through TS bit (TIMx\_SMCR register)

Note: CC3S bits are writable only when the channel is OFF (CC3E = '0' in TIMx\_CCER).

# 13.4.9 TIM1&TIM8 capture/compare enable register (TIMx\_CCER)

Address offset: 0x20 Reset value: 0x0000

15	14	13	12	11	10	9	8	1	ь	5	4	3	2	ı	U
Reser	wod	CC4P	CC4E	CC3NP	CC3NE	CC3P	CC3E	CC2NP	CC2NE	CC2P	CC2E	CC1NP	CC1NE	CC1P	CC1E
пезеі	veu	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 15:14 Reserved, must be kept at reset value.

Bit 13 CC4P: Capture/Compare 4 output polarity

refer to CC1P description

Bit 12 CC4E: Capture/Compare 4 output enable

refer to CC1E description

Bit 11 CC3NP: Capture/Compare 3 complementary output polarity

refer to CC1NP description

Bit 10 CC3NE: Capture/Compare 3 complementary output enable

refer to CC1NE description

Bit 9 CC3P: Capture/Compare 3 output polarity

refer to CC1P description

Bit 8 CC3E: Capture/Compare 3 output enable

refer to CC1E description

- Bit 7 **CC2NP**: Capture/Compare 2 complementary output polarity refer to CC1NP description
- Bit 6 **CC2NE**: Capture/Compare 2 complementary output enable refer to CC1NE description
- Bit 5 **CC2P**: Capture/Compare 2 output polarity refer to CC1P description
- Bit 4 **CC2E**: Capture/Compare 2 output enable refer to CC1E description
- Bit 3 CC1NP: Capture/Compare 1 complementary output polarity

### CC1 channel configured as output:

0: OC1N active high.

1: OC1N active low.

### CC1 channel configured as input:

This bit is used in conjunction with CC1P to define the polarity of TI1FP1 and TI2FP1. Refer to CC1P description.

Note: On channels having a complementary output, this bit is preloaded. If the CCPC bit is set in the TIMx\_CR2 register then the CC1NP active bit takes the new value from the preloaded bit only when a Commutation event is generated.

Note: This bit is not writable as soon as LOCK level 2 or 3 has been programmed (LOCK bits in TIMx\_BDTR register) and CC1S="00" (the channel is configured in output).

- Bit 2 **CC1NE**: Capture/Compare 1 complementary output enable
  - 0: Off OC1N is not active. OC1N level is then function of MOE, OSSI, OSSR, OIS1, OIS1N and CC1E bits.
  - 1: On OC1N signal is output on the corresponding output pin depending on MOE, OSSI, OSSR, OIS1, OIS1N and CC1E bits.

Note: On channels having a complementary output, this bit is preloaded. If the CCPC bit is set in the TIMx\_CR2 register then the CC1NE active bit takes the new value from the preloaded bit only when a Commutation event is generated.



### Bit 1 CC1P: Capture/Compare 1 output polarity

### CC1 channel configured as output:

0: OC1 active high

1: OC1 active low

### CC1 channel configured as input:

CC1NP/CC1P bits select the active polarity of TI1FP1 and TI2FP1 for trigger or capture operations.

00: non-inverted/rising edge

The circuit is sensitive to TIxFP1 rising edge (capture or trigger operations in reset, external clock or trigger mode), TIxFP1 is not inverted (trigger operation in gated mode or encoder mode).

01: inverted/falling edge

The circuit is sensitive to TIxFP1 falling edge (capture or trigger operations in reset, external clock or trigger mode), TIxFP1 is inverted (trigger operation in gated mode or encoder mode).

10: reserved, do not use this configuration.

11: non-inverted/both edges

The circuit is sensitive to both TIxFP1 rising and falling edges (capture or trigger operations in reset, external clock or trigger mode), TIxFP1 is not inverted (trigger operation in gated mode). This configuration must not be used in encoder mode.

Note: On channels having a complementary output, this bit is preloaded. If the CCPC bit is set in the TIMx\_CR2 register then the CC1P active bit takes the new value from the preloaded bit only when a Commutation event is generated.

Note: This bit is not writable as soon as LOCK level 2 or 3 has been programmed (LOCK bits in TIMx BDTR register).

### Bit 0 CC1E: Capture/Compare 1 output enable

### CC1 channel configured as output:

0: Off - OC1 is not active. OC1 level is then function of MOE, OSSI, OSSR, OIS1, OIS1N and CC1NE bits.

1: On - OC1 signal is output on the corresponding output pin depending on MOE, OSSI, OSSR, OIS1, OIS1N and CC1NE bits.

### CC1 channel configured as input:

This bit determines if a capture of the counter value can actually be done into the input capture/compare register 1 (TIMx\_CCR1) or not.

0: Capture disabled.

1: Capture enabled.

Note: On channels having a complementary output, this bit is preloaded. If the CCPC bit is set in the TIMx\_CR2 register then the CC1E active bit takes the new value from the preloaded bit only when a Commutation event is generated.

Table 59. Output control bits for complementary OCx and OCxN channels with break feature

		Control I	oits		Outp	ut states <sup>(1)</sup>
MOE bit	OSSI bit	OSSR bit	CCxE bit	CCxNE bit	OCx output state	OCxN output state
		0	0	0	Output Disabled (not driven by the timer) OCx=0, OCx_EN=0	Output Disabled (not driven by the timer) OCxN=0, OCxN_EN=0
		0	0	1	Output Disabled (not driven by the timer) OCx=0, OCx_EN=0	OCxREF + Polarity OCxN=OCxREF xor CCxNP, OCxN_EN=1
		0	1	0	OCxREF + Polarity OCx=OCxREF xor CCxP, OCx_EN=1	Output Disabled (not driven by the timer) OCxN=0, OCxN_EN=0
1	X	0	1	1	OCREF + Polarity + dead-time OCx_EN=1	Complementary to OCREF (not OCREF) + Polarity + dead-time OCxN_EN=1
'	^	1	0	0	Output Disabled (not driven by the timer) OCx=CCxP, OCx_EN=0	Output Disabled (not driven by the timer) OCxN=CCxNP, OCxN_EN=0
		1	0	1	Off-State (output enabled with inactive state) OCx=CCxP, OCx_EN=1	OCxREF + Polarity OCxN=OCxREF xor CCxNP, OCxN_EN=1
		1	1	0	OCxREF + Polarity OCx=OCxREF xor CCxP, OCx_EN=1	Off-State (output enabled with inactive state) OCxN=CCxNP, OCxN_EN=1
		1	1	1	OCREF + Polarity + dead-time OCx_EN=1	Complementary to OCREF (not OCREF) + Polarity + dead-time OCxN_EN=1
	0		0	0	Output Disabled (not driven by the timer) OCx=CCxP, OCx_EN=0	Output Disabled (not driven by the timer) OCxN=CCxNP, OCxN_EN=0
	0		0	1	Output Disabled (not driven by t	he timer)
	0		1	0	Asynchronously: OCx=CCxP, O OCxN_EN=0	Cx_EN=0, OCxN=CCxNP,
0	0	X	1	1	Then if the clock is present: OC	x=OISx and OCxN=OISxN after a and OISxN do not correspond to OCX
	1	^	0	0	Output Disabled (not driven by the timer) OCx=CCxP, OCx_EN=0	Output Disabled (not driven by the timer) OCxN=CCxNP, OCxN_EN=0
	1		0	1	Off-State (output enabled with in	nactive state)
	1		1	0	Asynchronously: OCx=CCxP, O OCxN_EN=1	•
	1		1	1	Then if the clock is present: OC	x=OISx and OCxN=OISxN after a and OISxN do not correspond to OCX

1. When both outputs of a channel are not used (CCxE = CCxNE = 0), the OISx, OISxN, CCxP and CCxNP bits must be kept cleared.

Note:

The state of the external I/O pins connected to the complementary OCx and OCxN channels depends on the OCx and OCxN channel state and the GPIO registers.

# 13.4.10 TIM1&TIM8 counter (TIMx\_CNT)

Address offset: 0x24 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	CNT[15:0]														
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 15:0 CNT[15:0]: Counter value

# 13.4.11 TIM1&TIM8 prescaler (TIMx\_PSC)

Address offset: 0x28 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
PSC[15:0]															
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

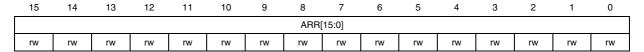
Bits 15:0 PSC[15:0]: Prescaler value

The counter clock frequency (CK\_CNT) is equal to  $f_{CK\_PSC}$  / (PSC[15:0] + 1).

PSC contains the value to be loaded in the active prescaler register at each update event (including when the counter is cleared through UG bit of TIMx\_EGR register or through trigger controller when configured in "reset mode").

# 13.4.12 TIM1&TIM8 auto-reload register (TIMx\_ARR)

Address offset: 0x2C Reset value: 0x0000



Bits 15:0 ARR[15:0]: Auto-reload value

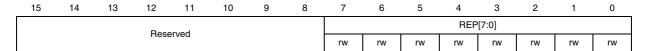
ARR is the value to be loaded in the actual auto-reload register.

Refer to Section 13.3.1: Time-base unit on page 301 for more details about ARR update and behavior.

The counter is blocked while the auto-reload value is null.

# 13.4.13 TIM1&TIM8 repetition counter register (TIMx\_RCR)

Address offset: 0x30 Reset value: 0x0000



Bits 15:8 Reserved, must be kept at reset value.

## Bits 7:0 REP[7:0]: Repetition counter value

These bits allow the user to set-up the update rate of the compare registers (i.e. periodic transfers from preload to active registers) when preload registers are enable, as well as the update interrupt generation rate, if this interrupt is enable.

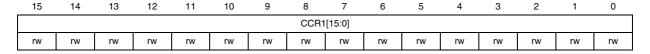
Each time the REP\_CNT related downcounter reaches zero, an update event is generated and it restarts counting from REP value. As REP\_CNT is reloaded with REP value only at the repetition update event U\_RC, any write to the TIMx\_RCR register is not taken in account until the next repetition update event.

It means in PWM mode (REP+1) corresponds to:

- the number of PWM periods in edge-aligned mode
- the number of half PWM period in center-aligned mode.

# 13.4.14 TIM1&TIM8 capture/compare register 1 (TIMx\_CCR1)

Address offset: 0x34 Reset value: 0x0000



Bits 15:0 CCR1[15:0]: Capture/Compare 1 value

### If channel CC1 is configured as output:

CCR1 is the value to be loaded in the actual capture/compare 1 register (preload value). It is loaded permanently if the preload feature is not selected in the TIMx\_CCMR1 register (bit OC1PE). Else the preload value is copied in the active capture/compare 1 register when an update event occurs.

The active capture/compare register contains the value to be compared to the counter TIMx\_CNT and signaled on OC1 output.

### If channel CC1 is configured as input:

CCR1 is the counter value transferred by the last input capture 1 event (IC1).

# 13.4.15 TIM1&TIM8 capture/compare register 2 (TIMx\_CCR2)

Address offset: 0x38 Reset value: 0x0000

_	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	CCR2[15:0]															
	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 15:0 CCR2[15:0]: Capture/Compare 2 value

#### If channel CC2 is configured as output:

CCR2 is the value to be loaded in the actual capture/compare 2 register (preload value). It is loaded permanently if the preload feature is not selected in the TIMx\_CCMR2 register (bit OC2PE). Else the preload value is copied in the active capture/compare 2 register when an update event occurs.

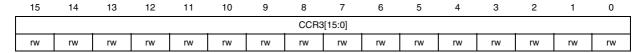
The active capture/compare register contains the value to be compared to the counter TIMx CNT and signalled on OC2 output.

### If channel CC2 is configured as input:

CCR2 is the counter value transferred by the last input capture 2 event (IC2).

# 13.4.16 TIM1&TIM8 capture/compare register 3 (TIMx\_CCR3)

Address offset: 0x3C Reset value: 0x0000



Bits 15:0 CCR3[15:0]: Capture/Compare value

#### If channel CC3 is configured as output:

CCR3 is the value to be loaded in the actual capture/compare 3 register (preload value). It is loaded permanently if the preload feature is not selected in the TIMx\_CCMR3 register (bit OC3PE). Else the preload value is copied in the active capture/compare 3 register when an update event occurs.

The active capture/compare register contains the value to be compared to the counter TIMx\_CNT and signalled on OC3 output.

## If channel CC3 is configured as input:

CCR3 is the counter value transferred by the last input capture 3 event (IC3).

# 13.4.17 TIM1&TIM8 capture/compare register 4 (TIMx CCR4)

Address offset: 0x40 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	CCR4[15:0]														
rw r															

Bits 15:0 CCR4[15:0]: Capture/Compare value

#### If channel CC4 is configured as output:

CCR4 is the value to be loaded in the actual capture/compare 4 register (preload value). It is loaded permanently if the preload feature is not selected in the TIMx\_CCMR4 register (bit OC4PE). Else the preload value is copied in the active capture/compare 4 register when an update event occurs.

The active capture/compare register contains the value to be compared to the counter TIMx CNT and signalled on OC4 output.

#### If channel CC4 is configured as input:

CCR4 is the counter value transferred by the last input capture 4 event (IC4).

# 13.4.18 TIM1&TIM8 break and dead-time register (TIMx\_BDTR)

Address offset: 0x44 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
MOE	AOE	BKP	BKE	OSSR	OSSI	LOCI	K[1:0]				DTG	[7:0]	_	_	
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Note:

As the bits AOE, BKP, BKE, OSSI, OSSR and DTG[7:0] can be write-locked depending on the LOCK configuration, it can be necessary to configure all of them during the first write access to the TIMx\_BDTR register.

#### Bit 15 MOE: Main output enable

This bit is cleared asynchronously by hardware as soon as the break input is active. It is set by software or automatically depending on the AOE bit. It is acting only on the channels which are configured in output.

- 0: OC and OCN outputs are disabled or forced to idle state.
- 1: OC and OCN outputs are enabled if their respective enable bits are set (CCxE, CCxNE in  $TIMx\_CCER$  register).

See OC/OCN enable description for more details (Section 13.4.9: TIM1&TIM8 capture/compare enable register (TIMx\_CCER) on page 354).

### Bit 14 AOE: Automatic output enable

- 0: MOE can be set only by software
- 1: MOE can be set by software or automatically at the next update event (if the break input is not be active)

Note: This bit can not be modified as long as LOCK level 1 has been programmed (LOCK bits in TIMx\_BDTR register).

#### Bit 13 BKP: Break polarity

0: Break input BRK is active low

1: Break input BRK is active high

Note: This bit can not be modified as long as LOCK level 1 has been programmed (LOCK bits in TIMx\_BDTR register).

Note: Any write operation to this bit takes a delay of 1 APB clock cycle to become effective.

#### Bit 12 BKE: Break enable

0: Break inputs (BRK and CSS clock failure event) disabled

1; Break inputs (BRK and CSS clock failure event) enabled

Note: This bit cannot be modified when LOCK level 1 has been programmed (LOCK bits in TIMx\_BDTR register).

Note: Any write operation to this bit takes a delay of 1 APB clock cycle to become effective.

#### Bit 11 OSSR: Off-state selection for Run mode

This bit is used when MOE=1 on channels having a complementary output which are configured as outputs. OSSR is not implemented if no complementary output is implemented in the timer.

See OC/OCN enable description for more details (Section 13.4.9: TIM1&TIM8 capture/compare enable register (TIMx CCER) on page 354).

0: When inactive, OC/OCN outputs are disabled (OC/OCN enable output signal=0).

1: When inactive, OC/OCN outputs are enabled with their inactive level as soon as CCxE=1 or CCxNE=1. Then, OC/OCN enable output signal=1

Note: This bit can not be modified as soon as the LOCK level 2 has been programmed (LOCK bits in TIMx\_BDTR register).

#### Bit 10 OSSI: Off-state selection for Idle mode

This bit is used when MOE=0 on channels configured as outputs.

See OC/OCN enable description for more details (Section 13.4.9: TIM1&TIM8 capture/compare enable register (TIMx\_CCER) on page 354).

0: When inactive, OC/OCN outputs are disabled (OC/OCN enable output signal=0).

1: When inactive, OC/OCN outputs are forced first with their idle level as soon as CCxE=1 or CCxNE=1. OC/OCN enable output signal=1)

Note: This bit can not be modified as soon as the LOCK level 2 has been programmed (LOCK bits in TIMx\_BDTR register).

#### Bits 9:8 LOCK[1:0]: Lock configuration

These bits offer a write protection against software errors.

00: LOCK OFF - No bit is write protected.

01: LOCK Level 1 = DTG bits in TIMx\_BDTR register, OISx and OISxN bits in TIMx\_CR2 register and BKE/BKP/AOE bits in TIMx\_BDTR register can no longer be written.

10: LOCK Level 2 = LOCK Level 1 + CC Polarity bits (CCxP/CCxNP bits in TIMx\_CCER register, as long as the related channel is configured in output through the CCxS bits) as well as OSSR and OSSI bits can no longer be written.

11: LOCK Level 3 = LOCK Level 2 + CC Control bits (OCxM and OCxPE bits in

TIMx\_CCMRx registers, as long as the related channel is configured in output through the CCxS bits) can no longer be written.

Note: The LOCK bits can be written only once after the reset. Once the TIMx\_BDTR register has been written, their content is frozen until the next reset.

#### Bits 7:0 DTG[7:0]: Dead-time generator setup

This bit-field defines the duration of the dead-time inserted between the complementary outputs. DT correspond to this duration.

 $DTG[7:5]=0xx \Rightarrow DT=DTG[7:0]x t_{dtg}$  with  $t_{dtg}=t_{DTS}$ .

 $DTG[7:5]=10x \Rightarrow DT=(64+DTG[5:0])xt_{dtg}$  with  $T_{dtg}=2xt_{DTS}$ .

 $DTG[7:5]=110 \Rightarrow DT=(32+DTG[4:0])xt_{dtg}$  with  $T_{dtg}=8xt_{DTS}$ .

 $DTG[7:5]=111 => DT=(32+DTG[4:0])xt_{dtg}$  with  $T_{dtg}=16xt_{DTS}$ .

Example if T<sub>DTS</sub>=125ns (8MHz), dead-time possible values are:

0 to 15875 ns by 125 ns steps,

16 us to 31750 ns by 250 ns steps,

32 us to 63us by 1 us steps,

64 us to 126 us by 2 us steps

Note: This bit-field can not be modified as long as LOCK level 1, 2 or 3 has been programmed (LOCK bits in TIMx\_BDTR register).

# 13.4.19 TIM1&TIM8 DMA control register (TIMx\_DCR)

Address offset: 0x48
Reset value: 0x0000

	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
	Reserved			DBL[4:0]					Reserved			DBA[4:0]					
				rw	rw	rw	rw	rw		neserveu		rw	rw	rw	rw	rw	

Bits 15:13 Reserved, must be kept at reset value.

#### Bits 12:8 DBL[4:0]: DMA burst length

This 5-bit vector defines the number of DMA transfers (the timer detects a burst transfer when a read or a write access to the TIMx\_DMAR register address is performed).

the TIMx\_DMAR address)

00000: 1 transfer

00001: 2 transfers

00010: 3 transfers

•••

10001: 18 transfers

Bits 7:5 Reserved, must be kept at reset value.

#### Bits 4:0 DBA[4:0]: DMA base address

This 5-bits vector defines the base-address for DMA transfers (when read/write access are done through the TIMx\_DMAR address). DBA is defined as an offset starting from the address of the TIMx\_CR1 register.

Example:

00000: TIMx\_CR1,

00001: TIMx\_CR2,

00010: TIMx\_SMCR,

...

**Example:** Let us consider the following transfer: DBL = 7 transfers and DBA = TIMx\_CR1. In this case the transfer is done to/from 7 registers starting from the TIMx\_CR1 address.

# 13.4.20 TIM1&TIM8 DMA address for full transfer (TIMx\_DMAR)

Address offset: 0x4C Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	DMAB[15:0]														
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 15:0 DMAB[15:0]: DMA register for burst accesses

A read or write operation to the DMAR register accesses the register located at the address (TIMx\_CR1 address) + (DBA + DMA index) x 4

where TIMx\_CR1 address is the address of the control register 1, DBA is the DMA base address configured in TIMx\_DCR register, DMA index is automatically controlled by the DMA transfer, and ranges from 0 to DBL (DBL configured in TIMx\_DCR).

### **Example of how to use the DMA burst feature**

In this example the timer DMA burst feature is used to update the contents of the CCRx registers (x = 2, 3, 4) with the DMA transferring half words into the CCRx registers.

This is done in the following steps:

- 1. Configure the corresponding DMA channel as follows:
  - DMA channel peripheral address is the DMAR register address
  - DMA channel memory address is the address of the buffer in the RAM containing the data to be transferred by DMA into CCRx registers.
  - Number of data to transfer = 3 (See note below).
  - Circular mode disabled.
- Configure the DCR register by configuring the DBA and DBL bit fields as follows: DBL = 3 transfers, DBA = 0xE.
- 3. Enable the TIMx update DMA request (set the UDE bit in the DIER register).
- 4. Enable TIMx
- Enable the DMA channel

Note:

This example is for the case where every CCRx register to be updated once. If every CCRx register is to be updated twice for example, the number of data to transfer should be 6. Let's take the example of a buffer in the RAM containing data1, data2, data3, data4, data5 and data6. The data is transferred to the CCRx registers as follows: on the first update DMA request, data1 is transferred to CCR2, data2 is transferred to CCR3, data3 is transferred to CCR4 and on the second update DMA request, data4 is transferred to CCR2, data5 is transferred to CCR3 and data6 is transferred to CCR4.

# 13.4.21 TIM1&TIM8 register map

TIM1&TIM8 registers are mapped as 16-bit addressable registers as described in the table below:

Table 60. TIM1&TIM8 register map and reset values

Offset	Register	100 20 20 20 20 30 30 30 30 30 30 30 30 30 30 30 30 30	12 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1					
0x00	TIMx_CR1  Reset value	Reserved	CKD					
0x04	TIMx_CR2	Reserved	O O O O O O O O O O O O O O O O O O O					
	Hesel value							
0x08	TIMx_SMCR Reset value	Reserved	L U U ETPS [1:0] ETF[3:0]					
0x0C	TIMx_DIER	Reserved	TDE COMDE CC3DE CC3DE CC2DE CC1DE UDE BIE TIE COMIE CC3IE					
	Reset value		0 0 0 0 0 0 0 0 0 0 0 0 0 0 0					
0x10	TIMx_SR  Reset value	Reserved	CC40F   CC20F   CC20					
0x14	TIMx_EGR	Reserved	COM CC26 CC26 CC26 CC16 CC16 CC16 CC16 CC16					
	TIMx_CCMR1							
	Output Compare mode  Reset value	Reserved	U					
0x18	TIMx_CCMR1 Input Capture mode Reset value	Reserved	IC2					
2.10	TIMx_CCMR2 Output Compare mode Reset value	Reserved	OC4M H H CC4S S [2:0] O S [1:0] S [2:0] O S [1:0] O O O O O O O O O O O O O O O O O O O					
0x1C	TIMx_CCMR2 Input Capture mode Reset value	Reserved	IC4F[3:0]					
0x20	TIMx_CCER  Reset value	Reserved	CC4P CC3NP CC3NP CC3NP CC2NP C					
0x24	TIMx_CNT	Reserved	CNT[15:0]					
	Reset value		0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0					
0x28	TIMx_PSC	Reserved	PSC[15:0]					
	Reset value		0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0					
0x2C	TIMx_ARR	Reserved	ARR[15:0]					
	Reset value		0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0					
0x30	TIMx_RCR	Reserved	REP[7:0]					
	Reset value		0 0 0 0 0 0 0 0					

Table 60. TIM1&TIM8 register map and reset values (continued)

Offset	Register	33 33 34 35 38 38 39 39 31 4 38 39 31 4 39 31 4 39 31 4 39 31 4 39 31 4 31 4	0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 0						
0x34	TIMx_CCR1	Reserved	CCR1[15:0]						
	Reset value		0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0						
0x38	TIMx_CCR2	Reserved	CCR2[15:0]						
	Reset value		0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0						
0x3C	TIMx_CCR3	Reserved	CCR3[15:0]						
	Reset value		0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0						
0x40	TIMx_CCR4	Reserved	CCR4[15:0]						
	Reset value		0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0						
0x44	TIMx_BDTR	Reserved	M M M M M M M M M M M M M M M M M M M						
	Reset value		0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0						
0x48	TIMx_DCR	Reserved	DBL[4:0] Reserved DBA[4:0]						
	Reset value		0 0 0 0 0 0 0 0						
0x4C	TIMx_DMAR	Reserved	DMAB[15:0]						
	Reset value		0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0						

Refer to *Table 2 on page 50* for the register boundary addresses.

# 14 General-purpose timers (TIM2 to TIM5)

This section applies to the whole STM32F4xx family, unless otherwise specified.

# 14.1 TIM2 to TIM5 introduction

The general-purpose timers consist of a 16-bit or 32-bit auto-reload counter driven by a programmable prescaler.

They may be used for a variety of purposes, including measuring the pulse lengths of input signals (*input capture*) or generating output waveforms (*output compare and PWM*).

Pulse lengths and waveform periods can be modulated from a few microseconds to several milliseconds using the timer prescaler and the RCC clock controller prescalers.

The timers are completely independent, and do not share any resources. They can be synchronized together as described in *Section 14.3.15*.

# 14.2 TIM2 to TIM5 main features

General-purpose TIMx timer features include:

- 16-bit (TIM3 and TIM4) or 32-bit (TIM2 and TIM5) up, down, up/down auto-reload counter.
- 16-bit programmable prescaler used to divide (also "on the fly") the counter clock frequency by any factor between 1 and 65536.
- Up to 4 independent channels for:
  - Input capture
  - Output compare
  - PWM generation (Edge- and Center-aligned modes)
  - One-pulse mode output
- Synchronization circuit to control the timer with external signals and to interconnect several timers.
- Interrupt/DMA generation on the following events:
  - Update: counter overflow/underflow, counter initialization (by software or internal/external trigger)
  - Trigger event (counter start, stop, initialization or count by internal/external trigger)
  - Input capture
  - Output compare
- Supports incremental (quadrature) encoder and hall-sensor circuitry for positioning purposes
- Trigger input for external clock or cycle-by-cycle current management

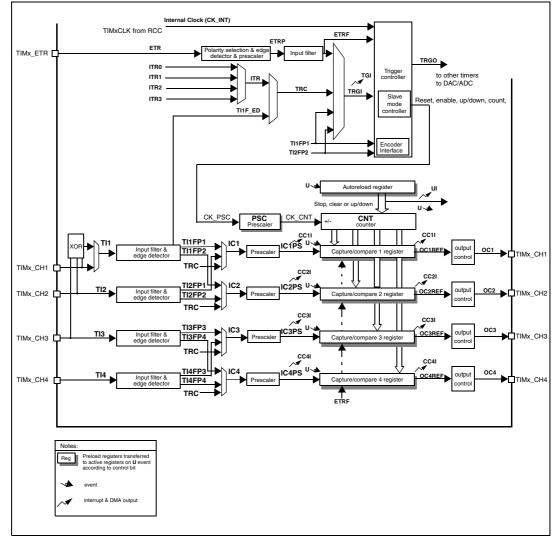


Figure 113. General-purpose timer block diagram

# 14.3 TIM2 to TIM5 functional description

### 14.3.1 Time-base unit

The main block of the programmable timer is a 16-bit/32-bit counter with its related autoreload register. The counter can count up. The counter clock can be divided by a prescaler.

The counter, the auto-reload register and the prescaler register can be written or read by software. This is true even when the counter is running.

The time-base unit includes:

- Counter Register (TIMx\_CNT)
- Prescaler Register (TIMx\_PSC):
- Auto-Reload Register (TIMx\_ARR)

The auto-reload register is preloaded. Writing to or reading from the auto-reload register accesses the preload register. The content of the preload register are transferred into the shadow register permanently or at each update event (UEV), depending on the auto-reload preload enable bit (ARPE) in TIMx\_CR1 register. The update event is sent when the counter reaches the overflow (or underflow when downcounting) and if the UDIS bit equals 0 in the TIMx\_CR1 register. It can also be generated by software. The generation of the update event is described in detail for each configuration.

The counter is clocked by the prescaler output CK\_CNT, which is enabled only when the counter enable bit (CEN) in TIMx\_CR1 register is set (refer also to the slave mode controller description to get more details on counter enabling).

Note that the actual counter enable signal CNT EN is set 1 clock cycle after CEN.

### **Prescaler description**

The prescaler can divide the counter clock frequency by any factor between 1 and 65536. It is based on a 16-bit counter controlled through a 16-bit/32-bit register (in the TIMx\_PSC register). It can be changed on the fly as this control register is buffered. The new prescaler ratio is taken into account at the next update event.

*Figure 114* and *Figure 115* give some examples of the counter behavior when the prescaler ratio is changed on the fly:

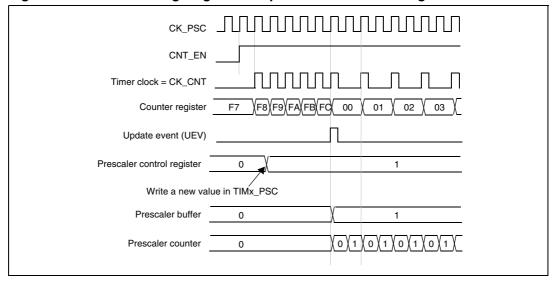


Figure 114. Counter timing diagram with prescaler division change from 1 to 2

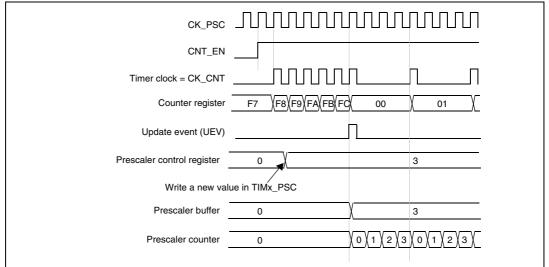


Figure 115. Counter timing diagram with prescaler division change from 1 to 4

### 14.3.2 Counter modes

### **Upcounting mode**

In upcounting mode, the counter counts from 0 to the auto-reload value (content of the TIMx ARR register), then restarts from 0 and generates a counter overflow event.

An Update event can be generated at each counter overflow or by setting the UG bit in the TIMx\_EGR register (by software or by using the slave mode controller).

The UEV event can be disabled by software by setting the UDIS bit in TIMx\_CR1 register. This is to avoid updating the shadow registers while writing new values in the preload registers. Then no update event occurs until the UDIS bit has been written to 0. However, the counter restarts from 0, as well as the counter of the prescaler (but the prescale rate does not change). In addition, if the URS bit (update request selection) in TIMx\_CR1 register is set, setting the UG bit generates an update event UEV but without setting the UIF flag (thus no interrupt or DMA request is sent). This is to avoid generating both update and capture interrupts when clearing the counter on the capture event.

When an update event occurs, all the registers are updated and the update flag (UIF bit in TIMx SR register) is set (depending on the URS bit):

- The buffer of the prescaler is reloaded with the preload value (content of the TIMx\_PSC register)
- The auto-reload shadow register is updated with the preload value (TIMx\_ARR)

The following figures show some examples of the counter behavior for different clock frequencies when TIMx\_ARR=0x36.

Figure 116. Counter timing diagram, internal clock divided by 1

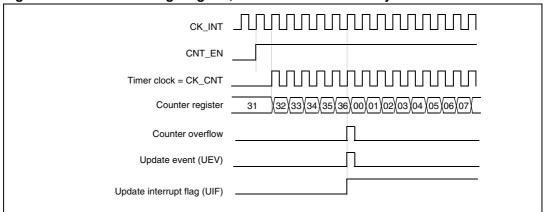


Figure 117. Counter timing diagram, internal clock divided by 2

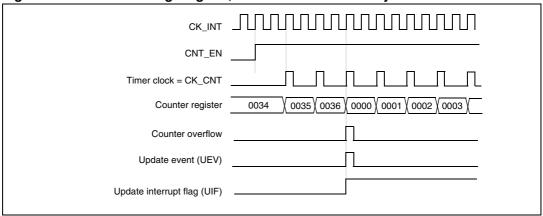


Figure 118. Counter timing diagram, internal clock divided by 4

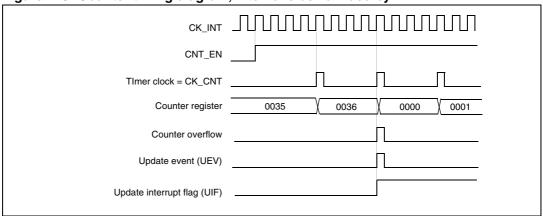


Figure 119. Counter timing diagram, internal clock divided by N

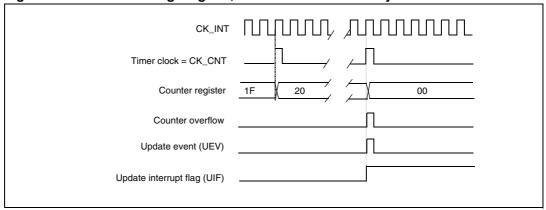
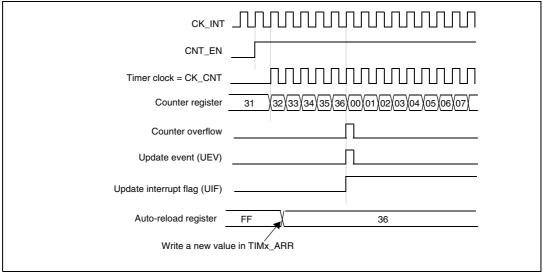


Figure 120. Counter timing diagram, Update event when ARPE=0 (TIMx\_ARR not preloaded)



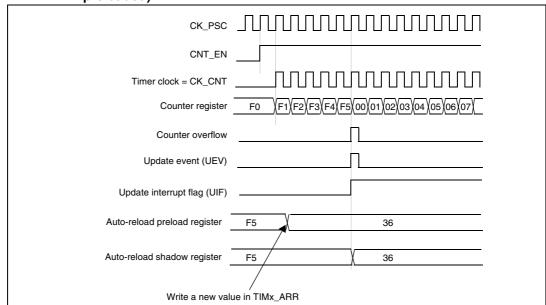


Figure 121. Counter timing diagram, Update event when ARPE=1 (TIMx\_ARR preloaded)

### **Downcounting mode**

In downcounting mode, the counter counts from the auto-reload value (content of the TIMx\_ARR register) down to 0, then restarts from the auto-reload value and generates a counter underflow event.

An Update event can be generate at each counter underflow or by setting the UG bit in the TIMx\_EGR register (by software or by using the slave mode controller)

The UEV update event can be disabled by software by setting the UDIS bit in TIMx\_CR1 register. This is to avoid updating the shadow registers while writing new values in the preload registers. Then no update event occurs until UDIS bit has been written to 0. However, the counter restarts from the current auto-reload value, whereas the counter of the prescaler restarts from 0 (but the prescale rate doesn't change).

In addition, if the URS bit (update request selection) in TIMx\_CR1 register is set, setting the UG bit generates an update event UEV but without setting the UIF flag (thus no interrupt or DMA request is sent). This is to avoid generating both update and capture interrupts when clearing the counter on the capture event.

When an update event occurs, all the registers are updated and the update flag (UIF bit in TIMx\_SR register) is set (depending on the URS bit):

- The buffer of the prescaler is reloaded with the preload value (content of the TIMx\_PSC register).
- The auto-reload active register is updated with the preload value (content of the TIMx\_ARR register). Note that the auto-reload is updated before the counter is reloaded, so that the next period is the expected one.

The following figures show some examples of the counter behavior for different clock frequencies when TIMx ARR=0x36.

Figure 122. Counter timing diagram, internal clock divided by 1

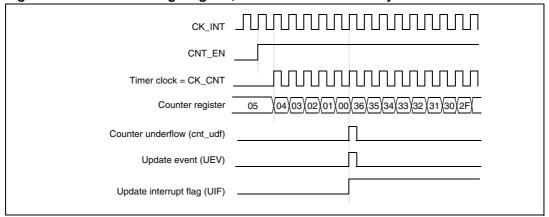


Figure 123. Counter timing diagram, internal clock divided by 2

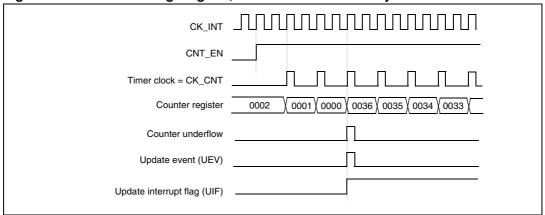
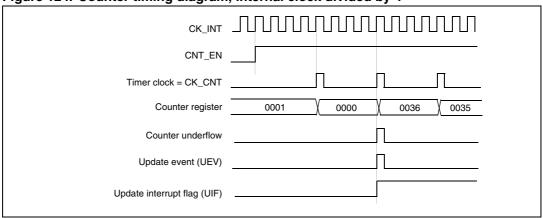


Figure 124. Counter timing diagram, internal clock divided by 4



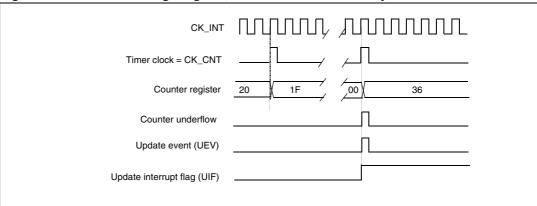
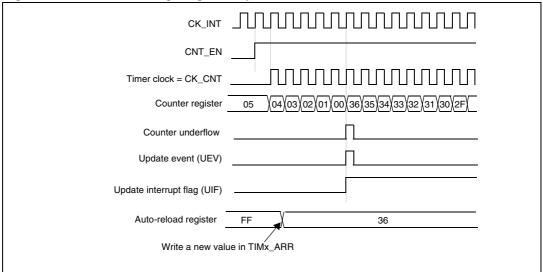


Figure 125. Counter timing diagram, internal clock divided by N





### Center-aligned mode (up/down counting)

In center-aligned mode, the counter counts from 0 to the auto-reload value (content of the  $TIMx\_ARR$  register) – 1, generates a counter overflow event, then counts from the auto-reload value down to 1 and generates a counter underflow event. Then it restarts counting from 0.

Center-aligned mode is active when the CMS bits in TIMx\_CR1 register are not equal to '00'. The Output compare interrupt flag of channels configured in output is set when: the counter counts down (Center aligned mode 1, CMS = "01"), the counter counts up (Center aligned mode 2, CMS = "10") the counter counts up and down (Center aligned mode 3, CMS = "11").

In this mode, the direction bit (DIR from TIMx\_CR1 register) cannot be written. It is updated by hardware and gives the current direction of the counter.

The update event can be generated at each counter overflow and at each counter underflow or by setting the UG bit in the TIMx\_EGR register (by software or by using the slave mode controller) also generates an update event. In this case, the counter restarts counting from 0, as well as the counter of the prescaler.

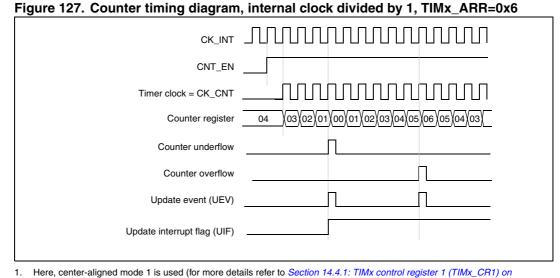
The UEV update event can be disabled by software by setting the UDIS bit in TIMx\_CR1 register. This is to avoid updating the shadow registers while writing new values in the preload registers. Then no update event occurs until the UDIS bit has been written to 0. However, the counter continues counting up and down, based on the current auto-reload value.

In addition, if the URS bit (update request selection) in TIMx\_CR1 register is set, setting the UG bit generates an update event UEV but without setting the UIF flag (thus no interrupt or DMA request is sent). This is to avoid generating both update and capture interrupt when clearing the counter on the capture event.

When an update event occurs, all the registers are updated and the update flag (UIF bit in TIMx\_SR register) is set (depending on the URS bit):

- The buffer of the prescaler is reloaded with the preload value (content of the TIMx\_PSC register).
- The auto-reload active register is updated with the preload value (content of the TIMx\_ARR register). Note that if the update source is a counter overflow, the auto-reload is updated before the counter is reloaded, so that the next period is the expected one (the counter is loaded with the new value).

The following figures show some examples of the counter behavior for different clock frequencies.



page 404).

Figure 128. Counter timing diagram, internal clock divided by 2

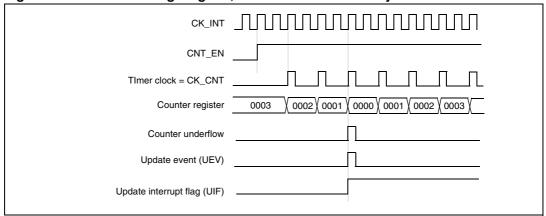
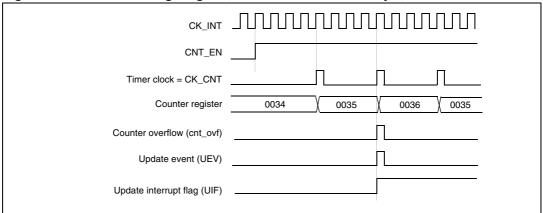
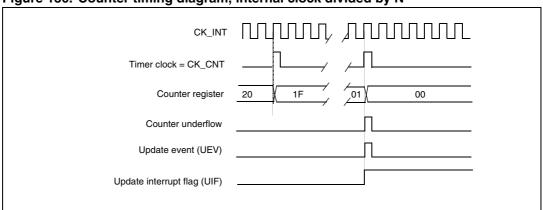


Figure 129. Counter timing diagram, internal clock divided by 4, TIMx\_ARR=0x36



1. Center-aligned mode 2 or 3 is used with an UIF on overflow.

Figure 130. Counter timing diagram, internal clock divided by N



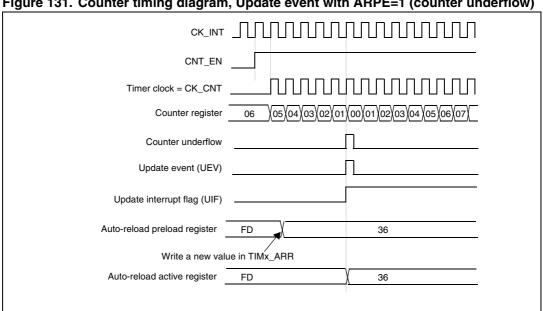
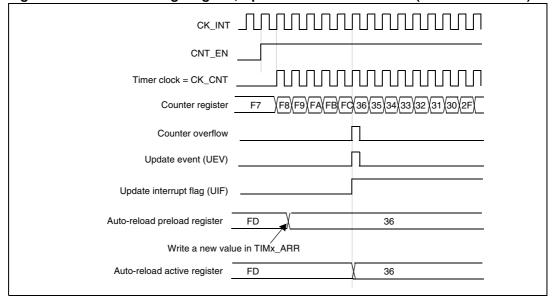


Figure 131. Counter timing diagram, Update event with ARPE=1 (counter underflow)





#### 14.3.3 Clock selection

The counter clock can be provided by the following clock sources:

- Internal clock (CK\_INT)
- External clock mode1: external input pin (Tlx)
- External clock mode2: external trigger input (ETR) available on TIM2, TIM3 and TIM4 only.
- Internal trigger inputs (ITRx): using one timer as prescaler for another timer, for example, you can configure Timer to act as a prescaler for Timer 2. Refer to : Using one timer as prescaler for another on page 398 for more details.

### Internal clock source (CK\_INT)

If the slave mode controller is disabled (SMS=000 in the TIMx\_SMCR register), then the CEN, DIR (in the TIMx\_CR1 register) and UG bits (in the TIMx\_EGR register) are actual control bits and can be changed only by software (except UG which remains cleared automatically). As soon as the CEN bit is written to 1, the prescaler is clocked by the internal clock CK\_INT.

*Figure 133* shows the behavior of the control circuit and the upcounter in normal mode, without prescaler.

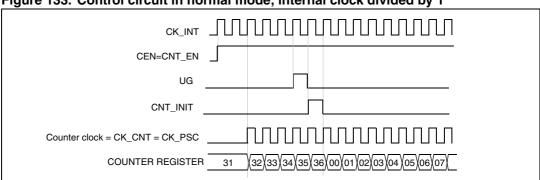


Figure 133. Control circuit in normal mode, internal clock divided by 1

### External clock source mode 1

This mode is selected when SMS=111 in the TIMx\_SMCR register. The counter can count at each rising or falling edge on a selected input.

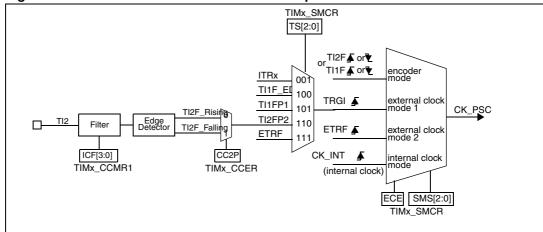


Figure 134. TI2 external clock connection example

For example, to configure the upcounter to count in response to a rising edge on the TI2 input, use the following procedure:

- Configure channel 2 to detect rising edges on the TI2 input by writing CC2S= '01 in the TIMx\_CCMR1 register.
- 2. Configure the input filter duration by writing the IC2F[3:0] bits in the TIMx\_CCMR1 register (if no filter is needed, keep IC2F=0000).

Note: The capture prescaler is not used for triggering, so you don't need to configure it.

- Select rising edge polarity by writing CC2P=0 and CC2NP=0 in the TIMx\_CCER register.
- 4. Configure the timer in external clock mode 1 by writing SMS=111 in the TIMx\_SMCR register.
- 5. Select TI2 as the input source by writing TS=110 in the TIMx\_SMCR register.
- 6. Enable the counter by writing CEN=1 in the TIMx\_CR1 register.

When a rising edge occurs on TI2, the counter counts once and the TIF flag is set.

The delay between the rising edge on TI2 and the actual clock of the counter is due to the resynchronization circuit on TI2 input.

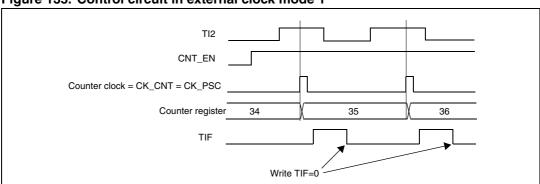


Figure 135. Control circuit in external clock mode 1

#### External clock source mode 2

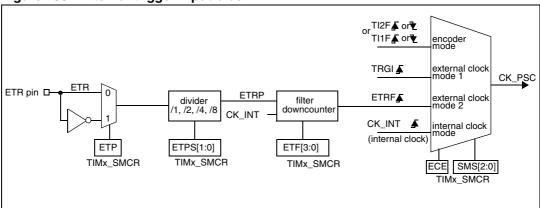
This mode is selected by writing ECE=1 in the TIMx\_SMCR register.

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The counter can count at each rising or falling edge on the external trigger input ETR.

Figure 136 gives an overview of the external trigger input block.

Figure 136. External trigger input block



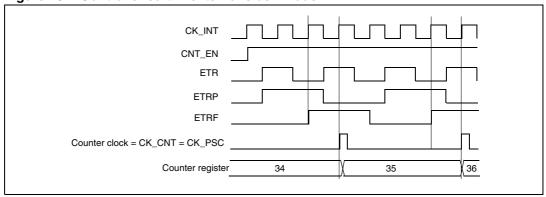
For example, to configure the upcounter to count each 2 rising edges on ETR, use the following procedure:

- 1. As no filter is needed in this example, write ETF[3:0]=0000 in the TIMx\_SMCR register.
- 2. Set the prescaler by writing ETPS[1:0]=01 in the TIMx\_SMCR register
- 3. Select rising edge detection on the ETR pin by writing ETP=0 in the TIMx\_SMCR register
- 4. Enable external clock mode 2 by writing ECE=1 in the TIMx\_SMCR register.
- 5. Enable the counter by writing CEN=1 in the TIMx\_CR1 register.

The counter counts once each 2 ETR rising edges.

The delay between the rising edge on ETR and the actual clock of the counter is due to the resynchronization circuit on the ETRP signal.

Figure 137. Control circuit in external clock mode 2



# 14.3.4 Capture/compare channels

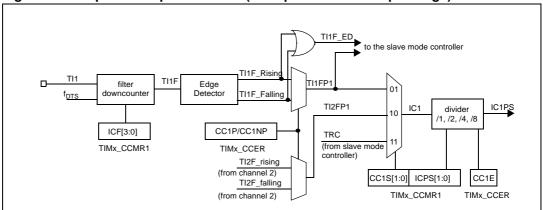
Each Capture/Compare channel is built around a capture/compare register (including a shadow register), a input stage for capture (with digital filter, multiplexing and prescaler) and an output stage (with comparator and output control).

The following figure gives an overview of one Capture/Compare channel.

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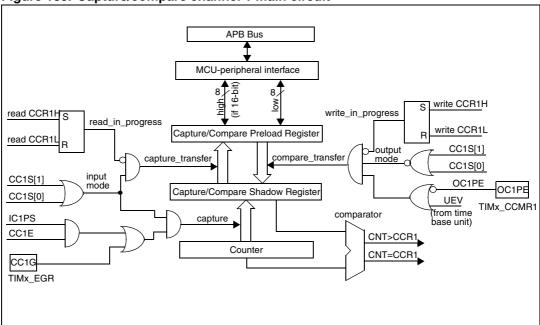
The input stage samples the corresponding TIx input to generate a filtered signal TIxF. Then, an edge detector with polarity selection generates a signal (TIxFPx) which can be used as trigger input by the slave mode controller or as the capture command. It is prescaled before the capture register (ICxPS).

Figure 138. Capture/compare channel (example: channel 1 input stage)



The output stage generates an intermediate waveform which is then used for reference: OCxRef (active high). The polarity acts at the end of the chain.

Figure 139. Capture/compare channel 1 main circuit



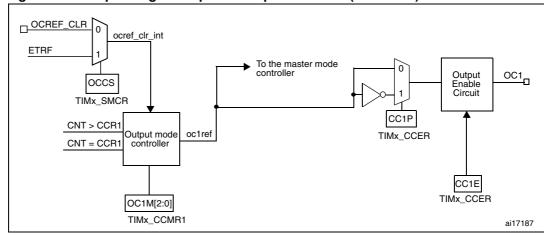


Figure 140. Output stage of capture/compare channel (channel 1)

The capture/compare block is made of one preload register and one shadow register. Write and read always access the preload register.

In capture mode, captures are actually done in the shadow register, which is copied into the preload register.

In compare mode, the content of the preload register is copied into the shadow register which is compared to the counter.

## 14.3.5 Input capture mode

In Input capture mode, the Capture/Compare Registers (TIMx\_CCRx) are used to latch the value of the counter after a transition detected by the corresponding ICx signal. When a capture occurs, the corresponding CCXIF flag (TIMx\_SR register) is set and an interrupt or a DMA request can be sent if they are enabled. If a capture occurs while the CCxIF flag was already high, then the over-capture flag CCxOF (TIMx\_SR register) is set. CCxIF can be cleared by software by writing it to 0 or by reading the captured data stored in the TIMx\_CCRx register. CCxOF is cleared when you write it to 0.

The following example shows how to capture the counter value in TIMx\_CCR1 when TI1 input rises. To do this, use the following procedure:

- Select the active input: TIMx\_CCR1 must be linked to the TI1 input, so write the CC1S bits to 01 in the TIMx\_CCMR1 register. As soon as CC1S becomes different from 00, the channel is configured in input and the TIMx\_CCR1 register becomes read-only.
- Program the input filter duration you need with respect to the signal you connect to the timer (by programming the ICxF bits in the TIMx\_CCMRx register if the input is one of the TIx inputs). Let's imagine that, when toggling, the input signal is not stable during at must 5 internal clock cycles. We must program a filter duration longer than these 5 clock cycles. We can validate a transition on TI1 when 8 consecutive samples with the

- new level have been detected (sampled at f<sub>DTS</sub> frequency). Then write IC1F bits to 0011 in the TIMx\_CCMR1 register.
- Select the edge of the active transition on the TI1 channel by writing the CC1P and CC1NP bits to 00 in the TIMx\_CCER register (rising edge in this case).
- Program the input prescaler. In our example, we wish the capture to be performed at each valid transition, so the prescaler is disabled (write IC1PS bits to 00 in the TIMx\_CCMR1 register).
- Enable capture from the counter into the capture register by setting the CC1E bit in the TIMx CCER register.
- If needed, enable the related interrupt request by setting the CC1IE bit in the TIMx\_DIER register, and/or the DMA request by setting the CC1DE bit in the TIMx\_DIER register.

When an input capture occurs:

- The TIMx\_CCR1 register gets the value of the counter on the active transition.
- CC1IF flag is set (interrupt flag). CC1OF is also set if at least two consecutive captures occurred whereas the flag was not cleared.
- An interrupt is generated depending on the CC1IE bit.
- A DMA request is generated depending on the CC1DE bit.

In order to handle the overcapture, it is recommended to read the data before the overcapture flag. This is to avoid missing an overcapture which could happen after reading the flag and before reading the data.

Note:

IC interrupt and/or DMA requests can be generated by software by setting the corresponding CCxG bit in the TIMx EGR register.

### 14.3.6 PWM input mode

This mode is a particular case of input capture mode. The procedure is the same except:

- Two ICx signals are mapped on the same TIx input.
- These 2 ICx signals are active on edges with opposite polarity.
- One of the two TIxFP signals is selected as trigger input and the slave mode controller is configured in reset mode.

For example, you can measure the period (in TIMx CCR1 register) and the duty cycle (in TIMx CCR2 register) of the PWM applied on TI1 using the following procedure (depending on CK INT frequency and prescaler value):

- Select the active input for TIMx CCR1: write the CC1S bits to 01 in the TIMx CCMR1 register (TI1 selected).
- Select the active polarity for TI1FP1 (used both for capture in TIMx CCR1 and counter clear): write the CC1P to '0' and the CC1NP bit to '0' (active on rising edge).
- Select the active input for TIMx CCR2: write the CC2S bits to 10 in the TIMx CCMR1 register (TI1 selected).
- Select the active polarity for TI1FP2 (used for capture in TIMx CCR2): write the CC2P bit to '1' and the CC2NP bit to '0' (active on falling edge).
- Select the valid trigger input: write the TS bits to 101 in the TIMx\_SMCR register (TI1FP1 selected).
- Configure the slave mode controller in reset mode: write the SMS bits to 100 in the TIMx\_SMCR register.
- Enable the captures: write the CC1E and CC2E bits to '1 in the TIMx\_CCER register.

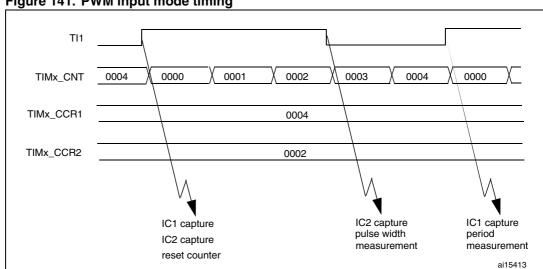


Figure 141. PWM input mode timing

#### 14.3.7 Forced output mode

In output mode (CCxS bits = 00 in the TIMx\_CCMRx register), each output compare signal (OCxREF and then OCx) can be forced to active or inactive level directly by software. independently of any comparison between the output compare register and the counter.

To force an output compare signal (ocxref/OCx) to its active level, you just need to write 101 in the OCxM bits in the corresponding TIMx\_CCMRx register. Thus ocxref is forced high (OCxREF is always active high) and OCx get opposite value to CCxP polarity bit.

e.g.: CCxP=0 (OCx active high) => OCx is forced to high level.

ocxref signal can be forced low by writing the OCxM bits to 100 in the TIMx\_CCMRx register.

Anyway, the comparison between the TIMx\_CCRx shadow register and the counter is still performed and allows the flag to be set. Interrupt and DMA requests can be sent accordingly. This is described in the Output Compare Mode section.

## 14.3.8 Output compare mode

This function is used to control an output waveform or indicating when a period of time has elapsed.

When a match is found between the capture/compare register and the counter, the output compare function:

- Assigns the corresponding output pin to a programmable value defined by the output compare mode (OCxM bits in the TIMx\_CCMRx register) and the output polarity (CCxP bit in the TIMx\_CCER register). The output pin can keep its level (OCXM=000), be set active (OCxM=001), be set inactive (OCxM=010) or can toggle (OCxM=011) on match.
- Sets a flag in the interrupt status register (CCxIF bit in the TIMx\_SR register).
- Generates an interrupt if the corresponding interrupt mask is set (CCXIE bit in the TIMx\_DIER register).
- Sends a DMA request if the corresponding enable bit is set (CCxDE bit in the TIMx\_DIER register, CCDS bit in the TIMx\_CR2 register for the DMA request selection).

The TIMx\_CCRx registers can be programmed with or without preload registers using the OCxPE bit in the TIMx\_CCMRx register.

In output compare mode, the update event UEV has no effect on ocxref and OCx output. The timing resolution is one count of the counter. Output compare mode can also be used to output a single pulse (in One-pulse mode).

#### Procedure:

- 1. Select the counter clock (internal, external, prescaler).
- 2. Write the desired data in the TIMx\_ARR and TIMx\_CCRx registers.
- 3. Set the CCxIE and/or CCxDE bits if an interrupt and/or a DMA request is to be generated.
- 4. Select the output mode. For example, you must write OCxM=011, OCxPE=0, CCxP=0 and CCxE=1 to toggle OCx output pin when CNT matches CCRx, CCRx preload is not used, OCx is enabled and active high.
- 5. Enable the counter by setting the CEN bit in the TIMx\_CR1 register.

The TIMx\_CCRx register can be updated at any time by software to control the output waveform, provided that the preload register is not enabled (OCxPE=0, else TIMx\_CCRx shadow register is updated only at the next update event UEV). An example is given in *Figure 142*.

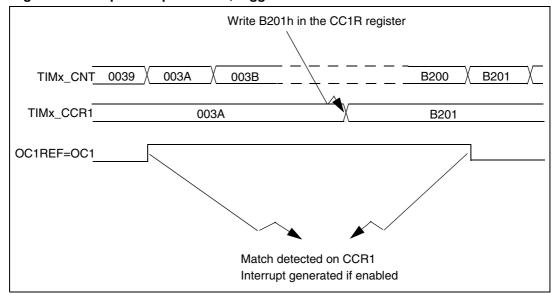


Figure 142. Output compare mode, toggle on OC1.

#### 14.3.9 PWM mode

Pulse width modulation mode allows you to generate a signal with a frequency determined by the value of the TIMx\_ARR register and a duty cycle determined by the value of the TIMx CCRx register.

The PWM mode can be selected independently on each channel (one PWM per OCx output) by writing 110 (PWM mode 1) or '111 (PWM mode 2) in the OCxM bits in the TIMx\_CCMRx register. You must enable the corresponding preload register by setting the OCxPE bit in the TIMx\_CCMRx register, and eventually the auto-reload preload register by setting the ARPE bit in the TIMx\_CR1 register.

As the preload registers are transferred to the shadow registers only when an update event occurs, before starting the counter, you have to initialize all the registers by setting the UG bit in the TIMx\_EGR register.

OCx polarity is software programmable using the CCxP bit in the TIMx\_CCER register. It can be programmed as active high or active low. OCx output is enabled by the CCxE bit in the TIMx\_CCER register. Refer to the TIMx\_CCERx register description for more details.

In PWM mode (1 or 2), TIMx\_CNT and TIMx\_CCRx are always compared to determine whether TIMx\_CCRx ≤ TIMx\_CNT or TIMx\_CNT ≤ TIMx\_CCRx (depending on the direction of the counter). However, to comply with the ETRF (OCREF can be cleared by an external event through the ETR signal until the next PWM period), the OCREF signal is asserted only:

- When the result of the comparison changes, or
- When the output compare mode (OCxM bits in TIMx\_CCMRx register) switches from the "frozen" configuration (no comparison, OCxM='000) to one of the PWM modes (OCxM='110 or '111).

This forces the PWM by software while the timer is running.

The timer is able to generate PWM in edge-aligned mode or center-aligned mode depending on the CMS bits in the TIMx\_CR1 register.

### PWM edge-aligned mode

Upcounting configuration

Upcounting is active when the DIR bit in the TIMx\_CR1 register is low. Refer to *Section : Upcounting mode on page 370*.

In the following example, we consider PWM mode 1. The reference PWM signal OCxREF is high as long as TIMx\_CNT <TIMx\_CCRx else it becomes low. If the compare value in TIMx\_CCRx is greater than the auto-reload value (in TIMx\_ARR) then OCxREF is held at '1. If the compare value is 0 then OCxREF is held at '0. *Figure 143* shows some edge-aligned PWM waveforms in an example where TIMx\_ARR=8.

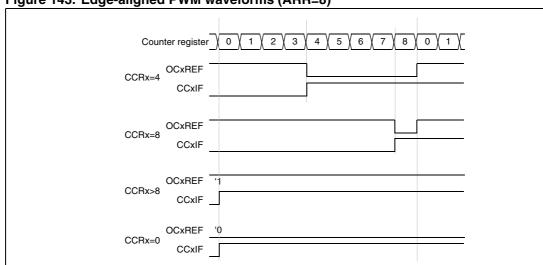


Figure 143. Edge-aligned PWM waveforms (ARR=8)

### **Downcounting configuration**

Downcounting is active when DIR bit in TIMx\_CR1 register is high. Refer to *Section : Downcounting mode on page 373*.

In PWM mode 1, the reference signal ocxref is low as long as TIMx\_CNT>TIMx\_CCRx else it becomes high. If the compare value in TIMx\_CCRx is greater than the auto-reload value in TIMx\_ARR, then ocxref is held at '1.0% PWM is not possible in this mode.

#### PWM center-aligned mode

Center-aligned mode is active when the CMS bits in TIMx\_CR1 register are different from '00 (all the remaining configurations having the same effect on the ocxref/OCx signals). The compare flag is set when the counter counts up, when it counts down or both when it counts up and down depending on the CMS bits configuration. The direction bit (DIR) in the TIMx\_CR1 register is updated by hardware and must not be changed by software. Refer to Section: Center-aligned mode (up/down counting) on page 375.

Figure 144 shows some center-aligned PWM waveforms in an example where:

- TIMx ARR=8,
- PWM mode is the PWM mode 1,
- The flag is set when the counter counts down corresponding to the center-aligned mode 1 selected for CMS=01 in TIMx\_CR1 register.

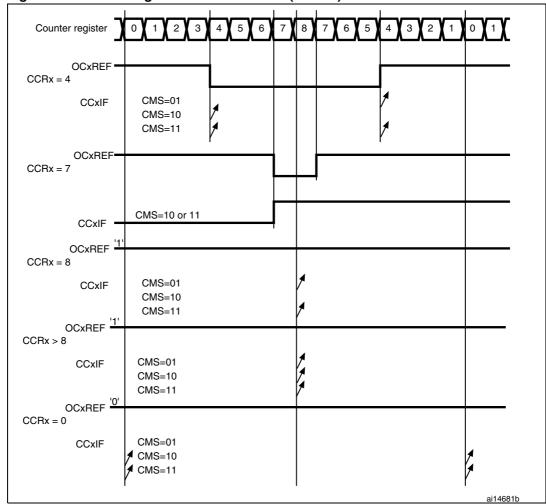


Figure 144. Center-aligned PWM waveforms (ARR=8)

Hints on using center-aligned mode:

- When starting in center-aligned mode, the current up-down configuration is used. It
  means that the counter counts up or down depending on the value written in the DIR bit
  in the TIMx\_CR1 register. Moreover, the DIR and CMS bits must not be changed at the
  same time by the software.
- Writing to the counter while running in center-aligned mode is not recommended as it can lead to unexpected results. In particular:
  - The direction is not updated if you write a value in the counter that is greater than the auto-reload value (TIMx\_CNT>TIMx\_ARR). For example, if the counter was counting up, it continues to count up.
  - The direction is updated if you write 0 or write the TIMx\_ARR value in the counter but no Update Event UEV is generated.
- The safest way to use center-aligned mode is to generate an update by software (setting the UG bit in the TIMx\_EGR register) just before starting the counter and not to write the counter while it is running.

## 14.3.10 One-pulse mode

One-pulse mode (OPM) is a particular case of the previous modes. It allows the counter to be started in response to a stimulus and to generate a pulse with a programmable length after a programmable delay.

Starting the counter can be controlled through the slave mode controller. Generating the waveform can be done in output compare mode or PWM mode. You select One-pulse mode by setting the OPM bit in the TIMx\_CR1 register. This makes the counter stop automatically at the next update event UEV.

A pulse can be correctly generated only if the compare value is different from the counter initial value. Before starting (when the timer is waiting for the trigger), the configuration must be:

- In upcounting: CNT<CCRx≤ ARR (in particular, 0<CCRx),</li>
- In downcounting: CNT>CCRx.

#### Figure 145. Example of one-pulse mode.

For example you may want to generate a positive pulse on OC1 with a length of  $t_{PULSE}$  and after a delay of  $t_{DELAY}$  as soon as a positive edge is detected on the TI2 input pin.

Let's use TI2FP2 as trigger 1:

- Map TI2FP2 on TI2 by writing IC2S=01 in the TIMx\_CCMR1 register.
- TI2FP2 must detect a rising edge, write CC2P=0 and CC2NP='0' in the TIMx\_CCER register.
- Configure TI2FP2 as trigger for the slave mode controller (TRGI) by writing TS=110 in the TIMx\_SMCR register.
- TI2FP2 is used to start the counter by writing SMS to '110 in the TIMx\_SMCR register (trigger mode).

The OPM waveform is defined by writing the compare registers (taking into account the clock frequency and the counter prescaler).

- The t<sub>DFLAY</sub> is defined by the value written in the TIMx\_CCR1 register.
- The t<sub>PULSE</sub> is defined by the difference between the auto-reload value and the compare value (TIMx\_ARR - TIMx\_CCR1).
- Let's say you want to build a waveform with a transition from '0 to '1 when a compare match occurs and a transition from '1 to '0 when the counter reaches the auto-reload value. To do this you enable PWM mode 2 by writing OC1M=111 in the TIMx\_CCMR1 register. You can optionally enable the preload registers by writing OC1PE=1 in the TIMx\_CCMR1 register and ARPE in the TIMx\_CR1 register. In this case you have to write the compare value in the TIMx\_CCR1 register, the auto-reload value in the TIMx\_ARR register, generate an update by setting the UG bit and wait for external trigger event on TI2. CC1P is written to '0 in this example.

In our example, the DIR and CMS bits in the TIMx\_CR1 register should be low.

You only want 1 pulse (Single mode), so you write '1 in the OPM bit in the TIMx\_CR1 register to stop the counter at the next update event (when the counter rolls over from the auto-reload value back to 0). When OPM bit in the TIMx\_CR1 register is set to '0', so the Repetitive Mode is selected.

#### Particular case: OCx fast enable:

In One-pulse mode, the edge detection on TIx input set the CEN bit which enables the counter. Then the comparison between the counter and the compare value makes the output toggle. But several clock cycles are needed for these operations and it limits the minimum delay t<sub>DFI AY</sub> min we can get.

If you want to output a waveform with the minimum delay, you can set the OCxFE bit in the TIMx\_CCMRx register. Then OCxRef (and OCx) is forced in response to the stimulus, without taking in account the comparison. Its new level is the same as if a compare match had occurred. OCxFE acts only if the channel is configured in PWM1 or PWM2 mode.

## 14.3.11 Clearing the OCxREF signal on an external event

The OCxREF signal for a given channel can be driven Low by applying a High level to the ETRF input (OCxCE enable bit of the corresponding TIMx\_CCMRx register set to '1'). The OCxREF signal remains Low until the next update event, UEV, occurs.

This function can only be used in output compare and PWM modes, and does not work in forced mode.

- 1. For example, the ETR signal can be connected to the output of a comparator to be used for current handling. In this case, ETR must be configured as follows:The external trigger prescaler should be kept off: bits ETPS[1:0] in the TIMx\_SMCR register are cleared to 00.
- 2. The external clock mode 2 must be disabled: bit ECE in the TIM1\_SMCR register is cleared to 0.
- 3. The external trigger polarity (ETP) and the external trigger filter (ETF) can be configured according to the application's needs.

*Figure 146* shows the behavior of the OCxREF signal when the ETRF input becomes high, for both values of the OCxCE enable bit. In this example, the timer TIMx is programmed in PWM mode.

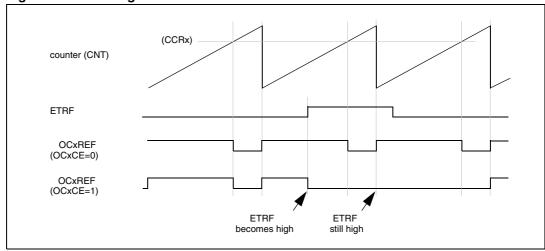


Figure 146. Clearing TIMx OCxREF

 In case of a PWM with a 100% duty cycle (if CCRx>ARR), OCxREF is enabled again at the next counter overflow.

#### 14.3.12 Encoder interface mode

To select Encoder Interface mode write SMS='001 in the TIMx\_SMCR register if the counter is counting on TI2 edges only, SMS=010 if it is counting on TI1 edges only and SMS=011 if it is counting on both TI1 and TI2 edges.

Select the TI1 and TI2 polarity by programming the CC1P and CC2P bits in the TIMx\_CCER register. When needed, you can program the input filter as well.

The two inputs TI1 and TI2 are used to interface to an incremental encoder. Refer to *Table 61*. The counter is clocked by each valid transition on TI1FP1 or TI2FP2 (TI1 and TI2 after input filter and polarity selection, TI1FP1=TI1 if not filtered and not inverted, TI2FP2=TI2 if not filtered and not inverted) assuming that it is enabled (CEN bit in TIMx\_CR1 register written to '1). The sequence of transitions of the two inputs is evaluated and generates count pulses as well as the direction signal. Depending on the sequence the counter counts up or down, the DIR bit in the TIMx\_CR1 register is modified by hardware accordingly. The DIR bit is calculated at each transition on any input (TI1 or TI2), whatever the counter is counting on TI1 only, TI2 only or both TI1 and TI2.

Encoder interface mode acts simply as an external clock with direction selection. This means that the counter just counts continuously between 0 and the auto-reload value in the TIMx\_ARR register (0 to ARR or ARR down to 0 depending on the direction). So you must configure TIMx\_ARR before starting. In the same way, the capture, compare, prescaler, trigger output features continue to work as normal.

In this mode, the counter is modified automatically following the speed and the direction of the incremental encoder and its content, therefore, always represents the encoder's position. The count direction correspond to the rotation direction of the connected sensor. The table summarizes the possible combinations, assuming TI1 and TI2 don't switch at the same time.

Active edge	Level on opposite	TI1FP1	signal	TI2FP2 signal		
Active edge	signal (TI1FP1 for TI2, TI2FP2 for TI1)	Rising	Falling	Rising	Falling	
Counting on	High	Down	Up	No Count	No Count	
TI1 only	Low	Up	Down	No Count	No Count	
Counting on	High	No Count	No Count	Up	Down	
TI2 only	Low	No Count	No Count	Down	Up	
Counting on	High	Down	Up	Up	Down	
TI1 and TI2	Low	Up	Down	Down	Up	

Table 61. Counting direction versus encoder signals

An external incremental encoder can be connected directly to the MCU without external interface logic. However, comparators are normally be used to convert the encoder's differential outputs to digital signals. This greatly increases noise immunity. The third encoder output which indicate the mechanical zero position, may be connected to an external interrupt input and trigger a counter reset.

Figure 147 gives an example of counter operation, showing count signal generation and direction control. It also shows how input jitter is compensated where both edges are

selected. This might occur if the sensor is positioned near to one of the switching points. For this example we assume that the configuration is the following:

- CC1S= '01' (TIMx\_CCMR1 register, TI1FP1 mapped on TI1)
- CC2S= '01' (TIMx CCMR2 register, TI2FP2 mapped on TI2)
- CC1P= '0', CC1NP = '0', IC1F = '0000' (TIMx\_CCER register, TI1FP1 noninverted, TI1FP1=TI1)
- CC2P= '0', CC2NP = '0', IC2F = '0000' (TIMx\_CCER register, TI2FP2 noninverted, TI2FP2=TI2)
- SMS= '011' (TIMx\_SMCR register, both inputs are active on both rising and falling edges)
- CEN = 1 (TIMx\_CR1 register, Counter is enabled)

configuration as above except CC1P=1).

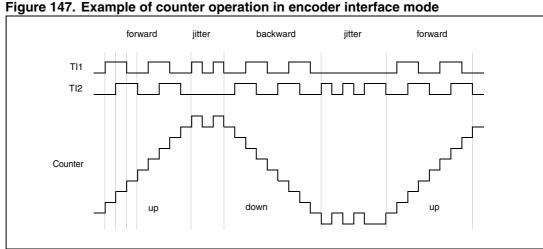


Figure 148 gives an example of counter behavior when TI1FP1 polarity is inverted (same

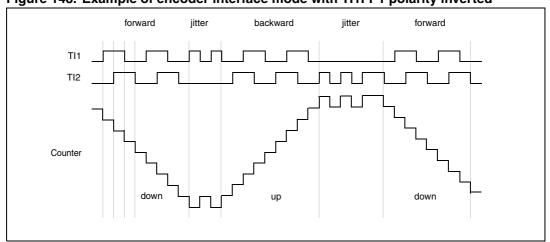


Figure 148. Example of encoder interface mode with TI1FP1 polarity inverted

The timer, when configured in Encoder Interface mode provides information on the sensor's current position. You can obtain dynamic information (speed, acceleration, deceleration) by measuring the period between two encoder events using a second timer configured in

capture mode. The output of the encoder which indicates the mechanical zero can be used for this purpose. Depending on the time between two events, the counter can also be read at regular times. You can do this by latching the counter value into a third input capture register if available (then the capture signal must be periodic and can be generated by another timer). when available, it is also possible to read its value through a DMA request generated by a Real-Time clock.

### 14.3.13 Timer input XOR function

The TI1S bit in the TIM\_CR2 register, allows the input filter of channel 1 to be connected to the output of a XOR gate, combining the three input pins TIMx\_CH1 to TIMx\_CH3.

The XOR output can be used with all the timer input functions such as trigger or input capture.

## 14.3.14 Timers and external trigger synchronization

The TIMx Timers can be synchronized with an external trigger in several modes: Reset mode, Gated mode and Trigger mode.

#### Slave mode: Reset mode

The counter and its prescaler can be reinitialized in response to an event on a trigger input. Moreover, if the URS bit from the TIMx\_CR1 register is low, an update event UEV is generated. Then all the preloaded registers (TIMx\_ARR, TIMx\_CCRx) are updated.

In the following example, the upcounter is cleared in response to a rising edge on TI1 input:

- Configure the channel 1 to detect rising edges on TI1. Configure the input filter duration (in this example, we don't need any filter, so we keep IC1F=0000). The capture prescaler is not used for triggering, so you don't need to configure it. The CC1S bits select the input capture source only, CC1S = 01 in the TIMx\_CCMR1 register. Write CC1P=0 and CC1NP=0 in TIMx\_CCER register to validate the polarity (and detect rising edges only).
- Configure the timer in reset mode by writing SMS=100 in TIMx\_SMCR register. Select TI1 as the input source by writing TS=101 in TIMx\_SMCR register.
- Start the counter by writing CEN=1 in the TIMx\_CR1 register.

The counter starts counting on the internal clock, then behaves normally until TI1 rising edge. When TI1 rises, the counter is cleared and restarts from 0. In the meantime, the trigger flag is set (TIF bit in the TIMx\_SR register) and an interrupt request, or a DMA request can be sent if enabled (depending on the TIE and TDE bits in TIMx\_DIER register).

The following figure shows this behavior when the auto-reload register TIMx\_ARR=0x36. The delay between the rising edge on TI1 and the actual reset of the counter is due to the resynchronization circuit on TI1 input.

TI1

UG

Counter clock = CK\_CNT = CK\_PSC

Counter register 30\31\32\33\34\35\36\00\01\02\03\00\01\02\03\00\

TIF

Figure 149. Control circuit in reset mode

### Slave mode: Gated mode

The counter can be enabled depending on the level of a selected input.

In the following example, the upcounter counts only when TI1 input is low:

- Configure the channel 1 to detect low levels on TI1. Configure the input filter duration (in this example, we don't need any filter, so we keep IC1F=0000). The capture prescaler is not used for triggering, so you don't need to configure it. The CC1S bits select the input capture source only, CC1S=01 in TIMx\_CCMR1 register. Write CC1P=1 in TIMx\_CCER register to validate the polarity (and detect low level only).
- Configure the timer in gated mode by writing SMS=101 in TIMx\_SMCR register. Select TI1 as the input source by writing TS=101 in TIMx\_SMCR register.
- Enable the counter by writing CEN=1 in the TIMx\_CR1 register (in gated mode, the counter doesn't start if CEN=0, whatever is the trigger input level).

The counter starts counting on the internal clock as long as TI1 is low and stops as soon as TI1 becomes high. The TIF flag in the TIMx\_SR register is set both when the counter starts or stops.

The delay between the rising edge on TI1 and the actual stop of the counter is due to the resynchronization circuit on TI1 input.

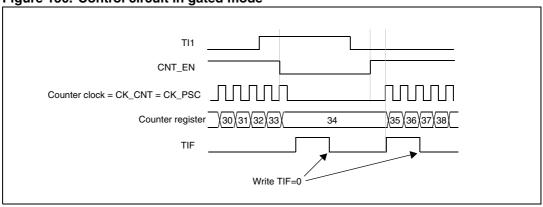


Figure 150. Control circuit in gated mode

1. The configuration "CCxP=CCxNP=1" (detection of both rising and falling edges) does not have any effect in gated mode because gated mode acts on a level and not on an edge.

## Slave mode: Trigger mode

The counter can start in response to an event on a selected input.

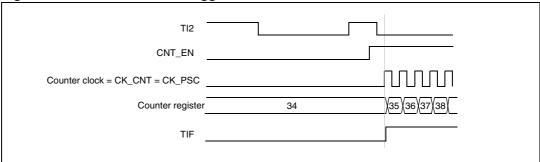
In the following example, the upcounter starts in response to a rising edge on TI2 input:

- Configure the channel 2 to detect rising edges on TI2. Configure the input filter duration (in this example, we don't need any filter, so we keep IC2F=0000). The capture prescaler is not used for triggering, so you don't need to configure it. CC2S bits are selecting the input capture source only, CC2S=01 in TIMx\_CCMR1 register. Write CC2P=1 in TIMx\_CCER register to validate the polarity (and detect low level only).
- Configure the timer in trigger mode by writing SMS=110 in TIMx\_SMCR register. Select TI2 as the input source by writing TS=110 in TIMx\_SMCR register.

When a rising edge occurs on TI2, the counter starts counting on the internal clock and the TIF flag is set.

The delay between the rising edge on TI2 and the actual start of the counter is due to the resynchronization circuit on TI2 input.

Figure 151. Control circuit in trigger mode



### Slave mode: External Clock mode 2 + trigger mode

The external clock mode 2 can be used in addition to another slave mode (except external clock mode 1 and encoder mode). In this case, the ETR signal is used as external clock input, and another input can be selected as trigger input when operating in reset mode, gated mode or trigger mode. It is recommended not to select ETR as TRGI through the TS bits of TIMx\_SMCR register.

In the following example, the upcounter is incremented at each rising edge of the ETR signal as soon as a rising edge of TI1 occurs:

- Configure the external trigger input circuit by programming the TIMx\_SMCR register as follows:
  - ETF = 0000: no filter
  - ETPS = 00: prescaler disabled
  - ETP = 0: detection of rising edges on ETR and ECE=1 to enable the external clock mode 2.
- 2. Configure the channel 1 as follows, to detect rising edges on TI:
  - IC1F = 0000: no filter.
  - The capture prescaler is not used for triggering and does not need to be configured.
  - CC1S = 01 in TIMx\_CCMR1 register to select only the input capture source
  - CC1P = 0 in TIMx\_CCER register to validate the polarity (and detect rising edge only).
- 3. Configure the timer in trigger mode by writing SMS=110 in TIMx\_SMCR register. Select TI1 as the input source by writing TS=101 in TIMx\_SMCR register.

A rising edge on TI1 enables the counter and sets the TIF flag. The counter then counts on ETR rising edges.

The delay between the rising edge of the ETR signal and the actual reset of the counter is due to the resynchronization circuit on ETRP input.

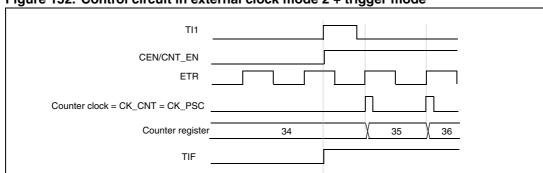


Figure 152. Control circuit in external clock mode 2 + trigger mode

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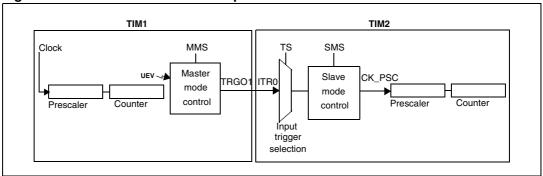
# 14.3.15 Timer synchronization

The TIMx timers are linked together internally for timer synchronization or chaining. When one Timer is configured in Master Mode, it can reset, start, stop or clock the counter of another Timer configured in Slave Mode.

Figure 153: Master/Slave timer example presents an overview of the trigger selection and the master mode selection blocks.

### Using one timer as prescaler for another

Figure 153. Master/Slave timer example



For example, you can configure Timer 1 to act as a prescaler for Timer 2. Refer to *Figure 153*. To do this:

- Configure Timer 1 in master mode so that it outputs a periodic trigger signal on each update event UEV. If you write MMS=010 in the TIM1\_CR2 register, a rising edge is output on TRGO1 each time an update event is generated.
- To connect the TRGO1 output of Timer 1 to Timer 2, Timer 2 must be configured in slave mode using ITR0 as internal trigger. You select this through the TS bits in the TIM2\_SMCR register (writing TS=000).
- Then you put the slave mode controller in external clock mode 1 (write SMS=111 in the TIM2\_SMCR register). This causes Timer 2 to be clocked by the rising edge of the periodic Timer 1 trigger signal (which correspond to the timer 1 counter overflow).
- Finally both timers must be enabled by setting their respective CEN bits (TIMx\_CR1 register).

Note:

If OCx is selected on Timer 1 as trigger output (MMS=1xx), its rising edge is used to clock the counter of timer 2.

## Using one timer to enable another timer

In this example, we control the enable of Timer 2 with the output compare 1 of Timer 1. Refer to *Figure 153* for connections. Timer 2 counts on the divided internal clock only when

Note:

OC1REF of Timer 1 is high. Both counter clock frequencies are divided by 3 by the prescaler compared to CK\_INT ( $f_{CK}$  CNT =  $f_{CK}$  INT/3).

- Configure Timer 1 master mode to send its Output Compare 1 Reference (OC1REF) signal as trigger output (MMS=100 in the TIM1\_CR2 register).
- Configure the Timer 1 OC1REF waveform (TIM1 CCMR1 register).
- Configure Timer 2 to get the input trigger from Timer 1 (TS=000 in the TIM2\_SMCR register).

The counter 2 clock is not synchronized with counter 1, this mode only affects the Timer 2

- Configure Timer 2 in gated mode (SMS=101 in TIM2\_SMCR register).
- Enable Timer 2 by writing '1 in the CEN bit (TIM2 CR1 register).
- Start Timer 1 by writing '1 in the CEN bit (TIM1\_CR1 register).

counter enable signal.

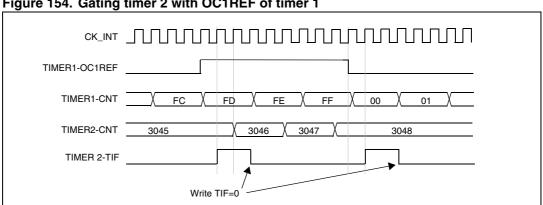


Figure 154. Gating timer 2 with OC1REF of timer 1

In the example in Figure 154, the Timer 2 counter and prescaler are not initialized before being started. So they start counting from their current value. It is possible to start from a given value by resetting both timers before starting Timer 1. You can then write any value you want in the timer counters. The timers can easily be reset by software using the UG bit in the TIMx\_EGR registers.

In the next example, we synchronize Timer 1 and Timer 2. Timer 1 is the master and starts from 0. Timer 2 is the slave and starts from 0xE7. The prescaler ratio is the same for both

timers. Timer 2 stops when Timer 1 is disabled by writing '0 to the CEN bit in the TIM1\_CR1 register:

- Configure Timer 1 master mode to send its Output Compare 1 Reference (OC1REF) signal as trigger output (MMS=100 in the TIM1\_CR2 register).
- Configure the Timer 1 OC1REF waveform (TIM1 CCMR1 register).
- Configure Timer 2 to get the input trigger from Timer 1 (TS=000 in the TIM2\_SMCR register).
- Configure Timer 2 in gated mode (SMS=101 in TIM2\_SMCR register).
- Reset Timer 1 by writing '1 in UG bit (TIM1\_EGR register).
- Reset Timer 2 by writing '1 in UG bit (TIM2\_EGR register).
- Initialize Timer 2 to 0xE7 by writing '0xE7' in the timer 2 counter (TIM2\_CNTL).
- Enable Timer 2 by writing '1 in the CEN bit (TIM2\_CR1 register).
- Start Timer 1 by writing '1 in the CEN bit (TIM1\_CR1 register).
- Stop Timer 1 by writing '0 in the CEN bit (TIM1\_CR1 register).

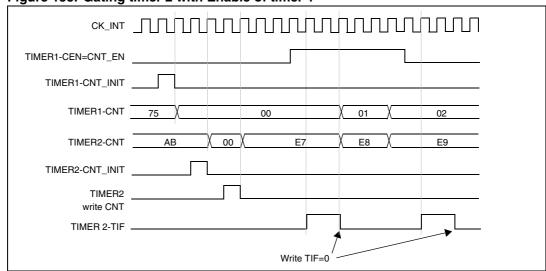


Figure 155. Gating timer 2 with Enable of timer 1

### Using one timer to start another timer

In this example, we set the enable of Timer 2 with the update event of Timer 1. Refer to *Figure 153* for connections. Timer 2 starts counting from its current value (which can be nonzero) on the divided internal clock as soon as the update event is generated by Timer 1. When Timer 2 receives the trigger signal its CEN bit is automatically set and the counter counts until we write '0 to the CEN bit in the TIM2\_CR1 register. Both counter clock frequencies are divided by 3 by the prescaler compared to CK\_INT ( $f_{CK-CNT} = f_{CK-INT}/3$ ).

- Configure Timer 1 master mode to send its Update Event (UEV) as trigger output (MMS=010 in the TIM1\_CR2 register).
- Configure the Timer 1 period (TIM1\_ARR registers).
- Configure Timer 2 to get the input trigger from Timer 1 (TS=000 in the TIM2\_SMCR register).
- Configure Timer 2 in trigger mode (SMS=110 in TIM2\_SMCR register).
- Start Timer 1 by writing '1 in the CEN bit (TIM1 CR1 register).

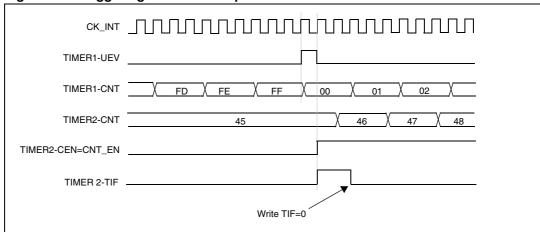


Figure 156. Triggering timer 2 with update of timer 1

As in the previous example, you can initialize both counters before starting counting. *Figure 157* shows the behavior with the same configuration as in *Figure 156* but in trigger mode instead of gated mode (SMS=110 in the TIM2\_SMCR register).

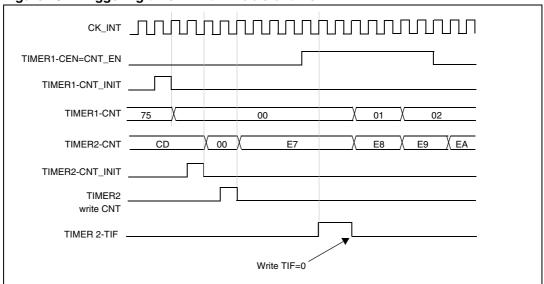


Figure 157. Triggering timer 2 with Enable of timer 1

### Using one timer as prescaler for another timer

For example, you can configure Timer 1 to act as a prescaler for Timer 2. Refer to *Figure 153* for connections. To do this:

- Configure Timer 1 master mode to send its Update Event (UEV) as trigger output (MMS=010 in the TIM1\_CR2 register). then it outputs a periodic signal on each counter overflow.
- Configure the Timer 1 period (TIM1 ARR registers).
- Configure Timer 2 to get the input trigger from Timer 1 (TS=000 in the TIM2\_SMCR register).
- Configure Timer 2 in external clock mode 1 (SMS=111 in TIM2\_SMCR register).
- Start Timer 2 by writing '1 in the CEN bit (TIM2\_CR1 register).
- Start Timer 1 by writing '1 in the CEN bit (TIM1\_CR1 register).

## Starting 2 timers synchronously in response to an external trigger

In this example, we set the enable of timer 1 when its TI1 input rises, and the enable of Timer 2 with the enable of Timer 1. Refer to *Figure 153* for connections. To ensure the counters are aligned, Timer 1 must be configured in Master/Slave mode (slave with respect to TI1, master with respect to Timer 2):

- Configure Timer 1 master mode to send its Enable as trigger output (MMS=001 in the TIM1 CR2 register).
- Configure Timer 1 slave mode to get the input trigger from TI1 (TS=100 in the TIM1\_SMCR register).
- Configure Timer 1 in trigger mode (SMS=110 in the TIM1\_SMCR register).
- Configure the Timer 1 in Master/Slave mode by writing MSM=1 (TIM1\_SMCR register).
- Configure Timer 2 to get the input trigger from Timer 1 (TS=000 in the TIM2\_SMCR register).
- Configure Timer 2 in trigger mode (SMS=110 in the TIM2 SMCR register).

When a rising edge occurs on TI1 (Timer 1), both counters starts counting synchronously on the internal clock and both TIF flags are set.

Note:

In this example both timers are initialized before starting (by setting their respective UG bits). Both counters starts from 0, but you can easily insert an offset between them by writing any of the counter registers (TIMx\_CNT). You can see that the master/slave mode insert a delay between CNT EN and CK PSC on timer 1.

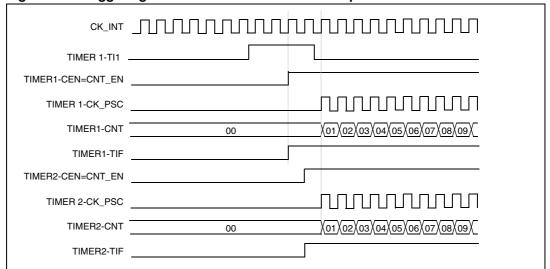


Figure 158. Triggering timer 1 and 2 with timer 1 TI1 input

# 14.3.16 **Debug mode**

When the microcontroller enters debug mode (Cortex<sup>™</sup>-M4F core - halted), the TIMx counter either continues to work normally or stops, depending on DBG\_TIMx\_STOP configuration bit in DBGMCU module. For more details, refer to *Section 32.16.2: Debug support for timers, watchdog, bxCAN and I2C*.

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# 14.4 TIM2 to TIM5 registers

Refer to Section 1.1 on page 46 for a list of abbreviations used in register descriptions.

The 32-bit peripheral registers have to be written by words (32 bits). All other peripheral registers have to be written by half-words (16 bits) or words (32 bits). Read accesses can be done by bytes (8 bits), half-words (16 bits) or words (32 bits).

# 14.4.1 TIMx control register 1 (TIMx\_CR1)

Address offset: 0x00 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		Rese	n rod			CKD	[1:0]	ARPE	CI	ИS	DIR	OPM	URS	UDIS	CEN
		nese	erveu			rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 15:10 Reserved, must be kept at reset value.

#### Bits 9:8 CKD: Clock division

This bit-field indicates the division ratio between the timer clock (CK\_INT) frequency and sampling clock used by the digital filters (ETR, TIx),

00:  $t_{DTS} = t_{CK\_INT}$ 01:  $t_{DTS} = 2 \times t_{CK\_INT}$ 10:  $t_{DTS} = 4 \times t_{CK\_INT}$ 

11: Reserved

Bit 7 ARPE: Auto-reload preload enable

0: TIMx\_ARR register is not buffered

1: TIMx\_ARR register is buffered

Bits 6:5 CMS: Center-aligned mode selection

00: Edge-aligned mode. The counter counts up or down depending on the direction bit (DIR)

01: Center-aligned mode 1. The counter counts up and down alternatively. Output compare interrupt flags of channels configured in output (CCxS=00 in TIMx\_CCMRx register) are set only when the counter is counting down.

10: Center-aligned mode 2. The counter counts up and down alternatively. Output compare interrupt flags of channels configured in output (CCxS=00 in TIMx\_CCMRx register) are set only when the counter is counting up.

11: Center-aligned mode 3. The counter counts up and down alternatively. Output compare interrupt flags of channels configured in output (CCxS=00 in TIMx\_CCMRx register) are set both when the counter is counting up or down.

Note: It is not allowed to switch from edge-aligned mode to center-aligned mode as long as the counter is enabled (CEN=1)

### Bit 4 DIR: Direction

0: Counter used as upcounter

1: Counter used as downcounter

Note: This bit is read only when the timer is configured in Center-aligned mode or Encoder mode.

## Bit 3 OPM: One-pulse mode

0: Counter is not stopped at update event

1: Counter stops counting at the next update event (clearing the bit CEN)

### Bit 2 URS: Update request source

This bit is set and cleared by software to select the UEV event sources.

0: Any of the following events generate an update interrupt or DMA request if enabled. These events can be:

- Counter overflow/underflow
- Setting the UG bit
- Update generation through the slave mode controller
- 1: Only counter overflow/underflow generates an update interrupt or DMA request if enabled.

### Bit 1 UDIS: Update disable

This bit is set and cleared by software to enable/disable UEV event generation.

- 0: UEV enabled. The Update (UEV) event is generated by one of the following events:
  - Counter overflow/underflow
  - Setting the UG bit
  - Update generation through the slave mode controller

Buffered registers are then loaded with their preload values.

1: UEV disabled. The Update event is not generated, shadow registers keep their value (ARR, PSC, CCRx). However the counter and the prescaler are reinitialized if the UG bit is set or if a hardware reset is received from the slave mode controller.

#### Bit 0 CEN: Counter enable

- 0: Counter disabled
- 1: Counter enabled

Note: External clock, gated mode and encoder mode can work only if the CEN bit has been previously set by software. However trigger mode can set the CEN bit automatically by hardware.

CEN is cleared automatically in one-pulse mode, when an update event occurs.

# 14.4.2 TIMx control register 2 (TIMx\_CR2)

Address offset: 0x04 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
			Poss	erved				TI1S		MMS[2:0]		CCDS		Reserved		Ì
			nese	erveu				rw	rw	rw	rw	rw		neserveu		l

Bits 15:8 Reserved, must be kept at reset value.

Bit 7 TI1S: TI1 selection

0: The TIMx\_CH1 pin is connected to TI1 input

1: The TIMx\_CH1, CH2 and CH3 pins are connected to the TI1 input (XOR combination)

#### Bits 6:4 MMS: Master mode selection

These bits allow to select the information to be sent in master mode to slave timers for synchronization (TRGO). The combination is as follows:

000: **Reset** - the UG bit from the TIMx\_EGR register is used as trigger output (TRGO). If the reset is generated by the trigger input (slave mode controller configured in reset mode) then the signal on TRGO is delayed compared to the actual reset.

001: **Enable** - the Counter enable signal, CNT\_EN, is used as trigger output (TRGO). It is useful to start several timers at the same time or to control a window in which a slave timer is enabled. The Counter Enable signal is generated by a logic OR between CEN control bit and the trigger input when configured in gated mode.

When the Counter Enable signal is controlled by the trigger input, there is a delay on TRGO, except if the master/slave mode is selected (see the MSM bit description in TIMx\_SMCR register).

010: **Update** - The update event is selected as trigger output (TRGO). For instance a master timer can then be used as a prescaler for a slave timer.

011: **Compare Pulse** - The trigger output send a positive pulse when the CC1IF flag is to be set (even if it was already high), as soon as a capture or a compare match occurred. (TRGO)

100: Compare - OC1REF signal is used as trigger output (TRGO)

101: Compare - OC2REF signal is used as trigger output (TRGO)

110: Compare - OC3REF signal is used as trigger output (TRGO)

111: Compare - OC4REF signal is used as trigger output (TRGO)

Bit 3 CCDS: Capture/compare DMA selection

0: CCx DMA request sent when CCx event occurs

1: CCx DMA requests sent when update event occurs

Bits 2:0 Reserved, must be kept at reset value.

# 14.4.3 TIMx slave mode control register (TIMx\_SMCR)

Address offset: 0x08 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ETP	ECE	ETPS	S[1:0]		ETF	[3:0]		MSM		TS[2:0]		Res.		SMS[2:0]	
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	nes.	rw	rw	rw

#### Bit 15 ETP: External trigger polarity

This bit selects whether ETR or ETR is used for trigger operations

0: ETR is noninverted, active at high level or rising edge

1: ETR is inverted, active at low level or falling edge

#### Bit 14 ECE: External clock enable

This bit enables External clock mode 2.

- 0: External clock mode 2 disabled
- 1: External clock mode 2 enabled. The counter is clocked by any active edge on the ETRF signal.
- 1: Setting the ECE bit has the same effect as selecting external clock mode 1 with TRGI connected to ETRF (SMS=111 and TS=111).
- 2: It is possible to simultaneously use external clock mode 2 with the following slave modes: reset mode, gated mode and trigger mode. Nevertheless, TRGI must not be connected to ETRF in this case (TS bits must not be 111).
- 3: If external clock mode 1 and external clock mode 2 are enabled at the same time, the external clock input is ETRF.

### Bits 13:12 ETPS: External trigger prescaler

External trigger signal ETRP frequency must be at most 1/4 of CK\_INT frequency. A prescaler can be enabled to reduce ETRP frequency. It is useful when inputting fast external clocks.

00: Prescaler OFF

01: ETRP frequency divided by 2

10: ETRP frequency divided by 4

11: ETRP frequency divided by 8

### Bits 11:8 ETF[3:0]: External trigger filter

This bit-field then defines the frequency used to sample ETRP signal and the length of the digital filter applied to ETRP. The digital filter is made of an event counter in which N events are needed to validate a transition on the output:

0000: No filter, sampling is done at fDTS

0001: f<sub>SAMPLING</sub>=f<sub>CK INT</sub>, N=2

0010: f<sub>SAMPLING</sub>=f<sub>CK\_INT</sub>, N=4

0011: f<sub>SAMPLING</sub>=f<sub>CK INT</sub>, N=8

0100: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/2, N=6

0101: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/2, N=8

0110:  $f_{SAMPLING} = f_{DTS}/4$ , N=6

0111:  $f_{SAMPLING} = f_{DTS}/4$ , N=8

1000: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/8, N=6

1001: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/8, N=8

1010: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/16, N=5

1011:  $f_{SAMPLING} = f_{DTS}/16$ , N=6

1100: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/16, N=8

1101: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/32, N=5

1110: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/32, N=6

1111: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/32, N=8

#### Bit 7 MSM: Master/Slave mode

0: No action

1: The effect of an event on the trigger input (TRGI) is delayed to allow a perfect synchronization between the current timer and its slaves (through TRGO). It is useful if we want to synchronize several timers on a single external event.

#### Bits 6:4 TS: Trigger selection

This bit-field selects the trigger input to be used to synchronize the counter.

000: Internal Trigger 0 (ITR0)

001: Internal Trigger 1 (ITR1).

010: Internal Trigger 2 (ITR2).

011: Internal Trigger 3 (ITR3).

100: TI1 Edge Detector (TI1F\_ED)

101: Filtered Timer Input 1 (TI1FP1)

110: Filtered Timer Input 2 (TI2FP2)

111: External Trigger input (ETRF)

See *Table 62: TIMx internal trigger connection on page 408* for more details on ITRx meaning for each Timer.

Note: These bits must be changed only when they are not used (e.g. when SMS=000) to avoid wrong edge detections at the transition.

Bit 3 Reserved, must be kept at reset value.

#### Bits 2:0 SMS: Slave mode selection

When external signals are selected the active edge of the trigger signal (TRGI) is linked to the polarity selected on the external input (see Input Control register and Control Register description.

000: Slave mode disabled - if CEN = 1 then the prescaler is clocked directly by the internal clock.

001: Encoder mode 1 - Counter counts up/down on TI2FP2 edge depending on TI1FP1 level.

010: Encoder mode 2 - Counter counts up/down on TI1FP1 edge depending on TI2FP2 level.

011: Encoder mode 3 - Counter counts up/down on both TI1FP1 and TI2FP2 edges depending on the level of the other input.

100: Reset Mode - Rising edge of the selected trigger input (TRGI) reinitializes the counter and generates an update of the registers.

101: Gated Mode - The counter clock is enabled when the trigger input (TRGI) is high. The counter stops (but is not reset) as soon as the trigger becomes low. Both start and stop of the counter are controlled.

110: Trigger Mode - The counter starts at a rising edge of the trigger TRGI (but it is not reset). Only the start of the counter is controlled.

111: External Clock Mode 1 - Rising edges of the selected trigger (TRGI) clock the counter.

Note: The gated mode must not be used if TI1F\_ED is selected as the trigger input (TS=100). Indeed, TI1F\_ED outputs 1 pulse for each transition on TI1F, whereas the gated mode checks the level of the trigger signal.

Table 62. TIMx internal trigger connection

Slave TIM	ITR0 (TS = 000)	ITR1 (TS = 001)	ITR2 (TS = 010)	ITR3 (TS = 011)
TIM2	TIM1	TIM8	TIM3	TIM4
ТІМЗ	TIM1	TIM2	TIM5	TIM4

Table 62. TIMx internal trigger connection

Slave TIM	ITR0 (TS = 000)	ITR1 (TS = 001)	ITR2 (TS = 010)	ITR3 (TS = 011)
TIM4	TIM1	TIM2	TIM3	TIM8
TIM5	TIM2	TIM3	TIM4	TIM8

# 14.4.4 TIMx DMA/Interrupt enable register (TIMx\_DIER)

Address offset: 0x0C Reset value: 0x0000

	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ſ	Res.	TDE	Res	CC4DE	CC3DE	CC2DE	CC1DE	UDE	Res.	TIE	Res	CC4IE	CC3IE	CC2IE	CC1IE	UIE
	nes.	rw	nes	rw	rw	rw	rw	rw	nes.	rw	nes	rw	rw	rw	rw	rw

Bit 15 Reserved, must be kept at reset value.

Bit 14 TDE: Trigger DMA request enable

0: Trigger DMA request disabled.

1: Trigger DMA request enabled.

Bit 13 Reserved, always read as 0

Bit 12 CC4DE: Capture/Compare 4 DMA request enable

0: CC4 DMA request disabled.

1: CC4 DMA request enabled.

Bit 11 CC3DE: Capture/Compare 3 DMA request enable

0: CC3 DMA request disabled.

1: CC3 DMA request enabled.

Bit 10 CC2DE: Capture/Compare 2 DMA request enable

0: CC2 DMA request disabled.

1: CC2 DMA request enabled.

Bit 9 CC1DE: Capture/Compare 1 DMA request enable

0: CC1 DMA request disabled.

1: CC1 DMA request enabled.

Bit 8 **UDE**: Update DMA request enable

0: Update DMA request disabled.

1: Update DMA request enabled.

Bit 7 Reserved, must be kept at reset value.

Bit 6 TIE: Trigger interrupt enable

0: Trigger interrupt disabled.

1: Trigger interrupt enabled.

Bit 5 Reserved, must be kept at reset value.

Bit 4 **CC4IE**: Capture/Compare 4 interrupt enable

0: CC4 interrupt disabled.

1: CC4 interrupt enabled.

Bit 3 CC3IE: Capture/Compare 3 interrupt enable

0: CC3 interrupt disabled

1: CC3 interrupt enabled

Bit 2 CC2IE: Capture/Compare 2 interrupt enable

0: CC2 interrupt disabled1: CC2 interrupt enabled

Bit 1 **CC1IE**: Capture/Compare 1 interrupt enable

0: CC1 interrupt disabled1: CC1 interrupt enabled

Bit 0 **UIE**: Update interrupt enable

0: Update interrupt disabled1: Update interrupt enabled

# 14.4.5 TIMx status register (TIMx\_SR)

Address offset: 0x10 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Reserved		CC4OF	CC3OF	CC2OF	CC1OF	Rese	rvod	TIF	Res	CC4IF	CC3IF	CC2IF	CC1IF	UIF
	neserveu		rc_w0	rc_w0	rc_w0	rc_w0	nese	erveu	rc_w0	nes	rc_w0	rc_w0	rc_w0	rc_w0	rc_w0

Bit 15:13 Reserved, must be kept at reset value.

Bit 12 CC40F: Capture/Compare 4 overcapture flag

refer to CC1OF description

Bit 11 CC3OF: Capture/Compare 3 overcapture flag

refer to CC1OF description

Bit 10 CC2OF: Capture/compare 2 overcapture flag

refer to CC1OF description

Bit 9 CC10F: Capture/Compare 1 overcapture flag

This flag is set by hardware only when the corresponding channel is configured in input capture mode. It is cleared by software by writing it to '0.

- 0: No overcapture has been detected
- 1: The counter value has been captured in TIMx\_CCR1 register while CC1IF flag was already set
- Bits 8:7 Reserved, must be kept at reset value.
  - Bit 6 TIF: Trigger interrupt flag

This flag is set by hardware on trigger event (active edge detected on TRGI input when the slave mode controller is enabled in all modes but gated mode. It is set when the counter starts or stops when gated mode is selected. It is cleared by software.

- 0: No trigger event occurred
- 1: Trigger interrupt pending
- Bit 5 Reserved, must be kept at reset value.
- Bit 4 CC4IF: Capture/Compare 4 interrupt flag

refer to CC1IF description

Bit 3 CC3IF: Capture/Compare 3 interrupt flag

refer to CC1IF description

Bit 2 CC2IF: Capture/Compare 2 interrupt flag

refer to CC1IF description

### Bit 1 CC1IF: Capture/compare 1 interrupt flag

### If channel CC1 is configured as output:

This flag is set by hardware when the counter matches the compare value, with some exception in center-aligned mode (refer to the CMS bits in the TIMx\_CR1 register description). It is cleared by software.

0: No match

1: The content of the counter TIMx\_CNT matches the content of the TIMx\_CCR1 register. When the contents of TIMx\_CCR1 are greater than the contents of TIMx\_ARR, the CC1IF bit goes high on the counter overflow (in upcounting and up/down-counting modes) or underflow (in downcounting mode)

### If channel CC1 is configured as input:

This bit is set by hardware on a capture. It is cleared by software or by reading the TIMx CCR1 register.

0: No input capture occurred

1: The counter value has been captured in TIMx\_CCR1 register (An edge has been detected on IC1 which matches the selected polarity)

#### Bit 0 **UIF**: Update interrupt flag

- This bit is set by hardware on an update event. It is cleared by software.
   No update occurred.
  - 1: Update interrupt pending. This bit is set by hardware when the registers are updated:
- At overflow or underflow (for TIM2 to TIM5) and if UDIS=0 in the TIMx\_CR1 register.
- When CNT is reinitialized by software using the UG bit in TIMx\_EGR register, if URS=0 and UDIS=0 in the TIMx\_CR1 register.

When CNT is reinitialized by a trigger event (refer to the synchro control register description), if URS=0 and UDIS=0 in the TIMx\_CR1 register.



# 14.4.6 TIMx event generation register (TIMx\_EGR)

Address offset: 0x14 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
				Dagamaa					TG	Res.	CC4G	CC3G	CC2G	CC1G	UG
				Reserved					w	nes.	W	W	w	W	w

Bits 15:7 Reserved, must be kept at reset value.

### Bit 6 TG: Trigger generation

This bit is set by software in order to generate an event, it is automatically cleared by hardware.

0: No action

1: The TIF flag is set in TIMx\_SR register. Related interrupt or DMA transfer can occur if enabled.

Bit 5 Reserved, must be kept at reset value.

Bit 4 CC4G: Capture/compare 4 generation

refer to CC1G description

Bit 3 CC3G: Capture/compare 3 generation

refer to CC1G description

Bit 2 CC2G: Capture/compare 2 generation

refer to CC1G description

Bit 1 CC1G: Capture/compare 1 generation

This bit is set by software in order to generate an event, it is automatically cleared by hardware.

0: No action

1: A capture/compare event is generated on channel 1:

### If channel CC1 is configured as output:

CC1IF flag is set, Corresponding interrupt or DMA request is sent if enabled.

### If channel CC1 is configured as input:

The current value of the counter is captured in TIMx\_CCR1 register. The CC1IF flag is set, the corresponding interrupt or DMA request is sent if enabled. The CC1OF flag is set if the CC1IF flag was already high.

### Bit 0 UG: Update generation

This bit can be set by software, it is automatically cleared by hardware.

0: No action

1: Re-initialize the counter and generates an update of the registers. Note that the prescaler counter is cleared too (anyway the prescaler ratio is not affected). The counter is cleared if the center-aligned mode is selected or if DIR=0 (upcounting), else it takes the auto-reload value (TIMx\_ARR) if DIR=1 (downcounting).

# 14.4.7 TIMx capture/compare mode register 1 (TIMx\_CCMR1)

Address offset: 0x18 Reset value: 0x0000

The channels can be used in input (capture mode) or in output (compare mode). The direction of a channel is defined by configuring the corresponding CCxS bits. All the other bits of this register have a different function in input and in output mode. For a given bit, OCxx describes its function when the channel is configured in output, ICxx describes its function when the channel is configured in input. So you must take care that the same bit can have a different meaning for the input stage and for the output stage.

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
OC2CE	(	OC2M[2:0	]	OC2PE	OC2FE	CCas	S[1:0]	OC1CE	(	OC1M[2:0	]	OC1PE	OC1FE	CC19	S[1:0]
	IC2F	[3:0]		IC2PS	C[1:0]	0028	5[1.0]		IC1F	[3:0]		IC1PS	SC[1:0]	0	5[1.0]
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

# **Output compare mode**

Bit 15 OC2CE: Output compare 2 clear enable

Bits 14:12 OC2M[2:0]: Output compare 2 mode

Bit 11 OC2PE: Output compare 2 preload enable

Bit 10 OC2FE: Output compare 2 fast enable

Bits 9:8 CC2S[1:0]: Capture/Compare 2 selection

This bit-field defines the direction of the channel (input/output) as well as the used input.

00: CC2 channel is configured as output

01: CC2 channel is configured as input, IC2 is mapped on TI2

10: CC2 channel is configured as input, IC2 is mapped on TI1

11: CC2 channel is configured as input, IC2 is mapped on TRC. This mode is working only if an internal trigger input is selected through the TS bit (TIMx\_SMCR register)

Note: CC2S bits are writable only when the channel is OFF (CC2E = 0 in TIMx\_CCER).

Bit 7 OC1CE: Output compare 1 clear enable

OC1CE: Output Compare 1 Clear Enable

0: OC1Ref is not affected by the ETRF input

1: OC1Ref is cleared as soon as a High level is detected on ETRF input

#### Bits 6:4 OC1M: Output compare 1 mode

These bits define the behavior of the output reference signal OC1REF from which OC1 and OC1N are derived. OC1REF is active high whereas OC1 and OC1N active level depends on CC1P and CC1NP bits.

000: Frozen - The comparison between the output compare register TIMx\_CCR1 and the counter TIMx\_CNT has no effect on the outputs.(this mode is used to generate a timing base).

001: Set channel 1 to active level on match. OC1REF signal is forced high when the counter TIMx CNT matches the capture/compare register 1 (TIMx CCR1).

010: Set channel 1 to inactive level on match. OC1REF signal is forced low when the counter TIMx\_CNT matches the capture/compare register 1 (TIMx\_CCR1).

011: Toggle - OC1REF toggles when TIMx\_CNT=TIMx\_CCR1.

100: Force inactive level - OC1REF is forced low.

101: Force active level - OC1REF is forced high.

110: PWM mode 1 - In upcounting, channel 1 is active as long as TIMx\_CNT<TIMx\_CCR1 else inactive. In downcounting, channel 1 is inactive (OC1REF='0) as long as

TIMx\_CNT>TIMx\_CCR1 else active (OC1REF=1).

111: PWM mode 2 - In upcounting, channel 1 is inactive as long as

TIMx\_CNT<TIMx\_CCR1 else active. In downcounting, channel 1 is active as long as TIMx\_CNT>TIMx\_CCR1 else inactive.

Note: In PWM mode 1 or 2, the OCREF level changes only when the result of the comparison changes or when the output compare mode switches from "frozen" mode to "PWM" mode.

### Bit 3 OC1PE: Output compare 1 preload enable

- 0: Preload register on TIMx\_CCR1 disabled. TIMx\_CCR1 can be written at anytime, the new value is taken in account immediately.
- 1: Preload register on TIMx\_CCR1 enabled. Read/Write operations access the preload register. TIMx\_CCR1 preload value is loaded in the active register at each update event.
- Note: 1: These bits can not be modified as long as LOCK level 3 has been programmed (LOCK bits in TIMx\_BDTR register) and CC1S=00 (the channel is configured in output).
  - 2: The PWM mode can be used without validating the preload register only in one-pulse mode (OPM bit set in TIMx\_CR1 register). Else the behavior is not guaranteed.

### Bit 2 OC1FE: Output compare 1 fast enable

This bit is used to accelerate the effect of an event on the trigger in input on the CC output. 0: CC1 behaves normally depending on counter and CCR1 values even when the trigger is ON. The minimum delay to activate CC1 output when an edge occurs on the trigger input is 5 clock cycles.

1: An active edge on the trigger input acts like a compare match on CC1 output. Then, OC is set to the compare level independently from the result of the comparison. Delay to sample the trigger input and to activate CC1 output is reduced to 3 clock cycles. OCFE acts only if the channel is configured in PWM1 or PWM2 mode.

### Bits 1:0 CC1S: Capture/Compare 1 selection

This bit-field defines the direction of the channel (input/output) as well as the used input.

00: CC1 channel is configured as output.

01: CC1 channel is configured as input, IC1 is mapped on TI1.

10: CC1 channel is configured as input, IC1 is mapped on TI2.

11: CC1 channel is configured as input, IC1 is mapped on TRC. This mode is working only if an internal trigger input is selected through TS bit (TIMx\_SMCR register)

Note: CC1S bits are writable only when the channel is OFF (CC1E = 0 in TIMx\_CCER).

### Input capture mode

```
Bits 15:12 IC2F: Input capture 2 filter
```

Bits 11:10 IC2PSC[1:0]: Input capture 2 prescaler

#### Bits 9:8 CC2S: Capture/compare 2 selection

This bit-field defines the direction of the channel (input/output) as well as the used input.

00: CC2 channel is configured as output.

01: CC2 channel is configured as input, IC2 is mapped on TI2.

10: CC2 channel is configured as input, IC2 is mapped on TI1.

11: CC2 channel is configured as input, IC2 is mapped on TRC. This mode is working only if an internal trigger input is selected through TS bit (TIMx\_SMCR register)

Note: CC2S bits are writable only when the channel is OFF (CC2E = 0 in TIMx\_CCER).

### Bits 7:4 IC1F: Input capture 1 filter

This bit-field defines the frequency used to sample TI1 input and the length of the digital filter applied to TI1. The digital filter is made of an event counter in which N events are needed to validate a transition on the output:

```
0000: No filter, sampling is done at f<sub>DTS</sub>
                                                                       1000: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/8, N=6
                                                                       1001: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/8, N=8
0001: f<sub>SAMPLING</sub>=f<sub>CK INT</sub>, N=2
0010: f<sub>SAMPLING</sub>=f<sub>CK_INT</sub>, N=4
                                                                       1010: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/16, N=5
0011: f<sub>SAMPLING</sub>=f<sub>CK INT</sub>, N=8
                                                                       1011: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/16, N=6
0100: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/2, N=6
                                                                       1100: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/16, N=8
0101: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/2, N=8
                                                                       1101: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/32, N=5
0110: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/4, N=6
                                                                       1110: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/32, N=6
0111: f_{SAMPLING} = f_{DTS}/4, N=8
                                                                       1111: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/32, N=8
```

Note: In current silicon revision,  $f_{DTS}$  is replaced in the formula by CK\_INT when ICxF[3:0]= 1, 2 or 3.

#### Bits 3:2 IC1PSC: Input capture 1 prescaler

This bit-field defines the ratio of the prescaler acting on CC1 input (IC1).

The prescaler is reset as soon as CC1E=0 (TIMx\_CCER register).

00: no prescaler, capture is done each time an edge is detected on the capture input

01: capture is done once every 2 events

10: capture is done once every 4 events

11: capture is done once every 8 events

### Bits 1:0 CC1S: Capture/Compare 1 selection

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This bit-field defines the direction of the channel (input/output) as well as the used input.

00: CC1 channel is configured as output

01: CC1 channel is configured as input, IC1 is mapped on TI1

10: CC1 channel is configured as input, IC1 is mapped on TI2

11: CC1 channel is configured as input, IC1 is mapped on TRC. This mode is working only if an internal trigger input is selected through TS bit (TIMx\_SMCR register)

Note: CC1S bits are writable only when the channel is OFF (CC1E = 0 in TIMx\_CCER).

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# 14.4.8 TIMx capture/compare mode register 2 (TIMx\_CCMR2)

Address offset: 0x1C Reset value: 0x0000

Refer to the above CCMR1 register description.

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	U
OC4CE	(	DC4M[2:0	]	OC4PE	OC4FE	CC4S	2[1:0]	OC3CE	(	OC3M[2:0	]	OC3PE	OC3FE	CC3S	10.13
	IC4F	[3:0]		IC4PS	C[1:0]	0043	5[1.0]		IC3F	[3:0]		IC3PS	C[1:0]	CCS	5[1.0]
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

### **Output compare mode**

Bit 15 **OC4CE:** Output compare 4 clear enable

Bits 14:12 OC4M: Output compare 4 mode

Bit 11 **OC4PE**: Output compare 4 preload enable
Bit 10 **OC4FE**: Output compare 4 fast enable

Bits 9:8 CC4S: Capture/Compare 4 selection

This bit-field defines the direction of the channel (input/output) as well as the used input.

00: CC4 channel is configured as output

01: CC4 channel is configured as input, IC4 is mapped on TI4

10: CC4 channel is configured as input, IC4 is mapped on TI3

11: CC4 channel is configured as input, IC4 is mapped on TRC. This mode is working only if an internal trigger input is selected through TS bit (TIMx\_SMCR register)

Note: CC4S bits are writable only when the channel is OFF (CC4E = 0 in TIMx\_CCER).

Bit 7 OC3CE: Output compare 3 clear enable

Bits 6:4 OC3M: Output compare 3 mode

Bit 3 OC3PE: Output compare 3 preload enable

Bit 2 OC3FE: Output compare 3 fast enable

Bits 1:0 CC3S: Capture/Compare 3 selection

This bit-field defines the direction of the channel (input/output) as well as the used input.

00: CC3 channel is configured as output

01: CC3 channel is configured as input, IC3 is mapped on TI3

10: CC3 channel is configured as input, IC3 is mapped on TI4

11: CC3 channel is configured as input, IC3 is mapped on TRC. This mode is working only if an internal trigger input is selected through TS bit (TIMx\_SMCR register)

Note: CC3S bits are writable only when the channel is OFF (CC3E = 0 in TIMx\_CCER).

## Input capture mode

Bits 15:12 IC4F: Input capture 4 filter

Bits 11:10 IC4PSC: Input capture 4 prescaler

Bits 9:8 CC4S: Capture/Compare 4 selection

This bit-field defines the direction of the channel (input/output) as well as the used input.

00: CC4 channel is configured as output

01: CC4 channel is configured as input, IC4 is mapped on TI4

10: CC4 channel is configured as input, IC4 is mapped on TI3

11: CC4 channel is configured as input, IC4 is mapped on TRC. This mode is working only if an internal trigger input is selected through TS bit (TIMx\_SMCR register)

Note: CC4S bits are writable only when the channel is OFF (CC4E = 0 in TIMx\_CCER).

Bits 7:4 IC3F: Input capture 3 filter

Bits 3:2 IC3PSC: Input capture 3 prescaler

Bits 1:0 CC3S: Capture/Compare 3 selection

This bit-field defines the direction of the channel (input/output) as well as the used input.

00: CC3 channel is configured as output

01: CC3 channel is configured as input, IC3 is mapped on TI3

10: CC3 channel is configured as input, IC3 is mapped on TI4

11: CC3 channel is configured as input, IC3 is mapped on TRC. This mode is working only if an internal trigger input is selected through TS bit (TIMx\_SMCR register)

Note: CC3S bits are writable only when the channel is OFF (CC3E = 0 in TIMx\_CCER).

# 14.4.9 TIMx capture/compare enable register (TIMx\_CCER)

Address offset: 0x20 Reset value: 0x0000

15	14	13	12	11	10	9	8	/	6	5	4	3	2	1	0
CC4NP	Res.	CC4P	CC4E	CC3NP	Res.	CC3P	CC3E	CC2NP	Res.	CC2P	CC2E	CC1NP	Res.	CC1P	CC1E
rw	nes.	rw	rw												

Bit 15 CC4NP: Capture/Compare 4 output Polarity.

Refer to CC1NP description

Bit 14 Reserved, must be kept at reset value.

Bit 13 CC4P: Capture/Compare 4 output Polarity.

refer to CC1P description

Bit 12 CC4E: Capture/Compare 4 output enable.

refer to CC1E description

Bit 11 CC3NP: Capture/Compare 3 output Polarity.

refer to CC1NP description

Bit 10 Reserved, must be kept at reset value.

Bit 9 CC3P: Capture/Compare 3 output Polarity.

refer to CC1P description

Bit 8 CC3E: Capture/Compare 3 output enable.

refer to CC1E description

Bit 7 **CC2NP**: Capture/Compare 2 output Polarity. refer to CC1NP description

Bit 6 Reserved, must be kept at reset value.

Bit 5 CC2P: Capture/Compare 2 output Polarity.

refer to CC1P description

Bit 4 CC2E: Capture/Compare 2 output enable.

refer to CC1E description

Bit 3 CC1NP: Capture/Compare 1 output Polarity.

#### CC1 channel configured as output:

CC1NP must be kept cleared in this case.

### CC1 channel configured as input:

This bit is used in conjunction with CC1P to define TI1FP1/TI2FP1 polarity. refer to CC1P description.

Bit 2 Reserved, must be kept at reset value.

### Bit 1 CC1P: Capture/Compare 1 output Polarity.

### CC1 channel configured as output:

0: OC1 active high

1: OC1 active low

### CC1 channel configured as input:

CC1NP/CC1P bits select TI1FP1 and TI2FP1 polarity for trigger or capture operations.

00: noninverted/rising edge

Circuit is sensitive to TIxFP1 rising edge (capture, trigger in reset, external clock or trigger mode), TIxFP1 is not inverted (trigger in gated mode, encoder mode).

01: inverted/falling edge

Circuit is sensitive to TIxFP1 falling edge (capture, trigger in reset, external clock or trigger mode), TIxFP1 is inverted (trigger in gated mode, encoder mode).

10: reserved, do not use this configuration.

11: noninverted/both edges

Circuit is sensitive to both TIxFP1 rising and falling edges (capture, trigger in reset, external clock or trigger mode), TIxFP1 is not inverted (trigger in gated mode). This configuration must not be used for encoder mode.

Bit 0 CC1E: Capture/Compare 1 output enable.

### CC1 channel configured as output:

0: Off - OC1 is not active

1: On - OC1 signal is output on the corresponding output pin

### CC1 channel configured as input:

This bit determines if a capture of the counter value can actually be done into the input capture/compare register 1 (TIMx\_CCR1) or not.

0: Capture disabled

1: Capture enabled

Table 63. Output control bit for standard OCx channels

CCxE bit	OCx output state
0	Output Disabled (OCx=0, OCx_EN=0)
1	OCx=OCxREF + Polarity, OCx_EN=1

Note:

The state of the external IO pins connected to the standard OCx channels depends on the OCx channel state and the GPIO registers.

# 14.4.10 TIMx counter (TIMx\_CNT)

Address offset: 0x24 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							CNT	[15:0]							
rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bits 15:0 CNT[15:0]: Counter value

# 14.4.11 TIMx prescaler (TIMx\_PSC)

Address offset: 0x28 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							PSC	[15:0]							
rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bits 15:0 PSC[15:0]: Prescaler value

The counter clock frequency CK\_CNT is equal to  $f_{\mbox{CK\_PSC}}\,/\,(\mbox{PSC}[15:0]$  + 1).

PSC contains the value to be loaded in the active prescaler register at each update event.

# 14.4.12 TIMx auto-reload register (TIMx\_ARR)

Address offset: 0x2C

Reset value: 0x0000

15	14	13	12	11	10	9	8	/	6	5	4	3	2	1	0
							ARR	[15:0]	_	_			_	_	
rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bits 15:0 ARR[15:0]: Auto-reloadvalue

ARR is the value to be loaded in the actual auto-reload register.

Refer to the Section 14.3.1: Time-base unit on page 368 for more details about ARR update and behavior.

The counter is blocked while the auto-reload value is null.

# 14.4.13 TIMx capture/compare register 1 (TIMx\_CCR1)

Address offset: 0x34

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
						CCR1[3	1:16] (dep	ending o	n timers)						
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw						
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							CCR1	[15:0]							
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw						

Bits 31:16 CCR1[31:16]: High Capture/Compare 1 value (on TIM2 and TIM5).

Bits 15:0 CCR1[15:0]: Low Capture/Compare 1 value

# If channel CC1 is configured as output:

CCR1 is the value to be loaded in the actual capture/compare 1 register (preload value). It is loaded permanently if the preload feature is not selected in the TIMx\_CCMR1 register (bit OC1PE). Else the preload value is copied in the active capture/compare 1 register when an update event occurs.

The active capture/compare register contains the value to be compared to the counter TIMx\_CNT and signaled on OC1 output.

### If channel CC1is configured as input:

CCR1 is the counter value transferred by the last input capture 1 event (IC1).

# 14.4.14 TIMx capture/compare register 2 (TIMx\_CCR2)

Address offset: 0x38

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
						CCR2[3	1:16] (dep	ending o	n timers)						
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw						
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		-	-		-	-	CCR2	[15:0]			-	-	-		
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw						

Bits 31:16 CCR2[31:16]: High Capture/Compare 2 value (on TIM2 and TIM5).

Bits 15:0 CCR2[15:0]: Low Capture/Compare 2 value

### If channel CC2 is configured as output:

CCR2 is the value to be loaded in the actual capture/compare 2 register (preload value). It is loaded permanently if the preload feature is not selected in the TIMx\_CCMR register (bit OC2PE). Else the preload value is copied in the active capture/compare 2 register when an update event occurs.

The active capture/compare register contains the value to be compared to the counter TIMx\_CNT and signalled on OC2 output.

### If channel CC2 is configured as input:

CCR2 is the counter value transferred by the last input capture 2 event (IC2).

# 14.4.15 TIMx capture/compare register 3 (TIMx\_CCR3)

Address offset: 0x3C

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
						CCR3[3	1:16] (dep	ending o	n timers)						
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw						
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							CCR	8[15:0]							
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw						

Bits 31:16 CCR3[31:16]: High Capture/Compare 3 value (on TIM2 and TIM5).

Bits 15:0 CCR3[15:0]: Low Capture/Compare value

# If channel CC3 is configured as output:

CCR3 is the value to be loaded in the actual capture/compare 3 register (preload value). It is loaded permanently if the preload feature is not selected in the TIMx\_CCMR register (bit OC3PE). Else the preload value is copied in the active capture/compare 3 register when an update event occurs.

The active capture/compare register contains the value to be compared to the counter TIMx\_CNT and signalled on OC3 output.

### If channel CC3is configured as input:

CCR3 is the counter value transferred by the last input capture 3 event (IC3).

# 14.4.16 TIMx capture/compare register 4 (TIMx\_CCR4)

Address offset: 0x40

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
						CCR4[3	1:16] (dep	ending o	n timers)						
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw						
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	_			_			CCR4	[15:0]		_				_	
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw						

Bits 31:16 CCR4[31:16]: High Capture/Compare 4 value (on TIM2 and TIM5).

### Bits 15:0 CCR4[15:0]: Low Capture/Compare value

1. if CC4 channel is configured as output (CC4S bits):

CCR4 is the value to be loaded in the actual capture/compare 4 register (preload value). It is loaded permanently if the preload feature is not selected in the TIMx\_CCMR register (bit OC4PE). Else the preload value is copied in the active capture/compare 4 register when an update event occurs.

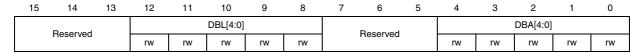
The active capture/compare register contains the value to be compared to the counter TIMx\_CNT and signalled on OC4 output.

if CC4 channel is configured as input (CC4S bits in TIMx\_CCMR4 register):
 CCR4 is the counter value transferred by the last input capture 4 event (IC4).

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# 14.4.17 TIMx DMA control register (TIMx\_DCR)

Address offset: 0x48 Reset value: 0x0000



Bits 15:13 Reserved, must be kept at reset value.

### Bits 12:8 DBL[4:0]: DMA burst length

This 5-bit vector defines the number of DMA transfers (the timer recognizes a burst transfer when a read or a write access is done to the TIMx\_DMAR address).

00000: 1 transfer, 00001: 2 transfers, 00010: 3 transfers,

•••

10001: 18 transfers.

Bits 7:5 Reserved, must be kept at reset value.

### Bits 4:0 DBA[4:0]: DMA base address

This 5-bit vector defines the base-address for DMA transfers (when read/write access are done through the TIMx\_DMAR address). DBA is defined as an offset starting from the address of the TIMx\_CR1 register.

Example:

00000: TIMx\_CR1, 00001: TIMx\_CR2, 00010: TIMx\_SMCR,

...

**Example:** Let us consider the following transfer: DBL = 7 transfers & DBA = TIMx\_CR1. In this case the transfer is done to/from 7 registers starting from the TIMx\_CR1 address.

# 14.4.18 TIMx DMA address for full transfer (TIMx\_DMAR)

Address offset: 0x4C

Reset value: 0x0000

15	14	13	12	11	10	9	8	1	6	5	4	3	2	1	Ü
							DMAE	3[15:0]							
rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bits 15:0 DMAB[15:0]: DMA register for burst accesses

A read or write operation to the DMAR register accesses the register located at the address (TIMx\_CR1 address) + (DBA + DMA index) x 4

where TIMx\_CR1 address is the address of the control register 1, DBA is the DMA base address configured in TIMx\_DCR register, DMA index is automatically controlled by the DMA transfer, and ranges from 0 to DBL (DBL configured in TIMx\_DCR).

## Example of how to use the DMA burst feature

In this example the timer DMA burst feature is used to update the contents of the CCRx registers (x = 2, 3, 4) with the DMA transferring half words into the CCRx registers.

This is done in the following steps:

- 1. Configure the corresponding DMA channel as follows:
  - DMA channel peripheral address is the DMAR register address
  - DMA channel memory address is the address of the buffer in the RAM containing the data to be transferred by DMA into CCRx registers.
  - Number of data to transfer = 3 (See note below).
  - Circular mode disabled.
- 2. Configure the DCR register by configuring the DBA and DBL bit fields as follows: DBL = 3 transfers, DBA = 0xE.
- 3. Enable the TIMx update DMA request (set the UDE bit in the DIER register).
- 4. Enable TIMx
- 5. Enable the DMA channel

Note:

This example is for the case where every CCRx register to be updated once. If every CCRx register is to be updated twice for example, the number of data to transfer should be 6. Let's take the example of a buffer in the RAM containing data1, data2, data3, data4, data5 and data6. The data is transferred to the CCRx registers as follows: on the first update DMA request, data1 is transferred to CCR2, data2 is transferred to CCR3, data3 is transferred to CCR4 and on the second update DMA request, data4 is transferred to CCR2, data5 is transferred to CCR3 and data6 is transferred to CCR4.

# 14.4.19 TIM2 option register (TIM2\_OR)

Address offset: 0x50 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Reserved			ITR1_	_RMP					Rese	vr.rod				
	nest	erveu		rw	rw					nese	erveu				

Bits 15:12 Reserved, must be kept at reset value.

Bits 11:10 ITR1\_RMP: Internal trigger 1 remap

Set and cleared by software.

00: TIM8\_TRGOUT

01: PTP trigger output is connected to TIM2\_ITR110: OTG FS SOF is connected to the TIM2\_ITR1 input11: OTG HS SOF is connected to the TIM2\_ITR1 input

Bits 9:0 Reserved, must be kept at reset value.

# 14.4.20 TIM5 option register (TIM5\_OR)

Address offset: 0x50 Reset value: 0x0000



Bits 15:8 Reserved, must be kept at reset value.

Bits 7:6 TI4\_RMP: Timer Input 4 remap

Set and cleared by software.

00: TIM5 Channel4 is connected to the GPIO: Refer to the Alternate function mapping table in the datasheets.

01: the LSI internal clock is connected to the TIM5\_CH4 input for calibration purposes10: the LSE internal clock is connected to the TIM5\_CH4 input for calibration purposes11: the RTC output event is connected to the TIM5\_CH4 input for calibration purposes

Bits 5:0 Reserved, must be kept at reset value.

# 14.4.21 TIMx register map

TIMx registers are mapped as described in the table below:

Table 64. TIM2 to TIM5 register map and reset values

Offset	Register	31	30	29	28	27	26	25	24		23	22	21	17	20	19	18	17	16	15	14	13	12	11	10	σ	α	7	. 9	נט	, <	4	3	2	1	0
0x00	TIMx_CR1												Res	ser	ved												CKD 1:0]	ARPE	[	MS 1:0]	ald	֓֞֞֞֓֓֓֓֓֓֓֓֓֓֓֓֓֓֓֓֓֓֓֓֓֓֓֓֓֓֓֓֓֓֓֓֓	ОРМ	URS	NDIS	CEN
	Reset value																									0	0	0	0	0	(	)	0	0	0	0
0x04	TIMx_CR2													Re	eser	vec	t											TI1S		IMS		D] [	CCDS	Re	serv	/ed
	Reset value																											0	0	0	(	)	0			
0x08	TIMx_SMCR								Res	er	vec	d								ETP	ECE		PS :0]		ETF	[3:	0]	MSM	-	TS[2	2:0]		Reserved	SN	1S[2	:0]
	Reset value																			0	0	0	0	0	0	0		0	0	0	(	) )	Re	0	0	0
0x0C	TIMx_DIER								Re	ese	erv	ec.	i								TDE	COMDE	CC4DE	CC3DE	CC2DE	CC1DE	UDE	Reserved	H	Reserved	CCAIF	) () 1	CC3IE	CC2IE	CC11E	NIE
	Reset value																				0	0	0	0	0	0		Ä	0	æ		0	0	0	0	0
0x10	TIMx_SR									F	Res	se	rve	t									CC40F	CC30F	CC2OF	CC10F		Reserved	ΗH	Reserved	CCAIF	1100	CC3IF	CC2IF	CC1IF	UIF
	Reset value																						0	0	0	0		Ä_	0	Ä	(	)	0	0	0	0
0x14	TIMx_EGR													ı	Res	erv	ed												TG	SS	CC4G			CC2G	O	ne
	Reset value																												0	æ	(	-	0	0	0	0
	TIMx_CCMR1 Output Compare mode								Res	er	vec	d								OC2CE		OC2 [2:0	]	OC2PE	OC2FE		C2S 1:0]			OC <sup>-</sup> [2:0	0]	1		OC1FE	[1	:1S :0]
0x18	Reset value																			0	0	0	0	0		0	0	0	0	0	(	)	0	0	0	0
	TIMx_CCMR1 Input Capture mode								Res	er	vec	d									IC2I	=[3:0	)]	Р	C2 SC :0]		C2S 1:0]	3	IC1	F[3:	.0]		IC PS [1:	SC	CC [1	:1S :0]
	Reset value																			0	0	0	0	0	0	0	0	0	0	0	(	)	0	0	0	0

Table 64. TIM2 to TIM5 register map and reset values (continued)

Offset	Register	+ 1	0 1 2 3 4 5 6 6 8 8 9 9 1 1 1 2 1 3 1 4 5 1 5 1 5 1 5 1 5 1 5 1 5 1 5 1 5 1
	TIMx_CCMR2		
	Output Compare	Reserved	OC4M
	mode Reset value		
0x1C	TIMx_CCMR2		10/1
	Input Capture	Reserved	
	mode	r teser veu	
	Reset value		0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
	TIMx_CCER		CCAP CCAP CCAP CCAP CCAP CCAP CCAP CCAP
0x20		Reserved	Reserved CC4F CC3NP Reserved CC3NP CC3P CC3P CC3P CC2P CC2P CC2NP Reserved CC2NP CC2P CC2NP CC2P CC2NP CC2P CC2NP CC2P CC2
	Reset value		0 000 0 0 0 0 0 0 0
		CNT[31:16]	
0x24	TIMx_CNT	(TIM2 and TIM5 only, reserved on the other timers)	CNT[15:0]
	Reset value	   0   0   0   0   0   0   0   0   0	
0x28	TIMx_PSC	Reserved	PSC[15:0]
	Reset value		0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
		ARR[31:16]	
0x2C	TIMx_ARR	(TIM2 and TIM5 only, reserved on the other timers)	ARR[15:0]
	Reset value	   0   0   0   0   0   0   0   0   0	
0x30		Reserved	
		CCR1[31:16]	
0x34	TIMx_CCR1	(TIM2 and TIM5 only, reserved on the other timers)	CCR1[15:0]
	Reset value		
	TIM-: CODO	CCR2[31:16]	0000(45.0)
0x38	TIMx_CCR2	(TIM2 and TIM5 only, reserved on the other timers)	CCR2[15:0]
	Reset value	0   0   0   0   0   0   0   0   0   0	
	TIM-: CODO	CCR3[31:16]	0000(45-0)
0x3C	TIMx_CCR3	(TIM2 and TIM5 only, reserved on the other timers)	CCR3[15:0]
	Reset value	0   0   0   0   0   0   0   0   0   0	
	TIM-: 00D4	CCR4[31:16]	000445-03
0x40	TIMx_CCR4	(TIM2 and TIM5 only, reserved on the other timers)	CCR4[15:0]
	Reset value	0   0   0   0   0   0   0   0   0   0	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
0x44		Reserved	
0,44			
	TIMx_DCR		DBL[4:0] DBA[4:0]
0x48	TIMX_DON	Reserved	Reserved BBA[4.0]
	Reset value		0 0 0 0 0 0 0 0 0 0 0
0x4C	TIMx_DMAR	Reserved	DMAB[15:0]
0.40		nesei veu	
	Reset value		0   0   0   0   0   0   0   0   0   0
0x50	TIM2_OR	Reserved	Reserved RMP Reserved
	Reset value		0 0
0x50	TIM5_OR	Reserved	Reserved IT4_R MP Reserved
0,00	Reset value	i iodol vod	0 0

Refer to *Table 2 on page 50* for the register boundary addresses.

# 15 General-purpose timers (TIM9 to TIM14)

This section applies to the whole STM32F4xx family unless otherwise specified.

## 15.1 TIM9 to TIM14 introduction

The TIM9 to TIM14 general-purpose timers consist of a 16-bit auto-reload counter driven by a programmable prescaler.

They may be used for a variety of purposes, including measuring the pulse lengths of input signals (input capture) or generating output waveforms (output compare, PWM).

Pulse lengths and waveform periods can be modulated from a few microseconds to several milliseconds using the timer prescaler and the RCC clock controller prescalers.

The TIM9 to TIM14 timers are completely independent, and do not share any resources. They can be synchronized together as described in *Section 15.4.12*.

### 15.2 TIM9 to TIM14 main features

### 15.2.1 TIM9/TIM12 main features

The features of the TIM9 to TIM14 general-purpose timers include:

- 16-bit auto-reload upcounter (in medium density devices)
- 16-bit programmable prescaler used to divide the counter clock frequency by any factor between 1 and 65536 (can be changed "on the fly")
- Up to 2 independent channels for:
  - Input capture
  - Output compare
  - PWM generation (edge-aligned mode)
  - One-pulse mode output
- Synchronization circuit to control the timer with external signals and to interconnect several timers together
- Interrupt generation on the following events:
  - Update: counter overflow, counter initialization (by software or internal trigger)
  - Trigger event (counter start, stop, initialization or count by internal trigger)
  - Input capture
  - Output compare

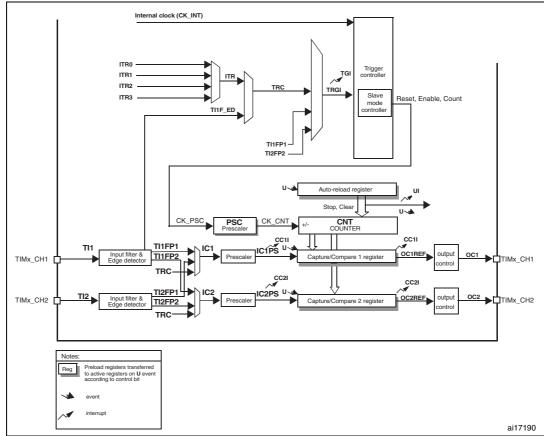


Figure 159. General-purpose timer block diagram (TIM9 and TIM12)

# 15.3 TIM10/TIM11 and TIM13/TIM14 main features

The features of general-purpose timers TIM10/TIM11 and TIM13/TIM14 include:

- 16-bit auto-reload upcounter
- 16-bit programmable prescaler used to divide the counter clock frequency by any factor between 1 and 65536 (can be changed "on the fly")
- independent channel for:
  - Input capture
  - Output compare
  - PWM generation (edge-aligned mode)
  - One-pulse mode output
- Interrupt generation on the following events:
  - Update: counter overflow, counter initialization (by software)
  - Input capture
  - Output compare

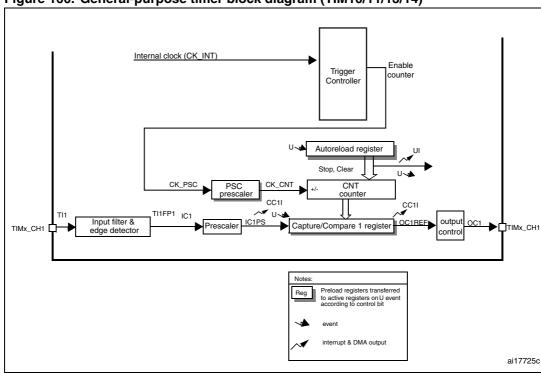


Figure 160. General-purpose timer block diagram (TIM10/11/13/14)

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# 15.4 TIM9 to TIM14 functional description

### 15.4.1 Time-base unit

The main block of the timer is a 16-bit counter with its related auto-reload register. The counters counts up.

The counter clock can be divided by a prescaler.

The counter, the auto-reload register and the prescaler register can be written or read by software. This is true even when the counter is running.

The time-base unit includes:

- Counter register (TIMx\_CNT)
- Prescaler register (TIMx\_PSC)
- Auto-reload register (TIMx\_ARR)

The auto-reload register is preloaded. Writing to or reading from the auto-reload register accesses the preload register. The content of the preload register are transferred into the shadow register permanently or at each update event (UEV), depending on the auto-reload preload enable bit (ARPE) in TIMx\_CR1 register. The update event is sent when the counter reaches the overflow and if the UDIS bit equals 0 in the TIMx\_CR1 register. It can also be generated by software. The generation of the update event is described in details for each configuration.

The counter is clocked by the prescaler output CK\_CNT, which is enabled only when the counter enable bit (CEN) in TIMx\_CR1 register is set (refer also to the slave mode controller description to get more details on counter enabling).

Note that the counter starts counting 1 clock cycle after setting the CEN bit in the TIMx\_CR1 register.

### **Prescaler description**

The prescaler can divide the counter clock frequency by any factor between 1 and 65536. It is based on a 16-bit counter controlled through a 16-bit register (in the TIMx\_PSC register). It can be changed on the fly as this control register is buffered. The new prescaler ratio is taken into account at the next update event.

*Figure 161* and *Figure 162* give some examples of the counter behavior when the prescaler ratio is changed on the fly.

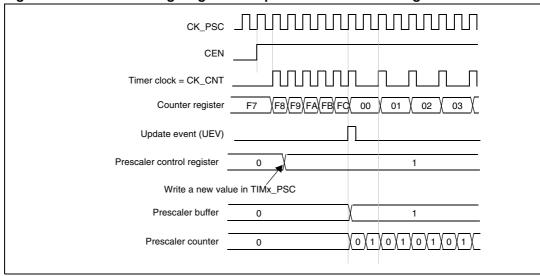
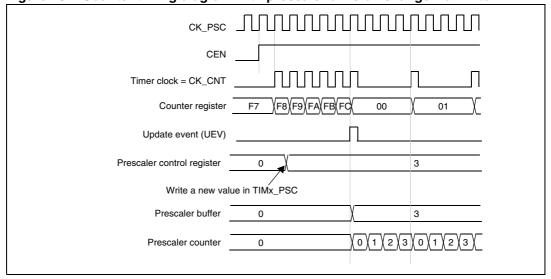


Figure 161. Counter timing diagram with prescaler division change from 1 to 2





# 15.4.2 Counter modes

### **Upcounting mode**

In upcounting mode, the counter counts from 0 to the auto-reload value (content of the TIMx\_ARR register), then restarts from 0 and generates a counter overflow event.

Setting the UG bit in the TIMx\_EGR register (by software or by using the slave mode controller on TIM9 and TIM12) also generates an update event.

The UEV event can be disabled by software by setting the UDIS bit in the TIMx\_CR1 register. This is to avoid updating the shadow registers while writing new values in the preload registers. Then no update event occurs until the UDIS bit has been written to 0. However, the counter restarts from 0, as well as the counter of the prescaler (but the prescale rate does not change). In addition, if the URS bit (update request selection) in TIMx\_CR1 register is set, setting the UG bit generates an update event UEV but without

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setting the UIF flag (thus no interrupt is sent). This is to avoid generating both update and capture interrupts when clearing the counter on the capture event.

When an update event occurs, all the registers are updated and the update flag (UIF bit in TIMx SR register) is set (depending on the URS bit):

- The auto-reload shadow register is updated with the preload value (TIMx\_ARR),
- The buffer of the prescaler is reloaded with the preload value (content of the TIMx\_PSC register).

The following figures show some examples of the counter behavior for different clock frequencies when TIMx\_ARR=0x36.

Figure 163. Counter timing diagram, internal clock divided by 1

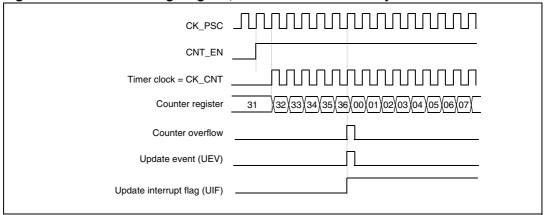


Figure 164. Counter timing diagram, internal clock divided by 2

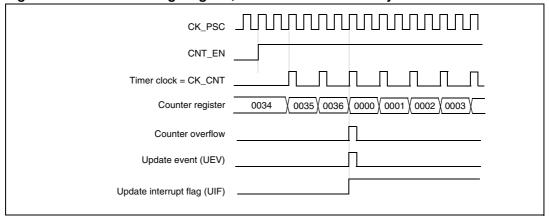


Figure 165. Counter timing diagram, internal clock divided by 4

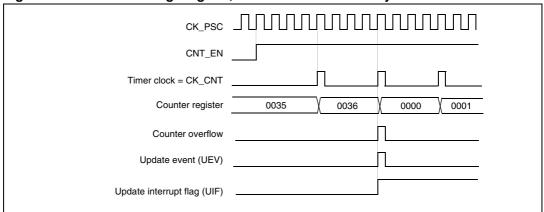


Figure 166. Counter timing diagram, internal clock divided by N

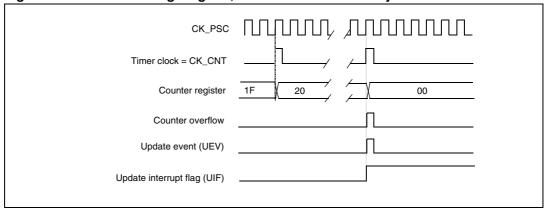
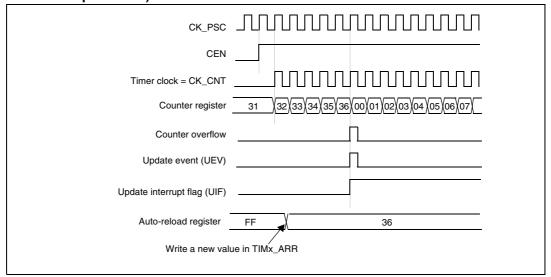


Figure 167. Counter timing diagram, update event when ARPE=0 (TIMx\_ARR not preloaded)



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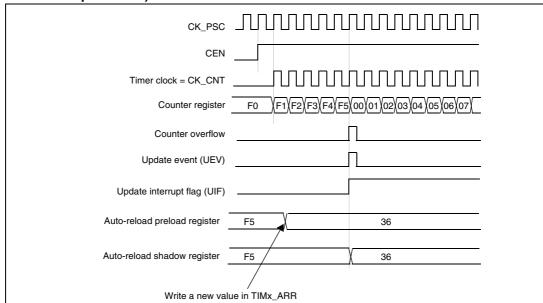


Figure 168. Counter timing diagram, update event when ARPE=1 (TIMx\_ARR preloaded)

### 15.4.3 Clock selection

The counter clock can be provided by the following clock sources:

- Internal clock (CK\_INT)
- External clock mode1 (for TIM9 and TIM12): external input pin (Tlx)
- Internal trigger inputs (ITRx) (for TIM9 and TIM12): connecting the trigger output from another timer. Refer to Section: Using one timer as prescaler for another for more details.

### Internal clock source (CK\_INT)

The internal clock source is the default clock source for TIM10/TIM11 and TIM13/TIM14.

For TIM9 and TIM12, the internal clock source is selected when the slave mode controller is disabled (SMS='000'). The CEN bit in the TIMx\_CR1 register and the UG bit in the TIMx\_EGR register are then used as control bits and can be changed only by software (except for UG which remains cleared). As soon as the CEN bit is programmed to 1, the prescaler is clocked by the internal clock CK\_INT.

*Figure 169* shows the behavior of the control circuit and the upcounter in normal mode, without prescaler.

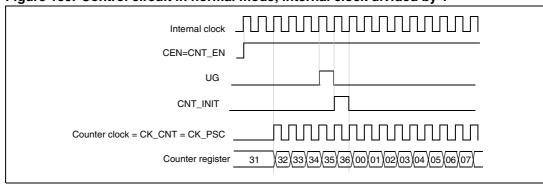
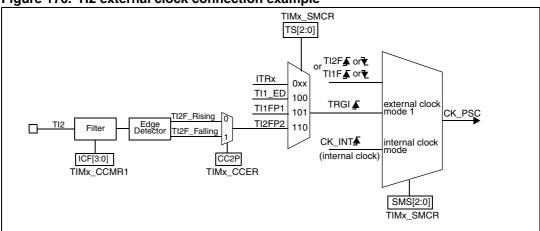


Figure 169. Control circuit in normal mode, internal clock divided by 1

### External clock source mode 1(TIM9 and TIM12)

This mode is selected when SMS='111' in the TIMx\_SMCR register. The counter can count at each rising or falling edge on a selected input.

Figure 170. TI2 external clock connection example



For example, to configure the upcounter to count in response to a rising edge on the TI2 input, use the following procedure:

- 1. Configure channel 2 to detect rising edges on the TI2 input by writing CC2S = '01' in the TIMx\_CCMR1 register.
- 2. Configure the input filter duration by writing the IC2F[3:0] bits in the TIMx\_CCMR1 register (if no filter is needed, keep IC2F='0000').
- 3. Select the rising edge polarity by writing CC2P='0' and CC2NP='0' in the TIMx\_CCER register.
- 4. Configure the timer in external clock mode 1 by writing SMS='111' in the TIMx\_SMCR register.
- 5. Select TI2 as the trigger input source by writing TS='110' in the TIMx SMCR register.
- 6. Enable the counter by writing CEN='1' in the TIMx\_CR1 register.

Note: The capture prescaler is not used for triggering, so you don't need to configure it.

When a rising edge occurs on TI2, the counter counts once and the TIF flag is set.

The delay between the rising edge on TI2 and the actual clock of the counter is due to the resynchronization circuit on TI2 input.

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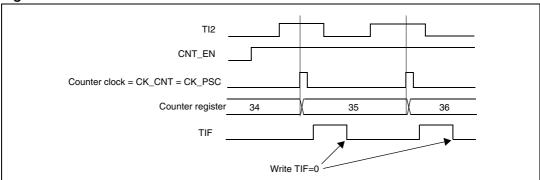


Figure 171. Control circuit in external clock mode 1

### 15.4.4 Capture/compare channels

Each Capture/Compare channel is built around a capture/compare register (including a shadow register), a input stage for capture (with digital filter, multiplexing and prescaler) and an output stage (with comparator and output control).

Figure 172 to Figure 174 give an overview of one capture/compare channel.

The input stage samples the corresponding TIx input to generate a filtered signal TIxF. Then, an edge detector with polarity selection generates a signal (TIxFPx) which can be used as trigger input by the slave mode controller or as the capture command. It is prescaled before the capture register (ICxPS).

TI1F\_Risir Edge TI1FP1 01 TI1F\_Falling downcounte TI2FP1 IC1PS IC1 divider 10 /1, /2, /4, /8 CC1P/CC1NP ICF[3:0] TRC (from slave mode controller) TIMx\_CCMR1 TIMx\_CCER TI2F\_rising (from channel 2) CC1E CC1S[1:0] ICPS[1:0] TI2F\_falling (from channel 2) TIMx CCMR1 TIMx\_CCER

Figure 172. Capture/compare channel (example: channel 1 input stage)

The output stage generates an intermediate waveform which is then used for reference: OCxRef (active high). The polarity acts at the end of the chain.

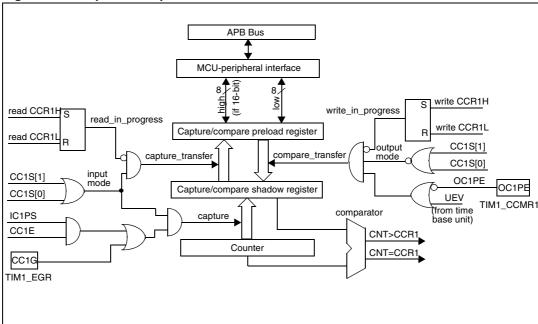
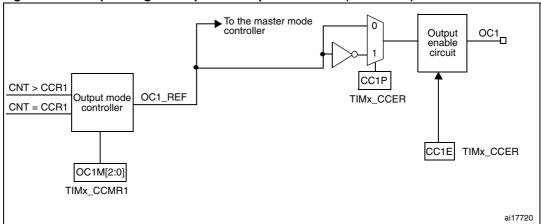


Figure 173. Capture/compare channel 1 main circuit

Figure 174. Output stage of capture/compare channel (channel 1)



The capture/compare block is made of one preload register and one shadow register. Write and read always access the preload register.

In capture mode, captures are actually done in the shadow register, which is copied into the preload register.

In compare mode, the content of the preload register is copied into the shadow register which is compared to the counter.

### 15.4.5 Input capture mode

In Input capture mode, the Capture/Compare Registers (TIMx\_CCRx) are used to latch the value of the counter after a transition detected by the corresponding ICx signal. When a capture occurs, the corresponding CCXIF flag (TIMx\_SR register) is set and an interrupt or a DMA request can be sent if they are enabled. If a capture occurs while the CCxIF flag was already high, then the over-capture flag CCxOF (TIMx\_SR register) is set. CCxIF can be

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cleared by software by writing it to '0' or by reading the captured data stored in the TIMx\_CCRx register. CCxOF is cleared when you write it to '0'.

The following example shows how to capture the counter value in TIMx\_CCR1 when TI1 input rises. To do this, use the following procedure:

- Select the active input: TIMx\_CCR1 must be linked to the TI1 input, so write the CC1S bits to '01' in the TIMx\_CCMR1 register. As soon as CC1S becomes different from '00', the channel is configured in input mode and the TIMx\_CCR1 register becomes readonly.
- 2. Program the input filter duration you need with respect to the signal you connect to the timer (by programming the ICxF bits in the TIMx\_CCMRx register if the input is one of the TIx inputs). Let's imagine that, when toggling, the input signal is not stable during at must 5 internal clock cycles. We must program a filter duration longer than these 5 clock cycles. We can validate a transition on TI1 when 8 consecutive samples with the new level have been detected (sampled at f<sub>DTS</sub> frequency). Then write IC1F bits to '0011' in the TIMx\_CCMR1 register.
- 3. Select the edge of the active transition on the TI1 channel by programming CC1P and CC1NP bits to '00' in the TIMx CCER register (rising edge in this case).
- 4. Program the input prescaler. In our example, we wish the capture to be performed at each valid transition, so the prescaler is disabled (write IC1PS bits to '00' in the TIMx\_CCMR1 register).
- 5. Enable capture from the counter into the capture register by setting the CC1E bit in the TIMx\_CCER register.
- 6. If needed, enable the related interrupt request by setting the CC1IE bit in the TIMx\_DIER register.

When an input capture occurs:

- The TIMx\_CCR1 register gets the value of the counter on the active transition.
- CC1IF flag is set (interrupt flag). CC1OF is also set if at least two consecutive captures occurred whereas the flag was not cleared.
- An interrupt is generated depending on the CC1IE bit.

In order to handle the overcapture, it is recommended to read the data before the overcapture flag. This is to avoid missing an overcapture which could happen after reading the flag and before reading the data.

Note:

IC interrupt requests can be generated by software by setting the corresponding CCxG bit in the TIMx\_EGR register.

### 15.4.6 PWM input mode (only for TIM9/12)

This mode is a particular case of input capture mode. The procedure is the same except:

- Two ICx signals are mapped on the same TIx input.
- These 2 ICx signals are active on edges with opposite polarity.
- One of the two TIxFP signals is selected as trigger input and the slave mode controller is configured in reset mode.

For example, you can measure the period (in TIMx\_CCR1 register) and the duty cycle (in TIMx\_CCR2 register) of the PWM applied on TI1 using the following procedure (depending on CK\_INT frequency and prescaler value):

- Select the active input for TIMx CCR1: write the CC1S bits to '01' in the TIMx CCMR1 register (TI1 selected).
- 2. Select the active polarity for TI1FP1 (used both for capture in TIMx\_CCR1 and counter clear): program the CC1P and CC1NP bits to '00' (active on rising edge).
- Select the active input for TIMx CCR2: write the CC2S bits to '10' in the TIMx CCMR1 register (TI1 selected).
- Select the active polarity for TI1FP2 (used for capture in TIMx CCR2): program the CC2P and CC2NP bits to '11' (active on falling edge).
- 5. Select the valid trigger input: write the TS bits to '101' in the TIMx SMCR register (TI1FP1 selected).
- Configure the slave mode controller in reset mode: write the SMS bits to '100' in the TIMx\_SMCR register.
- Enable the captures: write the CC1E and CC2E bits to '1' in the TIMx\_CCER register. 7.

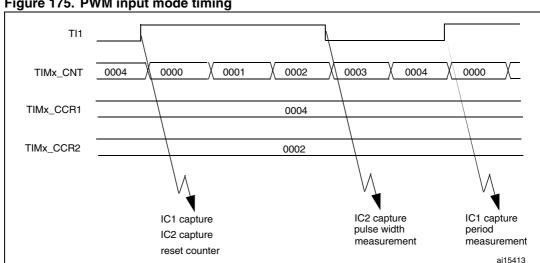


Figure 175. PWM input mode timing

The PWM input mode can be used only with the TIMx\_CH1/TIMx\_CH2 signals due to the fact that only TI1FP1 and TI2FP2 are connected to the slave mode controller.

#### 15.4.7 Forced output mode

In output mode (CCxS bits = '00' in the TIMx CCMRx register), each output compare signal (OCxREF and then OCx) can be forced to active or inactive level directly by software, independently of any comparison between the output compare register and the counter.

To force an output compare signal (OCXREF/OCx) to its active level, you just need to write '101' in the OCxM bits in the corresponding TIMx CCMRx register. Thus OCXREF is forced high (OCxREF is always active high) and OCx get opposite value to CCxP polarity bit.

For example: CCxP='0' (OCx active high) => OCx is forced to high level.

The OCxREF signal can be forced low by writing the OCxM bits to '100' in the TIMx CCMRx register.

Anyway, the comparison between the TIMx CCRx shadow register and the counter is still performed and allows the flag to be set. Interrupt requests can be sent accordingly. This is described in the output compare mode section below.

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### 15.4.8 Output compare mode

This function is used to control an output waveform or indicating when a period of time has elapsed.

When a match is found between the capture/compare register and the counter, the output compare function:

- Assigns the corresponding output pin to a programmable value defined by the output compare mode (OCxM bits in the TIMx\_CCMRx register) and the output polarity (CCxP bit in the TIMx\_CCER register). The output pin can keep its level (OCXM='000'), be set active (OCxM='001'), be set inactive (OCxM='010') or can toggle (OCxM='011') on match.
- 2. Sets a flag in the interrupt status register (CCxIF bit in the TIMx\_SR register).
- Generates an interrupt if the corresponding interrupt mask is set (CCXIE bit in the TIMx\_DIER register).

The TIMx\_CCRx registers can be programmed with or without preload registers using the OCxPE bit in the TIMx\_CCMRx register.

In output compare mode, the update event UEV has no effect on OCxREF and OCx output. The timing resolution is one count of the counter. Output compare mode can also be used to output a single pulse (in One-pulse mode).

#### Procedure:

- 1. Select the counter clock (internal, external, prescaler).
- 2. Write the desired data in the TIMx\_ARR and TIMx\_CCRx registers.
- 3. Set the CCxIE bit if an interrupt request is to be generated.
- 4. Select the output mode. For example:
  - Write OCxM = '011' to toggle OCx output pin when CNT matches CCRx
  - Write OCxPE = '0' to disable preload register
  - Write CCxP = '0' to select active high polarity
  - Write CCxE = '1' to enable the output
- 5. Enable the counter by setting the CEN bit in the TIMx\_CR1 register.

The TIMx\_CCRx register can be updated at any time by software to control the output waveform, provided that the preload register is not enabled (OCxPE='0', else TIMx\_CCRx shadow register is updated only at the next update event UEV). An example is given in *Figure 176*.

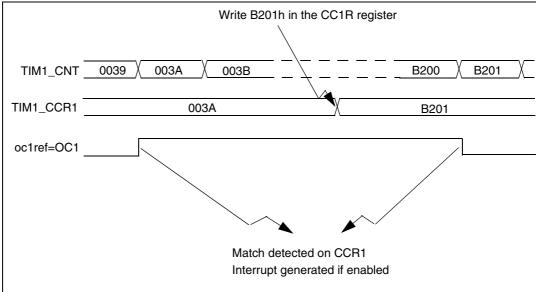


Figure 176. Output compare mode, toggle on OC1.

### 15.4.9 **PWM mode**

Pulse Width Modulation mode allows you to generate a signal with a frequency determined by the value of the TIMx\_ARR register and a duty cycle determined by the value of the TIMx\_CCRx register.

The PWM mode can be selected independently on each channel (one PWM per OCx output) by writing '110' (PWM mode 1) or '111' (PWM mode 2) in the OCxM bits in the TIMx\_CCMRx register. You must enable the corresponding preload register by setting the OCxPE bit in the TIMx\_CCMRx register, and eventually the auto-reload preload register (in upcounting or center-aligned modes) by setting the ARPE bit in the TIMx\_CR1 register.

As the preload registers are transferred to the shadow registers only when an update event occurs, before starting the counter, you have to initialize all the registers by setting the UG bit in the TIMx\_EGR register.

The OCx polarity is software programmable using the CCxP bit in the TIMx\_CCER register. It can be programmed as active high or active low. The OCx output is enabled by the CCxE bit in the TIMx\_CCER register. Refer to the TIMx\_CCERx register description for more details.

In PWM mode (1 or 2), TIMx\_CNT and TIMx\_CCRx are always compared to determine whether TIMx\_CNT  $\leq$  TIMx\_CCRx.

The timer is able to generate PWM in edge-aligned mode only since the counter is upcounting.

### PWM edge-aligned mode

In the following example, we consider PWM mode 1. The reference PWM signal OCxREF is high as long as TIMx\_CNT < TIMx\_CCRx else it becomes low. If the compare value in TIMx\_CCRx is greater than the auto-reload value (in TIMx\_ARR) then OCxREF is held at '1'. If the compare value is 0 then OCxRef is held at '0'. *Figure 177* shows some edgealigned PWM waveforms in an example where TIMx\_ARR=8.

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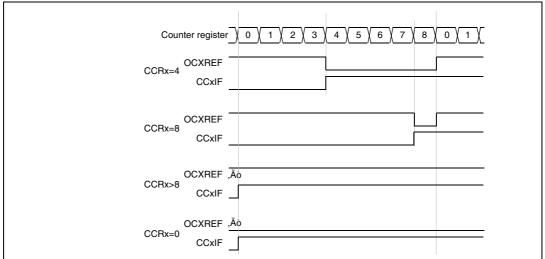


Figure 177. Edge-aligned PWM waveforms (ARR=8)

### 15.4.10 One-pulse mode

One-pulse mode (OPM) is a particular case of the previous modes. It allows the counter to be started in response to a stimulus and to generate a pulse with a programmable length after a programmable delay.

Starting the counter can be controlled through the slave mode controller. Generating the waveform can be done in output compare mode or PWM mode. You select One-pulse mode by setting the OPM bit in the TIMx\_CR1 register. This makes the counter stop automatically at the next update event UEV.

A pulse can be correctly generated only if the compare value is different from the counter initial value. Before starting (when the timer is waiting for the trigger), the configuration must be as follows:

CNT < CCRx≤ ARR (in particular, 0 < CCRx)

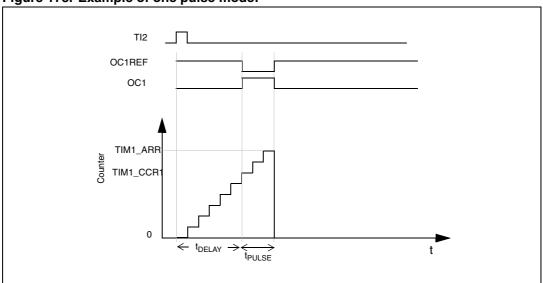


Figure 178. Example of one pulse mode.

For example you may want to generate a positive pulse on OC1 with a length of  $t_{PULSE}$  and after a delay of  $t_{DELAY}$  as soon as a positive edge is detected on the TI2 input pin.

Use TI2FP2 as trigger 1:

- 1. Map TI2FP2 to TI2 by writing CC2S='01' in the TIMx\_CCMR1 register.
- 2. TI2FP2 must detect a rising edge, write CC2P='0' and CC2NP = '0' in the TIMx\_CCER register.
- 3. Configure TI2FP2 as trigger for the slave mode controller (TRGI) by writing TS='110' in the TIMx\_SMCR register.
- 4. TI2FP2 is used to start the counter by writing SMS to '110' in the TIMx\_SMCR register (trigger mode).

The OPM waveform is defined by writing the compare registers (taking into account the clock frequency and the counter prescaler).

- The t<sub>DELAY</sub> is defined by the value written in the TIMx\_CCR1 register.
- The t<sub>PULSE</sub> is defined by the difference between the auto-reload value and the compare value (TIMx\_ARR - TIMx\_CCR1).
- Let's say you want to build a waveform with a transition from '0' to '1' when a compare match occurs and a transition from '1' to '0' when the counter reaches the auto-reload value. To do this you enable PWM mode 2 by writing OC1M='111' in the TIMx\_CCMR1 register. You can optionally enable the preload registers by writing OC1PE='1' in the TIMx\_CCMR1 register and ARPE in the TIMx\_CR1 register. In this case you have to write the compare value in the TIMx\_CCR1 register, the auto-reload value in the TIMx\_ARR register, generate an update by setting the UG bit and wait for external trigger event on TI2. CC1P is written to '0' in this example.

You only want 1 pulse (Single mode), so you write '1 in the OPM bit in the TIMx\_CR1 register to stop the counter at the next update event (when the counter rolls over from the auto-reload value back to 0). When OPM bit in the TIMx\_CR1 register is set to '0', so the Repetitive Mode is selected.

### Particular case: OCx fast enable

In One-pulse mode, the edge detection on TIx input set the CEN bit which enables the counter. Then the comparison between the counter and the compare value makes the output toggle. But several clock cycles are needed for these operations and it limits the minimum delay  $t_{DELAY}$  min we can get.

If you want to output a waveform with the minimum delay, you can set the OCxFE bit in the TIMx\_CCMRx register. Then OCxRef (and OCx) are forced in response to the stimulus, without taking in account the comparison. Its new level is the same as if a compare match had occurred. OCxFE acts only if the channel is configured in PWM1 or PWM2 mode.

### 15.4.11 TIM9/12 external trigger synchronization

The TIM9/12 timers can be synchronized with an external trigger in several modes: Reset mode, Gated mode and Trigger mode.

### Slave mode: Reset mode

The counter and its prescaler can be reinitialized in response to an event on a trigger input. Moreover, if the URS bit from the TIMx\_CR1 register is low, an update event UEV is generated. Then all the preloaded registers (TIMx\_ARR, TIMx\_CCRx) are updated.

In the following example, the upcounter is cleared in response to a rising edge on TI1 input:

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- Configure the channel 1 to detect rising edges on TI1. Configure the input filter duration (in this example, we don't need any filter, so we keep IC1F='0000'). The capture prescaler is not used for triggering, so you don't need to configure it. The CC1S bits select the input capture source only, CC1S = '01' in the TIMx CCMR1 register. Program CC1P and CC1NP to '00' in TIMx\_CCER register to validate the polarity (and detect rising edges only).
- Configure the timer in reset mode by writing SMS='100' in TIMx SMCR register. Select TI1 as the input source by writing TS='101' in TIMx SMCR register.
- Start the counter by writing CEN='1' in the TIMx CR1 register.

The counter starts counting on the internal clock, then behaves normally until TI1 rising edge. When TI1 rises, the counter is cleared and restarts from 0. In the meantime, the trigger flag is set (TIF bit in the TIMx\_SR register) and an interrupt request can be sent if enabled (depending on the TIE bit in TIMx\_DIER register).

The following figure shows this behavior when the auto-reload register TIMx ARR=0x36. The delay between the rising edge on TI1 and the actual reset of the counter is due to the resynchronization circuit on TI1 input.

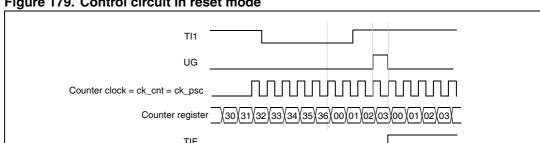


Figure 179. Control circuit in reset mode

#### Slave mode: Gated mode

The counter can be enabled depending on the level of a selected input.

In the following example, the upcounter counts only when TI1 input is low:

- Configure the channel 1 to detect low levels on TI1. Configure the input filter duration (in this example, we don't need any filter, so we keep IC1F='0000'). The capture prescaler is not used for triggering, so you don't need to configure it. The CC1S bits select the input capture source only, CC1S='01' in TIMx\_CCMR1 register. Program CC1P='1' and CC1NP= '0' in TIMx\_CCER register to validate the polarity (and detect low level only).
- Configure the timer in gated mode by writing SMS='101' in TIMx\_SMCR register. Select TI1 as the input source by writing TS='101' in TIMx SMCR register.
- Enable the counter by writing CEN='1' in the TIMx\_CR1 register (in gated mode, the counter doesn't start if CEN='0', whatever is the trigger input level).

The counter starts counting on the internal clock as long as TI1 is low and stops as soon as TI1 becomes high. The TIF flag in the TIMx\_SR register is set both when the counter starts or stops.

The delay between the rising edge on TI1 and the actual stop of the counter is due to the resynchronization circuit on TI1 input.

Figure 180. Control circuit in gated mode

### Slave mode: Trigger mode

The counter can start in response to an event on a selected input.

In the following example, the upcounter starts in response to a rising edge on TI2 input:

- Configure the channel 2 to detect rising edges on TI2. Configure the input filter duration (in this example, we don't need any filter, so we keep IC2F='0000'). The capture prescaler is not used for triggering, so you don't need to configure it. The CC2S bits are configured to select the input capture source only, CC2S='01' in TIMx\_CCMR1 register. Program CC2P='1' and CC2NP='0' in TIMx\_CCER register to validate the polarity (and detect low level only).
- Configure the timer in trigger mode by writing SMS='110' in TIMx\_SMCR register.
   Select TI2 as the input source by writing TS='110' in TIMx\_SMCR register.

When a rising edge occurs on TI2, the counter starts counting on the internal clock and the TIF flag is set.

The delay between the rising edge on TI2 and the actual start of the counter is due to the resynchronization circuit on TI2 input.

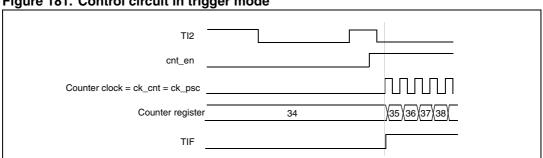


Figure 181. Control circuit in trigger mode

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### 15.4.12 Timer synchronization (TIM9/12)

The TIM timers are linked together internally for timer synchronization or chaining. Refer to *Section 14.3.15: Timer synchronization on page 398* for details.

### **15.4.13** Debug mode

When the microcontroller enters debug mode (Cortex<sup>TM</sup>-M4F core halted), the TIMx counter either continues to work normally or stops, depending on DBG\_TIMx\_STOP configuration bit in DBG module. For more details, refer to *Section 32.16.2: Debug support for timers, watchdog, bxCAN and I2C*.

## 15.5 TIM9 and TIM12 registers

Refer to Section 1.1 for a list of abbreviations used in register descriptions.

The peripheral registers have to be written by half-words (16 bits) or words (32 bits). Read accesses can be done by bytes (8 bits), half-words (16 bits) or words (32 bits).

### 15.5.1 TIM9/12 control register 1 (TIMx\_CR1)

Address offset: 0x00
Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Reserved					CKE	[1:0]	ARPE		Reserved		ОРМ	URS	UDIS	CEN
	Reserved					rw	rw	rw		neserveu		rw	rw	rw	rw

Bits 15:10 Reserved, must be kept at reset value.

### Bits 9:8 CKD: Clock division

This bit-field indicates the division ratio between the timer clock (CK\_INT) frequency and sampling clock used by the digital filters (TIx),

00:  $t_{DTS} = t_{CK\_INT}$ 

01:  $t_{DTS} = 2 \times t_{CK \ INT}$ 

10:  $t_{DTS} = 4 \times t_{CK\_INT}$ 

11: Reserved

### Bit 7 ARPE: Auto-reload preload enable

0: TIMx\_ARR register is not buffered.

1: TIMx\_ARR register is buffered.

Bits 6:4 Reserved, must be kept at reset value.

### Bit 3 OPM: One-pulse mode

0: Counter is not stopped on the update event

1: Counter stops counting on the next update event (clearing the CEN bit).

### Bit 2 URS: Update request source

This bit is set and cleared by software to select the UEV event sources.

0: Any of the following events generates an update interrupt if enabled:

- Counter overflow
- Setting the UG bit
- 1: Only counter overflow generates an update interrupt if enabled.

### Bit 1 UDIS: Update disable

This bit is set and cleared by software to enable/disable update event (UEV) generation.

- 0: UEV enabled. An UEV is generated by one of the following events:
  - Counter overflow
  - Setting the UG bit

Buffered registers are then loaded with their preload values.

1: UEV disabled. No UEV is generated, shadow registers keep their value (ARR, PSC, CCRx). The counter and the prescaler are reinitialized if the UG bit is set.

### Bit 0 CEN: Counter enable

- 0: Counter disabled
- 1: Counter enabled

CEN is cleared automatically in one-pulse mode, when an update event occurs.

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### 15.5.2 TIM9/12 control register 2 (TIMx\_CR2)

Address offset: 0x04 Reset value: 0x0000



Bits 15:7 Reserved, must be kept at reset value.

#### Bits 6:4 MMS: Master mode selection

These bits are used to select the information to be sent in Master mode to slave timers for synchronization (TRGO). The combination is as follows:

000: **Reset** - the UG bit in the TIMx\_EGR register is used as the trigger output (TRGO). If the reset is generated by the trigger input (slave mode controller configured in reset mode) then the signal on TRGO is delayed compared to the actual reset.

001: **Enable** - the Counter enable signal, CNT\_EN, is used as the trigger output (TRGO). It is useful to start several timers at the same time or to control a window in which a slave timer is enabled. The Counter Enable signal is generated by a logic OR between the CEN control bit and the trigger input when configured in Gated mode.

When the Counter Enable signal is controlled by the trigger input, there is a delay on TRGO, except if the master/slave mode is selected (see the MSM bit description in the TIMx\_SMCR register).

010: **Update** - The update event is selected as the trigger output (TRGO). For instance a master timer can be used as a prescaler for a slave timer.

011: **Compare pulse** - The trigger output sends a positive pulse when the CC1IF flag is to be set (even if it was already high), as soon as a capture or a compare match occurs. (TRGO).

100: **Compare** - OC1REF signal is used as the trigger output (TRGO).

101: Compare - OC2REF signal is used as the trigger output (TRGO).

110: Reserved

111: Reserved

Bits 3:0 Reserved, must be kept at reset value.

#### TIM9/12 slave mode control register (TIMx\_SMCR) 15.5.3

Address offset: 0x08 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			Poor	erved				MSM		TS[2:0]		Res.		SMS[2:0]	
			nese	erveu				rw	rw	rw	rw	nes.	rw	rw	rw

Bits 15:8 Reserved, must be kept at reset value.

#### Bit 7 MSM: Master/Slave mode

0: No action

1: The effect of an event on the trigger input (TRGI) is delayed to allow a perfect synchronization between the current timer and its slaves (through TRGO). It is useful in order to synchronize several timers on a single external event.

### Bits 6:4 TS: Trigger selection

This bitfield selects the trigger input to be used to synchronize the counter.

000: Internal Trigger 0 (ITR0)

001: Internal Trigger 1 (ITR1)

010: Internal Trigger 2 (ITR2)

011: Internal Trigger 3 (ITR3)

100: TI1 Edge Detector (TI1F\_ED)

101: Filtered Timer Input 1 (TI1FP1)

110: Filtered Timer Input 2 (TI2FP2)

111: Reserved.

See Table 65: TIMx internal trigger connection on page 450 for more details on the meaning of ITRx for each timer.

Note: These bits must be changed only when they are not used (e.g. when SMS='000') to avoid wrong edge detections at the transition.

Bit 3 Reserved, must be kept at reset value.

### Bits 2:0 SMS: Slave mode selection

When external signals are selected, the active edge of the trigger signal (TRGI) is linked to the polarity selected on the external input (see Input control register and Control register descriptions.

000: Slave mode disabled - if CEN = 1 then the prescaler is clocked directly by the internal clock

001: Reserved

010: Reserved

011: Reserved

100: Reset mode - Rising edge of the selected trigger input (TRGI) reinitializes the counter and generates an update of the registers

101: Gated mode - The counter clock is enabled when the trigger input (TRGI) is high. The counter stops (but is not reset) as soon as the trigger becomes low. Counter starts and stops are both controlled

110: Trigger mode - The counter starts on a rising edge of the trigger TRGI (but it is not reset). Only the start of the counter is controlled

111: External clock mode 1 - Rising edges of the selected trigger (TRGI) clock the counter

Note: The Gated mode must not be used if TI1F\_ED is selected as the trigger input (TS='100'). Indeed, TI1F\_ED outputs 1 pulse for each transition on TI1F, whereas the Gated mode checks the level of the trigger signal.

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Table 65. TIMx internal trigger connection

Slave TIM	ITR0 (TS =' 000')	ITR1 (TS = '001')	ITR2 (TS = '010')	ITR3 (TS = '011')
TIM2	TIM1	TIM8	TIM3	TIM4
TIM3	TIM1	TIM2	TIM5	TIM4
TIM4	TIM1	TIM2	TIM3	TIM8
TIM5	TIM2	TIM3	TIM4	TIM8
TIM9	TIM2	TIM3	TIM10	TIM11
TIM12	TIM4	TIM5	TIM13	TIM14

### 15.5.4 TIM9/12 Interrupt enable register (TIMx\_DIER)

Address offset: 0x0C Reset value: 0x0000



Bit 15:7 Reserved, must be kept at reset value.

Bit 6 TIE: Trigger interrupt enable

0: Trigger interrupt disabled.

1: Trigger interrupt enabled.

Bit 5:3 Reserved, must be kept at reset value.

Bit 2 CC2IE: Capture/Compare 2 interrupt enable

0: CC2 interrupt disabled.

1: CC2 interrupt enabled.

Bit 1 CC1IE: Capture/Compare 1 interrupt enable

0: CC1 interrupt disabled.1: CC1 interrupt enabled.

Bit 0 **UIE**: Update interrupt enable

0: Update interrupt disabled.

1: Update interrupt enabled.

### 15.5.5 TIM9/12 status register (TIMx\_SR)

Address offset: 0x10 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
		Reserved			CC2OF	CC1OF		erved	TIF		Reserved		CC2IF	CC1IF	UIF	
		neserveu			rc_w0	rc_w0	nest	erveu	rc_w0		neserveu		rc_w0	rc_w0	rc_w0	

Bit 15:11 Reserved, must be kept at reset value.

Bit 10 CC2OF: Capture/compare 2 overcapture flag

refer to CC1OF description

Bit 9 CC10F: Capture/Compare 1 overcapture flag

This flag is set by hardware only when the corresponding channel is configured in input capture mode. It is cleared by software by writing it to '0'.

0: No overcapture has been detected.

1: The counter value has been captured in TIMx\_CCR1 register while CC1IF flag was already set

Bits 8:7 Reserved, must be kept at reset value.

Bit 6 TIF: Trigger interrupt flag

This flag is set by hardware on trigger event (active edge detected on TRGI input when the slave mode controller is enabled in all modes but gated mode. It is set when the counter starts or stops when gated mode is selected. It is cleared by software.

0: No trigger event occurred.

1: Trigger interrupt pending.

Bit 5:3 Reserved, must be kept at reset value.

Bit 2 CC2IF: Capture/Compare 2 interrupt flag

refer to CC1IF description

Bit 1 CC1IF: Capture/compare 1 interrupt flag

### If channel CC1 is configured as output:

This flag is set by hardware when the counter matches the compare value. It is cleared by software.

0: No match.

1: The content of the counter TIMx\_CNT matches the content of the TIMx\_CCR1 register. When the contents of TIMx\_CCR1 are greater than the contents of TIMx\_ARR, the CC1IF bit goes high on the counter overflow.

### If channel CC1 is configured as input:

This bit is set by hardware on a capture. It is cleared by software or by reading the TIMx\_CCR1 register.

0: No input capture occurred.

1: The counter value has been captured in TIMx\_CCR1 register (an edge has been detected on IC1 which matches the selected polarity).

#### Bit 0 **UIF**: Update interrupt flag

This bit is set by hardware on an update event. It is cleared by software.

- 0: No update occurred.
- 1: Update interrupt pending. This bit is set by hardware when the registers are updated:
  - -At overflow and if UDIS='0' in the TIMx\_CR1 register.
  - -When CNT is reinitialized by software using the UG bit in TIMx\_EGR register, if URS='0' and UDIS='0' in the TIMx\_CR1 register.
  - -When CNT is reinitialized by a trigger event (refer to the synchro control register description), if URS='0' and UDIS='0' in the TIMx\_CR1 register.

### 15.5.6 TIM9/12 event generation register (TIMx\_EGR)

Address offset: 0x14 Reset value: 0x0000

15	14	13	12	11	10	9	8	/	ь	5	4	3	2	1	U
				Reserved					TG		Reserved		CC2G	CC1G	UG
				neserveu					W		neserveu		w	W	w

### Bits 15:7 Reserved, must be kept at reset value.

### Bit 6 TG: Trigger generation

This bit is set by software in order to generate an event, it is automatically cleared by hardware.

0: No action

1: The TIF flag is set in the TIMx\_SR register. Related interrupt can occur if enabled

#### Bits 5:3 Reserved, must be kept at reset value.

### Bit 2 CC2G: Capture/compare 2 generation

refer to CC1G description

### Bit 1 CC1G: Capture/compare 1 generation

This bit is set by software to generate an event, it is automatically cleared by hardware.

0: No action

1: A capture/compare event is generated on channel 1:

### If channel CC1 is configured as output:

the CC1IF flag is set, the corresponding interrupt is sent if enabled.

#### If channel CC1 is configured as input:

The current counter value is captured in the TIMx\_CCR1 register. The CC1IF flag is set, the corresponding interrupt is sent if enabled. The CC1OF flag is set if the CC1IF flag was already high.

### Bit 0 UG: Update generation

This bit can be set by software, it is automatically cleared by hardware.

0: No action

1: Re-initializes the counter and generates an update of the registers. The prescaler counter is also cleared and the prescaler ratio is not affected. The counter is cleared.

### 15.5.7 TIM9/12 capture/compare mode register 1 (TIMx\_CCMR1)

Address offset: 0x18 Reset value: 0x0000

The channels can be used in input (capture mode) or in output (compare mode). The direction of a channel is defined by configuring the corresponding CCxS bits. All the other bits in this register have different functions in input and output modes. For a given bit, OCxx describes its function when the channel is configured in output mode, ICxx describes its function when the channel is configured in input mode. So you must take care that the same bit can have different meanings for the input stage and the output stage.

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Res.	(	OC2M[2:0	]	OC2PE	OC2FE	CCas	S[1:0]	Res.	(	OC1M[2:0	]	OC1PE	OC1FE		S[1:0]
	IC2F[3:0]			IC2PS	C[1:0]	0028	5[1.0]		IC1F	[3:0]		IC1PS	SC[1:0]	0010	5[1.0]
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

### **Output compare mode**

Bit 15 Reserved, must be kept at reset value.

Bits 14:12 OC2M[2:0]: Output compare 2 mode

Bit 11 OC2PE: Output compare 2 preload enable

Bit 10 OC2FE: Output compare 2 fast enable

Bits 9:8 CC2S[1:0]: Capture/Compare 2 selection

This bitfield defines the direction of the channel (input/output) as well as the used input.

00: CC2 channel is configured as output

01: CC2 channel is configured as input, IC2 is mapped on TI2

10: CC2 channel is configured as input, IC2 is mapped on TI1

11: CC2 channel is configured as input, IC2 is mapped on TRC. This mode works only if an internal trigger input is selected through the TS bit (TIMx\_SMCR register

Note: The CC2S bits are writable only when the channel is OFF (CC2E = 0 in TIMx\_CCER).

Bit 7 Reserved, must be kept at reset value.

#### Bits 6:4 OC1M: Output compare 1 mode

These bits define the behavior of the output reference signal OC1REF from which OC1 and OC1N are derived. OC1REF is active high whereas the active levels of OC1 and OC1N depend on the CC1P and CC1NP bits, respectively.

000: Frozen - The comparison between the output compare register TIMx\_CCR1 and the counter TIMx\_CNT has no effect on the outputs.(this mode is used to generate a timing base).

001: Set channel 1 to active level on match. The OC1REF signal is forced high when the TIMx CNT counter matches the capture/compare register 1 (TIMx CCR1).

010: Set channel 1 to inactive level on match. The OC1REF signal is forced low when the TIMx\_CNT counter matches the capture/compare register 1 (TIMx\_CCR1).

011: Toggle - OC1REF toggles when TIMx\_CNT=TIMx\_CCR1

100: Force inactive level - OC1REF is forced low

101: Force active level - OC1REF is forced high

110: PWM mode 1 - In upcounting, channel 1 is active as long as TIMx\_CNT<TIMx\_CCR1 else it is inactive. In downcounting, channel 1 is inactive (OC1REF='0) as long as TIMx\_CNT>TIMx\_CCR1, else it is active (OC1REF='1')

111: PWM mode 2 - In upcounting, channel 1 is inactive as long as

TIMx\_CNT<TIMx\_CCR1 else it is active. In downcounting, channel 1 is active as long as TIMx\_CNT>TIMx\_CCR1 else it is inactive.

Note: In PWM mode 1 or 2, the OCREF level changes only when the result of the comparison changes or when the output compare mode switches from "frozen" mode to "PWM" mode.

#### Bit 3 OC1PE: Output compare 1 preload enable

0: Preload register on TIMx\_CCR1 disabled. TIMx\_CCR1 can be written at anytime, the new value is taken into account immediately

1: Preload register on TIMx\_CCR1 enabled. Read/Write operations access the preload register. TIMx\_CCR1 preload value is loaded into the active register at each update event

Note: The PWM mode can be used without validating the preload register only in one-pulse mode (OPM bit set in the TIMx\_CR1 register). Else the behavior is not guaranteed.

### Bit 2 OC1FE: Output compare 1 fast enable

This bit is used to accelerate the effect of an event on the trigger in input on the CC output. 0: CC1 behaves normally depending on the counter and CCR1 values even when the trigger is ON. The minimum delay to activate the CC1 output when an edge occurs on the trigger input is 5 clock cycles

1: An active edge on the trigger input acts like a compare match on the CC1 output. Then, OC is set to the compare level independently of the result of the comparison. Delay to sample the trigger input and to activate CC1 output is reduced to 3 clock cycles. OC1FE acts only if the channel is configured in PWM1 or PWM2 mode.

### Bits 1:0 CC1S: Capture/Compare 1 selection

This bitfield defines the direction of the channel (input/output) as well as the used input. 00: CC1 channel is configured as output

01: CC1 channel is configured as input, IC1 is mapped on TI1

10: CC1 channel is configured as input, IC1 is mapped on TI2

11: CC1 channel is configured as input, IC1 is mapped on TRC. This mode works only if an internal trigger input is selected through the TS bit (TIMx\_SMCR register)

Note: The CC1S bits are writable only when the channel is OFF (CC1E = 0 in TIMx\_CCER).

### Input capture mode

```
Bits 15:12 IC2F: Input capture 2 filter
```

Bits 11:10 IC2PSC[1:0]: Input capture 2 prescaler

#### Bits 9:8 CC2S: Capture/compare 2 selection

This bitfield defines the direction of the channel (input/output) as well as the used input.

00: CC2 channel is configured as output

01: CC2 channel is configured as input, IC2 is mapped on TI2

10: CC2 channel is configured as input, IC2 is mapped on TI1

11: CC2 channel is configured as input, IC2 is mapped on TRC. This mode works only if an internal trigger input is selected through the TS bit (TIMx\_SMCR register)

Note: The CC2S bits are writable only when the channel is OFF (CC2E = 0 in TIMx\_CCER).

### Bits 7:4 IC1F: Input capture 1 filter

This bitfield defines the frequency used to sample the TI1 input and the length of the digital filter applied to TI1. The digital filter is made of an event counter in which N events are needed to validate a transition on the output:

```
0000: No filter, sampling is done at f<sub>DTS</sub>
                                                                       1000: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/8, N=6
                                                                        1001: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/8, N=8
0001: f<sub>SAMPLING</sub>=f<sub>CK INT</sub>, N=2
0010: f<sub>SAMPLING</sub>=f<sub>CK INT</sub>, N=4
                                                                       1010: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/16, N=5
0011: f<sub>SAMPLING</sub>=f<sub>CK INT</sub>, N=8
                                                                       1011: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/16, N=6
0100: f_{SAMPLING} = f_{DTS}/2, N=6
                                                                       1100: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/16, N=8
0101: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/2, N=8
                                                                       1101: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/32, N=5
0110: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/4, N=6
                                                                       1110: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/32, N=6
0111: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/4, N=8
                                                                       1111: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/32, N=8
```

Note: In the current silicon revision,  $f_{DTS}$  is replaced in the formula by CK\_INT when ICxF[3:0]=1, 2 or 3.

#### Bits 3:2 IC1PSC: Input capture 1 prescaler

This bitfield defines the ratio of the prescaler acting on the CC1 input (IC1).

The prescaler is reset as soon as CC1E='0' (TIMx\_CCER register).

00: no prescaler, capture is done each time an edge is detected on the capture input

01: capture is done once every 2 events

10: capture is done once every 4 events

11: capture is done once every 8 events

### Bits 1:0 CC1S: Capture/Compare 1 selection

This bitfield defines the direction of the channel (input/output) as well as the used input.

00: CC1 channel is configured as output

01: CC1 channel is configured as input, IC1 is mapped on TI1

10: CC1 channel is configured as input, IC1 is mapped on TI2

11: CC1 channel is configured as input, IC1 is mapped on TRC. This mode is working only if an internal trigger input is selected through TS bit (TIMx\_SMCR register)

Note: The CC1S bits are writable only when the channel is OFF (CC1E = 0 in TIMx\_CCER).

### 15.5.8 TIM9/12 capture/compare enable register (TIMx CCER)

Address offset: 0x20 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			Rese	nuad				CC2NP	Res.	CC2P	CC2E	CC1NP	Res.	CC1P	CC1E
			nese	erveu				rw	nes.	rw	rw	rw	nes.	rw	rw

Bits 15:8 Reserved, must be kept at reset value.

Bit 7 CC2NP: Capture/Compare 2 output Polarity

refer to CC1NP description

Bits 6 Reserved, must be kept at reset value.

Bit 5 CC2P: Capture/Compare 2 output Polarity

refer to CC1P description

Bit 4 CC2E: Capture/Compare 2 output enable

refer to CC1E description

Bit 3 CC1NP: Capture/Compare 1 complementary output Polarity

CC1 channel configured as output: CC1NP must be kept cleared

CC1 channel configured as input: CC1NP is used in conjunction with CC1P to define TI1FP1/TI2FP1 polarity (refer to CC1P description).

Bits 2 Reserved, must be kept at reset value.

Bit 1 CC1P: Capture/Compare 1 output Polarity.

### CC1 channel configured as output:

0: OC1 active high.

1: OC1 active low.

### CC1 channel configured as input:

CC1NP/CC1P bits select TI1FP1 and TI2FP1 polarity for trigger or capture operations.

00: noninverted/rising edge

Circuit is sensitive to TlxFP1 rising edge (capture, trigger in reset, external clock or trigger mode), TlxFP1 is not inverted (trigger in gated mode, encoder mode).

01: inverted/falling edge

Circuit is sensitive to TIxFP1 falling edge (capture, trigger in reset, external clock or trigger mode), TIxFP1 is inverted (trigger in gated mode, encoder mode).

10: reserved, do not use this configuration.

Note: 11: noninverted/both edges

Circuit is sensitive to both TIxFP1 rising and falling edges (capture, trigger in reset, external clock or trigger mode), TIxFP1 is not inverted (trigger in gated mode). This configuration must not be used for encoder mode.

Bit 0 CC1E: Capture/Compare 1 output enable.

### CC1 channel configured as output:

0: Off - OC1 is not active.

1: On - OC1 signal is output on the corresponding output pin.

### CC1 channel configured as input:

This bit determines if a capture of the counter value can actually be done into the input capture/compare register 1 (TIMx\_CCR1) or not.

0: Capture disabled.

1: Capture enabled.



Table 66. Output control bit for standard OCx channels

CCxE bit	OCx output state
0	Output disabled (OCx='0', OCx_EN='0')
1	OCx=OCxREF + Polarity, OCx_EN='1'

Note:

The states of the external I/O pins connected to the standard OCx channels depend on the state of the OCx channel and on the GPIO registers.

### 15.5.9 TIM9/12 counter (TIMx\_CNT)

Address offset: 0x24

Reset value: 0x0000 0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							CNT[	[15:0]							
rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bits 15:0 CNT[15:0]: Counter value

### 15.5.10 TIM9/12 prescaler (TIMx\_PSC)

Address offset: 0x28 Reset value: 0x0000

	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Γ				•	•	•		PSC	[15:0]	•	•	•			•	
	rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bits 15:0 PSC[15:0]: Prescaler value

The counter clock frequency CK\_CNT is equal to  $f_{CK}$  PSC / (PSC[15:0] + 1).

PSC contains the value to be loaded into the active prescaler register at each update event.

### 15.5.11 TIM9/12 auto-reload register (TIMx\_ARR)

Address offset: 0x2C

Reset value: 0x0000 0000



Bits 15:0 ARR[15:0]: Auto-reload value

ARR is the value to be loaded into the actual auto-reload register.

Refer to the Section 15.4.1: Time-base unit on page 430 for more details about ARR update and behavior.

The counter is blocked while the auto-reload value is null.

### 15.5.12 TIM9/12 capture/compare register 1 (TIMx\_CCR1)

Address offset: 0x34 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							CCR1	[15:0]							
rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bits 15:0 CCR1[15:0]: Capture/Compare 1 value

#### If channel CC1 is configured as output:

CCR1 is the value to be loaded into the actual capture/compare 1 register (preload value). It is loaded permanently if the preload feature is not selected in the TIMx\_CCMR1 register (OC1PE bit). Else the preload value is copied into the active capture/compare 1 register when an update event occurs.

The active capture/compare register contains the value to be compared to the TIMx\_CNT counter and signaled on the OC1 output.

### If channel CC1is configured as input:

CCR1 is the counter value transferred by the last input capture 1 event (IC1).

## 15.5.13 TIM9/12 capture/compare register 2 (TIMx\_CCR2)

Address offset: 0x38 Reset value: 0x0000

	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Ī					_			CCR2	[15:0]							
	rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bits 15:0 CCR2[15:0]: Capture/Compare 2 value

#### If channel CC2 is configured as output:

CCR2 is the value to be loaded into the actual capture/compare 2 register (preload value). It is loaded permanently if the preload feature is not selected in the TIMx\_CCMR2 register (OC2PE bit). Else the preload value is copied into the active capture/compare 2 register when an update event occurs.

The active capture/compare register contains the value to be compared to the TIMx\_CNT counter and signalled on the OC2 output.

### If channel CC2 is configured as input:

CCR2 is the counter value transferred by the last input capture 2 event (IC2).

### 15.5.14 TIM9/12 register map

TIM9/12 registers are mapped as 16-bit addressable registers as described below:

Table 67. TIM9/12 register map and reset values

Offset	Register	10 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	6 8	7	9	4	3	1	0
0x00	TIMx_CR1	Reserved	CKD [1:0]	ARPE	Reserve		URS	SIGN	CEN
	Reset value		0 0	0			0 0	0	0
0x04	TIMx_CR2	Reserved		·	MMS[2:	0]	Res	erve	∍d
	Reset value				0 0	0			

Table 67. TIM9/12 register map and reset values (continued)

Offset	Register	33 30 30 30 30 30 30 30 30 30 30 30 30 3	112 11 11 11 11 11 11 11 11 11 11 11 11
0x08	TIMx_SMCR Reset value	Reserved	No   0   0   0   0   0   0   0   0   0
0x0C	TIMx_DIER  Reset value	Reserved	
0x10	TIMx_SR  Reset value	Reserved	C   C   C   C   C   C   C   C   C   C
0x14	TIMx_EGR Reset value	Reserved	5 Reserved 5 5 5 5 0 0 0
0x18	TIMx_CCMR1 Output Compare mode Reset value	Reserved	OC2M
0.10	TIMx_CCMR1 Input Capture mode Reset value	Reserved	IC2F[3:0]
0x1C		Reserved	
0x20	TIMx_CCER Reset value	Reserved	o CC2NP C CC2NP C CC2P C CC2P C CC2P C CC2P C CC1NP Reserved C CC1NP C CC1NP
0x24	TIMx_CNT Reset value	Reserved	CNT[15:0]
0x28	TIMx_PSC Reset value	Reserved	PSC[15:0]
0x2C	TIMx_ARR Reset value	Reserved	ARR[15:0]
0x30		Reserved	
0x34	TIMx_CCR1 Reset value	Reserved	CCR1[15:0]
0x38	TIMx_CCR2  Reset value	Reserved	CCR2[15:0]
0x3C to 0x4C	1.0001 14.140	Reserved	

Refer to *Table 2 on page 50* for the register boundary addresses.

# 15.6 TIM10/11/13/14 registers

The peripheral registers have to be written by half-words (16 bits) or words (32 bits). Read accesses can be done by bytes (8 bits), half-words (16 bits) or words (32 bits).

### 15.6.1 TIM10/11/13/14 control register 1 (TIMx\_CR1)

Address offset: 0x00 Reset value: 0x0000



Bits 15:10 Reserved, must be kept at reset value.

#### Bits 9:8 CKD: Clock division

This bit-field indicates the division ratio between the timer clock (CK\_INT) frequency and sampling clock used by the digital filters (ETR, TIx),

00:  $t_{DTS} = t_{CK\_INT}$ 01:  $t_{DTS} = 2 \times t_{CK\_INT}$ 10:  $t_{DTS} = 4 \times t_{CK\_INT}$ 11: Reserved

### Bit 7 ARPE: Auto-reload preload enable

0: TIMx\_ARR register is not buffered1: TIMx\_ARR register is buffered

Bits 6:3 Reserved, must be kept at reset value.

## Bit 2 **URS**: Update request source

This bit is set and cleared by software to select the update interrupt (UEV) sources.

0: Any of the following events generate an UEV if enabled:

- Counter overflow
- Setting the UG bit
- 1: Only counter overflow generates an UEV if enabled.

### Bit 1 UDIS: Update disable

This bit is set and cleared by software to enable/disable update interrupt (UEV) event generation.

0: UEV enabled. An UEV is generated by one of the following events:

- Counter overflow
- Setting the UG bit.

Buffered registers are then loaded with their preload values.

1: UEV disabled. No UEV is generated, shadow registers keep their value (ARR, PSC, CCRx). The counter and the prescaler are reinitialized if the UG bit is set.

Bit 0 CEN: Counter enable

0: Counter disabled

1: Counter enabled

### 10/11/13/14

### 15.6.2 TIM10/11/13/14 status register (TIMx\_SR)

Address offset: 0x10

#### Reset value: 0x0000 13 12 14 10 O 15 11 9 6 5 2 CC10F CC1IF UIF Reserved Reserved rc\_w0 rc\_w0 rc\_w0

#### Bit 15:10 Reserved, must be kept at reset value.

#### Bit 9 CC10F: Capture/Compare 1 overcapture flag

This flag is set by hardware only when the corresponding channel is configured in input capture mode. It is cleared by software by writing it to '0'.

- 0: No overcapture has been detected.
- 1: The counter value has been captured in TIMx\_CCR1 register while CC1IF flag was already set
- Bits 8:2 Reserved, must be kept at reset value.

### Bit 1 CC1IF: Capture/compare 1 interrupt flag

### If channel CC1 is configured as output:

This flag is set by hardware when the counter matches the compare value. It is cleared by software.

- 0: No match.
- 1: The content of the counter TIMx\_CNT matches the content of the TIMx\_CCR1 register. When the contents of TIMx\_CCR1 are greater than the contents of TIMx\_ARR, the CC1IF bit goes high on the counter overflow.

### If channel CC1 is configured as input:

This bit is set by hardware on a capture. It is cleared by software or by reading the  $TIMx\_CCR1$  register.

- 0: No input capture occurred.
- 1: The counter value has been captured in TIMx\_CCR1 register (an edge has been detected on IC1 which matches the selected polarity).

#### Bit 0 **UIF**: Update interrupt flag

This bit is set by hardware on an update event. It is cleared by software.

- 0: No update occurred.
- 1: Update interrupt pending. This bit is set by hardware when the registers are updated:
  - At overflow and if UDIS='0' in the TIMx\_CR1 register.
  - When CNT is reinitialized by software using the UG bit in TIMx\_EGR register, if URS='0' and UDIS='0' in the TIMx\_CR1 register.

### 15.6.3 TIM10/11/13/14 event generation register (TIMx\_EGR)

Address offset: 0x14
Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
													CC1G	UG	
Reserved													w	w	

Bits 15:2 Reserved, must be kept at reset value.

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### Bit 1 CC1G: Capture/compare 1 generation

This bit is set by software in order to generate an event, it is automatically cleared by hardware.

0: No action

1: A capture/compare event is generated on channel 1:

### If channel CC1 is configured as output:

CC1IF flag is set, Corresponding interrupt or is sent if enabled.

### If channel CC1 is configured as input:

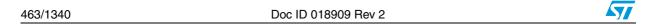
The current value of the counter is captured in TIMx\_CCR1 register. The CC1IF flag is set, the corresponding interrupt is sent if enabled. The CC1OF flag is set if the CC1IF flag was already high.

### Bit 0 UG: Update generation

This bit can be set by software, it is automatically cleared by hardware.

0: No action

1: Re-initialize the counter and generates an update of the registers. Note that the prescaler counter is cleared too (anyway the prescaler ratio is not affected). The counter is cleared.



# 15.6.4 TIM10/11/13/14 capture/compare mode register 1 (TIMx\_CCMR1)

Address offset: 0x18 Reset value: 0x0000

The channels can be used in input (capture mode) or in output (compare mode). The direction of a channel is defined by configuring the corresponding CCxS bits. All the other bits of this register have a different function in input and in output mode. For a given bit, OCxx describes its function when the channel is configured in output, ICxx describes its function when the channel is configured in input. So you must take care that the same bit can have a different meaning for the input stage and for the output stage.

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
				Reserved	ı	(	OC1M[2:0	)]	OC1PE	OC1FE	CC1S[1:0]				
	Reserved									[3:0]		IC1PS	C[1:0]	0013[1.0]	
										rw	rw	rw	rw	rw	rw

### **Output compare mode**

Bits 15:7 Reserved, must be kept at reset value.

#### Bits 6:4 **OC1M**: Output compare 1 mode

These bits define the behavior of the output reference signal OC1REF from which OC1 is derived. OC1REF is active high whereas OC1 active level depends on CC1P bit.

000: Frozen. The comparison between the output compare register TIMx\_CCR1 and the counter TIMx\_CNT has no effect on the outputs.

001: Set channel 1 to active level on match. OC1REF signal is forced high when the counter TIMx\_CNT matches the capture/compare register 1 (TIMx\_CCR1).

010: Set channel 1 to inactive level on match. OC1REF signal is forced low when the counter TIMx\_CNT matches the capture/compare register 1 (TIMx\_CCR1).

011: Toggle - OC1REF toggles when TIMx\_CNT = TIMx\_CCR1.

100: Force inactive level - OC1REF is forced low.

101: Force active level - OC1REF is forced high.

110: PWM mode 1 - Channel 1 is active as long as TIMx\_CNT < TIMx\_CCR1 else inactive.

111: PWM mode 2 - Channel 1 is inactive as long as TIMx\_CNT < TIMx\_CCR1 else active.

Note: In PWM mode 1 or 2, the OCREF level changes when the result of the comparison changes or when the output compare mode switches from frozen to PWM mode.

### Bit 3 **OC1PE**: Output compare 1 preload enable

0: Preload register on TIMx\_CCR1 disabled. TIMx\_CCR1 can be written at anytime, the new value is taken in account immediately.

1: Preload register on TIMx\_CCR1 enabled. Read/Write operations access the preload register. TIMx\_CCR1 preload value is loaded in the active register at each update event.

Note: The PWM mode can be used without validating the preload register only in one pulse mode (OPM bit set in TIMx\_CR1 register). Else the behavior is not guaranteed.

#### Bit 2 OC1FE: Output compare 1 fast enable

This bit is used to accelerate the effect of an event on the trigger in input on the CC output.

0: CC1 behaves normally depending on counter and CCR1 values even when the trigger is ON. The minimum delay to activate CC1 output when an edge occurs on the trigger input is 5 clock cycles.

1: An active edge on the trigger input acts like a compare match on CC1 output. OC is then set to the compare level independently of the result of the comparison. Delay to sample the trigger input and to activate CC1 output is reduced to 3 clock cycles. OC1FE acts only if the channel is configured in PWM1 or PWM2 mode.

### Bits 1:0 CC1S: Capture/Compare 1 selection

This bit-field defines the direction of the channel (input/output) as well as the used input.

00: CC1 channel is configured as output.

01: CC1 channel is configured as input, IC1 is mapped on TI1.

10:

11:

Note: CC1S bits are writable only when the channel is OFF (CC1E = 0 in TIMx\_CCER).

### Input capture mode

Bits 15:8 Reserved, must be kept at reset value.

#### Bits 7:4 IC1F: Input capture 1 filter

This bit-field defines the frequency used to sample TI1 input and the length of the digital filter applied to TI1. The digital filter is made of an event counter in which N events are needed to validate a transition on the output:

```
0000: No filter, sampling is done at f<sub>DTS</sub>1000: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/8, N=6
```

0001: f<sub>SAMPLING</sub>=f<sub>CK INT</sub>, N=21001: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/8, N=8

0010: f<sub>SAMPLING</sub>=f<sub>CK INT</sub>, N=41010: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/16, N=5

0011: f<sub>SAMPLING</sub>=f<sub>CK\_INT</sub>, N=81011: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/16, N=6

0100: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/2, N=61100: f<sub>SAMPLING</sub>=f<sub>DTS</sub>/16, N=8

0101:  $f_{SAMPLING} = f_{DTS}/2$ , N=81101:  $f_{SAMPLING} = f_{DTS}/32$ , N=5

0110:  $f_{SAMPLING} = f_{DTS}/4$ , N=61110:  $f_{SAMPLING} = f_{DTS}/32$ , N=6 0111:  $f_{SAMPLING} = f_{DTS}/4$ , N=81111:  $f_{SAMPLING} = f_{DTS}/32$ , N=8

Note: In current silicon revision,  $f_{DTS}$  is replaced in the formula by CK\_INT when ICxF[3:0]= 1, 2 or 3.

### Bits 3:2 IC1PSC: Input capture 1 prescaler

This bit-field defines the ratio of the prescaler acting on CC1 input (IC1).

The prescaler is reset as soon as CC1E='0' (TIMx\_CCER register).

00: no prescaler, capture is done each time an edge is detected on the capture input

01: capture is done once every 2 events

10: capture is done once every 4 events

11: capture is done once every 8 events

#### Bits 1:0 CC1S: Capture/Compare 1 selection

This bit-field defines the direction of the channel (input/output) as well as the used input.

00: CC1 channel is configured as output

01: CC1 channel is configured as input, IC1 is mapped on TI1

10: Reserved

11: Reserved

Note: CC1S bits are writable only when the channel is OFF (CC1E = 0 in  $TIMx\_CCER$ ).

# 15.6.5 TIM10/11/13/14 capture/compare enable register (TIMx\_CCER)

Address offset: 0x20 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
					Rese	ruad						CC1NP	Res.	CC1P	CC1E
					nese	erveu						rw	nes.	rw	rw

Bits 15:4 Reserved, must be kept at reset value.

Bit 3 **CC1NP**: Capture/Compare 1 complementary output Polarity.

CC1 channel configured as output: CC1NP must be kept cleared.

CC1 channel configured as input: CC1NP bit is used in conjunction with CC1P to define TI1FP1 polarity (refer to CC1P description).

Bit 2 Reserved, must be kept at reset value.

Bit 1 CC1P: Capture/Compare 1 output Polarity.

### CC1 channel configured as output:

0: OC1 active high

1: OC1 active low

#### CC1 channel configured as input:

The CC1P bit selects TI1FP1 and TI2FP1 polarity for trigger or capture operations.

00: noninverted/rising edge

Circuit is sensitive to TI1FP1 rising edge (capture mode), TI1FP1 is not inverted.

01: inverted/falling edge

Circuit is sensitive to TI1FP1 falling edge (capture mode), TI1FP1 is inverted.

10: reserved, do not use this configuration.

11: noninverted/both edges

Circuit is sensitive to both TI1FP1 rising and falling edges (capture mode), TI1FP1 is not inverted.

Bit 0 CC1E: Capture/Compare 1 output enable.

### CC1 channel configured as output:

0: Off - OC1 is not active

1: On - OC1 signal is output on the corresponding output pin

### CC1 channel configured as input:

This bit determines if a capture of the counter value can actually be done into the input capture/compare register 1 (TIMx\_CCR1) or not.

0: Capture disabled

1: Capture enabled

Table 68. Output control bit for standard OCx channels

CCxE bit	OCx output state
0	Output Disabled (OCx='0', OCx_EN='0')
1	OCx=OCxREF + Polarity, OCx_EN='1'

Note: The state of the external I/O pins connected to the standard OCx channels depends on the OCx channel state and the GPIO registers.

### 15.6.6 TIM10/11/13/14 counter (TIMx\_CNT)

Address offset: 0x24 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	CNT[15:0]														
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 15:0 CNT[15:0]: Counter value

### 15.6.7 TIM10/11/13/14 prescaler (TIMx\_PSC)

Address offset: 0x28 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							PSC[	[15:0]							
rw	rw	rw	rw	rw	rw	rw	rw	rw							

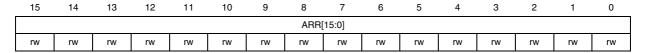
Bits 15:0 PSC[15:0]: Prescaler value

The counter clock frequency CK\_CNT is equal to  $f_{\text{CK\_PSC}}$  / (PSC[15:0] + 1).

PSC contains the value to be loaded in the active prescaler register at each update event.

### 15.6.8 TIM10/11/13/14 auto-reload register (TIMx\_ARR)

Address offset: 0x2C Reset value: 0x0000



Bits 15:0 ARR[15:0]: Auto-reload value

ARR is the value to be loaded in the actual auto-reload register.

Refer to Section 15.4.1: Time-base unit on page 430 for more details about ARR update and behavior.

The counter is blocked while the auto-reload value is null.

### 15.6.9 TIM10/11/13/14 capture/compare register 1 (TIMx\_CCR1)

Address offset: 0x34 Reset value: 0x0000

	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		CCR1[15:0]														
Ī	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

### Bits 15:0 CCR1[15:0]: Capture/Compare 1 value

#### If channel CC1 is configured as output:

CCR1 is the value to be loaded in the actual capture/compare 1 register (preload value). It is loaded permanently if the preload feature is not selected in the TIMx\_CCMR1 register (bit OC1PE). Else the preload value is copied in the active capture/compare 1 register when an update event occurs.

The active capture/compare register contains the value to be compared to the counter TIMx\_CNT and signaled on OC1 output.

#### If channel CC1is configured as input:

CCR1 is the counter value transferred by the last input capture 1 event (IC1).

### 15.6.10 TIM11 option register 1 (TIM11\_OR)

Address offset: 0x50 Reset value: 0x0000



Bits 15:2 Reserved, must be kept at reset value.

Bits 1:0 TI1\_RMP[1:0]: TIM11 Input 1 remapping capability

Set and cleared by software.

00,01,11: TIM11 Channel1 is connected to the GPIO (refer to the Alternate function mapping table in the datasheets).

10:  $HSE\_RTC$  clock (HSE divided by programmable prescaler) is connected to the  $TIM11\_CH1$  input for measurement purposes

# 15.6.11 TIM10/11/13/14 register map

TIMx registers are mapped as 16-bit addressable registers as described in the tables below:

Table 69. TIM10/11/13/14 register map and reset values

Offset	Register	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	1 2 3 4 5 6 7 8 9 9 10 11 12 13 14 15	0
0x00	TIMx_CR1 Reset value	Reserved	CKD   H   Reserved   S   S   O   O   O	
0x08	TIMx_SMCR Reset value	Rese	erved	J
0x0C	TIMx_DIER Reset value	Reserve	o O O O O O O	o UIE
0x10	TIMx_SR Reset value	Reserved	HOTO Reserved OO 0	
0x14	TIMx_EGR Reset value	Reserve	ā	ng
0x18	TIMx_CCMR1 Output compare mode Reset value TIMx_CCMR1 Input capture mode	Reserved Reserved	OC1M [2:0]	C1S 1:0] 0 C1S 1:0]
0x1C	Reset value	Reserved	0 0 0 0 0 0 0	0
0x20	TIMx_CCER Reset value	Reserved	o CC1NP Reserved	
0x24	TIMx_CNT Reset value	Reserved	CNT[15:0]	0
0x28	TIMx_PSC  Reset value	Reserved	PSC[15:0]	0
0x2C	TIMx_ARR Reset value	Reserved	ARR[15:0]	0
0x30		Reserved		
0x34	TIMx_CCR1 Reset value	Reserved	CCR1[15:0]	0
0x38 to 0x4C		Reserved		,
0x50	TIMx_OR Reset value	Reserve	ed C	O TI1_RMP

Refer to *Table 2 on page 50* for the register boundary addresses.

# 16 Basic timers (TIM6&TIM7)

This section applies to the whole STM32F4xx family, unless otherwise specified.

## 16.1 TIM6&TIM7 introduction

The basic timers TIM6 and TIM7 consist of a 16-bit auto-reload counter driven by a programmable prescaler.

They may be used as generic timers for time-base generation but they are also specifically used to drive the digital-to-analog converter (DAC). In fact, the timers are internally connected to the DAC and are able to drive it through their trigger outputs.

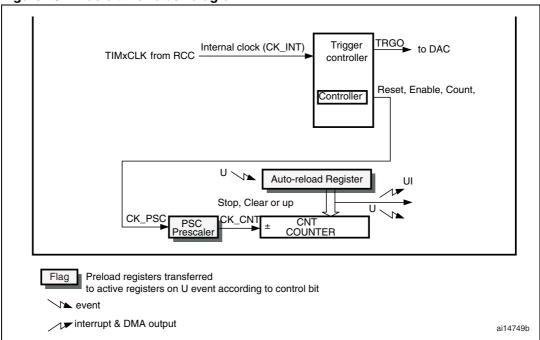
The timers are completely independent, and do not share any resources.

# 16.2 TIM6&TIM7 main features

Basic timer (TIM6&TIM7) features include:

- 16-bit auto-reload upcounter
- 16-bit programmable prescaler used to divide (also "on the fly") the counter clock frequency by any factor between 1 and 65536
- Synchronization circuit to trigger the DAC
- Interrupt/DMA generation on the update event: counter overflow

Figure 182. Basic timer block diagram



# 16.3 TIM6&TIM7 functional description

### 16.3.1 Time-base unit

The main block of the programmable timer is a 16-bit upcounter with its related auto-reload register. The counter clock can be divided by a prescaler.

The counter, the auto-reload register and the prescaler register can be written or read by software. This is true even when the counter is running.

The time-base unit includes:

- Counter Register (TIMx\_CNT)
- Prescaler Register (TIMx\_PSC)
- Auto-Reload Register (TIMx\_ARR)

The auto-reload register is preloaded. The preload register is accessed each time an attempt is made to write or read the auto-reload register. The contents of the preload register are transferred into the shadow register permanently or at each update event UEV, depending on the auto-reload preload enable bit (ARPE) in the TIMx\_CR1 register. The update event is sent when the counter reaches the overflow value and if the UDIS bit equals 0 in the TIMx\_CR1 register. It can also be generated by software. The generation of the update event is described in detail for each configuration.

The counter is clocked by the prescaler output CK\_CNT, which is enabled only when the counter enable bit (CEN) in the TIMx\_CR1 register is set.

Note that the actual counter enable signal CNT\_EN is set 1 clock cycle after CEN.

## **Prescaler description**

The prescaler can divide the counter clock frequency by any factor between 1 and 65536. It is based on a 16-bit counter controlled through a 16-bit register (in the TIMx\_PSC register). It can be changed on the fly as the TIMx\_PSC control register is buffered. The new prescaler ratio is taken into account at the next update event.

*Figure 183* and *Figure 184* give some examples of the counter behavior when the prescaler ratio is changed on the fly.

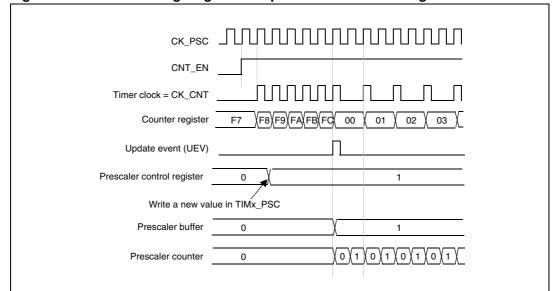
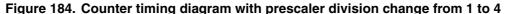
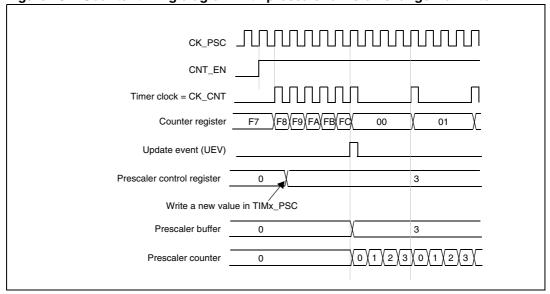


Figure 183. Counter timing diagram with prescaler division change from 1 to 2





# 16.3.2 Counting mode

The counter counts from 0 to the auto-reload value (contents of the TIMx\_ARR register), then restarts from 0 and generates a counter overflow event.

An update event can be generate at each counter overflow or by setting the UG bit in the TIMx\_EGR register (by software or by using the slave mode controller).

The UEV event can be disabled by software by setting the UDIS bit in the TIMx\_CR1 register. This avoids updating the shadow registers while writing new values into the preload registers. In this way, no update event occurs until the UDIS bit has been written to 0, however, the counter and the prescaler counter both restart from 0 (but the prescale rate does not change). In addition, if the URS (update request selection) bit in the TIMx\_CR1

register is set, setting the UG bit generates an update event UEV, but the UIF flag is not set (so no interrupt or DMA request is sent).

When an update event occurs, all the registers are updated and the update flag (UIF bit in the TIMx\_SR register) is set (depending on the URS bit):

- The buffer of the prescaler is reloaded with the preload value (contents of the TIMx\_PSC register)
- The auto-reload shadow register is updated with the preload value (TIMx\_ARR)

The following figures show some examples of the counter behavior for different clock frequencies when  $TIMx\_ARR = 0x36$ .

Figure 185. Counter timing diagram, internal clock divided by 1

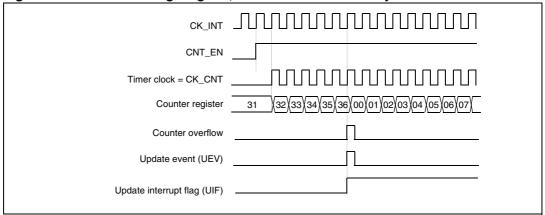
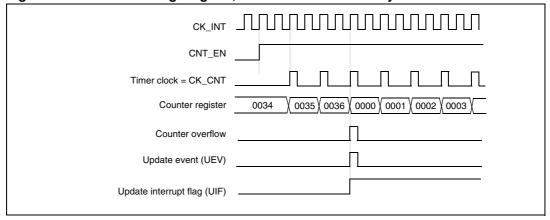


Figure 186. Counter timing diagram, internal clock divided by 2



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Figure 187. Counter timing diagram, internal clock divided by 4

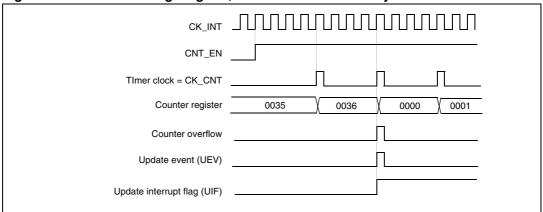


Figure 188. Counter timing diagram, internal clock divided by N

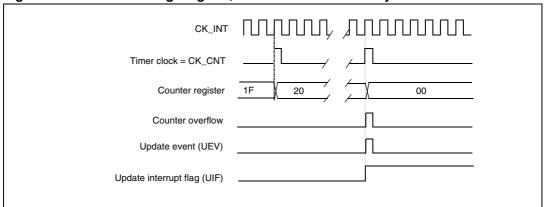
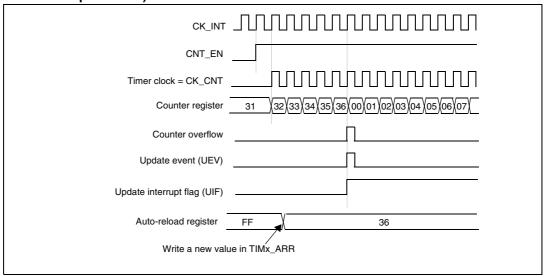


Figure 189. Counter timing diagram, update event when ARPE = 0 (TIMx\_ARR not preloaded)



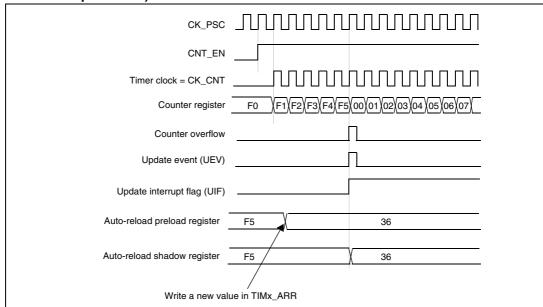


Figure 190. Counter timing diagram, update event when ARPE=1 (TIMx\_ARR preloaded)

### 16.3.3 Clock source

The counter clock is provided by the Internal clock (CK\_INT) source.

The CEN (in the TIMx\_CR1 register) and UG bits (in the TIMx\_EGR register) are actual control bits and can be changed only by software (except for UG that remains cleared automatically). As soon as the CEN bit is written to 1, the prescaler is clocked by the internal clock CK\_INT.

*Figure 191* shows the behavior of the control circuit and the upcounter in normal mode, without prescaler.

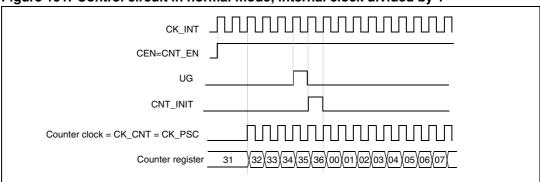


Figure 191. Control circuit in normal mode, internal clock divided by 1

# 16.3.4 Debug mode

When the microcontroller enters the debug mode (Cortex<sup>TM</sup>-M4F core - halted), the TIMx counter either continues to work normally or stops, depending on the DBG\_TIMx\_STOP configuration bit in the DBG module. For more details, refer to *Section 32.16.2: Debug support for timers, watchdog, bxCAN and I2C*.

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# 16.4 TIM6&TIM7 registers

Refer to Section 1.1 on page 46 for a list of abbreviations used in register descriptions.

The peripheral registers have to be written by half-words (16 bits) or words (32 bits). Read accesses can be done by bytes (8 bits), half-words (16 bits) or words (32 bits).

# 16.4.1 TIM6&TIM7 control register 1 (TIMx\_CR1)

Address offset: 0x00 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
			Rese	erved				ARPE		Reserved		ОРМ	URS	UDIS	CEN	
			11000	J. 704				rw		110001700		rw	rw	rw	rw	

Bits 15:8 Reserved, must be kept at reset value.

Bit 7 ARPE: Auto-reload preload enable

0: TIMx\_ARR register is not buffered.

1: TIMx\_ARR register is buffered.

Bits 6:4 Reserved, must be kept at reset value.

Bit 3 **OPM**: One-pulse mode

0: Counter is not stopped at update event

1: Counter stops counting at the next update event (clearing the CEN bit).

Bit 2 URS: Update request source

This bit is set and cleared by software to select the UEV event sources.

0: Any of the following events generates an update interrupt or DMA request if enabled. These events can be:

- Counter overflow/underflow
- Setting the UG bit
- Update generation through the slave mode controller
- 1: Only counter overflow/underflow generates an update interrupt or DMA request if enabled.

### Bit 1 UDIS: Update disable

This bit is set and cleared by software to enable/disable UEV event generation.

0: UEV enabled. The Update (UEV) event is generated by one of the following events:

- Counter overflow/underflow
- Setting the UG bit
- Update generation through the slave mode controller

Buffered registers are then loaded with their preload values.

1: UEV disabled. The Update event is not generated, shadow registers keep their value (ARR, PSC). However the counter and the prescaler are reinitialized if the UG bit is set or if a hardware reset is received from the slave mode controller.

Bit 0 CEN: Counter enable

0: Counter disabled

1: Counter enabled

Note: Gated mode can work only if the CEN bit has been previously set by software. However trigger mode can set the CEN bit automatically by hardware.

CEN is cleared automatically in one-pulse mode, when an update event occurs.

# 16.4.2 TIM6&TIM7 control register 2 (TIMx\_CR2)

Address offset: 0x04 Reset value: 0x0000



Bits 15:7 Reserved, must be kept at reset value.

#### Bits 6:4 MMS: Master mode selection

These bits are used to select the information to be sent in master mode to slave timers for synchronization (TRGO). The combination is as follows:

000: **Reset** - the UG bit from the TIMx\_EGR register is used as a trigger output (TRGO). If reset is generated by the trigger input (slave mode controller configured in reset mode) then the signal on TRGO is delayed compared to the actual reset.

001: **Enable** - the Counter enable signal, CNT\_EN, is used as a trigger output (TRGO). It is useful to start several timers at the same time or to control a window in which a slave timer is enabled. The Counter Enable signal is generated by a logic OR between CEN control bit and the trigger input when configured in gated mode.

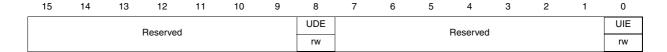
When the Counter Enable signal is controlled by the trigger input, there is a delay on TRGO, except if the master/slave mode is selected (see the MSM bit description in the TIMx\_SMCR register).

010: **Update** - The update event is selected as a trigger output (TRGO). For instance a master timer can then be used as a prescaler for a slave timer.

Bits 3:0 Reserved, must be kept at reset value.

## 16.4.3 TIM6&TIM7 DMA/Interrupt enable register (TIMx DIER)

Address offset: 0x0C Reset value: 0x0000



Bit 15:9 Reserved, must be kept at reset value.

Bit 8 **UDE**: Update DMA request enable

0: Update DMA request disabled.1: Update DMA request enabled.

Bit 7:1 Reserved, must be kept at reset value.

Bit 0 **UIE**: Update interrupt enable

0: Update interrupt disabled.

1: Update interrupt enabled.

# 16.4.4 TIM6&TIM7 status register (TIMx\_SR)

Address offset: 0x10 Reset value: 0x0000



Bits 15:1 Reserved, must be kept at reset value.

Bit 0 UIF: Update interrupt flag

This bit is set by hardware on an update event. It is cleared by software.

- 0: No update occurred.
- 1: Update interrupt pending. This bit is set by hardware when the registers are updated:
  - -At overflow or underflow and if UDIS = 0 in the TIMx CR1 register.
  - -When CNT is reinitialized by software using the UG bit in the TIMx\_EGR register, if URS = 0 and UDIS = 0 in the TIMx\_CR1 register.

# 16.4.5 TIM6&TIM7 event generation register (TIMx\_EGR)

Address offset: 0x14 Reset value: 0x0000



Bits 15:1 Reserved, must be kept at reset value.

Bit 0 UG: Update generation

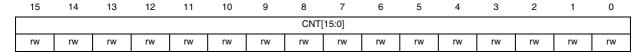
This bit can be set by software, it is automatically cleared by hardware.

0: No action.

1: Re-initializes the timer counter and generates an update of the registers. Note that the prescaler counter is cleared too (but the prescaler ratio is not affected).

# 16.4.6 TIM6&TIM7 counter (TIMx\_CNT)

Address offset: 0x24 Reset value: 0x0000



Bits 15:0 CNT[15:0]: Counter value

# 16.4.7 TIM6&TIM7 prescaler (TIMx\_PSC)

Address offset: 0x28

Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							PSC	[15:0]							
rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bits 15:0 PSC[15:0]: Prescaler value

The counter clock frequency CK\_CNT is equal to  $f_{CK_PSC}$  / (PSC[15:0] + 1).

PSC contains the value to be loaded into the active prescaler register at each update event.

# 16.4.8 TIM6&TIM7 auto-reload register (TIMx\_ARR)

Address offset: 0x2C

Reset value: 0x0000

15	14	13	12		10	9	O	,	U	5	4	3	2		U
							ARR	[15:0]							
rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bits 15:0 ARR[15:0]: Auto-reload value

ARR is the value to be loaded into the actual auto-reload register.

Refer to Section 16.3.1: Time-base unit on page 471 for more details about ARR update and behavior.

The counter is blocked while the auto-reload value is null.

# 16.4.9 TIM6&TIM7 register map

TIMx registers are mapped as 16-bit addressable registers as described in the table below:

Table 70. TIM6&TIM7 register map and reset values

Offset	Register	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	9 0 10		9 2 4	8 2	- 0							
0x00	TIMx_CR1	Reserved		ARPE	Reserved		UDIS							
	Reset value			0	Re	0 0	0 0							
0x04	TIMx_CR2	Reserved			MMS[2:0]	Reserved								
	Reset value				0 0 0	Be								
0x08		Reserved												
0x0C	TIMx_DIER	Reserved		UDE	Reserved		UIE							
	Reset value			0	Res		0							
0x10	TIMx_SR	Reserv	ved				当							
	Reset value						0							
0x14	TIMx_EGR	Reserv	ved				ng							
	Reset value						0							
0x18		Reserved												
0x1C		Reserved												
0x20		Reserved												
0x24	TIMx_CNT	Reserved	C	NT[15	:0]									
	Reset value		0 0 0 0 0 0 0	0 0	0 0 0	0 0	0   0							
0x28	TIMx_PSC	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0												
	Reset value		0   0   0   0   0   0   0	0 0	0 0 0	0 0	0 0							
0x2C	TIMx_ARR	Reserved	Д	RR[15	:0]									
	Reset value		0 0 0 0 0 0 0 0 0	0 0	0 0 0	0 0	0 0							

Refer to *Table 2 on page 50* for the register boundary addresses.

# 17 Independent watchdog (IWDG)

This section applies to the whole STM32F4xx family, unless otherwise specified.

## 17.1 IWDG introduction

The devices have two embedded watchdog peripherals which offer a combination of high safety level, timing accuracy and flexibility of use. Both watchdog peripherals (Independent and Window) serve to detect and resolve malfunctions due to software failure, and to trigger system reset or an interrupt (window watchdog only) when the counter reaches a given timeout value.

The independent watchdog (IWDG) is clocked by its own dedicated low-speed clock (LSI) and thus stays active even if the main clock fails. The window watchdog (WWDG) clock is prescaled from the APB1 clock and has a configurable time-window that can be programmed to detect abnormally late or early application behavior.

The IWDG is best suited to applications which require the watchdog to run as a totally independent process outside the main application, but have lower timing accuracy constraints. The WWDG is best suited to applications which require the watchdog to react within an accurate timing window. For further information on the window watchdog, refer to Section 18 on page 486.

## 17.2 IWDG main features

- Free-running downcounter
- clocked from an independent RC oscillator (can operate in Standby and Stop modes)
- Reset (if watchdog activated) when the downcounter value of 0x000 is reached

# 17.3 IWDG functional description

Figure 192 shows the functional blocks of the independent watchdog module.

When the independent watchdog is started by writing the value 0xCCCC in the Key register (IWDG\_KR), the counter starts counting down from the reset value of 0xFFF. When it reaches the end of count value (0x000) a reset signal is generated (IWDG reset).

Whenever the key value 0xAAAA is written in the IWDG\_KR register, the IWDG\_RLR value is reloaded in the counter and the watchdog reset is prevented.

### 17.3.1 Hardware watchdog

If the "Hardware watchdog" feature is enabled through the device option bits, the watchdog is automatically enabled at power-on, and will generate a reset unless the Key register is written by the software before the counter reaches end of count.

### 17.3.2 Register access protection

Write access to the IWDG\_PR and IWDG\_RLR registers is protected. To modify them, you must first write the code 0x5555 in the IWDG\_KR register. A write access to this register

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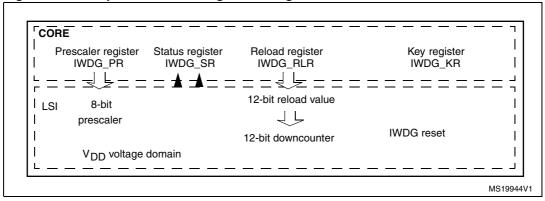
with a different value will break the sequence and register access will be protected again. This implies that it is the case of the reload operation (writing 0xAAAA).

A status register is available to indicate that an update of the prescaler or the down-counter reload value is on going.

## 17.3.3 Debug mode

When the microcontroller enters debug mode (Cortex<sup>™</sup>-M4F core halted), the IWDG counter either continues to work normally or stops, depending on DBG\_IWDG\_STOP configuration bit in DBG module. For more details, refer to *Section 32.16.2: Debug support for timers, watchdog, bxCAN and I2C*.

Figure 192. Independent watchdog block diagram



Note:

The watchdog function is implemented in the  $V_{DD}$  voltage domain that is still functional in Stop and Standby modes.

Table 71. Min/max IWDG timeout period at 32 kHz (LSI) (1)

Prescaler divider	PR[2:0] bits	Min timeout (ms) RL[11:0]= 0x000	Max timeout (ms) RL[11:0]= 0xFFF
/4	0	0.125	512
/8	1	0.25	1024
/16	2	0.5	2048
/32	3	1	4096
/64	4	2	8192
/128	5	4	16384
/256	6	8	32768

These timings are given for a 32 kHz clock but the microcontroller's internal RC frequency can vary from 30 to 60 kHz. Moreover, given an exact RC oscillator frequency, the exact timings still depend on the phasing of the APB interface clock versus the LSI clock so that there is always a full RC period of uncertainty.

# 17.4 IWDG registers

Refer to Section 1.1 on page 46 for a list of abbreviations used in register descriptions.

The peripheral registers can be accessed by half-words (16 bits) or words (32 bits).

# 17.4.1 Key register (IWDG\_KR)

Address offset: 0x00

Reset value: 0x0000 0000 (reset by Standby mode)

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Paganyad								KEY	[15:0	)]						
Reserved	w	w	w	w	w	w	w	w	w	w	w	w	w	w	w	w

Bits 31:16 Reserved, must be kept at reset value.

Bits 15:0 KEY[15:0]: Key value (write only, read 0000h)

These bits must be written by software at regular intervals with the key value AAAAh, otherwise the watchdog generates a reset when the counter reaches 0.

Writing the key value 5555h to enable access to the IWDG\_PR and IWDG\_RLR registers (see *Section 17.3.2*)

Writing the key value CCCCh starts the watchdog (except if the hardware watchdog option is selected)

# 17.4.2 Prescaler register (IWDG\_PR)

Address offset: 0x04

Reset value: 0x0000 0000

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

| PR[2:0] | rw | rw | rw |

Bits 31:3 Reserved, must be kept at reset value.

### Bits 2:0 PR[2:0]: Prescaler divider

These bits are write access protected see Section 17.3.2. They are written by software to select the prescaler divider feeding the counter clock. PVU bit of IWDG\_SR must be reset in order to be able to change the prescaler divider.

000: divider /4 001: divider /8 010: divider /16 011: divider /32 100: divider /64 101: divider /128 110: divider /256 111: divider /256

Note: Reading this register returns the prescaler value from the VDD voltage domain. This value may not be up to date/valid if a write operation to this register is ongoing. For this reason the value read from this register is valid only when the PVU bit in the IWDG\_SR register is reset.

# 17.4.3 Reload register (IWDG RLR)

Address offset: 0x08

Reset value: 0x0000 0FFF (reset by Standby mode)

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved

Bits 31:12 Reserved, must be kept at reset value.

#### Bits11:0 RL[11:0]: Watchdog counter reload value

These bits are write access protected see *Section 17.3.2*. They are written by software to define the value to be loaded in the watchdog counter each time the value AAAAh is written in the IWDG\_KR register. The watchdog counter counts down from this value. The timeout period is a function of this value and the clock prescaler. Refer to *Table 71*.

rw rw rw rw rw rw

The RVU bit in the IWDG\_SR register must be reset in order to be able to change the reload value.

Note: Reading this register returns the reload value from the VDD voltage domain. This value may not be up to date/valid if a write operation to this register is ongoing on this register. For this reason the value read from this register is valid only when the RVU bit in the IWDG\_SR register is reset.

# 17.4.4 Status register (IWDG\_SR)

Address offset: 0x0C

Reset value: 0x0000 0000 (not reset by Standby mode)

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	RVU	PVU
neserveu	r	r

Bits 31:2 Reserved, must be kept at reset value.

#### Bit 1 RVU: Watchdog counter reload value update

This bit is set by hardware to indicate that an update of the reload value is ongoing. It is reset by hardware when the reload value update operation is completed in the  $V_{DD}$  voltage domain (takes up to 5 RC 40 kHz cycles).

Reload value can be updated only when RVU bit is reset.

### Bit 0 PVU: Watchdog prescaler value update

This bit is set by hardware to indicate that an update of the prescaler value is ongoing. It is reset by hardware when the prescaler update operation is completed in the  $V_{DD}$  voltage domain (takes up to 5 RC 40 kHz cycles).

Prescaler value can be updated only when PVU bit is reset.

Note:

If several reload values or prescaler values are used by application, it is mandatory to wait until RVU bit is reset before changing the reload value and to wait until PVU bit is reset before changing the prescaler value. However, after updating the prescaler and/or the reload value it is not necessary to wait until RVU or PVU is reset before continuing code execution (even in case of low-power mode entry, the write operation is taken into account and will complete)

## 17.4.5 IWDG register map

The following table gives the IWDG register map and reset values.

Table 72. IWDG register map and reset values

Offset	Register	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	6	8	7	9	2	4	ဗ	7	-	0
0x00	IWDG_KR							В	000	rvec	1													K	EY[	15:0	0]				-	_	
0,000	Reset value							.,	.030	i vec	4							0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x04	IWDG_PR														Res	e e ru	hay							•							PR	[2:0	)]
0,04	Reset value														110	JC1 V	cu														0	0	0
0x08	IWDG_RLR									R	esei	rvor	1													ı	RL[1	11:0	]		•	•	_
0,00	Reset value									• • •	0301	1 400	•									1	1	1	1	1	1	1	1	1	1	1	1
0x0C	IWDG_SR															lese	rvo	Ч													01/1	3,00	PVU
0,00	Reset value															1636	ii VE	u															О

Refer to Table 2 on page 50 for the register boundary addresses.

# 18 Window watchdog (WWDG)

This section applies to the whole STM32F4xx family, unless otherwise specified.

## 18.1 WWDG introduction

The window watchdog is used to detect the occurrence of a software fault, usually generated by external interference or by unforeseen logical conditions, which causes the application program to abandon its normal sequence. The watchdog circuit generates an MCU reset on expiry of a programmed time period, unless the program refreshes the contents of the downcounter before the T6 bit becomes cleared. An MCU reset is also generated if the 7-bit downcounter value (in the control register) is refreshed before the downcounter has reached the window register value. This implies that the counter must be refreshed in a limited window.

## 18.2 WWDG main features

- Programmable free-running downcounter
- Conditional reset
  - Reset (if watchdog activated) when the downcounter value becomes less than 0x40
  - Reset (if watchdog activated) if the downcounter is reloaded outside the window (see Figure 194)
- Early wakeup interrupt (EWI): triggered (if enabled and the watchdog activated) when the downcounter is equal to 0x40.

# 18.3 WWDG functional description

If the watchdog is activated (the WDGA bit is set in the WWDG\_CR register) and when the 7-bit downcounter (T[6:0] bits) rolls over from 0x40 to 0x3F (T6 becomes cleared), it initiates a reset. If the software reloads the counter while the counter is greater than the value stored in the window register, then a reset is generated.

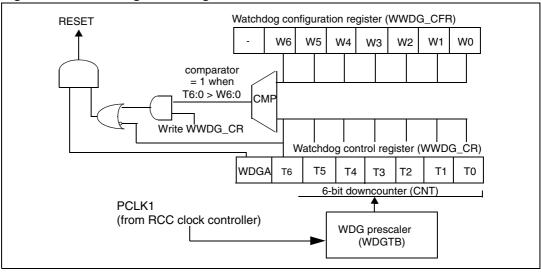


Figure 193. Watchdog block diagram

The application program must write in the WWDG\_CR register at regular intervals during normal operation to prevent an MCU reset. This operation must occur only when the counter value is lower than the window register value. The value to be stored in the WWDG\_CR register must be between 0xFF and 0xC0:

### **Enabling the watchdog**

The watchdog is always disabled after a reset. It is enabled by setting the WDGA bit in the WWDG\_CR register, then it cannot be disabled again except by a reset.

### Controlling the downcounter

This downcounter is free-running: It counts down even if the watchdog is disabled. When the watchdog is enabled, the T6 bit must be set to prevent generating an immediate reset.

The T[5:0] bits contain the number of increments which represents the time delay before the watchdog produces a reset. The timing varies between a minimum and a maximum value due to the unknown status of the prescaler when writing to the WWDG\_CR register (see *Figure 194*). The Configuration register (WWDG\_CFR) contains the high limit of the window: To prevent a reset, the downcounter must be reloaded when its value is lower than the window register value and greater than 0x3F. *Figure 194* describes the window watchdog process.

Note:

The T6 bit can be used to generate a software reset (the WDGA bit is set and the T6 bit is cleared).

### Advanced watchdog interrupt feature

The Early Wakeup Interrupt (EWI) can be used if specific safety operations or data logging must be performed before the actual reset is generated. The EWI interrupt is enabled by setting the EWI bit in the WWDG\_CFR register. When the downcounter reaches the value 0x40, an EWI interrupt is generated and the corresponding interrupt service routine (ISR) can be used to trigger specific actions (such as communications or data logging), before resetting the device.

In some applications, the EWI interrupt can be used to manage a software system check and/or system recovery/graceful degradation, without generating a WWDG reset. In this

case, the corresponding interrupt service routine (ISR) should reload the WWDG counter to avoid the WWDG reset, then trigger the required actions.

The EWI interrupt is cleared by writing '0' to the EWIF bit in the WWDG\_SR register.

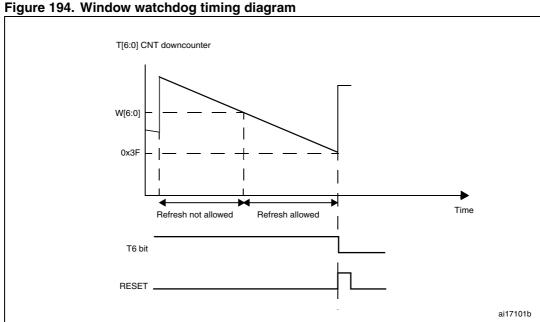
Note:

When the EWI interrupt cannot be served, e.g. due to a system lock in a higher priority task, the WWDG reset will eventually be generated.

#### 18.4 How to program the watchdog timeout

You can use the formula in Figure 194 to calculate the WWDG timeout.

Warning: When writing to the WWDG\_CR register, always write 1 in the T6 bit to avoid generating an immediate reset.



The formula to calculate the timeout value is given by:

$$t_{WWDG} = t_{PCLK1} \times 4096 \times 2^{WDGTB} \times (t[5:0] + 1)$$
 (ms)

where:

twwDG: WWDG timeout

t<sub>PCI K1</sub>: APB1 clock period measured in ms

Refer to the table below for the minimum and maximum values of the T<sub>WWDG</sub>.

Table 73. Timeout values at 30 MHz (f<sub>PCLK1</sub>)

Prescaler	WDGTB	Min timeout (μs) T[5:0] = 0x00	Max timeout (ms) T[5:0] = 0x3F
1	0	136.53	8.74
2	1	273.07	17.48
4	2	546.13	34.95
8	3	1092.27	69.91

# 18.5 Debug mode

When the microcontroller enters debug mode (Cortex<sup>™</sup>-M4F core halted), the WWDG counter either continues to work normally or stops, depending on DBG\_WWDG\_STOP configuration bit in DBG module. For more details, refer to *Section 32.16.2: Debug support for timers, watchdog, bxCAN and I2C*.

# 18.6 WWDG registers

Refer to Section 1.1 on page 46 for a list of abbreviations used in register descriptions.

The peripheral registers can be accessed by half-words (16 bits) or words (32 bits).

# 18.6.1 Control register (WWDG\_CR)

Address offset: 0x00

Reset value: 0x0000 007F

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Re	served							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			Pag	erved				WDGA				T[6:0]			
			nes	erved				rs				rw			

Bits 31:8 Reserved, must be kept at reset value.

### Bit 7 WDGA: Activation bit

This bit is set by software and only cleared by hardware after a reset. When WDGA = 1, the watchdog can generate a reset.

0: Watchdog disabled

1: Watchdog enabled

### Bits 6:0 T[6:0]: 7-bit counter (MSB to LSB)

These bits contain the value of the watchdog counter. It is decremented every (4096 x  $2^{WDGTB}$ ) PCLK1 cycles. A reset is produced when it rolls over from 0x40 to 0x3F (T6 becomes cleared).

# 18.6.2 Configuration register (WWDG\_CFR)

Address offset: 0x04

Reset value: 0x0000 007F

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Res	served							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		Rese	erved			EWI	WDGT	ΓΒ[1:0]				W[6:0]			
						rs	n	W				rw			

Bit 31:10 Reserved, must be kept at reset value.

### Bit 9 EWI: Early wakeup interrupt

When set, an interrupt occurs whenever the counter reaches the value 0x40. This interrupt is only cleared by hardware after a reset.

#### Bits 8:7 WDGTB[1:0]: Timer base

The time base of the prescaler can be modified as follows:

00: CK Counter Clock (PCLK1 div 4096) div 1

01: CK Counter Clock (PCLK1 div 4096) div 2

10: CK Counter Clock (PCLK1 div 4096) div 4 11: CK Counter Clock (PCLK1 div 4096) div 8

Bits 6:0 W[6:0]: 7-bit window value

These bits contain the window value to be compared to the downcounter.

# 18.6.3 Status register (WWDG\_SR)

Address offset: 0x08 Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Res	served							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							Reserve	, d							EWIF
										rc_w0					

Bit 31:1Reserved, must be kept at reset value.

# Bit 0 EWIF: Early wakeup interrupt flag

This bit is set by hardware when the counter has reached the value 0x40. It must be cleared by software by writing '0. A write of '1 has no effect. This bit is also set if the interrupt is not enabled.

# 18.6.4 WWDG register map

The following table gives the WWDG register map and reset values.

Table 74. WWDG register map and reset values

Offset	Register	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	6	8	7	9	5	4	3	2	_	0
0x00	WWDG_CR												lese													WDGA			1	[6:0]			
	Reset value																									0	1	1	1	1	1	1	1
0x04	WWDG_CFR										R	Rese	rved	ŀ										EWI	WDGTB1	WDGTB0			٧	V[6:0	]	·	
	Reset value																							0	0	0	1	1	1	1	1	1	1
0x08	WWDG_SR															Res	serv	ed															EWIF
	Reset value																																0

Refer to Table 2 on page 50 for the register boundary addresses.

# 19 Cryptographic processor (CRYP)

This section applies to the whole STM32F4xx family, unless otherwise specified.

## 19.1 CRYP introduction

The cryptographic processor can be used to both encipher and decipher data using the DES, Triple-DES or AES (128, 192, or 256) algorithms. It is a fully compliant implementation of the following standards:

- The data encryption standard (DES) and Triple-DES (TDES) as defined by Federal Information Processing Standards Publication (FIPS PUB 46-3, 1999 October 25). It follows the American National Standards Institute (ANSI) X9.52 standard.
- The advanced encryption standard (AES) as defined by Federal Information Processing Standards Publication (FIPS PUB 197, 2001 November 26)

The CRYP processor performs data encryption and decryption using DES and TDES algorithms in Electronic codebook (ECB) or Cipher block chaining (CBC) mode.

The CRYP peripheral is a 32-bit AHB2 peripheral. It supports DMA transfer for incoming and processed data, and has input and output FIFOs (each 8 words deep).

# 19.2 CRYP main features

- Suitable for AES, DES and TDES enciphering and deciphering operations
- AES
  - Supports the ECB, CBC, CTR chaining algorithms
  - Supports 128-, 192- and 256-bit keys
  - 4 x 32-bit initialization vectors (IV) used in the CBC, CTR modes

Table 75. Number of cycles required to process each 128-bit block

Algorithm / Key size	ECB	CBC	CTR
128b	14	14	14
192b	16	16	16
256b	18	18	18

### DES/TDES

- Direct implementation of simple DES algorithms (a single key, K1, is used)
- Supports the ECB and CBC chaining algorithms
- Supports 64-, 128- and 192-bit keys (including parity)
- 2 x 32-bit initialization vectors (IV) used in the CBC mode
- 16 HCLK cycles to process one 64-bit block in DES
- 48 HCLK cycles to process one 64-bit block in TDES

### Common to DES/TDES and AES

- IN and OUT FIFO (each with an 8-word depth, a 32-bit width, corresponding to 4 DES blocks or 2 AES blocks)
- Automatic data flow control with support of direct memory access (DMA) (using 2 channels, one for incoming data the other for processed data)
- Data swapping logic to support 1-, 8-, 16- or 32-bit data

# 19.3 CRYP functional description

The cryptographic processor implements a Triple-DES (TDES, that also supports DES) core and an AES cryptographic core. *Section 19.3.1* and *Section 19.3.2* provide details on these cores.

Since the TDES and the AES algorithms use block ciphers, incomplete input data blocks have to be padded prior to encryption (extra bits should be appended to the trailing end of the data string). After decryption, the padding has to be discarded. The hardware does not manage the padding operation, the software has to handle it.

Figure 195 shows the block diagram of the cryptographic processor.

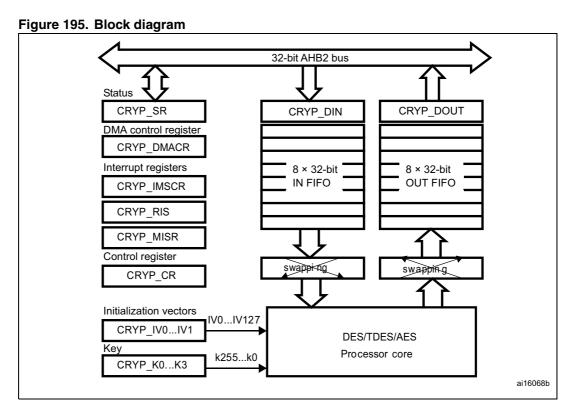


Figure 196.

### 19.3.1 DES/TDES cryptographic core

The DES/Triple-DES cryptographic core consists of three components:

- The DES algorithm (DEA)
- Multiple keys (1 for the DES algorithm, 1 to 3 for the TDES algorithm)
- The initialization vector (used in the CBC mode)

The basic processing involved in the TDES is as follows: an input block is read in the DEA and encrypted using the first key, K1 (K0 is not used in TDES mode). The output is then decrypted using the second key, K2, and encrypted using the third key, K3. The key depends on the algorithm which is used:

- DES mode: Key = [K1]
- TDES mode: Key = [K3 K2 K1]

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where Kx=[KxR KxL], R = right, L = left

According to the mode implemented, the resultant output block is used to calculate the ciphertext.

Note that the outputs of the intermediate DEA stages is never revealed outside the cryptographic boundary.

The TDES allows three different keying options:

Three independent keys

The first option specifies that all the keys are independent, that is, K1, K2 and K3 are independent. FIPS PUB 46-3 – 1999 (and ANSI X9.52 – 1998) refers to this option as the Keying Option 1 and, to the TDES as 3-key TDES.

Two independent keys

The second option specifies that K1 and K2 are independent and K3 is equal to K1, that is, K1 and K2 are independent, K3 = K1. FIPS PUB 46-3 - 1999 (and ANSI X9.52 - 1998) refers to this second option as the Keying Option 2 and, to the TDES as 2-key TDES.

Three equal keys

The third option specifies that K1, K2 and K3 are equal, that is, K1 = K2 = K3. FIPS PUB 46-3 - 1999 (and ANSI X9.52 - 1998) refers to the third option as the Keying Option 3. This "1-key" TDES is equivalent to single DES.

FIPS PUB 46-3 – 1999 (and ANSI X9.52-1998) provides a thorough explanation of the processing involved in the four operation modes supplied by the TDEA (TDES algorithm): TDES-ECB encryption, TDES-ECB decryption, TDES-CBC encryption and TDES-CBC decryption.

This reference manual only gives a brief explanation of each mode.

## DES and TDES Electronic codebook (DES/TDES-ECB) mode

DES/TDES-ECB mode encryption

Figure 197 illustrates the encryption in DES and TDES Electronic codebook (DES/TDES-ECB) mode. A 64-bit plaintext data block (P) is used after bit/byte/half-word swapping (refer to Section 19.3.3: Data type on page 507) as the input block (I). The input block is processed through the DEA in the encrypt state using K1. The output of this process is fed back directly to the input of the DEA where the DES is performed in the decrypt state using K2. The output of this process is fed back directly to the input of the DEA where the DES is performed in the encrypt state using K3. The resultant 64-bit output block (O) is used, after bit/byte/half-word swapping, as ciphertext (C) and it is pushed into the OUT FIFO.

DES/TDES-ECB mode decryption

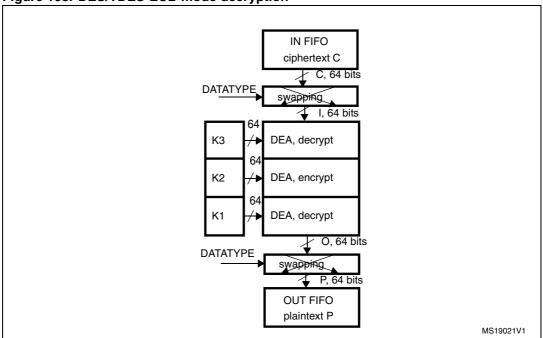
Figure 198 illustrates the DES/TDES-ECB decryption. A 64-bit ciphertext block (C) is used, after bit/byte/half-word swapping, as the input block (I). The keying sequence is reversed compared to that used in the encryption process. The input block is processed through the DEA in the decrypt state using K3. The output of this process is fed back directly to the input of the DEA where the DES is performed in the encrypt state using K2. The new result is directly fed to the input of the DEA where the DES is performed in the decrypt state using K1. The resultant 64-bit output block (O), after bit/byte/half-word swapping, produces the plaintext (P).

IN FIFO plaintext P P, 64 bits DATATYPE swapping DEA, encrypt 64 K2 DEA, decrypt ΚЗ DEA, encrypt O, 64 bits DATATYPE swapping **OUT FIFO** ciphertext C ai16069b

Figure 197. DES/TDES-ECB mode encryption

1. K: key; C: cipher text; I: input block; O: output block; P: plain text.

Figure 198. DES/TDES-ECB mode decryption



1. K: key; C: cipher text; I: input block; O: output block; P: plain text.

# DES and TDES Cipher block chaining (DES/TDES-CBC) mode

DES/TDES-CBC mode encryption

Figure 199 illustrates the DES and Triple-DES Cipher block chaining (DES/TDES-CBC) mode encryption. This mode begins by dividing a plaintext message into 64-bit data blocks. In TCBC encryption, the first input block (I1), obtained after bit/byte/half-word swapping (refer to Section 19.3.3: Data type on page 507), is formed by exclusive-ORing the first plaintext data block (P<sub>1</sub>) with a 64-bit initialization vector IV (I<sub>1</sub> = IV  $\oplus$  P<sub>1</sub>). The input block is processed through the DEA in the encrypt state using K1. The output of this process is fed back directly to the input of the DEA, which performs the DES in the decrypt state using K2. The output of this process is fed directly to the input of the DEA, which performs the DES in the encrypt state using K3. The resultant 64-bit output block (O<sub>1</sub>) is used directly as the ciphertext (C<sub>1</sub>), that is, C<sub>1</sub> = O<sub>1</sub>. This first ciphertext block is then exclusive-ORed with the second plaintext data block to produce the second input block,  $(I_2) = (C_1 \oplus P_2)$ . Note that  $I_2$  and  $P_2$  now refer to the second block. The second input block is processed through the TDEA to produce the second ciphertext block. This encryption process continues to "chain" successive cipher and plaintext blocks together until the last plaintext block in the message is encrypted. If the message does not consist of an integral number of data blocks, then the final partial data block should be encrypted in a manner specified for the application.

DES/TDES-CBC mode decryption

In DES/TDES-CBC decryption (see *Figure 200*), the first ciphertext block (C<sub>1</sub>) is used directly as the input block (I<sub>1</sub>). The keying sequence is reversed compared to that used for the encrypt process. The input block is processed through the DEA in the decrypt state using K3. The output of this process is fed directly to the input of the DEA where the DES is processed in the encrypt state using K2. This resulting value is directly fed to the input of the DEA where the DES is processed in the decrypt state using K1. The resulting output block is exclusive-ORed with the IV (which must be the same as that used during encryption) to produce the first plaintext block (P<sub>1</sub> = O<sub>1</sub>  $\oplus$  IV). The second ciphertext block is then used as the next input block and is processed through the TDEA. The resulting output block is exclusive-ORed with the first ciphertext block to produce the second plaintext data block (P<sub>2</sub> = O<sub>2</sub>  $\oplus$  C<sub>1</sub>). (Note that P<sub>2</sub> and O<sub>2</sub> refer to the second block of data.) The TCBC decryption process continues in this manner until the last complete ciphertext block has been decrypted. Ciphertext representing a partial data block must be decrypted in a manner specified for the application.

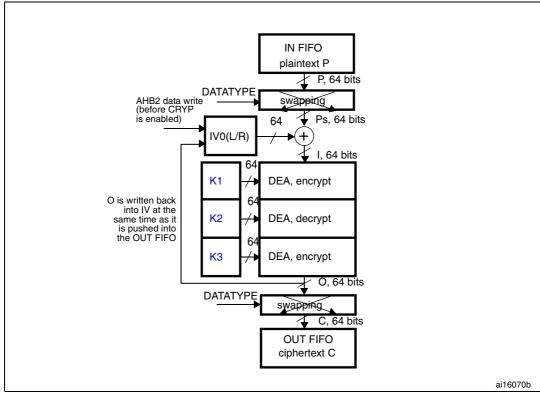


Figure 199. DES/TDES-CBC mode encryption

 K: key; C: cipher text; I: input block; O: output block; Ps: plain text before swapping (when decoding) or after swapping (when encoding); P: plain text; IV: initialization vectors.

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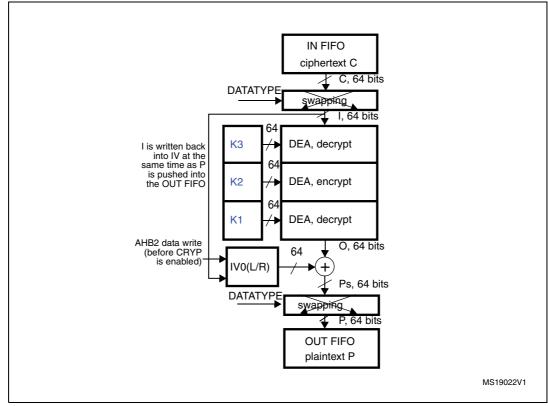


Figure 200. DES/TDES-CBC mode decryption

 K: key; C: cipher text; I: input block; O: output block; Ps: plain text before swapping (when decoding) or after swapping (when encoding); P: plain text; IV: initialization vectors.

# 19.3.2 AES cryptographic core

The AES cryptographic core consists of three components:

- The AES algorithm (AEA: advanced encryption algorithm)
- Multiple keys
- Initialization vector(s) or Nonce

The AES utilizes keys of 3 possible lengths: 128, 192 or 256 bits and, depending on the operation mode used, zero or one 128-bit initialization vector (IV).

The basic processing involved in the AES is as follows: an input block of 128 bits is read from the input FIFO and sent to the AEA to be encrypted using the key (K0...3). The key format depends on the key size:

- If Key size = 128: Key = [K3 K2]
- If Key size = 192: Key = [K3 K2 K1]
- If Key size = 256: Key = [K3 K2 K1 K0]

where Kx=[KxR KxL],R=right, L=left

According to the mode implemented, the resultant output block is used to calculate the ciphertext.

FIPS PUB 197 (November 26, 2001) provides a thorough explanation of the processing involved in the four operation modes supplied by the AES core: AES-ECB encryption, AES-

ECB decryption, AES-CBC encryption and AES-CBC decryption. This reference manual only gives a brief explanation of each mode.

### AES Electronic codebook (AES-ECB) mode

**AES-ECB** mode encryption

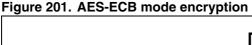
Figure 201 illustrates the AES Electronic codebook (AES-ECB) mode encryption. In AES-ECB encryption, a 128- bit plaintext data block (P) is used after bit/byte/halfword swapping (refer to Section 19.3.3: Data type on page 507) as the input block (I). The input block is processed through the AEA in the encrypt state using the 128, 192 or 256-bit key. The resultant 128-bit output block (O) is used after bit/byte/half-word swapping as ciphertext (C). It is then pushed into the OUT FIFO.

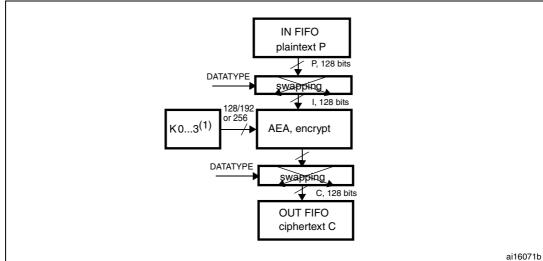
**AES-ECB** mode decryption

Figure 202 illustrates the AES Electronic codebook (AES-ECB) mode encryption.

To perform an AES decryption in the ECB mode, the secret key has to be prepared (it is necessary to execute the complete key schedule for encryption) by collecting the last round key, and using it as the first round key for the decryption of the ciphertext. This preparation function is computed by the AES core. Refer to Section 19.3.6: Procedure to perform an encryption or a decryption for more details on how to prepare the key.

In AES-ECB decryption, a 128-bit ciphertext block (C) is used after bit/byte/half-word swapping as the input block (I). The keying sequence is reversed compared to that of the encryption process. The resultant 128-bit output block (O), after bit/byte or halfword swapping, produces the plaintext (P).





- 1. K: key; C: cipher text; I: input block; O: output block; P: plain text.
- 2. If Key size = 128: Key = [K3 K2]. If Key size = 192: Key = [K3 K2 K1] If Key size = 256: Key = [K3 K2 K1 K0].

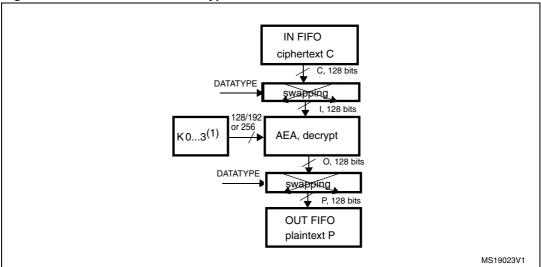


Figure 202. AES-ECB mode decryption

- 1. K: key; C: cipher text; I: input block; O: output block; P: plain text.
- 2. If Key size = 128 => Key = [K3 K2]. If Key size = 192 => Key = [K3 K2 K1] If Key size = 256 => Key = [K3 K2 K1 K0].

### AES Cipher block chaining (AES-CBC) mode

AES-CBC mode encryption

The AES Cipher block chaining (AES-CBC) mode decryption is shown on *Figure 203*. In AES-CBC encryption, the first input block (I<sub>1</sub>) obtained after bit/byte/half-word swapping (refer to *Section 19.3.3: Data type on page 507*) is formed by exclusive-ORing the first plaintext data block (P<sub>1</sub>) with a 128-bit initialization vector IV (I<sub>1</sub> = IV  $\oplus$  P<sub>1</sub>). The input block is processed through the AEA in the encrypt state using the 128-, 192- or 256-bit key (K0...K3). The resultant 128-bit output block (O<sub>1</sub>) is used directly as ciphertext (C<sub>1</sub>), that is, C<sub>1</sub> = O<sub>1</sub>. This first ciphertext block is then exclusive-ORed with the second plaintext data block to produce the second input block, (I<sub>2</sub>) = (C<sub>1</sub>  $\oplus$  P<sub>2</sub>). Note that I<sub>2</sub> and P<sub>2</sub> now refer to the second block. The second input block is processed through the AEA to produce the second ciphertext block. This encryption process continues to "chain" successive cipher and plaintext blocks together until the last plaintext block in the message is encrypted. If the message does not consist of an integral number of data blocks, then the final partial data block should be encrypted in a manner specified for the application.

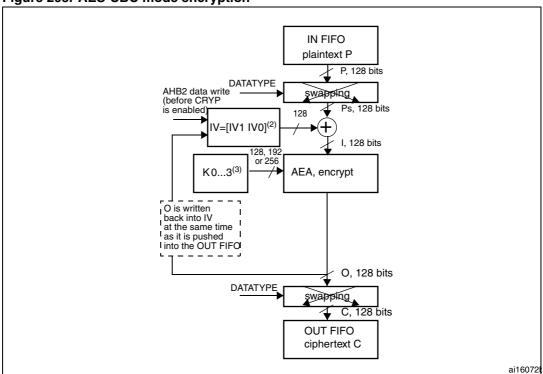
In the CBC mode, like in the ECB mode, the secret key must be prepared to perform an AES decryption. Refer to *Section 19.3.6: Procedure to perform an encryption or a decryption on page 512* for more details on how to prepare the key.

AES-CBC mode decryption

In AES-CBC decryption (see *Figure 204*), the first 128-bit ciphertext block (C<sub>1</sub>) is used directly as the input block (I<sub>1</sub>). The input block is processed through the AEA in the decrypt state using the 128-, 192- or 256-bit key. The resulting output block is exclusive-ORed with the 128-bit initialization vector IV (which must be the same as that used during encryption) to produce the first plaintext block (P<sub>1</sub> = O<sub>1</sub>  $\oplus$  IV). The second ciphertext block is then used as the next input block and is processed through the AEA. The resulting output block is exclusive-ORed with the first ciphertext block to produce the second plaintext data block (P<sub>2</sub> = O<sub>2</sub>  $\oplus$  C<sub>1</sub>). (Note that P<sub>2</sub> and O<sub>2</sub> refer to the second

block of data.) The AES-CBC decryption process continues in this manner until the last complete ciphertext block has been decrypted. Ciphertext representing a partial data block must be decrypted in a manner specified for the application.





- 1. K: key; C: cipher text; I: input block; O: output block; Ps: plain text before swapping (when decoding) or after swapping (when encoding); P: plain text; IV: Initialization vectors.
- 2. IVx=[IVxR IVxL], R=right, L=left.
- 3. If Key size = 128 => Key = [K3 K2]. If Key size = 192 => Key = [K3 K2 K1] If Key size = 256 => Key = [K3 K2 K1 K0].

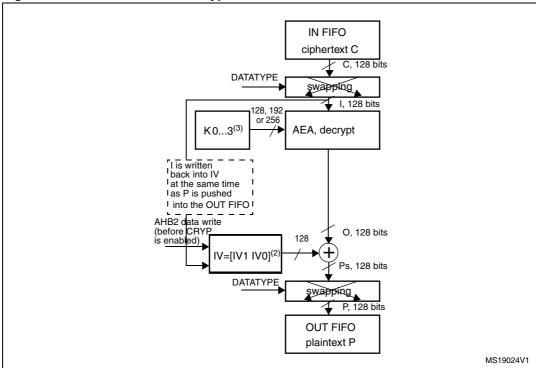


Figure 204. AES-CBC mode decryption

- 1. K: key; C: cipher text; I: input block; O: output block; Ps: plain text before swapping (when decoding) or after swapping (when encoding); P: plain text; IV: Initialization vectors.
- 2. IVx=[IVxR IVxL], R=right, L=left.
- 3. If Key size = 128 => Key = [K3 K2]. If Key size = 192 => Key = [K3 K2 K1] If Key size = 256 => Key = [K3 K2 K1 K0].

### **AES counter mode (AES-CTR) mode**

The AES counter mode uses the AES block as a key stream generator. The generated keys are then XORed with the plaintext to obtain the cipher. For this reason, it makes no sense to speak of different CTR encryption/decryption, since the two operations are exactly the same.

### In fact, given:

- Plaintext: P[0], P[1], ..., P[n] (128 bits each)
- A key K to be used (the size does not matter)
- An initial counter block (call it ICB but it has the same functionality as the IV of CBC)

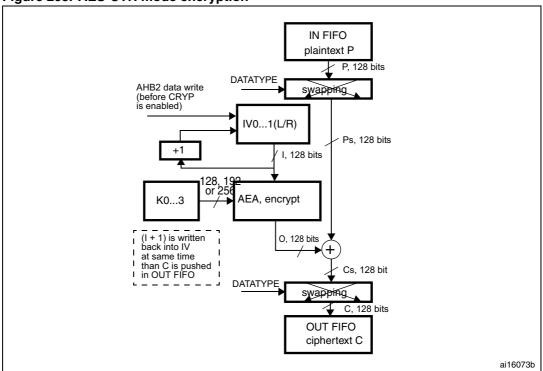
The cipher is computed as follows:

```
C[i] = enck(iv[i]) xor P[i], where:
  iv[0] = ICB and iv[i+1] = func(iv[i]), where func is an update function
  applied to the previous iv block; func is basically an increment of one of the fields
  composing the iv block.
```

Given that the ICB for decryption is the same as the one for encryption, the key stream generated during decryption is the same as the one generated during encryption. Then, the ciphertext is XORed with the key stream in order to retrieve the original plaintext. The decryption operation therefore acts exactly in the same way as the encryption operation.

Figure 205 and Figure 206 illustrate AES-CTR encryption and decryption, respectively.

Figure 205. AES-CTR mode encryption



K: key; C: cipher text; I: input Block; o: output block; Ps: plain text before swapping (when decoding) or after swapping (when encoding); Cs: cipher text after swapping (when decoding) or before swapping (when encoding); P: plain text; IV: Initialization vectors.

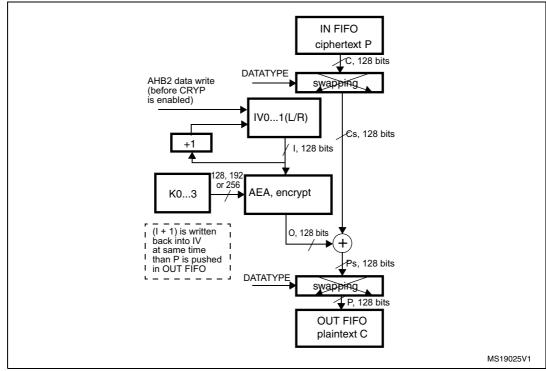
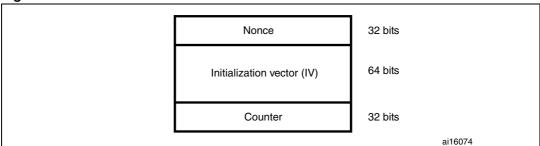


Figure 206. AES-CTR mode decryption

1. K: key; C: cipher text; I: input Block; o: output block; Ps: plain text before swapping (when decoding) or after swapping (when encoding); Cs: cipher text after swapping (when decoding) or before swapping (when encoding); P: plain text; IV: Initialization vectors.

Figure 207 shows the structure of the IV block as defined by the standard [2]. It is composed of three distinct fields.

Figure 207. Initial counter block structure for the Counter mode



- Nonce is a 32-bit, single-use value. A new nonce should be assigned to each different communication.
- The initialization vector (IV) is a 64-bit value and the standard specifies that the encryptor must choose IV so as to ensure that a given value is used only once for a given key
- The counter is a 32-bit big-endian integer that is incremented each time a block has been encrypted. The initial value of the counter should be set to '1'.

The block increments the least significant 32 bits, while it leaves the other (most significant) 96 bits unchanged.

## 19.3.3 Data type

Data enter the CRYP processor 32 bits (word) at a time as they are written into the CRYP\_DIN register. The principle of the DES is that streams of data are processed 64 bits by 64 bits and, for each 64-bit block, the bits are numbered from M1 to M64, with M1 the leftmost bit and M64 the right-most bit of the block. The same principle is used for the AES, but with a 128-bit block size.

The system memory organization is little-endian: whatever the data type (bit, byte, 16-bit half-word, 32-bit word) used, the least-significant data occupy the lowest address locations. A bit, byte, or half-word swapping operation (depending on the kind of data to be encrypted) therefore has to be performed on the data read from the IN FIFO before they enter the CRYP processor. The same swapping operation should be performed on the CRYP data before they are written into the OUT FIFO. For example, the operation would be byte swapping for an ASCII text stream.

The kind of data to be processed is configured with the DATATYPE bitfield in the CRYP control register (CRYP\_CR).

Table 76. Data types

DATATYPE in CRYP_CR	Swapping performed	System memory data (plaintext or cypher)
00b	No swapping	Example: TDES block value <b>0xABCD77206973FE01</b> is represented in system memory as:  TDES block size = 64bit = 2x 32 bit <b>0xABCD7720</b> 6973FE01 <b>oxABCD7720</b> @ <b>ox6973FE01 ox6973FE01</b>
01b	Half-word (16-bit) swapping	Example: TDES block value <b>0xABCD77206973FE01</b> is represented in system memory as:  TDES block size = 64bit = 2x 32 bit <b>0xABCD 7720 6973 FE01</b> OxFE01 6973  @ +4

Table 76. Data types

DATATYPE in CRYP_CR	Swapping performed	System memory data (plaintext or cypher)
		Example: TDES block value <b>0xABCD77206973FE01</b> is represented in system memory as:
10b	Byte (8-bit) swapping	system memory  TDES block size = 64bit = 2x 32 bit  0x AB CD 77 20 69 73 FE 01  0x 01 FE 73 69  0x 44
		TDES block value <b>0x4E6F772069732074</b> is represented in system memory as:  TDES Bloc size = 64bit = 2x 32 bit system memory
11b	Bit swapping	0x4E 6F 77 20 69 73 20 74

Figure 208 shows how the 64-bit data block M1...64 is constructed from two consecutive 32-bit words popped off the IN FIFO by the CRYP processor, according to the DATATYPE value. The same schematic can easily be extended to form the 128-bit block for the AES cryptographic algorithm (for the AES, the block length is four 32-bit words, but swapping only takes place at word level, so it is identical to the one described here for the TDES).

Note:

The same swapping is performed between the IN FIFO and the CRYP data block, and between the CRYP data block and the OUT FIFO.

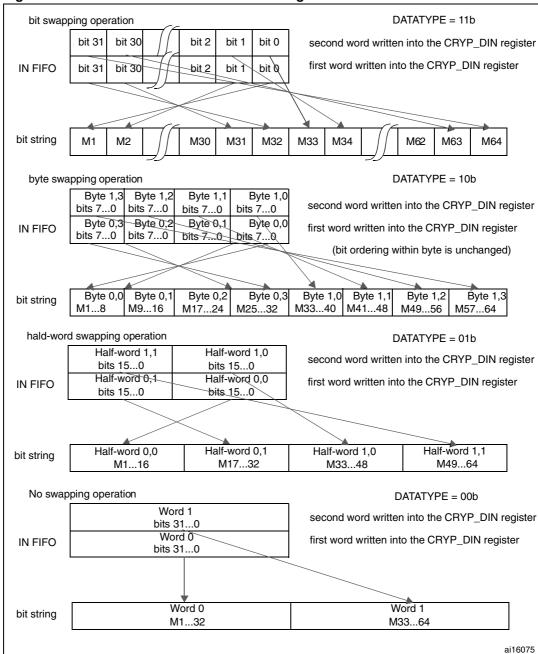


Figure 208. 64-bit block construction according to DATATYPE

## 19.3.4 Initialization vectors - CRYP\_IV0...1(L/R)

Initialization vectors are considered as two 64-bit data items. They therefore do not have the same data format and representation in system memory as plaintext or cypher data, and they are not affected by the DATATYPE value.

Initialization vectors are defined by two consecutive 32-bit words, CRYP\_IVL (left part, noted as bits IV1...32) and CRYP\_IVR (right part, noted as bits IV33...64).

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During the DES or TDES CBC encryption, the CRYP\_IV0(L/R) bits are XORed with the 64-bit data block popped off the IN FIFO after swapping (according to the DATATYPE value), that is, with the M1...64 bits of the data block. When the output of the DEA3 block is available, it is copied back into the CRYP\_IV0(L/R) vector, and this new content is XORed with the next 64-bit data block popped off the IN FIFO, and so on.

During the DES or TDES CBC decryption, the CRYP\_IV0(L/R) bits are XORed with the 64-bit data block (that is, with the M1...64 bits) delivered by the TDEA1 block before swapping (according to the DATATYPE value), and pushed into the OUT FIFO. When the XORed result is swapped and pushed into the OUT FIFO, the CRYP\_IV0(L/R) value is replaced by the output of the IN FIFO, then the IN FIFO is popped, and a new 64-bit data block can be processed.

During the AES CBC encryption, the CRYP\_IV0...1(L/R) bits are XORed with the 128-bit data block popped off the IN FIFO after swapping (according to the DATATYPE value). When the output of the AES core is available, it is copied back into the CRYP\_IV0...1(L/R) vector, and this new content is XORed with the next 128-bit data block popped off the IN FIFO, and so on.

During the AES CBC decryption, the CRYP\_IV0...1(L/R) bits are XORed with the 128-bit data block delivered by the AES core before swapping (according to the DATATYPE value) and pushed into the OUT FIFO. When the XORed result is swapped and pushed into the OUT FIFO, the CRYP\_IV0...1(L/R) value is replaced by the output of the IN FIFO, then the IN FIFO is popped, and a new 128-bit data block can be processed.

During the AES CTR encryption or decryption, the CRYP\_IV0...1(L/R) bits are encrypted by the AES core. Then the result of the encryption is XORed with the 128-bit data block popped off the IN FIFO after swapping (according to the DATATYPE value). When the XORed result is swapped and pushed into the OUT FIFO, the counter part of the CRYP\_IV0...1(L/R) value (32 LSB) is incremented.

Any write operation to the CRYP\_IV0...1(L/R) registers when bit BUSY = 1b in the CRYP\_SR register is disregarded (CRYP\_IV0...1(L/R) register content not modified). Thus, you must check that bit BUSY = 0b before modifying initialization vectors.

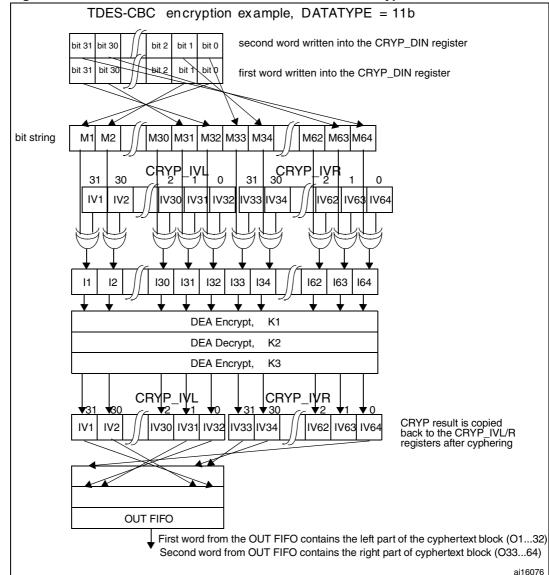


Figure 209. Initialization vectors use in the TDES-CBC encryption

## 19.3.5 CRYP busy state

When there is enough data in the input FIFO (at least 2 words for the DES or TDES algorithm mode, 4 words for the AES algorithm mode) and enough free-space in the output FIFO (at least 2 (DES/TDES) or 4 (AES) word locations), and when the bit CRYPEN = 1 in the CRYP\_CR register, then the cryptographic processor automatically starts an encryption or decryption process (according to the value of the ALGODIR bit in the CRYP\_CR register).

This process takes 48 AHB2 clock cycles for the Triple-DES algorithm, 16 AHB2 clock cycles for the simple DES algorithm, and 14, 16 or 18 AHB2 clock cycles for the AES with key lengths of 128, 192 or 256 bits, respectively. During the whole process, the BUSY bit in the CRYP\_SR register is set to '1'. At the end of the process, two (DES/TDES) or four (AES) words are written by the CRYP Core into the output FIFO, and the BUSY bit is cleared. In the CBC, CTR mode, the initialization vectors CRYP\_IVx(L/R)R (x = 0..3) are updated as well.

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A write operation to the key registers (CRYP\_Kx(L/R)R, x = 0..3), the initialization registers (CRYP\_IVx(L/R)R, x = 0..3), or to bits [9:2] in the CRYP\_CR register are ignored when the cryptographic processor is busy (bit BUSY = 1b in the CRYP\_SR register), and the registers are not modified. It is thus not possible to modify the configuration of the cryptographic processor while it is processing a block of data. It is however possible to clear the CRYPEN bit while BUSY = 1, in which case the ongoing DES, TDES or AES processing is completed and the two/four word results are written into the output FIFO, and then, only then, the BUSY bit is cleared.

Note:

When a block is being processed in the DES or TDES mode, if the output FIFO becomes full and if the input FIFO contains at least one new block, then the new block is popped off the input FIFO and the BUSY bit remains high until there is enough space to store this new block into the output FIFO.

#### 19.3.6 Procedure to perform an encryption or a decryption

#### Initialization

- 1. Initialize the peripheral (the order of operations is not important except for the key preparation for AES-ECB or AES-CBC decryption. The key size and the key value must be entered before preparing the key and the algorithm must be configured once the key has been prepared):
  - a) Configure the key size (128-, 192- or 256-bit, in the AES only) with the KEYSIZE bits in the CRYP\_CR register
  - b) Write the symmetric key into the CRYP\_KxL/R registers (2 to 8 registers to be written depending on the algorithm)
  - c) Configure the data type (1-, 8-, 16- or 32-bit), with the DATATYPE bits in the CRYP\_CR register
  - d) In case of decryption in AES-ECB or AES-CBC, you must prepare the key: configure the key preparation mode by setting the ALGOMODE bits to '111' in the CRYP\_CR register. Then write the CRYPEN bit to '1': the BUSY bit is set. Wait until BUSY returns to 0 (CRYPEN is automatically cleared as well): the key is prepared for decryption
  - e) Configure the algorithm and chaining (the DES/TDES in ECB/CBC, the AES in ECB/CBC/CTR) with the ALGOMODE bits in the CRYP\_CR register
  - f) Configure the direction (encryption/decryption), with the ALGODIR bit in the CRYP\_CR register
  - g) Write the initialization vectors into the CRYP\_IVxL/R register (in CBC or CTR modes only)
- 2. Flush the IN and OUT FIFOs by writing the FFLUSH bit to 1 in the CRYP\_CR register

#### Processing when the DMA is used to transfer the data from/to the memory

- 1. Configure the DMA controller to transfer the input data from the memory. The transfer length is the length of the message. As message padding is not managed by the peripheral, the message length must be an entire number of blocks. The data are transferred in burst mode. The burst length is 4 words in the AES and 2 or 4 words in the DES/TDES. The DMA should be configured to set an interrupt on transfer completion of the output data to indicate that the processing is finished.
- 2. Enable the cryptographic processor by writing the CRYPEN bit to 1. Enable the DMA requests by setting the DIEN and DOEN bits in the CRYP\_DMACR register.

3. All the transfers and processing are managed by the DMA and the cryptographic processor. The DMA interrupt indicates that the processing is complete. Both FIFOs are normally empty and BUSY = 0.

#### Processing when the data are transferred by the CPU during interrupts

- 1. Enable the interrupts by setting the INIM and OUTIM bits in the CRYP\_IMSCR register.
- 2. Enable the cryptographic processor by setting the CRYPEN bit in the CRYP\_CR register.
- 3. In the interrupt managing the input data: load the input message into the IN FIFO. You can load 2 or 4 words at a time, or load data until the FIFO is full. When the last word of the message has been entered into the FIFO, disable the interrupt by clearing the INIM bit.
- 4. In the interrupt managing the output data: read the output message from the OUT FIFO. You can read 1 block (2 or 4 words) at a time or read data until the FIFO is empty. When the last word has been read, INIM=0, BUSY=0 and both FIFOs are empty (IFEM=1 and OFNE=0). You can disable the interrupt by clearing the OUTIM bit and, the peripheral by clearing the CRYPEN bit.

#### Processing without using the DMA nor interrupts

- 1. Enable the cryptographic processor by setting the CRYPEN bit in the CRYP\_CR register.
- 2. Write the first blocks in the input FIFO (2 to 8 words).
- 3. Repeat the following sequence until the complete message has been processed:
  - a) Wait for OFNE=1, then read the OUT-FIFO (1 block or until the FIFO is empty)
  - b) Wait for IFNF=1, then write the IN FIFO (1 block or until the FIFO is full)
- 4. At the end of the processing, BUSY=0 and both FIFOs are empty (IFEM=1 and OFNE=0). You can disable the peripheral by clearing the CRYPEN bit.

## 19.3.7 Context swapping

If a context switching is needed because a new task launched by the OS requires this resource, the following tasks have to be performed for full context restoration (example when the DMA is used):

#### Case of the AES and DES

- 1. Context saving
  - Stop DMA transfers on the IN FIFO by clearing the DIEN bit in the CRYP\_DMACR register.
  - b) Wait until both the IN and OUT FIFOs are empty (IFEM=1 and OFNE=0 in the CRYP\_SR register) and the BUSY bit is cleared.
  - Stop DMA transfers on the OUT FIFO by writing the DOEN bit to 0 in the CRYP\_DMACR register and clear the CRYPEN bit.
  - d) Save the current configuration (bits [9:2] and bits 19 in the CRYP\_CR register) and, if not in ECB mode, the initialization vectors. The key value must already be available in the memory. When needed, save the DMA status (pointers for IN and OUT messages, number of remaining bytes, etc.).
- Configure and execute the other processing.

#### Context restoration

- a) Configure the processor as in *Section 19.3.6: Procedure to perform an encryption or a decryption on page 512, Initialization* with the saved configuration. For the AES-ECB or AES-CBC decryption, the key must be prepared again.
- b) If needed, reconfigure the DMA controller to transfer the rest of the message.
- c) Enable the processor by setting the CRYPEN bit and, the DMA requests by setting the DIEN and DOEN bits.

#### Case of the TDES

Context swapping can be done in the TDES in the same way as in the AES. But as the input FIFO can contain up to 4 unprocessed blocks and as the processing duration per block is higher, it can be faster in certain cases to interrupt the processing without waiting for the IN FIFO to be empty.

- Context saving
  - Stop DMA transfers on the IN FIFO by clearing the DIEN bit in the CRYP\_DMACR register.
  - Disable the processor by clearing the CRYPEN bit (the processing will stop at the end of the current block).
  - c) Wait until the OUT FIFO is empty (OFNE=0 in the CRYP\_SR register) and the BUSY bit is cleared.
  - d) Stop DMA transfers on the OUT FIFO by writing the DOEN bit to 0 in the CRYP\_DMACR register.
  - e) Save the current configuration (bits [9:2] and bits 19 in the CRYP\_CR register) and, if not in ECB mode, the initialization vectors. The key value must already be available in the memory. When needed, save the DMA status (pointers for IN and OUT messages, number of remaining bytes, etc.). Read back the data loaded in the IN FIFO that have not been processed and save them in the memory until the FIFO is empty.
- 2. Configure and execute the other processing.
- 3. Context restoration
  - a) Configure the processor as in *Section 19.3.6: Procedure to perform an encryption or a decryption on page 512, Initialization* with the saved configuration. For the AES-ECB or AES-CBC decryption, the key must be prepared again.
  - b) Write the data that were saved during context saving into the IN FIFO.
  - c) If needed, reconfigure the DMA controller to transfer the rest of the message.
  - Enable the processor by setting the CRYPEN bit and, the DMA requests by setting the DIEN and DOEN bits.

# 19.4 CRYP interrupts

There are two individual maskable interrupt sources generated by the CRYP. These two sources are combined into a single interrupt signal, which is the only interrupt signal from the CRYP that drives the NVIC (nested vectored interrupt controller). This combined interrupt, which is an OR function of the individual masked sources, is asserted if any of the individual interrupts listed below is asserted and enabled.

You can enable or disable the interrupt sources individually by changing the mask bits in the CRYP\_IMSCR register. Setting the appropriate mask bit to '1' enables the interrupt.

The status of the individual interrupt sources can be read either from the CRYP\_RISR register, for raw interrupt status, or from the CRYP\_MISR register, for the masked interrupt status.

#### **Output FIFO service interrupt - OUTMIS**

The output FIFO service interrupt is asserted when there is one or more (32-bit word) data items in the output FIFO. This interrupt is cleared by reading data from the output FIFO until there is no valid (32-bit) word left (that is, the interrupt follows the state of the OFNE (output FIFO not empty) flag).

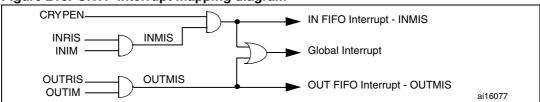
The output FIFO service interrupt OUTMIS is NOT enabled with the CRYP enable bit. Consequently, disabling the CRYP will not force the OUTMIS signal low if the output FIFO is not empty.

#### Input FIFO service interrupt - INMIS

The input FIFO service interrupt is asserted when there are less than four words in the input FIFO. It is cleared by performing write operations to the input FIFO until it holds four or more words.

The input FIFO service interrupt INMIS is enabled with the CRYP enable bit. Consequently, when CRYP is disabled, the INMIS signal is low even if the input FIFO is empty.

Figure 210. CRYP interrupt mapping diagram



#### 19.5 CRYP DMA interface

The cryptographic processor provides an interface to connect to the DMA controller. The DMA operation is controlled through the CRYP DMA control register, CRYP\_DMACR.

The burst and single transfer request signals are not mutually exclusive. They can both be asserted at the same time. For example, when there are 6 words available in the OUT FIFO, the burst transfer request and the single transfer request are asserted. After a burst transfer of 4 words, the single transfer request only is asserted to transfer the last 2 available words. This is useful for situations where the number of words left to be received in the stream is less than a burst.

Each request signal remains asserted until the relevant DMA clear signal is asserted. After the request clear signal is deasserted, a request signal can become active again, depending on the above described conditions. All request signals are deasserted if the CRYP peripheral is disabled or the DMA enable bit is cleared (DIEN bit for the IN FIFO and DOEN bit for the OUT FIFO in the CRYP\_DMACR register).

Note:

The DMA controller must be configured to perform burst of 4 words or less. Otherwise some data could be lost.

In order to let the DMA controller empty the OUT FIFO before filling up the IN FIFO, the OUTDMA channel should have a higher priority than the INDMA channel.

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# 19.6 CRYP registers

The cryptographic core is associated with several control and status registers, eight key registers and four initialization vectors registers.

# 19.6.1 CRYP control register (CRYP\_CR)

Address offset: 0x00

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
CRYPEN	FFLUSH		Rese	rvod		KEY	SIZE	DATA	TYPE	ALG	OMODE	[2:0]	ALGODIR	Res.	Res.
rw	w		nese	ii veu		rw	rw	rw	rw	rw	rw	rw	rw		

Bit 31:18 Reserved, must be kept at reset value

Bit 17:16 Reserved, must be kept at reset value

Bit 15 CRYPEN: Cryptographic processor enable

0: CRYP processor is disabled1: CRYP processor is enabled

Note: The CRYPEN bit is automatically cleared by hardware when the key preparation process ends (ALGOMODE=111b) or GCM\_CCM init Phase

Bit 14 FFLUSH: FIFO flush

When CRYPEN = 0, writing this bit to 1 flushes the IN and OUT FIFOs (that is read and write pointers of the FIFOs are reset. Writing this bit to 0 has no effect. When CRYPEN = 1, writing this bit to 0 or 1 has no effect.

Reading this bit always returns 0.

Bits 13:10 Reserved, must be kept at reset value

Bits 9:8 **KEYSIZE[1:0]:** Key size selection (AES mode only)

This bitfield defines the bit-length of the key used for the AES cryptographic core. This bitfield is 'don't care' in the DES or TDES modes.

00: 128 bit key length 01: 192 bit key length 10: 256 bit key length

11: Reserved, do not use this value

#### Bits 7:6 DATATYPE[1:0]: Data type selection

This bitfield defines the format of data entered in the CRYP\_DIN register (refer to *Section 19.3.3: Data type*).

00: 32-bit data. No swapping of each word. First word pushed into the IN FIFO (or popped off the OUT FIFO) forms bits 1...32 of the data block, the second word forms bits 33...64.

01: 16-bit data, or half-word. Each word pushed into the IN FIFO (or popped off the OUT FIFO) is considered as 2 half-words, which are swapped with each other.

10: 8-bit data, or bytes. Each word pushed into the IN FIFO (or popped off the OUT FIFO) is considered as 4 bytes, which are swapped with each other.

11: bit data, or bit-string. Each word pushed into the IN FIFO (or popped off the OUT FIFO) is considered as 32 bits (1st bit of the string at position 0), which are swapped with each other.

#### Bits 5:3 ALGOMODE[2:0]: Algorithm mode

000: TDES-ECB (triple-DES Electronic codebook): no feedback between blocks of data. Initialization vectors (CRYP\_IV0(L/R)) are not used, three key vectors (K1, K2, and K3) are used (K0 is not used).

001: TDES-CBC (triple-DES Cipher block chaining): output block is XORed with the subsequent input block before its entry into the algorithm. Initialization vectors (CRYP\_IV0L/R) must be initialized, three key vectors (K1, K2, and K3) are used (K0 is not used).

010: DES-ECB (simple DES Electronic codebook): no feedback between blocks of data. Initialization vectors (CRYP\_IV0L/R) are not used, only one key vector (K1) is used (K0, K2, K3 are not used).

011: DES-CBC (simple DES Cipher block chaining): output block is XORed with the subsequent input block before its entry into the algorithm. Initialization vectors (CRYP\_IV0L/R) must be initialized. Only one key vector (K1) is used (K0, K2, K3 are not used).

100: AES-ECB (AES Electronic codebook): no feedback between blocks of data. Initialization vectors (CRYP\_IV0L/R...1L/R) are not used. All four key vectors (K0...K3) are used.

101: AES-CBC (AES Cipher block chaining): output block is XORed with the subsequent input block before its entry into the algorithm. Initialization vectors (CRYP\_IV0L/R...1L/R) must be initialized. All four key vectors (K0...K3) are used. 110: AES-CTR (AES counter mode): output block is XORed with the subsequent input block before its entry into the algorithm. Initialization vectors (CRYP\_IV0L/R...1L/R) must be initialized. All four key vectors (K0...K3) are used. CTR decryption does not differ from CTR encryption, since the core always encrypts the current counter block to produce the key stream that will be XORed with the plaintext or cipher in input. Thus, ALGODIR is don't care when ALGOMODE = 110b, and the key must NOT be unrolled (prepared) for decryption.

111: AES key preparation for decryption mode. Writing this value when CRYPEN = 1 immediately starts an AES round for key preparation. The secret key must have previously been loaded into the K0...K3 registers. The BUSY bit in the CRYP\_SR register is set during the key preparation. After key processing, the resulting key is copied back into the K0...K3 registers, and the BUSY bit is cleared.

#### Bit 2 ALGODIR: Algorithm direction

0: Encrypt1: Decrypt

Bit 1:0 Reserved, must be kept at reset value

Note:

Writing to the KEYSIZE, DATATYPE, ALGOMODE and ALGODIR bits while BUSY=1 has no effect. These bits can only be configured when BUSY=0.

The FFLUSH bit has to be set only when BUSY=0. If not, the FIFO is flushed, but the block being processed may be pushed into the output FIFO just after the flush operation, resulting in a nonempty FIFO condition.

## 19.6.2 CRYP status register (CRYP\_SR)

Address offset: 0x04

Reset value: 0x0000 0003

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
					Dagamiaa	ı					BUSY	OFFU	OFNE	IFNF	IFEM
					Reserved	ı					r	r	r	r	r

Bit 31:5 Reserved, must be kept at reset value

#### Bit 4 BUSY: Busy bit

- 0: The CRYP Core is not processing any data. The reason is either that:
  - the CRYP core is disabled (CRYPEN=0 in the CRYP\_CR register) and the last processing has completed, or
  - The CRYP core is waiting for enough data in the input FIFO or enough free space in the output FIFO (that is in each case at least 2 words in the DES, 4 words in the AES).
- 1: The CRYP core is currently processing a block of data or a key preparation (for AES decryption).

Bit 3 OFFU: Output FIFO full

0: Output FIFO is not full

1: Output FIFO is full

Bits 2 OFNE: Output FIFO not empty

0: Output FIFO is empty

1: Output FIFO is not empty

Bit 1 IFNF: Input FIFO not full

0: Input FIFO is full1: Input FIFO is not full

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Bits 0 **IFEM:** Input FIFO empty 0: Input FIFO is not empty

1: Input FIFO is empty

## 19.6.3 CRYP data input register (CRYP\_DIN)

Address offset: 0x08

Reset value: 0x0000 0000

The CRYP\_DIN register is the data input register. It is 32-bit wide. It is used to enter up to four 64-bit (TDES) or two 128-bit (AES) plaintext (when encrypting) or ciphertext (when decrypting) blocks into the input FIFO, one 32-bit word at a time.

The first word written into the FIFO is the MSB of the input block. The LSB of the input block is written at the end. Disregarding the data swapping, this gives:

- In the DES/TDES modes: a block is a sequence of bits numbered from bit 1 (leftmost bit) to bit 64 (rightmost bit). Bit 1 corresponds to the MSB (bit 31) of the first word entered into the FIFO, bit 64 corresponds to the LSB (bit 0) of the second word entered into the FIFO.
- In the AES mode: a block is a sequence of bits numbered from 0 (leftmost bit) to 127 (rightmost bit). Bit 0 corresponds to the MSB (bit 31) of the first word written into the FIFO, bit 127 corresponds to the LSB (bit 0) of the 4th word written into the FIFO.

To fit different data sizes, the data written in the CRYP\_DIN register can be swapped before being processed by configuring the DATATYPE bits in the CRYP\_CR register. Refer to *Section 19.3.3: Data type on page 507* for more details.

When CRYP\_DIN register is written to, the data are pushed into the input FIFO. When at least two 32-bit words in the DES/TDES mode (or four 32-bit words in the AES mode) have been pushed into the input FIFO, and when at least 2 words are free in the output FIFO, the CRYP engine starts an encrypting or decrypting process. This process takes two 32-bit words in the DES/TDES mode (or four 32-bit words in the AES mode) from the input FIFO and delivers two 32-bit words (or 4, respectively) to the output FIFO per process round.

When CRYP DIN register is read:

- If CRYPEN = 0, the FIFO is popped, and then the data present in the Input FIFO are returned, from the oldest one (first reading) to the newest one (last reading). The IFEM flag must be checked before each read operation to make sure that the FIFO is not empty.
- if CRYPEN = 1, an undefined value is returned.

After the CRYP\_DIN register has been read once or several times, the FIFO must be flushed by setting the FFLUSH bit prior to processing new data.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							DAT	AIN							
rw	rw	rw	rw	rw	rw	rw	rw	rw							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							DAT	AIN							
rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bit 31:0 DATAIN: Data input

Read = returns Input FIFO content if CRYPEN = 0, else returns an undefined value.

Write = Input FIFO is written.

## 19.6.4 CRYP data output register (CRYP\_DOUT)

Address offset: 0x0C

Reset value: 0x0000 0000

The CRYP\_DOUT register is the data output register. It is read-only and 32-bit wide. It is used to retrieve up to four 64-bit (TDES mode) or two 128-bit (AES mode) ciphertext (when encrypting) or plaintext (when decrypting) blocks from the output FIFO, one 32-bit word at a time.

Like for the input data, the MSB of the output block is the first word read from the output FIFO. The LSB of the output block is read at the end. Disregarding data swapping, this gives:

- In the DES/TDES modes: Bit 1 (leftmost bit) corresponds to the MSB (bit 31) of the first word read from the FIFO, bit 64 (rightmost bit) corresponds to the LSB (bit 0) of the second word read from the FIFO.
- In the AES mode: Bit 0 (leftmost bit) corresponds to the MSB (bit 31) of the first word read from the FIFO, bit 127 (rightmost bit) corresponds to the LSB (bit 0) of the 4th word read from the FIFO.

To fit different data sizes, the data can be swapped after processing by configuring the DATATYPE bits in the CRYP\_CR register. Refer to *Section 19.3.3: Data type on page 507* for more details.

When CRYP\_DOUT register is read, the last data entered into the output FIFO (pointed to by the read pointer) is returned.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							DATA	OUT							
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							DATA	OUT							
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bit 31:0 DATAOUT: Data output

Read = returns output FIFO content.

Write = No effect.

## 19.6.5 CRYP DMA control register (CRYP\_DMACR)

Address offset: 0x10

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
						Poor	erved							DOEN	DIEN
						nese	erved							rw	rw

Bit 31:2 Reserved, must be kept at reset value

Bit 1 DOEN: DMA output enable

0: DMA for outgoing data transfer is disabled1: DMA for outgoing data transfer is enabled

Bit 0 DIEN: DMA input enable

0: DMA for incoming data transfer is disabled1: DMA for incoming data transfer is enabled

## 19.6.6 CRYP interrupt mask set/clear register (CRYP\_IMSCR)

Address offset: 0x14

Reset value: 0x0000 0000

The CRYP\_IMSCR register is the interrupt mask set or clear register. It is a read/write register. On a read operation, this register gives the current value of the mask on the relevant interrupt. Writing 1 to the particular bit sets the mask, enabling the interrupt to be read. Writing 0 to this bit clears the corresponding mask. All the bits are cleared to 0 when the peripheral is reset.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Res	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
						D	d							OUTIM	INIM
						Hese	erved							rw	rw

Bit 31:2 Reserved, must be kept at reset value

Bit 1 **OUTIM:** Output FIFO service interrupt mask

0: Output FIFO service interrupt is masked

1: Output FIFO service interrupt is not masked

Bit 0 INIM: Input FIFO service interrupt mask

0: Input FIFO service interrupt is masked

1: Input FIFO service interrupt is not masked

## 19.6.7 CRYP raw interrupt status register (CRYP\_RISR)

Address offset: 0x18

Reset value: 0x0000 0001

The CRYP\_RISR register is the raw interrupt status register. It is a read-only register. On a read, this register gives the current raw status of the corresponding interrupt prior to

masking. A write has no effect.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Res	served							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
						Page	nuad							OUTRIS	INRIS
						nese	erved							r	r

Bit 31:2 Reserved, must be kept at reset value

Bit 1 OUTRIS: Output FIFO service raw interrupt status

Gives the raw interrupt state prior to masking of the output FIFO service interrupt.

0: Raw interrupt not pending

1: Raw interrupt pending

Bit 0 INRIS: Input FIFO service raw interrupt status

Gives the raw interrupt state prior to masking of the Input FIFO service interrupt.

0: Raw interrupt not pending

1: Raw interrupt pending

## 19.6.8 CRYP masked interrupt status register (CRYP\_MISR)

Address offset: 0x1C

Reset value: 0x0000 0000

The CRYP\_MISR register is the masked interrupt status register. It is a read-only register. On a read, this register gives the current masked status of the corresponding interrupt prior to masking. A write has no effect.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Res	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
						Page	erved							OUTMIS	INMIS
						Hese	erved							r	r

Bit 31:2 Reserved, must be kept at reset value

Bit 1 OUTMIS: Output FIFO service masked interrupt status

Gives the interrupt state after masking of the output FIFO service interrupt.

0: Interrupt not pending

1: Interrupt pending

Bit 0 INMIS: Input FIFO service masked interrupt status

Gives the interrupt state after masking of the input FIFO service interrupt.

0: Interrupt not pending

1: Interrupt pending when CRYPEN = 1

## 19.6.9 CRYP key registers (CRYP\_K0...3(L/R)R)

Address offset: 0x20 to 0x3C Reset value: 0x0000 0000

These registers contain the cryptographic keys.

In the TDES mode, keys are 64-bit binary values (number from left to right, that is the leftmost bit is bit 1), named K1, K2 and K3 (K0 is not used), each key consists of 56 information bits and 8 parity bits. The parity bits are reserved for error detection purposes and are not used by the current block. Thus, bits 8, 16, 24, 32, 40, 48, 56 and 64 of each 64-bit key value Kx[1:64] are not used.

In the AES mode, the key is considered as a single 128-, 192- or 256-bit long bit sequence,  $k_0k_1k_2...k_{127/191/255}$  ( $k_0$  being the leftmost bit). The AES key is entered into the registers as follows:

- for AES-128: k<sub>0</sub>..k<sub>127</sub> corresponds to b<sub>127</sub>..b<sub>0</sub> (b<sub>255</sub>..b<sub>128</sub> are not used),
- for AES-192:  $k_0..k_{191}$  corresponds to  $b_{191}..b_0$  ( $b_{255}..b_{192}$  are not used),
- for AES-256: k<sub>0</sub>..k<sub>255</sub> corresponds to b<sub>255</sub>..b<sub>0</sub>.

In any case b<sub>0</sub> is the rightmost bit.

#### CRYP\_K0LR (address offset: 0x20)

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
b255	b254	b253	b252	b251	b250	b249	b248	b247	b246	b245	b244	b243	b242	b241	b240
w	w	w	w	w	w	w	W	w	w	w	W	W	w	w	w
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
15 b239	14 b238	13 b237	12 b236	11 b235	10 b234	9 b233	8 b232	7 b231	6 b230	5 b229	4 b228	3 b227	2 b226	1 b225	0 b224

### CRYP\_K0RR (address offset: 0x24)

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
b223	b222	b221	b220	b219	b218	b217	b216	b215	b214	b213	b212	b211	b210	b209	b208
w	w	W	W	W	W	W	W	W	W	W	W	W	W	w	W
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
15 b207	14 b206	13 b205	12 b204	11 b203	10 b202	9 b201	8 b200	7 b199	6 b198	5 b197	4 b196	3 b195	2 b194	1 b193	0 b192

#### CRYP\_K1LR (address offset: 0x28)

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
k1.1	k1.2	k1.3	k1.4	k1.5	k1.6	k1.7	k1.8	k1.9	k1.10	k1.11	k1.12	k1.13	k1.14	k1.15	k1.16
b191	b190	b189	b188	b187	b186	b185	b184	b183	b182	b181	b180	b179	b178	b177	b176
W	W	W	W	W	W	W	W	w	w	W	W	W	w	W	W

	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	k1.17 b175	k1.18 b174	k1.19 b173	k1.20 b172	k1.21 b171	k1.22 b170	k1.23 b169	k1.24 b168	k1.25 b167	k1.26 b166	k1.27 b165	k1.28 b164	k1.29 b163	k1.30 b162	k1.31 b161	k1.32 b160
Ī	W	w	w	w	w	w	w	w	W	W	W	W	w	w	W	w

## CRYP\_K1RR (address offset: 0x2C)

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
k1.33 b159	k1.34 b158	k1.35 b157	k1.36 b156	k1.37 b155	k1.38 b154	k1.39 b153	k1.40 b152	k1.41 b151	k1.42 b150	k1.43 b149	k1.44 b148	k1.45 b147	k1.46 b146	k1.47 b145	k1.48 b144
w	w	W	W	W	W	W	W	W	W	W	W	W	W	w	w
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
15 k1.49 b143	14 k1.50 b142	13 k1.51 b141	12 k1.52 b140	11 k1.53 b139	10 k1.54 b138	9 k1.55 b137	8 k1.56 b136	7 k1.57 b135	6 k1.58 b134	5 k1.59 b133	4 k1.60 b132	3 k1.61 b131	2 k1.62 b130	1 k1.63 b129	0 k1.64 b128

# CRYP\_K2LR (address offset: 0x30)

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
k2.1 b127	k2.2 b126	k2.3 b125	k2.4 b124	k2.5 b123	k2.6 b122	k2.7 b121	k2.8 b120	k2.9 b119	k2.10 b118	k2.11 b117	k2.12 b116	k2.13 b115	k2.14 b114	k2.15 b113	k2.16 b112
w	w	w	w	w	w	w	w	w	w	w	w	w	w	w	w
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
15 k2.17 b111	14 k2.18 b110	13 k2.19 b109	12 k2.20 b108	11 k2.21 b107	10 k2.22 b106	9 k2.23 b105	8 k2.24 b104	7 k2.25 b103	6 k2.26 b102	5 k2.27 b101	4 k2.28 b100	3 k2.29 b99	2 k2.30 b98	1 k2.31 b97	0 k2.32 b96

# CRYP\_K2RR (address offset: 0x34)

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
k2.33 b95	k2.34 b94	k2.35 b93	k2.36 b92	k2.37 b91	k2.38 b90	k2.39 b89	k2.40 b88	k2.41 b87	k2.42 b86	k2.43 b85	k2.44 b84	k2.45 b83	k2.46 b82	k2.47 b81	k2.48 b80
W	w	W	W	W	W	W	W	W	W	W	W	W	W	W	W
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
k2.49 b79	k2.50 b78	k2.51 b77	k2.52 b76	k2.53 b75	k2.54 b74	k2.55 b73	k2.56 b72	k2.57 b71	k2.58 b70	k2.59 b69	k2.60 b68	k2.61 b67	k2.62 b66	k2.63 b65	k2.64 b64
W	w	w	w	w	W	w	w	w	w	w	w	w	W	w	w

## CRYP\_K3LR (address offset: 0x38)

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
k3.1	k3.2	k3.3	k3.4	k3.5	k3.6	k3.7	k3.8	k3.9	k3.10	k3.11	k3.12	k3.13	k3.14	k3.15	k3.16
b63	b62	b61	b60	b59	b58	b57	b56	b55	b54	b53	b52	b51	b50	b49	b48
w	w	w	w	w	w	w	w	w	w	w	w	w	w	w	w
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
k3.17	k3.18	k3.19	k3.20	k3.21	k3.22	k3.23	k3.24	k3.25	k3.26	k3.27	k3.28	k3.29	k3.30	k3.31	k3.32
b47	b46	b45	b44	b43	b42	b41	b40	b39	b38	b37	b36	b35	b34	b33	b32
w	w	w	w	w	w	w	w	w	w	w	w	w	w	w	w

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### CRYP\_K3RR (address offset: 0x3C)

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
k3.33 b31	k3.34 b30	k3.35 b29	k3.36 b28	k3.37 b27	k3.38 b26	k3.39 b25	k3.40 b24	k3.41 b23	k3.42 b22	k3.43 b21	k3.44 b20	k3.45 b19	k3.46 b18	k3.47 b17	k3.48 b16
w	w	w	w	w	w	w	w	w	w	w	w	w	w	w	w
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
15 k3.49 b15	14 k3.50 b14	13 k3.51 b13	12 k3.52 b12	11 k3.53 b11	10 k3.54 b10	9 k3.55 b9	8 k3.56 b8	7 k3.57 b7	6 k3.58 b6	5 k3.59 b5	4 k3.60 b4	3 k3.61 b3	2 k3.62 b2	1 k3.63 b1	0 k3.64 b0

Note:

Write accesses to these registers are disregarded when the cryptographic processor is busy (bit BUSY = 1 in the CRYP\_SR register).

### 19.6.10 CRYP initialization vector registers (CRYP\_IV0...1(L/R)R)

Address offset: 0x40 to 0x4C Reset value: 0x0000 0000

The CRYP\_IV0...1(L/R)R are the left-word and right-word registers for the initialization vector (64 bits for DES/TDES and 128 bits for AES) and are used in the CBC (Cipher block chaining) and Counter (CTR) modes. After each computation round of the TDES or AES Core, the CRYP\_IV0...1(L/R)R registers are updated as described in Section: DES and TDES Cipher block chaining (DES/TDES-CBC) mode on page 498, Section: AES Cipher block chaining (AES-CBC) mode on page 502 and Section: AES counter mode (AES-CTR) mode on page 504.

IV0 is the leftmost bit whereas IV63 (DES, TDES) or IV127 (AES) are the rightmost bits of the initialization vector. IV1(L/R)R is used only in the AES.

#### CRYP\_IV0LR (address offset: 0x40)

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
IV0	IV1	IV2	IV3	IV4	IV5	IV6	IV7	IV8	IV9	IV10	IV11	IV12	IV13	IV14	IV15
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
15 IV16	14 IV17	13 IV18	12 IV19	11 IV20	10 IV21	9 IV22	8 IV23	7 IV24	6 IV25	5 IV26	4 IV27	3 IV28	2 IV29	1 IV30	0 IV31

### CRYP\_IV0RR (address offset: 0x44)

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
IV32	IV33	IV34	IV35	IV36	IV37	IV38	IV39	IV40	IV41	IV42	IV43	IV44	IV45	IV46	IV47
rw															
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
IV48	IV49	IV50	IV51	IV52	IV53	IV54	IV55	IV56	IV57	IV58	IV59	IV60	IV61	IV62	IV63
rw															

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## CRYP\_IV1LR (address offset: 0x48)

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
IV64	IV65	IV66	IV67	IV68	IV69	IV70	IV71	IV72	IV73	IV74	IV75	IV76	IV77	IV78	IV79
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
15 IV80	14 IV81	13 IV82	12 IV83	11 IV84	10 IV85	9 IV86	8 IV87	7 IV88	6 IV89	5 IV90	4 IV91	3 IV92	2 IV93	1 IV94	0 IV95

# CRYP\_IV1RR (address offset: 0x4C)

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
IV96	IV97	IV98	IV99	IV100	IV101	IV102	IV103	IV104	IV105	IV106	IV107	IV108	IV109	IV110	IV111
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
15 IV112	14 IV113	13 IV114	12 IV115	11 IV116	10 IV117	9 IV118	8 IV119	7 IV120	6 IV121	5 IV122	4 IV123	3 IV124	2 IV125	1 IV126	0 IV127

Note: In DES/3DES modes, only CRYP\_IV0(L/R) is used.

Write access to these registers are disregarded when the cryptographic processor is busy (bit BUSY = 1 in the CRYP\_SR register).

# 19.6.11 CRYP register map

Table 77. CRYP register map and reset values

04	Register name	Register size	
Offset	reset value	3 4 4 2 9 9 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	- 0
0x00 0x00	CRYP_CR Reset value	CRYPEN  CRYPEN  FFLUSH  KEYSIZE  DATATYPE  DATATYPE  ALGODIR	Res.
			ш 5
0x04	CRYP_SR		I FEM
	Reset value CRYP_DIN	DATAIN	1 1
0x08	Reset value		0 0
0x0C	CRYP_DOUT	DATAOUT	
	Reset value		0 0
0x10	CRYP_DMACR	Reserved	DOEN
	Reset value		0 0
0x14	CRYP_IMSCR	Reserved	0 OUTIM
	Reset value		0 0
0x18	CRYP_RISR		OUTRIS
	Reset value		0 1
0x1C	CRYP_MISR	Reserved	SIWNIO
	Reset value CRYP_K0LR	CRYP_K0LR	0 0
0x20	Reset value	_	0 0
0x24	CRYP_K0RR Reset value	CRYP_KORR	0 0
	CRYP_K3LR	 CRYP_K3LR	
0x38	Reset value		0 0
0x3C	CRYP_K3RR	CRYP_K3RR	
0,00	Reset value		0 0
0x40	CRYP_IV0LR Reset value	CRYP_IVOLR   0   0   0   0   0   0   0   0   0   0	0 0
	CRYP_IV0RR	CRYP_IVORR	0 0
0x44	Reset value		0 0
0x48	CRYP_IV1LR	CRYP_IV1LR	
	Reset value		0 0
0x4C	CRYP_IV1RR Reset value	CRYP_IV1RR	0 0

Refer to *Table 2 on page 50* for the register boundary addresses.

# 20 Random number generator (RNG)

This section applies to the whole STM32F4xx family, unless otherwise specified.

### 20.1 RNG introduction

The RNG processor is a random number generator, based on a continuous analog noise, that provides a random 32-bit value to the host when read.

The RNG passed the FIPS PUB 140-2 (2001 October 10) tests with a success ratio of 99%.

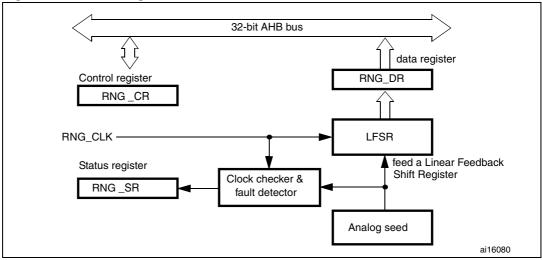
#### 20.2 RNG main features

- It delivers 32-bit random numbers, produced by an analog generator
- 40 periods of the PLL48CLK clock signal between two consecutive random numbers
- Monitoring of the RNG entropy to flag abnormal behavior (generation of stable values, or of a stable sequence of values)
- It can be disabled to reduce power-consumption

# 20.3 RNG functional description

Figure 211 shows the RNG block diagram.

Figure 211. Block diagram



The random number generator implements an analog circuit. This circuit generates seeds that feed a linear feedback shift register (RNG\_LFSR) in order to produce 32-bit random numbers.

The analog circuit is made of several ring oscillators whose outputs are XORed to generate the seeds. The RNG\_LFSR is clocked by a dedicated clock (PLL48CLK) at a constant frequency, so that the quality of the random number is independent of the HCLK frequency.

The contents of the RNG\_LFSR are transferred into the data register (RNG\_DR) when a significant number of seeds have been introduced into the RNG\_LFSR.

In parallel, the analog seed and the dedicated PLL48CLK clock are monitored. Status bits (in the RNG\_SR register) indicate when an abnormal sequence occurs on the seed or when the frequency of the PLL48CLK clock is too low. An interrupt can be generated when an error is detected.

### 20.3.1 Operation

To run the RNG, follow the steps below:

- 1. Enable the interrupt if needed (to do so, set the IE bit in the RNG\_CR register). An interrupt is generated when a random number is ready or when an error occurs.
- 2. Enable the random number generation by setting the RNGEN bit in the RNG\_CR register. This activates the analog part, the RNG\_LFSR and the error detector.
- 3. At each interrupt, check that no error occurred (the SEIS and CEIS bits should be '0' in the RNG\_SR register) and that a random number is ready (the DRDY bit is '1' in the RNG\_SR register). The contents of the RNG\_DR register can then be read.

As required by the FIPS PUB (Federal Information Processing Standard Publication) 140-2, the first random number generated after setting the RNGEN bit should not be used, but saved for comparison with the next generated random number. Each subsequent generated random number has to be compared with the previously generated number. The test fails if any two compared numbers are equal (continuous random number generator test).

## 20.3.2 Error management

#### If the CEIS bit is read as '1' (clock error)

In the case of a clock, the RNG is no more able to generate random numbers because the PLL48CLK clock is not correct. Check that the clock controller is correctly configured to provide the RNG clock and clear the CEIS bit. The RNG can work when the CECS bit is '0'. The clock error has no impact on the previously generated random numbers, and the RNG\_DR register contents can be used.

#### If the SEIS bit is read as '1' (seed error)

In the case of a seed error, the generation of random numbers is interrupted for as long as the SECS bit is '1'. If a number is available in the RNG\_DR register, it must not be used because it may not have enough entropy.

What you should do is clear the SEIS bit, then clear and set the RNGEN bit to reinitialize and restart the RNG.

# 20.4 RNG registers

The RNG is associated with a control register, a data register and a status register. They have to be accessed by words (32 bits).

## 20.4.1 RNG control register (RNG\_CR)

Address offset: 0x00

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	Reserved														
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Descript									IE	RNGEN	Boorwood			
Reserved									rw	rw	Reserved				

Bits 31:4 Reserved, must be kept at reset value

Bit 3 IE: Interrupt enable

0: RNG Interrupt is disabled

1: RNG Interrupt is enabled. An interrupt is pending as soon as DRDY=1 or SEIS=1 or

CEIS=1 in the RNG\_SR register.

Bit 2 RNGEN: Random number generator enable

0: Random number generator is disabled1: random Number Generator is enabled.

Bits 1:0 Reserved, must be kept at reset value

## 20.4.2 RNG status register (RNG\_SR)

Address offset: 0x04

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
Reserved															
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Description								SEIS	CEIS	Decembed		SECS	CECS	DRDY
	Reserved									rc_w0	Reserved		r	r	r

Bits 31:3 Reserved, must be kept at reset value

Bit 6 SEIS: Seed error interrupt status

This bit is set at the same time as SECS, it is cleared by writing it to 0.

0: No faulty sequence detected

1: One of the following faulty sequences has been detected:

- More than 64 consecutive bits at the same value (0 or 1)
- More than 32 consecutive alternances of 0 and 1 (0101010101...01)

An interrupt is pending if IE = 1 in the RNG\_CR register.

Bit 5 CEIS: Clock error interrupt status

This bit is set at the same time as CECS, it is cleared by writing it to 0.

0: The PLL48CLK clock was correctly detected

1: The PLL48CLK was not correctly detected (f $_{\rm PLL48CLK}$ < f $_{\rm HCLK}$ /16)

An interrupt is pending if IE = 1 in the RNG\_CR register.

Bits 4:3 Reserved, must be kept at reset value

#### Bit 2 SECS: Seed error current status

0: No faulty sequence has currently been detected. If the SEIS bit is set, this means that a faulty sequence was detected and the situation has been recovered.

- 1: One of the following faulty sequences has been detected:
  - More than 64 consecutive bits at the same value (0 or 1)
  - More than 32 consecutive alternances of 0 and 1 (0101010101...01)

#### Bit 1 CECS: Clock error current status

0: The PLL48CLK clock has been correctly detected. If the CEIS bit is set, this means that a clock error was detected and the situation has been recovered

1: The PLL48CLK was not correctly detected ( $f_{PLL48CLK} < f_{HCLK}/16$ ).

#### Bit 0 DRDY: Data ready

0: The RNG\_DR register is not yet valid, no random data is available

1: The RNG\_DR register contains valid random data

Note: An interrupt is pending if IE = 1 in the RNG\_CR register.

Once the RNG\_DR register has been read, this bit returns to 0 until a new valid value is computed.

## 20.4.3 RNG data register (RNG\_DR)

Address offset: 0x08

Reset value: 0x0000 0000

The RNG\_DR register is a read-only register that delivers a 32-bit random value when read. After being read, this register delivers a new random value after a maximum time of 40 periods of the PLL48CLK clock. The software must check that the DRDY bit is set before reading the RNDATA value.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
RNDATA															
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
RNDATA															
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:0 **RNDATA:** Random data 32-bit random data.

# 20.4.4 RNG register map

Table 78 gives the RNG register map and reset values.

Table 78. RNG register map and reset map

Offset	Register name	Register size													
011001	reset value	1 2 3 4 5 6 6 7 8 8 6 7 9 8 7 9 8 7 9 8 7 9 8 7 9 8 7 9 8 7 9 9 7 9 9 7 9 9 7 9 9 7 9 9 7 9 9 7 9 9 7 9 9 9 7 9 9 9 7 9 9 9 7 9 9 9 7 9	0												
0x00	RNG_CR 0x0000000	Reserved H Reserved													
0x04	RNG_SR 0x0000000	Reserved SEIS CEIS SECS SECS SECS SECS SECS SECS	YOHO												
0x08	RNG_DR 0x0000000	RNDATA[31:0]													

# 21 Hash processor (HASH)

This section applies to the whole STM32F4xx family, unless otherwise specified.

### 21.1 HASH introduction

The hash processor is a fully compliant implementation of the secure hash algorithm (SHA-1), the MD5 (message-digest algorithm 5) hash algorithm and the HMAC (keyed-hash message authentication code) algorithm suitable for a variety of applications. It computes a message digest (160 bits for the SHA-1 algorithm, 128 bits for the MD5 algorithm) for messages of up to  $(2^{64}-1)$  bits, while HMAC algorithms provide a way of authenticating messages by means of hash functions. HMAC algorithms consist in calling the SHA-1, or MD5 hash function twice.

#### 21.2 HASH main features

- Suitable for data authentication applications, compliant with:
  - FIPS PUB 180-2 (Federal Information Processing Standards Publication 180-2)
  - Secure Hash Standard specifications (SHA-1)
  - IETF RFC 1321 (Internet Engineering Task Force Request For Comments number 1321) specifications (MD5)
- Fast computation of SHA-1, and MD5
- AHB slave peripheral
- 32-bit data words for input data, supporting word, half-word, byte and bit bit-string representations, with little-endian data representation only.
- Automatic swapping to comply with the big-endian SHA1 computation standard with little-endian input bit-string representation
- $\bullet$  Automatic padding to complete the input bit string to fit modulo 512 (16  $\times$  32 bits) message digest computing
- 5x 32-bit words (H0 to H5) for output message digest, reload able to continue interrupted message digest computation.
- Corresponding 32-bit words of the digest from consecutive message blocks are added to each other to form the digest of the whole message
- Automatic data flow control with support for direct memory access (DMA)

Note:

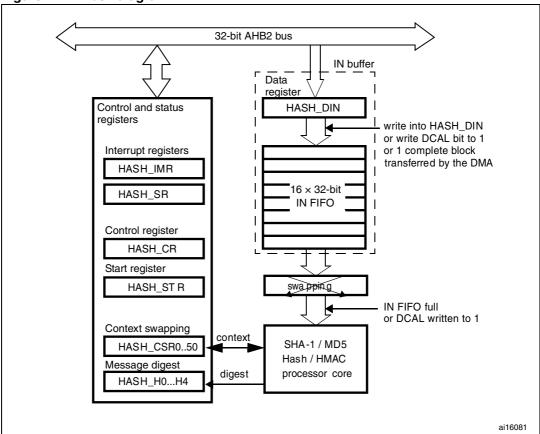
Padding, as defined in the SHA-1 algorithm, consists in adding a bit at bx1 followed by N bits at bx0 to get a total length congruent to 448 modulo 512. After this, the message is completed with a 64-bit integer which is the binary representation of the original message length.

For this hash processor, the quanta for entering the message is a 32-bit word, so an additional information must be provided at the end of the message entry, which is the number of valid bits in the last 32-bit word entered.

# 21.3 HASH functional description

Figure 1 shows the block diagram of the hash processor.

Figure 212. Block diagram



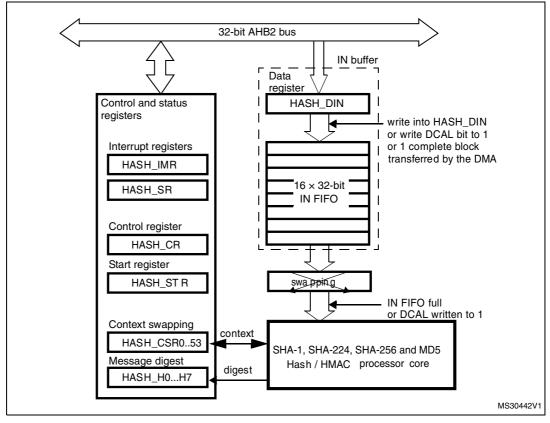


Figure 213. Block diagram for

The FIPS PUB 180-2 standard and the IETF RFC 1321 publication specify the SHA-1, and MD5 secure hash algorithms, respectively, for computing a condensed representation of a message or data file. When a message of any length below 2<sup>64</sup> bits is provided on input, the SHA-1, and MD5 produce respective a 160-bit, and 128-bit output string, respectively, called a message digest. The message digest can then be processed with a digital signature algorithm in order to generate or verify the signature for the message. Signing the message digest rather than the message often improves the efficiency of the process because the message digest is usually much smaller in size than the message. The verifier of a digital signature has to use the same hash algorithm as the one used by the creator of the digital signature.

The SHA-1, and MD5 are qualified as "secure" because it is computationally infeasible to find a message that corresponds to a given message digest, or to find two different messages that produce the same message digest. Any change to a message in transit will, with very high probability, result in a different message digest, and the signature will fail to verify. For more detail on the SHA-1 algorithm, please refer to the FIPS PUB 180-2 (Federal Information Processing Standards Publication 180-2), 2002 august 1.

The current implementation of this standard works with little-endian input data convention. For example, the C string "abc" must be represented in memory as the 24-bit hexadecimal value 0x434241.

A message or data file to be processed by the hash processor should be considered a bit string. The length of the message is the number of bits in the message (the empty message has length 0). You can consider that 32 bits of this bit string forms a 32-bit word. Note that the FIPS PUB 180-1 standard uses the convention that bit strings grow from left to right, and

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bits can be grouped as bytes (8 bits) or words (32 bits) (but some implementations also use half-words (16 bits), and implicitly, uses the big-endian byte (half-word) ordering. This convention is mainly important for padding (see *Section 1.3.4: Message padding on page 12*).

### 21.3.1 Duration of the processing

The computation of an intermediate block of a message takes:

- 66 HCLK clock cycles in SHA-1
- 50 HCLK clock cycles in MD5

to which you must add the time needed to load the 16 words of the block into the processor (at least 16 clock cycles for a 512-bit block).

The time needed to process the last block of a message (or of a key in HMAC) can be longer. This time depends on the length of the last block and the size of the key (in HMAC mode). Compared to the processing of an intermediate block, it can be increased by a factor of:

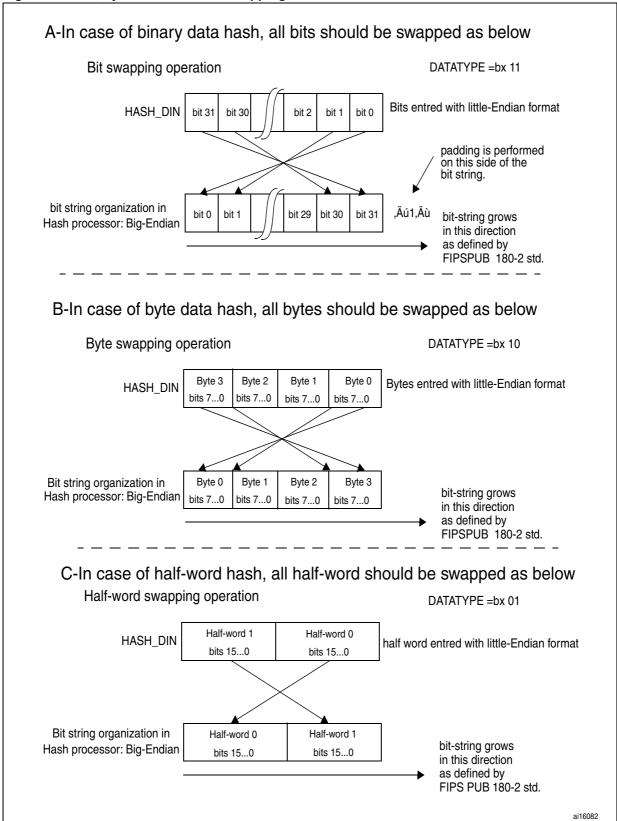
- 1 to 2.5 for a hash message
- around 2.5 for an HMAC input-key
- 1 to 2.5 for an HMAC message
- around 2.5 for an HMAC output key in case of a short key
- 3.5 to 5 for an HMAC output key in case of a long key

### 21.3.2 Data type

Data are entered into the hash processor 32 bits (word) at a time, by writing them into the HASH\_DIN register. But the original bit-string can be organized in bytes, half-words or words, or even be represented as bits. As the system memory organization is little-endian and SHA1 computation is big-endian, depending on the way the original bit string is grouped, a bit, byte, or half-word swapping operation is performed automatically by the hash processor.

The kind of data to be processed is configured with the DATATYPE bitfield in the HASH control register (HASH\_CR).

Figure 214. Bit, byte and half-word swapping



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The least significant bit of the message has to be at position 0 (right) in the first word entered into the hash processor, the 32nd bit of the bit string has to be at position 0 in the second word entered into the hash processor and so on.

## 21.3.3 Message digest computing

The HASH sequentially processes blocks of 512 bits when computing the message digest. Thus, each time  $16 \times 32$ -bit words (= 512 bits) have been written by the DMA or the CPU, into the hash processor, the HASH automatically starts computing the message digest. This operation is known as a partial digest computation.

The message to be processed is entered into the peripheral by 32-bit words written into the HASH\_DIN register. The current contents of the HASH\_DIN register are transferred to the input FIFO (IN FIFO) each time the register is written with new data. HASH\_DIN and the input FIFO form a FIFO of a 17-word length (named the IN buffer).

The processing of a block can start only once the last value of the block has entered the IN FIFO. The peripheral must get the information as to whether the HASH\_DIN register contains the last bits of the message or not. Two cases may occur:

- When the DMA is not used:
  - In case of a partial digest computation, this is done by writing an additional word into the HASH\_DIN register (actually the first word of the next block). Then the software must wait until the processor is ready again (when DINIS=1) before writing new data into HASH\_DIN.
  - In case of a final digest computation (last block entered), this is done by writing the DCAL bit to 1.
- When the DMA is used:

The contents of the HASH\_DIN register are interpreted automatically with the information sent by the DMA controller.

 The contents of the HASH\_DIN register are interpreted automatically with the information sent by the DMA controller.

This process —data entering + partial digest computation— continues until the last bits of the original message are written to the HASH\_DIN register. As the length (number of bits) of a message can be any integer value, the last word written into the HASH processor may have a valid number of bits between 1 and 32. This number of valid bits in the last word, NBLW, has to be written into the HASH\_STR register, so that message padding is correctly performed before the final message digest computation.

Once this is done, writing into HASH\_STR with bit DCAL = 1 starts the processing of the last entered block of message by the hash processor. This processing consists in:

- Automatically performing the message padding operation: the purpose of this operation
  is to make the total length of a padded message a multiple of 512. The HASH
  sequentially processes blocks of 512 bits when computing the message digest
- Computing the final message digest

When the DMA is enabled, it provides the information to the hash processor when it is transferring the last data word. Then the padding and digest computation are performed automatically as if DCAL had been written to 1.

## 21.3.4 Message padding

Message padding consists in appending a "1" followed by m "0"s followed by a 64-bit integer to the end of the original message to produce a padded message block of length 512. The "1" is added to the last word written into the HASH\_DIN register at the bit position defined by the NBLW bitfield, and the remaining upper bits are cleared ("0"s).

Example: let us assume that the original message is the ASCII binary-coded form of "abc", of length L = 24:

```
byte 0 byte 1 byte 2 byte 3 01100001 01100010 01100011 UUUUUUUUU <-- 1st word written to HASH DIN -->
```

NBLW has to be loaded with the value 24: a "1" is appended at bit location 24 in the bit string (starting counting from left to right in the above bit string), which corresponds to bit 31 in the HASH\_DIN register (little-endian convention):

```
01100001 01100010 01100011 1UUUUUUU
```

Since L = 24, the number of bits in the above bit string is 25, and 423 "0"s are appended, making now 448. This gives (in hexadecimal, big-endian format):

The L value, in two-word representation (that is 00000000 0000018) is appended. Hence the final padded message in hexadecimal:

If the HASH is programmed to use the little-endian byte input format, the above message has to be entered by doing the following steps:

- 1. 0xUU636261 is written into the HASH\_DIN register (where 'U' means don't care)
- 2. 0x18 is written into the HASH\_STR register (the number of valid bits in the last word written into the HASH\_DIN register is 24, as the original message length is 24 bits)
- 3. 0x10 is written into the HASH\_STR register to start the message padding and digest computation. When NBLW ≠ 0x00, the message padding puts a "1" into the HASH\_DIN register at the bit position defined by the NBLW value, and inserts "0"s at bit locations [31:(NBLW+1)]. When NBLW == 0x00, the message padding inserts one new word with value 0x0000 0001. Then an all zero word (0x0000 0000) is added and the message length in a two-word representation, to get a block of 16 x 32-bit words.
- 4. The HASH computing is performed, and the message digest is then available in the HASH\_Hx registers (x = 0...4) for the SHA-1 algorithm. For example:

```
H0 = 0xA9993E36

H1 = 0x4706816A

H2 = 0xBA3E2571

H3 = 0x7850C26C

H4 = 0x9CD0D89D
```

### 21.3.5 Hash operation

The hash function (SHA-1, and MD5) is selected when the INIT bit is written to '1' in the HASH\_CR register while the MODE bit is at '0' in HASH\_CR. The algorithm (SHA-1, or MD5) is selected at the same time (that is when the INIT bit is set) using the ALGO bits.

The message can then be sent by writing it word by word into the HASH\_DIN register. When a block of 512 bits —that is 16 words— has been written, a partial digest computation starts upon writing the first data of the next block. The hash processor remains busy for 66 cycles for the SHA-1 algorithm, or 50 cycles for the MD5 algorithm.

The process can then be repeated until the last word of the message. If DMA transfers are used, refer to the *Procedure where the data are loaded by DMA* section. Otherwise, if the message length is not an exact multiple of 512 bits, then the HASH\_STR register has to be written to launch the computation of the final digest.

Once computed, the digest can be read from the HASH\_H0...HASH\_H4 registers (for the MD5 algorithm, HASH\_H4 is not relevant).

## 21.3.6 HMAC operation

The HMAC algorithm is used for message authentication, by irreversibly binding the message being processed to a key chosen by the user. For HMAC specifications, refer to "HMAC: keyed-hashing for message authentication, H. Krawczyk, M. Bellare, R. Canetti, February 1997.

Basically, the algorithm consists of two nested hash operations:

#### where:

- pad is a sequence of zeroes needed to extend the key to the length of the underlying hash function data block (that is 512 bits for both the SHA-1, and MD5 hash algorithms)
- represents the concatenation operator

To compute the HMAC, four different phases are required:

- The block is initialized by writing the INIT bit to '1' with the MODE bit at '1' and the ALGO bits set to the value corresponding to the desired algorithm. The LKEY bit must also be set during this phase if the key being used is longer than 64 bytes (in this case, the HMAC specifications specify that the hash of the key should be used in place of the real key).
- 2. The key (to be used for the inner hash function) is then given to the core. This operation follows the same mechanism as the one used to send the message in the hash operation (that is, by writing into HASH\_DIN and, finally, into HASH\_STR).
- Once the last word has been entered and computation has started, the hash processor elaborates the key. It is then ready to accept the message text using the same mechanism as the one used to send the message in the hash operation.
- 4. After the first hash round, the hash processor returns "ready" to indicate that it is ready to receive the key to be used for the outer hash function (normally, this key is the same as the one used for the inner hash function). When the last word of the key is entered and computation starts, the HMAC result is made available in the HASH\_H0...HASH\_H4 registers.

Note: 1 The computation latency of the HMAC primitive depends on the lengths of the keys and message. You could the HMAC as two nested underlying hash functions with the same key length (long or short).

### 21.3.7 Context swapping

It is possible to interrupt a hash/HMAC process to perform another processing with a higher priority, and to complete the interrupted process later on, when the higher-priority task is complete. To do so, the context of the interrupted task must be saved from the hash registers to memory, and then be restored from memory to the hash registers.

The procedures where the data flow is controlled by software or by DMA are described below.

#### Procedure where the data are loaded by software

The context can be saved only when no block is currently being processed. That is, you must wait for DINIS = 1 (the last block has been processed and the input FIFO is empty) or NBW  $\neq$  0 (the FIFO is not full and no processing is ongoing).

Context saving:

Store the contents of the following registers into memory:

- HASH\_IMR
- HASH STR
- HASH CR
- HASH\_CSR0 to HASH\_CSR50.
- Context restoring:

The context can be restored when the high-priority task is complete. Please follow the order of the sequence below.

- Write the following registers with the values saved in memory: HASH\_IMR, HASH\_STR and HASH\_CR
- b) Initialize the hash processor by setting the INIT bit in the HASH\_CR register
- Write the HASH\_CSR0 to HASH\_CSR50 registers with the values saved in memory

You can now restart the processing from the point where it has been interrupted.

### Procedure where the data are loaded by DMA

In this case it is not possible to predict if a DMA transfer is in progress or if the process is ongoing. Thus, you must stop the DMA transfers, then wait until the HASH is ready in order to interrupt the processing of a message.

- Interrupting a processing:
  - Clear the DMAE bit to disable the DMA interface
  - Wait until the current DMA transfer is complete (wait for DMAES = 0 in the HASH\_SR register). Note that the block may or not have been totally transferred to the HASH.
  - Disable the corresponding channel in the DMA controller
  - Wait until the hash processor is ready (no block is being processed), that is wait for DINIS = 1
- The context saving and context restoring phases are the same as above (see *Procedure where the data are loaded by software*).

Reconfigure the DMA controller so that it transfers the end of the message. You can now restart the processing from the point where it was interrupted by setting the DMAE bit.

Note: 1 If context swapping does not involve HMAC operations, the HASH\_CSR38 to HASH\_CSR50 registers do not have to be saved and restored.

2 If context swapping occurs between two blocks (the last block was completely processed and the next block has not yet been pushed into the IN FIFO, NBW = 000 in the HASH\_CR register), the HASH\_CSR22 to HASH\_CSR37 registers do not have to be saved and restored.

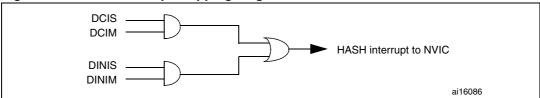
# 21.3.8 HASH interrupt

There are two individual maskable interrupt sources generated by the HASH processor. They are connected to the same interrupt vector.

You can enable or disable the interrupt sources individually by changing the mask bits in the HASH\_IMR register. Setting the appropriate mask bit to 1 enables the interrupt.

The status of the individual interrupt sources can be read from the HASH SR register.

Figure 215. HASH interrupt mapping diagram



# 21.4 HASH registers

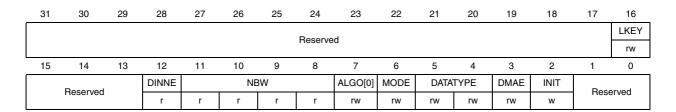
The HASH core is associated with several control and status registers and five message digest registers.

All these registers are accessible through word accesses only, else an AHB error is generated.

### 21.4.1 HASH control register (HASH\_CR)

Address offset: 0x00

Reset value: 0x0000 0000



Bits 31:17 Reserved, forced by hardware to 0.

#### Bit 16 LKEY: Long key selection

This bit selects between short key ( $\leq$  64 bytes) or long key (> 64 bytes) in HMAC mode

0: Short key (≤ 64 bytes)

1: Long key (> 64 bytes)

Note: This selection is only taken into account when the INIT bit is set and MODE = 1. Changing this bit during a computation has no effect.

Bits 15:13 Reserved, forced by hardware to 0.

#### Bit 12 DINNE: DIN not empty

This bit is set when the HASH\_DIN register holds valid data (that is after being written at least once). It is cleared when either the INIT bit (initialization) or the DCAL bit (completion of the previous message processing) is written to 1.

0: No data are present in the data input buffer

1: The input buffer contains at least one word of data

### Bits 11:8 NBW: Number of words already pushed

This bitfield reflects the number of words in the message that have already been pushed into the IN FIFO.

NBW increments (+1) when a write access is performed to the HASH\_DIN register while DINNE = 1.

It goes to 0000 when the INIT bit is written to 1 or when a digest calculation starts (DCAL written to 1 or DMA end of transfer).

#### If the DMA is not used:

0000 and DINNE=0: no word has been pushed into the DIN buffer (the buffer is empty, both the HASH\_DIN register and the IN FIFO are empty)

0000 and DINNE=1: 1 word has been pushed into the DIN buffer (The HASH\_DIN register contains 1 word, the IN FIFO is empty)

0001: 2 words have been pushed into the DIN buffer (the HASH\_DIN register and the IN FIFO contain 1 word each)

•••

1111: 16 words have been pushed into the DIN buffer

 If the DMA is used, NBW is the exact number of words that have been pushed into the IN FIFO.

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#### Bit 7 ALGO[1:0]: Algorithm selection

These bits selects the SHA-1 or the MD5 algorithm:

0: SHA-1 algorithm selected

1: MD5 algorithm selected

Note: This selection is only taken into account when the INIT bit is set. Changing this bit during a computation has no effect.

#### Bit 6 MODE: Mode selection

This bit selects the HASH or HMAC mode for the selected algorithm:

0: Hash mode selected

1: HMAC mode selected. LKEY must be set if the key being used is longer than 64 bytes.

Note: This selection is only taken into account when the INIT bit is set. Changing this bit during a computation has no effect.

#### Bits 5:4 **DATATYPE:** Data type selection

Defines the format of the data entered into the HASH\_DIN register:

00: 32-bit data. The data written into HASH\_DIN are directly used by the HASH processing, without reordering.

01: 16-bit data, or half-word. The data written into HASH\_DIN are considered as 2 half-words, and are swapped before being used by the HASH processing.

10: 8-bit data, or bytes. The data written into HASH\_DIN are considered as 4 bytes, and are swapped before being used by the HASH processing.

11: bit data, or bit-string. The data written into HASH\_DIN are considered as 32 bits (1st bit of the sting at position 0), and are swapped before being used by the HASH processing (1st bit of the string at position 31).

#### Bit 3 DMAE: DMA enable

0: DMA transfers disabled

1: DMA transfers enabled. A DMA request is sent as soon as the HASH core is ready to receive data.

Note: 1: This bit is cleared by hardware when the DMA asserts the DMA terminal count signal (while transferring the last data of the message). This bit is not cleared when the INIT bit is written to 1.

2: If this bit is written to 0 while a DMA transfer has already been requested to the DMA, DMAE is cleared but the current transfer is not aborted. Instead, the DMA interface remains internally enabled until the transfer is complete or INIT is written to 1.

#### Bit 2 INIT: Initialize message digest calculation

Writing this bit to 1 resets the hash processor core, so that the HASH is ready to compute the message digest of a new message.

Writing this bit to 0 has no effect.

Reading this bit always return 0.

Bit 1:0 Reserved, must be kept cleared.

### 21.4.2 HASH data input register (HASH\_DIN)

Address offset: 0x04

Reset value: 0x0000 0000

HASH\_DIN is the data input register. It is 32-bit wide. It is used to enter the message by blocks of 512 bits. When the HASH\_DIN register is written to, the value presented on the AHB databus is 'pushed' into the HASH core and the register takes the new value presented on the AHB databus. The DATATYPE bits must previously have been configured in the HASH\_CR register to get a correct message representation.

When a block of 16 words has been written to the HASH\_DIN register, an intermediate digest calculation is launched:

- by writing new data into the HASH\_DIN register (the first word of the next block) if the DMA is not used (intermediate digest calculation)
- automatically if the DMA is used

When the last block has been written to the HASH\_DIN register, the final digest calculation (including padding) is launched:

- by writing the DCAL bit to 1 in the HASH\_STR register (final digest calculation)
- automatically if the DMA is used.

When a digest calculation (intermediate or final) is in progress, any new write access to the HASH\_DIN register is extended (by wait-state insertion on the AHB bus) until the HASH calculation completes.

When the HASH\_DIN register is read, the last word written in this location is accessed (zero after reset).

rw	DAT rw	AIN rw	rw												
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
rw	rw	rw	rw	rw	rw	rw	rw	rw							
							DAT	AIN							
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16

Bit 31:0 DATAIN: Data input

Read = returns the current register content.

Write = the current register content is pushed into the IN FIFO, and the register takes the new value presented on the AHB databus.

#### 21.4.3 **HASH start register (HASH\_STR)**

Address offset: 0x08

Reset value: 0x0000 0000

The HASH STR register has two functions:

- It is used to define the number of valid bits in the last word of the message entered in the hash processor (that is the number of valid least significant bits in the last data written into the HASH\_DIN register)
- It is used to start the processing of the last block in the message by writing the DCAL bit to 1

3	1	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
								Resei	rved							
1	5	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			De	eserved				DCAL		Reserved				NBLW		
			ne	eserveu				w		neserveu		rw	rw	rw	rw	rw

Bits 31:9 Reserved, forced by hardware to 0.

#### Bit 8 DCAL: Digest calculation

Writing this bit to 1 starts the message padding, using the previously written value of NBLW, and starts the calculation of the final message digest with all data words written to the IN FIFO since the INIT bit was last written to 1.

Reading this bit returns 0.

Note

Bits 7:5 Reserved, forced by hardware to 0.

Bits 4:0 NBLW: Number of valid bits in the last word of the message in the bit string organization of hash processor

> When these bits are written and DCAL is at '0', they take the value on the AHB databus:

0x00: All 32 bits of the last data written in the bit string organization of hash processor (after data swapping) are valid.

0x01: Only bit [31] of the last data written in the bit string organization of hash processor (after data swapping) are valid

0x02: Only bits [31:30] of the last data written in the bit string organization of hash processor (after data swapping) are valid

0x03: Only bits [31:29] of the last data written in the bit string organization of hash processor (after data swapping) are valid

0x1F: Only bits [0] of the last data written in the bit string organization of hash processor (after data swapping) are valid

When these bits are written and DCAL is at '1', the bitfield is not changed. Reading them returns the last value written to NBLW.

Note: These bits must be configured before setting the DCAL bit, else they are not

taken into account. Especially, it is not possible to configure NBLW and set DCAL at the same time.

### 21.4.4 HASH digest registers (HASH\_HR0..4)

Address offset: 0x0C to 0x1C Reset value: 0x0000 0000

These registers contain the message digest result named as:

- 1. H0, H1, H2, H3 and H4, respectively, in the SHA1 algorithm description
- 2. A, B, C and D, respectively, in the MD5 algorithm description Note that in this case, the HASH\_H4 is not used, and is read as zero.
- 3. H0 to H6, respectively, in the SHA224 algorithm description,
- 4. H0 to H7, respectively, in the SHA256 algorithm description,

If a read access to one of these registers occurs while the HASH core is calculating an intermediate digest or a final message digest (that is when the DCAL bit has been written to 1), then the read is stalled until the completion of the HASH calculation.

Note:

### HASH\_HR0

Address offset: 0x0C

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Н	0							
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							Н	0							
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

### HASH HR1

Address offset: 0x10

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
				_	_		Н	1			_	_		_	
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			_		_		Н	1	_	_		_			_
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

### HASH\_HR2

Address offset: 0x14

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Н	12							
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							Н	12							
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

### HASH\_HR3

Address offset: 0x18

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Н	13							
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							Н	13							
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

### HASH\_HR4

Address offset: 0x1C

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Н	14							
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							Н	14							
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Note:

When starting a digest computation for a new bit stream (by writing the INIT bit to 1), these registers assume their reset values.

# 21.4.5 HASH interrupt enable register (HASH\_IMR)

Address offset: 0x20

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
						Dass	- w. co.d							DCIE	DINIE
						Hese	erved							rw	rw

Bits 31:2 Reserved, forced by hardware to 0.

Bit 1 DCIE: Digest calculation completion interrupt enable

0: Digest calculation completion interrupt disabled

1: Digest calculation completion interrupt enabled.

Bit 0 DINIE: Data input interrupt enable

0: Data input interrupt disabled

1: Data input interrupt enabled

### 21.4.6 HASH status register (HASH\_SR)

Address offset: 0x24

Reset value: 0x0000 0001

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
					Rese	vr.cd						BUSY	DMAS	DCIS	DINIS
					Hese	ervea						r	r	rc_w0	rc_w0

Bits 31:4 Reserved, forced by hardware to 0.

#### Bit 3 BUSY: Busy bit

0: No block is currently being processed

1: The hash core is processing a block of data

#### Bit 2 DMAS: DMA Status

This bit provides information on the DMA interface activity. It is set with DMAE and cleared when DMAE=0 and no DMA transfer is ongoing. No interrupt is associated with this bit.

- 0: DMA interface is disabled (DMAE=0) and no transfer is ongoing
- 1: DMA interface is enabled (DMAE=1) or a transfer is ongoing

#### Bit 1 DCIS: Digest calculation completion interrupt status

This bit is set by hardware when a digest becomes ready (the whole message has been processed). It is cleared by writing it to 0 or by writing the INIT bit to 1 in the HASH\_CR register.

- 0: No digest available in the HASH\_Hx registers
- 1: Digest calculation complete, a digest is available in the HASH\_Hx registers. An interrupt is generated if the DCIE bit is set in the HASH\_IMR register.

#### Bit 0 DINIS: Data input interrupt status

This bit is set by hardware when the input buffer is ready to get a new block (16 locations are free). It is cleared by writing it to 0 or by writing the HASH\_DIN register.

- 0: Less than 16 locations are free in the input buffer
- 1: A new block can be entered into the input buffer. An interrupt is generated if the DINIE bit is set in the HASH\_IMR register.

### 21.4.7 HASH context swap registers (HASH\_CSR0...53)

Address offset: 0x0F8 to 0x1CC

a) For HASH\_CSR0 register: Reset value: 0x0000 0002.

b) For others registers: Reset value: 0x0000 0000.

These registers contain the complete internal register states of the hash processor, and are useful when a context swap has to be done because a high-priority task has to use the hash processor while it is already in use by another task.

When such an event occurs, the HASH\_CSRx registers have to be read and the read values have to be saved somewhere in the system memory space. Then the hash processor can be used by the preemptive task, and when hash computation is finished, the saved context can be read from memory and written back into these HASH\_CSRx registers.

### HASH\_CSRx

Address offset: 0x0F8 to 0x1C0

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							C	Sx							
rw															
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							C	Sx							
rw															

# 21.4.8 HASH register map

Table 9 gives the summary HASH register map and reset values.

Table 79. HASH register map and reset values

Offset	Register name	Ť													R	eg	ist	er s	siz	е													٦
Oliset	reset value	31	30	53	28	27	56	22	24	23	22	21	20	19	18	17	16	15	14	13	12	7	10	6	8	7	9	2	4	က	2	_	0
0x00	HASH_CR		7.0					Res						•			LKEY		Reserved .		DINNE			ВW		ALGO[0]	MODE	DATATYPE		DMAE	INIT	Reserved	
	Reset value																0		ш_		0	0	0	0	0	0	0	0	0	0	0		
0x04	HASH_DIN																DAT	AIN															
0.01	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0
0x08	HASH_STR											Re	serv	ed											DCAL		Reserved			N	IBLV	V	
	Reset value																								0		Be		0	0	0	0 (	J
0x0C	HASH_HR0																Н																
	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0
0x10	HASH_HR1					-					-						Н																
	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	ე
0x14	HASH_HR2	١.,						-									Н																
	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	ე
0x18	HASH_HR3					-					-						Н																
	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0
0x1C	HASH_HR4					-					-						Н																
	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		0
0x20	HASH_IMR																rve								ı								DINIE
	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		0
0x24	HASH_SR													R	ese	rve	d													BUSY	DMAS	DCIS	_
	Reset value																													0	0	0	1
0xF8	HASH_CSR0																CS	R0															
	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1 (	0
0x1C0	HASH_CSR50																CSF	350															
ox : o	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0
												Re	eser	ved																			
0x310	HASH_HR0																Н	0															
0.310	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0
0x314	HASH_HR1							•							•		Н	1	•						•								
03314	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0
0v219	HASH_HR2						-						!	-			Н	2	•														T
0x318	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0v210	HASH_HR3						!	1						1			Н	3							•		•	•					ヿ
0x31C	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0
0,000	HASH_HR4					1		1			1				1		Н	4															┪
0x320	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0
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# 22 Real-time clock (RTC)

This section applies to the whole STM32F4xx family, unless otherwise specified.

### 22.1 Introduction

The real-time clock (RTC) is an independent BCD timer/counter. The RTC provides a time-of-day clock/calendar, two programmable alarm interrupts, and a periodic programmable wakeup flag with interrupt capability. The RTC also includes an automatic wakeup unit to manage low power modes.

Two 32-bit registers contain the seconds, minutes, hours (12- or 24-hour format), day (day of week), date (day of month), month, and year, expressed in binary coded decimal format (BCD). The sub-seconds value is also available in binary format.

Compensations for 28-, 29- (leap year), 30-, and 31-day months are performed automatically. Daylight saving time compensation can also be performed.

Additional 32-bit registers contain the programmable alarm subseconds, seconds, minutes, hours, day, and date.

A digital calibration feature is available to compensate for any deviation in crystal oscillator accuracy.

After power-on reset, all RTC registers are protected against possible parasitic write accesses.

As long as the supply voltage remains in the operating range, the RTC never stops, regardless of the device status (Run mode, low power mode or under reset).

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### 22.2 RTC main features

The RTC unit main features are the following (see Figure 216: RTC block diagram):

- Calendar with subseconds, seconds, minutes, hours (12 or 24 format), day (day of week), date (day of month), month, and year.
- Daylight saving compensation programmable by software.
- Two programmable alarms with interrupt function. The alarms can be triggered by any combination of the calendar fields.
- Automatic wakeup unit generating a periodic flag that triggers an automatic wakeup interrupt.
- Reference clock detection: a more precise second source clock (50 or 60 Hz) can be used to enhance the calendar precision.
- Accurate synchronization with an external clock using the subsecond shift feature.
- Maskable interrupts/events:
  - Alarm A
  - Alarm B
  - Wakeup interrupt
  - Timestamp
  - Tamper detection
- Digital calibration circuit (periodic counter correction)
  - 5 ppm accuracy
  - 0.95 ppm accuracy, obtained in a calibration window of several seconds
- Timestamp function for event saving (1 event)
- Tamper detection:
  - 2 tamper events with configurable filter and internal pull-up.
- 20 backup registers (80 bytes). The backup registers are reset when a tamper detection event occurs.
- RTC alternate function outputs (RTC\_AFO):
  - AFO\_CALIB: 512 Hz or 1Hz clock output (with an LSE frequency of 32.768 kHz).
     It is routed to the device RTC\_AF1 pin.
  - AFO\_ALARM: Alarm A or Alarm B or wakeup (only one can be selected). It is routed to the device RTC\_AF1 pin.
- RTC alternate function inputs (RTC AFI):
  - AFI\_TAMPER1: tamper event detection. It is routed to the device RTC\_AF1 and RTC\_AF2 pins.
  - AFI\_TAMPER2: tamper2 event detection. It is routed to the device RTC\_TAMPER2 pin.
  - AFI\_TIMESTAMP: timestamp event detection. It is routed to the device RTC\_AF1 and RTC\_AF2 pins.

Note: Refer to Section 6.4.1: GPIO port mode register (GPIOx\_MODER) (x = A..I) for more details on how to select RTC alternate functions (RTC\_AF1 and RTC\_AF2).

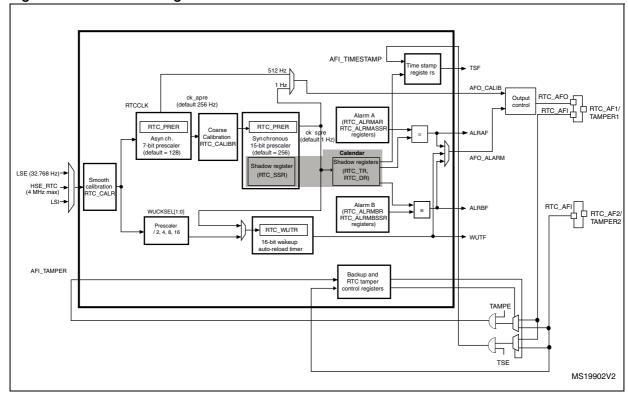


Figure 216. RTC block diagram

1. On STM32F4xx devices, the RTC\_AF1 and RTC\_AF2 alternate functions are connected to PC13 and PI8, respectively.

# 22.3 RTC functional description

### 22.3.1 Clock and prescalers

The RTC clock source (RTCCLK) is selected through the clock controller among the LSE clock, the LSI oscillator clock, and the HSE clock. For more information on the RTC clock source configuration, refer to *Section 5: Reset and clock control for (RCC)*.

A programmable prescaler stage generates a 1 Hz clock which is used to update the calendar. To minimize power consumption, the prescaler is split into 2 programmable prescalers (see *Figure 216: RTC block diagram*):

- A 7-bit asynchronous prescaler configured through the PREDIV\_A bits of the RTC\_PRER register.
- A 15-bit synchronous prescaler configured through the PREDIV\_S bits of the RTC\_PRER register.

Note:

When both prescalers are used, it is recommended to configure the asynchronous prescaler to a high value to minimize consumption.

The asynchronous prescaler division factor is set to 128, and the synchronous division factor to 256, to obtain an internal clock frequency of 1 Hz (ck\_spre) with an LSE frequency of 32.768 kHz.

The minimum division factor is 1 and the maximum division factor is  $2^{22}$ .

This corresponds to a maximum input frequency of around 1 MHz.

f<sub>ck apre</sub> is given by the following formula:

$$f_{CK\_APRE} = \frac{f_{RTCCLK}}{PREDIV\_A + 1}$$

The ck\_apre clock is used to clock the binary RTC\_SSR subseconds downcounter. When it reaches 0, RTC\_SSR is reloaded with the content of PREDIV\_S.

f<sub>ck spre</sub> is given by the following formula:

$$f_{CK\_SPRE} = \frac{f_{RTCCLK}}{(PREDIV\_S + 1) \times (PREDIV\_A + 1)}$$

The ck\_spre clock can be used either to update the calendar or as timebase for the 16-bit wakeup auto-reload timer. To obtain short timeout periods, the 16-bit wakeup auto-reload timer can also run with the RTCCLK divided by the programmable 4-bit asynchronous prescaler (see *Section 22.3.4: Periodic auto-wakeup* for details).

#### 22.3.2 Real-time clock and calendar

The RTC calendar time and date registers are accessed through shadow registers which are synchronized with PCLK1 (APB1 clock). They can also be accessed directly in order to avoid waiting for the synchronization duration.

- RTC\_SSR for the subseconds
- RTC TR for the time
- RTC\_DR for the date

Every two RTCCLK periods, the current calendar value is copied into the shadow registers, and the RSF bit of RTC\_ISR register is set (see *Section 22.6.4*). The copy is not performed in Stop and Standby mode. When exiting these modes, the shadow registers are updated after up to 2 RTCCLK periods.

When the application reads the calendar registers, it accesses the content of the shadow registers. it is possible to make a direct access to the calendar registers by setting the BYPSHAD control bit in the RTC\_CR register. By default, this bit is cleared, and the user accesses the shadow registers.

When reading the RTC\_SSR, RTC\_TR or RTC\_DR registers in BYPSHAD=0 mode, the frequency of the APB clock (f<sub>APB</sub>) must be at least 7 times the frequency of the RTC clock (f<sub>RTCCLK</sub>).

The shadow registers are reset by system reset.

#### 22.3.3 Programmable alarms

The RTC unit provides two programmable alarms, Alarm A and Alarm B.

The programmable alarm functions are enabled through the ALRAIE and ALRBIE bits in the RTC\_CR register. The ALRAF and ALRBF flags are set to 1 if the calendar subseconds, seconds, minutes, hours, date or day match the values programmed in the alarm registers RTC\_ALRMASSR/RTC\_ALRMAR and RTC\_ALRMBSSR/RTC\_ALRMBR, respectively. Each calendar field can be independently selected through the MSKx bits of the RTC\_ALRMAR and RTC\_ALRMBR registers, and through the MASKSSx bits of the RTC\_ALRMASSR and RTC\_ALRMBSSR registers. The alarm interrupts are enabled through the ALRAIE and ALRBIE bits in the RTC\_CR register.

Alarm A and Alarm B (if enabled by bits OSEL[0:1] in RTC\_CR register) can be routed to the AFO\_ALARM output. AFO\_ALARM polarity can be configured through bit POL the RTC\_CR register.

#### Caution:

If the seconds field is selected (MSK0 bit reset in RTC\_ALRMAR or RTC\_ALRMBR), the synchronous prescaler division factor set in the RTC\_PRER register must be at least 3 to ensure correct behavior.

### 22.3.4 Periodic auto-wakeup

The periodic wakeup flag is generated by a 16-bit programmable auto-reload down-counter. The wakeup timer range can be extended to 17 bits.

The wakeup function is enabled through the WUTE bit in the RTC\_CR register.

The wakeup timer clock input can be:

- RTC clock (RTCCLK) divided by 2, 4, 8, or 16.
   When RTCCLK is LSE(32.768kHz), this allows to configure the wakeup interrupt period from 122 μs to 32 s, with a resolution down to 61μs.
- ck\_spre (usually 1 Hz internal clock)

When ck\_spre frequency is 1Hz, this allows to achieve a wakeup time from 1 s to around 36 hours with one-second resolution. This large programmable time range is divided in 2 parts:

- from 1s to 18 hours when WUCKSEL [2:1] = 10
- and from around 18h to 36h when WUCKSEL[2:1] = 11. In this last case 216 is added to the 16-bit counter current value. When the initialization sequence is complete (see *Programming the wakeup timer on page 558*), the timer starts counting down. When the wakeup function is enabled, the down-counting remains active in low power modes. In addition, when it reaches 0, the WUTF flag is set in the RTC\_ISR register, and the wakeup counter is automatically reloaded with its reload value (RTC\_WUTR register value).

The WUTF flag must then be cleared by software.

When the periodic wakeup interrupt is enabled by setting the WUTIE bit in the RTC\_CR2 register, it can exit the device from low power modes.

The periodic wakeup flag can be routed to the AFO\_ALARM output provided it has been enabled through bits OSEL[0:1] of RTC\_CR register. AFO\_ALARM polarity can be configured through the POL bit in the RTC\_CR register.

System reset, as well as low power modes (Sleep, Stop and Standby) have no influence on the wakeup timer.

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### 22.3.5 RTC initialization and configuration

#### **RTC** register access

The RTC registers are 32-bit registers. The APB interface introduces 2 wait-states in RTC register accesses except on read accesses to calendar shadow registers when BYPSHAD=0.

#### RTC register write protection

After system reset, the RTC registers are protected against parasitic write access with the DBP bit of the PWR power control register (PWR\_CR). The DBP bit must be set to enable RTC registers write access.

After power-on reset, all the RTC registers are write-protected. Writing to the RTC registers is enabled by writing a key into the Write Protection register, RTC\_WPR.

The following steps are required to unlock the write protection on all the RTC registers except for RTC\_ISR[13:8], RTC\_TAFCR, and RTC\_BKPxR.

- Write '0xCA' into the RTC WPR register.
- 2. Write '0x53' into the RTC WPR register.

Writing a wrong key reactivates the write protection.

The protection mechanism is not affected by system reset.

#### Calendar initialization and configuration

To program the initial time and date calendar values, including the time format and the prescaler configuration, the following sequence is required:

- 1. Set INIT bit to 1 in the RTC\_ISR register to enter initialization mode. In this mode, the calendar counter is stopped and its value can be updated.
- 2. Poll INITF bit of in the RTC\_ISR register. The initialization phase mode is entered when INITF is set to 1. It takes around 2 RTCCLK clock cycles (due to clock synchronization).
- 3. To generate a 1 Hz clock for the calendar counter, program first the synchronous prescaler factor in RTC\_PRER register, and then program the asynchronous prescaler factor. Even if only one of the two fields needs to be changed, 2 separate write accesses must be performed to the RTC\_PRER register.
- 4. Load the initial time and date values in the shadow registers (RTC\_TR and RTC\_DR), and configure the time format (12 or 24 hours) through the FMT bit in the RTC\_CR register.
- 5. Exit the initialization mode by clearing the INIT bit. The actual calendar counter value is then automatically loaded and the counting restarts after 4 RTCCLK clock cycles.

When the initialization sequence is complete, the calendar starts counting.

Note:

After a system reset, the application can read the INITS flag in the RTC\_ISR register to check if the calendar has been initialized or not. If this flag equals 0, the calendar has not been initialized since the year field is set at its power-on reset default value (0x00).

To read the calendar after initialization, the software must first check that the RSF flag is set in the RTC\_ISR register.

### Daylight saving time

The daylight saving time management is performed through bits SUB1H, ADD1H, and BKP of the RTC\_CR register.

Using SUB1H or ADD1H, the software can subtract or add one hour to the calendar in one single operation without going through the initialization procedure.

In addition, the software can use the BKP bit to memorize this operation.

#### Programming the alarm

A similar procedure must be followed to program or update the programmable alarm (Alarm A or Alarm B):

- 1. Clear ALRAE or ALRBIE in RTC\_CR to disable Alarm A or Alarm B.
- Poll ALRAWF or ALRBWF in RTC\_ISR until it is set to make sure the access to alarm registers is allowed. This takes around 2 RTCCLK clock cycles (due to clock synchronization).
- Program the Alarm A or Alarm B registers (RTC\_ALRMASSR/RTC\_ALRMAR or RTC\_ALRMBSSR/RTC\_ALRMBR).
- 4. Set ALRAE or ALRBIE in the RTC\_CR register to enable Alarm A or Alarm B again.

Note:

Each change of the RTC\_CR register is taken into account after around 2 RTCCLK clock cycles due to clock synchronization.

#### Programming the wakeup timer

The following sequence is required to configure or change the wakeup timer auto-reload value (WUT[15:0] in RTC\_WUTR):

- 1. Clear WUTE in RTC CR to disable the wakeup timer.
- 2. Poll WUTWF until it is set in RTC\_ISR to make sure the access to wakeup auto-reload counter and to WUCKSEL[2:0] bits is allowed. It takes around 2 RTCCLK clock cycles (due to clock synchronization).
- 3. Program the wakeup auto-reload value WUT[15:0], and the wakeup clock selection (WUCKSEL[2:0] bits in RTC\_CR). Set WUTE in RTC\_CR to enable the timer again. The wakeup timer restarts down-counting.

#### 22.3.6 Reading the calendar

#### When BYPSHAD control bit is cleared in the RTC CR register

To read the RTC calendar registers (RTC\_SSR, RTC\_TR and RTC\_DR) properly, the APB1 clock frequency ( $f_{PCLK1}$ ) must be equal to or greater than seven times the  $f_{RTCCLK}$  RTC clock frequency. This ensures a secure behavior of the synchronization mechanism.

If the APB1 clock frequency is less than seven times the RTC clock frequency, the software must read the calendar time and date registers twice. If the second read of the RTC\_TR gives the same result as the first read, this ensures that the data is correct. Otherwise a third read access must be done. In any case the APB1 clock frequency must never be lower than the RTC clock frequency.

The RSF bit is set in RTC\_ISR register each time the calendar registers are copied into the RTC\_SSR, RTC\_TR and RTC\_DR shadow registers. The copy is performed every two RTCCLK cycles. To ensure consistency between the 3 values, reading either RTC\_SSR or

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RTC\_TR locks the values in the higher-order calendar shadow registers until RTC\_DR is read. In case the software makes read accesses to the calendar in a time interval smaller than 2 RTCCLK periods: RSF must be cleared by software after the first calendar read, and then the software must wait until RSF is set before reading again the RTC\_SSR, RTC\_TR and RTC\_DR registers.

After waking up from low power mode (Stop or Standby), RSF must be cleared by software. The software must then wait until it is set again before reading the RTC\_SSR, RTC\_TR and RTC\_DR registers.

The RSF bit must be cleared after wakeup and not before entering low power mode.

Note:

After a system reset, the software must wait until RSF is set before reading the RTC\_SSR, RTC\_TR and RTC\_DR registers. Indeed, a system reset resets the shadow registers to their default values.

After an initialization (refer to Calendar initialization and configuration on page 557): the software must wait until RSF is set before reading the RTC\_SSR, RTC\_TR and RTC\_DR registers.

After synchronization (refer to Section 22.3.8: RTC synchronization): the software must wait until RSF is set before reading the RTC\_SSR, RTC\_TR and RTC\_DR registers.

# When the BYPSHAD control bit is set in the RTC\_CR register (bypass shadow registers)

Reading the calendar registers gives the values from the calendar counters directly, thus eliminating the need to wait for the RSF bit to be set. This is especially useful after exiting from low power modes (STOP or Standby), since the shadow registers are not updated during these modes.

When the BYPSHAD bit is set to 1, the results of the different registers might not be coherent with each other if an RTCCLK edge occurs between two read accesses to the registers. Additionally, the value of one of the registers may be incorrect if an RTCCLK edge occurs during the read operation. The software must read all the registers twice, and then compare the results to confirm that the data is coherent and correct. Alternatively, the software can just compare the two results of the least-significant calendar register.

Note:

While BYPSHAD=1, instructions which read the calendar registers require one extra APB cycle to complete.

#### 22.3.7 Resetting the RTC

The calendar shadow registers (RTC\_SSR, RTC\_TR and RTC\_DR) and the RTC status register (RTC\_ISR) are reset to their default values by all available system reset sources.

On the contrary, the following registers are reset to their default values by a power-on reset and are not affected by a system reset: the RTC current calendar registers, the RTC control register (RTC\_CR), the prescaler register (RTC\_PRER), the RTC calibration registers (RTC\_CALIBR or RTC\_CALR), the RTC shift register (RTC\_SHIFTR), the RTC timestamp registers (RTC\_TSSSR, RTC\_TSTR and RTC\_TSDR), the RTC tamper and alternate function configuration register (RTC\_TAFCR), the RTC backup registers (RTC\_BKPxR), the wakeup timer register (RTC\_WUTR), the Alarm A and Alarm B registers (RTC\_ALRMASSR/RTC\_ALRMAR and RTC\_ALRMBSSR/RTC\_ALRMBR).

In addition, the RTC keeps on running under system reset if the reset source is different from the power-on reset one. When a power-on reset occurs, the RTC is stopped and all the RTC registers are set to their reset values.

### 22.3.8 RTC synchronization

The RTC can be synchronized to a remote clock with a high degree of precision. After reading the sub-second field (RTC\_SSR or RTC\_TSSSR), a calculation can be made of the precise offset between the times being maintained by the remote clock and the RTC. The RTC can then be adjusted to eliminate this offset by "shifting" its clock by a fraction of a second using RTC\_SHIFTR.

RTC\_SSR contains the value of the synchronous prescaler's counter. This allows one to calculate the exact time being maintained by the RTC down to a resolution of 1 / (PREDIV\_S + 1) seconds. As a consequence, the resolution can be improved by increasing the synchronous prescaler value (PREDIV\_S[14:0]. The maximum resolution allowed (30.52 µs with a 32768 Hz clock) is obtained with PREDIV S set to 0x7FFF.

However, increasing PREDIV\_S means that PREDIV\_A must be decreased in order to maintain the synchronous prescaler's output at 1 Hz. In this way, the frequency of the asynchronous prescaler's output increases, which may increase the RTC dynamic consumption.

The RTC can be finely adjusted using the RTC shift control register (RTC\_SHIFTR). Writing to RTC\_SHIFTR can shift (either delay or advance) the clock by up to a second with a resolution of 1 / (PREDIV\_S + 1) seconds. The shift operation consists of adding the SUBFS[14:0] value to the synchronous prescaler counter SS[15:0]: this will delay the clock. If at the same time the ADD1S bit is set, this results in adding one second and at the same time subtracting a fraction of second, so this will advance the clock.

Caution:

Before initiating a shift operation, the user must check that SS[15] = 0 in order to ensure that no overflow will occur.

As soon as a shift operation is initiated by a write to the RTC\_SHIFTR register, the SHPF flag is set by hardware to indicate that a shift operation is pending. This bit is cleared by hardware as soon as the shift operation has completed.

Caution:

This synchronization feature is not compatible with the reference clock detection feature: firmware must not write to RTC\_SHIFTR when REFCKON=1.

#### 22.3.9 RTC reference clock detection

The RTC calendar update can be synchronized to a reference clock RTC\_REFIN, usually the mains (50 or 60 Hz). The RTC\_REFIN reference clock should have a higher precision than the 32.768 kHz LSE clock. When the RTC\_REFIN detection is enabled (REFCKON bit of RTC\_CR set to 1), the calendar is still clocked by the LSE, and RTC\_REFIN is used to compensate for the imprecision of the calendar update frequency (1 Hz).

Each 1 Hz clock edge is compared to the nearest reference clock edge (if one is found within a given time window). In most cases, the two clock edges are properly aligned. When the 1 Hz clock becomes misaligned due to the imprecision of the LSE clock, the RTC shifts the 1 Hz clock a bit so that future 1 Hz clock edges are aligned. Thanks to this mechanism, the calendar becomes as precise as the reference clock.

The RTC detects if the reference clock source is present by using the 256 Hz clock (ck\_apre) generated from the 32.768 kHz quartz. The detection is performed during a time window around each of the calendar updates (every 1 s). The window equals 7 ck\_apre periods when detecting the first reference clock edge. A smaller window of 3 ck\_apre periods is used for subsequent calendar updates.

Each time the reference clock is detected in the window, the asynchronous prescaler which outputs the ck\_apre clock is forced to reload. This has no effect when the reference clock and the 1 Hz clock are aligned because the prescaler is being reloaded at the same moment. When the clocks are not aligned, the reload shifts future 1 Hz clock edges a little for them to be aligned with the reference clock.

If the reference clock halts (no reference clock edge occurred during the 3 ck\_apre window), the calendar is updated continuously based solely on the LSE clock. The RTC then waits for the reference clock using a large 7 ck\_apre period detection window centered on the ck spre edge.

When the reference clock detection is enabled, PREDIV\_A and PREDIV\_S must be set to their default values:

- PREDIV\_A = 0x007F
- PREVID\_S = 0x00FF

Note: The reference clock detection is not available in Standby mode.

**Caution:** The reference clock detection feature cannot be used in conjunction with the coarse digital calibration: RTC\_CALIBR must be kept at 0x0000 0000 when REFCKON=1.

### 22.3.10 RTC coarse digital calibration

Two digital calibration methods are available: coarse and smooth calibration. To perform coarse calibration refer to *Section 22.6.7: RTC calibration register (RTC CALIBR)*.

The two calibration methods are not intended to be used together, the application must select one of the two methods. Coarse calibration is provided for compatibly reasons. To perform smooth calibration refer to Section 22.3.11: RTC smooth digital calibration and the Section 22.6.16: RTC calibration register (RTC\_CALR)

The coarse digital calibration can be used to compensate crystal inaccuracy by adding (positive calibration) or masking (negative calibration) clock cycles at the output of the asynchronous prescaler (ck apre).

Positive and negative calibration are selected by setting the DCS bit in RTC\_CALIBR register to '0' and '1', respectively.

When positive calibration is enabled (DCS = '0'), 2 ck\_apre cycles are added every minute (around 15360 ck\_apre cycles) for 2xDC minutes. This causes the calendar to be updated sooner, thereby adjusting the effective RTC frequency to be a bit higher.

When negative calibration is enabled (DCS = '1'), 1 ck\_apre cycle is removed every minute (around 15360 ck\_apre cycles) for 2xDC minutes. This causes the calendar to be updated later, thereby adjusting the effective RTC frequency to be a bit lower.

DC is configured through bits DC[4:0] of RTC\_CALIBR register. This number ranges from 0 to 31 corresponding to a time interval (2xDC) ranging from 0 to 62.

The coarse digital calibration can be configured only in initialization mode, and starts when the INIT bit is cleared. The full calibration cycle lasts 64 minutes. The first 2xDC minutes of the 64 -minute cycle are modified as just described.

Negative calibration can be performed with a resolution of about 2 ppm while positive calibration can be performed with a resolution of about 4 ppm. The maximum calibration ranges from –63 ppm to 126 ppm.

The calibration can be performed either on the LSE or on the HSE clock.

**Caution:** Digital calibration may not work correctly if PREDIV A < 6.

#### Case of RTCCLK=32.768 kHz and PREDIV\_A+1=128

The following description assumes that ck\_apre frequency is 256 Hz obtained with an LSE clock nominal frequency of 32.768 kHz, and PREDIV\_A set to 127 (default value).

The ck\_spre clock frequency is only modified during the first 2xDC minutes of the 64-minute cycle. For example, when DC equals 1, only the first 2 minutes are modified. This means that the first 2xDC minutes of each 64-minute cycle have, once per minute, one second either shortened by 256 or lengthened by 128 RTCCLK cycles, given that each ck\_apre cycle represents 128 RTCCLK cycles (with PREDIV A+1=128).

Therefore each calibration step has the effect of adding 512 or subtracting 256 oscillator cycles for every 125829120 RTCCLK cycles (64min x 60 s/min x 32768 cycles/s). This is equivalent to +4.069 ppm or-2.035 ppm per calibration step. As a result, the calibration resolution is +10.5 or -5.27 seconds per month, and the total calibration ranges from +5.45 to -2.72 minutes per month.

In order to measure the clock deviation, a 512 Hz clock is output for calibration. Refer to *Section 22.3.14: Calibration clock output*.

### 22.3.11 RTC smooth digital calibration

RTC frequency can be digitally calibrated with a resolution of about 0.954 ppm with a range from -487.1 ppm to +488.5 ppm. The correction of the frequency is performed using series of small adjustments (adding and/or subtracting individual RTCCLK pulses). These adjustments are fairly well distributed so that the RTC is well calibrated even when observed over short durations of time.

The smooth digital calibration is performed during a cycle of about  $2^{20}$  RTCCLK pulses, or 32 seconds when the input frequency is 32768 Hz.

The smooth calibration register (RTC\_CALR) specifies the number of RTCCLK clock cycles to be masked during the 32-second cycle:

- Setting the bit CALM[0] to 1 causes exactly one pulse to be masked during the 32second cycle.
- Setting CALM[1] to 1 causes two additional cycles to be masked
- Setting SMC[2] to 1 causes four additional cycles to be masked
- and so on up to SMC[8] set to 1 which causes 256 clocks to be masked.

While CALM allows the RTC frequency to be reduced by up to 487.1 ppm with fine resolution, the bit CALP can be used to increase the frequency by 488.5 ppm. Setting CALP to '1' effectively inserts an extra RTCCLK pulse every 2<sup>11</sup> RTCCLK cycles, which means that 512 clocks are added during every 32-second cycle.

Using CALM together with CALP, an offset ranging from -511 to +512 RTCCLK cycles can be added during the 32-second cycle, which translates to a calibration range of -487.1 ppm to +488.5 ppm with a resolution of about 0.954 ppm.

The formula to calculate the effective calibrated frequency (FCAL) given the input frequency (FRTCCLK) is as follows:

$$F_{CAL} = F_{RTCCLK} x [1 + (CALP x 512 - CALM) / (2^{20} + CALM - CALP x 512)]$$

#### Calibration when PREDIV A<3

The CALP bit can not be set to 1 when the asynchronous prescaler value (PREDIV\_A bits in RTC\_PRER register) is less than 3. If CALP was already set to 1 and PREDIV\_A bits are set to a value less than 3, CALP is ignored and the calibration operates as if CALP was equal to 0.

To perform a calibration with PREDIV\_A less than 3, the synchronous prescaler value (PREDIV\_S) should be reduced so that each second is accelerated by 8 RTCCLK clock cycles, which is equivalent to adding 256 clock cycles every 32 seconds. As a result, between 255 and 256 clock pulses (corresponding to a calibration range from 243.3 to 244.1 ppm) can effectively be added during each 32-second cycle using only the CALM bits.

With a nominal RTCCLK frequency of 32768 Hz, when PREDIV\_A equals 1 (division factor of 2), PREDIV\_S should be set to 16379 rather than 16383 (4 less). The only other interesting case is when PREDIV\_A equals 0, PREDIV\_S should be set to 32759 rather than 32767 (8 less).

If PREDIV\_S is reduced in this way, the formula given the effective frequency of the calibrated input clock is as follows:

$$F_{CAL} = F_{RTCCLK} x [1 + (256 - CALM) / (2^{20} + CALM - 256)]$$

In this case, CALM[7:0] equals 0x100 (the midpoint of the CALM range) is the correct setting if RTCCLK is exactly 32768.00 Hz.

#### Verifying the RTC calibration

RTC precision is performed by measuring the precise frequency of RTCCLK and calculating the correct CALM value and CALP values. An optional 1 Hz output is provided to allow applications to measure and verify the RTC precision.

Measuring the precise frequency of the RTC over a limited interval can result in a measurement error of up to 2 RTCCLK clock cycles over the measurement period, depending on how the digital calibration cycle is aligned with the measurement period.

However, this measurement error can be eliminated if the measurement period is the same length as the calibration cycle period. In this case, the only error observed is the error due to the resolution of the digital calibration.

- By default, the calibration cycle period is 32 seconds.
   Using this mode and measuring the accuracy of the 1 Hz output over exactly 32 seconds guarantees that the measure is within 0.477 ppm (0.5 RTCCLK cycles over 32 seconds, due to the limitation of the calibration resolution).
- CALW16 bit of the RTC\_CALR register can be set to 1 to force a 16- second calibration cycle period.
  - In this case, the RTC precision can be measured during 16 seconds with a maximum error of 0.954 ppm (0.5 RTCCLK cycles over 16 seconds). However, since the calibration resolution is reduced, the long term RTC precision is also reduced to 0.954 ppm: CALM[0] bit is stuck at 0 when CALW16 is set to 1.
- CALW8 bit of the RTC\_CALR register can be set to 1 to force a 8- second calibration cycle period.
  - In this case, the RTC precision can be measured during 8 seconds with a maximum error of 1.907 ppm (0.5 RTCCLK cycles over 8s). The long term RTC precision is also reduced to 1.907 ppm: CALM[1:0] bits are stuck at 00 when CALW8 is set to 1.

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### Re-calibration on-the-fly

The calibration register (RTC\_CALR) can be updated on-the-fly while RTC\_ISR/INITF=0, by using the follow process:

- 1. Poll the RTC\_ISR/RECALPF (re-calibration pending flag).
- 2. If it is set to 0, write a new value to RTC\_CALR, if necessary. RECALPF is then automatically set to 1
- 3. Within three ck\_apre cycles after the write operation to RTC\_CALR, the new calibration settings take effect.

### 22.3.12 Timestamp function

Timestamp is enabled by setting the TSE bit of RTC\_CR register to 1.

The calendar is saved in the timestamp registers (RTC\_TSSSR, RTC\_TSTR, RTC\_TSDR) when a timestamp event is detected on the pin to which the TIMESTAMP alternate function is mapped. When a timestamp event occurs, the timestamp flag bit (TSF) in RTC\_ISR register is set.

By setting the TSIE bit in the RTC\_CR register, an interrupt is generated when a timestamp event occurs.

If a new timestamp event is detected while the timestamp flag (TSF) is already set, the timestamp overflow flag (TSOVF) flag is set and the timestamp registers (RTC\_TSTR and RTC\_TSDR) maintain the results of the previous event.

Note:

TSF is set 2 ck\_apre cycles after the timestamp event occurs due to synchronization process.

There is no delay in the setting of TSOVF. This means that if two timestamp events are close together, TSOVF can be seen as '1' while TSF is still '0'. As a consequence, it is recommended to poll TSOVF only after TSF has been set.

#### Caution:

If a timestamp event occurs immediately after the TSF bit is supposed to be cleared, then both TSF and TSOVF bits are set. To avoid masking a timestamp event occurring at the same moment, the application must not write '0' into TSF bit unless it has already read it to '1'.

Optionally, a tamper event can cause a timestamp to be recorded. See the description of the TAMPTS control bit in *Section 22.6.17: RTC tamper and alternate function configuration register (RTC\_TAFCR)*. If the timestamp event is on the same pin as a tamper event configured in filtered mode (TAMPFLT set to a non-zero value), the timestamp on tamper detection event mode must be selected by setting TAMPTS='1' in RTC\_TAFCR register.

#### **TIMESTAMP** alternate function

The TIMESTAMP alternate function can be mapped to either RTC\_AF1 or RTC\_AF2 depending on the value of the TSINSEL bit in the RTC\_TAFCR register (see *Section 22.6.17 on page 585*). Mapping the timestamp event on RTC\_AF2 is not allowed if RTC\_AF1 is used as TAMPER in filtered mode (TAMPFLT set to a non-zero value).

#### 22.3.13 Tamper detection

Two tamper detection inputs are available. They can be configured either for edge detection, or for level detection with filtering.

### RTC backup registers

The backup registers (RTC\_BKPxR) are twenty 32-bit registers for storing 80 bytes of user application data. They are implemented in the backup domain that remains powered-on by  $V_{BAT}$  when the  $V_{DD}$  power is switched off. They are not reset by system reset, power-on reset, or when the device wakes up from Standby mode.

The backup registers are reset when a tamper detection event occurs (see *Section 22.6.20: RTC backup registers (RTC BKPxR)* and *Tamper detection initialization on page 565.* 

#### **Tamper detection initialization**

Each tamper detection input is associated with a flag TAMP1F/TAMP2F in the RTC\_ISR2 register. Each input can be enabled by setting the corresponding TAMP1E/TAMP2E bits to 1 in the RTC\_TAFCR register.

A tamper detection event resets all backup registers (RTC BKPxR).

By setting the TAMPIE bit in the RTC\_TAFCR register, an interrupt is generated when a tamper detection event occurs.

#### Timestamp on tamper event

With TAMPTS set to '1, any tamper event causes a timestamp to occur. In this case, either the TSF bit or the TSOVF bit are set in RTC\_ISR, in the same manner as if a normal timestamp event occurs. The affected tamper flag register (TAMP1F, TAMP2F) is set at the same time that TSF or TSOVF is set.

#### Edge detection on tamper inputs

If the TAMPFLT bits are "00", the TAMPER pins generate tamper detection events(AFI\_TAMPER[2:1]) when either a rising edge is observed or an falling edge is observed depending on the corresponding TAMPxTRG bit. The internal pull-up resistors on the TAMPER inputs are deactivated when edge detection is selected.

#### Caution:

To avoid losing tamper detection events, the signal used for edge detection is logically ANDed with TAMPxE in order to detect a tamper detection event in case it occurs before the TAMPERx pin is enabled.

- When TAMPxTRG = 0: if the TAMPERx alternate function is already high before tamper detection is enabled (TAMPxE bit set to 1), a tamper event is detected as soon as TAMPERx is enabled, even if there was no rising edge on TAMPERx after TAMPxE was set
- When TAMPxTRG = 1: if the TAMPERx alternate function is already low before tamper detection is enabled, a tamper event is detected as soon as TAMPERx is enabled (even if there was no falling edge on TAMPERx after TAMPxE was set.

After a tamper event has been detected and cleared, the TAMPERx alternate function should be disabled and then re-enabled (TAMPxE set to 1) before re-programming the backup registers (RTC\_BKPxR). This prevents the application from writing to the backup registers while the TAMPERx value still indicates a tamper detection. This is equivalent to a level detection on the TAMPERx alternate function.

Note:

Tamper detection is still active when  $V_{DD}$  power is switched off. To avoid unwanted resetting of the backup registers, the pin to which the TAMPER alternate function is mapped should be externally tied to the correct level.

### Level detection with filtering on tamper inputs

Level detection with filtering is performed by setting TAMPFLT to a non-zero value. A tamper detection event is generated when either 2, 4, or 8 (depending on TAMPFLT) consecutive samples are observed at the level designated by the TAMPxTRG bits (TAMP1TRG/TAMP2TRG).

The TAMPER inputs are pre-charged through the I/O internal pull-up resistance before its state is sampled, unless disabled by setting TAMPPUDIS to 1, The duration of the precharge is determined by the TAMPPRCH bits, allowing for larger capacitances on the tamper inputs.

The trade-off between tamper detection latency and power consumption through the pull-up can be optimized by using TAMPFREQ to determine the frequency of the sampling for level detection.

Note: Refer to the datasheets for the electrical characteristics of the pull-up resistors.

#### **TAMPER** alternate function detection

The TAMPER1 alternate function can be mapped either to RTC\_AF1(PC13) or RTC\_AF2 (PI8) depending on the value of TAMP1INSEL bit in RTC\_TAFCR register (see Section 22.6.17: RTC tamper and alternate function configuration register (RTC\_TAFCR)). TAMPE bit must be cleared when TAMP1INSEL is modified to avoid unwanted setting of TAMPF.

The TAMPER 2 alternate functions is mapped to RTC\_TAMPER2 pin.

### 22.3.14 Calibration clock output

When the COE bit is set to 1 in the RTC\_CR register, a reference clock is provided on the RTC\_CALIB device output. If the COSEL bit in the RTC\_CR register is reset and PREDIV\_A = 0x7F, the RTC\_CALIB frequency is f<sub>RTCCLK</sub>/64. This corresponds to a calibration output at 512 Hz for an RTCCLK frequency at 32.768 kHz.

The RTC\_CALIB output is not impacted by the calibration value programmed in RTC\_CALIBR register. The RTC\_CALIB duty cycle is irregular: there is a light jitter on falling edges. It is therefore recommended to use rising edges.

If COSEL is set and "PREDIV\_S+1" is a non-zero multiple of 256 (i.e: PREDIV\_S[7:0] = 0xFF), the RTC\_CALIB frequency is fRTCCLK/(256 \* (PREDIV\_A+1)). This corresponds to a calibration output at 1 Hz for prescaler default values (PREDIV\_A = 0x7F, PREDIV\_S = 0xFF), with an RTCCLK frequency at 32.768 kHz.

#### Calibration alternate function output

When the COE bit in the RTC\_CR register is set to 1, the calibration alternate function (AFO\_CALIB) is enabled on RTC\_AF1.

#### 22.3.15 Alarm output

Three functions can be selected on Alarm output: ALRAF, ALRBF and WUTF. These functions reflect the contents of the corresponding flags in the RTC\_ISR register.

The OSEL[1:0] control bits in the RTC\_CR register are used to activate the alarm alternate function output (AFO\_ALARM) in RTC\_AF1, and to select the function which is output on AFO\_ALARM.

The polarity of the output is determined by the POL control bit in RTC\_CR so that the opposite of the selected flag bit is output when POL is set to 1.

### Alarm alternate function output

AFO\_ALARM can be configured in output open drain or output push-pull using the control bit ALARMOUTTYPE in the RTC\_TAFCR register.

Note:

Once AFO\_ALARM is enabled, it has priority over AFO\_CALIB (COE bit is don't care on RTC\_AF1).

When AFO\_CALIB or AFO\_ALARM is selected, RTC\_AF1 is automatically configured in output alternate function.

# 22.4 RTC and low power modes

Table 80. Effect of low power modes on RTC

Mode	Description
Sleep	No effect RTC interrupts cause the device to exit the Sleep mode.
Stop	The RTC remains active when the RTC clock source is LSE or LSI. RTC alarm, RTC tamper event, RTC time stamp event, and RTC Wakeup cause the device to exit the Stop mode.
Standby	The RTC remains active when the RTC clock source is LSE or LSI. RTC alarm, RTC tamper event, RTC time stamp event, and RTC Wakeup cause the device to exit the Standby mode.

# 22.5 RTC interrupts

All RTC interrupts are connected to the EXTI controller.

To enable the RTC Alarm interrupt, the following sequence is required:

- 1. Configure and enable the EXTI Line 17 in interrupt mode and select the rising edge sensitivity.
- 2. Configure and enable the RTC\_Alarm IRQ channel in the NVIC.
- 3. Configure the RTC to generate RTC alarms (Alarm A or Alarm B).

To enable the RTC Wakeup interrupt, the following sequence is required:

- Configure and enable the EXTI Line 22 in interrupt mode and select the rising edge sensitivity.
- 2. Configure and enable the RTC\_WKUP IRQ channel in the NVIC.
- 3. Configure the RTC to generate the RTC wakeup timer event.

To enable the RTC Tamper interrupt, the following sequence is required:

- Configure and enable the EXTI Line 21 in interrupt mode and select the rising edge sensitivity.
- 2. Configure and Enable the TAMP\_STAMP IRQ channel in the NVIC.
- 3. Configure the RTC to detect the RTC tamper event.

To enable the RTC TimeStamp interrupt, the following sequence is required:

- Configure and enable the EXTI Line 21 in interrupt mode and select the rising edge sensitivity.
- 2. Configure and Enable the TAMP\_STAMP IRQ channel in the NVIC.
- 3. Configure the RTC to detect the RTC timestamp event.

Table 81. Interrupt control bits

Interrupt event	Event flag	Enable control bit	Exit the Sleep mode	Exit the Stop mode	Exit the Standby mode
Alarm A	ALRAF	ALRAIE	yes	yes <sup>(1)</sup>	yes <sup>(1)</sup>
Alarm B	ALRBF	ALRBIE	yes	yes <sup>(1)</sup>	yes <sup>(1)</sup>
Wakeup	WUTF	WUTIE	yes	yes <sup>(1)</sup>	yes <sup>(1)</sup>
TimeStamp	TSF	TSIE	yes	yes <sup>(1)</sup>	yes <sup>(1)</sup>
Tamper1 detection	TAMP1F	TAMPIE	yes	yes <sup>(1)</sup>	yes <sup>(1)</sup>
Tamper2 detection <sup>(2)</sup>	TAMP2F	TAMPIE	yes	yes <sup>(1)</sup>	yes <sup>(1)</sup>

<sup>1.</sup> Wakeup from STOP and Standby modes is possible only when the RTC clock source is LSE or LSI.

<sup>2.</sup> If RTC\_TAMPER2 pin is present. Refer to device datasheet pinout.

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# 22.6 RTC registers

Refer to *Section 1.1* of the reference manual for a list of abbreviations used in register descriptions.

The peripheral registers can be accessed by words (32 bits).

### 22.6.1 RTC time register (RTC\_TR)

The RTC\_TR is the calendar time shadow register. This register must be written in initialization mode only. Refer to *Calendar initialization and configuration on page 557* and *Reading the calendar on page 558*.

Address offset: 0x00

Power-on reset value: 0x0000 0000

System reset: 0x0000 0000 when BYPSHAD = 0. Not affected when BYPSHAD = 1

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
				Reserved					PM	HT[	1:0]		HU[	3:0]	
				neserved					rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Reserv	eserv MNT[2:0] MNU[3:0] Reserv						Reserv		ST[2:0]			SU[	3:0]		
ed	rw	rw	rw	rw	rw	rw	rw	ed	rw	rw	rw	rw	rw	rw	rw

Bits 31-24 Reserved

Bit 23 Reserved, must be kept at reset value.

Bit 22 PM: AM/PM notation

0: AM or 24-hour format

1: PM

Bits 21:20 HT[1:0]: Hour tens in BCD format

Bit 16:16 HU[3:0]: Hour units in BCD format

Bit 15 Reserved, must be kept at reset value.

Bits 14:12 MNT[2:0]: Minute tens in BCD format

Bit 11:8 MNU[3:0]: Minute units in BCD format

Bit 7 Reserved, must be kept at reset value.

Bits 6:4 ST[2:0]: Second tens in BCD format

Bit 3:0 SU[3:0]: Second units in BCD format

Note: This register is write protected. The write access procedure is described in RTC register write protection on page 557.

### 22.6.2 RTC date register (RTC\_DR)

The RTC\_DR is the calendar date shadow register. This register must be written in initialization mode only. Refer to *Calendar initialization and configuration on page 557* and *Reading the calendar on page 558*.

Address offset: 0x04

Power-on reset value: 0x0000 2101

System reset: 0x0000 2101 when BYPSHAD = 0. Not affected when BYPSHAD = 1.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
			Rese	vrvod					YT[	3:0]			YU[	[3:0]	
			nese	rveu				rw	rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	WDU[2:0]	]	MT		MU	[3:0]		Book	erved	DT[	1:0]		DU[	[3:0]	
rw	rw	rw	rw	rw	rw	rw	rw	nese	i veu	rw	rw	rw	rw	rw	rw

Bits 31-24 Reserved

Bits 23:20 YT[3:0]: Year tens in BCD format

Bits 19:16 YU[3:0]: Year units in BCD format

Bits 15:13 WDU[2:0]: Week day units

000: forbidden 001: Monday

111: Sunday

Bit 12 MT: Month tens in BCD format

Bits 11:8 MU: Month units in BCD format

Bits 7:6 Reserved, must be kept at reset value.

Bits 5:4 DT[1:0]: Date tens in BCD format

Bits 3:0 DU[3:0]: Date units in BCD format

Note:

This register is write protected. The write access procedure is described in RTC register write protection on page 557.

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### 22.6.3 RTC control register (RTC\_CR)

Address offset: 0x08

Power-on value: 0x0000 0000 System reset: not affected

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	Reserved								OSE	L[1:0]	POL	COSEL	BKP	SUB1H	ADD1H
	neserved								rw	rw	rw	rw	rw	W	w
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
TSIE	WUTIE	ALRBIE	ALRAIE	TSE	WUTE	ALRBI E	ALRAE	DCE	FMT	BYPS HAD	REFCKON	TSEDGE	W	WUCKSEL[2:	
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:24 Reserved, must be kept at reset value.

Bit 23 COE: Calibration output enable

This bit enables the AFO\_CALIB RTC output

0: Calibration output disabled

1: Calibration output enabled

Bits 22:21 OSEL[1:0]: Output selection

These bits are used to select the flag to be routed to AFO\_ALARM RTC output

00: Output disabled

01: Alarm A output enabled 10: Alarm B output enabled

11: Wakeup output enabled

Bit 20 POL: Output polarity

This bit is used to configure the polarity of AFO\_ALARM RTC output

0: The pin is high when ALRAF/ALRBF/WUTF is asserted (depending on OSEL[1:0])

1: The pin is low when ALRAF/ALRBF/WUTF is asserted (depending on OSEL[1:0]).

Bit 19 COSEL: Calibration output selection

When COE=1, this bit selects which signal is output on RTC\_CALIB.

0: Calibration output is 512 Hz

1: Calibration output is 1 Hz

These frequencies are valid for RTCCLK at 32.768 kHz and prescalers at their default values (PREDIV\_A=127 and PREDIV\_S=255). Refer to Section 22.3.14: Calibration clock output

#### Bit 18 BKP: Backup

This bit can be written by the user to memorize whether the daylight saving time change has been performed or not.

#### Bit 17 **SUB1H**: Subtract 1 hour (winter time change)

When this bit is set outside initialization mode, 1 hour is subtracted to the calendar time if the current hour is not 0. This bit is always read as 0.

Setting this bit has no effect when current hour is 0.

0: No effect

1: Subtracts 1 hour to the current time. This can be used for winter time change.

#### Bit 16 ADD1H: Add 1 hour (summer time change)

When this bit is set outside initialization mode, 1 hour is added to the calendar time. This bit is always read as 0.

- 0: No effect
- 1: Adds 1 hour to the current time. This can be used for summer time change
- Bit 15 TSIE: Timestamp interrupt enable
  - 0: Timestamp Interrupt disable
  - 1: Timestamp Interrupt enable
- Bit 14 WUTIE: Wakeup timer interrupt enable
  - 0: Wakeup timer interrupt disabled
  - 1: Wakeup timer interrupt enabled
- Bit 13 ALRBIE: Alarm B interrupt enable
  - 0: Alarm B Interrupt disable
  - 1: Alarm B Interrupt enable
- Bit 12 ALRAIE: Alarm A interrupt enable
  - 0: Alarm A interrupt disabled
  - 1: Alarm A interrupt enabled
- Bit 11 TSE: Time stamp enable
  - 0: Time stamp disable
  - 1: Time stamp enable
- Bit 10 WUTE: Wakeup timer enable
  - 0: Wakeup timer disabled
  - 1: Wakeup timer enabled
- Bit 9 ALRBIE: Alarm B enable
  - 0: Alarm B disabled
  - 1: Alarm B enabled
- Bit 8 ALRAE: Alarm A enable
  - 0: Alarm A disabled
  - 1: Alarm A enabled
- Bit 7 DCE: Coarse digital calibration enable
  - 0: Digital calibration disabled
  - 1: Digital calibration enabled
  - PREDIV\_A must be 6 or greater
- Bit 6 FMT: Hour format
  - 0: 24 hour/day format
  - 1: AM/PM hour format
- Bit 5 BYPSHAD: Bypass the shadow registers
  - 0: Calendar values (when reading from RTC\_SSR, RTC\_TR, and RTC\_DR) are taken from the shadow registers, which are updated once every two RTCCLK cycles.
  - 1: Calendar values (when reading from RTC\_SSR, RTC\_TR, and RTC\_DR) are taken directly from the calendar counters.

Note: If the frequency of the APB1 clock is less than seven times the frequency of RTCCLK, BYPSHAD must be set to '1'.

- Bit 4 **REFCKON**: Reference clock detection enable (50 or 60 Hz)
  - 0: Reference clock detection disabled
  - 1: Reference clock detection enabled
  - Note: PREDIV\_S must be 0x00FF.

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Bit 3 TSEDGE: Timestamp event active edge

0: TIMESTAMP rising edge generates a timestamp event

1: TIMESTAMP falling edge generates a timestamp event

TSE must be reset when TSEDGE is changed to avoid unwanted TSF setting

#### Bits 2:0 WUCKSEL[2:0]: Wakeup clock selection

000: RTC/16 clock is selected 001: RTC/8 clock is selected 010: RTC/4 clock is selected

011: RTC/2 clock is selected

10x: ck\_spre (usually 1 Hz) clock is selected

11x: ck\_spre (usually 1 Hz) clock is selected and 2<sup>16</sup> is added to the WUT counter value (see note below)

Note:

 $WUT = Wakeup \ unit \ counter \ value. \ WUT = (0x0000 \ to \ 0xFFFF) + 0x10000 \ added \ when \ WUCKSEL[2:1 = 11].$ 

Bits 7, 6 and 4 of this register can be written in initialization mode only (RTC\_ISR/INITF = 1).

Bits 2 to 0 of this register can be written only when RTC\_CR WUTE bit = 0 and RTC\_ISR WUTWF bit = 1.

It is recommended not to change the hour during the calendar hour increment as it could mask the incrementation of the calendar hour.

ADD1H and SUB1H changes are effective in the next second.

This register is write protected. The write access procedure is described in RTC register write protection on page 557.

# 22.6.4 RTC initialization and status register (RTC\_ISR)

Address offset: 0x0C

Reset value: 0x0000 0007

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
															RECAL PF
							Reserved								r
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Res.	TAMP 2F	TAMP 1F	TSOVF	TSF	WUTF	ALRBF	ALRAF	INIT	INITF	RSF	INITS	SHPF	WUT WF	ALRB WF	ALRA WF
	rc_w0	rc_w0	rc_w0	rc_w0	rc_w0	rc_w0	rc_w0	rw	r	rc_w0	r	rc_w0	r	r	r

Bits 31:17 Reserved

#### Bit 16 **RECALPF**: Recalibration pending Flag

The RECALPF status flag is automatically set to '1' when software writes to the RTC\_CALR register, indicating that the RTC\_CALR register is blocked. When the new calibration settings are taken into account, this bit returns to '0'. Refer to Section: Re-calibration on-the-fly.

Bit 15 Reserved, must be kept at reset value.

#### Bit 14 TAMP2F: TAMPER2 detection flag

This flag is set by hardware when a tamper detection event is detected on tamper input 2. It is cleared by software writing 0.

#### Bit 13 TAMP1F: Tamper detection flag

This flag is set by hardware when a tamper detection event is detected.

It is cleared by software writing 0.

#### Bit 12 TSOVF: Timestamp overflow flag

This flag is set by hardware when a timestamp event occurs while TSF is already set. This flag is cleared by software by writing 0. It is recommended to check and then clear TSOVF only after clearing the TSF bit. Otherwise, an overflow might not be noticed if a timestamp event occurs immediately before the TSF bit is cleared.

#### Bit 11 TSF: Timestamp flag

This flag is set by hardware when a timestamp event occurs.

This flag is cleared by software by writing 0.

#### Bit 10 WUTF: Wakeup timer flag

This flag is set by hardware when the wakeup auto-reload counter reaches 0.

This flag is cleared by software by writing 0.

This flag must be cleared by software at least 1.5 RTCCLK periods before WUTF is set to 1 again.

#### Bit 9 ALRBF: Alarm B flag

This flag is set by hardware when the time/date registers (RTC\_TR and RTC\_DR) match the Alarm B register (RTC\_ALRMBR).

This flag is cleared by software by writing 0.

#### Bit 8 ALRAF: Alarm A flag

This flag is set by hardware when the time/date registers (RTC\_TR and RTC\_DR) match the Alarm A register (RTC\_ALRMAR).

This flag is cleared by software by writing 0.

#### Bit 7 INIT: Initialization mode

0: Free running mode

1: Initialization mode used to program time and date register (RTC\_TR and RTC\_DR), and prescaler register (RTC\_PRER). Counters are stopped and start counting from the new value when INIT is reset.

#### Bit 6 INITF: Initialization flag

When this bit is set to 1, the RTC is in initialization state, and the time, date and prescaler registers can be updated.

- 0: Calendar registers update is not allowed
- 1: Calendar registers update is allowed.

#### Bit 5 RSF: Registers synchronization flag

This bit is set by hardware each time the calendar registers are copied into the shadow registers (RTC\_SSRx, RTC\_TRx and RTC\_DRx). This bit is cleared by hardware in initialization mode, while a shift operation is pending (SHPF=1), or when in bypass shadow register mode (BYPSHAD=1). This bit can also be cleared by software.

- 0: Calendar shadow registers not yet synchronized
- 1: Calendar shadow registers synchronized

#### Bit 4 INITS: Initialization status flag

This bit is set by hardware when the calendar year field is different from 0 (power-on reset state).

- 0: Calendar has not been initialized
- 1: Calendar has been initialized

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#### Bit 3 SHPF: Shift operation pending

0: No shift operation is pending

1: A shift operation is pending

This flag is set by hardware as soon as a shift operation is initiated by a write to the RTC\_SHIFTR. It is cleared by hardware when the corresponding shift operation has been executed. Writing to SHPF has no effect.

#### Bit 2 WUTWF: Wakeup timer write flag

This bit is set by hardware when the wakeup timer values can be changed, after the WUTE bit has been set to 0 in RTC\_CR.

- 0: Wakeup timer configuration update not allowed
- 1: Wakeup timer configuration update allowed

#### Bit 1 ALRBWF: Alarm B write flag

This bit is set by hardware when Alarm B values can be changed, after the ALRBIE bit has been set to 0 in RTC\_CR.

It is cleared by hardware in initialization mode.

0: Alarm B update not allowed

1: Alarm B update allowed.

#### Bit 0 ALRAWF: Alarm A write flag

This bit is set by hardware when Alarm A values can be changed, after the ALRAE bit has been set to 0 in RTC\_CR.

It is cleared by hardware in initialization mode.

0: Alarm A update not allowed

1: Alarm A update allowed

Note: The ALRAF, ALRBF, WUTF and TSF bits are cleared 2 APB clock cycles after programming them to 0.

This register is write protected (except for RTC\_ISR[13:8] bits). The write access procedure is described in RTC register write protection on page 557.

### 22.6.5 RTC prescaler register (RTC\_PRER)

Address offset: 0x10

Power-on reset value: 0x007F 00FF

System reset: not affected

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
				Reserved	1						PR	EDIV_A[6	6:0]		
				nesei veu					rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Res.							PRE	EDIV_S[1	4:0]						
nes.	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:24 Reserved

Bit 23 Reserved, must be kept at reset value.

Bits 22:16 PREDIV\_A[6:0]: Asynchronous prescaler factor

This is the asynchronous division factor:

ck\_apre frequency = RTCCLK frequency/(PREDIV\_A+1)

Note: PREDIV\_A [6:0]= 000000 is a prohibited value.

Bit 15 Reserved, must be kept at reset value.

Bits 14:0 PREDIV\_S[14:0]: Synchronous prescaler factor

This is the synchronous division factor:

ck\_spre frequency = ck\_apre frequency/(PREDIV\_S+1)

Note:

This register must be written in initialization mode only. The initialization must be performed in two separate write accesses. Refer to Calendar initialization and configuration on page 557

This register is write protected. The write access procedure is described in RTC register write protection on page 557.

### 22.6.6 RTC wakeup timer register (RTC\_WUTR)

Address offset: 0x14

Power-on reset value: 0x0000 FFFF

System reset: not affected

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							WUT	[15:0]							
rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bits 31:16 Reserved

Bits 15:0 WUT[15:0]: Wakeup auto-reload value bits

When the wakeup timer is enabled (WUTE set to 1), the WUTF flag is set every (WUT[15:0] + 1) ck\_wut cycles. The ck\_wut period is selected through WUCKSEL[2:0] bits of the RTC\_CR register

When WUCKSEL[2] = 1, the wakeup timer becomes 17-bits and WUCKSEL[1] effectively becomes WUT[16] the most-significant bit to be reloaded into the timer.

Note: The first assertion of WUTF occurs (WUT+1) ck\_wut cycles after WUTE is set. Setting WUT[15:0] to 0x0000 with WUCKSEL[2:0] =011 (RTCCLK/2) is forbidden.

Note:

This register can be written only when WUTWF is set to 1 in RTC\_ISR.

This register is write protected. The write access procedure is described in RTC register write protection on page 557.

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## 22.6.7 RTC calibration register (RTC\_CALIBR)

Address offset: 0x18

Power-on reset value: 0x0000 0000

System reset: not affected

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			Rese	arvod				DCS	Rese	rvod			DC[4:0	]	
			nese	erveu				rw	пезе	i veu	rw	rw	rw	rw	rw

Bits 31:8 Reserved

Bit 7 DCS: Digital calibration sign

0: Positive calibration: calendar update frequency is increased1: Negative calibration: calendar update frequency is decreased

Bits 6:5 Reserved, must be kept at reset value.

Bits 4:0 **DC[4:0]**: Digital calibration

DCS = 0 (positive calibration)

00000: + 0 ppm

00001: + 4 ppm (rounded value)

00010: + 8 ppm (rounded value)

..

11111: + 126 ppm (rounded value)

DCS = 1 (negative calibration)

00000: -0 ppm

00001: -2 ppm (rounded value)

00010: -4 ppm (rounded value)

••

11111: -63 ppm (rounded value)

Refer to Case of RTCCLK=32.768 kHz and PREDIV\_A+1=128 on page 562 for the exact step value.

Note: This register can be written in initialization mode only (RTC\_ISR/INITF = '1').

This register is write protected. The write access procedure is described in RTC register write protection on page 557.

## 22.6.8 RTC alarm A register (RTC\_ALRMAR)

Address offset: 0x1C

Power-on reset value: 0x0000 0000

System reset: not affected

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
MSK4	WDSEL	DT[	[1:0]		DU[	[3:0]		MSK3	PM	HT[	1:0]		HU	[3:0]	
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
MSK2		MNT[2:0]	]		MNU	J[3:0]		MSK1		ST[2:0]			SU[	[3:0]	
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bit 31 MSK4: Alarm A date mask

0: Alarm A set if the date/day match

1: Date/day don't care in Alarm A comparison

Bit 30 WDSEL: Week day selection

0: DU[3:0] represents the date units

1: DU[3:0] represents the week day. DT[1:0] is don't care.

Bits 29:28 DT[1:0]: Date tens in BCD format.

Bits 27:24 DU[3:0]: Date units or day in BCD format.

Bit 23 MSK3: Alarm A hours mask

0: Alarm A set if the hours match

1: Hours don't care in Alarm A comparison

Bit 22 PM: AM/PM notation

0: AM or 24-hour format

1: PM

Bits 21:20 HT[1:0]: Hour tens in BCD format.

Bits 19:16 HU[3:0]: Hour units in BCD format.

Bit 15 MSK2: Alarm A minutes mask

0: Alarm A set if the minutes match

1: Minutes don't care in Alarm A comparison

Bits 14:12 MNT[2:0]: Minute tens in BCD format.

Bits 11:8 MNU[3:0]: Minute units in BCD format.

Bit 7 MSK1: Alarm A seconds mask

0: Alarm A set if the seconds match

1: Seconds don't care in Alarm A comparison

Bits 6:4 ST[2:0]: Second tens in BCD format.

Bits 3:0 SU[3:0]: Second units in BCD format.

Note: This register can be written only when ALRAWF is set to 1 in RTC\_ISR, or in initialization mode.

This register is write protected. The write access procedure is described in RTC register write protection on page 557.

## 22.6.9 RTC alarm B register (RTC\_ALRMBR)

Address offset: 0x20

Power-on reset value: 0x0000 0000

System reset: not affected

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	
MSK4	WDSEL	DI	Γ[1:0]		DU	[3:0]		MSK3	PM	НТ	[1:0]		Н	U[3:0]		
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	
15	14	10	12	11	10	9	8	7	6	E	4	3	2	4	0	
15	14	13	12	11	10	9	8	- /	6	5	4	3		<u> </u>	0	
MSK2		MNT[2:0	)]		MNU	J[3:0]		MSK1		ST[2:0]			SU[3:0]			
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	

Bit 31 MSK4: Alarm B date mask

0: Alarm B set if the date and day match

1: Date and day don't care in Alarm B comparison

Bit 30 WDSEL: Week day selection

0: DU[3:0] represents the date units

1: DU[3:0] represents the week day. DT[1:0] is don't care.

Bits 29:28 DT[1:0]: Date tens in BCD format

Bits 27:24 DU[3:0]: Date units or day in BCD format

Bit 23 MSK3: Alarm B hours mask

0: Alarm B set if the hours match

1: Hours don't care in Alarm B comparison

Bit 22 PM: AM/PM notation

0: AM or 24-hour format

1: PM

Bits 21:20 HT[1:0]: Hour tens in BCD format

Bits 19:16 HU[3:0]: Hour units in BCD format

Bit 15 MSK2: Alarm B minutes mask

0: Alarm B set if the minutes match

1: Minutes don't care in Alarm B comparison

Bits 14:12 MNT[2:0]: Minute tens in BCD format

Bits 11:8 MNU[3:0]: Minute units in BCD format

Bit 7 MSK1: Alarm B seconds mask

0: Alarm B set if the seconds match

1: Seconds don't care in Alarm B comparison

Bits 6:4 ST[2:0]: Second tens in BCD format

Bits 3:0 SU[3:0]: Second units in BCD format

Note: This register can be written only when ALRBWF is set to 1 in RTC\_ISR, or in initialization mode.

This register is write protected. The write access procedure is described in RTC register write protection on page 557.

## 22.6.10 RTC write protection register (RTC\_WPR)

Address offset: 0x24

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
 15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			Rese	ruod							KE	ΞY			
			nese	rveu				W	w	w	w	W	w	w	w

Bits 31:8 Reserved, must be kept at reset value.

Bits 7:0 KEY: Write protection key

This byte is written by software.

Reading this byte always returns 0x00.

Refer to *RTC register write protection* for a description of how to unlock RTC register write protection.

## 22.6.11 RTC sub second register (RTC\_SSR)

Address offset: 0x28

Power-on reset value: 0x0000 0000

System reset: 0x0000 0000 when BYPSHAD = 0. Not affected when BYPSHAD = 1.

	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
								Rese	erved							
	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r
																_
_	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		-			-	-	-	SS[	15:0]	-	-					
	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits31:16 Reserved

Bits 15:0 SS: Sub second value

SS[15:0] is the value in the synchronous prescaler's counter. The fraction of a second is given by the formula below:

Second fraction = ( PREDIV\_S - SS ) / ( PREDIV\_S + 1 )

Note: SS can be larger than PREDIV\_S only after a shift operation. In that case, the correct time/date is one second less than as indicated by RTC\_TR/RTC\_DR.

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## 22.6.12 RTC shift control register (RTC\_SHIFTR)

Address offset: 0x2C

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
ADD1S								Reserved	ł						
w	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Rsvd							S	UBFS[14	:0]						
r	W	W	w	w	W	w	w	W	w	W	w	W	w	w	w

#### Bit 31 ADD1S: Add one second

0: No effect

1: Add one second to the clock/calendar

This bit is write only and is always read as zero. Writing to this bit has no effect when a shift operation is pending (when SHPF=1, in RTC\_ISR).

This function is intended to be used with SUBFS (see description below) in order to effectively add a fraction of a second to the clock in an atomic operation.

Bits 31:15 Reserved

#### Bits 14:0 SUBFS: Subtract a fraction of a second

These bits are write only and is always read as zero. Writing to this bit has no effect when a shift operation is pending (when SHPF=1, in RTC\_ISR).

The value which is written to SUBFS is added to the synchronous prescaler's counter. Since this counter counts down, this operation effectively subtracts from (delays) the clock by:

Delay (seconds) = SUBFS / ( PREDIV\_S + 1 )

A fraction of a second can effectively be added to the clock (advancing the clock) when the ADD1S function is used in conjunction with SUBFS, effectively advancing the clock by:

Advance (seconds) = (1 - (SUBFS / (PREDIV\_S + 1))).

Note: Writing to SUBFS causes RSF to be cleared. Software can then wait until RSF=1 to be sure that the shadow registers have been updated with the shifted time.

Refer to Section 22.3.8: RTC synchronization.

Note: This register is write protected. The write access procedure is described in RTC register write protection on page 557

## 22.6.13 RTC time stamp time register (RTC\_TSTR)

Address offset: 0x30

Power-on reset value: 0x0000 0000

System reset: not affected

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
				Reserved	1				PM	НТІ	[1:0]		Н	J[3:0]	
				neserved	,				r	r	r	r	r	r	r
								_		_					
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Reserv-		MNT[2:0	)]		MNU	J[3:0]		Reserv-		ST[2:0]			SU	[3:0]	
ed	r	r	r	r	r	r	r	ed	r	r	r	r	r	r	r

Bits 31:23 Reserved, must be kept at reset value.

Bit 22 PM: AM/PM notation

0: AM or 24-hour format

1: PM

Bits 21:20 HT[1:0]: Hour tens in BCD format.

Bits 19:16 HU[3:0]: Hour units in BCD format.

Bit 15 Reserved, must be kept at reset value.

Bits 14:12 MNT[2:0]: Minute tens in BCD format.

Bits 11:8 MNU[3:0]: Minute units in BCD format.

Bit 7 Reserved, must be kept at reset value.

Bits 6:4 ST[2:0]: Second tens in BCD format.

Bits 3:0 SU[3:0]: Second units in BCD format.

Note:

The content of this register is valid only when TSF is set to 1 in RTC\_ISR. It is cleared when TSF bit is reset.

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#### 22.6.14 RTC time stamp date register (RTC\_TSDR)

Address offset: 0x34

Power-on reset value: 0x0000 0000

System reset: not affected

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Res	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	WDU[1:0	]	MT		MU	[3:0]		Book	erved	DT[	1:0]		DU	[3:0]	
r	r	r	r	r	r	r	r	nese	erveu	r	r	r	r	r	r

Bits 31:16 Reserved, must be kept at reset value.

Bits 15:13 WDU[1:0]: Week day units

Bit 12 MT: Month tens in BCD format

Bits 11:8 MU[3:0]: Month units in BCD format

Bits 7:6 Reserved, must be kept at reset value.

Bits 5:4 DT[1:0]: Date tens in BCD format

Bit 3:0 DU[3:0]: Date units in BCD format

Note: The content of this register is valid only when TSF is set to 1 in RTC\_ISR. It is cleared when

TSF bit is reset.

#### 22.6.15 RTC timestamp sub second register (RTC\_TSSSR)

Address offset: 0x38

Power-on reset value: 0x0000 0000

System reset: not affected

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r
45	4.4	40	10	44	10	0		7		_	4	0	0		
15	14	13	12	11	10	9	8	/	6	5	4	3	2	ı	0
							SS[	15:0]							
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:16 Reserved

Bits 15:0 SS: Sub second value

SS[15:0] is the value of the synchronous prescaler's counter when the timestamp event occurred.

Note: The content of this register is valid only when RTC\_ISR/TSF is set. It is cleared when the

RTC\_ISR/TSF bit is reset.

## 22.6.16 RTC calibration register (RTC\_CALR)

Address offset: 0x3C

Power-on reset value: 0x0000 0000

System reset: not affected

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
CALP	CALW8	CALW16		Res	erved					(	CALM[8:	0]	•	•	
rw	rw	rw	r	r	r	r	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bit 31:16 Reserved

Bit 15 CALP: Increase frequency of RTC by 488.5 ppm

0: No RTCCLK pulses are added.

1: One RTCCLK pulse is effectively inserted every 2<sup>11</sup> pulses (frequency increased by 488.5 ppm).

This feature is intended to be used in conjunction with CALM, which lowers the frequency of the calendar with a fine resolution. if the input frequency is 32768 Hz, the number of RTCCLK pulses added during a 32-second window is calculated as follows: (512 \* CALP) - CALM.

Refer to Section 22.3.11: RTC smooth digital calibration.

Bit 14 CALW8: Use an 8-second calibration cycle period

When CALW8 is set to '1', the 8-second calibration cycle period is selected.

CALM[1:0] are stuck at "00" when CALW8='1'.

Refer to Section 22.3.11: RTC smooth digital calibration.

Bit 13 CALW16: Use a 16-second calibration cycle period

When CALW16 is set to '1', the 16-second calibration cycle period is selected. This bit must not be set to '1' if CALW8=1.

Note: CALM[0] is stuck at '0' when CALW16='1'.

Refer to Section 22.3.11: RTC smooth digital calibration.

Bits 12:9 Reserved

Bits 8:0 CALM[8:0]: Calibration minus

The frequency of the calendar is reduced by masking CALM out of  $2^{20}$  RTCCLK pulses (32 seconds if the input frequency is 32768 Hz). This decreases the frequency of the calendar with a resolution of 0.9537 ppm.

To increase the frequency of the calendar, this feature should be used in conjunction with CALP

See Section 22.3.11: RTC smooth digital calibration on page 562.

Note: This register is write protected. The write access procedure is described in RTC register write protection on page 557

Real-time clock (RTC)

RM0090

# 22.6.17 RTC tamper and alternate function configuration register (RTC\_TAFCR)

Address offset: 0x40

Power-on reset value: 0x0000 0000

System reset: not affected

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
						Reserved	l						ALARMOUT TYPE	TSIN SEL	TAMP1 INSEL
													rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
TAMP- PUDIS		ИР- H[1:0]	TAMPE	ELT[1:0]	TAN	/IPFREQ[	2:0]	TAMPT S	Rese	erved	TAMP2 -TRG	TAMP2 E	TAMPIE	TAMP1 TRG	TAMP1 E
rw	rw	rw	rw	rw	rw	rw	rw	rw			rw	rw	rw	rw	rw

Bit 31:19 Reserved. Always read as 0.

Bit 18 ALARMOUTTYPE: AFO\_ALARM output type

0: ALARM\_AF0 is an open-drain output

1: ALARM\_AF0 is a push-pull output

Bit 17 TSINSEL: TIMESTAMP mapping

0: RTC\_AF1 used as TIMESTAMP

1: RTC\_AF2 used as TIMESTAMP

Bit 16 TAMP1INSEL: TAMPER1 mapping

0: RTC\_AF1 used as TAMPER

1: RTC\_AF2 used as TAMPER

Note: TAMP1E must be reset when TAMP1INSEL is changed to avoid unwanted setting of

TAMP1F.

Bit 15 TAMPPUDIS: TAMPER pull-up disable

This bit determines if each of the tamper pins are pre-charged before each sample.

0: Precharge tamper pins before sampling (enable internal pull-up)

1: Disable precharge of tamper pins

Note:

## Bits 14:13 TAMPPRCH[1:0]: Tamper precharge duration

These bit determines the duration of time during which the pull-up/is activated before each sample. TAMPPRCH is valid for each of the tamper inputs.

0x0: 1 RTCCLK cycle 0x1: 2 RTCCLK cycles 0x2: 4 RTCCLK cycles 0x3: 8 RTCCLK cycles

## Bits 12:11 TAMPFLT[1:0]: Tamper filter count

These bits determines the number of consecutive samples at the specified level (TAMP\*TRG) necessary to activate a Tamper event. TAMPFLT is valid for each of the tamper inputs.

0x0: Tamper is activated on edge of tamper input transitions to the active level (no internal pull-up on tamper input).

0x1: Tamper is activated after 2 consecutive samples at the active level.

0x2: Tamper is activated after 4 consecutive samples at the active level.

0x3: Tamper is activated after 8 consecutive samples at the active level.

## Bits 10:8 TAMPFREQ[2:0]: Tamper sampling frequency

Determines the frequency at which each of the tamper inputs are sampled.

0x0: RTCCLK / 32768 (1 Hz when RTCCLK = 32768 Hz)

0x1: RTCCLK / 16384 (2 Hz when RTCCLK = 32768 Hz)

0x2: RTCCLK / 8192 (4 Hz when RTCCLK = 32768 Hz)

0x3: RTCCLK / 4096 (8 Hz when RTCCLK = 32768 Hz)

0x4: RTCCLK / 2048 (16 Hz when RTCCLK = 32768 Hz)

0x5: RTCCLK / 1024 (32 Hz when RTCCLK = 32768 Hz)

0x6: RTCCLK / 512 (64 Hz when RTCCLK = 32768 Hz)

0x7: RTCCLK / 256 (128 Hz when RTCCLK = 32768 Hz)

#### Bit 7 TAMPTS: Activate timestamp on tamper detection event

0: Tamper detection event does not cause a timestamp to be saved

1: Save timestamp on tamper detection event

TAMPTS is valid even if TSE=0 in the RTC\_CR register.

#### Bit 6:5 Reserved. Always read as 0.

#### Bit 4 TAMP2TRG: Active level for tamper 2

if TAMPFLT != 00:

0: TAMPER2 staying low triggers a tamper detection event.

1: TAMPER2 staying high triggers a tamper detection event.

if TAMPFLT = 00:

0: TAMPER2 rising edge triggers a tamper detection event.

1: TAMPER2 falling edge triggers a tamper detection event.

### Bit 3 TAMP2E: Tamper 2 detection enable

0: Tamper 2 detection disabled

1: Tamper 2 detection enabled

### Bit 2 TAMPIE: Tamper interrupt enable

0: Tamper interrupt disabled

1: Tamper interrupt enabled

## Bit 1 TAMP1TRG: Active level for tamper 1

if TAMPFLT != 00

0: TAMPER1 staying low triggers a tamper detection event.

1: TAMPER1 staying high triggers a tamper detection event.

if TAMPFLT = 00:

0: TAMPER1 rising edge triggers a tamper detection event.

1: TAMPER1 falling edge triggers a tamper detection event.

**Caution:** When TAMPFLT = 0, TAMP1E must be reset when TAMP1TRG is changed to avoid spuriously setting TAMP1F.

Bit 0 TAMP1E: Tamper 1 detection enable

0: Tamper 1 detection disabled1: Tamper 1 detection enabled

## 22.6.18 RTC alarm A sub second register (RTC\_ALRMASSR)

Address offset: 0x44

Power-on reset value: 0x0000 0000

System reset: not affected

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	Rese	ved			MASKS	S[3:0]					Res	erved			
r	r	r	r	rw	rw	rw	rw	r	r	r	r	r	r	r	r
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Reserved								SS[14:0]							
r	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	W	rw	rw

Bit 31:28 Reserved

Bit 27:24 MASKSS[3:0]: Mask the most-significant bits starting at this bit

0: No comparison on sub seconds for Alarm A. The alarm is set when the seconds unit is incremented (assuming that the rest of the fields match).

1: SS[14:1] are don't care in Alarm A comparison. Only SS[0] is compared.

2: SS[14:2] are don't care in Alarm A comparison. Only SS[1:0] are compared.

3: SS[14:3] are don't care in Alarm A comparison. Only SS[2:0] are compared.

...

12: SS[14:12] are don't care in Alarm A comparison. SS[11:0] are compared.

13: SS[14:13] are don't care in Alarm A comparison. SS[12:0] are compared.

14: SS[14] is don't care in Alarm A comparison. SS[13:0] are compared.

15: All 15 SS bits are compared and must match to activate alarm.

The overflow bits of the synchronous counter (bits 15) is never compared. This bit can be different from 0 only after a shift operation.

Bit 23:15 Reserved

Bit 14:0 SS[14:0]: Sub seconds value

This value is compared with the contents of the synchronous prescaler's counter to determine if Alarm A is to be activated. Only bits 0 up MASKSS-1 are compared.

Note: This register can be written only when ALRAE is reset in RTC\_CR register, or in initialization mode.

This register is write protected. The write access procedure is described in RTC register write protection on page 557

## 22.6.19 RTC alarm B sub second register (RTC\_ALRMBSSR)

Address offset: 0x48

Power-on reset value: 0x0000 0000

System reset: not affected

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16				
Reserved MASKSS[3:0]									Reserved										
r	r	r	r	rw	rw	rw	rw	r	r	r	r	r	r	r	r				
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0				
Reserved								SS[14:0]											
r	rw	rw	rw	rw	rw	w	rw	rw											

Bit 31:28 Reserved

Bit 27:24 MASKSS[3:0]: Mask the most-significant bits starting at this bit

0x0: No comparison on sub seconds for Alarm B. The alarm is set when the seconds unit is incremented (assuming that the rest of the fields match).

0x1: SS[14:1] are don't care in Alarm B comparison. Only SS[0] is compared.

0x2: SS[14:2] are don't care in Alarm B comparison. Only SS[1:0] are compared.

0x3: SS[14:3] are don't care in Alarm B comparison. Only SS[2:0] are compared.

...

0xC: SS[14:12] are don't care in Alarm B comparison. SS[11:0] are compared.

0xD: SS[14:13] are don't care in Alarm B comparison. SS[12:0] are compared.

0xE: SS[14] is don't care in Alarm B comparison. SS[13:0] are compared.

0xF: All 15 SS bits are compared and must match to activate alarm.

The overflow bits of the synchronous counter (bits 15) is never compared. This bit can be different from 0 only after a shift operation.

Bit 23:15 Reserved

Bit 14:0 SS[14:0]: Sub seconds value

This value is compared with the contents of the synchronous prescaler's counter to determine if Alarm B is to be activated. Only bits 0 up to MASKSS-1 are compared.

Note: This register can be written only when ALRBIE is reset in RTC\_CR register, or in initialization mode.

This register is write protected. The write access procedure is described in Section: RTC register write protection

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## 22.6.20 RTC backup registers (RTC\_BKPxR)

Address offset: 0x50 to 0x9C

Power-on reset value: 0x0000 0000

System reset: not affected

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	BKP[31:16]														
rw	rw r														
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							BKP[	15:0]							
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	W	rw	rw

Bits 31:0 BKP[31:0]

The application can write or read data to and from these registers.

They are powered-on by  $V_{BAT}$  when  $V_{DD}$  is switched off, so that they are not reset by System reset, and their contents remain valid when the device operates in low-power mode. This register is reset on a tamper detection event, as long as TAMPxF=1. or when the Flash readout protection is disabled.

## 22.6.21 RTC register map

Table 82. RTC register map and reset values

		gister illa									,																						
Offset	Register	31 30 29 28	27 26	25 24	23	22	21	20	18	17	16	15	14	13	12	11	10	6	œ	7	9	2	4	ဗ	7	1	0						
0x00	RTC_TR	Re	served			PM	H7 [1:0	0]		J[3:0		Reserved		NT[2			ИNU		)]	Reserved		T[2:				3:0]							
	Reset value					0	0	0 (	0 0	0	0	æ	0	0	0	0	0	0	0	ď	0	0	0	0	0	0	0						
0x04	RTC_DR	Res	erved			YT[	3:0]		YL	[3:0]	]	WE	) DU[2	2:0]	MT		MU[	3:0]		Becerved	200	D [1:			DU[	3:0]							
	Reset value											0	0	1	0	0	0	0	1	ď	-	0	0	0	0	0	1						
0x08	RTC_CR	Reso	Reserved O					POL	BKP	SUB1H	ADD1H	TSIE	WUTIE	ALRBIE	ALRAIE	TSE	WUTE	ALRBIE	ALRAE	DCE	FMT	BYPSHAD	REFCKON	TSEDGE		CKSI [2:0]							
	Reset value							0 (	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0						
0x0C	RTC_ISR			Re	Reserved											TSF	WUTF	ALRBF	ALRAF	LN	INITF	RSF	INITS	SHPF WUTWF		ALRBWF	ALRAWF						
	Reset value																		0	0	0	0	0	0	0	0	0	0	0	0	1	0	1
0x10	RTC_PRER	Re	served			PR	EDI	/_A[6	6:0]		Reserved						PR	EDI	V_8	S[14	:0]												
	Reset value					1	1	1	1 1	1	1	0         0         0         0         0         0         1							1	1	1												
0x14	RTC_WUTR			Rese	erve	ved							W	/UT[15:0]																			
	Reset value											1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1						
0x18	RTC_CALIBR								Reserved								DCS	Reserved			D	C[4:	0]										
	Reset value		_																	0	B	2	0	0	0	0	0						
0x1C												MNT[2:0] MNU[3:0]					MSK1				SU[3:0]												
	Reset value	0 0 0 0						0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0						

Table 82. RTC register map and reset values (continued)

Offset	Register	30 31	28 27 26	25 24	23	22	21	19	1 0	16	15	14	13	1	10	8	7	9 2	4	ю c	7	0
0x20	RTC_ALRMBR  Reset value	0 0 MSK4	0]   DU(		o MSK3	O PM	HT [1:0]		U[3:	0]	o MSK2		Γ[2:0]		/NU[3:	-	o MSK2	ST[2			J[3:0	-
		0 0 0 0 1	0 0 0	0 0	U	U	0 0	[0]	o j c	)   0	JU	0	0   0	0	0 0	0	U	0 0	-		0	0
0x24	RTC_WPR						Rese	erved											KEY	[7:0]		
	Reset value																0	0 0	0	0 (	0	0
0x28	RTC_SSR			Rese	erve	d										SS[	15:0]					
	Reset value	(0.									0	0	0 0	0	0 0	0	0	0 0	0	0 (	0	0
0x2C	RTC_SHIFTR	ADD1S		F	Rese	eserved										SUB	FS[14:0]					
	Reset value	0								0 0	0	0 0	0	0	0 0	0	0 (	0	0			
0x30	RTC_TSTR		Reserved			PM	HT[1:0]	Н	U[3:	0]	Reserved	S.C.T.I.V.	[0:3]   NIM   [7:0]   O   O	N	/NU[3:	0]	Reserved	ST[2	:0]	SI	J[3:0	)]
	Reset value					0	0 0	0 (	0 (	0	æ	0	0 0	0	0 0	0	æ	0 0	0	0 (	0	0
0x38	RTC_TSSSR			Rese	erve	d										SS[	[15:0]					
	Reset value										0						0 0 0 0 0 0 0 0					
0x3C	RTC_ CALR			Rese	erve	d					CALP	O CALW18					CALM[8:0]					
	Reset value										0	0	0	à	-	0	0	0 0	0	0 (	0	0
0x40	RTC_TAFCR		Res	served	O ALARMOUTTYPE O TSINSEL O TAMPINSEL					TAMPPUDIS	TAMPPRCH[1:0]		TAMPFLT[1:0]	TAMPFREQ[2:0]		TAMPTS	Reserved	TAMP2TRG	TAMP2E	TAMP1ETRG	TAMP1E	
	Reset value								0 0	0	0	0	0 0	0	0 0	0	0		0	0 (	0	0
0x44	RTC_ ALRMASSR	Reserved					Re	serve	d							SS	S[14	:0]				
	Reset value		0 0	0 0								0	0 0	0	0 0	0	0	0 0	0	0 (	0	0
0x48	RTC_ ALRMBSSR	Reserved	MASKS	SS[3:0]			Re	serve	d							SS	S[14	:0]				
	Reset value		0 0	0 0							00000000000000000											
	RTC_BKP0R									BKP	[31:	0]										
0x50 to 0x9C	Reset value	0 0 0	0 0 0	0 0	0	0	0 0	0 (	0 (	0	0	0	0 0	0	0 0	0	0	0 0	0	0 (	0	0
io oxac	to RTC_BKP19R									BKP	[31:	0]										
	Reset value	0   0   0   0   0   0   0   0   0   0										0	0 0	0	0	0 0	0	0 (	0	0		

Refer to *Table 2 on page 50* for the register boundary addresses.

## 23 Inter-integrated circuit (I<sup>2</sup>C) interface

This section applies to the whole STM32F4xx family, unless otherwise specified.

## 23.1 I<sup>2</sup>C introduction

I<sup>2</sup>C (inter-integrated circuit) bus Interface serves as an interface between the microcontroller and the serial I<sup>2</sup>C bus. It provides multimaster capability, and controls all I<sup>2</sup>C bus-specific sequencing, protocol, arbitration and timing. It supports standard and fast speed modes. It is also SMBus 2.0 compatible.

It may be used for a variety of purposes, including CRC generation and verification, SMBus (system management bus) and PMBus (power management bus).

Depending on specific device implementation DMA capability can be available for reduced CPU overload.

## 23.2 I<sup>2</sup>C main features

- Parallel-bus/I<sup>2</sup>C protocol converter
- Multimaster capability: the same interface can act as Master or Slave
- I<sup>2</sup>C Master features:
  - Clock generation
  - Start and Stop generation
- I<sup>2</sup>C Slave features:
  - Programmable I<sup>2</sup>C Address detection
  - Dual Addressing Capability to acknowledge 2 slave addresses
  - Stop bit detection
- Generation and detection of 7-bit/10-bit addressing and General Call
- Supports different communication speeds:
  - Standard Speed (up to 100 kHz)
  - Fast Speed (up to 400 kHz)
- Status flags:
  - Transmitter/Receiver mode flag
  - End-of-Byte transmission flag
  - I<sup>2</sup>C busy flag
- Error flags:
  - Arbitration lost condition for master mode
  - Acknowledgement failure after address/ data transmission
  - Detection of misplaced start or stop condition
  - Overrun/Underrun if clock stretching is disabled
- 2 Interrupt vectors:
  - 1 Interrupt for successful address/ data communication
  - 1 Interrupt for error condition

- Optional clock stretching
- 1-byte buffer with DMA capability
- Configurable PEC (packet error checking) generation or verification:
  - PEC value can be transmitted as last byte in Tx mode
  - PEC error checking for last received byte
- SMBus 2.0 Compatibility:
  - 25 ms clock low timeout delay
  - 10 ms master cumulative clock low extend time
  - 25 ms slave cumulative clock low extend time
  - Hardware PEC generation/verification with ACK control
  - Address Resolution Protocol (ARP) supported
- PMBus Compatibility

Note:

Some of the above features may not be available in certain products. The user should refer to the product data sheet, to identify the specific features supported by the  $l^2C$  interface implementation.

## 23.3 I<sup>2</sup>C functional description

In addition to receiving and transmitting data, this interface converts it from serial to parallel format and vice versa. The interrupts are enabled or disabled by software. The interface is connected to the I<sup>2</sup>C bus by a data pin (SDA) and by a clock pin (SCL). It can be connected with a standard (up to 100 kHz) or fast (up to 400 kHz) I<sup>2</sup>C bus.

### 23.3.1 Mode selection

The interface can operate in one of the four following modes:

- Slave transmitter
- Slave receiver
- Master transmitter
- Master receiver

By default, it operates in slave mode. The interface automatically switches from slave to master, after it generates a START condition and from master to slave, if an arbitration loss or a Stop generation occurs, allowing multimaster capability.

#### **Communication flow**

In Master mode, the I<sup>2</sup>C interface initiates a data transfer and generates the clock signal. A serial data transfer always begins with a start condition and ends with a stop condition. Both start and stop conditions are generated in master mode by software.

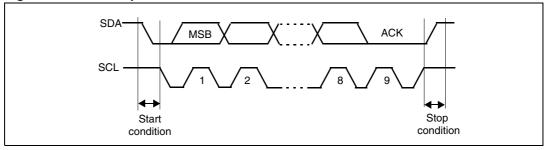
In Slave mode, the interface is capable of recognizing its own addresses (7 or 10-bit), and the General Call address. The General Call address detection may be enabled or disabled by software.

Data and addresses are transferred as 8-bit bytes, MSB first. The first byte(s) following the start condition contain the address (one in 7-bit mode, two in 10-bit mode). The address is always transmitted in Master mode.

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A 9th clock pulse follows the 8 clock cycles of a byte transfer, during which the receiver must send an acknowledge bit to the transmitter. Refer to *Figure 217*.

Figure 217. I<sup>2</sup>C bus protocol



Acknowledge may be enabled or disabled by software. The I<sup>2</sup>C interface addresses (dual addressing 7-bit/ 10-bit and/or general call address) can be selected by software.

The block diagram of the I<sup>2</sup>C interface is shown in *Figure 218*.

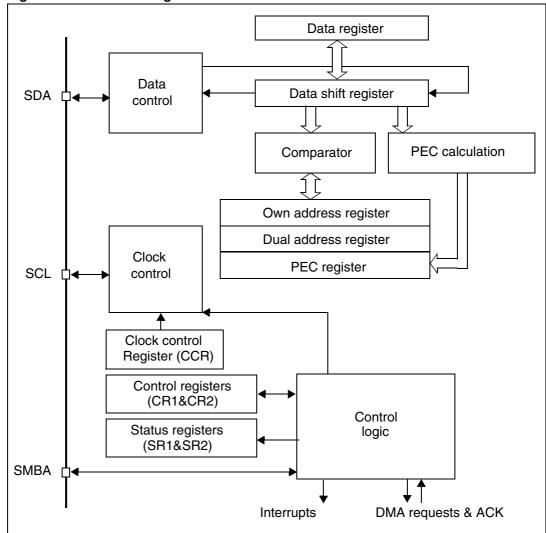


Figure 218. I<sup>2</sup>C block diagram

1. SMBA is an optional signal in SMBus mode. This signal is not applicable if SMBus is disabled.

## 23.3.2 I<sup>2</sup>C slave mode

By default the I<sup>2</sup>C interface operates in Slave mode. To switch from default Slave mode to Master mode a Start condition generation is needed.

The peripheral input clock must be programmed in the I2C\_CR2 register in order to generate correct timings. The peripheral input clock frequency must be at least:

- 2 MHz in Standard mode
- 4 MHz in Fast mode

As soon as a start condition is detected, the address is received from the SDA line and sent to the shift register. Then it is compared with the address of the interface (OAR1) and with OAR2 (if ENDUAL=1) or the General Call address (if ENGC = 1).

Note:

In 10-bit addressing mode, the comparison includes the header sequence (11110xx0), where xx denotes the two most significant bits of the address.

**Header or address not matched**: the interface ignores it and waits for another Start condition.

**Header matched** (10-bit mode only): the interface generates an acknowledge pulse if the ACK bit is set and waits for the 8-bit slave address.

**Address matched**: the interface generates in sequence:

- An acknowledge pulse if the ACK bit is set
- The ADDR bit is set by hardware and an interrupt is generated if the ITEVFEN bit is set.
- If ENDUAL=1, the software has to read the DUALF bit to check which slave address has been acknowledged.

In 10-bit mode, after receiving the address sequence the slave is always in Receiver mode. It will enter Transmitter mode on receiving a repeated Start condition followed by the header sequence with matching address bits and the least significant bit set (11110xx1).

The TRA bit indicates whether the slave is in Receiver or Transmitter mode.

#### Slave transmitter

Following the address reception and after clearing ADDR, the slave sends bytes from the DR register to the SDA line via the internal shift register.

The slave stretches SCL low until ADDR is cleared and DR filled with the data to be sent (see *Figure 219* Transfer sequencing EV1 EV3).

When the acknowledge pulse is received:

 The TxE bit is set by hardware with an interrupt if the ITEVFEN and the ITBUFEN bits are set.

If TxE is set and some data were not written in the I2C\_DR register before the end of the next data transmission, the BTF bit is set and the interface waits until BTF is cleared by a read to I2C\_SR1 followed by a write to the I2C\_DR register, stretching SCL low.



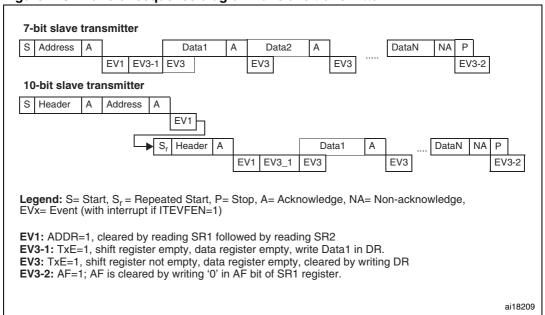


Figure 219. Transfer sequence diagram for slave transmitter

- 1. The EV1 and EV3\_1 events stretch SCL low until the end of the corresponding software sequence.
- The EV3 event stretches SCL low if the software sequence is not completed before the end of the next byte transmission

#### Slave receiver

Following the address reception and after clearing ADDR, the slave receives bytes from the SDA line into the DR register via the internal shift register. After each byte the interface generates in sequence:

- An acknowledge pulse if the ACK bit is set
- The RxNE bit is set by hardware and an interrupt is generated if the ITEVFEN and ITBUFEN bit is set.

If RxNE is set and the data in the DR register is not read before the end of the next data reception, the BTF bit is set and the interface waits until BTF is cleared by a read from the I2C\_DR register, stretching SCL low (see *Figure 220* Transfer sequencing).

7-bit slave receiver S Address A Data1 DataN Data2 EV1 EV2 EV2 EV2 EV4 10-bit slave receiver S Header Address Data1 DataN EV1 EV2 EV2 EV4 **Legend:** S= Start, S<sub>r</sub> = Repeated Start, P= Stop, A= Acknowledge, EVx= Event (with interrupt if ITEVFEN=1) EV1: ADDR=1, cleared by reading SR1 followed by reading SR2 **EV2:** RxNE=1 cleared by reading DR register. EV4: STOPF=1, cleared by reading SR1 register followed by writing to the CR1 register ai18208

Figure 220. Transfer sequence diagram for slave receiver

- 1. The EV1 event stretches SCL low until the end of the corresponding software sequence.
- 2. The EV2 event stretches SCL low if the software sequence is not completed before the end of the next byte reception.
- After checking the SR1 register content, the user should perform the complete clearing sequence for each flag found set.
   Thus, for ADDR and STOPF flags, the following sequence is required inside the I2C interrupt routine: READ SR1 if (ADDR == 1) {READ SR1: READ SR2}

if (ADDR == 1) {READ SR1; READ SR2} if (STOPF == 1) {READ SR1; WRITE CR1}

The purpose is to make sure that both ADDR and STOPF flags are cleared if both are found set.

## Closing slave communication

After the last data byte is transferred a Stop Condition is generated by the master. The interface detects this condition and sets:

The STOPF bit and generates an interrupt if the ITEVFEN bit is set.

The STOPF bit is cleared by a read of the SR1 register followed by a write to the CR1 register (see *Figure 220: Transfer sequence diagram for slave receiver* EV4).

## 23.3.3 I<sup>2</sup>C master mode

In Master mode, the  $I^2C$  interface initiates a data transfer and generates the clock signal. A serial data transfer always begins with a Start condition and ends with a Stop condition. Master mode is selected as soon as the Start condition is generated on the bus with a START bit.

The following is the required sequence in master mode.

- Program the peripheral input clock in I2C\_CR2 Register in order to generate correct timings
- Configure the clock control registers
- Configure the rise time register
- Program the I2C\_CR1 register to enable the peripheral
- Set the START bit in the I2C CR1 register to generate a Start condition

The peripheral input clock frequency must be at least:

- 2 MHz in Standard mode
- 4 MHz in Fast mode

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#### Start condition

Setting the START bit causes the interface to generate a Start condition and to switch to Master mode (M/SL bit set) when the BUSY bit is cleared.

Note:

In master mode, setting the START bit causes the interface to generate a ReStart condition at the end of the current byte transfer.

Once the Start condition is sent:

The SB bit is set by hardware and an interrupt is generated if the ITEVFEN bit is set.

Then the master waits for a read of the SR1 register followed by a write in the DR register with the Slave address (see *Figure 221* and *Figure 222* Transfer sequencing EV5).

#### Slave address transmission

Then the slave address is sent to the SDA line via the internal shift register.

- In 10-bit addressing mode, sending the header sequence causes the following event:
  - The ADD10 bit is set by hardware and an interrupt is generated if the ITEVFEN bit is set.

Then the master waits for a read of the SR1 register followed by a write in the DR register with the second address byte (see *Figure 221* and *Figure 222* Transfer sequencing).

 The ADDR bit is set by hardware and an interrupt is generated if the ITEVFEN bit is set.

Then the master waits for a read of the SR1 register followed by a read of the SR2 register (see *Figure 221* and *Figure 222* Transfer sequencing).

In 7-bit addressing mode, one address byte is sent.

As soon as the address byte is sent,

 The ADDR bit is set by hardware and an interrupt is generated if the ITEVFEN bit is set.

Then the master waits for a read of the SR1 register followed by a read of the SR2 register (see *Figure 221* and *Figure 222* Transfer sequencing).

The master can decide to enter Transmitter or Receiver mode depending on the LSB of the slave address sent.

- In 7-bit addressing mode,
  - To enter Transmitter mode, a master sends the slave address with LSB reset.
  - To enter Receiver mode, a master sends the slave address with LSB set.
- In 10-bit addressing mode,
  - To enter Transmitter mode, a master sends the header (11110xx0) and then the slave address, (where xx denotes the two most significant bits of the address).
  - To enter Receiver mode, a master sends the header (11110xx0) and then the slave address. Then it should send a repeated Start condition followed by the header (11110xx1), (where xx denotes the two most significant bits of the address).

The TRA bit indicates whether the master is in Receiver or Transmitter mode.

#### Master transmitter

Following the address transmission and after clearing ADDR, the master sends bytes from the DR register to the SDA line via the internal shift register.

The master waits until the first data byte is written into I2C\_DR (see *Figure 221* Transfer sequencing EV8\_1).

When the acknowledge pulse is received, the TxE bit is set by hardware and an interrupt is generated if the ITEVFEN and ITBUFEN bits are set.

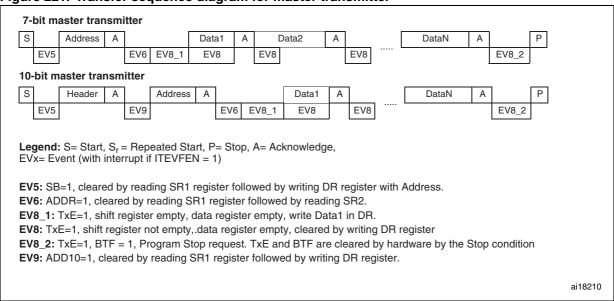
If TxE is set and a data byte was not written in the DR register before the end of the last data transmission, BTF is set and the interface waits until BTF is cleared by a write to I2C\_DR, stretching SCL low.

## Closing the communication

After the last byte is written to the DR register, the STOP bit is set by software to generate a Stop condition (see *Figure 221* Transfer sequencing EV8\_2). The interface automatically goes back to slave mode (M/SL bit cleared).

Note: Stop condition should be programmed during EV8\_2 event, when either TxE or BTF is set.

Figure 221. Transfer sequence diagram for master transmitter



- 1. The EV5, EV6, EV9, EV8\_1 and EV8\_2 events stretch SCL low until the end of the corresponding software sequence.
- 2. The EV8 event stretches SCL low if the software sequence is not complete before the end of the next byte transmission.

#### Master receiver

Following the address transmission and after clearing ADDR, the I<sup>2</sup>C interface enters Master Receiver mode. In this mode the interface receives bytes from the SDA line into the DR register via the internal shift register. After each byte the interface generates in sequence:

- An acknowledge pulse if the ACK bit is set
- 2. The RxNE bit is set and an interrupt is generated if the ITEVFEN and ITBUFEN bits are set (see *Figure 222* Transfer sequencing EV7).

If the RxNE bit is set and the data in the DR register is not read before the end of the last data reception, the BTF bit is set by hardware and the interface waits until BTF is cleared by a read in the DR register, stretching SCL low.

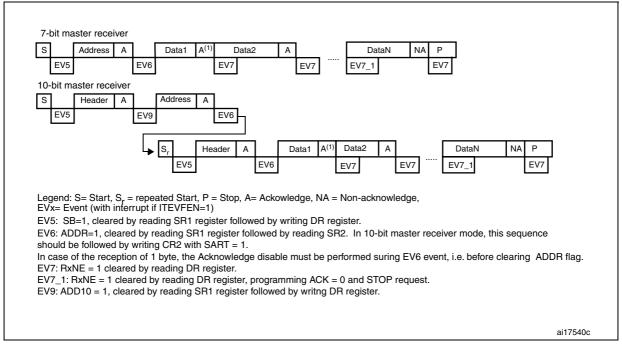
## Closing the communication

The master sends a NACK for the last byte received from the slave. After receiving this NACK, the slave releases the control of the SCL and SDA lines. Then the master can send a Stop/Restart condition.

- 1. To generate the nonacknowledge pulse after the last received data byte, the ACK bit must be cleared just after reading the second last data byte (after second last RxNE event).
- 2. In order to generate the Stop/Restart condition, software must set the STOP/START bit after reading the second last data byte (after the second last RxNE event).
- 3. In case a single byte has to be received, the Acknowledge disable is made during EV6 (before ADDR flag is cleared) and the STOP condition generation is made after EV6.

After the Stop condition generation, the interface goes automatically back to slave mode (M/SL bit cleared).

Figure 222. Transfer sequence diagram for master receiver



- 1. If a single byte is received, it is NA.
- 2. The EV5, EV6 and EV9 events stretch SCL low until the end of the corresponding software sequence.
- The EV7 event stretches SCL low if the software sequence is not completed before the end of the next byte reception.
- 4. The EV7\_1 software sequence must be completed before the ACK pulse of the current byte transfer.

The procedures described below are recommended if the EV7-1 software sequence is not completed before the ACK pulse of the current byte transfer.

These procedures must be followed to make sure:

- The ACK bit is set low on time before the end of the last data reception
- The STOP bit is set high after the last data reception without reception of supplementary data.

### For 2-byte reception:

- Wait until ADDR = 1 (SCL stretched low until the ADDR flag is cleared)
- Set ACK low, set POS high
- Clear ADDR flag
- Wait until BTF = 1 (Data 1 in DR, Data2 in shift register, SCL stretched low until a data 1 is read)
- Set STOP high
- Read data 1 and 2

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#### For N >2 -byte reception, from N-2 data reception

- Wait until BTF = 1 (data N-2 in DR, data N-1 in shift register, SCL stretched low until data N-2 is read)
- Set ACK low
- Read data N-2
- Wait until BTF = 1 (data N-1 in DR, data N in shift register, SCL stretched low until a data N-1 is read)
- Set STOP high
- Read data N-1 and N

#### 23.3.4 Error conditions

The following are the error conditions which may cause communication to fail.

## **Bus error (BERR)**

This error occurs when the I<sup>2</sup>C interface detects an external Stop or Start condition during an address or a data transfer. In this case:

- the BERR bit is set and an interrupt is generated if the ITERREN bit is set
- in Slave mode: data are discarded and the lines are released by hardware:
  - in case of a misplaced Start, the slave considers it is a restart and waits for an address, or a Stop condition
  - in case of a misplaced Stop, the slave behaves like for a Stop condition and the lines are released by hardware
- In Master mode: the lines are not released and the state of the current transmission is not affected. It is up to the software to abort or not the current transmission

#### Acknowledge failure (AF)

This error occurs when the interface detects a nonacknowledge bit. In this case:

- the AF bit is set and an interrupt is generated if the ITERREN bit is set
- a transmitter which receives a NACK must reset the communication:
  - If Slave: lines are released by hardware
  - If Master: a Stop or repeated Start condition must be generated by software

## **Arbitration lost (ARLO)**

This error occurs when the I<sup>2</sup>C interface detects an arbitration lost condition. In this case,

- the ARLO bit is set by hardware (and an interrupt is generated if the ITERREN bit is set)
- the I<sup>2</sup>C Interface goes automatically back to slave mode (the M/SL bit is cleared). When the I<sup>2</sup>C loses the arbitration, it is not able to acknowledge its slave address in the same transfer, but it can acknowledge it after a repeated Start from the winning master.
- lines are released by hardware

## Overrun/underrun error (OVR)

An overrun error can occur in slave mode when clock stretching is disabled and the I<sup>2</sup>C interface is receiving data. The interface has received a byte (RxNE=1) and the data in DR has not been read, before the next byte is received by the interface. In this case,

- The last received byte is lost.
- In case of Overrun error, software should clear the RxNE bit and the transmitter should re-transmit the last received byte.

Underrun error can occur in slave mode when clock stretching is disabled and the I<sup>2</sup>C interface is transmitting data. The interface has not updated the DR with the next byte (TxE=1), before the clock comes for the next byte. In this case,

- The same byte in the DR register will be sent again
- The user should make sure that data received on the receiver side during an underrun error are discarded and that the next bytes are written within the clock low time specified in the I<sup>2</sup>C bus standard.

For the first byte to be transmitted, the DR must be written after ADDR is cleared and before the first SCL rising edge. If not possible, the receiver must discard the first data.

## 23.3.5 SDA/SCL line control

- If clock stretching is enabled:
  - Transmitter mode: If TxE=1 and BTF=1: the interface holds the clock line low before transmission to wait for the microcontroller to write the byte in the Data Register (both buffer and shift register are empty).
  - Receiver mode: If RxNE=1 and BTF=1: the interface holds the clock line low after reception to wait for the microcontroller to read the byte in the Data Register (both buffer and shift register are full).
- If clock stretching is disabled in Slave mode:
  - Overrun Error in case of RxNE=1 and no read of DR has been done before the next byte is received. The last received byte is lost.
  - Underrun Error in case TxE=1 and no write into DR has been done before the next byte must be transmitted. The same byte will be sent again.
  - Write Collision not managed.

### 23.3.6 SMBus

#### Introduction

The System Management Bus (SMBus) is a two-wire interface through which various devices can communicate with each other and with the rest of the system. It is based on I<sup>2</sup>C principles of operation. SMBus provides a control bus for system and power management related tasks. A system may use SMBus to pass messages to and from devices instead of toggling individual control lines.

The System Management Bus Specification refers to three types of devices. A *slave* is a device that is receiving or responding to a command. A *master* is a device that issues commands, generates the clocks, and terminates the transfer. A *host* is a specialized master that provides the main interface to the system's CPU. A host must be a master-slave and must support the SMBus host notify protocol. Only one host is allowed in a system.

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## Similarities between SMBus and I<sup>2</sup>C

- 2 wire bus protocol (1 Clk, 1 Data) + SMBus Alert line optional
- Master-slave communication, Master provides clock
- Multi master capability
- SMBus data format similar to I<sup>2</sup>C 7-bit addressing format (Figure 217).

## Differences between SMBus and I<sup>2</sup>C

The following table describes the differences between SMBus and I<sup>2</sup>C.

Table 83. SMBus vs. I<sup>2</sup>C

SMBus	I <sup>2</sup> C
Max. speed 100 kHz	Max. speed 400 kHz
Min. clock speed 10 kHz	No minimum clock speed
35 ms clock low timeout	No timeout
Logic levels are fixed	Logic levels are V <sub>DD</sub> dependent
Different address types (reserved, dynamic etc.)	7-bit, 10-bit and general call slave address types
Different bus protocols (quick command, process call etc.)	No bus protocols

## SMBus application usage

With System Management Bus, a device can provide manufacturer information, tell the system what its model/part number is, save its state for a suspend event, report different types of errors, accept control parameters, and return its status. SMBus provides a control bus for system and power management related tasks.

#### **Device identification**

Any device that exists on the System Management Bus as a slave has a unique address called the Slave Address. For the list of reserved slave addresses, refer to the SMBus specification version. 2.0 (http://smbus.org/specs/).

## **Bus protocols**

The SMBus specification supports up to 9 bus protocols. For more details of these protocols and SMBus address types, refer to SMBus specification version. 2.0 (http://smbus.org/specs/). These protocols should be implemented by the user software.

## Address resolution protocol (ARP)

SMBus slave address conflicts can be resolved by dynamically assigning a new unique address to each slave device. The Address Resolution Protocol (ARP) has the following attributes:

- Address assignment uses the standard SMBus physical layer arbitration mechanism
- Assigned addresses remain constant while device power is applied; address retention through device power loss is also allowed
- No additional SMBus packet overhead is incurred after address assignment. (i.e. subsequent accesses to assigned slave addresses have the same overhead as accesses to fixed address devices.)
- Any SMBus master can enumerate the bus

## Unique device identifier (UDID)

In order to provide a mechanism to isolate each device for the purpose of address assignment, each device must implement a unique device identifier (UDID).

For the details on 128 bit UDID and more information on ARP, refer to SMBus specification version 2.0 (http://smbus.org/specs/).

#### SMBus alert mode

SMBus Alert is an optional signal with an interrupt line for devices that want to trade their ability to master for a pin. SMBA is a wired-AND signal just as the SCL and SDA signals are. SMBA is used in conjunction with the SMBus General Call Address. Messages invoked with the SMBus are 2 bytes long.

A slave-only device can signal the host through SMBA that it wants to talk by setting ALERT bit in I2C\_CR1 register. The host processes the interrupt and simultaneously accesses all SMBA devices through the Alert Response Address (known as ARA having a value 0001 100X). Only the device(s) which pulled SMBA low will acknowledge the Alert Response Address. This status is identified using SMBALERT Status flag in I2C\_SR1 register. The host performs a modified Receive Byte operation. The 7 bit device address provided by the slave transmit device is placed in the 7 most significant bits of the byte. The eighth bit can be a zero or one.

If more than one device pulls SMBA low, the highest priority (lowest address) device will win communication rights via standard arbitration during the slave address transfer. After acknowledging the slave address the device must disengage its SMBA pull-down. If the host still sees SMBA low when the message transfer is complete, it knows to read the ARA again.

A host which does not implement the SMBA signal may periodically access the ARA.

For more details on SMBus Alert mode, refer to SMBus specification version 2.0 (http://smbus.org/specs/).

#### Timeout error

There are differences in the timing specifications between I<sup>2</sup>C and SMBus. SMBus defines a clock low timeout, TIMEOUT of 35 ms. Also SMBus specifies TLOW: SEXT as the cumulative clock low extend time for a slave device. SMBus specifies TLOW: MEXT as the cumulative clock low extend time for a master device. For more details on these timeouts, refer to SMBus specification version 2.0 (http://smbus.org/specs/).

The status flag Timeout or Tlow Error in I2C\_SR1 shows the status of this feature.

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#### How to use the interface in SMBus mode

To switch from I<sup>2</sup>C mode to SMBus mode, the following sequence should be performed.

- Set the SMBus bit in the I2C\_CR1 register
- Configure the SMBTYPE and ENARP bits in the I2C\_CR1 register as required for the application

If you want to configure the device as a master, follow the Start condition generation procedure in *Section 23.3.3: I2C master mode*. Otherwise, follow the sequence in *Section 23.3.2: I2C slave mode*.

The application has to control the various SMBus protocols by software.

- SMB Device Default Address acknowledged if ENARP=1 and SMBTYPE=0
- SMB Host Header acknowledged if ENARP=1 and SMBTYPE=1
- SMB Alert Response Address acknowledged if SMBALERT=1

## 23.3.7 DMA requests

DMA requests (when enabled) are generated only for data transfer. DMA requests are generated by Data Register becoming empty in transmission and Data Register becoming full in reception. The DMA must be initialized and enabled before the I2C data transfer. The DMAEN bit must be set in the I2C\_CR2 register before the ADDR event. In master mode or in slave mode when clock stretching is enabled, the DMAEN bit can also be set during the ADDR event, before clearing the ADDR flag. The DMA request must be served before the end of the current byte transfer. When the number of data transfers which has been programmed for the corresponding DMA channel is reached, the DMA controller sends an End of Transfer EOT signal to the I<sup>2</sup>C interface and generates a Transfer Complete interrupt if enabled:

- Master transmitter: In the interrupt routine after the EOT interrupt, disable DMA requests then wait for a BTF event before programming the Stop condition.
- Master receiver
  - When the number of bytes to be received is equal to or greater than two, the DMA controller sends a hardware signal, EOT\_1, corresponding to the last but one data byte (number\_of\_bytes 1). If, in the I2C\_CR2 register, the LAST bit is set, I<sup>2</sup>C automatically sends a NACK after the next byte following EOT\_1. The user can generate a Stop condition in the DMA Transfer Complete interrupt routine if enabled.
  - When a single byte must be received: the NACK must be programmed during EV6 event, i.e. program ACK=0 when ADDR=1, before clearing ADDR flag. Then the user can program the STOP condition either after clearing ADDR flag, or in the DMA Transfer Complete interrupt routine.

## **Transmission using DMA**

DMA mode can be enabled for transmission by setting the DMAEN bit in the I2C\_CR2 register. Data will be loaded from a Memory area configured using the DMA peripheral (refer to the DMA specification) to the I2C\_DR register whenever the TxE bit is set. To map a DMA channel for I<sup>2</sup>C transmission, perform the following sequence. Here x is the channel number.

- 1. Set the I2C\_DR register address in the DMA\_CPARx register. The data will be moved to this address from the memory after each TxE event.
- 2. Set the memory address in the DMA\_CMARx register. The data will be loaded into I2C\_DR from this memory after each TxE event.
- 3. Configure the total number of bytes to be transferred in the DMA\_CNDTRx register. After each TxE event, this value will be decremented.
- 4. Configure the channel priority using the PL[0:1] bits in the DMA CCRx register
- 5. Set the DIR bit and, in the DMA\_CCRx register, configure interrupts after half transfer or full transfer depending on application requirements.
- 6. Activate the channel by setting the EN bit in the DMA\_CCRx register.

When the number of data transfers which has been programmed in the DMA Controller registers is reached, the DMA controller sends an End of Transfer EOT/ EOT\_1 signal to the I<sup>2</sup>C interface and the DMA generates an interrupt, if enabled, on the DMA channel interrupt vector.

Note: Do not enable the ITBUFEN bit in the I2C\_CR2 register if DMA is used for transmission.

## **Reception using DMA**

DMA mode can be enabled for reception by setting the DMAEN bit in the I2C\_CR2 register. Data will be loaded from the I2C\_DR register to a Memory area configured using the DMA peripheral (refer to the DMA specification) whenever a data byte is received. To map a DMA channel for I<sup>2</sup>C reception, perform the following sequence. Here x is the channel number.

- 1. Set the I2C\_DR register address in DMA\_CPARx register. The data will be moved from this address to the memory after each RxNE event.
- 2. Set the memory address in the DMA\_CMARx register. The data will be loaded from the I2C\_DR register to this memory area after each RxNE event.
- 3. Configure the total number of bytes to be transferred in the DMA\_CNDTRx register. After each RxNE event, this value will be decremented.
- 4. Configure the channel priority using the PL[0:1] bits in the DMA\_CCRx register
- 5. Reset the DIR bit and configure interrupts in the DMA\_CCRx register after half transfer or full transfer depending on application requirements.
- 6. Activate the channel by setting the EN bit in the DMA\_CCRx register.

When the number of data transfers which has been programmed in the DMA Controller registers is reached, the DMA controller sends an End of Transfer EOT/ EOT\_1 signal to the I<sup>2</sup>C interface and DMA generates an interrupt, if enabled, on the DMA channel interrupt vector.

Note: Do not enable the ITBUFEN bit in the I2C\_CR2 register if DMA is used for reception.



## 23.3.8 Packet error checking

A PEC calculator has been implemented to improve the reliability of communication. The PEC is calculated by using the  $C(x) = x^8 + x^2 + x + 1$  CRC-8 polynomial serially on each bit.

- PEC calculation is enabled by setting the ENPEC bit in the I2C\_CR1 register. PEC is a CRC-8 calculated on all message bytes including addresses and R/W bits.
  - In transmission: set the PEC transfer bit in the I2C\_CR1 register after the TxE event corresponding to the last byte. The PEC will be transferred after the last transmitted byte.
  - In reception: set the PEC bit in the I2C\_CR1 register after the RxNE event corresponding to the last byte so that the receiver sends a NACK if the next received byte is not equal to the internally calculated PEC. In case of Master-Receiver, a NACK must follow the PEC whatever the check result. The PEC must be set before the ACK of the CRC reception in slave mode. It must be set when the ACK is set low in master mode.
- A PECERR error flag/interrupt is also available in the I2C\_SR1 register.
- If DMA and PEC calculation are both enabled:-
  - In transmission: when the I<sup>2</sup>C interface receives an EOT signal from the DMA controller, it automatically sends a PEC after the last byte.
  - In reception: when the I<sup>2</sup>C interface receives an EOT\_1 signal from the DMA controller, it will automatically consider the next byte as a PEC and will check it. A DMA request is generated after PEC reception.
- To allow intermediate PEC transfers, a control bit is available in the I2C\_CR2 register (LAST bit) to determine if it is really the last DMA transfer or not. If it is the last DMA request for a master receiver, a NACK is automatically sent after the last received byte.
- PEC calculation is corrupted by an arbitration loss.

## 23.4 I<sup>2</sup>C interrupts

The table below gives the list of I<sup>2</sup>C interrupt requests.

Table 84. I<sup>2</sup>C Interrupt requests

Interrupt event	Event flag	Enable control bit
Start bit sent (Master)	SB	
Address sent (Master) or Address matched (Slave)	ADDR	
10-bit header sent (Master)	ADD10	ITEVFEN
Stop received (Slave)	STOPF	
Data byte transfer finished	BTF	
Receive buffer not empty	RxNE	ITEVFEN and ITBUFEN
Transmit buffer empty	TxE	TIEVPEN AND TIBUPEN

Table 84. I<sup>2</sup>C Interrupt requests (continued)

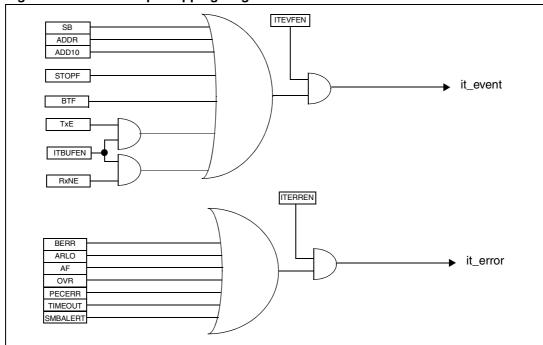
Interrupt event	Event flag	Enable control bit
Bus error	BERR	
Arbitration loss (Master)	ARLO	
Acknowledge failure	AF	
Overrun/Underrun	OVR	ITERREN
PEC error	PECERR	
Timeout/Tlow error	TIMEOUT	
SMBus Alert	SMBALERT	

Note:

SB, ADDR, ADD10, STOPF, BTF, RxNE and TxE are logically ORed on the same interrupt channel.

BERR, ARLO, AF, OVR, PECERR, TIMEOUT and SMBALERT are logically ORed on the same interrupt channel.

Figure 223. I<sup>2</sup>C interrupt mapping diagram



## 23.5 I<sup>2</sup>C debug mode

When the microcontroller enters the debug mode (Cortex<sup>TM</sup>-M4F core halted), the SMBUS timeout either continues to work normally or stops, depending on the DBG\_I2Cx\_SMBUS\_TIMEOUT configuration bits in the DBG module. For more details, refer to Section 32.16.2: Debug support for timers, watchdog, bxCAN and I2C on page 1316.

## 23.6 I<sup>2</sup>C registers

Refer to Section 1.1 on page 46 for a list of abbreviations used in register descriptions.

The peripheral registers can be accessed by half-words (16 bits) or words (32 bits).

## 23.6.1 I<sup>2</sup>C Control register 1 (I2C\_CR1)

Address offset: 0x00 Reset value: 0x0000

15	14	13	12	11	10	9	8	/	ь	5	4	3	2	1	U	
SWRST	Res.	ALERT	PEC	POS	ACK	STOP	START	NO STRETCH	ENGC	ENPEC	ENARP	SMB TYPE	Res.	SMBUS	PE	
rw		rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw		rw	rw	

#### Bit 15 SWRST: Software reset

When set, the I2C is under reset state. Before resetting this bit, make sure the I2C lines are released and the bus is free.

- 0: I2C Peripheral not under reset
- 1: I<sup>2</sup>C Peripheral under reset state

Note: This bit can be used to reinitialize the peripheral after an error or a locked state. As an example, if the BUSY bit is set and remains locked due to a glitch on the bus, the SWRST bit can be used to exit from this state.

## Bit 14 Reserved, must be kept at reset value

### Bit 13 ALERT: SMBus alert

This bit is set and cleared by software, and cleared by hardware when PE=0.

- 0: Releases SMBA pin high. Alert Response Address Header followed by NACK.
- 1: Drives SMBA pin low. Alert Response Address Header followed by ACK.

#### Bit 12 PEC: Packet error checking

This bit is set and cleared by software, and cleared by hardware when PEC is transferred or by a START or Stop condition or when PE=0.

- 0: No PEC transfer
- 1: PEC transfer (in Tx or Rx mode)

Note: PEC calculation is corrupted by an arbitration loss.

#### Bit 11 POS: Acknowledge/PEC Position (for data reception)

This bit is set and cleared by software and cleared by hardware when PE=0.

0: ACK bit controls the (N)ACK of the current byte being received in the shift register. The PEC bit indicates that current byte in shift register is a PEC.

1: ACK bit controls the (N)ACK of the next byte which will be received in the shift register. The PEC bit indicates that the next byte in the shift register is a PEC

Note: The POS bit must be used only in 2-byte reception configuration in master mode. It must be configured before data reception starts, as described in the 2-byte reception procedure recommended in Section: Master receiver on page 600.

#### Bit 10 ACK: Acknowledge enable

This bit is set and cleared by software and cleared by hardware when PE=0.

- 0: No acknowledge returned
- 1: Acknowledge returned after a byte is received (matched address or data)

#### Bit 9 STOP: Stop generation

The bit is set and cleared by software, cleared by hardware when a Stop condition is detected, set by hardware when a timeout error is detected.

- In Master Mode:
- 0: No Stop generation.
- 1: Stop generation after the current byte transfer or after the current Start condition is sent.
- In Slave mode:
- 0: No Stop generation.
- 1: Release the SCL and SDA lines after the current byte transfer.

#### Bit 8 START: Start generation

This bit is set and cleared by software and cleared by hardware when start is sent or PE=0.

- In Master Mode:
- 0: No Start generation
- 1: Repeated start generation
- In Slave mode:
- 0: No Start generation
- 1: Start generation when the bus is free

### Bit 7 NOSTRETCH: Clock stretching disable (Slave mode)

This bit is used to disable clock stretching in slave mode when ADDR or BTF flag is set, until it is reset by software.

- 0: Clock stretching enabled
- 1: Clock stretching disabled

#### Bit 6 ENGC: General call enable

- 0: General call disabled. Address 00h is NACKed.
- 1: General call enabled, Address 00h is ACKed.

## Bit 5 ENPEC: PEC enable

- 0: PEC calculation disabled
- 1: PEC calculation enabled

#### Bit 4 ENARP: ARP enable

- 0: ARP disable
- 1: ARP enable

SMBus Device default address recognized if SMBTYPE=0

SMBus Host address recognized if SMBTYPE=1

Bit 3 SMBTYPE: SMBus type

0: SMBus Device

1: SMBus Host

Bit 2 Reserved, must be kept at reset value

Bit 1 SMBUS: SMBus mode

0: I<sup>2</sup>C mode

1: SMBus mode

Bit 0 PE: Peripheral enable

0: Peripheral disable

1: Peripheral enable

Note: If this bit is reset while a communication is on going, the peripheral is disabled at the end of the current communication, when back to IDLE state.

All bit resets due to PE=0 occur at the end of the communication.

In master mode, this bit must not be reset before the end of the communication.

Note:

When the STOP, START or PEC bit is set, the software must not perform any write access to I2C\_CR1 before this bit is cleared by hardware. Otherwise there is a risk of setting a second STOP, START or PEC request.

## 23.6.2 I<sup>2</sup>C Control register 2 (I2C\_CR2)

Address offset: 0x04 Reset value: 0x0000

15	14	13	12	11	10	9	8	/	6	5	4	3	2	1	0
	Reserved		LAST	DMA EN	ITBUF EN	ITEVT EN	ITERR EN	Rese	rved			FREG	Q[5:0]		
			rw	rw	rw	rw	rw			rw	rw	rw	rw	rw	rw

Bits 15:13 Reserved, must be kept at reset value

Bit 12 LAST: DMA last transfer

0: Next DMA EOT is not the last transfer

1: Next DMA EOT is the last transfer

Note: This bit is used in master receiver mode to permit the generation of a NACK on the last received data.

Bit 11 DMAEN: DMA requests enable

0: DMA requests disabled

1: DMA request enabled when TxE=1 or RxNE =1

Bit 10 ITBUFEN: Buffer interrupt enable

0: TxE = 1 or RxNE = 1 does not generate any interrupt.

1:TxE = 1 or RxNE = 1 generates Event Interrupt (whatever the state of DMAEN)

#### Bit 9 ITEVTEN: Event interrupt enable

- 0: Event interrupt disabled
- 1: Event interrupt enabled

This interrupt is generated when:

- -SB = 1 (Master)
- -ADDR = 1 (Master/Slave)
- -ADD10= 1 (Master)
- -STOPF = 1 (Slave)
- -BTF = 1 with no TxE or RxNE event
- -TxE event to 1 if ITBUFEN = 1
- -RxNE event to 1if ITBUFEN = 1

# Bit 8 ITERREN: Error interrupt enable

- 0: Error interrupt disabled
- 1: Error interrupt enabled

This interrupt is generated when:

- BERR = 1
- ARLO = 1
- AF = 1
- OVR = 1
- PECERR = 1
- TIMEOUT = 1
- SMBALERT = 1

Bits 7:6 Reserved, must be kept at reset value

#### Bits 5:0 FREQ[5:0]: Peripheral clock frequency

The peripheral clock frequency must be configured using the input APB clock frequency (I2C peripheral connected to APB). The minimum allowed frequency is 2 MHz, the maximum frequency is limited by the maximum APB frequency (42 MHz) and an intrinsic limitation of 46 MHz.

0b000000: Not allowed 0b000001: Not allowed 0b000010: 2 MHz

...

0b101010: 42MHz

Higher than 0b101010: Not allowed

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# 23.6.3 I<sup>2</sup>C Own address register 1 (I2C\_OAR1)

Address offset: 0x08 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ADD MODE			Reserved	l		ADD	[9:8]			A	\DD[7:1]				ADD0
rw						rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

## Bit 15 **ADDMODE** Addressing mode (slave mode)

0: 7-bit slave address (10-bit address not acknowledged)1: 10-bit slave address (7-bit address not acknowledged)

Bit 14 Should always be kept at 1 by software.

Bits 13:10 Reserved, must be kept at reset value

Bits 9:8 ADD[9:8]: Interface address

7-bit addressing mode: don't care

10-bit addressing mode: bits9:8 of address

Bits 7:1 ADD[7:1]: Interface address

bits 7:1 of address

Bit 0 ADD0: Interface address

7-bit addressing mode: don't care 10-bit addressing mode: bit 0 of address

# 23.6.4 I<sup>2</sup>C Own address register 2 (I2C\_OAR2)

Address offset: 0x0C Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			Rese	nuod						P	ADD2[7:1]	]			ENDUAL
			nese	rveu				rw	rw	rw	rw	rw	rw	rw	rw

Bits 15:8 Reserved, must be kept at reset value

Bits 7:1 ADD2[7:1]: Interface address

bits 7:1 of address in dual addressing mode

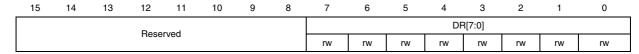
Bit 0 ENDUAL: Dual addressing mode enable

0: Only OAR1 is recognized in 7-bit addressing mode

1: Both OAR1 and OAR2 are recognized in 7-bit addressing mode

# 23.6.5 I<sup>2</sup>C Data register (I2C\_DR)

Address offset: 0x10 Reset value: 0x0000



Bits 15:8 Reserved, must be kept at reset value

# Bits 7:0 DR[7:0] 8-bit data register

Byte received or to be transmitted to the bus.

- -Transmitter mode: Byte transmission starts automatically when a byte is written in the DR register. A continuous transmit stream can be maintained if the next data to be transmitted is put in DR once the transmission is started (TxE=1)
- -Receiver mode: Received byte is copied into DR (RxNE=1). A continuous transmit stream can be maintained if DR is read before the next data byte is received (RxNE=1).

Note: In slave mode, the address is not copied into DR.

Note: Write collision is not managed (DR can be written if TxE=0).

Note: If an ARLO event occurs on ACK pulse, the received byte is not copied into DR and so

cannot be read.

# 23.6.6 I<sup>2</sup>C Status register 1 (I2C\_SR1)

Address offset: 0x14 Reset value: 0x0000

15	14	13	12	11	10	9	8	/	ь	5	4	3	2	ı	U
SMB ALERT	TIME OUT	Res.	PEC ERR	OVR	AF	ARLO	BERR	TxE	RxNE	Res.	STOPF	ADD10	BTF	ADDR	SB
rc_w0	rc_w0		rc_w0	rc_w0	rc_w0	rc_w0	rc_w0	r	r		r	r	r	r	r

### Bit 15 SMBALERT: SMBus alert

In SMBus host mode:

0: no SMBALERT

1: SMBALERT event occurred on pin

In SMBus slave mode:

0: no SMBALERT response address header

- 1: SMBALERT response address header to SMBALERT LOW received
- Cleared by software writing 0, or by hardware when PE=0.

#### Bit 14 TIMEOUT: Timeout or Tlow error

0: No timeout error

1: SCL remained LOW for 25 ms (Timeout)

or

Master cumulative clock low extend time more than 10 ms (Tlow:mext)

or

Slave cumulative clock low extend time more than 25 ms (Tlow:sext)

- When set in slave mode: slave resets the communication and lines are released by hardware
- When set in master mode: Stop condition sent by hardware
- Cleared by software writing 0, or by hardware when PE=0.

Note: This functionality is available only in SMBus mode.

#### Bit 13 Reserved, must be kept at reset value

#### Bit 12 PECERR: PEC Error in reception

- 0: no PEC error: receiver returns ACK after PEC reception (if ACK=1)
- 1: PEC error: receiver returns NACK after PEC reception (whatever ACK)
- -Cleared by software writing 0, or by hardware when PE=0.
- -Note: When the received CRC is wrong, PECERR is not set in slave mode if the PEC control bit is not set before the end of the CRC reception. Nevertheless, reading the PEC value determines whether the received CRC is right or wrong.

#### Bit 11 OVR: Overrun/Underrun

- 0: No overrun/underrun
- 1: Overrun or underrun
- -Set by hardware in slave mode when NOSTRETCH=1 and:
- -In reception when a new byte is received (including ACK pulse) and the DR register has not been read yet. New received byte is lost.
- -In transmission when a new byte should be sent and the DR register has not been written yet. The same byte is sent twice.
- -Cleared by software writing 0, or by hardware when PE=0.

Note: If the DR write occurs very close to SCL rising edge, the sent data is unspecified and a hold timing error occurs

#### Bit 10 AF: Acknowledge failure

- 0: No acknowledge failure
- 1: Acknowledge failure
- -Set by hardware when no acknowledge is returned.
- -Cleared by software writing 0, or by hardware when PE=0.

#### Bit 9 ARLO: Arbitration lost (master mode)

- 0: No Arbitration Lost detected
- 1: Arbitration Lost detected
- Set by hardware when the interface loses the arbitration of the bus to another master
- -Cleared by software writing 0, or by hardware when PE=0.
- After an ARLO event the interface switches back automatically to Slave mode (M/SL=0).

Note: In SMBUS, the arbitration on the data in slave mode occurs only during the data phase, or the acknowledge transmission (not on the address acknowledge).

#### Bit 8 BERR: Bus error

- 0: No misplaced Start or Stop condition
- 1: Misplaced Start or Stop condition
- -Set by hardware when the interface detects an SDA rising or falling edge while SCL is high, occurring in a non-valid position during a byte transfer.
- -Cleared by software writing 0, or by hardware when PE=0.

#### Bit 7 TxE: Data register empty (transmitters)

- 0: Data register not empty
- 1: Data register empty
- -Set when DR is empty in transmission. TxE is not set during address phase.
- -Cleared by software writing to the DR register or by hardware after a start or a stop condition or when PE=0.

TxE is not set if either a NACK is received, or if next byte to be transmitted is PEC (PEC=1)

Note: TxE is not cleared by writing the first data being transmitted, or by writing data when BTF is set, as in both cases the data register is still empty.

Bit 6 RxNE: Data register not empty (receivers)

- 0: Data register empty
- 1: Data register not empty
- -Set when data register is not empty in receiver mode. RxNE is not set during address phase.
- -Cleared by software reading or writing the DR register or by hardware when PE=0.

RxNE is not set in case of ARLO event.

Note: RxNE is not cleared by reading data when BTF is set, as the data register is still full.

Bit 5 Reserved, must be kept at reset value

Bit 4 STOPF: Stop detection (slave mode)

- 0: No Stop condition detected
- 1: Stop condition detected
- -Set by hardware when a Stop condition is detected on the bus by the slave after an acknowledge (if ACK=1).
- -Cleared by software reading the SR1 register followed by a write in the CR1 register, or by hardware when PE=0

Note: The STOPF bit is not set after a NACK reception.

It is recommended to perform the complete clearing sequence (READ SR1 then WRITE CR1) after the STOPF is set. Refer to Figure 220: Transfer sequence diagram for slave receiver on page 597.

Bit 3 ADD10: 10-bit header sent (Master mode)

- 0: No ADD10 event occurred.
- 1: Master has sent first address byte (header).
- -Set by hardware when the master has sent the first byte in 10-bit address mode.
- -Cleared by software reading the SR1 register followed by a write in the DR register of the second address byte, or by hardware when PE=0.

Note: ADD10 bit is not set after a NACK reception

Bit 2 BTF: Byte transfer finished

- 0: Data byte transfer not done
- 1: Data byte transfer succeeded
- -Set by hardware when NOSTRETCH=0 and:
  - -In reception when a new byte is received (including ACK pulse) and DR has not been read yet (RxNE=1).
  - -In transmission when a new byte should be sent and DR has not been written yet (TxE=1).
  - -Cleared by software by either a read or write in the DR register or by hardware after a start or a stop condition in transmission or when PE=0.

Note: The BTF bit is not set after a NACK reception

The BTF bit is not set if next byte to be transmitted is the PEC (TRA=1 in I2C\_SR2 register and PEC=1 in I2C\_CR1 register)

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# Bit 1 ADDR: Address sent (master mode)/matched (slave mode)

This bit is cleared by software reading SR1 register followed reading SR2, or by hardware when PE=0.

#### Address matched (Slave)

- 0: Address mismatched or not received.
- 1: Received address matched.
- -Set by hardware as soon as the received slave address matched with the OAR registers content or a general call or a SMBus Device Default Address or SMBus Host or SMBus Alert is recognized. (when enabled depending on configuration).

Note: In slave mode, it is recommended to perform the complete clearing sequence (READ SR1 then READ SR2) after ADDR is set. Refer to Figure 220: Transfer sequence diagram for slave receiver on page 597.

#### Address sent (Master)

- 0: No end of address transmission
- 1: End of address transmission
- -For 10-bit addressing, the bit is set after the ACK of the 2nd byte.
- -For 7-bit addressing, the bit is set after the ACK of the byte.

Note: ADDR is not set after a NACK reception

#### Bit 0 SB: Start bit (Master mode)

- 0: No Start condition
- 1: Start condition generated.
- -Set when a Start condition generated.
- -Cleared by software by reading the SR1 register followed by writing the DR register, or by hardware when PE=0

# 23.6.7 I<sup>2</sup>C Status register 2 (I2C\_SR2)

Address offset: 0x18 Reset value: 0x0000

Note:

Reading I2C\_SR2 after reading I2C\_SR1 clears the ADDR flag, even if the ADDR flag was set after reading I2C\_SR1. Consequently, I2C\_SR2 must be read only when ADDR is found set in I2C\_SR1 or when the STOPF bit is cleared.

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			PEC	[7:0]				DUALF	SMB HOST	SMBDE FAULT	GEN CALL	Res.	TRA	BUSY	MSL
r	r	r	r	r	r	r	r	r	r	r	r		r	r	r

#### Bits 15:8 PEC[7:0] Packet error checking register

This register contains the internal PEC when ENPEC=1.

#### Bit 7 **DUALF**: Dual flag (Slave mode)

- 0: Received address matched with OAR1
- 1: Received address matched with OAR2
- -Cleared by hardware after a Stop condition or repeated Start condition, or when PE=0.

### Bit 6 SMBHOST: SMBus host header (Slave mode)

- 0: No SMBus Host address
- 1: SMBus Host address received when SMBTYPE=1 and ENARP=1.
- -Cleared by hardware after a Stop condition or repeated Start condition, or when PE=0.

#### Bit 5 SMBDEFAULT: SMBus device default address (Slave mode)

- 0: No SMBus Device Default address
- 1: SMBus Device Default address received when ENARP=1
- -Cleared by hardware after a Stop condition or repeated Start condition, or when PE=0.

### Bit 4 GENCALL: General call address (Slave mode)

- 0: No General Call
- 1: General Call Address received when ENGC=1
- -Cleared by hardware after a Stop condition or repeated Start condition, or when PE=0.

#### Bit 3 Reserved, must be kept at reset value

#### Bit 2 TRA: Transmitter/receiver

- 0: Data bytes received
- 1: Data bytes transmitted

This bit is set depending on the R/W bit of the address byte, at the end of total address phase.

It is also cleared by hardware after detection of Stop condition (STOPF=1), repeated Start condition, loss of bus arbitration (ARLO=1), or when PE=0.

#### Bit 1 BUSY: Bus busy

- 0: No communication on the bus
- 1: Communication ongoing on the bus
- -Set by hardware on detection of SDA or SCL low
- -cleared by hardware on detection of a Stop condition.

It indicates a communication in progress on the bus. This information is still updated when the interface is disabled (PE=0).

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Bit 0 MSL: Master/slave

- 0: Slave Mode
- 1: Master Mode
- -Set by hardware as soon as the interface is in Master mode (SB=1).
- -Cleared by hardware after detecting a Stop condition on the bus or a loss of arbitration (ARLO=1), or by hardware when PE=0.

Note: Reading I2C\_SR2 after reading I2C\_SR1 clears the ADDR flag, even if the ADDR flag was

set after reading I2C\_SR1. Consequently, I2C\_SR2 must be read only when ADDR is found

set in I2C SR1 or when the STOPF bit is cleared.

# 23.6.8 I<sup>2</sup>C Clock control register (I2C\_CCR)

Address offset: 0x1C Reset value: 0x0000

Note:  $f_{PCLK1}$  must be at least 2 MHz to achieve standard mode  $I^2C$  frequencies. It must be at least

4 MHz to achieve fast mode I<sup>2</sup>C frequencies. It must be a multiple of 10MHz to reach the

400 kHz maximum I<sup>2</sup>C fast mode clock.

The CCR register must be configured only when the I2C is disabled (PE = 0).

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
F/S	DUTY	Reser	nuod.		_	_	_		CC	CR[11:0]			_		
rw	rw	nesei	veu	rw	rw	rw	rw	rw	rw						

Bit 15 F/S: I2C master mode selection

0: Standard Mode I2C

1: Fast Mode I2C

Bit 14 DUTY: Fast mode duty cycle

0: Fast Mode  $t_{low}/t_{high} = 2$ 

1: Fast Mode t<sub>low</sub>/t<sub>high</sub> = 16/9 (see CCR)

Bits 13:12 Reserved, must be kept at reset value

# Bits 11:0 CCR[11:0]: Clock control register in Fast/Standard mode (Master mode)

Controls the SCL clock in master mode.

Standard mode or SMBus:

$$T_{high} = CCR * T_{PCLK1}$$
  
 $T_{low} = CCR * T_{PCLK1}$ 

Fast mode:

If DUTY = 0:

 $T_{high} = CCR * T_{PCLK1}$ 

 $T_{low} = 2 * CCR * T_{PCLK1}$ 

If DUTY = 1: (to reach 400 kHz)

 $T_{high} = 9 * CCR * T_{PCLK1}$ 

 $T_{low} = 16 * CCR * T_{PCLK1}$ 

For instance: in standard mode, to generate a 100 kHz SCL frequency:

If FREQR = 08,  $T_{PCLK1}$  = 125 ns so CCR must be programmed with 0x28

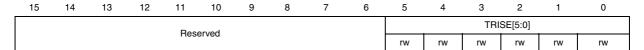
 $(0x28 \le 40d x 125 ns = 5000 ns.)$ 

Note: 1. The minimum allowed value is 0x04, except in FAST DUTY mode where the minimum allowed value is 0x01

- . These timings are without filters.
- . The CCR register must be configured only when the  $l^2C$  is disabled (PE = 0).

# 23.6.9 I<sup>2</sup>C TRISE register (I2C\_TRISE)

Address offset: 0x20 Reset value: 0x0002



Bits 15:6 Reserved, must be kept at reset value

Bits 5:0 TRISE[5:0]: Maximum rise time in Fast/Standard mode (Master mode)

These bits must be programmed with the maximum SCL rise time given in the  $I^2C$  bus specification, incremented by 1.

For instance: in standard mode, the maximum allowed SCL rise time is 1000 ns.

If, in the I2C\_CR2 register, the value of FREQ[5:0] bits is equal to 0x08 and  $T_{PCLK1}$  = 125 ns therefore the TRISE[5:0] bits must be programmed with 09h.

(1000 ns / 125 ns = 8 + 1)

The filter value can also be added to TRISE[5:0].

If the result is not an integer, TRISE[5:0] must be programmed with the integer part, in order to respect the t<sub>HIGH</sub> parameter.

Note: TRISE[5:0] must be configured only when the I2C is disabled (PE = 0).

# 23.6.10 I<sup>2</sup>C register map

The table below provides the I<sup>2</sup>C register map and reset values.

Table 85. I<sup>2</sup>C register map and reset values

Offset	Register	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	15	14	13	12	11	10	6	8	7	9	2	4	3	7	-	0
0x00	I2C_CR1  Reset value	Reserved	o SWRST	Reserved	o ALERT	o PEC	o POS	o ACK	o STOP	o START	ONOSTRETCH	o ENGC	o ENPEC	o ENARP	o SMBTYPE	Reserved	o SMBUS	o PE
0x04	I2C_CR2  Reset value	Reserved	<u> </u>		J			O ITBUFEN o	OITEVTEN	OITERREN	Doggana		0	F	REC	Q[5:0	)]	0
0x08	I2C_OAR1	Reserved	ADDMODE	Re	ser					)[9:8]			ΑС	)D[7	':1]			ADD0
0x0C	I2C_OAR2  Reset value	Reserved	<u> </u>						0	0	0	0	O AD	0 D2[	7:1] 0	0	0	o ENDUAL o
0x10	I2C_DR  Reset value	Reserved									0	0			7:0]		0	0
0x14	I2C_SR1	Reserved	O SMBALERT	TIMEOUT	Reserved		OVR	AF	ARLO	BERR	TxE	RXNE		STOPF	ADD10	BTF	ADDR	SB
0x18	Reset value	Reserved	0	0		0 PE	0 C[7	(:0]	0	0	DUALF		S		Reserved		_	MSL
0x1C	I2C_CCR Reset value	Reserved	o F/S o	o DUTY o	Dovraga	0	0	0	0	0	0	0 R[1	1:0]	0	0	0	0	0
0x20	I2C_TRISE  Reset value	Reserved									'		0	TI	RISI 0	E[5:		0

Refer to *Table 2 on page 50* for the register boundary addresses table.

# 24 Universal synchronous asynchronous receiver transmitter (USART)

This section applies to the whole STM32F4xx family, unless otherwise specified.

# 24.1 USART introduction

The universal synchronous asynchronous receiver transmitter (USART) offers a flexible means of full-duplex data exchange with external equipment requiring an industry standard NRZ asynchronous serial data format. The USART offers a very wide range of baud rates using a fractional baud rate generator.

It supports synchronous one-way communication and half-duplex single wire communication. It also supports the LIN (local interconnection network), Smartcard Protocol and IrDA (infrared data association) SIR ENDEC specifications, and modem operations (CTS/RTS). It allows multiprocessor communication.

High speed data communication is possible by using the DMA for multibuffer configuration.

# 24.2 USART main features

- Full duplex, asynchronous communications
- NRZ standard format (Mark/Space)
- Configurable oversampling method by 16 or by 8 to give flexibility between speed and clock tolerance
- Fractional baud rate generator systems
  - Common programmable transmit and receive baud rate (refer to the datasheets for the value of the baud rate at the maximum APB frequency.
- Programmable data word length (8 or 9 bits)
- Configurable stop bits support for 1 or 2 stop bits
- LIN Master Synchronous Break send capability and LIN slave break detection capability
  - 13-bit break generation and 10/11 bit break detection when USART is hardware configured for LIN
- Transmitter clock output for synchronous transmission
- IrDA SIR encoder decoder
  - Support for 3/16 bit duration for normal mode
- Smartcard emulation capability
  - The Smartcard interface supports the asynchronous protocol Smartcards as defined in the ISO 7816-3 standards
  - 0.5, 1.5 stop bits for Smartcard operation
- Single-wire half-duplex communication
- Configurable multibuffer communication using DMA (direct memory access)
  - Buffering of received/transmitted bytes in reserved SRAM using centralized DMA
- Separate enable bits for transmitter and receiver

- Transfer detection flags:
  - Receive buffer full
  - Transmit buffer empty
  - End of transmission flags
- Parity control:
  - Transmits parity bit
  - Checks parity of received data byte
- Four error detection flags:
  - Overrun error
  - Noise detection
  - Frame error
  - Parity error
- Ten interrupt sources with flags:
  - CTS changes
  - LIN break detection
  - Transmit data register empty
  - Transmission complete
  - Receive data register full
  - Idle line received
  - Overrun error
  - Framing error
  - Noise error
  - Parity error
- Multiprocessor communication enter into mute mode if address match does not occur
- Wake up from mute mode (by idle line detection or address mark detection)
- Two receiver wakeup modes: Address bit (MSB, 9<sup>th</sup> bit), Idle line

# 24.3 USART functional description

The interface is externally connected to another device by three pins (see *Figure 224*). Any USART bidirectional communication requires a minimum of two pins: Receive Data In (RX) and Transmit Data Out (TX):

**RX:** Receive Data Input is the serial data input. Oversampling techniques are used for data recovery by discriminating between valid incoming data and noise.

**TX:** Transmit Data Output. When the transmitter is disabled, the output pin returns to its I/O port configuration. When the transmitter is enabled and nothing is to be transmitted, the TX pin is at high level. In single-wire and smartcard modes, this I/O is used to transmit and receive the data (at USART level, data are then received on SW\_RX).

Through these pins, serial data is transmitted and received in normal USART mode as frames comprising:

- An Idle Line prior to transmission or reception
- A start bit
- A data word (8 or 9 bits) least significant bit first
- 0.5,1, 1.5, 2 Stop bits indicating that the frame is complete
- This interface uses a fractional baud rate generator with a 12-bit mantissa and 4-bit fraction
- A status register (USART\_SR)
- Data Register (USART\_DR)
- A baud rate register (USART\_BRR) 12-bit mantissa and 4-bit fraction.
- A Guardtime Register (USART\_GTPR) in case of Smartcard mode.

Refer to Section 24.6: USART registers on page 663 for the definitions of each bit.

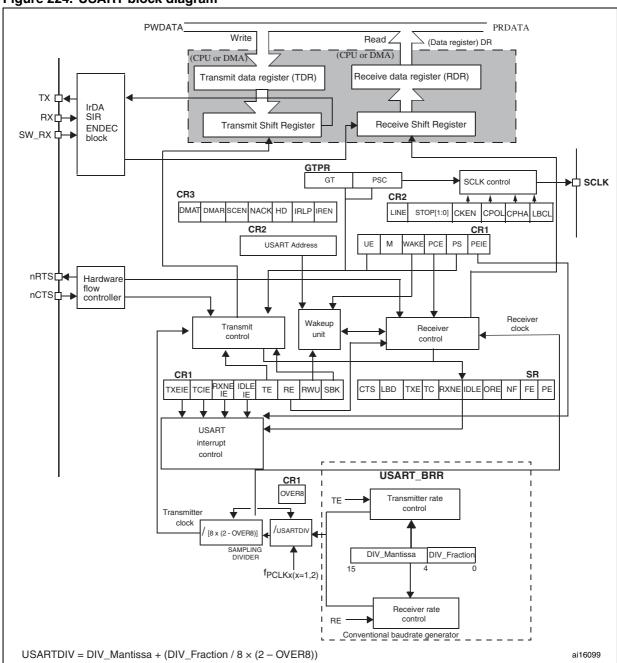
The following pin is required to interface in synchronous mode:

SCLK: Transmitter clock output. This pin outputs the transmitter data clock for synchronous transmission corresponding to SPI master mode (no clock pulses on start bit and stop bit, and a software option to send a clock pulse on the last data bit). In parallel data can be received synchronously on RX. This can be used to control peripherals that have shift registers (e.g. LCD drivers). The clock phase and polarity are software programmable. In smartcard mode, SCLK can provide the clock to the smartcard.

The following pins are required in Hardware flow control mode:

- nCTS: Clear To Send blocks the data transmission at the end of the current transfer when high
- nRTS: Request to send indicates that the USART is ready to receive a data (when low).

Figure 224. USART block diagram



# 24.3.1 USART character description

Word length may be selected as being either 8 or 9 bits by programming the M bit in the USART\_CR1 register (see *Figure 225*).

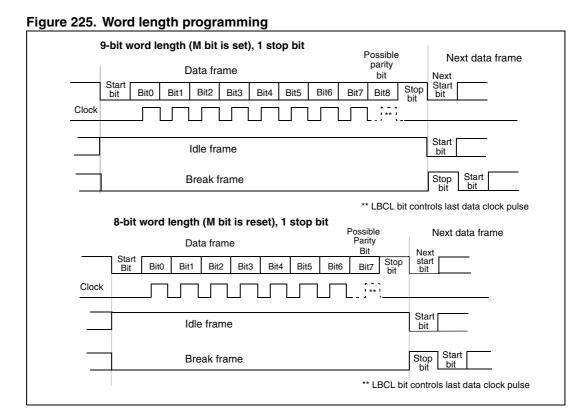
The TX pin is in low state during the start bit. It is in high state during the stop bit.

An *Idle character* is interpreted as an entire frame of "1"s followed by the start bit of the next frame which contains data (The number of "1" 's will include the number of stop bits).

A *Break character* is interpreted on receiving "0"s for a frame period. At the end of the break frame the transmitter inserts either 1 or 2 stop bits (logic "1" bit) to acknowledge the start bit.

Transmission and reception are driven by a common baud rate generator, the clock for each is generated when the enable bit is set respectively for the transmitter and receiver.

The details of each block is given below.



#### 24.3.2 Transmitter

The transmitter can send data words of either 8 or 9 bits depending on the M bit status. When the transmit enable bit (TE) is set, the data in the transmit shift register is output on the TX pin and the corresponding clock pulses are output on the SCLK pin.

#### Character transmission

During an USART transmission, data shifts out least significant bit first on the TX pin. In this mode, the USART\_DR register consists of a buffer (TDR) between the internal bus and the transmit shift register (see *Figure 224*).

Every character is preceded by a start bit which is a logic level low for one bit period. The character is terminated by a configurable number of stop bits.

The following stop bits are supported by USART: 0.5, 1, 1.5 and 2 stop bits.

Note:

The TE bit should not be reset during transmission of data. Resetting the TE bit during the transmission will corrupt the data on the TX pin as the baud rate counters will get frozen. The current data being transmitted will be lost.

An idle frame will be sent after the TE bit is enabled.

# Configurable stop bits

The number of stop bits to be transmitted with every character can be programmed in Control register 2, bits 13,12.

- 1 stop bit: This is the default value of number of stop bits.
- 2 Stop bits: This will be supported by normal USART, single-wire and modem modes.
- 0.5 stop bit: To be used when receiving data in Smartcard mode.
- 1.5 stop bits: To be used when transmitting and receiving data in Smartcard mode.

An idle frame transmission will include the stop bits.

A break transmission will be 10 low bits followed by the configured number of stop bits (when m=0) and 11 low bits followed by the configured number of stop bits (when m=1). It is not possible to transmit long breaks (break of length greater than 10/11 low bits).

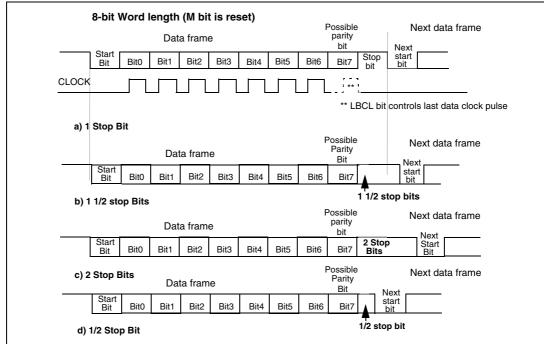


Figure 226. Configurable stop bits

#### Procedure:

- 1. Enable the USART by writing the UE bit in USART\_CR1 register to 1.
- 2. Program the M bit in USART CR1 to define the word length.
- 3. Program the number of stop bits in USART CR2.
- 4. Select DMA enable (DMAT) in USART\_CR3 if Multi buffer Communication is to take place. Configure the DMA register as explained in multibuffer communication.
- 5. Select the desired baud rate using the USART\_BRR register.
- 6. Set the TE bit in USART\_CR1 to send an idle frame as first transmission.
- 7. Write the data to send in the USART\_DR register (this clears the TXE bit). Repeat this for each data to be transmitted in case of single buffer.
- 8. After writing the last data into the USART\_DR register, wait until TC=1. This indicates that the transmission of the last frame is complete. This is required for instance when the USART is disabled or enters the Halt mode to avoid corrupting the last transmission.

## Single byte communication

Clearing the TXE bit is always performed by a write to the data register.

The TXE bit is set by hardware and it indicates:

- The data has been moved from TDR to the shift register and the data transmission has started.
- The TDR register is empty.
- The next data can be written in the USART\_DR register without overwriting the previous data.

This flag generates an interrupt if the TXEIE bit is set.

When a transmission is taking place, a write instruction to the USART\_DR register stores the data in the TDR register and which is copied in the shift register at the end of the current transmission.

When no transmission is taking place, a write instruction to the USART\_DR register places the data directly in the shift register, the data transmission starts, and the TXE bit is immediately set.

If a frame is transmitted (after the stop bit) and the TXE bit is set, the TC bit goes high. An interrupt is generated if the TCIE bit is set in the USART\_CR1 register.

After writing the last data into the USART\_DR register, it is mandatory to wait for TC=1 before disabling the USART or causing the microcontroller to enter the low power mode (see *Figure 227: TC/TXE behavior when transmitting*).

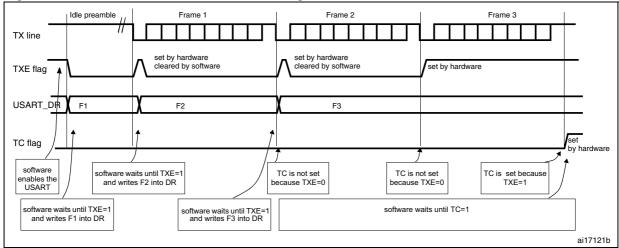
The TC bit is cleared by the following software sequence:

- 1. A read from the USART\_SR register
- 2. A write to the USART\_DR register

Note:

The TC bit can also be cleared by writing a '0 to it. This clearing sequence is recommended only for Multibuffer communication.





# **Break characters**

Setting the SBK bit transmits a break character. The break frame length depends on the M bit (see *Figure 225*).

If the SBK bit is set to '1 a break character is sent on the TX line after completing the current character transmission. This bit is reset by hardware when the break character is completed (during the stop bit of the break character). The USART inserts a logic 1 bit at the end of the last break frame to guarantee the recognition of the start bit of the next frame.

Note:

If the software resets the SBK bit before the commencement of break transmission, the break character will not be transmitted. For two consecutive breaks, the SBK bit should be set after the stop bit of the previous break.

## Idle characters

Setting the TE bit drives the USART to send an idle frame before the first data frame.

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# 24.3.3 Receiver

The USART can receive data words of either 8 or 9 bits depending on the M bit in the USART\_CR1 register.

#### Start bit detection

The start bit detection sequence is the same when oversampling by 16 or by 8.

In the USART, the start bit is detected when a specific sequence of samples is recognized. This sequence is: 1 1 1 0 X 0 X 0 X 0 0 0 0.

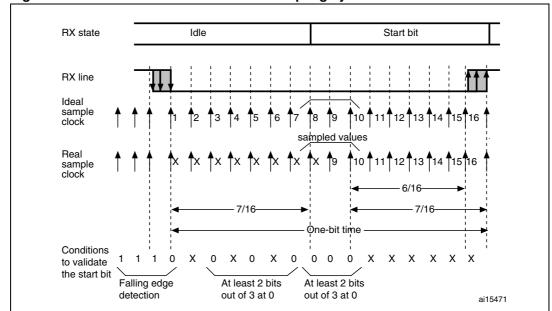


Figure 228. Start bit detection when oversampling by 16 or 8

Note:

If the sequence is not complete, the start bit detection aborts and the receiver returns to the idle state (no flag is set) where it waits for a falling edge.

The start bit is confirmed (RXNE flag set, interrupt generated if RXNEIE=1) if the 3 sampled bits are at 0 (first sampling on the 3rd, 5th and 7th bits finds the 3 bits at 0 and second sampling on the 8th, 9th and 10th bits also finds the 3 bits at 0).

The start bit is validated (RXNE flag set, interrupt generated if RXNEIE=1) but the NE noise flag is set if, for both samplings, at least 2 out of the 3 sampled bits are at 0 (sampling on the 3rd, 5th and 7th bits and sampling on the 8th, 9th and 10th bits). If this condition is not met, the start detection aborts and the receiver returns to the idle state (no flag is set).

If, for one of the samplings (sampling on the 3rd, 5th and 7th bits or sampling on the 8th, 9th and 10th bits), 2 out of the 3 bits are found at 0, the start bit is validated but the NE noise flag bit is set.

# **Character reception**

During an USART reception, data shifts in least significant bit first through the RX pin. In this mode, the USART\_DR register consists of a buffer (RDR) between the internal bus and the received shift register.

#### Procedure:

- 1. Enable the USART by writing the UE bit in USART CR1 register to 1.
- 2. Program the M bit in USART\_CR1 to define the word length.
- 3. Program the number of stop bits in USART\_CR2.
- 4. Select DMA enable (DMAR) in USART\_CR3 if multibuffer communication is to take place. Configure the DMA register as explained in multibuffer communication. STEP 3
- Select the desired baud rate using the baud rate register USART\_BRR
- Set the RE bit USART\_CR1. This enables the receiver which begins searching for a start bit.

#### When a character is received

- The RXNE bit is set. It indicates that the content of the shift register is transferred to the RDR. In other words, data has been received and can be read (as well as its associated error flags).
- An interrupt is generated if the RXNEIE bit is set.
- The error flags can be set if a frame error, noise or an overrun error has been detected during reception.
- In multibuffer, RXNE is set after every byte received and is cleared by the DMA read to the Data Register.
- In single buffer mode, clearing the RXNE bit is performed by a software read to the USART\_DR register. The RXNE flag can also be cleared by writing a zero to it. The RXNE bit must be cleared before the end of the reception of the next character to avoid an overrun error.

Note:

The RE bit should not be reset while receiving data. If the RE bit is disabled during reception, the reception of the current byte will be aborted.

#### **Break character**

When a break character is received, the USART handles it as a framing error.

## Idle character

When an idle frame is detected, there is the same procedure as a data received character plus an interrupt if the IDLEIE bit is set.

# **Overrun error**

An overrun error occurs when a character is received when RXNE has not been reset. Data can not be transferred from the shift register to the RDR register until the RXNE bit is cleared.

The RXNE flag is set after every byte received. An overrun error occurs if RXNE flag is set when the next data is received or the previous DMA request has not been serviced. When an overrun error occurs:

• The ORE bit is set.

- The RDR content will not be lost. The previous data is available when a read to USART\_DR is performed.
- The shift register will be overwritten. After that point, any data received during overrun
  is lost
- An interrupt is generated if either the RXNEIE bit is set or both the EIE and DMAR bits are set.
- The ORE bit is reset by a read to the USART\_SR register followed by a USART\_DR register read operation.

Note: The ORE bit, when set, indicates that at least 1 data has been lost. There are two possibilities:

- if RXNE=1, then the last valid data is stored in the receive register RDR and can be read,
- if RXNE=0, then it means that the last valid data has already been read and thus there is nothing to be read in the RDR. This case can occur when the last valid data is read in the RDR at the same time as the new (and lost) data is received. It may also occur when the new data is received during the reading sequence (between the USART\_SR register read access and the USART\_DR read access).

# Selecting the proper oversampling method

The receiver implements different user-configurable oversampling techniques (except in synchronous mode) for data recovery by discriminating between valid incoming data and noise.

The oversampling method can be selected by programming the OVER8 bit in the USART\_CR1 register and can be either 16 or 8 times the baud rate clock (*Figure 229* and *Figure 230*).

Depending on the application:

- select oversampling by 8 (OVER8=1) to achieve higher speed (up to f<sub>PCLK</sub>/8). In this
  case the maximum receiver tolerance to clock deviation is reduced (refer to
  Section 24.3.5: USART receiver tolerance to clock deviation on page 645)
- select oversampling by 16 (OVER8=0) to increase the tolerance of the receiver to clock deviations. In this case, the maximum speed is limited to maximum f<sub>PCLK</sub>/16

Programming the ONEBIT bit in the USART\_CR3 register selects the method used to evaluate the logic level. There are two options:

- the majority vote of the three samples in the center of the received bit. In this case, when the 3 samples used for the majority vote are not equal, the NF bit is set
- a single sample in the center of the received bit Depending on the application:
  - select the three samples' majority vote method (ONEBIT=0) when operating in a noisy environment and reject the data when a noise is detected (refer to *Figure 86*) because this indicates that a glitch occurred during the sampling.
  - select the single sample method (ONEBIT=1) when the line is noise-free to increase the receiver's tolerance to clock deviations (see Section 24.3.5: USART receiver tolerance to clock deviation on page 645). In this case the NF bit will never be set.



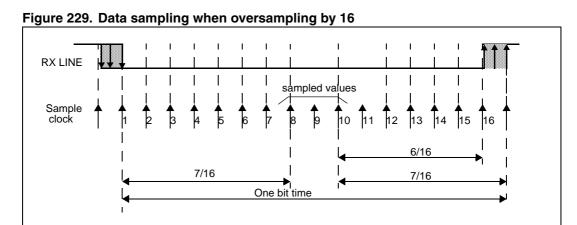
When noise is detected in a frame:

- The NF bit is set at the rising edge of the RXNE bit.
- The invalid data is transferred from the Shift register to the USART\_DR register.
- No interrupt is generated in case of single byte communication. However this bit rises at the same time as the RXNE bit which itself generates an interrupt. In case of multibuffer communication an interrupt will be issued if the EIE bit is set in the USART CR3 register.

The NF bit is reset by a USART\_SR register read operation followed by a USART\_DR register read operation.

Note:

Oversampling by 8 is not available in the Smartcard, IrDA and LIN modes. In those modes, the OVER8 bit is forced to '0 by hardware.



**RX LINE** sampled values Sample clock(x8) 2/8 3/8 3/8 One bit time

Figure 230. Data sampling when oversampling by 8

Table 86. Noise detection from sampled data

Sampled value	NE status	Received bit value
000	0	0
001	1	0
010	1	0
011	1	1
100	1	0
101	1	1

Table 86. Noise detection from sampled data (continued)

Sampled value	NE status	Received bit value
110	1	1
111	0	1

# Framing error

A framing error is detected when:

The stop bit is not recognized on reception at the expected time, following either a desynchronization or excessive noise.

When the framing error is detected:

- The FE bit is set by hardware
- The invalid data is transferred from the Shift register to the USART\_DR register.
- No interrupt is generated in case of single byte communication. However this bit rises
  at the same time as the RXNE bit which itself generates an interrupt. In case of
  multibuffer communication an interrupt will be issued if the EIE bit is set in the
  USART CR3 register.

The FE bit is reset by a USART\_SR register read operation followed by a USART\_DR register read operation.

# Configurable stop bits during reception

The number of stop bits to be received can be configured through the control bits of Control Register 2 - it can be either 1 or 2 in normal mode and 0.5 or 1.5 in Smartcard mode.

- 1. **0.5 stop bit (reception in Smartcard mode)**: No sampling is done for 0.5 stop bit. As a consequence, no framing error and no break frame can be detected when 0.5 stop bit is selected.
- 2. 1 stop bit. Sampling for 1 stop Bit is done on the 8th, 9th and 10th samples.
- 3. 1.5 stop bits (Smartcard mode): When transmitting in smartcard mode, the device must check that the data is correctly sent. Thus the receiver block must be enabled (RE =1 in the USART\_CR1 register) and the stop bit is checked to test if the smartcard has detected a parity error. In the event of a parity error, the smartcard forces the data signal low during the sampling NACK signal-, which is flagged as a framing error. Then, the FE flag is set with the RXNE at the end of the 1.5 stop bit. Sampling for 1.5 stop bits is done on the 16th, 17th and 18th samples (1 baud clock period after the beginning of the stop bit). The 1.5 stop bit can be decomposed into 2 parts: one 0.5 baud clock period during which nothing happens, followed by 1 normal stop bit period during which sampling occurs halfway through. Refer to Section 24.3.11: Smartcard on page 654 for more details.
- 4. **2** stop bits: Sampling for 2 stop bits is done on the 8th, 9th and 10th samples of the first stop bit. If a framing error is detected during the first stop bit the framing error flag will be set. The second stop bit is not checked for framing error. The RXNE flag will be set at the end of the first stop bit.

# 24.3.4 Fractional baud rate generation

The baud rate for the receiver and transmitter (Rx and Tx) are both set to the same value as programmed in the Mantissa and Fraction values of USARTDIV.

# Equation 1: Baud rate for standard USART (SPI mode included)

$$Tx/Rx \text{ baud } = \frac{f_{CK}}{8 \times (2 - OVER8) \times USARTDIV}$$

# Equation 2: Baud rate in Smartcard, LIN and IrDA modes

$$Tx/Rx \text{ baud } = \frac{f_{CK}}{16 \times USARTDIV}$$

USARTDIV is an unsigned fixed point number that is coded on the USART\_BRR register.

- When OVER8=0, the fractional part is coded on 4 bits and programmed by the DIV\_fraction[3:0] bits in the USART\_BRR register
- When OVER8=1, the fractional part is coded on 3 bits and programmed by the DIV\_fraction[2:0] bits in the USART\_BRR register, and bit DIV\_fraction[3] must be kept cleared.

Note:

The baud counters are updated to the new value in the baud registers after a write operation to USART\_BRR. Hence the baud rate register value should not be changed during communication.

# How to derive USARTDIV from USART\_BRR register values when OVER8=0

## Example 1:

If DIV\_Mantissa = 0d27 and DIV\_Fraction = 0d12 (USART\_BRR = 0x1BC), then

Mantissa (USARTDIV) = 0d27

Fraction (USARTDIV) = 12/16 = 0d0.75

Therefore USARTDIV = 0d27.75

# Example 2:

To program USARTDIV = 0d25.62

This leads to:

DIV Fraction = 16\*0d0.62 = 0d9.92

The nearest real number is 0d10 = 0xA

 $DIV_Mantissa = mantissa (0d25.620) = 0d25 = 0x19$ 

Then, USART\_BRR = 0x19A hence USARTDIV = 0d25.625

# Example 3:

To program USARTDIV = 0d50.99

This leads to:

DIV Fraction = 16\*0d0.99 = 0d15.84

The nearest real number is  $0d16 = 0x10 \Rightarrow 0verflow of DIV_frac[3:0] \Rightarrow 0d16 \Rightarrow 0$ 

DIV\_Mantissa = mantissa (0d50.990 + carry) = 0d51 = 0x33

Then, USART\_BRR = 0x330 hence USARTDIV = 0d51.000

# How to derive USARTDIV from USART\_BRR register values when OVER8=1

### Example 1:

If DIV\_Mantissa = 0x27 and DIV\_Fraction[2:0]= 0d6 (USART\_BRR = 0x1B6), then

Mantissa (USARTDIV) = 0d27

Fraction (USARTDIV) = 6/8 = 0d0.75

Therefore USARTDIV = 0d27.75

# Example 2:

To program USARTDIV = 0d25.62

This leads to:

DIV\_Fraction = 8\*0d0.62 = 0d4.96

The nearest real number is 0d5 = 0x5

DIV\_Mantissa = mantissa (0d25.620) = 0d25 = 0x19

Then, USART\_BRR = 0x195 => USARTDIV = 0d25.625

# Example 3:

To program USARTDIV = 0d50.99

This leads to:

 $DIV_Fraction = 8*0d0.99 = 0d7.92$ 

The nearest real number is  $0d8 = 0x8 => overflow of the DIV_frac[2:0] => carry must be added up to the mantissa$ 

DIV\_Mantissa = mantissa (0d50.990 + carry) = 0d51 = 0x33

Then, USART\_BRR = 0x0330 => USARTDIV = 0d51.000

Table 87. Error calculation for programmed baud rates at  $f_{PCLK} = 8$  MHz or  $f_{PCLK} = 12$  MHz, oversampling by  $16^{(1)}$ 

	Oversampling by 16 (OVER8=0)													
Ba	ud rate7		f <sub>PCLK</sub> = 8 MH	f <sub>P</sub>	<sub>CLK</sub> = 12 MHz									
S.No	Desired	Actual	Value programmed in the baud rate register	% Error = (Calculated - Desired) B.rate / Desired B.rate	Actual	Value programmed in the baud rate register	% Error							
1	1.2 KBps	1.2 KBps	416.6875	0	1.2 KBps	625	0							
2	2.4 KBps	2.4 KBps	208.3125	0.01	2.4 KBps	312.5	0							
3	9.6 KBps	9.604 KBps	52.0625	0.04	9.6 KBps	78.125	0							
4	4 19.2 KBps 19.185 KBps		26.0625	0.08	19.2 KBps	39.0625	0							

Table 87. Error calculation for programmed baud rates at  $f_{PCLK}$  = 8 MHz or  $f_{PCLK}$  = 12 MHz, oversampling by  $16^{(1)}$  (continued)

	Oversampling by 16 (OVER8=0)												
Ва	ud rate7		f <sub>PCLK</sub> = 8 MH	z	f <sub>P</sub> (	<sub>CLK</sub> = 12 MHz							
S.No	Desired	Actual	Value programmed in the baud rate register	% Error = (Calculated - Desired) B.rate / Desired B.rate	Actual	Value programmed in the baud rate register	% Error						
5	38.4 KBps	38.462 KBps	13	0.16	38.339 KBps	19.5625	0.16						
6	57.6 KBps	57.554 KBps	8.6875	0.08	57.692 KBps	13	0.16						
7	115.2 KBps	115.942 KBps	4.3125	0.64	115.385 KBps	6.5	0.16						
8	230.4 KBps	228.571 KBps	2.1875	0.79	230.769 KBps	3.25	0.16						
9	460.8 KBps	470.588 KBps	1.0625	2.12	461.538 KBps	1.625	0.16						
10	921.6 KBps	NA	NA	NA	NA	NA	NA						
11	2 MBps	NA	NA	NA	NA	NA	NA						
12	3 MBps	NA	NA	NA	NA	NA	NA						

<sup>1.</sup> The lower the CPU clock the lower the accuracy for a particular baud rate. The upper limit of the achievable baud rate can be fixed with these data.

Table 88. Error calculation for programmed baud rates at  $f_{PCLK}$  = 8 MHz or  $f_{PCLK}$  =12 MHz, oversampling by  $8^{(1)}$ 

	Oversampling by 8 (OVER8 = 1)													
Ва	aud rate	f	PCLK = 8 MHz		f <sub>PCL</sub>	<sub>K</sub> = 12 MHz								
S.No	Desired	Actual	Value programmed in the baud rate register	% Error = (Calculated - Desired) B.rate / Desired B.rate	Actual	Value programmed in the baud rate register	% Error							
1	1.2 KBps	1.2 KBps	833.375	0	1.2 KBps	1250	0							
2	2.4 KBps	2.4 KBps	416.625	0.01	2.4 KBps	625	0							
3	9.6 KBps	9.604 KBps	104.125	0.04	9.6 KBps	156.25	0							
4	19.2 KBps	19.185 KBps	52.125	80.0	19.2 KBps	78.125	0							
5	38.4 KBps	38.462 KBps	26	0.16	38.339 KBps	39.125	0.16							
6	57.6 KBps	57.554 KBps	17.375	0.08	57.692 KBps	26	0.16							
7	115.2 KBps	115.942 KBps	8.625	0.64	115.385 KBps	13	0.16							
8	230.4 KBps	228.571 KBps	4.375	0.79	230.769 KBps	6.5	0.16							
9	460.8 KBps	470.588 KBps	2.125	2.12	461.538 KBps	3.25	0.16							
10	921.6 KBps	888.889 KBps	1.125	3.55	923.077 KBps	1.625	0.16							

Table 88. Error calculation for programmed baud rates at  $f_{PCLK}$  = 8 MHz or  $f_{PCLK}$  =12 MHz, oversampling by  $8^{(1)}$  (continued)

	Oversampling by 8 (OVER8 = 1)													
Ва	aud rate	f	PCLK = 8 MHz	f <sub>PCLK</sub> = 12 MHz										
S.No	Desired	Actual	Value programmed in the baud rate register	% Error = (Calculated - Desired) B.rate / Desired B.rate	Actual	Value programmed in the baud rate register	% Error							
11 2 MBps		NA	NA	NA	NA	NA	NA							
12	3 MBps	NA	NA	NA	NA	NA	NA							

<sup>1.</sup> The lower the CPU clock the lower the accuracy for a particular baud rate. The upper limit of the achievable baud rate can be fixed with these data.

Table 89. Error calculation for programmed baud rates at  $f_{PCLK}$  = 16 MHz or  $f_{PCLK}$  = 24 MHz, oversampling by  $16^{(1)}$ 

		inplining by 10					
		Ov	ersampling b	y 16 (OVER8 = 0	)		
В	Baud rate f <sub>PCLK</sub> = 16 MHz					<sub>CLK</sub> = 24 MHz	
S.No	Desired	Actual	Value programmed in the baud rate register	% Error = (Calculated - Desired) B.rate / Desired B.rate	Actual	Value programmed in the baud rate register	% Error
1	1.2 KBps	1.2 KBps	833.3125	0	1.2	1250	0
2	2.4 KBps	2.4 KBps	416.6875	0	2.4	625	0
3	9.6 KBps	9.598 KBps	104.1875	0.02	9.6	156.25	0
4	19.2 KBps	19.208 KBps	52.0625	0.04	19.2	78.125	0
5	38.4 KBps	38.369 KBps	26.0625	0.08	38.4	39.0625	0
6	57.6 KBps	57.554 KBps	17.375	0.08	57.554	26.0625	0.08
7	115.2 KBps	115.108 KBps	8.6875	0.08	115.385	13	0.16
8	230.4 KBps	231.884 KBps	4.3125	0.64	230.769	6.5	0.16
9	460.8 KBps	457.143 KBps	2.1875	0.79	461.538	3.25	0.16
10	921.6 KBps	941.176 KBps	1.0625	2.12	923.077	1.625	0.16
11	2 MBps	NA	NA	NA	NA	NA	NA
12	3 MBps	NA	NA	NA	NA	NA	NA

<sup>1.</sup> The lower the CPU clock the lower the accuracy for a particular baud rate. The upper limit of the achievable baud rate can be fixed with these data.

Table 90. Error calculation for programmed baud rates at  $f_{PCLK}$  = 16 MHz or  $f_{PCLK}$  = 24 MHz, oversampling by  $8^{(1)}$ 

		ampining by 0	0	b 0 (OVEDO 4)							
	Oversampling by 8 (OVER8=1)										
В	aud rate		f <sub>PCLK</sub> = 16 MH	lz	f <sub>PC</sub>	<sub>LK</sub> = 24 MHz					
S.No	Desired	Actual	Value programmed in the baud rate register	% Error = (Calculated - Desired) B.rate / Desired B.rate	Actual	Value programmed in the baud rate register	% Error				
1	1.2 KBps	1.2 KBps	1666.625	0	1.2 KBps	2500	0				
2	2.4 KBps	2.4 KBps	833.375	0	2.4 KBps	1250	0				
3	9.6 KBps	9.598 KBps	208.375	0.02	9.6 KBps	312.5	0				
4	19.2 KBps	19.208 KBps	104.125	0.04	19.2 KBps	156.25	0				
5	38.4 KBps	38.369 KBps	52.125	0.08	38.4 KBps	78.125	0				
6	57.6 KBps	57.554 KBps	34.75	0.08	57.554 KBps	52.125	0.08				
7	115.2 KBps	115.108 KBps	17.375	0.08	115.385 KBps	26	0.16				
8	230.4 KBps	231.884 KBps	8.625	0.64	230.769 KBps	13	0.16				
9	460.8 KBps	457.143 KBps	4.375	0.79	461.538 KBps	6.5	0.16				
10	921.6 KBps	941.176 KBps	2.125	2.12	923.077 KBps	3.25	0.16				
11	2 MBps	2000 KBps	1	0	2000 KBps	1.5	0				
12	3 MBps	NA	NA	NA	3000 KBps	1	0				

<sup>1.</sup> The lower the CPU clock the lower the accuracy for a particular baud rate. The upper limit of the achievable baud rate can be fixed with these data.

Table 91. Error calculation for programmed baud rates at  $f_{PCLK}$  = 8 MHz or  $f_{PCLK}$  = 16 MHz, oversampling by  $16^{(1)}$ 

	Oversampling by 16 (OVER8=0)										
Е	Baud rate f <sub>PCLK</sub> = 8 MHz				f <sub>PC</sub>	<sub>LK</sub> = 16 MHz					
S.No	Desired	Actual	Value programmed in the baud rate register	% Error = (Calculated - Desired)B.Rate /Desired B.Rate	Actual	Value programmed in the baud rate register	% Error				
1.	2.4 KBps	2.400 KBps	208.3125	0.00%	2.400 KBps	416.6875	0.00%				
2.	9.6 KBps	9.604 KBps	52.0625	0.04%	9.598 KBps	104.1875	0.02%				
3.	19.2 KBps	19.185 KBps	26.0625	0.08%	19.208 KBps	52.0625	0.04%				
4.	57.6 KBps	57.554 KBps	8.6875	0.08%	57.554 KBps	17.3750	0.08%				
5.	115.2 KBps	115.942 KBps	4.3125	0.64%	115.108 KBps	8.6875	0.08%				
6.	230.4 KBps	228.571 KBps	2.1875	0.79%	231.884 KBps	4.3125	0.64%				
7.	460.8 KBps	470.588 KBps	1.0625	2.12%	457.143 KBps	2.1875	0.79%				
8.	896 KBps	NA	NA	NA	888.889 KBps	1.1250	0.79%				

Table 91. Error calculation for programmed baud rates at  $f_{PCLK}$  = 8 MHz or  $f_{PCLK}$  = 16 MHz, oversampling by  $16^{(1)}$  (continued)

	Oversampling by 16 (OVER8=0)										
E	Baud rate		f <sub>PCLK</sub> = 8 MI	Hz	f <sub>PC</sub>	<sub>LK</sub> = 16 MHz					
S.No	Desired	Actual	Value programmed in the baud rate register	% Error = (Calculated - Desired)B.Rate /Desired B.Rate	Actual	Value programmed in the baud rate register	% Error				
9.	921.6 KBps	NA	NA	NA	941.176 KBps	1.0625	2.12%				
10.	1.792 MBps	NA	NA	NA	NA	NA	NA				
11.	1.8432 MBps	NA	NA	NA	NA	NA	NA				
12.	3.584 MBps	NA	NA	NA	NA	NA	NA				
13.	3.6864 MBps	NA	NA	NA	NA	NA	NA				
14.	7.168 MBps	NA	NA	NA	NA	NA	NA				
15.	7.3728 MBps	NA	NA	NA	NA	NA	NA				

<sup>1.</sup> The lower the CPU clock the lower the accuracy for a particular baud rate. The upper limit of the achievable baud rate can be fixed with these data.

Table 92. Error calculation for programmed baud rates at  $f_{PCLK}$  = 8 MHz or  $f_{PCLK}$  = 16 MHz, oversampling by  $8^{(1)}$ 

Oversampling by 8 (OVER8=1)									
Baud rate		f <sub>PCLK</sub> = 8 MH	z	f <sub>PCLK</sub> = 16 MHz					
Desired	Actual	Value programmed in the baud rate register	% Error = (Calculated - Desired)B.Rate /Desired B.Rate	Actual	Value programmed in the baud rate register	% Error			
2.4 KBps	2.400 KBps	416.625	0.01%	2.400 KBps	833.375	0.00%			
9.6 KBps	9.604 KBps	104.125	0.04%	9.598 KBps	208.375	0.02%			
19.2 KBps	19.185 KBps	52.125	0.08%	19.208 KBps	104.125	0.04%			
57.6 KBps	57.557 KBps	17.375	0.08%	57.554 KBps	34.750	0.08%			
115.2 KBps	115.942 KBps	8.625	0.64%	115.108 KBps	17.375	0.08%			
230.4 KBps	228.571 KBps	4.375	0.79%	231.884 KBps	8.625	0.64%			
460.8 KBps	470.588 KBps	2.125	2.12%	457.143 KBps	4.375	0.79%			
896 KBps	888.889 KBps	1.125	0.79%	888.889 KBps	2.250	0.79%			
921.6 KBps	888.889 KBps	1.125	3.55%	941.176 KBps	2.125	2.12%			
1.792 MBps	NA	NA	NA	1.7777 MBps	1.125	0.79%			
1.8432 MBps	NA	NA	NA	1.7777 MBps	1.125	3.55%			
3.584 MBps	NA	NA	NA	NA	NA	NA			
3.6864 MBps	NA	NA	NA	NA	NA	NA			
7.168 MBps	NA	NA	NA	NA	NA	NA			
7.3728 MBps	NA	NA	NA	NA	NA	NA			
	Desired  2.4 KBps 9.6 KBps 19.2 KBps 57.6 KBps 115.2 KBps 230.4 KBps 460.8 KBps 921.6 KBps 1.792 MBps 1.8432 MBps 3.584 MBps 3.6864 MBps 7.168 MBps	Desired         Actual           2.4 KBps         2.400 KBps           9.6 KBps         9.604 KBps           19.2 KBps         19.185 KBps           57.6 KBps         57.557 KBps           115.2 KBps         115.942 KBps           230.4 KBps         228.571 KBps           460.8 KBps         470.588 KBps           896 KBps         888.889 KBps           921.6 KBps         888.889 KBps           1.792 MBps         NA           1.8432 MBps         NA           3.6864 MBps         NA           7.168 MBps         NA	Desired         Actual         Value programmed in the baud rate register           2.4 KBps         2.400 KBps         416.625           9.6 KBps         9.604 KBps         104.125           19.2 KBps         19.185 KBps         52.125           57.6 KBps         57.557 KBps         17.375           115.2 KBps         115.942 KBps         8.625           230.4 KBps         228.571 KBps         4.375           460.8 KBps         470.588 KBps         2.125           896 KBps         888.889 KBps         1.125           921.6 KBps         888.889 KBps         1.125           1.792 MBps         NA         NA           1.8432 MBps         NA         NA           3.6864 MBps         NA         NA           7.168 MBps         NA         NA	Desired   Actual   Value programmed in the baud rate register   Calculated - Desired)   B.Rate	Desired         Actual         Value programmed in the baud rate register         % Error = (Calculated - Desired)B.Rate         Actual           2.4 KBps         2.400 KBps         416.625         0.01%         2.400 KBps           9.6 KBps         9.604 KBps         104.125         0.04%         9.598 KBps           19.2 KBps         19.185 KBps         52.125         0.08%         19.208 KBps           57.6 KBps         57.557 KBps         17.375         0.08%         57.554 KBps           115.2 KBps         115.942 KBps         8.625         0.64%         115.108 KBps           230.4 KBps         228.571 KBps         4.375         0.79%         231.884 KBps           460.8 KBps         470.588 KBps         2.125         2.12%         457.143 KBps           896 KBps         888.889 KBps         1.125         0.79%         888.889 KBps           921.6 KBps         888.889 KBps         1.125         3.55%         941.176 KBps           1.792 MBps         NA         NA         NA         1.7777 MBps           1.8432 MBps         NA         NA         NA         NA           3.6864 MBps         NA         NA         NA         NA           7.168 MBps         NA         NA	Desired   Actual   Value programmed in the baud rate register   Calculated - Desired   Desired   B.Rate   Actual   Programmed in the baud rate register   Calculated - Desired B.Rate   Desired B.Rate   Actual   Programmed in the baud rate register   Programmed in the baud rate register   Desired B.Rate   Actual   Programmed in the baud rate register   Programmed in the baud rate register   Programmed in the baud rate register   Desired B.Rate   Actual   Programmed in the baud rate register   Programmed in the baud rate register   Desired B.Rate   Desired B.Rat			

1. The lower the CPU clock the lower the accuracy for a particular baud rate. The upper limit of the achievable baud rate can be fixed with these data.

Table 93. Error calculation for programmed baud rates at  $f_{PCLK}$  = 30 MHz or  $f_{PCLK}$  = 60 MHz, oversampling by  $16^{(1)(2)}$ 

	Oversampling by 16 (OVER8=0)									
В	aud rate		f <sub>PCLK</sub> = 30 MI	-lz	f <sub>P</sub> (	<sub>CLK</sub> = 60 MHz				
S.No	Desired	Actual	Value programmed in the baud rate register	% Error = (Calculated - Desired)B.Rate /Desired B.Rate	Actual	Value programmed in the baud rate register	% Error			
1.	2.4 KBps	2.400 KBps	781.2500	0.00%	2.400 KBps	1562.5000	0.00%			
2.	9.6 KBps	9.600 KBps	195.3125	0.00%	9.600 KBps	390.6250	0.00%			
3.	19.2 KBps	19.194 KBps	97.6875	0.03%	19.200 KBps	195.3125	0.00%			
4.	57.6 KBps	57.582KBps	32.5625	0.03%	57.582 KBps	65.1250	0.03%			
5.	115.2 KBps	115.385 KBps	16.2500	0.16%	115.163 KBps	32.5625	0.03%			
6.	230.4 KBps	230.769 KBps	8.1250	0.16%	230.769KBps	16.2500	0.16%			
7.	460.8 KBps	461.538 KBps	4.0625	0.16%	461.538 KBps	8.1250	0.16%			
8.	896 KBps	909.091 KBps	2.0625	1.46%	895.522 KBps	4.1875	0.05%			
9.	921.6 KBps	909.091 KBps	2.0625	1.36%	923.077 KBps	4.0625	0.16%			
10.	1.792 MBps	1.1764 MBps	1.0625	1.52%	1.8182 MBps	2.0625	1.36%			
11.	1.8432 MBps	1.8750 MBps	1.0000	1.73%	1.8182 MBps	2.0625	1.52%			
12.	3.584 MBps	NA	NA	NA	3.2594 MBps	1.0625	1.52%			
13.	3.6864 MBps	NA	NA	NA	3.7500 MBps	1.0000	1.73%			
14.	7.168 MBps	NA	NA	NA	NA	NA	NA			
15.	7.3728 MBps	NA	NA	NA	NA	NA	NA			

<sup>1.</sup> The lower the CPU clock the lower the accuracy for a particular baud rate. The upper limit of the achievable baud rate can be fixed with these data.

Table 94. Error calculation for programmed baud rates at  $f_{PCLK}$  = 30 MHz or  $f_{PCLK}$  = 60 MHz, oversampling by  $8^{(1)}$  (2)

	Oversampling by 8 (OVER8=1)									
E	Baud rate f <sub>PCLK</sub> = 30 MHz			f <sub>PCLK</sub> =60 MHz						
S.No	Desired	Actual	Value programmed in the baud rate register	Desired)B.Rate	Actual	Value programmed in the baud rate register	% Error			
1.	2.4 KBps	2.400 KBps	1562.5000	0.00%	2.400 KBps	3125.0000	0.00%			
2.	9.6 KBps	9.600 KBps	390.6250	0.00%	9.600 KBps	781.2500	0.00%			

<sup>2.</sup> Only USART1 and USART6 are clocked with PCLK2. Other USARTs are clocked with PCLK1. Refer to the device datasheets for the maximum values for PCLK1 and PCLK2.

Table 94. Error calculation for programmed baud rates at  $f_{PCLK}$  = 30 MHz or  $f_{PCLK}$  = 60 MHz, oversampling by  $8^{(1)}$  (continued)

	Oversampling by 8 (OVER8=1)									
E	Baud rate		f <sub>PCLK</sub> = 30 MH	Iz	f <sub>PCLK</sub> =60 MHz					
S.No	Desired	Actual	Value programmed in the baud rate register	% Error = (Calculated - Desired)B.Rate /Desired B.Rate	Actual	Value programmed in the baud rate register	% Error			
3.	19.2 KBps	19.194 KBps	195.3750	0.03%	19.200 KBps	390.6250	0.00%			
4.	57.6 KBps	57.582 KBps	65.1250	0.16%	57.582 KBps	130.2500	0.03%			
5.	115.2 KBps	115.385 KBps	32.5000	0.16%	115.163 KBps	65.1250	0.03%			
6.	230.4 KBps	230.769 KBps	16.2500	0.16%	230.769 KBps	32.5000	0.16%			
7.	460.8 KBps	461.538 KBps	8.1250	0.16%	461.538 KBps	16.2500	0.16%			
8.	896 KBps	909.091 KBps	4.1250	1.46%	895.522 KBps	8.3750	0.05%			
9.	921.6 KBps	909.091 KBps	4.1250	1.36%	923.077 KBps	8.1250	0.16%			
10.	1.792 MBps	1.7647 MBps	2.1250	1.52%	1.8182 MBps	4.1250	1.46%			
11.	1.8432 MBps	1.8750 MBps	2.0000	1.73%	1.8182 MBps	4.1250	1.36%			
12.	3.584 MBps	3.7500 MBps	1.0000	4.63%	3.5294 MBps	2.1250	1.52%			
13.	3.6864 MBps	3.7500 MBps	1.0000	1.73%	3.7500 MBps	2.0000	1.73%			
14.	7.168 MBps	NA	NA	NA	7.5000 MBps	1.0000	4.63%			
15.	7.3728 MBps	NA	NA	NA	7.5000 MBps	1.0000	1.73%			

<sup>1.</sup> The lower the CPU clock the lower the accuracy for a particular baud rate. The upper limit of the achievable baud rate can be fixed with these data.

Table 95. Error calculation for programmed baud rates at  $f_{PCLK}$  = 42 MHz or  $f_{PCLK}$  = 84 Hz, oversampling by  $16^{(1)(2)}$ 

	Oversampling by 16 (OVER8=0)										
Ва	aud rate		f <sub>PCLK</sub> = 42 MH	-lz	f <sub>PC</sub>	<sub>LK</sub> = 84 MHz					
S.No	Desired	Actual	Value programmed in the baud rate register	% Error = (Calculated - Desired)B.Rate /Desired B.Rate		Value programmed in the baud rate register	% Error				
1	1.2 KBps	1.2 KBps	2187.5	0	1.2 KBps	4375	0				
2	2.4 KBps	2.4 KBps	1093.75	0	2.4 KBps	2187.5	0				
3	9.6 KBps	9.6 KBps	273.4375	0	9.6 KBps	546.875	0				
4	19.2 KBps	19.195 KBps	136.75	0.02	19.2 KBps	273.4375	0				
5	38.4 KBps	38.391 KBps	68.375	0.02	38.391 KBps	136.75	0.02				
6	57.6 KBps	57.613 KBps	45.5625	0.02	57.613 KBps	91.125	0.02				
7	115.2 KBps	115.068 KBps	22.8125	0.11	115.226 KBps	45.5625	0.02				

<sup>2.</sup> Only USART1 and USART6 are clocked with PCLK2. Other USARTs are clocked with PCLK1. Refer to the device datasheets for the maximum values for PCLK1 and PCLK2.

Table 95. Error calculation for programmed baud rates at  $f_{PCLK} = 42$  MHz or  $f_{PCLK} = 84$  Hz, oversampling by  $16^{(1)(2)}$  (continued)

	Oversampling by 16 (OVER8=0)									
Ва	aud rate		f <sub>PCLK</sub> = 42 MH	łz	f <sub>PCLK</sub> = 84 MHz					
S.No	Desired	Actual	Value programmed in the baud rate register	% Error = (Calculated - Desired)B.Rate /Desired B.Rate	Actual	Value programmed in the baud rate register	% Error			
8	230.4 KBps	230.769 KBps	11.375	0.16	230.137 KBps	22.8125	0.11			
9	460.8 KBps	461.538 KBps	5.6875	0.16	461.538 KBps	11.375	0.16			
10	921.6 KBps	913.043 KBps	2.875	0.93	923.076 KBps	5.6875	0.93			
11	1.792 MBps	1.826 MBps	1.4375	1.9	1.787 MBps	2.9375	0.27			
12	1.8432 MBps	1.826 MBps	1.4375	0.93	1.826 MBps	2.875	0.93			
13	3.584 MBps	N.A	N.A	N.A	3.652 MBps	1.4375	1.9			
14	3.6864 MBps	N.A	N.A	N.A	3.652 MBps	1.4375	0.93			
15	7.168 MBps	N.A	N.A	N.A	N.A	N.A	N.A			
16	7.3728 MBps	N.A	N.A	N.A	N.A	N.A	N.A			
18	9 MBps	N.A	N.A	N.A	N.A	N.A	N.A			
20	10.5 MBps	N.A	N.A	N.A	N.A	N.A	N.A			

<sup>1.</sup> The lower the CPU clock the lower the accuracy for a particular baud rate. The upper limit of the achievable baud rate can be fixed with these data.

Table 96. Error calculation for programmed baud rates at  $f_{PCLK}$  = 42 MHz or  $f_{PCLK}$  = 84 MHz, oversampling by  $8^{(1)(2)}$ 

	Oversampling by 8 (OVER8=1)									
Е	Baud rate f <sub>PCLK</sub> = 42 MHz			f <sub>PCLK</sub> = 84 MHz						
S.No	Desired	Actual	Value % Error = programmed (Calculated - in the baud rate register /Desired B.Rate		Actual	Value programmed in the baud rate register	% Error			
1.	1.2 KBps	1.2 KBps	4375	0	1.2 KBps	8750	0			
2.	2.4 KBps	2.4 KBps	2187.5	0	2.4 KBps	4375	0			
3.	9.6 KBps	9.6 KBps	546.875	0	9.6 KBps	1093.75	0			
4.	19.2 KBps	19.195 KBps	273.5	0.02	19.2 KBps	546.875	0			
5.	38.4 KBps	38.391 KBps	136.75	0.02	38.391 KBps	273.5	0.02			
6.	57.6 KBps	57.613 KBps	91.125	0.02	57.613 KBps	182.25	0.02			
7.	115.2 KBps	115.068 KBps	45.625	0.11	115.226 KBps	91.125	0.02			
8.	230.4 KBps	230.769 KBps	22.75	0.11	230.137 KBps	45.625	0.11			
9.	460.8 KBps	461.538 KBps	11.375	0.16	461.538 KBps	22.75	0.16			

<sup>2.</sup> Only USART1 and USART6 are clocked with PCLK2. Other USARTs are clocked with PCLK1. Refer to the device datasheets for the maximum values for PCLK1 and PCLK2.

0.93

2.34

3.57

3.7

0

Oversampling by 8 (OVER8=1)  $f_{PCLK} = 42 \text{ MHz}$  $f_{PCLK} = 84 \text{ MHz}$ **Baud rate** Value % Error = Value (Calculated programmed programmed **Desired** S.No **Actual Actual** % Error Desired)B.Rate in the baud in the baud /Desired B.Rate rate register rate register 10. 921.6 KBps 913.043 KBps 5.75 0.93 923.076 KBps 11.375 0.93 11. 1.792 MBps 1.826 MBps 2.875 1.9 1.787Mbps 5.875 0.27 12. 1.8432 MBps 1.826 MBps 2.875 0.93 1.826 MBps 5.75 0.93 13. 3.584 MBps 3.5 MBps 1.5 2.34 3.652 MBps 2.875 1.9

Table 96. Error calculation for programmed baud rates at  $f_{PCLK}$  = 42 MHz or  $f_{PCLK}$  = 84 MHz, oversampling by  $8^{(1)(2)}$  (continued)

1.375

N.A

N.A

N.A

N.A

#### 24.3.5 USART receiver tolerance to clock deviation

3.82 MBps

N.A

N.A

N.A

N.A

14.

15.

16.

18.

20.

3.6864 MBps

7.168 MBps

7.3728 MBps

9 MBps

10.5 MBps

The USART's asynchronous receiver works correctly only if the total clock system deviation is smaller than the USART receiver's tolerance. The causes which contribute to the total deviation are:

3.57

N.A

N.A

N.A

N.A

3.652 MBps

7 MBps

7.636 MBps

9.333 MBps

10.5 MBps

2.875

1.5

1.375

1.125

1

- DTRA: Deviation due to the transmitter error (which also includes the deviation of the transmitter's local oscillator)
- DQUANT: Error due to the baud rate quantization of the receiver
- DREC: Deviation of the receiver's local oscillator
- DTCL: Deviation due to the transmission line (generally due to the transceivers which can introduce an asymmetry between the low-to-high transition timing and the high-tolow transition timing)

DTRA + DQUANT + DREC + DTCL < USART receiver's tolerance

The USART receiver's tolerance to properly receive data is equal to the maximum tolerated deviation and depends on the following choices:

- 10- or 11-bit character length defined by the M bit in the USART CR1 register
- oversampling by 8 or 16 defined by the OVER8 bit in the USART\_CR1 register
- use of fractional baud rate or not
- use of 1 bit or 3 bits to sample the data, depending on the value of the ONEBIT bit in the USART\_CR3 register

<sup>1.</sup> The lower the CPU clock the lower the accuracy for a particular baud rate. The upper limit of the achievable baud rate can be fixed with these data

Only USART1 and USART6 are clocked with PCLK2. Other USARTs are clocked with PCLK1. Refer to the device datasheets for the maximum values for PCLK1 and PCLK2.

OVER8 bit = 0OVER8 bit = 1 M bit ONEBIT=0 ONEBIT=1 ONEBIT=0 ONEBIT=1 3.75% 4.375% 3.75% n 2 50% 3.41% 1 3.41% 3.97% 2.27%

Table 97. USART receiver's tolerance when DIV fraction is 0

Table 98. USART receiver's tolerance when DIV\_Fraction is different from 0

M bit	OVER	3 bit = 0	OVER8 bit = 1		
IVI DIL	ONEBIT=0	ONEBIT=1	ONEBIT=0	ONEBIT=1	
0	3.33%	3.88%	2%	3%	
1	3.03%	3.53%	1.82%	2.73%	

Note:

The figures specified in Table and Table 98 may slightly differ in the special case when the received frames contain some Idle frames of exactly 10-bit times when M=0 (11-bit times when M=1).

# 24.3.6 Multiprocessor communication

There is a possibility of performing multiprocessor communication with the USART (several USARTs connected in a network). For instance one of the USARTs can be the master, its TX output is connected to the RX input of the other USART. The others are slaves, their respective TX outputs are logically ANDed together and connected to the RX input of the master.

In multiprocessor configurations it is often desirable that only the intended message recipient should actively receive the full message contents, thus reducing redundant USART service overhead for all non addressed receivers.

The non addressed devices may be placed in mute mode by means of the muting function. In mute mode:

- None of the reception status bits can be set.
- All the receive interrupts are inhibited.
- The RWU bit in USART\_CR1 register is set to 1. RWU can be controlled automatically by hardware or written by the software under certain conditions.

The USART can enter or exit from mute mode using one of two methods, depending on the WAKE bit in the USART\_CR1 register:

- Idle Line detection if the WAKE bit is reset,
- Address Mark detection if the WAKE bit is set.

# Idle line detection (WAKE=0)

The USART enters mute mode when the RWU bit is written to 1.

It wakes up when an Idle frame is detected. Then the RWU bit is cleared by hardware but the IDLE bit is not set in the USART\_SR register. RWU can also be written to 0 by software.

An example of mute mode behavior using Idle line detection is given in *Figure 231*.

RX Data 1 Data 2 Data 3 Data 4 IDLE Data 5 Data 6

RWU Mute Mode Normal Mode

RWU written to 1 Idle frame detected

Figure 231. Mute mode using Idle line detection

# Address mark detection (WAKE=1)

In this mode, bytes are recognized as addresses if their MSB is a '1 else they are considered as data. In an address byte, the address of the targeted receiver is put on the 4 LSB. This 4-bit word is compared by the receiver with its own address which is programmed in the ADD bits in the USART\_CR2 register.

The USART enters mute mode when an address character is received which does not match its programmed address. In this case, the RWU bit is set by hardware. The RXNE flag is not set for this address byte and no interrupt nor DMA request is issued as the USART would have entered mute mode.

It exits from mute mode when an address character is received which matches the programmed address. Then the RWU bit is cleared and subsequent bytes are received normally. The RXNE bit is set for the address character since the RWU bit has been cleared.

The RWU bit can be written to as 0 or 1 when the receiver buffer contains no data (RXNE=0 in the USART\_SR register). Otherwise the write attempt is ignored.

An example of mute mode behavior using address mark detection is given in Figure 232.

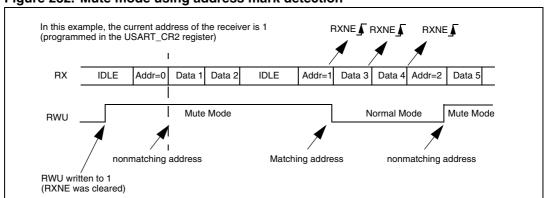


Figure 232. Mute mode using address mark detection

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# 24.3.7 Parity control

Parity control (generation of parity bit in transmission and parity checking in reception) can be enabled by setting the PCE bit in the USART\_CR1 register. Depending on the frame length defined by the M bit, the possible USART frame formats are as listed in *Table 99*.

Table 99. Frame formats

M bit	PCE bit	USART frame <sup>(1)</sup>
0	0	SB   8 bit data   STB
0	1	SB   7-bit data   PB   STB
1	0	SB   9-bit data   STB
1	1	SB   8-bit data PB   STB

<sup>1.</sup> Legends: SB: start bit, STB: stop bit, PB: parity bit.

# **Even parity**

The parity bit is calculated to obtain an even number of "1s" inside the frame made of the 7 or 8 LSB bits (depending on whether M is equal to 0 or 1) and the parity bit.

E.g.: data=00110101; 4 bits set => parity bit will be 0 if even parity is selected (PS bit in USART\_CR1 = 0).

# **Odd parity**

The parity bit is calculated to obtain an odd number of "1s" inside the frame made of the 7 or 8 LSB bits (depending on whether M is equal to 0 or 1) and the parity bit.

E.g.: data=00110101; 4 bits set => parity bit will be 1 if odd parity is selected (PS bit in USART CR1 = 1).

#### Parity checking in reception

If the parity check fails, the PE flag is set in the USART\_SR register and an interrupt is generated if PEIE is set in the USART\_CR1 register. The PE flag is cleared by a software sequence (a read from the status register followed by a read or write access to the USART\_DR data register).

Note:

In case of wakeup by an address mark: the MSB bit of the data is taken into account to identify an address but not the parity bit. And the receiver does not check the parity of the address data (PE is not set in case of a parity error).

# Parity generation in transmission

If the PCE bit is set in USART\_CR1, then the MSB bit of the data written in the data register is transmitted but is changed by the parity bit (even number of "1s" if even parity is selected (PS=0) or an odd number of "1s" if odd parity is selected (PS=1)).

Note:

The software routine that manages the transmission can activate the software sequence which clears the PE flag (a read from the status register followed by a read or write access to the data register). When operating in half-duplex mode, depending on the software, this can cause the PE flag to be unexpectedly cleared.

# 24.3.8 LIN (local interconnection network) mode

The LIN mode is selected by setting the LINEN bit in the USART\_CR2 register. In LIN mode, the following bits must be kept cleared:

- CLKEN in the USART\_CR2 register,
- STOP[1:0], SCEN, HDSEL and IREN in the USART\_CR3 register.

### LIN transmission

The same procedure explained in *Section 24.3.2* has to be applied for LIN Master transmission than for normal USART transmission with the following differences:

- Clear the M bit to configure 8-bit word length.
- Set the LINEN bit to enter LIN mode. In this case, setting the SBK bit sends 13 '0 bits as a break character. Then a bit of value '1 is sent to allow the next start detection.

## LIN reception

A break detection circuit is implemented on the USART interface. The detection is totally independent from the normal USART receiver. A break can be detected whenever it occurs, during Idle state or during a frame.

When the receiver is enabled (RE=1 in USART\_CR1), the circuit looks at the RX input for a start signal. The method for detecting start bits is the same when searching break characters or data. After a start bit has been detected, the circuit samples the next bits exactly like for the data (on the 8th, 9th and 10th samples). If 10 (when the LBDL = 0 in USART\_CR2) or 11 (when LBDL=1 in USART\_CR2) consecutive bits are detected as '0, and are followed by a delimiter character, the LBD flag is set in USART\_SR. If the LBDIE bit=1, an interrupt is generated. Before validating the break, the delimiter is checked for as it signifies that the RX line has returned to a high level.

If a '1 is sampled before the 10 or 11 have occurred, the break detection circuit cancels the current detection and searches for a start bit again.

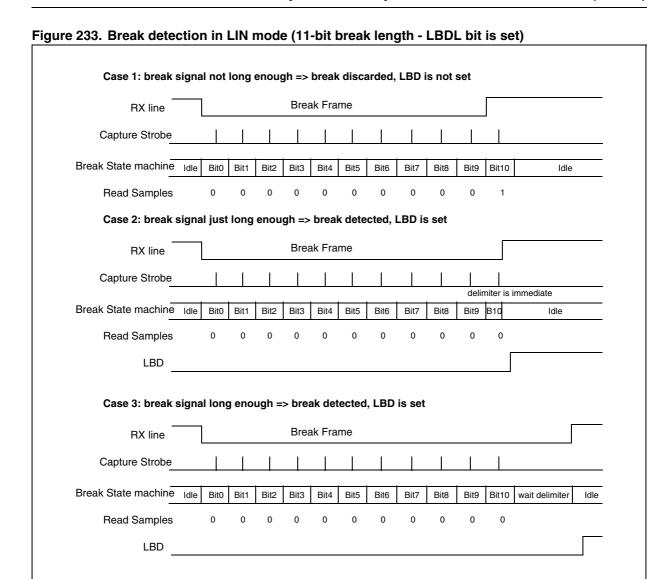
If the LIN mode is disabled (LINEN=0), the receiver continues working as normal USART, without taking into account the break detection.

If the LIN mode is enabled (LINEN=1), as soon as a framing error occurs (i.e. stop bit detected at '0, which will be the case for any break frame), the receiver stops until the break detection circuit receives either a '1, if the break word was not complete, or a delimiter character if a break has been detected.

The behavior of the break detector state machine and the break flag is shown on the Figure 233: Break detection in LIN mode (11-bit break length - LBDL bit is set) on page 650.

Examples of break frames are given on *Figure 234: Break detection in LIN mode vs. Framing error detection on page 651.* 

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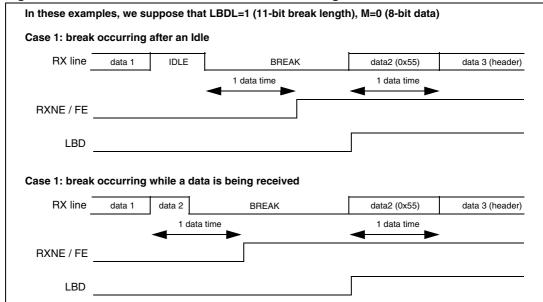


Figure 234. Break detection in LIN mode vs. Framing error detection

# 24.3.9 USART synchronous mode

The synchronous mode is selected by writing the CLKEN bit in the USART\_CR2 register to 1. In synchronous mode, the following bits must be kept cleared:

- LINEN bit in the USART\_CR2 register,
- SCEN, HDSEL and IREN bits in the USART\_CR3 register.

The USART allows the user to control a bidirectional synchronous serial communications in master mode. The SCLK pin is the output of the USART transmitter clock. No clock pulses are sent to the SCLK pin during start bit and stop bit. Depending on the state of the LBCL bit in the USART\_CR2 register clock pulses will or will not be generated during the last valid data bit (address mark). The CPOL bit in the USART\_CR2 register allows the user to select the clock polarity, and the CPHA bit in the USART\_CR2 register allows the user to select the phase of the external clock (see *Figure 235*, *Figure 236* & *Figure 237*).

During the Idle state, preamble and send break, the external SCLK clock is not activated.

In synchronous mode the USART transmitter works exactly like in asynchronous mode. But as SCLK is synchronized with TX (according to CPOL and CPHA), the data on TX is synchronous.

In this mode the USART receiver works in a different manner compared to the asynchronous mode. If RE=1, the data is sampled on SCLK (rising or falling edge, depending on CPOL and CPHA), without any oversampling. A setup and a hold time must be respected (which depends on the baud rate: 1/16 bit time).

Note:

The SCLK pin works in conjunction with the TX pin. Thus, the clock is provided only if the transmitter is enabled (TE=1) and a data is being transmitted (the data register USART\_DR

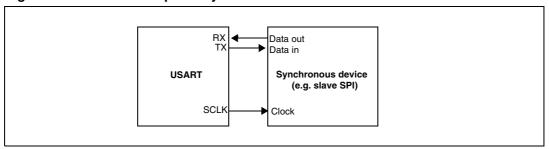
has been written). This means that it is not possible to receive a synchronous data without transmitting data.

The LBCL, CPOL and CPHA bits have to be selected when both the transmitter and the receiver are disabled (TE=RE=0) to ensure that the clock pulses function correctly. These bits should not be changed while the transmitter or the receiver is enabled.

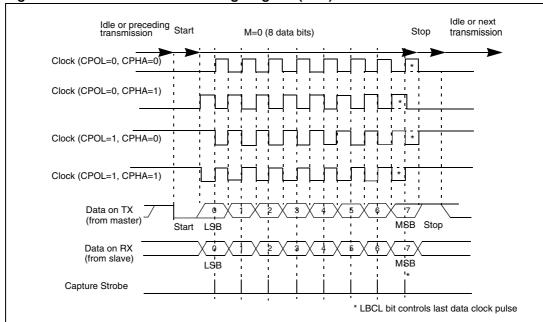
It is advised that TE and RE are set in the same instruction in order to minimize the setup and the hold time of the receiver.

The USART supports master mode only: it cannot receive or send data related to an input clock (SCLK is always an output).

Figure 235. USART example of synchronous transmission







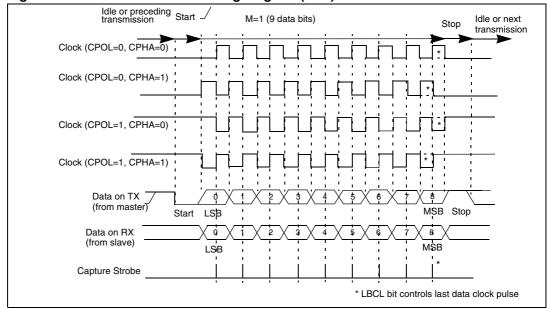
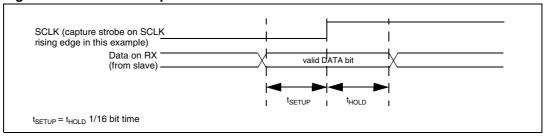


Figure 237. USART data clock timing diagram (M=1)

Figure 238. RX data setup/hold time



Note:

The function of SCLK is different in Smartcard mode. Refer to the Smartcard mode chapter for more details.

# 24.3.10 Single-wire half-duplex communication

The single-wire half-duplex mode is selected by setting the HDSEL bit in the USART\_CR3 register. In this mode, the following bits must be kept cleared:

- LINEN and CLKEN bits in the USART\_CR2 register,
- SCEN and IREN bits in the USART\_CR3 register.

The USART can be configured to follow a single-wire half-duplex protocol where the TX and RX lines are internally connected. The selection between half- and full-duplex communication is made with a control bit 'HALF DUPLEX SEL' (HDSEL in USART\_CR3).

As soon as HDSEL is written to 1:

- the TX and RX lines are internally connected
- the RX pin is no longer used
- the TX pin is always released when no data is transmitted. Thus, it acts as a standard I/O in idle or in reception. It means that the I/O must be configured so that TX is configured as floating input (or output high open-drain) when not driven by the USART.

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Apart from this, the communications are similar to what is done in normal USART mode. The conflicts on the line must be managed by the software (by the use of a centralized arbiter, for instance). In particular, the transmission is never blocked by hardware and continue to occur as soon as a data is written in the data register while the TE bit is set.

### 24.3.11 Smartcard

The Smartcard mode is selected by setting the SCEN bit in the USART\_CR3 register. In smartcard mode, the following bits must be kept cleared:

- LINEN bit in the USART\_CR2 register,
- HDSEL and IREN bits in the USART\_CR3 register.

Moreover, the CLKEN bit may be set in order to provide a clock to the smartcard.

The Smartcard interface is designed to support asynchronous protocol Smartcards as defined in the ISO 7816-3 standard. The USART should be configured as:

- 8 bits plus parity: where M=1 and PCE=1 in the USART CR1 register
- 1.5 stop bits when transmitting and receiving: where STOP=11 in the USART\_CR2 register.

Note:

It is also possible to choose 0.5 stop bit for receiving but it is recommended to use 1.5 stop bits for both transmitting and receiving to avoid switching between the two configurations.

Figure 239 shows examples of what can be seen on the data line with and without parity error.

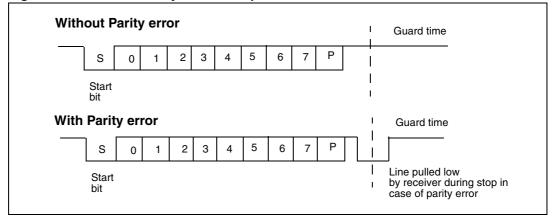


Figure 239. ISO 7816-3 asynchronous protocol

When connected to a Smartcard, the TX output of the USART drives a bidirectional line that is also driven by the Smartcard. The TX pin must be configured as open-drain.

Smartcard is a single wire half duplex communication protocol.

- Transmission of data from the transmit shift register is guaranteed to be delayed by a minimum of 1/2 baud clock. In normal operation a full transmit shift register will start shifting on the next baud clock edge. In Smartcard mode this transmission is further delayed by a guaranteed 1/2 baud clock.
- If a parity error is detected during reception of a frame programmed with a 0.5 or 1.5 stop bit period, the transmit line is pulled low for a baud clock period after the completion of the receive frame. This is to indicate to the Smartcard that the data transmitted to USART has not been correctly received. This NACK signal (pulling transmit line low for 1 baud clock) will cause a framing error on the transmitter side

(configured with 1.5 stop bits). The application can handle re-sending of data according to the protocol. A parity error is 'NACK'ed by the receiver if the NACK control bit is set, otherwise a NACK is not transmitted.

- The assertion of the TC flag can be delayed by programming the Guard Time register. In normal operation, TC is asserted when the transmit shift register is empty and no further transmit requests are outstanding. In Smartcard mode an empty transmit shift register triggers the guard time counter to count up to the programmed value in the Guard Time register. TC is forced low during this time. When the guard time counter reaches the programmed value TC is asserted high.
- The de-assertion of TC flag is unaffected by Smartcard mode.
- If a framing error is detected on the transmitter end (due to a NACK from the receiver), the NACK will not be detected as a start bit by the receive block of the transmitter.
   According to the ISO protocol, the duration of the received NACK can be 1 or 2 baud clock periods.
- On the receiver side, if a parity error is detected and a NACK is transmitted the receiver will not detect the NACK as a start bit.

Note: A break character is not significant in Smartcard mode. A 0x00 data with a framing error will be treated as data and not as a break.

No Idle frame is transmitted when toggling the TE bit. The Idle frame (as defined for the other configurations) is not defined by the ISO protocol.

*Figure 240* details how the NACK signal is sampled by the USART. In this example the USART is transmitting a data and is configured with 1.5 stop bits. The receiver part of the USART is enabled in order to check the integrity of the data and the NACK signal.

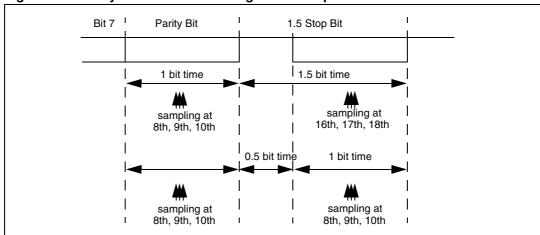


Figure 240. Parity error detection using the 1.5 stop bits

The USART can provide a clock to the smartcard through the SCLK output. In smartcard mode, SCLK is not associated to the communication but is simply derived from the internal peripheral input clock through a 5-bit prescaler. The division ratio is configured in the prescaler register USART\_GTPR. SCLK frequency can be programmed from  $f_{CK}/2$  to  $f_{CK}/62$ , where  $f_{CK}$  is the peripheral input clock.

### 24.3.12 IrDA SIR ENDEC block

The IrDA mode is selected by setting the IREN bit in the USART\_CR3 register. In IrDA mode, the following bits must be kept cleared:

- LINEN, STOP and CLKEN bits in the USART\_CR2 register,
- SCEN and HDSEL bits in the USART\_CR3 register.

The IrDA SIR physical layer specifies use of a Return to Zero, Inverted (RZI) modulation scheme that represents logic 0 as an infrared light pulse (see *Figure 241*).

The SIR Transmit encoder modulates the Non Return to Zero (NRZ) transmit bit stream output from USART. The output pulse stream is transmitted to an external output driver and infrared LED. USART supports only bit rates up to 115.2Kbps for the SIR ENDEC. In normal mode the transmitted pulse width is specified as 3/16 of a bit period.

The SIR receive decoder demodulates the return-to-zero bit stream from the infrared detector and outputs the received NRZ serial bit stream to USART. The decoder input is normally HIGH (marking state) in the Idle state. The transmit encoder output has the opposite polarity to the decoder input. A start bit is detected when the decoder input is low.

- IrDA is a half duplex communication protocol. If the Transmitter is busy (i.e. the USART is sending data to the IrDA encoder), any data on the IrDA receive line will be ignored by the IrDA decoder and if the Receiver is busy (USART is receiving decoded data from the USART), data on the TX from the USART to IrDA will not be encoded by IrDA. While receiving data, transmission should be avoided as the data to be transmitted could be corrupted.
- A '0 is transmitted as a high pulse and a '1 is transmitted as a '0. The width of the pulse is specified as 3/16th of the selected bit period in normal mode (see *Figure 242*).
- The SIR decoder converts the IrDA compliant receive signal into a bit stream for USART.
- The SIR receive logic interprets a high state as a logic one and low pulses as logic zeros.
- The transmit encoder output has the opposite polarity to the decoder input. The SIR output is in low state when Idle.
- The IrDA specification requires the acceptance of pulses greater than 1.41 us. The acceptable pulse width is programmable. Glitch detection logic on the receiver end filters out pulses of width less than 2 PSC periods (PSC is the prescaler value programmed in the IrDA low-power Baud Register, USART\_GTPR). Pulses of width less than 1 PSC period are always rejected, but those of width greater than one and less than two periods may be accepted or rejected, those greater than 2 periods will be accepted as a pulse. The IrDA encoder/decoder doesn't work when PSC=0.
- The receiver can communicate with a low-power transmitter.
- In IrDA mode, the STOP bits in the USART\_CR2 register must be configured to "1 stop bit".

## IrDA low-power mode

### Transmitter:

In low-power mode the pulse width is not maintained at 3/16 of the bit period. Instead, the width of the pulse is 3 times the low-power baud rate which can be a minimum of 1.42 MHz. Generally this value is 1.8432 MHz (1.42 MHz < PSC< 2.12 MHz). A low-power mode programmable divisor divides the system clock to achieve this value.

### Receiver:

Receiving in low-power mode is similar to receiving in normal mode. For glitch detection the USART should discard pulses of duration shorter than 1/PSC. A valid low is accepted only if its duration is greater than 2 periods of the IrDA low-power Baud clock (PSC value in USART\_GTPR).

Note:

A pulse of width less than two and greater than one PSC period(s) may or may not be rejected.

The receiver set up time should be managed by software. The IrDA physical layer specification specifies a minimum of 10 ms delay between transmission and reception (IrDA is a half duplex protocol).

Figure 241. IrDA SIR ENDEC- block diagram

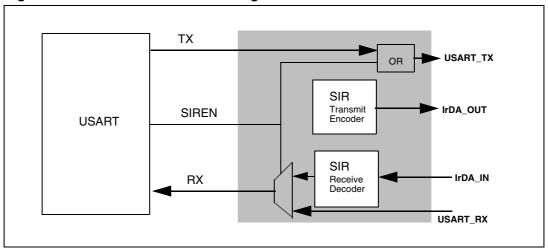
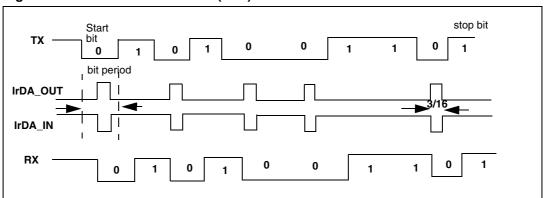


Figure 242. IrDA data modulation (3/16) -Normal mode



## 24.3.13 Continuous communication using DMA

The USART is capable of continuous communication using the DMA. The DMA requests for Rx buffer and Tx buffer are generated independently.

Note:

You should refer to product specs for availability of the DMA controller. If DMA is not available in the product, you should use the USART as explained in Section 24.3.2 or 24.3.3. In the USART\_SR register, you can clear the TXE/RXNE flags to achieve continuous communication.

### **Transmission using DMA**

DMA mode can be enabled for transmission by setting DMAT bit in the USART\_CR3 register. Data is loaded from a SRAM area configured using the DMA peripheral (refer to the DMA specification) to the USART\_DR register whenever the TXE bit is set. To map a DMA channel for USART transmission, use the following procedure (x denotes the channel number):

- Write the USART\_DR register address in the DMA control register to configure it as the destination of the transfer. The data will be moved to this address from memory after each TXE event.
- Write the memory address in the DMA control register to configure it as the source of the transfer. The data will be loaded into the USART\_DR register from this memory area after each TXE event.
- 3. Configure the total number of bytes to be transferred to the DMA control register.
- 4. Configure the channel priority in the DMA register
- Configure DMA interrupt generation after half/ full transfer as required by the application.
- 6. Clear the TC bit in the SR register by writing 0 to it.
- Activate the channel in the DMA register.

When the number of data transfers programmed in the DMA Controller is reached, the DMA controller generates an interrupt on the DMA channel interrupt vector.

In transmission mode, once the DMA has written all the data to be transmitted (the TCIF flag is set in the DMA\_ISR register), the TC flag can be monitored to make sure that the USART communication is complete. This is required to avoid corrupting the last transmission before disabling the USART or entering the Stop mode. The software must wait until TC=1. The TC flag remains cleared during all data transfers and it is set by hardware at the last frame's end of transmission.

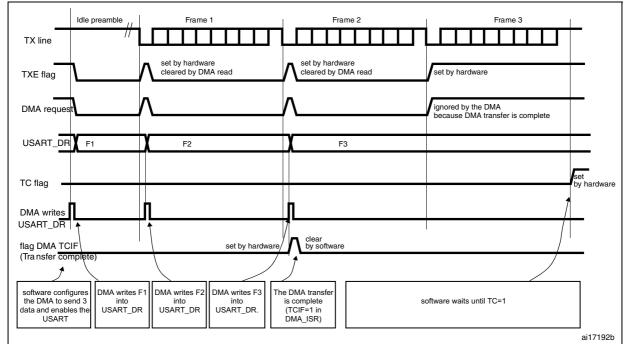


Figure 243. Transmission using DMA

# **Reception using DMA**

DMA mode can be enabled for reception by setting the DMAR bit in USART\_CR3 register. Data is loaded from the USART\_DR register to a SRAM area configured using the DMA peripheral (refer to the DMA specification) whenever a data byte is received. To map a DMA channel for USART reception, use the following procedure:

- Write the USART\_DR register address in the DMA control register to configure it as the source of the transfer. The data will be moved from this address to the memory after each RXNE event.
- 2. Write the memory address in the DMA control register to configure it as the destination of the transfer. The data will be loaded from USART\_DR to this memory area after each RXNE event.
- 3. Configure the total number of bytes to be transferred in the DMA control register.
- 4. Configure the channel priority in the DMA control register
- 5. Configure interrupt generation after half/ full transfer as required by the application.
- 6. Activate the channel in the DMA control register.

When the number of data transfers programmed in the DMA Controller is reached, the DMA controller generates an interrupt on the DMA channel interrupt vector. The DMAR bit should be cleared by software in the USART\_CR3 register during the interrupt subroutine.

Note: If DMA is used for reception, do not enable the RXNEIE bit.

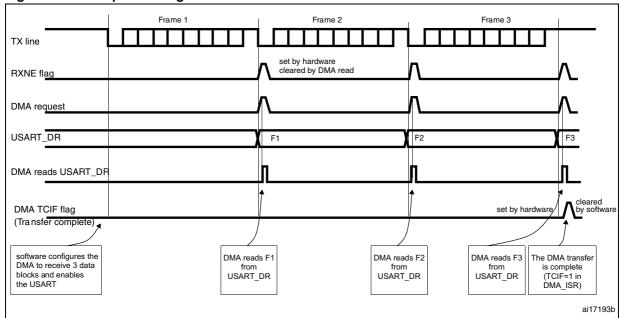


Figure 244. Reception using DMA

# Error flagging and interrupt generation in multibuffer communication

In case of multibuffer communication if any error occurs during the transaction the error flag will be asserted after the current byte. An interrupt will be generated if the interrupt enable flag is set. For framing error, overrun error and noise flag which are asserted with RXNE in case of single byte reception, there will be separate error flag interrupt enable bit (EIE bit in the USART\_CR3 register), which if set will issue an interrupt after the current byte with either of these errors.

### 24.3.14 Hardware flow control

It is possible to control the serial data flow between 2 devices by using the nCTS input and the nRTS output. The *Figure 245* shows how to connect 2 devices in this mode:

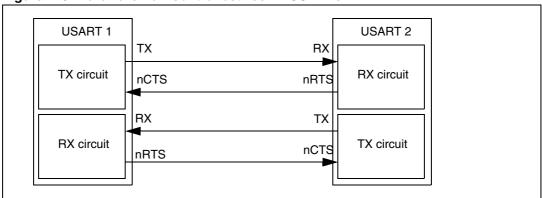


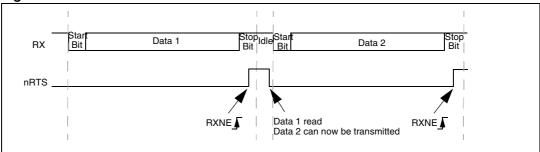
Figure 245. Hardware flow control between 2 USARTs

RTS and CTS flow control can be enabled independently by writing respectively RTSE and CTSE bits to 1 (in the USART\_CR3 register).

#### RTS flow control

If the RTS flow control is enabled (RTSE=1), then nRTS is asserted (tied low) as long as the USART receiver is ready to receive a new data. When the receive register is full, nRTS is deasserted, indicating that the transmission is expected to stop at the end of the current frame. *Figure 246* shows an example of communication with RTS flow control enabled.

Figure 246. RTS flow control



### CTS flow control

If the CTS flow control is enabled (CTSE=1), then the transmitter checks the nCTS input before transmitting the next frame. If nCTS is asserted (tied low), then the next data is transmitted (assuming that a data is to be transmitted, in other words, if TXE=0), else the transmission does not occur. When nCTS is deasserted during a transmission, the current transmission is completed before the transmitter stops.

When CTSE=1, the CTSIF status bit is automatically set by hardware as soon as the nCTS input toggles. It indicates when the receiver becomes ready or not ready for communication. An interrupt is generated if the CTSIE bit in the USART\_CR3 register is set. The figure below shows an example of communication with CTS flow control enabled.

CTS CTS A nCTS Transmit data register Data 3 TDR Data 2 empty empty Data 2 Data 1 Data 3 Bit Bit Bit Writing data 3 in TDR Transmission of Data 3 is delayed until nCTS = 0

Figure 247. CTS flow control

Note:

**Special behavior of break frames:** when the CTS flow is enabled, the transmitter does not check the nCTS input state to send a break.

#### **USART** interrupts 24.4

Table 100. USART interrupt requests

Interrupt event	Event flag	Enable control bit
Transmit Data Register Empty	TXE	TXEIE
CTS flag	CTS	CTSIE
Transmission Complete	TC	TCIE
Received Data Ready to be Read	RXNE	RXNEIE
Overrun Error Detected	ORE	HAINEIE
Idle Line Detected	IDLE	IDLEIE
Parity Error	PE	PEIE
Break Flag	LBD	LBDIE
Noise Flag, Overrun error and Framing Error in multibuffer communication	NF or ORE or FE	EIE

The USART interrupt events are connected to the same interrupt vector (see Figure 248).

- During transmission: Transmission Complete, Clear to Send or Transmit Data Register empty interrupt.
- While receiving: Idle Line detection, Overrun error, Receive Data register not empty, Parity error, LIN break detection, Noise Flag (only in multi buffer communication) and Framing Error (only in multi buffer communication).

These events generate an interrupt if the corresponding Enable Control Bit is set.

**USART** IDLE IDLEIE interrupt **RXNEIE** ORE RXNEIE RXNE PEIE LBD LBDIE EIE DMAR

Figure 248. USART interrupt mapping diagram

# 24.5 USART mode configuration

Table 101. USART mode configuration<sup>(1)</sup>

USART modes	USART1	USART2	USART3	UART4	UART5	USART6
Asynchronous mode	Х	Х	Х	Х	Х	Х
Hardware flow control	Х	Х	Х	NA	NA	Х
Multibuffer communication (DMA)	Х	Х	Х	Х	Х	Х
Multiprocessor communication	Х	Х	Х	Х	Х	Х
Synchronous	Х	Х	Х	NA	NA	Х
Smartcard	Х	Х	Х	NA	NA	Х
Half-duplex (single-wire mode)	Х	Х	Х	Х	Х	Х
IrDA	Х	Х	Х	Х	Х	Х
LIN	Х	Х	Х	Х	Х	Х

<sup>1.</sup> X = supported; NA = not applicable.

# 24.6 USART registers

Refer to Section 1.1 on page 46 for a list of abbreviations used in register descriptions.

The peripheral registers can be accessed by half-words (16 bits) or words (32 bits).

# 24.6.1 Status register (USART\_SR)

Address offset: 0x00

Reset value: 0x00C0 0000

31	30	29	28	27	26	26 25		23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		Rese	n rod			CTS	LBD	TXE	TC	RXNE	IDLE	ORE	NF	FE	PE
		nese	erveu			rc_w0	rc_w0	r	rc_w0	rc_w0	r	r	r	r	r

Bits 31:10 Reserved, must be kept at reset value

Bit 9 CTS: CTS flag

This bit is set by hardware when the nCTS input toggles, if the CTSE bit is set. It is cleared by software (by writing it to 0). An interrupt is generated if CTSIE=1 in the USART\_CR3 register.

0: No change occurred on the nCTS status line

1: A change occurred on the nCTS status line

Note: This bit is not available for UART4 & UART5.

### Bit 8 LBD: LIN break detection flag

This bit is set by hardware when the LIN break is detected. It is cleared by software (by writing it to 0). An interrupt is generated if LBDIE = 1 in the USART\_CR2 register.

0: LIN Break not detected

1: LIN break detected

Note: An interrupt is generated when LBD=1 if LBDIE=1

#### Bit 7 TXE: Transmit data register empty

This bit is set by hardware when the content of the TDR register has been transferred into the shift register. An interrupt is generated if the TXEIE bit =1 in the USART\_CR1 register. It is cleared by a write to the USART\_DR register.

- 0: Data is not transferred to the shift register
- 1: Data is transferred to the shift register)

Note: This bit is used during single buffer transmission.

#### Bit 6 TC: Transmission complete

This bit is set by hardware if the transmission of a frame containing data is complete and if TXE is set. An interrupt is generated if TCIE=1 in the USART\_CR1 register. It is cleared by a software sequence (a read from the USART\_SR register followed by a write to the USART\_DR register). The TC bit can also be cleared by writing a '0' to it. This clearing sequence is recommended only for multibuffer communication.

- 0: Transmission is not complete
- 1: Transmission is complete

#### Bit 5 RXNE: Read data register not empty

This bit is set by hardware when the content of the RDR shift register has been transferred to the USART\_DR register. An interrupt is generated if RXNEIE=1 in the USART\_CR1 register. It is cleared by a read to the USART\_DR register. The RXNE flag can also be cleared by writing a zero to it. This clearing sequence is recommended only for multibuffer communication.

- 0: Data is not received
- 1: Received data is ready to be read.

### Bit 4 IDLE: IDLE line detected

This bit is set by hardware when an Idle Line is detected. An interrupt is generated if the IDLEIE=1 in the USART\_CR1 register. It is cleared by a software sequence (an read to the USART\_SR register followed by a read to the USART\_DR register).

- 0: No Idle Line is detected
- 1: Idle Line is detected

Note: The IDLE bit will not be set again until the RXNE bit has been set itself (i.e. a new idle line occurs).

## Bit 3 ORE: Overrun error

This bit is set by hardware when the word currently being received in the shift register is ready to be transferred into the RDR register while RXNE=1. An interrupt is generated if RXNEIE=1 in the USART\_CR1 register. It is cleared by a software sequence (an read to the USART\_SR register followed by a read to the USART\_DR register).

- 0: No Overrun error
- 1: Overrun error is detected

Note: When this bit is set, the RDR register content will not be lost but the shift register will be overwritten. An interrupt is generated on ORE flag in case of Multi Buffer communication if the EIE bit is set.

#### Bit 2 NF: Noise detected flag

This bit is set by hardware when noise is detected on a received frame. It is cleared by a software sequence (an read to the USART\_SR register followed by a read to the USART\_DR register).

0: No noise is detected

1: Noise is detected

Note: This bit does not generate interrupt as it appears at the same time as the RXNE bit which itself generates an interrupting interrupt is generated on NF flag in case of Multi Buffer communication if the EIE bit is set.

Note: When the line is noise-free, the NF flag can be disabled by programming the ONEBIT bit to 1 to increase the USART tolerance to deviations (Refer to Section 24.3.5: USART receiver tolerance to clock deviation on page 645).

### Bit 1 FE: Framing error

This bit is set by hardware when a de-synchronization, excessive noise or a break character is detected. It is cleared by a software sequence (an read to the USART\_SR register followed by a read to the USART\_DR register).

0: No Framing error is detected

1: Framing error or break character is detected

Note: This bit does not generate interrupt as it appears at the same time as the RXNE bit which itself generates an interrupt. If the word currently being transferred causes both frame error and overrun error, it will be transferred and only the ORE bit will be set.

An interrupt is generated on FE flag in case of Multi Buffer communication if the EIE bit is set.

### Bit 0 PE: Parity error

This bit is set by hardware when a parity error occurs in receiver mode. It is cleared by a software sequence (a read from the status register followed by a read or write access to the USART\_DR data register). The software must wait for the RXNE flag to be set before clearing the PE bit.

An interrupt is generated if PEIE = 1 in the USART\_CR1 register.

0: No parity error1: Parity error

# 24.6.2 Data register (USART\_DR)

Address offset: 0x04

Reset value: 0xXXXX XXXX

Bits 31:9 Reserved, must be kept at reset value

### Bits 8:0 DR[8:0]: Data value

Contains the Received or Transmitted data character, depending on whether it is read from or written to.

The Data register performs a double function (read and write) since it is composed of two registers, one for transmission (TDR) and one for reception (RDR)

The TDR register provides the parallel interface between the internal bus and the output shift register (see Figure 1).

The RDR register provides the parallel interface between the input shift register and the internal bus.

When transmitting with the parity enabled (PCE bit set to 1 in the USART\_CR1 register), the value written in the MSB (bit 7 or bit 8 depending on the data length) has no effect because it is replaced by the parity.

When receiving with the parity enabled, the value read in the MSB bit is the received parity bit.

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# 24.6.3 Baud rate register (USART\_BRR)

Note: The baud counters stop counting if the TE or RE bits are disabled respectively.

Address offset: 0x08

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
				ı	DIV_Mant	tissa[11:0	]						DIV_Frac	ction[3:0]	
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:16 Reserved, must be kept at reset value

### Bits 15:4 DIV\_Mantissa[11:0]: mantissa of USARTDIV

These 12 bits define the mantissa of the USART Divider (USARTDIV)

### Bits 3:0 DIV\_Fraction[3:0]: fraction of USARTDIV

These 4 bits define the fraction of the USART Divider (USARTDIV). When OVER8=1, the DIV\_Fraction3 bit is not considered and must be kept cleared.

# 24.6.4 Control register 1 (USART\_CR1)

Address offset: 0x0C

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	ved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
OVER8	Reserved	UE	М	WAKE	PCE	PS	PEIE	TXEIE	TCIE	RXNEIE	IDLEIE	TE	RE	RWU	SBK
rw	Res.	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:16 Reserved, must be kept at reset value

Bit 15 OVER8: Oversampling mode

0: oversampling by 16 1: oversampling by 8

Note: Oversampling by 8 is not available in the Smartcard, IrDA and LIN modes: when SCEN=1,IREN=1 or LINEN=1 then OVER8 is forced to '0 by hardware.

Bit 14 Reserved, must be kept at reset value

### Bit 13 UE: USART enable

When this bit is cleared the USART prescalers and outputs are stopped and the end of the current

byte transfer in order to reduce power consumption. This bit is set and cleared by software.

0: USART prescaler and outputs disabled

1: USART enabled

## Bit 12 M: Word length

This bit determines the word length. It is set or cleared by software.

0: 1 Start bit, 8 Data bits, n Stop bit

1: 1 Start bit, 9 Data bits, n Stop bit

Note: The M bit must not be modified during a data transfer (both transmission and reception)

#### Bit 11 WAKE: Wakeup method

This bit determines the USART wakeup method, it is set or cleared by software.

0: Idle Line

1: Address Mark

### Bit 10 PCE: Parity control enable

This bit selects the hardware parity control (generation and detection). When the parity control is enabled, the computed parity is inserted at the MSB position (9th bit if M=1; 8th bit if M=0) and parity is checked on the received data. This bit is set and cleared by software. Once it is set, PCE is active after the current byte (in reception and in transmission).

0: Parity control disabled

1: Parity control enabled

#### Bit 9 PS: Parity selection

This bit selects the odd or even parity when the parity generation/detection is enabled (PCE bit set). It is set and cleared by software. The parity will be selected after the current byte.

0: Even parity

1: Odd parity

### Bit 8 PEIE: PE interrupt enable

This bit is set and cleared by software.

0: Interrupt is inhibited

1: An USART interrupt is generated whenever PE=1 in the USART\_SR register

### Bit 7 TXEIE: TXE interrupt enable

This bit is set and cleared by software.

0: Interrupt is inhibited

1: An USART interrupt is generated whenever TXE=1 in the USART\_SR register

### Bit 6 TCIE: Transmission complete interrupt enable

This bit is set and cleared by software.

0: Interrupt is inhibited

1: An USART interrupt is generated whenever TC=1 in the USART\_SR register

### Bit 5 RXNEIE: RXNE interrupt enable

This bit is set and cleared by software.

0: Interrupt is inhibited

1: An USART interrupt is generated whenever ORE=1 or RXNE=1 in the USART\_SR register

### Bit 4 IDLEIE: IDLE interrupt enable

This bit is set and cleared by software.

0: Interrupt is inhibited

1: An USART interrupt is generated whenever IDLE=1 in the USART\_SR register

### Bit 3 TE: Transmitter enable

This bit enables the transmitter. It is set and cleared by software.

0: Transmitter is disabled

1: Transmitter is enabled

Note: 1: During transmission, a "0" pulse on the TE bit ("0" followed by "1") sends a preamble (idle line) after the current word, except in smartcard mode.

2: When TE is set there is a 1 bit-time delay before the transmission starts.

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#### Bit 2 RE: Receiver enable

This bit enables the receiver. It is set and cleared by software.

- 0: Receiver is disabled
- 1: Receiver is enabled and begins searching for a start bit

### Bit 1 RWU: Receiver wakeup

This bit determines if the USART is in mute mode or not. It is set and cleared by software and can be cleared by hardware when a wakeup sequence is recognized.

- 0: Receiver in active mode
- 1: Receiver in mute mode

Note: 1: Before selecting Mute mode (by setting the RWU bit) the USART must first receive a data byte, otherwise it cannot function in Mute mode with wakeup by Idle line detection.

2: In Address Mark Detection wakeup configuration (WAKE bit=1) the RWU bit cannot be modified by software while the RXNE bit is set.

### Bit 0 SBK: Send break

This bit set is used to send break characters. It can be set and cleared by software. It should be set by software, and will be reset by hardware during the stop bit of break.

- 0: No break character is transmitted
- 1: Break character will be transmitted

# 24.6.5 Control register 2 (USART\_CR2)

Address offset: 0x10

Reset value: 0x0000 0000

31	30	29	29 28 27		26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Res.	LINEN	STO	P[1:0]	1:0] CLKEN		СРНА	LBCL	Res.	LBDIE	LBDL	Res.		ADD	[3:0]	
nes.	rw	rw	rw	rw	rw	rw	rw		rw	rw	rw	rw	rw	rw	rw

Bits 31:15 Reserved, must be kept at reset value

### Bit 14 LINEN: LIN mode enable

This bit is set and cleared by software.

0: LIN mode disabled

1: LIN mode enabled

The LIN mode enables the capability to send LIN Synch Breaks (13 low bits) using the SBK bit in the USART\_CR1 register, and to detect LIN Sync breaks.

### Bits 13:12 STOP: STOP bits

These bits are used for programming the stop bits.

00: 1 Stop bit 01: 0.5 Stop bit 10: 2 Stop bits 11: 1.5 Stop bit

Note: The 0.5 Stop bit and 1.5 Stop bit are not available for UART4 & UART5.

### Bit 11 CLKEN: Clock enable

This bit allows the user to enable the SCLK pin.

0: SCLK pin disabled

1: SCLK pin enabled

This bit is not available for UART4 & UART5.

### Bit 10 CPOL: Clock polarity

This bit allows the user to select the polarity of the clock output on the SCLK pin in synchronous mode. It works in conjunction with the CPHA bit to produce the desired clock/data relationship

0: Steady low value on SCLK pin outside transmission window.

1: Steady high value on SCLK pin outside transmission window.

This bit is not available for UART4 & UART5.

### Bit 9 CPHA: Clock phase

This bit allows the user to select the phase of the clock output on the SCLK pin in synchronous mode. It works in conjunction with the CPOL bit to produce the desired clock/data relationship (see figures 236 to 237)

0: The first clock transition is the first data capture edge

1: The second clock transition is the first data capture edge

Note: This bit is not available for UART4 & UART5.

Bit 8 LBCL: Last bit clock pulse

This bit allows the user to select whether the clock pulse associated with the last data bit transmitted (MSB) has to be output on the SCLK pin in synchronous mode.

0: The clock pulse of the last data bit is not output to the SCLK pin

1: The clock pulse of the last data bit is output to the SCLK pin

Note: 1: The last bit is the 8th or 9th data bit transmitted depending on the 8 or 9 bit format selected by the M bit in the USART\_CR1 register.

2: This bit is not available for UART4 & UART5.

Bit 7 Reserved, must be kept at reset value

Bit 6 LBDIE: LIN break detection interrupt enable

Break interrupt mask (break detection using break delimiter).

0: Interrupt is inhibited

1: An interrupt is generated whenever LBD=1 in the USART\_SR register

Bit 5 LBDL: lin break detection length

This bit is for selection between 11 bit or 10 bit break detection.

0: 10-bit break detection1: 11-bit break detection

Bit 4 Reserved, must be kept at reset value

Bits 3:0 ADD[3:0]: Address of the USART node

This bit-field gives the address of the USART node.

This is used in multiprocessor communication during mute mode, for wake up with address mark detection.

Note: These 3 bits (CPOL, CPHA, LBCL) should not be written while the transmitter is enabled.

# 24.6.6 Control register 3 (USART\_CR3)

Address offset: 0x14

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	1/	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Reserved			ONEBIT	CTSIE	CTSE	RTSE	DMAT	DMAR	SCEN	NACK	HDSEL	IRLP	IREN	EIE	
			rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	

Bits 31:12 Reserved, must be kept at reset value

Bit 11 **ONEBIT**: One sample bit method enable

This bit allows the user to select the sample method. When the one sample bit method is selected the noise detection flag (NF) is disabled.

0: Three sample bit method

1: One sample bit method

Bit 10 CTSIE: CTS interrupt enable

0: Interrupt is inhibited

1: An interrupt is generated whenever CTS=1 in the USART\_SR register

Note: This bit is not available for UART4 & UART5.

#### Bit 9 CTSE: CTS enable

0: CTS hardware flow control disabled

1: CTS mode enabled, data is only transmitted when the nCTS input is asserted (tied to 0). If the nCTS input is deasserted while a data is being transmitted, then the transmission is completed before stopping. If a data is written into the data register while nCTS is asserted, the transmission is postponed until nCTS is asserted.

Note: This bit is not available for UART4 & UART5.

#### Bit 8 RTSE: RTS enable

0: RTS hardware flow control disabled

1: RTS interrupt enabled, data is only requested when there is space in the receive buffer. The transmission of data is expected to cease after the current character has been transmitted. The nRTS output is asserted (tied to 0) when a data can be received.

Note: This bit is not available for UART4 & UART5.

### Bit 7 DMAT: DMA enable transmitter

This bit is set/reset by software

- 1: DMA mode is enabled for transmission.
- 0: DMA mode is disabled for transmission.

#### Bit 6 DMAR: DMA enable receiver

This bit is set/reset by software

- 1: DMA mode is enabled for reception
- 0: DMA mode is disabled for reception

### Bit 5 SCEN: Smartcard mode enable

This bit is used for enabling Smartcard mode.

- 0: Smartcard Mode disabled
- 1: Smartcard Mode enabled

Note: This bit is not available for UART4 & UART5.

#### Bit 4 NACK: Smartcard NACK enable

- 0: NACK transmission in case of parity error is disabled
- 1: NACK transmission during parity error is enabled

Note: This bit is not available for UART4 & UART5.

### Bit 3 HDSEL: Half-duplex selection

Selection of Single-wire Half-duplex mode

- 0: Half duplex mode is not selected
- 1: Half duplex mode is selected

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### Bit 2 IRLP: IrDA low-power

This bit is used for selecting between normal and low-power IrDA modes

0: Normal mode

1: Low-power mode

#### Bit 1 IREN: IrDA mode enable

This bit is set and cleared by software.

0: IrDA disabled

1: IrDA enabled

### Bit 0 EIE: Error interrupt enable

Error Interrupt Enable Bit is required to enable interrupt generation in case of a framing error, overrun error or noise flag (FE=1 or ORE=1 or NF=1 in the USART\_SR register) in case of Multi Buffer Communication (DMAR=1 in the USART\_CR3 register).

0: Interrupt is inhibited

1: An interrupt is generated whenever DMAR=1 in the USART\_CR3 register and FE=1 or ORE=1 or NF=1 in the USART\_SR register.

# 24.6.7 Guard time and prescaler register (USART\_GTPR)

Address offset: 0x18

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	11 10		8	7	7 6		4	3	2	1	0
			GT[	7:0]							PSC	[7:0]			
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:16 Reserved, must be kept at reset value

Bits 15:8 GT[7:0]: Guard time value

This bit-field gives the Guard time value in terms of number of baud clocks.

This is used in Smartcard mode. The Transmission Complete flag is set after this guard time value

Note: This bit is not available for UART4 & UART5.

### Bits 7:0 PSC[7:0]: Prescaler value

### - In IrDA Low-power mode:

PSC[7:0] = IrDA Low-Power Baud Rate

Used for programming the prescaler for dividing the system clock to achieve the low-power frequency:

The source clock is divided by the value given in the register (8 significant bits):

0000000: Reserved - do not program this value

00000001: divides the source clock by 1 00000010: divides the source clock by 2  $\,$ 

...

- In normal IrDA mode: PSC must be set to 00000001.

#### - In smartcard mode:

PSC[4:0]: Prescaler value

Used for programming the prescaler for dividing the system clock to provide the smartcard clock.

The value given in the register (5 significant bits) is multiplied by 2 to give the division factor of the source clock frequency:

00000: Reserved - do not program this value

00001: divides the source clock by 2 00010: divides the source clock by 4 00011: divides the source clock by 6

...

Note: 1: Bits [7:5] have no effect if Smartcard mode is used.

2: This bit is not available for UART4 & UART5.

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# 24.6.8 USART register map

The table below gives the USART register map and reset values.

Table 102. USART register map and reset values

Offset	Register	31	30	29	28	1	26	25	C 7	24	73	22	21	17	20	18	18	17	16	15		1	13	12	11	10	6	α	۱ د	,	9	2	4	က	7	-	0
0x00	USART_SR												Res	ser	ved												CTS	LBD	TXF	1 (	IC	RXNE	IDLE	ORE	NF	FE	PE
	Reset value																										0	0	) 1		1	0	0	0	0	0	0
0x04	USART_DR												R	les	erve	ed																D	R[8	:0]			
	Reset value																											0	) (	)	0	0	0	0	0	0	0
0x08	USART_BRR			Reserved DIV_Mantis															ntiss	a[1	5:4					DI		ract :0]	ion								
	Reset value																			0	(	)	0	0	0	0	0	0	) (	) [	0	0	0	0	0	0	0
0x0C	USART_CR1								Re	eserv	ved	d								OVER8	Received	2000	NE	Μ	WAKE	PCE	PS	PEIE	TXFIF	1011	ICIE	RXNEIE	IDLEIE	TE	RE	RWU	SBK
	Reset value																			0	ď	=	0	0	0	0	0	0	) (	)	0	0	0	0	0	0	0
0x10	USART_CR2								ı	Rese	erve	ec	d						LINEN				ST [1	OP :0]	CLKEN	CPOL	CPHA	LBCL	Reserved	7	LBUIE	LBDL	Reserved	,	ADD	)[3:0	[0
	Reset value																					)	0	0	0	0	0	0	ď		0	0	Ä	0	0	0	0
0x14	USART_CR3			Reserved																	ONEBIT						CTSE	RTSE	DMAT		DIMAR	SCEN	NACK	HDSEL	IRLP	IREN	EIE
	Reset value																			0 (								0	) (	)	0	0	0	0	0	0	0
0x18	USART_GTPR		Reserved													GT[7:0]													ı	PSC	(7:0	)]					
	Reset value																			0	(	)	0	0	0	0	0	0	) (	)	0	0	0	0	0	0	0

Refer to *Table 2 on page 50* for the register boundary addresses.

#### Serial peripheral interface (SPI) 25

This section applies to the whole STM32F4xx family, unless otherwise specified.

#### **SPI** introduction 25.1

The SPI interface provides two main functions, supporting either the SPI protocol or the I<sup>2</sup>S audio protocol. By default, it is the SPI function that is selected. It is possible to switch the interface from SPI to I<sup>2</sup>S by software.

The serial peripheral interface (SPI) allows half/full-duplex, synchronous, serial communication with external devices. The interface can be configured as the master and in this case it provides the communication clock (SCK) to the external slave device. The interface is also capable of operating in multimaster configuration.

It may be used for a variety of purposes, including simplex synchronous transfers on two lines with a possible bidirectional data line or reliable communication using CRC checking.

The I<sup>2</sup>S is also a synchronous serial communication interface. It can address four different audio standards including the I<sup>2</sup>S Philips standard, the MSB- and LSB-justified standards, and the PCM standard. It can operate as a slave or a master device in full-duplex mode (using 4 pins) or in half-duplex mode (using 3 pins). Master clock can be provided by the interface to an external slave component when the I<sup>2</sup>S is configured as the communication master.

#### Warning:

Since some SPI1 and SPI3/I2S3 pins may be mapped onto some pins used by the JTAG interface (SPI1\_NSS onto JTDI, SPI3\_NSS/I2S3\_WS onto JTDI and SPI3\_SCK/I2S3\_CK onto JTDO), you may either:

- map SPI/I2S onto other pins
- disable the JTAG and use the SWD interface prior to configuring the pins listed as SPI I/Os (when debugging the application) or
- disable both JTAG/SWD interfaces (for standalone applications).

For more information on the configuration of the JTAG/SWD interface pins, please refer to Section 6.3.2: I/O pin multiplexer and mapping.

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# 25.2 SPI and I<sup>2</sup>S main features

### 25.2.1 SPI features

- Full-duplex synchronous transfers on three lines
- Simplex synchronous transfers on two lines with or without a bidirectional data line
- 8- or 16-bit transfer frame format selection
- Master or slave operation
- Multimaster mode capability
- 8 master mode baud rate prescalers (f<sub>PCLK</sub>/2 max.)
- Slave mode frequency (f<sub>PCLK</sub>/2 max)
- Faster communication for both master and slave
- NSS management by hardware or software for both master and slave: dynamic change of master/slave operations
- Programmable clock polarity and phase
- Programmable data order with MSB-first or LSB-first shifting
- Dedicated transmission and reception flags with interrupt capability
- SPI bus busy status flag
- SPI TI mode
- Hardware CRC feature for reliable communication:
  - CRC value can be transmitted as last byte in Tx mode
  - Automatic CRC error checking for last received byte
- Master mode fault, overrun and CRC error flags with interrupt capability
- 1-byte transmission and reception buffer with DMA capability: Tx and Rx requests

## 25.2.2 I<sup>2</sup>S features

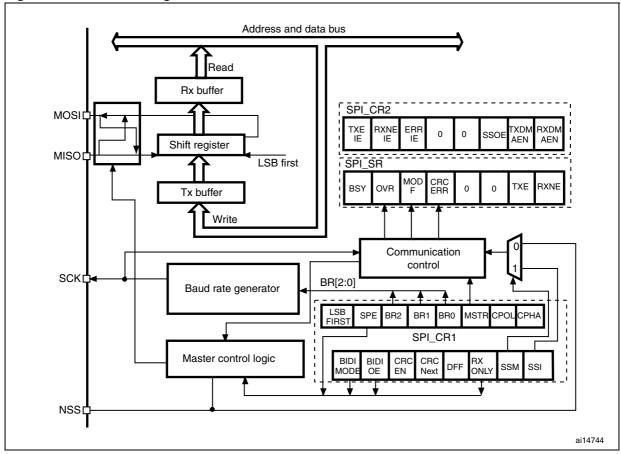
- Full duplex communication
- Half-duplex communication (only transmitter or receiver)
- Master or slave operations
- 8-bit programmable linear prescaler to reach accurate audio sample frequencies (from 8 kHz to 192 kHz)
- Data format may be 16-bit, 24-bit or 32-bit
- Packet frame is fixed to 16-bit (16-bit data frame) or 32-bit (16-bit, 24-bit, 32-bit data frame) by audio channel
- Programmable clock polarity (steady state)
- Underrun flag in slave transmission mode, overrun flag in reception mode (master and slave), and Frame Error flag in reception and transmission mode (slave only)
- 16-bit register for transmission and reception with one data register for both channel sides
- Supported I<sup>2</sup>S protocols:
  - I<sup>2</sup>S Phillps standard
  - MSB-justified standard (left-justified)
  - LSB-justified standard (right-justified)
  - PCM standard (with short and long frame synchronization on 16-bit channel frame or 16-bit data frame extended to 32-bit channel frame)
- Data direction is always MSB first
- DMA capability for transmission and reception (16-bit wide)
- Master clock may be output to drive an external audio component. Ratio is fixed at  $256 \times F_S$  (where  $F_S$  is the audio sampling frequency)
- Both I<sup>2</sup>S (I2S2 and I2S3) have a dedicated PLL (PLLI2S) to generate an even more accurate clock.
- I<sup>2</sup>S (I2S2 and I2S3) clock can be derived from an external clock mapped on the I2S\_CKIN pin.

# 25.3 SPI functional description

### 25.3.1 General description

The block diagram of the SPI is shown in Figure 249.

Figure 249. SPI block diagram



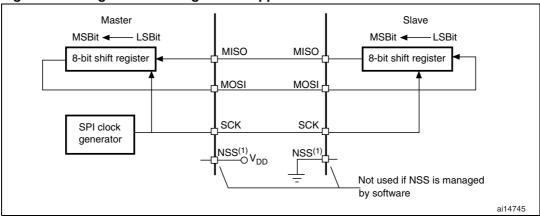
Usually, the SPI is connected to external devices through 4 pins:

- MISO: Master In / Slave Out data. This pin can be used to transmit data in slave mode and receive data in master mode.
- MOSI: Master Out / Slave In data. This pin can be used to transmit data in master mode and receive data in slave mode.
- SCK: Serial Clock output for SPI masters and input for SPI slaves.
- NSS: Slave select. This is an optional pin to select a slave device. This pin acts as a 'chip select' to let the SPI master communicate with slaves individually and to avoid contention on the data lines. Slave NSS inputs can be driven by standard IO ports on the master device. The NSS pin may also be used as an output if enabled (SSOE bit) and driven low if the SPI is in master configuration. In this manner, all NSS pins from devices connected to the Master NSS pin see a low level and become slaves when they are configured in NSS hardware mode. When configured in master mode with NSS configured as an input (MSTR=1 and SSOE=0) and if NSS is pulled low, the SPI

enters the master mode fault state: the MSTR bit is automatically cleared and the device is configured in slave mode (refer to *Section 25.3.10: Error flags on page 699*).

A basic example of interconnections between a single master and a single slave is illustrated in *Figure 250*.

Figure 250. Single master/ single slave application



1. Here, the NSS pin is configured as an input.

The MOSI pins are connected together and the MISO pins are connected together. In this way data is transferred serially between master and slave (most significant bit first).

The communication is always initiated by the master. When the master device transmits data to a slave device via the MOSI pin, the slave device responds via the MISO pin. This implies full-duplex communication with both data out and data in synchronized with the same clock signal (which is provided by the master device via the SCK pin).

### Slave select (NSS) pin management

Hardware or software slave select management can be set using the SSM bit in the SPI\_CR1 register.

- Software NSS management (SSM = 1)
   The slave select information is driven internally by the value of the SSI bit in the SPI\_CR1 register. The external NSS pin remains free for other application uses.
- Hardware NSS management (SSM = 0)
   Two configurations are possible depending on the NSS output configuration (SSOE bit in register SPI\_CR1).
  - NSS output enabled (SSM = 0, SSOE = 1)
     This configuration is used only when the device operates in master mode. The NSS signal is driven low when the master starts the communication and is kept low until the SPI is disabled.
  - NSS output disabled (SSM = 0, SSOE = 0)
     This configuration allows multimaster capability for devices operating in master mode. For devices set as slave, the NSS pin acts as a classical NSS input: the slave is selected when NSS is low and deselected when NSS high.

# Clock phase and clock polarity

Four possible timing relationships may be chosen by software, using the CPOL and CPHA bits in the SPI\_CR1 register. The CPOL (clock polarity) bit controls the steady state value of the clock when no data is being transferred. This bit affects both master and slave modes. If CPOL is reset, the SCK pin has a low-level idle state. If CPOL is set, the SCK pin has a high-level idle state.

If the CPHA (clock phase) bit is set, the second edge on the SCK pin (falling edge if the CPOL bit is reset, rising edge if the CPOL bit is set) is the MSBit capture strobe. Data are latched on the occurrence of the second clock transition. If the CPHA bit is reset, the first edge on the SCK pin (falling edge if CPOL bit is set, rising edge if CPOL bit is reset) is the MSBit capture strobe. Data are latched on the occurrence of the first clock transition.

The combination of the CPOL (clock polarity) and CPHA (clock phase) bits selects the data capture clock edge.

*Figure 251*, shows an SPI transfer with the four combinations of the CPHA and CPOL bits. The diagram may be interpreted as a master or slave timing diagram where the SCK pin, the MISO pin, the MOSI pin are directly connected between the master and the slave device.

Note: Prior to changing the CPOL/CPHA bits the SPI must be disabled by resetting the SPE bit.

Master and slave must be programmed with the same timing mode.

The idle state of SCK must correspond to the polarity selected in the SPI\_CR1 register (by pulling up SCK if CPOL=1 or pulling down SCK if CPOL=0).

The Data Frame Format (8- or 16-bit) is selected through the DFF bit in SPI\_CR1 register, and determines the data length during transmission/reception.

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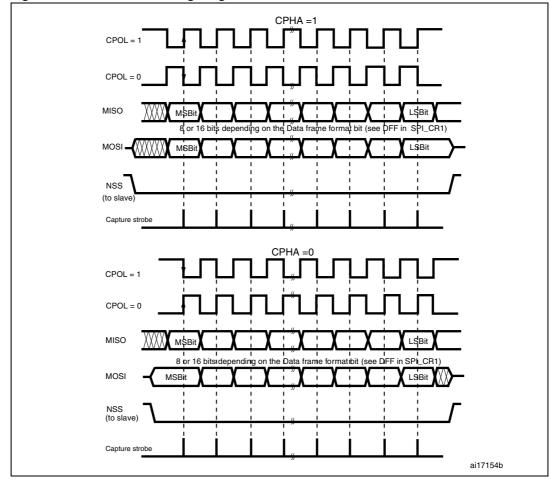


Figure 251. Data clock timing diagram

1. These timings are shown with the LSBFIRST bit reset in the SPI\_CR1 register.

## **Data frame format**

Data can be shifted out either MSB-first or LSB-first depending on the value of the LSBFIRST bit in the SPI\_CR1 Register.

Each data frame is 8 or 16 bits long depending on the size of the data programmed using the DFF bit in the SPI\_CR1 register. The selected data frame format is applicable for transmission and/or reception.

# 25.3.2 Configuring the SPI in slave mode

In the slave configuration, the serial clock is received on the SCK pin from the master device. The value set in the BR[2:0] bits in the SPI\_CR1 register, does not affect the data transfer rate.

Note:

It is recommended to enable the SPI slave before the master sends the clock. If not, undesired data transmission might occur. The data register of the slave needs to be ready before the first edge of the communication clock or before the end of the ongoing communication. It is mandatory to have the polarity of the communication clock set to the steady state value before the slave and the master are enabled.

Follow the procedure below to configure the SPI in slave mode:

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### **Procedure**

- Set the DFF bit to define 8- or 16-bit data frame format
- Select the CPOL and CPHA bits to define one of the four relationships between the
  data transfer and the serial clock (see *Figure 251*). For correct data transfer, the CPOL
  and CPHA bits must be configured in the same way in the slave device and the master
  device. This step is not required when the TI mode is selected through the FRF bit in
  the SPI\_CR2 register.
- 3. The frame format (MSB-first or LSB-first depending on the value of the LSBFIRST bit in the SPI\_CR1 register) must be the same as the master device. This step is not required when TI mode is selected.
- 4. In Hardware mode (refer to *Slave select (NSS) pin management on page 679*), the NSS pin must be connected to a low level signal during the complete byte transmit sequence. In NSS software mode, set the SSM bit and clear the SSI bit in the SPI\_CR1 register. This step is not required when TI mode is selected.
- 5. Set the FRF bit in the SPI\_CR2 register to select the TI mode protocol for serial communications.
- 6. Clear the MSTR bit and set the SPE bit (both in the SPI\_CR1 register) to assign the pins to alternate functions.

In this configuration the MOSI pin is a data input and the MISO pin is a data output.

### **Transmit sequence**

The data byte is parallel-loaded into the Tx buffer during a write cycle.

The transmit sequence begins when the slave device receives the clock signal and the most significant bit of the data on its MOSI pin. The remaining bits (the 7 bits in 8-bit data frame format, and the 15 bits in 16-bit data frame format) are loaded into the shift-register. The TXE flag in the SPI\_SR register is set on the transfer of data from the Tx Buffer to the shift register and an interrupt is generated if the TXEIE bit in the SPI\_CR2 register is set.

### Receive sequence

For the receiver, when data transfer is complete:

- The Data in shift register is transferred to Rx Buffer and the RXNE flag (SPI\_SR register) is set
- An Interrupt is generated if the RXNEIE bit is set in the SPI\_CR2 register.

After the last sampling clock edge the RXNE bit is set, a copy of the data byte received in the shift register is moved to the Rx buffer. When the SPI\_DR register is read, the SPI peripheral returns this buffered value.

Clearing of the RXNE bit is performed by reading the SPI\_DR register.

### SPI TI protocol in slave mode

In slave mode, the SPI interface is compatible with the TI protocol. The FRF bit of the SPI\_CR2 register can be used to configure the slave SPI serial communications to be compliant with this protocol.

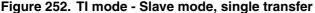
The clock polarity and phase are forced to conform to the TI protocol requirements whatever the values set in the SPI\_CR1 register. NSS management is also specific to the TI protocol which makes the configuration of NSS management through the SPI\_CR1 and SPI\_CR2 registers (such as SSM, SSI, SSOE) transparent for the user.

In Slave mode (*Figure 252: TI mode - Slave mode, single transfer* and *Figure 253: TI mode - Slave mode, continuous transfer*), the SPI baud rate prescaler is used to control the moment when the MISO pin state changes to HiZ. Any baud rate can be used thus allowing to determine this moment with optimal flexibility. However, the baud rate is generally set to the external master clock baud rate. The time for the MISO signal to become HiZ (t<sub>release</sub>) depends on internal resynchronizations and on the baud rate value set in through BR[2:0] of SPI\_CR1 register. It is given by the formula:

$$\frac{t_{baud\_rate}}{2} + 4 \times t_{pclk} < t_{release} < \frac{t_{baud\_rate}}{2} + 6 \times t_{pclk}$$

Note: This feature is not available for Motorola SPI communications (FRF bit set to 0).

To detect TI frame errors in Slave transmitter only mode by using the Error interrupt (ERRIE = 1), the SPI must be configured in 2-line unidirectional mode by setting BIDIMODE and BIDIOE to 1 in the SPI\_CR1 register. When BIDIMODE is set to 0, OVR is set to 1 because the data register is never read and error interrupt are always generated, while when BIDIMODE is set to 1, data are not received and OVR is never set.



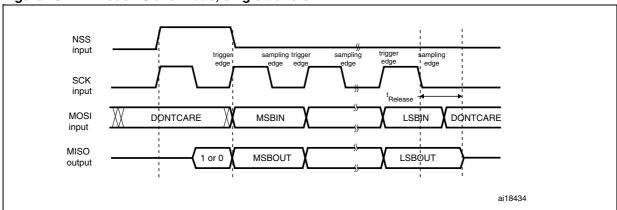
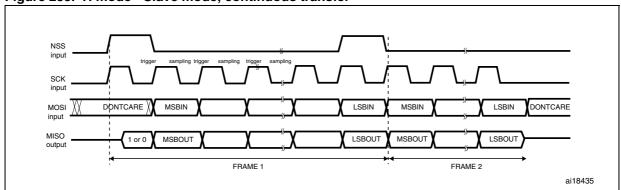


Figure 253. TI mode - Slave mode, continuous transfer



# 25.3.3 Configuring the SPI in master mode

In the master configuration, the serial clock is generated on the SCK pin.

### **Procedure**

- 1. Select the BR[2:0] bits to define the serial clock baud rate (see SPI\_CR1 register).
- 2. Select the CPOL and CPHA bits to define one of the four relationships between the data transfer and the serial clock (see *Figure 251*). This step is not required when the TI mode is selected.
- 3. Set the DFF bit to define 8- or 16-bit data frame format
- 4. Configure the LSBFIRST bit in the SPI\_CR1 register to define the frame format. This step is not required when the TI mode is selected.
- 5. If the NSS pin is required in input mode, in hardware mode, connect the NSS pin to a high-level signal during the complete byte transmit sequence. In NSS software mode, set the SSM and SSI bits in the SPI\_CR1 register. If the NSS pin is required in output mode, the SSOE bit only should be set. This step is not required when the TI mode is selected.
- 6. Set the FRF bit in SPI\_CR2 to select the TI protocol for serial communications.
- 7. The MSTR and SPE bits must be set (they remain set only if the NSS pin is connected to a high-level signal).

In this configuration the MOSI pin is a data output and the MISO pin is a data input.

### **Transmit sequence**

The transmit sequence begins when a byte is written in the Tx Buffer.

The data byte is parallel-loaded into the shift register (from the internal bus) during the first bit transmission and then shifted out serially to the MOSI pin MSB first or LSB first depending on the LSBFIRST bit in the SPI\_CR1 register. The TXE flag is set on the transfer of data from the Tx Buffer to the shift register and an interrupt is generated if the TXEIE bit in the SPI\_CR2 register is set.

### Receive sequence

For the receiver, when data transfer is complete:

- The data in the shift register is transferred to the RX Buffer and the RXNE flag is set
- An interrupt is generated if the RXNEIE bit is set in the SPI\_CR2 register

At the last sampling clock edge the RXNE bit is set, a copy of the data byte received in the shift register is moved to the Rx buffer. When the SPI\_DR register is read, the SPI peripheral returns this buffered value.

Clearing the RXNE bit is performed by reading the SPI\_DR register.

A continuous transmit stream can be maintained if the next data to be transmitted is put in the Tx buffer once the transmission is started. Note that TXE flag should be '1 before any attempt to write the Tx buffer is made.

Note:

When a master is communicating with SPI slaves which need to be de-selected between transmissions, the NSS pin must be configured as GPIO or another GPIO must be used and toggled by software.

### SPI TI protocol in master mode

In master mode, the SPI interface is compatible with the TI protocol. The FRF bit of the SPI\_CR2 register can be used to configure the master SPI serial communications to be compliant with this protocol.

The clock polarity and phase are forced to conform to the TI protocol requirements whatever the values set in the SPI\_CR1 register. NSS management is also specific to the TI protocol which makes the configuration of NSS management through the SPI\_CR1 and SPI\_CR2 registers (SSM, SSI, SSOE) transparent for the user.

Figure 254: TI mode - master mode, single transfer and Figure 255: TI mode - master mode, continuous transfer) show the SPI master communication waveforms when the TI mode is selected in master mode.



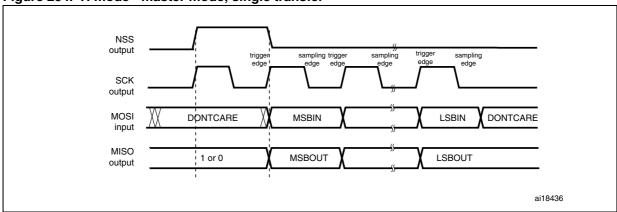
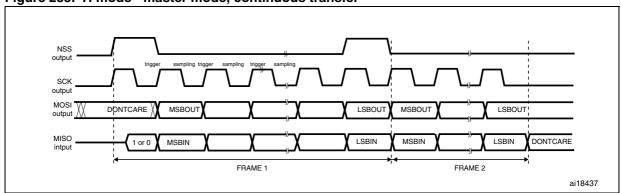


Figure 255. TI mode - master mode, continuous transfer



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### 25.3.4 Configuring the SPI for half-duplex communication

The SPI is capable of operating in half-duplex mode in 2 configurations.

- 1 clock and 1 bidirectional data wire
- 1 clock and 1 data wire (receive-only or transmit-only)

### 1 clock and 1 bidirectional data wire (BIDIMODE=1)

This mode is enabled by setting the BIDIMODE bit in the SPI\_CR1 register. In this mode SCK is used for the clock and MOSI in master or MISO in slave mode is used for data communication. The transfer direction (Input/Output) is selected by the BIDIOE bit in the SPI\_CR1 register. When this bit is 1, the data line is output otherwise it is input.

### 1 clock and 1 unidirectional data wire (BIDIMODE=0)

In this mode, the application can use the SPI either in transmit-only mode or in receive-only mode.

- Transmit-only mode is similar to full-duplex mode (BIDIMODE=0, RXONLY=0): the data are transmitted on the transmit pin (MOSI in master mode or MISO in slave mode) and the receive pin (MISO in master mode or MOSI in slave mode) can be used as a general-purpose IO. In this case, the application just needs to ignore the Rx buffer (if the data register is read, it does not contain the received value).
- In receive-only mode, the application can disable the SPI output function by setting the RXONLY bit in the SPI\_CR2 register. In this case, it frees the transmit IO pin (MOSI in master mode or MISO in slave mode), so it can be used for other purposes.

To start the communication in receive-only mode, configure and enable the SPI:

- In master mode, the communication starts immediately and stops when the SPE bit is cleared and the current reception stops. There is no need to read the BSY flag in this mode. It is always set when an SPI communication is ongoing.
- In slave mode, the SPI continues to receive as long as the NSS is pulled down (or the SSI bit is cleared in NSS software mode) and the SCK is running.

### 25.3.5 Data transmission and reception procedures

### **Rx and Tx buffers**

In reception, data are received and then stored into an internal Rx buffer while In transmission, data are first stored into an internal Tx buffer before being transmitted.

A read access of the SPI\_DR register returns the Rx buffered value whereas a write access to the SPI\_DR stores the written data into the Tx buffer.

### Start sequence in master mode

- In full-duplex (BIDIMODE=0 and RXONLY=0)
  - The sequence begins when data are written into the SPI\_DR register (Tx buffer).
  - The data are then parallel loaded from the Tx buffer into the 8-bit shift register during the first bit transmission and then shifted out serially to the MOSI pin.
  - At the same time, the received data on the MISO pin is shifted in serially to the 8bit shift register and then parallel loaded into the SPI\_DR register (Rx buffer).
- In unidirectional receive-only mode (BIDIMODE=0 and RXONLY=1)
  - The sequence begins as soon as SPE=1
  - Only the receiver is activated and the received data on the MISO pin are shifted in serially to the 8-bit shift register and then parallel loaded into the SPI\_DR register (Rx buffer).
- In bidirectional mode, when transmitting (BIDIMODE=1 and BIDIOE=1)
  - The sequence begins when data are written into the SPI\_DR register (Tx buffer).
  - The data are then parallel loaded from the Tx buffer into the 8-bit shift register during the first bit transmission and then shifted out serially to the MOSI pin.
  - No data are received.
- In bidirectional mode, when receiving (BIDIMODE=1 and BIDIOE=0)
  - The sequence begins as soon as SPE=1 and BIDIOE=0.
  - The received data on the MOSI pin are shifted in serially to the 8-bit shift register and then parallel loaded into the SPI\_DR register (Rx buffer).
  - The transmitter is not activated and no data are shifted out serially to the MOSI pin.

### Start sequence in slave mode

- In full-duplex mode (BIDIMODE=0 and RXONLY=0)
  - The sequence begins when the slave device receives the clock signal and the first bit of the data on its MOSI pin. The 7 remaining bits are loaded into the shift register.
  - At the same time, the data are parallel loaded from the Tx buffer into the 8-bit shift register during the first bit transmission, and then shifted out serially to the MISO pin. The software must have written the data to be sent before the SPI master device initiates the transfer.
- In unidirectional receive-only mode (BIDIMODE=0 and RXONLY=1)
  - The sequence begins when the slave device receives the clock signal and the first bit of the data on its MOSI pin. The 7 remaining bits are loaded into the shift register.
  - The transmitter is not activated and no data are shifted out serially to the MISO pin.
- In bidirectional mode, when transmitting (BIDIMODE=1 and BIDIOE=1)
  - The sequence begins when the slave device receives the clock signal and the first bit in the Tx buffer is transmitted on the MISO pin.
  - The data are then parallel loaded from the Tx buffer into the 8-bit shift register during the first bit transmission and then shifted out serially to the MISO pin. The software must have written the data to be sent before the SPI master device initiates the transfer.
  - No data are received.

- In bidirectional mode, when receiving (BIDIMODE=1 and BIDIOE=0)
  - The sequence begins when the slave device receives the clock signal and the first bit of the data on its MISO pin.
  - The received data on the MISO pin are shifted in serially to the 8-bit shift register and then parallel loaded into the SPI\_DR register (Rx buffer).
  - The transmitter is not activated and no data are shifted out serially to the MISO pin.

### Handling data transmission and reception

The TXE flag (Tx buffer empty) is set when the data are transferred from the Tx buffer to the shift register. It indicates that the internal Tx buffer is ready to be loaded with the next data. An interrupt can be generated if the TXEIE bit in the SPI\_CR2 register is set. Clearing the TXE bit is performed by writing to the SPI\_DR register.

Note:

The software must ensure that the TXE flag is set to 1 before attempting to write to the Tx buffer. Otherwise, it overwrites the data previously written to the Tx buffer.

The RXNE flag (Rx buffer not empty) is set on the last sampling clock edge, when the data are transferred from the shift register to the Rx buffer. It indicates that data are ready to be read from the SPI\_DR register. An interrupt can be generated if the RXNEIE bit in the SPI\_CR2 register is set. Clearing the RXNE bit is performed by reading the SPI\_DR register.

For some configurations, the BSY flag can be used during the last data transfer to wait until the completion of the transfer.

# Full-duplex transmit and receive procedure in master or slave mode (BIDIMODE=0 and RXONLY=0)

The software has to follow this procedure to transmit and receive data (see *Figure 256* and *Figure 257*):

- 1. Enable the SPI by setting the SPE bit to 1.
- 2. Write the first data item to be transmitted into the SPI\_DR register (this clears the TXE flag).
- 3. Wait until TXE=1 and write the second data item to be transmitted. Then wait until RXNE=1 and read the SPI\_DR to get the first received data item (this clears the RXNE bit). Repeat this operation for each data item to be transmitted/received until the n-1 received data.
- Wait until RXNE=1 and read the last received data.
- 5. Wait until TXE=1 and then wait until BSY=0 before disabling the SPI.

This procedure can also be implemented using dedicated interrupt subroutines launched at each rising edges of the RXNE or TXE flag.

Figure 256. TXE/RXNE/BSY behavior in Master / full-duplex mode (BIDIMODE=0 and RXONLY=0) in the case of continuous transfers

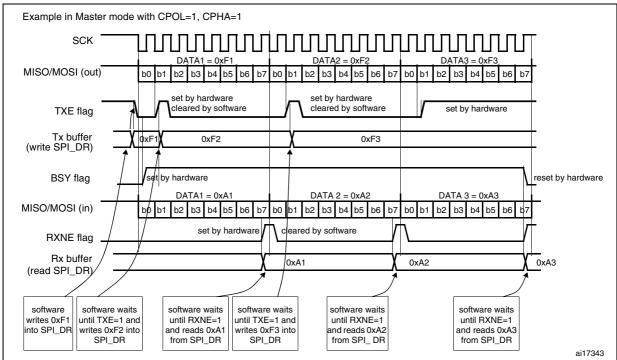
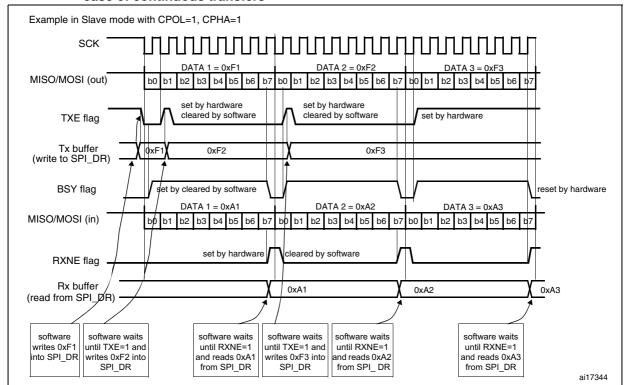


Figure 257. TXE/RXNE/BSY behavior in Slave / full-duplex mode (BIDIMODE=0, RXONLY=0) in the case of continuous transfers



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### Transmit-only procedure (BIDIMODE=0 RXONLY=0)

In this mode, the procedure can be reduced as described below and the BSY bit can be used to wait until the completion of the transmission (see *Figure 258* and *Figure 259*).

- 1. Enable the SPI by setting the SPE bit to 1.
- 2. Write the first data item to send into the SPI\_DR register (this clears the TXE bit).
- 3. Wait until TXE=1 and write the next data item to be transmitted. Repeat this step for each data item to be transmitted.
- 4. After writing the last data item into the SPI\_DR register, wait until TXE=1, then wait until BSY=0, this indicates that the transmission of the last data is complete.

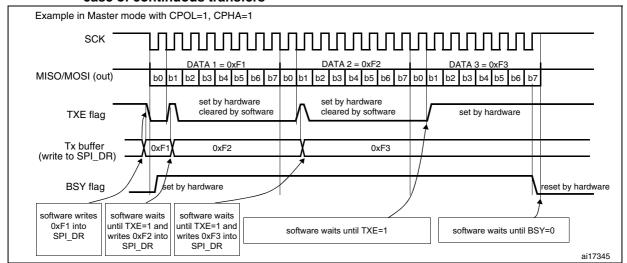
This procedure can be also implemented using dedicated interrupt subroutines launched at each rising edge of the TXE flag.

Note:

During discontinuous communications, there is a 2 APB clock period delay between the write operation to SPI\_DR and the BSY bit setting. As a consequence, in transmit-only mode, it is mandatory to wait first until TXE is set and then until BSY is cleared after writing the last data.

After transmitting two data items in transmit-only mode, the OVR flag is set in the SPI\_SR register since the received data are never read.

Figure 258. TXE/BSY behavior in Master transmit-only mode (BIDIMODE=0 and RXONLY=0) in the case of continuous transfers



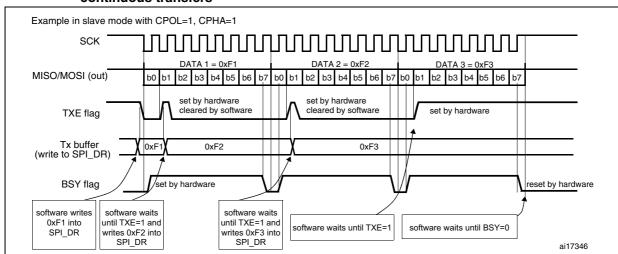


Figure 259. TXE/BSY in Slave transmit-only mode (BIDIMODE=0 and RXONLY=0) in the case of continuous transfers

### Bidirectional transmit procedure (BIDIMODE=1 and BIDIOE=1)

In this mode, the procedure is similar to the procedure in Transmit-only mode except that the BIDIMODE and BIDIOE bits both have to be set in the SPI\_CR2 register before enabling the SPI.

### Unidirectional receive-only procedure (BIDIMODE=0 and RXONLY=1)

In this mode, the procedure can be reduced as described below (see *Figure 260*):

- Set the RXONLY bit in the SPI CR2 register.
- 2. Enable the SPI by setting the SPE bit to 1:
  - a) In master mode, this immediately activates the generation of the SCK clock, and data are serially received until the SPI is disabled (SPE=0).
  - b) In slave mode, data are received when the SPI master device drives NSS low and generates the SCK clock.
- 3. Wait until RXNE=1 and read the SPI\_DR register to get the received data (this clears the RXNE bit). Repeat this operation for each data item to be received.

This procedure can also be implemented using dedicated interrupt subroutines launched at each rising edge of the RXNE flag.

Note:

If it is required to disable the SPI after the last transfer, follow the recommendation described in Section 25.3.8: Disabling the SPI on page 696.

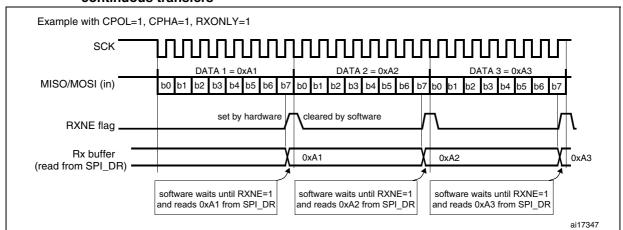


Figure 260. RXNE behavior in receive-only mode (BIDIRMODE=0 and RXONLY=1) in the case of continuous transfers

### Bidirectional receive procedure (BIDIMODE=1 and BIDIOE=0)

In this mode, the procedure is similar to the Receive-only mode procedure except that the BIDIMODE bit has to be set and the BIDIOE bit cleared in the SPI\_CR2 register before enabling the SPI.

### Continuous and discontinuous transfers

When transmitting data in master mode, if the software is fast enough to detect each rising edge of TXE (or TXE interrupt) and to immediately write to the SPI\_DR register before the ongoing data transfer is complete, the communication is said to be continuous. In this case, there is no discontinuity in the generation of the SPI clock between each data item and the BSY bit is never cleared between each data transfer.

On the contrary, if the software is not fast enough, this can lead to some discontinuities in the communication. In this case, the BSY bit is cleared between each data transmission (see *Figure 261*).

In Master receive-only mode (RXONLY=1), the communication is always continuous and the BSY flag is always read at 1.

In slave mode, the continuity of the communication is decided by the SPI master device. In any case, even if the communication is continuous, the BSY flag goes low between each transfer for a minimum duration of one SPI clock cycle (see *Figure 259*).

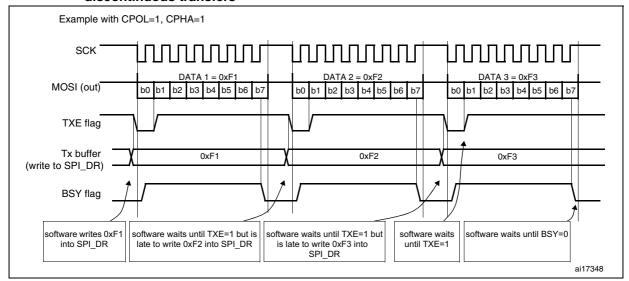


Figure 261. TXE/BSY behavior when transmitting (BIDIRMODE=0 and RXONLY=0) in the case of discontinuous transfers

### 25.3.6 CRC calculation

A CRC calculator has been implemented for communication reliability. Separate CRC calculators are implemented for transmitted data and received data. The CRC is calculated using a programmable polynomial serially on each bit. It is calculated on the sampling clock edge defined by the CPHA and CPOL bits in the SPI\_CR1 register.

Note:

This SPI offers two kinds of CRC calculation standard which depend directly on the data frame format selected for the transmission and/or reception: 8-bit data (CR8) and 16-bit data (CRC16).

CRC calculation is enabled by setting the CRCEN bit in the SPI\_CR1 register. This action resets the CRC registers (SPI\_RXCRCR and SPI\_TXCRCR). In full duplex or transmitter only mode, when the transfers are managed by the software (CPU mode), it is necessary to write the bit CRCNEXT immediately after the last data to be transferred is written to the SPI\_DR. At the end of this last data transfer, the SPI\_TXCRCR value is transmitted.

In receive only mode and when the transfers are managed by software (CPU mode), it is necessary to write the CRCNEXT bit after the second last data has been received. The CRC is received just after the last data reception and the CRC check is then performed.

At the end of data and CRC transfers, the CRCERR flag in the SPI\_SR register is set if corruption occurs during the transfer.

If data are present in the TX buffer, the CRC value is transmitted only after the transmission of the data byte. During CRC transmission, the CRC calculator is switched off and the register value remains unchanged.

SPI communication using the CRC is possible through the following procedure:

- 1. Program the CPOL, CPHA, LSBFirst, BR, SSM, SSI and MSTR values.
- 2. Program the polynomial in the SPI\_CRCPR register.
- 3. Enable the CRC calculation by setting the CRCEN bit in the SPI\_CR1 register. This also clears the SPI\_RXCRCR and SPI\_TXCRCR registers.
- 4. Enable the SPI by setting the SPE bit in the SPI\_CR1 register.
- 5. Start the communication and sustain the communication until all but one byte or half-word have been transmitted or received.
  - In full duplex or transmitter-only mode, when the transfers are managed by software, when writing the last byte or half word to the Tx buffer, set the CRCNEXT bit in the SPI\_CR1 register to indicate that the CRC will be transmitted after the transmission of the last byte.
  - In receiver only mode, set the bit CRCNEXT just after the reception of the second to last data to prepare the SPI to enter in CRC Phase at the end of the reception of the last data. CRC calculation is frozen during the CRC transfer.
- 6. After the transfer of the last byte or half word, the SPI enters the CRC transfer and check phase. In full duplex mode or receiver-only mode, the received CRC is compared to the SPI\_RXCRCR value. If the value does not match, the CRCERR flag in SPI\_SR is set and an interrupt can be generated when the ERRIE bit in the SPI\_CR2 register is set.

Note:

When the SPI is in slave mode, be careful to enable CRC calculation only when the clock is stable, that is, when the clock is in the steady state. If not, a wrong CRC calculation may be done. In fact, the CRC is sensitive to the SCK slave input clock as soon as CRCEN is set, and this, whatever the value of the SPE bit.

With high bitrate frequencies, be careful when transmitting the CRC. As the number of used CPU cycles has to be as low as possible in the CRC transfer phase, it is forbidden to call software functions in the CRC transmission sequence to avoid errors in the last data and CRC reception. In fact, CRCNEXT bit has to be written before the end of the transmission/reception of the last data.

For high bit rate frequencies, it is advised to use the DMA mode to avoid the degradation of the SPI speed performance due to CPU accesses impacting the SPI bandwidth.

When the devices are configured as slaves and the NSS hardware mode is used, the NSS pin needs to be kept low between the data phase and the CRC phase.

When the SPI is configured in slave mode with the CRC feature enabled, CRC calculation takes place even if a high level is applied on the NSS pin. This may happen for example in case of a multislave environment where the communication master addresses slaves alternately.

Between a slave deselection (high level on NSS) and a new slave selection (low level on NSS), the CRC value should be cleared on both master and slave sides in order to resynchronize the master and slave for their respective CRC calculation.

To clear the CRC, follow the procedure below:

- Disable SPI (SPE = 0)
- 2. Clear the CRCEN bit
- 3. Set the CRCEN bit
- 4. Enable the SPI (SPE = 1)

## 25.3.7 Status flags

Three status flags are provided for the application to completely monitor the state of the SPI bus.

### Tx buffer empty flag (TXE)

When it is set, this flag indicates that the Tx buffer is empty and the next data to be transmitted can be loaded into the buffer. The TXE flag is cleared when writing to the SPI DR register.

### Rx buffer not empty (RXNE)

When set, this flag indicates that there are valid received data in the Rx buffer. It is cleared when SPI DR is read.

### **BUSY flag**

This BSY flag is set and cleared by hardware (writing to this flag has no effect). The BSY flag indicates the state of the communication layer of the SPI.

When BSY is set, it indicates that the SPI is busy communicating. There is one exception in master mode / bidirectional receive mode (MSTR=1 and BDM=1 and BDOE=0) where the BSY flag is kept low during reception.

The BSY flag is useful to detect the end of a transfer if the software wants to disable the SPI and enter Halt mode (or disable the peripheral clock). This avoids corrupting the last transfer. For this, the procedure described below must be strictly respected.

The BSY flag is also useful to avoid write collisions in a multimaster system.

The BSY flag is set when a transfer starts, with the exception of master mode / bidirectional receive mode (MSTR=1 and BDM=1 and BDOE=0).

It is cleared:

- when a transfer is finished (except in master mode if the communication is continuous)
- when the SPI is disabled
- when a master mode fault occurs (MODF=1)

When communication is not continuous, the BSY flag is low between each communication.

When communication is continuous:

- in master mode, the BSY flag is kept high during all the transfers
- in slave mode, the BSY flag goes low for one SPI clock cycle between each transfer

Note:

Do not use the BSY flag to handle each data transmission or reception. It is better to use the TXE and RXNE flags instead.

### 25.3.8 Disabling the SPI

When a transfer is terminated, the application can stop the communication by disabling the SPI peripheral. This is done by clearing the SPE bit.

For some configurations, disabling the SPI and entering the Halt mode while a transfer is ongoing can cause the current transfer to be corrupted and/or the BSY flag might become unreliable.

To avoid any of those effects, it is recommended to respect the following procedure when disabling the SPI:

### In master or slave full-duplex mode (BIDIMODE=0, RXONLY=0)

- 1. Wait until RXNE=1 to receive the last data
- 2. Wait until TXE=1
- 3. Then wait until BSY=0
- 4. Disable the SPI (SPE=0) and, eventually, enter the Halt mode (or disable the peripheral clock)

# In master or slave unidirectional transmit-only mode (BIDIMODE=0, RXONLY=0) or bidirectional transmit mode (BIDIMODE=1, BIDIOE=1)

After the last data is written into the SPI\_DR register:

- 1. Wait until TXE=1
- 2. Then wait until BSY=0
- 3. Disable the SPI (SPE=0) and, eventually, enter the Halt mode (or disable the peripheral clock)

# In master unidirectional receive-only mode (MSTR=1, BIDIMODE=0, RXONLY=1) or bidirectional receive mode (MSTR=1, BIDIMODE=1, BIDIOE=0)

This case must be managed in a particular way to ensure that the SPI does not initiate a new transfer. The sequence below is valid only for SPI Motorola configuration (FRF bit set to 0):

- 1. Wait for the second to last occurrence of RXNE=1 (n-1)
- 2. Then wait for one SPI clock cycle (using a software loop) before disabling the SPI (SPE=0)
- 3. Then wait for the last RXNE=1 before entering the Halt mode (or disabling the peripheral clock)

When the SPI is configured in TI mode (Bit FRF set to 1), the following procedure has to be respected to avoid generating an undesired pulse on NSS when the SPI is disabled:

- 1. Wait for the second to last occurrence of RXNE = 1 (n-1).
- 2. Disable the SPI (SPE = 0) in the following window frame using a software loop:
  - After at least one SPI clock cycle,
  - Before the beginning of the LSB data transfer.

Note:

In master bidirectional receive mode (MSTR=1 and BDM=1 and BDOE=0), the BSY flag is kept low during transfers.

# In slave receive-only mode (MSTR=0, BIDIMODE=0, RXONLY=1) or bidirectional receive mode (MSTR=0, BIDIMODE=1, BIDOE=0)

- 1. You can disable the SPI (write SPE=1) at any time: the current transfer will complete before the SPI is effectively disabled
- 2. Then, if you want to enter the Halt mode, you must first wait until BSY = 0 before entering the Halt mode (or disabling the peripheral clock).

### 25.3.9 SPI communication using DMA (direct memory addressing)

To operate at its maximum speed, the SPI needs to be fed with the data for transmission and the data received on the Rx buffer should be read to avoid overrun. To facilitate the transfers, the SPI features a DMA capability implementing a simple request/acknowledge protocol.

A DMA access is requested when the enable bit in the SPI\_CR2 register is enabled. Separate requests must be issued to the Tx and Rx buffers (see *Figure 262* and *Figure 263*):

- In transmission, a DMA request is issued each time TXE is set to 1. The DMA then writes to the SPI\_DR register (this clears the TXE flag).
- In reception, a DMA request is issued each time RXNE is set to 1. The DMA then reads the SPI\_DR register (this clears the RXNE flag).

When the SPI is used only to transmit data, it is possible to enable only the SPI Tx DMA channel. In this case, the OVR flag is set because the data received are not read.

When the SPI is used only to receive data, it is possible to enable only the SPI Rx DMA channel.

In transmission mode, when the DMA has written all the data to be transmitted (flag TCIF is set in the DMA\_ISR register), the BSY flag can be monitored to ensure that the SPI communication is complete. This is required to avoid corrupting the last transmission before disabling the SPI or entering the Stop mode. The software must first wait until TXE=1 and then until BSY=0.

Note:

During discontinuous communications, there is a 2 APB clock period delay between the write operation to SPI\_DR and the BSY bit setting. As a consequence, it is mandatory to wait first until TXE=1 and then until BSY=0 after writing the last data.

Figure 262. Transmission using DMA

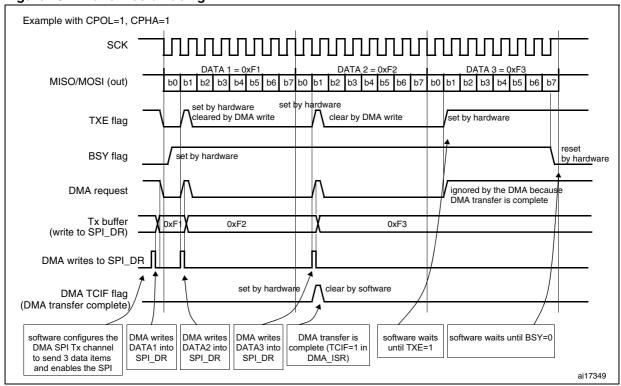
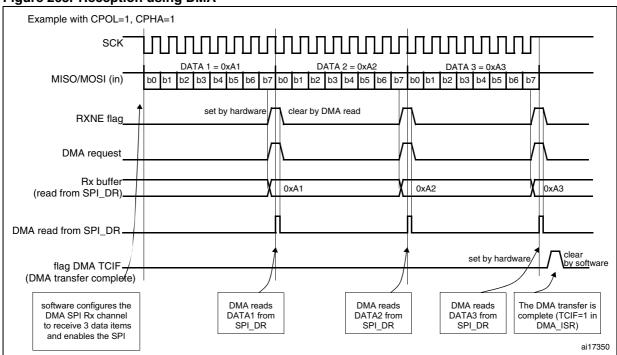


Figure 263. Reception using DMA



### DMA capability with CRC

When SPI communication is enabled with CRC communication and DMA mode, the transmission and reception of the CRC at the end of communication are automatic that is without using the bit CRCNEXT. After the CRC reception, the CRC must be read in the SPI DR register to clear the RXNE flag.

At the end of data and CRC transfers, the CRCERR flag in SPI\_SR is set if corruption occurs during the transfer.

### 25.3.10 Error flags

### Master mode fault (MODF)

Master mode fault occurs when the master device has its NSS pin pulled low (in NSS hardware mode) or SSI bit low (in NSS software mode), this automatically sets the MODF bit. Master mode fault affects the SPI peripheral in the following ways:

- The MODF bit is set and an SPI interrupt is generated if the ERRIE bit is set.
- The SPE bit is cleared. This blocks all output from the device and disables the SPI interface.
- The MSTR bit is cleared, thus forcing the device into slave mode.

Use the following software sequence to clear the MODF bit:

- 1. Make a read or write access to the SPI\_SR register while the MODF bit is set.
- Then write to the SPI\_CR1 register.

To avoid any multiple slave conflicts in a system comprising several MCUs, the NSS pin must be pulled high during the MODF bit clearing sequence. The SPE and MSTR bits can be restored to their original state after this clearing sequence.

As a security, hardware does not allow the setting of the SPE and MSTR bits while the MODF bit is set.

In a slave device the MODF bit cannot be set. However, in a multimaster configuration, the device can be in slave mode with this MODF bit set. In this case, the MODF bit indicates that there might have been a multimaster conflict for system control. An interrupt routine can be used to recover cleanly from this state by performing a reset or returning to a default state.

### Overrun condition

An overrun condition occurs when the master device has sent data bytes and the slave device has not cleared the RXNE bit resulting from the previous data byte transmitted. When an overrun condition occurs:

the OVR bit is set and an interrupt is generated if the ERRIE bit is set.

In this case, the receiver buffer contents will not be updated with the newly received data from the master device. A read from the SPI\_DR register returns this byte. All other subsequently transmitted bytes are lost.

Clearing the OVR bit is done by a read from the SPI\_DR register followed by a read access to the SPI\_SR register.

### **CRC** error

This flag is used to verify the validity of the value received when the CRCEN bit in the SPI CR1 register is set. The CRCERR flag in the SPI SR register is set if the value received in the shift register does not match the receiver SPI RXCRCR value.

### TI mode frame format error

A TI mode frame format error is detected when an NSS pulse occurs during an ongoing communication when the SPI is acting in slave mode and configured to conform to the TI mode protocol. When this error occurs, the FRE flag is set in the SPI\_SR register. The SPI is not disabled when an error occurs, the NSS pulse is ignored, and the SPI waits for the next NSS pulse before starting a new transfer. The data may be corrupted since the error detection may result in the lost of two data bytes.

The FRE flag is cleared when SPI SR register is read. If the bit ERRIE is set, an interrupt is generated on the NSS error detection. In this case, the SPI should be disabled because data consistency is no more guaranteed and communications should be reinitiated by the master when the slave SPI is re-enabled.

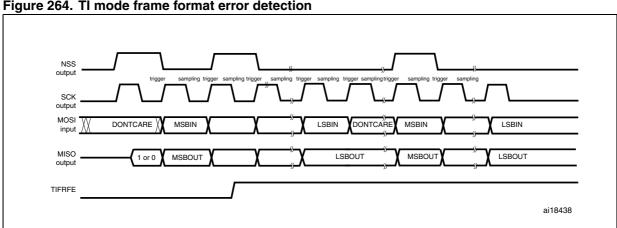


Figure 264. TI mode frame format error detection

#### 25.3.11 **SPI** interrupts

Table 103. SPI interrupt requests

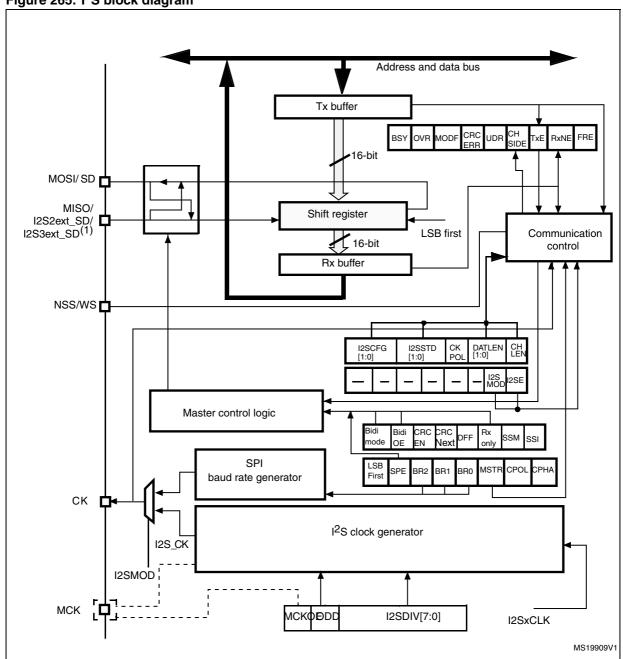
Interrupt event	Event flag	Enable Control bit		
Transmit buffer empty flag	TXE	TXEIE		
Receive buffer not empty flag	RXNE	RXNEIE		
Master Mode fault event	MODF			
Overrun error	OVR	ERRIE		
CRC error flag	CRCERR			
TI frame format error	FRE	ERRIE		

# 25.4 I<sup>2</sup>S functional description

# 25.4.1 I<sup>2</sup>S general description

The block diagram of the I<sup>2</sup>S is shown in *Figure 265*.

Figure 265. I<sup>2</sup>S block diagram



1.  $I2S2ext\_SD$  and  $I2S3ext\_SD$  are the extended SD pins that control the  $I^2S$  full duplex mode.

The SPI could function as an audio I<sup>2</sup>S interface when the I<sup>2</sup>S capability is enabled (by setting the I2SMOD bit in the SPI\_I2SCFGR register). This interface uses almost the same pins, flags and interrupts as the SPI.

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The I<sup>2</sup>S shares three common pins with the SPI:

- SD: Serial Data (mapped on the MOSI pin) to transmit or receive the two timemultiplexed data channels (in half-duplex mode only).
- WS: Word Select (mapped on the NSS pin) is the data control signal output in master mode and input in slave mode.
- CK: Serial Clock (mapped on the SCK pin) is the serial clock output in master mode and serial clock input in slave mode.
- I2S2ext\_SD and I2S3ext\_SD: additional pins (mapped on the MISO pin) to control the I<sup>2</sup>S full duplex mode.

An additional pin could be used when a master clock output is needed for some external audio devices:

 MCK: Master Clock (mapped separately) is used, when the I<sup>2</sup>S is configured in master mode (and when the MCKOE bit in the SPI\_I2SPR register is set), to output this additional clock generated at a preconfigured frequency rate equal to 256 × F<sub>S</sub>, where F<sub>S</sub> is the audio sampling frequency.

The I<sup>2</sup>S uses its own clock generator to produce the communication clock when it is set in master mode. This clock generator is also the source of the master clock output. Two additional registers are available in I<sup>2</sup>S mode. One is linked to the clock generator configuration SPI\_I2SPR and the other one is a generic I<sup>2</sup>S configuration register SPI\_I2SCFGR (audio standard, slave/master mode, data format, packet frame, clock polarity, etc.).

The SPI\_CR1 register and all CRC registers are not used in the I<sup>2</sup>S mode. Likewise, the SSOE bit in the SPI\_CR2 register and the MODF and CRCERR bits in the SPI\_SR are not used.

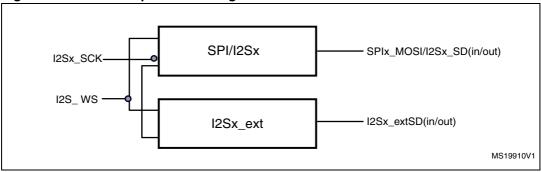
The I<sup>2</sup>S uses the same SPI register for data transfer (SPI\_DR) in 16-bit wide mode.

### **25.4.2 I2S** full duplex

To support I2S full duplex mode, two extra I<sup>2</sup>S instances called extended I2Ss (I2S2\_ext, I2S3\_ext) are available in addition to I2S2 and I2S3 (see *Figure 266*). The first I2S full-duplex interface is consequently based on I2S2 and I2S2\_ext, and the second one on I2S3 and I2S3\_ext.

Note: I2S2\_ext an I2S3\_ext are used only in full-duplex mode.

Figure 266. I2S full duplex block diagram



1. Where x can be 2 or 3.

I2Sx can operate in master mode. As a result:

- Only I2Sx can output SCK and WS in half duplex mode
- Only I2Sx can deliver SCK and WS to I2S2\_ext and I2S3\_ext in full duplex mode.

The extended I2Ss (I2Sx\_ext) can be used only in full duplex mode. The I2Sx\_ext operate always in slave mode.

Both I2Sx and I2Sx\_ext can be configured as transmitters or receivers.

### 25.4.3 Supported audio protocols

The four-line bus has to handle only audio data generally time-multiplexed on two channels: the right channel and the left channel. However there is only one 16-bit register for the transmission and the reception. So, it is up to the software to write into the data register the adequate value corresponding to the considered channel side, or to read the data from the data register and to identify the corresponding channel by checking the CHSIDE bit in the SPI\_SR register. Channel Left is always sent first followed by the channel right (CHSIDE has no meaning for the PCM protocol).

Four data and packet frames are available. Data may be sent with a format of:

- 16-bit data packed in 16-bit frame
- 16-bit data packed in 32-bit frame
- 24-bit data packed in 32-bit frame
- 32-bit data packed in 32-bit frame

When using 16-bit data extended on 32-bit packet, the first 16 bits (MSB) are the significant bits, the 16-bit LSB is forced to 0 without any need for software action or DMA request (only one read/write operation).

The 24-bit and 32-bit data frames need two CPU read or write operations to/from the SPI\_DR or two DMA operations if the DMA is preferred for the application. For 24-bit data frame specifically, the 8 nonsignificant bits are extended to 32 bits with 0-bits (by hardware). For all data formats and communication standards, the most significant bit is always sent first (MSB first).

The I<sup>2</sup>S interface supports four audio standards, configurable using the I2SSTD[1:0] and PCMSYNC bits in the SPI\_I2SCFGR register.

### I<sup>2</sup>S Philips standard

For this standard, the WS signal is used to indicate which channel is being transmitted. It is activated one CK clock cycle before the first bit (MSB) is available.

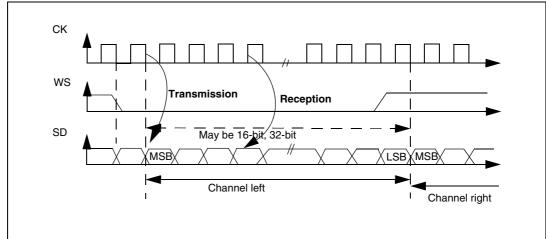


Figure 267. I<sup>2</sup>S Philips protocol waveforms (16/32-bit full accuracy, CPOL = 0)

Data are latched on the falling edge of CK (for the transmitter) and are read on the rising edge (for the receiver). The WS signal is also latched on the falling edge of CK.

CK

WS

Transmission

Reception

SD

Channel left 32-bit

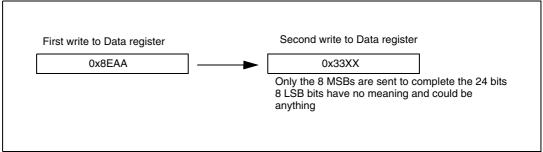
Channel right

Figure 268. I<sup>2</sup>S Philips standard waveforms (24-bit frame with CPOL = 0)

This mode needs two write or read operations to/from the SPI\_DR.

In transmission mode: if 0x8EAA33 has to be sent (24-bit):

Figure 269. Transmitting 0x8EAA33



In reception mode: if data 0x8EAA33 is received:

Figure 270. Receiving 0x8EAA33

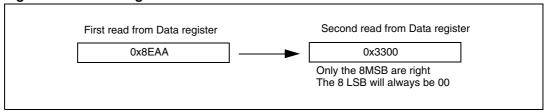
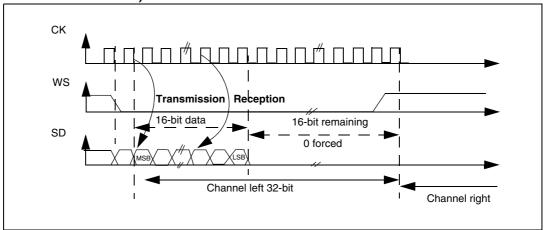


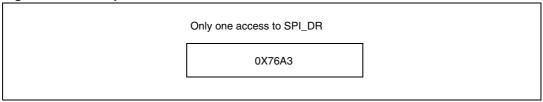
Figure 271.  $I^2$ S Philips standard (16-bit extended to 32-bit packet frame with CPOL = 0)



When 16-bit data frame extended to 32-bit channel frame is selected during the  $I^2S$  configuration phase, only one access to SPI\_DR is required. The 16 remaining bits are forced by hardware to 0x0000 to extend the data to 32-bit format.

If the data to transmit or the received data are 0x76A3 (0x76A30000 extended to 32-bit), the operation shown in *Figure 272* is required.

Figure 272. Example



For transmission, each time an MSB is written to SPI\_DR, the TXE flag is set and its interrupt, if allowed, is generated to load SPI\_DR with the new value to send. This takes place even if 0x0000 have not yet been sent because it is done by hardware.

For reception, the RXNE flag is set and its interrupt, if allowed, is generated when the first 16 MSB half-word is received.

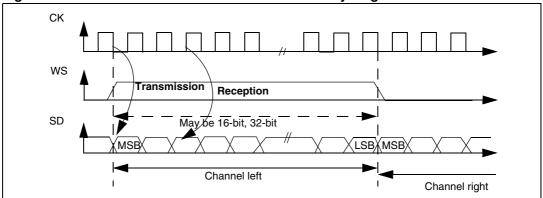
In this way, more time is provided between two write or read operations, which prevents underrun or overrun conditions (depending on the direction of the data transfer).

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### MSB justified standard

For this standard, the WS signal is generated at the same time as the first data bit, which is the MSBit.

Figure 273. MSB Justified 16-bit or 32-bit full-accuracy length with CPOL = 0



Data are latched on the falling edge of CK (for transmitter) and are read on the rising edge (for the receiver).

Figure 274. MSB Justified 24-bit frame length with CPOL = 0

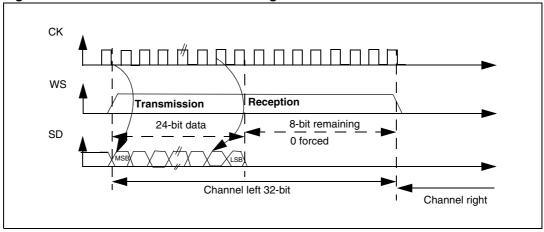
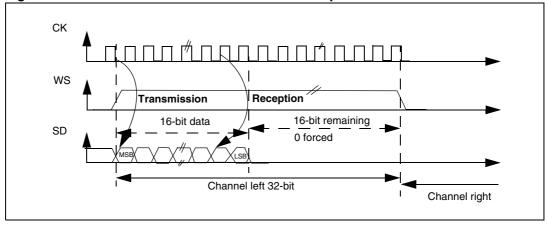


Figure 275. MSB Justified 16-bit extended to 32-bit packet frame with CPOL = 0



### LSB justified standard

This standard is similar to the MSB justified standard (no difference for the 16-bit and 32-bit full-accuracy frame formats).

Figure 276. LSB justified 16-bit or 32-bit full-accuracy with CPOL = 0

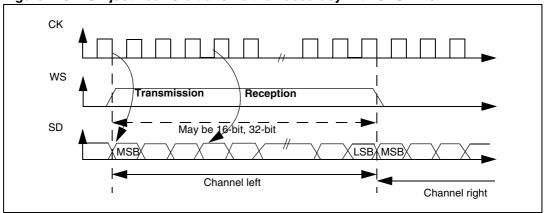
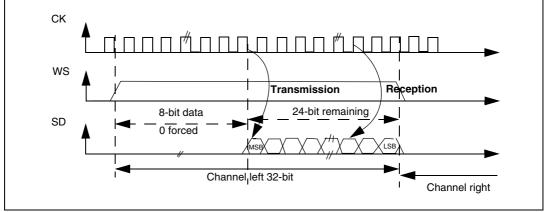


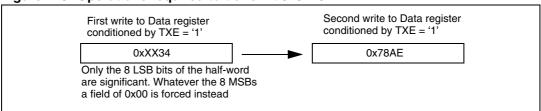
Figure 277. LSB Justified 24-bit frame length with CPOL = 0



In transmission mode:

If data 0x3478AE have to be transmitted, two write operations to the SPI\_DR register are required from software or by DMA. The operations are shown below.

Figure 278. Operations required to transmit 0x3478AE



In reception mode:

If data 0x3478AE are received, two successive read operations from SPI\_DR are required on each RXNE event.

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Figure 279. Operations required to receive 0x3478AE

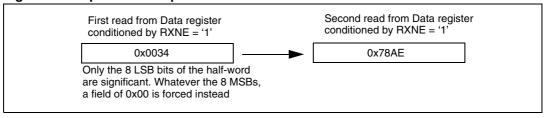
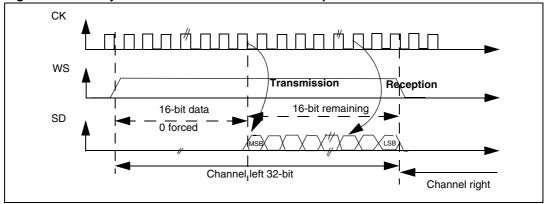


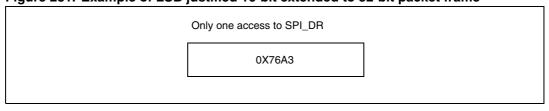
Figure 280. LSB justified 16-bit extended to 32-bit packet frame with CPOL = 0



When 16-bit data frame extended to 32-bit channel frame is selected during the I<sup>2</sup>S configuration phase, Only one access to SPI\_DR is required. The 16 remaining bits are forced by hardware to 0x0000 to extend the data to 32-bit format. In this case it corresponds to the half-word MSB.

If the data to transmit or the received data are 0x76A3 (0x0000 76A3 extended to 32-bit), the operation shown in *Figure 281* is required.

Figure 281. Example of LSB justified 16-bit extended to 32-bit packet frame



In transmission mode, when TXE is asserted, the application has to write the data to be transmitted (in this case 0x76A3). The 0x000 field is transmitted first (extension on 32-bit). TXE is asserted again as soon as the effective data (0x76A3) is sent on SD.

In reception mode, RXNE is asserted as soon as the significant half-word is received (and not the 0x0000 field).

In this way, more time is provided between two write or read operations to prevent underrun or overrun conditions.

### **PCM** standard

For the PCM standard, there is no need to use channel-side information. The two PCM modes (short and long frame) are available and configurable using the PCMSYNC bit in SPI\_I2SCFGR.

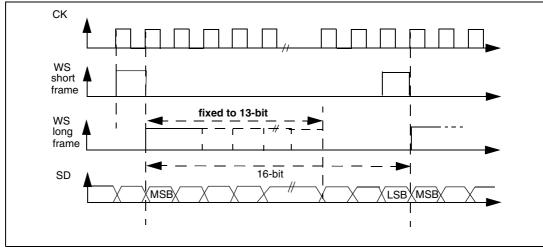


Figure 282. PCM standard waveforms (16-bit)

For long frame synchronization, the WS signal assertion time is fixed 13 bits in master mode.

For short frame synchronization, the WS synchronization signal is only one cycle long.

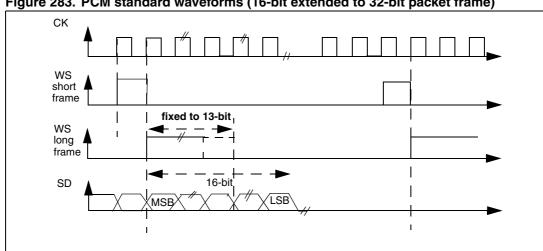


Figure 283. PCM standard waveforms (16-bit extended to 32-bit packet frame)

Note:

For both modes (master and slave) and for both synchronizations (short and long), the number of bits between two consecutive pieces of data (and so two synchronization signals) needs to be specified (DATLEN and CHLEN bits in the SPI\_I2SCFGR register) even in slave mode.

#### 25.4.4 **Clock generator**

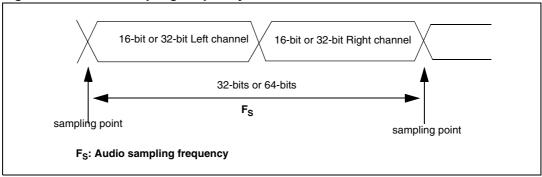
The I<sup>2</sup>S bitrate determines the dataflow on the I<sup>2</sup>S data line and the I<sup>2</sup>S clock signal frequency.

 $I^2S$  bitrate = number of bits per channel × number of channels × sampling audio frequency For a 16-bit audio, left and right channel, the I<sup>2</sup>S bitrate is calculated as follows:

$$I^2S$$
 bitrate =  $16 \times 2 \times F_S$ 

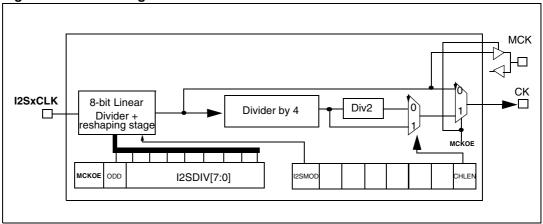
It will be:  $I^2S$  bitrate = 32 x 2 x  $F_S$  if the packet length is 32-bit wide.

Figure 284. Audio sampling frequency definition



When the master mode is configured, a specific action needs to be taken to properly program the linear divider in order to communicate with the desired audio frequency.

Figure 285. I<sup>2</sup>S clock generator architecture



1. Where x could be 2 or 3.

*Figure 284* presents the communication clock architecture. To achieve high-quality audio performance, the I2SxCLK clock source can be either the PLLI2S output (through R division factor) or an external clock (mapped to I2S\_CKIN pin).

The audio sampling frequency can be 192 kHz, 96 kHz, or 48 kHz. In order to reach the desired frequency, the linear divider needs to be programmed according to the formulas below:

When the master clock is generated (MCKOE in the SPI\_I2SPR register is set):

 $F_S = I2SxCLK / [(16*2)*((2*I2SDIV)+ODD)*8)]$  when the channel frame is 16-bit wide

 $F_S = I2SxCLK / [(32*2)*((2*I2SDIV)+ODD)*4)]$  when the channel frame is 32-bit wide

When the master clock is disabled (MCKOE bit cleared):

 $F_S = I2SxCLK / [(16*2)*((2*I2SDIV)+ODD))]$  when the channel frame is 16-bit wide

 $F_S = I2SxCLK / [(32*2)*((2*I2SDIV)+ODD))]$  when the channel frame is 32-bit wide

*Table 104* provides example precision values for different clock configurations.

Note: Other configurations are possible that allow optimum clock precision.

Table 104. Audio frequency precision (for PLLM VCO = 1 MHz or 2 MHz) $^{(1)}$ 

Master clock	Target f <sub>S</sub> (Hz)	Data format	PLLI2SN	PLLI2SR	I2SDIV	I2SODD	Real f <sub>S</sub> (Hz)	Error
	9000	16-bit	192	2	187	1	8000	0.0000%
	8000	32-bit	192	3	62	1	8000	0.0000%
	16000	16-bit	192	3	62	1	16000	0.0000%
	16000	32-bit	256	2	62	1	16000	0.0000%
	32000	16-bit	256	2	62	1	32000	0.0000%
	32000	32-bit	256	5	12	1	32000	0.0000%
	48000	16-bit	192	5	12	1	48000	0.0000%
Disabled	46000	32-bit	384	5	12	1	48000	0.0000%
Disabled	96000	16-bit	384	5	12	1	96000	0.0000%
		32-bit	424	3	11	1	96014.49219	0.0151%
	22050	16-bit	290	3	68	1	22049.87695	0.0006%
		32-bit	302	2	53	1	22050.23438	0.0011%
	44100	16-bit	302	2	53	1	44100.46875	0.0011%
	44100	32-bit	429	4	19	0	44099.50781	0.0011%
	100000	16-bit	424	3	11	1	192028.9844	0.0151%
	192000	32-bit	258	3	3	1	191964.2813	0.0186%
	8000	don't care	256	5	12	1	8000	0.0000%
	16000	don't care	213	2	13	0	16000.60059	0.0038%
	32000	don't care	213	2	6	1	32001.20117	0.0038%
Enabled	48000	don't care	258	3	3	1	47991.07031	0.0186%
	96000	don't care	344	2	3	1	95982.14063	0.0186%
	22050	don't care	429	4	9	1	22049.75391	0.0011%
	44100	don't care	271	2	6	0	44108.07422	0.0183%

This table gives only example values for different clock configurations. Other configurations allowing optimum clock precision are possible.

# 25.4.5 I<sup>2</sup>S master mode

The I<sup>2</sup>S can be configured as follows:

- In master mode for transmission or reception (half-duplex mode using I2Sx)
- In master mode transmission and reception (full duplex mode using I2Sx and I2Sx\_ext).

This means that the serial clock is generated on the CK pin as well as the Word Select signal WS. Master clock (MCK) may be output or not, thanks to the MCKOE bit in the SPI\_I2SPR register.

### **Procedure**

- Select the I2SDIV[7:0] bits in the SPI\_I2SPR register to define the serial clock baud rate to reach the proper audio sample frequency. The ODD bit in the SPI\_I2SPR register also has to be defined.
- Select the CKPOL bit to define the steady level for the communication clock. Set the MCKOE bit in the SPI\_I2SPR register if the master clock MCK needs to be provided to the external DAC/ADC audio component (the I2SDIV and ODD values should be computed depending on the state of the MCK output, for more details refer to Section 25.4.4: Clock generator).
- 3. Set the I2SMOD bit in SPI\_I2SCFGR to activate the I<sup>2</sup>S functionalities and choose the I<sup>2</sup>S standard through the I2SSTD[1:0] and PCMSYNC bits, the data length through the DATLEN[1:0] bits and the number of bits per channel by configuring the CHLEN bit. Select also the I<sup>2</sup>S master mode and direction (Transmitter or Receiver) through the I2SCFG[1:0] bits in the SPI\_I2SCFGR register.
- 4. If needed, select all the potential interruption sources and the DMA capabilities by writing the SPI\_CR2 register.
- 5. The I2SE bit in SPI\_I2SCFGR register must be set.

WS and CK are configured in output mode. MCK is also an output, if the MCKOE bit in SPI\_I2SPR is set.

### **Transmission sequence**

The transmission sequence begins when a half-word is written into the Tx buffer.

Assumedly, the first data written into the Tx buffer correspond to the channel Left data. When data are transferred from the Tx buffer to the shift register, TXE is set and data corresponding to the channel Right have to be written into the Tx buffer. The CHSIDE flag indicates which channel is to be transmitted. It has a meaning when the TXE flag is set because the CHSIDE flag is updated when TXE goes high.

A full frame has to be considered as a Left channel data transmission followed by a Right channel data transmission. It is not possible to have a partial frame where only the left channel is sent.

The data half-word is parallel loaded into the 16-bit shift register during the first bit transmission, and then shifted out, serially, to the MOSI/SD pin, MSB first. The TXE flag is set after each transfer from the Tx buffer to the shift register and an interrupt is generated if the TXEIE bit in the SPI\_CR2 register is set.

For more details about the write operations depending on the I<sup>2</sup>S standard mode selected, refer to *Section 25.4.3: Supported audio protocols*).

To ensure a continuous audio data transmission, it is mandatory to write the SPI\_DR with the next data to transmit before the end of the current transmission.

To switch off the  $I^2S$ , by clearing I2SE, it is mandatory to wait for TXE = 1 and BSY = 0.

### **Reception sequence**

The operating mode is the same as for the transmission mode except for the point 3 (refer to the procedure described in *Section 25.4.5: I2S master mode*), where the configuration should set the master reception mode through the I2SCFG[1:0] bits.

Whatever the data or channel length, the audio data are received by 16-bit packets. This means that each time the Rx buffer is full, the RXNE flag is set and an interrupt is generated

if the RXNEIE bit is set in SPI\_CR2 register. Depending on the data and channel length configuration, the audio value received for a right or left channel may result from one or two receptions into the Rx buffer.

Clearing the RXNE bit is performed by reading the SPI\_DR register.

CHSIDE is updated after each reception. It is sensitive to the WS signal generated by the  $I^2S$  cell.

For more details about the read operations depending on the I<sup>2</sup>S standard mode selected, refer to *Section 25.4.3: Supported audio protocols*.

If data are received while the previously received data have not been read yet, an overrun is generated and the OVR flag is set. If the ERRIE bit is set in the SPI\_CR2 register, an interrupt is generated to indicate the error.

To switch off the  $I^2S$ , specific actions are required to ensure that the  $I^2S$  completes the transfer cycle properly without initiating a new data transfer. The sequence depends on the configuration of the data and channel lengths, and on the audio protocol mode selected. In the case of:

- 16-bit data length extended on 32-bit channel length (DATLEN = 00 and CHLEN = 1) using the LSB justified mode (I2SSTD = 10)
  - a) Wait for the second to last RXNE = 1 (n 1)
  - b) Then wait 17 I<sup>2</sup>S clock cycles (using a software loop)
  - c) Disable the  $I^2S$  (I2SE = 0)
- 16-bit data length extended on 32-bit channel length (DATLEN = 00 and CHLEN = 1) in MSB justified, I<sup>2</sup>S or PCM modes (I2SSTD = 00, I2SSTD = 01 or I2SSTD = 11, respectively)
  - a) Wait for the last RXNE
  - b) Then wait 1 I<sup>2</sup>S clock cycle (using a software loop)
  - c) Disable the  $I^2S$  (I2SE = 0)
- For all other combinations of DATLEN and CHLEN, whatever the audio mode selected through the I2SSTD bits, carry out the following sequence to switch off the I<sup>2</sup>S:
  - a) Wait for the second to last RXNE = 1 (n 1)
  - b) Then wait one I<sup>2</sup>S clock cycle (using a software loop)
  - c) Disable the  $I^2S$  (I2SE = 0)

Note: The BSY flag is kept low during transfers.

## 25.4.6 I<sup>2</sup>S slave mode

The I<sup>2</sup>S can be configured as follows:

- In slave mode for transmission or reception (half-duplex mode using I2Sx)
- In slave mode transmission and reception (full duplex mode using I2Sx and I2Sx\_ext).

The operating mode is following mainly the same rules as described for the I<sup>2</sup>S master configuration. In slave mode, there is no clock to be generated by the I<sup>2</sup>S interface. The clock and WS signals are input from the external master connected to the I<sup>2</sup>S interface. There is then no need, for the user, to configure the clock.

The configuration steps to follow are listed below:

- Set the I2SMOD bit in the SPI\_I2SCFGR register to reach the I<sup>2</sup>S functionalities and choose the I<sup>2</sup>S standard through the I2SSTD[1:0] bits, the data length through the DATLEN[1:0] bits and the number of bits per channel for the frame configuring the CHLEN bit. Select also the mode (transmission or reception) for the slave through the I2SCFG[1:0] bits in SPI\_I2SCFGR register.
- 2. If needed, select all the potential interrupt sources and the DMA capabilities by writing the SPI CR2 register.
- 3. The I2SE bit in SPI I2SCFGR register must be set.

### **Transmission sequence**

The transmission sequence begins when the external master device sends the clock and when the NSS\_WS signal requests the transfer of data. The slave has to be enabled before the external master starts the communication. The I<sup>2</sup>S data register has to be loaded before the master initiates the communication.

For the I<sup>2</sup>S, MSB justified and LSB justified modes, the first data item to be written into the data register corresponds to the data for the left channel. When the communication starts, the data are transferred from the Tx buffer to the shift register. The TXE flag is then set in order to request the right channel data to be written into the I<sup>2</sup>S data register.

The CHSIDE flag indicates which channel is to be transmitted. Compared to the master transmission mode, in slave mode, CHSIDE is sensitive to the WS signal coming from the external master. This means that the slave needs to be ready to transmit the first data before the clock is generated by the master. WS assertion corresponds to left channel transmitted first.

Note:

The I2SE has to be written at least two PCLK cycles before the first clock of the master comes on the CK line.

The data half-word is parallel-loaded into the 16-bit shift register (from the internal bus) during the first bit transmission, and then shifted out serially to the MOSI/SD pin MSB first. The TXE flag is set after each transfer from the Tx buffer to the shift register and an interrupt is generated if the TXEIE bit in the SPI\_CR2 register is set.

Note that the TXE flag should be checked to be at 1 before attempting to write the Tx buffer.

For more details about the write operations depending on the I<sup>2</sup>S standard mode selected, refer to *Section 25.4.3: Supported audio protocols*.

To secure a continuous audio data transmission, it is mandatory to write the SPI\_DR register with the next data to transmit before the end of the current transmission. An underrun flag is set and an interrupt may be generated if the data are not written into the SPI\_DR register before the first clock edge of the next data communication. This indicates to the software that the transferred data are wrong. If the ERRIE bit is set into the SPI\_CR2 register, an interrupt is generated when the UDR flag in the SPI\_SR register goes high. In this case, it is mandatory to switch off the I<sup>2</sup>S and to restart a data transfer starting from the loft channel.

To switch off the  $I^2S$ , by clearing the I2SE bit, it is mandatory to wait for TXE = 1 and BSY = 0.

### **Reception sequence**

The operating mode is the same as for the transmission mode except for the point 1 (refer to the procedure described in *Section 25.4.6: I2S slave mode*), where the configuration should set the master reception mode using the I2SCFG[1:0] bits in the SPI\_I2SCFGR register.

Whatever the data length or the channel length, the audio data are received by 16-bit packets. This means that each time the RX buffer is full, the RXNE flag in the SPI\_SR register is set and an interrupt is generated if the RXNEIE bit is set in the SPI\_CR2 register. Depending on the data length and channel length configuration, the audio value received for a right or left channel may result from one or two receptions into the RX buffer.

The CHSIDE flag is updated each time data are received to be read from SPI\_DR. It is sensitive to the external WS line managed by the external master component.

Clearing the RXNE bit is performed by reading the SPI\_DR register.

For more details about the read operations depending the I<sup>2</sup>S standard mode selected, refer to *Section 25.4.3: Supported audio protocols*.

If data are received while the precedent received data have not yet been read, an overrun is generated and the OVR flag is set. If the bit ERRIE is set in the SPI\_CR2 register, an interrupt is generated to indicate the error.

To switch off the  $I^2S$  in reception mode, I2SE has to be cleared immediately after receiving the last RXNE = 1.

Note:

The external master components should have the capability of sending/receiving data in 16-bit or 32-bit packets via an audio channel.

### 25.4.7 Status flags

Three status flags are provided for the application to fully monitor the state of the I<sup>2</sup>S bus.

### **Busy flag (BSY)**

The BSY flag is set and cleared by hardware (writing to this flag has no effect). It indicates the state of the communication layer of the I<sup>2</sup>S.

When BSY is set, it indicates that the  $I^2S$  is busy communicating. There is one exception in master receive mode (I2SCFG = 11) where the BSY flag is kept low during reception.

The BSY flag is useful to detect the end of a transfer if the software needs to disable the I<sup>2</sup>S. This avoids corrupting the last transfer. For this, the procedure described below must be strictly respected.

The BSY flag is set when a transfer starts, except when the I<sup>2</sup>S is in master receiver mode.

The BSY flag is cleared:

- when a transfer completes (except in master transmit mode, in which the communication is supposed to be continuous)
- when the I<sup>2</sup>S is disabled

When communication is continuous:

- In master transmit mode, the BSY flag is kept high during all the transfers
- In slave mode, the BSY flag goes low for one I<sup>2</sup>S clock cycle between each transfer

Note:

Do not use the BSY flag to handle each data transmission or reception. It is better to use the TXE and RXNE flags instead.

### Tx buffer empty flag (TXE)

When set, this flag indicates that the Tx buffer is empty and the next data to be transmitted can then be loaded into it. The TXE flag is reset when the Tx buffer already contains data to be transmitted. It is also reset when the I<sup>2</sup>S is disabled (I2SE bit is reset).

### RX buffer not empty (RXNE)

When set, this flag indicates that there are valid received data in the RX Buffer. It is reset when SPI\_DR register is read.

### Channel Side flag (CHSIDE)

In transmission mode, this flag is refreshed when TXE goes high. It indicates the channel side to which the data to transfer on SD has to belong. In case of an underrun error event in slave transmission mode, this flag is not reliable and I<sup>2</sup>S needs to be switched off and switched on before resuming the communication.

In reception mode, this flag is refreshed when data are received into SPI\_DR. It indicates from which channel side data have been received. Note that in case of error (like OVR) this flag becomes meaningless and the I<sup>2</sup>S should be reset by disabling and then enabling it (with configuration if it needs changing).

This flag has no meaning in the PCM standard (for both Short and Long frame modes).

When the OVR or UDR flag in the SPI\_SR is set and the ERRIE bit in SPI\_CR2 is also set, an interrupt is generated. This interrupt can be cleared by reading the SPI\_SR status register (once the interrupt source has been cleared).

## 25.4.8 Error flags

There are three error flags for the I<sup>2</sup>S cell.

### **Underrun flag (UDR)**

In slave transmission mode this flag is set when the first clock for data transmission appears while the software has not yet loaded any value into SPI\_DR. It is available when the I2SMOD bit in SPI\_I2SCFGR is set. An interrupt may be generated if the ERRIE bit in SPI\_CR2 is set

The UDR bit is cleared by a read operation on the SPI\_SR register.

### Overrun flag (OVR)

This flag is set when data are received and the previous data have not yet been read from SPI\_DR. As a result, the incoming data are lost. An interrupt may be generated if the ERRIE bit is set in SPI\_CR2.

In this case, the receive buffer contents are not updated with the newly received data from the transmitter device. A read operation to the SPI\_DR register returns the previous correctly received data. All other subsequently transmitted half-words are lost.

Clearing the OVR bit is done by a read operation on the SPI\_DR register followed by a read access to the SPI\_SR register.

### Frame error flag (FRE)

This flag can be set by hardware only if the I2S is configured in Slave mode. It is set if the external master is changing the WS line at a moment when the slave is not expected this

change. If the synchronization is lost, to recover from this state and resynchronize the external master device with the I2S slave device, follow the steps below:

- 1. Disable the I2S
- 2. Re-enable it when the correct level is detected on the WS line (WS line is high in I2S mode, or low for MSB- or LSB-justified or PCM modes).

Desynchronization between the master and slave device may be due to noisy environment on the SCK communication clock or on the WS frame synchronization line. An error interrupt can be generated if the ERRIE bit is set. The desynchronization flag (FRE) is cleared by software when the status register is read.

# 25.4.9 I<sup>2</sup>S interrupts

Table 105 provides the list of I<sup>2</sup>S interrupts.

Table 105. I<sup>2</sup>S interrupt requests

Interrupt event	Event flag	Enable Control bit			
Transmit buffer empty flag	TXE	TXEIE			
Receive buffer not empty flag	RXNE	RXNEIE			
Overrun error	OVR	ERRIE			
Underrun error	UDR				
Frame error flag	FRE	ERRIE			

### 25.4.10 DMA features

DMA is working in exactly the same way as for the SPI mode. There is no difference on the  $I^2S$ . Only the CRC feature is not available in  $I^2S$  mode since there is no data transfer protection system.

# 25.5 SPI and I<sup>2</sup>S registers

Refer to Section 1.1 on page 46 for a list of abbreviations used in register descriptions.

The peripheral registers can be accessed by half-words (16 bits) or words (32 bits).

# 25.5.1 SPI control register 1 (SPI\_CR1) (not used in I<sup>2</sup>S mode)

Address offset: 0x00
Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
BIDI MODE	BIDI OE	CRC EN	CRC NEXT	DFF	RX ONLY	SSM	SSI	LSB FIRST	SPE	BR [2:0]		MSTR	CPOL	СРНА	
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bit 15 BIDIMODE: Bidirectional data mode enable

0: 2-line unidirectional data mode selected1: 1-line bidirectional data mode selected

Note: Not used in I<sup>2</sup>S mode

Bit 14 BIDIOE: Output enable in bidirectional mode

This bit combined with the BIDImode bit selects the direction of transfer in bidirectional mode

0: Output disabled (receive-only mode)

1: Output enabled (transmit-only mode)

Note: In master mode, the MOSI pin is used and in slave mode, the MISO pin is used.

### Not used in I<sup>2</sup>S mode

Bit 13 CRCEN: Hardware CRC calculation enable

0: CRC calculation disabled

1: CRC calculation enabled

Note: This bit should be written only when SPI is disabled (SPE = '0') for correct operation

Not used in PS mode

Bit 12 CRCNEXT: CRC transfer next

0: Data phase (no CRC phase)

1: Next transfer is CRC (CRC phase)

Note: When the SPI is configured in full duplex or transmitter only modes, CRCNEXT must be written as soon as the last data is written to the SPI\_DR register.

When the SPI is configured in receiver only mode, CRCNEXT must be set after the second last data reception.

This bit should be kept cleared when the transfers are managed by DMA.

Not used in  $l^2S$  mode

Bit 11 DFF: Data frame format

0: 8-bit data frame format is selected for transmission/reception

1: 16-bit data frame format is selected for transmission/reception

Note: This bit should be written only when SPI is disabled (SPE = '0') for correct operation

Not used in I<sup>2</sup>S mode

### Bit 10 RXONLY: Receive only

This bit combined with the BIDImode bit selects the direction of transfer in 2-line unidirectional mode. This bit is also useful in a multislave system in which this particular slave is not accessed, the output from the accessed slave is not corrupted.

0: Full duplex (Transmit and receive)

1: Output disabled (Receive-only mode)

Note: Not used in I<sup>2</sup>S mode

### Bit 9 SSM: Software slave management

When the SSM bit is set, the NSS pin input is replaced with the value from the SSI bit.

0: Software slave management disabled

1: Software slave management enabled

Note: Not used in I<sup>2</sup>S mode and SPI TI mode

### Bit 8 SSI: Internal slave select

This bit has an effect only when the SSM bit is set. The value of this bit is forced onto the NSS pin and the IO value of the NSS pin is ignored.

Note: Not used in I<sup>2</sup>S mode and SPI TI mode

### Bit 7 LSBFIRST: Frame format

0: MSB transmitted first

1: LSB transmitted first

Note: This bit should not be changed when communication is ongoing.

### Not used in PS mode and SPI TI mode

### Bit 6 SPE: SPI enable

0: Peripheral disabled

1: Peripheral enabled

Note: 1- Not used in PS mode.

Note: 2- When disabling the SPI, follow the procedure described in Section 25.3.8: Disabling the

SPI.

### Bits 5:3 BR[2:0]: Baud rate control

000: f<sub>PCLK</sub>/2100: f<sub>PCLK</sub>/32 001: f<sub>PCLK</sub>/4101: f<sub>PCLK</sub>/64 010: f<sub>PCLK</sub>/8110: f<sub>PCLK</sub>/128 011: f<sub>PCLK</sub>/16111: f<sub>PCLK</sub>/256

Note: These bits should not be changed when communication is ongoing.

### Not used in $l^2S$ mode

### Bit 2 MSTR: Master selection

0: Slave configuration

1: Master configuration

Note: This bit should not be changed when communication is ongoing.

### Not used in I<sup>2</sup>S mode

### Bit1 CPOL: Clock polarity

0: CK to 0 when idle

1: CK to 1 when idle

Note: This bit should not be changed when communication is ongoing.

# Not used in $l^2$ S mode and SPI TI mode

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Bit 0 CPHA: Clock phase

0: The first clock transition is the first data capture edge1: The second clock transition is the first data capture edge

Note: This bit should not be changed when communication is ongoing.

Note: Not used in I'S mode and SPI TI mode

## 25.5.2 SPI control register 2 (SPI\_CR2)

Address offset: 0x04 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Reserved								TXEIE	RXNEIE	ERRIE	FRF	Res.	SSOE	TXDMAEN	RXDMAEN
								rw	rw	rw	rw		rw	rw	rw

Bits 15:8 Reserved, must be kept at reset value.

Bit 7 TXEIE: Tx buffer empty interrupt enable

0: TXE interrupt masked

1: TXE interrupt not masked. Used to generate an interrupt request when the TXE flag is set.

Bit 6 RXNEIE: RX buffer not empty interrupt enable

0: RXNE interrupt masked

1: RXNE interrupt not masked. Used to generate an interrupt request when the RXNE flag is set

Bit 5 ERRIE: Error interrupt enable

This bit controls the generation of an interrupt when an error condition occurs (CRCERR, OVR, MODF in SPI mode and UDR, OVR and FRE).

0: Error interrupt is masked

1: Error interrupt is enabled

Bit 4 FRF: Frame format

0: SPI Motorola mode

1 SPI TI mode

Note: Not used in I<sup>2</sup>S mode

Bit 3 Reserved. Forced to 0 by hardware.

Bit 2 SSOE: SS output enable

0: SS output is disabled in master mode and the cell can work in multimaster configuration

1: SS output is enabled in master mode and when the cell is enabled. The cell cannot work in a multimaster environment.

Note: Not used in I<sup>2</sup>S mode and SPI TI mode

Bit 1 TXDMAEN: Tx buffer DMA enable

When this bit is set, the DMA request is made whenever the TXE flag is set.

0: Tx buffer DMA disabled

1: Tx buffer DMA enabled

Bit 0 RXDMAEN: Rx buffer DMA enable

When this bit is set, the DMA request is made whenever the RXNE flag is set.

0: Rx buffer DMA disabled

1: Rx buffer DMA enabled

# 25.5.3 SPI status register (SPI\_SR)

Address offset: 0x08 Reset value: 0x0002

15	14	13	12	11	10	9	8	1	6	5	4	3	2	1	0	
			Reserved	1			FRE	BSY	OVR	MODF	CRC ERR	UDR	CHSID E	TXE	RXNE	
							r	r	r	r	rc_w0	r	r	r	r	

Bits 15:9 Reserved. Forced to 0 by hardware.

Bit 8 FRE: Frame format error

0: No frame format error

1: A frame format error occurred

This flag is set by hardware and cleared by software when the SPIx\_SR register is read.

Note: This flag is used when the SPI operates in TI slave mode or I2S slave mode (refer to Section 25.3.10).

Bit 7 BSY: Busy flag

0: SPI (or I2S) not busy

1: SPI (or I2S) is busy in communication or Tx buffer is not empty

This flag is set and cleared by hardware.

Note: BSY flag must be used with caution: refer to Section 25.3.7: Status flags and Section 25.3.8: Disabling the SPI.

Bit 6 OVR: Overrun flag

0: No overrun occurred

1: Overrun occurred

This flag is set by hardware and reset by a software sequence. Refer to *Section 25.4.8 on page 716* for the software sequence.

Bit 5 MODF: Mode fault

0: No mode fault occurred

1: Mode fault occurred

This flag is set by hardware and reset by a software sequence. Refer to *Section 25.3.10 on page 699* for the software sequence.

Note: Not used in I<sup>2</sup>S mode

Bit 4 CRCERR: CRC error flag

0: CRC value received matches the SPI\_RXCRCR value

1: CRC value received does not match the SPI\_RXCRCR value

This flag is set by hardware and cleared by software writing 0.

Note: Not used in I<sup>2</sup>S mode

Bit 3 UDR: Underrun flag

0: No underrun occurred

1: Underrun occurred

This flag is set by hardware and reset by a software sequence. Refer to *Section 25.4.8 on page 716* for the software sequence.

Note: Not used in SPI mode

Bit 2 CHSIDE: Channel side

0: Channel Left has to be transmitted or has been received

1: Channel Right has to be transmitted or has been received

Note: Not used for SPI mode. No meaning in PCM mode

Bit 1 TXE: Transmit buffer empty

0: Tx buffer not empty

1: Tx buffer empty

Bit 0 RXNE: Receive buffer not empty

0: Rx buffer empty1: Rx buffer not empty

# 25.5.4 SPI data register (SPI\_DR)

Address offset: 0x0C Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							DR[	15:0]							
rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bits 15:0 DR[15:0]: Data register

Data received or to be transmitted.

The data register is split into 2 buffers - one for writing (Transmit Buffer) and another one for reading (Receive buffer). A write to the data register will write into the Tx buffer and a read from the data register will return the value held in the Rx buffer.

#### Notes for the SPI mode:

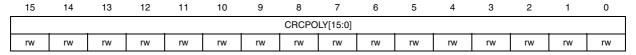
Depending on the data frame format selection bit (DFF in SPI\_CR1 register), the data sent or received is either 8-bit or 16-bit. This selection has to be made before enabling the SPI to ensure correct operation.

For an 8-bit data frame, the buffers are 8-bit and only the LSB of the register (SPI\_DR[7:0]) is used for transmission/reception. When in reception mode, the MSB of the register (SPI\_DR[15:8]) is forced to 0.

For a 16-bit data frame, the buffers are 16-bit and the entire register, SPI\_DR[15:0] is used for transmission/reception.

# 25.5.5 SPI CRC polynomial register (SPI\_CRCPR) (not used in I<sup>2</sup>S mode)

Address offset: 0x10 Reset value: 0x0007



Bits 15:0 CRCPOLY[15:0]: CRC polynomial register

This register contains the polynomial for the CRC calculation.

The CRC polynomial (0007h) is the reset value of this register. Another polynomial can be configured as required.

Note: Not used for the  $l^2S$  mode.

# 25.5.6 SPI RX CRC register (SPI\_RXCRCR) (not used in I<sup>2</sup>S mode)

Address offset: 0x14 Reset value: 0x0000

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							RXCR	C[15:0]							
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

#### Bits 15:0 RXCRC[15:0]: Rx CRC register

When CRC calculation is enabled, the RxCRC[15:0] bits contain the computed CRC value of the subsequently received bytes. This register is reset when the CRCEN bit in SPI\_CR1 register is written to 1. The CRC is calculated serially using the polynomial programmed in the SPI\_CRCPR register.

Only the 8 LSB bits are considered when the data frame format is set to be 8-bit data (DFF bit of SPI\_CR1 is cleared). CRC calculation is done based on any CRC8 standard.

The entire 16-bits of this register are considered when a 16-bit data frame format is selected (DFF bit of the SPI\_CR1 register is set). CRC calculation is done based on any CRC16 standard.

Note: A read to this register when the BSY Flag is set could return an incorrect value. Not used for  $l^2S$  mode.

# 25.5.7 SPI TX CRC register (SPI\_TXCRCR) (not used in I<sup>2</sup>S mode)

Address offset: 0x18
Reset value: 0x0000

	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ſ								TXCR	C[15:0]							
Ī	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

# Bits 15:0 TXCRC[15:0]: Tx CRC register

When CRC calculation is enabled, the TxCRC[7:0] bits contain the computed CRC value of the subsequently transmitted bytes. This register is reset when the CRCEN bit of SPI\_CR1 is written to 1. The CRC is calculated serially using the polynomial programmed in the SPI\_CRCPR register.

Only the 8 LSB bits are considered when the data frame format is set to be 8-bit data (DFF bit of SPI\_CR1 is cleared). CRC calculation is done based on any CRC8 standard.

The entire 16-bits of this register are considered when a 16-bit data frame format is selected (DFF bit of the SPI\_CR1 register is set). CRC calculation is done based on any CRC16 standard.

Note: A read to this register when the BSY flag is set could return an incorrect value. Not used for I<sup>2</sup>S mode.

# 25.5.8 SPI\_I<sup>2</sup>S configuration register (SPI\_I2SCFGR)

Address offset: 0x1C Reset value: 0x0000

15 14 12 11 7 5 3 0 PCMSY DATLEN I2SMOD 12SE 12SCFG I2SSTD CKPOL CHLEN Reserved Reserved rw rw rw rw rw rw rw rw

Bits 15:12 Reserved, must be kept at reset value.

Bit 11 **I2SMOD**: I2S mode selection 0: SPI mode is selected 1: I2S mode is selected

Note: This bit should be configured when the SPI or I<sup>2</sup>S is disabled

Bit 10 I2SE: I2S Enable

0: I<sup>2</sup>S peripheral is disabled 1: I<sup>2</sup>S peripheral is enabled Note: Not used in SPI mode

Bit 9:8 I2SCFG: I2S configuration mode

00: Slave - transmit01: Slave - receive10: Master - transmit11: Master - receive

Note: This bit should be configured when the  $l^2S$  is disabled.

Not used in SPI mode

Bit 7 PCMSYNC: PCM frame synchronization

0: Short frame synchronization1: Long frame synchronization

Note: This bit has a meaning only if I2SSTD = 11 (PCM standard is used)

Not used in SPI mode

Bit 6 Reserved: forced at 0 by hardware

Bit 5:4 I2SSTD: I2S standard selection

00: I<sup>2</sup>S Philips standard.

01: MSB justified standard (left justified)

10: LSB justified standard (right justified)

11: PCM standard

For more details on I<sup>2</sup>S standards, refer to Section 25.4.3 on page 703. Not used in SPI mode.

Note: For correct operation, these bits should be configured when the  $l^2S$  is disabled.

Bit 3 CKPOL: Steady state clock polarity

0: I<sup>2</sup>S clock steady state is low level

1: I<sup>2</sup>S clock steady state is high level

Note: For correct operation, this bit should be configured when the  $\hat{F}S$  is disabled.

Not used in SPI mode

#### Bit 2:1 DATLEN: Data length to be transferred

00: 16-bit data length

01: 24-bit data length

10: 32-bit data length

11: Not allowed

Note: For correct operation, these bits should be configured when the  $l^2S$  is disabled.

Not used in SPI mode.

#### Bit 0 CHLEN: Channel length (number of bits per audio channel)

0: 16-bit wide

1: 32-bit wide

The bit write operation has a meaning only if DATLEN = 00 otherwise the channel length is fixed to 32-bit by hardware whatever the value filled in. Not used in SPI mode.

Note: For correct operation, this bit should be configured when the  ${}^{\rho}S$  is disabled.

#### SPI\_I<sup>2</sup>S prescaler register (SPI\_I2SPR) 25.5.9

Address offset: 0x20

Reset value: 0000 0010 (0x0002)

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		Page	nuod			MCKOE	ODD				I2S	DIV			
		Rese	riveu			rw	rw				r	N			

#### Bits 15:10 Reserved, must be kept at reset value.

#### Bit 9 MCKOE: Master clock output enable

0: Master clock output is disabled

1: Master clock output is enabled

Note: This bit should be configured when the  $l^2S$  is disabled. It is used only when the  $l^2S$  is in master mode.

Not used in SPI mode.

#### Bit 8 **ODD**: Odd factor for the prescaler

0: real divider value is = I2SDIV \*2

1: real divider value is = (I2SDIV \* 2)+1

Refer to Section 25.4.4 on page 709. Not used in SPI mode.

Note: This bit should be configured when the  $l^2S$  is disabled. It is used only when the  $l^2S$  is in master mode.

# Bit 7:0 I2SDIV: I2S Linear prescaler

I2SDIV [7:0] = 0 or I2SDIV [7:0] = 1 are forbidden values.

Refer to Section 25.4.4 on page 709. Not used in SPI mode.

Note: These bits should be configured when the I<sup>2</sup>S is disabled. It is used only when the I<sup>2</sup>S is in master mode.

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# 25.5.10 SPI register map

The table provides shows the SPI register map and reset values.

Table 106. SPI register map and reset values

Offset	Register	30 29 27 27 26	25 24	23	21 20	19	17	15	14	13	12	11	10	<b>D</b> 0	<b>^</b>	. 9	2	4	3	7		0
0x00	SPI_CR1		Reser					BIDIMODE		CRCEN	CRCNEXT	DFF	RXONLY		LSE		ВІ	R [2	:0]	MSTR	CPOL	CPITA
	Reset value							0	0	0	0	0	0	0 0	0	0	0	0	0	0	0	0
0x04	SPI_CR2				Rese	rved									TXEIE	RXNEIE	ERRIE	FRF	Reserved	SSOE	TXDMAEN	OHXUMAEN
	Reset value														0	0	0	0			0	0
0x08	SPI_SR			ſ	Reserv	ed								FRF	BSY	OVR	MODF	CRCERR	UDR	CHSIDE	TXE	HAINE
	Reset value													(	0	0	0	0	0	0	1	0
0x0C	SPI_DR		Reser	hav											[15:							
0,00	Reset value		110301	vcu				0	0	0	0	0		0 0				0	0	0	0	0
0x10	SPI_CRCPR Reset value		Reser	ved				0	0	0	0	0		RCP 0   0				0	0	11	11	1
0.44	SPI_RXCRCR							Ė		<u> </u>	لت			RxCl					اتا		ــــــــــــــــــــــــــــــــــــــ	Ė
0x14	Reset value		Reser	vea				0	0	0	0	0	0	0 0	0	0	0	0	0	0	0	0
0x18	SPI_TXCRCR		Reser	wod										ГхСГ	RC[1	5:0]						T
0.10	Reset value		116361	veu				0	0	0	0	0	0	0 0	0	0	0	0	0	0	0	0
0x1C	SPI_I2SCFGR			Reser	ved							ISSMOD	ISSE	12SCFG	PCMSYNC	Reserved	USCATO	200	CKPOL	DATLEN	i	CHLEIN
	Reset value											0	0	0 (	0	ľ	0	0	0	0	0	0
0x20	SPI_I2SPR			Re	eserved	i							10/1014					I2S				
	Reset value													0 0	0	0	0	0	0	0	1	0

Refer to *Table 2 on page 50* for the register boundary addresses.

# 26 Secure digital input/output interface (SDIO)

This section applies to the whole STM32F4xx family, unless otherwise specified.

# 26.1 SDIO main features

The SD/SDIO MMC card host interface (SDIO) provides an interface between the APB2 peripheral bus and MultiMediaCards (MMCs), SD memory cards, SDIO cards and CE-ATA devices.

The MultiMediaCard system specifications are available through the MultiMediaCard Association website at *www.mmca.org*, published by the MMCA technical committee.

SD memory card and SD I/O card system specifications are available through the SD card Association website at *www.sdcard.org*.

CE-ATA system specifications are available through the CE-ATA workgroup website at www.ce-ata.org.

The SDIO features include the following:

- Full compliance with *MultiMediaCard System Specification Version 4.2*. Card support for three different databus modes: 1-bit (default), 4-bit and 8-bit
- Full compatibility with previous versions of MultiMediaCards (forward compatibility)
- Full compliance with SD Memory Card Specifications Version 2.0
- Full compliance with SD I/O Card Specification Version 2.0: card support for two different databus modes: 1-bit (default) and 4-bit
- Full support of the CE-ATA features (full compliance with CE-ATA digital protocol Rev1.1)
- Data transfer up to 48 MHz for the 8 bit mode
- Data and command output enable signals to control external bidirectional drivers.

Note:

The SDIO does not have an SPI-compatible communication mode.

The SD memory card protocol is a superset of the MultiMediaCard protocol as defined in the MultiMediaCard system specification V2.11. Several commands required for SD memory devices are not supported by either SD I/O-only cards or the I/O portion of combo cards. Some of these commands have no use in SD I/O devices, such as erase commands, and thus are not supported in the SDIO. In addition, several commands are different between SD memory cards and SD I/O cards and thus are not supported in the SDIO. For details refer to SD I/O card Specification Version 1.0. CE-ATA is supported over the MMC electrical interface using a protocol that utilizes the existing MMC access primitives. The interface electrical and signaling definition is as defined in the MMC reference.

The MultiMediaCard/SD bus connects cards to the controller.

The current version of the SDIO supports only one SD/SDIO/MMC4.2 card at any one time and a stack of MMC4.1 or previous.

# 26.2 SDIO bus topology

Communication over the bus is based on command and data transfers.

The basic transaction on the MultiMediaCard/SD/SD I/O bus is the command/response transaction. These types of bus transaction transfer their information directly within the command or response structure. In addition, some operations have a data token.

Data transfers to/from SD/SDIO memory cards are done in data blocks. Data transfers to/from MMC are done data blocks or streams. Data transfers to/from the CE-ATA Devices are done in data blocks.

Figure 286. SDIO "no response" and "no data" operations

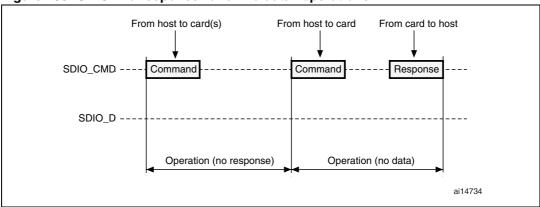
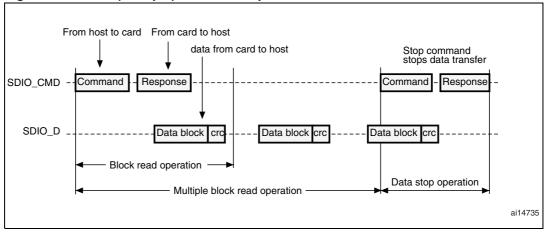


Figure 287. SDIO (multiple) block read operation



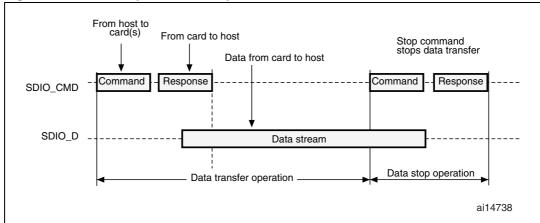
From host to card From card to host Stop command stops data transfer Data from host to card Response Response Command Command SDIO CMD SDIO\_D Data block Busy Data block crc Busy Optional cards Busy Data stop operation Needed for CE-ATA Block write operation Multiple block write operation

Figure 288. SDIO (multiple) block write operation

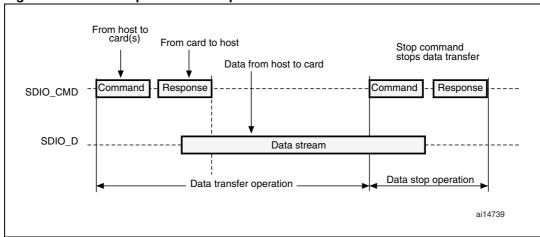
Note:

The SDIO will not send any data as long as the Busy signal is asserted (SDIO\_D0 pulled low).

Figure 289. SDIO sequential read operation





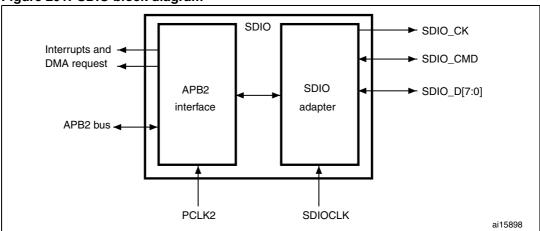


# 26.3 SDIO functional description

The SDIO consists of two parts:

- The SDIO adapter block provides all functions specific to the MMC/SD/SD I/O card such as the clock generation unit, command and data transfer.
- The APB2 interface accesses the SDIO adapter registers, and generates interrupt and DMA request signals.

Figure 291. SDIO block diagram



By default SDIO\_D0 is used for data transfer. After initialization, the host can change the databus width.

If a MultiMediaCard is connected to the bus, SDIO\_D0, SDIO\_D[3:0] or SDIO\_D[7:0] can be used for data transfer. MMC V3.31 or previous, supports only 1 bit of data so only SDIO\_D0 can be used.

If an SD or SD I/O card is connected to the bus, data transfer can be configured by the host to use SDIO\_D0 or SDIO\_D[3:0]. All data lines are operating in push-pull mode.

# SDIO\_CMD has two operational modes:

- Open-drain for initialization (only for MMCV3.31 or previous)
- Push-pull for command transfer (SD/SD I/O card MMC4.2 use push-pull drivers also for initialization)

**SDIO\_CK** is the clock to the card: one bit is transferred on both command and data lines with each clock cycle. The clock frequency can vary between 0 MHz and 20 MHz (for a MultiMediaCard V3.31), between 0 and 48 MHz for a MultiMediaCard V4.0/4.2, or between 0 and 25 MHz (for an SD/SD I/O card).

The SDIO uses two clock signals:

- SDIO adapter clock (SDIOCLK = 48 MHz)
- APB2 bus clock (PCLK2)

PCLK2 and SDIO\_CK clock frequencies must respect the following condition:

 $Frequenc(PCLK2) \ge 3/8 \times Frequency(SDIO\_CK)$ 

The signals shown in *Table 107* are used on the MultiMediaCard/SD/SD I/O card bus.

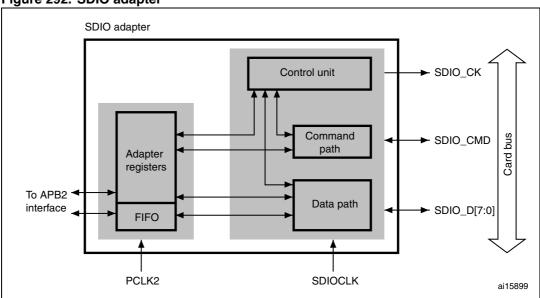
Table 107. SDIO I/O definitions

Pin	Direction	Description
SDIO_CK	Output	MultiMediaCard/SD/SDIO card clock. This pin is the clock from host to card.
SDIO_CMD	Bidirectional	MultiMediaCard/SD/SDIO card command. This pin is the bidirectional command/response signal.
SDIO_D[7:0]	Bidirectional	MultiMediaCard/SD/SDIO card data. These pins are the bidirectional databus.

# 26.3.1 SDIO adapter

Figure 292 shows a simplified block diagram of an SDIO adapter.

Figure 292. SDIO adapter



The SDIO adapter is a multimedia/secure digital memory card bus master that provides an interface to a multimedia card stack or to a secure digital memory card. It consists of five subunits:

- Adapter register block
- Control unit
- Command path
- Data path
- Data FIFO

Note:

The adapter registers and FIFO use the APB2 bus clock domain (PCLK2). The control unit, command path and data path use the SDIO adapter clock domain (SDIOCLK).

# Adapter register block

The adapter register block contains all system registers. This block also generates the signals that clear the static flags in the multimedia card. The clear signals are generated when 1 is written into the corresponding bit location in the SDIO Clear register.



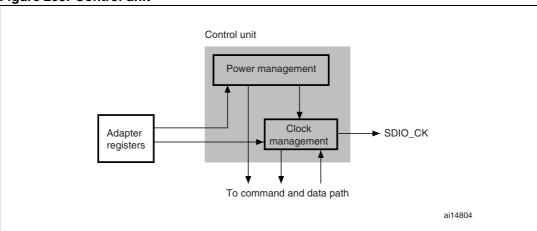
#### **Control unit**

The control unit contains the power management functions and the clock divider for the memory card clock.

There are three power phases:

- power-off
- power-up
- power-on

Figure 293. Control unit



The control unit is illustrated in *Figure 293*. It consists of a power management subunit and a clock management subunit.

The power management subunit disables the card bus output signals during the power-off and power-up phases.

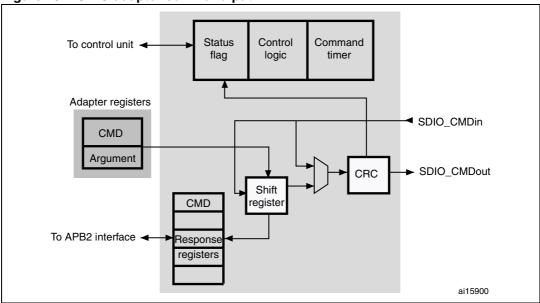
The clock management subunit generates and controls the SDIO\_CK signal. The SDIO\_CK output can use either the clock divide or the clock bypass mode. The clock output is inactive:

- after reset
- during the power-off or power-up phases
- if the power saving mode is enabled and the card bus is in the Idle state (eight clock periods after both the command and data path subunits enter the Idle phase)

# **Command path**

The command path unit sends commands to and receives responses from the cards.

Figure 294. SDIO adapter command path



- Command path state machine (CPSM)
  - When the command register is written to and the enable bit is set, command transfer starts. When the command has been sent, the command path state machine (CPSM) sets the status flags and enters the Idle state if a response is not required. If a response is required, it waits for the response (see *Figure 295 on page 734*). When the response is received, the received CRC code and the internally generated code are compared, and the appropriate status flags are set.

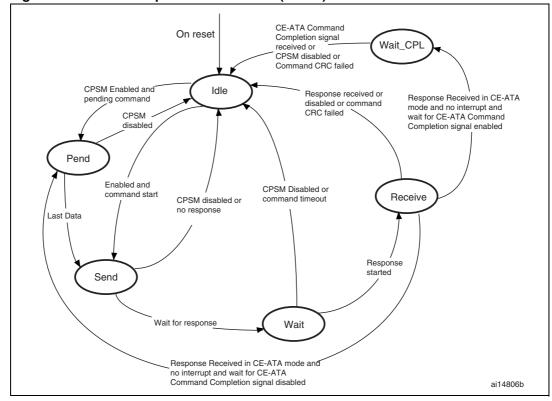


Figure 295. Command path state machine (CPSM)

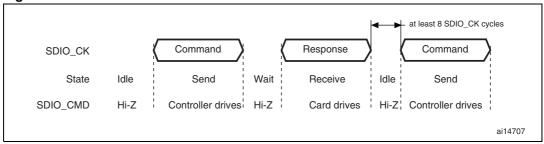
When the Wait state is entered, the command timer starts running. If the timeout is reached before the CPSM moves to the Receive state, the timeout flag is set and the Idle state is entered.

Note: The command timeout has a fixed value of 64 SDIO\_CK clock periods.

If the interrupt bit is set in the command register, the timer is disabled and the CPSM waits for an interrupt request from one of the cards. If a pending bit is set in the command register, the CPSM enters the Pend state, and waits for a CmdPend signal from the data path subunit. When CmdPend is detected, the CPSM moves to the Send state. This enables the data counter to trigger the stop command transmission.

Note: The CPSM remains in the Idle state for at least eight SDIO\_CK periods to meet the  $N_{CC}$  and  $N_{RC}$  timing constraints.  $N_{CC}$  is the minimum delay between two host commands, and  $N_{RC}$  is the minimum delay between the host command and the card response.

Figure 296. SDIO command transfer



#### Command format

Command: a command is a token that starts an operation. Command are sent from the host either to a single card (addressed command) or to all connected cards (broadcast command are available for MMC V3.31 or previous). Commands are transferred serially on the CMD line. All commands have a fixed length of 48 bits. The general format for a command token for MultiMediaCards, SD-Memory cards and SDIO-Cards is shown in *Table 108*. CE-ATA commands are an extension of MMC commands V4.2, and so have the same format.

The command path operates in a half-duplex mode, so that commands and responses can either be sent or received. If the CPSM is not in the Send state, the SDIO\_CMD output is in the Hi-Z state, as shown in *Figure 296 on page 735*. Data on SDIO\_CMD are synchronous with the rising edge of SDIO\_CK. *Table* shows the command format.

Table 108. Command format

Bit position	Width	Value	Description
47	1	0	Start bit
46	1	1	Transmission bit
[45:40]	6	-	Command index
[39:8]	32	-	Argument
[7:1]	7	-	CRC7
0	1	1	End bit

 Response: a response is a token that is sent from an addressed card (or synchronously from all connected cards for MMC V3.31 or previous), to the host as an answer to a previously received command. Responses are transferred serially on the CMD line.

The SDIO supports two response types. Both use CRC error checking:

- 48 bit short response
- 136 bit long response

Note:

If the response does not contain a CRC (CMD1 response), the device driver must ignore the CRC failed status.

Bit position Width Value **Description** 47 0 Start bit 46 0 Transmission bit [45:40] 6 Command index [39:8] 32 Argument 7 [7:1] CRC7(or 1111111) 0 1 1 End bit

Table 109. Short response format

Table 110. Long response format

Bit position	Width	Value	Description
135	1	0	Start bit
134	1	0	Transmission bit
[133:128]	6	111111	Reserved
[127:1]	127	-	CID or CSD (including internal CRC7)
0	1	1	End bit

The command register contains the command index (six bits sent to a card) and the command type. These determine whether the command requires a response, and whether the response is 48 or 136 bits long (see *Section 26.9.4 on page 771*). The command path implements the status flags shown in *Table 111*:

Table 111. Command path status flags

Flag	Description
CMDREND	Set if response CRC is OK.
CCRCFAIL	Set if response CRC fails.
CMDSENT	Set when command (that does not require response) is sent
CTIMEOUT	Response timeout.
CMDACT	Command transfer in progress.

The CRC generator calculates the CRC checksum for all bits before the CRC code. This includes the start bit, transmitter bit, command index, and command argument (or card status). The CRC checksum is calculated for the first 120 bits of CID or CSD for the long response format. Note that the start bit, transmitter bit and the six reserved bits are not used in the CRC calculation.

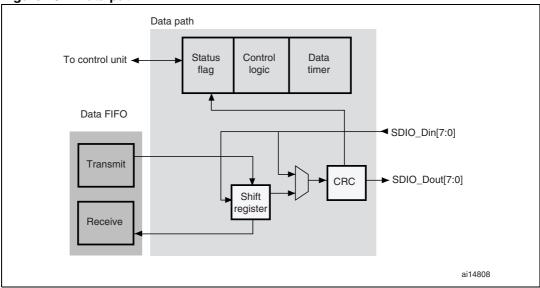
The CRC checksum is a 7-bit value:

CRC[6:0] = Remainder [(M(x) \* 
$$x^7$$
) / G(x)]  
G(x) =  $x^7 + x^3 + 1$   
M(x) = (start bit) \*  $x^{39} + ... +$  (last bit before CRC) \*  $x^0$ , or  
M(x) = (start bit) \*  $x^{119} + ... +$  (last bit before CRC) \*  $x^0$ 

# Data path

The data path subunit transfers data to and from cards. *Figure 297* shows a block diagram of the data path.

Figure 297. Data path



The card databus width can be programmed using the clock control register. If the 4-bit wide bus mode is enabled, data is transferred at four bits per clock cycle over all four data signals (SDIO\_D[3:0]). If the 8-bit wide bus mode is enabled, data is transferred at eight bits per clock cycle over all eight data signals (SDIO\_D[7:0]). If the wide bus mode is not enabled, only one bit per clock cycle is transferred over SDIO\_D0.

Depending on the transfer direction (send or receive), the data path state machine (DPSM) moves to the Wait S or Wait R state when it is enabled:

- Send: the DPSM moves to the Wait\_S state. If there is data in the transmit FIFO, the DPSM moves to the Send state, and the data path subunit starts sending data to a card.
- Receive: the DPSM moves to the Wait\_R state and waits for a start bit. When it
  receives a start bit, the DPSM moves to the Receive state, and the data path subunit
  starts receiving data from a card.

Data path state machine (DPSM)

The DPSM operates at SDIO\_CK frequency. Data on the card bus signals is synchronous to the rising edge of SDIO\_CK. The DPSM has six states, as shown in *Figure 298: Data path state machine (DPSM)*.

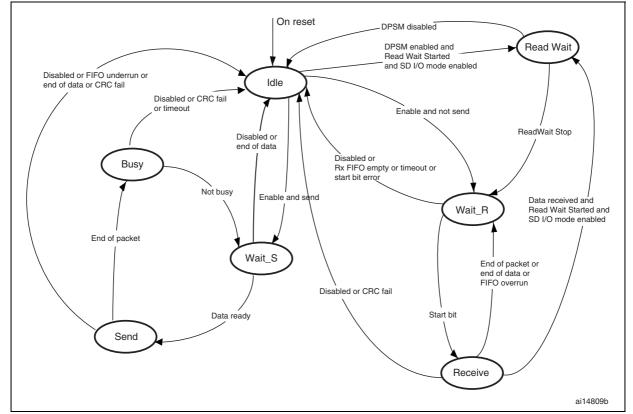


Figure 298. Data path state machine (DPSM)

- Idle: the data path is inactive, and the SDIO\_D[7:0] outputs are in Hi-Z. When the data
  control register is written and the enable bit is set, the DPSM loads the data counter
  with a new value and, depending on the data direction bit, moves to either the Wait\_S
  or the Wait\_R state.
- Wait\_R: if the data counter equals zero, the DPSM moves to the Idle state when the
  receive FIFO is empty. If the data counter is not zero, the DPSM waits for a start bit on
  SDIO\_D. The DPSM moves to the Receive state if it receives a start bit before a
  timeout, and loads the data block counter. If it reaches a timeout before it detects a
  start bit, or a start bit error occurs, it moves to the Idle state and sets the timeout status
  flag.
- Receive: serial data received from a card is packed in bytes and written to the data FIFO. Depending on the transfer mode bit in the data control register, the data transfer mode can be either block or stream:
  - In block mode, when the data block counter reaches zero, the DPSM waits until it receives the CRC code. If the received code matches the internally generated CRC code, the DPSM moves to the Wait\_R state. If not, the CRC fail status flag is set and the DPSM moves to the Idle state.
  - In stream mode, the DPSM receives data while the data counter is not zero. When
    the counter is zero, the remaining data in the shift register is written to the data
    FIFO, and the DPSM moves to the Wait\_R state.

If a FIFO overrun error occurs, the DPSM sets the FIFO error flag and moves to the Idle state:

Wait\_S: the DPSM moves to the Idle state if the data counter is zero. If not, it waits until
the data FIFO empty flag is deasserted, and moves to the Send state.

Note:

The DPSM remains in the Wait\_S state for at least two clock periods to meet the  $N_{WR}$  timing requirements, where  $N_{WR}$  is the number of clock cycles between the reception of the card response and the start of the data transfer from the host.

- Send: the DPSM starts sending data to a card. Depending on the transfer mode bit in the data control register, the data transfer mode can be either block or stream:
  - In block mode, when the data block counter reaches zero, the DPSM sends an internally generated CRC code and end bit, and moves to the Busy state.
  - In stream mode, the DPSM sends data to a card while the enable bit is high and the data counter is not zero. It then moves to the Idle state.

If a FIFO underrun error occurs, the DPSM sets the FIFO error flag and moves to the Idle state.

- Busy: the DPSM waits for the CRC status flag:
  - If it does not receive a positive CRC status, it moves to the Idle state and sets the CRC fail status flag.
  - If it receives a positive CRC status, it moves to the Wait\_S state if SDIO\_D0 is not low (the card is not busy).

If a timeout occurs while the DPSM is in the Busy state, it sets the data timeout flag and moves to the Idle state.

The data timer is enabled when the DPSM is in the Wait\_R or Busy state, and generates the data timeout error:

- When transmitting data, the timeout occurs if the DPSM stays in the Busy state for longer than the programmed timeout period
- When receiving data, the timeout occurs if the end of the data is not true, and if the DPSM stays in the Wait\_R state for longer than the programmed timeout period.
- Data: data can be transferred from the card to the host or vice versa. Data is transferred via the data lines. They are stored in a FIFO of 32 words, each word is 32 bits wide.

Table 112. Data token format

Description	Start bit	Data	CRC16	End bit
Block Data	0	-	yes	1
Stream Data	0	-	no	1

# **Data FIFO**

The data FIFO (first-in-first-out) subunit is a data buffer with a transmit and receive unit.

The FIFO contains a 32-bit wide, 32-word deep data buffer, and transmit and receive logic. Because the data FIFO operates in the APB2 clock domain (PCLK2), all signals from the subunits in the SDIO clock domain (SDIOCLK) are resynchronized.

Depending on the TXACT and RXACT flags, the FIFO can be disabled, transmit enabled, or receive enabled. TXACT and RXACT are driven by the data path subunit and are mutually exclusive:

- The transmit FIFO refers to the transmit logic and data buffer when TXACT is asserted
- The receive FIFO refers to the receive logic and data buffer when RXACT is asserted

#### Transmit FIFO:

Data can be written to the transmit FIFO through the APB2 interface when the SDIO is enabled for transmission.

The transmit FIFO is accessible via 32 sequential addresses. The transmit FIFO contains a data output register that holds the data word pointed to by the read pointer. When the data path subunit has loaded its shift register, it increments the read pointer and drives new data out.

If the transmit FIFO is disabled, all status flags are deasserted. The data path subunit asserts TXACT when it transmits data.

Table 113. Transmit FIFO status flags

Flag	Description
TXFIFOF	Set to high when all 32 transmit FIFO words contain valid data.
TXFIFOE	Set to high when the transmit FIFO does not contain valid data.
TXFIFOHE	Set to high when 8 or more transmit FIFO words are empty. This flag can be used as a DMA request.
TXDAVL	Set to high when the transmit FIFO contains valid data. This flag is the inverse of the TXFIFOE flag.
TXUNDERR	Set to high when an underrun error occurs. This flag is cleared by writing to the SDIO Clear register.

### Receive FIFO

When the data path subunit receives a word of data, it drives the data on the write databus. The write pointer is incremented after the write operation completes. On the read side, the contents of the FIFO word pointed to by the current value of the read pointer is driven onto the read databus. If the receive FIFO is disabled, all status flags are deasserted, and the read and write pointers are reset. The data path subunit asserts RXACT when it receives data. *Table 114* lists the receive FIFO status flags. The receive FIFO is accessible via 32 sequential addresses.

Flag Description **RXFIFOF** Set to high when all 32 receive FIFO words contain valid data **RXFIFOE** Set to high when the receive FIFO does not contain valid data. Set to high when 8 or more receive FIFO words contain valid data. This flag can be **RXFIFOHF** used as a DMA request. Set to high when the receive FIFO is not empty. This flag is the inverse of the **RXDAVL** RXFIFOE flag. Set to high when an overrun error occurs. This flag is cleared by writing to the SDIO **RXOVERR** Clear register.

Receive FIFO status flags Table 114.

#### 26.3.2 **SDIO APB2 interface**

The APB2 interface generates the interrupt and DMA requests, and accesses the SDIO adapter registers and the data FIFO. It consists of a data path, register decoder, and interrupt/DMA logic.

# **SDIO** interrupts

The interrupt logic generates an interrupt request signal that is asserted when at least one of the selected status flags is high. A mask register is provided to allow selection of the conditions that will generate an interrupt. A status flag generates the interrupt request if a corresponding mask flag is set.

# SDIO/DMA interface - procedure for data transfers between the SDIO and memory

In the example shown, the transfer is from the SDIO host controller to an MMC (512 bytes using CMD24 (WRITE BLOCK). The SDIO FIFO is filled by data stored in a memory using the DMA controller.

- Do the card identification process
- Increase the SDIO CK frequency
- Select the card by sending CMD7 3
- 4. Configure the DMA2 as follows:
  - Enable DMA2 controller and clear any pending interrupts.
  - Program the DMA2\_Stream3 or DMA2\_Stream6 Channel4 source address register with the memory location's base address and DMA2\_Stream3 or DMA2 Stream6 Channel4 destination address register with the SDIO FIFO register address.
  - Program DMA2 Stream3 or DMA2 Stream6 Channel4 control register (memory increment, not peripheral increment, peripheral and source width is word size).
  - Program DMA2 Stream3 or DMA2 Stream6 Channel4 to select the peripheral as flow controller (set PFCTRL bit in DMA\_S3CR or DMA\_S6CR configuration register).
  - Configure the incremental burst transfer to 4 beats (at least from peripheral side) in DMA2 Stream3 or DMA2 Stream6 Channel4.
  - f) Enable DMA2\_Stream3 or DMA2\_Stream6 Channel4



- Send CMD24 (WRITE BLOCK) as follows:
  - a) Program the SDIO data length register (SDIO data timer register should be already programmed before the card identification process).
  - b) Program the SDIO argument register with the address location of the card where data is to be transferred.
  - c) Program the SDIO command register: CmdIndex with 24 (WRITE\_BLOCK); WaitResp with '1' (SDIO card host waits for a response); CPSMEN with '1' (SDIO card host enabled to send a command). Other fields are at their reset value.
  - d) Wait for SDIO\_STA[6] = CMDREND interrupt, then program the SDIO data control register: DTEN with '1' (SDIO card host enabled to send data); DTDIR with '0' (from controller to card); DTMODE with '0' (block data transfer); DMAEN with '1' (DMA enabled); DBLOCKSIZE with 0x9 (512 bytes). Other fields are don't care.
  - e) Wait for SDIO\_STA[10] = DBCKEND.
- 6. Check that no channels are still enabled by polling the DMA Enabled Channel Status register.

# 26.4 Card functional description

#### 26.4.1 Card identification mode

While in card identification mode the host resets all cards, validates the operation voltage range, identifies cards and sets a relative card address (RCA) for each card on the bus. All data communications in the card identification mode use the command line (CMD) only.

#### 26.4.2 Card reset

The GO\_IDLE\_STATE command (CMD0) is the software reset command and it puts the MultiMediaCard and SD memory in the Idle state. The IO\_RW\_DIRECT command (CMD52) resets the SD I/O card. After power-up or CMD0, all cards output bus drivers are in the high-impedance state and the cards are initialized with a default relative card address (RCA=0x0001) and with a default driver stage register setting (lowest speed, highest driving current capability).

# 26.4.3 Operating voltage range validation

All cards can communicate with the SDIO card host using any operating voltage within the specification range. The supported minimum and maximum  $V_{DD}$  values are defined in the operation conditions register (OCR) on the card.

Cards that store the card identification number (CID) and card specific data (CSD) in the payload memory are able to communicate this information only under data-transfer  $V_{DD}$  conditions. When the SDIO card host module and the card have incompatible  $V_{DD}$  ranges, the card is not able to complete the identification cycle and cannot send CSD data. For this purpose, the special commands, SEND\_OP\_COND (CMD1), SD\_APP\_OP\_COND (ACMD41 for SD Memory), and IO\_SEND\_OP\_COND (CMD5 for SD I/O), are designed to provide a mechanism to identify and reject cards that do not match the  $V_{DD}$  range desired by the SDIO card host. The SDIO card host sends the required  $V_{DD}$  voltage window as the operand of these commands. Cards that cannot perform data transfer in the specified range disconnect from the bus and go to the inactive state.

By using these commands without including the voltage range as the operand, the SDIO card host can query each card and determine the common voltage range before placing out-of-range cards in the inactive state. This query is used when the SDIO card host is able to select a common voltage range or when the user requires notification that cards are not usable.

# 26.4.4 Card identification process

The card identification process differs for MultiMediaCards and SD cards. For MultiMediaCard cards, the identification process starts at clock rate  $F_{od}$ . The SDIO\_CMD line output drivers are open-drain and allow parallel card operation during this process. The registration process is accomplished as follows:

- 1. The bus is activated.
- 2. The SDIO card host broadcasts SEND\_OP\_COND (CMD1) to receive operation conditions.
- 3. The response is the wired AND operation of the operation condition registers from all cards.
- 4. Incompatible cards are placed in the inactive state.
- 5. The SDIO card host broadcasts ALL\_SEND\_CID (CMD2) to all active cards.
- 6. The active cards simultaneously send their CID numbers serially. Cards with outgoing CID bits that do not match the bits on the command line stop transmitting and must wait for the next identification cycle. One card successfully transmits a full CID to the SDIO card host and enters the Identification state.
- 7. The SDIO card host issues SET\_RELATIVE\_ADDR (CMD3) to that card. This new address is called the relative card address (RCA); it is shorter than the CID and addresses the card. The assigned card changes to the Standby state, it does not react to further identification cycles, and its output switches from open-drain to push-pull.
- 8. The SDIO card host repeats steps 5 through 7 until it receives a timeout condition.

For the SD card, the identification process starts at clock rate  $F_{od}$ , and the SDIO\_CMD line output drives are push-pull drivers instead of open-drain. The registration process is accomplished as follows:

- 1. The bus is activated.
- 2. The SDIO card host broadcasts SD APP OP COND (ACMD41).
- 3. The cards respond with the contents of their operation condition registers.
- 4. The incompatible cards are placed in the inactive state.
- 5. The SDIO card host broadcasts ALL SEND CID (CMD2) to all active cards.
- 6. The cards send back their unique card identification numbers (CIDs) and enter the Identification state.
- 7. The SDIO card host issues SET\_RELATIVE\_ADDR (CMD3) to an active card with an address. This new address is called the relative card address (RCA); it is shorter than the CID and addresses the card. The assigned card changes to the Standby state. The SDIO card host can reissue this command to change the RCA. The RCA of the card is the last assigned value.
- 8. The SDIO card host repeats steps 5 through 7 with all active cards.

For the SD I/O card, the registration process is accomplished as follows:

- 1. The bus is activated.
- 2. The SDIO card host sends IO SEND OP COND (CMD5).
- 3. The cards respond with the contents of their operation condition registers.
- 4. The incompatible cards are set to the inactive state.
- 5. The SDIO card host issues SET\_RELATIVE\_ADDR (CMD3) to an active card with an address. This new address is called the relative card address (RCA); it is shorter than the CID and addresses the card. The assigned card changes to the Standby state. The SDIO card host can reissue this command to change the RCA. The RCA of the card is the last assigned value.

#### 26.4.5 Block write

During block write (CMD24 - 27) one or more blocks of data are transferred from the host to the card with a CRC appended to the end of each block by the host. A card supporting block write is always able to accept a block of data defined by WRITE\_BL\_LEN. If the CRC fails, the card indicates the failure on the SDIO\_D line and the transferred data are discarded and not written, and all further transmitted blocks (in multiple block write mode) are ignored.

If the host uses partial blocks whose accumulated length is not block aligned and, block misalignment is not allowed (CSD parameter WRITE\_BLK\_MISALIGN is not set), the card will detect the block misalignment error before the beginning of the first misaligned block. (ADDRESS\_ERROR error bit is set in the status register). The write operation will also be aborted if the host tries to write over a write-protected area. In this case, however, the card will set the WP\_VIOLATION bit.

Programming of the CID and CSD registers does not require a previous block length setting. The transferred data is also CRC protected. If a part of the CSD or CID register is stored in ROM, then this unchangeable part must match the corresponding part of the receive buffer. If this match fails, then the card reports an error and does not change any register contents. Some cards may require long and unpredictable times to write a block of data. After receiving a block of data and completing the CRC check, the card begins writing and holds the SDIO\_D line low if its write buffer is full and unable to accept new data from a new WRITE\_BLOCK command. The host may poll the status of the card with a SEND\_STATUS command (CMD13) at any time, and the card will respond with its status. The READY\_FOR\_DATA status bit indicates whether the card can accept new data or whether the write process is still in progress. The host may deselect the card by issuing CMD7 (to select a different card), which will place the card in the Disconnect state and release the SDIO\_D line(s) without interrupting the write operation. When selecting the card again, it will reactivate busy indication by pulling SDIO\_D to low if programming is still in progress and the write buffer is unavailable.

### 26.4.6 Block read

In Block read mode the basic unit of data transfer is a block whose maximum size is defined in the CSD (READ\_BL\_LEN). If READ\_BL\_PARTIAL is set, smaller blocks whose start and end addresses are entirely contained within one physical block (as defined by READ\_BL\_LEN) may also be transmitted. A CRC is appended to the end of each block, ensuring data transfer integrity. CMD17 (READ\_SINGLE\_BLOCK) initiates a block read and after completing the transfer, the card returns to the Transfer state.

CMD18 (READ MULTIPLE BLOCK) starts a transfer of several consecutive blocks.

The host can abort reading at any time, within a multiple block operation, regardless of its type. Transaction abort is done by sending the stop transmission command.

If the card detects an error (for example, out of range, address misalignment or internal error) during a multiple block read operation (both types) it stops the data transmission and remains in the data state. The host must than abort the operation by sending the stop transmission command. The read error is reported in the response to the stop transmission command.

If the host sends a stop transmission command after the card transmits the last block of a multiple block operation with a predefined number of blocks, it is responded to as an illegal command, since the card is no longer in the data state. If the host uses partial blocks whose accumulated length is not block-aligned and block misalignment is not allowed, the card detects a block misalignment error condition at the beginning of the first misaligned block (ADDRESS\_ERROR error bit is set in the status register).

# 26.4.7 Stream access, stream write and stream read (MultiMediaCard only)

In stream mode, data is transferred in bytes and no CRC is appended at the end of each block.

# Stream write (MultiMediaCard only)

WRITE\_DAT\_UNTIL\_STOP (CMD20) starts the data transfer from the SDIO card host to the card, beginning at the specified address and continuing until the SDIO card host issues a stop command. When partial blocks are allowed (CSD parameter WRITE\_BL\_PARTIAL is set), the data stream can start and stop at any address within the card address space, otherwise it can only start and stop at block boundaries. Because the amount of data to be transferred is not determined in advance, a CRC cannot be used. When the end of the memory range is reached while sending data and no stop command is sent by the SD card host, any additional transferred data are discarded.

The maximum clock frequency for a stream write operation is given by the following equation fields of the card-specific data register:

$$\label{eq:maximumspeed} \begin{aligned} \text{Maximumspeed= MIN(TRANSPEED,} & \frac{(8 \times 2^{writebllen})(-\text{NSAC})}{\text{TAAC} \times \text{R2WFACTOR}}) \end{aligned}$$

- Maximumspeed = maximum write frequency
- TRANSPEED = maximum data transfer rate
- writebllen = maximum write data block length
- NSAC = data read access time 2 in CLK cycles
- TAAC = data read access time 1
- R2WFACTOR = write speed factor

If the host attempts to use a higher frequency, the card may not be able to process the data and stop programming, set the OVERRUN error bit in the status register, and while ignoring all further data transfer, wait (in the receive data state) for a stop command. The write operation is also aborted if the host tries to write over a write-protected area. In this case, however, the card sets the WP VIOLATION bit.

## Stream read (MultiMediaCard only)

READ DAT UNTIL STOP (CMD11) controls a stream-oriented data transfer.

This command instructs the card to send its data, starting at a specified address, until the SDIO card host sends STOP\_TRANSMISSION (CMD12). The stop command has an execution delay due to the serial command transmission and the data transfer stops after the end bit of the stop command. When the end of the memory range is reached while sending data and no stop command is sent by the SDIO card host, any subsequent data sent are considered undefined.

The maximum clock frequency for a stream read operation is given by the following equation and uses fields of the card specific data register.

$$\label{eq:maximumspeed} \begin{aligned} & \text{Maximumspeed= MIN(TRANSPEED,} \\ & \underbrace{\frac{(8 \times 2^{readbllen})(-\text{NSAC})}{\text{TAAC} \times \text{R2WFACTOR}}}) \end{aligned}$$

- Maximumspeed = maximum read frequency
- TRANSPEED = maximum data transfer rate
- readbllen = maximum read data block length
- writebllen = maximum write data block length
- NSAC = data read access time 2 in CLK cycles
- TAAC = data read access time 1
- R2WFACTOR = write speed factor

If the host attempts to use a higher frequency, the card is not able to sustain data transfer. If this happens, the card sets the UNDERRUN error bit in the status register, aborts the transmission and waits in the data state for a stop command.

# 26.4.8 Erase: group erase and sector erase

The erasable unit of the MultiMediaCard is the erase group. The erase group is measured in write blocks, which are the basic writable units of the card. The size of the erase group is a card-specific parameter and defined in the CSD.

The host can erase a contiguous range of Erase Groups. Starting the erase process is a three-step sequence.

First the host defines the start address of the range using the ERASE\_GROUP\_START (CMD35) command, next it defines the last address of the range using the ERASE\_GROUP\_END (CMD36) command and, finally, it starts the erase process by issuing the ERASE (CMD38) command. The address field in the erase commands is an Erase Group address in byte units. The card ignores all LSBs below the Erase Group size, effectively rounding the address down to the Erase Group boundary.

If an erase command is received out of sequence, the card sets the ERASE\_SEQ\_ERROR bit in the status register and resets the whole sequence.

If an out-of-sequence (neither of the erase commands, except SEND\_STATUS) command received, the card sets the ERASE\_RESET status bit in the status register, resets the erase sequence and executes the last command.

If the erase range includes write protected blocks, they are left intact and only unprotected blocks are erased. The WP\_ERASE\_SKIP status bit in the status register is set.

The card indicates that an erase is in progress by holding SDIO\_D low. The actual erase time may be quite long, and the host may issue CMD7 to deselect the card.

### 26.4.9 Wide bus selection or deselection

Wide bus (4-bit bus width) operation mode is selected or deselected using SET\_BUS\_WIDTH (ACMD6). The default bus width after power-up or GO\_IDLE\_STATE (CMD0) is 1 bit. SET\_BUS\_WIDTH (ACMD6) is only valid in a transfer state, which means that the bus width can be changed only after a card is selected by SELECT/DESELECT CARD (CMD7).

# 26.4.10 Protection management

Three write protection methods for the cards are supported in the SDIO card host module:

- 1. internal card write protection (card responsibility)
- mechanical write protection switch (SDIO card host module responsibility only)
- 3. password-protected card lock operation

# Internal card write protection

Card data can be protected against write and erase. By setting the permanent or temporary write-protect bits in the CSD, the entire card can be permanently write-protected by the manufacturer or content provider. For cards that support write protection of groups of sectors by setting the WP\_GRP\_ENABLE bit in the CSD, portions of the data can be protected, and the write protection can be changed by the application. The write protection is in units of WP\_GRP\_SIZE sectors as specified in the CSD. The SET\_WRITE\_PROT and CLR\_WRITE\_PROT commands control the protection of the addressed group. The SEND\_WRITE\_PROT command is similar to a single block read command. The card sends a data block containing 32 write protection bits (representing 32 write protect groups starting at the specified address) followed by 16 CRC bits. The address field in the write protect commands is a group address in byte units.

The card ignores all LSBs below the group size.

# Mechanical write protect switch

A mechanical sliding tab on the side of the card allows the user to set or clear the write protection on a card. When the sliding tab is positioned with the window open, the card is write-protected, and when the window is closed, the card contents can be changed. A matched switch on the socket side indicates to the SDIO card host module that the card is write-protected. The SDIO card host module is responsible for protecting the card. The position of the write protect switch is unknown to the internal circuitry of the card.

### Password protect

The password protection feature enables the SDIO card host module to lock and unlock a card with a password. The password is stored in the 128-bit PWD register and its size is set in the 8-bit PWD\_LEN register. These registers are nonvolatile so that a power cycle does not erase them. Locked cards respond to and execute certain commands. This means that the SDIO card host module is allowed to reset, initialize, select, and query for status, however it is not allowed to access data on the card. When the password is set (as indicated by a nonzero value of PWD\_LEN), the card is locked automatically after power-up. As with the CSD and CID register write commands, the lock/unlock commands are available in the transfer state only. In this state, the command does not include an address argument and the card must be selected before using it. The card lock/unlock commands have the structure and bus transaction types of a regular single-block write command. The transferred data block includes all of the required information for the command (the

password setting mode, the PWD itself, and card lock/unlock). The command data block size is defined by the SDIO card host module before it sends the card lock/unlock command, and has the structure shown in *Table 128*.

The bit settings are as follows:

- ERASE: setting it forces an erase operation. All other bits must be zero, and only the command byte is sent
- LOCK\_UNLOCK: setting it locks the card. LOCK\_UNLOCK can be set simultaneously with SET\_PWD, however not with CLR\_PWD
- CLR\_PWD: setting it clears the password data
- SET\_PWD: setting it saves the password data to memory
- PWD\_LEN: it defines the length of the password in bytes
- PWD: the password (new or currently used, depending on the command)

The following sections list the command sequences to set/reset a password, lock/unlock the card, and force an erase.

# Setting the password

- 1. Select a card (SELECT/DESELECT CARD, CMD7), if none is already selected.
- Define the block length (SET\_BLOCKLEN, CMD16) to send, given by the 8-bit card lock/unlock mode, the 8-bit PWD\_LEN, and the number of bytes of the new password. When a password replacement is done, the block size must take into account that both the old and the new passwords are sent with the command.
- 3. Send LOCK/UNLOCK (CMD42) with the appropriate data block size on the data line including the 16-bit CRC. The data block indicates the mode (SET\_PWD = 1), the length (PWD\_LEN), and the password (PWD) itself. When a password replacement is done, the length value (PWD\_LEN) includes the length of both passwords, the old and the new one, and the PWD field includes the old password (currently used) followed by the new password.
- 4. When the password is matched, the new password and its size are saved into the PWD and PWD\_LEN fields, respectively. When the old password sent does not correspond (in size and/or content) to the expected password, the LOCK\_UNLOCK\_FAILED error bit is set in the card status register, and the password is not changed.

The password length field (PWD\_LEN) indicates whether a password is currently set. When this field is nonzero, there is a password set and the card locks itself after power-up. It is possible to lock the card immediately in the current power session by setting the LOCK\_UNLOCK bit (while setting the password) or sending an additional command for card locking.

#### Resetting the password

- Select a card (SELECT/DESELECT CARD, CMD7), if none is already selected.
- 2. Define the block length (SET\_BLOCKLEN, CMD16) to send, given by the 8-bit card lock/unlock mode, the 8-bit PWD\_LEN, and the number of bytes in the currently used password.
- Send LOCK/UNLOCK (CMD42) with the appropriate data block size on the data line including the 16-bit CRC. The data block indicates the mode (CLR\_PWD = 1), the length (PWD\_LEN) and the password (PWD) itself. The LOCK\_UNLOCK bit is ignored.
- 4. When the password is matched, the PWD field is cleared and PWD\_LEN is set to 0. When the password sent does not correspond (in size and/or content) to the expected

password, the LOCK\_UNLOCK\_FAILED error bit is set in the card status register, and the password is not changed.

## Locking a card

- 1. Select a card (SELECT/DESELECT CARD, CMD7), if none is already selected.
- 2. Define the block length (SET\_BLOCKLEN, CMD16) to send, given by the 8-bit card lock/unlock mode (byte 0 in *Table 128*), the 8-bit PWD\_LEN, and the number of bytes of the current password.
- 3. Send LOCK/UNLOCK (CMD42) with the appropriate data block size on the data line including the 16-bit CRC. The data block indicates the mode (LOCK\_UNLOCK = 1), the length (PWD\_LEN), and the password (PWD) itself.
- 4. When the password is matched, the card is locked and the CARD\_IS\_LOCKED status bit is set in the card status register. When the password sent does not correspond (in size and/or content) to the expected password, the LOCK\_UNLOCK\_FAILED error bit is set in the card status register, and the lock fails.

It is possible to set the password and to lock the card in the same sequence. In this case, the SDIO card host module performs all the required steps for setting the password (see *Setting the password on page 748*), however it is necessary to set the LOCK\_UNLOCK bit in Step 3 when the new password command is sent.

When the password is previously set (PWD\_LEN is not 0), the card is locked automatically after power on reset. An attempt to lock a locked card or to lock a card that does not have a password fails and the LOCK\_UNLOCK\_FAILED error bit is set in the card status register.

## Unlocking the card

- Select a card (SELECT/DESELECT CARD, CMD7), if none is already selected.
- Define the block length (SET\_BLOCKLEN, CMD16) to send, given by the 8-bit cardlock/unlock mode (byte 0 in *Table 128*), the 8-bit PWD\_LEN, and the number of bytes of the current password.
- 3. Send LOCK/UNLOCK (CMD42) with the appropriate data block size on the data line including the 16-bit CRC. The data block indicates the mode (LOCK\_UNLOCK = 0), the length (PWD\_LEN), and the password (PWD) itself.
- 4. When the password is matched, the card is unlocked and the CARD\_IS\_LOCKED status bit is cleared in the card status register. When the password sent is not correct in size and/or content and does not correspond to the expected password, the LOCK\_UNLOCK\_FAILED error bit is set in the card status register, and the card remains locked.

The unlocking function is only valid for the current power session. When the PWD field is not clear, the card is locked automatically on the next power-up.

An attempt to unlock an unlocked card fails and the LOCK\_UNLOCK\_FAILED error bit is set in the card status register.

# Forcing erase

If the user has forgotten the password (PWD content), it is possible to access the card after clearing all the data on the card. This forced erase operation erases all card data and all password data.

- 1. Select a card (SELECT/DESELECT CARD, CMD7), if none is already selected.
- 2. Set the block length (SET\_BLOCKLEN, CMD16) to 1 byte. Only the 8-bit card lock/unlock byte (byte 0 in *Table 128*) is sent.
- 3. Send LOCK/UNLOCK (CMD42) with the appropriate data byte on the data line including the 16-bit CRC. The data block indicates the mode (ERASE = 1). All other bits must be zero.
- 4. When the ERASE bit is the only bit set in the data field, all card contents are erased, including the PWD and PWD\_LEN fields, and the card is no longer locked. When any other bits are set, the LOCK\_UNLOCK\_FAILED error bit is set in the card status register and the card retains all of its data, and remains locked.

An attempt to use a force erase on an unlocked card fails and the LOCK\_UNLOCK\_FAILED error bit is set in the card status register.

# 26.4.11 Card status register

The response format R1 contains a 32-bit field named card status. This field is intended to transmit the card status information (which may be stored in a local status register) to the host. If not specified otherwise, the status entries are always related to the previously issued command.

*Table 115* defines the different entries of the status. The type and clear condition fields in the table are abbreviated as follows:

#### Type:

- E: error bit
- S: status bit
- R: detected and set for the actual command response
- X: detected and set during command execution. The SDIO card host must poll the card by issuing the status command to read these bits.

#### Clear condition:

- A: according to the card current state
- B: always related to the previous command. Reception of a valid command clears it (with a delay of one command)
- C: clear by read

Table 115. Card status

Bits	Identifier	Туре	Value	Description	Clear condition
31	ADDRESS_ OUT_OF_RANGE	ERX	'0'= no error '1'= error	The command address argument was out of the allowed range for this card.  A multiple block or stream read/write operation is (although started in a valid address) attempting to read or write beyond the card capacity.	С
30	ADDRESS_MISALIGN		'0'= no error '1'= error	The commands address argument (in accordance with the currently set block length) positions the first data block misaligned to the card physical blocks.  A multiple block read/write operation (although started with a valid address/block-length combination) is attempting to read or write a data block which is not aligned with the physical blocks of the card.	С
29	BLOCK_LEN_ERROR		'0'= no error '1'= error	Either the argument of a SET_BLOCKLEN command exceeds the maximum value allowed for the card, or the previously defined block length is illegal for the current command (e.g. the host issues a write command, the current block length is smaller than the maximum allowed value for the card and it is not allowed to write partial blocks)	С
28	ERASE_SEQ_ERROR		'0'= no error '1'= error	An error in the sequence of erase commands occurred.	С
27	ERASE_PARAM	ΕX	'0'= no error '1'= error	An invalid selection of erase groups for erase occurred.	С
26	WP_VIOLATION	ΕX	'0'= no error '1'= error	Attempt to program a write-protected block.	O
25	CARD_IS_LOCKED	SR	'0' = card unlocked '1' = card locked	When set, signals that the card is locked by the host	А
24	LOCK_UNLOCK_ FAILED	ΕX	'0'= no error '1'= error	Set when a sequence or password error has been detected in lock/unlock card command	С
23	COM_CRC_ERROR	ER	'0'= no error '1'= error	The CRC check of the previous command failed.	В
22	ILLEGAL_COMMAND	ΕR	'0'= no error '1'= error	Command not legal for the card state	В
21	CARD_ECC_FAILED	ΕX	'0'= success '1'= failure	Card internal ECC was applied but failed to correct the data.	С
20	CC_ERROR	ER	'0'= no error '1'= error	(Undefined by the standard) A card error occurred, which is not related to the host command.	С

Table 115. Card status (continued)

Bits	Identifier	Туре	Value	Description	Clear condition
19	ERROR	ΕX	'0'= no error '1'= error	(Undefined by the standard) A generic card error related to the (and detected during) execution of the last host command (e.g. read or write failures).	С
18	Reserved				
17	Reserved				
16	CID/CSD_OVERWRITE	ΕX	'0'= no error '1'= error	Can be either of the following errors:  - The CID register has already been written and cannot be overwritten  - The read-only section of the CSD does not match the card contents  - An attempt to reverse the copy (set as original) or permanent WP (unprotected) bits was made	С
15	WP_ERASE_SKIP	ΕX	'0'= not protected '1'= protected	Set when only partial address space was erased due to existing write	С
14	CARD_ECC_DISABLED	sx	'0'= enabled '1'= disabled	The command has been executed without using the internal ECC.	А
13	ERASE_RESET		'0'= cleared '1'= set	An erase sequence was cleared before executing because an out of erase sequence command was received (commands other than CMD35, CMD36, CMD38 or CMD13)	С
12:9	CURRENT_STATE	SR	0 = Idle 1 = Ready 2 = Ident 3 = Stby 4 = Tran 5 = Data 6 = Rcv 7 = Prg 8 = Dis 9 = Btst 10-15 = reserved	The state of the card when receiving the command. If the command execution causes a state change, it will be visible to the host in the response on the next command. The four bits are interpreted as a binary number between 0 and 15.	В
8	READY_FOR_DATA	SR	'0'= not ready '1' = ready	Corresponds to buffer empty signalling on the bus	
7	SWITCH_ERROR	ΕX	'0'= no error '1'= switch error	If set, the card did not switch to the expected mode as requested by the SWITCH command	В
6	Reserved				
5	APP_CMD	SR	'0' = Disabled '1' = Enabled	The card will expect ACMD, or an indication that the command has been interpreted as ACMD	С
4	Reserved for SD I/O Card	•			•

Table 115. Card status (continued)

Bits	Identifier	Туре	Value	Description	Clear condition
3	AKE_SEQ_ERROR	ER	'0'= no error '1'= error	Error in the sequence of the authentication process	С
2	Reserved for application specific commands				
1	Deserved for manufacturer test made				
0	Treserved for manufacture	Reserved for manufacturer test mode			

# 26.4.12 SD status register

The SD status contains status bits that are related to the SD memory card proprietary features and may be used for future application-specific usage. The size of the SD Status is one data block of 512 bits. The contents of this register are transmitted to the SDIO card host if ACMD13 is sent (CMD55 followed with CMD13). ACMD13 can be sent to a card in transfer state only (card is selected).

*Table 116* defines the different entries of the SD status register. The type and clear condition fields in the table are abbreviated as follows:

# Type:

- E: error bit
- S: status bit
- R: detected and set for the actual command response
- X: detected and set during command execution. The SDIO card Host must poll the card by issuing the status command to read these bits

#### Clear condition:

- A: according to the card current state
- B: always related to the previous command. Reception of a valid command clears it (with a delay of one command)
- C: clear by read

Table 116. SD status

Bits	Identifier	Туре	Value	Description	Clear condition
511: 510	DAT_BUS_WIDTH	SR	'00'= 1 (default) '01'= reserved '10'= 4 bit width '11'= reserved	Shows the currently defined databus width that was defined by SET_BUS_WIDTH command	A
509	SECURED_MODE	SR	'0'= Not in the mode '1'= In Secured Mode	Card is in Secured Mode of operation (refer to the "SD Security Specification").	А
508: 496	Reserved				•

Table 116. SD status (continued)

Bits	Identifier	Туре	Value	Description	Clear condition
495: 480	SD_CARD_TYPE	SR	'00xxh'= SD Memory Cards as defined in Physical Spec Ver1.01-2.00 ('x'= don't care). The following cards are currently defined: '0000'= Regular SD RD/WR Card. '0001'= SD ROM Card	In the future, the 8 LSBs will be used to define different variations of an SD memory card (each bit will define different SD types). The 8 MSBs will be used to define SD Cards that do not comply with current SD physical layer specification.	Α
479: 448	SIZE_OF_PROTE CT ED_AREA	SR	Size of protected area (See below)	(See below)	А
447: 440	SPEED_CLASS	SR	Speed Class of the card (See below)	(See below)	А
439: 432	PERFORMANCE_ MOVE	SR	Performance of move indicated by 1 [MB/s] step. (See below)	(See below)	A
431:428	AU_SIZE	SR	Size of AU (See below)	(See below)	А
427:424	Reserved				
423:408	ERASE_SIZE	SR	Number of AUs to be erased at a time	(See below)	А
407:402	ERASE_TIMEOUT	SR	Timeout value for erasing areas specified by UNIT_OF_ERASE_AU	(See below)	A
401:400	ERASE_OFFSET	SR	Fixed offset value added to erase time.	(See below)	А
399:312	Reserved				
311:0	Reserved for Manufacturer				

# SIZE\_OF\_PROTECTED\_AREA

Setting this field differs between standard- and high-capacity cards. In the case of a standard-capacity card, the capacity of protected area is calculated as follows:

Protected area = SIZE\_OF\_PROTECTED\_AREA\_\* MULT \* BLOCK\_LEN. SIZE\_OF\_PROTECTED\_AREA is specified by the unit in MULT\*BLOCK\_LEN.

In the case of a high-capacity card, the capacity of protected area is specified in this field:

Protected area = SIZE\_OF\_PROTECTED\_AREA

SIZE\_OF\_PROTECTED\_AREA is specified by the unit in bytes.

# SPEED\_CLASS

This 8-bit field indicates the speed class and the value can be calculated by  $P_W/2$  (where  $P_W$  is the write performance).

Table 117. Speed class code field

SPEED_CLASS	Value definition
00h	Class 0
01h	Class 2
02h	Class 4
03h	Class 6
04h – FFh	Reserved

# PERFORMANCE\_MOVE

This 8-bit field indicates Pm (performance move) and the value can be set by 1 [MB/sec] steps. If the card does not move used RUs (recording units), Pm should be considered as infinity. Setting the field to FFh means infinity.

Table 118. Performance move field

PERFORMANCE_MOVE	Value definition
00h	Not defined
01h	1 [MB/sec]
02h	02h 2 [MB/sec]
FEh	254 [MB/sec]
FFh	Infinity

# AU\_SIZE

This 4-bit field indicates the AU size and the value can be selected in the power of 2 base from 16 KB.

Table 119. AU\_SIZE field

AU_SIZE	Value definition
00h	Not defined
01h	16 KB
02h	32 KB
03h	64 KB
04h	128 KB
05h	256 KB
06h	512 KB
07h	1 MB
08h	2 MB

Table 119. AU SIZE field (continued)

AU_SIZE	Value definition
09h	4 MB
Ah – Fh	Reserved

The maximum AU size, which depends on the card capacity, is defined in *Table 120*. The card can be set to any AU size between RU size and maximum AU size.

Table 120. Maximum AU size

Capacity	16 MB-64 MB	128 MB-256 MB	512 MB	1 GB-32 GB
Maximum AU Size	512 KB	1 MB	2 MB	4 MB

# **ERASE\_SIZE**

This 16-bit field indicates Nerase. When Nerase numbers of AUs are erased, the timeout value is specified by ERASE\_TIMEOUT (Refer to *ERASE\_TIMEOUT*). The host should determine the proper number of AUs to be erased in one operation so that the host can show the progress of the erase operation. If this field is set to 0, the erase timeout calculation is not supported.

Table 121. Erase size field

ERASE_SIZE	Value definition
0000h	Erase timeout calculation is not supported.
0001h	1 AU
0002h	2 AU
0003h	3 AU
FFFFh	65535 AU

### **ERASE\_TIMEOUT**

This 6-bit field indicates Terase and the value indicates the erase timeout from offset when multiple AUs are being erased as specified by ERASE\_SIZE. The range of ERASE\_TIMEOUT can be defined as up to 63 seconds and the card manufacturer can choose any combination of ERASE\_SIZE and ERASE\_TIMEOUT depending on the implementation. Determining ERASE\_TIMEOUT determines the ERASE\_SIZE.

Table 122. Erase timeout field

ERASE_TIMEOUT	Value definition
00	Erase timeout calculation is not supported.
01	1 [sec]
02	2 [sec]
03	3 [sec]

Table 122. Erase timeout field (continued)

ERASE_TIMEOUT	Value definition
63	63 [sec]

### **ERASE OFFSET**

This 2-bit field indicates Toffset and one of four values can be selected. This field is meaningless if the ERASE\_SIZE and ERASE\_TIMEOUT fields are set to 0.

Table 123. Erase offset field

ERASE_OFFSET	Value definition
0h	0 [sec]
1h	1 [sec]
2h	2 [sec]
3h	3 [sec]

#### 26.4.13 SD I/O mode

### SD I/O interrupts

To allow the SD I/O card to interrupt the MultiMediaCard/SD module, an interrupt function is available on a pin on the SD interface. Pin 8, used as SDIO\_D1 when operating in the 4-bit SD mode, signals the cards interrupt to the MultiMediaCard/SD module. The use of the interrupt is optional for each card or function within a card. The SD I/O interrupt is level-sensitive, which means that the interrupt line must be held active (low) until it is either recognized and acted upon by the MultiMediaCard/SD module or deasserted due to the end of the interrupt period. After the MultiMediaCard/SD module has serviced the interrupt, the interrupt status bit is cleared via an I/O write to the appropriate bit in the SD I/O card's internal registers. The interrupt output of all SD I/O cards is active low and the application must provide external pull-up resistors on all data lines (SDIO\_D[3:0]). The MultiMediaCard/SD module samples the level of pin 8 (SDIO\_D/IRQ) into the interrupt detector only during the interrupt period. At all other times, the MultiMediaCard/SD module ignores this value.

The interrupt period is applicable for both memory and I/O operations. The definition of the interrupt period for operations with single blocks is different from the definition for multiple-block data transfers.

## SD I/O suspend and resume

Within a multifunction SD I/O or a card with both I/O and memory functions, there are multiple devices (I/O and memory) that share access to the MMC/SD bus. To share access to the MMC/SD module among multiple devices, SD I/O and combo cards optionally implement the concept of suspend/resume. When a card supports suspend/resume, the MMC/SD module can temporarily halt a data transfer operation to one function or memory (suspend) to free the bus for a higher-priority transfer to a different function or memory. After this higher-priority transfer is complete, the original transfer is resumed (restarted) where it left off. Support of suspend/resume is optional on a per-card basis. To perform the

suspend/resume operation on the MMC/SD bus, the MMC/SD module performs the following steps:

- 1. Determines the function currently using the SDIO\_D [3:0] line(s)
- 2. Requests the lower-priority or slower transaction to suspend
- 3. Waits for the transaction suspension to complete
- 4. Begins the higher-priority transaction
- 5. Waits for the completion of the higher priority transaction
- 6. Restores the suspended transaction

#### SD I/O ReadWait

The optional ReadWait (RW) operation is defined only for the SD 1-bit and 4-bit modes. The ReadWait operation allows the MMC/SD module to signal a card that it is reading multiple registers (IO\_RW\_EXTENDED, CMD53) to temporarily stall the data transfer while allowing the MMC/SD module to send commands to any function within the SD I/O device. To determine when a card supports the ReadWait protocol, the MMC/SD module must test capability bits in the internal card registers. The timing for ReadWait is based on the interrupt period.

## 26.4.14 Commands and responses

### Application-specific and general commands

The SD card host module system is designed to provide a standard interface for a variety of applications types. In this environment, there is a need for specific customer/application features. To implement these features, two types of generic commands are defined in the standard: application-specific commands (ACMD) and general commands (GEN\_CMD).

When the card receives the APP\_CMD (CMD55) command, the card expects the next command to be an application-specific command. ACMDs have the same structure as regular MultiMediaCard commands and can have the same CMD number. The card recognizes it as ACMD because it appears after APP\_CMD (CMD55). When the command immediately following the APP\_CMD (CMD55) is not a defined application-specific command, the standard command is used. For example, when the card has a definition for SD\_STATUS (ACMD13), and receives CMD13 immediately following APP\_CMD (CMD55), this is interpreted as SD\_STATUS (ACMD13). However, when the card receives CMD7 immediately following APP\_CMD (CMD55) and the card does not have a definition for ACMD7, this is interpreted as the standard (SELECT/DESELECT\_CARD) CMD7.

To use one of the manufacturer-specific ACMDs the SD card Host must perform the following steps:

- Send APP\_CMD (CMD55)
   The card responds to the MultiMediaCard/SD module, indicating that the APP\_CMD bit is set and an ACMD is now expected.
- Send the required ACMD
   The card responds to the MultiMediaCard/SD module, indicating that the APP\_CMD bit is set and that the accepted command is interpreted as an ACMD. When a nonACMD is sent, it is handled by the card as a normal MultiMediaCard command and the APP\_CMD bit in the card status register stays clear.

When an invalid command is sent (neither ACMD nor CMD) it is handled as a standard MultiMediaCard illegal command error.

The bus transaction for a GEN\_CMD is the same as the single-block read or write commands (WRITE\_BLOCK, CMD24 or READ\_SINGLE\_BLOCK, CMD17). In this case, the argument denotes the direction of the data transfer rather than the address, and the data block has vendor-specific format and meaning.

The card must be selected (in transfer state) before sending GEN\_CMD (CMD56). The data block size is defined by SET\_BLOCKLEN (CMD16). The response to GEN\_CMD (CMD56) is in R1b format.

### **Command types**

Both application-specific and general commands are divided into the four following types:

- broadcast command (BC): sent to all cards; no responses returned.
- broadcast command with response (BCR): sent to all cards; responses received from all cards simultaneously.
- addressed (point-to-point) command (AC): sent to the card that is selected; does
  not include a data transfer on the SDIO\_D line(s).
- addressed (point-to-point) data transfer command (ADTC): sent to the card that is selected; includes a data transfer on the SDIO\_D line(s).

#### **Command formats**

See Table 108 on page 735 for command formats.

#### Commands for the MultiMediaCard/SD module

Table 124. Block-oriented write commands

CMD index	Туре	Argument	Response format	Abbreviation	Description
CMD23	ac	[31:16] set to 0 [15:0] number of blocks	R1	SET_BLOCK_COUNT	Defines the number of blocks which are going to be transferred in the multiple-block read or write command that follows.
CMD24	adtc	[31:0] data address	R1	WRITE_BLOCK	Writes a block of the size selected by the SET_BLOCKLEN command.
CMD25	adtc	[31:0] data address	R1	WRITE_MULTIPLE_BLOCK	Continuously writes blocks of data until a STOP_TRANSMISSION follows or the requested number of blocks has been received.
CMD26	adtc	[31:0] stuff bits	R1	PROGRAM_CID	Programming of the card identification register. This command must be issued only once per card. The card contains hardware to prevent this operation after the first programming. Normally this command is reserved for manufacturer.
CMD27	adtc	[31:0] stuff bits	R1	PROGRAM_CSD	Programming of the programmable bits of the CSD.

Table 125. Block-oriented write protection commands

CMD index	Туре	Argument	Response format	Abbreviation	Description	
CMD28	ac	[31:0] data address	R1b	SET_WRITE_PROT	If the card has write protection features, this command sets the write protection bit of the addressed group. The properties of write protection are coded in the card-specific data (WP_GRP_SIZE).	
CMD29	ac	[31:0] data address	R1b	CLR_WRITE_PROT	If the card provides write protection features, this command clears the write protection bit of the addressed group.	
CMD30	adtc	[31:0] write protect data address	R1	SEND_WRITE_PROT	If the card provides write protection features, this command asks the card to send the status of the write protection bits.	
CMD31	Reserved					

## Table 126. Erase commands

CMD index	Туре	Argument	Response format	Abbreviation	Description
CMD32  CMD34	Reserved. These command indexes cannot be used in order to maintain backward compatibility with older versions of the MultiMediaCard.				
CMD35	ac	[31:0] data address	R1	ERASE_GROUP_START	Sets the address of the first erase group within a range to be selected for erase.
CMD36	ac	[31:0] data address	R1	ERASE_GROUP_END	Sets the address of the last erase group within a continuous range to be selected for erase.
CMD37	Reserved. This command index cannot be used in order to maintain backward compatibility with older versions of the MultiMediaCards				
CMD38	ac	[31:0] stuff bits	R1	ERASE	Erases all previously selected write blocks.

Table 127. I/O mode commands

CMD index	Туре	Argument	Response format	Abbreviation	Description
CMD39	ac	[31:16] RCA [15:15] register write flag [14:8] register address [7:0] register data	R4	FAST_IO	Used to write and read 8-bit (register) data fields. The command addresses a card and a register and provides the data for writing if the write flag is set. The R4 response contains data read from the addressed register. This command accesses application-dependent registers that are not defined in the MultiMediaCard standard.

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Table 127. I/O mode commands (continued)

CMD index	Туре	Argument	Response format	Abbreviation	Description	
CMD40	bcr	[31:0] stuff bits	R5	GO_IRQ_STATE	Places the system in the interrupt mode.	
CMD41	Reserved					

#### Table 128. Lock card

CMD index	Туре	Argument	Response format	Abbreviation	Description
CMD42	adtc	[31:0] stuff bits	R1b	LOCK_UNLOCK	Sets/resets the password or locks/unlocks the card. The size of the data block is set by the SET_BLOCK_LEN command.
CMD43  CMD54	Reserv	ved			

Table 129. Application-specific commands

CMD index	Туре	Argument	Response format	Abbreviation	Description		
CMD55	ac	[31:16] RCA [15:0] stuff bits	R1	APP_CMD	Indicates to the card that the next command bits is an application specific command rather than a standard command		
CMD56	adtc	[31:1] stuff bits [0]: RD/WR			Used either to transfer a data block to the card or to get a data block from the card for general purpose/application-specific commands. The size of the data block shall be set by the SET_BLOCK_LEN command.		
CMD57  CMD59	Reserved.						
CMD60  CMD63	Reser	Reserved for manufacturer.					

# 26.5 Response formats

All responses are sent via the MCCMD command line SDIO\_CMD. The response transmission always starts with the left bit of the bit string corresponding to the response code word. The code length depends on the response type.

A response always starts with a start bit (always 0), followed by the bit indicating the direction of transmission (card = 0). A value denoted by x in the tables below indicates a variable entry. All responses, except for the R3 response type, are protected by a CRC. Every command code word is terminated by the end bit (always 1).

There are five types of responses. Their formats are defined as follows:

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## 26.5.1 R1 (normal response command)

Code length = 48 bits. The 45:40 bits indicate the index of the command to be responded to, this value being interpreted as a binary-coded number (between 0 and 63). The status of the card is coded in 32 bits.

Table 130. R1 response

Bit position	Width (bits	Value	Description
47	1	0	Start bit
46	1	0	Transmission bit
[45:40]	6	Х	Command index
[39:8]	32	Х	Card status
[7:1]	7	Х	CRC7
0	1	1	End bit

### 26.5.2 R1b

It is identical to R1 with an optional busy signal transmitted on the data line. The card may become busy after receiving these commands based on its state prior to the command reception.

## 26.5.3 R2 (CID, CSD register)

Code length = 136 bits. The contents of the CID register are sent as a response to the CMD2 and CMD10 commands. The contents of the CSD register are sent as a response to CMD9. Only the bits [127...1] of the CID and CSD are transferred, the reserved bit [0] of these registers is replaced by the end bit of the response. The card indicates that an erase is in progress by holding MCDAT low. The actual erase time may be quite long, and the host may issue CMD7 to deselect the card.

Table 131. R2 response

Bit position	Width (bits	Value	Description
135	1	0	Start bit
134	1	0	Transmission bit
[133:128]	6	'111111'	Command index
[127:1]	127	Х	Card status
0	1	1	End bit

# **26.5.4** R3 (OCR register)

Code length: 48 bits. The contents of the OCR register are sent as a response to CMD1. The level coding is as follows: restricted voltage windows = low, card busy = low.

Table 132. R3 response

Bit position	Width (bits	Value	Description
47	1	0	Start bit
46	1	0	Transmission bit
[45:40]	6	'111111'	Reserved
[39:8]	32	Х	OCR register
[7:1]	7	'1111111'	Reserved
0	1	1	End bit

## 26.5.5 R4 (Fast I/O)

Code length: 48 bits. The argument field contains the RCA of the addressed card, the register address to be read out or written to, and its content.

Table 133. R4 response

Bit position	l	Width (bits	Value	Description
47		1	0	Start bit
46		1	0	Transmission bit
[45:40]		6	'100111'	CMD39
	[31:16]	16	Х	RCA
[39:8] Argument field	[15:8]	8	Х	register address
	[7:0]	8	Х	read register contents
[7:1]		7	Х	CRC7
0		1	1	End bit

### 26.5.6 R4b

For SD I/O only: an SDIO card receiving the CMD5 will respond with a unique SDIO response R4. The format is:

Table 134. R4b response

Bit position	Width (bits	Value	Description
47	1	0	Start bit
46	1	0	Transmission bit
[45:40]	6	х	Reserved

Bit position		Width (bits	Value	Description
	39	16	Х	Card is ready
	[38:36]	3	X	Number of I/O functions
[39:8] Argument field	35	1	Х	Present memory
	[34:32]	3	Х	Stuff bits
	[31:8]	24	Х	I/O ORC
[7:1]		7	Х	Reserved
0		1	1	End bit

Table 134. R4b response (continued)

Once an SD I/O card has received a CMD5, the I/O portion of that card is enabled to respond normally to all further commands. This I/O enable of the function within the I/O card will remain set until a reset, power cycle or CMD52 with write to I/O reset is received by the card. Note that an SD memory-only card may respond to a CMD5. The proper response for a memory-only card would be *Present memory* = 1 and *Number of I/O functions* = 0. A memory-only card built to meet the SD Memory Card specification version 1.0 would detect the CMD5 as an illegal command and not respond. The I/O aware host will send CMD5. If the card responds with response R4, the host determines the card's configuration based on the data contained within the R4 response.

## 26.5.7 R5 (interrupt request)

Only for MultiMediaCard. Code length: 48 bits. If the response is generated by the host, the RCA field in the argument will be 0x0.

Table 135. R5 response

Bit position		Width (bits	Value	Description
47		1	0	Start bit
46		1	0	Transmission bit
[45:40]		6	'101000'	CMD40
[39:8] Argument field	[31:16]	16	х	RCA [31:16] of winning card or of the host
[59.6] Argument neid	[15:0]	16	x	Not defined. May be used for IRQ data
[7:1]		7	Х	CRC7
0		1	1	End bit

#### 26.5.8 R6

Only for SD I/O. The normal response to CMD3 by a memory device. It is shown in *Table 136*.

Table 136. R6 response

Bit positio	n	Width (bits)	Value	Description
47		1	0	Start bit
46		1	0	Transmission bit
[45:40]		6	'101000'	CMD40
[39:8] Argument	[31:16]	16	Х	RCA [31:16] of winning card or of the host
field	[15:0]	16	Х	Not defined. May be used for IRQ data
[7:1]		7	Х	CRC7
0		1	1	End bit

The card [23:8] status bits are changed when CMD3 is sent to an I/O-only card. In this case, the 16 bits of response are the SD I/O-only values:

- Bit [15] COM\_CRC\_ERROR
- Bit [14] ILLEGAL\_COMMAND
- Bit [13] ERROR
- Bits [12:0] Reserved

## 26.6 SDIO I/O card-specific operations

The following features are SD I/O-specific operations:

- SDIO read wait operation by SDIO D2 signalling
- SDIO read wait operation by stopping the clock
- SDIO suspend/resume operation (write and read suspend)
- SDIO interrupts

The SDIO supports these operations only if the SDIO\_DCTRL[11] bit is set, except for read suspend that does not need specific hardware implementation.

### 26.6.1 SDIO I/O read wait operation by SDIO\_D2 signalling

It is possible to start the readwait interval before the first block is received: when the data path is enabled (SDIO\_DCTRL[0] bit set), the SDIO-specific operation is enabled (SDIO\_DCTRL[11] bit set), read wait starts (SDIO\_DCTRL[10] =0 and SDI\_DCTRL[8] =1) and data direction is from card to SDIO (SDIO\_DCTRL[1] = 1), the DPSM directly moves from Idle to Readwait. In Readwait the DPSM drives SDIO\_D2 to 0 after 2 SDIO\_CK clock cycles. In this state, when you set the RWSTOP bit (SDIO\_DCTRL[9]), the DPSM remains in Wait for two more SDIO\_CK clock cycles to drive SDIO\_D2 to 1 for one clock cycle (in accordance with SDIO specification). The DPSM then starts waiting again until it receives data from the card. The DPSM will not start a readwait interval while receiving a block even if read wait start is set: the readwait interval will start after the CRC is received. The RWSTOP bit has to be cleared to start a new read wait operation. During the readwait interval, the SDIO can detect SDIO interrupts on SDIO\_D1.

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## 26.6.2 SDIO read wait operation by stopping SDIO\_CK

If the SDIO card does not support the previous read wait method, the SDIO can perform a read wait by stopping SDIO\_CK (SDIO\_DCTRL is set just like in the method presented in *Section 26.6.1*, but SDIO\_DCTRL[10] =1): DSPM stops the clock two SDIO\_CK cycles after the end bit of the current received block and starts the clock again after the read wait start bit is set.

As SDIO\_CK is stopped, any command can be issued to the card. During a read/wait interval, the SDIO can detect SDIO interrupts on SDIO\_D1.

## 26.6.3 SDIO suspend/resume operation

While sending data to the card, the SDIO can suspend the write operation. the SDIO\_CMD[11] bit is set and indicates to the CPSM that the current command is a suspend command. The CPSM analyzes the response and when the ACK is received from the card (suspend accepted), it acknowledges the DPSM that goes Idle after receiving the CRC token of the current block.

The hardware does not save the number of the remaining block to be sent to complete the suspended operation (resume).

The write operation can be suspended by software, just by disabling the DPSM (SDIO\_DCTRL[0] =0) when the ACK of the suspend command is received from the card. The DPSM enters then the Idle state.

To suspend a read: the DPSM waits in the Wait\_r state as the function to be suspended sends a complete packet just before stopping the data transaction. The application continues reading RxFIFO until the FIF0 is empty, and the DPSM goes Idle automatically.

## 26.6.4 SDIO interrupts

SDIO interrupts are detected on the SDIO\_D1 line once the SDIO\_DCTRL[11] bit is set.

## 26.7 CE-ATA specific operations

The following features are CE-ATA specific operations:

- sending the command completion signal disable to the CE-ATA device
- receiving the command completion signal from the CE-ATA device
- signaling the completion of the CE-ATA command to the CPU, using the status bit and/or interrupt.

The SDIO supports these operations only for the CE-ATA CMD61 command, that is, if  $SDIO\_CMD[14]$  is set.

## 26.7.1 Command completion signal disable

Command completion signal disable is sent 8 bit cycles after the reception of a **short** response if the 'enable CMD completion' bit, SDIO\_CMD[12], is not set and the 'not interrupt Enable' bit, SDIO\_CMD[13], is set.

The CPSM enters the Pend state, loading the command shift register with the disable sequence "00001" and, the command counter with 43. Eight cycles after, a trigger moves

the CPSM to the Send state. When the command counter reaches 48, the CPSM becomes Idle as no response is awaited.

## 26.7.2 Command completion signal enable

If the 'enable CMD completion' bit SDIO\_CMD[12] is set and the 'not interrupt Enable' bit SDIO\_CMD[13] is set, the CPSM waits for the command completion signal in the Waitcpl state.

When '0' is received on the CMD line, the CPSM enters the Idle state. No new command can be sent for 7 bit cycles. Then, for the last 5 cycles (out of the 7) the CMD line is driven to '1' in push-pull mode.

## 26.7.3 CE-ATA interrupt

The command completion is signaled to the CPU by the status bit SDIO\_STA[23]. This static bit can be cleared with the clear bit SDIO\_ICR[23].

The SDIO\_STA[23] status bit can generate an interrupt on each interrupt line, depending on the mask bit SDIO\_MASKx[23].

## 26.7.4 Aborting CMD61

If the command completion disable signal has not been sent and CMD61 needs to be aborted, the command state machine must be disabled. It then becomes Idle, and the CMD12 command can be sent. No command completion disable signal is sent during the operation.

### 26.8 HW flow control

The HW flow control functionality is used to avoid FIFO underrun (TX mode) and overrun (RX mode) errors.

The behavior is to stop SDIO\_CK and freeze SDIO state machines. The data transfer is stalled while the FIFO is unable to transmit or receive data. Only state machines clocked by SDIOCLK are frozen, the APB2 interface is still alive. The FIFO can thus be filled or emptied even if flow control is activated.

To enable HW flow control, the SDIO\_CLKCR[14] register bit must be set to 1. After reset Flow Control is disabled.

# 26.9 SDIO registers

The device communicates to the system via 32-bit-wide control registers accessible via APB2.

The peripheral registers have to be accessed by words (32 bits).

## 26.9.1 SDIO power control register (SDIO\_POWER)

Address offset: 0x00

Reset value: 0x0000 0000

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved

Reserved

Reserved

Bits 31:2 Reserved, must be kept at reset value

Bits 1:0 PWRCTRL: Power supply control bits.

These bits are used to define the current functional state of the card clock:

00: Power-off: the clock to card is stopped.

01: Reserved

10: Reserved power-up

11: Power-on: the card is clocked.

Note: At least seven HCLK clock periods are needed between two write accesses to this register.

Note: After a data write, data cannot be written to this register for three SDIOCLK (48 MHz) clock

periods plus two PCLK2 clock periods.

## 26.9.2 SDI clock control register (SDIO\_CLKCR)

Address offset: 0x04

Reset value: 0x0000 0000

The SDIO\_CLKCR register controls the SDIO\_CK output clock.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0 Ш, NEGEDGE YPASS **PWRSAV** CLKEN WID HWFC **CLKDIV** BUS Reserved 6 rw rw rw rw rw rw rw rw rw rw

#### Bits 31:15 Reserved, must be kept at reset value

#### Bit 14 HWFC\_EN: HW Flow Control enable

0b: HW Flow Control is disabled

1b: HW Flow Control is enabled

When HW Flow Control is enabled, the meaning of the TXFIFOE and RXFIFOF interrupt signals, please see SDIO Status register definition in Section 26.9.11.

#### Bit 13 NEGEDGE:SDIO\_CK dephasing selection bit

0b: SDIO\_CK generated on the rising edge of the master clock SDIOCLK 1b: SDIO\_CK generated on the falling edge of the master clock SDIOCLK

#### Bits 12:11 WIDBUS: Wide bus mode enable bit

00: Default bus mode: SDIO D0 used 01: 4-wide bus mode: SDIO\_D[3:0] used 10: 8-wide bus mode: SDIO\_D[7:0] used

### Bit 10 BYPASS: Clock divider bypass enable bit

0: Disable bypass: SDIOCLK is divided according to the CLKDIV value before driving the SDIO\_CK output signal.

1: Enable bypass: SDIOCLK directly drives the SDIO\_CK output signal.

### Bit 9 PWRSAV: Power saving configuration bit

For power saving, the SDIO\_CK clock output can be disabled when the bus is idle by setting PWRSAV:

0: SDIO\_CK clock is always enabled

1: SDIO\_CK is only enabled when the bus is active

#### Bit 8 CLKEN: Clock enable bit

0: SDIO CK is disabled 1: SDIO\_CK is enabled

## Bits 7:0 CLKDIV: Clock divide factor

This field defines the divide factor between the input clock (SDIOCLK) and the output clock (SDIO\_CK): SDIO\_CK frequency = SDIOCLK / [CLKDIV + 2].

#### Note:

While the SD/SDIO card or MultiMediaCard is in identification mode, the SDIO\_CK frequency must be less than 400 kHz.

The clock frequency can be changed to the maximum card bus frequency when relative card addresses are assigned to all cards.

After a data write, data cannot be written to this register for three SDIOCLK (48 MHz) clock periods plus two PCLK2 clock periods. SDIO CK can also be stopped during the read wait interval for SD I/O cards: in this case the SDIO\_CLKCR register does not control SDIO\_CK.

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# 26.9.3 SDIO argument register (SDIO\_ARG)

Address offset: 0x08

Reset value: 0x0000 0000

The SDIO\_ARG register contains a 32-bit command argument, which is sent to a card as

part of a command message.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

															CMD	ARG	ì														
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw															

## Bits 31:0 CMDARG: Command argument

Command argument sent to a card as part of a command message. If a command contains an argument, it must be loaded into this register before writing a command to the command register.

## 26.9.4 SDIO command register (SDIO\_CMD)

Address offset: 0x0C

Reset value: 0x0000 0000

The SDIO\_CMD register contains the command index and command type bits. The command index is sent to a card as part of a command message. The command type bits control the command path state machine (CPSM).

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	CE-ATACMD	NIEN	ENCMDcompl	SDIOSuspend	CPSMEN	WAITPEND	WAITINT	OS DE LA	Ü			CMDINDEX	CINIDING		
	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:15 Reserved, must be kept at reset value

Bit 14 ATACMD: CE-ATA command

If ATACMD is set, the CPSM transfers CMD61.

Bit 13 nIEN: not Interrupt Enable

if this bit is 0, interrupts in the CE-ATA device are enabled.

Bit 12 ENCMDcompl: Enable CMD completion

If this bit is set, the command completion signal is enabled.

Bit 11 SDIOSuspend: SD I/O suspend command

If this bit is set, the command to be sent is a suspend command (to be used only with SDIO card).

Bit 10 CPSMEN: Command path state machine (CPSM) Enable bit

If this bit is set, the CPSM is enabled.

Bit 9 WAITPEND: CPSM Waits for ends of data transfer (CmdPend internal signal).

If this bit is set, the CPSM waits for the end of data transfer before it starts sending a command.

Bit 8 WAITINT: CPSM waits for interrupt request

If this bit is set, the CPSM disables command timeout and waits for an interrupt request.

Bits 7:6 WAITRESP: Wait for response bits

They are used to configure whether the CPSM is to wait for a response, and if yes, which kind of response.

00: No response, expect CMDSENT flag

01: Short response, expect CMDREND or CCRCFAIL flag

10: No response, expect CMDSENT flag

11: Long response, expect CMDREND or CCRCFAIL flag

Bit 5:0 CMDINDEX: Command index

The command index is sent to the card as part of a command message.

Note: After a data write, data cannot be written to this register for three SDIOCLK (48 MHz) clock periods plus two PCLK2 clock periods.

MultiMediaCards can send two kinds of response: short responses, 48 bits long, or long responses, 136 bits long. SD card and SD I/O card can send only short responses, the

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argument can vary according to the type of response: the software will distinguish the type of response according to the sent command. CE-ATA devices send only short responses.

## 26.9.5 SDIO command response register (SDIO\_RESPCMD)

Address offset: 0x10

Reset value: 0x0000 0000

The SDIO\_RESPCMD register contains the command index field of the last command response received. If the command response transmission does not contain the command index field (long or OCR response), the RESPCMD field is unknown, although it must contain 111111b (the value of the reserved field from the response).

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved		F	RESP	CMI	)	
neserveu	r	r	r	r	r	r

Bits 31:6 Reserved, must be kept at reset value

Bits 5:0 RESPCMD: Response command index

Read-only bit field. Contains the command index of the last command response received.

## 26.9.6 SDIO response 1..4 register (SDIO RESPx)

Address offset:  $(0x10 + (4 \times x))$ ; x = 1..4

Reset value: 0x0000 0000

The SDIO\_RESP1/2/3/4 registers contain the status of a card, which is part of the received response.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
														CA	RDS	TAT	JSx														
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:0 CARDSTATUSx: see Table 137.

The Card Status size is 32 or 127 bits, depending on the response type.

Table 137. Response type and SDIO\_RESPx registers

Register	Short response	Long response
SDIO_RESP1	Card Status[31:0]	Card Status [127:96]
SDIO_RESP2	Unused	Card Status [95:64]
SDIO_RESP3	Unused	Card Status [63:32]
SDIO_RESP4	Unused	Card Status [31:1]0b

The most significant bit of the card status is received first. The SDIO\_RESP3 register LSB is always 0b.

## 26.9.7 SDIO data timer register (SDIO\_DTIMER)

Address offset: 0x24

Reset value: 0x0000 0000

The SDIO\_DTIMER register contains the data timeout period, in card bus clock periods.

A counter loads the value from the SDIO\_DTIMER register, and starts decrementing when the data path state machine (DPSM) enters the Wait\_R or Busy state. If the timer reaches 0 while the DPSM is in either of these states, the timeout status flag is set.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

														I	DATA	TIME	=														
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw															

Bits 31:0 DATATIME: Data timeout period

Data timeout period expressed in card bus clock periods.

Note: A data transfer must be written to the data timer register and the data length register before

being written to the data control register.

## 26.9.8 SDIO data length register (SDIO\_DLEN)

Address offset: 0x28

Reset value: 0x0000 0000

The SDIO\_DLEN register contains the number of data bytes to be transferred. The value is loaded into the data counter when data transfer starts.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

DATALENGTH

Boonwood												DATA	LEN	IGTH											
Reserved	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw											

Bits 31:25 Reserved, must be kept at reset value

Bits 24:0 **DATALENGTH:** Data length value

Number of data bytes to be transferred.

Note:

For a block data transfer, the value in the data length register must be a multiple of the block size (see SDIO\_DCTRL). A data transfer must be written to the data timer register and the data length register before being written to the data control register.

For an SDIO multibyte transfer the value in the data length register must be between 1 and 512.

## 26.9.9 SDIO data control register (SDIO\_DCTRL)

Address offset: 0x2C

Reset value: 0x0000 0000

The SDIO\_DCTRL register control the data path state machine (DPSM).

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved

Bits 31:12 Reserved, must be kept at reset value

Bit 11 SDIOEN: SD I/O enable functions

If this bit is set, the DPSM performs an SD I/O-card-specific operation.

Bit 10 RWMOD: Read wait mode

0: Read Wait control stopping SDIO\_D21: Read Wait control using SDIO\_CK

Bit 9 RWSTOP: Read wait stop

0: Read wait in progress if RWSTART bit is set 1: Enable for read wait stop if RWSTART bit is set

Bit 8 RWSTART: Read wait start

If this bit is set, read wait operation starts.

Bits 7:4 DBLOCKSIZE: Data block size

Define the data block length when the block data transfer mode is selected:

0000: (0 decimal) lock length =  $2^0 = 1$  byte

0001: (1 decimal) lock length =  $2^1$  = 2 bytes

0010: (2 decimal) lock length =  $2^2 = 4$  bytes

0011: (3 decimal) lock length =  $2^3$  = 8 bytes

0100: (4 decimal) lock length =  $2^4$  = 16 bytes

0101: (5 decimal) lock length =  $2^5$  = 32 bytes

0110: (6 decimal) lock length =  $2^6$  = 64 bytes

0111: (7 decimal) lock length =  $2^7$  = 128 bytes

1000: (8 decimal) lock length =  $2^8$  = 256 bytes

1001: (9 decimal) lock length =  $2^9$  = 512 bytes

1010: (10 decimal) lock length =  $2^{10}$  = 1024 bytes

1011: (11 decimal) lock length =  $2^{11}$  = 2048 bytes

1100: (12 decimal) lock length =  $2^{12}$  = 4096 bytes

1101: (13 decimal) lock length =  $2^{13}$  = 8192 bytes

1110: (14 decimal) lock length =  $2^{14}$  = 16384 bytes

1111: (15 decimal) reserved

Bit 3 DMAEN: DMA enable bit

0: DMA disabled.

1: DMA enabled.

Bit 2 DTMODE: Data transfer mode selection 1: Stream or SDIO multibyte data transfer.

0: Block data transfer

1: Stream or SDIO multibyte data transfer

Bit 1 DTDIR: Data transfer direction selection

0: From controller to card.

1: From card to controller.

Bit 0 DTEN: Data transfer enabled bit

Data transfer starts if 1b is written to the DTEN bit. Depending on the direction bit, DTDIR, the DPSM moves to the Wait\_S, Wait\_R state or Readwait if RW Start is set immediately at the beginning of the transfer. It is not necessary to clear the enable bit after the end of a data transfer but the SDIO\_DCTRL must be updated to enable a new data transfer

Note:

After a data write, data cannot be written to this register for three SDIOCLK (48 MHz) clock periods plus two PCLK2 clock periods.

The meaning of the DTMODE bit changes according to the value of the SDIOEN bit. When SDIOEN=0 and DTMODE=1, the MultiMediaCard stream mode is enabled, and when SDIOEN=1 and DTMODE=1, the peripheral enables an SDIO multibyte transfer.

## 26.9.10 SDIO data counter register (SDIO\_DCOUNT)

Address offset: 0x30

Reset value: 0x0000 0000

The SDIO\_DCOUNT register loads the value from the data length register (see SDIO\_DLEN) when the DPSM moves from the Idle state to the Wait\_R or Wait\_S state. As data is transferred, the counter decrements the value until it reaches 0. The DPSM then moves to the Idle state and the data status end flag, DATAEND, is set.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved												DAT	ACO	UNT											
neserveu	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:25 Reserved, must be kept at reset value

Bits 24:0 DATACOUNT: Data count value

When this bit is read, the number of remaining data bytes to be transferred is returned. Write has no effect.

Note: This register should be read only when the data transfer is complete.

## 26.9.11 SDIO status register (SDIO\_STA)

Address offset: 0x34

Reset value: 0x0000 0000

The SDIO\_STA register is a read-only register. It contains two types of flag:

- Static flags (bits [23:22,10:0]): these bits remain asserted until they are cleared by writing to the SDIO Interrupt Clear register (see SDIO\_ICR)
- Dynamic flags (bits [21:11]): these bits change state depending on the state of the underlying logic (for example, FIFO full and empty flags are asserted and deasserted as data while written to the FIFO)

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			Ċ	חפאפו אפת				CEATAEND	SDIOIT	RXDAVL	TXDAVL	RXFIFOE	TXFIFOE	RXFIFOF	TXFIFOF	RXFIFOHF	TXFIFOHE	RXACT	TXACT	CMDACT	DBCKEND	STBITERR	DATAEND	CMDSENT	CMDREND	RXOVERR	TXUNDERR	DTIMEOUT	СТІМЕОПТ	DCRCFAIL	CCRCFAIL
			Re	es.				r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:24 Reserved, must be kept at reset value

Bit 23 CEATAEND: CE-ATA command completion signal received for CMD61

Bit 22 SDIOIT: SDIO interrupt received

Bit 21 **RXDAVL:** Data available in receive FIFO Bit 20 **TXDAVL:** Data available in transmit FIFO

Bit 19 **RXFIFOE:** Receive FIFO empty Bit 18 **TXFIFOE:** Transmit FIFO empty

When HW Flow Control is enabled, TXFIFOE signals becomes activated when the FIFO contains 2 words.

Bit 17 RXFIFOF: Receive FIFO full

When HW Flow Control is enabled, RXFIFOF signals becomes activated 2 words before the FIFO is full.

Bit 16 TXFIFOF: Transmit FIFO full

Bit 15 RXFIFOHF: Receive FIFO half full: there are at least 8 words in the FIFO

Bit 14 TXFIFOHE: Transmit FIFO half empty: at least 8 words can be written into the FIFO

Bit 13 **RXACT:** Data receive in progress Bit 12 **TXACT:** Data transmit in progress

Bit 11 CMDACT: Command transfer in progress

Bit 10 DBCKEND: Data block sent/received (CRC check passed)

Bit 9 STBITERR: Start bit not detected on all data signals in wide bus mode

Bit 8 DATAEND: Data end (data counter, SDIDCOUNT, is zero)

Bit 7 CMDSENT: Command sent (no response required)

Bit 6 CMDREND: Command response received (CRC check passed)

Bit 5 RXOVERR: Received FIFO overrun error

Bit 4 TXUNDERR: Transmit FIFO underrun error

Bit 3 DTIMEOUT: Data timeout

Bit 2 CTIMEOUT: Command response timeout

The Command TimeOut period has a fixed value of 64 SDIO\_CK clock periods.

Bit 1 DCRCFAIL: Data block sent/received (CRC check failed)

Bit 0 CCRCFAIL: Command response received (CRC check failed)

## 26.9.12 SDIO interrupt clear register (SDIO\_ICR)

Address offset: 0x38

Reset value: 0x0000 0000

The SDIO\_ICR register is a write-only register. Writing a bit with 1b clears the corresponding bit in the SDIO\_STA Status register.

3	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
				Rese	erved	I			CEATAENDC	SDIOITC					Re	eserv	ed					DBCKENDC	STBITERRC	DATAENDC	CMDSENTC	CMDRENDC	RXOVERRC	TXUNDERRC	DTIMEOUTC	СТІМЕОUTC	DCRCFAILC	CCRCFAILC
									rw	rw												rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:24 Reserved, must be kept at reset value

Bit 23 CEATAENDC: CEATAEND flag clear bit

Set by software to clear the CEATAEND flag.

0: CEATAEND not cleared1: CEATAEND cleared

Bit 22 SDIOITC: SDIOIT flag clear bit

Set by software to clear the SDIOIT flag.

0: SDIOIT not cleared1: SDIOIT cleared

Bits 21:11 Reserved, must be kept at reset value

Bit 10 DBCKENDC: DBCKEND flag clear bit

Set by software to clear the DBCKEND flag.

0: DBCKEND not cleared

1: DBCKEND cleared

Bit 9 STBITERRC: STBITERR flag clear bit

Set by software to clear the STBITERR flag.

0: STBITERR not cleared

1: STBITERR cleared

Bit 8 DATAENDC: DATAEND flag clear bit

Set by software to clear the DATAEND flag.

0: DATAEND not cleared

1: DATAEND cleared

### Bit 7 CMDSENTC: CMDSENT flag clear bit

Set by software to clear the CMDSENT flag.

0: CMDSENT not cleared

1: CMDSENT cleared

### Bit 6 CMDRENDC: CMDREND flag clear bit

Set by software to clear the CMDREND flag.

0: CMDREND not cleared

1: CMDREND cleared

### Bit 5 RXOVERRC: RXOVERR flag clear bit

Set by software to clear the RXOVERR flag.

0: RXOVERR not cleared

1: RXOVERR cleared

#### Bit 4 TXUNDERRC: TXUNDERR flag clear bit

Set by software to clear TXUNDERR flag.

0: TXUNDERR not cleared

1: TXUNDERR cleared

#### Bit 3 DTIMEOUTC: DTIMEOUT flag clear bit

Set by software to clear the DTIMEOUT flag.

0: DTIMEOUT not cleared

1: DTIMEOUT cleared

## Bit 2 CTIMEOUTC: CTIMEOUT flag clear bit

Set by software to clear the CTIMEOUT flag.

0: CTIMEOUT not cleared

1: CTIMEOUT cleared

## Bit 1 DCRCFAILC: DCRCFAIL flag clear bit

Set by software to clear the DCRCFAIL flag.

0: DCRCFAIL not cleared

1: DCRCFAIL cleared

### Bit 0 CCRCFAILC: CCRCFAIL flag clear bit

Set by software to clear the CCRCFAIL flag.

0: CCRCFAIL not cleared

1: CCRCFAIL cleared

## 26.9.13 SDIO mask register (SDIO\_MASK)

Address offset: 0x3C

Reset value: 0x0000 0000

The interrupt mask register determines which status flags generate an interrupt request by

setting the corresponding bit to 1b.

31 30 29 28 27 26 25 24 23 22 21 18 17 16 15 14 3 0 20 19 13 12 11 10 9 8 6 5 4 TXUNDERRIE DTIMEOUTIE CEATAENDIE CMDRENDIE RXOVERRIE RXFIFOHFIE TXFIFOHEIE DBCKENDIE STBITERRIE CMDSENTIE RXFIFOEIE DATAENDIE RXDAVLIE TXFIFOEIE RXFIFOFIE **TXFIFOFIE** CMDACTIE TXDAVLIE SDIOITIE **RXACTIE** TXACTIE Reserved

Bits 31:24 Reserved, must be kept at reset value

### Bit 23 **CEATAENDIE:** CE-ATA command completion signal received interrupt enable

Set and cleared by software to enable/disable the interrupt generated when receiving the CE-ATA command completion signal.

- 0: CE-ATA command completion signal received interrupt disabled
- 1: CE-ATA command completion signal received interrupt enabled

#### Bit 22 SDIOITIE: SDIO mode interrupt received interrupt enable

Set and cleared by software to enable/disable the interrupt generated when receiving the SDIO mode interrupt.

- 0: SDIO Mode Interrupt Received interrupt disabled
- 1: SDIO Mode Interrupt Received interrupt enabled

### Bit 21 RXDAVLIE: Data available in Rx FIFO interrupt enable

Set and cleared by software to enable/disable the interrupt generated by the presence of data available in Rx FIFO.

- 0: Data available in Rx FIFO interrupt disabled
- 1: Data available in Rx FIFO interrupt enabled

#### Bit 20 TXDAVLIE: Data available in Tx FIFO interrupt enable

Set and cleared by software to enable/disable the interrupt generated by the presence of data available in Tx FIFO.

- 0: Data available in Tx FIFO interrupt disabled
- 1: Data available in Tx FIFO interrupt enabled

#### Bit 19 RXFIFOEIE: Rx FIFO empty interrupt enable

Set and cleared by software to enable/disable interrupt caused by Rx FIFO empty.

- 0: Rx FIFO empty interrupt disabled
- 1: Rx FIFO empty interrupt enabled

## Bit 18 TXFIFOEIE: Tx FIFO empty interrupt enable

Set and cleared by software to enable/disable interrupt caused by Tx FIFO empty.

- 0: Tx FIFO empty interrupt disabled
- 1: Tx FIFO empty interrupt enabled

#### Bit 17 RXFIFOFIE: Rx FIFO full interrupt enable

Set and cleared by software to enable/disable interrupt caused by Rx FIFO full.

- 0: Rx FIFO full interrupt disabled
- 1: Rx FIFO full interrupt enabled

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#### Bit 16 TXFIFOFIE: Tx FIFO full interrupt enable

Set and cleared by software to enable/disable interrupt caused by Tx FIFO full.

0: Tx FIFO full interrupt disabled

1: Tx FIFO full interrupt enabled

#### Bit 15 RXFIFOHFIE: Rx FIFO half full interrupt enable

Set and cleared by software to enable/disable interrupt caused by Rx FIFO half full.

0: Rx FIFO half full interrupt disabled

1: Rx FIFO half full interrupt enabled

#### Bit 14 **TXFIFOHEIE:** Tx FIFO half empty interrupt enable

Set and cleared by software to enable/disable interrupt caused by Tx FIFO half empty.

0: Tx FIFO half empty interrupt disabled

1: Tx FIFO half empty interrupt enabled

#### Bit 13 RXACTIE: Data receive acting interrupt enable

Set and cleared by software to enable/disable interrupt caused by data being received (data receive acting).

0: Data receive acting interrupt disabled

1: Data receive acting interrupt enabled

#### Bit 12 TXACTIE: Data transmit acting interrupt enable

Set and cleared by software to enable/disable interrupt caused by data being transferred (data transmit acting).

0: Data transmit acting interrupt disabled

1: Data transmit acting interrupt enabled

#### Bit 11 CMDACTIE: Command acting interrupt enable

Set and cleared by software to enable/disable interrupt caused by a command being transferred (command acting).

0: Command acting interrupt disabled

1: Command acting interrupt enabled

#### Bit 10 DBCKENDIE: Data block end interrupt enable

Set and cleared by software to enable/disable interrupt caused by data block end.

0: Data block end interrupt disabled

1: Data block end interrupt enabled

#### Bit 9 STBITERRIE: Start bit error interrupt enable

Set and cleared by software to enable/disable interrupt caused by start bit error.

0: Start bit error interrupt disabled

1: Start bit error interrupt enabled

#### Bit 8 DATAENDIE: Data end interrupt enable

Set and cleared by software to enable/disable interrupt caused by data end.

0: Data end interrupt disabled

1: Data end interrupt enabled

### Bit 7 **CMDSENTIE:** Command sent interrupt enable

Set and cleared by software to enable/disable interrupt caused by sending command.

0: Command sent interrupt disabled

1: Command sent interrupt enabled

#### Bit 6 **CMDRENDIE:** Command response received interrupt enable

Set and cleared by software to enable/disable interrupt caused by receiving command response.

- 0: Command response received interrupt disabled
- 1: command Response Received interrupt enabled

#### Bit 5 RXOVERRIE: Rx FIFO overrun error interrupt enable

Set and cleared by software to enable/disable interrupt caused by Rx FIFO overrun error.

- 0: Rx FIFO overrun error interrupt disabled
- 1: Rx FIFO overrun error interrupt enabled

### Bit 4 TXUNDERRIE: Tx FIFO underrun error interrupt enable

Set and cleared by software to enable/disable interrupt caused by Tx FIFO underrun error.

- 0: Tx FIFO underrun error interrupt disabled
- 1: Tx FIFO underrun error interrupt enabled

### Bit 3 DTIMEOUTIE: Data timeout interrupt enable

Set and cleared by software to enable/disable interrupt caused by data timeout.

- 0: Data timeout interrupt disabled
- 1: Data timeout interrupt enabled

#### Bit 2 CTIMEOUTIE: Command timeout interrupt enable

Set and cleared by software to enable/disable interrupt caused by command timeout.

- 0: Command timeout interrupt disabled
- 1: Command timeout interrupt enabled

## Bit 1 DCRCFAILIE: Data CRC fail interrupt enable

Set and cleared by software to enable/disable interrupt caused by data CRC failure.

- 0: Data CRC fail interrupt disabled
- 1: Data CRC fail interrupt enabled

## Bit 0 CCRCFAILIE: Command CRC fail interrupt enable

Set and cleared by software to enable/disable interrupt caused by command CRC failure.

- 0: Command CRC fail interrupt disabled
- 1: Command CRC fail interrupt enabled

## 26.9.14 SDIO FIFO counter register (SDIO\_FIFOCNT)

Address offset: 0x48

Reset value: 0x0000 0000

The SDIO\_FIFOCNT register contains the remaining number of words to be written to or read from the FIFO. The FIFO counter loads the value from the data length register (see SDIO\_DLEN) when the data transfer enable bit, DTEN, is set in the data control register (SDIO\_DCTRL register) and the DPSM is at the Idle state. If the data length is not word-aligned (multiple of 4), the remaining 1 to 3 bytes are regarded as a word.

 $31 \quad 30 \quad 29 \quad 28 \quad 27 \quad 26 \quad 25 \quad 24 \quad 23 \quad 22 \quad 21 \quad 20 \quad 19 \quad 18 \quad 17 \quad 16 \quad 15 \quad 14 \quad 13 \quad 12 \quad 11 \quad 10 \quad 9 \quad 8 \quad 7 \quad 6 \quad 5 \quad 4 \quad 3 \quad 2 \quad 1 \quad 0$ 

Reserved												FOC												
neserved	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:24 Reserved, must be kept at reset value

Bits 23:0 FIFOCOUNT: Remaining number of words to be written to or read from the FIFO.

## 26.9.15 SDIO data FIFO register (SDIO\_FIFO)

Address offset: 0x80

Reset value: 0x0000 0000

The receive and transmit FIFOs can be read or written as 32-bit wide registers. The FIFOs contain 32 entries on 32 sequential addresses. This allows the CPU to use its load and store multiple operands to read from/write to the FIFO.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

															FIF0	Data															
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw															

bits 31:0 FIFOData: Receive and transmit FIFO data

The FIFO data occupies 32 entries of 32-bit words, from address: SDIO base + 0x080 to SDIO base + 0xFC.

## 26.9.16 SDIO register map

The following table summarizes the SDIO registers.

Table 138. SDIO register map

Offset	Register	31 30 29 27 27 26 25	24 23 22 22 23 20 19 19 17 16	13 13 17 17	1	10	8	7	5	က	1	0
0x00	SDIO_POWER		Reserved								HOOK	PWRCIRL
0x04	SDIO_CLKCR		Reserved		WIDBUS	PWRSAV	CLKEN			CLKUI		
0x08	SDIO_ARG		CMD	ARG								
0x0C	SDIO_CMD			CMDINDEX								
0x10	SDIO_RESPCM D			F	RESPC	MD						
0x14	SDIO_RESP1		CARDS									
0x18	SDIO_RESP2		CARDS									
0x1C	SDIO_RESP3		CARDS									
0x20	SDIO_RESP4		CARDS									
0x24	SDIO_DTIMER		DATA									
0x28	SDIO_DLEN	Reserved		DATALE	NGTH							
0x2C	SDIO_DCTRL		Reserved		SDIOEN	RWSTOP	RWSTART	HZISXJO IBU		DMAEN	DTDIR	DTEN
0x30	SDIO_DCOUNT	Reserved		DATACO	TNUC	•				•	•	
0x34	SDIO_STA	Reserved	CEATAEND SDIOIT RXDAVL TXDAVL RXFIFOE TXFIFOE TXFIFOE TXFIFOF	TXFIFOHE RXACT TXACT	CMDACT	STBITERR	DATAEND	CMDSENT	RXOVERR TXUNDERR	DTIMEOUT	DCRCFAIL	CCRCFAIL

Table 138. SDIO register map (continued)

Offset	Register	31 30 29 27 27 26 25 25	23	21	20		18			14	13	11	10	6	8	7	9	5	4	3	2	-	0
0x38	SDIO_ICR	Reserved	CEATAENDC					Reserved					DBCKENDC	STBITERRC	DATAENDC	CMDSENTC	CMDRENDC	RXOVERRC	TXUNDERRC	DTIMEOUTC	IMEO	RCF,	CCRCFAILC
0x3C	SDIO_MASK	Reserved	CEATAENDIE	RXDAVLIE	TXDAVLIE	RXFIFOEIE	RXFIFOFIE	TXFIFOFIE	RXFIFOHFIE	TXFIFOHEIE	RXACTIE	CMDACTIE	DBCKENDIE	STBITERRIE	DATAENDIE	CMDSENTIE	CMDRENDIE	RXOVERRIE	TXUNDERRIE	DTIMEOUTIE			CCRCFAILIE
0x48	SDIO_FIFOCNT	Reserved									FIFO	COU	NT										
0x80	SDIO_FIFO							FIF	0Dat	а													

Refer to Table 2 on page 50 for the register boundary addresses.

# 27 Controller area network (bxCAN)

This section applies to the whole STM32F4xx family, unless otherwise specified.

## 27.1 bxCAN introduction

The **Basic Extended CAN** peripheral, named **bxCAN**, interfaces the CAN network. It supports the CAN protocols version 2.0A and B. It has been designed to manage a high number of incoming messages efficiently with a minimum CPU load. It also meets the priority requirements for transmit messages.

For safety-critical applications, the CAN controller provides all hardware functions for supporting the CAN Time Triggered Communication option.

## 27.2 bxCAN main features

- Supports CAN protocol version 2.0 A, B Active
- Bit rates up to 1 Mbit/s
- Supports the Time Triggered Communication option

#### **Transmission**

- Three transmit mailboxes
- Configurable transmit priority
- Time Stamp on SOF transmission

## Reception

- Two receive FIFOs with three stages
- Scalable filter banks:
  - 28 filter banks shared between CAN1 and CAN2
- Identifier list feature
- Configurable FIFO overrun
- Time Stamp on SOF reception

#### Time-triggered communication option

- Disable automatic retransmission mode
- 16-bit free running timer
- Time Stamp sent in last two data bytes

### Management

- Maskable interrupts
- Software-efficient mailbox mapping at a unique address space

#### **Dual CAN**

- CAN1: Master bxCAN for managing the communication between a Slave bxCAN and the 512-byte SRAM memory
- CAN2: Slave bxCAN, with no direct access to the SRAM memory.
- The two bxCAN cells share the 512-byte SRAM memory (see Figure 300: Dual CAN block diagram)

#### **Dual CAN**

- CAN1: Master bxCAN for managing the communication between a Slave bxCAN and the 512-byte SRAM memory
- CAN2: Slave bxCAN, with no direct access to the SRAM memory.
- The two bxCAN cells share the 512-byte SRAM memory (see Figure 300 on page 787)

Note:

In low, medium-, high- and XL-density devices the USB and CAN share a dedicated 512-byte SRAM memory for data transmission and reception, and so they cannot be used concurrently (the shared SRAM is accessed through CAN and USB exclusively). The USB and CAN can be used in the same application but not at the same time.

## 27.3 bxCAN general description

In today's CAN applications, the number of nodes in a network is increasing and often several networks are linked together via gateways. Typically the number of messages in the system (and thus to be handled by each node) has significantly increased. In addition to the application messages, Network Management and Diagnostic messages have been introduced.

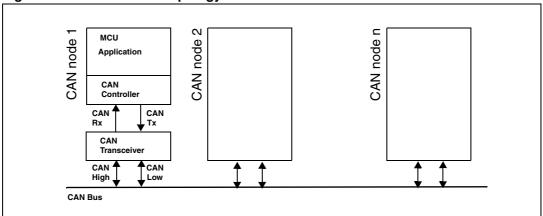
An enhanced filtering mechanism is required to handle each type of message.

Furthermore, application tasks require more CPU time, therefore real-time constraints caused by message reception have to be reduced.

 A receive FIFO scheme allows the CPU to be dedicated to application tasks for a long time period without losing messages.

The standard HLP (Higher Layer Protocol) based on standard CAN drivers requires an efficient interface to the CAN controller.

Figure 299. CAN network topology



#### 27.3.1 CAN 2.0B active core

The bxCAN module handles the transmission and the reception of CAN messages fully autonomously. Standard identifiers (11-bit) and extended identifiers (29-bit) are fully supported by hardware.

### 27.3.2 Control, status and configuration registers

The application uses these registers to:

- Configure CAN parameters, e.g. baud rate
- Request transmissions
- Handle receptions
- Manage interrupts
- Get diagnostic information

#### 27.3.3 Tx mailboxes

Three transmit mailboxes are provided to the software for setting up messages. The transmission Scheduler decides which mailbox has to be transmitted first.

## 27.3.4 Acceptance filters

The bxCAN provides 28 scalable/configurable identifier filter banks for selecting the incoming messages the software needs and discarding the others.

The bxCAN provides 28 scalable/configurable identifier filter banks for selecting the incoming messages the software needs and discarding the others. In other devices there are 14 scalable/configurable identifier filter banks.

#### **Receive FIFO**

Two receive FIFOs are used by hardware to store the incoming messages. Three complete messages can be stored in each FIFO. The FIFOs are managed completely by hardware.

CAM (Master) with 512 bytes SRAM Master Receive FIFO 0 Master Control Master Status Tx Status Rx FIFO 0 Status Scheduler Rx FIFO 1 Status Control/Status/Configuration Acceptance Filters Interrupt Enable CAN 2.0B Active Core 26 Error Status Memory Filter Access Bit Timing Controller Master Filters (0 to 27) Filter Master Filter Mode Transmission Filter Scale Scheduler Slave Receive FIFO 1 Receive FFO 0 Save Tx Mailboxes Filter FIFOAssign Filter Activation Mailbox 0 Mailbox 0 Mailbox 0 CAN2 (Save) Master Control Control/Status/Configuration Master Status Tx Status Rx FIFO 0 Status CAN 2.0B Active Core Rx FIFO 1 Status Interrupt Enable Error Status Note: CAN 2 start filter bank number n is conf gurable by writing to Bit Timing the CAN2SB[5:0] bits in the CAN\_FMRregister. ai16094b

Figure 300. Dual CAN block diagram

#### 27.4 bxCAN operating modes

bxCAN has three main operating modes: initialization, normal and Sleep. After a hardware reset, bxCAN is in Sleep mode to reduce power consumption and an internal pullup is active on CANTX. The software requests bxCAN to enter initialization or Sleep mode by setting the INRQ or SLEEP bits in the CAN\_MCR register. Once the mode has been entered, bxCAN confirms it by setting the INAK or SLAK bits in the CAN\_MSR register and the internal pull-up is disabled. When neither INAK nor SLAK are set, bxCAN is in normal

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mode. Before entering **normal** mode bxCAN always has to **synchronize** on the CAN bus. To synchronize, bxCAN waits until the CAN bus is idle, this means 11 consecutive recessive bits have been monitored on CANRX.

### 27.4.1 Initialization mode

The software initialization can be done while the hardware is in Initialization mode. To enter this mode the software sets the INRQ bit in the CAN\_MCR register and waits until the hardware has confirmed the request by setting the INAK bit in the CAN\_MSR register.

To leave Initialization mode, the software clears the INQR bit. bxCAN has left Initialization mode once the INAK bit has been cleared by hardware.

While in Initialization Mode, all message transfers to and from the CAN bus are stopped and the status of the CAN bus output CANTX is recessive (high).

Entering Initialization Mode does not change any of the configuration registers.

To initialize the CAN Controller, software has to set up the Bit Timing (CAN\_BTR) and CAN options (CAN\_MCR) registers.

To initialize the registers associated with the CAN filter banks (mode, scale, FIFO assignment, activation and filter values), software has to set the FINIT bit (CAN\_FMR). Filter initialization also can be done outside the initialization mode.

Note: When FINIT=1, CAN reception is deactivated.

The filter values also can be modified by deactivating the associated filter activation bits (in the CAN\_FA1R register).

If a filter bank is not used, it is recommended to leave it non active (leave the corresponding FACT bit cleared).

#### 27.4.2 Normal mode

Once the initialization is complete, the software must request the hardware to enter Normal mode to be able to synchronize on the CAN bus and start reception and transmission.

The request to enter Normal mode is issued by clearing the INRQ bit in the CAN\_MCR register. The bxCAN enters Normal mode and is ready to take part in bus activities when it has synchronized with the data transfer on the CAN bus. This is done by waiting for the occurrence of a sequence of 11 consecutive recessive bits (Bus Idle state). The switch to Normal mode is confirmed by the hardware by clearing the INAK bit in the CAN\_MSR register.

The initialization of the filter values is independent from Initialization Mode but must be done while the filter is not active (corresponding FACTx bit cleared). The filter scale and mode configuration must be configured before entering Normal Mode.

### 27.4.3 Sleep mode (low power)

To reduce power consumption, bxCAN has a low-power mode called Sleep mode. This mode is entered on software request by setting the SLEEP bit in the CAN\_MCR register. In this mode, the bxCAN clock is stopped, however software can still access the bxCAN mailboxes.

If software requests entry to **initialization** mode by setting the INRQ bit while bxCAN is in **Sleep** mode, it must also clear the SLEEP bit.

bxCAN can be woken up (exit Sleep mode) either by software clearing the SLEEP bit or on detection of CAN bus activity.

On CAN bus activity detection, hardware automatically performs the wakeup sequence by clearing the SLEEP bit if the AWUM bit in the CAN MCR register is set. If the AWUM bit is cleared, software has to clear the SLEEP bit when a wakeup interrupt occurs, in order to exit from Sleep mode.

Note:

If the wakeup interrupt is enabled (WKUIE bit set in CAN IER register) a wakeup interrupt will be generated on detection of CAN bus activity, even if the bxCAN automatically performs the wakeup sequence.

After the SLEEP bit has been cleared, Sleep mode is exited once bxCAN has synchronized with the CAN bus, refer to Figure 301: bxCAN operating modes. The Sleep mode is exited once the SLAK bit has been cleared by hardware.

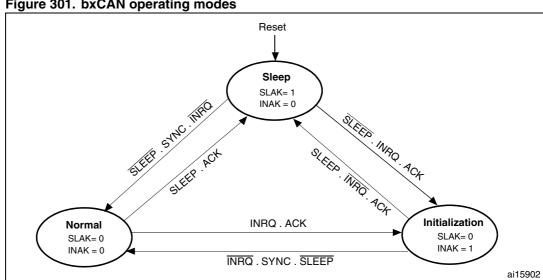


Figure 301. bxCAN operating modes

- ACK = The wait state during which hardware confirms a request by setting the INAK or SLAK bits in the CAN\_MSR register
- 2. SYNC = The state during which bxCAN waits until the CAN bus is idle, meaning 11 consecutive recessive bits have been monitored on CANRX

#### 27.5 Test mode

Test mode can be selected by the SILM and LBKM bits in the CAN\_BTR register. These bits must be configured while bxCAN is in Initialization mode. Once test mode has been selected, the INRQ bit in the CAN MCR register must be reset to enter Normal mode.

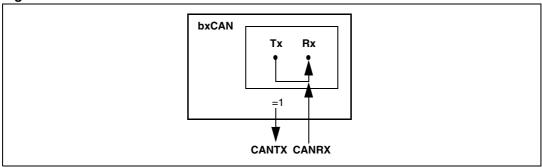
#### 27.5.1 Silent mode

The bxCAN can be put in Silent mode by setting the SILM bit in the CAN\_BTR register.

In Silent mode, the bxCAN is able to receive valid data frames and valid remote frames, but it sends only recessive bits on the CAN bus and it cannot start a transmission. If the bxCAN has to send a dominant bit (ACK bit, overload flag, active error flag), the bit is rerouted internally so that the CAN Core monitors this dominant bit, although the CAN bus may

remain in recessive state. Silent mode can be used to analyze the traffic on a CAN bus without affecting it by the transmission of dominant bits (Acknowledge Bits, Error Frames).

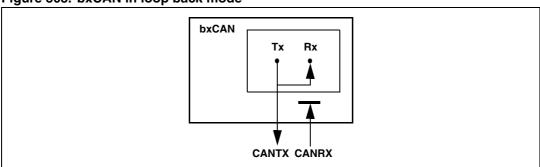
Figure 302. bxCAN in silent mode



## 27.5.2 Loop back mode

The bxCAN can be set in Loop Back Mode by setting the LBKM bit in the CAN\_BTR register. In Loop Back Mode, the bxCAN treats its own transmitted messages as received messages and stores them (if they pass acceptance filtering) in a Receive mailbox.

Figure 303. bxCAN in loop back mode

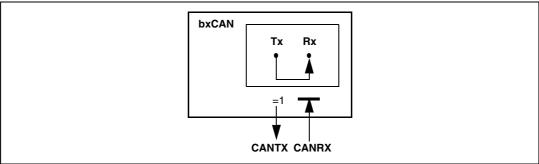


This mode is provided for self-test functions. To be independent of external events, the CAN Core ignores acknowledge errors (no dominant bit sampled in the acknowledge slot of a data / remote frame) in Loop Back Mode. In this mode, the bxCAN performs an internal feedback from its Tx output to its Rx input. The actual value of the CANRX input pin is disregarded by the bxCAN. The transmitted messages can be monitored on the CANTX pin.

## 27.5.3 Loop back combined with silent mode

It is also possible to combine Loop Back mode and Silent mode by setting the LBKM and SILM bits in the CAN\_BTR register. This mode can be used for a "Hot Selftest", meaning the bxCAN can be tested like in Loop Back mode but without affecting a running CAN system connected to the CANTX and CANRX pins. In this mode, the CANRX pin is disconnected from the bxCAN and the CANTX pin is held recessive.

Figure 304. bxCAN in combined mode



## 27.6 Debug mode

When the microcontroller enters the debug mode (Cortex<sup>™</sup>-M4F core halted), the bxCAN continues to work normally or stops, depending on:

- the DBG\_CAN1\_STOP bit for CAN1 or the DBG\_CAN2\_STOP bit for CAN2 in the DBG module. For more details, refer to Section 32.16.2: Debug support for timers, watchdog, bxCAN and I2C.
- the DBF bit in CAN\_MCR. For more details, refer to Section 27.9.2: CAN control and status registers.

## 27.7 bxCAN functional description

### 27.7.1 Transmission handling

In order to transmit a message, the application must select one **empty** transmit mailbox, set up the identifier, the data length code (DLC) and the data before requesting the transmission by setting the corresponding TXRQ bit in the CAN\_TIxR register. Once the mailbox has left **empty** state, the software no longer has write access to the mailbox registers. Immediately after the TXRQ bit has been set, the mailbox enters **pending** state and waits to become the highest priority mailbox, see *Transmit Priority*. As soon as the mailbox has the highest priority it will be **scheduled** for transmission. The transmission of the message of the scheduled mailbox will start (enter **transmit** state) when the CAN bus becomes idle. Once the mailbox has been successfully transmitted, it will become **empty** again. The hardware indicates a successful transmission by setting the RQCP and TXOK bits in the CAN\_TSR register.

If the transmission fails, the cause is indicated by the ALST bit in the CAN\_TSR register in case of an Arbitration Lost, and/or the TERR bit, in case of transmission error detection.

### **Transmit priority**

## By identifier

When more than one transmit mailbox is pending, the transmission order is given by the identifier of the message stored in the mailbox. The message with the lowest identifier value has the highest priority according to the arbitration of the CAN protocol. If the identifier values are equal, the lower mailbox number will be scheduled first.

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#### By transmit request order

The transmit mailboxes can be configured as a transmit FIFO by setting the TXFP bit in the CAN\_MCR register. In this mode the priority order is given by the transmit request order.

This mode is very useful for segmented transmission.

#### **Abort**

A transmission request can be aborted by the user setting the ABRQ bit in the CAN\_TSR register. In **pending** or **scheduled** state, the mailbox is aborted immediately. An abort request while the mailbox is in **transmit** state can have two results. If the mailbox is transmitted successfully the mailbox becomes **empty** with the TXOK bit set in the CAN\_TSR register. If the transmission fails, the mailbox becomes **scheduled**, the transmission is aborted and becomes **empty** with TXOK cleared. In all cases the mailbox will become **empty** again at least at the end of the current transmission.

#### Nonautomatic retransmission mode

This mode has been implemented in order to fulfil the requirement of the Time Triggered Communication option of the CAN standard. To configure the hardware in this mode the NART bit in the CAN\_MCR register must be set.

In this mode, each transmission is started only once. If the first attempt fails, due to an arbitration loss or an error, the hardware will not automatically restart the message transmission.

At the end of the first transmission attempt, the hardware considers the request as completed and sets the RQCP bit in the CAN\_TSR register. The result of the transmission is indicated in the CAN\_TSR register by the TXOK, ALST and TERR bits.

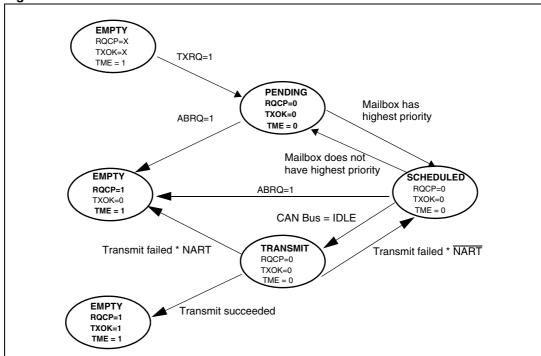


Figure 305. Transmit mailbox states

#### 27.7.2 Time triggered communication mode

In this mode, the internal counter of the CAN hardware is activated and used to generate the Time Stamp value stored in the CAN RDTxR/CAN TDTxR registers, respectively (for Rx and Tx mailboxes). The internal counter is incremented each CAN bit time (refer to Section 27.7.7: Bit timing). The internal counter is captured on the sample point of the Start Of Frame bit in both reception and transmission.

#### 27.7.3 Reception handling

For the reception of CAN messages, three mailboxes organized as a FIFO are provided. In order to save CPU load, simplify the software and guarantee data consistency, the FIFO is managed completely by hardware. The application accesses the messages stored in the FIFO through the FIFO output mailbox.

## Valid message

A received message is considered as valid when it has been received correctly according to the CAN protocol (no error until the last but one bit of the EOF field) and It passed through the identifier filtering successfully, see Section 27.7.4: Identifier filtering.

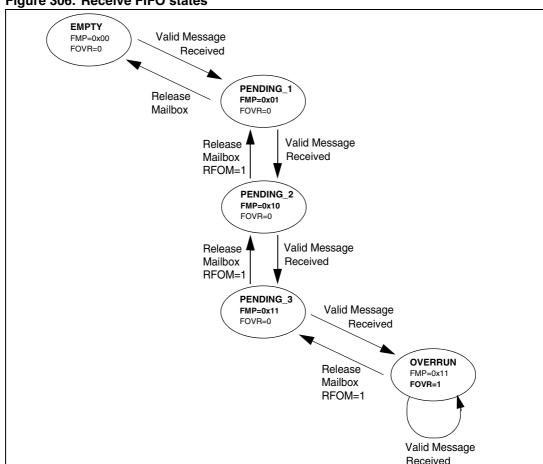


Figure 306. Receive FIFO states

## FIFO management

Starting from the **empty** state, the first valid message received is stored in the FIFO which becomes **pending\_1**. The hardware signals the event setting the FMP[1:0] bits in the CAN\_RFR register to the value 01b. The message is available in the FIFO output mailbox. The software reads out the mailbox content and releases it by setting the RFOM bit in the CAN\_RFR register. The FIFO becomes **empty** again. If a new valid message has been received in the meantime, the FIFO stays in **pending\_1** state and the new message is available in the output mailbox.

If the application does not release the mailbox, the next valid message will be stored in the FIFO which enters **pending\_2** state (FMP[1:0] = 10b). The storage process is repeated for the next valid message putting the FIFO into **pending\_3** state (FMP[1:0] = 11b). At this point, the software must release the output mailbox by setting the RFOM bit, so that a mailbox is free to store the next valid message. Otherwise the next valid message received will cause a loss of message.

Refer also to Section 27.7.5: Message storage

#### Overrun

Once the FIFO is in **pending\_3** state (i.e. the three mailboxes are full) the next valid message reception will lead to an **overrun** and a message will be lost. The hardware signals the overrun condition by setting the FOVR bit in the CAN\_RFR register. Which message is lost depends on the configuration of the FIFO:

- If the FIFO lock function is disabled (RFLM bit in the CAN\_MCR register cleared) the last message stored in the FIFO will be overwritten by the new incoming message. In this case the latest messages will be always available to the application.
- If the FIFO lock function is enabled (RFLM bit in the CAN\_MCR register set) the most recent message will be discarded and the software will have the three oldest messages in the FIFO available.

## **Reception related interrupts**

Once a message has been stored in the FIFO, the FMP[1:0] bits are updated and an interrupt request is generated if the FMPIE bit in the CAN IER register is set.

When the FIFO becomes full (i.e. a third message is stored) the FULL bit in the CAN\_RFR register is set and an interrupt is generated if the FFIE bit in the CAN\_IER register is set.

On overrun condition, the FOVR bit is set and an interrupt is generated if the FOVIE bit in the CAN\_IER register is set.

## 27.7.4 Identifier filtering

In the CAN protocol the identifier of a message is not associated with the address of a node but related to the content of the message. Consequently a transmitter broadcasts its message to all receivers. On message reception a receiver node decides - depending on the identifier value - whether the software needs the message or not. If the message is needed, it is copied into the SRAM. If not, the message must be discarded without intervention by the software.

To fulfill this requirement, the bxCAN Controller provides 28 configurable and scalable filter banks (27-0) to the application. In other devices the bxCAN Controller provides 14 configurable and scalable filter banks (13-0) to the application in order to receive only the messages the software needs. This hardware filtering saves CPU resources which would be

otherwise needed to perform filtering by software. Each filter bank x consists of two 32-bit registers, CAN\_FxR0 and CAN\_FxR1.

#### Scalable width

To optimize and adapt the filters to the application needs, each filter bank can be scaled independently. Depending on the filter scale a filter bank provides:

- One 32-bit filter for the STDID[10:0], EXTID[17:0], IDE and RTR bits.
- Two 16-bit filters for the STDID[10:0], RTR, IDE and EXTID[17:15] bits.

Refer to Figure 307.

Furthermore, the filters can be configured in mask mode or in identifier list mode.

#### Mask mode

In **mask** mode the identifier registers are associated with mask registers specifying which bits of the identifier are handled as "must match" or as "don't care".

#### Identifier list mode

In **identifier list** mode, the mask registers are used as identifier registers. Thus instead of defining an identifier and a mask, two identifiers are specified, doubling the number of single identifiers. All bits of the incoming identifier must match the bits specified in the filter registers.

## Filter bank scale and mode configuration

The filter banks are configured by means of the corresponding CAN\_FMR register. To configure a filter bank it must be deactivated by clearing the FACT bit in the CAN\_FAR register. The filter scale is configured by means of the corresponding FSCx bit in the CAN\_FS1R register, refer to *Figure 307*. The **identifier list** or **identifier mask** mode for the corresponding Mask/Identifier registers is configured by means of the FBMx bits in the CAN\_FMR register.

To filter a group of identifiers, configure the Mask/Identifier registers in mask mode.

To select single identifiers, configure the Mask/Identifier registers in identifier list mode.

Filters not used by the application should be left deactivated.

Each filter within a filter bank is numbered (called the *Filter Number*) from 0 to a maximum dependent on the mode and the scale of each of the filter banks.

Concerning the filter configuration, refer to Figure 307.

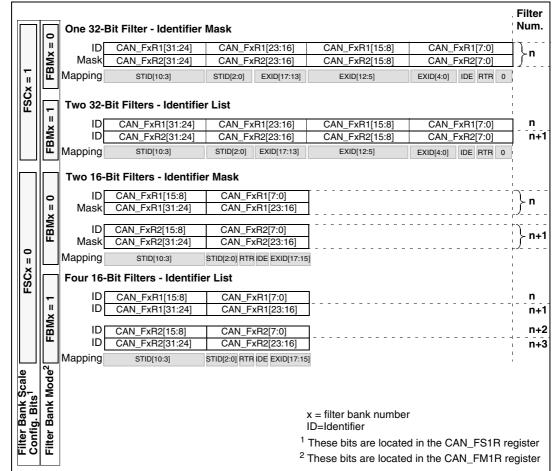


Figure 307. Filter bank scale configuration - register organization

#### Filter match index

Once a message has been received in the FIFO it is available to the application. Typically, application data is copied into SRAM locations. To copy the data to the right location the application has to identify the data by means of the identifier. To avoid this, and to ease the access to the SRAM locations, the CAN controller provides a Filter Match Index.

This index is stored in the mailbox together with the message according to the filter priority rules. Thus each received message has its associated filter match index.

The Filter Match index can be used in two ways:

- Compare the Filter Match index with a list of expected values.
- Use the Filter Match Index as an index on an array to access the data destination location.

For nonmasked filters, the software no longer has to compare the identifier.

If the filter is masked the software reduces the comparison to the masked bits only.

The index value of the filter number does not take into account the activation state of the filter banks. In addition, two independent numbering schemes are used, one for each FIFO. Refer to *Figure 308* for an example.

ID=Identifier



Figure 308. Example of filter numbering

## Filter priority rules

Depending on the filter combination it may occur that an identifier passes successfully through several filters. In this case the filter match value stored in the receive mailbox is chosen according to the following priority rules:

- A 32-bit filter takes priority over a 16-bit filter.
- For filters of equal scale, priority is given to the Identifier List mode over the Identifier Mask mode
- For filters of equal scale and mode, priority is given by the filter number (the lower the number, the higher the priority).

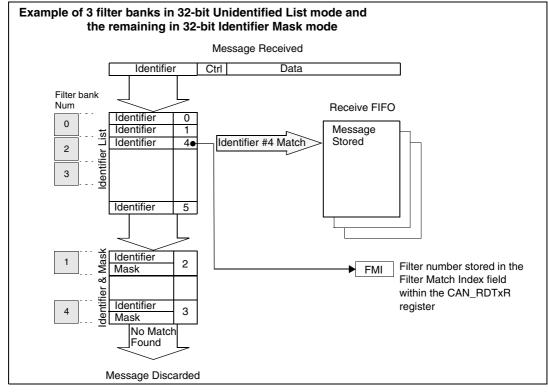


Figure 309. Filtering mechanism - example

The example above shows the filtering principle of the bxCAN. On reception of a message, the identifier is compared first with the filters configured in identifier list mode. If there is a match, the message is stored in the associated FIFO and the index of the matching filter is stored in the Filter Match Index. As shown in the example, the identifier matches with Identifier #2 thus the message content and FMI 2 is stored in the FIFO.

If there is no match, the incoming identifier is then compared with the filters configured in mask mode.

If the identifier does not match any of the identifiers configured in the filters, the message is discarded by hardware without disturbing the software.

## 27.7.5 Message storage

The interface between the software and the hardware for the CAN messages is implemented by means of mailboxes. A mailbox contains all information related to a message; identifier, data, control, status and time stamp information.

#### **Transmit mailbox**

The software sets up the message to be transmitted in an empty transmit mailbox. The status of the transmission is indicated by hardware in the CAN\_TSR register.

Table 139. Transmit mailbox mapping

Offset to transmit mailbox base address	Register name
0	CAN_TIxR
4	CAN_TDTxR
8	CAN_TDLxR
12	CAN_TDHxR

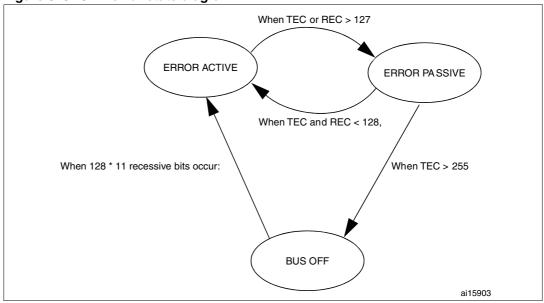
#### Receive mailbox

When a message has been received, it is available to the software in the FIFO output mailbox. Once the software has handled the message (e.g. read it) the software must release the FIFO output mailbox by means of the RFOM bit in the CAN\_RFR register to make the next incoming message available. The filter match index is stored in the MFMI field of the CAN\_RDTxR register. The 16-bit time stamp value is stored in the TIME[15:0] field of CAN\_RDTxR.

Table 140. Receive mailbox mapping

Offset to receive mailbox base address (bytes)	Register name
0	CAN_RixR
4	CAN_RDTxR
8	CAN_RDLxR
12	CAN_RDHxR

Figure 310. CAN error state diagram



## 27.7.6 Error management

The error management as described in the CAN protocol is handled entirely by hardware using a Transmit Error Counter (TEC value, in CAN\_ESR register) and a Receive Error Counter (REC value, in the CAN\_ESR register), which get incremented or decremented according to the error condition. For detailed information about TEC and REC management, please refer to the CAN standard.

Both of them may be read by software to determine the stability of the network. Furthermore, the CAN hardware provides detailed information on the current error status in CAN\_ESR register. By means of the CAN\_IER register (ERRIE bit, etc.), the software can configure the interrupt generation on error detection in a very flexible way.

## **Bus-Off recovery**

The Bus-Off state is reached when TEC is greater than 255, this state is indicated by BOFF bit in CAN\_ESR register. In Bus-Off state, the bxCAN is no longer able to transmit and receive messages.

Depending on the ABOM bit in the CAN\_MCR register bxCAN will recover from Bus-Off (become error active again) either automatically or on software request. But in both cases the bxCAN has to wait at least for the recovery sequence specified in the CAN standard (128 occurrences of 11 consecutive recessive bits monitored on CANRX).

If ABOM is set, the bxCAN will start the recovering sequence automatically after it has entered Bus-Off state.

If ABOM is cleared, the software must initiate the recovering sequence by requesting bxCAN to enter and to leave initialization mode.

Note:

In initialization mode, bxCAN does not monitor the CANRX signal, therefore it cannot complete the recovery sequence. **To recover, bxCAN must be in normal mode**.

## **27.7.7 Bit timing**

The bit timing logic monitors the serial bus-line and performs sampling and adjustment of the sample point by synchronizing on the start-bit edge and resynchronizing on the following edges.

Its operation may be explained simply by splitting nominal bit time into three segments as follows:

- **Synchronization segment (SYNC\_SEG)**: a bit change is expected to occur within this time segment. It has a fixed length of one time quantum (1 x t<sub>CAN</sub>).
- Bit segment 1 (BS1): defines the location of the sample point. It includes the PROP\_SEG and PHASE\_SEG1 of the CAN standard. Its duration is programmable between 1 and 16 time quanta but may be automatically lengthened to compensate for positive phase drifts due to differences in the frequency of the various nodes of the network.
- Bit segment 2 (BS2): defines the location of the transmit point. It represents the PHASE\_SEG2 of the CAN standard. Its duration is programmable between 1 and 8 time quanta but may also be automatically shortened to compensate for negative phase drifts.

The resynchronization Jump Width (SJW) defines an upper bound to the amount of lengthening or shortening of the bit segments. It is programmable between 1 and 4 time quanta.

A valid edge is defined as the first transition in a bit time from dominant to recessive bus level provided the controller itself does not send a recessive bit.

If a valid edge is detected in BS1 instead of SYNC\_SEG, BS1 is extended by up to SJW so that the sample point is delayed.

Conversely, if a valid edge is detected in BS2 instead of SYNC\_SEG, BS2 is shortened by up to SJW so that the transmit point is moved earlier.

As a safeguard against programming errors, the configuration of the Bit Timing Register (CAN\_BTR) is only possible while the device is in Standby mode.

Note:

For a detailed description of the CAN bit timing and resynchronization mechanism, please refer to the ISO 11898 standard.

#### Figure 311. Bit timing

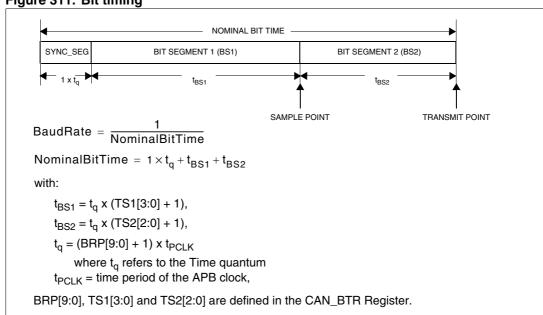
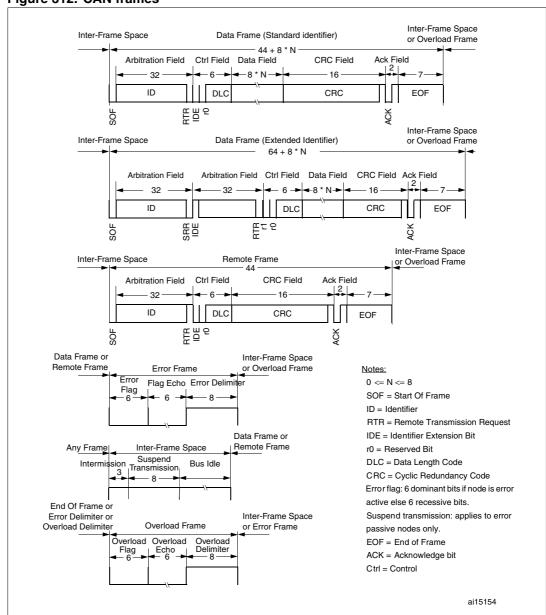


Figure 312. CAN frames



## 27.8 bxCAN interrupts

Four interrupt vectors are dedicated to bxCAN. Each interrupt source can be independently enabled or disabled by means of the CAN Interrupt Enable Register (CAN\_IER).

Figure 313. Event flags and interrupt generation CAN\_IER TRANSMIT INTERRUPT TMEIE & CAN\_TSR FMPIE0 & FIFO 0 FMP0 INTERRUPT FFIE0 & + CAN\_RF0R FULL0 FOVIE0 & FOVR0 **^** FIFO 1 FMP1 INTERRUPT FFIE1 & CAN\_RF1R -FULL1 FOVIE1 & FOVR1 ERRIE **EWGIE** & EWGF **EPVIE** & CAN ESR -EPVF & BOFIE ERRI CAN MSR & BOFF STATUS CHANGE ERROR LECIE & INTERRUPT 1≤LEC≤6 WKUIE & WKUI CAN MSR

- The transmit interrupt can be generated by the following events:
  - Transmit mailbox 0 becomes empty, RQCP0 bit in the CAN\_TSR register set.

&

- Transmit mailbox 1 becomes empty, RQCP1 bit in the CAN\_TSR register set.
- Transmit mailbox 2 becomes empty, RQCP2 bit in the CAN\_TSR register set.
- The **FIFO 0 interrupt** can be generated by the following events:
  - Reception of a new message, FMP0 bits in the CAN\_RF0R register are not '00'.
  - FIFO0 full condition, FULL0 bit in the CAN\_RF0R register set.
  - FIFO0 overrun condition, FOVR0 bit in the CAN\_RF0R register set.
- The FIFO 1 interrupt can be generated by the following events:
  - Reception of a new message, FMP1 bits in the CAN\_RF1R register are not '00'.
  - FIFO1 full condition, FULL1 bit in the CAN\_RF1R register set.

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- FIFO1 overrun condition, FOVR1 bit in the CAN RF1R register set.
- The error and status change interrupt can be generated by the following events:
  - Error condition, for more details on error conditions please refer to the CAN Error Status register (CAN\_ESR).
  - Wakeup condition, SOF monitored on the CAN Rx signal.
  - Entry into Sleep mode.

## 27.9 CAN registers

The peripheral registers have to be accessed by words (32 bits).

## 27.9.1 Register access protection

Erroneous access to certain configuration registers can cause the hardware to temporarily disturb the whole CAN network. Therefore the CAN\_BTR register can be modified by software only while the CAN hardware is in initialization mode.

Although the transmission of incorrect data will not cause problems at the CAN network level, it can severely disturb the application. A transmit mailbox can be only modified by software while it is in empty state, refer to *Figure 305: Transmit mailbox states*.

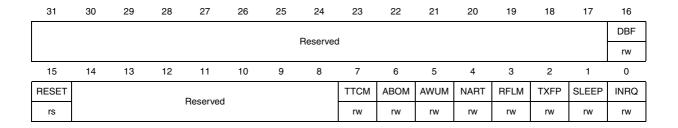
The filter values can be modified either deactivating the associated filter banks or by setting the FINIT bit. Moreover, the modification of the filter configuration (scale, mode and FIFO assignment) in CAN\_FMxR, CAN\_FSxR and CAN\_FFAR registers can only be done when the filter initialization mode is set (FINIT=1) in the CAN\_FMR register.

## 27.9.2 CAN control and status registers

Refer to Section 1.1 for a list of abbreviations used in register descriptions.

#### CAN master control register (CAN\_MCR)

Address offset: 0x00 Reset value: 0x0001 0002



Bits 31:17 Reserved, must be kept at reset value.

Bit 16 DBF: Debug freeze

0: CAN working during debug

1: CAN reception/transmission frozen during debug. Reception FIFOs can still be accessed/controlled normally.

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#### Bit 15 RESET: bxCAN software master reset

- 0: Normal operation.
- 1: Force a master reset of the bxCAN -> Sleep mode activated after reset (FMP bits and CAN\_MCR register are initialized to the reset values). This bit is automatically reset to 0.

#### Bits 14:8 Reserved, must be kept at reset value.

#### Bit 7 TTCM: Time triggered communication mode

- 0: Time Triggered Communication mode disabled.
- 1: Time Triggered Communication mode enabled

Note: For more information on Time Triggered Communication mode, please refer to Section 27.7.2: Time triggered communication mode.

#### Bit 6 ABOM: Automatic bus-off management

This bit controls the behavior of the CAN hardware on leaving the Bus-Off state.

- 0: The Bus-Off state is left on software request, once 128 occurrences of 11 recessive bits have been monitored and the software has first set and cleared the INRQ bit of the CAN\_MCR register.
- 1: The Bus-Off state is left automatically by hardware once 128 occurrences of 11 recessive bits have been monitored.

For detailed information on the Bus-Off state please refer to *Section 27.7.6: Error management.* 

#### Bit 5 AWUM: Automatic wakeup mode

This bit controls the behavior of the CAN hardware on message reception during Sleep mode.

- 0: The Sleep mode is left on software request by clearing the SLEEP bit of the CAN\_MCR register.
- 1: The Sleep mode is left automatically by hardware on CAN message detection.

The SLEEP bit of the CAN\_MCR register and the SLAK bit of the CAN\_MSR register are cleared by hardware.

## Bit 4 NART: No automatic retransmission

- 0: The CAN hardware will automatically retransmit the message until it has been successfully transmitted according to the CAN standard.
- 1: A message will be transmitted only once, independently of the transmission result (successful, error or arbitration lost).

## Bit 3 RFLM: Receive FIFO locked mode

- 0: Receive FIFO not locked on overrun. Once a receive FIFO is full the next incoming message will overwrite the previous one.
- 1: Receive FIFO locked against overrun. Once a receive FIFO is full the next incoming message will be discarded.

## Bit 2 TXFP: Transmit FIFO priority

This bit controls the transmission order when several mailboxes are pending at the same time

- 0: Priority driven by the identifier of the message
- 1: Priority driven by the request order (chronologically)

#### Bit 1 SLEEP: Sleep mode request

This bit is set by software to request the CAN hardware to enter the Sleep mode. Sleep mode will be entered as soon as the current CAN activity (transmission or reception of a CAN frame) has been completed.

This bit is cleared by software to exit Sleep mode.

This bit is cleared by hardware when the AWUM bit is set and a SOF bit is detected on the CAN Rx signal.

This bit is set after reset - CAN starts in Sleep mode.

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## Bit 0 INRQ: Initialization request

The software clears this bit to switch the hardware into normal mode. Once 11 consecutive recessive bits have been monitored on the Rx signal the CAN hardware is synchronized and ready for transmission and reception. Hardware signals this event by clearing the INAK bit in the CAN\_MSR register.

Software sets this bit to request the CAN hardware to enter initialization mode. Once software has set the INRQ bit, the CAN hardware waits until the current CAN activity (transmission or reception) is completed before entering the initialization mode. Hardware signals this event by setting the INAK bit in the CAN\_MSR register.

## CAN master status register (CAN\_MSR)

Address offset: 0x04 Reset value: 0x0000 0C02

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15 14 13 12 11 10 9 8 7 6 5 4											4	3	2	1	0
	Booo	nuod		RX	SAMP	RXM	TXM		Reserved		SLAKI	WKUI	ERRI	SLAK	INAK
	Reserved.			r	r	r	r		neserveu		rc_w1	rc_w1	rc_w1	r	r

Bits 31:12 Reserved, must be kept at reset value.

Bit 11 RX: CAN Rx signal

Monitors the actual value of the CAN\_RX Pin.

Bit 10 SAMP: Last sample point

The value of RX on the last sample point (current received bit value).

Bit 9 RXM: Receive mode

The CAN hardware is currently receiver.

Bit 8 TXM: Transmit mode

The CAN hardware is currently transmitter.

Bits 7:5 Reserved, must be kept at reset value.

Bit 4 SLAKI: Sleep acknowledge interrupt

When SLKIE=1, this bit is set by hardware to signal that the bxCAN has entered Sleep Mode. When set, this bit generates a status change interrupt if the SLKIE bit in the CAN\_IER register is set.

This bit is cleared by software or by hardware, when SLAK is cleared.

Note: When SLKIE=0, no polling on SLAKI is possible. In this case the SLAK bit can be polled.

Bit 3 WKUI: Wakeup interrupt

This bit is set by hardware to signal that a SOF bit has been detected while the CAN hardware was in Sleep mode. Setting this bit generates a status change interrupt if the WKUIE bit in the CAN\_IER register is set.

This bit is cleared by software.

#### Bit 2 ERRI: Error interrupt

This bit is set by hardware when a bit of the CAN\_ESR has been set on error detection and the corresponding interrupt in the CAN\_IER is enabled. Setting this bit generates a status change interrupt if the ERRIE bit in the CAN\_IER register is set.

This bit is cleared by software.

#### Bit 1 SLAK: Sleep acknowledge

This bit is set by hardware and indicates to the software that the CAN hardware is now in Sleep mode. This bit acknowledges the Sleep mode request from the software (set SLEEP bit in CAN\_MCR register).

This bit is cleared by hardware when the CAN hardware has left Sleep mode (to be synchronized on the CAN bus). To be synchronized the hardware has to monitor a sequence of 11 consecutive recessive bits on the CAN RX signal.

Note: The process of leaving Sleep mode is triggered when the SLEEP bit in the CAN\_MCR register is cleared. Please refer to the AWUM bit of the CAN\_MCR register description for detailed information for clearing SLEEP bit

#### Bit 0 INAK: Initialization acknowledge

This bit is set by hardware and indicates to the software that the CAN hardware is now in initialization mode. This bit acknowledges the initialization request from the software (set INRQ bit in CAN\_MCR register).

This bit is cleared by hardware when the CAN hardware has left the initialization mode (to be synchronized on the CAN bus). To be synchronized the hardware has to monitor a sequence of 11 consecutive recessive bits on the CAN RX signal.

## CAN transmit status register (CAN\_TSR)

Address offset: 0x08 Reset value: 0x1C00 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
LOW2	LOW1	LOW0	TME2	TME1	TME0	CODI	E[1:0]	ABRQ 2		Reserved		TERR 2	ALST2	TXOK 2	RQCP 2
r	r	r	r	r	r	r	r	rs				rc_w1	rc_w1	rc_w1	rc_w1
15	14	13	12	11	10	9	8	7	6 5 4		4	3	2	1	0
ABRQ 1		Reserved			ALST1	TXOK 1	RQCP 1	ABRQ 0		Reserved		TERR 0	ALST0	TXOK 0	RQCP 0
rs	Res.			rc_w1	rc_w1	rc_w1	rc_w1	rs				rc_w1	rc_w1	rc_w1	rc_w1

## Bit 31 LOW2: Lowest priority flag for mailbox 2

This bit is set by hardware when more than one mailbox are pending for transmission and mailbox 2 has the lowest priority.

#### Bit 30 LOW1: Lowest priority flag for mailbox 1

This bit is set by hardware when more than one mailbox are pending for transmission and mailbox 1 has the lowest priority.

#### Bit 29 LOW0: Lowest priority flag for mailbox 0

This bit is set by hardware when more than one mailbox are pending for transmission and mailbox 0 has the lowest priority.

Note: The LOW[2:0] bits are set to zero when only one mailbox is pending.

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Bit 28 TME2: Transmit mailbox 2 empty

This bit is set by hardware when no transmit request is pending for mailbox 2.

Bit 27 TME1: Transmit mailbox 1 empty

This bit is set by hardware when no transmit request is pending for mailbox 1.

Bit 26 TME0: Transmit mailbox 0 empty

This bit is set by hardware when no transmit request is pending for mailbox 0.

Bits 25:24 CODE[1:0]: Mailbox code

In case at least one transmit mailbox is free, the code value is equal to the number of the next transmit mailbox free.

In case all transmit mailboxes are pending, the code value is equal to the number of the transmit mailbox with the lowest priority.

Bit 23 ABRQ2: Abort request for mailbox 2

Set by software to abort the transmission request for the corresponding mailbox.

Cleared by hardware when the mailbox becomes empty.

Setting this bit has no effect when the mailbox is not pending for transmission.

Bits 22:20 Reserved, must be kept at reset value.

Bit 19 TERR2: Transmission error of mailbox 2

This bit is set when the previous TX failed due to an error.

Bit 18 ALST2: Arbitration lost for mailbox 2

This bit is set when the previous TX failed due to an arbitration lost.

Bit 17 TXOK2: Transmission OK of mailbox 2

The hardware updates this bit after each transmission attempt.

0: The previous transmission failed

1: The previous transmission was successful

This bit is set by hardware when the transmission request on mailbox 2 has been completed successfully. Please refer to *Figure 305*.

Bit 16 RQCP2: Request completed mailbox2

Set by hardware when the last request (transmit or abort) has been performed.

Cleared by software writing a "1" or by hardware on transmission request (TXRQ2 set in CAN\_TMID2R register).

Clearing this bit clears all the status bits (TXOK2, ALST2 and TERR2) for Mailbox 2.

Bit 15 ABRQ1: Abort request for mailbox 1

Set by software to abort the transmission request for the corresponding mailbox.

Cleared by hardware when the mailbox becomes empty.

Setting this bit has no effect when the mailbox is not pending for transmission.

Bits 14:12 Reserved, must be kept at reset value.

Bit 11 TERR1: Transmission error of mailbox1

This bit is set when the previous TX failed due to an error.

Bit 10 ALST1: Arbitration lost for mailbox1

This bit is set when the previous TX failed due to an arbitration lost.

Bit 9 TXOK1: Transmission OK of mailbox1

The hardware updates this bit after each transmission attempt.

0: The previous transmission failed

1: The previous transmission was successful

This bit is set by hardware when the transmission request on mailbox 1 has been completed successfully. Please refer to *Figure 305* 



#### Bit 8 RQCP1: Request completed mailbox1

Set by hardware when the last request (transmit or abort) has been performed.

Cleared by software writing a "1" or by hardware on transmission request (TXRQ1 set in CAN\_TI1R register).

Clearing this bit clears all the status bits (TXOK1, ALST1 and TERR1) for Mailbox 1.

#### Bit 7 ABRQ0: Abort request for mailbox0

Set by software to abort the transmission request for the corresponding mailbox.

Cleared by hardware when the mailbox becomes empty.

Setting this bit has no effect when the mailbox is not pending for transmission.

#### Bits 6:4 Reserved, must be kept at reset value.

#### Bit 3 TERR0: Transmission error of mailbox0

This bit is set when the previous TX failed due to an error.

#### Bit 2 ALST0: Arbitration lost for mailbox0

This bit is set when the previous TX failed due to an arbitration lost.

#### Bit 1 TXOK0: Transmission OK of mailbox0

The hardware updates this bit after each transmission attempt.

0: The previous transmission failed

1: The previous transmission was successful

This bit is set by hardware when the transmission request on mailbox 1 has been completed successfully. Please refer to *Figure 305* 

#### Bit 0 RQCP0: Request completed mailbox0

Set by hardware when the last request (transmit or abort) has been performed.

Cleared by software writing a "1" or by hardware on transmission request (TXRQ0 set in CAN\_TIOR register).

Clearing this bit clears all the status bits (TXOK0, ALST0 and TERR0) for Mailbox 0.

## CAN receive FIFO 0 register (CAN\_RF0R)

Address offset: 0x0C Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Res	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
				Poor	erved					RFOM0	FOVR0	FULL0	Res.	FMP	0[1:0]
				nese	erveu					rs	rc_w1	rc_w1	nes.	r	r

#### Bit 31:6 Reserved, must be kept at reset value.

#### Bit 5 RFOM0: Release FIFO 0 output mailbox

Set by software to release the output mailbox of the FIFO. The output mailbox can only be released when at least one message is pending in the FIFO. Setting this bit when the FIFO is empty has no effect. If at least two messages are pending in the FIFO, the software has to release the output mailbox to access the next message.

Cleared by hardware when the output mailbox has been released.

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#### Bit 4 FOVR0: FIFO 0 overrun

This bit is set by hardware when a new message has been received and passed the filter while the FIFO was full.

This bit is cleared by software.

#### Bit 3 FULLO: FIFO 0 full

Set by hardware when three messages are stored in the FIFO.

This bit is cleared by software.

Bit 2 Reserved, must be kept at reset value.

#### Bits 1:0 FMP0[1:0]: FIFO 0 message pending

These bits indicate how many messages are pending in the receive FIFO. FMP is increased each time the hardware stores a new message in to the FIFO. FMP is

decreased each time the software releases the output mailbox by setting the RFOM0 bit.

## CAN receive FIFO 1 register (CAN\_RF1R)

Address offset: 0x10 Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Res	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
				Pos	erved					RFOM1	FOVR1	FULL1	Res.	FMP <sup>-</sup>	1[1:0]
				nest	ei veu					rs	rc_w1	rc_w1	nes.	r	r

## Bits 31:6 Reserved, must be kept at reset value.

## Bit 5 RFOM1: Release FIFO 1 output mailbox

Set by software to release the output mailbox of the FIFO. The output mailbox can only be released when at least one message is pending in the FIFO. Setting this bit when the FIFO is empty has no effect. If at least two messages are pending in the FIFO, the software has to release the output mailbox to access the next message.

Cleared by hardware when the output mailbox has been released.

#### Bit 4 FOVR1: FIFO 1 overrun

This bit is set by hardware when a new message has been received and passed the filter while the FIFO was full.

This bit is cleared by software.

#### Bit 3 FULL1: FIFO 1 full

Set by hardware when three messages are stored in the FIFO.

This bit is cleared by software.

Bit 2 Reserved, must be kept at reset value.

## Bits 1:0 FMP1[1:0]: FIFO 1 message pending

These bits indicate how many messages are pending in the receive FIFO1.

FMP1 is increased each time the hardware stores a new message in to the FIFO1. FMP is decreased each time the software releases the output mailbox by setting the RFOM1 bit.

## CAN interrupt enable register (CAN\_IER)

Address offset: 0x14 Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
						Dane								SLKIE	WKUIE
		Reserved													rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ERRIE	LEC										FMP IE0	TME IE			
rw				rw	rw	rw	rw		rw	rw	rw	rw	rw	rw	rw

Bits 31:18 Reserved, must be kept at reset value.

Bit 17 **SLKIE**: Sleep interrupt enable

0: No interrupt when SLAKI bit is set.

1: Interrupt generated when SLAKI bit is set.

Bit 16 WKUIE: Wakeup interrupt enable

0: No interrupt when WKUI is set.

1: Interrupt generated when WKUI bit is set.

Bit 15 ERRIE: Error interrupt enable

0: No interrupt will be generated when an error condition is pending in the CAN\_ESR.

1: An interrupt will be generation when an error condition is pending in the CAN\_ESR.

Bits 14:12 Reserved, must be kept at reset value.

Bit 11 LECIE: Last error code interrupt enable

0: ERRI bit will not be set when the error code in LEC[2:0] is set by hardware on error detection.

1: ERRI bit will be set when the error code in LEC[2:0] is set by hardware on error detection.

Bit 10 BOFIE: Bus-off interrupt enable

0: ERRI bit will not be set when BOFF is set.

1: ERRI bit will be set when BOFF is set.

Bit 9 EPVIE: Error passive interrupt enable

0: ERRI bit will not be set when EPVF is set.

1: ERRI bit will be set when EPVF is set.

Bit 8 **EWGIE**: Error warning interrupt enable

0: ERRI bit will not be set when EWGF is set.

1: ERRI bit will be set when EWGF is set.

Bit 7 Reserved, must be kept at reset value.

Bit 6 FOVIE1: FIFO overrun interrupt enable

0: No interrupt when FOVR is set.

1: Interrupt generation when FOVR is set.

Bit 5 FFIE1: FIFO full interrupt enable

0: No interrupt when FULL bit is set.

1: Interrupt generated when FULL bit is set.

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#### Bit 4 FMPIE1: FIFO message pending interrupt enable

0: No interrupt generated when state of FMP[1:0] bits are not 00b.

1: Interrupt generated when state of FMP[1:0] bits are not 00b.

#### Bit 3 FOVIE0: FIFO overrun interrupt enable

0: No interrupt when FOVR bit is set.

1: Interrupt generated when FOVR bit is set.

#### Bit 2 FFIE0: FIFO full interrupt enable

0: No interrupt when FULL bit is set.

1: Interrupt generated when FULL bit is set.

#### Bit 1 FMPIE0: FIFO message pending interrupt enable

0: No interrupt generated when state of FMP[1:0] bits are not 00b.

1: Interrupt generated when state of FMP[1:0] bits are not 00b.

#### Bit 0 TMEIE: Transmit mailbox empty interrupt enable

0: No interrupt when RQCPx bit is set.

1: Interrupt generated when RQCPx bit is set. *Note: Refer to Section 27.8: bxCAN interrupts.* 

## **CAN error status register (CAN\_ESR)**

Address offset: 0x18 Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
			REC	[7:0]							TEC	[7:0]			
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
				Reserved	1					LEC[2:0]		Res.	BOFF	EPVF	EWGF
				neserved	1				rw	rw	rw	nes.	r	r	r

## Bits 31:24 REC[7:0]: Receive error counter

The implementing part of the fault confinement mechanism of the CAN protocol. In case of an error during reception, this counter is incremented by 1 or by 8 depending on the error condition as defined by the CAN standard. After every successful reception the counter is decremented by 1 or reset to 120 if its value was higher than 128. When the counter value exceeds 127, the CAN controller enters the error passive state.

# Bits 23:16 **TEC[7:0]**: Least significant byte of the 9-bit transmit error counter

The implementing part of the fault confinement mechanism of the CAN protocol.

Bits 15:7 Reserved, must be kept at reset value.

## Bits 6:4 LEC[2:0]: Last error code

This field is set by hardware and holds a code which indicates the error condition of the last error detected on the CAN bus. If a message has been transferred (reception or transmission) without error, this field will be cleared to '0'.

The LEC[2:0] bits can be set to value 0b111 by software. They are updated by hardware to indicate the current communication status.

000: No Error 001: Stuff Error 010: Form Error

011: Acknowledgment Error 100: Bit recessive Error 101: Bit dominant Error

110: CRC Error 111: Set by software

Bit 3 Reserved, must be kept at reset value.

#### Bit 2 BOFF: Bus-off flag

This bit is set by hardware when it enters the bus-off state. The bus-off state is entered on TEC overflow, greater than 255, refer to *Section 27.7.6 on page 800*.

#### Bit 1 EPVF: Error passive flag

This bit is set by hardware when the Error Passive limit has been reached (Receive Error Counter or Transmit Error Counter>127).

#### Bit 0 EWGF: Error warning flag

This bit is set by hardware when the warning limit has been reached (Receive Error Counter or Transmit Error Counter≥96).

## CAN bit timing register (CAN\_BTR)

Address offset: 0x1C Reset value: 0x0123 0000

This register can only be accessed by the software when the CAN hardware is in initialization mode.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
SILM	LBKM		Rese	nyod		SJW	/[1:0]	Res.		TS2[2:0]			TS1	[3:0]	
rw	rw		11636	rveu		rw	rw		rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		Rese	nuod							BRP	[9:0]				
		nese	rveu			rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bit 31 SILM: Silent mode (debug)

0: Normal operation1: Silent Mode

Bit 30 LBKM: Loop back mode (debug)

0: Loop Back Mode disabled1: Loop Back Mode enabled

Bits 29:26 Reserved, must be kept at reset value.

Bits 25:24 SJW[1:0]: Resynchronization jump width

These bits define the maximum number of time quanta the CAN hardware is allowed to lengthen or shorten a bit to perform the resynchronization.

 $t_{RJW} = t_{CAN} \times (SJW[1:0] + 1)$ 

Bit 23 Reserved, must be kept at reset value.

Bits 22:20 TS2[2:0]: Time segment 2

These bits define the number of time quanta in Time Segment 2.

 $t_{BS2} = t_{CAN} \times (TS2[2:0] + 1)$ 

Bits 19:16 TS1[3:0]: Time segment 1

These bits define the number of time quanta in Time Segment 1

 $t_{BS1} = t_{CAN} x (TS1[3:0] + 1)$ 

For more information on bit timing, please refer to Section 27.7.7: Bit timing on page 800.

Bits 15:10 Reserved, must be kept at reset value.

Bits 9:0 BRP[9:0]: Baud rate prescaler

These bits define the length of a time quanta.

 $t_q = (BRP[9:0]+1) \times t_{PCLK}$ 

## 27.9.3 CAN mailbox registers

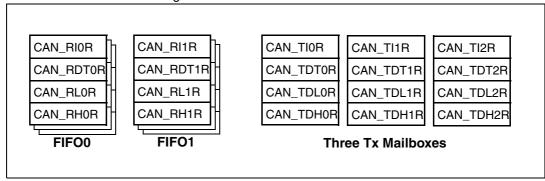
This chapter describes the registers of the transmit and receive mailboxes. Refer to Section 27.7.5: Message storage on page 798 for detailed register mapping.

Transmit and receive mailboxes have the same registers except:

- The FMI field in the CAN RDTxR register.
- A receive mailbox is always write protected.
- A transmit mailbox is write-enabled only while empty, corresponding TME bit in the CAN\_TSR register set.

There are 3 TX Mailboxes and 2 RX Mailboxes. Each RX Mailbox allows access to a 3 level depth FIFO, the access being offered only to the oldest received message in the FIFO.

Each mailbox consist of 4 registers.



## CAN TX mailbox identifier register (CAN\_TIxR) (x=0..2)

Address offsets: 0x180, 0x190, 0x1A0

Reset value: 0xXXXX XXXX (except bit 0, TXRQ = 0)

All TX registers are write protected when the mailbox is pending transmission (TMEx reset).

This register also implements the TX request control (bit 0) - reset value 0.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
				STID[1	0:0]/EXID	[28:18]						E	XID[17:1	3]	
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
					E	EXID[12:0	]						IDE	RTR	TXRQ
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

#### Bits 31:21 STID[10:0]/EXID[28:18]: Standard identifier or extended identifier

The standard identifier or the MSBs of the extended identifier (depending on the IDE bit value).

## Bit 20:3 EXID[17:0]: Extended identifier

The LSBs of the extended identifier.

Bit 2 IDE: Identifier extension

This bit defines the identifier type of message in the mailbox.

0: Standard identifier.

1: Extended identifier.

Bit 1 RTR: Remote transmission request

0: Data frame

1: Remote frame

Bit 0 TXRQ: Transmit mailbox request

Set by software to request the transmission for the corresponding mailbox.

Cleared by hardware when the mailbox becomes empty.

# CAN mailbox data length control and time stamp register (CAN\_TDTxR) (x=0..2)

All bits of this register are write protected when the mailbox is not in empty state.

Address offsets: 0x184, 0x194, 0x1A4

Reset value: 0xXXXX XXXX

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							TIME	[15:0]							
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			Dogoryod	1			TGT		Rese	ruod			DLC	[3:0]	
	Reserved								nese	erveu		rw	rw	rw	rw

## Bits 31:16 TIME[15:0]: Message time stamp

This field contains the 16-bit timer value captured at the SOF transmission.

Bits 15:9 Reserved, must be kept at reset value.

#### Bit 8 TGT: Transmit global time

This bit is active only when the hardware is in the Time Trigger Communication mode, TTCM bit of the CAN\_MCR register is set.

0: Time stamp TIME[15:0] is not sent.

1: Time stamp TIME[15:0] value is sent in the last two data bytes of the 8-byte message: TIME[7:0] in data byte 7 and TIME[15:8] in data byte 6, replacing the data written in CAN\_TDHxR[31:16] register (DATA6[7:0] and DATA7[7:0]). DLC must be programmed as 8 in order these two bytes to be sent over the CAN bus.

Bits 7:4 Reserved, must be kept at reset value.

## Bits 3:0 DLC[3:0]: Data length code

This field defines the number of data bytes a data frame contains or a remote frame request. A message can contain from 0 to 8 data bytes, depending on the value in the DLC field.

## CAN mailbox data low register (CAN\_TDLxR) (x=0..2)

All bits of this register are write protected when the mailbox is not in empty state.

Address offsets: 0x188, 0x198, 0x1A8

Reset value: 0xXXXX XXXX

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
			DATA	3[7:0]							DATA	2[7:0]			
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			DATA	1[7:0]							DATA	.0[7:0]			
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:24 DATA3[7:0]: Data byte 3

Data byte 3 of the message.

Bits 23:16 DATA2[7:0]: Data byte 2

Data byte 2 of the message.

Bits 15:8 DATA1[7:0]: Data byte 1

Data byte 1 of the message.

Bits 7:0 DATA0[7:0]: Data byte 0

Data byte 0 of the message.

A message can contain from 0 to 8 data bytes and starts with byte 0.

## CAN mailbox data high register (CAN\_TDHxR) (x=0..2)

All bits of this register are write protected when the mailbox is not in empty state.

Address offsets: 0x18C, 0x19C, 0x1AC

Reset value: 0xXXXX XXXX

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
			DATA	7[7:0]							DATA	6[7:0]			
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			DATA	5[7:0]							DATA	4[7:0]			
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:24 DATA7[7:0]: Data byte 7

Data byte 7 of the message.

Note: If TGT of this message and TTCM are active, DATA7 and DATA6 will be replaced by the

TIME stamp value.

Bits 23:16 DATA6[7:0]: Data byte 6

Data byte 6 of the message.

Bits 15:8 DATA5[7:0]: Data byte 5

Data byte 5 of the message.

Bits 7:0 **DATA4[7:0]**: Data byte 4

Data byte 4 of the message.

## CAN receive FIFO mailbox identifier register (CAN\_RIxR) (x=0..1)

Address offsets: 0x1B0, 0x1C0 Reset value: 0xXXXX XXXX

All RX registers are write protected.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
				STID[1	0:0]/EXID	[28:18]						Е	XID[17:1	3]	
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
					E	EXID[12:0	]						IDE	RTR	Res.
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	nes.

## Bits 31:21 STID[10:0]/EXID[28:18]: Standard identifier or extended identifier

The standard identifier or the MSBs of the extended identifier (depending on the IDE bit value).

Bits 20:3 EXID[17:0]: Extended identifier

The LSBs of the extended identifier.

Bit 2 IDE: Identifier extension

This bit defines the identifier type of message in the mailbox.

 $0: Standard\ identifier.$ 

1: Extended identifier.

Bit 1 RTR: Remote transmission request

0: Data frame

1: Remote frame

Bit 0 Reserved, must be kept at reset value.

# CAN receive FIFO mailbox data length control and time stamp register (CAN\_RDTxR) (x=0..1)

Address offsets: 0x1B4, 0x1C4 Reset value: 0xXXXX XXXX

All RX registers are write protected.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							TIME	[15:0]							
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			FMI	[7:0]					Rese	ruod			DLC	[3:0]	
r	r	r	r	r	r	r	r		nese	ii veu		r	r	r	r

## Bits 31:16 TIME[15:0]: Message time stamp

This field contains the 16-bit timer value captured at the SOF detection.

#### Bits 15:8 FMI[7:0]: Filter match index

This register contains the index of the filter the message stored in the mailbox passed through. For more details on identifier filtering please refer to *Section 27.7.4: Identifier filtering on page 794* - **Filter Match Index** paragraph.

Bits 7:4 Reserved, must be kept at reset value.

## Bits 3:0 DLC[3:0]: Data length code

This field defines the number of data bytes a data frame contains (0 to 8). It is 0 in the case of a remote frame request.



## CAN receive FIFO mailbox data low register (CAN\_RDLxR) (x=0..1)

All bits of this register are write protected when the mailbox is not in empty state.

Address offsets: 0x1B8, 0x1C8 Reset value: 0xXXXX XXXX

All RX registers are write protected.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
			DATA	3[7:0]							DATA	2[7:0]			
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			DATA	1[7:0]							DATA	0[7:0]			
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:24 DATA3[7:0]: Data Byte 3

Data byte 3 of the message.

Bits 23:16 DATA2[7:0]: Data Byte 2

Data byte 2 of the message.

Bits 15:8 DATA1[7:0]: Data Byte 1

Data byte 1 of the message.

Bits 7:0 DATA0[7:0]: Data Byte 0

Data byte 0 of the message.

A message can contain from 0 to 8 data bytes and starts with byte 0.

## CAN receive FIFO mailbox data high register (CAN\_RDHxR) (x=0..1)

Address offsets: 0x1BC, 0x1CC Reset value: 0xXXXX XXXX

All RX registers are write protected.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
			DATA	7[7:0]							DATA	6[7:0]			
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			DATA	5[7:0]							DATA	4[7:0]			
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:24 DATA7[7:0]: Data Byte 7

Data byte 3 of the message.

Bits 23:16 DATA6[7:0]: Data Byte 6

Data byte 2 of the message.

Bits 15:8 DATA5[7:0]: Data Byte 5

Data byte 1 of the message.

Bits 7:0 **DATA4[7:0]**: Data Byte 4 Data byte 0 of the message.

## 27.9.4 CAN filter registers

## **CAN filter master register (CAN\_FMR)**

Address offset: 0x200 Reset value: 0x2A1C 0E01

All bits of this register are set and cleared by software.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Rese	nuad			CAN2S	SB[5:0]						Reserved	ı			FINIT
Hese	erved	rw	rw	rw	rw	rw	rw				neserved	l			rw

Bits 31:14 Reserved, must be kept at reset value.

Bits 13:8 CAN2SB[5:0]: CAN2 start bank

These bits are set and cleared by software. They define the start bank for the CAN2 interface (Slave) in the range 0 to 27.

Note: When CAN2SB[5:0] = 28d, all the filters to CAN1 can be used.
When CAN2SB[5:0] is set to 0, no filters are assigned to CAN1.

Bits 7:1 Reserved, must be kept at reset value.

Bit 0 FINIT: Filter init mode

Initialization mode for filter banks

0: Active filters mode.

1: Initialization mode for the filters.

## **CAN filter mode register (CAN\_FM1R)**

Address offset: 0x204 Reset value: 0x0000 0000

This register can be written only when the filter initialization mode is set (FINIT=1) in the

CAN\_FMR register.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	Poor	erved		FBM27	FBM26	FBM25	FBM24	FBM23	FBM22	FBM21	FBM20	FBM19	FBM18	FBM17	FBM16
	nest	erveu		rw											
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
FBM15	FBM14	FBM13	FBM12	FBM11	FBM10	FBM9	FBM8	FBM7	FBM6	FBM5	FBM4	FBM3	FBM2	FBM1	FBM0
rw															

Note: Please refer to Figure 307: Filter bank scale configuration - register organization on page 796

Bits 31:28 Reserved, must be kept at reset value.

Bits 27:0 FBMx: Filter mode

Mode of the registers of Filter x.

0: Two 32-bit registers of filter bank x are in Identifier Mask mode.1: Two 32-bit registers of filter bank x are in Identifier List mode.

## CAN filter scale register (CAN\_FS1R)

Address offset: 0x20C Reset value: 0x0000 0000

This register can be written only when the filter initialization mode is set (FINIT=1) in the CAN\_FMR register.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	Rese	nuod		FSC27	FSC26	FSC25	FSC24	FSC23	FSC22	FSC21	FSC20	FSC19	FSC18	FSC17	FSC16
	nese	erveu		rw											
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
FSC15	FSC14	FSC13	FSC12	FSC11	FSC10	FSC9	FSC8	FSC7	FSC6	FSC5	FSC4	FSC3	FSC2	FSC1	FSC0
rw															

Bits 31:28 Reserved, must be kept at reset value.

Bits 27:0 FSCx: Filter scale configuration

These bits define the scale configuration of Filters 13-0.

0: Dual 16-bit scale configuration1: Single 32-bit scale configuration

Note: Please refer to Figure 307: Filter bank scale configuration - register organization on

page 796.

## **CAN filter FIFO assignment register (CAN\_FFA1R)**

Address offset: 0x214 Reset value: 0x0000 0000

This register can be written only when the filter initialization mode is set (FINIT=1) in the

CAN\_FMR register.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	Pose	erved		FFA27	FFA26	FFA25	FFA24	FFA23	FFA22	FFA21	FFA20	FFA19	FFA18	FFA17	FFA16
	nese	erveu		rw											
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
FFA15	FFA14	FFA13	FFA12	FFA11	FFA10	FFA9	FFA8	FFA7	FFA6	FFA5	FFA4	FFA3	FFA2	FFA1	FFA0
rw															

Bits 31:28 Reserved, must be kept at reset value.

Bits 27:0 FFAx: Filter FIFO assignment for filter x

The message passing through this filter will be stored in the specified FIFO.

0: Filter assigned to FIFO 0 1: Filter assigned to FIFO 1

## CAN filter activation register (CAN\_FA1R)

Address offset: 0x21C Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	Pos	erved		FACT27	FACT26	FACT25	FACT24	FACT23	FACT22	FACT21	FACT20	FACT19	FACT18	FACT17	FACT16
	nese	erveu		rw											
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
FACT15	FACT14	FACT13	FACT12	FACT11	FACT10	FACT9	FACT8	FACT7	FACT6	FACT5	FACT4	FACT3	FACT2	FACT1	FACT0
rw															

Bits 31:28 Reserved, must be kept at reset value.

Bits 27:0 FACTx: Filter active

The software sets this bit to activate Filter x. To modify the Filter x registers (CAN\_FxR[0:7]), the FACTx bit must be cleared or the FINIT bit of the CAN\_FMR register must be set.

0: Filter x is not active

1: Filter x is active

## Filter bank i register x (CAN\_FiRx) (i=0..27, x=1, 2)

Address offsets: 0x240..0x31C Reset value: 0xXXXX XXXX

There are 28 filter banks, i=0 .. 27. Each filter bank i is composed of two 32-bit registers,

CAN\_FiR[2:1].

This register can only be modified when the FACTx bit of the CAN\_FAxR register is cleared or when the FINIT bit of the CAN\_FMR register is set.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
FB31	FB30	FB29	FB28	FB27	FB26	FB25	FB24	FB23	FB22	FB21	FB20	FB19	FB18	FB17	FB16
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
15 FB15	14 FB14	13 FB13	12 FB12	11 FB11	10 FB10	9 FB9	8 FB8	7 FB7	6 FB6	5 FB5	4 FB4	3 FB3	2 FB2	1 FB1	0 FB0

In all configurations:

Bits 31:0 FB[31:0]: Filter bits

#### Identifier

Each bit of the register specifies the level of the corresponding bit of the expected identifier.

- 0: Dominant bit is expected
- 1: Recessive bit is expected

#### Mask

Each bit of the register specifies whether the bit of the associated identifier register must match with the corresponding bit of the expected identifier or not.

- 0: Don't care, the bit is not used for the comparison
- 1: Must match, the bit of the incoming identifier must have the same level has specified in the corresponding identifier register of the filter.

Note:

Depending on the scale and mode configuration of the filter the function of each register can differ. For the filter mapping, functions description and mask registers association, refer to Section 27.7.4: Identifier filtering on page 794.

A Mask/Identifier register in **mask mode** has the same bit mapping as in **identifier list** mode.

For the register mapping/addresses of the filter banks please refer to the Table 141 on page 825.

## 27.9.5 bxCAN register map

Refer to *Table 2 on page 50* for the register boundary addresses. The registers from offset 0x200 to 31C are present only in CAN1.

Table 141. bxCAN register map and reset values

Offset	Register	24 25 28 29 24 25 25 25 25 25 25 25 25 25 25 25 25 25	23	20 119 17	16	9 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	8 ~	0 7 8 8 6 0
0x000	CAN_MCR Reset value	Reser			DBF O RESET	Reserved	O TTCM	0 0 AWUM O O AWUM O O O AWUM O O O O O O O O O O O O O O O O O O O
0x004	CAN_MSR  Reset value		Reserved	d		<u> </u>	Σ	Served 0 0 0 1 0
0x008	CAN_TSR	LOW[2:0] TME[2:0] 0	Reserv Reserv		RQCP2 ABRQ1		ABRQ0	Beserved TERRO TXOKO TXOKO
0x00C	CAN_RF0R		0	0 0 0	0 0	0 0 0	0 0	O   PFOM0   O   FOVR0   O   FULL0   O   Puserved   O   O   O   O   O   O   O   O   O
0x010	CAN_RF1R  Reset value			Reserved				O   PFOM1   O   O   POVR1   O   O   POVR1   O   O   O   O   O   O   O   O   O
0x014	CAN_IER  Reset value	Reserve	d	o SLKIE	o WKUIE o ERRIE	Reserved S O O O	o EWGIE Reserved	FOVIE1   C   FOVIE1   C   FMPIE1   C   FMPIE1   C   FMPIE0   C   FMP
0x018	CAN_ESR  Reset value	REC[7:0]	0 0 0 0	FEC[7:0]	0	Reserved	-	C   C   C   C   C   C   C   C   C   C
0x01C	CAN_BTR  Reset value	Reserved 0 0 0	TS2[2:	:0] TS1[3:0	1	Reserved	0   0	BRP[9:0]
0x020- 0x17F				Reserved	i	,		
0x180	CAN_TIOR  Reset value	STID[10:0]/EXID[28:	18]	x   x   x   x	хх	EXID[17:0]	x   x	x x x x x x x 0 RTR DE
0x184	CAN_TDT0R  Reset value		[15:0]		V	Reserved	·	eserved DLC[3:0]
0x188	CAN_TDL0R	DATA3[7:0]	X   X   X     D/	X X X X X ATA2[7:0]	X	DATA1[7:0]	X	x   x   x   x   x   DATA0[7:0]
0x18C	Reset value  CAN_TDH0R	DATA7[7:0]		X X X X X X ATA6[7:0]	x x	DATA5[7:0]	x x	X X X X X X X X X DATA4[7:0]
0x190	Reset value  CAN_TI1R	X   X   X   X   X   X   X   X   X   STID[10:0]/EXID[28:	18]	1 1	хх	EXID[17:0]		A A A A A A A A A A A A A A A A A A A
	Reset value	x x x x x x x x x	x x x	x x x x	Х	X   X   X   X   X   X	ХХ	x x x x x x 0

Table 141. bxCAN register map and reset values (continued)

Offset	Register	2 2 2 2 3 3 3 3 4 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5	16 17 8 19 20 21 22 23 19 29 19 29 19 29 19 29 19 29 19 29 19 29 29 29 29 29 29 29 29 29 29 29 29 29		0 1 2 3 4 5 6 7
0.11000					
0x194	CAN_TDT1R	TIME	[15:0]	Reserved E	Reserved DLC[3:0]
	Reset value	x x x x x x x x x	x x x x x x x x x	x	x x x x
0x198	CAN_TDL1R	DATA3[7:0]	DATA2[7:0]	DATA1[7:0]	DATA0[7:0]
	Reset value			x   x   x   x   x   x   x   x	
0x19C					
	CAN_TDH1R	DATA7[7:0]	DATA6[7:0]	DATA5[7:0]	DATA4[7:0]
	Reset value	x x x x x x x x x	x x x x x x x x x	x x x x x x x x x	
0x1A0	CAN_TI2R	STID[10:0]/EXID[28:	EXID[17:0]		IDE RTR
	Reset value	x x x x x x x x x		X   X   X   X   X   X   X   X   X   X	x x x x x x x 0
0x1A4	CAN_TDT2R	TIME	[15:0]	Reserved L5	DLC[3:0]
	Reset value		x   x   x   x   x   x   x   x	Reserved P	Reserved x x x x
0x1A8					
	CAN_TDL2R	DATA3[7:0]	DATA2[7:0]	DATA1[7:0]	DATA0[7:0]
	Reset value	x x x x x x x x x	x x x x x x x x x	x x x x x x x x x	x x x x x x x x x
0x1AC	CAN_TDH2R	DATA7[7:0]	DATA6[7:0]	DATA5[7:0]	DATA4[7:0]
	Reset value			x x x x x x x x x	
0x1B0	CAN_RIOR	STID[10:0]/EXID[28:	18]	EXID[17:0]	IDE RTR
UXIBU	Reset value	x x x x x x x x	x x x x x x x x	x x x x x x x x x	X
0x1B4	CAN_RDT0R	TIME	[15:0]	FMI[7:0]	DI C[3:0]
					Heserved
	Reset value	x   x   x   x   x   x   x   x	x   x   x   x   x   x   x   x	x   x   x   x   x   x   x   x	x   x   x   x
0x1B8	CAN_RDL0R	DATA3[7:0]	DATA2[7:0]	DATA1[7:0]	DATA0[7:0]
	Reset value	x x x x x x x x x	x x x x x x x x x	x x x x x x x x x x	x x x x x x x x x
0x1BC	CAN_RDH0R	DATA7[7:0]	DATA6[7:0]	DATA5[7:0]	DATA4[7:0]
	Reset value	   x   x   x   x   x   x   x   x	   x   x   x   x   x   x   x   x	x   x   x   x   x   x   x   x	
0x1C0	CAN_RI1R	STID[10:0]/EXID[28:	18]	EXID[17:0]	IDE RTR
0.100	Reset value	x x x x x x x x	18] x x x x x x x x x x	x x x x x x x x x	x x x x x x x 8 8 8
0x1C4	CAN_RDT1R		[15:0]	FMI[7:0]	DI C[3:0]
					Heserved
	Reset value	x   x   x   x   x   x   x   x	x   x   x   x   x   x   x   x	x x x x x x x x x	
0x1C8	CAN_RDL1R	DATA3[7:0]	DATA2[7:0]	DATA1[7:0]	DATA0[7:0]
	Reset value	x x x x x x x x x	x x x x x x x x x	x x x x x x x x x	x x x x x x x x x
0x1CC	CAN_RDH1R	DATA7[7:0]	DATA6[7:0]	DATA5[7:0]	DATA4[7:0]
	Reset value	x   x   x   x   x   x   x   x		x   x   x   x   x   x   x   x	
0x1D0-					1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
0x1FF			Reserved		

Table 141. bxCAN register map and reset values (continued)

Offset	Register	2082812181818181818181818181818181818181	8 2 9 4 8 2 0			
0x200	CAN_FMR	Reserved CAN2SB[5:0]	Reserved $\mathbb{Z}$			
	Reset value	0 0 1 1 1 1	0 1			
0x204	CAN_FM1R	Reserved FBM[27:0]				
	Reset value		0 0 0 0 0 0 0 0 0 0 0			
0x208		Reserved				
0x20C	CAN_FS1R	Reserved FSC[27:0]				
	Reset value	Reset value 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0				
0x210	Reserved					
0x214	CAN_FFA1R	Reserved FFA[27:0]				
	Reset value	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0				
0x218		Reserved				
0x21C	CAN_FA1R	Reserved FACT[27:0]				
	Reset value					
0x220		Reserved				
0x224- 0x23F		Reserved				
0x240	CAN_F0R1	FB[31:0]				
	Reset value	X   X   X   X   X   X   X   X   X   X				
0x244	CAN_F0R2	FB[31:0]				
	Reset value	X   X   X   X   X   X   X   X   X   X				
0x248	CAN_F1R1	FB[31:0]				
	Reset value					
0x24C	CAN_F1R2	FB[31:0]				
	Reset value					
		:				
0x318	CAN_F27R1	FB[31:0]				
	Reset value	x   x   x   x   x   x   x   x   x   x	X   X   X   X   X   X   X   X   X   X			
0x31C	CAN_F27R2	FB[31:0]				
	Reset value	X	X   X   X   X   X   X   X   X   X   X			

# 28 Ethernet (ETH): media access control (MAC) with DMA controller

This section applies to the whole STM32F4xx family, unless otherwise specified.

## 28.1 Ethernet introduction

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The Ethernet peripheral enables the STM32F4xx to transmit and receive data over Ethernet in compliance with the IEEE 802.3-2002 standard.

The Ethernet provides a configurable, flexible peripheral to meet the needs of various applications and customers. It supports two industry standard interfaces to the external physical layer (PHY): the default media independent interface (MII) defined in the IEEE 802.3 specifications and the reduced media independent interface (RMII). It can be used in number of applications such as switches, network interface cards, etc.

The Ethernet is compliant with the following standards:

- IEEE 802.3-2002 for Ethernet MAC
- IEEE 1588-2008 standard for precision networked clock synchronization
- AMBA 2.0 for AHB Master/Slave ports
- RMII specification from RMII consortium

## 28.2 Ethernet main features

The Ethernet (ETH) peripheral includes the following features, listed by category:

#### 28.2.1 MAC core features

- Supports 10/100 Mbit/s data transfer rates with external PHY interfaces
- IEEE 802.3-compliant MII interface to communicate with an external Fast Ethernet PHY
- Supports both full-duplex and half-duplex operations
  - Supports CSMA/CD Protocol for half-duplex operation
  - Supports IEEE 802.3x flow control for full-duplex operation
  - Optional forwarding of received pause control frames to the user application in fullduplex operation
  - Back-pressure support for half-duplex operation
  - Automatic transmission of zero-quanta pause frame on deassertion of flow control input in full-duplex operation
- Preamble and start-of-frame data (SFD) insertion in Transmit, and deletion in Receive paths
- Automatic CRC and pad generation controllable on a per-frame basis
- Options for automatic pad/CRC stripping on receive frames
- Programmable frame length to support Standard frames with sizes up to 16 KB
- Programmable interframe gap (40-96 bit times in steps of 8)
- Supports a variety of flexible address filtering modes:
  - Up to four 48-bit perfect (DA) address filters with masks for each byte
  - Up to three 48-bit SA address comparison check with masks for each byte
  - 64-bit Hash filter (optional) for multicast and unicast (DA) addresses
  - Option to pass all multicast addressed frames
  - Promiscuous mode support to pass all frames without any filtering for network monitoring
  - Passes all incoming packets (as per filter) with a status report
- Separate 32-bit status returned for transmission and reception packets
- Supports IEEE 802.1Q VLAN tag detection for reception frames
- Separate transmission, reception, and control interfaces to the Application
- Supports mandatory network statistics with RMON/MIB counters (RFC2819/RFC2665)
- MDIO interface for PHY device configuration and management
- Detection of LAN wakeup frames and AMD Magic Packet™ frames
- Receive feature for checksum off-load for received IPv4 and TCP packets encapsulated by the Ethernet frame
- Enhanced receive feature for checking IPv4 header checksum and TCP, UDP, or ICMP checksum encapsulated in IPv4 or IPv6 datagrams
- Support Ethernet frame time stamping as described in IEEE 1588-2008. Sixty-four-bit time stamps are given in each frame's transmit or receive status
- Two sets of FIFOs: a 2-KB Transmit FIFO with programmable threshold capability, and a 2-KB Receive FIFO with a configurable threshold (default of 64 bytes)
- Receive Status vectors inserted into the Receive FIFO after the EOF transfer enables multiple-frame storage in the Receive FIFO without requiring another FIFO to store those frames' Receive Status



- Option to filter all error frames on reception and not forward them to the application in Store-and-Forward mode
- Option to forward under-sized good frames
- Supports statistics by generating pulses for frames dropped or corrupted (due to overflow) in the Receive FIFO
- Supports Store and Forward mechanism for transmission to the MAC core
- Automatic generation of PAUSE frame control or back pressure signal to the MAC core based on Receive FIFO-fill (threshold configurable) level
- Handles automatic retransmission of Collision frames for transmission.
- Discards frames on late collision, excessive collisions, excessive deferral and underrun conditions
- Software control to flush Tx FIFO
- Calculates and inserts IPv4 header checksum and TCP, UDP, or ICMP checksum in frames transmitted in Store-and-Forward mode
- Supports internal loopback on the MII for debugging

#### 28.2.2 DMA features

- Supports all AHB burst types in the AHB Slave Interface
- Software can select the type of AHB burst (fixed or indefinite burst) in the AHB Master interface.
- Option to select address-aligned bursts from AHB master port
- Optimization for packet-oriented DMA transfers with frame delimiters
- Byte-aligned addressing for data buffer support
- Dual-buffer (ring) or linked-list (chained) descriptor chaining
- Descriptor architecture, allowing large blocks of data transfer with minimum CPU intervention:
- each descriptor can transfer up to 8 KB of data
- Comprehensive status reporting for normal operation and transfers with errors
- Individual programmable burst size for Transmit and Receive DMA Engines for optimal host bus utilization
- Programmable interrupt options for different operational conditions
- Per-frame Transmit/Receive complete interrupt control
- Round-robin or fixed-priority arbitration between Receive and Transmit engines
- Start/Stop modes
- Current Tx/Rx Buffer pointer as status registers
- Current Tx/Rx Descriptor pointer as status registers

### 28.2.3 PTP features

- Received and transmitted frames time stamping
- Coarse and fine correction methods
- Trigger interrupt when system time becomes greater than target time
- Pulse per second output (product alternate function output)

# 28.3 Ethernet pins

Table 142 shows the MAC signals and the corresponding MII/RMII signal mapping. All MAC signals are mapped onto AF11, some signals are mapped onto different I/O pins, and should be configured in Alternate function mode (for more details, refer to Section 6.3.2: I/O pin multiplexer and mapping).

Table 142. Alternate function mapping

D t	AF11					
Port	ETH					
PA0-WKUP	ETH_MII_CRS					
PA1	ETH_MII _RX_CLK / ETH_RMII _REF_CLK					
PA2	ETH _MDIO					
PA3	ETH _MII_COL					
PA7	ETH_MII_RX_DV / ETH_RMII_CRS_DV					
PB0	ETH _MII_RXD2					
PB1	ETH _MII_RXD3					
PB5	ETH _PPS_OUT					
PB8	ETH _MII_TXD3					
PB10	ETH_ MII_RX_ER					
PB11	ETH _MII_TX_EN / ETH _RMII_TX_EN					
PB12	ETH _MII_TXD0 / ETH _RMII_TXD0					
PB13	ETH _MII_TXD1 / ETH _RMII_TXD1					
PC1	ETH _MDC					
PC2	ETH _MII_TXD2					
PC3	ETH _MII_TX_CLK					
PC4	ETH_MII_RXD0 / ETH_RMII_RXD0					
PC5	ETH _MII_RXD1/ ETH _RMII_RXD1					
PE2	ETH_MII_TXD3					
PG8	ETH_PPS_OUT					
PG11	ETH _MII_TX_EN / ETH _RMII_TX_EN					
PG13	ETH _MII_TXD0 / ETH _RMII_TXD0					
PG14	ETH _MII_TXD1 / ETH _RMII_TXD1					
PH2	ETH _MII_CRS					
PH3	ETH _MII_COL					
PH6	ETH _MII_RXD2					
PH7	ETH _MII_RXD3					
PI10	ETH _MII_RX_ER					
1	ı					

## 28.4 Ethernet functional description: SMI, MII and RMII

The Ethernet peripheral consists of a MAC 802.3 (media access control) with a dedicated DMA controller. It supports both default media-independent interface (MII) and reduced media-independent interface (RMII) through one selection bit (refer to SYSCFG\_PMC register).

The DMA controller interfaces with the Core and memories through the AHB Master and Slave interfaces. The AHB Master Interface controls data transfers while the AHB Slave interface accesses Control and Status Registers (CSR) space.

The Transmit FIFO (Tx FIFO) buffers data read from system memory by the DMA before transmission by the MAC Core. Similarly, the Receive FIFO (Rx FIFO) stores the Ethernet frames received from the line until they are transferred to system memory by the DMA.

The Ethernet peripheral also includes an SMI to communicate with external PHY. A set of configuration registers permit the user to select the desired mode and features for the MAC and the DMA controller.

Note: The AHB clock frequency must be at least 25 MHz when the Ethernet is used.

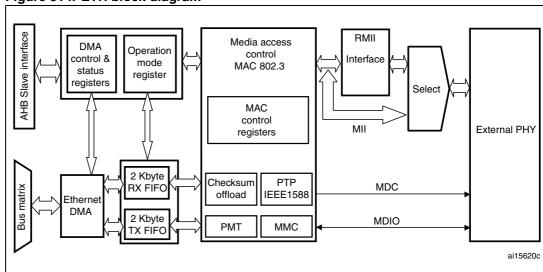


Figure 314. ETH block diagram

1. For AHB connections please refer to Figure 1: System architecture.

## 28.4.1 Station management interface: SMI

The station management interface (SMI) allows the application to access any PHY registers through a 2-wire clock and data lines. The interface supports accessing up to 32 PHYs.

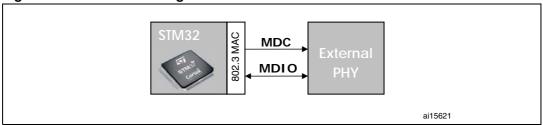
The application can select one of the 32 PHYs and one of the 32 registers within any PHY and send control data or receive status information. Only one register in one PHY can be addressed at any given time.

Both the MDC clock line and the MDIO data line are implemented as alternate function I/O in the microcontroller:

• MDC: a periodic clock that provides the timing reference for the data transfer at the maximum frequency of 2.5 MHz. The minimum high and low times for MDC must be

- 160 ns each, and the minimum period for MDC must be 400 ns. In idle state the SMI management interface drives the MDC clock signal low.
- MDIO: data input/output bitstream to transfer status information to/from the PHY device synchronously with the MDC clock signal

Figure 315. SMI interface signals



#### SMI frame format

The frame structure related to a read or write operation is shown in *Table 13*, the order of bit transmission must be from left to right.

Table 143. Management frame format

	Management frame fields							
	Preamble (32 bits) Start Operation PADDR RAD					ТА	Data (16 bits)	Idle
Read	1 1	01	10	ppppp	rrrrr	Z0	ddddddddddddd	Z
Write	1 1	01	01	ppppp	rrrrr	10	ddddddddddddd	Z

The management frame consists of eight fields:

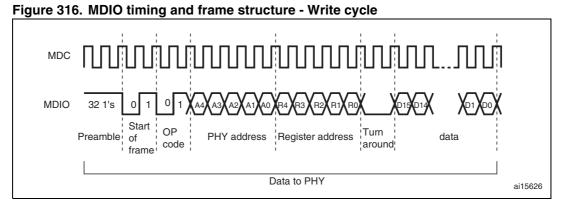
- **Preamble**: each transaction (read or write) can be initiated with the preamble field that corresponds to 32 contiguous logic one bits on the MDIO line with 32 corresponding cycles on MDC. This field is used to establish synchronization with the PHY device.
- **Start**: the start of frame is defined by a <01> pattern to verify transitions on the line from the default logic one state to zero and back to one.
- **Operation**: defines the type of transaction (read or write) in progress.
- PADDR: the PHY address is 5 bits, allowing 32 unique PHY addresses. The MSB bit of the address is the first transmitted and received.
- RADDR: the register address is 5 bits, allowing 32 individual registers to be addressed within the selected PHY device. The MSB bit of the address is the first transmitted and received.
- TA: the turn-around field defines a 2-bit pattern between the RADDR and DATA fields to avoid contention during a read transaction. For a read transaction the MAC controller drives high-impedance on the MDIO line for the 2 bits of TA. The PHY device must drive a high-impedance state on the first bit of TA, a zero bit on the second one. For a write transaction, the MAC controller drives a <10> pattern during the TA field. The PHY device must drive a high-impedance state for the 2 bits of TA.
- Data: the data field is 16-bit. The first bit transmitted and received must be bit 15 of the ETH\_MIID register.

• **Idle**: the MDIO line is driven in high-impedance state. All three-state drivers must be disabled and the PHY's pull-up resistor keeps the line at logic one.

## **SMI** write operation

When the application sets the MII Write and Busy bits (in *Ethernet MAC MII address register (ETH\_MACMIIAR)*), the SMI initiates a write operation into the PHY registers by transferring the PHY address, the register address in PHY, and the write data (in *Ethernet MAC MII data register (ETH\_MACMIIDR)*). The application should not change the MII Address register contents or the MII Data register while the transaction is ongoing. Write operations to the MII Address register or the MII Data Register during this period are ignored (the Busy bit is high), and the transaction is completed without any error. After the Write operation has completed, the SMI indicates this by resetting the Busy bit.

*Figure 316* shows the frame format for the write operation.



### **SMI** read operation

When the user sets the MII Busy bit in the Ethernet MAC MII address register (ETH\_MACMIIAR) with the MII Write bit at 0, the SMI initiates a read operation in the PHY registers by transferring the PHY address and the register address in PHY. The application should not change the MII Address register contents or the MII Data register while the transaction is ongoing. Write operations to the MII Address register or MII Data Register during this period are ignored (the Busy bit is high) and the transaction is completed without any error. After the read operation has completed, the SMI resets the Busy bit and then updates the MII Data register with the data read from the PHY.

*Figure 317* shows the frame format for the read operation.

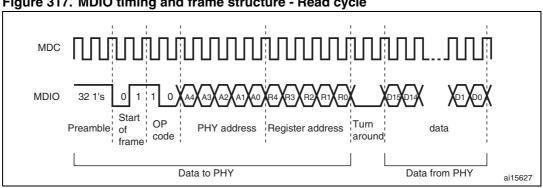


Figure 317. MDIO timing and frame structure - Read cycle

#### SMI clock selection

The MAC initiates the Management Write/Read operation. The SMI clock is a divided clock whose source is the application clock (AHB clock). The divide factor depends on the clock range setting in the MII Address register.

Table 144 shows how to set the clock ranges.

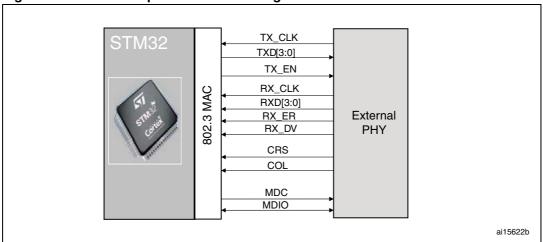
Table 144. Clock range

Selection	HCLK clock	MDC clock
000	60-100 MHz	AHB clock / 42
001	100-168 MHz	AHB clock / 62
010	20-35 MHz	AHB clock / 16
011	35-60 MHz	AHB clock / 26
100, 101, 110, 111	Reserved	-

## 28.4.2 Media-independent interface: MII

The media-independent interface (MII) defines the interconnection between the MAC sublayer and the PHY for data transfer at 10 Mbit/s and 100 Mbit/s.

Figure 318. Media independent interface signals



- MII\_TX\_CLK: continuous clock that provides the timing reference for the TX data transfer. The nominal frequency is: 2.5 MHz at 10 Mbit/s speed; 25 MHz at 100 Mbit/s speed.
- MII\_RX\_CLK: continuous clock that provides the timing reference for the RX data transfer. The nominal frequency is: 2.5 MHz at 10 Mbit/s speed; 25 MHz at 100 Mbit/s speed.
- MII\_TX\_EN: transmission enable indicates that the MAC is presenting nibbles on the MII for transmission. It must be asserted synchronously (MII\_TX\_CLK) with the first nibble of the preamble and must remain asserted while all nibbles to be transmitted are presented to the MII.
- MII\_TXD[3:0]: transmit data is a bundle of 4 data signals driven synchronously by the MAC sublayer and qualified (valid data) on the assertion of the MII\_TX\_EN signal.

- MII\_TXD[0] is the least significant bit, MII\_TXD[3] is the most significant bit. While MII\_TX\_EN is deasserted the transmit data must have no effect upon the PHY.
- MII\_CRS: carrier sense is asserted by the PHY when either the transmit or receive medium is non idle. It shall be deasserted by the PHY when both the transmit and receive media are idle. The PHY must ensure that the MII\_CS signal remains asserted throughout the duration of a collision condition. This signal is not required to transition synchronously with respect to the TX and RX clocks. In full duplex mode the state of this signal is don't care for the MAC sublayer.
- MII\_COL: collision detection must be asserted by the PHY upon detection of a collision
  on the medium and must remain asserted while the collision condition persists. This
  signal is not required to transition synchronously with respect to the TX and RX clocks.
  In full duplex mode the state of this signal is don't care for the MAC sublayer.
- MII\_RXD[3:0]: reception data is a bundle of 4 data signals driven synchronously by the PHY and qualified (valid data) on the assertion of the MII\_RX\_DV signal. MII\_RXD[0] is the least significant bit, MII\_RXD[3] is the most significant bit. While MII\_RX\_EN is deasserted and MII\_RX\_ER is asserted, a specific MII\_RXD[3:0] value is used to transfer specific information from the PHY (see *Table 146*).
- MII\_RX\_DV: receive data valid indicates that the PHY is presenting recovered and decoded nibbles on the MII for reception. It must be asserted synchronously (MII\_RX\_CLK) with the first recovered nibble of the frame and must remain asserted through the final recovered nibble. It must be deasserted prior to the first clock cycle that follows the final nibble. In order to receive the frame correctly, the MII\_RX\_DV signal must encompass the frame, starting no later than the SFD field.
- MII\_RX\_ER: receive error must be asserted for one or more clock periods (MII\_RX\_CLK) to indicate to the MAC sublayer that an error was detected somewhere in the frame. This error condition must be qualified by MII\_RX\_DV assertion as described in *Table 146*.

Table 145. TX interface signal encoding

MII_TX_EN	MII_TXD[3:0]	Description
0	0000 through 1111	Normal inter-frame
1	0000 through 1111	Normal data transmission

Table 146. RX interface signal encoding

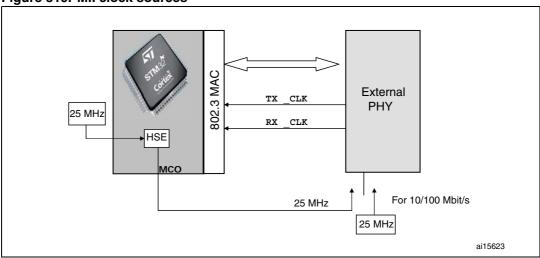
MII_RX_DV	MII_RX_ERR	MII_RXD[3:0]	Description	
0	0	0000 through 1111	Normal inter-frame	
0	1	0000	Normal inter-frame	
0	1	0001 through 1101	Reserved	
0	1	1110	False carrier indication	
0	1	1111	Reserved	
1	0	0000 through 1111	Normal data reception	
1	1	0000 through 1111	Data reception with errors	



#### MII clock sources

To generate both TX\_CLK and RX\_CLK clock signals, the external PHY must be clocked with an external 25 MHz as shown in *Figure 319*. Instead of using an external 25 MHz quartz to provide this clock, the STM32F4xxmicrocontroller can output this signal on its MCO pin. In this case, the PLL multiplier has to be configured so as to get the desired frequency on the MCO pin, from the 25 MHz external quartz.

Figure 319. MII clock sources



## 28.4.3 Reduced media-independent interface: RMII

The reduced media-independent interface (RMII) specification reduces the pin count between the microcontroller Ethernet peripheral and the external Ethernet in 10/100 Mbit/s. According to the IEEE 802.3u standard, an MII contains 16 pins for data and control. The RMII specification is dedicated to reduce the pin count to 7 pins (a 62.5% decrease in pin count).

The RMII is instantiated between the MAC and the PHY. This helps translation of the MAC's MII into the RMII. The RMII block has the following characteristics:

- It supports 10-Mbit/s and 100-Mbit/s operating rates
- The clock reference must be doubled to 50 MHz
- The same clock reference must be sourced externally to both MAC and external Ethernet PHY
- It provides independent 2-bit wide (dibit) transmit and receive data paths

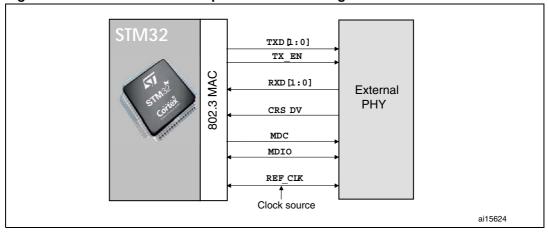
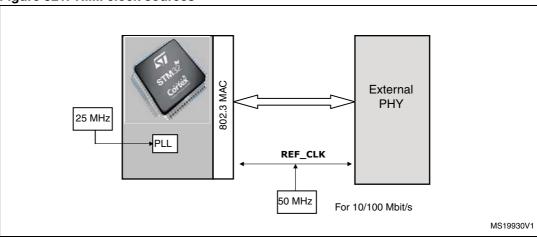


Figure 320. Reduced media-independent interface signals

#### RMII clock sources

Either clock the PHY from an external 50 MHz clock or use a PHY with an embedded PLL to generate the 50 MHz frequency.

Figure 321. RMII clock sources



## 28.4.4 MII/RMII selection

The mode, MII or RMII, is selected using the configuration bit 23, MII\_RMII\_SEL, in the SYSCFG\_PMC register. The application has to set the MII/RMII mode while the Ethernet controller is under reset or before enabling the clocks.

### MII/RMII internal clock scheme

The clock scheme required to support both the MII and RMII, as well as 10 and 100 Mbit/s operations is described in *Figure 322*.

MAC 25 MHz or 2.5 MHz GPIO and AF MII TX CLK as AF 25 MHz TX controller (25 MHz or 2.5 MHz) or 2.5 MHz 0 MII Sync. divider 50 MHz 1 RMII<sup>(1)</sup> /2 for 100 Mb/s /20 for 10 Mb/s MII\_RX\_CLK as AF 25 MHz RX 25 MHz or 2.5 MH (25 MHz or 2.5 MHz) 📮 GPIO and AF or 2.5 MHz controller RMII\_REF\_CK as AF (50 MHz) RMII HCI K HCI K must be greater than 25 MHz

Figure 322. Clock scheme

1. The MII/RMII selection is controlled through bit 23, MII\_RMII\_SEL, in the SYSCFG\_PMC register.

To save a pin, the two input clock signals, RMII\_REF\_CK and MII\_RX\_CLK, are multiplexed on the same GPIO pin.

# 28.5 Ethernet functional description: MAC 802.3

The IEEE 802.3 International Standard for local area networks (LANs) employs the CSMA/CD (carrier sense multiple access with collision detection) as the access method.

The Ethernet peripheral consists of a MAC 802.3 (media access control) controller with media independent interface (MII) and a dedicated DMA controller.

The MAC block implements the LAN CSMA/CD sublayer for the following families of systems: 10 Mbit/s and 100 Mbit/s of data rates for baseband and broadband systems. Half-and full-duplex operation modes are supported. The collision detection access method is applied only to the half-duplex operation mode. The MAC control frame sublayer is supported.

The MAC sublayer performs the following functions associated with a data link control procedure:

- Data encapsulation (transmit and receive)
  - Framing (frame boundary delimitation, frame synchronization)
  - Addressing (handling of source and destination addresses)
  - Error detection
- Media access management
  - Medium allocation (collision avoidance)
  - Contention resolution (collision handling)

Basically there are two operating modes of the MAC sublayer:

- Half-duplex mode: the stations contend for the use of the physical medium, using the CSMA/CD algorithms.
- Full duplex mode: simultaneous transmission and reception without contention resolution (CSMA/CD algorithm are unnecessary) when all the following conditions are met:
  - physical medium capability to support simultaneous transmission and reception
  - exactly 2 stations connected to the LAN
  - both stations configured for full-duplex operation

#### 28.5.1 MAC 802.3 frame format

The MAC block implements the MAC sublayer and the optional MAC control sublayer (10/100 Mbit/s) as specified by the IEEE 802.3-2002 standard.

Two frame formats are specified for data communication systems using the CSMA/CD MAC:

- Basic MAC frame format
- Tagged MAC frame format (extension of the basic MAC frame format)

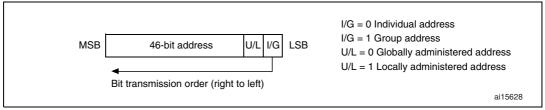
*Figure 324* and *Figure 325* describe the frame structure (untagged and tagged) that includes the following fields:

- Preamble: 7-byte field used for synchronization purposes (PLS circuitry)
   Hexadecimal value: 55-55-55-55-55-55
  - Bit pattern: 01010101 01010101 01010101 01010101 01010101 01010101 01010101 (right-to-left bit transmission)
- Start frame delimiter (SFD): 1-byte field used to indicate the start of a frame.
   Hexadecimal value: D5
  - Bit pattern: 11010101 (right-to-left bit transmission)
- Destination and Source Address fields: 6-byte fields to indicate the destination and source station addresses as follows (see Figure 323):
  - Each address is 48 bits in length
  - The first LSB bit (I/G) in the destination address field is used to indicate an individual (I/G = 0) or a group address (I/G = 1). A group address could identify none, one or more, or all the stations connected to the LAN. In the source address the first bit is reserved and reset to 0.
  - The second bit (U/L) distinguishes between locally (U/L = 1) or globally (U/L = 0) administered addresses. For broadcast addresses this bit is also 1.
  - Each byte of each address field must be transmitted least significant bit first.

The address designation is based on the following types:

- Individual address: this is the physical address associated with a particular station on the network.
- Group address. A multidestination address associated with one or more stations on a given network. There are two kinds of multicast address:
  - Multicast-group address: an address associated with a group of logically related stations.
  - Broadcast address: a distinguished, predefined multicast address (all 1's in the destination address field) that always denotes all the stations on a given LAN.

Figure 323. Address field format



- QTag Prefix: 4-byte field inserted between the Source address field and the MAC Client Length/Type field. This field is an extension of the basic frame (untagged) to obtain the tagged MAC frame. The untagged MAC frames do not include this field. The extensions for tagging are as follows:
  - 2-byte constant Length/Type field value consistent with the Type interpretation (greater than 0x0600) equal to the value of the 802.1Q Tag Protocol Type (0x8100 hexadecimal). This constant field is used to distinguish tagged and untagged MAC frames.
  - 2-byte field containing the Tag control information field subdivided as follows: a 3-bit user priority, a canonical format indicator (CFI) bit and a 12-bit VLAN Identifier.
     The length of the tagged MAC frame is extended by 4 bytes by the QTag Prefix.
- MAC client length/type: 2-byte field with different meaning (mutually exclusive), depending on its value:
  - If the value is less than or equal to maxValidFrame (0d1500) then this field indicates the number of MAC client data bytes contained in the subsequent data field of the 802.3 frame (length interpretation).
  - If the value is greater than or equal to MinTypeValue (0d1536 decimal, 0x0600) then this field indicates the nature of the MAC client protocol (Type interpretation) related to the Ethernet frame.

Regardless of the interpretation of the length/type field, if the length of the data field is less than the minimum required for proper operation of the protocol, a PAD field is added after the data field but prior to the FCS (frame check sequence) field. The length/type field is transmitted and received with the higher-order byte first.

For length/type field values in the range between maxValidLength and minTypeValue (boundaries excluded), the behavior of the MAC sublayer is not specified: they may or may not be passed by the MAC sublayer.

- Data and PAD fields: n-byte data field. Full data transparency is provided, it means that
  any arbitrary sequence of byte values may appear in the data field. The size of the PAD,
  if any, is determined by the size of the data field. Max and min length of the data and
  PAD field are:
  - Maximum length = 1500 bytes
  - Minimum length for untagged MAC frames = 46 bytes
  - Minimum length for tagged MAC frames = 42 bytes

When the data field length is less than the minimum required, the PAD field is added to match the minimum length (42 bytes for tagged frames, 46 bytes for untagged frames).

• Frame check sequence: 4-byte field that contains the cyclic redundancy check (CRC) value. The CRC computation is based on the following fields: source address,

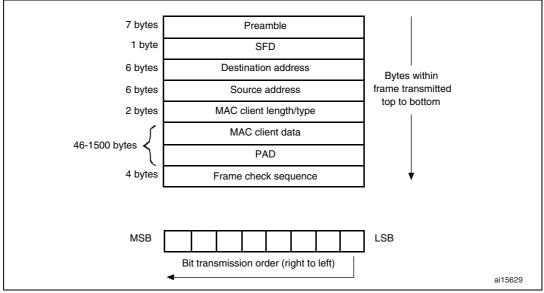
destination address, QTag prefix, length/type, LLC data and PAD (that is, all fields except the preamble, SFD). The generating polynomial is the following:

$$G(x) = x^{32} + x^{26} + x^{23} + x^{22} + x^{16} + x^{12} + x^{11} + x^{10} + x^{8} + x^{7} + x^{5} + x^{4} + x^{2} + x + 1$$

The CRC value of a frame is computed as follows:

- The first 2 bits of the frame are complemented
- The n-bits of the frame are the coefficients of a polynomial M(x) of degree (n-1). The first bit of the destination address corresponds to the  $x^{n-1}$  term and the last bit of the data field corresponds to the  $x^0$  term
- M(x) is multiplied by x<sup>32</sup> and divided by G(x), producing a remainder R(x) of degree
   31
- The coefficients of R(x) are considered as a 32-bit sequence
- The bit sequence is complemented and the result is the CRC
- The 32-bits of the CRC value are placed in the frame check sequence. The  $x^{32}$  term is the first transmitted, the  $x^0$  term is the last one

Figure 324. MAC frame format



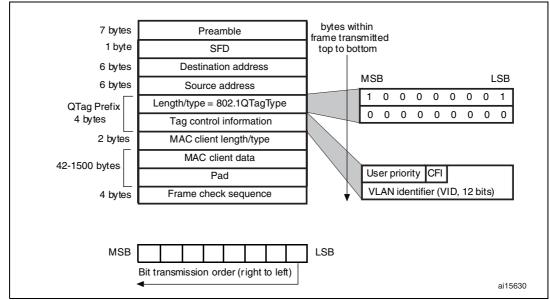


Figure 325. Tagged MAC frame format

Each byte of the MAC frame, except the FCS field, is transmitted low-order bit first.

An invalid MAC frame is defined by one of the following conditions:

- The frame length is inconsistent with the expected value as specified by the length/type field. If the length/type field contains a type value, then the frame length is assumed to be consistent with this field (no invalid frame)
- The frame length is not an integer number of bytes (extra bits)
- The CRC value computed on the incoming frame does not match the included FCS

#### 28.5.2 **MAC frame transmission**

The DMA controls all transactions for the transmit path. Ethernet frames read from the system memory are pushed into the FIFO by the DMA. The frames are then popped out and transferred to the MAC core. When the end-of-frame is transferred, the status of the transmission is taken from the MAC core and transferred back to the DMA. The Transmit FIFO has a depth of 2 Kbyte. FIFO-fill level is indicated to the DMA so that it can initiate a data fetch in required bursts from the system memory, using the AHB interface. The data from the AHB Master interface is pushed into the FIFO.

When the SOF is detected, the MAC accepts the data and begins transmitting to the MII. The time required to transmit the frame data to the MII after the application initiates transmission is variable, depending on delay factors like IFG delay, time to transmit preamble/SFD, and any back-off delays for Half-duplex mode. After the EOF is transferred to the MAC core, the core completes normal transmission and then gives the status of transmission back to the DMA. If a normal collision (in Half-duplex mode) occurs during transmission, the MAC core makes the transmit status valid, then accepts and drops all further data until the next SOF is received. The same frame should be retransmitted from SOF on observing a Retry request (in the Status) from the MAC. The MAC issues an underflow status if the data are not provided continuously during the transmission. During the normal transfer of a frame, if the MAC receives an SOF without getting an EOF for the previous frame, then the SOF is ignored and the new frame is considered as the continuation of the previous frame.



There are two modes of operation for popping data towards the MAC core:

- In Threshold mode, as soon as the number of bytes in the FIFO crosses the configured threshold level (or when the end-of-frame is written before the threshold is crossed), the data is ready to be popped out and forwarded to the MAC core. The threshold level is configured using the TTC bits of ETH DMABMR.
- In Store-and-forward mode, only after a complete frame is stored in the FIFO, the frame
  is popped towards the MAC core. If the Tx FIFO size is smaller than the Ethernet frame
  to be transmitted, then the frame is popped towards the MAC core when the Tx FIFO
  becomes almost full.

The application can flush the Transmit FIFO of all contents by setting the FTF (ETH\_DMAOMR register [20]) bit. This bit is self-clearing and initializes the FIFO pointers to the default state. If the FTF bit is set during a frame transfer to the MAC core, then transfer is stopped as the FIFO is considered to be empty. Hence an underflow event occurs at the MAC transmitter and the corresponding Status word is forwarded to the DMA.

### **Automatic CRC and pad generation**

When the number of bytes received from the application falls below 60 (DA+SA+LT+Data), zeros are appended to the transmitting frame to make the data length exactly 46 bytes to meet the minimum data field requirement of IEEE 802.3. The MAC can be programmed not to append any padding. The cyclic redundancy check (CRC) for the frame check sequence (FCS) field is calculated and appended to the data being transmitted. When the MAC is programmed to not append the CRC value to the end of Ethernet frames, the computed CRC is not transmitted. An exception to this rule is that when the MAC is programmed to append pads for frames (DA+SA+LT+Data) less than 60 bytes, CRC will be appended at the end of the padded frames.

The CRC generator calculates the 32-bit CRC for the FCS field of the Ethernet frame. The encoding is defined by the following polynomial.

$$G(x) = x^{32} + x^{26} + x^{23} + x^{22} + x^{16} + x^{12} + x^{11} + x^{10} + x^{8} + x^{7} + x^{5} + x^{4} + x^{2} + x + 1$$

### **Transmit protocol**

The MAC controls the operation of Ethernet frame transmission. It performs the following functions to meet the IEEE 802.3/802.3z specifications. It:

- generates the preamble and SFD
- generates the jam pattern in Half-duplex mode
- controls the Jabber timeout
- controls the flow for Half-duplex mode (back pressure)
- generates the transmit frame status
- contains time stamp snapshot logic in accordance with IEEE 1588

When a new frame transmission is requested, the MAC sends out the preamble and SFD, followed by the data. The preamble is defined as 7 bytes of 0b10101010 pattern, and the SFD is defined as 1 byte of 0b10101011 pattern. The collision window is defined as 1 slot time (512 bit times for 10/100 Mbit/s Ethernet). The jam pattern generation is applicable only to Half-duplex mode, not to Full-duplex mode.

In MII mode, if a collision occurs at any time from the beginning of the frame to the end of the CRC field, the MAC sends a 32-bit jam pattern of 0x5555 5555 on the MII to inform all

other stations that a collision has occurred. If the collision is seen during the preamble transmission phase, the MAC completes the transmission of the preamble and SFD and then sends the jam pattern.

A jabber timer is maintained to cut off the transmission of Ethernet frames if more than 2048 (default) bytes have to be transferred. The MAC uses the deferral mechanism for flow control (back pressure) in Half-duplex mode. When the application requests to stop receiving frames, the MAC sends a JAM pattern of 32 bytes whenever it senses the reception of a frame, provided that transmit flow control is enabled. This results in a collision and the remote station backs off. The application requests flow control by setting the BPA bit (bit 0) in the ETH\_MACFCR register. If the application requests a frame to be transmitted, then it is scheduled and transmitted even when back pressure is activated. Note that if back pressure is kept activated for a long time (and more than 16 consecutive collision events occur) then the remote stations abort their transmissions due to excessive collisions. If IEEE 1588 time stamping is enabled for the transmit frame, this block takes a snapshot of the system time when the SFD is put onto the transmit MII bus.

#### Transmit scheduler

The MAC is responsible for scheduling the frame transmission on the MII. It maintains the interframe gap between two transmitted frames and follows the truncated binary exponential backoff algorithm for Half-duplex mode. The MAC enables transmission after satisfying the IFG and backoff delays. It maintains an idle period of the configured interframe gap (IFG bits in the ETH\_MACCR register) between any two transmitted frames. If frames to be transmitted arrive sooner than the configured IFG time, the MII waits for the enable signal from the MAC before starting the transmission on it. The MAC starts its IFG counter as soon as the carrier signal of the MII goes inactive. At the end of the programmed IFG value, the MAC enables transmission in Full-duplex mode. In Half-duplex mode and when IFG is configured for 96 bit times, the MAC follows the rule of deference specified in Section 4.2.3.2.1 of the IEEE 802.3 specification. The MAC resets its IFG counter if a carrier is detected during the first two-thirds (64-bit times for all IFG values) of the IFG interval. If the carrier is detected during the final one third of the IFG interval, the MAC continues the IFG count and enables the transmitter after the IFG interval. The MAC implements the truncated binary exponential backoff algorithm when it operates in Half-duplex mode.

### **Transmit flow control**

When the Transmit Flow Control Enable bit (TFE bit in ETH\_MACFCR) is set, the MAC generates Pause frames and transmits them as necessary, in Full-duplex mode. The Pause frame is appended with the calculated CRC, and is sent. Pause frame generation can be initiated in two ways.

A pause frame is sent either when the application sets the FCB bit in the ETH\_MACFCR register or when the receive FIFO is full (packet buffer).

- If the application has requested flow control by setting the FCB bit in ETH\_MACFCR, the MAC generates and transmits a single Pause frame. The value of the pause time in the generated frame contains the programmed pause time value in ETH\_MACFCR. To extend the pause or end the pause prior to the time specified in the previously transmitted Pause frame, the application must request another Pause frame transmission after programming the Pause Time value (PT in ETH\_MACFCR register) with the appropriate value.
- If the application has requested flow control when the receive FIFO is full, the MAC generates and transmits a Pause frame. The value of the pause time in the generated frame is the programmed pause time value in ETH\_MACFCR. If the receive FIFO

remains full at a configurable number of slot-times (PLT bits in ETH\_MACFCR) before this Pause time runs out, a second Pause frame is transmitted. The process is repeated as long as the receive FIFO remains full. If this condition is no more satisfied prior to the sampling time, the MAC transmits a Pause frame with zero pause time to indicate to the remote end that the receive buffer is ready to receive new data frames.

### Single-packet transmit operation

The general sequence of events for a transmit operation is as follows:

- If the system has data to be transferred, the DMA controller fetches them from the memory through the AHB Master interface and starts forwarding them to the FIFO. It continues to receive the data until the end of frame is transferred.
- When the threshold level is crossed or a full packet of data is received into the FIFO, the frame data are popped and driven to the MAC core. The DMA continues to transfer data from the FIFO until a complete packet has been transferred to the MAC. Upon completion of the frame, the DMA controller is notified by the status coming from the MAC.

### Transmit operation—Two packets in the buffer

- Because the DMA must update the descriptor status before releasing it to the Host, there can be at the most two frames inside a transmit FIFO. The second frame is fetched by the DMA and put into the FIFO only if the OSF (operate on second frame) bit is set. If this bit is not set, the next frame is fetched from the memory only after the MAC has completely processed the frame and the DMA has released the descriptors.
- 2. If the OSF bit is set, the DMA starts fetching the second frame immediately after completing the transfer of the first frame to the FIFO. It does not wait for the status to be updated. In the meantime, the second frame is received into the FIFO while the first frame is being transmitted. As soon as the first frame has been transferred and the status is received from the MAC, it is pushed to the DMA. If the DMA has already completed sending the second packet to the FIFO, the second transmission must wait for the status of the first packet before proceeding to the next frame.

### Retransmission during collision

While a frame is being transferred to the MAC, a collision event may occur on the MAC line interface in Half-duplex mode. The MAC would then indicate a retry attempt by giving the status even before the end of frame is received. Then the retransmission is enabled and the frame is popped out again from the FIFO. After more than 96 bytes have been popped towards the MAC core, the FIFO controller frees up that space and makes it available to the DMA to push in more data. This means that the retransmission is not possible after this threshold is crossed or when the MAC core indicates a late collision event.

### **Transmit FIFO flush operation**

The MAC provides a control to the software to flush the Transmit FIFO through the use of Bit 20 in the Operation mode register. The Flush operation is immediate and the Tx FIFO and the corresponding pointers are cleared to the initial state even if the Tx FIFO is in the middle of transferring a frame to the MAC Core. This results in an underflow event in the MAC transmitter, and the frame transmission is aborted. The status of such a frame is marked with both underflow and frame flush events (TDES0 bits 13 and 1). No data are coming to the FIFO from the application (DMA) during the Flush operation. Transfer transmit status words are transferred to the application for the number of frames that is flushed (including partial frames). Frames that are completely flushed have the Frame flush status bit (TDES0 13) set. The Flush operation is completed when the application (DMA) has accepted all of



the Status words for the frames that were flushed. The Transmit FIFO Flush control register bit is then cleared. At this point, new frames from the application (DMA) are accepted. All data presented for transmission after a Flush operation are discarded unless they start with an SOF marker.

#### Transmit status word

At the end of the Ethernet frame transfer to the MAC core and after the core has completed the transmission of the frame, the transmit status is given to the application. The detailed description of the Transmit Status is the same as for bits [23:0] in TDES0. If IEEE 1588 time stamping is enabled, a specific frames' 64-bit time stamp is returned, along with the transmit status.

#### Transmit checksum offload

Communication protocols such as TCP and UDP implement checksum fields, which helps determine the integrity of data transmitted over a network. Because the most widespread use of Ethernet is to encapsulate TCP and UDP over IP datagrams, the Ethernet controller has a transmit checksum offload feature that supports checksum calculation and insertion in the transmit path, and error detection in the receive path. This section explains the operation of the checksum offload feature for transmitted frames.

Note:

The checksum for TCP, UDP or ICMP is calculated over a complete frame, then inserted into its corresponding header field. Due to this requirement, this function is enabled only when the Transmit FIFO is configured for Store-and-forward mode (that is, when the TSF bit is set in the ETH\_ETH\_DMAOMR register). If the core is configured for Threshold (cut-through) mode, the Transmit checksum offload is bypassed.

You must make sure the Transmit FIFO is deep enough to store a complete frame before that frame is transferred to the MAC Core transmitter. If the FIFO depth is less than the input Ethernet frame size, the payload (TCP/UDP/ICMP) checksum insertion function is bypassed and only the frame's IPv4 Header checksum is modified, even in Store-and-forward mode.

The transmit checksum offload supports two types of checksum calculation and insertion. This checksum can be controlled for each frame by setting the CIC bits (Bits 28:27 in TDES1, described in *TDES1: Transmit descriptor Word1 on page 880*).

See IETF specifications RFC 791, RFC 793, RFC 768, RFC 792, RFC 2460 and RFC 4443 for IPv4, TCP, UDP, ICMP, IPv6 and ICMPv6 packet header specifications, respectively.

### IP header checksum

In IPv4 datagrams, the integrity of the header fields is indicated by the 16-bit header checksum field (the eleventh and twelfth bytes of the IPv4 datagram). The checksum offload detects an IPv4 datagram when the Ethernet frame's Type field has the value 0x0800 and the IP datagram's Version field has the value 0x4. The input frame's checksum field is ignored during calculation and replaced by the calculated value. IPv6 headers do not have a checksum field; thus, the checksum offload does not modify IPv6 header fields. The result of this IP header checksum calculation is indicated by the IP Header Error status bit in the Transmit status (Bit 16). This status bit is set whenever the values of the Ethernet Type field and the IP header's Version field are not consistent, or when the Ethernet frame does not have enough data, as indicated by the

IP header Length field. In other words, this bit is set when an IP header error is asserted under the following circumstances:

- a) For IPv4 datagrams:
- The received Ethernet type is 0x0800, but the IP header's Version field does not equal 0x4
- The IPv4 Header Length field indicates a value less than 0x5 (20 bytes)
- The total frame length is less than the value given in the IPv4 Header Length field
- b) For IPv6 datagrams:
- The Ethernet type is 0x86DD but the IP header Version field does not equal 0x6
- The frame ends before the IPv6 header (40 bytes) or extension header (as given in the corresponding Header Length field in an extension header) has been completely received. Even when the checksum offload detects such an IP header error, it inserts an IPv4 header checksum if the Ethernet Type field indicates an IPv4 payload.

### • TCP/UDP/ICMP checksum

The TCP/UDP/ICMP checksum processes the IPv4 or IPv6 header (including extension headers) and determines whether the encapsulated payload is TCP, UDP or ICMP.

Note that:

- a) For non-TCP, -UDP, or -ICMP/ICMPv6 payloads, this checksum is bypassed and nothing further is modified in the frame.
- b) Fragmented IP frames (IPv4 or IPv6), IP frames with security features (such as an authentication header or encapsulated security payload), and IPv6 frames with routing headers are bypassed and not processed by the checksum.

The checksum is calculated for the TCP, UDP, or ICMP payload and inserted into its corresponding field in the header. It can work in the following two modes:

- In the first mode, the TCP, UDP, or ICMPv6 pseudo-header is not included in the checksum calculation and is assumed to be present in the input frame's checksum field. The checksum field is included in the checksum calculation, and then replaced by the final calculated checksum.
- In the second mode, the checksum field is ignored, the TCP, UDP, or ICMPv6 pseudo-header data are included into the checksum calculation, and the checksum field is overwritten with the final calculated value.

Note that: for ICMP-over-IPv4 packets, the checksum field in the ICMP packet must always be 0x0000 in both modes, because pseudo-headers are not defined for such packets. If it does not equal 0x0000, an incorrect checksum may be inserted into the packet.

The result of this operation is indicated by the payload checksum error status bit in the Transmit Status vector (bit 12). The payload checksum error status bit is set when either of the following is detected:

- the frame has been forwarded to the MAC transmitter in Store-and-forward mode without the end of frame being written to the FIFO
- the packet ends before the number of bytes indicated by the payload length field in the IP header is received.

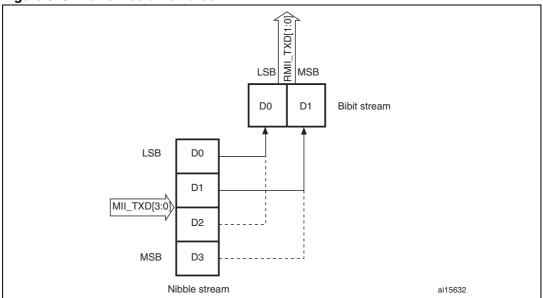
When the packet is longer than the indicated payload length, the bytes are ignored as stuff bytes, and no error is reported. When the first type of error is detected, the TCP,

UDP or ICMP header is not modified. For the second error type, still, the calculated checksum is inserted into the corresponding header field.

### MII/RMII transmit bit order

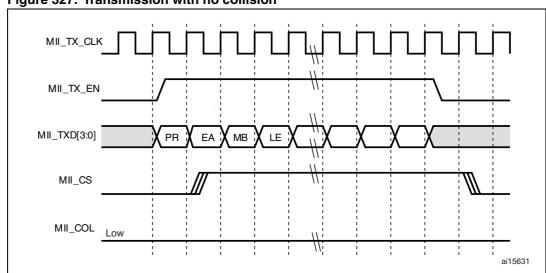
Each nibble from the MII is transmitted on the RMII a dibit at a time with the order of dibit transmission shown in *Figure 326*. Lower order bits (D1 and D0) are transmitted first followed by higher order bits (D2 and D3).

Figure 326. Transmission bit order



### MII/RMII transmit timing diagrams

Figure 327. Transmission with no collision



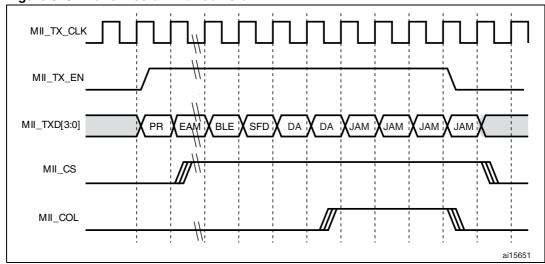


Figure 328. Transmission with collision

Figure 329 shows a frame transmission in MII and RMII.

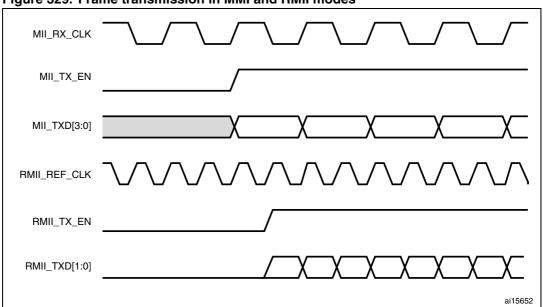


Figure 329. Frame transmission in MMI and RMII modes

## 28.5.3 MAC frame reception

The MAC received frames are pushes into the Rx FIFO. The status (fill level) of this FIFO is indicated to the DMA once it crosses the configured receive threshold (RTC in the ETH\_DMAOMR register) so that the DMA can initiate pre-configured burst transfers towards the AHB interface.

In the default Cut-through mode, when 64 bytes (configured with the RTC bits in the ETH\_DMAOMR register) or a full packet of data are received into the FIFO, the data are popped out and the DMA is notified of its availability. Once the DMA has initiated the transfer to the AHB interface, the data transfer continues from the FIFO until a complete packet has

been transferred. Upon completion of the EOF frame transfer, the status word is popped out and sent to the DMA controller.

In Rx FIFO Store-and-forward mode (configured by the RSF bit in the ETH\_DMAOMR register), a frame is read out only after being written completely into the Receive FIFO. In this mode, all error frames are dropped (if the core is configured to do so) such that only valid frames are read out and forwarded to the application. In Cut-through mode, some error frames are not dropped, because the error status is received at the end of the frame, by which time the start of that frame has already been read out of the FIFO.

A receive operation is initiated when the MAC detects an SFD on the MII. The core strips the preamble and SFD before proceeding to process the frame. The header fields are checked for the filtering and the FCS field used to verify the CRC for the frame. The frame is dropped in the core if it fails the address filter.

### Receive protocol

The received frame preamble and SFD are stripped. Once the SFD has been detected, the MAC starts sending the Ethernet frame data to the receive FIFO, beginning with the first byte following the SFD (destination address). If IEEE 1588 time stamping is enabled, a snapshot of the system time is taken when any frame's SFD is detected on the MII. Unless the MAC filters out and drops the frame, this time stamp is passed on to the application.

If the received frame length/type field is less than 0x600 and if the MAC is programmed for the auto CRC/pad stripping option, the MAC sends the data of the frame to RxFIFO up to the count specified in the length/type field, then starts dropping bytes (including the FCS field). If the Length/Type field is greater than or equal to 0x600, the MAC sends all received Ethernet frame data to Rx FIFO, regardless of the value on the programmed auto-CRC strip option. The MAC watchdog timer is enabled by default, that is, frames above 2048 bytes (DA + SA + LT + Data + pad + FCS) are cut off. This feature can be disabled by programming the watchdog disable (WD) bit in the MAC configuration register. However, even if the watchdog timer is disabled, frames greater than 16 KB in size are cut off and a watchdog timeout status is given.

### Receive CRC: automatic CRC and pad stripping

The MAC checks for any CRC error in the receiving frame. It calculates the 32-bit CRC for the received frame that includes the Destination address field through the FCS field. The encoding is defined by the following polynomial.

$$G(x) = x^{32} + x^{26} + x^{23} + x^{22} + x^{16} + x^{12} + x^{11} + x^{10} + x^{8} + x^{7} + x^{5} + x^{4} + x^{2} + x + 1$$

Regardless of the auto-pad/CRC strip, the MAC receives the entire frame to compute the CRC check for the received frame.

### Receive checksum offload

Both IPv4 and IPv6 frames in the received Ethernet frames are detected and processed for data integrity. You can enable the receive checksum offload by setting the IPCO bit in the ETH\_MACCR register. The MAC receiver identifies IPv4 or IPv6 frames by checking for value 0x0800 or 0x86DD, respectively, in the received Ethernet frame Type field. This identification applies to VLAN-tagged frames as well. The receive checksum offload calculates IPv4 header checksums and checks that they match the received IPv4 header checksums. The IP Header Error bit is set for any mismatch between the indicated payload



type (Ethernet Type field) and the IP header version, or when the received frame does not have enough bytes, as indicated by the IPv4 header's Length field (or when fewer than 20 bytes are available in an IPv4 or IPv6 header). The receive checksum offload also identifies a TCP, UDP or ICMP payload in the received IP datagrams (IPv4 or IPv6) and calculates the checksum of such payloads properly, as defined in the TCP, UDP or ICMP specifications. It includes the TCP/UDP/ICMPv6 pseudo-header bytes for checksum calculation and checks whether the received checksum field matches the calculated value. The result of this operation is given as a Payload Checksum Error bit in the receive status word. This status bit is also set if the length of the TCP, UDP or ICMP payload does not match the expected payload length given in the IP header. As mentioned in TCP/UDP/ICMP checksum on page 848, the receive checksum offload bypasses the payload of fragmented IP datagrams, IP datagrams with security features, IPv6 routing headers, and payloads other than TCP, UDP or ICMP. This information (whether the checksum is bypassed or not) is given in the receive status, as described in the RDES0: Receive descriptor Word0 section. In this configuration, the core does not append any payload checksum bytes to the received Ethernet frames.

As mentioned in *RDES0: Receive descriptor Word0 on page 886*, the meaning of certain register bits changes as shown in *Table 147*.

Table 147. Frame statuses

Bit 18: Ethernet frame	Bit 27: Header checksum error	Bit 28: Payload checksum error	Frame status
0	0	0	The frame is an IEEE 802.3 frame (Length field value is less than 0x0600).
1	0	0	IPv4/IPv6 Type frame in which no checksum error is detected.
1	0	1	IPv4/IPv6 Type frame in which a payload checksum error (as described for PCE) is detected
1	1	0	IPv4/IPv6 Type frame in which IP header checksum error (as described for IPCO HCE) is detected.
1	1	1	IPv4/IPv6 Type frame in which both PCE and IPCO HCE are detected.
0	0	1	IPv4/IPv6 Type frame in which there is no IP HCE and the payload check is bypassed due to unsupported payload.
0	1	1	Type frame which is neither IPv4 or IPv6 (checksum offload bypasses the checksum check completely)
0	1	0	Reserved

#### Receive frame controller

If the RA bit is reset in the MAC CSR frame filter register, the MAC performs frame filtering based on the destination/source address (the application still needs to perform another level of filtering if it decides not to receive any bad frames like runt, CRC error frames, etc.). On detecting a filter-fail, the frame is dropped and not transferred to the application. When the filtering parameters are changed dynamically, and in case of (DA-SA) filter-fail, the rest of

the frame is dropped and the Rx Status Word is immediately updated (with zero frame length, CRC error and Runt Error bits set), indicating the filter fail. In Ethernet power down mode, all received frames are dropped, and are not forwarded to the application.

#### Receive flow control

The MAC detects the receiving Pause frame and pauses the frame transmission for the delay specified within the received Pause frame (only in Full-duplex mode). The Pause frame detection function can be enabled or disabled with the RFCE bit in ETH MACFCR. Once receive flow control has been enabled, the received frame destination address begins to be monitored for any match with the multicast address of the control frame (0x0180 C200 0001). If a match is detected (the destination address of the received frame matches the reserved control frame destination address), the MAC then decides whether or not to transfer the received control frame to the application, based on the level of the PCF bit in ETH\_MACFFR.

The MAC also decodes the type, opcode, and Pause Timer fields of the receiving control frame. If the byte count of the status indicates 64 bytes, and if there is no CRC error, the MAC transmitter pauses the transmission of any data frame for the duration of the decoded Pause time value, multiplied by the slot time (64 byte times for both 10/100 Mbit/s modes). Meanwhile, if another Pause frame is detected with a zero Pause time value, the MAC resets the Pause time and manages this new pause request.

If the received control frame matches neither the type field (0x8808), the opcode (0x00001), nor the byte length (64 bytes), or if there is a CRC error, the MAC does not generate a Pause.

In the case of a pause frame with a multicast destination address, the MAC filters the frame based on the address match.

For a pause frame with a unicast destination address, the MAC filtering depends on whether the DA matched the contents of the MAC address 0 register and whether the UPDF bit in ETH\_MACFCR is set (detecting a pause frame even with a unicast destination address). The PCF register bits (bits [7:6] in ETH\_MACFFR) control filtering for control frames in addition to address filtering.

### Receive operation multiframe handling

Since the status is available immediately following the data, the FIFO is capable of storing any number of frames into it, as long as it is not full.

### Error handling

If the Rx FIFO is full before it receives the EOF data from the MAC, an overflow is declared and the whole frame is dropped, and the overflow counter in the (ETH\_DMAMFBOCR register) is incremented. The status indicates a partial frame due to overflow. The Rx FIFO can filter error and undersized frames, if enabled (using the FEF and FUGF bits in ETH DMAOMR).

If the Receive FIFO is configured to operate in Store-and-forward mode, all error frames can be filtered and dropped.

In Cut-through mode, if a frame's status and length are available when that frame's SOF is read from the Rx FIFO, then the complete erroneous frame can be dropped. The DMA can flush the error frame being read from the FIFO, by enabling the receive frame flash bit. The data transfer to the application (DMA) is then stopped and the rest of the frame is internally read and dropped. The next frame transfer can then be started, if available.



#### Receive status word

At the end of the Ethernet frame reception, the MAC outputs the receive status to the application (DMA). The detailed description of the receive status is the same as for bits[31:0] in RDES0, given in *RDES0: Receive descriptor Word0 on page 886*.

### Frame length interface

In case of switch applications, data transmission and reception between the application and MAC happen as complete frame transfers. The application layer should be aware of the length of the frames received from the ingress port in order to transfer the frame to the egress port. The MAC core provides the frame length of each received frame inside the status at the end of each frame reception.

Note: A frame length value of 0 is given for partial frames written into the Rx FIFO due to overflow.

#### MII/RMII receive bit order

Each nibble is transmitted to the MII from the dibit received from the RMII in the nibble transmission order shown in *Figure 330*. The lower-order bits (D0 and D1) are received first, followed by the higher-order bits (D2 and D3).

LSB D0 D1 Di-bit stream

LSB D0 D1

MII\_RXD[3:0] D2

Figure 330. Receive bit order

Nibble stream

ai15633

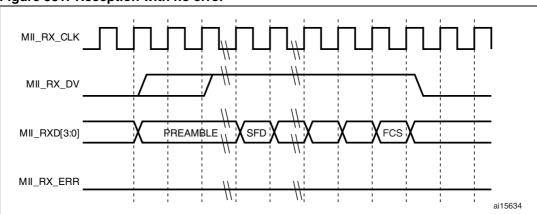


Figure 331. Reception with no error



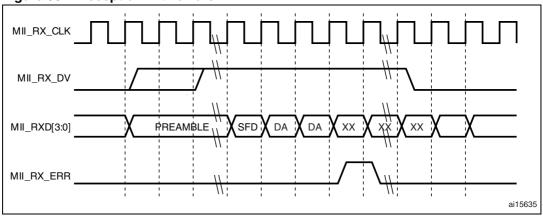
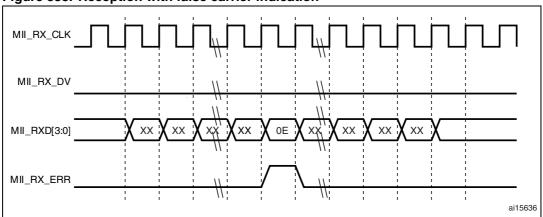


Figure 333. Reception with false carrier indication



#### 28.5.4 **MAC** interrupts

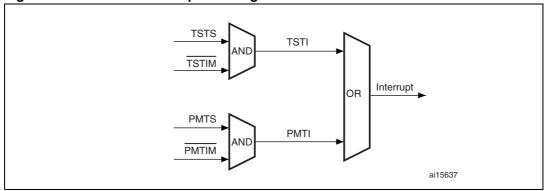
Interrupts can be generated from the MAC core as a result of various events.

The ETH\_MACSR register describes the events that can cause an interrupt from the MAC core. You can prevent each event from asserting the interrupt by setting the corresponding mask bits in the Interrupt Mask register.



The interrupt register bits only indicate the block from which the event is reported. You have to read the corresponding status registers and other registers to clear the interrupt. For example, bit 3 of the Interrupt register, set high, indicates that the Magic packet or Wake-on-LAN frame is received in Power-down mode. You must read the ETH\_MACPMTCSR Register to clear this interrupt event.

Figure 334. MAC core interrupt masking scheme



## 28.5.5 MAC filtering

### Address filtering

Address filtering checks the destination and source addresses on all received frames and the address filtering status is reported accordingly. Address checking is based on different parameters (Frame filter register) chosen by the application. The filtered frame can also be identified: multicast or broadcast frame.

Address filtering uses the station's physical (MAC) address and the Multicast Hash table for address checking purposes.

### Unicast destination address filter

The MAC supports up to 4 MAC addresses for unicast perfect filtering. If perfect filtering is selected (HU bit in the Frame filter register is reset), the MAC compares all 48 bits of the received unicast address with the programmed MAC address for any match. Default MacAddr0 is always enabled, other addresses MacAddr1–MacAddr3 are selected with an individual enable bit. Each byte of these other addresses (MacAddr1–MacAddr3) can be masked during comparison with the corresponding received DA byte by setting the corresponding Mask Byte Control bit in the register. This helps group address filtering for the DA. In Hash filtering mode (when HU bit is set), the MAC performs imperfect filtering for unicast addresses using a 64-bit Hash table. For hash filtering, the MAC uses the 6 upper CRC (see note 1 below) bits of the received destination address to index the content of the Hash table. A value of 000000 selects bit 0 in the selected register, and a value of 111111 selects bit 63 in the Hash Table register. If the corresponding bit (indicated by the 6-bit CRC) is set to 1, the unicast frame is said to have passed the Hash filter; otherwise, the frame has failed the Hash filter.

Note:

This CRC is a 32-bit value coded by the following polynomial (for more details refer to Section 28.5.3: MAC frame reception):

$$G(x) = x^{32} + x^{26} + x^{23} + x^{22} + x^{16} + x^{12} + x^{11} + x^{10} + x^{8} + x^{7} + x^{5} + x^{4} + x^{2} + x + 1$$

#### Multicast destination address filter

The MAC can be programmed to pass all multicast frames by setting the PAM bit in the Frame filter register. If the PAM bit is reset, the MAC performs the filtering for multicast addresses based on the HM bit in the Frame filter register. In Perfect filtering mode, the multicast address is compared with the programmed MAC destination address registers (1–3). Group address filtering is also supported. In Hash filtering mode, the MAC performs imperfect filtering using a 64-bit Hash table. For hash filtering, the MAC uses the 6 upper CRC (see note 1 below) bits of the received multicast address to index the content of the Hash table. A value of 000000 selects bit 0 in the selected register and a value of 111111 selects bit 63 in the Hash Table register. If the corresponding bit is set to 1, then the multicast frame is said to have passed the Hash filter; otherwise, the frame has failed the Hash filter.

Note:

This CRC is a 32-bit value coded by the following polynomial (for more details refer to Section 28.5.3: MAC frame reception):

$$G(x) \ = \ x^{32} + x^{26} + x^{23} + x^{22} + x^{16} + x^{12} + x^{11} + x^{10} + x^8 + x^7 + x^5 + x^4 + x^2 + x + 1$$

### Hash or perfect address filter

The DA filter can be configured to pass a frame when its DA matches either the Hash filter or the Perfect filter by setting the HPF bit in the Frame filter register and setting the corresponding HU or HM bits. This configuration applies to both unicast and multicast frames. If the HPF bit is reset, only one of the filters (Hash or Perfect) is applied to the received frame.

#### Broadcast address filter

The MAC does not filter any broadcast frames in the default mode. However, if the MAC is programmed to reject all broadcast frames by setting the BFD bit in the Frame filter register, any broadcast frames are dropped.

### Unicast source address filter

The MAC can also perform perfect filtering based on the source address field of the received frames. By default, the MAC compares the SA field with the values programmed in the SA registers. The MAC address registers [1:3] can be configured to contain SA instead of DA for comparison, by setting bit 30 in the corresponding register. Group filtering with SA is also supported. The frames that fail the SA filter are dropped by the MAC if the SAF bit in the Frame filter register is set. Otherwise, the result of the SA filter is given as a status bit in the Receive Status word (see *RDES0: Receive descriptor Word0*).

When the SAF bit is set, the result of the SA and DA filters is AND'ed to decide whether the frame needs to be forwarded. This means that either of the filter fail result will drop the frame. Both filters have to pass the frame for the frame to be forwarded to the application.

### Inverse filtering operation

For both destination and source address filtering, there is an option to invert the filter-match result at the final output. These are controlled by the DAIF and SAIF bits in the Frame filter register, respectively. The DAIF bit is applicable for both Unicast and Multicast DA frames. The result of the unicast/multicast destination address filter is inverted in this mode. Similarly, when the SAIF bit is set, the result of the unicast SA filter is inverted. *Table 148* 

and *Table 149* summarize destination and source address filtering based on the type of frame received.

Table 148. Destination address filtering

Frame type	РМ	HPF	HU	DAIF	нм	PAM	DB	DA filter operation
	1	Х	Х	Х	Х	Х	Χ	Pass
Broadcast	0	Х	Х	Х	Х	Х	0	Pass
	0	Х	Х	Х	Х	Х	1	Fail
	1	Х	Х	Х	Х	Х	Χ	Pass all frames
	0	Х	0	0	Х	Х	Χ	Pass on perfect/group filter match
	0	Х	0	1	Х	Х	Χ	Fail on perfect/Group filter match
Unicast	0	0	1	0	Х	Х	Χ	Pass on hash filter match
	0	0	1	1	Х	Х	Х	Fail on hash filter match
	0	1	1	0	Х	Х	х	Pass on hash or perfect/Group filter match
	0	1	1	1	Х	Х	Х	Fail on hash or perfect/Group filter match
	1	Х	Х	Х	Х	Х	Х	Pass all frames
	Χ	Х	Х	Х	Х	1	Х	Pass all frames
	0	Х	Х	0	0	0	Х	Pass on Perfect/Group filter match and drop PAUSE control frames if PCF = 0x
	0	0	Х	0	1	0	Х	Pass on hash filter match and drop PAUSE control frames if PCF = 0x
Multicast	0	1	х	0	1	0	х	Pass on hash or perfect/Group filter match and drop PAUSE control frames if PCF = 0x
	0	Х	Х	1	0	0	х	Fail on perfect/Group filter match and drop PAUSE control frames if PCF = 0x
	0	0	Х	1	1	0	Х	Fail on hash filter match and drop PAUSE control frames if PCF = 0x
	0	1	х	1	1	0	Х	Fail on hash or perfect/Group filter match and drop PAUSE control frames if PCF = 0x

SAIF SAF Frame type PΜ **SA filter operation** Χ Χ Pass all frames Pass status on perfect/Group filter match but do not drop 0 0 0 frames that fail Unicast 0 Fail status on perfect/group filter match but do not drop frame 0 0 Pass on perfect/group filter match and drop frames that fail 0 Fail on perfect/group filter match and drop frames that fail

Table 149. Source address filtering

## 28.5.6 MAC loopback mode

The MAC supports loopback of transmitted frames onto its receiver. By default, the MAC loopback function is disabled, but this feature can be enabled by programming the Loopback bit in the MAC ETH MACCR register.

## 28.5.7 MAC management counters: MMC

The MAC management counters (MMC) maintain a set of registers for gathering statistics on the received and transmitted frames. These include a control register for controlling the behavior of the registers, two 32-bit registers containing generated interrupts (receive and transmit), and two 32-bit registers containing masks for the Interrupt register (receive and transmit). These registers are accessible from the application. Each register is 32 bits wide.

Section 28.8: Ethernet register descriptions describes the various counters and lists the addresses of each of the statistics counters. This address is used for read/write accesses to the desired transmit/receive counter.

The Receive MMC counters are updated for frames that pass address filtering. Dropped frames statistics are not updated unless the dropped frames are runt frames of less than 6 bytes (DA bytes are not received fully).

#### Good transmitted and received frames

Transmitted frames are considered "good" if transmitted successfully. In other words, a transmitted frame is good if the frame transmission is not aborted due to any of the following errors:

- + Jabber Timeout
- + No Carrier/Loss of Carrier
- + Late Collision
- + Frame Underflow
- + Excessive Deferral
- + Excessive Collision

Received frames are considered "good" if none of the following errors exists:

- + CRC error
- + Runt Frame (shorter than 64 bytes)
- + Alignment error (in 10/100 Mbit/s only)
- + Length error (non-Type frames only)
- + Out of Range (non-Type frames only, longer than maximum size)
- + MII RXER Input error

The maximum frame size depends on the frame type, as follows:

- + Untagged frame maxsize = 1518
- + VLAN Frame maxsize = 1522

## 28.5.8 Power management: PMT

This section describes the power management (PMT) mechanisms supported by the MAC. PMT supports the reception of network (remote) wakeup frames and Magic Packet frames. PMT generates interrupts for wakeup frames and Magic Packets received by the MAC. The PMT block is enabled with remote wakeup frame enable and Magic Packet enable. These enable bits (WFE and MPE) are in the ETH\_MACPMTCSR register and are programmed by the application. When the power down mode is enabled in the PMT, then all received frames are dropped by the MAC and they are not forwarded to the application. The MAC comes out of the power down mode only when either a Magic Packet or a Remote wakeup frame is received and the corresponding detection is enabled.

### Remote wakeup frame filter register

There are eight wakeup frame filter registers. To write on each of them, load the wakeup frame filter register value by value. The wanted values of the wakeup frame filter are loaded by sequentially loading eight times the wakeup frame filter register. The read operation is identical to the write operation. To read the eight values, you have to read eight times the wakeup frame filter register to reach the last register. Each read/write points the wakeup frame filter register to the next filter register.

Wakeup frame filter reg0 Filter 0 Byte Mask Wakeup frame filter reg1 Filter 1 Byte Mask Wakeup frame filter reg2 Filter 2 Byte Mask Wakeup frame filter reg3 Filter 3 Byte Mask Filter 2 Filter 3 Filter 1 Wakeup frame filter reg4 RSVD RSVD RSVD RSVD Command Command Command Command Wakeup frame filter reg5 Filter 3 Offset Filter 2 Offset Filter 1 Offset Filter 0 Offset Wakeup frame filter reg6 Filter 1 CRC - 16 Filter 0 CRC - 16 Wakeup frame filter reg7 Filter 3 CRC - 16 Filter 2 CRC - 16 ai15647

Figure 335. Wakeup frame filter register

#### Filter i Byte Mask

This register defines which bytes of the frame are examined by filter i (0, 1, 2, and 3) in order to determine whether or not the frame is a wakeup frame. The MSB (thirty-first bit) must be zero. Bit j [30:0] is the Byte Mask. If bit j (byte number) of the Byte Mask is set, then Filter i Offset + j of the incoming frame is processed by the CRC block; otherwise Filter i Offset + j is ignored.

#### Filter i Command

This 4-bit command controls the filter i operation. Bit 3 specifies the address type, defining the pattern's destination address type. When the bit is set, the pattern applies to only multicast frames. When the bit is reset, the pattern applies only to unicast frames. Bit 2 and bit 1 are reserved. Bit 0 is the enable bit for filter i; if bit 0 is not set, filter i is disabled.

### Filter i Offset

This register defines the offset (within the frame) from which the frames are examined by filter i. This 8-bit pattern offset is the offset for the filter i first byte to be examined. The minimum allowed is 12, which refers to the 13th byte of the frame (offset value 0 refers to the first byte of the frame).

#### Filter i CRC-16

This register contains the CRC\_16 value calculated from the pattern, as well as the byte mask programmed to the wakeup filter register block.

### Remote wakeup frame detection

When the MAC is in sleep mode and the remote wakeup bit is enabled in the ETH\_MACPMTCSR register, normal operation is resumed after receiving a remote wakeup frame. The application writes all eight wakeup filter registers, by performing a sequential write to the wakeup frame filter register address. The application enables remote wakeup by writing a 1 to bit 2 in the ETH\_MACPMTCSR register. PMT supports four programmable filters that provide different receive frame patterns. If the incoming frame passes the address filtering of Filter Command, and if Filter CRC-16 matches the incoming examined pattern, then the wakeup frame is received. Filter\_offset (minimum value 12, which refers to the 13th byte of the frame) determines the offset from which the frame is to be examined. Filter Byte Mask determines which bytes of the frame must be examined. The thirty-first bit of Byte Mask must be set to zero. The wakeup frame is checked only for length error, FCS error, dribble bit error, MII error, collision, and to ensure that it is not a runt frame. Even if the

wakeup frame is more than 512 bytes long, if the frame has a valid CRC value, it is considered valid. Wakeup frame detection is updated in the ETH\_MACPMTCSR register for every remote wakeup frame received. If enabled, a PMT interrupt is generated to indicate the reception of a remote wakeup frame.

### Magic packet detection

The Magic Packet frame is based on a method that uses Advanced Micro Device's Magic Packet technology to power up the sleeping device on the network. The MAC receives a specific packet of information, called a Magic Packet, addressed to the node on the network. Only Magic Packets that are addressed to the device or a broadcast address are checked to determine whether they meet the wakeup requirements. Magic Packets that pass address filtering (unicast or broadcast) are checked to determine whether they meet the remote Wake-on-LAN data format of 6 bytes of all ones followed by a MAC address appearing 16 times. The application enables Magic Packet wakeup by writing a 1 to bit 1 in the ETH MACPMTCSR register. The PMT block constantly monitors each frame addressed to the node for a specific Magic Packet pattern. Each received frame is checked for a 0xFFFF FFFF pattern following the destination and source address field. The PMT block then checks the frame for 16 repetitions of the MAC address without any breaks or interruptions. In case of a break in the 16 repetitions of the address, the 0xFFFF FFFF FFFF pattern is scanned for again in the incoming frame. The 16 repetitions can be anywhere in the frame, but must be preceded by the synchronization stream (0xFFFF FFFF FFFF). The device also accepts a multicast frame, as long as the 16 duplications of the MAC address are detected. If the MAC address of a node is 0x0011 2233 4455, then the MAC scans for the data sequence:

Magic Packet detection is updated in the ETH\_MACPMTCSR register for received Magic Packet. If enabled, a PMT interrupt is generated to indicate the reception of a Magic Packet.

#### System consideration during power-down

The Ethernet PMT block is able to detect frames while the system is in the Stop mode, provided that the EXTI line 19 is enabled.

The MAC receiver state machine should remain enabled during the power-down mode. This means that the RE bit has to remain set in the ETH\_MACCR register because it is involved in magic packet/ wake-on-LAN frame detection. The transmit state machine should however be turned off during the power-down mode by clearing the TE bit in the ETH\_MACCR register. Moreover, the Ethernet DMA should be disabled during the power-down mode, because it is not necessary to copy the magic packet/wake-on-LAN frame into the SRAM. To disable the Ethernet DMA, clear the ST bit and the SR bit (for the transmit DMA and the receive DMA, respectively) in the ETH\_DMAOMR register.

The recommended power-down and wakeup sequences are as follows:

- Disable the transmit DMA and wait for any previous frame transmissions to complete.
   These transmissions can be detected when the transmit interrupt ETH\_DMASR register[0] is received.
- 2. Disable the MAC transmitter and MAC receiver by clearing the RE and TE bits in the ETH\_MACCR configuration register.
- 3. Wait for the receive DMA to have emptied all the frames in the Rx FIFO.
- 4. Disable the receive DMA.
- 5. Configure and enable the EXTI line 19 to generate either an event or an interrupt.
- 6. If you configure the EXTI line 19 to generate an interrupt, you also have to correctly configure the ETH\_WKUP\_IRQ Handler function, which should clear the pending bit of the EXTI line 19.
- 7. Enable Magic packet/Wake-on-LAN frame detection by setting the MFE/ WFE bit in the ETH\_MACPMTCSR register.
- Enable the MAC power-down mode, by setting the PD bit in the ETH\_MACPMTCSR register.
- 9. Enable the MAC Receiver by setting the RE bit in the ETH\_MACCR register.
- 10. Enter the system's Stop mode (for more details refer to Section 5.3.4: Stop mode):
- On receiving a valid wakeup frame, the Ethernet peripheral exits the power-down mode.
- 12. Read the ETH\_MACPMTCSR to clear the power management event flag, enable the MAC transmitter state machine, and the receive and transmit DMA.
- 13. Configure the system clock: enable the HSE and set the clocks.

## 28.5.9 Precision time protocol (IEEE1588 PTP)

The IEEE 1588 standard defines a protocol that allows precise clock synchronization in measurement and control systems implemented with technologies such as network communication, local computing and distributed objects. The protocol applies to systems that communicate by local area networks supporting multicast messaging, including (but not limited to) Ethernet. This protocol is used to synchronize heterogeneous systems that include clocks of varying inherent precision, resolution and stability. The protocol supports system-wide synchronization accuracy in the submicrosecond range with minimum network and local clock computing resources. The message-based protocol, known as the precision time protocol (PTP), is transported over UDP/IP. The system or network is classified into Master and Slave nodes for distributing the timing/clock information. The protocol's technique for synchronizing a slave node to a master node by exchanging PTP messages is described in *Figure 336*.

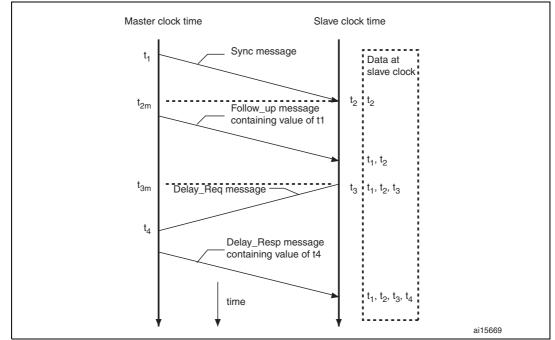


Figure 336. Networked time synchronization

- The master broadcasts PTP Sync messages to all its nodes. The Sync message contains the master's reference time information. The time at which this message leaves the master's system is t<sub>1</sub>. For Ethernet ports, this time has to be captured at the MII.
- 2. A slave receives the Sync message and also captures the exact time, t<sub>2</sub>, using its timing reference.
- 3. The master then sends the slave a Follow\_up message, which contains the t<sub>1</sub> information for later use.
- 4. The slave sends the master a Delay\_Req message, noting the exact time, t<sub>3</sub>, at which this frame leaves the MII.
- 5. The master receives this message and captures the exact time, t<sub>4</sub>, at which it enters its system.
- 6. The master sends the t<sub>4</sub> information to the slave in the Delay\_Resp message.
- 7. The slave uses the four values of t<sub>1</sub>, t<sub>2</sub>, t<sub>3</sub>, and t<sub>4</sub> to synchronize its local timing reference to the master's timing reference.

Most of the protocol implementation occurs in the software, above the UDP layer. As described above, however, hardware support is required to capture the exact time when specific PTP packets enter or leave the Ethernet port at the MII. This timing information has to be captured and returned to the software for a proper, high-accuracy implementation of PTP.

### Reference timing source

To get a snapshot of the time, the core requires a reference time in 64-bit format (split into two 32-bit channels, with the upper 32 bits providing time in seconds, and the lower 32 bits indicating time in nanoseconds) as defined in the IEEE 1588 specification.

The PTP reference clock input is used to internally generate the reference time (also called the System Time) and to capture time stamps. The frequency of this reference clock must

be greater than or equal to the resolution of time stamp counter. The synchronization accuracy target between the master node and the slaves is around 100 ns.

The generation, update and modification of the System Time are described in the *Section : System Time correction methods.* 

The accuracy depends on the PTP reference clock input period, the characteristics of the oscillator (drift) and the frequency of the synchronization procedure.

Due to the synchronization from the Tx and Rx clock input domain to the PTP reference clock domain, the uncertainty on the time stamp latched value is 1 reference clock period. If we add the uncertainty due to resolution, we will add half the period for time stamping.

### Transmission of frames with the PTP feature

When a frame's SFD is output on the MII, a time stamp is captured. Frames for which time stamp capture is required are controllable on a per-frame basis. In other words, each transmitted frame can be marked to indicate whether a time stamp must be captured or not for that frame. The transmitted frames are not processed to identify PTP frames. Frame control is exercised through the control bits in the transmit descriptor. Captured time stamps are returned to the application in the same way as the status is provided for frames. The time stamp is sent back along with the Transmit status of the frame, inside the corresponding transmit descriptor, thus connecting the time stamp automatically to the specific PTP frame. The 64-bit time stamp information is written back to the TDES2 and TDES3 fields, with TDES2 holding the time stamp's 32 least significant bits.

### Reception of frames with the PTP feature

When the IEEE 1588 time stamping feature is enabled, the Ethernet MAC captures the time stamp of all frames received on the MII. The MAC provides the time stamp as soon as the frame reception is complete. Captured time stamps are returned to the application in the same way as the frame status is provided. The time stamp is sent back along with the Receive status of the frame, inside the corresponding receive descriptor. The 64-bit time stamp information is written back to the RDES2 and RDES3 fields, with RDES2 holding the time stamp's 32 least significant bits.

### System Time correction methods

The 64-bit PTP time is updated using the PTP input reference clock, HCLK. This PTP time is used as a source to take snapshots (time stamps) of the Ethernet frames being transmitted or received at the MII. The System Time counter can be initialized or corrected using either the Coarse or the Fine correction method.

In the Coarse correction method, the initial value or the offset value is written to the Time stamp update register (refer to *Section 28.8.3: IEEE 1588 time stamp registers on page 921*). For initialization, the System Time counter is written with the value in the Time stamp update registers, whereas for system time correction, the offset value (Time stamp update register) is added to or subtracted from the system time.

In the Fine correction method, the slave clock (reference clock) frequency drift with respect to the master clock (as defined in IEEE 1588) is corrected over a period of time, unlike in the Coarse correction method where it is corrected in a single clock cycle. The longer correction time helps maintain linear time and does not introduce drastic changes (or a large jitter) in the reference time between PTP Sync message intervals. In this method, an accumulator sums up the contents of the Addend register as shown in *Figure 337*. The arithmetic carry that the accumulator generates is used as a pulse to increment the system time counter.

The accumulator and the addend are 32-bit registers. Here, the accumulator acts as a high-precision frequency multiplier or divider. *Figure 337* shows this algorithm.

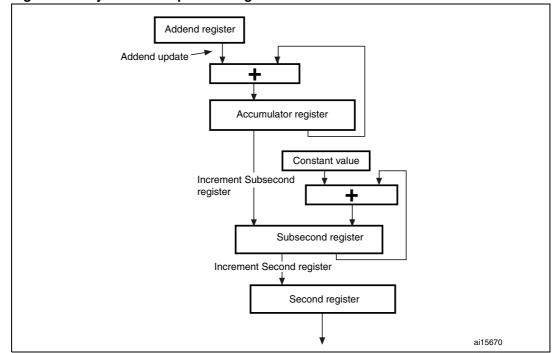


Figure 337. System time update using the Fine correction method

The system time update logic requires a 50 MHz clock frequency to achieve 20 ns accuracy. The frequency division is the ratio of the reference clock frequency to the required clock frequency. Hence, if the reference clock (HCLK) is, let us say, 66 MHz, the ratio is calculated as 66 MHz/50 MHz = 1.32. Hence, the default addend value to be set in the register is  $2^{32}/1.32$ , which is equal to 0xC1F0 7C1F.

If the reference clock drifts lower, to 65 MHz for example, the ratio is 65/50 or 1.3 and the value to set in the addend register is  $2^{32}/1.30$  equal to 0xC4EC 4EC4. If the clock drifts higher, to 67 MHz for example, the addend register must be set to 0xBF0 B7672. When the clock drift is zero, the default addend value of 0xC1F0 7C1F ( $2^{32}/1.32$ ) should be programmed.

In *Figure 337*, the constant value used to increment the subsecond register is 0d43. This makes an accuracy of 20 ns in the system time (in other words, it is incremented by 20 ns steps).

The software has to calculate the drift in frequency based on the Sync messages, and to update the Addend register accordingly. Initially, the slave clock is set with FreqCompensationValue0 in the Addend register. This value is as follows:

FreqCompensationValue0 = 2<sup>32</sup> / FreqDivisionRatio

If MasterToSlaveDelay is initially assumed to be the same for consecutive Sync messages, the algorithm described below must be applied. After a few Sync cycles, frequency lock occurs. The slave clock can then determine a precise MasterToSlaveDelay value and resynchronize with the master using the new value.

The algorithm is as follows:

- At time MasterSyncTime (n) the master sends the slave clock a Sync message. The slave receives this message when its local clock is SlaveClockTime (n) and computes MasterClockTime (n) as:
  - MasterClockTime (n) = MasterSyncTime (n) + MasterToSlaveDelay (n)
- The master clock count for current Sync cycle, MasterClockCount (n) is given by:
   MasterClockCount (n) = MasterClockTime (n) − MasterClockTime (n − 1) (assuming that MasterToSlaveDelay is the same for Sync cycles n and n − 1)
- The slave clock count for current Sync cycle, SlaveClockCount (n) is given by:
   SlaveClockCount (n) = SlaveClockTime (n) − SlaveClockTime (n − 1)
- The difference between master and slave clock counts for current Sync cycle, ClockDiffCount (n) is given by:
  - ClockDiffCount (n) = MasterClockCount (n) SlaveClockCount (n)
- The frequency-scaling factor for slave clock, FreqScaleFactor (n) is given by:
   FreqScaleFactor (n) = (MasterClockCount (n) + ClockDiffCount (n)) /
   SlaveClockCount (n)
- The frequency compensation value for Addend register, FreqCompensationValue (n) is given by:

 $FreqCompensationValue (n) = FreqScaleFactor (n) \times FreqCompensationValue (n - 1)$ 

In theory, this algorithm achieves lock in one Sync cycle; however, it may take several cycles, due to changing network propagation delays and operating conditions.

This algorithm is self-correcting: if for any reason the slave clock is initially set to a value from the master that is incorrect, the algorithm corrects it at the cost of more Sync cycles.

### Programming steps for system time generation initialization

The time stamping feature can be enabled by setting bit 0 in the Time stamp control register (ETH\_\_PTPTSCR). However, it is essential to initialize the time stamp counter after this bit is set to start time stamp operation. The proper sequence is the following:

- 1. Mask the Time stamp trigger interrupt by setting bit 9 in the MACIMR register.
- 2. Program Time stamp register bit 0 to enable time stamping.
- 3. Program the Subsecond increment register based on the PTP clock frequency.
- 4. If you are using the Fine correction method, program the Time stamp addend register and set Time stamp control register bit 5 (addend register update).
- 5. Poll the Time stamp control register until bit 5 is cleared.
- 6. To select the Fine correction method (if required), program Time stamp control register bit 1.
- 7. Program the Time stamp high update and Time stamp low update registers with the appropriate time value.
- 8. Set Time stamp control register bit 2 (Time stamp init).
- 9. The Time stamp counter starts operation as soon as it is initialized with the value written in the Time stamp update register.
- 10. Enable the MAC receiver and transmitter for proper time stamping.

Note: If time stamp operation is disabled by clearing bit 0 in the ETH\_PTPTSCR register, the above steps must be repeated to restart the time stamp operation.



### Programming steps for system time update in the Coarse correction method

To synchronize or update the system time in one process (coarse correction method), perform the following steps:

- 1. Write the offset (positive or negative) in the Time stamp update high and low registers.
- 2. Set bit 3 (TSSTU) in the Time stamp control register.
- 3. The value in the Time stamp update registers is added to or subtracted from the system time when the TSSTU bit is cleared.

### Programming steps for system time update in the Fine correction method

To synchronize or update the system time to reduce system-time jitter (fine correction method), perform the following steps:

- 1. With the help of the algorithm explained in *Section : System Time correction methods*, calculate the rate by which you want to speed up or slow down the system time increments.
- 2. Update the time stamp.
- 3. Wait the time you want the new value of the Addend register to be active. You can do this by activating the Time stamp trigger interrupt after the system time reaches the target value.
- 4. Program the required target time in the Target time high and low registers. Unmask the Time stamp interrupt by clearing bit 9 in the ETH MACIMR register.
- 5. Set Time stamp control register bit 4 (TSARU).
- 6. When this trigger causes an interrupt, read the ETH\_MACSR register.
- 7. Reprogram the Time stamp addend register with the old value and set ETH\_TPTSCR bit 5 again.

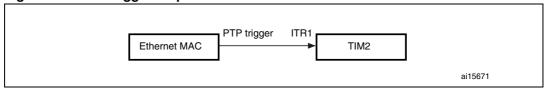
### PTP trigger internal connection with TIM2

The MAC provides a trigger interrupt when the system time becomes greater than the target time. Using an interrupt introduces a known latency plus an uncertainty in the command execution time.

In order to avoid this uncertainty, a PTP trigger output signal is set high when the system time is greater than the target time. It is internally connected to the TIM2 input trigger. With this signal, the input capture feature, the output compare feature and the waveforms of the timer can be used, triggered by the synchronized PTP system time. No uncertainty is introduced since the clock of the timer (PCLK1: TIM2 APB1 clock) and PTP reference clock (HCLK) are synchronous.

This PTP trigger signal is connected to the TIM2 ITR1 input selectable by software. The connection is enabled through bits 11 and 10 in the TIM2 option register (TIM2\_OR). *Figure 338* shows the connection.

Figure 338. PTP trigger output to TIM2 ITR1 connection



### PTP pulse-per-second output signal

This PTP pulse output is used to check the synchronization between all nodes in the network. To be able to test the difference between the local slave clock and the master reference clock, both clocks were given a pulse-per-second (PPS) output signal that may be connected to an oscilloscope if necessary. The deviation between the two signals can therefore be measured. The pulse width of the PPS output is 125 ms.

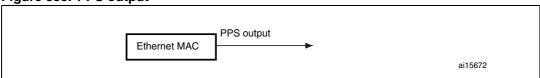
The PPS output is enabled through bits 11 and 10 in the TIM2 option register (TIM2\_OR).

The default frequency of the PPS output is 1 Hz. PPSFREQ[3:0] (in ETH\_PTPPPSCR) can be used to set the frequency of the PPS output to 2<sup>PPSFREQ</sup> Hz.

When set to 1 Hz, the PPS pulse width is 125 ms with binary rollover (TSSSR=0, bit 9 in ETH\_PTPTSCR) and 100 ms with digital rollover (TSSSR=1). When set to 2 Hz and higher, the duty cycle of the PPS output is 50% with binary rollover.

With digital rollover (TSSSR=1), it is recommended not to use the PPS output with a frequency other than 1 Hz as it would have irregular waveforms (though its average frequency would always be correct during any one-second window).





# 28.6 Ethernet functional description: DMA controller operation

The DMA has independent transmit and receive engines, and a CSR space. The transmit engine transfers data from system memory into the Tx FIFO while the receive engine transfers data from the Rx FIFO into system memory. The controller utilizes descriptors to efficiently move data from source to destination with minimum CPU intervention. The DMA is designed for packet-oriented data transfers such as frames in Ethernet. The controller can be programmed to interrupt the CPU in cases such as frame transmit and receive transfer completion, and other normal/error conditions. The DMA and the STM32F4xx communicate through two data structures:

- Control and status registers (CSR)
- Descriptor lists and data buffers.

Control and status registers are described in detail in *Section 28.8 on page 897*. Descriptors are described in detail in *Section on page 877*.

The DMA transfers the received data frames to the receive buffer in the STM32F4xx memory, and transmits data frames from the transmit buffer in the STM32F4xx memory. Descriptors that reside in the STM32F4xx memory act as pointers to these buffers. There are two descriptor lists: one for reception, and one for transmission. The base address of each list is written into DMA Registers 3 and 4, respectively. A descriptor list is forward-linked (either implicitly or explicitly). The last descriptor may point back to the first entry to create a ring structure. Explicit chaining of descriptors is accomplished by configuring the second address chained in both the receive and transmit descriptors (RDES1[14] and TDES0[20]). The descriptor lists reside in the Host's physical memory space. Each descriptor can point to a maximum of two buffers. This enables the use of two

physically addressed buffers, instead of two contiguous buffers in memory. A data buffer resides in the Host's physical memory space, and consists of an entire frame or part of a frame, but cannot exceed a single frame. Buffers contain only data. The buffer status is maintained in the descriptor. Data chaining refers to frames that span multiple data buffers. However, a single descriptor cannot span multiple frames. The DMA skips to the next frame buffer when the end of frame is detected. Data chaining can be enabled or disabled. The descriptor ring and chain structure is shown in *Figure 340*.

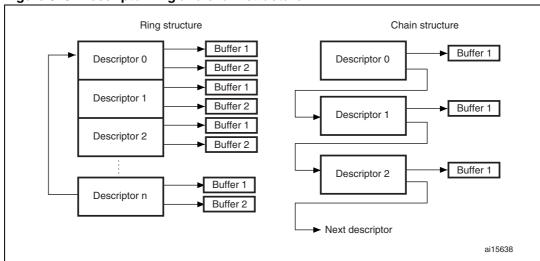


Figure 340. Descriptor ring and chain structure

## 28.6.1 Initialization of a transfer using DMA

Initialization for the MAC is as follows:

- 1. Write to ETH DMABMR to set STM32F4xx bus access parameters.
- Write to the ETH\_DMAIER register to mask unnecessary interrupt causes.
- 3. The software driver creates the transmit and receive descriptor lists. Then it writes to both the ETH\_DMARDLAR and ETH\_DMATDLAR registers, providing the DMA with the start address of each list.
- 4. Write to MAC Registers 1, 2, and 3 to choose the desired filtering options.
- 5. Write to the MAC ETH\_MACCR register to configure and enable the transmit and receive operating modes. The PS and DM bits are set based on the auto-negotiation result (read from the PHY).
- 6. Write to the ETH\_DMAOMR register to set bits 13 and 1 and start transmission and reception.
- 7. The transmit and receive engines enter the running state and attempt to acquire descriptors from the respective descriptor lists. The receive and transmit engines then begin processing receive and transmit operations. The transmit and receive processes are independent of each other and can be started or stopped separately.

### 28.6.2 Host bus burst access

The DMA attempts to execute fixed-length burst transfers on the AHB master interface if configured to do so (FB bit in ETH\_DMABMR). The maximum burst length is indicated and limited by the PBL field (ETH\_DMABMR [13:8]). The receive and transmit descriptors are

always accessed in the maximum possible burst size (limited by PBL) for the 16 bytes to be read

The Transmit DMA initiates a data transfer only when there is sufficient space in the Transmit FIFO to accommodate the configured burst or the number of bytes until the end of frame (when it is less than the configured burst length). The DMA indicates the start address and the number of transfers required to the AHB Master Interface. When the AHB Interface is configured for fixed-length burst, then it transfers data using the best combination of INCR4, INCR8, INCR16 and SINGLE transactions. Otherwise (no fixed-length burst), it transfers data using INCR (undefined length) and SINGLE transactions.

The Receive DMA initiates a data transfer only when sufficient data for the configured burst is available in Receive FIFO or when the end of frame (when it is less than the configured burst length) is detected in the Receive FIFO. The DMA indicates the start address and the number of transfers required to the AHB master interface. When the AHB interface is configured for fixed-length burst, then it transfers data using the best combination of INCR4, INCR8, INCR16 and SINGLE transactions. If the end of frame is reached before the fixed-burst ends on the AHB interface, then dummy transfers are performed in order to complete the fixed-length burst. Otherwise (FB bit in ETH\_DMABMR is reset), it transfers data using INCR (undefined length) and SINGLE transactions.

When the AHB interface is configured for address-aligned beats, both DMA engines ensure that the first burst transfer the AHB initiates is less than or equal to the size of the configured PBL. Thus, all subsequent beats start at an address that is aligned to the configured PBL. The DMA can only align the address for beats up to size 16 (for PBL > 16), because the AHB interface does not support more than INCR16.

### 28.6.3 Host data buffer alignment

The transmit and receive data buffers do not have any restrictions on start address alignment. In our system with 32-bit memory, the start address for the buffers can be aligned to any of the four bytes. However, the DMA always initiates transfers with address aligned to the bus width with dummy data for the byte lanes not required. This typically happens during the transfer of the beginning or end of an Ethernet frame.

### Example of buffer read:

If the Transmit buffer address is 0x0000 0FF2, and 15 bytes need to be transferred, then the DMA will read five full words from address 0x0000 0FF0, but when transferring data to the Transmit FIFO, the extra bytes (the first two bytes) will be dropped or ignored. Similarly, the last 3 bytes of the last transfer will also be ignored. The DMA always ensures it transfers a full 32-bit data items to the Transmit FIFO, unless it is the end of frame.

### Example of buffer write:

If the Receive buffer address is 0x0000 0FF2, and 16 bytes of a received frame need to be transferred, then the DMA will write five full 32-bit data items from address 0x0000 0FF0. But the first 2 bytes of the first transfer and the last 2 bytes of the third transfer will have dummy data.

### 28.6.4 Buffer size calculations

The DMA does not update the size fields in the transmit and receive descriptors. The DMA updates only the status fields (xDES0) of the descriptors. The driver has to calculate the sizes. The transmit DMA transfers the exact number of bytes (indicated by buffer size field in TDES1) towards the MAC core. If a descriptor is marked as first (FS bit in TDES0 is set),

then the DMA marks the first transfer from the buffer as the start of frame. If a descriptor is marked as last (LS bit in TDES0), then the DMA marks the last transfer from that data buffer as the end of frame. The receive DMA transfers data to a buffer until the buffer is full or the end of frame is received. If a descriptor is not marked as last (LS bit in RDES0), then the buffer(s) that correspond to the descriptor are full and the amount of valid data in a buffer is accurately indicated by the buffer size field minus the data buffer pointer offset when the descriptor's FS bit is set. The offset is zero when the data buffer pointer is aligned to the databus width. If a descriptor is marked as last, then the buffer may not be full (as indicated by the buffer size in RDES1). To compute the amount of valid data in this final buffer, the driver must read the frame length (FL bits in RDES0[29:16]) and subtract the sum of the buffer sizes of the preceding buffers in this frame. The receive DMA always transfers the start of next frame with a new descriptor.

Note:

Even when the start address of a receive buffer is not aligned to the system databus width the system should allocate a receive buffer of a size aligned to the system bus width. For example, if the system allocates a 1024 byte (1 KB) receive buffer starting from address 0x1000, the software can program the buffer start address in the receive descriptor to have a 0x1002 offset. The receive DMA writes the frame to this buffer with dummy data in the first two locations (0x1000 and 0x1001). The actual frame is written from location 0x1002. Thus, the actual useful space in this buffer is 1022 bytes, even though the buffer size is programmed as 1024 bytes, due to the start address offset.

### 28.6.5 DMA arbiter

The arbiter inside the DMA takes care of the arbitration between transmit and receive channel accesses to the AHB master interface. Two types of arbitrations are possible: round-robin, and fixed-priority. When round-robin arbitration is selected (DA bit in ETH\_DMABMR is reset), the arbiter allocates the databus in the ratio set by the PM bits in ETH\_DMABMR, when both transmit and receive DMAs request access simultaneously. When the DA bit is set, the receive DMA always gets priority over the transmit DMA for data access.

### 28.6.6 Error response to DMA

For any data transfer initiated by a DMA channel, if the slave replies with an error response, that DMA stops all operations and updates the error bits and the fatal bus error bit in the Status register (ETH\_DMASR register). That DMA controller can resume operation only after soft- or hard-resetting the peripheral and re-initializing the DMA.

### 28.6.7 Tx DMA configuration

### TxDMA operation: default (non-OSF) mode

The transmit DMA engine in default mode proceeds as follows:

- 1. The user sets up the transmit descriptor (TDES0-TDES3) and sets the OWN bit (TDES0[31]) after setting up the corresponding data buffer(s) with Ethernet frame data.
- 2. Once the ST bit (ETH\_DMAOMR register[13]) is set, the DMA enters the Run state.
- 3. While in the Run state, the DMA polls the transmit descriptor list for frames requiring transmission. After polling starts, it continues in either sequential descriptor ring order or chained order. If the DMA detects a descriptor flagged as owned by the CPU, or if an error condition occurs, transmission is suspended and both the Transmit Buffer

- Unavailable (ETH\_DMASR register[2]) and Normal Interrupt Summary (ETH\_DMASR register[16]) bits are set. The transmit engine proceeds to Step 9.
- If the acquired descriptor is flagged as owned by DMA (TDES0[31] is set), the DMA decodes the transmit data buffer address from the acquired descriptor.
- 5. The DMA fetches the transmit data from the STM32F4xx memory and transfers the data.
- 6. If an Ethernet frame is stored over data buffers in multiple descriptors, the DMA closes the intermediate descriptor and fetches the next descriptor. Steps 3, 4, and 5 are repeated until the end of Ethernet frame data is transferred.
- 7. When frame transmission is complete, if IEEE 1588 time stamping was enabled for the frame (as indicated in the transmit status) the time stamp value is written to the transmit descriptor (TDES2 and TDES3) that contains the end-of-frame buffer. The status information is then written to this transmit descriptor (TDES0). Because the OWN bit is cleared during this step, the CPU now owns this descriptor. If time stamping was not enabled for this frame, the DMA does not alter the contents of TDES2 and TDES3.
- 8. Transmit Interrupt (ETH\_DMASR register [0]) is set after completing the transmission of a frame that has Interrupt on Completion (TDES1[31]) set in its last descriptor. The DMA engine then returns to Step 3.
- 9. In the Suspend state, the DMA tries to re-acquire the descriptor (and thereby returns to Step 3) when it receives a transmit poll demand, and the Underflow Interrupt Status bit is cleared.

Figure 341 shows the TxDMA transmission flow in default mode.



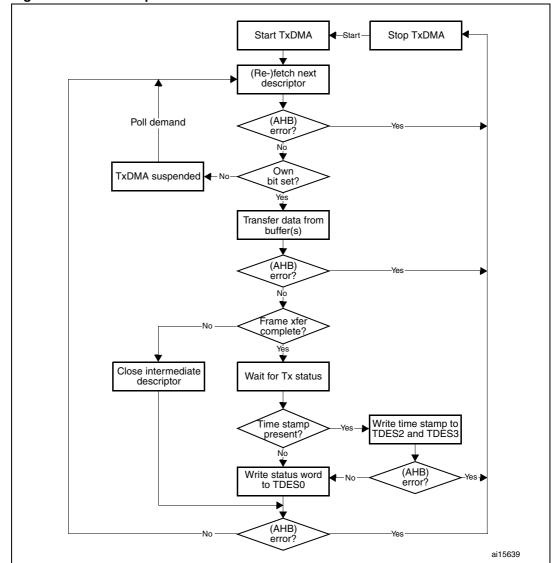


Figure 341. TxDMA operation in Default mode

### **TxDMA operation: OSF mode**

While in the Run state, the transmit process can simultaneously acquire two frames without closing the Status descriptor of the first (if the OSF bit is set in ETH\_DMAOMR register[2]). As the transmit process finishes transferring the first frame, it immediately polls the transmit descriptor list for the second frame. If the second frame is valid, the transmit process transfers this frame before writing the first frame's status information. In OSF mode, the Run-state transmit DMA operates according to the following sequence:

- The DMA operates as described in steps 1–6 of the TxDMA (default mode).
- 2. Without closing the previous frame's last descriptor, the DMA fetches the next descriptor.
- 3. If the DMA owns the acquired descriptor, the DMA decodes the transmit buffer address in this descriptor. If the DMA does not own the descriptor, the DMA goes into Suspend mode and skips to Step 7.
- 4. The DMA fetches the Transmit frame from the STM32F4xx memory and transfers the frame until the end of frame data are transferred, closing the intermediate descriptors if this frame is split across multiple descriptors.
- 5. The DMA waits for the transmission status and time stamp of the previous frame. When the status is available, the DMA writes the time stamp to TDES2 and TDES3, if such time stamp was captured (as indicated by a status bit). The DMA then writes the status, with a cleared OWN bit, to the corresponding TDES0, thus closing the descriptor. If time stamping was not enabled for the previous frame, the DMA does not alter the contents of TDES2 and TDES3.
- 6. If enabled, the Transmit interrupt is set, the DMA fetches the next descriptor, then proceeds to Step 3 (when Status is normal). If the previous transmission status shows an underflow error, the DMA goes into Suspend mode (Step 7).
- 7. In Suspend mode, if a pending status and time stamp are received by the DMA, it writes the time stamp (if enabled for the current frame) to TDES2 and TDES3, then writes the status to the corresponding TDES0. It then sets relevant interrupts and returns to Suspend mode.
- 8. The DMA can exit Suspend mode and enter the Run state (go to Step 1 or Step 2 depending on pending status) only after receiving a Transmit Poll demand (ETH\_DMATPDR register).

Figure 342 shows the basic flowchart in OSF mode.



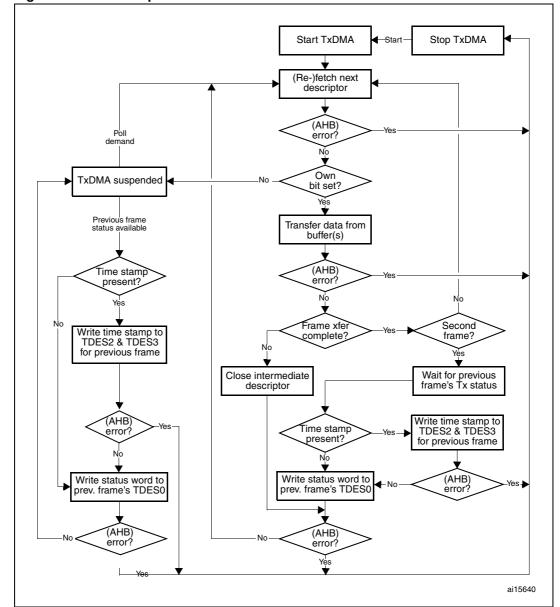


Figure 342. TxDMA operation in OSF mode

### **Transmit frame processing**

The transmit DMA expects that the data buffers contain complete Ethernet frames, excluding preamble, pad bytes, and FCS fields. The DA, SA, and Type/Len fields contain valid data. If the transmit descriptor indicates that the MAC core must disable CRC or pad insertion, the buffer must have complete Ethernet frames (excluding preamble), including the CRC bytes. Frames can be data-chained and span over several buffers. Frames have to be delimited by the first descriptor (TDES0[28]) and the last descriptor (TDES0[29]). As the transmission starts, TDES0[28] has to be set in the first descriptor. When this occurs, the frame data are transferred from the memory buffer to the Transmit FIFO. Concurrently, if the last descriptor (TDES0[29]) of the current frame is cleared, the transmit process attempts to acquire the next descriptor. The transmit process expects TDES0[28] to be cleared in this descriptor. If TDES0[29] is cleared, it indicates an intermediary buffer. If TDES0[29] is set, it

indicates the last buffer of the frame. After the last buffer of the frame has been transmitted, the DMA writes back the final status information to the transmit descriptor 0 (TDES0) word of the descriptor that has the last segment set in transmit descriptor 0 (TDES0[29]). At this time, if Interrupt on Completion (TDES0[30]) is set, Transmit Interrupt (in ETH\_DMASR register [0]) is set, the next descriptor is fetched, and the process repeats. Actual frame transmission begins after the Transmit FIFO has reached either a programmable transmit threshold (ETH\_DMAOMR register[16:14]), or a full frame is contained in the FIFO. There is also an option for the Store and forward mode (ETH\_DMAOMR register[21]). Descriptors are released (OWN bit TDES0[31] is cleared) when the DMA finishes transferring the frame.

### Transmit polling suspended

Transmit polling can be suspended by either of the following conditions:

- The DMA detects a descriptor owned by the CPU (TDES0[31]=0) and the Transmit buffer unavailable flag is set (ETH\_DMASR register[2]). To resume, the driver must give descriptor ownership to the DMA and then issue a Poll Demand command.
- A frame transmission is aborted when a transmit error due to underflow is detected. The appropriate Transmit Descriptor 0 (TDES0) bit is set. If the second condition occurs, both the Abnormal Interrupt Summary (in ETH\_DMASR register [15]) and Transmit Underflow bits (in ETH\_DMASR register[5]) are set, and the information is written to Transmit Descriptor 0, causing the suspension. If the DMA goes into Suspend state due to the first condition, then both the Normal Interrupt Summary (ETH\_DMASR register [16]) and Transmit Buffer Unavailable (ETH\_DMASR register[2]) bits are set. In both cases, the position in the transmit list is retained. The retained position is that of the descriptor following the last descriptor closed by the DMA. The driver must explicitly issue a Transmit Poll Demand command after rectifying the suspension cause.

### **Normal Tx DMA descriptors**

The normal transmit descriptor structure consists of four 32-bit words as shown in *Figure 343*. The bit descriptions of TDES0, TDES1, TDES2 and TDES3 are given below.

Note that enhanced descriptors must be used if time stamping is activated (ETH\_PTPTSCR bit 0, TSE=1) or if IPv4 checksum offload is activated (ETH\_MACCR bit 10, IPCO=1).

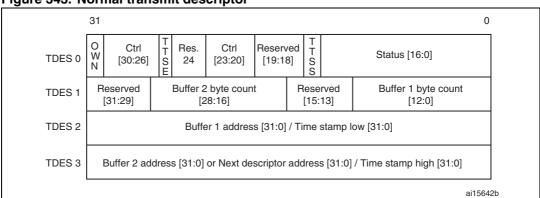


Figure 343. Normal transmit descriptor

### TDES0: Transmit descriptor Word0

The application software has to program the control bits [30:26]+[23:20] plus the OWN bit [31] during descriptor initialization. When the DMA updates the descriptor (or writes it back), it resets all the control bits plus the OWN bit, and reports only the status bits.

#### 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 31 30 29 28 27 26 25 TT ΤE TC TT ΙH IP E LC LC O IC LS FS DP ES JT FF NC EC VF СС ED UF DB DC CIC R SS Ε SF Н Α Res Res. Ν rw rw

#### Bit 31 OWN: Own bit

When set, this bit indicates that the descriptor is owned by the DMA. When this bit is reset, it indicates that the descriptor is owned by the CPU. The DMA clears this bit either when it completes the frame transmission or when the buffers allocated in the descriptor are read completely. The ownership bit of the frame's first descriptor must be set after all subsequent descriptors belonging to the same frame have been set.

#### Bit 30 IC: Interrupt on completion

When set, this bit sets the Transmit Interrupt (Register 5[0]) after the present frame has been transmitted.

#### Bit 29 LS: Last segment

When set, this bit indicates that the buffer contains the last segment of the frame.

#### Bit 28 FS: First segment

When set, this bit indicates that the buffer contains the first segment of a frame.

#### Bit 27 DC: Disable CRC

When this bit is set, the MAC does not append a cyclic redundancy check (CRC) to the end of the transmitted frame. This is valid only when the first segment (TDES0[28]) is set.

### Bit 26 DP: Disable pad

When set, the MAC does not automatically add padding to a frame shorter than 64 bytes. When this bit is reset, the DMA automatically adds padding and CRC to a frame shorter than 64 bytes, and the CRC field is added despite the state of the DC (TDES0[27]) bit. This is valid only when the first segment (TDES0[28]) is set.

#### Bit 25 TTSE: Transmit time stamp enable

When TTSE is set and when TSE is set (ETH\_PTPTSCR bit 0), IEEE1588 hardware time stamping is activated for the transmit frame described by the descriptor. This field is only valid when the First segment control bit (TDES0[28]) is set.

Bit 24 Reserved, must be kept at reset value.

#### Bits 23:22 CIC: Checksum insertion control

These bits control the checksum calculation and insertion. Bit encoding is as shown below:

00: Checksum Insertion disabled

01: Only IP header checksum calculation and insertion are enabled

10: IP header checksum and payload checksum calculation and insertion are enabled, but pseudo-header checksum is not calculated in hardware

11: IP Header checksum and payload checksum calculation and insertion are enabled, and pseudo-header checksum is calculated in hardware.

### Bit 21 TER: Transmit end of ring

When set, this bit indicates that the descriptor list reached its final descriptor. The DMA returns to the base address of the list, creating a descriptor ring.



#### Bit 20 TCH: Second address chained

When set, this bit indicates that the second address in the descriptor is the next descriptor address rather than the second buffer address. When TDES0[20] is set, TBS2 (TDES1[28:16]) is a "don't care" value. TDES0[21] takes precedence over TDES0[20].

#### Bits 19:18 Reserved, must be kept at reset value.

#### Bit 17 TTSS: Transmit time stamp status

This field is used as a status bit to indicate that a time stamp was captured for the described transmit frame. When this bit is set, TDES2 and TDES3 have a time stamp value captured for the transmit frame. This field is only valid when the descriptor's Last segment control bit (TDES0[29]) is set.

Note that when enhanced descriptors are enabled (EDFE=1 in ETH\_DMABMR), TTSS=1 indicates that TDES6 and TDES7 have the time stamp value.

#### Bit 16 IHE: IP header error

When set, this bit indicates that the MAC transmitter detected an error in the IP datagram header. The transmitter checks the header length in the IPv4 packet against the number of header bytes received from the application and indicates an error status if there is a mismatch. For IPv6 frames, a header error is reported if the main header length is not 40 bytes. Furthermore, the Ethernet length/type field value for an IPv4 or IPv6 frame must match the IP header version received with the packet. For IPv4 frames, an error status is also indicated if the Header Length field has a value less than 0x5.

### Bit 15 ES: Error summary

Indicates the logical OR of the following bits:

TDES0[14]: Jabber timeout

TDES0[13]: Frame flush

TDES0[11]: Loss of carrier

TDES0[10]: No carrier

TDES0[9]: Late collision

TDES0[8]: Excessive collision

TDES0[2]:Excessive deferral

TDES0[1]: Underflow error

TDES0[16]: IP header error

TDES0[12]: IP payload error

### Bit 14 JT: Jabber timeout

When set, this bit indicates the MAC transmitter has experienced a jabber timeout. This bit is only set when the MAC configuration register's JD bit is not set.

### Bit 13 FF: Frame flushed

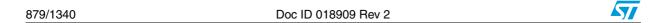
When set, this bit indicates that the DMA/MTL flushed the frame due to a software Flush command given by the CPU.

### Bit 12 IPE: IP payload error

When set, this bit indicates that MAC transmitter detected an error in the TCP, UDP, or ICMP IP datagram payload. The transmitter checks the payload length received in the IPv4 or IPv6 header against the actual number of TCP, UDP or ICMP packet bytes received from the application and issues an error status in case of a mismatch.

### Bit 11 LCA: Loss of carrier

When set, this bit indicates that a loss of carrier occurred during frame transmission (that is, the MII\_CRS signal was inactive for one or more transmit clock periods during frame transmission). This is valid only for the frames transmitted without collision when the MAC operates in Half-duplex mode.



#### Bit 10 NC: No carrier

When set, this bit indicates that the Carrier Sense signal form the PHY was not asserted during transmission.

#### Bit 9 LCO: Late collision

When set, this bit indicates that frame transmission was aborted due to a collision occurring after the collision window (64 byte times, including preamble, in MII mode). This bit is not valid if the Underflow Error bit is set.

#### Bit 8 EC: Excessive collision

When set, this bit indicates that the transmission was aborted after 16 successive collisions while attempting to transmit the current frame. If the RD (Disable retry) bit in the MAC Configuration register is set, this bit is set after the first collision, and the transmission of the frame is aborted

### Bit 7 VF: VLAN frame

When set, this bit indicates that the transmitted frame was a VLAN-type frame.

#### Bits 6:3 CC: Collision count

This 4-bit counter value indicates the number of collisions occurring before the frame was transmitted. The count is not valid when the Excessive collisions bit (TDES0[8]) is set.

#### Bit 2 ED: Excessive deferral

When set, this bit indicates that the transmission has ended because of excessive deferral of over 24 288 bit times if the Deferral check (DC) bit in the MAC Control register is set high.

#### Bit 1 UF: Underflow error

When set, this bit indicates that the MAC aborted the frame because data arrived late from the RAM memory. Underflow error indicates that the DMA encountered an empty transmit buffer while transmitting the frame. The transmission process enters the Suspended state and sets both Transmit underflow (Register 5[5]) and Transmit interrupt (Register 5[0]).

### Bit 0 DB: Deferred bit

When set, this bit indicates that the MAC defers before transmission because of the presence of the carrier. This bit is valid only in Half-duplex mode.

### TDES1: Transmit descriptor Word1

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved						Т	BS2				_			Reserved						TI	BS1						
neserveu	rw	rw	rw	rw	rw	rw	rw	neserveu	rw	rw	rw	rw	rw	rw	rw												

31:29 Reserved, must be kept at reset value.

### 28:16 TBS2: Transmit buffer 2 size

These bits indicate the second data buffer size in bytes. This field is not valid if TDES0[20] is set

15:13 Reserved, must be kept at reset value.

#### 12:0 TBS1: Transmit buffer 1 size

These bits indicate the first data buffer byte size, in bytes. If this field is 0, the DMA ignores this buffer and uses Buffer 2 or the next descriptor, depending on the value of TCH (TDES0[20]).



### TDES2: Transmit descriptor Word2

TDES2 contains the address pointer to the first buffer of the descriptor or it contains time stamp data.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

TBAP1/TBAP/TTSL

rw

Bits 31:0 TBAP1: Transmit buffer 1 address pointer / Transmit frame time stamp low

These bits have two different functions: they indicate to the DMA the location of data in memory, and after all data are transferred, the DMA can then use these bits to pass back time stamp data.

**TBAP:** When the software makes this descriptor available to the DMA (at the moment that the OWN bit is set to 1 in TDES0), these bits indicate the physical address of Buffer 1. There is no limitation on the buffer address alignment. See *Host data buffer alignment on page 871* for further details on buffer address alignment.

TTSL: Before it clears the OWN bit in TDES0, the DMA updates this field with the 32 least significant bits of the time stamp captured for the corresponding transmit frame (overwriting the value for TBAP1). This field has the time stamp only if time stamping is activated for this frame (see TTSE, TDES0 bit 25) and if the Last segment control bit (LS) in the descriptor is set.

### TDES3: Transmit descriptor Word3

TDES3 contains the address pointer either to the second buffer of the descriptor or the next descriptor, or it contains time stamp data.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

TBAP2/TBAP2/TTSH

rw

Bits 31:0 **TBAP2:** Transmit buffer 2 address pointer (Next descriptor address) / Transmit frame time stamp high

These bits have two different functions: they indicate to the DMA the location of data in memory, and after all data are transferred, the DMA can then use these bits to pass back time stamp data.

**TBAP2:** When the software makes this descriptor available to the DMA (at the moment when the OWN bit is set to 1 in TDES0), these bits indicate the physical address of Buffer 2 when a descriptor ring structure is used. If the Second address chained (TDES1 [24]) bit is set, this address contains the pointer to the physical memory where the next descriptor is present. The buffer address pointer must be aligned to the bus width only when TDES1 [24] is set. (LSBs are ignored internally.)

**TTSH:** Before it clears the OWN bit in TDES0, the DMA updates this field with the 32 most significant bits of the time stamp captured for the corresponding transmit frame (overwriting the value for TBAP2). This field has the time stamp only if time stamping is activated for this frame (see TDES0 bit 25, TTSE) and if the Last segment control bit (LS) in the descriptor is set.

### **Enhanced Tx DMA descriptors**

Enhanced descriptors (enabled with EDFE=1, ETHDMABMR bit 7), must be used if time stamping is activated (TSE=1, ETH\_PTPTSCR bit 0) or if IPv4 checksum offload is activated (IPCO=1, ETH\_MACCR bit 10).

Enhanced descriptors comprise eight 32-bit words, twice the size of normal descriptors. TDES0, TDES1, TDES2 and TDES3 have the same definitions as for normal transmit descriptors (refer to *Normal Tx DMA descriptors*). TDES6 and TDES7 hold the time stamp. TDES4, TDES5, TDES6 and TDES7 are defined below.

When the Enhanced descriptor mode is selected, the software needs to allocate 32-bytes (8 DWORDS) of memory for every descriptor. When time stamping or IPv4 checksum offload are not being used, the enhanced descriptor format may be disabled and the software can use normal descriptors with the default size of 16 bytes.

31 0 Ctrl Res. Ctrl Reserved w Status [16:0] TDES 0 [30:26] S 24 [23:20] [19:18] Buffer 1 byte count Reserved Buffer 2 byte count Reserved TDES 1 [31:29] [15:13] [28:16] [12:0] TDES 2 Buffer 1 address [31:0] TDES 3 Buffer 2 address [31:0] or Next descriptor address [31:0] TDFS 4 Reserved TDFS 5 Reserved

Time stamp low [31:0]

Time stamp high [31:0]

Figure 344. Enhanced transmit descriptor

TDES4: Transmit descriptor Word4

Reserved

TDFS 6

TDES 7

TDES5: Transmit descriptor Word5

Reserved

TDES6: Transmit descriptor Word6

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

TTSL

rw

Bits 31:0 TTSL: Transmit frame time stamp low

This field is updated by DMA with the 32 least significant bits of the time stamp captured for the corresponding transmit frame. This field has the time stamp only if the Last segment control bit (LS) in the descriptor is set.



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### TDES7: Transmit descriptor Word7

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
															TTSI	1															
															rw																

Bits 31:0 TTSH: Transmit frame time stamp high

This field is updated by DMA with the 32 most significant bits of the time stamp captured for the corresponding transmit frame. This field has the time stamp only if the Last segment control bit (LS) in the descriptor is set.

### 28.6.8 Rx DMA configuration

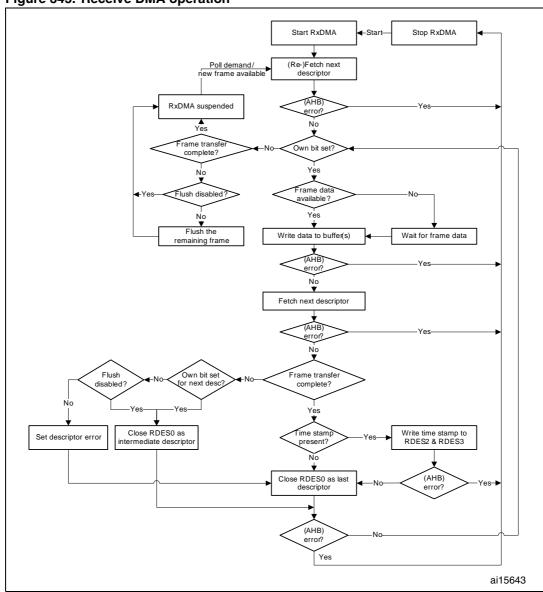
The Receive DMA engine's reception sequence is illustrated in *Figure 345* and described below:

- 1. The CPU sets up Receive descriptors (RDES0-RDES3) and sets the OWN bit (RDES0[31]).
- 2. Once the SR (ETH\_DMAOMR register[1]) bit is set, the DMA enters the Run state. While in the Run state, the DMA polls the receive descriptor list, attempting to acquire free descriptors. If the fetched descriptor is not free (is owned by the CPU), the DMA enters the Suspend state and jumps to Step 9.
- 3. The DMA decodes the receive data buffer address from the acquired descriptors.
- 4. Incoming frames are processed and placed in the acquired descriptor's data buffers.
- 5. When the buffer is full or the frame transfer is complete, the Receive engine fetches the next descriptor.
- 6. If the current frame transfer is complete, the DMA proceeds to step 7. If the DMA does not own the next fetched descriptor and the frame transfer is not complete (EOF is not yet transferred), the DMA sets the Descriptor error bit in RDES0 (unless flushing is disabled). The DMA closes the current descriptor (clears the OWN bit) and marks it as intermediate by clearing the Last segment (LS) bit in the RDES1 value (marks it as last descriptor if flushing is not disabled), then proceeds to step 8. If the DMA owns the next descriptor but the current frame transfer is not complete, the DMA closes the current descriptor as intermediate and returns to step 4.
- 7. If IEEE 1588 time stamping is enabled, the DMA writes the time stamp (if available) to the current descriptor's RDES2 and RDES3. It then takes the received frame's status and writes the status word to the current descriptor's RDES0, with the OWN bit cleared and the Last segment bit set.
- 8. The Receive engine checks the latest descriptor's OWN bit. If the CPU owns the descriptor (OWN bit is at 0) the Receive buffer unavailable bit (in ETH\_DMASR register[7]) is set and the DMA Receive engine enters the Suspended state (step 9). If the DMA owns the descriptor, the engine returns to step 4 and awaits the next frame.
- Before the Receive engine enters the Suspend state, partial frames are flushed from the Receive FIFO (you can control flushing using bit 24 in the ETH\_DMAOMR register).
- 10. The Receive DMA exits the Suspend state when a Receive Poll demand is given or the start of next frame is available from the Receive FIFO. The engine proceeds to step 2 and re-fetches the next descriptor.

The DMA does not acknowledge accepting the status until it has completed the time stamp write-back and is ready to perform status write-back to the descriptor. If software has enabled time stamping through CSR, when a valid time stamp value is not available for the

frame (for example, because the receive FIFO was full before the time stamp could be written to it), the DMA writes all ones to RDES2 and RDES3. Otherwise (that is, if time stamping is not enabled), RDES2 and RDES3 remain unchanged.





### Receive descriptor acquisition

The receive engine always attempts to acquire an extra descriptor in anticipation of an incoming frame. Descriptor acquisition is attempted if any of the following conditions is/are satisfied:

- The receive Start/Stop bit (ETH\_DMAOMR register[1]) has been set immediately after the DMA has been placed in the Run state.
- The data buffer of the current descriptor is full before the end of the frame currently being transferred
- The controller has completed frame reception, but the current receive descriptor has not yet been closed.
- The receive process has been suspended because of a CPU-owned buffer (RDES0[31] = 0) and a new frame is received.
- A Receive poll demand has been issued.

### Receive frame processing

The MAC transfers the received frames to the STM32F4xx memory only when the frame passes the address filter and the frame size is greater than or equal to the configurable threshold bytes set for the Receive FIFO, or when the complete frame is written to the FIFO in Store-and-forward mode. If the frame fails the address filtering, it is dropped in the MAC block itself (unless Receive All ETH\_MACFFR [31] bit is set). Frames that are shorter than 64 bytes, because of collision or premature termination, can be purged from the Receive FIFO. After 64 (configurable threshold) bytes have been received, the DMA block begins transferring the frame data to the receive buffer pointed to by the current descriptor. The DMA sets the first descriptor (RDES0[9]) after the DMA AHB Interface becomes ready to receive a data transfer (if DMA is not fetching transmit data from the memory), to delimit the frame. The descriptors are released when the OWN (RDES0[31]) bit is reset to 0, either as the data buffer fills up or as the last segment of the frame is transferred to the receive buffer. If the frame is contained in a single descriptor, both the last descriptor (RDES0[8]) and first descriptor (RDES0[9]) bits are set. The DMA fetches the next descriptor, sets the last descriptor (RDES0[8]) bit, and releases the RDES0 status bits in the previous frame descriptor. Then the DMA sets the receive interrupt bit (ETH\_DMASR register [6]). The same process repeats unless the DMA encounters a descriptor flagged as being owned by the CPU. If this occurs, the receive process sets the receive buffer unavailable bit (ETH DMASR register[7]) and then enters the Suspend state. The position in the receive list is retained.

### Receive process suspended

If a new receive frame arrives while the receive process is in Suspend state, the DMA refetches the current descriptor in the STM32F4xx memory. If the descriptor is now owned by the DMA, the receive process re-enters the Run state and starts frame reception. If the descriptor is still owned by the host, by default, the DMA discards the current frame at the top of the Rx FIFO and increments the missed frame counter. If more than one frame is stored in the Rx FIFO, the process repeats. The discarding or flushing of the frame at the top of the Rx FIFO can be avoided by setting the DMA Operation mode register bit 24 (DFRF). In such conditions, the receive process sets the receive buffer unavailable status bit and returns to the Suspend state.

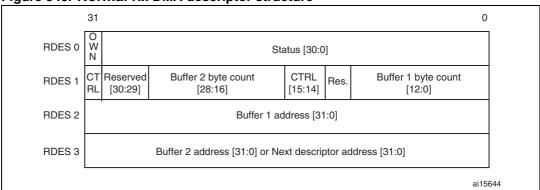


### **Normal Rx DMA descriptors**

The normal receive descriptor structure consists of four 32-bit words (16 bytes). These are shown in *Figure 346*. The bit descriptions of RDES0, RDES1, RDES2 and RDES3 are given below.

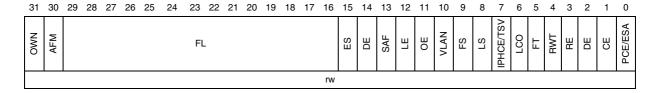
Note that enhanced descriptors must be used if time stamping is activated (TSE=1, ETH\_PTPTSCR bit 0) or if IPv4 checksum offload is activated (IPCO=1, ETH\_MACCR bit 10).

Figure 346. Normal Rx DMA descriptor structure



### RDES0: Receive descriptor Word0

RDES0 contains the received frame status, the frame length and the descriptor ownership information.



### Bit 31 OWN: Own bit

When set, this bit indicates that the descriptor is owned by the DMA of the MAC Subsystem. When this bit is reset, it indicates that the descriptor is owned by the Host. The DMA clears this bit either when it completes the frame reception or when the buffers that are associated with this descriptor are full.

### Bit 30 AFM: Destination address filter fail

When set, this bit indicates a frame that failed the DA filter in the MAC Core.

#### Bits 29:16 FL: Frame length

These bits indicate the byte length of the received frame that was transferred to host memory (including CRC). This field is valid only when last descriptor (RDES0[8]) is set and descriptor error (RDES0[14]) is reset.

This field is valid when last descriptor (RDES0[8]) is set. When the last descriptor and error summary bits are not set, this field indicates the accumulated number of bytes that have been transferred for the current frame.

#### Bit 15 ES: Error summary

Indicates the logical OR of the following bits:

RDES0[1]: CRC error RDES0[3]: Receive error RDES0[4]: Watchdog timeout RDES0[6]: Late collision

RDES0[7]: Giant frame (This is not applicable when RDES0[7] indicates an IPV4 header

checksum error.)

RDES0[11]: Overflow error RDES0[14]: Descriptor error.

This field is valid only when the last descriptor (RDES0[8]) is set.

#### Bit 14 DE: Descriptor error

When set, this bit indicates a frame truncation caused by a frame that does not fit within the current descriptor buffers, and that the DMA does not own the next descriptor. The frame is truncated. This field is valid only when the last descriptor (RDES0[8]) is set.

#### Bit 13 SAF: Source address filter fail

When set, this bit indicates that the SA field of frame failed the SA filter in the MAC Core.

#### Bit 12 LE: Length error

When set, this bit indicates that the actual length of the received frame does not match the value in the Length/ Type field. This bit is valid only when the Frame type (RDES0[5]) bit is reset.

### Bit 11 OE: Overflow error

When set, this bit indicates that the received frame was damaged due to buffer overflow.

### Bit 10 VLAN: VLAN tag

When set, this bit indicates that the frame pointed to by this descriptor is a VLAN frame tagged by the MAC core.

### Bit 9 FS: First descriptor

When set, this bit indicates that this descriptor contains the first buffer of the frame. If the size of the first buffer is 0, the second buffer contains the beginning of the frame. If the size of the second buffer is also 0, the next descriptor contains the beginning of the frame.

#### Bit 8 LS: Last descriptor

When set, this bit indicates that the buffers pointed to by this descriptor are the last buffers of the frame.

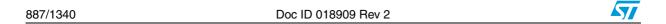
### Bit 7 IPHCE/TSV: IPv header checksum error / time stamp valid

If IPHCE is set, it indicates an error in the IPv4 or IPv6 header. This error can be due to inconsistent Ethernet Type field and IP header Version field values, a header checksum mismatch in IPv4, or an Ethernet frame lacking the expected number of IP header bytes. This bit can take on special meaning as specified in *Table 150*.

If enhanced descriptor format is enabled (EDFE=1, bit 7 of ETH\_DMABMR), this bit takes on the TSV function (otherwise it is IPHCE). When TSV is set, it indicates that a snapshot of the timestamp is written in descriptor words 6 (RDES6) and 7 (RDES7). TSV is valid only when the Last descriptor bit (RDES0[8]) is set.

### Bit 6 LCO: Late collision

When set, this bit indicates that a late collision has occurred while receiving the frame in Halfduplex mode.



#### Bit 5 FT: Frame type

When set, this bit indicates that the Receive frame is an Ethernet-type frame (the LT field is greater than or equal to 0x0600). When this bit is reset, it indicates that the received frame is an IEEE802.3 frame. This bit is not valid for Runt frames less than 14 bytes. When the normal descriptor format is used (ETH\_DMABMR EDFE=0), FT can take on special meaning as specified in *Table 150*.

### Bit 4 RWT: Receive watchdog timeout

When set, this bit indicates that the Receive watchdog timer has expired while receiving the current frame and the current frame is truncated after the watchdog timeout.

#### Bit 3 RE: Receive error

When set, this bit indicates that the RX\_ERR signal is asserted while RX\_DV is asserted during frame reception.

### Bit 2 DE: Dribble bit error

When set, this bit indicates that the received frame has a non-integer multiple of bytes (odd nibbles). This bit is valid only in MII mode.

#### Bit 1 CE: CRC error

When set, this bit indicates that a cyclic redundancy check (CRC) error occurred on the received frame. This field is valid only when the last descriptor (RDES0[8]) is set.

#### Bit 0 PCE/ESA: Payload checksum error / extended status available

When set, it indicates that the TCP, UDP or ICMP checksum the core calculated does not match the received encapsulated TCP, UDP or ICMP segment's Checksum field. This bit is also set when the received number of payload bytes does not match the value indicated in the Length field of the encapsulated IPv4 or IPv6 datagram in the received Ethernet frame. This bit can take on special meaning as specified in *Table 150*.

If the enhanced descriptor format is enabled (EDFE=1, bit 7 in ETH\_DMABMR), this bit takes on the ESA function (otherwise it is PCE). When ESA is set, it indicates that the extended status is available in descriptor word 4 (RDES4). ESA is valid only when the last descriptor bit (RDES0[8]) is set.

Bits 5, 7, and 0 reflect the conditions discussed in *Table 150*.

Table 150. Receive descriptor 0 - encoding for bits 7, 5 and 0 (normal descriptor format only, EDFE=0)

Bit 5: frame type	Bit 7: IPC checksum error	Bit 0: payload checksum error	Frame status
0	0	0	IEEE 802.3 Type frame (Length field value is less than 0x0600.)
1	0	0	IPv4/IPv6 Type frame, no checksum error detected
1	0	1	IPv4/IPv6 Type frame with a payload checksum error (as described for PCE) detected
1	1	0	IPv4/IPv6 Type frame with an IP header checksum error (as described for IPC CE) detected
1	1	1	IPv4/IPv6 Type frame with both IP header and payload checksum errors detected
0	0	1	IPv4/IPv6 Type frame with no IP header checksum error and the payload check bypassed, due to an unsupported payload

Table 150. Receive descriptor 0 - encoding for bits 7, 5 and 0 (normal descriptor format only, EDFE=0) (continued)

Bit 5: frame type	Bit 7: IPC checksum error	Bit 0: payload checksum error	Frame status
0	1	1	A Type frame that is neither IPv4 or IPv6 (the checksum offload engine bypasses checksum completely.)
0	1	0	Reserved

### RDES1: Receive descriptor Word1

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
DIC	RE	3S2						R	BS2							RER	RCH	erved						F	RBS						
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	Res	rw	rw	rw	rw	rw	rw	rw						

### Bit 31 DIC: Disable interrupt on completion

When set, this bit prevents setting the Status register's RS bit (CSR5[6]) for the received frame ending in the buffer indicated by this descriptor. This, in turn, disables the assertion of the interrupt to Host due to RS for that frame.

Bits 30:29 Reserved, must be kept at reset value.

#### Bits 28:16 RBS2: Receive buffer 2 size

These bits indicate the second data buffer size, in bytes. The buffer size must be a multiple of 4, 8, or 16, depending on the bus widths (32, 64 or 128, respectively), even if the value of RDES3 (buffer2 address pointer) is not aligned to bus width. If the buffer size is not an appropriate multiple of 4, 8 or 16, the resulting behavior is undefined. This field is not valid if RDES1 [14] is set.

#### Bit 15 RER: Receive end of ring

When set, this bit indicates that the descriptor list reached its final descriptor. The DMA returns to the base address of the list, creating a descriptor ring.

#### Bit 14 RCH: Second address chained

When set, this bit indicates that the second address in the descriptor is the next descriptor address rather than the second buffer address. When this bit is set, RBS2 (RDES1[28:16]) is a "don't care" value. RDES1[15] takes precedence over RDES1[14].

Bit 13 Reserved, must be kept at reset value.

### Bits 12:0 RBS1: Receive buffer 1 size

Indicates the first data buffer size in bytes. The buffer size must be a multiple of 4, 8 or 16, depending upon the bus widths (32, 64 or 128), even if the value of RDES2 (buffer1 address pointer) is not aligned. When the buffer size is not a multiple of 4, 8 or 16, the resulting behavior is undefined. If this field is 0, the DMA ignores this buffer and uses Buffer 2 or next descriptor depending on the value of RCH (bit 14).

### RDES2: Receive descriptor Word2

RDES2 contains the address pointer to the first data buffer in the descriptor, or it contains time stamp data.

Bits 31:0 RBAP1 / RTSL: Receive buffer 1 address pointer / Receive frame time stamp low

These bits take on two different functions: the application uses them to indicate to the DMA where to store the data in memory, and then after transferring all the data the DMA may use these bits to pass back time stamp data.

RBAP1: When the software makes this descriptor available to the DMA (at the moment that the OWN bit is set to 1 in RDES0), these bits indicate the physical address of Buffer 1. There are no limitations on the buffer address alignment except for the following condition: the DMA uses the configured value for its address generation when the RDES2 value is used to store the start of frame. Note that the DMA performs a write operation with the RDES2[3/2/1:0] bits as 0 during the transfer of the start of frame but the frame data is shifted as per the actual Buffer address pointer. The DMA ignores RDES2[3/2/1:0] (corresponding to bus width of 128/64/32) if the address pointer is to a buffer where the middle or last part of the frame is stored.

**RTSL:** Before it clears the OWN bit in RDES0, the DMA updates this field with the 32 least significant bits of the time stamp captured for the corresponding receive frame (overwriting the value for RBAP1). This field has the time stamp only if time stamping is activated for this frame and if the Last segment control bit (LS) in the descriptor is set.

### RDES3: Receive descriptor Word3

RDES3 contains the address pointer either to the second data buffer in the descriptor or to the next descriptor, or it contains time stamp data.

Bits 31:0 RBAP2 / RTSH: Receive buffer 2 address pointer (next descriptor address) / Receive frame time stamp high

These bits take on two different functions: the application uses them to indicate to the DMA the location of where to store the data in memory, and then after transferring all the data the DMA may use these bits to pass back time stamp data.

**RBAP1:** When the software makes this descriptor available to the DMA (at the moment that the OWN bit is set to 1 in RDES0), these bits indicate the physical address of buffer 2 when a descriptor ring structure is used. If the second address chained (RDES1 [24]) bit is set, this address contains the pointer to the physical memory where the next descriptor is present. If RDES1 [24] is set, the buffer (next descriptor) address pointer must be bus width-aligned (RDES3[3, 2, or 1:0] = 0, corresponding to a bus width of 128, 64 or 32. LSBs are ignored internally.) However, when RDES1 [24] is reset, there are no limitations on the RDES3 value, except for the following condition: the DMA uses the configured value for its buffer address generation when the RDES3 value is used to store the start of frame. The DMA ignores RDES3[3, 2, or 1:0] (corresponding to a bus width of 128, 64 or 32) if the address pointer is to a buffer where the middle or last part of the frame is stored.

**RTSH:** Before it clears the OWN bit in RDES0, the DMA updates this field with the 32 most significant bits of the time stamp captured for the corresponding receive frame (overwriting the value for RBAP2). This field has the time stamp only if time stamping is activated and if the Last segment control bit (LS) in the descriptor is set.

### Enhanced Rx DMA descriptors format with IEEE1588 time stamp

Enhanced descriptors (enabled with EDFE=1, ETHDMABMR bit 7), must be used if time stamping is activated (TSE=1, ETH\_PTPTSCR bit 0) or if IPv4 checksum offload is activated (IPCO=1, ETH\_MACCR bit 10).

Enhanced descriptors comprise eight 32-bit words, twice the size of normal descriptors. RDES0, RDES1, RDES2 and RDES3 have the same definitions as for normal receive descriptors (refer to *Normal Rx DMA descriptors*). RDES4 contains extended status while RDES6 and RDES7 hold the time stamp. RDES4, RDES5, RDES6 and RDES7 are defined below.

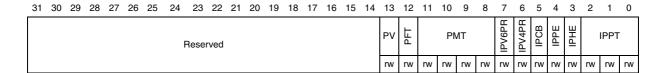
When the Enhanced descriptor mode is selected, the software needs to allocate 32 bytes (8 DWORDS) of memory for every descriptor. When time stamping or IPv4 checksum offload are not being used, the enhanced descriptor format may be disabled and the software can use normal descriptors with the default size of 16 bytes.

31 0 O W RDES 0 Status [30:0] Reserved Buffer 2 byte count CTRI Buffer 1 byte count CT RDES 1 Res [30:29] [28:16] [15:14] [12:0] RDES 2 Buffer 1 address [31:0] RDES 3 Buffer 2 address [31:0] or Next descriptor address [31:0] RDES 4 Extended Status [31:0] RDES 5 Reserved RDES 6 Time stamp low [31:0] RDES 7 Time stamp high [31:0] ai17104

Figure 347. Enhanced receive descriptor field format with IEEE1588 time stamp enabled

### RDES4: Receive descriptor Word4

The extended status, shown below, is valid only when there is status related to IPv4 checksum or time stamp available as indicated by bit 0 in RDES0.



Bits 31:14 Reserved, must be kept at reset value.

Bit 13 PV: PTP version

When set, indicates that the received PTP message uses the IEEE 1588 version 2 format. When cleared, it uses version 1 format. This is valid only if the message type is non-zero.

Bit 12 PFT: PTP frame type

When set, this bit indicates that the PTP message is sent directly over Ethernet. When this bit is cleared and the message type is non-zero, it indicates that the PTP message is sent over UDP-IPv4 or UDP-IPv6. The information on IPv4 or IPv6 can be obtained from bits 6 and 7.

### Bits 11:8 PMT: PTP message type

These bits are encoded to give the type of the message received.

- 0000: No PTP message received
- 0001: SYNC (all clock types)
- 0010: Follow\_Up (all clock types)
- 0011: Delay\_Req (all clock types)
- 0100: Delay\_Resp (all clock types)
- 0101: Pdelay\_Req (in peer-to-peer transparent clock) or Announce (in ordinary or boundary clock)
- 0110: Pdelay\_Resp (in peer-to-peer transparent clock) or Management (in ordinary or boundary clock)
- 0111: Pdelay\_Resp\_Follow\_Up (in peer-to-peer transparent clock) or Signaling (for ordinary or boundary clock)
- 1xxx Reserved

#### Bit 7 IPV6PR: IPv6 packet received

When set, this bit indicates that the received packet is an IPv6 packet.

### Bit 6 IPV4PR: IPv4 packet received

When set, this bit indicates that the received packet is an IPv4 packet.

#### Bit 5 IPCB: IP checksum bypassed

When set, this bit indicates that the checksum offload engine is bypassed.

### Bit 4 IPPE: IP payload error

When set, this bit indicates that the 16-bit IP payload checksum (that is, the TCP, UDP, or ICMP checksum) that the core calculated does not match the corresponding checksum field in the received segment. It is also set when the TCP, UDP, or ICMP segment length does not match the payload length value in the IP Header field.

### Bit 3 IPHE: IP header error

When set, this bit indicates either that the 16-bit IPv4 header checksum calculated by the core does not match the received checksum bytes, or that the IP datagram version is not consistent with the Ethernet Type value.

### Bits 2:0 IPPT: IP payload type

if IPv4 checksum offload is activated (IPCO=1, ETH\_MACCR bit 10), these bits indicate the type of payload encapsulated in the IP datagram. These bits are '00' if there is an IP header error or fragmented IP.

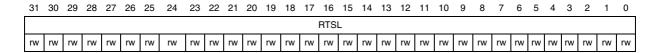
- 000: Unknown or did not process IP payload
- 001: UDP
- 010: TCP
- 011: ICMP
- 1xx: Reserved

### • RDES5: Receive descriptor Word5

Reserved.

### **RDES6: Receive descriptor Word6**

The table below describes the fields that have different meaning for RDES6 when the receive descriptor is closed and time stamping is enabled.



Bits 31:0 RTSL: Receive frame time stamp low

The DMA updates this field with the 32 least significant bits of the time stamp captured for the corresponding receive frame. The DMA updates this field only for the last descriptor of the receive frame indicated by last descriptor status bit (RDES0[8]). When this field and the RTSH field in RDES7 show all ones, the time stamp must be treated as corrupt.

### **RDES7: Receive descriptor Word7**

The table below describes the fields that have a different meaning for RDES7 when the receive descriptor is closed and time stamping is enabled.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
															RTSI	+															
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw															

Bits 31:0 RTSH: Receive frame time stamp high

The DMA updates this field with the 32 most significant bits of the time stamp captured for the corresponding receive frame. The DMA updates this field only for the last descriptor of the receive frame indicated by last descriptor status bit (RDES0[8]).

When this field and RDES7's RTSL field show all ones, the time stamp must be treated as corrupt.

#### 28.6.9 **DMA** interrupts

Interrupts can be generated as a result of various events. The ETH\_DMASR register contains all the bits that might cause an interrupt. The ETH DMAIER register contains an enable bit for each of the events that can cause an interrupt.

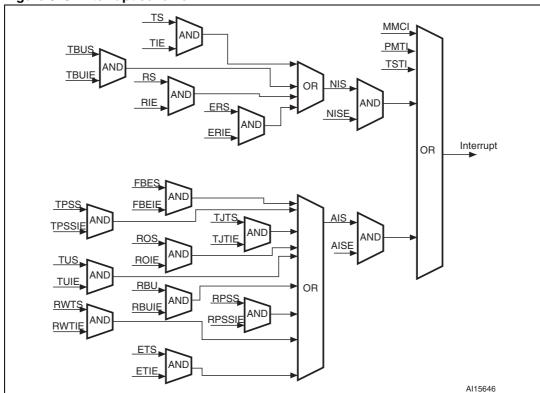
There are two groups of interrupts, Normal and Abnormal, as described in the ETH DMASR register. Interrupts are cleared by writing a 1 to the corresponding bit position. When all the enabled interrupts within a group are cleared, the corresponding summary bit is cleared. If the MAC core is the cause for assertion of the interrupt, then any of the TSTS or PMTS bits in the ETH DMASR register is set high.

Interrupts are not queued and if the interrupt event occurs before the driver has responded to it, no additional interrupts are generated. For example, the Receive Interrupt bit (ETH\_DMASR register [6]) indicates that one or more frames were transferred to the STM32F4xx buffer. The driver must scan all descriptors, from the last recorded position to the first one owned by the DMA.

An interrupt is generated only once for simultaneous, multiple events. The driver must scan the ETH DMASR register for the cause of the interrupt. The interrupt is not generated again unless a new interrupting event occurs, after the driver has cleared the appropriate bit in the ETH\_DMASR register. For example, the controller generates a Receive interrupt (ETH\_DMASR register[6]) and the driver begins reading the ETH\_DMASR register. Next,

receive buffer unavailable (ETH\_DMASR register[7]) occurs. The driver clears the Receive interrupt. Even then, a new interrupt is generated, due to the active or pending Receive buffer unavailable interrupt.

Figure 348. Interrupt scheme



# 28.7 Ethernet interrupts

The Ethernet controller has two interrupt vectors: one dedicated to normal Ethernet operations and the other, used only for the Ethernet wakeup event (with wakeup frame or Magic Packet detection) when it is mapped on EXTI Ilne19.

The first Ethernet vector is reserved for interrupts generated by the MAC and the DMA as listed in the *MAC interrupts* and *DMA interrupts* sections.

The second vector is reserved for interrupts generated by the PMT on wakeup events. The mapping of a wakeup event on EXTI line19 causes the STM32F4xx to exit the low power mode, and generates an interrupt.

When an Ethernet wakeup event mapped on EXTI Line19 occurs and the MAC PMT interrupt is enabled and the EXTI Line19 interrupt, with detection on rising edge, is also enabled, both interrupts are generated.

A watchdog timer (see ETH\_DMARSWTR register) is given for flexible control of the RS bit (ETH\_DMASR register). When this watchdog timer is programmed with a non-zero value, it gets activated as soon as the RxDMA completes a transfer of a received frame to system memory without asserting the Receive Status because it is not enabled in the corresponding Receive descriptor (RDES1[31]). When this timer runs out as per the programmed value, the RS bit is set and the interrupt is asserted if the corresponding RIE is enabled in the

ETH\_DMAIER register. This timer is disabled before it runs out, when a frame is transferred to memory and the RS is set because it is enabled for that descriptor.

Note:

Reading the PMT control and status register automatically clears the Wakeup Frame Received and Magic Packet Received PMT interrupt flags. However, since the registers for these flags are in the CLK\_RX domain, there may be a significant delay before this update is visible by the firmware. The delay is especially long when the RX clock is slow (in 10 Mbit mode) and when the AHB bus is high-frequency.

Since interrupt requests from the PMT to the CPU are based on the same registers in the CLK\_RX domain, the CPU may spuriously call the interrupt routine a second time even after reading PMT\_CSR. Thus, it may be necessary that the firmware polls the Wakeup Frame Received and Magic Packet Received bits and exits the interrupt service routine only when they are found to be at '0'.

# 28.8 Ethernet register descriptions

The peripheral registers can be accessed by bytes (8-bit), half-words (16-bit) or words (32-bits).

### 28.8.1 MAC register description

### Ethernet MAC configuration register (ETH\_MACCR)

Address offset: 0x0000
Reset value: 0x0000 8000

The MAC configuration register is the operation mode register of the MAC. It establishes receive and transmit operating modes.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	. 8	7	6	5	4	3	. 2	1	0	
		Rese	ervec	i		CSTF	served	WD	۵r	2	D 2 D		IFG		CSD	served	FES	ROD	МЛ	DM	IPCO	RD	served	APCS	ā	DL.	DC	TE	BE	5	served	
						rw	Re	rw	rw	ä			rw		rw	Re	rw	rw	rw	rw	rw	rw	Re	rw	r	w	rw	rw	rw	_	r C	

Bits 31:246 Reserved, must be kept at reset value.

### Bits 25 CSTF: CRC stripping for Type frames

When set, the last 4 bytes (FCS) of all frames of Ether type (type field greater than 0x0600) will be stripped and dropped before forwarding the frame to the application.

Bits 24 Reserved, must be kept at reset value.

### Bit 23 WD: Watchdog disable

When this bit is set, the MAC disables the watchdog timer on the receiver, and can receive frames of up to 16 384 bytes.

When this bit is reset, the MAC allows no more than 2 048 bytes of the frame being received and cuts off any bytes received after that.

#### Bit 22 JD: Jabber disable

When this bit is set, the MAC disables the jabber timer on the transmitter, and can transfer frames of up to 16 384 bytes.

When this bit is reset, the MAC cuts off the transmitter if the application sends out more than 2 048 bytes of data during transmission.

Bits 21:20 Reserved, must be kept at reset value.

### Bits 19:17 IFG: Interframe gap

These bits control the minimum interframe gap between frames during transmission.

000: 96 bit times 001: 88 bit times 010: 80 bit times

111: 40 bit times

Note: In Half-duplex mode, the minimum IFG can be configured for 64 bit times (IFG = 100) only. Lower values are not considered.

#### Bit 16 CSD: Carrier sense disable

When set high, this bit makes the MAC transmitter ignore the MII CRS signal during frame transmission in Half-duplex mode. No error is generated due to Loss of Carrier or No Carrier during such transmission.

When this bit is low, the MAC transmitter generates such errors due to Carrier Sense and even aborts the transmissions.

#### Bit 15 Reserved, must be kept at reset value.

#### Bit 14 FES: Fast Ethernet speed

Indicates the speed in Fast Ethernet (MII) mode:

0: 10 Mbit/s 1: 100 Mbit/s

#### Bit 13 ROD: Receive own disable

When this bit is set, the MAC disables the reception of frames in Half-duplex mode. When this bit is reset, the MAC receives all packets that are given by the PHY while transmitting.

This bit is not applicable if the MAC is operating in Full-duplex mode.

#### Bit 12 LM: Loopback mode

When this bit is set, the MAC operates in loopback mode at the MII. The MII receive clock input (RX\_CLK) is required for the loopback to work properly, as the transmit clock is not looped-back internally.

### Bit 11 DM: Duplex mode

When this bit is set, the MAC operates in a Full-duplex mode where it can transmit and receive simultaneously.

#### Bit 10 IPCO: IPv4 checksum offload

When set, this bit enables IPv4 checksum checking for received frame payloads' TCP/UDP/ICMP headers. When this bit is reset, the checksum offload function in the receiver is disabled and the corresponding PCE and IP HCE status bits (see *Table 147 on page 852*) are always cleared.

### Bit 9 RD: Retry disable

When this bit is set, the MAC attempts only 1 transmission. When a collision occurs on the MII, the MAC ignores the current frame transmission and reports a Frame Abort with excessive collision error in the transmit frame status.

When this bit is reset, the MAC attempts retries based on the settings of BL.

Note: This bit is applicable only in the Half-duplex mode.

Bit 8 Reserved, must be kept at reset value.

### Bit 7 APCS: Automatic pad/CRC stripping

When this bit is set, the MAC strips the Pad/FCS field on incoming frames only if the length's field value is less than or equal to 1 500 bytes. All received frames with length field greater than or equal to 1 501 bytes are passed on to the application without stripping the Pad/FCS field.

When this bit is reset, the MAC passes all incoming frames unmodified.

#### Bits 6:5 BL: Back-off limit

The Back-off limit determines the random integer number (r) of slot time delays (4 096 bit times for 1000 Mbit/s and 512 bit times for 10/100 Mbit/s) the MAC waits before rescheduling a transmission attempt during retries after a collision.

Note: This bit is applicable only to Half-duplex mode.

```
00: k = min (n, 10)
01: k = min (n, 8)
10: k = min (n, 4)
11: k = min (n, 1),
```

where n = retransmission attempt. The random integer r takes the value in the range  $0 \le r < 2^k$ 

#### Bit 4 DC: Deferral check

When this bit is set, the deferral check function is enabled in the MAC. The MAC issues a Frame Abort status, along with the excessive deferral error bit set in the transmit frame status when the transmit state machine is deferred for more than 24 288 bit times in 10/100-Mbit/s mode. Deferral begins when the transmitter is ready to transmit, but is prevented because of an active CRS (carrier sense) signal on the MII. Defer time is not cumulative. If the transmitter defers for 10 000 bit times, then transmits, collides, backs off, and then has to defer again after completion of back-off, the deferral timer resets to 0 and restarts. When this bit is reset, the deferral check function is disabled and the MAC defers until the CRS signal goes inactive. This bit is applicable only in Half-duplex mode.

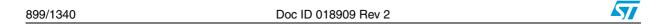
### Bit 3 TE: Transmitter enable

When this bit is set, the transmit state machine of the MAC is enabled for transmission on the MII. When this bit is reset, the MAC transmit state machine is disabled after the completion of the transmission of the current frame, and does not transmit any further frames.

### Bit 2 RE: Receiver enable

When this bit is set, the receiver state machine of the MAC is enabled for receiving frames from the MII. When this bit is reset, the MAC receive state machine is disabled after the completion of the reception of the current frame, and will not receive any further frames from the MII.

Bits 1:0 Reserved, must be kept at reset value.



### **Ethernet MAC frame filter register (ETH\_MACFFR)**

Address offset: 0x0004 Reset value: 0x0000 0000

The MAC frame filter register contains the filter controls for receiving frames. Some of the controls from this register go to the address check block of the MAC, which performs the first level of address filtering. The second level of filtering is performed on the incoming frame, based on other controls such as pass bad frames and pass control frames.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

RA	Reserved	HPF	SAF	SAIF	ח	7	BFD	PAM	DAIF	MH	ПΗ	PM
rw		rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

#### Bit 31 RA: Receive all

When this bit is set, the MAC receiver passes all received frames on to the application, irrespective of whether they have passed the address filter. The result of the SA/DA filtering is updated (pass or fail) in the corresponding bits in the receive status word. When this bit is reset, the MAC receiver passes on to the application only those frames that have passed the SA/DA address filter.

#### Bits 30:11 Reserved, must be kept at reset value.

#### Bit 10 HPF: Hash or perfect filter

When this bit is set and if the HM or HU bit is set, the address filter passes frames that match either the perfect filtering or the hash filtering.

When this bit is cleared and if the HU or HM bit is set, only frames that match the Hash filter are passed.

### Bit 9 SAF: Source address filter

The MAC core compares the SA field of the received frames with the values programmed in the enabled SA registers. If the comparison matches, then the SAMatch bit in the RxStatus word is set high. When this bit is set high and the SA filter fails, the MAC drops the frame. When this bit is reset, the MAC core forwards the received frame to the application. It also forwards the updated SA Match bit in RxStatus depending on the SA address comparison.

#### Bit 8 SAIF: Source address inverse filtering

When this bit is set, the address check block operates in inverse filtering mode for the SA address comparison. The frames whose SA matches the SA registers are marked as failing the SA address filter.

When this bit is reset, frames whose SA does not match the SA registers are marked as failing the SA address filter.

#### Bits 7:6 PCF: Pass control frames

These bits control the forwarding of all control frames (including unicast and multicast PAUSE frames). Note that the processing of PAUSE control frames depends only on RFCE in Flow Control Register[2].

00: MAC prevents all control frames from reaching the application

01: MAC forwards all control frames to application except Pause control frames

10: MAC forwards all control frames to application even if they fail the address filter

11: MAC forwards control frames that pass the address filter.

These bits control the forwarding of all control frames (including unicast and multicast PAUSE frames). Note that the processing of PAUSE control frames depends only on RFCE in Flow Control Register[2].

00 or 01: MAC prevents all control frames from reaching the application

10: MAC forwards all control frames to application even if they fail the address filter

11: MAC forwards control frames that pass the address filter.

#### Bit 5 BFD: Broadcast frames disable

When this bit is set, the address filters filter all incoming broadcast frames.

When this bit is reset, the address filters pass all received broadcast frames.

#### Bit 4 PAM: Pass all multicast

When set, this bit indicates that all received frames with a multicast destination address (first bit in the destination address field is '1') are passed.

When reset, filtering of multicast frame depends on the HM bit.

#### Bit 3 DAIF: Destination address inverse filtering

When this bit is set, the address check block operates in inverse filtering mode for the DA address comparison for both unicast and multicast frames.

When reset, normal filtering of frames is performed.

#### Bit 2 HM: Hash multicast

When set, MAC performs destination address filtering of received multicast frames according to the hash table.

When reset, the MAC performs a perfect destination address filtering for multicast frames, that is, it compares the DA field with the values programmed in DA registers.

#### Bit 1 HU: Hash unicast

When set, MAC performs destination address filtering of unicast frames according to the hash table.

When reset, the MAC performs a perfect destination address filtering for unicast frames, that is, it compares the DA field with the values programmed in DA registers.

#### Bit 0 PM: Promiscuous mode

When this bit is set, the address filters pass all incoming frames regardless of their destination or source address. The SA/DA filter fails status bits in the receive status word are always cleared when PM is set.

# Ethernet MAC hash table high register (ETH MACHTHR)

Address offset: 0x0008 Reset value: 0x0000 0000

The 64-bit Hash table is used for group address filtering. For hash filtering, the contents of the destination address in the incoming frame are passed through the CRC logic, and the upper 6 bits in the CRC register are used to index the contents of the Hash table. This CRC

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is a 32-bit value coded by the following polynomial (for more details refer to *Section 28.5.3: MAC frame reception*):

$$G(x) = x^{32} + x^{26} + x^{23} + x^{22} + x^{16} + x^{12} + x^{11} + x^{10} + x^8 + x^7 + x^5 + x^4 + x^2 + x + 1$$

The most significant bit determines the register to be used (hash table high/hash table low), and the other 5 bits determine which bit within the register. A hash value of 0b0 0000 selects bit 0 in the selected register, and a value of 0b1 1111 selects bit 31 in the selected register.

For example, if the DA of the incoming frame is received as 0x1F52 419C B6AF (0x1F is the first byte received on the MII interface), then the internally calculated 6-bit Hash value is 0x2C and the HTH register bit[12] is checked for filtering. If the DA of the incoming frame is received as 0xA00A 9800 0045, then the calculated 6-bit Hash value is 0x07 and the HTL register bit[7] is checked for filtering.

If the corresponding bit value in the register is 1, the frame is accepted. Otherwise, it is rejected. If the PAM (pass all multicast) bit is set in the ETH\_MACFFR register, then all multicast frames are accepted regardless of the multicast hash values.

The Hash table high register contains the higher 32 bits of the multicast Hash table.

Bits 31:0 HTH: Hash table high

This field contains the upper 32 bits of Hash table.

# Ethernet MAC hash table low register (ETH\_MACHTLR)

Address offset: 0x000C Reset value: 0x0000 0000

The Hash table low register contains the lower 32 bits of the multi-cast Hash table.

Bits 31:0 HTL: Hash table low

This field contains the lower 32 bits of the Hash table.

# Ethernet MAC MII address register (ETH MACMIIAR)

Address offset: 0x0010 Reset value: 0x0000 0000

The MII address register controls the management cycles to the external PHY through the management interface.

3	1	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
																			PA					MR			ved		CR		MW	МВ
								Rese	ervec	i							rw	Reser	rw	rw	rw	rw	rc_ w1									



Bits 31:16 Reserved, must be kept at reset value.

Bits 15:11 PA: PHY address

This field tells which of the 32 possible PHY devices are being accessed.

Bits 10:6 MR: MII register

These bits select the desired MII register in the selected PHY device.

Bit 5 Reserved, must be kept at reset value.

#### Bits 4:2 CR: Clock range

The CR clock range selection determines the HCLK frequency and is used to decide the frequency of the MDC clock:

Selection	HCLK	MDC Clock
000	60-100 MHz	HCLK/42
001	100-150 MHz	HCLK/62
010	20-35 MHz	HCLK/16
011	35-60 MHz	HCLK/26
100	150-168 MHz	HCLK/102
101, 110, 111	Reserved	_

#### Bit 1 MW: MII write

When set, this bit tells the PHY that this will be a Write operation using the MII Data register. If this bit is not set, this will be a Read operation, placing the data in the MII Data register.

#### Bit 0 MB: MII busy

This bit should read a logic 0 before writing to ETH\_MACMIIAR and ETH\_MACMIIDR. This bit must also be reset to 0 during a Write to ETH\_MACMIIAR. During a PHY register access, this bit is set to 0b1 by the application to indicate that a read or write access is in progress. ETH\_MACMIIDR (MII Data) should be kept valid until this bit is cleared by the MAC during a PHY Write operation. The ETH\_MACMIIDR is invalid until this bit is cleared by the MAC during a PHY Read operation. The ETH\_MACMIIAR (MII Address) should not be written to until this bit is cleared.

# Ethernet MAC MII data register (ETH\_MACMIIDR)

Address offset: 0x0014
Reset value: 0x0000 0000

The MAC MII Data register stores write data to be written to the PHY register located at the address specified in ETH\_MACMIIAR. ETH\_MACMIIDR also stores read data from the PHY register located at the address specified by ETH MACMIIAR.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved								М	D							
	rw															

Bits 31:16 Reserved, must be kept at reset value.

#### Bits 15:0 MD: MII data

This contains the 16-bit data value read from the PHY after a Management Read operation, or the 16-bit data value to be written to the PHY before a Management Write operation.

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# Ethernet MAC flow control register (ETH MACFCR)

Address offset: 0x0018 Reset value: 0x0000 0000

The Flow control register controls the generation and reception of the control (Pause Command) frames by the MAC. A write to a register with the Busy bit set to '1' causes the MAC to generate a pause control frame. The fields of the control frame are selected as specified in the 802.3x specification, and the Pause Time value from this register is used in the Pause Time field of the control frame. The Busy bit remains set until the control frame is transferred onto the cable. The Host must make sure that the Busy bit is cleared before writing to the register.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

							Р	Т								Reserved	ZQPD	erved	Pl	LT	OPFD	RFCE	TFCE	FCB/ BPA
rw	.1000.100	rw	Res	rw	rw	rw	rw	rw	rc_w1 /rw															

#### Bits 31:16 PT: Pause time

This field holds the value to be used in the Pause Time field in the transmit control frame. If the Pause Time bits is configured to be double-synchronized to the MII clock domain, then consecutive write operations to this register should be performed only after at least 4 clock cycles in the destination clock domain.

Bits 15:8 Reserved, must be kept at reset value.

#### Bit 7 ZQPD: Zero-quanta pause disable

When set, this bit disables the automatic generation of Zero-quanta pause control frames on the deassertion of the flow-control signal from the FIFO layer.

When this bit is reset, normal operation with automatic Zero-quanta pause control frame generation is enabled.

Bit 6 Reserved, must be kept at reset value.

# Bits 5:4 PLT: Pause low threshold

This field configures the threshold of the Pause timer at which the Pause frame is automatically retransmitted. The threshold values should always be less than the Pause Time configured in bits[31:16]. For example, if PT = 100H (256 slot-times), and PLT = 01, then a second PAUSE frame is automatically transmitted if initiated at 228 (256 – 28) slot-times after the first PAUSE frame is transmitted.

Selection Threshold

Pause time minus 4 slot times
Pause time minus 28 slot times
Pause time minus 144 slot times
Pause time minus 256 slot times

Slot time is defined as time taken to transmit 512 bits (64 bytes) on the MII interface.

#### Bit 3 **UPFD:** Unicast pause frame detect

When this bit is set, the MAC detects the Pause frames with the station's unicast address specified in the ETH\_MACA0HR and ETH\_MACA0LR registers, in addition to detecting Pause frames with the unique multicast address.

When this bit is reset, the MAC detects only a Pause frame with the unique multicast address specified in the 802.3x standard.

#### Bit 2 RFCE: Receive flow control enable

When this bit is set, the MAC decodes the received Pause frame and disables its transmitter for a specified (Pause Time) time.

When this bit is reset, the decode function of the Pause frame is disabled.

#### Bit 1 TFCE: Transmit flow control enable

In Full-duplex mode, when this bit is set, the MAC enables the flow control operation to transmit Pause frames. When this bit is reset, the flow control operation in the MAC is disabled, and the MAC does not transmit any Pause frames.

In Half-duplex mode, when this bit is set, the MAC enables the back-pressure operation. When this bit is reset, the back pressure feature is disabled.

#### Bit 0 FCB/BPA: Flow control busy/back pressure activate

This bit initiates a Pause Control frame in Full-duplex mode and activates the back pressure function in Half-duplex mode if TFCE bit is set.

In Full-duplex mode, this bit should be read as 0 before writing to the Flow control register. To initiate a Pause control frame, the Application must set this bit to 1. During a transfer of the Control frame, this bit continues to be set to signify that a frame transmission is in progress. After completion of the Pause control frame transmission, the MAC resets this bit to 0. The Flow control register should not be written to until this bit is cleared.

In Half-duplex mode, when this bit is set (and TFCE is set), back pressure is asserted by the MAC core. During back pressure, when the MAC receives a new frame, the transmitter starts sending a JAM pattern resulting in a collision. When the MAC is configured to Full-duplex mode, the BPA is automatically disabled.

# **Ethernet MAC VLAN tag register (ETH\_MACVLANTR)**

Address offset: 0x001C Reset value: 0x0000 0000

The VLAN tag register contains the IEEE 802.1Q VLAN Tag to identify the VLAN frames. The MAC compares the 13<sup>th</sup> and 14<sup>th</sup> bytes of the receiving frame (Length/Type) with 0x8100, and the following 2 bytes are compared with the VLAN tag; if a match occurs, the received VLAN bit in the receive frame status is set. The legal length of the frame is increased from 1518 bytes to 1522 bytes.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	/	6	5	4	3	2	1	0
						Re	eserv	/ed							VLANTC								VLA	NTI							
															rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

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Bits 31:17 Reserved, must be kept at reset value.

# Bit 16 VLANTC: 12-bit VLAN tag comparison

When this bit is set, a 12-bit VLAN identifier, rather than the complete 16-bit VLAN tag, is used for comparison and filtering. Bits[11:0] of the VLAN tag are compared with the corresponding field in the received VLAN-tagged frame.

When this bit is reset, all 16 bits of the received VLAN frame's fifteenth and sixteenth bytes are used for comparison.

# Bits 15:0 VLANTI: VLAN tag identifier (for receive frames)

This contains the 802.1Q VLAN tag to identify VLAN frames, and is compared to the fifteenth and sixteenth bytes of the frames being received for VLAN frames. Bits[15:13] are the user priority, Bit[12] is the canonical format indicator (CFI) and bits[11:0] are the VLAN tag's VLAN identifier (VID) field. When the VLANTC bit is set, only the VID (bits[11:0]) is used for comparison.

If VLANTI (VLANTI[11:0] if VLANTC is set) is all zeros, the MAC does not check the fifteenth and sixteenth bytes for VLAN tag comparison, and declares all frames with a Type field value of 0x8100 as VLAN frames.

# Ethernet MAC remote wakeup frame filter register (ETH MACRWUFFR)

Address offset: 0x0028 Reset value: 0x0000 0000

This is the address through which the remote wakeup frame filter registers are written/read by the application. The Wakeup frame filter register is actually a pointer to eight (not transparent) such wakeup frame filter registers. Eight sequential write operations to this address with the offset (0x0028) will write all wakeup frame filter registers. Eight sequential read operations from this address with the offset (0x0028) will read all wakeup frame filter registers. This register contains the higher 16 bits of the 7<sup>th</sup> MAC address. Refer to *Remote wakeup frame filter register* section for additional information.

Figure 349. Ethernet MAC remote wakeup frame filter register (ETH\_MACRWUFFR)

Wakeup frame filter reg0				Filter 0 B	yte Mask			
Wakeup frame filter reg1				Filter 1 B	yte Mask			
Wakeup frame filter reg2				Filter 2 B	yte Mask			
Wakeup frame filter reg3				Filter 3 B	yte Mask			
Wakeup frame filter reg4	RSVD	Filter 3 Command	RSVD	Filter 2 Command	RSVD	Filter 1 Command	RSVD	Filter 0 Command
Wakeup frame filter reg5	Filter 3	3 Offset	Filter 2	2 Offset	Filter 1	Offset	Filter (	Offset
Wakeup frame filter reg6		Filter 1 C	CRC - 16			Filter 0 C	CRC - 16	
Wakeup frame filter reg7		Filter 3 C	CRC - 16			Filter 2 (	CRC - 16	
								ai15648

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# Ethernet MAC PMT control and status register (ETH\_MACPMTCSR)

Address offset: 0x002C Reset value: 0x0000 0000

The ETH\_MACPMTCSR programs the request wakeup events and monitors the wakeup

events.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

WFFRPR	Reserved	GU	eserved	WFR	MPR	eserved	WFE	MPE	PD
rs	Res.	rw	Œ	rc_r	rc_r	Œ	rw	rw	rs

#### Bit 31 WFFRPR: Wakeup frame filter register pointer reset

When set, it resets the Remote wakeup frame filter register pointer to 0b000. It is automatically cleared after 1 clock cycle.

Bits 30:10 Reserved, must be kept at reset value.

#### Bit 9 GU: Global unicast

When set, it enables any unicast packet filtered by the MAC (DAF) address recognition to be a wakeup frame.

Bits 8:7 Reserved, must be kept at reset value.

#### Bit 6 WFR: Wakeup frame received

When set, this bit indicates the power management event was generated due to reception of a wakeup frame. This bit is cleared by a read into this register.

#### Bit 5 MPR: Magic packet received

When set, this bit indicates the power management event was generated by the reception of a Magic Packet. This bit is cleared by a read into this register.

Bits 4:3 Reserved, must be kept at reset value.

#### Bit 2 WFE: Wakeup frame enable

When set, this bit enables the generation of a power management event due to wakeup frame reception.

# Bit 1 MPE: Magic Packet enable

When set, this bit enables the generation of a power management event due to Magic Packet reception.

# Bit 0 PD: Power down

When this bit is set, all received frames will be dropped. This bit is cleared automatically when a magic packet or wakeup frame is received, and Power-down mode is disabled. Frames received after this bit is cleared are forwarded to the application. This bit must only be set when either the Magic Packet Enable or Wakeup Frame Enable bit is set high.

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# Ethernet MAC debug register (ETH\_MACDBGR)

Address offset: 0x0034 Reset value: 0x0000 0000

This debug register gives the status of all the main modules of the transmit and receive data paths and the FIFOs. An all-zero status indicates that the MAC core is in Idle state (and FIFOs are empty) and no activity is going on in the data paths.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	TFF	TFNE	Reserved	TFWA	TEDO	L	MTP	G CL	_	MMTEA	Reserved	0	L	Reserved	ودوره	5	RFWRA	Reserved	MSFBWCS		MMRPEA
	ro	ro		ro	ro	ro	ro	ro	ro	ro		ro	ro		ro	ro	ro		ro	ro	ro

Bits 31:26 Reserved, must be kept at reset value.

#### Bit 25 TFF: Tx FIFO full

When high, it indicates that the Tx FIFO is full and hence no more frames will be accepted for transmission.

# Bit 24 TFNE: Tx FIFO not empty

When high, it indicates that the TxFIFO is not empty and has some data left for transmission.

Bit 23 Reserved, must be kept at reset value.

#### Bit 22 TFWA: Tx FIFO write active

When high, it indicates that the TxFIFO write controller is active and transferring data to the TxFIFO.

# Bits 21:20 TFRS: Tx FIFO read status

This indicates the state of the TxFIFO read controller:

00: Idle state

01: Read state (transferring data to the MAC transmitter)

10: Waiting for TxStatus from MAC transmitter

11: Writing the received TxStatus or flushing the TxFIFO

# Bit 19 MTP: MAC transmitter in pause

When high, it indicates that the MAC transmitter is in Pause condition (in full-duplex mode only) and hence will not schedule any frame for transmission

#### Bits 18:17 MTFCS: MAC transmit frame controller status

This indicates the state of the MAC transmit frame controller:

00. Idle

01: Waiting for Status of previous frame or IFG/backoff period to be over

10: Generating and transmitting a Pause control frame (in full duplex mode)

11: Transferring input frame for transmission

#### Bit 16 MMTEA: MAC MII transmit engine active

When high, it indicates that the MAC MII transmit engine is actively transmitting data and that it is not in the Idle state.

Bits 15:10 Reserved, must be kept at reset value.

#### Bits 9:8 RFFL: Rx FIFO fill level

This gives the status of the Rx FIFO fill-level:

00: RxFIFO empty

01: RxFIFO fill-level below flow-control de-activate threshold 10: RxFIFO fill-level above flow-control activate threshold

11: RxFIFO full

Bit 7 Reserved, must be kept at reset value.

#### Bits 6:5 RFRCS: Rx FIFO read controller status

It gives the state of the Rx FIFO read controller:

00: IDLE state

01: Reading frame data

10: Reading frame status (or time-stamp)

11: Flushing the frame data and status

# Bit 4 RFWRA: Rx FIFO write controller active

When high, it indicates that the Rx FIFO write controller is active and transferring a received frame to the FIFO.

Bit 3 Reserved, must be kept at reset value.

# Bits 2:1 MSFRWCS: MAC small FIFO read / write controllers status

When high, these bits indicate the respective active state of the small FIFO read and write controllers of the MAC receive frame controller module.

#### Bit 0 MMRPEA: MAC MII receive protocol engine active

When high, it indicates that the MAC MII receive protocol engine is actively receiving data and is not in the Idle state.

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# Ethernet MAC interrupt status register (ETH\_MACSR)

Address offset: 0x0038 Reset value: 0x0000 0000

The ETH\_MACSR register contents identify the events in the MAC that can generate an

interrupt.

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		Poss	vrvod			TSTS	Posc	erved	MMCTS	MMCRS	MMCS	PMTS		Reserved	
		Rese	riveu			rc_r	nese	ei veu	r	r	r	r		neserveu	

Bits 15:10 Reserved, must be kept at reset value.

# Bit 9 TSTS: Time stamp trigger status

This bit is set high when the system time value equals or exceeds the value specified in the Target time high and low registers. This bit is cleared when this register is read.

Bits 8:7 Reserved, must be kept at reset value.

#### Bit 6 MMCTS: MMC transmit status

This bit is set high whenever an interrupt is generated in the ETH\_MMCTIR Register. This bit is cleared when all the bits in this interrupt register (ETH\_MMCTIR) are cleared.

# Bit 5 MMCRS: MMC receive status

This bit is set high whenever an interrupt is generated in the ETH\_MMCRIR register. This bit is cleared when all the bits in this interrupt register (ETH\_MMCRIR) are cleared.

#### Bit 4 MMCS: MMC status

This bit is set high whenever any of bits 6:5 is set high. It is cleared only when both bits are low.

#### Bit 3 PMTS: PMT status

This bit is set whenever a Magic packet or Wake-on-LAN frame is received in Power-down mode (See bits 5 and 6 in the ETH\_MACPMTCSR register *Ethernet MAC PMT control and status register (ETH\_MACPMTCSR)* on page 907). This bit is cleared when both bits[6:5], of this last register, are cleared due to a read operation to the ETH\_MACPMTCSR register.

Bits 2:0 Reserved, must be kept at reset value.

# Ethernet MAC interrupt mask register (ETH MACIMR)

Address offset: 0x003C Reset value: 0x0000 0000

The ETH\_MACIMR register bits make it possible to mask the interrupt signal due to the

corresponding event in the ETH\_MACSR register.

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		Rese	arvod			TSTIM			Reserved			PMTIM		Reserved	
		nese	erveu			rw			neserveu			rw		neserveu	

Bits 15:10 Reserved, must be kept at reset value.

Bit 9 TSTIM: Time stamp trigger interrupt mask

When set, this bit disables the time stamp interrupt generation.

Bits 8:4 Reserved, must be kept at reset value.

Bit 3 PMTIM: PMT interrupt mask

When set, this bit disables the assertion of the interrupt signal due to the setting of the PMT Status bit in ETH\_MACSR.

Bits 2:0 Reserved, must be kept at reset value.

# Ethernet MAC address 0 high register (ETH MACA0HR)

Address offset: 0x0040 Reset value: 0x0010 FFFF

The MAC address 0 high register holds the upper 16 bits of the 6-byte first MAC address of the station. Note that the first DA byte that is received on the MII interface corresponds to the LS Byte (bits [7:0]) of the MAC address low register. For example, if 0x1122 3344 5566 is received (0x11 is the first byte) on the MII as the destination address, then the MAC address 0 register [47:0] is compared with 0x6655 4433 2211.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

МО	Reserved			-					MAC	A0H							
1		rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bit 31 MO: Always 1.

Bits 30:16 Reserved, must be kept at reset value.

Bits 15:0 MACA0H: MAC address0 high [47:32]

This field contains the upper 16 bits (47:32) of the 6-byte MAC address0. This is used by the MAC for filtering for received frames and for inserting the MAC address in the transmit flow control (Pause) frames.

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# Ethernet MAC address 0 low register (ETH MACA0LR)

Address offset: 0x0044

Reset value: 0xFFFF FFFF

The MAC address 0 low register holds the lower 32 bits of the 6-byte first MAC address of

the station.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

															MAC	A0L															
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw															

# Bits 31:0 MACA0L: MAC address0 low [31:0]

This field contains the lower 32 bits of the 6-byte MAC address0. This is used by the MAC for filtering for received frames and for inserting the MAC address in the transmit flow control (Pause) frames.

# Ethernet MAC address 1 high register (ETH\_MACA1HR)

Address offset: 0x0048

Reset value: 0x0000 FFFF

The MAC address 1 high register holds the upper 16 bits of the 6-byte second MAC address of the station.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

ΑE	SA			ME	3C			Reserved								MAC	A1H							
rw	neserveu	rw	rw	rw	rw	rw	rw	rw	rw	rw														

# Bit 31 AE: Address enable

When this bit is set, the address filters use the MAC address1 for perfect filtering. When this bit is cleared, the address filters ignore the address for filtering.

#### Bit 30 SA: Source address

When this bit is set, the MAC address1[47:0] is used for comparison with the SA fields of the received frame.

When this bit is cleared, the MAC address1[47:0] is used for comparison with the DA fields of the received frame.

# Bits 29:24 MBC: Mask byte control

These bits are mask control bits for comparison of each of the MAC address1 bytes. When they are set high, the MAC core does not compare the corresponding byte of received DA/SA with the contents of the MAC address1 registers. Each bit controls the masking of the bytes as follows:

- Bit 29: ETH\_MACA1HR [15:8]
- Bit 28: ETH\_MACA1HR [7:0]
- Bit 27: ETH\_MACA1LR [31:24]

...

- Bit 24: ETH\_MACA1LR [7:0]

Bits 23:16 Reserved, must be kept at reset value.

Bits 15:0 MACA1H: MAC address1 high [47:32]

This field contains the upper 16 bits (47:32) of the 6-byte second MAC address.

# Ethernet MAC address1 low register (ETH\_MACA1LR)

Address offset: 0x004C

Reset value: 0xFFFF FFFF

The MAC address 1 low register holds the lower 32 bits of the 6-byte second MAC address of the station.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
															111																

															MAC	A1L															
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw															

#### Bits 31:0 MACA1L: MAC address1 low [31:0]

This field contains the lower 32 bits of the 6-byte MAC address1. The content of this field is undefined until loaded by the application after the initialization process.

# Ethernet MAC address 2 high register (ETH MACA2HR)

Address offset: 0x0050

Reset value: 0x0000 FFFF

The MAC address 2 high register holds the upper 16 bits of the 6-byte second MAC address of the station.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
AE	SA			M	вс						Rese	VI. (0.0											MAC	A2H							
rw				nese	rvec				rw	rw	rw	rw	rw	rw	rw	rw	rw														

### Bit 31 AE: Address enable

When this bit is set, the address filters use the MAC address2 for perfect filtering. When reset, the address filters ignore the address for filtering.

#### Bit 30 SA: Source address

When this bit is set, the MAC address 2 [47:0] is used for comparison with the SA fields of the received frame.

When this bit is reset, the MAC address 2 [47:0] is used for comparison with the DA fields of the received frame.

#### Bits 29:24 MBC: Mask byte control

These bits are mask control bits for comparison of each of the MAC address2 bytes. When set high, the MAC core does not compare the corresponding byte of received DA/SA with the contents of the MAC address 2 registers. Each bit controls the masking of the bytes as follows:

- Bit 29: ETH\_MACA2HR [15:8]
- Bit 28: ETH\_MACA2HR [7:0]
- Bit 27: ETH\_MACA2LR [31:24]

٠..

- Bit 24: ETH\_MACA2LR [7:0]

Bits 23:16Reserved, must be kept at reset value.

Bits 15:0 MACA2H: MAC address2 high [47:32]

This field contains the upper 16 bits (47:32) of the 6-byte MAC address2.

# Ethernet MAC address 2 low register (ETH\_MACA2LR)

Address offset: 0x0054

Reset value: 0xFFFF FFFF

The MAC address 2 low register holds the lower 32 bits of the 6-byte second MAC address

of the station.

# 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

															MAC	A2L															
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw															

#### Bits 31:0 MACA2L: MAC address2 low [31:0]

This field contains the lower 32 bits of the 6-byte second MAC address2. The content of this field is undefined until loaded by the application after the initialization process.

# Ethernet MAC address 3 high register (ETH\_MACA3HR)

Address offset: 0x0058

Reset value: 0x0000 FFFF

The MAC address 3 high register holds the upper 16 bits of the 6-byte second MAC address

of the station.

#### 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

A	Æ	SA			М	вс			Reserved								MAC	АЗН							
r	w	rw	neserveu	rw	rw	rw	rw	rw	rw	rw	rw	rw													

# Bit 31 AE: Address enable

When this bit is set, the address filters use the MAC address3 for perfect filtering. When this bit is cleared, the address filters ignore the address for filtering.

#### Bit 30 SA: Source address

When this bit is set, the MAC address 3 [47:0] is used for comparison with the SA fields of the received frame.

When this bit is cleared, the MAC address 3[47:0] is used for comparison with the DA fields of the received frame.

# Bits 29:24 MBC: Mask byte control

These bits are mask control bits for comparison of each of the MAC address3 bytes. When these bits are set high, the MAC core does not compare the corresponding byte of received DA/SA with the contents of the MAC address 3 registers. Each bit controls the masking of the bytes as follows:

- Bit 29: ETH\_MACA3HR [15:8]
- Bit 28: ETH\_MACA3HR [7:0]
- Bit 27: ETH\_MACA3LR [31:24]

...

- Bit 24: ETH\_MACA3LR [7:0]

Bits 23:16 Reserved, must be kept at reset value.

Bits 15:0 MACA3H: MAC address3 high [47:32]

This field contains the upper 16 bits (47:32) of the 6-byte MAC address3.

# Ethernet MAC address 3 low register (ETH\_MACA3LR)

Address offset: 0x005C

Reset value: 0xFFFF FFFF

The MAC address 3 low register holds the lower 32 bits of the 6-byte second MAC address

of the station.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

															MAC	A3L															
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw															

Bits 31:0 MACA3L: MAC address3 low [31:0]

This field contains the lower 32 bits of the 6-byte second MAC address3. The content of this field is undefined until loaded by the application after the initialization process.

# 28.8.2 MMC register description

# Ethernet MMC control register (ETH\_MMCCR)

Address offset: 0x0100
Reset value: 0x0000 0000

The Ethernet MMC Control register establishes the operating mode of the management

counters.

 $31 \quad 30 \quad 29 \quad 28 \quad 27 \quad 26 \quad 25 \quad 24 \quad 23 \quad 22 \quad 21 \quad 20 \quad 19 \quad 18 \quad 17 \quad 16 \quad 15 \quad 14 \quad 13 \quad 12 \quad 11 \quad 10 \quad 9 \quad 8 \quad 7 \quad 6 \quad 5 \quad 4 \quad 3 \quad 2 \quad 1 \quad 0$ 

Reserved	MCFHP	MCP	MCF	ROR	CSR	CR
	rw	rw	rw	rw	rw	rw

Bits 31:6 Reserved, must be kept at reset value.

# Bit 5 MCFHP: MMC counter Full-Half preset

When MCFHP is low and bit4 is set, all MMC counters get preset to almost-half value. All frame-counters get preset to 0x7FFF\_FFF0 (half - 16)

When MCFHP is high and bit4 is set, all MMC counters get preset to almost-full value. All frame-counters get preset to 0xFFFF\_FFF0 (full - 16)

#### Bit 4 MCP: MMC counter preset

When set, all counters will be initialized or preset to almost full or almost half as per Bit5 above. This bit will be cleared automatically after 1 clock cycle. This bit along with bit5 is useful for debugging and testing the assertion of interrupts due to MMC counter becoming half-full or full.

#### Bit 3 MCF: MMC counter freeze

When set, this bit freezes all the MMC counters to their current value. (None of the MMC counters are updated due to any transmitted or received frame until this bit is cleared to 0. If any MMC counter is read with the Reset on Read bit set, then that counter is also cleared in this mode.)

#### Bit 2 ROR: Reset on read

When this bit is set, the MMC counters is reset to zero after read (self-clearing after reset). The counters are cleared when the least significant byte lane (bits [7:0]) is read.

#### Bit 1 CSR: Counter stop rollover

When this bit is set, the counter does not roll over to zero after it reaches the maximum value.

#### Bit 0 CR: Counter reset

When it is set, all counters are reset. This bit is cleared automatically after 1 clock cycle.

# Ethernet MMC receive interrupt register (ETH\_MMCRIR)

Address offset: 0x0104
Reset value: 0x0000 0000

The Ethernet MMC receive interrupt register maintains the interrupts generated when receive statistic counters reach half their maximum values. (MSB of the counter is set.) It is a 32-bit wide register. An interrupt bit is cleared when the respective MMC counter that

caused the interrupt is read. The least significant byte lane (bits [7:0]) of the respective counter must be read in order to clear the interrupt bit.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
						Rese	ervec	d						RGUFS				ı	Rese	ervec	i				RFAES	RFCES		Re	eserv	ed	
														rc_r											rc_r	rc_r					

Bits 31:18 Reserved, must be kept at reset value.

Bit 17 RGUFS: Received Good Unicast Frames Status

This bit is set when the received, good unicast frames, counter reaches half the maximum value.

Bits 16:7 Reserved, must be kept at reset value.

Bit 6 RFAES: Received frames alignment error status

This bit is set when the received frames, with alignment error, counter reaches half the maximum value.

Bit 5 RFCES: Received frames CRC error status

This bit is set when the received frames, with CRC error, counter reaches half the maximum value.

Bits 4:0 Reserved, must be kept at reset value.

# Ethernet MMC transmit interrupt register (ETH\_MMCTIR)

Address offset: 0x0108
Reset value: 0x0000 0000

The Ethernet MMC transmit Interrupt register maintains the interrupts generated when transmit statistic counters reach half their maximum values. (MSB of the counter is set.) It is a 32-bit wide register. An interrupt bit is cleared when the respective MMC counter that caused the interrupt is read. The least significant byte lane (bits [7:0]) of the respective counter must be read in order to clear the interrupt bit.

31 30 29	28 27 26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Reserved	d				TGFS		Re	eserv	/ed		TGFMSCS	TGFSCS							Rese	rved	I					
						rc_r						rc_r	rc_r														

Bits 31:22 Reserved, must be kept at reset value.

Bit 21 TGFS: Transmitted good frames status

This bit is set when the transmitted, good frames, counter reaches half the maximum value.

Bits 20:16 Reserved, must be kept at reset value.

Bit 15 TGFMSCS: Transmitted good frames more single collision status

This bit is set when the transmitted, good frames after more than a single collision, counter reaches half the maximum value.

Bit 14 TGFSCS: Transmitted good frames single collision status

This bit is set when the transmitted, good frames after a single collision, counter reaches half the maximum value.

Bits 13:0 Reserved, must be kept at reset value.

# Ethernet MMC receive interrupt mask register (ETH\_MMCRIMR)

Address offset: 0x010C Reset value: 0x0000 0000

The Ethernet MMC receive interrupt mask register maintains the masks for interrupts generated when the receive statistic counters reach half their maximum value. (MSB of the counter is set.) It is a 32-bit wide register.

3	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
						Rese	erved	d						RGUFM					Rese	erved	I				RFAEM	RFCEM		Re	serv	ed	
														rw											rw	rw					

- Bits 31:18 Reserved, must be kept at reset value.
  - Bit 17 RGUFM: Received good unicast frames mask

Setting this bit masks the interrupt when the received, good unicast frames, counter reaches half the maximum value.

- Bits 16:7 Reserved, must be kept at reset value.
  - Bit 6 RFAEM: Received frames alignment error mask

Setting this bit masks the interrupt when the received frames, with alignment error, counter reaches half the maximum value.

Bit 5 RFCEM: Received frame CRC error mask

Setting this bit masks the interrupt when the received frames, with CRC error, counter reaches half the maximum value.

Bits 4:0 Reserved, must be kept at reset value.

# Ethernet MMC transmit interrupt mask register (ETH\_MMCTIMR)

Address offset: 0x0110
Reset value: 0x0000 0000

The Ethernet MMC transmit interrupt mask register maintains the masks for interrupts generated when the transmit statistic counters reach half their maximum value. (MSB of the counter is set). It is a 32-bit wide register.

 $31 \quad 30 \quad 29 \quad 28 \quad 27 \quad 26 \quad 25 \quad 24 \quad 23 \quad 22 \quad 21 \quad 20 \quad 19 \quad 18 \quad 17 \quad 16 \quad 15 \quad 14 \quad 13 \quad 12 \quad 11 \quad 10 \quad 9 \quad 8 \quad 7 \quad 6 \quad 5 \quad 4 \quad 3 \quad 2 \quad 1 \quad 0$ 

Reserved	TGFM	Reserved	TGFMSCM	TGFSCM	Reserved
	rw		rw	rw	

Bits 31:22 Reserved, must be kept at reset value.

Bit 21 TGFM: Transmitted good frames mask

Setting this bit masks the interrupt when the transmitted, good frames, counter reaches half the maximum value.

Bits 20:16 Reserved, must be kept at reset value.

Bit 15 TGFMSCM: Transmitted good frames more single collision mask

Setting this bit masks the interrupt when the transmitted good frames after more than a single collision counter reaches half the maximum value.

Bit 14 TGFSCM: Transmitted good frames single collision mask

Setting this bit masks the interrupt when the transmitted good frames after a single collision counter reaches half the maximum value.

Bits 13:0 Reserved, must be kept at reset value.

# Ethernet MMC transmitted good frames after a single collision counter register (ETH\_MMCTGFSCCR)

Address offset: 0x014C Reset value: 0x0000 0000

This register contains the number of successfully transmitted frames after a single collision in Half-duplex mode.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

															TGF	SCC															
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:0 **TGFSCC:** Transmitted good frames single collision counter Transmitted good frames after a single collision counter.

# Ethernet MMC transmitted good frames after more than a single collision counter register (ETH\_MMCTGFMSCCR)

Address offset: 0x0150
Reset value: 0x0000 0000

This register contains the number of successfully transmitted frames after more than a single collision in Half-duplex mode.

Bits 31:0 **TGFMSCC:** Transmitted good frames more single collision counter Transmitted good frames after more than a single collision counter

# Ethernet MMC transmitted good frames counter register (ETH\_MMCTGFCR)

Address offset: 0x0168 Reset value: 0x0000 0000

This register contains the number of good frames transmitted.

Bits 31:0 TGFC: Transmitted good frames counter

# Ethernet MMC received frames with CRC error counter register (ETH\_MMCRFCECR)

Address offset: 0x0194 Reset value: 0x0000 0000

This register contains the number of frames received with CRC error.

Bits 31:0 RFCEC: Received frames CRC error counter Received frames with CRC error counter

# Ethernet MMC received frames with alignment error counter register (ETH\_MMCRFAECR)

Address offset: 0x0198
Reset value: 0x0000 0000

This register contains the number of frames received with alignment (dribble) error.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
															RFA	AEC															
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:0 **RFAEC:** Received frames alignment error counter Received frames with alignment error counter

# MMC received good unicast frames counter register (ETH\_MMCRGUFCR)

Address offset: 0x01C4
Reset value: 0x0000 0000

This register contains the number of good unicast frames received.

3	1 3	0 29	2	8 2	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
																RGI	JFC															
ı	r	r	r		r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:0 RGUFC: Received good unicast frames counter

# 28.8.3 IEEE 1588 time stamp registers

This section describes the registers required to support precision network clock synchronization functions under the IEEE 1588 standard.

# Ethernet PTP time stamp control register (ETH\_PTPTSCR)

Address offset: 0x0700

Reset value: 0x0000 00002000

This register controls the time stamp generation and update logic.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
					Re	eserv	/ed						TSPFFMAE	TNOOT		TSSMRME	TSSEME	TSSIPV4FE	TSSIPV6FE	TSSPTPOEFE	TSPTPPSV2E	TSSSR	TSSARFE		חפאפו	TTSARU	TSITE	UTSST	ILSST	LSFCU	TSE
													rw	r	N	rw	rw	rw	rw	rw	rw	rw	rw			rw	rw	rw	rw	rw	rw

- Bits 31:19 Reserved, must be kept at reset value.
  - Bit 18 **TSPFFMAE:** Time stamp PTP frame filtering MAC address enable

When set, this bit uses the MAC address (except for MAC address 0) to filter the PTP frames when PTP is sent directly over Ethernet.

Bits 17:16 TSCNT: Time stamp clock node type

The following are the available types of clock node:

00: Ordinary clock

01: Boundary clock

10: End-to-end transparent clock

11: Peer-to-peer transparent clock

Bit 15 TSSMRME: Time stamp snapshot for message relevant to master enable

When this bit is set, the snapshot is taken for messages relevant to the master node only. When this bit is cleared the snapshot is taken for messages relevant to the slave node only. This is valid only for the ordinary clock and boundary clock nodes.

Bit 14 TSSEME: Time stamp snapshot for event message enable

When this bit is set, the time stamp snapshot is taken for event messages only (SYNC, Delay\_Req, Pdelay\_Req or Pdelay\_Resp). When this bit is cleared the snapshot is taken for all other messages except for Announce, Management and Signaling.

Bit 13 TSSIPV4FE: Time stamp snapshot for IPv4 frames enable

When this bit is set, the time stamp snapshot is taken for IPv4 frames.

Bit 12 TSSIPV6FE: Time stamp snapshot for IPv6 frames enable

When this bit is set, the time stamp snapshot is taken for IPv6 frames.

Bit 11 TSSPTPOEFE: Time stamp snapshot for PTP over ethernet frames enable

When this bit is set, the time stamp snapshot is taken for frames which have PTP messages in Ethernet frames (PTP over Ethernet) also. By default snapshots are taken for UDP-IPEthernet PTP packets.

Bit 10 TSPTPPSV2E: Time stamp PTP packet snooping for version2 format enable

When this bit is set, the PTP packets are snooped using the version 2 format. When the bit is cleared, the PTP packets are snooped using the version 1 format.

Note: IEEE 1588 Version 1 and Version 2 formats as indicated in IEEE standard 1588-2008 (Revision of IEEE STD. 1588-2002).

Bit 9 TSSSR: Time stamp subsecond rollover: digital or binary rollover control

When this bit is set, the Time stamp low register rolls over when the subsecond counter reaches the value 0x3B9A C9FF (999 999 999 in decimal), and increments the Time Stamp (high) seconds.

When this bit is cleared, the rollover value of the subsecond register reaches 0x7FFF FFF. The subsecond increment has to be programmed correctly depending on the PTP's reference clock frequency and this bit value.

Bit 8 TSSARFE: Time stamp snapshot for all received frames enable

When this bit is set, the time stamp snapshot is enabled for all frames received by the core.

- Bits 7:6 Reserved, must be kept at reset value.
  - Bit 5 TSARU: Time stamp addend register update

When this bit is set, the Time stamp addend register's contents are updated to the PTP block for fine correction. This bit is cleared when the update is complete. This register bit must be read as zero before you can set it.



#### Bit 4 TSITE: Time stamp interrupt trigger enable

When this bit is set, a time stamp interrupt is generated when the system time becomes greater than the value written in the Target time register. When the Time stamp trigger interrupt is generated, this bit is cleared.

#### Bit 3 TSSTU: Time stamp system time update

When this bit is set, the system time is updated (added to or subtracted from) with the value specified in the Time stamp high update and Time stamp low update registers. Both the TSSTU and TSSTI bits must be read as zero before you can set this bit. Once the update is completed in hardware, this bit is cleared.

#### Bit 2 TSSTI: Time stamp system time initialize

When this bit is set, the system time is initialized (overwritten) with the value specified in the Time stamp high update and Time stamp low update registers. This bit must be read as zero before you can set it. When initialization is complete, this bit is cleared.

#### Bit 1 TSFCU: Time stamp fine or coarse update

When set, this bit indicates that the system time stamp is to be updated using the Fine Update method. When cleared, it indicates the system time stamp is to be updated using the Coarse method.

#### Bit 0 TSE: Time stamp enable

When this bit is set, time stamping is enabled for transmit and receive frames. When this bit is cleared, the time stamp function is suspended and time stamps are not added for transmit and receive frames. Because the maintained system time is suspended, you must always initialize the time stamp feature (system time) after setting this bit high.

The table below indicates the messages for which a snapshot is taken depending on the clock, enable master and enable snapshot for event message register settings.

Table 151. Time stamp snapshot dependency on registers bits

TSCNT (bits 17:16)	TSSMRME (bit 15) <sup>(1)</sup>	TSSEME (bit 14)	Messages for which snapshots are taken
00 or 01	X <sup>(2)</sup>	0	SYNC, Follow_Up, Delay_Req, Delay_Resp
00 or 01	1	1	Delay_Req
00 or 01	0	1	SYNC
10	N/A	0	SYNC, Follow_Up, Delay_Req, Delay_Resp
10	N/A	1	SYNC, Follow_Up
11	N/A	0	SYNC, Follow_Up, Delay_Req, Delay_Resp, Pdelay_Req, Pdelay_Resp
11	N/A	1	SYNC, Pdelay_Req, Pdelay_Resp

<sup>1.</sup> N/A = not applicable.

# Ethernet PTP subsecond increment register (ETH PTPSSIR)

Address offset: 0x0704
Reset value: 0x0000 0000

This register contains the 8-bit value by which the subsecond register is incremented. In Coarse update mode (TSFCU bit in ETH\_PTPTSCR), the value in this register is added to

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<sup>2.</sup> X = don't care.

the system time every clock cycle of HCLK. In Fine update mode, the value in this register is added to the system time whenever the accumulator gets an overflow.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved

Reserved

Reserved

Bits 31:8 Reserved, must be kept at reset value.

#### Bits 7:0 STSSI: System time subsecond increment

The value programmed in this register is added to the contents of the subsecond value of the system time in every update.

For example, to achieve 20 ns accuracy, the value is: 20 / 0.467 = ~43 (or 0x2A).

# **Ethernet PTP time stamp high register (ETH\_PTPTSHR)**

Address offset: 0x0708

Reset value: 0x0000 0000

This register contains the most significant (higher) 32 time bits. This read-only register contains the seconds system time value. The Time stamp high register, along with Time stamp low register, indicates the current value of the system time maintained by the MAC. Though it is updated on a continuous basis.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
															S	ГS															
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:0 STS: System time second

The value in this field indicates the current value in seconds of the System Time maintained by the core.

# Ethernet PTP time stamp low register (ETH\_PTPTSLR)

Address offset: 0x070C Reset value: 0x0000 0000

This register contains the least significant (lower) 32 time bits. This read-only register contains the subsecond system time value.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
STPNS															(	STSS	6														
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

#### Bit 31 STPNS: System time positive or negative sign

This bit indicates a positive or negative time value. When set, the bit indicates that time representation is negative. When cleared, it indicates that time representation is positive. Because the system time should always be positive, this bit is normally zero.

#### Bits 30:0 STSS: System time subseconds

The value in this field has the subsecond time representation, with 0.46 ns accuracy.

# Ethernet PTP time stamp high update register (ETH\_PTPTSHUR)

Address offset: 0x0710
Reset value: 0x0000 0000

This register contains the most significant (higher) 32 bits of the time to be written to, added to, or subtracted from the System Time value. The Time stamp high update register, along with the Time stamp update low register, initializes or updates the system time maintained by the MAC. You have to write both of these registers before setting the TSSTI or TSSTU bits in the Time stamp control register.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
															TS	SUS															

															TS	US															
rw																															

Bits 31:0 TSUS: Time stamp update second

The value in this field indicates the time, in seconds, to be initialized or added to the system time.

# Ethernet PTP time stamp low update register (ETH\_PTPTSLUR)

Address offset: 0x0714
Reset value: 0x0000 0000

This register contains the least significant (lower) 32 bits of the time to be written to, added

to, or subtracted from the System Time value.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

TSUPNS															Т	SUS	S														
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

#### Bit 31 **TSUPNS:** Time stamp update positive or negative sign

This bit indicates positive or negative time value. When set, the bit indicates that time representation is negative. When cleared, it indicates that time representation is positive. When TSSTI is set (system time initialization) this bit should be zero. If this bit is set when TSSTU is set, the value in the Time stamp update registers is subtracted from the system time. Otherwise it is added to the system time.

#### Bits 30:0 TSUSS: Time stamp update subseconds

The value in this field indicates the subsecond time to be initialized or added to the system time. This value has an accuracy of 0.46 ns (in other words, a value of 0x0000\_0001 is 0.46 ns).

# Ethernet PTP time stamp addend register (ETH\_PTPTSAR)

Address offset: 0x0718
Reset value: 0x0000 0000

This register is used by the software to readjust the clock frequency linearly to match the master clock frequency. This register value is used only when the system time is configured for Fine update mode (TSFCU bit in ETH\_PTPTSCR). This register content is added to a 32-bit accumulator in every clock cycle and the system time is updated whenever the accumulator overflows.

Bits 31:0 TSA: Time stamp addend

This register indicates the 32-bit time value to be added to the Accumulator register to achieve time synchronization.

# **Ethernet PTP target time high register (ETH\_PTPTTHR)**

Address offset: 0x071C Reset value: 0x0000 0000

This register contains the higher 32 bits of time to be compared with the system time for interrupt event generation. The Target time high register, along with Target time low register, is used to schedule an interrupt event (TSARU bit in ETH\_PTPTSCR) when the system time exceeds the value programmed in these registers.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
															TT	SH															
rw																															

Bits 31:0 TTSH: Target time stamp high

This register stores the time in seconds. When the time stamp value matches or exceeds both Target time stamp registers, the MAC, if enabled, generates an interrupt.

# Ethernet PTP target time low register (ETH\_PTPTTLR)

Address offset: 0x0720
Reset value: 0x0000 0000

This register contains the lower 32 bits of time to be compared with the system time for interrupt event generation.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
															TT	SL															
rw																															

Bits 31:0 TTSL: Target time stamp low

This register stores the time in (signed) nanoseconds. When the value of the time stamp matches or exceeds both Target time stamp registers, the MAC, if enabled, generates an interrupt.

# Ethernet PTP time stamp status register (ETH\_PTPTSSR)

Address offset: 0x0728 Reset value: 0x0000 0000

This register contains the time stamp status register.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
														Res	ervec	i														TSTTR	TSSO
																														ro	ro

Bits 31:2 Reserved, must be kept at reset value.

#### Bit 1 TSTTR: Time stamp target time reached

When set, this bit indicates that the value of the system time is greater than or equal to the value specified in the Target time high and low registers

#### Bit 0 TSSO: Time stamp second overflow

When set, this bit indicates that the second value of the time stamp has overflowed beyond 0xFFFF FFFF.

# Ethernet PTP PPS control register (ETH\_PTPPPSCR)

Address offset: 0x072C

Reset value: 0x0000 0000

This register controls the frequency of the PPS output.

# 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0 Reserved Reserved

#### Bits 31:4 Reserved, must be kept at reset value.

#### Bits 3:0 PPSFREQ: PPS frequency selection

The PPS output frequency is set to 2PPSFREQ Hz.

0000: 1 Hz with a pulse width of 125 ms for binary rollover and, of 100 ms for digital rollover

0001: 2 Hz with 50% duty cycle for binary rollover (digital rollover not recommended)

0010: 4 Hz with 50% duty cycle for binary rollover (digital rollover not recommended)

0011: 8 Hz with 50% duty cycle for binary rollover (digital rollover not recommended)

0100: 16 Hz with 50% duty cycle for binary rollover (digital rollover not recommended)

...

1111: 32768 Hz with 50% duty cycle for binary rollover (digital rollover not recommended)

Note: If digital rollover is used (TSSSR=1, bit 9 in ETH\_PTPTSCR), it is recommended not to use the PPS output with a frequency other than 1 Hz. Otherwise, with digital rollover, the PPS output has irregular waveforms at higher frequencies (though its average frequency will always be correct during any one-second window).

# 28.8.4 DMA register description

This section defines the bits for each DMA register. Non-32 bit accesses are allowed as long as the address is word-aligned.

# Ethernet DMA bus mode register (ETH\_DMABMR)

Address offset: 0x1000
Reset value: 0x0000 2101

The bus mode register establishes the bus operating modes for the DMA.

31	1	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		Re	eserv	/ed		MB	AAB	FPM	USP			RI	DР			FB	Р	М			PI	BL			)FE			DSL			PΑ	SR
						rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	Ш	rw	rw	rw	rw	rw	rw	rs



#### Bits 31:27 Reserved, must be kept at reset value.

#### Bit 26 MB: Mixed burst

When this bit is set high and the FB bit is low, the AHB master interface starts all bursts of a length greater than 16 with INCR (undefined burst). When this bit is cleared, it reverts to fixed burst transfers (INCRx and SINGLE) for burst lengths of 16 and below.

#### Bit 25 AAB: Address-aligned beats

When this bit is set high and the FB bit equals 1, the AHB interface generates all bursts aligned to the start address LS bits. If the FB bit equals 0, the first burst (accessing the data buffer's start address) is not aligned, but subsequent bursts are aligned to the address.

# Bit 24 FPM: 4xPBL mode

When set high, this bit multiplies the PBL value programmed (bits [22:17] and bits [13:8]) four times. Thus the DMA transfers data in a maximum of 4, 8, 16, 32, 64 and 128 beats depending on the PBL value.

#### Bit 23 USP: Use separate PBL

When set high, it configures the RxDMA to use the value configured in bits [22:17] as PBL while the PBL value in bits [13:8] is applicable to TxDMA operations only. When this bit is cleared, the PBL value in bits [13:8] is applicable for both DMA engines.

#### Bits 22:17 RDP: Rx DMA PBL

These bits indicate the maximum number of beats to be transferred in one RxDMA transaction. This is the maximum value that is used in a single block read/write operation. The RxDMA always attempts to burst as specified in RDP each time it starts a burst transfer on the host bus. RDP can be programmed with permissible values of 1, 2, 4, 8, 16, and 32. Any other value results in undefined behavior.

These bits are valid and applicable only when USP is set high.

#### Bit 16 FB: Fixed burst

This bit controls whether the AHB Master interface performs fixed burst transfers or not. When set, the AHB uses only SINGLE, INCR4, INCR8 or INCR16 during start of normal burst transfers. When reset, the AHB uses SINGLE and INCR burst transfer operations.

# Bits 15:14 PM: Rx Tx priority ratio

#### RxDMA requests are given priority over TxDMA requests in the following ratio:

00: 1:1

01: 2:1

10: 3:1

11: 4:1

This is valid only when the DA bit is cleared.

#### Bits 13:8 PBL: Programmable burst length

These bits indicate the maximum number of beats to be transferred in one DMA transaction. This is the maximum value that is used in a single block read/write operation. The DMA always attempts to burst as specified in PBL each time it starts a burst transfer on the host bus. PBL can be programmed with permissible values of 1, 2, 4, 8, 16, and 32. Any other value results in undefined behavior. When USP is set, this PBL value is applicable for TxDMA transactions only.

The PBL values have the following limitations:

- The maximum number of beats (PBL) possible is limited by the size of the Tx FIFO and Rx
- The FIFO has a constraint that the maximum beat supported is half the depth of the FIFO.
- If the PBL is common for both transmit and receive DMA, the minimum Rx FIFO and Tx FIFO depths must be considered.
- Do not program out-of-range PBL values, because the system may not behave properly.

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#### Bit 7 EDFE: Enhanced descriptor format enable

When this bit is set, the enhanced descriptor format is enabled and the descriptor size is increased to 32 bytes (8 DWORDS). This is required when time stamping is activated (TSE=1, ETH\_PTPTSCR bit 0) or if IPv4 checksum offload is activated (IPCO=1, ETH\_MACCR bit 10).

#### Bits 6:2 DSL: Descriptor skip length

This bit specifies the number of words to skip between two unchained descriptors. The address skipping starts from the end of current descriptor to the start of next descriptor. When DSL value equals zero, the descriptor table is taken as contiguous by the DMA, in Ring mode.

#### Bit 1 DA: DMA Arbitration

0: Round-robin with Rx:Tx priority given in bits [15:14]

1: Rx has priority over Tx

#### Bit 0 SR: Software reset

When this bit is set, the MAC DMA controller resets all MAC Subsystem internal registers and logic. It is cleared automatically after the reset operation has completed in all of the core clock domains. Read a 0 value in this bit before re-programming any register of the core.

# Ethernet DMA transmit poll demand register (ETH\_DMATPDR)

Address offset: 0x1004
Reset value: 0x0000 0000

This register is used by the application to instruct the DMA to poll the transmit descriptor list. The transmit poll demand register enables the Transmit DMA to check whether or not the current descriptor is owned by DMA. The Transmit Poll Demand command is given to wake up the TxDMA if it is in Suspend mode. The TxDMA can go into Suspend mode due to an underflow error in a transmitted frame or due to the unavailability of descriptors owned by transmit DMA. You can issue this command anytime and the TxDMA resets it once it starts re-fetching the current descriptor from host memory.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0

01	00	20	20	 20	20	24	20	 - 1	20	10	10	.,	10	10	17	10	12	 10	J	U	,	U	0	-	U	_	•	U
													TF	PD														
													rw_	_wt														

# Bits 31:0 TPD: Transmit poll demand

When these bits are written with any value, the DMA reads the current descriptor pointed to by the ETH\_DMACHTDR register. If that descriptor is not available (owned by Host), transmission returns to the Suspend state and ETH\_DMASR register bit 2 is asserted. If the descriptor is available, transmission resumes.

# **EHERNET DMA receive poll demand register (ETH\_DMARPDR)**

Address offset: 0x1008 Reset value: 0x0000 0000

This register is used by the application to instruct the DMA to poll the receive descriptor list. The Receive poll demand register enables the receive DMA to check for new descriptors. This command is given to wake up the RxDMA from Suspend state. The RxDMA can go into Suspend state only due to the unavailability of descriptors owned by it.



31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

RPD rw\_wt

Bits 31:0 RPD: Receive poll demand

When these bits are written with any value, the DMA reads the current descriptor pointed to by the ETH\_DMACHRDR register. If that descriptor is not available (owned by Host), reception returns to the Suspended state and ETH\_DMASR register bit 7 is not asserted. If the descriptor is available, the Receive DMA returns to active state.

# Ethernet DMA receive descriptor list address register (ETH\_DMARDLAR)

Address offset: 0x100C

Reset value: 0x0000 0000

The Receive descriptor list address register points to the start of the receive descriptor list. The descriptor lists reside in the STM32F4xx's physical memory space and must be word-aligned. The DMA internally converts it to bus-width aligned address by making the corresponding LS bits low. Writing to the ETH\_DMARDLAR register is permitted only when reception is stopped. When stopped, the ETH\_DMARDLAR register must be written to before the receive Start command is given.

 $31 \quad 30 \quad 29 \quad 28 \quad 27 \quad 26 \quad 25 \quad 24 \quad 23 \quad 22 \quad 21 \quad 20 \quad 19 \quad 18 \quad 17 \quad 16 \quad 15 \quad 14 \quad 13 \quad 12 \quad 11 \quad 10 \quad 9 \quad 8 \quad 7 \quad 6 \quad 5 \quad 4 \quad 3 \quad 2 \quad 1 \quad 0$ 

															SF	RL															
rw																															

Bits 31:0 SRL: Start of receive list

This field contains the base address of the first descriptor in the receive descriptor list. The LSB bits [1/2/3:0] for 32/64/128-bit bus width) are internally ignored and taken as all-zero by the DMA. Hence these LSB bits are read only.

# Ethernet DMA transmit descriptor list address register (ETH\_DMATDLAR)

Address offset: 0x1010
Reset value: 0x0000 0000

The Transmit descriptor list address register points to the start of the transmit descriptor list. The descriptor lists reside in the STM32F4xx's physical memory space and must be word-aligned. The DMA internally converts it to bus-width-aligned address by taking the corresponding LSB to low. Writing to the ETH\_DMATDLAR register is permitted only when transmission has stopped. Once transmission has stopped, the ETH\_DMATDLAR register can be written before the transmission Start command is given.

 $31 \quad 30 \quad 29 \quad 28 \quad 27 \quad 26 \quad 25 \quad 24 \quad 23 \quad 22 \quad 21 \quad 20 \quad 19 \quad 18 \quad 17 \quad 16 \quad 15 \quad 14 \quad 13 \quad 12 \quad 11 \quad 10 \quad 9 \quad 8 \quad 7 \quad 6 \quad 5 \quad 4 \quad 3 \quad 2 \quad 1 \quad 0$ 

Ī																S	TL															
	rw																															

#### Bits 31:0 STL: Start of transmit list

This field contains the base address of the first descriptor in the transmit descriptor list. The LSB bits [1/2/3:0] for 32/64/128-bit bus width) are internally ignored and taken as all-zero by the DMA. Hence these LSB bits are read-only.

# Ethernet DMA status register (ETH\_DMASR)

Address offset: 0x1014
Reset value: 0x0000 0000

The Status register contains all the status bits that the DMA reports to the application. The ETH\_DMASR register is usually read by the software driver during an interrupt service routine or polling. Most of the fields in this register cause the host to be interrupted. The ETH\_DMASR register bits are not cleared when read. Writing 1 to (unreserved) bits in ETH\_DMASR register[16:0] clears them and writing 0 has no effect. Each field (bits [16:0]) can be masked by masking the appropriate bit in the ETH\_DMAIER register.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
70,20	>	TSTS	PMTS	MMCS	erved		EBS			TPS			RPS		SIN	AIS	ERS	FBES	9	neved	ETS	RWTS	RPSS	RBUS	RS	SUL	ROS	STLT	TBUS	SSAL	TS
000	2	r	r	r	Res	r	r	r	r	r	r	r	r	r	rc- w1	rc- w1	rc- w1	rc- w1		S E E	rc- w1	٠.	-	rc- w1			rc- w1		rc- w1		

Bits 31:30 Reserved, must be kept at reset value.

#### Bit 29 TSTS: Time stamp trigger status

This bit indicates an interrupt event in the MAC core's Time stamp generator block. The software must read the MAC core's status register, clearing its source (bit 9), to reset this bit to 0. When this bit is high an interrupt is generated if enabled.

## Bit 28 PMTS: PMT status

This bit indicates an event in the MAC core's PMT. The software must read the corresponding registers in the MAC core to get the exact cause of interrupt and clear its source to reset this bit to 0. The interrupt is generated when this bit is high if enabled.

#### Bit 27 MMCS: MMC status

This bit reflects an event in the MMC of the MAC core. The software must read the corresponding registers in the MAC core to get the exact cause of interrupt and clear the source of interrupt to make this bit as 0. The interrupt is generated when this bit is high if enabled.

Bit 26 Reserved, must be kept at reset value.

# Bits 25:23 EBS: Error bits status

These bits indicate the type of error that caused a bus error (error response on the AHB interface). Valid only with the fatal bus error bit (ETH\_DMASR register [13]) set. This field does not generate an interrupt.

Bit 23

1 Error during data transfer by TxDMA

0 Error during data transfer by RxDMA

Bit 24

1 Error during read transfer

0 Error during write transfer

Bit 25 1 Error during descriptor access

0 Error during data buffer access

#### Bits 22:20 TPS: Transmit process state

These bits indicate the Transmit DMA FSM state. This field does not generate an interrupt.

000: Stopped; Reset or Stop Transmit Command issued

001: Running; Fetching transmit transfer descriptor

010: Running; Waiting for status

011: Running; Reading Data from host memory buffer and queuing it to transmit buffer (Tx FIFO)

100, 101: Reserved for future use

110: Suspended; Transmit descriptor unavailable or transmit buffer underflow

111: Running; Closing transmit descriptor

#### Bits 19:17 RPS: Receive process state

These bits indicate the Receive DMA FSM state. This field does not generate an interrupt.

000: Stopped: Reset or Stop Receive Command issued

001: Running: Fetching receive transfer descriptor

010: Reserved for future use

011: Running: Waiting for receive packet

100: Suspended: Receive descriptor unavailable

101: Running: Closing receive descriptor

110: Reserved for future use

111: Running: Transferring the receive packet data from receive buffer to host memory

#### Bit 16 NIS: Normal interrupt summary

The normal interrupt summary bit value is the logical OR of the following when the corresponding interrupt bits are enabled in the ETH\_DMAIER register:

- ETH\_DMASR [0]: Transmit interrupt
- ETH\_DMASR [2]: Transmit buffer unavailable
- ETH\_DMASR [6]: Receive interrupt
- ETH\_DMASR [14]: Early receive interrupt

Only unmasked bits affect the normal interrupt summary bit.

This is a sticky bit and it must be cleared (by writing a 1 to this bit) each time a corresponding bit that causes NIS to be set is cleared.

#### Bit 15 AIS: Abnormal interrupt summary

The abnormal interrupt summary bit value is the logical OR of the following when the corresponding interrupt bits are enabled in the ETH\_DMAIER register:

- ETH\_DMASR [1]:Transmit process stopped
- ETH\_DMASR [3]:Transmit jabber timeout
- ETH\_DMASR [4]: Receive FIFO overflow
- ETH\_DMASR [5]: Transmit underflow
- ETH\_DMASR [7]: Receive buffer unavailable
- ETH\_DMASR [8]: Receive process stopped
- ETH\_DMASR [9]: Receive watchdog timeout
- ETH\_DMASR [10]: Early transmit interrupt
- ETH\_DMASR [13]: Fatal bus error

Only unmasked bits affect the abnormal interrupt summary bit.

This is a sticky bit and it must be cleared each time a corresponding bit that causes AIS to be set is cleared.

#### Bit 14 ERS: Early receive status

This bit indicates that the DMA had filled the first data buffer of the packet. Receive Interrupt ETH\_DMASR [6] automatically clears this bit.



#### Bit 13 FBES: Fatal bus error status

This bit indicates that a bus error occurred, as detailed in [25:23]. When this bit is set, the corresponding DMA engine disables all its bus accesses.

#### Bits 12:11 Reserved, must be kept at reset value.

#### Bit 10 ETS: Early transmit status

This bit indicates that the frame to be transmitted was fully transferred to the Transmit FIFO.

#### Bit 9 RWTS: Receive watchdog timeout status

This bit is asserted when a frame with a length greater than 2 048 bytes is received.

#### Bit 8 RPSS: Receive process stopped status

This bit is asserted when the receive process enters the Stopped state.

#### Bit 7 RBUS: Receive buffer unavailable status

This bit indicates that the next descriptor in the receive list is owned by the host and cannot be acquired by the DMA. Receive process is suspended. To resume processing receive descriptors, the host should change the ownership of the descriptor and issue a Receive Poll Demand command. If no Receive Poll Demand is issued, receive process resumes when the next recognized incoming frame is received. ETH\_DMASR [7] is set only when the previous receive descriptor was owned by the DMA.

#### Bit 6 RS: Receive status

This bit indicates the completion of the frame reception. Specific frame status information has been posted in the descriptor. Reception remains in the Running state.

#### Bit 5 TUS: Transmit underflow status

This bit indicates that the transmit buffer had an underflow during frame transmission. Transmission is suspended and an underflow error TDES0[1] is set.

#### Bit 4 ROS: Receive overflow status

This bit indicates that the receive buffer had an overflow during frame reception. If the partial frame is transferred to the application, the overflow status is set in RDES0[11].

# Bit 3 TJTS: Transmit jabber timeout status

This bit indicates that the transmit jabber timer expired, meaning that the transmitter had been excessively active. The transmission process is aborted and placed in the Stopped state. This causes the transmit jabber timeout TDES0[14] flag to be asserted.

#### Bit 2 TBUS: Transmit buffer unavailable status

This bit indicates that the next descriptor in the transmit list is owned by the host and cannot be acquired by the DMA. Transmission is suspended. Bits [22:20] explain the transmit process state transitions. To resume processing transmit descriptors, the host should change the ownership of the bit of the descriptor and then issue a Transmit Poll Demand command.

# Bit 1 TPSS: Transmit process stopped status

This bit is set when the transmission is stopped.

# Bit 0 TS: Transmit status

This bit indicates that frame transmission is finished and TDES1[31] is set in the first descriptor.

# Ethernet DMA operation mode register (ETH\_DMAOMR)

Address offset: 0x1018
Reset value: 0x0000 0000



The operation mode register establishes the Transmit and Receive operating modes and commands. The ETH\_DMAOMR register should be the last CSR to be written as part of DMA initialization.

26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 7 6 8 5 4 3 DTCEFD FUGF DFRF Reserved Reserved RSF Reserved Ħ 2 监 RTC OSF SR TS. ST Reserved Reserved Reserved rs rw rw rw rw rw rw rw rw rw

Bits 31:27 Reserved, must be kept at reset value.

#### Bit 26 DTCEFD: Dropping of TCP/IP checksum error frames disable

When this bit is set, the core does not drop frames that only have errors detected by the receive checksum offload engine. Such frames do not have any errors (including FCS error) in the Ethernet frame received by the MAC but have errors in the encapsulated payload only. When this bit is cleared, all error frames are dropped if the FEF bit is reset.

#### Bit 25 RSF: Receive store and forward

When this bit is set, a frame is read from the Rx FIFO after the complete frame has been written to it, ignoring RTC bits. When this bit is cleared, the Rx FIFO operates in Cut-through mode, subject to the threshold specified by the RTC bits.

#### Bit 24 **DFRF:** Disable flushing of received frames

When this bit is set, the RxDMA does not flush any frames due to the unavailability of receive descriptors/buffers as it does normally when this bit is cleared. (See *Receive process suspended on page 885*)

Bits 23:22 Reserved, must be kept at reset value.

#### Bit 21 TSF: Transmit store and forward

When this bit is set, transmission starts when a full frame resides in the Transmit FIFO. When this bit is set, the TTC values specified by the ETH\_DMAOMR register bits [16:14] are ignored.

When this bit is cleared, the TTC values specified by the ETH\_DMAOMR register bits [16:14] are taken into account.

This bit should be changed only when transmission is stopped.

#### Bit 20 FTF: Flush transmit FIFO

When this bit is set, the transmit FIFO controller logic is reset to its default values and thus all data in the Tx FIFO are lost/flushed. This bit is cleared internally when the flushing operation is complete. The Operation mode register should not be written to until this bit is cleared.

Bits 19:17 Reserved, must be kept at reset value.

#### Bits 16:14 TTC: Transmit threshold control

These three bits control the threshold level of the Transmit FIFO. Transmission starts when the frame size within the Transmit FIFO is larger than the threshold. In addition, full frames with a length less than the threshold are also transmitted. These bits are used only when the TSF bit (Bit 21) is cleared.

000: 64 001: 128 010: 192 011: 256 100: 40 101: 32 110: 24 111: 16

#### Bit 13 ST: Start/stop transmission

When this bit is set, transmission is placed in the Running state, and the DMA checks the transmit list at the current position for a frame to be transmitted. Descriptor acquisition is attempted either from the current position in the list, which is the transmit list base address set by the ETH\_DMATDLAR register, or from the position retained when transmission was stopped previously. If the current descriptor is not owned by the DMA, transmission enters the Suspended state and the transmit buffer unavailable bit (ETH\_DMASR [2]) is set. The Start Transmission command is effective only when transmission is stopped. If the command is issued before setting the DMA ETH\_DMATDLAR register, the DMA behavior is unpredictable.

When this bit is cleared, the transmission process is placed in the Stopped state after completing the transmission of the current frame. The next descriptor position in the transmit list is saved, and becomes the current position when transmission is restarted. The Stop Transmission command is effective only when the transmission of the current frame is complete or when the transmission is in the Suspended state.

#### Bits 12:8 Reserved, must be kept at reset value.

#### Bit 7 FEF: Forward error frames

When this bit is set, all frames except runt error frames are forwarded to the DMA. When this bit is cleared, the Rx FIFO drops frames with error status (CRC error, collision error, giant frame, watchdog timeout, overflow). However, if the frame's start byte (write) pointer is already transferred to the read controller side (in Threshold mode), then the frames are not dropped. The Rx FIFO drops the error frames if that frame's start byte is not transferred (output) on the ARI bus.

# Bit 6 FUGF: Forward undersized good frames

When this bit is set, the Rx FIFO forwards undersized frames (frames with no error and length less than 64 bytes) including pad-bytes and CRC).

When this bit is cleared, the Rx FIFO drops all frames of less than 64 bytes, unless such a frame has already been transferred due to lower value of receive threshold (e.g., RTC = 01).

Bit 5 Reserved, must be kept at reset value.



#### Bits 4:3 RTC: Receive threshold control

These two bits control the threshold level of the Receive FIFO. Transfer (request) to DMA starts when the frame size within the Receive FIFO is larger than the threshold. In addition, full frames with a length less than the threshold are transferred automatically.

Note: Note that value of 11 is not applicable if the configured Receive FIFO size is 128 bytes.

Note: These bits are valid only when the RSF bit is zero, and are ignored when the RSF bit is set to 1.

00: 64 01: 32 10: 96 11: 128

#### Bit 2 OSF: Operate on second frame

When this bit is set, this bit instructs the DMA to process a second frame of Transmit data even before status for first frame is obtained.

#### Bit 1 SR: Start/stop receive

When this bit is set, the receive process is placed in the Running state. The DMA attempts to acquire the descriptor from the receive list and processes incoming frames. Descriptor acquisition is attempted from the current position in the list, which is the address set by the DMA ETH\_DMARDLAR register or the position retained when the receive process was previously stopped. If no descriptor is owned by the DMA, reception is suspended and the receive buffer unavailable bit (ETH\_DMASR [7]) is set. The Start Receive command is effective only when reception has stopped. If the command was issued before setting the DMA ETH\_DMARDLAR register, the DMA behavior is unpredictable.

When this bit is cleared, RxDMA operation is stopped after the transfer of the current frame. The next descriptor position in the receive list is saved and becomes the current position when the receive process is restarted. The Stop Receive command is effective only when the Receive process is in either the Running (waiting for receive packet) or the Suspended state

Bit 0 Reserved, must be kept at reset value.



## Ethernet DMA interrupt enable register (ETH\_DMAIER)

Address offset: 0x101C Reset value: 0x0000 0000

The Interrupt enable register enables the interrupts reported by ETH\_DMASR. Setting a bit to 1 enables a corresponding interrupt. After a hardware or software reset, all interrupts are disabled.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	NISE	AISE	ERIE	FBEIE	served	ETIE	RWTIE	RPSIE	RBUIE	RIE	TUIE	ROIE	TJTIE	TBUIE	TPSIE	TE
	rw	rw	rw	rw	Re	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:17 Reserved, must be kept at reset value.

#### Bit 16 NISE: Normal interrupt summary enable

When this bit is set, a normal interrupt is enabled. When this bit is cleared, a normal interrupt is disabled. This bit enables the following bits:

- ETH\_DMASR [0]: Transmit Interrupt
- ETH\_DMASR [2]: Transmit buffer unavailable
- ETH\_DMASR [6]: Receive interrupt
- ETH\_DMASR [14]: Early receive interrupt

#### Bit 15 AISE: Abnormal interrupt summary enable

When this bit is set, an abnormal interrupt is enabled. When this bit is cleared, an abnormal interrupt is disabled. This bit enables the following bits:

- ETH\_DMASR [1]: Transmit process stopped
- ETH\_DMASR [3]: Transmit jabber timeout
- ETH\_DMASR [4]: Receive overflow
- ETH\_DMASR [5]: Transmit underflow
- ETH\_DMASR [7]: Receive buffer unavailable
- ETH\_DMASR [8]: Receive process stopped
- ETH\_DMASR [9]: Receive watchdog timeout
- ETH\_DMASR [10]: Early transmit interrupt
- ETH\_DMASR [13]: Fatal bus error

#### Bit 14 **ERIE:** Early receive interrupt enable

When this bit is set with the normal interrupt summary enable bit (ETH\_DMAIER register[16]), the early receive interrupt is enabled.

When this bit is cleared, the early receive interrupt is disabled.

#### Bit 13 **FBEIE:** Fatal bus error interrupt enable

When this bit is set with the abnormal interrupt summary enable bit (ETH\_DMAIER register[15]), the fatal bus error interrupt is enabled.

When this bit is cleared, the fatal bus error enable interrupt is disabled.

## Bits 12:11 Reserved, must be kept at reset value.

#### Bit 10 ETIE: Early transmit interrupt enable

When this bit is set with the abnormal interrupt summary enable bit (ETH\_DMAIER register [15]), the early transmit interrupt is enabled.

When this bit is cleared, the early transmit interrupt is disabled.

#### Bit 9 **RWTIE:** receive watchdog timeout interrupt enable

When this bit is set with the abnormal interrupt summary enable bit (ETH\_DMAIER register[15]), the receive watchdog timeout interrupt is enabled.

When this bit is cleared, the receive watchdog timeout interrupt is disabled.

#### Bit 8 RPSIE: Receive process stopped interrupt enable

When this bit is set with the abnormal interrupt summary enable bit (ETH\_DMAIER register[15]), the receive stopped interrupt is enabled. When this bit is cleared, the receive stopped interrupt is disabled.

#### Bit 7 **RBUIE:** Receive buffer unavailable interrupt enable

When this bit is set with the abnormal interrupt summary enable bit (ETH\_DMAIER register[15]), the receive buffer unavailable interrupt is enabled.

When this bit is cleared, the receive buffer unavailable interrupt is disabled.

#### Bit 6 RIE: Receive interrupt enable

When this bit is set with the normal interrupt summary enable bit (ETH\_DMAIER register[16]), the receive interrupt is enabled.

When this bit is cleared, the receive interrupt is disabled.

### Bit 5 TUIE: Underflow interrupt enable

When this bit is set with the abnormal interrupt summary enable bit (ETH\_DMAIER register[15]), the transmit underflow interrupt is enabled.

When this bit is cleared, the underflow interrupt is disabled.

#### Bit 4 ROIE: Overflow interrupt enable

When this bit is set with the abnormal interrupt summary enable bit (ETH\_DMAIER register[15]), the receive overflow interrupt is enabled.

When this bit is cleared, the overflow interrupt is disabled.

#### Bit 3 **TJTIE:** Transmit jabber timeout interrupt enable

When this bit is set with the abnormal interrupt summary enable bit (ETH\_DMAIER register[15]), the transmit jabber timeout interrupt is enabled.

When this bit is cleared, the transmit jabber timeout interrupt is disabled.

#### Bit 2 TBUIE: Transmit buffer unavailable interrupt enable

When this bit is set with the normal interrupt summary enable bit (ETH\_DMAIER register[16]), the transmit buffer unavailable interrupt is enabled.

When this bit is cleared, the transmit buffer unavailable interrupt is disabled.

#### Bit 1 **TPSIE:** Transmit process stopped interrupt enable

When this bit is set with the abnormal interrupt summary enable bit (ETH\_DMAIER register[15]), the transmission stopped interrupt is enabled.

When this bit is cleared, the transmission stopped interrupt is disabled.

#### Bit 0 TIE: Transmit interrupt enable

When this bit is set with the normal interrupt summary enable bit (ETH\_DMAIER register[16]), the transmit interrupt is enabled.

When this bit is cleared, the transmit interrupt is disabled.

The Ethernet interrupt is generated only when the TSTS or PMTS bits of the DMA Status register is asserted with their corresponding interrupt are unmasked, or when the NIS/AIS Status bit is asserted and the corresponding Interrupt Enable bits (NISE/AISE) are enabled.

## Ethernet DMA missed frame and buffer overflow counter register (ETH DMAMFBOCR)

Address offset: 0x1020 Reset value: 0x0000 0000

The DMA maintains two counters to track the number of missed frames during reception. This register reports the current value of the counter. The counter is used for diagnostic purposes. Bits [15:0] indicate missed frames due to the STM32F4xx buffer being unavailable (no receive descriptor was available). Bits [27:17] indicate missed frames due to Rx FIFO overflow conditions and runt frames (good frames of less than 64 bytes).

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	OFOC						MFA						OMFC								MF	-C							
	rc_	rc_	rc_	rc_	rc_	rc_	rc_	rc_	rc_	rc_	rc_	rc_	rc_	rc_	rc_	rc_	rc_	rc_	rc_	rc_	rc_	rc_	rc_	rc_	rc_	rc_	rc_	rc_	rc_
	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:29 Reserved, must be kept at reset value.

Bit 28 OFOC: Overflow bit for FIFO overflow counter

Bits 27:17 MFA: Missed frames by the application

Indicates the number of frames missed by the application

Bit 16 OMFC: Overflow bit for missed frame counter

Bits 15:0 MFC: Missed frames by the controller

Indicates the number of frames missed by the Controller due to the host receive buffer being unavailable. This counter is incremented each time the DMA discards an incoming frame.

#### Ethernet DMA receive status watchdog timer register (ETH\_DMARSWTR)

Address offset: 0x1024 Reset value: 0x0000 0000

This register, when written with a non-zero value, enables the watchdog timer for the receive

status (RS, ETH\_DMASR[6]).

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 **BSWTC** 

Reserved rw rw rw rw rw

Bits 31:8 Reserved, must be kept at reset value.

Bits 7:0 RSWTC: Receive status (RS) watchdog timer count

Indicates the number of HCLK clock cycles multiplied by 256 for which the watchdog timer is set. The watchdog timer gets triggered with the programmed value after the RxDMA completes the transfer of a frame for which the RS status bit is not set due to the setting of RDES1[31] in the corresponding descriptor. When the watchdog timer runs out, the RS bit is set and the timer is stopped. The watchdog timer is reset when the RS bit is set high due to automatic setting of RS as per RDES1[31] of any received frame.

## Ethernet DMA current host transmit descriptor register (ETH\_DMACHTDR)

Address offset: 0x1048 Reset value: 0x0000 0000

The Current host transmit descriptor register points to the start address of the current

transmit descriptor read by the DMA.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
															НΤΙ	DAP															
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:0 HTDAP: Host transmit descriptor address pointer

Cleared on reset. Pointer updated by DMA during operation.

## Ethernet DMA current host receive descriptor register (ETH\_DMACHRDR)

Address offset: 0x104C Reset value: 0x0000 0000

The Current host receive descriptor register points to the start address of the current receive descriptor read by the DMA.

,	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
																HRI	DAP															
	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:0 HRDAP: Host receive descriptor address pointer

Cleared on Reset. Pointer updated by DMA during operation.

# Ethernet DMA current host transmit buffer address register (ETH\_DMACHTBAR)

Address offset: 0x1050
Reset value: 0x0000 0000

The Current host transmit buffer address register points to the current transmit buffer address being read by the DMA.

Bits 31:0 HTBAP: Host transmit buffer address pointer

Cleared on reset. Pointer updated by DMA during operation.

# Ethernet DMA current host receive buffer address register (ETH\_DMACHRBAR)

Address offset: 0x1054
Reset value: 0x0000 0000

The current host receive buffer address register points to the current receive buffer address

being read by the DMA.

Bits 31:0 HRBAP: Host receive buffer address pointer

Cleared on reset. Pointer updated by DMA during operation.

## 28.8.5 Ethernet register maps

Table 152 gives the ETH register map and reset values.

Table 152. Ethernet register map and reset values

Offset	Register	31	30	28	27	26	25	24	23	22	21 20	19	18	17	16	15	14	13	12	11	10	6	8	7	9	2	4	3	2	- 0	0
0x00	ETH_MACCR		Res	serve	d			eserved	QM '	ar '	Reserved		IFG		CSD	Reserved		ROD	LM		=	, RD	Reserved	, APCS	ā				, RE	Reserved	
	Reset value						0		0	0	В	0	0	0	0	Ж	0	0	0	0	0	_	Ж	0	0	0	0	0	0	ш.	
0x04	ETH_MACFF R	RA								F	Reserve	ed										SAF		PCE			PAM			⊋ ₹	
	Reset value	0																			0	0	0	0	0	0	0	0	0	0	0
0x08	ETH_MACHT HR													H	łΤΗ	[31:	0]														
	Reset value	0	0 0	0	0	0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x0C	ETH_MACHT LR													ŀ	HTL[	31:	0]														
	Reset value	0	0 0	0	0	0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	-	0
0x10	ETH_MACMII AR						F	Rese	erve	d								PA					MR				CR	;			M B
	Reset value															0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	ETH_MACMII																						N/	1D							
0x14	DR						F	Rese	erve	d																					
	Reset value															0	0	0	0	0	0	0	0	0	0	0	0	0	0		0
0x18	ETH_MACFC R							Р	Т									F	Rese	erve	d			ZQPD	Reserved	Pl	T	UPFD	RFCE	TFCE CONTRACT	700/071
	Reset value	0	0 0	0	0	0	0	0	0	0	0 0	0	0	0	0									0	ш	0	0	0	0	0 (	0
0x1C	ETH_MACVL ANTR	-					Re	ser	/ed						VLANTC								VLA	ANT	I						
	Reset value														0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x28	ETH_MACRW UFFR		F	rame	e filt	er re	g0\I	Fran	ne f	ilter	reg1\F	rame	filt	er re	g2\F	ran	ne fi	lter	reg3	\Fra	me	filte	r re	g4\	.\Fra	ame	filte	r reç	g7		
	Reset value														(	0															
0x2C	ETH_MACPM TCSR	WFFRPR									Rese	ved										gn '	Becerved	000	WFR	MPR	Reserved			MPE	
	Reset value	0																				0	Ц	-	0	0	4	•	0	0	0

Table 152. Ethernet register map and reset values (continued)

Offset	Register	- 0	തെ	7		4	က	N	- 6	െ	8	ì	ပ္		4 (	, (2)	1-1	0			Τ.	_	Τ			
Oliset	negistei	30	29 28	2	25	24	ĸ	22	21	19	18	17	16	12	7	13	7	9	6	8 1	٠ (	ט	4	3		0
0x34	ETH_MACDB GR Reset value		Reserved		O	TFNEGU	Reserved	o TFWA	TFRS	O MTP		S S S E N S	O MMTEA		Re	eserve	ed		O FF	Reserved		RFRCS	o RFWRA	Reserved	MSFRWCS	o MMRPEA
	neset value				U	U		U	0 0	, 0	U	U	0					-	-						U	0 0
0x38	ETH_MACSR														Re	eserve	ed		2 2	Reserved	OTOMAN		MMCS	PMTS	Res	erved
	Reset value																		0	Щ.	'	0 0	0	0		
0x3C	ETH_MACIM R														Re	eserve	ed		O ISTIM	F	Res	erved	I	o PMTIM	Res	erved
	ETH_MACA0	0				Д.		امما											_	IACA	01.1					
0x40	HR	MO					eser												IV	IACA	υп					
	Reset value	1 0	0 0	0	0 0	0	0	0	0 0	0	0	0	0	1	1	1 1	1	1	1	1 1		1 1	1	1	1	1 1
0.44	ETH_MACA0 LR												MAC	A0L	-											
0x44	Reset value	1 1 1	1   1	111	1   1	11	11	11	1 1 1	11	11	11	1	1	1 1	1   1	111	1 [	1 [	1   1	-	1   1	11	11	1	1   1
	ETH_MACA1	-		MDO!	0.01	<u> </u>		<u> </u>		1	<u> </u>	<u> </u>				- 1 -	1 ' 1				411	.   .			- 1	
0x48	HR	AE SA		MBC[	6:0]				Re	serve	ed								IV	IACA	IH					
	Reset value	0 0	0 0	0	0 0	0								1	1	1 1	1	1	1	1 1		1 1	1	1	1	1 1
0x4C	ETH_MACA1 LR												MAC	A1L	-											
0.00	Reset value	1 1 1	1   1	111	1   1	11	11	11	1 1 1	11	11	11	1	1	1 1	1   1	111	1 [	1 [	1   1	-	1   1	11	11	1	1   1
	ETH MACA2	<del> </del>	<u> </u>		<u>.   .</u>	<u> </u>	Ė	L'	L'L	<u> </u>	÷	<u> </u>	٠	• 1	•	<u> </u>	1				01.1		<u> </u>	<u> </u>	• 1	
0x50	HR	AE SA		MB	C				Re	serve	ed								IV	IACA	2H					
	Reset value	0 0	0 0	0	0 0	0								1	1	1 1	1	1	1	1 1		1 1	1	1	1	1 1
054	ETH_MACA2												MAC	A2L												
0x54	LR Reset value	1 1 1	1   1	111	1 1 1	1 1	1 1	1 1	1 1 1	11	1 1	1 1	111	1 1	1 1	1   1	111	1 1	1 1	1   1	-	1 1 1	1 1	1 1	1	1   1
	ETH_MACA3	+++	'   '	T.T	<u>.   .</u>	' '	Ľ	<u>'</u>	L'.L'	<u>'                                    </u>	<u> </u>	<u>'</u>	1	•	•	.   .	1'1	- 1	<u>'  </u>	'   '		.   .	Η'	<u> </u>	•	.   .
0x58	HR	AE SA		MB	C				Re	serve	ed								IV	IACA:	ЗН					
	Reset value	0 0	0 0	0	0 0	0								1	1	1 1	1	1	1	1 1		1 1	1	1	1	1 1
050	ETH_MACA3 LR												MAC	A3L												
0x5C	Reset value	1 1 1	1   1	11	1 1 1	1 1	1 1	1 1	1 1 1	11	1 1	1 1	11	1 1	1	1   1	T 1 T	1 1	1 1	1   1	_	1 1 1	1 1	1 1	1 1	1   1
	11000t value		<u> </u>	+	• • •	<u> </u>	<u> </u>	<u> </u>		<u> </u>	<u> </u>	<u> </u>	<u> </u>		<u> </u>	<u> </u>	1.1		•	<u></u>		. д	+:	<u>t</u>	-	÷
0x100	ETH_MMCCR									Res	orvo	-d										MCFHP	MCP	MCF	ROR	CSR
0.00										nes	erve	au										M	Δ			
	Reset value											,												0	0	0 0
	ETH_MMCRI											ORGUFS									DEVEC	RFCES				
0x104	R				Res	erve	d					g				Res	erved	l			Į,			Re	serve	ed
	Reset value											10										0 0				
0x108	ETH_MMCTI R		ı	Reser	ved				TGFS	Re	eser	ved		OTGFMSCS	TGFSCS					Re	ser	ved				
	Reset value								0				ď	0	0											
	ETH MMCRI											Σ									2	≥ ≥				
0x10C	MR				Res	erve	h					ORGUFM				Res	erved	ı			DEAFIN	RECEM		Re	serve	ed
OX.100							•					N N														
	Reset value											0		= 1								0 0				
0x110	ETH_MMCTI MR		ı	Reser	ved				TGFM	Re	eser	ved		TGFMSCM	TGFSCM					Re	ser	ved				
	Reset value								0				-		0											
	ETH_MMCTG												TGF	SCC	:											
0x14C	FSCCR	0101	010	101	010	1 ^	1.0		1010	\ T \	1 0					0 1 0	101	<u> </u>	<u>Λ Ι</u>	<u>ΛΙ</u>		<u> </u>	Ι.	1.0	0	0 1 0
-	Reset value ETH_MMCTG	0 0	0 0	0	0 0	0	0	0	0 (	0	0			_		0 0	0	0	0	0 0	'   '	0 0	0	0	0	0 0
0x150	FMSCCR											Т	GFM	ISC	С											
	Reset value	0 0	0 0	0	0 0	0	0	0	0 (	0	0	0	0	0	0	0 0	0	0	0	0 0	) (	0 0	0	0	0	0 0
_	_		_							_	_												_			

Table 152. Ethernet register map and reset values (continued)

Iable				-			P				vaiu	_	· ·	,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,		~,												
Offset	- 5	31	29	28	27	26	27	23	22	21	20 19	18	17	16	14	13	12	7	10	6	8	7	9	2	4	ဗ	7	- 0
0100	ETH_MMCTG FCR													TGF	С													
0x168	Reset value	0 0	0   0	10	0	0	0 0	10	0	0	0 0	0	0	0 (	0 0	0	0	0	0	0	0	0	0	0	0	0	0   0	0   0
	ETH_MMCRF	<u> </u>			ľ	<u> </u>		ľ	٦٠	<u> </u>	0 0			RFCE	- 1	ľ		<u> </u>		•	•		•		١			
0x194	CECR											_																
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0x708	ETH_PTPTS HR												S	TS[3	1:0]													
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0x1014	ETH_DMASR	Reserved	TSTS	<b>PMTS</b>	o MMCS	Reserved	EBS			TPS		RPS		SIN N	ERS	FBES	Reserved		ETS	RWTS	RPSS	RBUS	RS	TUS	ROS	TJTS	TPSS	2 2
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Table 152. Ethernet register map and reset values (continued)

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0x101C	ETH_DMAIE R								serv	/ed							0		o ERIE	O FBEIE	Reserved		o ETIE	O NW III	ORPSIE	BUIE	O RIE	o TUE	□ BOIE		O TBUIE	OTPSIE	O TIE
0x1020	ETH_DMAMF BOCR	Re	serv	/ed	OFOC						MFA						OMFC								М								
0x1024	Reset value ETH_DMARS WTR Reset value ETH_DMACH				0	0	0	0	0	0	0	O F	0 Rese	0 erved	d 0	0	0	0	0	0	0	0	0	0	0	0	0	0 I	0 RSV 0	0 VTC 0	0	0	0
0x1048	TDR Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	HTI 0	OAP 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x104C	ETH_DMACH RDR Reset value	0	0	0	0	0	10	0	0	0	0	0	0	0	۸	٥	HRI 0	DAP 0		0	0 1	0 [	0 1	0 [	0	0	0	0	0	0	0	0	0
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	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

Refer to *Table 2 on page 50* for the register boundary addresses.

## 29 USB on-the-go full-speed (OTG\_FS)

This section applies to the whole STM32F4xx family, unless otherwise specified.

## 29.1 OTG FS introduction

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This section presents the architecture and the programming model of the OTG\_FS controller.

The following acronyms are used throughout the section:

FS full-speed LS Low-speed

MAC Media access controller

OTG On-the-go

PFC Packet FIFO controller

PHY Physical layer

USB Universal serial bus

UTMI USB 2.0 transceiver macrocell interface (UTMI)

References are made to the following documents:

- USB On-The-Go Supplement, Revision 1.3
- Universal Serial Bus Revision 2.0 Specification

The OTG\_FS is a dual-role device (DRD) controller that supports both device and host functions and is fully compliant with the *On-The-Go Supplement to the USB 2.0 Specification*. It can also be configured as a host-only or device-only controller, fully compliant with the *USB 2.0 Specification*. In host mode, the OTG\_FS supports full-speed (FS, 12 Mbits/s) and low-speed (LS, 1.5 Mbits/s) transfers whereas in device mode, it only supports full-speed (FS, 12 Mbits/s) transfers. The OTG\_FS supports both HNP and SRP. The only external device required is a charge pump for  $V_{BUS}$  in host mode.

## 29.2 OTG\_FS main features

The main features can be divided into three categories: general, host-mode and device-mode features.

#### 29.2.1 General features

The OTG\_FS interface general features are the following:

- It is USB-IF certified to the Universal Serial Bus Specification Rev 2.0
- It includes full support (PHY) for the optional On-The-Go (OTG) protocol detailed in the On-The-Go Supplement Rev 1.3 specification
  - Integrated support for A-B Device Identification (ID line)
  - Integrated support for host Negotiation Protocol (HNP) and Session Request Protocol (SRP)
  - It allows host to turn V<sub>BUS</sub> off to conserve battery power in OTG applications
  - It supports OTG monitoring of V<sub>BUS</sub> levels with internal comparators
  - It supports dynamic host-peripheral switch of role
- It is software-configurable to operate as:
  - SRP capable USB FS Peripheral (B-device)
  - SRP capable USB FS/LS host (A-device)
  - USB On-The-Go Full-Speed Dual Role device
- It supports FS SOF and LS Keep-alives with
  - SOF pulse PAD connectivity
  - SOF pulse internal connection to timer2 (TIM2)
  - Configurable framing period
  - Configurable end of frame interrupt
- It includes power saving features such as system stop during USB Suspend, switch-off of clock domains internal to the digital core, PHY and DFIFO power management
- It features a dedicated RAM of 1.25 Kbytes with advanced FIFO control:
  - Configurable partitioning of RAM space into different FIFOs for flexible and efficient use of RAM
  - Each FIFO can hold multiple packets
  - Dynamic memory allocation
  - Configurable FIFO sizes that are not powers of 2 to allow the use of contiguous memory locations
- It guarantees max USB bandwidth for up to one frame (1ms) without system intervention

#### 29.2.2 Host-mode features

The OTG\_FS interface main features and requirements in host-mode are the following:

- External charge pump for V<sub>BUS</sub> voltage generation.
- Up to 8 host channels (pipes): each channel is dynamically reconfigurable to allocate any type of USB transfer.
- Built-in hardware scheduler holding:
  - Up to 8 interrupt plus isochronous transfer requests in the periodic hardware queue
  - Up to 8 control plus bulk transfer requests in the non-periodic hardware queue
- Management of a shared RX FIFO, a periodic TX FIFO and a nonperiodic TX FIFO for efficient usage of the USB data RAM.

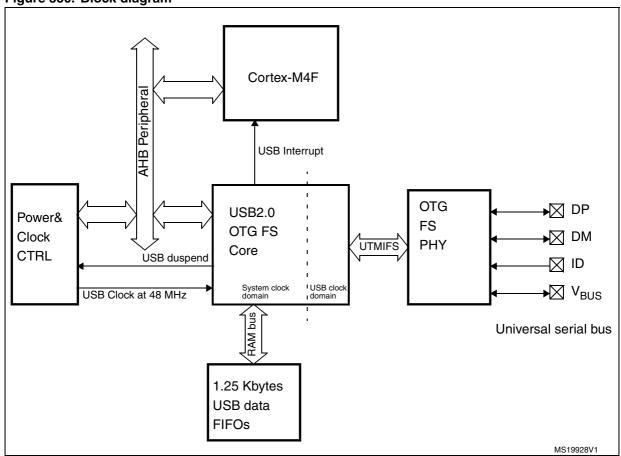
## 29.2.3 Peripheral-mode features

The OTG\_FS interface main features in peripheral-mode are the following:

- 1 bidirectional control endpoint0
- 3 IN endpoints (EPs) configurable to support Bulk, Interrupt or Isochronous transfers
- 3 OUT endpoints configurable to support Bulk, Interrupt or Isochronous transfers
- Management of a shared Rx FIFO and a Tx-OUT FIFO for efficient usage of the USB data RAM
- Management of up to 4 dedicated Tx-IN FIFOs (one for each active IN EP) to put less load on the application
- Support for the soft disconnect feature.

## 29.3 OTG\_FS functional description

Figure 350. Block diagram



### 29.3.1 OTG full-speed core

The USB OTG FS receives the 48 MHz  $\pm 0.25\%$  clock from the reset and clock controller (RCC), via an external quartz. The USB clock is used for driving the 48 MHz domain at full-speed (12 Mbit/s) and must be enabled prior to configuring the OTG FS core.

The CPU reads and writes from/to the OTG FS core registers through the AHB peripheral bus. It is informed of USB events through the single USB OTG interrupt line described in *Section 29.15: OTG\_FS interrupts*.

The CPU submits data over the USB by writing 32-bit words to dedicated OTG\_FS locations (push registers). The data are then automatically stored into Tx-data FIFOs configured within the USB data RAM. There is one Tx-FIFO push register for each in-endpoint (peripheral mode) or out-channel (host mode).

The CPU receives the data from the USB by reading 32-bit words from dedicated OTG\_FS addresses (pop registers). The data are then automatically retrieved from a shared Rx-FIFO configured within the 1.25 KB USB data RAM. There is one Rx-FIFO pop register for each out-endpoint or in-channel.

The USB protocol layer is driven by the serial interface engine (SIE) and serialized over the USB by the full-/low-speed transceiver module within the on-chip physical layer (PHY).

## 29.3.2 Full-speed OTG PHY

The embedded full-speed OTG PHY is controlled by the OTG FS core and conveys USB control & data signals through the full-speed subset of the UTMI+ Bus (UTMIFS). It provides the physical support to USB connectivity.

The full-speed OTG PHY includes the following components:

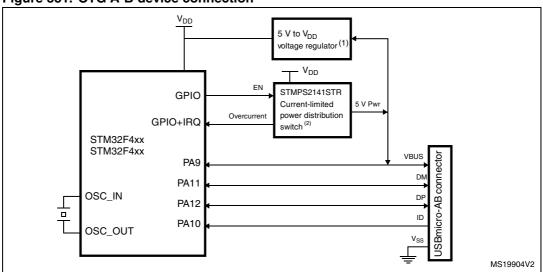
- FS/LS transceiver module used by both host and device. It directly drives transmission and reception on the single-ended USB lines.
- integrated ID pull-up resistor used to sample the ID line for A/B device identification.
- DP/DM integrated pull-up and pull-down resistors controlled by the OTG\_FS core depending on the current role of the device. As a peripheral, it enables the DP pull-up resistor to signal full-speed peripheral connections as soon as V<sub>BUS</sub> is sensed to be at a valid level (B-session valid). In host mode, pull-down resistors are enabled on both DP/DM. Pull-up and pull-down resistors are dynamically switched when the device's role is changed via the host negotiation protocol (HNP).
- Pull-up/pull-down resistor ECN circuit. The DP pull-up consists of 2 resistors controlled separately from the OTG\_FS as per the resistor Engineering Change Notice applied to USB Rev2.0. The dynamic trimming of the DP pull-up strength allows for better noise rejection and Tx/Rx signal quality.
- V<sub>BUS</sub> sensing comparators with hysteresis used to detect V<sub>BUS</sub> Valid, A-B Session Valid and session-end voltage thresholds. They are used to drive the session request protocol (SRP), detect valid startup and end-of-session conditions, and constantly monitor the V<sub>BUS</sub> supply during USB operations.
- V<sub>BUS</sub> pulsing method circuit used to charge/discharge V<sub>BUS</sub> through resistors during the SRP (weak drive).

Caution:

To guarantee a correct operation for the USB OTG FS peripheral, the AHB frequency should be higher than 14.2 MHz.

## 29.4 OTG dual role device (DRD)

Figure 351. OTG A-B device connection



- 1. External voltage regulator only needed when building a V<sub>BUS</sub> powered device
- STMPS2141STR needed only if the application has to support a V<sub>BUS</sub> powered device. A basic power switch can be used if 5 V are available on the application board.
- 3. V<sub>DD</sub> range is between 2 V and 3.6 V.

#### 29.4.1 ID line detection

The host or peripheral (the default) role is assumed depending on the ID input pin. The ID line status is determined on plugging in the USB, depending on which side of the USB cable is connected to the micro-AB receptacle.

- If the B-side of the USB cable is connected with a floating ID wire, the integrated pull-up resistor detects a high ID level and the default Peripheral role is confirmed. In this configuration the OTG\_FS complies with the standard FSM described by section 6.8.2: On-The-Go B-device of the On-The-Go Specification Rev1.3 supplement to the USB2.0.
- If the A-side of the USB cable is connected with a grounded ID, the OTG\_FS issues an ID line status change interrupt (CIDSCHG bit in OTG\_FS\_GINTSTS) for host software initialization, and automatically switches to the host role. In this configuration the OTG\_FS complies with the standard FSM described by section 6.8.1: On-The-Go Adevice of the On-The-Go Specification Rev1.3 supplement to the USB2.0.

#### 29.4.2 HNP dual role device

The HNP capable bit in the Global USB configuration register (HNPCAP bit in OTG\_FS\_GUSBCFG) enables the OTG\_FS core to dynamically change its role from A-host to A-peripheral and vice-versa, or from B-Peripheral to B-host and vice-versa according to the host negotiation protocol (HNP). The current device status can be read by the combined values of the Connector ID Status bit in the Global OTG control and status register (CIDSTS bit in OTG\_FS\_GOTGCTL) and the current mode of operation bit in the global interrupt and status register (CMOD bit in OTG\_FS\_GINTSTS).

The HNP program model is described in detail in *Section 29.17: OTG\_FS programming model*.

#### 29.4.3 SRP dual role device

The SRP capable bit in the global USB configuration register (SRPCAP bit in OTG\_FS\_GUSBCFG) enables the OTG\_FS core to switch off the generation of  $V_{BUS}$  for the A-device to save power. Note that the A-device is always in charge of driving  $V_{BUS}$  regardless of the host or peripheral role of the OTG\_FS.

the SRP A/B-device program model is described in detail in *Section 29.17: OTG\_FS programming model*.

## 29.5 USB peripheral

This section gives the functional description of the OTG\_FS in the USB peripheral mode. The OTG\_FS works as an USB peripheral in the following circumstances:

- OTG B-Peripheral
  - OTG B-device default state if B-side of USB cable is plugged in
- OTG A-Peripheral
  - OTG A-device state after the HNP switches the OTG\_FS to its peripheral role
- B-device
  - If the ID line is present, functional and connected to the B-side of the USB cable, and the HNP-capable bit in the Global USB Configuration register (HNPCAP bit in OTG\_FS\_GUSBCFG) is cleared (see On-The-Go Rev1.3 par. 6.8.3).
- Peripheral only (see Figure 352: USB peripheral-only connection)
  - The force device mode bit in the Global USB configuration register (FDMOD in OTG\_FS\_GUSBCFG) is set to 1, forcing the OTG\_FS core to work as a USB peripheral-only (see On-The-Go Rev1.3 par. 6.8.3). In this case, the ID line is ignored even if present on the USB connector.

Note:

To build a bus-powered device implementation in case of the B-device or peripheral-only configuration, an external regulator has to be added that generates the  $V_{DD}$  chip-supply from  $V_{BUS}$ .

The  $V_{BUS}$  pin can be freed by disabling the  $V_{BUS}$  sensing option. This is done by setting the NOVBUSSENS bit in the OTG\_FS\_GCCFG register. In this case the  $V_{BUS}$  is considered internally to be always at  $V_{BUS}$  valid level (5 V).

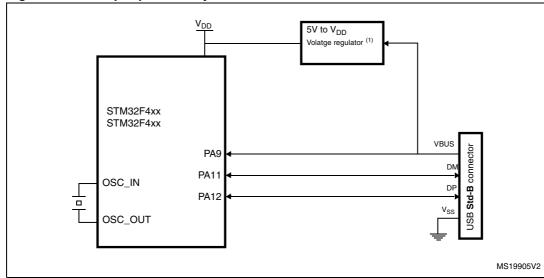


Figure 352. USB peripheral-only connection

- 1. Use a regulator to build a bus-powered device.
- 2.  $V_{DD}$  range between 2 V and 3.6 V.

## 29.5.1 SRP-capable peripheral

The SRP capable bit in the Global USB configuration register (SRPCAP bit in OTG\_FS\_GUSBCFG) enables the OTG\_FS to support the session request protocol (SRP). In this way, it allows the remote A-device to save power by switching off V<sub>BUS</sub> while the USB session is suspended.

The SRP peripheral mode program model is described in detail in the *B-device session request protocol* section.

### 29.5.2 Peripheral states

#### Powered state

The  $V_{BUS}$  input detects the B-Session valid voltage by which the USB peripheral is allowed to enter the powered state (see USB2.0 par9.1). The OTG\_FS then automatically connects the DP pull-up resistor to signal full-speed device connection to the host and generates the session request interrupt (SRQINT bit in OTG\_FS\_GINTSTS) to notify the powered state.

The  $V_{BUS}$  input also ensures that valid  $V_{BUS}$  levels are supplied by the host during USB operations. If a drop in  $V_{BUS}$  below B-session valid happens to be detected (for instance because of a power disturbance or if the host port has been switched off), the OTG\_FS automatically disconnects and the session end detected (SEDET bit in OTG\_FS\_GOTGINT) interrupt is generated to notify that the OTG\_FS has exited the powered state.

In the powered state, the OTG\_FS expects to receive some reset signaling from the host. No other USB operation is possible. When a reset signaling is received the reset detected interrupt (USBRST in OTG\_FS\_GINTSTS) is generated. When the reset signaling is complete, the enumeration done interrupt (ENUMDNE bit in OTG\_FS\_GINTSTS) is generated and the OTG\_FS enters the Default state.

#### Soft disconnect

The powered state can be exited by software with the soft disconnect feature. The DP pullup resistor is removed by setting the soft disconnect bit in the device control register (SDIS bit in OTG\_FS\_DCTL), causing a device disconnect detection interrupt on the host side even though the USB cable was not really removed from the host port.

#### **Default state**

In the Default state the OTG\_FS expects to receive a SET\_ADDRESS command from the host. No other USB operation is possible. When a valid SET\_ADDRESS command is decoded on the USB, the application writes the corresponding number into the device address field in the device configuration register (DAD bit in OTG\_FS\_DCFG). The OTG\_FS then enters the address state and is ready to answer host transactions at the configured USB address.

#### Suspended state

The OTG\_FS peripheral constantly monitors the USB activity. After counting 3 ms of USB idleness, the early suspend interrupt (ESUSP bit in OTG\_FS\_GINTSTS) is issued, and confirmed 3 ms later, if appropriate, by the suspend interrupt (USBSUSP bit in OTG\_FS\_GINTSTS). The device suspend bit is then automatically set in the device status register (SUSPSTS bit in OTG\_FS\_DSTS) and the OTG\_FS enters the suspended state.

The suspended state may optionally be exited by the device itself. In this case the application sets the remote wakeup signaling bit in the device control register (RWUSIG bit in OTG\_FS\_DCTL) and clears it after 1 to 15 ms.

When a resume signaling is detected from the host, the resume interrupt (WKUPINT bit in OTG FS GINTSTS) is generated and the device suspend bit is automatically cleared.

## 29.5.3 Peripheral endpoints

The OTG\_FS core instantiates the following USB endpoints:

- Control endpoint 0:
  - Bidirectional and handles control messages only
  - Separate set of registers to handle in and out transactions
  - Proper control (OTG\_FS\_DIEPCTL0/OTG\_FS\_DOEPCTL0), transfer configuration (OTG\_FS\_DIEPTSIZ0/OTG\_FS\_DIEPTSIZ0), and status-interrupt (OTG\_FS\_DIEPINTx/)OTG\_FS\_DOEPINT0) registers. The available set of bits inside the control and transfer size registers slightly differs from that of other endpoints
- 3 IN endpoints
  - Each of them can be configured to support the isochronous, bulk or interrupt transfer type
  - Each of them has proper control (OTG\_FS\_DIEPCTLx), transfer configuration (OTG\_FS\_DIEPTSIZx), and status-interrupt (OTG\_FS\_DIEPINTx) registers
  - The Device IN endpoints common interrupt mask register (OTG\_FS\_DIEPMSK) is available to enable/disable a single kind of endpoint interrupt source on all of the IN endpoints (EP0 included)
  - Support for incomplete isochronous IN transfer interrupt (IISOIXFR bit in OTG\_FS\_GINTSTS), asserted when there is at least one isochronous IN endpoint

on which the transfer is not completed in the current frame. This interrupt is asserted along with the end of periodic frame interrupt (OTG\_FS\_GINTSTS/EOPF).

- 3 OUT endpoints
  - Each of them can be configured to support the isochronous, bulk or interrupt transfer type
  - Each of them has a proper control (OTG\_FS\_DOEPCTLx), transfer configuration (OTG\_FS\_DOEPTSIZx) and status-interrupt (OTG\_FS\_DOEPINTx) register
  - Device Out endpoints common interrupt mask register (OTG\_FS\_DOEPMSK) is available to enable/disable a single kind of endpoint interrupt source on all of the OUT endpoints (EP0 included)
  - Support for incomplete isochronous OUT transfer interrupt (INCOMPISOOUT bit in OTG\_FS\_GINTSTS), asserted when there is at least one isochronous OUT endpoint on which the transfer is not completed in the current frame. This interrupt is asserted along with the end of periodic frame interrupt (OTG\_FS\_GINTSTS/EOPF).

#### **Endpoint control**

- The following endpoint controls are available to the application through the device endpoint-x IN/OUT control register (DIEPCTLx/DOEPCTLx):
  - Endpoint enable/disable
  - Endpoint activate in current configuration
  - Program USB transfer type (isochronous, bulk, interrupt)
  - Program supported packet size
  - Program Tx-FIFO number associated with the IN endpoint
  - Program the expected or transmitted data0/data1 PID (bulk/interrupt only)
  - Program the even/odd frame during which the transaction is received or transmitted (isochronous only)
  - Optionally program the NAK bit to always negative-acknowledge the host regardless of the FIFO status
  - Optionally program the STALL bit to always stall host tokens to that endpoint
  - Optionally program the SNOOP mode for OUT endpoint not to check the CRC field of received data

#### **Endpoint transfer**

The device endpoint-*x* transfer size registers (DIEPTSIZ*x*/DOEPTSIZ*x*) allow the application to program the transfer size parameters and read the transfer status. Programming must be done before setting the endpoint enable bit in the endpoint control register. Once the endpoint is enabled, these fields are read-only as the OTG FS core updates them with the current transfer status.

The following transfer parameters can be programmed:

- Transfer size in bytes
- Number of packets that constitute the overall transfer size

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### **Endpoint status/interrupt**

The device endpoint-*x* interrupt registers (DIEPINT*x*/DOPEPINT*x*) indicate the status of an endpoint with respect to USB- and AHB-related events. The application must read these registers when the OUT endpoint interrupt bit or the IN endpoint interrupt bit in the core interrupt register (OEPINT bit in OTG\_FS\_GINTSTS or IEPINT bit in OTG\_FS\_GINTSTS, respectively) is set. Before the application can read these registers, it must first read the device all endpoints interrupt (OTG\_FS\_DAINT) register to get the exact endpoint number for the device endpoint-*x* interrupt register. The application must clear the appropriate bit in this register to clear the corresponding bits in the DAINT and GINTSTS registers

The peripheral core provides the following status checks and interrupt generation:

- Transfer completed interrupt, indicating that data transfer was completed on both the application (AHB) and USB sides
- Setup stage has been done (control-out only)
- Associated transmit FIFO is half or completely empty (in endpoints)
- NAK acknowledge has been transmitted to the host (isochronous-in only)
- IN token received when Tx-FIFO was empty (bulk-in/interrupt-in only)
- Out token received when endpoint was not yet enabled
- Babble error condition has been detected
- Endpoint disable by application is effective
- Endpoint NAK by application is effective (isochronous-in only)
- More than 3 back-to-back setup packets were received (control-out only)
- Timeout condition detected (control-in only)
- Isochronous out packet has been dropped, without generating an interrupt

#### **29.6 USB** host

This section gives the functional description of the OTG\_FS in the USB host mode. The OTG\_FS works as a USB host in the following circumstances:

- OTG A-host
  - OTG A-device default state when the A-side of the USB cable is plugged in
- OTG B-host
  - OTG B-device after HNP switching to the host role
- A-device
  - If the ID line is present, functional and connected to the A-side of the USB cable, and the HNP-capable bit is cleared in the Global USB Configuration register (HNPCAP bit in OTG\_FS\_GUSBCFG). Integrated pull-down resistors are automatically set on the DP/DM lines.
- Host only (see figure Figure 353: USB host-only connection).
  - The force host mode bit in the global USB configuration register (FHMOD bit in OTG\_FS\_GUSBCFG) forces the OTG\_FS core to work as a USB host-only. In this case, the ID line is ignored even if present on the USB connector. Integrated pulldown resistors are automatically set on the DP/DM lines.

Note:

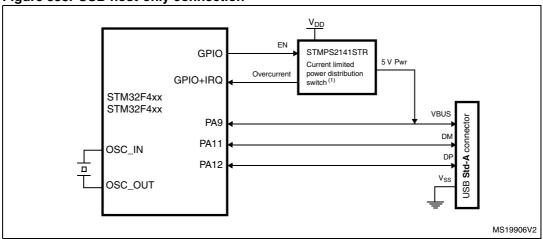
On-chip 5 V  $V_{BUS}$  generation is not supported. For this reason, a charge pump or, if 5 V are available on the application board, a basic power switch must be added externally to drive

the 5 V  $V_{BUS}$  line. The external charge pump can be driven by any GPIO output. This is required for the OTG A-host, A-device and host-only configurations.

The  $V_{BUS}$  input ensures that valid  $V_{BUS}$  levels are supplied by the charge pump during USB operations while the charge pump overcurrent output can be input to any GPIO pin configured to generate port interrupts. The overcurrent ISR must promptly disable the  $V_{BUS}$  generation.

The  $V_{BUS}$  pin can be freed by disabling the  $V_{BUS}$  sensing option. This is done by setting the NOVBUSSENS bit in the OTG\_FS\_GCCFG register. In this case the  $V_{BUS}$  is considered internally to be always at  $V_{BUS}$  valid level (5 V).





- STMPS2141STR needed only if the application has to support a V<sub>BUS</sub> powered device. A basic power switch can be used if 5 V are available on the application board.
- 2.  $V_{DD}$  range is between 2 V and 3.6 V.

### 29.6.1 SRP-capable host

SRP support is available through the SRP capable bit in the global USB configuration register (SRPCAP bit in OTG\_FS\_GUSBCFG). With the SRP feature enabled, the host can save power by switching off the  $V_{BUS}$  power while the USB session is suspended.

The SRP host mode program model is described in detail in the *A-device session request protocol*) section.

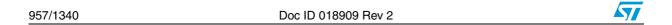
#### 29.6.2 USB host states

## Host port power

On-chip 5 V  $V_{BUS}$  generation is not supported. For this reason, a charge pump or, if 5 V are available on the application board, a basic power switch, must be added externally to drive the 5 V  $V_{BUS}$  line. The external charge pump can be driven by any GPIO output. When the application decides to power on  $V_{BUS}$  using the chosen GPIO, it must also set the port power bit in the host port control and status register (PPWR bit in OTG\_FS\_HPRT).

### V<sub>BUS</sub> valid

The  $V_{BUS}$  input ensures that valid  $V_{BUS}$  levels are supplied by the charge pump during USB operations.



Any unforeseen  $V_{BUS}$  voltage drop below the  $V_{BUS}$  valid threshold (4.25 V) leads to an OTG interrupt triggered by the session end detected bit (SEDET bit in OTG\_FS\_GOTGINT). The application is then required to remove the  $V_{BUS}$  power and clear the port power bit. The charge pump overcurrent flag can also be used to prevent electrical damage. Connect the overcurrent flag output from the charge pump to any GPIO input and configure it to generate a port interrupt on the active level. The overcurrent ISR must promptly disable the  $V_{BUS}$  generation and clear the port power bit.

#### Host detection of a peripheral connection

Even if USB peripherals or B-devices can be attached at any time, the OTG\_FS will not detect any bus connection until  $V_{BUS}$  is no longer sensed at a valid level (5 V).

When V<sub>BUS</sub> is at a valid level and a remote B-device is attached, the OTG\_FS core issues a host port interrupt triggered by the device connected bit in the host port control and status register (PCDET bit in OTG\_FS\_HPRT).

#### Host detection of peripheral a disconnection

The peripheral disconnection event triggers the disconnect detected interrupt (DISCINT bit in OTG\_FS\_GINTSTS).

#### **Host enumeration**

After detecting a peripheral connection the host must start the enumeration process by sending USB reset and configuration commands to the new peripheral.

Before starting to drive a USB reset, the application waits for the OTG interrupt triggered by the debounce done bit (DBCDNE bit in OTG\_FS\_GOTGINT), which indicates that the bus is stable again after the electrical debounce caused by the attachment of a pull-up resistor on DP (FS) or DM (LS).

The application drives a USB reset signaling (single-ended zero) over the USB by keeping the port reset bit set in the host port control and status register (PRST bit in OTG\_FS\_HPRT) for a minimum of 10 ms and a maximum of 20 ms. The application takes care of the timing count and then of clearing the port reset bit.

Once the USB reset sequence has completed, the host port interrupt is triggered by the port enable/disable change bit (PENCHNG bit in OTG\_FS\_HPRT). This informs the application that the speed of the enumerated peripheral can be read from the port speed field in the host port control and status register (PSPD bit in OTG\_FS\_HPRT) and that the host is starting to drive SOFs (FS) or Keep alives (LS). The host is now ready to complete the peripheral enumeration by sending peripheral configuration commands.

#### Host suspend

The application decides to suspend the USB activity by setting the port suspend bit in the host port control and status register (PSUSP bit in OTG\_FS\_HPRT). The OTG\_FS core stops sending SOFs and enters the suspended state.

The suspended state can be optionally exited on the remote device's initiative (remote wakeup). In this case the remote wakeup interrupt (WKUPINT bit in OTG\_FS\_GINTSTS) is generated upon detection of a remote wakeup signaling, the port resume bit in the host port control and status register (PRES bit in OTG\_FS\_HPRT) self-sets, and resume signaling is automatically driven over the USB. The application must time the resume window and then clear the port resume bit to exit the suspended state and restart the SOF.

If the suspended state is exited on the host initiative, the application must set the port resume bit to start resume signaling on the host port, time the resume window and finally clear the port resume bit.

#### 29.6.3 Host channels

The OTG\_FS core instantiates 8 host channels. Each host channel supports an USB host transfer (USB pipe). The host is not able to support more than 8 transfer requests at the same time. If more than 8 transfer requests are pending from the application, the host controller driver (HCD) must re-allocate channels when they become available from previous duty, that is, after receiving the transfer completed and channel halted interrupts.

Each host channel can be configured to support in/out and any type of periodic/nonperiodic transaction. Each host channel makes us of proper control (HCCHARx), transfer configuration (HCTSIZx) and status/interrupt (HCINTx) registers with associated mask (HCINTMSKx) registers.

#### Host channel control

- The following host channel controls are available to the application through the host channel-*x* characteristics register (HCCHAR*x*):
  - Channel enable/disable
  - Program the FS/LS speed of target USB peripheral
  - Program the address of target USB peripheral
  - Program the endpoint number of target USB peripheral
  - Program the transfer IN/OUT direction
  - Program the USB transfer type (control, bulk, interrupt, isochronous)
  - Program the maximum packet size (MPS)
  - Program the periodic transfer to be executed during odd/even frames

#### Host channel transfer

The host channel transfer size registers (HCTSIZx) allow the application to program the transfer size parameters, and read the transfer status. Programming must be done before setting the channel enable bit in the host channel characteristics register. Once the endpoint is enabled the packet count field is read-only as the OTG FS core updates it according to the current transfer status.

- The following transfer parameters can be programmed:
  - transfer size in bytes
  - number of packets making up the overall transfer size
  - initial data PID

#### Host channel status/interrupt

The host channel-*x* interrupt register (HCINT*x*) indicates the status of an endpoint with respect to USB- and AHB-related events. The application must read these register when the host channels interrupt bit in the core interrupt register (HCINT bit in OTG\_FS\_GINTSTS) is set. Before the application can read these registers, it must first read the host all channels interrupt (HCAINT) register to get the exact channel number for the host channel-*x* interrupt register. The application must clear the appropriate bit in this register to clear the

corresponding bits in the HAINT and GINTSTS registers. The mask bits for each interrupt source of each channel are also available in the OTG FS HCINTMSK-x register.

- The host core provides the following status checks and interrupt generation:
  - Transfer completed interrupt, indicating that the data transfer is complete on both the application (AHB) and USB sides
  - Channel has stopped due to transfer completed, USB transaction error or disable command from the application
  - Associated transmit FIFO is half or completely empty (IN endpoints)
  - ACK response received
  - NAK response received
  - STALL response received
  - USB transaction error due to CRC failure, timeout, bit stuff error, false EOP
  - Babble error
  - fraMe overrun
  - dAta toggle error

#### 29.6.4 Host scheduler

The host core features a built-in hardware scheduler which is able to autonomously re-order and manage the USB transaction requests posted by the application. At the beginning of each frame the host executes the periodic (isochronous and interrupt) transactions first, followed by the nonperiodic (control and bulk) transactions to achieve the higher level of priority granted to the isochronous and interrupt transfer types by the USB specification.

The host processes the USB transactions through request queues (one for periodic and one for nonperiodic). Each request queue can hold up to 8 entries. Each entry represents a pending transaction request from the application, and holds the IN or OUT channel number along with other information to perform a transaction on the USB. The order in which the requests are written to the queue determines the sequence of the transactions on the USB interface.

At the beginning of each frame, the host processes the periodic request queue first, followed by the nonperiodic request queue. The host issues an incomplete periodic transfer interrupt (IPXFR bit in OTG\_FS\_GINTSTS) if an isochronous or interrupt transaction scheduled for the current frame is still pending at the end of the current frame. The OTG HS core is fully responsible for the management of the periodic and nonperiodic request queues. The periodic transmit FIFO and queue status register (HPTXSTS) and nonperiodic transmit FIFO and queue status register (HNPTXSTS) are read-only registers which can be used by the application to read the status of each request queue. They contain:

- The number of free entries currently available in the periodic (nonperiodic) request queue (8 max)
- Free space currently available in the periodic (nonperiodic) Tx-FIFO (out-transactions)
- IN/OUT token, host channel number and other status information.

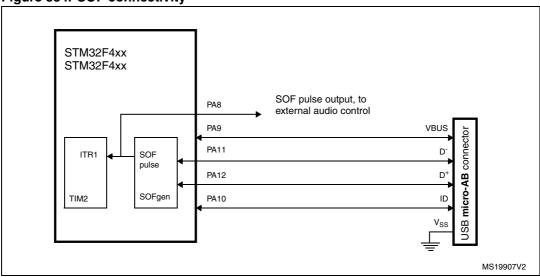
As request queues can hold a maximum of 8 entries each, the application can push to schedule host transactions in advance with respect to the moment they physically reach the SB for a maximum of 8 pending periodic transactions plus 8 pending nonperiodic transactions.

To post a transaction request to the host scheduler (queue) the application must check that there is at least 1 entry available in the periodic (nonperiodic) request queue by reading the

PTXQSAV bits in the OTG\_FS\_HNPTXSTS register or NPTQXSAV bits in the OTG\_FS\_HNPTXSTS register.

## 29.7 SOF trigger

Figure 354. SOF connectivity



The OTG FS core provides means to monitor, track and configure SOF framing in the host and peripheral, as well as an SOF pulse output connectivity feature.

Such utilities are especially useful for adaptive audio clock generation techniques, where the audio peripheral needs to synchronize to the isochronous stream provided by the PC, or the host needs to trim its framing rate according to the requirements of the audio peripheral.

#### 29.7.1 Host SOFs

In host mode the number of PHY clocks occurring between the generation of two consecutive SOF (FS) or Keep-alive (LS) tokens is programmable in the host frame interval register (HFIR), thus providing application control over the SOF framing period. An interrupt is generated at any start of frame (SOF bit in OTH\_FS\_GINTSTS). The current frame number and the time remaining until the next SOF are tracked in the host frame number register (HFNUM).

An SOF pulse signal, generated at any SOF starting token and with a width of 12 system clock cycles, can be made available externally on the SOF pin using the SOFOUTEN bit in the global control and configuration register. The SOF pulse is also internally connected to the input trigger of timer 2 (TIM2), so that the input capture feature, the output compare feature and the timer can be triggered by the SOF pulse. The TIM2 connection is enabled through the ITR1\_RMP bits of TIM2\_OR register.

### 29.7.2 Peripheral SOFs

In device mode, the start of frame interrupt is generated each time an SOF token is received on the USB (SOF bit in OTH\_FS\_GINTSTS). The corresponding frame number can be read from the device status register (FNSOF bit in OTG\_FS\_DSTS). An SOF pulse signal with a width of 12 system clock cycles is also generated and can be made available externally on

the SOF pin by using the SOF output enable bit in the global control and configuration register (SOFOUTEN bit in OTG\_FS\_GCCFG). The SOF pulse signal is also internally connected to the TIM2 input trigger, so that the input capture feature, the output compare feature and the timer can be triggered by the SOF pulse. The TIM2 connection is enabled through the ITR1\_RMP bits the TIM2 option register (TIM2\_OR).

The end of periodic frame interrupt (GINTSTS/EOPF) is used to notify the application when 80%, 85%, 90% or 95% of the time frame interval elapsed depending on the periodic frame interval field in the device configuration register (PFIVL bit in OTG\_FS\_DCFG). This feature can be used to determine if all of the isochronous traffic for that frame is complete.

## 29.8 Power options

The power consumption of the OTG PHY is controlled by three bits in the general core configuration register:

- PHY power down (GCCFG/PWRDWN)
   It switches on/off the full-speed transceiver module of the PHY. It must be preliminarily set to allow any USB operation.
- A-V<sub>BUS</sub> sensing enable (GCCFG/VBUSASEN)
   It switches on/off the V<sub>BUS</sub> comparators associated with A-device operations. It must be set when in A-device (USB host) mode and during HNP.
- B-V<sub>BUS</sub> sensing enable (GCCFG/VBUSASEN)
   It switches on/off the V<sub>BUS</sub> comparators associated with B-device operations. It must be set when in B-device (USB peripheral) mode and during HNP.

Power reduction techniques are available while in the USB suspended state, when the USB session is not yet valid or the device is disconnected.

- Stop PHY clock (STPPCLK bit in OTG\_FS\_PCGCCTL)
  - When setting the stop PHY clock bit in the clock gating control register, most of the 48 MHz clock domain internal to the OTG full-speed core is switched off by clock gating. The dynamic power consumption due to the USB clock switching activity is cut even if the 48 MHz clock input is kept running by the application
  - Most of the transceiver is also disabled, and only the part in charge of detecting the asynchronous resume or remote wakeup event is kept alive.
- Gate HCLK (GATEHCLK bit in OTG\_FS\_PCGCCTL)
  - When setting the Gate HCLK bit in the clock gating control register, most of the system clock domain internal to the OTG\_FS core is switched off by clock gating. Only the register read and write interface is kept alive. The dynamic power consumption due to the USB clock switching activity is cut even if the system clock is kept running by the application for other purposes.
- USB system stop
  - When the OTG\_FS is in the USB suspended state, the application may decide to drastically reduce the overall power consumption by a complete shut down of all the clock sources in the system. USB System Stop is activated by first setting the Stop PHY clock bit and then configuring the system deep sleep mode in the power control system module (PWR).

The OTG\_FS core automatically reactivates both system and USB clocks by asynchronous detection of remote wakeup (as an host) or resume (as a device) signaling on the USB.

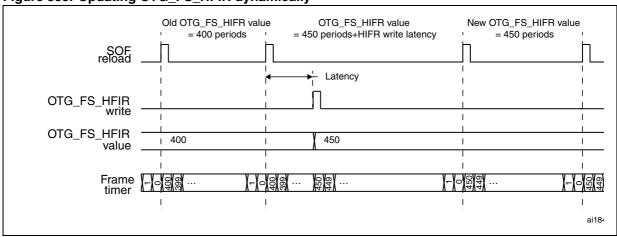
To save dynamic power, the USB data FIFO is clocked only when accessed by the OTG\_FS core.

## 29.9 Dynamic update of the OTG\_FS\_HFIR register

The USB core embeds a dynamic trimming capability of micro-SOF framing period in host mode allowing to synchronize an external device with the micro-SOF frames.

When the OTG\_HS\_HFIR register is changed within a current micro-SOF frame, the SOF period correction is applied in the next frame as described in *Figure 355*.





## 29.10 USB data FIFOs

The USB system features 1.25 Kbyte of dedicated RAM with a sophisticated FIFO control mechanism. The packet FIFO controller module in the OTG\_FS core organizes RAM space into Tx-FIFOs into which the application pushes the data to be temporarily stored before the USB transmission, and into a single Rx FIFO where the data received from the USB are temporarily stored before retrieval (popped) by the application. The number of instructed FIFOs and how these are organized inside the RAM depends on the device's role. In peripheral mode an additional Tx-FIFO is instructed for each active IN endpoint. Any FIFO size is software configured to better meet the application requirements.

## 29.11 Peripheral FIFO architecture

Single data ĔΙΕΩ DIEPTXF2[31:16] IN endpoint Tx FIFO #n Tx FIFO #n Dedicated Tx DFIFO push access packet FIFO #n contro DIEPTXFx[15:0] from AHB (optional) ■ MAC pop DIEPTXF2[15:0] IN endpoint Tx FIFO #1 Dedicated Tx DFIFO push access Tx FIFO #1 packet FIFO #1 contr DIEPTXF1[31:16] from AHB (optional) DIEPTXF1[15:0] MAC pop IN endpoint Tx FIFO #0 Dedicated Tx GNPTXFSIZ[31:16] Tx FIFO #0 packet DFIFO push access FIFO #0 contro from AHB (optional) - MAC pop GNPTXFSIZ[15:0] Any OUT endpoint DFIFO pop GRXFSIZ[31:16] Rx FIFO contro Rx packets access from AHB (Rx start -MAC push A1 = 0 address fixed to 0) ai15611

Figure 356. Device-mode FIFO address mapping and AHB FIFO access mapping

### 29.11.1 Peripheral Rx FIFO

The OTG peripheral uses a single receive FIFO that receives the data directed to all OUT endpoints. Received packets are stacked back-to-back until free space is available in the Rx-FIFO. The status of the received packet (which contains the OUT endpoint destination number, the byte count, the data PID and the validity of the received data) is also stored by the core on top of the data payload. When no more space is available, host transactions are NACKed and an interrupt is received on the addressed endpoint. The size of the receive FIFO is configured in the receive FIFO Size register (GRXFSIZ).

The single receive FIFO architecture makes it more efficient for the USB peripheral to fill in the receive RAM buffer:

- All OUT endpoints share the same RAM buffer (shared FIFO)
- The OTG FS core can fill in the receive FIFO up to the limit for any host sequence of OUT tokens

The application keeps receiving the Rx-FIFO non-empty interrupt (RXFLVL bit in OTG\_FS\_GINTSTS) as long as there is at least one packet available for download. It reads the packet information from the receive status read and pop register (GRXSTSP) and finally pops data off the receive FIFO by reading from the endpoint-related pop address.

## 29.11.2 Peripheral Tx FIFOs

The core has a dedicated FIFO for each IN endpoint. The application configures FIFO sizes by writing the non periodic transmit FIFO size register (OTG\_FS\_TX0FSIZ) for IN endpoint0 and the device IN endpoint transmit FIFOx registers (DIEPTXFx) for IN endpoint-x.

## 29.12 Host FIFO architecture

Figure 357. Host-mode FIFO address mapping and AHB FIFO access mapping Single data **FIFO** Any periodic channel Periodic Tx HPTXFSIZ[31:16] Periodic Tx packets DFIFO push access FIFO control from AHB (optional) — MAC pop HPTXFSIZ[15:0] Any non-periodic Periodic Tx packets NPTXFSIZ[31:16] channel DFIFO push Non-periodic access from AHB Tx FIFO contro NPTXFSIZ[15:0] MAC pop RXFSIZ[31:16] Rx packets Any channel DFIFO pop access from AHB Rx FIFO contro Rx start address fixed to 0 -MAC push A1 = 0ai15610

#### 29.12.1 Host Rx FIFO

The host uses one receiver FIFO for all periodic and nonperiodic transactions. The FIFO is used as a receive buffer to hold the received data (payload of the received packet) from the USB until it is transferred to the system memory. Packets received from any remote IN endpoint are stacked back-to-back until free space is available. The status of each received packet with the host channel destination, byte count, data PID and validity of the received data are also stored into the FIFO. The size of the receive FIFO is configured in the receive FIFO size register (GRXFSIZ).

The single receive FIFO architecture makes it highly efficient for the USB host to fill in the receive data buffer:

- All IN configured host channels share the same RAM buffer (shared FIFO)
- The OTG FS core can fill in the receive FIFO up to the limit for any sequence of IN tokens driven by the host software

The application receives the Rx FIFO not-empty interrupt as long as there is at least one packet available for download. It reads the packet information from the receive status read and pop register and finally pops the data off the receive FIFO.

#### 29.12.2 Host Tx FIFOs

The host uses one transmit FIFO for all non-periodic (control and bulk) OUT transactions and one transmit FIFO for all periodic (isochronous and interrupt) OUT transactions. FIFOs are used as transmit buffers to hold the data (payload of the transmit packet) to be transmitted over the USB. The size of the periodic (nonperiodic) Tx FIFO is configured in the host periodic (nonperiodic) transmit FIFO size (HPTXFSIZ/HNPTXFSIZ) register.

The two Tx FIFO implementation derives from the higher priority granted to the periodic type of traffic over the USB frame. At the beginning of each frame, the built-in host scheduler processes the periodic request queue first, followed by the nonperiodic request queue.

The two transmit FIFO architecture provides the USB host with separate optimization for periodic and nonperiodic transmit data buffer management:

- All host channels configured to support periodic (nonperiodic) transactions in the OUT direction share the same RAM buffer (shared FIFOs)
- The OTG FS core can fill in the periodic (nonperiodic) transmit FIFO up to the limit for any sequence of OUT tokens driven by the host software

The OTG\_FS core issues the periodic Tx FIFO empty interrupt (PTXFE bit in OTG\_FS\_GINTSTS) as long as the periodic Tx-FIFO is half or completely empty, depending on the value of the periodic Tx-FIFO empty level bit in the AHB configuration register (PTXFELVL bit in OTG\_FS\_GAHBCFG). The application can push the transmission data in advance as long as free space is available in both the periodic Tx FIFO and the periodic request queue. The host periodic transmit FIFO and queue status register (HPTXSTS) can be read to know how much space is available in both.

OTG\_FS core issues the non periodic Tx FIFO empty interrupt (NPTXFE bit in OTG\_FS\_GINTSTS) as long as the nonperiodic Tx FIFO is half or completely empty depending on the non periodic Tx FIFO empty level bit in the AHB configuration register (TXFELVL bit in OTG\_FS\_GAHBCFG). The application can push the transmission data as long as free space is available in both the nonperiodic Tx FIFO and nonperiodic request queue. The host nonperiodic transmit FIFO and queue status register (HNPTXSTS) can be read to know how much space is available in both.

#### 29.13 FIFO RAM allocation

#### **29.13.1** Device mode

Receive FIFO RAM allocation: the application should allocate RAM for SETUP Packets: 10 locations must be reserved in the receive FIFO to receive SETUP packets on control endpoint. The core does not use these locations, which are reserved for SETUP packets, to write any other data. One location is to be allocated for Global OUT NAK. Status information is written to the FIFO along with each received packet. Therefore, a minimum space of (Largest Packet Size / 4) + 1 must be allocated to receive packets. If multiple isochronous endpoints are enabled, then at least two (Largest Packet Size / 4) + 1 spaces must be allocated to receive back-to-back packets. Typically, two (Largest Packet Size / 4) + 1 spaces are recommended so that when the previous packet is being transferred to the CPU, the USB can receive the subsequent packet.

Along with the last packet for each endpoint, transfer complete status information is also pushed to the FIFO. Typically, one location for each OUT endpoint is recommended.

**Transmit FIFO RAM allocation:** the minimum RAM space required for each IN Endpoint Transmit FIFO is the maximum packet size for that particular IN endpoint.

Note:

More space allocated in the transmit IN Endpoint FIFO results in better performance on the USB.

#### 29.13.2 Host mode

#### **Receive FIFO RAM allocation**

Status information is written to the FIFO along with each received packet. Therefore, a minimum space of (Largest Packet Size / 4) + 1 must be allocated to receive packets. If multiple isochronous channels are enabled, then at least two (Largest Packet Size / 4) + 1 spaces must be allocated to receive back-to-back packets. Typically, two (Largest Packet Size / 4) + 1 spaces are recommended so that when the previous packet is being transferred to the CPU, the USB can receive the subsequent packet.

Along with the last packet in the host channel, transfer complete status information is also pushed to the FIFO. So one location must be allocated for this.

#### Transmit FIFO RAM allocation

The minimum amount of RAM required for the host Non-periodic Transmit FIFO is the largest maximum packet size among all supported non-periodic OUT channels.

Typically, two Largest Packet Sizes worth of space is recommended, so that when the current packet is under transfer to the USB, the CPU can get the next packet.

The minimum amount of RAM required for host periodic Transmit FIFO is the largest maximum packet size out of all the supported periodic OUT channels. If there is at least one Isochronous OUT endpoint, then the space must be at least two times the maximum packet size of that channel.

Note:

More space allocated in the Transmit Non-periodic FIFO results in better performance on the USB.

## 29.14 USB system performance

Best USB and system performance is achieved owing to the large RAM buffers, the highly configurable FIFO sizes, the quick 32-bit FIFO access through AHB push/pop registers and, especially, the advanced FIFO control mechanism. Indeed, this mechanism allows the

OTG\_FS to fill in the available RAM space at best regardless of the current USB sequence. With these features:

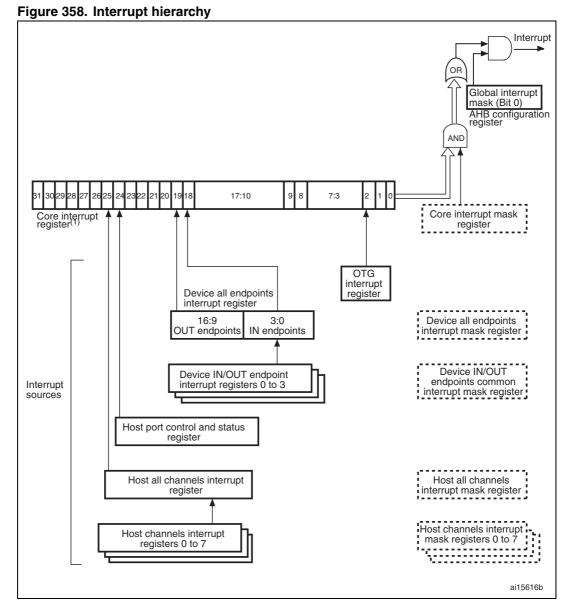
- The application gains good margins to calibrate its intervention in order to optimize the CPU bandwidth usage:
  - It can accumulate large amounts of transmission data in advance compared to when they are effectively sent over the USB
  - It benefits of a large time margin to download data from the single receive FIFO
- The USB Core is able to maintain its full operating rate, that is to provide maximum full-speed bandwidth with a great margin of autonomy versus application intervention:
  - It has a large reserve of transmission data at its disposal to autonomously manage the sending of data over the USB
  - It has a lot of empty space available in the receive buffer to autonomously fill it in with the data coming from the USB

As the OTG\_FS core is able to fill in the 1.25 Kbyte RAM buffer very efficiently, and as 1.25 Kbyte of transmit/receive data is more than enough to cover a full speed frame, the USB system is able to withstand the maximum full-speed data rate for up to one USB frame (1 ms) without any CPU intervention.

## 29.15 OTG\_FS interrupts

When the OTG\_FS controller is operating in one mode, either device or host, the application must not access registers from the other mode. If an illegal access occurs, a mode mismatch interrupt is generated and reflected in the Core interrupt register (MMIS bit in the OTG\_FS\_GINTSTS register). When the core switches from one mode to the other, the registers in the new mode of operation must be reprogrammed as they would be after a power-on reset.

Figure 358 shows the interrupt hierarchy.



 The core interrupt register bits are shown in OTG\_FS core interrupt register (OTG\_FS\_GINTSTS) on page 984.

## 29.16 OTG\_FS control and status registers

By reading from and writing to the control and status registers (CSRs) through the AHB slave interface, the application controls the OTG\_FS controller. These registers are 32 bits wide, and the addresses are 32-bit block aligned. The OTG\_FS registers must be accessed by words (32 bits).

CSRs are classified as follows:

- Core global registers
- Host-mode registers
- Host global registers
- Host port CSRs
- Host channel-specific registers
- Device-mode registers
- Device global registers
- Device endpoint-specific registers
- Power and clock-gating registers
- Data FIFO (DFIFO) access registers

Only the Core global, Power and clock-gating, Data FIFO access, and host port control and status registers can be accessed in both host and device modes. When the OTG\_FS controller is operating in one mode, either device or host, the application must not access registers from the other mode. If an illegal access occurs, a mode mismatch interrupt is generated and reflected in the Core interrupt register (MMIS bit in the OTG\_FS\_GINTSTS register). When the core switches from one mode to the other, the registers in the new mode of operation must be reprogrammed as they would be after a power-on reset.

## 29.16.1 **CSR** memory map

The host and device mode registers occupy different addresses. All registers are implemented in the AHB clock domain.

0000h Core global CSRs (1 Kbyte) 0400h Host mode CSRs (1 Kbyte) 0800h Device mode CSRs (1.5 Kbyte) 0E00h Power and clock gating CSRs (0.5 Kbyte) 1000h Device EP 0/Host channel 0 FIFO (4 Kbyte) 2000h Device EP1/Host channel 1 FIFO (4 Kbyte) 3000h **DFIFO** push/pop to this region Device EP  $(x-1)^{(1)}$ /Host channel  $(x-1)^{(1)}$  FIFO (4 Kbyte) Device EP x<sup>(1)</sup>/Host channel x<sup>(1)</sup> FIFO (4 Kbyte) Reserved 2 0000h DFIFO debug read/ Direct access to data FIFO RAM for debugging (128 Kbyte) region 3 FFFFh ai15615b

Figure 359. CSR memory map

## **Global CSR map**

These registers are available in both host and device modes.

Table 153. Core global control and status registers (CSRs)

Acronym	Address offset	Register name
OTG_FS_GOTGCTL	0x000	OTG_FS control and status register (OTG_FS_GOTGCTL) on page 975
OTG_FS_GOTGINT	0x004	OTG_FS interrupt register (OTG_FS_GOTGINT) on page 977
OTG_FS_GAHBCFG	0x008	OTG_FS AHB configuration register (OTG_FS_GAHBCFG) on page 979
OTG_FS_GUSBCFG	0x00C	OTG_FS USB configuration register (OTG_FS_GUSBCFG) on page 980
OTG_FS_GRSTCTL	0x010	OTG_FS reset register (OTG_FS_GRSTCTL) on page 982
OTG_FS_GINTSTS	0x014	OTG_FS core interrupt register (OTG_FS_GINTSTS) on page 984
OTG_FS_GINTMSK	0x018	OTG_FS interrupt mask register (OTG_FS_GINTMSK) on page 988

<sup>1.</sup> x = 3 in device mode and x = 7 in host mode.

Table 153. Core global control and status registers (CSRs) (continued)

Acronym	Address offset	Register name
OTG_FS_GRXSTSR	0x01C	OTG_FS Receive status debug read/OTG status read and pop registers
OTG_FS_GRXSTSP	0x020	(OTG_FS_GRXSTSR/OTG_FS_GRXSTSP) on page 991
OTG_FS_GRXFSIZ	0x024	OTG_FS Receive FIFO size register (OTG_FS_GRXFSIZ) on page 992
OTG_FS_HNPTXFSIZ/ OTG_FS_DIEPTXF0 <sup>(1)</sup>	0x028	OTG_FS Host non-periodic transmit FIFO size register (OTG_FS_HNPTXFSIZ)/Endpoint 0 Transmit FIFO size (OTG_FS_DIEPTXF0)
OTG_FS_HNPTXSTS	0x02C	OTG_FS non-periodic transmit FIFO/queue status register (OTG_FS_HNPTXSTS) on page 993
OTG_FS_GCCFG	0x038	OTG_FS general core configuration register (OTG_FS_GCCFG) on page 994
OTG_FS_CID	0x03C	OTG_FS core ID register (OTG_FS_CID) on page 995
OTG_FS_HPTXFSIZ	0x100	OTG_FS Host periodic transmit FIFO size register (OTG_FS_HPTXFSIZ) on page 996
OTG_FS_DIEPTXFx	0x104 0x124  0x138	OTG_FS device IN endpoint transmit FIFO size register (OTG_FS_DIEPTXFx) ( $x = 13$ , where $x$ is the FIFO_number) on page 997

<sup>1.</sup> The general rule is to use OTG\_FS\_HNPTXFSIZ for host mode and OTG\_FS\_DIEPTXF0 for device mode.

## **Host-mode CSR map**

These registers must be programmed every time the core changes to host mode.

Table 154. Host-mode control and status registers (CSRs)

Acronym	Offset address	Register name
OTG_FS_HCFG	0x400	OTG_FS Host configuration register (OTG_FS_HCFG) on page 997
OTG_FS_HFIR	0x404	OTG_FS Host frame interval register (OTG_FS_HFIR) on page 998
OTG_FS_HFNUM	0x408	OTG_FS Host frame number/frame time remaining register (OTG_FS_HFNUM) on page 999
OTG_FS_HPTXSTS	0x410	OTG_FS_Host periodic transmit FIFO/queue status register (OTG_FS_HPTXSTS) on page 999
OTG_FS_HAINT	0x414	OTG_FS Host all channels interrupt register (OTG_FS_HAINT) on page 1000
OTG_FS_HAINTMSK	0x418	OTG_FS Host all channels interrupt mask register (OTG_FS_HAINTMSK) on page 1001
OTG_FS_HPRT	0x440	OTG_FS Host port control and status register (OTG_FS_HPRT) on page 1001

Table 154. Host-mode control and status registers (CSRs) (continued)

Acronym	Offset address	Register name
OTG_FS_HCCHARx	0x500 0x520  0x6E0h	OTG_FS Host channel- $x$ characteristics register (OTG_FS_HCCHAR $x$ ) ( $x = 07$ , where $x = Channel_number$ ) on page 1004
OTG_FS_HCINTx	508h	$OTG\_FS$ Host channel-x interrupt register ( $OTG\_FS\_HCINTx$ ) ( $x = 07$ , where $x = Channel\_number$ ) on page 1005
OTG_FS_HCINTMSKx	50Ch	$OTG\_FS$ Host channel-x interrupt mask register ( $OTG\_FS\_HCINTMSKx$ ) ( $x = 07$ , where $x = Channel\_number$ ) on page 1006
OTG_FS_HCTSIZx	510h	$OTG\_FS$ Host channel-x transfer size register ( $OTG\_FS\_HCTSIZx$ ) ( $x = 07$ , where $x = Channel\_number$ ) on page 1007

# **Device-mode CSR map**

These registers must be programmed every time the core changes to device mode.

Table 155. Device-mode control and status registers

Acronym	Offset address	Register name
OTG_FS_DCFG	0x800	OTG_FS device configuration register (OTG_FS_DCFG) on page 1008
OTG_FS_DCTL	0x804	OTG_FS device control register (OTG_FS_DCTL) on page 1009
OTG_FS_DSTS	0x808	OTG_FS device status register (OTG_FS_DSTS) on page 1010
OTG_FS_DIEPMSK	0x810	OTG_FS device IN endpoint common interrupt mask register (OTG_FS_DIEPMSK) on page 1011
OTG_FS_DOEPMSK	0x814	OTG_FS device OUT endpoint common interrupt mask register (OTG_FS_DOEPMSK) on page 1012
OTG_FS_DAINT	0x818	OTG_FS device all endpoints interrupt register (OTG_FS_DAINT) on page 1013
OTG_FS_DAINTMSK	0x81C	OTG_FS all endpoints interrupt mask register (OTG_FS_DAINTMSK) on page 1014
OTG_FS_DVBUSDIS	0x828	OTG_FS device VBUS discharge time register (OTG_FS_DVBUSDIS) on page 1014
OTG_FS_DVBUSPULSE	0x82C	OTG_FS device VBUS pulsing time register (OTG_FS_DVBUSPULSE) on page 1015
OTG_FS_DIEPEMPMSK	0x834	OTG_FS device IN endpoint FIFO empty interrupt mask register: (OTG_FS_DIEPEMPMSK) on page 1015
OTG_FS_DIEPCTL0	0x900	OTG_FS device control IN endpoint 0 control register (OTG_FS_DIEPCTL0) on page 1016

Table 155. Device-mode control and status registers (continued)

Acronym	Offset address	Register name
OTG_FS_DIEPCTLx	0x920 0x940  0xAE0	OTG device endpoint-x control register (OTG_FS_DIEPCTLx) (x = 13, where x = Endpoint_number) on page 1017
OTG_FS_DIEPINTx	0x908	OTG_FS device endpoint-x interrupt register (OTG_FS_DIEPINTx) $(x = 03, where \ x = Endpoint_number)$ on page 1024
OTG_FS_DIEPTSIZ0	0x910	OTG_FS device IN endpoint 0 transfer size register (OTG_FS_DIEPTSIZ0) on page 1026
OTG_FS_DTXFSTSx	0x918	OTG_FS device IN endpoint transmit FIFO status register (OTG_FS_DTXFSTSx) (x = 03, where x = Endpoint_number) on page 1029
OTG_FS_DIEPTSIZx	0x930 0x950  0xAF0	OTG_FS device OUT endpoint-x transfer size register (OTG_FS_DOEPTSIZx) (x = 13, where x = Endpoint_number) on page 1029
OTG_FS_DOEPCTL0	0xB00	OTG_FS device control OUT endpoint 0 control register (OTG_FS_DOEPCTL0) on page 1020
OTG_FS_DOEPCTLx	0xB20 0xB40  0xCC0 0xCE0 0xCFD	OTG device endpoint-x control register (OTG_FS_DIEPCTLx) (x = 13, where x = Endpoint_number) on page 1017
OTG_FS_DOEPINTx	0xB08	OTG_FS device endpoint-x interrupt register (OTG_FS_DIEPINTx) (x = 03, where x = Endpoint_number) on page 1024
OTG_FS_DOEPTSIZx	0xB10	OTG_FS device OUT endpoint-x transfer size register (OTG_FS_DOEPTSIZx) ( $x = 13$ , where $x = \text{Endpoint\_number}$ ) on page 1029

# Data FIFO (DFIFO) access register map

These registers, available in both host and device modes, are used to read or write the FIFO space for a specific endpoint or a channel, in a given direction. If a host channel is of type IN, the FIFO can only be read on the channel. Similarly, if a host channel is of type OUT, the FIFO can only be written on the channel.

Table 156. Data FIFO (DFIFO) access register map

FIFO access register section	Address range	Access
Device IN Endpoint 0/Host OUT Channel 0: DFIFO Write Access Device OUT Endpoint 0/Host IN Channel 0: DFIFO Read Access	0x1000-0x1FFC	w r
Device IN Endpoint 1/Host OUT Channel 1: DFIFO Write Access Device OUT Endpoint 1/Host IN Channel 1: DFIFO Read Access	0x2000-0x2FFC	w r
Device IN Endpoint $x^{(1)}$ /Host OUT Channel $x^{(1)}$ : DFIFO Write Access Device OUT Endpoint $x^{(1)}$ /Host IN Channel $x^{(1)}$ : DFIFO Read Access	0xX000h-0xXFFCh	w r

<sup>1.</sup> Where x is 3 in device mode and 7 in host mode.

# Power and clock gating CSR map

There is a single register for power and clock gating. It is available in both host and device modes.

Table 157. Power and clock gating control and status registers

Register name	Acronym	Offset address: 0xE00-0xFFF
Power and clock gating control register	PCGCR	0xE00-0xE04
Reserved		0xE05-0xFFF

# 29.16.2 OTG\_FS global registers

These registers are available in both host and device modes, and do not need to be reprogrammed when switching between these modes.

Bit values in the register descriptions are expressed in binary unless otherwise specified.

# OTG\_FS control and status register (OTG\_FS\_GOTGCTL)

Address offset: 0x000

Reset value: 0x0000 0800

The OTG\_FS\_GOTGCTL register controls the behavior and reflects the status of the OTG function of the core.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	BSVLD	ASVLD	DBCT	CIDSTS	Reserved	DHNPEN	HSHNPEN	HNPRQ	HNGSCS	Reserved	SRQ	SRQSCS
	r	r	r	r		rw	rw	rw	r		rw	r

## Bits 31:20 Reserved, must be kept at reset value.

#### Bit 19 BSVLD: B-session valid

Indicates the device mode transceiver status.

- 0: B-session is not valid.
- 1: B-session is valid.
- In OTG mode, you can use this bit to determine if the device is connected or disconnected.

Note: Only accessible in device mode.

#### Bit 18 ASVLD: A-session valid

Indicates the host mode transceiver status.

0: A-session is not valid

1: A-session is valid

Note: Only accessible in host mode.

## Bit 17 DBCT: Long/short debounce time

Indicates the debounce time of a detected connection.

- 0: Long debounce time, used for physical connections (100 ms +  $2.5 \mu s$ )
- 1: Short debounce time, used for soft connections (2.5 µs)

Note: Only accessible in host mode.

## Bit 16 CIDSTS: Connector ID status

Indicates the connector ID status on a connect event.

0: The OTG\_FS controller is in A-device mode

1: The OTG\_FS controller is in B-device mode

Note: Accessible in both device and host modes.

## Bits 15:12 Reserved, must be kept at reset value.

## Bit 11 DHNPEN: Device HNP enabled

The application sets this bit when it successfully receives a SetFeature.SetHNPEnable command from the connected USB host.

0: HNP is not enabled in the application

1: HNP is enabled in the application

Note: Only accessible in device mode.

## Bit 10 HSHNPEN: host set HNP enable

The application sets this bit when it has successfully enabled HNP (using the SetFeature.SetHNPEnable command) on the connected device.

0: Host Set HNP is not enabled

1: Host Set HNP is enabled

Note: Only accessible in host mode.

## Bit 9 HNPRQ: HNP request

The application sets this bit to initiate an HNP request to the connected USB host. The application can clear this bit by writing a 0 when the host negotiation success status change bit in the OTG\_FS\_GOTGINT register (HNSSCHG bit in OTG\_FS\_GOTGINT) is set. The core clears this bit when the HNSSCHG bit is cleared.

0: No HNP request

1: HNP request

Note: Only accessible in device mode.

#### Bit 8 HNGSCS: Host negotiation success

The core sets this bit when host negotiation is successful. The core clears this bit when the HNP Request (HNPRQ) bit in this register is set.

0: Host negotiation failure

1: Host negotiation success

Note: Only accessible in device mode.

## Bits 7:2 Reserved, must be kept at reset value.

#### Bit 1 SRQ: Session request

The application sets this bit to initiate a session request on the USB. The application can clear this bit by writing a 0 when the host negotiation success status change bit in the OTG\_FS\_GOTGINT register (HNSSCHG bit in OTG\_FS\_GOTGINT) is set. The core clears this bit when the HNSSCHG bit is cleared.

If you use the USB 1.1 full-speed serial transceiver interface to initiate the session request, the application must wait until  $V_{BUS}$  discharges to 0.2 V, after the B-Session Valid bit in this register (BSVLD bit in OTG\_FS\_GOTGCTL) is cleared. This discharge time varies between different PHYs and can be obtained from the PHY vendor.

0: No session request

1: Session request

Note: Only accessible in device mode.

#### Bit 0 SRQSCS: Session request success

The core sets this bit when a session request initiation is successful.

0: Session request failure1: Session request success

Note: Only accessible in device mode.

# OTG\_FS interrupt register (OTG\_FS\_GOTGINT)

Address offset: 0x04

Reset value: 0x0000 0000

The application reads this register whenever there is an OTG interrupt and clears the bits in this register to clear the OTG interrupt.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
					Rese	erved	d					DBCDNE	АБТОСНВ	HNGDET			Re	eserv	ed′			HNSSCHG	SRSSCHG		Re	eserv	ed		SEDET	Re	es.

Bits 31:20 Reserved, must be kept at reset value.

## Bit 19 **DBCDNE**: Debounce done

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The core sets this bit when the debounce is completed after the device connect. The application can start driving USB reset after seeing this interrupt. This bit is only valid when the HNP Capable or SRP Capable bit is set in the OTG\_FS\_GUSBCFG register (HNPCAP bit or SRPCAP bit in OTG\_FS\_GUSBCFG, respectively).

Note: Only accessible in host mode.

#### Bit 18 ADTOCHG: A-device timeout change

The core sets this bit to indicate that the A-device has timed out while waiting for the B-device to connect.

Note: Accessible in both device and host modes.

# Bit 17 HNGDET: Host negotiation detected

The core sets this bit when it detects a host negotiation request on the USB.

Note: Accessible in both device and host modes.

Bits 16:10 Reserved, must be kept at reset value.

## Bit 9 HNSSCHG: Host negotiation success status change

The core sets this bit on the success or failure of a USB host negotiation request. The application must read the host negotiation success bit of the OTG\_FS\_GOTGCTL register (HNGSCS in OTG\_FS\_GOTGCTL) to check for success or failure.

Note: Accessible in both device and host modes.

Bits 7:3 Reserved, must be kept at reset value.

#### Bit 8 SRSSCHG: Session request success status change

The core sets this bit on the success or failure of a session request. The application must read the session request success bit in the OTG\_FS\_GOTGCTL register (SRQSCS bit in OTG\_FS\_GOTGCTL) to check for success or failure.

Note: Accessible in both device and host modes.

#### Bit 2 SEDET: Session end detected

The core sets this bit to indicate that the level of the voltage on  $V_{BUS}$  is no longer valid for a B-Peripheral session when  $V_{BUS} < 0.8 \text{ V}$ .

Bits 1:0 Reserved, must be kept at reset value.

# OTG FS AHB configuration register (OTG FS GAHBCFG)

Address offset: 0x008 Reset value: 0x0000 0000

This register can be used to configure the core after power-on or a change in mode. This register mainly contains AHB system-related configuration parameters. Do not change this register after the initial programming. The application must program this register before starting any transactions on either the AHB or the USB.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
										Re	eserv	ed											PTXFELVL	TXFELVL			Rese	erved	I		GINTMSK
																							rw	rw							rw

Bits 31:20 Reserved, must be kept at reset value.

## Bit 8 PTXFELVL: Periodic TxFIFO empty level

Indicates when the periodic TxFIFO empty interrupt bit in the OTG\_FS\_GINTSTS register (PTXFE bit in OTG\_FS\_GINTSTS) is triggered.

0: PTXFE (in OTG\_FS\_GINTSTS) interrupt indicates that the Periodic TxFIFO is half empty 1: PTXFE (in OTG\_FS\_GINTSTS) interrupt indicates that the Periodic TxFIFO is completely empty

Note: Only accessible in host mode.

## Bit 7 TXFELVL: TxFIFO empty level

In device mode, this bit indicates when IN endpoint Transmit FIFO empty interrupt (TXFE in OTG\_FS\_DIEPINTx.) is triggered.

0: the TXFE (in OTG\_FS\_DIEPINTx) interrupt indicates that the IN Endpoint TxFIFO is half

1: the TXFE (in OTG\_FS\_DIEPINTx) interrupt indicates that the IN Endpoint TxFIFO is completely empty

In host mode, this bit indicates when the nonperiodic Tx FIFO empty interrupt (NPTXFE bit in OTG\_FS\_GINTSTS) is triggered:

0: the NPTXFE (in OTG\_FS\_GINTSTS) interrupt indicates that the nonperiodic Tx FIFO is half empty

1: the NPTXFE (in OTG\_FS\_GINTSTS) interrupt indicates that the nonperiodic Tx FIFO is completely empty

## Bits 6:1 Reserved, must be kept at reset value.

# Bit 0 GINTMSK: Global interrupt mask

The application uses this bit to mask or unmask the interrupt line assertion to itself. Irrespective of this bit's setting, the interrupt status registers are updated by the core.

0: Mask the interrupt assertion to the application.

1: Unmask the interrupt assertion to the application.

Note: Accessible in both device and host modes.

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# OTG\_FS USB configuration register (OTG\_FS\_GUSBCFG)

Address offset: 0x00C Reset value: 0x0000 0A00

This register can be used to configure the core after power-on or a changing to host mode or device mode. It contains USB and USB-PHY related configuration parameters. The application must program this register before starting any transactions on either the AHB or the USB. Do not make changes to this register after the initial programming.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

CTXPKT	FDMOD	FHMOD	Reserved	TRDT	HNPCAP	SRPCAP	Res.	PHYSEL	Reserved	TOCAL
rw	rw	rw		rw	r/rw	r/rw		wo		rw

Bits 31:20 Reserved, must be kept at reset value.

#### Bit 31 CTXPKT: Corrupt Tx packet

This bit is for debug purposes only. Never set this bit to 1.

Note: Accessible in both device and host modes.

#### Bit 30 **FDMOD:** Force device mode

Writing a 1 to this bit forces the core to device mode irrespective of the OTG\_FS\_ID input pin.

0: Normal mode

1: Force device mode

After setting the force bit, the application must wait at least 25 ms before the change takes effect.

Note: Accessible in both device and host modes.

## Bit 29 FHMOD: Force host mode

Writing a 1 to this bit forces the core to host mode irrespective of the OTG\_FS\_ID input pin.

- 0: Normal mode
- 1: Force host mode

After setting the force bit, the application must wait at least 25 ms before the change takes effect.

Note: Accessible in both device and host modes.

Bits 28:14 Reserved, must be kept at reset value.

#### Bits 13:10 TRDT: USB turnaround time

Sets the turnaround time in PHY clocks.

To calculate the value of TRDT, use the following formula:

 $TRDT = 4 \times AHB \ clock + 1 \ PHY \ clock$ 

#### Examples:

- 1. if AHB clock = 72 MHz (PHY Clock is 48), the TRDT is set to 9.
- 2. if AHB clock = 48 MHz (PHY Clock is 48), the TRDT is set to 5.

Note: Only accessible in device mode.

## Bit 9 HNPCAP: HNP-capable

The application uses this bit to control the OTG\_FS controller's HNP capabilities.

- 0: HNP capability is not enabled.
- 1: HNP capability is enabled.

Note: Accessible in both device and host modes.

## Bit 8 SRPCAP: SRP-capable

The application uses this bit to control the OTG\_FS controller's SRP capabilities. If the core operates as a non-SRP-capable

B-device, it cannot request the connected A-device (host) to activate  $V_{\text{BUS}}$  and start a session.

0: SRP capability is not enabled.

1: SRP capability is enabled.

Note: Accessible in both device and host modes.

Bit 7 Reserved, must be kept at reset value.

Bit 6 PHYSEL: Full Speed serial transceiver select

This bit is always 1 with write-only access.

Bits5:3 Reserved, must be kept at reset value.

#### Bits 2:0 TOCAL: FS timeout calibration

The number of PHY clocks that the application programs in this field is added to the full-speed interpacket timeout duration in the core to account for any additional delays introduced by the PHY. This can be required, because the delay introduced by the PHY in generating the line state condition can vary from one PHY to another.

The USB standard timeout value for full-speed operation is 16 to 18 (inclusive) bit times. The application must program this field based on the speed of enumeration. The number of bit times added per PHY clock is 0.25 bit times.



# OTG\_FS reset register (OTG\_FS\_GRSTCTL)

Address offset: 0x10

Reset value: 0x2000 0000

The application uses this register to reset various hardware features inside the core.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

AHBIDL	Reserved	TXFNUM	TXFFLSH	RXFFLSH	leserved	FCRST	HSRST	CSRST
r		rw	rs	rs	н	rs	rs	rs

#### Bit 31 AHBIDL: AHB master idle

Indicates that the AHB master state machine is in the Idle condition.

Note: Accessible in both device and host modes.

Bits 30:11 Reserved, must be kept at reset value.

#### Bits 10:6 TXFNUM: TxFIFO number

This is the FIFO number that must be flushed using the TxFIFO Flush bit. This field must not be changed until the core clears the TxFIFO Flush bit.

#### 00000:

- Non-periodic TxFIFO flush in host mode
- Tx FIFO 0 flush in device mode

#### 00001:

- Periodic TxFIFO flush in host mode
- TXFIFO 1 flush in device mode

00010: TXFIFO 2 flush in device mode

...

00101: TXFIFO 15 flush in device mode

10000: Flush all the transmit FIFOs in device or host mode.

Note: Accessible in both device and host modes.

## Bit 5 TXFFLSH: TxFIFO flush

This bit selectively flushes a single or all transmit FIFOs, but cannot do so if the core is in the midst of a transaction.

The application must write this bit only after checking that the core is neither writing to the TxFIFO nor reading from the TxFIFO. Verify using these registers:

Read—NAK Effective Interrupt ensures the core is not reading from the FIFO

Write—AHBIDL bit in OTG\_FS\_GRSTCTL ensures the core is not writing anything to the FIFO.

Note: Accessible in both device and host modes.

#### Bit 4 RXFFLSH: RxFIFO flush

The application can flush the entire RxFIFO using this bit, but must first ensure that the core is not in the middle of a transaction.

The application must only write to this bit after checking that the core is neither reading from the RxFIFO nor writing to the RxFIFO.

The application must wait until the bit is cleared before performing any other operations. This bit requires 8 clocks (slowest of PHY or AHB clock) to clear.

Note: Accessible in both device and host modes.

Bit 3 Reserved, must be kept at reset value.

#### Bit 2 FCRST: Host frame counter reset

The application writes this bit to reset the frame number counter inside the core. When the frame counter is reset, the subsequent SOF sent out by the core has a frame number of 0.

Note: Only accessible in host mode.

## Bit 1 HSRST: HCLK soft reset

The application uses this bit to flush the control logic in the AHB Clock domain. Only AHB Clock Domain pipelines are reset.

FIFOs are not flushed with this bit.

All state machines in the AHB clock domain are reset to the Idle state after terminating the transactions on the AHB, following the protocol.

CSR control bits used by the AHB clock domain state machines are cleared.

To clear this interrupt, status mask bits that control the interrupt status and are generated by the AHB clock domain state machine are cleared.

Because interrupt status bits are not cleared, the application can get the status of any core events that occurred after it set this bit.

This is a self-clearing bit that the core clears after all necessary logic is reset in the core. This can take several clocks, depending on the core's current state.

Note: Accessible in both device and host modes.

#### Bit 0 CSRST: Core soft reset

Resets the HCLK and PCLK domains as follows:

Clears the interrupts and all the CSR register bits except for the following bits:

- RSTPDMODL bit in OTG\_FS\_PCGCCTL
- GAYEHCLK bit in OTG\_FS\_PCGCCTL
- PWRCLMP bit in OTG\_FS\_PCGCCTL
- STPPCLK bit in OTG\_FS\_PCGCCTL
- FSLSPCS bit in OTG\_FS\_HCFG
- DSPD bit in OTG\_FS\_DCFG

All module state machines (except for the AHB slave unit) are reset to the Idle state, and all the transmit FIFOs and the receive FIFO are flushed.

Any transactions on the AHB Master are terminated as soon as possible, after completing the last data phase of an AHB transfer. Any transactions on the USB are terminated immediately. The application can write to this bit any time it wants to reset the core. This is a self-clearing bit and the core clears this bit after all the necessary logic is reset in the core, which can take several clocks, depending on the current state of the core. Once this bit has been cleared, the software must wait at least 3 PHY clocks before accessing the PHY domain (synchronization delay). The software must also check that bit 31 in this register is set to 1 (AHB Master is Idle) before starting any operation.

Typically, the software reset is used during software development and also when you dynamically change the PHY selection bits in the above listed USB configuration registers. When you change the PHY, the corresponding clock for the PHY is selected and used in the PHY domain. Once a new clock is selected, the PHY domain has to be reset for proper operation.

Note: Accessible in both device and host modes.



# OTG\_FS core interrupt register (OTG\_FS\_GINTSTS)

Address offset: 0x014

Reset value: 0x0400 0020

This register interrupts the application for system-level events in the current mode (device mode or host mode).

Some of the bits in this register are valid only in host mode, while others are valid in device mode only. This register also indicates the current mode. To clear the interrupt status bits of the rc w1 type, the application must write 1 into the bit.

The FIFO status interrupts are read-only; once software reads from or writes to the FIFO while servicing these interrupts, FIFO interrupt conditions are cleared automatically.

The application must clear the OTG\_FS\_GINTSTS register at initialization before unmasking the interrupt bit to avoid any interrupts generated prior to initialization.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
WKUINT	SRQINT	DISCINT	CIDSCHG	Reserved	PTXFE	HCINT	HPRTINT	Reserved		IPXFR/INCOMPISOOUT	IISOIXFR	OEPINT	IEPINT	Contract of the contract of th	0000000	EOPF	ISOODRP	ENUMDNE	USBRST	USBSUSP	ESUSP	Beggge		GOUTNAKEFF	GINAKEFF	NPTXFE	RXFLVL	SOF	OTGINT	SIMM	СМОД
	rc_	_w1			r	r	r	Res	S.	rc_	w1	r	r					rc_	_w1					r	r	r	r	rc_w1	r	rc_w1	r

## Bit 31 WKUPINT: Resume/remote wakeup detected interrupt

In device mode, this interrupt is asserted when a resume is detected on the USB. In host mode, this interrupt is asserted when a remote wakeup is detected on the USB.

Note: Accessible in both device and host modes.

## Bit 30 SRQINT: Session request/new session detected interrupt

In host mode, this interrupt is asserted when a session request is detected from the device. In device mode, this interrupt is asserted when  $V_{BUS}$  is in the valid range for a B-peripheral device. Accessible in both device and host modes.

## Bit 29 **DISCINT:** Disconnect detected interrupt

Asserted when a device disconnect is detected.

Note: Only accessible in host mode.

## Bit 28 CIDSCHG: Connector ID status change

The core sets this bit when there is a change in connector ID status.

Note: Accessible in both device and host modes.

## Bit 27 Reserved, must be kept at reset value.

## Bit 26 PTXFE: Periodic TxFIFO empty

Asserted when the periodic transmit FIFO is either half or completely empty and there is space for at least one entry to be written in the periodic request queue. The half or completely empty status is determined by the periodic TxFIFO empty level bit in the OTG\_FS\_GAHBCFG register (PTXFELVL bit in OTG\_FS\_GAHBCFG).

Note: Only accessible in host mode.

#### Bit 25 **HCINT:** Host channels interrupt

The core sets this bit to indicate that an interrupt is pending on one of the channels of the core (in host mode). The application must read the OTG\_FS\_HAINT register to determine the exact number of the channel on which the interrupt occurred, and then read the corresponding OTG\_FS\_HCINTx register to determine the exact cause of the interrupt. The application must clear the appropriate status bit in the OTG\_FS\_HCINTx register to clear this bit.

Note: Only accessible in host mode.

#### Bit 24 HPRTINT: Host port interrupt

The core sets this bit to indicate a change in port status of one of the OTG\_FS controller ports in host mode. The application must read the OTG\_FS\_HPRT register to determine the exact event that caused this interrupt. The application must clear the appropriate status bit in the OTG\_FS\_HPRT register to clear this bit.

Note: Only accessible in host mode.

#### Bits 23:22 Reserved, must be kept at reset value.

#### Bit 21 IPXFR: Incomplete periodic transfer

In host mode, the core sets this interrupt bit when there are incomplete periodic transactions still pending, which are scheduled for the current frame.

**INCOMPISOOUT:** Incomplete isochronous OUT transfer

In device mode, the core sets this interrupt to indicate that there is at least one isochronous OUT endpoint on which the transfer is not completed in the current frame. This interrupt is asserted along with the End of periodic frame interrupt (EOPF) bit in this register.

#### Bit 20 IISOIXFR: Incomplete isochronous IN transfer

The core sets this interrupt to indicate that there is at least one isochronous IN endpoint on which the transfer is not completed in the current frame. This interrupt is asserted along with the End of periodic frame interrupt (EOPF) bit in this register.

Note: Only accessible in device mode.

## Bit 19 **OEPINT:** OUT endpoint interrupt

The core sets this bit to indicate that an interrupt is pending on one of the OUT endpoints of the core (in device mode). The application must read the OTG\_FS\_DAINT register to determine the exact number of the OUT endpoint on which the interrupt occurred, and then read the corresponding OTG\_FS\_DOEPINTx register to determine the exact cause of the interrupt. The application must clear the appropriate status bit in the corresponding OTG\_FS\_DOEPINTx register to clear this bit.

Note: Only accessible in device mode.

## Bit 18 IEPINT: IN endpoint interrupt

The core sets this bit to indicate that an interrupt is pending on one of the IN endpoints of the core (in device mode). The application must read the OTG\_FS\_DAINT register to determine the exact number of the IN endpoint on which the interrupt occurred, and then read the corresponding OTG\_FS\_DIEPINTx register to determine the exact cause of the interrupt. The application must clear the appropriate status bit in the corresponding OTG\_FS\_DIEPINTx register to clear this bit.

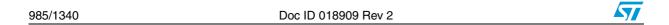
Note: Only accessible in device mode.

# Bits 17:16 Reserved, must be kept at reset value.

## Bit 15 EOPF: End of periodic frame interrupt

Indicates that the period specified in the periodic frame interval field of the OTG\_FS\_DCFG register (PFIVL bit in OTG\_FS\_DCFG) has been reached in the current frame.

Note: Only accessible in device mode.



## Bit 14 ISOODRP: Isochronous OUT packet dropped interrupt

The core sets this bit when it fails to write an isochronous OUT packet into the RxFIFO because the RxFIFO does not have enough space to accommodate a maximum size packet for the isochronous OUT endpoint.

Note: Only accessible in device mode.

#### Bit 13 ENUMDNE: Enumeration done

The core sets this bit to indicate that speed enumeration is complete. The application must read the OTG\_FS\_DSTS register to obtain the enumerated speed.

Note: Only accessible in device mode.

## Bit 12 USBRST: USB reset

The core sets this bit to indicate that a reset is detected on the USB.

Note: Only accessible in device mode.

#### Bit 11 USBSUSP: USB suspend

The core sets this bit to indicate that a suspend was detected on the USB. The core enters the Suspended state when there is no activity on the data lines for a period of 3 ms.

Note: Only accessible in device mode.

#### Bit 10 ESUSP: Early suspend

The core sets this bit to indicate that an Idle state has been detected on the USB for 3 ms.

Note: Only accessible in device mode.

#### Bits 9:8 Reserved, must be kept at reset value.

## Bit 7 GONAKEFF: Global OUT NAK effective

Indicates that the Set global OUT NAK bit in the OTG\_FS\_DCTL register (SGONAK bit in OTG\_FS\_DCTL), set by the application, has taken effect in the core. This bit can be cleared by writing the Clear global OUT NAK bit in the OTG\_FS\_DCTL register (CGONAK bit in OTG\_FS\_DCTL).

Note: Only accessible in device mode.

## Bit 6 GINAKEFF: Global IN non-periodic NAK effective

Indicates that the Set global non-periodic IN NAK bit in the OTG\_FS\_DCTL register (SGINAK bit in OTG\_FS\_DCTL), set by the application, has taken effect in the core. That is, the core has sampled the Global IN NAK bit set by the application. This bit can be cleared by clearing the Clear global non-periodic IN NAK bit in the OTG\_FS\_DCTL register (CGINAK bit in OTG\_FS\_DCTL).

This interrupt does not necessarily mean that a NAK handshake is sent out on the USB. The STALL bit takes precedence over the NAK bit.

Note: Only accessible in device mode.

## Bit 5 NPTXFE: Non-periodic TxFIFO empty

This interrupt is asserted when the non-periodic TxFIFO is either half or completely empty, and there is space for at least one entry to be written to the non-periodic transmit request queue. The half or completely empty status is determined by the non-periodic TxFIFO empty level bit in the OTG\_FS\_GAHBCFG register (TXFELVL bit in OTG\_FS\_GAHBCFG).

Note: Accessible in host mode only.

## Bit 4 RXFLVL: RxFIFO non-empty

Indicates that there is at least one packet pending to be read from the RxFIFO.

Note: Accessible in both host and device modes.

#### Bit 3 SOF: Start of frame

In host mode, the core sets this bit to indicate that an SOF (FS), or Keep-Alive (LS) is transmitted on the USB. The application must write a 1 to this bit to clear the interrupt. In device mode, in the core sets this bit to indicate that an SOF token has been received on the USB. The application can read the Device Status register to get the current frame number. This interrupt is seen only when the core is operating in FS.

Note: Accessible in both host and device modes.

## Bit 2 OTGINT: OTG interrupt

The core sets this bit to indicate an OTG protocol event. The application must read the OTG Interrupt Status (OTG\_FS\_GOTGINT) register to determine the exact event that caused this interrupt. The application must clear the appropriate status bit in the OTG\_FS\_GOTGINT register to clear this bit.

Note: Accessible in both host and device modes.

#### Bit 1 MMIS: Mode mismatch interrupt

The core sets this bit when the application is trying to access:

- A host mode register, when the core is operating in device mode
- A device mode register, when the core is operating in host mode

The register access is completed on the AHB with an OKAY response, but is ignored by the core internally and does not affect the operation of the core.

Note: Accessible in both host and device modes.

## Bit 0 CMOD: Current mode of operation

Indicates the current mode.

0: Device mode

1: Host mode

Note: Accessible in both host and device modes.

# OTG\_FS interrupt mask register (OTG\_FS\_GINTMSK)

Address offset: 0x018
Reset value: 0x0000 0000

This register works with the Core interrupt register to interrupt the application. When an interrupt bit is masked, the interrupt associated with that bit is not generated. However, the Core Interrupt (OTG\_FS\_GINTSTS) register bit corresponding to that interrupt is still set.

31	30	29	28	27	26	25	24	23 22	21	20	19	18	17	16	15	14	13	12	11	10	9 8	/	6	5	4	3	2	1	Ü
WUIM	SRQIM	DISCINT	CIDSCHGM	Reserved	PTXFEM	WIOH	PRTIM	Reserved	IPXFRM/IISOOXFRM	IISOIXFRM	OEPINT	IEPINT	EPMISM	Reserved	EOPFM	ISOODRPM	ENUMDNEM	USBRST	USBSUSPM	ESUSPM	Reserved	GONAKEFFM	GINAKEFFM	NPTXFEM	RXFLVLM	SOFM	OTGINT	MMISM	Reserved
rw	rw	rw	rw		rw	rw	r		rw	rw	rw	rw	rw		rw	rw	rw	rw	rw	rw		rw	rw	rw	rw	rw	rw	rw	

Bit 31 WUIM: Resume/remote wakeup detected interrupt mask

0: Masked interrupt1: Unmasked interrupt

Note: Accessible in both host and device modes.

Bit 30 SRQIM: Session request/new session detected interrupt mask

0: Masked interrupt1: Unmasked interrupt

Note: Accessible in both host and device modes.

Bit 29 **DISCINT:** Disconnect detected interrupt mask

0: Masked interrupt1: Unmasked interrupt

Note: Only accessible in device mode.

Bit 28 CIDSCHGM: Connector ID status change mask

0: Masked interrupt1: Unmasked interrupt

Note: Accessible in both host and device modes.

Bit 27 Reserved, must be kept at reset value.

Bit 26 PTXFEM: Periodic TxFIFO empty mask

0: Masked interrupt1: Unmasked interrupt

Note: Only accessible in host mode.

Bit 25 HCIM: Host channels interrupt mask

0: Masked interrupt1: Unmasked interrupt

Note: Only accessible in host mode.

Bit 24 PRTIM: Host port interrupt mask

0: Masked interrupt1: Unmasked interrupt

Note: Only accessible in host mode.

Bits 23:22 Reserved, must be kept at reset value.

Bit 21 IPXFRM: Incomplete periodic transfer mask

0: Masked interrupt1: Unmasked interrupt

Note: Only accessible in host mode.

IISOOXFRM: Incomplete isochronous OUT transfer mask

0: Masked interrupt1: Unmasked interrupt

Note: Only accessible in device mode.

Bit 20 IISOIXFRM: Incomplete isochronous IN transfer mask

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in device mode.

Bit 19 **OEPINT:** OUT endpoints interrupt mask

0: Masked interrupt1: Unmasked interrupt

Note: Only accessible in device mode.

Bit 18 IEPINT: IN endpoints interrupt mask

0: Masked interrupt1: Unmasked interrupt

Note: Only accessible in device mode.

Bit 17 EPMISM: Endpoint mismatch interrupt mask

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in device mode.

Bit 16 Reserved, must be kept at reset value.

Bit 15 **EOPFM:** End of periodic frame interrupt mask

0: Masked interrupt1: Unmasked interrupt

Note: Only accessible in device mode.

Bit 14 ISOODRPM: Isochronous OUT packet dropped interrupt mask

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in device mode.

Bit 13 ENUMDNEM: Enumeration done mask

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in device mode.

Bit 12 USBRST: USB reset mask

989/1340

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in device mode.

## Bit 11 USBSUSPM: USB suspend mask

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in device mode.

## Bit 10 ESUSPM: Early suspend mask

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in device mode.

## Bits 9:8 Reserved, must be kept at reset value.

## Bit 7 GONAKEFFM: Global OUT NAK effective mask

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in device mode.

#### Bit 6 GINAKEFFM: Global non-periodic IN NAK effective mask

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in device mode.

## Bit 5 NPTXFEM: Non-periodic TxFIFO empty mask

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in Host mode.

## Bit 4 RXFLVLM: Receive FIFO non-empty mask

0: Masked interrupt

1: Unmasked interrupt

Note: Accessible in both device and host modes.

#### Bit 3 SOFM: Start of frame mask

0: Masked interrupt

1: Unmasked interrupt

Note: Accessible in both device and host modes.

## Bit 2 OTGINT: OTG interrupt mask

0: Masked interrupt

1: Unmasked interrupt

Note: Accessible in both device and host modes.

## Bit 1 MMISM: Mode mismatch interrupt mask

0: Masked interrupt

1: Unmasked interrupt

Note: Accessible in both device and host modes.

Bit 0 Reserved, must be kept at reset value.

# OTG\_FS Receive status debug read/OTG status read and pop registers (OTG\_FS\_GRXSTSR/OTG\_FS\_GRXSTSP)

Address offset for Read: 0x01C Address offset for Pop: 0x020

Reset value: 0x0000 0000

A read to the Receive status debug read register returns the contents of the top of the Receive FIFO. A read to the Receive status read and pop register additionally pops the top data entry out of the RxFIFO.

The receive status contents must be interpreted differently in host and device modes. The core ignores the receive status pop/read when the receive FIFO is empty and returns a value of 0x0000 0000. The application must only pop the Receive Status FIFO when the Receive FIFO non-empty bit of the Core interrupt register (RXFLVL bit in OTG\_FS\_GINTSTS) is asserted.

## Host mode:

 $31 \ \ 30 \ \ 29 \ \ 28 \ \ 27 \ \ 26 \ \ 25 \ \ 24 \ \ 23 \ \ 22 \ \ 21 \ \ 20 \ \ 19 \ \ 18 \ \ 17 \ \ 16 \ \ 15 \ \ 14 \ \ 13 \ \ 12 \ \ 11 \ \ 10 \ \ 9 \ \ 8 \ \ 7 \ \ 6 \ \ 5 \ \ 4 \ \ 3 \ \ 2 \ \ 1 \ \ 0$ 

Reserved	PKTSTS	DPID	BCNT	CHNUM
i lesei veu	r	r	r	r

Bits 31:21 Reserved, must be kept at reset value.

Bits 20:17 PKTSTS: Packet status

Indicates the status of the received packet

0010: IN data packet received

0011: IN transfer completed (triggers an interrupt)0101: Data toggle error (triggers an interrupt)0111: Channel halted (triggers an interrupt)

Others: Reserved

Bits 16:15 DPID: Data PID

Indicates the Data PID of the received packet

00: DATA0 10: DATA1 01: DATA2 11: MDATA

Bits 14:4 BCNT: Byte count

Indicates the byte count of the received IN data packet.

Bits 3:0 CHNUM: Channel number

Indicates the channel number to which the current received packet belongs.

## **Device mode:**

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	FRMNUM	PKTSTS	DPID	BCNT	EPNUM
rieserveu	r	ŗ	r	r	r

Bits 31:25 Reserved, must be kept at reset value.

Bits 24:21 FRMNUM: Frame number

This is the least significant 4 bits of the frame number in which the packet is received on the USB. This field is supported only when isochronous OUT endpoints are supported.

Bits 20:17 PKTSTS: Packet status

Indicates the status of the received packet 0001: Global OUT NAK (triggers an interrupt)

0010: OUT data packet received

0011: OUT transfer completed (triggers an interrupt)0100: SETUP transaction completed (triggers an interrupt)

0110: SETUP data packet received

Others: Reserved

Bits 16:15 DPID: Data PID

Indicates the Data PID of the received OUT data packet

00: DATA0 10: DATA1 01: DATA2 11: MDATA

Bits 14:4 BCNT: Byte count

Indicates the byte count of the received data packet.

Bits 3:0 **EPNUM:** Endpoint number

Indicates the endpoint number to which the current received packet belongs.

## OTG\_FS Receive FIFO size register (OTG\_FS\_GRXFSIZ)

Address offset: 0x024

Reset value: 0x0000 0200

The application can program the RAM size that must be allocated to the RxFIFO.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	RXFD
i iesei veu	r/rw

Bits 31:16 Reserved, must be kept at reset value.

Bits 15:0 RXFD: RxFIFO depth

This value is in terms of 32-bit words.

Minimum value is 16 Maximum value is 256

The power-on reset value of this register is specified as the largest Rx data FIFO depth.

# OTG\_FS Host non-periodic transmit FIFO size register (OTG\_FS\_HNPTXFSIZ)/Endpoint 0 Transmit FIFO size (OTG\_FS\_DIEPTXF0)

Address offset: 0x028

Reset value: 0x0000 0200

#### 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

NPTXFD/TX0FD	NPTXFSA/TX0FSA
r/rw	r/rw

#### Host mode

Bits 31:16 NPTXFD: Non-periodic TxFIFO depth

This value is in terms of 32-bit words.

Minimum value is 16 Maximum value is 256

Bits 15:0 NPTXFSA: Non-periodic transmit RAM start address

This field contains the memory start address for non-periodic transmit FIFO RAM.

#### **Device mode**

Bits 31:16 TX0FD: Endpoint 0 TxFIFO depth

This value is in terms of 32-bit words.

Minimum value is 16 Maximum value is 256

Bits 15:0 TX0FSA: Endpoint 0 transmit RAM start address

This field contains the memory start address for the endpoint 0 transmit FIFO RAM.

# OTG\_FS non-periodic transmit FIFO/queue status register (OTG\_FS\_HNPTXSTS)

Address offset: 0x02C

Reset value: 0x0008 0200

Note: In Device mode, this register is not valid.

This read-only register contains the free space information for the non-periodic TxFIFO and

the non-periodic transmit request queue.

# $31 \ \ 30 \ \ 29 \ \ 28 \ \ 27 \ \ 26 \ \ 25 \ \ 24 \ \ 23 \ \ 22 \ \ 21 \ \ 20 \ \ 19 \ \ 18 \ \ 17 \ \ 16 \ \ 15 \ \ 14 \ \ 13 \ \ 12 \ \ 11 \ \ 10 \ \ 9 \ \ 8 \ \ 7 \ \ 6 \ \ 5 \ \ 4 \ \ 3 \ \ 2 \ \ 1 \ \ 0$

ved	NPTXQTOP	NPTQXSAV	NPTXFSAV
Reser	r	r	r

Bit 31 Reserved, must be kept at reset value.

Bits 30:24 NPTXQTOP: Top of the non-periodic transmit request queue

Entry in the non-periodic Tx request queue that is currently being processed by the MAC.

Bits 30:27: Channel/endpoint number

Bits 26:25:

00: IN/OUT token

01: Zero-length transmit packet (device IN/host OUT)

11: Channel halt command

Bit 24: Terminate (last entry for selected channel/endpoint)

## Bits 23:16 NPTQXSAV: Non-periodic transmit request queue space available

Indicates the amount of free space available in the non-periodic transmit request queue. This queue holds both IN and OUT requests in host mode. Device mode has only IN

00: Non-periodic transmit request queue is full

01: 1 location available 10: 2 locations available

bxn: n locations available  $(0 \le n \le 8)$ 

Others: Reserved

## Bits 15:0 NPTXFSAV: Non-periodic TxFIFO space available

Indicates the amount of free space available in the non-periodic TxFIFO.

Values are in terms of 32-bit words. 00: Non-periodic TxFIFO is full

01: 1 word available 10: 2 words available

0xn: n words available (where  $0 \le n \le 256$ )

Others: Reserved

# OTG\_FS general core configuration register (OTG\_FS\_GCCFG)

Address offset: 0x038

Reset value: 0x0000 0000

3	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											SNE	Z	Z	z		Z																

Reserved	NOVBUSSENS	SOFOUTEN	VBUSBSEN	VBUSASEN	Reserved	.PWRDWN	Reserved
	rw	rw	rw	rw		rw	

Bits 31:22 Reserved, must be kept at reset value.

## Bit 21 NOVBUSSENS: V<sub>BUS</sub> sensing disable option

When this bit is set,  $V_{BUS}$  is considered internally to be always at  $V_{BUS}$  valid level (5 V). This option removes the need for a dedicated  $V_{BUS}$  pad, and leave this pad free to be used for other purposes such as a shared functionality.  $V_{BUS}$  connection can be remapped on another general purpose input pad and monitored by software.

This option is only suitable for host-only or device-only applications.

0: V<sub>BUS</sub> sensing available by hardware

1: V<sub>BUS</sub> sensing not available by hardware.

#### Bit 20 SOFOUTEN: SOF output enable

0: SOF pulse not available on PAD

1: SOF pulse available on PAD

# Bit 19 $\,$ VBUSBSEN: Enable the $\,$ VBUS sensing "B" device

0: V<sub>BUS</sub> sensing "B" disabled

1: V<sub>BUS</sub> sensing "B" enabled

## Bit 18 VBUSASEN: Enable the V<sub>BUS</sub> sensing "A" device

0: V<sub>BUS</sub> sensing "A" disabled

1: V<sub>BUS</sub> sensing "A" enabled

Bit 17 Reserved, must be kept at reset value.

#### Bit 16 PWRDWN: Power down

Used to activate the transceiver in transmission/reception

0: Power down active

1: Power down deactivated ("Transceiver active")

Bits 15:0 Reserved, must be kept at reset value.

## OTG\_FS core ID register (OTG\_FS\_CID)

Address offset: 0x03C

Reset value:0x0000 1100

This is a read only register containing the Product ID.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

														PF	RODU	JCT_	_ID														
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw															

## Bits 31:0 PRODUCT\_ID: Product ID field

Application-programmable ID field.

# OTG\_FS Host periodic transmit FIFO size register (OTG\_FS\_HPTXFSIZ)

Address offset: 0x100

Reset value: 0x0200 0600

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

							PTX	FSIZ															PT>	(SA							
r/r												r/r																			
w	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W

Bits 31:16 PTXFD: Host periodic TxFIFO depth

This value is in terms of 32-bit words.

Minimum value is 16

Bits 15:0 PTXSA: Host periodic TxFIFO start address

The power-on reset value of this register is the sum of the largest Rx data FIFO depth and largest non-periodic Tx data FIFO depth.

# OTG\_FS device IN endpoint transmit FIFO size register (OTG\_FS\_DIEPTXFx) $(x = 1..3, where x is the FIFO_number)$

Address offset:  $0x104 + (FIFO_number - 1) \times 0x04$ 

Reset value: 0x02000400

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

	INEPTXFD 																					П	NEP.	TXS	4						
r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r	r/r
w	w	w	w	w	w	w	w	w	W	w	w	W	w	w	w	w	w	w	w	w	w	w	w	W	w	w	w	w	w	w	w

## Bits 31:16 INEPTXFD: IN endpoint TxFIFO depth

This value is in terms of 32-bit words.

Minimum value is 16

The power-on reset value of this register is specified as the largest IN endpoint FIFO number depth.

## Bits 15:0 INEPTXSA: IN endpoint FIFOx transmit RAM start address

This field contains the memory start address for IN endpoint transmit FIFOx. The address must be aligned with a 32-bit memory location.

# 29.16.3 Host-mode registers

Bit values in the register descriptions are expressed in binary unless otherwise specified.

Host-mode registers affect the operation of the core in the host mode. Host mode registers must not be accessed in device mode, as the results are undefined. Host mode registers can be categorized as follows:

## OTG\_FS Host configuration register (OTG\_FS\_HCFG)

Address offset: 0x400 Reset value: 0x0000 0000

This register configures the core after power-on. Do not make changes to this register after initializing the host.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
														Reserved															FSLSS		
														ш.															r	rw	rw

Bits 31:3 Reserved, must be kept at reset value.

## Bit 2 FSLSS: FS- and LS-only support

The application uses this bit to control the core's enumeration speed. Using this bit, the application can make the core enumerate as an FS host, even if the connected device supports HS traffic. Do not make changes to this field after initial programming.

1: FS/LS-only, even if the connected device can support HS (read-only)

Bits 1:0 FSLSPCS: FS/LS PHY clock select

When the core is in FS host mode 01: PHY clock is running at 48 MHz

Others: Reserved

When the core is in LS host mode

00: Reserved

01: Select 48 MHz PHY clock frequency 10: Select 6 MHz PHY clock frequency

11: Reserved

Note: The FSLSPCS must be set on a connection event according to the speed of the connected device (after changing this bit, a software reset must be performed).

## OTG\_FS Host frame interval register (OTG\_FS\_HFIR)

Address offset: 0x404

Reset value: 0x0000 EA60

This register stores the frame interval information for the current speed to which the OTG\_FS controller has enumerated.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							D																FR	IVL							
							Hese	erved	1							rw	rw	rw	rw	rw	rw	rw	rw								

Bits 31:16 Reserved, must be kept at reset value.

## Bits 15:0 FRIVL: Frame interval

The value that the application programs to this field specifies the interval between two consecutive SOFs (FS) or Keep-Alive tokens (LS). This field contains the number of PHY clocks that constitute the required frame interval. The application can write a value to this register only after the Port enable bit of the host port control and status register (PENA bit in OTG\_FS\_HPRT) has been set. If no value is programmed, the core calculates the value based on the PHY clock specified in the FS/LS PHY Clock Select field of the host configuration register (FSLSPCS in OTG\_FS\_HCFG). Do not change the value of this field after the initial configuration.

1 ms × (PHY clock frequency)

# OTG FS Host frame number/frame time remaining register (OTG FS HFNUM)

Address offset: 0x408

Reset value: 0x0000 3FFF

This register indicates the current frame number. It also indicates the time remaining (in terms of the number of PHY clocks) in the current frame.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							FTE	SEM															FRI	II IM							

							FTF	REM															FRN	MUI							
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

## Bits 31:16 FTREM: Frame time remaining

Indicates the amount of time remaining in the current frame, in terms of PHY clocks. This field decrements on each PHY clock. When it reaches zero, this field is reloaded with the value in the Frame interval register and a new SOF is transmitted on the USB.

#### Bits 15:0 FRNUM: Frame number

This field increments when a new SOF is transmitted on the USB, and is cleared to 0 when it reaches 0x3FFF.

# OTG\_FS\_Host periodic transmit FIFO/queue status register (OTG\_FS\_HPTXSTS)

Address offset: 0x410

Reset value: 0x0008 0100

This read-only register contains the free space information for the periodic TxFIFO and the periodic transmit request queue.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

		-	PTXC	TOF	•					F	PTXC	QSA\	/									F	TXF	SAV	L						
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	rw	rw	rw	rw	rw	rw	rw	rw	rw							

#### Bits 31:24 PTXQTOP: Top of the periodic transmit request queue

This indicates the entry in the periodic Tx request queue that is currently being processed by the MAC.

This register is used for debugging.

Bit 31: Odd/Even frame

0: send in even frame

1: send in odd frame

Bits 30:27: Channel/endpoint number

Bits 26:25: Type

00: IN/OUT

01: Zero-length packet

11: Disable channel command

Bit 24: Terminate (last entry for the selected channel/endpoint)

## Bits 23:16 PTXQSAV: Periodic transmit request queue space available

Indicates the number of free locations available to be written in the periodic transmit request queue. This queue holds both IN and OUT requests.

00: Periodic transmit request queue is full

01: 1 location available10: 2 locations available

bxn: n locations available  $(0 \le n \le 8)$ 

Others: Reserved

## Bits 15:0 PTXFSAVL: Periodic transmit data FIFO space available

Indicates the number of free locations available to be written to in the periodic TxFIFO.

Values are in terms of 32-bit words 0000: Periodic TxFIFO is full

0001: 1 word available 0010: 2 words available

bxn: n words available (where  $0 \le n \le PTXFD$ )

Others: Reserved

## OTG\_FS Host all channels interrupt register (OTG\_FS\_HAINT)

Address offset: 0x414
Reset value: 0x0000 000

When a significant event occurs on a channel, the host all channels interrupt register interrupts the application using the host channels interrupt bit of the Core interrupt register (HCINT bit in OTG\_FS\_GINTSTS). This is shown in *Figure 358*. There is one interrupt bit per channel, up to a maximum of 16 bits. Bits in this register are set and cleared when the application sets and clears bits in the corresponding host channel-x interrupt register.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved								HA	INT							
neserveu	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:16 Reserved, must be kept at reset value.

Bits 15:0 HAINT: Channel interrupts

One bit per channel: Bit 0 for Channel 0, bit 15 for Channel 15

# OTG\_FS Host all channels interrupt mask register (OTG\_FS\_HAINTMSK)

Address offset: 0x418 Reset value: 0x0000 0000

The host all channel interrupt mask register works with the host all channel interrupt register to interrupt the application when an event occurs on a channel. There is one interrupt mask bit per channel, up to a maximum of 16 bits.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Penanyad								HAII	MTV							
Reserved	rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bits 31:16 Reserved, must be kept at reset value.

Bits 15:0 HAINTM: Channel interrupt mask

0: Masked interrupt1: Unmasked interrupt

One bit per channel: Bit 0 for channel 0, bit 15 for channel 15

## OTG\_FS Host port control and status register (OTG\_FS\_HPRT)

Address offset: 0x440
Reset value: 0x0000 0000

This register is available only in host mode. Currently, the OTG host supports only one port.

A single register holds USB port-related information such as USB reset, enable, suspend, resume, connect status, and test mode for each port. It is shown in *Figure 358*. The rc\_w1 bits in this register can trigger an interrupt to the application through the host port interrupt bit of the core interrupt register (HPRTINT bit in OTG\_FS\_GINTSTS). On a Port Interrupt, the application must read this register and clear the bit that caused the interrupt. For the rc\_w1 bits, the application must write a 1 to the bit to clear the interrupt.

31 30 29 28 27 26 25 24 23 22 21 20	19 18 17	' 16 15 14 13	12 11 10	9 8	7 6	5 4	3 2 1 0
Reserved	PSPD	PTCTL	유 교	eserved	PSUSP	POCCHNG	PENCHNG PENA PCDET PCSTS
	r r	rw rw rw rw	rw r r	rw rw	rs rw	rc_ w1 r	rc_ rc_ rc_ rc_ r

Bits 31:19 Reserved, must be kept at reset value.

Bits 18:17 PSPD: Port speed

Indicates the speed of the device attached to this port.

01: Full speed10: Low speed11: Reserved

#### Bits 16:13 PTCTL: Port test control

The application writes a nonzero value to this field to put the port into a Test mode, and the corresponding pattern is signaled on the port.

0000: Test mode disabled 0001: Test\_J mode 0010: Test\_K mode

0011: Test\_SE0\_NAK mode 0100: Test\_Packet mode 0101: Test\_Force\_Enable Others: Reserved

## Bit 12 PPWR: Port power

The application uses this field to control power to this port, and the core clears this bit on an overcurrent condition.

0: Power off 1: Power on

## Bits 11:10 PLSTS: Port line status

Indicates the current logic level USB data lines

Bit 10: Logic level of OTG\_FS\_FS\_DP Bit 11: Logic level of OTG\_FS\_FS\_DM

Bit 9 Reserved, must be kept at reset value.

#### Bit 8 PRST: Port reset

When the application sets this bit, a reset sequence is started on this port. The application must time the reset period and clear this bit after the reset sequence is complete.

0: Port not in reset

1: Port in reset

The application must leave this bit set for a minimum duration of at least 10 ms to start a reset on the port. The application can leave it set for another 10 ms in addition to the required minimum duration, before clearing the bit, even though there is no maximum limit set by the USB standard.

## Bit 7 PSUSP: Port suspend

The application sets this bit to put this port in Suspend mode. The core only stops sending SOFs when this is set. To stop the PHY clock, the application must set the Port clock stop bit, which asserts the suspend input pin of the PHY.

The read value of this bit reflects the current suspend status of the port. This bit is cleared by the core after a remote wakeup signal is detected or the application sets the Port reset bit or Port resume bit in this register or the Resume/remote wakeup detected interrupt bit or Disconnect detected interrupt bit in the Core interrupt register (WKUINT or DISCINT in OTG\_FS\_GINTSTS, respectively).

0: Port not in Suspend mode

1: Port in Suspend mode

## Bit 6 PRES: Port resume

The application sets this bit to drive resume signaling on the port. The core continues to drive the resume signal until the application clears this bit.

If the core detects a USB remote wakeup sequence, as indicated by the Port resume/remote wakeup detected interrupt bit of the Core interrupt register (WKUINT bit in OTG\_FS\_GINTSTS), the core starts driving resume signaling without application intervention and clears this bit when it detects a disconnect condition. The read value of this bit indicates whether the core is currently driving resume signaling.

0: No resume driven

1: Resume driven

## Bit 5 POCCHNG: Port overcurrent change

The core sets this bit when the status of the Port overcurrent active bit (bit 4) in this register changes.

#### Bit 4 POCA: Port overcurrent active

Indicates the overcurrent condition of the port.

0: No overcurrent condition

1: Overcurrent condition

#### Bit 3 PENCHNG: Port enable/disable change

The core sets this bit when the status of the Port enable bit 2 in this register changes.

#### Bit 2 PENA: Port enable

A port is enabled only by the core after a reset sequence, and is disabled by an overcurrent condition, a disconnect condition, or by the application clearing this bit. The application cannot set this bit by a register write. It can only clear it to disable the port. This bit does not trigger any interrupt to the application.

0: Port disabled

1: Port enabled

# Bit 1 PCDET: Port connect detected

The core sets this bit when a device connection is detected to trigger an interrupt to the application using the host port interrupt bit in the Core interrupt register (HPRTINT bit in OTG\_FS\_GINTSTS). The application must write a 1 to this bit to clear the interrupt.

## Bit 0 PCSTS: Port connect status

0: No device is attached to the port

1: A device is attached to the port

# OTG\_FS Host channel-x characteristics register (OTG\_FS\_HCCHARx) (x = 0..7, where x = Channel\_number)

Address offset:  $0x500 + (Channel number \times 0x20)$ 

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
CHENA	CHDIS	ODDFRM				DAD	ı			МС	NT	EDTVD	_	LSDEV	Reserved	EPDIR		EPN	NUM						N	/IPSI	Z				
rs	rs	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	ш	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

#### Bit 31 CHENA: Channel enable

This field is set by the application and cleared by the OTG host.

0: Channel disabled

1: Channel enabled

## Bit 30 CHDIS: Channel disable

The application sets this bit to stop transmitting/receiving data on a channel, even before the transfer for that channel is complete. The application must wait for the Channel disabled interrupt before treating the channel as disabled.

## Bit 29 ODDFRM: Odd frame

This field is set (reset) by the application to indicate that the OTG host must perform a transfer in an odd frame. This field is applicable for only periodic (isochronous and interrupt) transactions.

0: Even frame

1: Odd frame

## Bits 28:22 DAD: Device address

This field selects the specific device serving as the data source or sink.

## Bits 21:20 MCNT: Multicount

This field indicates to the host the number of transactions that must be executed per frame for this periodic endpoint. For non-periodic transfers, this field is not used

00: Reserved. This field yields undefined results

01: 1 transaction

10: 2 transactions per frame to be issued for this endpoint

11: 3 transactions per frame to be issued for this endpoint

Note: This field must be set to at least 01.

## Bits 19:18 EPTYP: Endpoint type

Indicates the transfer type selected.

00: Control

01: Isochronous

10: Bulk

11: Interrupt

## Bit 17 LSDEV: Low-speed device

This field is set by the application to indicate that this channel is communicating to a low-speed device.

Bit 16 Reserved, must be kept at reset value.

Bit 15 EPDIR: Endpoint direction

Indicates whether the transaction is IN or OUT.

0: OUT 1: IN

Bits 14:11 EPNUM: Endpoint number

Indicates the endpoint number on the device serving as the data source or sink.

Bits 10:0 MPSIZ: Maximum packet size

Indicates the maximum packet size of the associated endpoint.

# OTG\_FS Host channel-x interrupt register (OTG\_FS\_HCINTx) (x = 0..7, where $x = Channel\_number$ )

Address offset: 0x508 + (Channel\_number × 0x20)

Reset value: 0x0000 0000

This register indicates the status of a channel with respect to USB- and AHB-related events. It is shown in *Figure 358*. The application must read this register when the host channels interrupt bit in the Core interrupt register (HCINT bit in OTG\_FS\_GINTSTS) is set. Before the application can read this register, it must first read the host all channels interrupt (OTG\_FS\_HAINT) register to get the exact channel number for the host channel-x interrupt register. The application must clear the appropriate bit in this register to clear the corresponding bits in the OTG\_FS\_HAINT and OTG\_FS\_GINTSTS registers.

Bits 31:11 Reserved, must be kept at reset value.

Bit 10 **DTERR:** Data toggle error Bit 9 **FRMOR:** Frame overrun

Bit 8 **BBERR**: Babble error

Bit 7 TXERR: Transaction error

Indicates one of the following errors occurred on the USB.

CRC check failure

Timeout Bit stuff error False EOP

Bit 6 Reserved, must be kept at reset value.

Bit 5 ACK: ACK response received/transmitted interrupt

Bit 4 NAK: NAK response received interrupt

Bit 3 STALL: STALL response received interrupt

Bit 2 Reserved, must be kept at reset value.

Bit 1 CHH: Channel halted

Indicates the transfer completed abnormally either because of any USB transaction error or in response to disable request by the application.

Bit 0 XFRC: Transfer completed

Transfer completed normally without any errors.

# OTG\_FS Host channel-x interrupt mask register (OTG\_FS\_HCINTMSKx) $(x = 0..7, where x = Channel_number)$

Address offset: 0x50C + (Channel\_number × 0x20)

Reset value: 0x0000 0000

This register reflects the mask for each channel status described in the previous section.

;	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
										Re	eserv	/ed										DTERRM	FRMORM	BBERRM	TXERRM	NYET	ACKM	NAKN	STALLM	eserved	СННМ	XFRCM
																						rw	rw	rw	rw	rw	rw	rw	rw	ш	rw	rw

Bits 31:11 Reserved, must be kept at reset value.

Bit 10 DTERRM: Data toggle error mask

0: Masked interrupt1: Unmasked interrupt

Bit 9 FRMORM: Frame overrun mask

0: Masked interrupt1: Unmasked interrupt

Bit 8 BBERRM: Babble error mask

0: Masked interrupt1: Unmasked interrupt

Bit 7 TXERRM: Transaction error mask

0: Masked interrupt1: Unmasked interrupt

Bit 6 NYET: response received interrupt mask

0: Masked interrupt1: Unmasked interrupt

Bit 5 ACKM: ACK response received/transmitted interrupt mask

0: Masked interrupt1: Unmasked interrupt

Bit 4 NAKM: NAK response received interrupt mask

0: Masked interrupt1: Unmasked interrupt

Bit 3 STALLM: STALL response received interrupt mask

0: Masked interrupt1: Unmasked interrupt

Bit 2 Reserved, must be kept at reset value.

Bit 1 CHHM: Channel halted mask

0: Masked interrupt1: Unmasked interrupt

Bit 0 XFRCM: Transfer completed mask

0: Masked interrupt1: Unmasked interrupt

# OTG\_FS Host channel-x transfer size register (OTG\_FS\_HCTSIZx) (x = 0...7, where x = Channel number)

Address offset: 0x510 + (Channel\_number × 0x20)

Reset value: 0x0000 0000

#### 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	DF	Ole					PKT	CNT													Х	FRS	ΙZ								
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bit 31 Reserved, must be kept at reset value.

Bits 30:29 DPID: Data PID

The application programs this field with the type of PID to use for the initial transaction. The host maintains this field for the rest of the transfer.

00: DATA0 01: DATA2 10: DATA1

11: MDATA (non-control)/SETUP (control)

## Bits 28:19 PKTCNT: Packet count

This field is programmed by the application with the expected number of packets to be transmitted (OUT) or received (IN).

The host decrements this count on every successful transmission or reception of an OUT/IN packet. Once this count reaches zero, the application is interrupted to indicate normal completion.

#### Bits 18:0 XFRSIZ: Transfer size

For an OUT, this field is the number of data bytes the host sends during the transfer. For an IN, this field is the buffer size that the application has reserved for the transfer. The application is expected to program this field as an integer multiple of the maximum packet size for IN transactions (periodic and non-periodic).

# 29.16.4 Device-mode registers

## OTG\_FS device configuration register (OTG\_FS\_DCFG)

Address offset: 0x800

Reset value: 0x0220 0000

This register configures the core in device mode after power-on or after certain control commands or enumeration. Do not make changes to this register after initial programming.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved

Bits 31:13 Reserved, must be kept at reset value.

#### Bits 12:11 PFIVL: Periodic frame interval

Indicates the time within a frame at which the application must be notified using the end of periodic frame interrupt. This can be used to determine if all the isochronous traffic for that frame is complete.

00: 80% of the frame interval 01: 85% of the frame interval 10: 90% of the frame interval 11: 95% of the frame interval

#### Bits 10:4 DAD: Device address

The application must program this field after every SetAddress control command.

Bit 3 Reserved, must be kept at reset value.

## Bit 2 NZLSOHSK: Non-zero-length status OUT handshake

The application can use this field to select the handshake the core sends on receiving a nonzero-length data packet during the OUT transaction of a control transfer's Status stage.

- 1: Send a STALL handshake on a nonzero-length status OUT transaction and do not send the received OUT packet to the application.
- 0: Send the received OUT packet to the application (zero-length or nonzero-length) and send a handshake based on the NAK and STALL bits for the endpoint in the Device endpoint control register.

#### Bits 1:0 DSPD: Device speed

Indicates the speed at which the application requires the core to enumerate, or the maximum speed the application can support. However, the actual bus speed is determined only after the chirp sequence is completed, and is based on the speed of the USB host to which the core is connected.

00: Reserved01: Reserved10: Reserved

11: Full speed (USB 1.1 transceiver clock is 48 MHz)

# OTG\_FS device control register (OTG\_FS\_DCTL)

Address offset: 0x804

Reset value: 0x0000 0000

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	POPRGDNE	CGONAK	SGONAK	CGINAK	SGINAK		TCTL		GONSTS	GINSTS	SDIS	RWUSIG
	rw	٧	W	w	w	rw	rw	rw	r	r	rw	rw

Bits 31:12 Reserved, must be kept at reset value.

## Bit 11 POPRGDNE: Power-on programming done

The application uses this bit to indicate that register programming is completed after a wakeup from power down mode.

#### Bit 10 CGONAK: Clear global OUT NAK

A write to this field clears the Global OUT NAK.

## Bit 9 SGONAK: Set global OUT NAK

A write to this field sets the Global OUT NAK.

The application uses this bit to send a NAK handshake on all OUT endpoints. The application must set the this bit only after making sure that the Global OUT NAK effective bit in the Core interrupt register (GONAKEFF bit in OTG\_FS\_GINTSTS) is cleared.

#### Bit 8 CGINAK: Clear global IN NAK

A write to this field clears the Global IN NAK.

# Bit 7 SGINAK: Set global IN NAK

A write to this field sets the Global non-periodic IN NAK. The application uses this bit to send a NAK handshake on all non-periodic IN endpoints.

The application must set this bit only after making sure that the Global IN NAK effective bit in the Core interrupt register (GINAKEFF bit in OTG\_FS\_GINTSTS) is cleared.

# Bits 6:4 TCTL: Test control

000: Test mode disabled

001: Test\_J mode

010: Test\_K mode

011: Test\_SE0\_NAK mode

100: Test\_Packet mode

101: Test\_Force\_Enable

Others: Reserved

### Bit 3 GONSTS: Global OUT NAK status

0: A handshake is sent based on the FIFO Status and the NAK and STALL bit settings.

1: No data is written to the RxFIFO, irrespective of space availability. Sends a NAK handshake on all packets, except on SETUP transactions. All isochronous OUT packets are dropped.

#### Bit 2 GINSTS: Global IN NAK status

0: A handshake is sent out based on the data availability in the transmit FIFO.

1: A NAK handshake is sent out on all non-periodic IN endpoints, irrespective of the data availability in the transmit FIFO.

## Bit 1 SDIS: Soft disconnect

The application uses this bit to signal the USB OTG core to perform a soft disconnect. As long as this bit is set, the host does not see that the device is connected, and the device does not receive signals on the USB. The core stays in the disconnected state until the application clears this bit.

0: Normal operation. When this bit is cleared after a soft disconnect, the core generates a device connect event to the USB host. When the device is reconnected, the USB host restarts device enumeration.

1: The core generates a device disconnect event to the USB host.

# Bit 0 RWUSIG: Remote wakeup signaling

When the application sets this bit, the core initiates remote signaling to wake up the USB host. The application must set this bit to instruct the core to exit the Suspend state. As specified in the USB 2.0 specification, the application must clear this bit 1 ms to 15 ms after setting it.

*Table 158* contains the minimum duration (according to device state) for which the Soft disconnect (SDIS) bit must be set for the USB host to detect a device disconnect. To accommodate clock jitter, it is recommended that the application add some extra delay to the specified minimum duration.

Table 158. Minimum duration for soft disconnect

Operating speed	Device state	Minimum duration
Full speed	Suspended	1 ms + 2.5 μs
Full speed	Idle	2.5 µs
Full speed	Not Idle or Suspended (Performing transactions)	2.5 µs

# OTG FS device status register (OTG FS DSTS)

Address offset: 0x808

Reset value: 0x0000 0010

This register indicates the status of the core with respect to USB-related events. It must be read on interrupts from the device all interrupts (OTG\_FS\_DAINT) register.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
				Rese	erved	ļ										FNS	SOF								Rese	erved		EERR	COMINA		SUSPSTS
										r	r	r	r	r	r	r	r	r	r	r	r	r	r					r	r	r	r

Bits 31:22 Reserved, must be kept at reset value.

Bits 21:8 FNSOF: Frame number of the received SOF

Bits 7:4 Reserved, must be kept at reset value.



#### Bit 3 **EERR:** Erratic error

The core sets this bit to report any erratic errors.

Due to erratic errors, the OTG\_FS controller goes into Suspended state and an interrupt is generated to the application with Early suspend bit of the OTG\_FS\_GINTSTS register (ESUSP bit in OTG\_FS\_GINTSTS). If the early suspend is asserted due to an erratic error, the application can only perform a soft disconnect recover.

## Bits 2:1 ENUMSPD: Enumerated speed

Indicates the speed at which the OTG\_FS controller has come up after speed detection through a chirp sequence.

01: Reserved 10: Reserved

11: Full speed (PHY clock is running at 48 MHz)

Others: reserved

# Bit 0 SUSPSTS: Suspend status

In device mode, this bit is set as long as a Suspend condition is detected on the USB. The core enters the Suspended state when there is no activity on the USB data lines for a period of 3 ms. The core comes out of the suspend:

- When there is an activity on the USB data lines
- When the application writes to the Remote wakeup signaling bit in the OTG\_FS\_DCTL register (RWUSIG bit in OTG\_FS\_DCTL).

# OTG\_FS device IN endpoint common interrupt mask register (OTG\_FS\_DIEPMSK)

Address offset: 0x810 Reset value: 0x0000 0000

This register works with each of the OTG\_FS\_DIEPINTx registers for all endpoints to generate an interrupt per IN endpoint. The IN endpoint interrupt for a specific status in the OTG\_FS\_DIEPINTx register can be masked by writing to the corresponding bit in this register. Status bits are masked by default.

(	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
												Re	eserv	ed												INEPNEM	INEPNMM	ITTXFEMSK	MOT	Reserved	MOde	XFRCM
																										rw	rw	rw	rw		rw	rw

Bits 31:7 Reserved, must be kept at reset value.

Bit 6 INEPNEM: IN endpoint NAK effective mask

0: Masked interrupt1: Unmasked interrupt

Bit 5 INEPNMM: IN token received with EP mismatch mask

0: Masked interrupt1: Unmasked interrupt

Bit 4 ITTXFEMSK: IN token received when TxFIFO empty mask

0: Masked interrupt1: Unmasked interrupt

Bit 3 TOM: Timeout condition mask (Non-isochronous endpoints)

0: Masked interrupt

1: Unmasked interrupt

Bit 2 Reserved, must be kept at reset value.

Bit 1 EPDM: Endpoint disabled interrupt mask

0: Masked interrupt1: Unmasked interrupt

Bit 0 XFRCM: Transfer completed interrupt mask

0: Masked interrupt1: Unmasked interrupt

# OTG\_FS device OUT endpoint common interrupt mask register (OTG\_FS\_DOEPMSK)

Address offset: 0x814

Reset value: 0x0000 0000

This register works with each of the OTG\_FS\_DOEPINTx registers for all endpoints to generate an interrupt per OUT endpoint. The OUT endpoint interrupt for a specific status in the OTG\_FS\_DOEPINTx register can be masked by writing into the corresponding bit in this register. Status bits are masked by default.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
												Re	eserv	red													OTEPDM	STUPM	Reserved	EPDM	XFRCM
																											rw	rw	ш	rw	rw

Bits 31:5 Reserved, must be kept at reset value.

Bit 4 OTEPDM: OUT token received when endpoint disabled mask

Applies to control OUT endpoints only.

0: Masked interrupt

1: Unmasked interrupt

Bit 3 STUPM: SETUP phase done mask

Applies to control endpoints only.

0: Masked interrupt

1: Unmasked interrupt

Bit 2 Reserved, must be kept at reset value.

Bit 1 EPDM: Endpoint disabled interrupt mask

0: Masked interrupt

1: Unmasked interrupt

Bit 0 XFRCM: Transfer completed interrupt mask

0: Masked interrupt

1: Unmasked interrupt

# OTG\_FS device all endpoints interrupt register (OTG\_FS\_DAINT)

Address offset: 0x818 Reset value: 0x0000 0000

When a significant event occurs on an endpoint, a OTG\_FS\_DAINT register interrupts the application using the Device OUT endpoints interrupt bit or Device IN endpoints interrupt bit of the OTG\_FS\_GINTSTS register (OEPINT or IEPINT in OTG\_FS\_GINTSTS, respectively). There is one interrupt bit per endpoint, up to a maximum of 16 bits for OUT endpoints and 16 bits for IN endpoints. For a bidirectional endpoint, the corresponding IN and OUT interrupt bits are used. Bits in this register are set and cleared when the application sets and clears bits in the corresponding Device Endpoint-x interrupt register (OTG\_FS\_DIEPINTx/OTG\_FS\_DOEPINTx).

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

							OEF	PINT															IEP	INT							
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:16 **OEPINT:** OUT endpoint interrupt bits

One bit per OUT endpoint:

Bit 16 for OUT endpoint 0, bit 18 for OUT endpoint 3.

Bits 15:0 IEPINT: IN endpoint interrupt bits

One bit per IN endpoint:

Bit 0 for IN endpoint 0, bit 3 for endpoint 3.

# OTG\_FS all endpoints interrupt mask register (OTG\_FS\_DAINTMSK)

Address offset: 0x81C Reset value: 0x0000 0000

The OTG\_FS\_DAINTMSK register works with the Device endpoint interrupt register to interrupt the application when an event occurs on a device endpoint. However, the

OTG\_FS\_DAINT register bit corresponding to that interrupt is still set.

 $31 \quad 30 \quad 29 \quad 28 \quad 27 \quad 26 \quad 25 \quad 24 \quad 23 \quad 22 \quad 21 \quad 20 \quad 19 \quad 18 \quad 17 \quad 16 \quad 15 \quad 14 \quad 13 \quad 12 \quad 11 \quad 10 \quad 9 \quad 8 \quad 7 \quad 6 \quad 5 \quad 4 \quad 3 \quad 2 \quad 1 \quad 0$ 

							OE	РМ															IEF	PM							
rw	rw	rw	rw	rw	rw	rw	rw	rw																							

Bits 31:16 **OEPM:** OUT EP interrupt mask bits

One per OUT endpoint:

Bit 16 for OUT EP 0, bit 18 for OUT EP 3

0: Masked interrupt1: Unmasked interrupt

Bits 15:0 IEPM: IN EP interrupt mask bits

One bit per IN endpoint:

Bit 0 for IN EP 0, bit 3 for IN EP 3

0: Masked interrupt1: Unmasked interrupt

# OTG\_FS device V<sub>BUS</sub> discharge time register (OTG\_FS\_DVBUSDIS)

Address offset: 0x0828 Reset value: 0x0000 17D7

This register specifies the V<sub>BUS</sub> discharge time after V<sub>BUS</sub> pulsing during SRP.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Peneryod								VBU	SDT							
Reserved	rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bits 31:16 Reserved, must be kept at reset value.

Bits 15:0 **VBUSDT:** Device  $V_{BUS}$  discharge time

Specifies the V<sub>BUS</sub> discharge time after V<sub>BUS</sub> pulsing during SRP. This value equals:

V<sub>BUS</sub> discharge time in PHY clocks / 1 024

Depending on your  $V_{BUS}$  load, this value may need adjusting.

# OTG\_FS device V<sub>BUS</sub> pulsing time register (OTG\_FS\_DVBUSPULSE)

Address offset: 0x082C Reset value: 0x0000 05B8

This register specifies the V<sub>BUS</sub> pulsing time during SRP.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved						DVB	USP					
neserveu	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:12 Reserved, must be kept at reset value.

Bits 11:0 **DVBUSP:** Device V<sub>BUS</sub> pulsing time

Specifies the V<sub>BUS</sub> pulsing time during SRP. This value equals:

V<sub>BUS</sub> pulsing time in PHY clocks / 1 024

# OTG\_FS device IN endpoint FIFO empty interrupt mask register: (OTG\_FS\_DIEPEMPMSK)

Address offset: 0x834 Reset value: 0x0000 0000

This register is used to control the IN endpoint FIFO empty interrupt generation

(TXFE\_OTG\_FS\_DIEPINTx).

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved							IN	IEPT	XFE	М						
neserveu	rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bits 31:16 Reserved, must be kept at reset value.

Bits 15:0 INEPTXFEM: IN EP Tx FIFO empty interrupt mask bits

These bits act as mask bits for OTG\_FS\_DIEPINTx.

TXFE interrupt one bit per IN endpoint:
Bit 0 for IN endpoint 0, bit 3 for IN endpoint 3

0: Masked interrupt1: Unmasked interrupt

# OTG\_FS device control IN endpoint 0 control register (OTG\_FS\_DIEPCTL0)

Address offset: 0x900

Reset value: 0x0000 0000

This section describes the OTG\_FS\_DIEPCTL0 register. Nonzero control endpoints use

registers for endpoints 1–3.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

EPENA	EPDIS	eserved	SNAK	CNAK		TXF	NUM	l	STALL	eserved	EP'	TYP	NAKSTS	eserved	USBAEP	Reserved	MP	PSIZ
r	r	Œ	w	٧	rw	rw	rw	rw	rs	н	r	r	r	R	r		rw	rw

#### Bit 31 **EPENA:** Endpoint enable

The application sets this bit to start transmitting data on the endpoint 0.

The core clears this bit before setting any of the following interrupts on this endpoint:

- Endpoint disabled
- Transfer completed

## Bit 30 EPDIS: Endpoint disable

The application sets this bit to stop transmitting data on an endpoint, even before the transfer for that endpoint is complete. The application must wait for the Endpoint disabled interrupt before treating the endpoint as disabled. The core clears this bit before setting the Endpoint disabled interrupt. The application must set this bit only if Endpoint enable is already set for this endpoint.

Bits 29:28 Reserved, must be kept at reset value.

# Bit 27 SNAK: Set NAK

A write to this bit sets the NAK bit for the endpoint.

Using this bit, the application can control the transmission of NAK handshakes on an endpoint. The core can also set this bit for an endpoint after a SETUP packet is received on that endpoint.

## Bit 26 CNAK: Clear NAK

# A write to this bit clears the NAK bit for the endpoint.

# Bits 25:22 TXFNUM: TxFIFO number

This value is set to the FIFO number that is assigned to IN endpoint 0.

# Bit 21 STALL: STALL handshake

The application can only set this bit, and the core clears it when a SETUP token is received for this endpoint. If a NAK bit, a Global IN NAK or Global OUT NAK is set along with this bit, the STALL bit takes priority.

Bit 20 Reserved, must be kept at reset value.

# Bits 19:18 EPTYP: Endpoint type

Hardcoded to '00' for control.

#### Bit 17 NAKSTS: NAK status

Indicates the following:

0: The core is transmitting non-NAK handshakes based on the FIFO status

1: The core is transmitting NAK handshakes on this endpoint.

When this bit is set, either by the application or core, the core stops transmitting data, even if there are data available in the TxFIFO. Irrespective of this bit's setting, the core always responds to SETUP data packets with an ACK handshake.

Bit 16 Reserved, must be kept at reset value.

## Bit 15 USBAEP: USB active endpoint

This bit is always set to 1, indicating that control endpoint 0 is always active in all configurations and interfaces.

Bits 14:2 Reserved, must be kept at reset value.

## Bits 1:0 MPSIZ: Maximum packet size

The application must program this field with the maximum packet size for the current logical endpoint.

00: 64 bytes 01: 32 bytes 10: 16 bytes 11: 8 bytes

# OTG device endpoint-x control register (OTG\_FS\_DIEPCTLx) (x = 1..3, where x = Endpoint number)

Address offset: 0x900 + (Endpoint\_number × 0x20)

Reset value: 0x0000 0000

The application uses this register to control the behavior of each logical endpoint other than endpoint 0.

31	30	29	28	21	20	25	24	23	22	21	20	19	18	17	10	15	14	13	12	11	10	9	Ö	/	О	Э	4	3	2		U
EPENA	EPDIS	SODDFRM	SDOPID/SEVNFRM	SNAK	CNAK		TXF	NUM	1	Stall	Reserved	dALGE		NAKSTS	EONUM/DPID	USBAEP		Rese	rved						N	/IPSI	Z				
rs	rs	w	W	W	w	rw	rw	rw	rw	rw/ rs		rw	rw	r	r	rw					rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

# Bit 31 EPENA: Endpoint enable

The application sets this bit to start transmitting data on an endpoint.

The core clears this bit before setting any of the following interrupts on this endpoint:

- SETUP phase done
- Endpoint disabled
- Transfer completed

# Bit 30 EPDIS: Endpoint disable

The application sets this bit to stop transmitting/receiving data on an endpoint, even before the transfer for that endpoint is complete. The application must wait for the Endpoint disabled interrupt before treating the endpoint as disabled. The core clears this bit before setting the Endpoint disabled interrupt. The application must set this bit only if Endpoint enable is already set for this endpoint.

## Bit 29 SODDFRM: Set odd frame

Applies to isochronous IN and OUT endpoints only.

Writing to this field sets the Even/Odd frame (EONUM) field to odd frame.

#### Bit 28 SD0PID: Set DATA0 PID

Applies to interrupt/bulk IN endpoints only.

Writing to this field sets the endpoint data PID (DPID) field in this register to DATA0.

#### SEVNFRM: Set even frame

Applies to isochronous IN endpoints only.

Writing to this field sets the Even/Odd frame (EONUM) field to even frame.

## Bit 27 SNAK: Set NAK

A write to this bit sets the NAK bit for the endpoint.

Using this bit, the application can control the transmission of NAK handshakes on an endpoint. The core can also set this bit for OUT endpoints on a Transfer completed interrupt, or after a SETUP is received on the endpoint.

#### Bit 26 CNAK: Clear NAK

A write to this bit clears the NAK bit for the endpoint.

#### Bits 25:22 TXFNUM: TxFIFO number

These bits specify the FIFO number associated with this endpoint. Each active IN endpoint must be programmed to a separate FIFO number.

This field is valid only for IN endpoints.

#### Bit 21 STALL: STALL handshake

Applies to non-control, non-isochronous IN endpoints only (access type is rw).

The application sets this bit to stall all tokens from the USB host to this endpoint. If a NAK bit, Global IN NAK, or Global OUT NAK is set along with this bit, the STALL bit takes priority. Only the application can clear this bit, never the core.

Applies to control endpoints only (access type is rs).

The application can only set this bit, and the core clears it, when a SETUP token is received for this endpoint. If a NAK bit, Global IN NAK, or Global OUT NAK is set along with this bit, the STALL bit takes priority. Irrespective of this bit's setting, the core always responds to SETUP data packets with an ACK handshake.

# Bit 20 Reserved, must be kept at reset value.

### Bits 19:18 EPTYP: Endpoint type

This is the transfer type supported by this logical endpoint.

00: Control 01: Isochronous

10: Bulk11: Interrupt

#### Bit 17 NAKSTS: NAK status

It indicates the following:

0: The core is transmitting non-NAK handshakes based on the FIFO status.

1: The core is transmitting NAK handshakes on this endpoint.

When either the application or the core sets this bit:

For non-isochronous IN endpoints: The core stops transmitting any data on an IN endpoint, even if there are data available in the TxFIFO.

For isochronous IN endpoints: The core sends out a zero-length data packet, even if there are data available in the TxFIFO.

Irrespective of this bit's setting, the core always responds to SETUP data packets with an ACK handshake.

## Bit 16 EONUM: Even/odd frame

Applies to isochronous IN endpoints only.

Indicates the frame number in which the core transmits/receives isochronous data for this endpoint. The application must program the even/odd frame number in which it intends to transmit/receive isochronous data for this endpoint using the SEVNFRM and SODDFRM fields in this register.

0: Even frame

1: Odd frame

**DPID:** Endpoint data PID

Applies to interrupt/bulk IN endpoints only.

Contains the PID of the packet to be received or transmitted on this endpoint. The application must program the PID of the first packet to be received or transmitted on this endpoint, after the endpoint is activated. The application uses the SD0PID register field to program either DATA0 or DATA1 PID.

0: DATA0 1: DATA1

## Bit 15 USBAEP: USB active endpoint

Indicates whether this endpoint is active in the current configuration and interface. The core clears this bit for all endpoints (other than EP 0) after detecting a USB reset. After receiving the SetConfiguration and SetInterface commands, the application must program endpoint registers accordingly and set this bit.

#### Bits 14:11 Reserved, must be kept at reset value.

## Bits 10:0 MPSIZ: Maximum packet size

The application must program this field with the maximum packet size for the current logical endpoint. This value is in bytes.

# OTG\_FS device control OUT endpoint 0 control register (OTG\_FS\_DOEPCTL0)

Address offset: 0xB00

Reset value: 0x0000 8000

This section describes the OTG\_FS\_DOEPCTL0 register. Nonzero control endpoints use

registers for endpoints 1–3.

## 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

į	EPENA	EPDIS	eserved	SNAK	CNAK	Reserved	Stall	SNPM	EPT	ΥP	NAKSTS	eserved	USBAEP	Reserved	MPS	SIZ	
,	N	r	Я	w	w		rs	rw	r	r	r	Ж	r		r	r	

## Bit 31 EPENA: Endpoint enable

The application sets this bit to start transmitting data on endpoint 0.

The core clears this bit before setting any of the following interrupts on this endpoint:

- SETUP phase done
- Endpoint disabled
- Transfer completed

### Bit 30 EPDIS: Endpoint disable

The application cannot disable control OUT endpoint 0.

Bits 29:28 Reserved, must be kept at reset value.

## Bit 27 SNAK: Set NAK

A write to this bit sets the NAK bit for the endpoint.

Using this bit, the application can control the transmission of NAK handshakes on an endpoint. The core can also set this bit on a Transfer completed interrupt, or after a SETUP is received on the endpoint.

#### Bit 26 CNAK: Clear NAK

A write to this bit clears the NAK bit for the endpoint.

Bits 25:22 Reserved, must be kept at reset value.

# Bit 21 STALL: STALL handshake

The application can only set this bit, and the core clears it, when a SETUP token is received for this endpoint. If a NAK bit or Global OUT NAK is set along with this bit, the STALL bit takes priority. Irrespective of this bit's setting, the core always responds to SETUP data packets with an ACK handshake.

## Bit 20 SNPM: Snoop mode

This bit configures the endpoint to Snoop mode. In Snoop mode, the core does not check the correctness of OUT packets before transferring them to application memory.

# Bits 19:18 EPTYP: Endpoint type

Hardcoded to 2'b00 for control.

#### Bit 17 NAKSTS: NAK status

Indicates the following:

0: The core is transmitting non-NAK handshakes based on the FIFO status.

1: The core is transmitting NAK handshakes on this endpoint.

When either the application or the core sets this bit, the core stops receiving data, even if there is space in the RxFIFO to accommodate the incoming packet. Irrespective of this bit's setting, the core always responds to SETUP data packets with an ACK handshake.

Bit 16 Reserved, must be kept at reset value.

## Bit 15 USBAEP: USB active endpoint

This bit is always set to 1, indicating that a control endpoint 0 is always active in all configurations and interfaces.

Bits 14:2 Reserved, must be kept at reset value.

## Bits 1:0 MPSIZ: Maximum packet size

The maximum packet size for control OUT endpoint 0 is the same as what is programmed in control IN endpoint 0.

00: 64 bytes 01: 32 bytes 10: 16 bytes 11: 8 bytes

# OTG\_FS device endpoint-x control register (OTG\_FS\_DOEPCTLx) (x = 1...3, where $x = Endpoint_number$ )

Address offset for OUT endpoints: 0xB00 + (Endpoint\_number × 0x20)

Reset value: 0x0000 0000

The application uses this register to control the behavior of each logical endpoint other than endpoint 0.

31	30	29	28	21	20	25	24	23	22	21	20	19	10	17	10	15	14	13	12	1.1	10	9	Ö	/	О	Э	4	3	2	- 1	U
EPENA	EPDIS	SODDFRM/SD1PID	SD0PID/SEVNFRM	SNAK	CNAK		Rese	erved	I	Stall	SNPM	EPTVP	- -	NAKSTS	EONUM/DPID	USBAEP		Rese	rved						N	/IPSI	Z				
rs	rs	w	w	w	w					rw/ rs	rw	rw	rw	r	r	rw					rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

# Bit 31 EPENA: Endpoint enable

Applies to IN and OUT endpoints.

The application sets this bit to start transmitting data on an endpoint.

The core clears this bit before setting any of the following interrupts on this endpoint:

- SETUP phase done
- Endpoint disabled
- Transfer completed

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## Bit 30 EPDIS: Endpoint disable

The application sets this bit to stop transmitting/receiving data on an endpoint, even before the transfer for that endpoint is complete. The application must wait for the Endpoint disabled interrupt before treating the endpoint as disabled. The core clears this bit before setting the Endpoint disabled interrupt. The application must set this bit only if Endpoint enable is already set for this endpoint.

#### Bit 29 SD1PID: Set DATA1 PID

Applies to interrupt/bulk IN and OUT endpoints only. Writing to this field sets the endpoint data PID (DPID) field in this register to DATA1.

**SODDFRM:** Set odd frame

Applies to isochronous IN and OUT endpoints only. Writing to this field sets the Even/Odd frame (EONUM) field to odd frame.

#### Bit 28 SD0PID: Set DATA0 PID

Applies to interrupt/bulk OUT endpoints only.

Writing to this field sets the endpoint data PID (DPID) field in this register to DATA0.

#### **SEVNFRM:** Set even frame

Applies to isochronous OUT endpoints only.

Writing to this field sets the Even/Odd frame (EONUM) field to even frame.

#### Bit 27 SNAK: Set NAK

A write to this bit sets the NAK bit for the endpoint.

Using this bit, the application can control the transmission of NAK handshakes on an endpoint. The core can also set this bit for OUT endpoints on a Transfer Completed interrupt, or after a SETUP is received on the endpoint.

## Bit 26 CNAK: Clear NAK

A write to this bit clears the NAK bit for the endpoint.

# Bits 25:22 Reserved, must be kept at reset value.

# Bit 21 STALL: STALL handshake

Applies to non-control, non-isochronous OUT endpoints only (access type is rw). The application sets this bit to stall all tokens from the USB host to this endpoint. If a NAK bit, Global IN NAK, or Global OUT NAK is set along with this bit, the STALL bit takes priority. Only the application can clear this bit, never the core.

Applies to control endpoints only (access type is rs).

The application can only set this bit, and the core clears it, when a SETUP token is received for this endpoint. If a NAK bit, Global IN NAK, or Global OUT NAK is set along with this bit, the STALL bit takes priority. Irrespective of this bit's setting, the core always responds to SETUP data packets with an ACK handshake.

## Bit 20 SNPM: Snoop mode

This bit configures the endpoint to Snoop mode. In Snoop mode, the core does not check the correctness of OUT packets before transferring them to application memory.

## Bits 19:18 EPTYP: Endpoint type

This is the transfer type supported by this logical endpoint.

00: Control

01: Isochronous

10: Bulk

11: Interrupt

#### Bit 17 NAKSTS: NAK status

Indicates the following:

0: The core is transmitting non-NAK handshakes based on the FIFO status.

1: The core is transmitting NAK handshakes on this endpoint.

When either the application or the core sets this bit:

The core stops receiving any data on an OUT endpoint, even if there is space in the RxFIFO to accommodate the incoming packet.

Irrespective of this bit's setting, the core always responds to SETUP data packets with an ACK handshake.

#### Bit 16 EONUM: Even/odd frame

Applies to isochronous IN and OUT endpoints only.

Indicates the frame number in which the core transmits/receives isochronous data for this endpoint. The application must program the even/odd frame number in which it intends to transmit/receive isochronous data for this endpoint using the SEVNFRM and SODDFRM fields in this register.

0: Even frame

1: Odd frame

**DPID:** Endpoint data PID

Applies to interrupt/bulk OUT endpoints only.

Contains the PID of the packet to be received or transmitted on this endpoint. The application must program the PID of the first packet to be received or transmitted on this endpoint, after the endpoint is activated. The application uses the SD0PID register field to program either DATA0 or DATA1 PID.

0: DATA0 1: DATA1

## Bit 15 USBAEP: USB active endpoint

Indicates whether this endpoint is active in the current configuration and interface. The core clears this bit for all endpoints (other than EP 0) after detecting a USB reset. After receiving the SetConfiguration and SetInterface commands, the application must program endpoint registers accordingly and set this bit.

## Bits 14:11 Reserved, must be kept at reset value.

# Bits 10:0 MPSIZ: Maximum packet size

The application must program this field with the maximum packet size for the current logical endpoint. This value is in bytes.



# OTG\_FS device endpoint-x interrupt register (OTG\_FS\_DIEPINTx) (x = 0..3, where x = Endpoint number)

Address offset:  $0x908 + (Endpoint number \times 0x20)$ 

Reset value: 0x0000 0080

This register indicates the status of an endpoint with respect to USB- and AHB-related events. It is shown in *Figure 358*. The application must read this register when the IN endpoints interrupt bit of the Core interrupt register (IEPINT in OTG FS GINTSTS) is set. Before the application can read this register, it must first read the device all endpoints interrupt (OTG\_FS\_DAINT) register to get the exact endpoint number for the Device endpoint-x interrupt register. The application must clear the appropriate bit in this register to clear the corresponding bits in the OTG FS DAINT and OTG FS GINTSTS registers.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	TXFE	INEPNE	erved	ITTXFE	70C	erved	EPDISD	XFRC	
	r	rc_ w1 /rw	Res	rc_ w1	rc_ w1	Res	rc_ w1	rc_ w1	

Bits 31:8 Reserved, must be kept at reset value.

### Bit 7 TXFE: Transmit FIFO empty

This interrupt is asserted when the TxFIFO for this endpoint is either half or completely empty. The half or completely empty status is determined by the TxFIFO Empty Level bit in the OTG\_FS\_GAHBCFG register (TXFELVL bit in OTG\_FS\_GAHBCFG).

# Bit 6 INEPNE: IN endpoint NAK effective

This bit can be cleared when the application clears the IN endpoint NAK by writing to the CNAK bit in OTG\_FS\_DIEPCTLx.

This interrupt indicates that the core has sampled the NAK bit set (either by the application or by the core). The interrupt indicates that the IN endpoint NAK bit set by the application has taken effect in the core.

This interrupt does not guarantee that a NAK handshake is sent on the USB. A STALL bit takes priority over a NAK bit.

Bit 5 Reserved, must be kept at reset value.

# Bit 4 ITTXFE: IN token received when TxFIFO is empty

Applies to non-periodic IN endpoints only.

Indicates that an IN token was received when the associated TxFIFO (periodic/non-periodic) was empty. This interrupt is asserted on the endpoint for which the IN token was received.

## Bit 3 TOC: Timeout condition

Applies only to Control IN endpoints.

Indicates that the core has detected a timeout condition on the USB for the last IN token on this endpoint.

Bit 2 Reserved, must be kept at reset value.

## Bit 1 **EPDISD:** Endpoint disabled interrupt

This bit indicates that the endpoint is disabled per the application's request.

#### Bit 0 XFRC: Transfer completed interrupt

This field indicates that the programmed transfer is complete on the AHB as well as on the USB, for this endpoint.

# OTG\_FS device endpoint-x interrupt register (OTG\_FS\_DOEPINTx) (x = 0...3, where $x = Endpoint_number$ )

Address offset: 0xB08 + (Endpoint\_number × 0x20)

Reset value: 0x0000 0080

This register indicates the status of an endpoint with respect to USB- and AHB-related events. It is shown in *Figure 358*. The application must read this register when the OUT Endpoints Interrupt bit of the OTG\_FS\_GINTSTS register (OEPINT bit in OTG\_FS\_GINTSTS) is set. Before the application can read this register, it must first read the OTG\_FS\_DAINT register to get the exact endpoint number for the OTG\_FS\_DOEPINTx register. The application must clear the appropriate bit in this register to clear the corresponding bits in the OTG\_FS\_DAINT and OTG\_FS\_GINTSTS registers.

3	1	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
												Rese	erved	I											Reserved	₹ g B2BSTUP	Reserved		a STUP	Reserved	QSIQA3 c' w1	L XFRC

Bits 31:7 Reserved, must be kept at reset value.

Bit 6 B2BSTUP: Back-to-back SETUP packets received

Applies to control OUT endpoint only.

This bit indicates that the core has received more than three back-to-back SETUP packets for this particular endpoint.

- Bit 5 Reserved, must be kept at reset value.
- Bit 4 OTEPDIS: OUT token received when endpoint disabled

Applies only to control OUT endpoints.

Indicates that an OUT token was received when the endpoint was not yet enabled. This interrupt is asserted on the endpoint for which the OUT token was received.

Bit 3 STUP: SETUP phase done

Applies to control OUT endpoint only.

Indicates that the SETUP phase for the control endpoint is complete and no more back-to-back SETUP packets were received for the current control transfer. On this interrupt, the application can decode the received SETUP data packet.

- Bit 2 Reserved, must be kept at reset value.
- Bit 1 **EPDISD:** Endpoint disabled interrupt

This bit indicates that the endpoint is disabled per the application's request.

Bit 0 XFRC: Transfer completed interrupt

This field indicates that the programmed transfer is complete on the AHB as well as on the USB, for this endpoint.

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# OTG\_FS device IN endpoint 0 transfer size register (OTG\_FS\_DIEPTSIZ0)

Address offset: 0x910 Reset value: 0x0000 0000

The application must modify this register before enabling endpoint 0. Once endpoint 0 is enabled using the endpoint enable bit in the device control endpoint 0 control registers (EPENA in OTG\_FS\_DIEPCTL0), the core modifies this register. The application can only read this register once the core has cleared the Endpoint enable bit.

Nonzero endpoints use the registers for endpoints 1-3.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	PKT	CNT	Decembed			Х	FRS	ΙZ		
neserved	rw	rw	Reserved	rw	rw	rw	rw	rw	rw	rw

Bits 31:21 Reserved, must be kept at reset value.

#### Bits 20:19 PKTCNT: Packet count

Indicates the total number of USB packets that constitute the Transfer Size amount of data for endpoint 0.

This field is decremented every time a packet (maximum size or short packet) is read from the TxFIFO.

Bits 18:7 Reserved, must be kept at reset value.

#### Bits 6:0 XFRSIZ: Transfer size

Indicates the transfer size in bytes for endpoint 0. The core interrupts the application only after it has exhausted the transfer size amount of data. The transfer size can be set to the maximum packet size of the endpoint, to be interrupted at the end of each packet.

The core decrements this field every time a packet from the external memory is written to the TxFIFO.

# OTG\_FS device OUT endpoint 0 transfer size register (OTG\_FS\_DOEPTSIZ0)

Address offset: 0xB10 Reset value: 0x0000 0000

The application must modify this register before enabling endpoint 0. Once endpoint 0 is enabled using the Endpoint enable bit in the OTG\_FS\_DOEPCTL0 registers (EPENA bit in OTG\_FS\_DOEPCTL0), the core modifies this register. The application can only read this register once the core has cleared the Endpoint enable bit.

Nonzero endpoints use the registers for endpoints 1-3.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

eserved	STUPC NT	Reserved	PKTCNT	Reserved			XF	RSI	Z		
ω.	rw rw		rw	r	w r	N I	w	rw	rw	rw	rw

Bit 31 Reserved, must be kept at reset value.

## Bits 30:29 STUPCNT: SETUP packet count

This field specifies the number of back-to-back SETUP data packets the endpoint can receive.

01: 1 packet 10: 2 packets 11: 3 packets

Bits 28:20 Reserved, must be kept at reset value.

#### Bit 19 PKTCNT: Packet count

This field is decremented to zero after a packet is written into the RxFIFO.

Bits 18:7 Reserved, must be kept at reset value.

## Bits 6:0 XFRSIZ: Transfer size

Indicates the transfer size in bytes for endpoint 0. The core interrupts the application only after it has exhausted the transfer size amount of data. The transfer size can be set to the maximum packet size of the endpoint, to be interrupted at the end of each packet.

The core decrements this field every time a packet is read from the RxFIFO and written to the external memory.

# OTG\_FS device endpoint-x transfer size register (OTG\_FS\_DIEPTSIZx) (x = 1..3, where x = Endpoint\_number)

Address offset:  $0x910 + (Endpoint number \times 0x20)$ 

Reset value: 0x0000 0000

The application must modify this register before enabling the endpoint. Once the endpoint is enabled using the Endpoint enable bit in the OTG\_FS\_DIEPCTLx registers (EPENA bit in OTG\_FS\_DIEPCTLx), the core modifies this register. The application can only read this register once the core has cleared the Endpoint enable bit.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

pe	MC	CNT					PKT	CNT													Х	FRS	ΙZ								
Reserve	rw/ r/r w	rw/ r/r w	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bit 31 Reserved, must be kept at reset value.

#### Bits 30:29 MCNT: Multi count

For periodic IN endpoints, this field indicates the number of packets that must be transmitted per frame on the USB. The core uses this field to calculate the data PID for isochronous IN endpoints.

01: 1 packet 10: 2 packets 11: 3 packets

#### Bit 28:19 PKTCNT: Packet count

Indicates the total number of USB packets that constitute the Transfer Size amount of data for this endpoint.

This field is decremented every time a packet (maximum size or short packet) is read from the TxFIFO.

# Bits 18:0 XFRSIZ: Transfer size

This field contains the transfer size in bytes for the current endpoint. The core only interrupts the application after it has exhausted the transfer size amount of data. The transfer size can be set to the maximum packet size of the endpoint, to be interrupted at the end of each packet.

The core decrements this field every time a packet from the external memory is written to the TxFIFO.

# OTG\_FS device IN endpoint transmit FIFO status register (OTG\_FS\_DTXFSTSx) (x = 0..3, where x = Endpoint\_number)

Address offset for IN endpoints:  $0x918 + (Endpoint\_number \times 0x20)$  This read-only register contains the free space information for the Device IN endpoint TxFIFO.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

| INEPTFSAV | Reserved | Reserved

31:16 Reserved, must be kept at reset value.

## 15:0 INEPTFSAV: IN endpoint TxFIFO space available

Indicates the amount of free space available in the Endpoint TxFIFO.

Values are in terms of 32-bit words:

0x0: Endpoint TxFIFO is full 0x1: 1 word available 0x2: 2 words available

0xn: n words available Others: Reserved

# OTG\_FS device OUT endpoint-x transfer size register (OTG\_FS\_DOEPTSIZx) $(x = 1..3, where x = Endpoint_number)$

Address offset: 0xB10 + (Endpoint\_number × 0x20)

Reset value: 0x0000 0000

The application must modify this register before enabling the endpoint. Once the endpoint is enabled using Endpoint Enable bit of the OTG\_FS\_DOEPCTLx registers (EPENA bit in OTG\_FS\_DOEPCTLx), the core modifies this register. The application can only read this register once the core has cleared the Endpoint enable bit.

 $31 \quad 30 \quad 29 \quad 28 \quad 27 \quad 26 \quad 25 \quad 24 \quad 23 \quad 22 \quad 21 \quad 20 \quad 19 \quad 18 \quad 17 \quad 16 \quad 15 \quad 14 \quad 13 \quad 12 \quad 11 \quad 10 \quad 9 \quad 8 \quad 7 \quad 6 \quad 5 \quad 4 \quad 3 \quad 2 \quad 1 \quad 0$ 

rved	RXDI TUP	PID/S CNT					PKT	CNT	•												X	FRS	ΙZ								
Rese	rw/r/ rw	rw/r/ rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bit 31 Reserved, must be kept at reset value.

Bits 30:29 RXDPID: Received data PID

Applies to isochronous OUT endpoints only.

This is the data PID received in the last packet for this endpoint.

00: DATA0 01: DATA2 10: DATA1 11: MDATA

# STUPCNT: SETUP packet count

Applies to control OUT Endpoints only.

This field specifies the number of back-to-back SETUP data packets the endpoint can receive.

01: 1 packet 10: 2 packets 11: 3 packets

## Bit 28:19 PKTCNT: Packet count

Indicates the total number of USB packets that constitute the Transfer Size amount of data for this endpoint.

This field is decremented every time a packet (maximum size or short packet) is written to the RxFIFO.

# Bits 18:0 XFRSIZ: Transfer size

This field contains the transfer size in bytes for the current endpoint. The core only interrupts the application after it has exhausted the transfer size amount of data. The transfer size can be set to the maximum packet size of the endpoint, to be interrupted at the end of each packet.

The core decrements this field every time a packet is read from the RxFIFO and written to the external memory.

# 29.16.5 OTG\_FS power and clock gating control register (OTG\_FS\_PCGCCTL)

Address offset: 0xE00
Reset value: 0x0000 0000

This register is available in host and device modes.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
												Res	erve	d													₹ PHYSUSP	por 2000	neseiven	GATEHCLK	<b>≥</b> STPPCLK

## Bit 31:5 Reserved, must be kept at reset value.

# Bit 4 PHYSUSP: PHY Suspended

Indicates that the PHY has been suspended. This bit is updated once the PHY is suspended after the application has set the STPPCLK bit (bit 0).

Bits 3:2 Reserved, must be kept at reset value.

#### Bit 1 GATEHCLK: Gate HCLK

The application sets this bit to gate HCLK to modules other than the AHB Slave and Master and wakeup logic when the USB is suspended or the session is not valid. The application clears this bit when the USB is resumed or a new session starts.

# Bit 0 STPPCLK: Stop PHY clock

The application sets this bit to stop the PHY clock when the USB is suspended, the session is not valid, or the device is disconnected. The application clears this bit when the USB is resumed or a new session starts.

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# 29.16.6 OTG\_FS register map

The table below gives the USB OTG register map and reset values.

Table 159. OTG\_FS register map and reset values

Offset	Register	31	30	53	58	27	56	25	54	23	7	20	6	18	17	16	15	14	3	12	11	10	6	8	7	9	2	4	က	2	_	0
0x000	OTG_FS_GOT GCTL							erve					o BSVLD	o ASVLD	o DBCT .	CIDSTS 1		Rese				O HSHNPEN	0	O HNGSCS		F	Rese	erve	d		o SRQ	o SRQSCS
0x004	OTG_FS_GOT GINT					F	Res	erve	ed				o DBCDNE	о АВТОСНG	o HNGDET				Reserved				PINSSCHG	o SRSSCHG		Re	ser	ved		o SEDET	Re	es.
0x008	OTG_FS_GAH BCFG										Re	eser	ved		•									o PTXFELVL	O TXFELVL		F	Rese	erve	:d		o GINTMSK
0x00C		CTXPKT	FDMOD	FHMOD						Re	ser	ved							0	TR	DT 0	1	O HNPCAP	0	Reserved	□ PHYSEL	Re	serv	ved	T(	OC#	-
0x010	OTG_FS_GRST CTL	¬ AHBIDL								F	Res	erve	ed									0	TX 0	FNU	JM	0	0	ORXFFLSH	o Reserved	o FCRST	o HSRST	o CSRST
0x014	OTG_FS_GINT STS	WKUINT	SRQINT	DISCINT	CIDSCHG	Reserved	PTXFE	HCINT	HPRTINT	Reserved	IPXFR/INCOMPISOOUT	IISOIXFR	OEPINT	IEPINT	possional	0000	EOPF	ISOODRP	ENOMDNE	USBRST	USBSUSP	ESUSP	Beserved		GOUTNAKEFF	GINAKEFF	NPTXFE	RXFLVL	SOF	OTGINT	MMIS	CMOD
	Reset value	0	0	0	0		1	0	0		0	0	0	0			0	0	0	0	0	0			0	0	1	0	0	0	0	0
0x018	OTG_FS_GINT MSK	WUIM	SRQIM	DISCINT	CIDSCHGM	Reserved	PTXFEM	HCIM	PRTIM	Reserved	IPXFRM/IISOOXFRM	IISOIXFRM	OEPINT	IEPINT	EPMISM	Reserved	EOPFM	ISOODRPM	ENUMDNEM	USBRST	USBSUSPM	ESUSPM	Beserved		GONAKEFFM	GINAKEFFM	NPTXFEM	RXFLVLM	SOFM	OTGINT	MMISM	Reserved
	Reset value	0	0	0	0		0	0	0		0	0	0	0	0		0	0	0	0	0	0			0	0	0	0	0	0	0	
	OTG_FS_GRXS TSR (host mode)					Re	ser	ved				0	PKT			DF	PID	0	10	10	I 0		BCN'		0	I 0	0	10		CHN		1
0x01C	OTG_FS_GRXS TSR (Device mode)			Re	ser	ved				RMNUI			PKT	ST	3		O PID	0	0	0	0		0 BCN		0	0	0	0		EPN		
	Reset value OTG_FS_GRXS TSR (host mode)					Re	ser	ved	0	0 0	0		0 PKT	1	1		0 PID	0	0	0	0		0 BCN	0 T	0	0	0	0	0	CHN		0
0x020	Reset value  OTG_FS_GRXS TSPR (Device mode)  Reset value			Re	eser	ved				RMNUI			0 PKT	ST			0 PID	0	0		0		BCN'		0	0		0		0 EPN		
0x024	OTG_FS_GRXF SIZ Reset value							ſ	Rese	erved			•				0	0	0	0	0	0	1	RX 0	FD 0	0	0	0	0	0	0	0

Table 159. OTG\_FS register map and reset values (continued)

		<u>_</u>	<u>,</u>	<del></del> .		I -	T 1 -	<del></del>				<del>,</del>		I			-	1 1	
Offset	Register	31	78 78 78	27 26	25 24	23	20	19	17	16	15	13	7	10	6 8 ~	9	4	N 0	- 0
0x028	OTG_FS_HNPT XFSIZ/ OTG_FS_DIEP			N	PTXFC	)/TX0F	D							N	PTXFSA/TX	0FSA			
	TXF0 Reset value	0 0 1	0 0	0   0	0 0	0 0	0 0	0 0	0	0	0	0   0   0	0 0	0	1 0 0	0 0	0 0	0 0 1	0 0
	OTG_FS_HNPT			(QTOP	-   -	- 1 -	NPTQ		1 - 1		- 1	-   -   -		_	NPTXFSA			1 - 1	- 1 -
0x02C	XSTS Reset value	Res.			0   0	0   0	1010		0	0	0	0 0 0	0 0	0		0 0	0 0	0 1 0 1	0 0
0x038	OTG_FS_ GCCFG	_   •	- 1 - 1	eserved		0 0	S	o VBUSBSEN o		o PWRDWN o	0	0   0   0	,   0	<u> </u>	Reserved		0 0	,   0	0 0
	OTG FS CID						0 0	0 0	PRO		JCT	ID							
0x03C	Reset value	0 0	0 0	0 0	0 0	0 0	0 0	0 0		0		0 0 1	0	0	0 1 0	0 0	0 0	0 0	0 0
0x100	OTG_FS_HPTX FSIZ Reset value	0   0	0   0	0   1	PTXI	FSIZ	1110	1 0	0 1	0	0	0   0   1	10	0	PTXSA	0   1	0 1 0	) [ 1 [	0   0
	OTG_FS_DIEP	9   9	0 0	<u> </u>	INEP		1 . 1 .		١٠١	•		0 0			INEPTXS.			<u> </u>	-1-
0x104	TXF1	010	0101	0101			1010	0 0		0	0 1	01017	110	1			0 1 4	1101	
	Reset value  OTG_FS_DIEP	0 0	0 0	0 0	1 0		0 0	0 0	0	0	0	0 0 0	0	l		0 0	0 (	0 0	0 0
0x108	TXF2				INEP.	TXFD									INEPTXS	A			
	Reset value	0 0	0 0	0 0	1 0	0 0	0 0	0 0	0	0	0	0 0 0	0	1	0 0 0	0 0	0 (	0	0 0
0x10C	OTG_FS_DIEP TXF3				INEP	TXFD									INEPTXS	Α			
0.100	Reset value	0 0	0 0	0 0	1 0	0 0	0 0	0 0	0	0	0	0 0 0	0 0	1	0 0 0	0 0	0 0	0 0	0 0
0x400	OTG_FS_HCFG			<u> </u>	,		<b>.</b>	Re	serv	ed								FSLSS	FSLSPCS
	Reset value  OTG_FS_HFIR														FRIVL			0	0 0
0x404	Reset value				Rese	erved					1	1   1   (	1	0	1 0 0	1 1	0 0	0 0	0 0
0.400	OTG_FS_HFNU M				FTF	REM									FRNUM				
0x408	Reset value	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0	0	0	0   1   1	11	1	1 1 1 1	1 1 1	1 1 1	111	1   1
	OTG_FS_HPTX		PTXQ				PTX								PTXFSAV	1			
0x410	STS Reset value	0   0			0   0	VIV		YIY	I V I	Υ	VI	<del>V I V I V</del>	/ I V	V			VI	/ T V T	VIV
	OTG_FS_HAIN	0   0	0 0	0 0	0   0	YY	YY		<u>'</u>	'	-		' '	'	YYYY	YY		'	YY
0x414	ТТ				Rese	erved									HAINT				
	Reset value OTG FS HAIN										0	0 0 0	0	0	0 0 0	0 0	0 0	0 0	0 0
0x418	TMSK				Rese	erved									HAINTM				
	Reset value										0	0 0 0	0	0	0 0 0	0 0	0 0		0 0
0x440	OTG_FS_HPRT			Res	erved				PD		PTC <sup>-</sup>	ď	TO Id	2	E    -	P(	PENCHNG	PEN	PCDE1
	Reset value	1 1	_				1	0	0	0		0 0 0	0	0	0 0 0	0 0	0 (	0 0	0 0
0x500	OTG_FS_HCC HAR0	CHENA	ODDFRM	D	AD		MCNT	EPTYP	LSDEV	Reserved	EPDIR	EPNU	М			MPSI	Z		
	Reset value			0 0	0 0	0 0	0 0	0 0	0	Œ	0	0 0 0	0	0	0 0 0	0 0	0 0	0	0 0
0x520	OTG_FS_HCC HAR1	CHENA	ODDFRM	D	AD		MCNT	EPTYP	LSDEV	Reserved	EPDIR	EPNU	М			MPSI	Z		
	Reset value	0 0	0 0	0 0	0 0	0 0	0 0	0 0	0	Œ	0	0 0 0	0	0	0 0 0	0 0	0 0	0 0	0 0
0x540	OTG_FS_HCC HAR2	CHENA	ODDFRM	D	AD		MCNT	EPTYP	LSDEV	Reserved	EPDIR	EPNU	М			MPSI	z		
	Reset value			0 0	0 0	0 0	0 0	0 0	0	Ř	0	0 0 0	0	0	0 0 0	0 0	0 0	0	0 0

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Table 159. OTG\_FS register map and reset values (continued)

Office   Register   Color   State		59. OIG_I	_		<del>-</del>							<u>_</u>	_				-				1		П		—т	$\neg \neg$	_
PROSERT VAILURE   D.   D.   D.   D.   D.   D.   D.   D	Offset	Register	31	_		28	72	22	24	2 2	2 2		17	16	15	4 5 2 1	10	6	8	7	9	2	4	က	7	- 0	>
OKASE    O	0x560	HAR3			ODDFRN									Reserved													
Preset value   O   O   O   O   O   O   O   O   O		Reset value		0		0 (	0 0	0	0	0 0	0 0	0 0	0	ш.	0	0 0 0 0	0	0	0	0	0	0	0	0	0	0 0	)
Preset value   O   O   O   O   O   O   O   O   O	0x580		CHENA	CHDIS	ODDFRM			DAD			MCNT	ЕРТҮР	LSDEV	eserved	EPDIR	EPNUM					N	IPSI	Z				
Ox5CO   OTG_FS_HCC   HAR6		Reset value	0	0	0	0 (	0 0	0	0	0 0	0 0	0 0	0	В	0	0 0 0 0	0	0	0	0	0	0	0	0	0	0 0	,
Ox5CO   OTG_FS_HCC   HAR6	0x5A0	HAR5			ODDFRM									Reserved													
Ox506   OTG_FS_HCIN T   Ox508   OTG_FS_HCIN T   Ox508   OTG_FS_HCIN T   Ox508   OTG_FS_HCIN T   Ox508   Ox50		Reset value	-			0 (	0 0	0	0	0 0	0 0	0 0	0	Ľ	<b>!</b>	0 0 0 0	0	0	0	0	0	0	0	0	0	0 0	_
Ox506   OTG_FS_HCIN T   Ox508   OTG_FS_HCIN T   Ox508   OTG_FS_HCIN T   Ox508   OTG_FS_HCIN T   Ox508   Ox50	0x5C0		CHENA	CHDIS	ODDFRN			DAD			MCNT	ЕРТҮР	LSDEV	eserved	EPDIR	EPNUM					N	IPSI	Z				
Reservalue		Reset value	0	0		0 (	0 0	0	0	0 0	0 0	0 0	0	В	0	0 0 0 0	0	0	0	0	0	0	0	0	0	0 0	,
Reservalue	0x5E0		CHENA	CHDIS	ODDFRM			DAD			MCNT	ЕРТҮР	LSDEV	eserved	EPDIR	EPNUM					N	IPSI	Z				
Ox528		Reset value	0	0	0	0 (	0 0	0	0	0 0	0 0	0 0	0	ш	0	0 0 0 0	_				0	0	0	0	0	0 0	,
Ox528	0x508	T0								Re	eserved							FRMOR	ввеяя		eserved			STALL	leserved		
OTG_FS_HCIN   T2		Reset value																				0	$\vdash$				-
OTG_FS_HCIN   T2	0x528	T1		Reserved  Reserved  Reserved  Reserved														FRMOF			eserved			STALL	eserved		
OTG_FS_HCIN   Reserved																						0	$\vdash$	- 1			
OTG_FS_HCIN   Reserved	0x548	T2								Re	eserved						DTERF	FRMOF	BBERF		Reserved			STALL	Reserved		
OTG_FS_HCIN T4	0x568	OTG_FS_HCIN		Reserved  Reserved																			$\vdash$	_	_		-
Reset value		Reset value																	0	0	Re	0	0	0	Re	0 0	,
Ox5A8   OTG_FS_HCIN   Feset value	0x588									Re	eserved						DTERR	FRMOR	BBERR	TXERR	served	ACK	NAK	STALL	served	XFRC	
O		Reset value															_		0		ď	0	0	-	_	0 0	)
Ox5C8   OTG_FS_HCIN   TO   O   O   O   O   O   O   O   O	0x5A8	T5		Reserved														FRMOF	BBERR	TXERR	eserved			STALL	eserved		
Name		Reset value															Ľ	Ŭ	•	•		0	Ť	-	_		
0x5E8         OTG_FS_HCIN T7         Reserved         HUNDRED HER HCIN TMSK1         Polytograph         HUNDRED HCIN TMSK1         Polytog	0x5C8	T6								Re	eserved							FRMOF	BBERF		eserved				eserved		
Reset value		Reset value																	-			0	1	$\vdash$	_	_	-1
0x50C         OTG_FS_HCIN TMSK0         Reserved         WH W	0x5E8	T7								Re	eserved							FRMOF	BBERF		eserved			STALL	eserved		
Reset value		Heset value															_				Щ	0	$\vdash$	-	ш		$\dashv$
0x52C OTG_FS_HCIN TMSK1 Reserved Reserv	0x50C	TMSK0								Re	eserved							FRMORN							Reserved		
		Reset value															_		-		0	0	0	0		0 0	Ц
	0x52C									Re	eserved									TXERRM	NYET	ACKM			Reserved		
		Reset value																		0	0	0	0	0	т_	0 0	Ī

Table 159. OTG\_FS register map and reset values (continued)

Offset	Register	31	30	28	22	2 2	24	23	22	7	20	19	9	17	16	15	41	13	12	1	9	6	8	7	9	2	4	3	2	_	0
0x54C	OTG_FS_HCIN TMSK2								Res			<u> </u>		•		•		•	•		_	0 FRMORM	o BBERRM	O TXERRM	o NYET	o ACKM	O NAKM	OSTALLM	Reserved	ОННМ	o XFRCM
0x56C	OTG_FS_HCIN TMSK3								Res	erve	ed										DTERRM	_	o BBERRM	o TXERRM			o NAKM	_	Reserved	ОННМ	o XFRCM
0x58C	OTG_FS_HCIN TMSK4		Reserved																				o BBERRM	O TXERRM	o NYET	o ACKM	O NAKM	OSTALLM	Reserved	ОННМ	o XFRCM
0x5AC	OTG_FS_HCIN TMSK5		Reserved													5 5 5 5												Reserved		o XFRCM	
0x5CC	OTG_FS_HCIN TMSK6		Reserved																		DTERRM	_	_	Σ		o ACKM		OSTALLM	Reserved	- +	OXFRCM
0x5EC	OTG_FS_HCIN TMSK7				O DTERRM O FRWORM O TXERRM O NYET O ACKM												NAKM	O STALLM	Reserved	CHHM	o XFRCM o										
0x510	OTG_FS_HCTS IZ0 Reset value	Reserved	DPID 0 0	0	0	0 0	PKT	CNT 0		0	0	0	XFRSIZ															0			
0x530	OTG_FS_HCTS IZ1 Reset value	Reserved	DPID 0 0	0	0	0 0	PKT(	CNT 0		0	0	0	0 0 0 0 0 0 0 0									RSI 0	Z 0	0	0	0	0	0	0	0	0
0x550	Reset value	Reserv	DPID 0 0	0	0	0 0	PKT(	O O		0	0	0	0	0	0	0	0	0	0	0		RSI 0	Z 0	0	0	0	0	0	0	0	0
0x570		Reserv	DPID 0 0	0	0	0 0	PKT(	O O		0	0	0	0	0		RSI 0	Z 0	0	0 0 0 0 0						0						
0x590	Reset value	d Reserv	DPID 0 0	0	0	0 0	PKT(	0	0	0	0	0	0	0	0	0	0	0	0	0	0		0	0	0	0	0	0	0	0	
0x5B0	IZ5 Reset value	Reserv	DPID 0 0	0	0	0 0	PKT(	0		0	0	0	0	0	0	0	0	0	0	0		RSI 0	0	0	0	0	0	0	0	0	0
0x5D0		le L	DPID 0 0	0	0	0 0	PKT(	CNT 0		0	0	0	0	0	0	0	0	0	0	0		RSI 0	Z 0	0	0	0	0	0	0	0	0
0x5F0	OTG_FS_HCTS IZ7 Reset value	e –	DPID 0 0	0	0	0 0	РКТ(	CNT 0		0	0	0	0	0	0	0	0	0	0	0		RSI 0	Z 0	0	0	0	0	0	0	0	0
0x800	OTG_FS_DCFG Reset value		成   1   1   1   1   1   1   1   1   1																O PEIVI		0	0	0	o DAD	0	0	0	Reserved	ONZLSOHSK	OSPD	

Table 159. OTG\_FS register map and reset values (continued)

Offset	Register	2224282383	12   20   14   14   14   14   14   14   14   1	1 2 1 3 1 5 1 5	0 1 2 3 4 5 6 7 8 9 10												
0x804	OTG_FS_DCTL		erved	POPRGDNE	CGONAK SGONAK CGINAK SGINAK TCTL GONSTS GINSTS SDIS RWUSIG												
	Reset value		T	0	0 0 0 0 0 0 0 0 0 0												
0x808	OTG_FS_DSTS	Reserved		FNSOF Reserved B B Q S S S S S S S S S S S S S S S S S													
	Reset value		0 0 0 0 0 0	0 0 0 0 0	0 0 0 0												
0x810	OTG_FS_DIEP MSK		Reserved	INEPNEM ITTYFEMSK TOM Reserved EPDM EPDM													
	Reset value			0 0 0 0 0													
0x814	OTG_FS_DOEP MSK Reset value		Reserved	o OTEPDM STUPM Reserved O KFRCM													
	OTG FS DAIN			IEPINT													
0x818	_т_	OEPINT															
	Reset value  OTG FS DAIN	0 0 0 0 0 0 0 0 0 0	0 0 0 0 0 0	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0													
0x81C	TMSK	OEPM		IEPM													
	Reset value	0 0 0 0 0 0 0 0 0 0	0 0 0 0 0 0	0 0 0 0 0	0 0 0 0 0 0 0 0 0 0 0 0												
0x828	OTG_FS_DVBU SDIS	Reserved			VBUSDT												
	Reset value			0 0 0 1 0	1 1 1 1 1 1 0 1 0 1 1 1												
0x82C	OTG_FS_DVBU SPULSE	Res	erved	DVBUSP													
	Reset value			0 1 0 1 1 0 1 1 1 0 0 0													
0x834	OTG_FS_DIEP EMPMSK	Reserved			INEPTXFEM												
	Reset value			0 0 0 0 0	0 0 0 0 0 0 0 0 0 0 0 0 0												
0x900	OTG_FS_DIEP CTL0	Peserved PENA NUMBER OF	Reserved  O A A O NAKSTS  Reserved	USBAEP	Reserved MPSI Z												
	TG_FS_DTXFS		101-1010101-	1													
0x918	TS0	Reserved			INEPTFSAV												
	Reset value			0 0 0 0 0	0 1 0 0 0 0 0 0 0 0 0 0												
0x920	OTG_FS_DIEP CTL1	EPENA EPDIS SODDFRM/SD1PID SD0PID/SEVNFRM SNAK CNAK TA	Stall Reserved EPTYP NAKSTS EONUM/DPID	Reserved	MPSIZ												
	Reset value	0 0 0 0 0 0 0 0 0 0	0 0 0 0	0	0 0 0 0 0 0 0 0 0 0 0 0												
0x938	TG_FS_DTXFS TS1	Reserved			INEPTFSAV												
OXOGO	Reset value	110001100		0 0 0 0 0	0 1 0 0 0 0 0 0 0 0 0 0												
0x940	OTG_FS_DIEP CTL2	O EPENA O SODDFRIM O SODDID/SEVNIFRIM O SNAK O CNAK O CNAK	Reserved  PROPING  PR	Reserved	MPSIZ												
	TG_FS_DTXFS		1 0 0 0 0														
0x958	TS2	Reserved			INEPTFSAV												
	Reset value			0 0 0 0 0	0 1 0 0 0 0 0 0 0 0 0												

Table 159. OTG\_FS register map and reset values (continued)

Offset	Register							io to lo o				_				0		1						
Oliset	negister	31	3	29	<b>1</b> 28	2	26	22 24 22 22 23	5	7	19	17	16	15	4 5 2 1	9 8	7 9 4	ю <sub>С</sub>	0					
0x960	OTG_FS_DIEP CTL3	EPENA	EPDIS	SODDFRM	SDOPID/SEVNFRM		CNAK	TXFNUM	Stall	Reserved		NAKSTS	EONUM/DPID	USBAEP	Reserved	MPSIZ								
	Reset value TG_FS_DTXFS													0 0 0 0	0 0	0 0								
0x978	TS3	Reserved													TFSAV									
	Reset value												0 0 0 0	0 0	0 0									
0xB00		EPENA	EPDIS	Bospryod	nexel ved	SNAK	CNAK	Reserved	Stall	SNPM	EPTY P	NAKSTS	Reserved	¬ USBAEP		Reserv		MPSI Z						
	Reset value	0	0	П		0	0		0	0	0 0	0	Н	1					0 0					
0xB20	OTG_FS_DOEP CTL1	EPENA	EPDIS	SODDFRM	SD0PID/SEVNFRM	SNAK	CNAK	Reserved	Stall	SNPM	ЕРТҮР	NAKSTS	EONUM/DPID	USBAEP	Reserved									
	Reset value	0	0	0	0	0	0		0	0	0 0	0	0	0		0 0 0	0 0 0 0	0 0	0 0					
0xB40	OTG_FS_DOEP CTL2	EPENA	EPDIS	SODDFRM	SD0PID/SEVNFRM	SNAK	CNAK	Reserved	Stall	SNPM	ЕРТҮР	NAKSTS	EONUM/DPID	USBAEP	Reserved		MPSIZ							
	Reset value	000000000000000000000000000000000000000													0 0 0 0	0 0 0 0 0 0 0								
0xB60	OTG_FS_DOEP CTL3	MR A GG G														MPSIZ								
	Reset value	0	0	0	0	0	0		0	0	0 0	0	0	0		0 0 0	0 0 0 0	-	0 0					
0x908	OTG_FS_DIEPI NT0								F	Rese	erved						TXFE INEPNE Reserved	D TOC Reserved	EPDISD XFRC					
	Reset value															1 0 4 0	0 œ	0 0						
0x928	OTG_FS_DIEPI NT1								F	Rese	erved						TXFE INEPNE Seserved ITTXFE	TOC	EPDISD XFRC					
	Reset value																1 0 4 0	0 &	0 0					
0x948	OTG_FS_DIEPI NT2								F	Rese	erved						TXFE O INEPNE Reserved O ITTXFE	D TOC	EPDISD XFRC					
	Reset value		1															<del>                                     </del>	0 0					
0x968	OTG_FS_DIEPI NT3								F	Rese	erved						TXFE INEPNE Reserved	Reserved	EPDISD XFRC					
	Reset value																1 0 4 0		0 0					
0xB08	OTG_FS_DOEP INT0								F	Rese	erved						Reserved  BZBSTUP  Reserved  OTEPDIS	~	o EPDISD o XFRC					
0xB28	OTG_FS_DOEP INT1								F	Rese	erved						Reserved BZBSTUP Reserved OTEPDIS	STUP	XFRC o					
	Reset value																0 0	0 "	0 0					

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Table 159. OTG\_FS register map and reset values (continued)

Offset	Register	31	90	53	28	27	56	22	24	23	52	21	20	19	8	17	16	15	14	13	12	11	10	6	8	7	9	2	4	3	2	-	0	
0xB48	OTG_FS_DOEP INT2											F	lese	erve	d											Reserved	o B2BSTUP	Reserved	OTEPDIS	OSTUP	Reserved	o EPDISD	o XFRC	
0xB68	OTG_FS_DOEP INT3		Reserved													ed (													OTEPDIS	o STUP	Reserved	o EPDISD	o XFRC	
0x910	OTG_FS_DIEP TSIZ0 Reset value	Reserved PKTC NT 0   0											Reserved												XFRSIZ									
0x930	OTG_FS_DIEP TSIZ1												XFRSIZ													0 0 0 0 0 0 0 0								
	Reset value  OTG_FS_DIEP		0	0	0	0	0	0	0	0	0	0	0 0 0 0 0 0 0 0 0 0										0	0	0	0	0	0	0	0	0			
0x950	TSIZ2  Reset value	Reserved	MCI 0	NT 0	0	0	0		КТ 0	ONT 0	0	0	0	0	0	0	0	0	0	0	0	0	0	RSI 0	2 0	0	0	0	0	0	0	0	0	
0x970	OTG_FS_DIEP TSIZ3			MCNT PKTCNT											XFRSIZ												U	•	Ů	Ü	Ů	٠	_	
0,370	Reset value	Reserved	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
0xB10	OTG_FS_DOEP TSIZ0	Reserved	STU					Res	erv	ed				PKTCNT	Reserved											XFRSIZ								
	Reset value	Ä	0											0									0 0 0 0 0 0 0							0				
0xB30	OTG_FS_DOEP TSIZ1	Reserved	RXDPID/	SIUPCIN				Р	KT	CNT					XFRSIZ																			
	Reset value	ш	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
0xB50	OTG_FS_DOEP TSIZ2	Reserved	RXDPID/	SIOPON				Р	KT	CNT					XFRSIZ																			
	Reset value	Н	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
0xB70	OTG_FS_DOEP TSIZ3	STUP												XFRSIZ																				
	Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	자 0	0	
0xE00	OTG_FS_PCG CCTL		Res													served													<b>JSUSYHA</b>	Beserved		GATEHCLK	STPPCLK	
	Reset value																																	

Refer to Table 2 on page 50 for the register boundary addresses.

# 29.17 OTG\_FS programming model

# 29.17.1 Core initialization

The application must perform the core initialization sequence. If the cable is connected during power-up, the current mode of operation bit in the OTG\_FS\_GINTSTS (CMOD bit in OTG\_FS\_GINTSTS) reflects the mode. The OTG\_FS controller enters host mode when an "A" plug is connected or device mode when a "B" plug is connected.

This section explains the initialization of the OTG\_FS controller after power-on. The application must follow the initialization sequence irrespective of host or device mode operation. All core global registers are initialized according to the core's configuration:

- 1. Program the following fields in the OTG\_FS\_GAHBCFG register:
  - Global interrupt mask bit GINTMSK = 1
  - RxFIFO non-empty (RXFLVL bit in OTG\_FS\_GINTSTS)
  - Periodic TxFIFO empty level
- 2. Program the following fields in the OTG\_FS\_GUSBCFG register:
  - HNP capable bit
  - SRP capable bit
  - FS timeout calibration field
  - USB turnaround time field
- The software must unmask the following bits in the OTG\_FS\_GINTMSK register:OTG interrupt mask

Mode mismatch interrupt mask

 The software can read the CMOD bit in OTG\_FS\_GINTSTS to determine whether the OTG\_FS controller is operating in host or device mode.

## 29.17.2 Host initialization

To initialize the core as host, the application must perform the following steps:

- 1. Program the HPRTINT in the OTG\_FS\_GINTMSK register to unmask
- 2. Program the OTG\_FS\_HCFG register to select full-speed host
- 3. Program the PPWR bit in OTG\_FS\_HPRT to 1. This drives V<sub>BUS</sub> on the USB.
- 4. Wait for the PCDET interrupt in OTG\_FS\_HPRT0. This indicates that a device is connecting to the port.
- 5. Program the PRST bit in OTG FS HPRT to 1. This starts the reset process.
- 6. Wait at least 10 ms for the reset process to complete.
- 7. Program the PRST bit in OTG\_FS\_HPRT to 0.
- Wait for the PENCHNG interrupt in OTG\_FS\_HPRT.
- 9. Read the PSPD bit in OTG\_FS\_HPRT to get the enumerated speed.
- 10. Program the HFIR register with a value corresponding to the selected PHY clock 1
- 11. Program the FSLSPCS field in the OTG\_FS\_HCFG register following the speed of the device detected in step 9. If FSLSPCS has been changed a port reset must be performed.
- 12. Program the OTG\_FS\_GRXFSIZ register to select the size of the receive FIFO.
- 13. Program the OTG\_FS\_HNPTXFSIZ register to select the size and the start address of the Non-periodic transmit FIFO for non-periodic transactions.
- 14. Program the OTG\_FS\_HPTXFSIZ register to select the size and start address of the periodic transmit FIFO for periodic transactions.

To communicate with devices, the system software must initialize and enable at least one channel.

# 29.17.3 Device initialization

The application must perform the following steps to initialize the core as a device on powerup or after a mode change from host to device.

- 1. Program the following fields in the OTG FS DCFG register:
  - Device speed
  - Non-zero-length status OUT handshake
- 2. Program the OTG\_FS\_GINTMSK register to unmask the following interrupts:
  - USB reset
  - Enumeration done
  - Early suspend
  - USB suspend
  - SOF
- 3. Program the VBUSBSEN bit in the OTG\_FS\_GCCFG register to enable V<sub>BUS</sub> sensing in "B" device mode and supply the 5 volts across the pull-up resistor on the DP line.
- 4. Wait for the USBRST interrupt in OTG\_FS\_GINTSTS. It indicates that a reset has been detected on the USB that lasts for about 10 ms on receiving this interrupt.

Wait for the ENUMDNE interrupt in OTG\_FS\_GINTSTS. This interrupt indicates the end of reset on the USB. On receiving this interrupt, the application must read the OTG\_FS\_DSTS

register to determine the enumeration speed and perform the steps listed in *Endpoint* initialization on enumeration completion on page 1057.

At this point, the device is ready to accept SOF packets and perform control transfers on control endpoint 0.

# 29.17.4 Host programming model

#### Channel initialization

The application must initialize one or more channels before it can communicate with connected devices. To initialize and enable a channel, the application must perform the following steps:

- Program the OTG\_FS\_GINTMSK register to unmask the following:
- 2. Channel interrupt
  - Non-periodic transmit FIFO empty for OUT transactions (applicable when operating in pipelined transaction-level with the packet count field programmed with more than one).
  - Non-periodic transmit FIFO half-empty for OUT transactions (applicable when operating in pipelined transaction-level with the packet count field programmed with more than one).
- 3. Program the OTG\_FS\_HAINTMSK register to unmask the selected channels' interrupts.
- 4. Program the OTG\_FS\_HCINTMSK register to unmask the transaction-related interrupts of interest given in the host channel interrupt register.
- Program the selected channel's OTG\_FS\_HCTSIZx register with the total transfer size, in bytes, and the expected number of packets, including short packets. The application must program the PID field with the initial data PID (to be used on the first OUT transaction or to be expected from the first IN transaction).
- 6. Program the OTG\_FS\_HCCHARx register of the selected channel with the device's endpoint characteristics, such as type, speed, direction, and so forth. (The channel can be enabled by setting the channel enable bit to 1 only when the application is ready to transmit or receive any packet).

# Halting a channel

The application can disable any channel by programming the OTG\_FS\_HCCHARx register with the CHDIS and CHENA bits set to 1. This enables the OTG\_FS host to flush the posted requests (if any) and generates a channel halted interrupt. The application must wait for the CHH interrupt in OTG\_FS\_HCINTx before reallocating the channel for other transactions. The OTG\_FS host does not interrupt the transaction that has already been started on the USB.

Before disabling a channel, the application must ensure that there is at least one free space available in the non-periodic request queue (when disabling a non-periodic channel) or the periodic request queue (when disabling a periodic channel). The application can simply flush the posted requests when the Request queue is full (before disabling the channel), by programming the OTG\_FS\_HCCHARx register with the CHDIS bit set to 1, and the CHENA bit cleared to 0.

The application is expected to disable a channel on any of the following conditions:

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- When an STALL, TXERR, BBERR or DTERR interrupt in OTG FS HCINTx is received for an IN or OUT channel. The application must be able to receive other interrupts (DTERR, Nak, Data, TXERR) for the same channel before receiving the halt.
- When a DISCINT (Disconnect Device) interrupt in OTG FS GINTSTS is received. (The application is expected to disable all enabled channels).
- 3. When the application aborts a transfer before normal completion.

# **Operational model**

The application must initialize a channel before communicating to the connected device. This section explains the sequence of operation to be performed for different types of USB transactions.

# Writing the transmit FIFO

The OTG\_FS host automatically writes an entry (OUT request) to the periodic/nonperiodic request queue, along with the last DWORD write of a packet. The application must ensure that at least one free space is available in the periodic/non-periodic request queue before starting to write to the transmit FIFO. The application must always write to the transmit FIFO in DWORDs. If the packet size is non-DWORD aligned, the application must use padding. The OTG\_FS host determines the actual packet size based on the programmed maximum packet size and transfer size.

Start Read GNPTXSTS/HPTXFSIZ registers for available FIFO and queue spaces Wait for NPTXFE/PTXFE 1 MPS of interrupt in LPS FIFO space OTG\_FS\_GINTSTS Yes available? Yes Write 1 packet data to transmit FIFO More packets to send? Νo MPS: Maximum packet size Done

Figure 360. Transmit FIFO write task

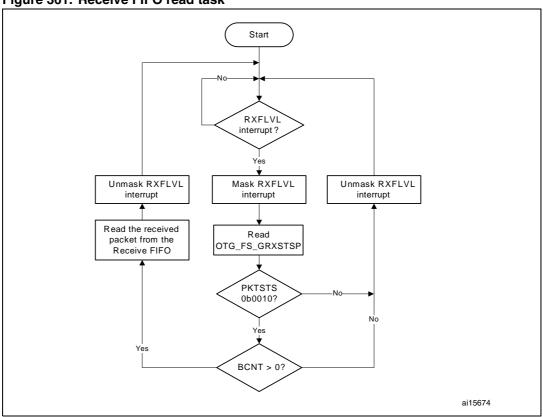
1041/1340 Doc ID 018909 Rev 2

LPS: Last packet size

# Reading the receive FIFO

The application must ignore all packet statuses other than IN data packet (bx0010).

Figure 361. Receive FIFO read task



# Bulk and control OUT/SETUP transactions

A typical bulk or control OUT/SETUP pipelined transaction-level operation is shown in *Figure 362*. See channel 1 (ch\_1). Two bulk OUT packets are transmitted. A control

SETUP transaction operates in the same way but has only one packet. The assumptions are:

- The application is attempting to send two maximum-packet-size packets (transfer size = 1, 024 bytes).
- The non-periodic transmit FIFO can hold two packets (128 bytes for FS).
- The non-periodic request queue depth = 4.

# Normal bulk and control OUT/SETUP operations

The sequence of operations in (channel 1) is as follows:

- a) Initialize channel 1
- b) Write the first packet for channel 1
- c) Along with the last Word write, the core writes an entry to the non-periodic request queue
- d) As soon as the non-periodic queue becomes non-empty, the core attempts to send an OUT token in the current frame
- e) Write the second (last) packet for channel 1
- f) The core generates the XFRC interrupt as soon as the last transaction is completed successfully
- g) In response to the XFRC interrupt, de-allocate the channel for other transfers
- h) Handling non-ACK responses

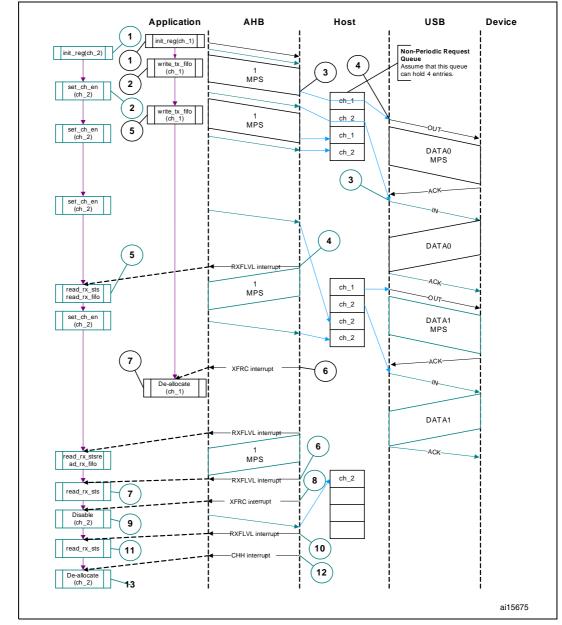


Figure 362. Normal bulk/control OUT/SETUP and bulk/control IN transactions

The channel-specific interrupt service routine for bulk and control OUT/SETUP transactions is shown in the following code samples.

- Interrupt service routine for bulk/control OUT/SETUP and bulk/control IN transactions
  - a) Bulk/Control OUT/SETUP

```
Unmask (NAK/TXERR/STALL/XFRC)
if (XFRC)
    {
    Reset Error Count
    Mask ACK
```

```
De-allocate Channel
  }
else if (STALL)
  Transfer Done = 1
  Unmask CHH
  Disable Channel
else if (NAK or TXERR )
  Rewind Buffer Pointers
  Unmask CHH
  Disable Channel
  if (TXERR)
    Increment Error Count
    Unmask ACK
  else
    Reset Error Count
else if (CHH)
  {
  Mask CHH
  if (Transfer Done or (Error_count == 3))
    De-allocate Channel
  else
    Re-initialize Channel
  else if (ACK)
    Reset Error Count
    Mask ACK
```

The application is expected to write the data packets into the transmit FIFO as and when the space is available in the transmit FIFO and the Request queue. The application can make use of the NPTXFE interrupt in OTG\_FS\_GINTSTS to find the transmit FIFO space.

# b) Bulk/Control IN

```
Unmask (TXERR/XFRC/BBERR/STALL/DTERR)
if (XFRC)
{
  Reset Error Count
   Unmask CHH
  Disable Channel
```

```
Reset Error Count
  Mask ACK
  }
else if (TXERR or BBERR or STALL)
  Unmask CHH
  Disable Channel
  if (TXERR)
    Increment Error Count
    Unmask ACK
else if (CHH)
  Mask CHH
  if (Transfer Done or (Error count == 3))
    De-allocate Channel
  else
    Re-initialize Channel
else if (ACK)
  Reset Error Count
  Mask ACK
  }
else if (DTERR)
  Reset Error Count
```

The application is expected to write the requests as and when the Request queue space is available and until the XFRC interrupt is received.

#### Bulk and control IN transactions

A typical bulk or control IN pipelined transaction-level operation is shown in *Figure 363*. See channel 2 (ch\_2). The assumptions are:

- The application is attempting to receive two maximum-packet-size packets (transfer size = 1 024 bytes).
- The receive FIFO can contain at least one maximum-packet-size packet and two status Words per packet (72 bytes for FS).
- The non-periodic request queue depth = 4.

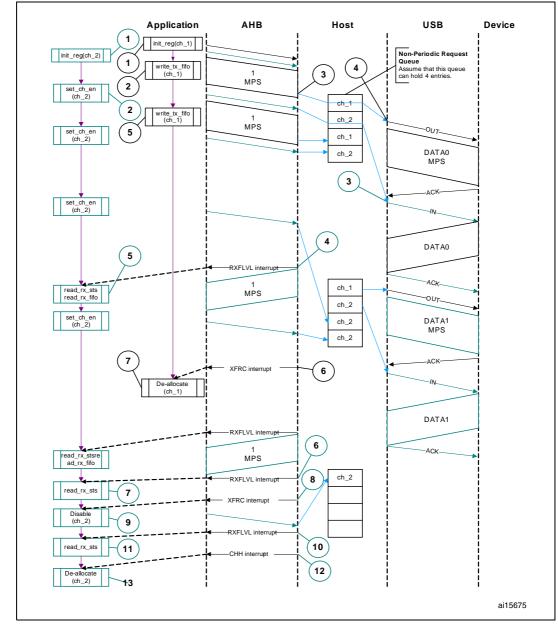


Figure 363. Bulk/control IN transactions

The sequence of operations is as follows:

- a) Initialize channel 2.
- b) Set the CHENA bit in HCCHAR2 to write an IN request to the non-periodic request queue.
- c) The core attempts to send an IN token after completing the current OUT transaction.
- d) The core generates an RXFLVL interrupt as soon as the received packet is written to the receive FIFO.
- e) In response to the RXFLVL interrupt, mask the RXFLVL interrupt and read the received packet status to determine the number of bytes received, then read the receive FIFO accordingly. Following this, unmask the RXFLVL interrupt.



- f) The core generates the RXFLVL interrupt for the transfer completion status entry in the receive FIFO.
- g) The application must read and ignore the receive packet status when the receive packet status is not an IN data packet (PKTSTS in GRXSTSR ≠ 0b0010).
- h) The core generates the XFRC interrupt as soon as the receive packet status is read.
- In response to the XFRC interrupt, disable the channel and stop writing the OTG\_FS\_HCCHAR2 register for further requests. The core writes a channel disable request to the non-periodic request queue as soon as the OTG\_FS\_HCCHAR2 register is written.
- j) The core generates the RXFLVL interrupt as soon as the halt status is written to the receive FIFO.
- k) Read and ignore the receive packet status.
- The core generates a CHH interrupt as soon as the halt status is popped from the receive FIFO.
- m) In response to the CHH interrupt, de-allocate the channel for other transfers.
- n) Handling non-ACK responses

#### Control transactions

Setup, Data, and Status stages of a control transfer must be performed as three separate transfers. Setup-, Data- or Status-stage OUT transactions are performed similarly to the bulk OUT transactions explained previously. Data- or Status-stage IN transactions are performed similarly to the bulk IN transactions explained previously. For all three stages, the application is expected to set the EPTYP field in OTG\_FS\_HCCHAR1 to Control. During the Setup stage, the application is expected to set the PID field in OTG\_FS\_HCTSIZ1 to SETUP.

### Interrupt OUT transactions

A typical interrupt OUT operation is shown in Figure 364. The assumptions are:

- The application is attempting to send one packet in every frame (up to 1 maximum packet size), starting with the odd frame (transfer size = 1 024 bytes)
- The periodic transmit FIFO can hold one packet (1 KB)
- Periodic request queue depth = 4

The sequence of operations is as follows:

- a) Initialize and enable channel 1. The application must set the ODDFRM bit in OTG\_FS\_HCCHAR1.
- b) Write the first packet for channel 1.
- c) Along with the last Word write of each packet, the OTG\_FS host writes an entry to the periodic request queue.
- d) The OTG\_FS host attempts to send an OUT token in the next (odd) frame.
- e) The OTG\_FS host generates an XFRC interrupt as soon as the last packet is transmitted successfully.
- f) In response to the XFRC interrupt, reinitialize the channel for the next transfer.

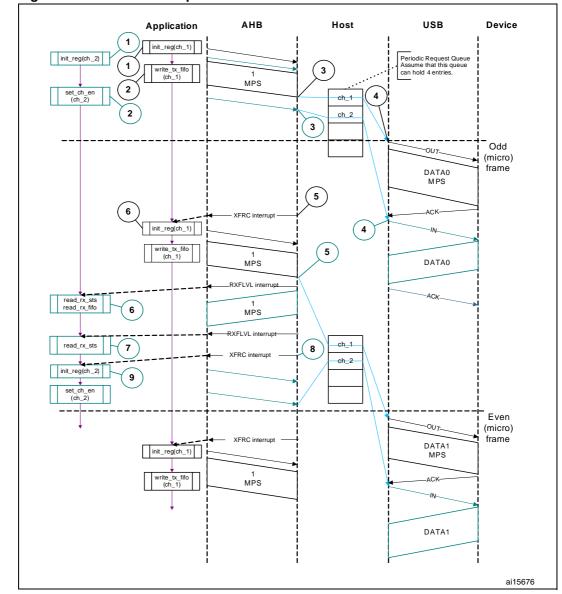


Figure 364. Normal interrupt OUT/IN transactions

# • Interrupt service routine for interrupt OUT/IN transactions

a) Interrupt OUT

```
Unmask (NAK/TXERR/STALL/XFRC/FRMOR)
if (XFRC)
    {
    Reset Error Count
    Mask ACK
    De-allocate Channel
    }
else
    if (STALL or FRMOR)
        {
        Mask ACK
        Unmask CHH
    }
}
```

```
Disable Channel
  if (STALL)
     {
    Transfer Done = 1
  }
else
  if (NAK or TXERR)
    Rewind Buffer Pointers
    Reset Error Count
    Mask ACK
    Unmask CHH
    Disable Channel
  else
    if (CHH)
    Mask CHH
     if (Transfer Done or (Error count == 3))
       De-allocate Channel
     else
       Re-initialize Channel (in next b_interval - 1 Frame)
  else
    if (ACK)
    Reset Error Count
    Mask ACK
 The application uses the NPTXFE interrupt in OTG_FS_GINTSTS to find the transmit
```

FIFO space.

## b) Interrupt IN

```
Unmask (NAK/TXERR/XFRC/BBERR/STALL/FRMOR/DTERR)
if (XFRC)
  Reset Error Count
  Mask ACK
  if (OTG_FS_HCTSIZx.PKTCNT == 0)
    De-allocate Channel
    }
  else
    Transfer Done = 1
    Unmask CHH
    Disable Channel
```

```
}
else
  if (STALL or FRMOR or NAK or DTERR or BBERR)
    Mask ACK
    Unmask CHH
    Disable Channel
    if (STALL or BBERR)
       Reset Error Count
       Transfer Done = 1
     else
       if (!FRMOR)
       Reset Error Count
else
  if (TXERR)
    Increment Error Count
    Unmask ACK
    Unmask CHH
    Disable Channel
else
  if (CHH)
    Mask CHH
     if (Transfer Done or (Error count == 3))
      De-allocate Channel
     else
      Re-initialize Channel (in next b_interval - 1 /Frame)
  }
else
  if (ACK)
    Reset Error Count
    Mask ACK
     }
```

## Interrupt IN transactions

The assumptions are:

- The application is attempting to receive one packet (up to 1 maximum packet size) in every frame, starting with odd (transfer size = 1 024 bytes).
- The receive FIFO can hold at least one maximum-packet-size packet and two status Words per packet (1 031 bytes).
- Periodic request queue depth = 4.

### Normal interrupt IN operation

The sequence of operations is as follows:

- a) Initialize channel 2. The application must set the ODDFRM bit in OTG\_FS\_HCCHAR2.
- b) Set the CHENA bit in OTG\_FS\_HCCHAR2 to write an IN request to the periodic request queue.
- c) The OTG\_FS host writes an IN request to the periodic request queue for each OTG\_FS\_HCCHAR2 register write with the CHENA bit set.
- d) The OTG\_FS host attempts to send an IN token in the next (odd) frame.
- e) As soon as the IN packet is received and written to the receive FIFO, the OTG\_FS host generates an RXFLVL interrupt.
- f) In response to the RXFLVL interrupt, read the received packet status to determine the number of bytes received, then read the receive FIFO accordingly. The application must mask the RXFLVL interrupt before reading the receive FIFO, and unmask after reading the entire packet.
- g) The core generates the RXFLVL interrupt for the transfer completion status entry in the receive FIFO. The application must read and ignore the receive packet status when the receive packet status is not an IN data packet (PKTSTS in GRXSTSR ≠ 0b0010).
- h) The core generates an XFRC interrupt as soon as the receive packet status is
- i) In response to the XFRC interrupt, read the PKTCNT field in OTG\_FS\_HCTSIZ2. If the PKTCNT bit in OTG\_FS\_HCTSIZ2 is not equal to 0, disable the channel before re-initializing the channel for the next transfer, if any). If PKTCNT bit in OTG\_FS\_HCTSIZ2 = 0, reinitialize the channel for the next transfer. This time, the application must reset the ODDFRM bit in OTG\_FS\_HCCHAR2.

### Isochronous OUT transactions

A typical isochronous OUT operation is shown in *Figure 365*. The assumptions are:

- The application is attempting to send one packet every frame (up to 1 maximum packet size), starting with an odd frame. (transfer size = 1 024 bytes).
- The periodic transmit FIFO can hold one packet (1 KB).
- Periodic request queue depth = 4.

The sequence of operations is as follows:

- Initialize and enable channel 1. The application must set the ODDFRM bit in OTG\_FS\_HCCHAR1.
- b) Write the first packet for channel 1.
- c) Along with the last Word write of each packet, the OTG\_FS host writes an entry to the periodic request queue.
- d) The OTG\_FS host attempts to send the OUT token in the next frame (odd).
- e) The OTG\_FS host generates the XFRC interrupt as soon as the last packet is transmitted successfully.
- f) In response to the XFRC interrupt, reinitialize the channel for the next transfer.
- g) Handling non-ACK responses

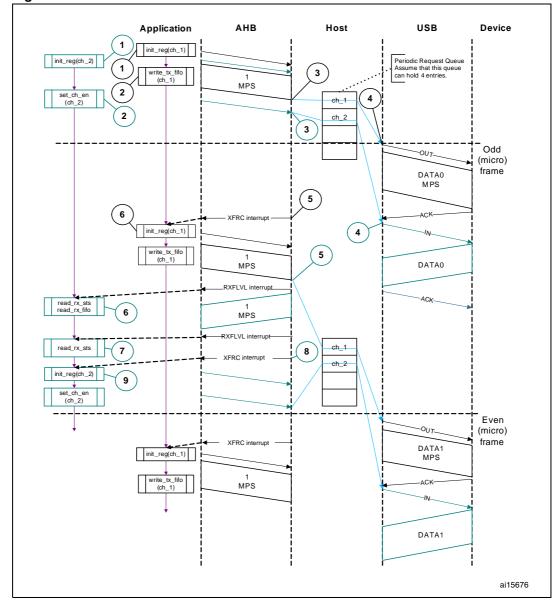


Figure 365. Normal isochronous OUT/IN transactions

Interrupt service routine for isochronous OUT/IN transactions

Code sample: Isochronous OUT

```
Unmask (FRMOR/XFRC)
if (XFRC)
  {
  De-allocate Channel
  }
else
  if (FRMOR)
    {
    Unmask CHH
    Disable Channel
  }
```

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```
else
  if (CHH)
    {
    Mask CHH
    De-allocate Channel
   Code sample: Isochronous IN
Unmask (TXERR/XFRC/FRMOR/BBERR)
if (XFRC or FRMOR)
  if (XFRC and (OTG_FS_HCTSIZx.PKTCNT == 0))
    Reset Error Count
    De-allocate Channel
    }
  else
    Unmask CHH
    Disable Channel
  }
  else
    if (TXERR or BBERR)
       Increment Error Count
       Unmask CHH
       Disable Channel
       }
    else
       if (CHH)
         Mask CHH
         if (Transfer Done or (Error_count == 3))
           De-allocate Channel
           }
         else
            Re-initialize Channel
```

#### Isochronous IN transactions

The assumptions are:

- The application is attempting to receive one packet (up to 1 maximum packet size) in every frame starting with the next odd frame (transfer size = 1 024 bytes).
- The receive FIFO can hold at least one maximum-packet-size packet and two status Word per packet (1 031 bytes).
- Periodic request queue depth = 4.

The sequence of operations is as follows:

- a) Initialize channel 2. The application must set the ODDFRM bit in OTG FS HCCHAR2.
- b) Set the CHENA bit in OTG\_FS\_HCCHAR2 to write an IN request to the periodic request queue.
- c) The OTG\_FS host writes an IN request to the periodic request queue for each OTG\_FS\_HCCHAR2 register write with the CHENA bit set.
- d) The OTG\_FS host attempts to send an IN token in the next odd frame.
- e) As soon as the IN packet is received and written to the receive FIFO, the OTG\_FS host generates an RXFLVL interrupt.
- f) In response to the RXFLVL interrupt, read the received packet status to determine the number of bytes received, then read the receive FIFO accordingly. The application must mask the RXFLVL interrupt before reading the receive FIFO, and unmask it after reading the entire packet.
- g) The core generates an RXFLVL interrupt for the transfer completion status entry in the receive FIFO. This time, the application must read and ignore the receive packet status when the receive packet status is not an IN data packet (PKTSTS bit in OTG\_FS\_GRXSTSR ≠ 0b0010).
- h) The core generates an XFRC interrupt as soon as the receive packet status is read.
- i) In response to the XFRC interrupt, read the PKTCNT field in OTG\_FS\_HCTSIZ2. If PKTCNT≠ 0 in OTG\_FS\_HCTSIZ2, disable the channel before re-initializing the channel for the next transfer, if any. If PKTCNT = 0 in OTG\_FS\_HCTSIZ2, reinitialize the channel for the next transfer. This time, the application must reset the ODDFRM bit in OTG\_FS\_HCCHAR2.

## Selecting the queue depth

Choose the periodic and non-periodic request queue depths carefully to match the number of periodic/non-periodic endpoints accessed.

The non-periodic request queue depth affects the performance of non-periodic transfers. The deeper the queue (along with sufficient FIFO size), the more often the core is able to pipeline non-periodic transfers. If the queue size is small, the core is able to put in new requests only when the queue space is freed up.

The core's periodic request queue depth is critical to perform periodic transfers as scheduled. Select the periodic queue depth, based on the number of periodic transfers scheduled in a microframe. If the periodic request queue depth is smaller than the periodic transfers scheduled in a microframe, a frame overrun condition occurs.

## Handling babble conditions

OTG\_FS controller handles two cases of babble: packet babble and port babble. Packet babble occurs if the device sends more data than the maximum packet size for

the channel. Port babble occurs if the core continues to receive data from the device at EOF2 (the end of frame 2, which is very close to SOF).

When OTG\_FS controller detects a packet babble, it stops writing data into the Rx buffer and waits for the end of packet (EOP). When it detects an EOP, it flushes already written data in the Rx buffer and generates a Babble interrupt to the application.

When OTG\_FS controller detects a port babble, it flushes the RxFIFO and disables the port. The core then generates a Port disabled interrupt (HPRTINT in OTG\_FS\_GINTSTS, PENCHNG in OTG\_FS\_HPRT). On receiving this interrupt, the application must determine that this is not due to an overcurrent condition (another cause of the Port Disabled interrupt) by checking POCA in OTG\_FS\_HPRT, then perform a soft reset. The core does not send any more tokens after it has detected a port babble condition.

# 29.17.5 Device programming model

# **Endpoint initialization on USB reset**

- 1. Set the NAK bit for all OUT endpoints
  - SNAK = 1 in OTG\_FS\_DOEPCTLx (for all OUT endpoints)
- 2. Unmask the following interrupt bits
  - INEP0 = 1 in OTG\_FS\_DAINTMSK (control 0 IN endpoint)
  - OUTEP0 = 1 in OTG\_FS\_DAINTMSK (control 0 OUT endpoint)
  - STUP = 1 in DOEPMSK
  - XFRC = 1 in DOEPMSK
  - XFRC = 1 in DIEPMSK
  - TOC = 1 in DIEPMSK
- 3. Set up the Data FIFO RAM for each of the FIFOs
  - Program the OTG\_FS\_GRXFSIZ register, to be able to receive control OUT data and setup data. If thresholding is not enabled, at a minimum, this must be equal to 1 max packet size of control endpoint 0 + 2 Words (for the status of the control OUT data packet) + 10 Words (for setup packets).
  - Program the OTG\_FS\_TX0FSIZ register (depending on the FIFO number chosen) to be able to transmit control IN data. At a minimum, this must be equal to 1 max packet size of control endpoint 0.
- 4. Program the following fields in the endpoint-specific registers for control OUT endpoint 0 to receive a SETUP packet
  - STUPCNT = 3 in OTG\_FS\_DOEPTSIZ0 (to receive up to 3 back-to-back SETUP packets)

At this point, all initialization required to receive SETUP packets is done.

### **Endpoint initialization on enumeration completion**

- On the Enumeration Done interrupt (ENUMDNE in OTG\_FS\_GINTSTS), read the OTG\_FS\_DSTS register to determine the enumeration speed.
- Program the MPSIZ field in OTG\_FS\_DIEPCTL0 to set the maximum packet size. This
  step configures control endpoint 0. The maximum packet size for a control endpoint
  depends on the enumeration speed.

At this point, the device is ready to receive SOF packets and is configured to perform control transfers on control endpoint 0.

## **Endpoint initialization on SetAddress command**

This section describes what the application must do when it receives a SetAddress command in a SETUP packet.

- Program the OTG\_FS\_DCFG register with the device address received in the SetAddress command
- 1. Program the core to send out a status IN packet

# Endpoint initialization on SetConfiguration/SetInterface command

This section describes what the application must do when it receives a SetConfiguration or SetInterface command in a SETUP packet.

- When a SetConfiguration command is received, the application must program the endpoint registers to configure them with the characteristics of the valid endpoints in the new configuration.
- 2. When a SetInterface command is received, the application must program the endpoint registers of the endpoints affected by this command.
- 3. Some endpoints that were active in the prior configuration or alternate setting are not valid in the new configuration or alternate setting. These invalid endpoints must be deactivated.
- 4. Unmask the interrupt for each active endpoint and mask the interrupts for all inactive endpoints in the OTG\_FS\_DAINTMSK register.
- 5. Set up the Data FIFO RAM for each FIFO.
- 6. After all required endpoints are configured; the application must program the core to send a status IN packet.

At this point, the device core is configured to receive and transmit any type of data packet.

# **Endpoint activation**

This section describes the steps required to activate a device endpoint or to configure an existing device endpoint to a new type.

- Program the characteristics of the required endpoint into the following fields of the OTG\_FS\_DIEPCTLx register (for IN or bidirectional endpoints) or the OTG\_FS\_DOEPCTLx register (for OUT or bidirectional endpoints).
  - Maximum packet size
  - USB active endpoint = 1
  - Endpoint start data toggle (for interrupt and bulk endpoints)
  - Endpoint type
  - TxFIFO number
- Once the endpoint is activated, the core starts decoding the tokens addressed to that endpoint and sends out a valid handshake for each valid token received for the endpoint.

# **Endpoint deactivation**

This section describes the steps required to deactivate an existing endpoint.

- In the endpoint to be deactivated, clear the USB active endpoint bit in the OTG\_FS\_DIEPCTLx register (for IN or bidirectional endpoints) or the OTG\_FS\_DOEPCTLx register (for OUT or bidirectional endpoints).
- 2. Once the endpoint is deactivated, the core ignores tokens addressed to that endpoint, which results in a timeout on the USB.

Note:

The application must meet the following conditions to set up the device core to handle traffic: NPTXFEM and RXFLVLM in the OTG FS GINTMSK register must be cleared.

# 29.17.6 Operational model

## **SETUP and OUT data transfers**

This section describes the internal data flow and application-level operations during data OUT transfers and SETUP transactions.

#### Packet read

This section describes how to read packets (OUT data and SETUP packets) from the receive FIFO.

- 1. On catching an RXFLVL interrupt (OTG\_FS\_GINTSTS register), the application must read the Receive status pop register (OTG\_FS\_GRXSTSP).
- 2. The application can mask the RXFLVL interrupt (in OTG\_FS\_GINTSTS) by writing to RXFLVL = 0 (in OTG\_FS\_GINTMSK), until it has read the packet from the receive FIFO.
- 3. If the received packet's byte count is not 0, the byte count amount of data is popped from the receive Data FIFO and stored in memory. If the received packet byte count is 0, no data is popped from the receive data FIFO.
- 4. The receive FIFO's packet status readout indicates one of the following:
  - a) Global OUT NAK pattern:
     PKTSTS = Global OUT NAK, BCNT = 0x000, EPNUM = Don't Care (0x0),
     DPID = Don't Care (0b00).
     These data indicate that the global OUT NAK bit has taken effect.
  - b) SETUP packet pattern:

    PKTSTS = SETUP, BCNT = 0x008, EPNUM = Control EP Num, DPID = D0.

    These data indicate that a SETUP packet for the specified endpoint is now available for reading from the receive FIFO.
  - c) Setup stage done pattern:
    PKTSTS = Setup Stage Done, BCNT = 0x0, EPNUM = Control EP Num,
    DPID = Don't Care (0b00).
    These data indicate that the Setup stage for the specified endpoint has completed
    and the Data stage has started. After this entry is popped from the receive FIFO,
    the core asserts a Setup interrupt on the specified control OUT endpoint.
  - d) Data OUT packet pattern:
     PKTSTS = DataOUT, BCNT = size of the received data OUT packet (0 ≤ BCNT ≤ 1 024), EPNUM = EPNUM on which the packet was received, DPID = Actual Data PID.
  - e) Data transfer completed pattern:

    PKTSTS = Data OUT Transfer Done, BCNT = 0x0, EPNUM = OUT EP Num

    on which the data transfer is complete, DPID = Don't Care (0b00).

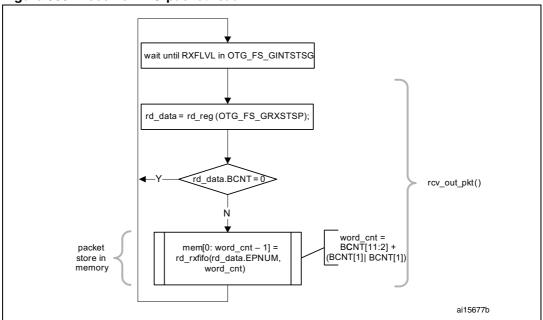
    These data indicate that an OUT data transfer for the specified OUT endpoint has

completed. After this entry is popped from the receive FIFO, the core asserts a Transfer Completed interrupt on the specified OUT endpoint.

- 5. After the data payload is popped from the receive FIFO, the RXFLVL interrupt (OTG FS GINTSTS) must be unmasked.
- 6. Steps 1–5 are repeated every time the application detects assertion of the interrupt line due to RXFLVL in OTG\_FS\_GINTSTS. Reading an empty receive FIFO can result in undefined core behavior.

Figure 366 provides a flowchart of the above procedure.

Figure 366. Receive FIFO packet read



## SETUP transactions

This section describes how the core handles SETUP packets and the application's sequence for handling SETUP transactions.

# Application requirements

1. To receive a SETUP packet, the STUPCNT field (OTG\_FS\_DOEPTSIZx) in a control OUT endpoint must be programmed to a non-zero value. When the application programs the STUPCNT field to a non-zero value, the core receives SETUP packets and writes them to the receive FIFO, irrespective of the NAK status and EPENA bit setting in OTG\_FS\_DOEPCTLx. The STUPCNT field is decremented every time the control endpoint receives a SETUP packet. If the STUPCNT field is not programmed to a proper value before receiving a SETUP packet, the core still receives the SETUP packet and decrements the STUPCNT field, but the application may not be able to

determine the correct number of SETUP packets received in the Setup stage of a control transfer.

- STUPCNT = 3 in OTG FS DOEPTSIZx
- 2. The application must always allocate some extra space in the Receive data FIFO, to be able to receive up to three SETUP packets on a control endpoint.
  - The space to be reserved is 10 Words. Three Words are required for the first SETUP packet, 1 Word is required for the Setup stage done Word and 6 Words are required to store two extra SETUP packets among all control endpoints.
  - 3 Words per SETUP packet are required to store 8 bytes of SETUP data and 4 bytes of SETUP status (Setup packet pattern). The core reserves this space in the receive data.
  - FIFO to write SETUP data only, and never uses this space for data packets.
- 3. The application must read the 2 Words of the SETUP packet from the receive FIFO.
- 4. The application must read and discard the Setup stage done Word from the receive FIFO.

### Internal data flow

- 5. When a SETUP packet is received, the core writes the received data to the receive FIFO, without checking for available space in the receive FIFO and irrespective of the endpoint's NAK and STALL bit settings.
  - The core internally sets the IN NAK and OUT NAK bits for the control IN/OUT endpoints on which the SETUP packet was received.
- 6. For every SETUP packet received on the USB, 3 Words of data are written to the receive FIFO, and the STUPCNT field is decremented by 1.
  - The first Word contains control information used internally by the core
  - The second Word contains the first 4 bytes of the SETUP command
  - The third Word contains the last 4 bytes of the SETUP command
- 7. When the Setup stage changes to a Data IN/OUT stage, the core writes an entry (Setup stage done Word) to the receive FIFO, indicating the completion of the Setup stage.
- 8. On the AHB side, SETUP packets are emptied by the application.
- 9. When the application pops the Setup stage done Word from the receive FIFO, the core interrupts the application with an STUP interrupt (OTG\_FS\_DOEPINTx), indicating it can process the received SETUP packet.
  - The core clears the endpoint enable bit for control OUT endpoints.

## Application programming sequence

- 1. Program the OTG FS DOEPTSIZx register.
  - STUPCNT = 3
- 2. Wait for the RXFLVL interrupt (OTG\_FS\_GINTSTS) and empty the data packets from the receive FIFO.
- Assertion of the STUP interrupt (OTG\_FS\_DOEPINTx) marks a successful completion of the SETUP Data Transfer.
  - On this interrupt, the application must read the OTG\_FS\_DOEPTSIZx register to determine the number of SETUP packets received and process the last received SETUP packet.

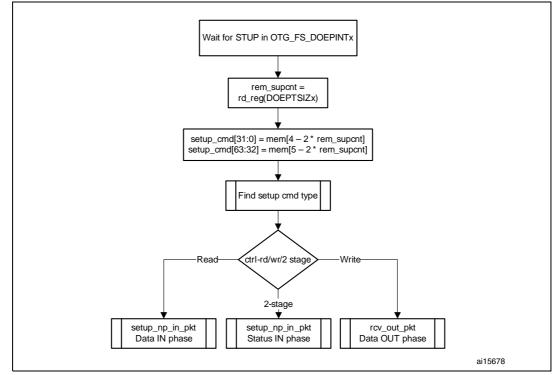


Figure 367. Processing a SETUP packet

## Handling more than three back-to-back SETUP packets

Per the USB 2.0 specification, normally, during a SETUP packet error, a host does not send more than three back-to-back SETUP packets to the same endpoint. However, the USB 2.0 specification does not limit the number of back-to-back SETUP packets a host can send to the same endpoint. When this condition occurs, the OTG\_FS controller generates an interrupt (B2BSTUP in OTG\_FS\_DOEPINTx).

## Setting the global OUT NAK

Internal data flow:

- When the application sets the Global OUT NAK (SGONAK bit in OTG\_FS\_DCTL), the
  core stops writing data, except SETUP packets, to the receive FIFO. Irrespective of the
  space availability in the receive FIFO, non-isochronous OUT tokens receive a NAK
  handshake response, and the core ignores isochronous OUT data packets
- 2. The core writes the Global OUT NAK pattern to the receive FIFO. The application must reserve enough receive FIFO space to write this data pattern.
- 3. When the application pops the Global OUT NAK pattern Word from the receive FIFO, the core sets the GONAKEFF interrupt (OTG\_FS\_GINTSTS).
- Once the application detects this interrupt, it can assume that the core is in Global OUT NAK mode. The application can clear this interrupt by clearing the SGONAK bit in OTG FS DCTL.

Application programming sequence

- 1. To stop receiving any kind of data in the receive FIFO, the application must set the Global OUT NAK bit by programming the following field:
  - SGONAK = 1 in OTG FS DCTL
- 2. Wait for the assertion of the GONAKEFF interrupt in OTG\_FS\_GINTSTS. When asserted, this interrupt indicates that the core has stopped receiving any type of data except SETUP packets.
- The application can receive valid OUT packets after it has set SGONAK in OTG\_FS\_DCTL and before the core asserts the GONAKEFF interrupt (OTG\_FS\_GINTSTS).
- 4. The application can temporarily mask this interrupt by writing to the GINAKEFFM bit in the OTG\_FS\_GINTMSK register.
  - GINAKEFFM = 0 in the OTG\_FS\_GINTMSK register
- Whenever the application is ready to exit the Global OUT NAK mode, it must clear the SGONAK bit in OTG\_FS\_DCTL. This also clears the GONAKEFF interrupt (OTG\_FS\_GINTSTS).
  - OTG\_FS\_DCTL = 1 in CGONAK
- 6. If the application has masked this interrupt earlier, it must be unmasked as follows:
  - GINAKEFFM = 1 in GINTMSK

### Disabling an OUT endpoint

The application must use this sequence to disable an OUT endpoint that it has enabled.

Application programming sequence:

- Before disabling any OUT endpoint, the application must enable Global OUT NAK mode in the core.
  - SGONAK = 1 in OTG\_FS\_DCTL
- 2. Wait for the GONAKEFF interrupt (OTG\_FS\_GINTSTS)
- 3. Disable the required OUT endpoint by programming the following fields:
  - EPDIS = 1 in OTG\_FS\_DOEPCTLx
  - SNAK = 1 in OTG\_FS\_DOEPCTLx
- 4. Wait for the EPDISD interrupt (OTG\_FS\_DOEPINTx), which indicates that the OUT endpoint is completely disabled. When the EPDISD interrupt is asserted, the core also clears the following bits:
  - EPDIS = 0 in OTG\_FS\_DOEPCTLx
  - EPENA = 0 in OTG\_FS\_DOEPCTLx
- The application must clear the Global OUT NAK bit to start receiving data from other non-disabled OUT endpoints.
  - SGONAK = 0 in OTG\_FS\_DCTL

#### Generic non-isochronous OUT data transfers

This section describes a regular non-isochronous OUT data transfer (control, bulk, or interrupt).

Application requirements:

- 1. Before setting up an OUT transfer, the application must allocate a buffer in the memory to accommodate all data to be received as part of the OUT transfer.
- 2. For OUT transfers, the transfer size field in the endpoint's transfer size register must be a multiple of the maximum packet size of the endpoint, adjusted to the Word boundary.
  - transfer size[EPNUM] =  $n \times (MPSIZ[EPNUM] + 4 (MPSIZ[EPNUM] \mod 4))$
  - packet count[EPNUM] = n
  - -n>0
- 3. On any OUT endpoint interrupt, the application must read the endpoint's transfer size register to calculate the size of the payload in the memory. The received payload size can be less than the programmed transfer size.
  - Payload size in memory = application programmed initial transfer size core updated final transfer size
  - Number of USB packets in which this payload was received = application programmed initial packet count – core updated final packet count

#### Internal data flow:

- The application must set the transfer size and packet count fields in the endpointspecific registers, clear the NAK bit, and enable the endpoint to receive the data.
- 2. Once the NAK bit is cleared, the core starts receiving data and writes it to the receive FIFO, as long as there is space in the receive FIFO. For every data packet received on the USB, the data packet and its status are written to the receive FIFO. Every packet (maximum packet size or short packet) written to the receive FIFO decrements the packet count field for that endpoint by 1.
  - OUT data packets received with bad data CRC are flushed from the receive FIFO automatically.
  - After sending an ACK for the packet on the USB, the core discards non-isochronous OUT data packets that the host, which cannot detect the ACK, resends. The application does not detect multiple back-to-back data OUT packets on the same endpoint with the same data PID. In this case the packet count is not decremented.
  - If there is no space in the receive FIFO, isochronous or non-isochronous data packets are ignored and not written to the receive FIFO. Additionally, nonisochronous OUT tokens receive a NAK handshake reply.
  - In all the above three cases, the packet count is not decremented because no data are written to the receive FIFO.
- 3. When the packet count becomes 0 or when a short packet is received on the endpoint, the NAK bit for that endpoint is set. Once the NAK bit is set, the isochronous or non-isochronous data packets are ignored and not written to the receive FIFO, and non-isochronous OUT tokens receive a NAK handshake reply.
- 4. After the data are written to the receive FIFO, the application reads the data from the receive FIFO and writes it to external memory, one packet at a time per endpoint.
- 5. At the end of every packet write on the AHB to external memory, the transfer size for the endpoint is decremented by the size of the written packet.

- 6. The OUT data transfer completed pattern for an OUT endpoint is written to the receive FIFO on one of the following conditions:
  - The transfer size is 0 and the packet count is 0
  - The last OUT data packet written to the receive FIFO is a short packet (0 ≤ packet size < maximum packet size)</li>
- 7. When either the application pops this entry (OUT data transfer completed), a transfer completed interrupt is generated for the endpoint and the endpoint enable is cleared.

### Application programming sequence:

- 1. Program the OTG\_FS\_DOEPTSIZx register for the transfer size and the corresponding packet count.
- 2. Program the OTG\_FS\_DOEPCTLx register with the endpoint characteristics, and set the EPENA and CNAK bits.
  - EPENA = 1 in OTG FS DOEPCTLx
  - CNAK = 1 in OTG\_FS\_DOEPCTLx
- 3. Wait for the RXFLVL interrupt (in OTG\_FS\_GINTSTS) and empty the data packets from the receive FIFO.
  - This step can be repeated many times, depending on the transfer size.
- 4. Asserting the XFRC interrupt (OTG\_FS\_DOEPINTx) marks a successful completion of the non-isochronous OUT data transfer.
- 5. Read the OTG\_FS\_DOEPTSIZx register to determine the size of the received data payload.

#### Generic isochronous OUT data transfer

This section describes a regular isochronous OUT data transfer.

## Application requirements:

- 1. All the application requirements for non-isochronous OUT data transfers also apply to isochronous OUT data transfers.
- For isochronous OUT data transfers, the transfer size and packet count fields must always be set to the number of maximum-packet-size packets that can be received in a single frame and no more. Isochronous OUT data transfers cannot span more than 1 frame.
- 3. The application must read all isochronous OUT data packets from the receive FIFO (data and status) before the end of the periodic frame (EOPF interrupt in OTG\_FS\_GINTSTS).
- 4. To receive data in the following frame, an isochronous OUT endpoint must be enabled after the EOPF (OTG\_FS\_GINTSTS) and before the SOF (OTG\_FS\_GINTSTS).

### Internal data flow:

- 1. The internal data flow for isochronous OUT endpoints is the same as that for non-isochronous OUT endpoints, but for a few differences.
- When an isochronous OUT endpoint is enabled by setting the Endpoint Enable and clearing the NAK bits, the Even/Odd frame bit must also be set appropriately. The core receives data on an isochronous OUT endpoint in a particular frame only if the following condition is met:
  - EONUM (in OTG\_FS\_DOEPCTLx) = SOFFN[0] (in OTG\_FS\_DSTS)
- 3. When the application completely reads an isochronous OUT data packet (data and status) from the receive FIFO, the core updates the RXDPID field in

OTG FS DOEPTSIZx with the data PID of the last isochronous OUT data packet read from the receive FIFO.

Application programming sequence:

- Program the OTG FS DOEPTSIZx register for the transfer size and the corresponding packet count
- 2. Program the OTG FS DOEPCTLx register with the endpoint characteristics and set the Endpoint Enable, ClearNAK, and Even/Odd frame bits.
  - EPENA = 1
  - CNAK = 1
  - EONUM = (0: Even/1: Odd)
- Wait for the RXFLVL interrupt (in OTG\_FS\_GINTSTS) and empty the data packets from the receive FIFO
  - This step can be repeated many times, depending on the transfer size.
- The assertion of the XFRC interrupt (in OTG FS DOEPINTx) marks the completion of 4. the isochronous OUT data transfer. This interrupt does not necessarily mean that the data in memory are good.
- This interrupt cannot always be detected for isochronous OUT transfers. Instead, the application can detect the IISOOXFRM interrupt in OTG\_FS\_GINTSTS.
- Read the OTG FS DOEPTSIZx register to determine the size of the received transfer and to determine the validity of the data received in the frame. The application must treat the data received in memory as valid only if one of the following conditions is met:
  - RXDPID = D0 (in OTG\_FS\_DOEPTSIZx) and the number of USB packets in which this payload was received = 1
  - RXDPID = D1 (in OTG\_FS\_DOEPTSIZx) and the number of USB packets in which this payload was received = 2
  - RXDPID = D2 (in OTG\_FS\_DOEPTSIZx) and the number of USB packets in which this payload was received = 3

The number of USB packets in which this payload was received = Application programmed initial packet count – Core updated final packet count

### Incomplete isochronous OUT data transfers

The application can discard invalid data packets.

This section describes the application programming sequence when isochronous OUT data packets are dropped inside the core.

### Internal data flow:

- For isochronous OUT endpoints, the XFRC interrupt (in OTG\_FS\_DOEPINTx) may not always be asserted. If the core drops isochronous OUT data packets, the application could fail to detect the XFRC interrupt (OTG\_FS\_DOEPINTx) under the following circumstances:
  - When the receive FIFO cannot accommodate the complete ISO OUT data packet, the core drops the received ISO OUT data
  - When the isochronous OUT data packet is received with CRC errors
  - When the isochronous OUT token received by the core is corrupted
  - When the application is very slow in reading the data from the receive FIFO
- When the core detects an end of periodic frame before transfer completion to all isochronous OUT endpoints, it asserts the incomplete Isochronous OUT data interrupt

(IISOOXFRM in OTG\_FS\_GINTSTS), indicating that an XFRC interrupt (in OTG\_FS\_DOEPINTx) is not asserted on at least one of the isochronous OUT endpoints. At this point, the endpoint with the incomplete transfer remains enabled, but no active transfers remain in progress on this endpoint on the USB.

Application programming sequence:

- Asserting the IISOOXFRM interrupt (OTG\_FS\_GINTSTS) indicates that in the current frame, at least one isochronous OUT endpoint has an incomplete transfer.
- 2. If this occurs because isochronous OUT data is not completely emptied from the endpoint, the application must ensure that the application empties all isochronous OUT data (data and status) from the receive FIFO before proceeding.
  - When all data are emptied from the receive FIFO, the application can detect the XFRC interrupt (OTG\_FS\_DOEPINTx). In this case, the application must reenable the endpoint to receive isochronous OUT data in the next frame.
- 3. When it receives an IISOOXFRM interrupt (in OTG\_FS\_GINTSTS), the application must read the control registers of all isochronous OUT endpoints (OTG\_FS\_DOEPCTLx) to determine which endpoints had an incomplete transfer in the current microframe. An endpoint transfer is incomplete if both the following conditions are met:
  - EONUM bit (in OTG\_FS\_DOEPCTLx) = SOFFN[0] (in OTG\_FS\_DSTS)
  - EPENA = 1 (in OTG FS DOEPCTLx)
- 4. The previous step must be performed before the SOF interrupt (in OTG\_FS\_GINTSTS) is detected, to ensure that the current frame number is not changed.
- For isochronous OUT endpoints with incomplete transfers, the application must discard the data in the memory and disable the endpoint by setting the EPDIS bit in OTG\_FS\_DOEPCTLx.
- 6. Wait for the EPDIS interrupt (in OTG\_FS\_DOEPINTx) and enable the endpoint to receive new data in the next frame.
  - Because the core can take some time to disable the endpoint, the application may not be able to receive the data in the next frame after receiving bad isochronous data.

## Stalling a non-isochronous OUT endpoint

This section describes how the application can stall a non-isochronous endpoint.

- 1. Put the core in the Global OUT NAK mode.
- 2. Disable the required endpoint
  - When disabling the endpoint, instead of setting the SNAK bit in OTG\_FS\_DOEPCTL, set STALL = 1 (in OTG\_FS\_DOEPCTL).
     The STALL bit always takes precedence over the NAK bit.
- 3. When the application is ready to end the STALL handshake for the endpoint, the STALL bit (in OTG\_FS\_DOEPCTLx) must be cleared.
- 4. If the application is setting or clearing a STALL for an endpoint due to a SetFeature.Endpoint Halt or ClearFeature.Endpoint Halt command, the STALL bit must be set or cleared before the application sets up the Status stage transfer on the control endpoint.

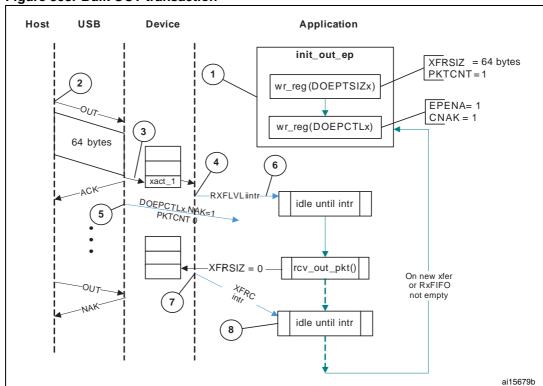
## **Examples**

This section describes and depicts some fundamental transfer types and scenarios.

Bulk OUT transaction

*Figure 368* depicts the reception of a single Bulk OUT Data packet from the USB to the AHB and describes the events involved in the process.

Figure 368. Bulk OUT transaction



After a SetConfiguration/SetInterface command, the application initializes all OUT endpoints by setting CNAK = 1 and EPENA = 1 (in OTG\_FS\_DOEPCTLx), and setting a suitable XFRSIZ and PKTCNT in the OTG\_FS\_DOEPTSIZx register.

- 1. host attempts to send data (OUT token) to an endpoint.
- 2. When the core receives the OUT token on the USB, it stores the packet in the RxFIFO because space is available there.
- 3. After writing the complete packet in the RxFIFO, the core then asserts the RXFLVL interrupt (in OTG\_FS\_GINTSTS).
- 4. On receiving the PKTCNT number of USB packets, the core internally sets the NAK bit for this endpoint to prevent it from receiving any more packets.
- 5. The application processes the interrupt and reads the data from the RxFIFO.
- 6. When the application has read all the data (equivalent to XFRSIZ), the core generates an XFRC interrupt (in OTG\_FS\_DOEPINTx).
- 7. The application processes the interrupt and uses the setting of the XFRC interrupt bit (in OTG\_FS\_DOEPINTx) to determine that the intended transfer is complete.

#### IN data transfers

#### Packet write

This section describes how the application writes data packets to the endpoint FIFO when dedicated transmit FIFOs are enabled.

- 1. The application can either choose the polling or the interrupt mode.
  - In polling mode, the application monitors the status of the endpoint transmit data FIFO by reading the OTG\_FS\_DTXFSTSx register, to determine if there is enough space in the data FIFO.
  - In interrupt mode, the application waits for the TXFE interrupt (in OTG\_FS\_DIEPINTx) and then reads the OTG\_FS\_DTXFSTSx register, to determine if there is enough space in the data FIFO.
  - To write a single non-zero length data packet, there must be space to write the entire packet in the data FIFO.
  - To write zero length packet, the application must not look at the FIFO space.
- 2. Using one of the above mentioned methods, when the application determines that there is enough space to write a transmit packet, the application must first write into the endpoint control register, before writing the data into the data FIFO. Typically, the application, must do a read modify write on the OTG\_FS\_DIEPCTLx register to avoid modifying the contents of the register, except for setting the Endpoint Enable bit.

The application can write multiple packets for the same endpoint into the transmit FIFO, if space is available. For periodic IN endpoints, the application must write packets only for one microframe. It can write packets for the next periodic transaction only after getting transfer complete for the previous transaction.

## Setting IN endpoint NAK

Internal data flow:

- 1. When the application sets the IN NAK for a particular endpoint, the core stops transmitting data on the endpoint, irrespective of data availability in the endpoint's transmit FIFO.
- 2. Non-isochronous IN tokens receive a NAK handshake reply
  - Isochronous IN tokens receive a zero-data-length packet reply
- 3. The core asserts the INEPNE (IN endpoint NAK effective) interrupt in OTG\_FS\_DIEPINTx in response to the SNAK bit in OTG\_FS\_DIEPCTLx.
- Once this interrupt is seen by the application, the application can assume that the endpoint is in IN NAK mode. This interrupt can be cleared by the application by setting the CNAK bit in OTG\_FS\_DIEPCTLx.

Application programming sequence:

- 1. To stop transmitting any data on a particular IN endpoint, the application must set the IN NAK bit. To set this bit, the following field must be programmed.
  - SNAK = 1 in OTG FS DIEPCTLx
- 2. Wait for assertion of the INEPNE interrupt in OTG\_FS\_DIEPINTx. This interrupt indicates that the core has stopped transmitting data on the endpoint.
- 3. The core can transmit valid IN data on the endpoint after the application has set the NAK bit, but before the assertion of the NAK Effective interrupt.
- 4. The application can mask this interrupt temporarily by writing to the INEPNEM bit in DIEPMSK.
  - INEPNEM = 0 in DIEPMSK
- 5. To exit Endpoint NAK mode, the application must clear the NAK status bit (NAKSTS) in OTG\_FS\_DIEPCTLx. This also clears the INEPNE interrupt (in OTG\_FS\_DIEPINTx).
  - CNAK = 1 in OTG FS DIEPCTLx
- 6. If the application masked this interrupt earlier, it must be unmasked as follows:
  - INEPNEM = 1 in DIEPMSK

### IN endpoint disable

Use the following sequence to disable a specific IN endpoint that has been previously enabled.

Application programming sequence:

- 1. The application must stop writing data on the AHB for the IN endpoint to be disabled.
- 2. The application must set the endpoint in NAK mode.
  - SNAK = 1 in OTG\_FS\_DIEPCTLx
- 3. Wait for the INEPNE interrupt in OTG\_FS\_DIEPINTx.
- Set the following bits in the OTG\_FS\_DIEPCTLx register for the endpoint that must be disabled.
  - EPDIS = 1 in OTG\_FS\_DIEPCTLx
  - SNAK = 1 in OTG\_FS\_DIEPCTLx
- 5. Assertion of the EPDISD interrupt in OTG\_FS\_DIEPINTx indicates that the core has completely disabled the specified endpoint. Along with the assertion of the interrupt, the core also clears the following bits:
  - EPENA = 0 in OTG FS DIEPCTLx
  - EPDIS = 0 in OTG\_FS\_DIEPCTLx
- 6. The application must read the OTG\_FS\_DIEPTSIZx register for the periodic IN EP, to calculate how much data on the endpoint were transmitted on the USB.
- 7. The application must flush the data in the Endpoint transmit FIFO, by setting the following fields in the OTG\_FS\_GRSTCTL register:
  - TXFNUM (in OTG\_FS\_GRSTCTL) = Endpoint transmit FIFO number
  - TXFFLSH in (OTG\_FS\_GRSTCTL) = 1

The application must poll the OTG\_FS\_GRSTCTL register, until the TXFFLSH bit is cleared by the core, which indicates the end of flush operation. To transmit new data on this endpoint, the application can re-enable the endpoint at a later point.

### Generic non-periodic IN data transfers

Application requirements:

- 1. Before setting up an IN transfer, the application must ensure that all data to be transmitted as part of the IN transfer are part of a single buffer.
- For IN transfers, the Transfer Size field in the Endpoint Transfer Size register denotes a
  payload that constitutes multiple maximum-packet-size packets and a single short
  packet. This short packet is transmitted at the end of the transfer.
  - To transmit a few maximum-packet-size packets and a short packet at the end of the transfer:

```
Transfer size[EPNUM] = x \times MPSIZ[EPNUM] + sp
If (sp > 0), then packet count[EPNUM] = x + 1.
Otherwise, packet count[EPNUM] = x
```

To transmit a single zero-length data packet:

```
Transfer size[EPNUM] = 0
Packet count[EPNUM] = 1
```

 To transmit a few maximum-packet-size packets and a zero-length data packet at the end of the transfer, the application must split the transfer into two parts. The first sends maximum-packet-size data packets and the second sends the zerolength data packet alone.

```
First transfer: transfer size[EPNUM] = x \times MPSIZ[epnum]; packet count = n;
Second transfer: transfer size[EPNUM] = 0; packet count = 1;
```

- 3. Once an endpoint is enabled for data transfers, the core updates the Transfer size register. At the end of the IN transfer, the application must read the Transfer size register to determine how much data posted in the transmit FIFO have already been sent on the USB.
- 4. Data fetched into transmit FIFO = Application-programmed initial transfer size coreupdated final transfer size
  - Data transmitted on USB = (application-programmed initial packet count Core updated final packet count) × MPSIZ[EPNUM]
  - Data yet to be transmitted on USB = (Application-programmed initial transfer size
     data transmitted on USB)

### Internal data flow:

- The application must set the transfer size and packet count fields in the endpointspecific registers and enable the endpoint to transmit the data.
- 2. The application must also write the required data to the transmit FIFO for the endpoint.
- 3. Every time a packet is written into the transmit FIFO by the application, the transfer size for that endpoint is decremented by the packet size. The data is fetched from the memory by the application, until the transfer size for the endpoint becomes 0. After writing the data into the FIFO, the "number of packets in FIFO" count is incremented (this is a 3-bit count, internally maintained by the core for each IN endpoint transmit FIFO. The maximum number of packets maintained by the core at any time in an IN endpoint FIFO is eight). For zero-length packets, a separate flag is set for each FIFO, without any data in the FIFO.
- 4. Once the data are written to the transmit FIFO, the core reads them out upon receiving an IN token. For every non-isochronous IN data packet transmitted with an ACK

- handshake, the packet count for the endpoint is decremented by one, until the packet count is zero. The packet count is not decremented on a timeout.
- 5. For zero length packets (indicated by an internal zero length flag), the core sends out a zero-length packet for the IN token and decrements the packet count field.
- 6. If there are no data in the FIFO for a received IN token and the packet count field for that endpoint is zero, the core generates an "IN token received when TxFIFO is empty" (ITTXFE) Interrupt for the endpoint, provided that the endpoint NAK bit is not set. The core responds with a NAK handshake for non-isochronous endpoints on the USB.
- 7. The core internally rewinds the FIFO pointers and no timeout interrupt is generated.
- 8. When the transfer size is 0 and the packet count is 0, the transfer complete (XFRC) interrupt for the endpoint is generated and the endpoint enable is cleared.

Application programming sequence:

- 1. Program the OTG\_FS\_DIEPTSIZx register with the transfer size and corresponding packet count.
- 2. Program the OTG\_FS\_DIEPCTLx register with the endpoint characteristics and set the CNAK and EPENA (Endpoint Enable) bits.
- 3. When transmitting non-zero length data packet, the application must poll the OTG\_FS\_DTXFSTSx register (where x is the FIFO number associated with that endpoint) to determine whether there is enough space in the data FIFO. The application can optionally use TXFE (in OTG\_FS\_DIEPINTx) before writing the data.
- Generic periodic IN data transfers

This section describes a typical periodic IN data transfer.

Application requirements:

- 1. Application requirements 1, 2, 3, and 4 of *Generic non-periodic IN data transfers on page 1071* also apply to periodic IN data transfers, except for a slight modification of requirement 2.
  - The application can only transmit multiples of maximum-packet-size data packets or multiples of maximum-packet-size packets, plus a short packet at the end. To transmit a few maximum-packet-size packets and a short packet at the end of the transfer, the following conditions must be met:

```
transfer size[EPNUM] = x \times MPSIZ[EPNUM] + sp
(where x is an integer \geq 0, and 0 \leq sp < MPSIZ[EPNUM])
If (sp > 0), packet count[EPNUM] = x + 1
Otherwise, packet count[EPNUM] = x;
MCNT[EPNUM] = packet count[EPNUM]
```

- The application cannot transmit a zero-length data packet at the end of a transfer.
   It can transmit a single zero-length data packet by itself. To transmit a single zero-length data packet:
- transfer size[EPNUM] = 0packet count[EPNUM] = 1MCNT[EPNUM] = packet count[EPNUM]

- 2. The application can only schedule data transfers one frame at a time.
  - $(MCNT 1) \times MPSIZ \le XFERSIZ \le MCNT \times MPSIZ$
  - PKTCNT = MCNT (in OTG\_FS\_DIEPTSIZx)
  - If XFERSIZ < MCNT × MPSIZ, the last data packet of the transfer is a short packet.
  - Note that: MCNT is in OTG\_FS\_DIEPTSIZx, MPSIZ is in OTG\_FS\_DIEPCTLx, PKTCNT is in OTG\_FS\_DIEPTSIZx and XFERSIZ is in OTG\_FS\_DIEPTSIZx
- 3. The complete data to be transmitted in the frame must be written into the transmit FIFO by the application, before the IN token is received. Even when 1 Word of the data to be transmitted per frame is missing in the transmit FIFO when the IN token is received, the core behaves as when the FIFO is empty. When the transmit FIFO is empty:
  - A zero data length packet would be transmitted on the USB for isochronous IN endpoints
  - A NAK handshake would be transmitted on the USB for interrupt IN endpoints

#### Internal data flow:

- The application must set the transfer size and packet count fields in the endpointspecific registers and enable the endpoint to transmit the data.
- 2. The application must also write the required data to the associated transmit FIFO for the endpoint.
- 3. Every time the application writes a packet to the transmit FIFO, the transfer size for that endpoint is decremented by the packet size. The data are fetched from application memory until the transfer size for the endpoint becomes 0.
- 4. When an IN token is received for a periodic endpoint, the core transmits the data in the FIFO, if available. If the complete data payload (complete packet, in dedicated FIFO mode) for the frame is not present in the FIFO, then the core generates an IN token received when TxFIFO empty interrupt for the endpoint.
  - A zero-length data packet is transmitted on the USB for isochronous IN endpoints
  - A NAK handshake is transmitted on the USB for interrupt IN endpoints
- 5. The packet count for the endpoint is decremented by 1 under the following conditions:
  - For isochronous endpoints, when a zero- or non-zero-length data packet is transmitted
  - For interrupt endpoints, when an ACK handshake is transmitted
  - When the transfer size and packet count are both 0, the transfer completed interrupt for the endpoint is generated and the endpoint enable is cleared.
- At the "Periodic frame Interval" (controlled by PFIVL in OTG\_FS\_DCFG), when the
  core finds non-empty any of the isochronous IN endpoint FIFOs scheduled for the
  current frame non-empty, the core generates an IISOIXFR interrupt in
  OTG\_FS\_GINTSTS.

Application programming sequence:

- Program the OTG\_FS\_DIEPCTLx register with the endpoint characteristics and set the CNAK and EPENA bits.
- Write the data to be transmitted in the next frame to the transmit FIFO.
- 3. Asserting the ITTXFE interrupt (in OTG\_FS\_DIEPINTx) indicates that the application has not yet written all data to be transmitted to the transmit FIFO.
- 4. If the interrupt endpoint is already enabled when this interrupt is detected, ignore the interrupt. If it is not enabled, enable the endpoint so that the data can be transmitted on the next IN token attempt.
- Asserting the XFRC interrupt (in OTG\_FS\_DIEPINTx) with no ITTXFE interrupt in OTG\_FS\_DIEPINTx indicates the successful completion of an isochronous IN transfer. A read to the OTG\_FS\_DIEPTSIZx register must give transfer size = 0 and packet count = 0, indicating all data were transmitted on the USB.
- 6. Asserting the XFRC interrupt (in OTG\_FS\_DIEPINTx), with or without the ITTXFE interrupt (in OTG\_FS\_DIEPINTx), indicates the successful completion of an interrupt IN transfer. A read to the OTG\_FS\_DIEPTSIZx register must give transfer size = 0 and packet count = 0, indicating all data were transmitted on the USB.
- 7. Asserting the incomplete isochronous IN transfer (IISOIXFR) interrupt in OTG\_FS\_GINTSTS with none of the aforementioned interrupts indicates the core did not receive at least 1 periodic IN token in the current frame.

### Incomplete isochronous IN data transfers

This section describes what the application must do on an incomplete isochronous IN data transfer.

#### Internal data flow:

- 1. An isochronous IN transfer is treated as incomplete in one of the following conditions:
  - a) The core receives a corrupted isochronous IN token on at least one isochronous IN endpoint. In this case, the application detects an incomplete isochronous IN transfer interrupt (IISOIXFR in OTG\_FS\_GINTSTS).
  - b) The application is slow to write the complete data payload to the transmit FIFO and an IN token is received before the complete data payload is written to the FIFO. In this case, the application detects an IN token received when TxFIFO empty interrupt in OTG\_FS\_DIEPINTx. The application can ignore this interrupt, as it eventually results in an incomplete isochronous IN transfer interrupt (IISOIXFR in OTG\_FS\_GINTSTS) at the end of periodic frame.
    - The core transmits a zero-length data packet on the USB in response to the received IN token.
- 2. The application must stop writing the data payload to the transmit FIFO as soon as possible.
- 3. The application must set the NAK bit and the disable bit for the endpoint.
- 4. The core disables the endpoint, clears the disable bit, and asserts the Endpoint Disable interrupt for the endpoint.

Application programming sequence:

- 1. The application can ignore the IN token received when TxFIFO empty interrupt in OTG\_FS\_DIEPINTx on any isochronous IN endpoint, as it eventually results in an incomplete isochronous IN transfer interrupt (in OTG\_FS\_GINTSTS).
- Assertion of the incomplete isochronous IN transfer interrupt (in OTG\_FS\_GINTSTS)
  indicates an incomplete isochronous IN transfer on at least one of the isochronous IN
  endpoints.
- 3. The application must read the Endpoint Control register for all isochronous IN endpoints to detect endpoints with incomplete IN data transfers.
- 4. The application must stop writing data to the Periodic Transmit FIFOs associated with these endpoints on the AHB.
- 5. Program the following fields in the OTG\_FS\_DIEPCTLx register to disable the endpoint:
  - SNAK = 1 in OTG\_FS\_DIEPCTLx
  - EPDIS = 1 in OTG\_FS\_DIEPCTLx
- 6. The assertion of the Endpoint Disabled interrupt in OTG\_FS\_DIEPINTx indicates that the core has disabled the endpoint.
  - At this point, the application must flush the data in the associated transmit FIFO or overwrite the existing data in the FIFO by enabling the endpoint for a new transfer in the next microframe. To flush the data, the application must use the OTG\_FS\_GRSTCTL register.

## Stalling non-isochronous IN endpoints

This section describes how the application can stall a non-isochronous endpoint.

Application programming sequence:

- 1. Disable the IN endpoint to be stalled. Set the STALL bit as well.
- 2. EPDIS = 1 in OTG\_FS\_DIEPCTLx, when the endpoint is already enabled
  - STALL = 1 in OTG\_FS\_DIEPCTLx
  - The STALL bit always takes precedence over the NAK bit
- Assertion of the Endpoint Disabled interrupt (in OTG\_FS\_DIEPINTx) indicates to the application that the core has disabled the specified endpoint.
- 4. The application must flush the non-periodic or periodic transmit FIFO, depending on the endpoint type. In case of a non-periodic endpoint, the application must re-enable the other non-periodic endpoints that do not need to be stalled, to transmit data.
- 5. Whenever the application is ready to end the STALL handshake for the endpoint, the STALL bit must be cleared in OTG\_FS\_DIEPCTLx.
- If the application sets or clears a STALL bit for an endpoint due to a
  SetFeature.Endpoint Halt command or ClearFeature.Endpoint Halt command, the
  STALL bit must be set or cleared before the application sets up the Status stage
  transfer on the control endpoint.

Special case: stalling the control OUT endpoint

The core must stall IN/OUT tokens if, during the data stage of a control transfer, the host sends more IN/OUT tokens than are specified in the SETUP packet. In this case, the application must enable the ITTXFE interrupt in OTG\_FS\_DIEPINTx and the OTEPDIS interrupt in OTG\_FS\_DOEPINTx during the data stage of the control transfer, after the core has transferred the amount of data specified in the SETUP packet. Then, when the

application receives this interrupt, it must set the STALL bit in the corresponding endpoint control register, and clear this interrupt.

# 29.17.7 Worst case response time

When the OTG\_FS controller acts as a device, there is a worst case response time for any tokens that follow an isochronous OUT. This worst case response time depends on the AHB clock frequency.

The core registers are in the AHB domain, and the core does not accept another token before updating these register values. The worst case is for any token following an isochronous OUT, because for an isochronous transaction, there is no handshake and the next token could come sooner. This worst case value is 7 PHY clocks when the AHB clock is the same as the PHY clock. When the AHB clock is faster, this value is smaller.

If this worst case condition occurs, the core responds to bulk/interrupt tokens with a NAK and drops isochronous and SETUP tokens. The host interprets this as a timeout condition for SETUP and retries the SETUP packet. For isochronous transfers, the Incomplete isochronous IN transfer interrupt (IISOIXFR) and Incomplete isochronous OUT transfer interrupt (IISOOXFR) inform the application that isochronous IN/OUT packets were dropped.

# Choosing the value of TRDT in OTG\_FS\_GUSBCFG

The value in TRDT (OTG\_FS\_GUSBCFG) is the time it takes for the MAC, in terms of PHY clocks after it has received an IN token, to get the FIFO status, and thus the first data from the PFC block. This time involves the synchronization delay between the PHY and AHB clocks. The worst case delay for this is when the AHB clock is the same as the PHY clock. In this case, the delay is 5 clocks.

Once the MAC receives an IN token, this information (token received) is synchronized to the AHB clock by the PFC (the PFC runs on the AHB clock). The PFC then reads the data from the SPRAM and writes them into the dual clock source buffer. The MAC then reads the data out of the source buffer (4 deep).

If the AHB is running at a higher frequency than the PHY, the application can use a smaller value for TRDT (in OTG\_FS\_GUSBCFG).

Figure 369 has the following signals:

- tkn\_rcvd: Token received information from MAC to PFC
- dynced\_tkn\_rcvd: Doubled sync tkn\_rcvd, from PCLK to HCLK domain
- spr read: Read to SPRAM
- spr addr: Address to SPRAM
- spr rdata: Read data from SPRAM
- srcbuf\_push: Push to the source buffer
- srcbuf rdata: Read data from the source buffer. Data seen by MAC

The application can use the following formula to calculate the value of TRDT:

 $4 \times AHB$  clock + 1 PHY clock = (2 clock sync + 1 clock memory address + 1 clock memory data from sync RAM) + (1 PHY clock (next PHY clock MAC can sample the 2 clock FIFO outputs)

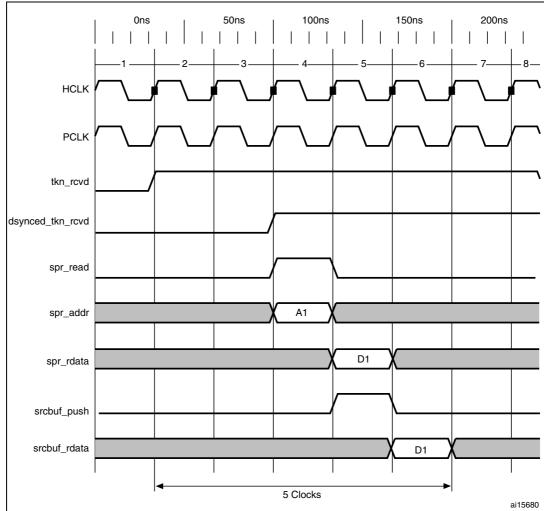


Figure 369. TRDT max timing case

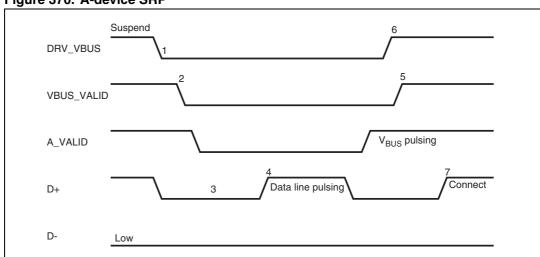
# 29.17.8 OTG programming model

The OTG\_FS controller is an OTG device supporting HNP and SRP. When the core is connected to an "A" plug, it is referred to as an A-device. When the core is connected to a "B" plug it is referred to as a B-device. In host mode, the OTG\_FS controller turns off  $V_{BUS}$  to conserve power. SRP is a method by which the B-device signals the A-device to turn on  $V_{BUS}$  power. A device must perform both data-line pulsing and  $V_{BUS}$  pulsing, but a host can detect either data-line pulsing or  $V_{BUS}$  pulsing for SRP. HNP is a method by which the B-device negotiates and switches to host role. In Negotiated mode after HNP, the B-device suspends the bus and reverts to the device role.

## A-device session request protocol

The application must set the SRP-capable bit in the Core USB configuration register. This enables the OTG FS controller to detect SRP as an A-device.

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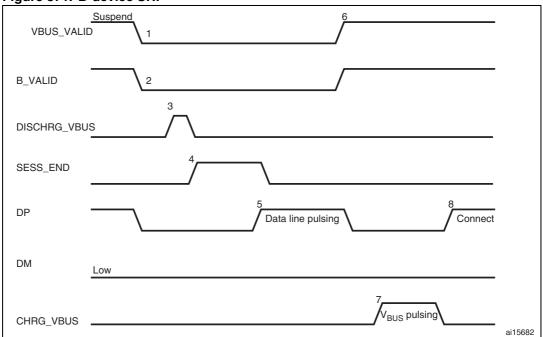
### Figure 370. A-device SRP

- DRV\_VBUS = V<sub>BUS</sub> drive signal to the PHY VBUS\_VALID = V<sub>BUS</sub> valid signal from PHY A\_VALID = A-peripheral V<sub>BUS</sub> level signal to PHY D+ = Data plus line D- = Data minus line
- To save power, the application suspends and turns off port power when the bus is idle
  by writing the port suspend and port power bits in the host port control and status
  register.
- 2. PHY indicates port power off by deasserting the VBUS\_VALID signal.
- 3. The device must detect SE0 for at least 2 ms to start SRP when  $V_{BUS}$  power is off.
- 4. To initiate SRP, the device turns on its data line pull-up resistor for 5 to 10 ms. The OTG\_FS controller detects data-line pulsing.
- 5. The device drives  $V_{BUS}$  above the A-device session valid (2.0 V minimum) for  $V_{BUS}$  pulsing.
  - The OTG\_FS controller interrupts the application on detecting SRP. The Session request detected bit is set in Global interrupt status register (SRQINT set in OTG\_FS\_GINTSTS).
- 6. The application must service the Session request detected interrupt and turn on the port power bit by writing the port power bit in the host port control and status register. The PHY indicates port power-on by asserting the VBUS\_VALID signal.
- 7. When the USB is powered, the device connects, completing the SRP process.

# **B-device session request protocol**

The application must set the SRP-capable bit in the Core USB configuration register. This enables the OTG\_FS controller to initiate SRP as a B-device. SRP is a means by which the OTG\_FS controller can request a new session from the host.



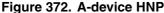


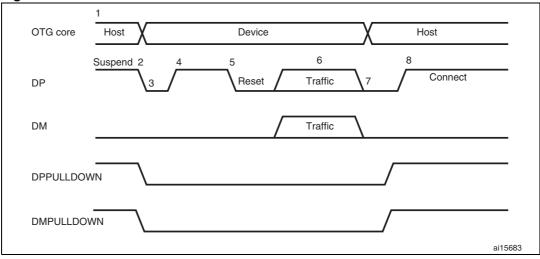
- I. VBUS\_VALID = V<sub>BUS</sub> valid signal from PHY B\_VALID = B-peripheral valid session to PHY DISCHRG\_VBUS = discharge signal to PHY SESS\_END = session end signal to PHY CHRG\_VBUS = charge V<sub>BUS</sub> signal to PHY DP = Data plus line DM = Data minus line
- 1. To save power, the host suspends and turns off port power when the bus is idle.
  - The OTG\_FS controller sets the early suspend bit in the Core interrupt register after 3 ms of bus idleness. Following this, the OTG\_FS controller sets the USB suspend bit in the Core interrupt register.
  - The OTG\_FS controller informs the PHY to discharge V<sub>BUS</sub>.
- 2. The PHY indicates the session's end to the device. This is the initial condition for SRP. The OTG\_FS controller requires 2 ms of SE0 before initiating SRP.
  - For a USB 1.1 full-speed serial transceiver, the application must wait until  $V_{BUS}$  discharges to 0.2 V after BSVLD (in OTG\_FS\_GOTGCTL) is deasserted. This

- discharge time can be obtained from the transceiver vendor and varies from one transceiver to another.
- 3. The USB OTG core informs the PHY to speed up  $V_{BUS}$  discharge.
- The application initiates SRP by writing the session request bit in the OTG Control and status register. The OTG\_FS controller perform data-line pulsing followed by V<sub>BUS</sub> pulsing.
- 5. The host detects SRP from either the data-line or  $V_{BUS}$  pulsing, and turns on  $V_{BUS}$ . The PHY indicates  $V_{BUS}$  power-on to the device.
- 6. The OTG\_FS controller performs V<sub>BUS</sub> pulsing. The host starts a new session by turning on V<sub>BUS</sub>, indicating SRP success. The OTG\_FS controller interrupts the application by setting the session request success status change bit in the OTG interrupt status register. The application reads the session request success bit in the OTG control and status register.
- 7. When the USB is powered, the OTG\_FS controller connects, completing the SRP process.

## A-device host negotiation protocol

HNP switches the USB host role from the A-device to the B-device. The application must set the HNP-capable bit in the Core USB configuration register to enable the OTG\_FS controller to perform HNP as an A-device.





- DPPULLDOWN = signal from core to PHY to enable/disable the pull-down on the DP line inside the PHY. DMPULLDOWN = signal from core to PHY to enable/disable the pull-down on the DM line inside the PHY.
- The OTG\_FS controller sends the B-device a SetFeature b\_hnp\_enable descriptor to enable HNP support. The B-device's ACK response indicates that the B-device supports HNP. The application must set host Set HNP Enable bit in the OTG Control

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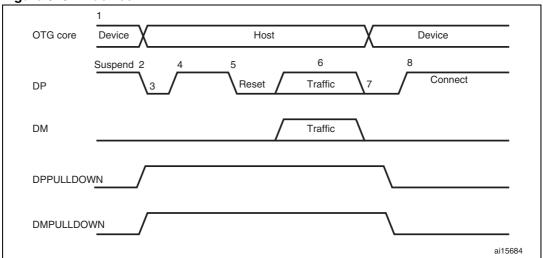
- and status register to indicate to the OTG\_FS controller that the B-device supports HNP
- 2. When it has finished using the bus, the application suspends by writing the Port suspend bit in the host port control and status register.
- 3. When the B-device observes a USB suspend, it disconnects, indicating the initial condition for HNP. The B-device initiates HNP only when it must switch to the host role; otherwise, the bus continues to be suspended.
  - The OTG\_FS controller sets the host negotiation detected interrupt in the OTG interrupt status register, indicating the start of HNP.
  - The OTG\_FS controller deasserts the DM pull down and DM pull down in the PHY to indicate a device role. The PHY enables the OTG\_FS\_DP pull-up resistor to indicate a connect for B-device.
  - The application must read the current mode bit in the OTG Control and status register to determine device mode operation.
- The B-device detects the connection, issues a USB reset, and enumerates the OTG\_FS controller for data traffic.
- 5. The B-device continues the host role, initiating traffic, and suspends the bus when done.
  - The OTG\_FS controller sets the early suspend bit in the Core interrupt register after 3 ms of bus idleness. Following this, the OTG\_FS controller sets the USB Suspend bit in the Core interrupt register.
- 6. In Negotiated mode, the OTG\_FS controller detects the suspend, disconnects, and switches back to the host role. The OTG\_FS controller asserts the DM pull down and DM pull down in the PHY to indicate its assumption of the host role.
- 7. The OTG\_FS controller sets the Connector ID status change interrupt in the OTG Interrupt Status register. The application must read the connector ID status in the OTG Control and Status register to determine the OTG\_FS controller operation as an Adevice. This indicates the completion of HNP to the application. The application must read the Current mode bit in the OTG control and status register to determine host mode operation.
- 8. The B-device connects, completing the HNP process.

## **B-device host negotiation protocol**

HNP switches the USB host role from B-device to A-device. The application must set the HNP-capable bit in the Core USB configuration register to enable the OTG\_FS controller to perform HNP as a B-device.

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- DPPULLDOWN = signal from core to PHY to enable/disable the pull-down on the DP line inside the PHY. DMPULLDOWN = signal from core to PHY to enable/disable the pull-down on the DM line inside the PHY.
- The A-device sends the SetFeature b\_hnp\_enable descriptor to enable HNP support.
   The OTG\_FS controller's ACK response indicates that it supports HNP. The application must set the device HNP enable bit in the OTG Control and status register to indicate HNP support.
  - The application sets the HNP request bit in the OTG Control and status register to indicate to the OTG\_FS controller to initiate HNP.
- 2. When it has finished using the bus, the A-device suspends by writing the Port suspend bit in the host port control and status register.
  - The OTG\_FS controller sets the Early suspend bit in the Core interrupt register after 3 ms of bus idleness. Following this, the OTG\_FS controller sets the USB suspend bit in the Core interrupt register.
  - The OTG\_FS controller disconnects and the A-device detects SE0 on the bus, indicating HNP. The OTG\_FS controller asserts the DP pull down and DM pull down in the PHY to indicate its assumption of the host role.
  - The A-device responds by activating its OTG\_FS\_DP pull-up resistor within 3 ms of detecting SE0. The OTG\_FS controller detects this as a connect.

The OTG\_FS controller sets the host negotiation success status change interrupt in the OTG Interrupt status register, indicating the HNP status. The application must read the host negotiation success bit in the OTG Control and status register to determine host

- negotiation success. The application must read the current Mode bit in the Core interrupt register (OTG\_FS\_GINTSTS) to determine host mode operation.
- 3. The application sets the reset bit (PRST in OTG\_FS\_HPRT) and the OTG\_FS controller issues a USB reset and enumerates the A-device for data traffic.
- 4. The OTG\_FS controller continues the host role of initiating traffic, and when done, suspends the bus by writing the Port suspend bit in the host port control and status register.
- 5. In Negotiated mode, when the A-device detects a suspend, it disconnects and switches back to the host role. The OTG\_FS controller deasserts the DP pull down and DM pull down in the PHY to indicate the assumption of the device role.
- 6. The application must read the current mode bit in the Core interrupt (OTG\_FS\_GINTSTS) register to determine the host mode operation.
- 7. The OTG\_FS controller connects, completing the HNP process.

# 30 USB on-the-go high-speed (OTG\_HS)

This section applies to the whole STM32F4xx family, unless otherwise specified.

# 30.1 OTG\_HS introduction

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This section presents the architecture and the programming model of the OTG\_HS controller.

The following acronyms are used throughout the section:

FS full-speed
HS High-speed
LS Low-speed

USB Universal serial bus

OTG On-the-go
PHY Physical layer

MAC Media access controller
PFC Packet FIFO controller

UTMI USB Transceiver Macrocell Interface

ULPI UTMI+ Low Pin Interface

References are made to the following documents:

- USB On-The-Go Supplement, Revision 1.3
- Universal Serial Bus Revision 2.0 Specification

The OTG\_HS is a dual-role device (DRD) controller that supports both peripheral and host functions and is fully compliant with the *On-The-Go Supplement to the USB 2.0 Specification*. It can also be configured as a host-only or peripheral-only controller, fully compliant with the *USB 2.0 Specification*. In host mode, the OTG\_HS supports high-speed (HS, 480 Mbits/s), full-speed (FS, 12 Mbits/s) and low-speed (LS, 1.5 Mbits/s) transfers whereas in peripheral mode, it only supports high-speed (HS, 480 Mbits/s) and full-speed (FS, 12 Mbits/s) transfers. The OTG\_HS supports both HNP and SRP. The only external device required is a charge pump for VBUS in OTG mode.

# 30.2 OTG\_HS main features

The main features can be divided into three categories: general, host-mode and peripheral-mode features.

#### 30.2.1 General features

The OTG\_HS interface main features are the following:

- It is USB-IF certified in compliance with the Universal Serial Bus Revision 2.0 Specification
- It supports 3 PHY interfaces
  - An on-chip full-speed PHY
  - An I<sup>2</sup>C Interface for external full-speed I<sup>2</sup>C PHY
  - An ULPI interface for external high-speed PHY.
- It supports the host negotiation protocol (HNP) and the session request protocol (SRP)
- It allows the host to turn V<sub>BUS</sub> off to save power in OTG applications, with no need for external components
- It allows to monitor V<sub>BUS</sub> levels using internal comparators
- It supports dynamic host-peripheral role switching
- It is software-configurable to operate as:
  - An SRP-capable USB HS/FS peripheral (B-device)
  - An SRP-capable USB HS/FS/low-speed host (A-device)
  - An USB OTG FS dual-role device
- It supports HS/FS SOFs as well as low-speed (LS) keep-alive tokens with:
  - SOF pulse PAD output capability
  - SOF pulse internal connection to timer 2 (TIM2)
  - Configurable framing period
  - Configurable end-of-frame interrupt
- It embeds an internal DMA with shareholding support and software selectable AHB burst type in DMA mode
- It has power saving features such as system clock stop during USB suspend, switching off of the digital core internal clock domains, PHY and DFIFO power management
- It features a dedicated 4-Kbyte data RAM with advanced FIFO management:
  - The memory partition can be configured into different FIFOs to allow flexible and efficient use of RAM
  - Each FIFO can contain multiple packets
  - Memory allocation is performed dynamically
  - The FIFO size can be configured to values that are not powers of 2 to allow the use of contiguous memory locations
- It ensures a maximum USB bandwidth of up to one frame without application intervention

#### 30.2.2 Host-mode features

The OTG\_HS interface features in host mode are the following:

- It requires an external charge pump to generate V<sub>BUS</sub>
- It has up to 12 host channels (pipes), each channel being dynamically reconfigurable to support any kind of USB transfer
- It features a built-in hardware scheduler holding:
  - Up to 8 interrupt plus isochronous transfer requests in the periodic hardware queue
  - Up to 8 control plus bulk transfer requests in the nonperiodic hardware queue
- It manages a shared RX FIFO, a periodic TX FIFO, and a nonperiodic TX FIFO for efficient usage of the USB data RAM
- It features dynamic trimming capability of SOF framing period in host mode.

# 30.2.3 Peripheral-mode features

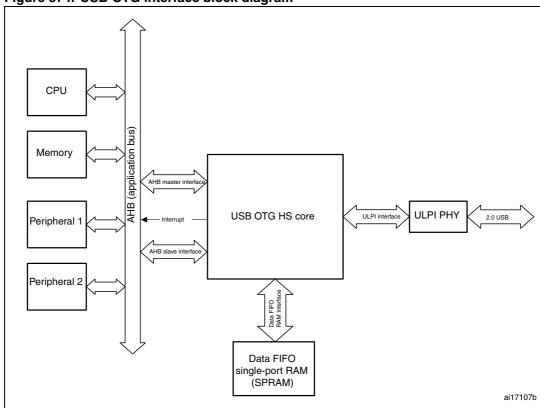
The OTG\_HS interface main features in peripheral mode are the following:

- It has 1 bidirectional control endpoint 0
- It has 5 IN endpoints (EP) configurable to support bulk, interrupt or isochronous transfers
- It has 5 OUT endpoints configurable to support bulk, interrupt or isochronous transfers
- It manages a shared Rx FIFO and a Tx-OUT FIFO for efficient usage of the USB data RAM
- It manages up to 6 dedicated Tx-IN FIFOs (one for each IN-configured EP) to reduce the application load
- It features soft disconnect capability

# 30.3 OTG\_HS functional description

Figure 374 shows the OTG\_HS interface block diagram.

Figure 374. USB OTG interface block diagram



1. The USB DMA cannot directly address the internal Flash memory.

# 30.3.1 High-speed OTG PHY

The USB OTG HS core embeds an ULPI interface to connect an external HS phy.

# 30.3.2 External Full-speed OTG PHY using the I2C interface

The USB OTG HS core embeds an I<sup>2</sup>C interface allowing to connect an external FS phy.

# 30.3.3 Embedded Full-speed OTG PHY

The full-speed OTG PHY includes the following components:

- FS/LS transceiver module used by both host and Device. It directly drives transmission and reception on the single-ended USB lines.
- Integrated ID pull-up resistor used to sample the ID line for A/B Device identification.
- DP/DM integrated pull-up and pull-down resistors controlled by the OTG\_HS core
  depending on the current role of the device. As a peripheral, it enables the DP pull-up
  resistor to signal full-speed peripheral connections as soon as V<sub>BUS</sub> is sensed to be at
  a valid level (B-session valid). In host mode, pull-down resistors are enabled on both

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DP/DM. Pull-up and pull-down resistors are dynamically switched when the peripheral role is changed via the host negotiation protocol (HNP).

- Pull-up/pull-down resistor ECN circuit
  - The DP pull-up consists of 2 resistors controlled separately from the OTG\_HS as per the resistor Engineering Change Notice applied to USB Rev2.0. The dynamic trimming of the DP pull-up strength allows to achieve a better noise rejection and Tx/Rx signal quality.
- V<sub>BUS</sub> sensing comparators with hysteresis used to detect VBUS\_VALID, A-B Session Valid and session-end voltage thresholds. They are used to drive the session request protocol (SRP), detect valid startup and end-of-session conditions, and constantly monitor the V<sub>BUS</sub> supply during USB operations.
- V<sub>BUS</sub> pulsing method circuit used to charge/discharge V<sub>BUS</sub> through resistors during the SRP (weak drive).

Caution:

To guarantee a correct operation for the USB OTG HS peripheral, the AHB frequency should be higher than 30 MHz.

## 30.4 OTG dual-role device

#### 30.4.1 ID line detection

The host or peripheral (the default) role depends on the level of the ID input line. It is determined when the USB cable is plugged in and depends on which side of the USB cable is connected to the micro-AB receptacle:

- If the B-side of the USB cable is connected with a floating ID wire, the integrated pull-up resistor detects a high ID level and the default peripheral role is confirmed. In this configuration the OTG\_HS conforms to the FSM standard described in section 6.8.2. On-The-Go B-device of the USB On-The-Go Supplement, Revision 1.3.
- If the A-side of the USB cable is connected with a grounded ID, the OTG\_HS issues an
  ID line status change interrupt (CIDSCHG bit in the OTG\_HS\_GINTSTS register) for
  host software initialization, and automatically switches to host role. In this configuration
  the OTG\_HS conforms to the FSM standard described by section 6.8.1: On-The-Go ADevice of the USB On-The-Go Supplement, Revision 1.3.

#### 30.4.2 HNP dual role device

The HNP capable bit in the Global USB configuration register (HNPCAP bit in the OTG\_HS\_GUSBCFG register) configures the OTG\_HS core to dynamically change from A-host to A-device role and vice-versa, or from B-device to B-host role and vice-versa, according to the host negotiation protocol (HNP). The current device status is defined by the combination of the Connector ID Status bit in the Global OTG control and status register (CIDSTS bit in OTG\_HS\_GOTGCTL) and the current mode of operation bit in the global interrupt and status register (CMOD bit in OTG\_HS\_GINTSTS).

The HNP programming model is described in detail in *Section 30.13: OTG\_HS programming model*.

#### 30.4.3 SRP dual-role device

The SRP capable bit in the global USB configuration register (SRPCAP bit in OTG\_HS\_GUSBCFG) configures the OTG\_HS core to switch  $V_{BUS}$  off for the A-device in

order to save power. The A-device is always in charge of driving V<sub>BUS</sub> regardless of the OTG\_HS role (host or peripheral). The SRP A/B-device program model is described in detail in *Section 30.13: OTG\_HS programming model*.

# 30.5 USB functional description in peripheral mode

The OTG\_HS operates as an USB peripheral in the following circumstances:

- OTG B-device
  - OTG B-device default state if the B-side of USB cable is plugged in
- OTG A-device
  - OTG A-device state after the HNP switches the OTG HS to peripheral role
- B-Device
  - If the ID line is present, functional and connected to the B-side of the USB cable, and the HNP-capable bit in the Global USB Configuration register (HNPCAP bit in OTG\_HS\_GUSBCFG) is cleared (see On-The-Go specification Revision 1.3 section 6.8.3).
- Peripheral only (see Figure 352: USB peripheral-only connection)
  The force peripheral mode bit in the Global USB configuration register (FDMOD in OTG\_HS\_GUSBCFG) is set to 1, forcing the OTG\_HS core to operate in USB peripheral-only mode (see On-The-Go specification Revision 1.3 section 6.8.3). In this case, the ID line is ignored even if it is available on the USB connector.

Note:

To build a bus-powered device architecture in the B-Device or peripheral-only configuration, an external regulator must be added to generate the  $V_{DD}$  supply voltage from  $V_{BUS}$ .

## 30.5.1 SRP-capable peripheral

The SRP capable bit in the Global USB configuration register (SRPCAP bit in OTG\_HS\_GUSBCFG) configures the OTG\_HS to support the session request protocol (SRP). As a result, it allows the remote A-device to save power by switching  $V_{BUS}$  off when the USB session is suspended.

The SRP peripheral mode program model is described in detail in *Section : B-device* session request protocol.

# 30.5.2 Peripheral states

## **Powered state**

The  $V_{BUS}$  input detects the B-session valid voltage used to put the USB peripheral in the Powered state (see USB2.0 specification section 9.1). The OTG\_HS then automatically connects the DP pull-up resistor to signal full-speed device connection to the host, and generates the session request interrupt (SRQINT bit in OTG\_HS\_GINTSTS) to notify the Powered state. The  $V_{BUS}$  input also ensures that valid  $V_{BUS}$  levels are supplied by the host during USB operations. If  $V_{BUS}$  drops below the B-session valid voltage (for example because power disturbances occurred or the host port has been switched off), the OTG\_HS automatically disconnects and the session end detected (SEDET bit in OTG\_HS\_GOTGINT) interrupt is generated to notify that the OTG\_HS has exited the Powered state.

In Powered state, the OTG\_HS expects a reset from the host. No other USB operations are possible. When a reset is received, the reset detected interrupt (USBRST in OTG\_HS\_GINTSTS) is generated. When the reset is complete, the enumeration done interrupt (ENUMDNE bit in OTG\_HS\_GINTSTS) is generated and the OTG\_HS enters the Default state.

### Soft disconnect

The Powered state can be exited by software by using the soft disconnect feature. The DP pull-up resistor is removed by setting the Soft disconnect bit in the device control register (SDIS bit in OTG\_HS\_DCTL), thus generating a device disconnect detection interrupt on the host side even though the USB cable was not really unplugged from the host port.

#### **Default state**

In Default state the OTG\_HS expects to receive a SET\_ADDRESS command from the host. No other USB operations are possible. When a valid SET\_ADDRESS command is decoded on the USB, the application writes the corresponding number into the device address field in the device configuration register (DAD bit in OTG\_HS\_DCFG). The OTG\_HS then enters the address state and is ready to answer host transactions at the configured USB address.

## Suspended state

The OTG\_HS peripheral constantly monitors the USB activity. When the USB remains idle for 3 ms, the early suspend interrupt (ESUSP bit in OTG\_HS\_GINTSTS) is issued. It is confirmed 3 ms later, if appropriate, by generating a suspend interrupt (USBSUSP bit in OTG\_HS\_GINTSTS). The device suspend bit is then automatically set in the device status register (SUSPSTS bit in OTG\_HS\_DSTS) and the OTG\_HS enters the Suspended state.

The device can also exit from the Suspended state by itself. In this case the application sets the remote wakeup signaling bit in the device control register (RWUSIG bit in OTG HS DCTL) and clears it after 1 to 15 ms.

When a resume signaling is detected from the host, the resume interrupt (WKUPINT bit in OTG\_HS\_GINTSTS) is generated and the device suspend bit is automatically cleared.

## 30.5.3 Peripheral endpoints

The OTG HS core instantiates the following USB endpoints:

Control endpoint 0

This endpoint is bidirectional and handles control messages only.

It has a separate set of registers to handle IN and OUT transactions, as well as dedicated control (OTG\_HS\_DIEPCTL0/OTG\_HS\_DOEPCTL0), transfer configuration (OTG\_HS\_DIEPTSIZ0/OTG\_HS\_DIEPTSIZ0), and status-interrupt

(OTG\_HS\_DIEPINTx/)OTG\_HS\_DOEPINT0) registers. The bits available inside the control and transfer size registers slightly differ from other endpoints.

# • 5 IN endpoints

- They can be configured to support the isochronous, bulk or interrupt transfer type.
- They feature dedicated control (OTG\_HS\_DIEPCTLx), transfer configuration (OTG\_HS\_DIEPTSIZx), and status-interrupt (OTG\_HS\_DIEPINTx) registers.
- The Device IN endpoints common interrupt mask register (OTG\_HS\_DIEPMSK)
  allows to enable/disable a single endpoint interrupt source on all of the
  IN endpoints (EP0 included).
- They support incomplete isochronous IN transfer interrupt (IISOIXFR bit in OTG\_HS\_GINTSTS). This interrupt is asserted when there is at least one isochronous IN endpoint for which the transfer is not completed in the current frame. This interrupt is asserted along with the end of periodic frame interrupt (OTG\_HS\_GINTSTS/EOPF).

### 5 OUT endpoints

- They can be configured to support the isochronous, bulk or interrupt transfer type.
- They feature dedicated control (OTG\_HS\_DOEPCTLx), transfer configuration (OTG\_HS\_DOEPTSIZx) and status-interrupt (OTG\_HS\_DOEPINTx) registers.
- The Device Out endpoints common interrupt mask register (OTG\_HS\_DOEPMSK) allows to enable/disable a single endpoint interrupt source on all OUT endpoints (EP0 included).
- They support incomplete isochronous OUT transfer interrupt (INCOMPISOOUT bit in OTG\_HS\_GINTSTS). This interrupt is asserted when there is at least one isochronous OUT endpoint on which the transfer is not completed in the current frame. This interrupt is asserted along with the end of periodic frame interrupt (OTG\_HS\_GINTSTS/EOPF).

## **Endpoint controls**

The following endpoint controls are available through the device endpoint-x IN/OUT control register (DIEPCTLx/DOEPCTLx):

- Endpoint enable/disable
- Endpoint activation in current configuration
- Program the USB transfer type (isochronous, bulk, interrupt)
- Program the supported packet size
- Program the Tx-FIFO number associated with the IN endpoint
- Program the expected or transmitted data0/data1 PID (bulk/interrupt only)
- Program the even/odd frame during which the transaction is received or transmitted (isochronous only)
- Optionally program the NAK bit to always send a negative acknowledge to the host regardless of the FIFO status
- Optionally program the STALL bit to always stall host tokens to that endpoint
- Optionally program the Snoop mode for OUT endpoint where the received data CRC is not checked

# **Endpoint transfer**

The device endpoint-x transfer size registers (DIEPTSIZx/DOEPTSIZx) allow the application to program the transfer size parameters and read the transfer status.

The programming operation must be performed before setting the endpoint enable bit in the endpoint control register.

Once the endpoint is enabled, these fields are read-only as the OTG FS core updates them with the current transfer status.

The following transfer parameters can be programmed:

- Transfer size in bytes
- Number of packets constituting the overall transfer size.

## **Endpoint status/interrupt**

The device endpoint-x interrupt registers (DIEPINTx/DOPEPINTx) indicate the status of an endpoint with respect to USB- and AHB-related events. The application must read these registers when the OUT endpoint interrupt bit or the IN endpoint interrupt bit in the core interrupt register (OEPINT bit in OTG\_HS\_GINTSTS or IEPINT bit in OTG\_HS\_GINTSTS, respectively) is set. Before the application can read these registers, it must first read the device all endpoints interrupt register (OTG\_HS\_DAINT) to get the exact endpoint number for the device endpoint-x interrupt register. The application must clear the appropriate bit in this register to clear the corresponding bits in the DAINT and GINTSTS registers.

The peripheral core provides the following status checks and interrupt generation:

- Transfer completed interrupt, indicating that data transfer has completed on both the application (AHB) and USB sides
- Setup stage done (control-out only)
- Associated transmit FIFO is half or completely empty (in endpoints)
- NAK acknowledge transmitted to the host (isochronous-in only)
- IN token received when Tx-FIFO was empty (bulk-in/interrupt-in only)
- OUT token received when endpoint was not yet enabled
- Babble error condition detected
- Endpoint disable by application is effective
- Endpoint NAK by application is effective (isochronous-in only)
- More than 3 back-to-back setup packets received (control-out only)
- Timeout condition detected (control-in only)
- Isochronous out packet dropped without generating an interrupt

# 30.6 USB functional description on host mode

This section gives the functional description of the OTG\_HS in the USB host mode. The OTG\_HS works as a USB host in the following circumstances:

- OTG A-host
  - OTG A-device default state when the A-side of the USB cable is plugged in
- OTG B-host
  - OTG B-device after HNP switching to the host role
- A-device

If the ID line is present, functional and connected to the A-side of the USB cable, and the HNP-capable bit is cleared in the Global USB Configuration register (HNPCAP bit in OTG\_HS\_GUSBCFG). Integrated pull-down resistors are automatically set on the DP/DM lines.

Host only (Figure 353: USB host-only connection).

The force host mode bit in the global USB configuration register (FHMOD bit in OTG\_HS\_GUSBCFG) forces the OTG\_HS core to operate in USB host-only mode. In this case, the ID line is ignored even if it is available on the USB connector. Integrated pull-down resistors are automatically set on the OTG\_HS\_FS\_DP/OTG\_HS\_FS\_DM lines.

Note:

On-chip 5 V  $V_{BUS}$  generation is not supported. As a result, a charge pump or a basic power switch (if a 5 V supply is available on the application board) must be added externally to drive the 5 V  $V_{BUS}$  line. The external charge pump can be driven by any GPIO output. This is required for the OTG A-host, A-device and host-only configurations.

The  $V_{BUS}$  input ensures that valid  $V_{BUS}$  levels are supplied by the charge pump during USB operations while the charge pump overcurrent output can be input to any GPIO pin configured to generate port interrupts. The overcurrent ISR must promptly disable the  $V_{BUS}$  generation.

# 30.6.1 SRP-capable host

SRP support is available through the SRP capable bit in the global USB configuration register (SRPCAP bit in OTG\_HS\_GUSBCFG). When the SRP feature is enabled, the host can save power by switching off the  $V_{BUS}$  power while the USB session is suspended. The SRP host mode program model is described in detail in *Section : A-device session request protocol*.

## 30.6.2 USB host states

#### Host port power

On-chip 5 V  $V_{BUS}$  generation is not supported. As a result, a charge pump or a basic power switch (if a 5 V supply voltage is available on the application board) must be added externally to drive the 5 V  $V_{BUS}$  line. The external charge pump can be driven by any GPIO output. When the application powers on  $V_{BUS}$  through the selected GPIO, it must also set the port power bit in the host port control and status register (PPWR bit in OTG\_HS\_HPRT).

# V<sub>BUS</sub> valid

The  $V_{BUS}$  input ensures that valid  $V_{BUS}$  levels are supplied by the charge pump during USB operations.

Any unforeseen  $V_{BUS}$  voltage drop below the  $V_{BUS}$  valid threshold (4.25 V) generates an OTG interrupt triggered by the session end detected bit (SEDET bit in OTG\_HS\_GOTGINT). The application must then switch the  $V_{BUS}$  power off and clear the port power bit. The charge pump overcurrent flag can also be used to prevent electrical damage. Connect the overcurrent flag output from the charge pump to any GPIO input, and configure it to generate a port interrupt on the active level. The overcurrent ISR must promptly disable the  $V_{BUS}$  generation and clear the port power bit.

## Detection of peripheral connection by the host

Even if USB peripherals or B-devices can be attached at any time, the OTG\_HS does not detect a bus connection until the end of the V<sub>BUS</sub> sensing (V<sub>BUS</sub> over 4.75 V).

When V<sub>BUS</sub> is at a valid level and a remote B-device is attached, the OTG\_HS core issues a host port interrupt triggered by the device connected bit in the host port control and status register (PCDET bit in OTG\_HS\_HPRT).

## Detection of peripheral disconnection by the host

The peripheral disconnection event triggers the disconnect detected interrupt (DISCINT bit in OTG\_HS\_GINTSTS).

#### Host enumeration

After detecting a peripheral connection, the host must start the enumeration process by issuing an USB reset and configuration commands to the new peripheral.

Before sending an USB reset, the application waits for the OTG interrupt triggered by the debounce done bit (DBCDNE bit in OTG\_HS\_GOTGINT), which indicates that the bus is stable again after the electrical debounce caused by the attachment of a pull-up resistor on OTG\_HS\_FS\_DP (full speed) or OTG\_HS\_FS\_DM (low speed).

The application issues an USB reset (single-ended zero) via the USB by keeping the port reset bit set in the Host port control and status register (PRST bit in OTG\_HS\_HPRT) for a minimum of 10 ms and a maximum of 20 ms. The application monitors the time and then clears the port reset bit.

Once the USB reset sequence has completed, the host port interrupt is triggered by the port enable/disable change bit (PENCHNG bit in OTG\_HS\_HPRT) to inform the application that the speed of the enumerated peripheral can be read from the port speed field in the host port control and status register (PSPD bit in OTG\_HS\_HPRT), and that the host is starting to drive SOFs (full speed) or keep-alive tokens (low speed). The host is then ready to complete the peripheral enumeration by sending peripheral configuration commands.

# Host suspend

The application can decide to suspend the USB activity by setting the port suspend bit in the host port control and status register (PSUSP bit in OTG\_HS\_HPRT). The OTG\_HS core stops sending SOFs and enters the Suspended state.

The Suspended state can be exited on the remote device initiative (remote wakeup). In this case the remote wakeup interrupt (WKUPINT bit in OTG\_HS\_GINTSTS) is generated upon detection of a remote wakeup event, the port resume bit in the host port control and status

register (PRES bit in OTG\_HS\_HPRT) is set, and a resume signaling is automatically issued on the USB. The application must monitor the resume window duration, and then clear the port resume bit to exit the Suspended state and restart the SOF.

If the Suspended state is exited on the host initiative, the application must set the port resume bit to start resume signaling on the host port, monitor the resume window duration and then clear the port resume bit.

#### 30.6.3 Host channels

The OTG\_HS core instantiates 12 host channels. Each host channel supports an USB host transfer (USB pipe). The host is not able to support more than 8 transfer requests simultaneously. If more than 8 transfer requests are pending from the application, the host controller driver (HCD) must re-allocate channels when they become available, that is, after receiving the transfer completed and channel halted interrupts.

Each host channel can be configured to support IN/OUT and any type of periodic/nonperiodic transaction. Each host channel has dedicated control (HCCHARx), transfer configuration (HCTSIZx) and status/interrupt (HCINTx) registers with associated mask (HCINTMSKx) registers.

## Host channel controls

The following host channel controls are available through the host channel-x characteristics register (HCCHARx):

- Channel enable/disable
- Program the HS/FS/LS speed of target USB peripheral
- Program the address of target USB peripheral
- Program the endpoint number of target USB peripheral
- Program the transfer IN/OUT direction
- Program the USB transfer type (control, bulk, interrupt, isochronous)
- Program the maximum packet size (MPS)
- Program the periodic transfer to be executed during odd/even frames

#### Host channel transfer

The host channel transfer size registers (HCTSIZx) allow the application to program the transfer size parameters, and read the transfer status.

The programming operation must be performed before setting the channel enable bit in the host channel characteristics register. Once the endpoint is enabled, the packet count field is read-only as the OTG HS core updates it according to the current transfer status.

The following transfer parameters can be programmed:

- Transfer size in bytes
- Number of packets constituting the overall transfer size
- Initial data PID

# Host channel status/interrupt

The host channel-x interrupt register (HCINTx) indicates the status of an endpoint with respect to USB- and AHB-related events. The application must read these register when the host channels interrupt bit in the core interrupt register (HCINT bit in OTG HS GINTSTS) is

set. Before the application can read these registers, it must first read the host all channels interrupt (HCAINT) register to get the exact channel number for the host channel-x interrupt register. The application must clear the appropriate bit in this register to clear the corresponding bits in the HAINT and GINTSTS registers. The mask bits for each interrupt source of each channel are also available in the OTG\_HS\_HCINTMSK-x register.

The host core provides the following status checks and interrupt generation:

- Transfer completed interrupt, indicating that the data transfer is complete on both the application (AHB) and USB sides
- Channel stopped due to transfer completed, USB transaction error or disable command from the application
- Associated transmit FIFO half or completely empty (IN endpoints)
- ACK response received
- NAK response received
- STALL response received
- USB transaction error due to CRC failure, timeout, bit stuff error, false EOP
- Babble error
- Frame overrun
- Data toggle error

#### 30.6.4 Host scheduler

The host core features a built-in hardware scheduler which is able to autonomously re-order and manage the USB the transaction requests posted by the application. At the beginning of each frame the host executes the periodic (isochronous and interrupt) transactions first, followed by the nonperiodic (control and bulk) transactions to achieve the higher level of priority granted to the isochronous and interrupt transfer types by the USB specification.

The host processes the USB transactions through request queues (one for periodic and one for nonperiodic). Each request queue can hold up to 8 entries. Each entry represents a pending transaction request from the application, and holds the IN or OUT channel number along with other information to perform a transaction on the USB. The order in which the requests are written to the queue determines the sequence of the transactions on the USB interface.

At the beginning of each frame, the host processes the periodic request queue first, followed by the nonperiodic request queue. The host issues an incomplete periodic transfer interrupt (IPXFR bit in OTG\_HS\_GINTSTS) if an isochronous or interrupt transaction scheduled for the current frame is still pending at the end of the current frame. The OTG HS core is fully responsible for the management of the periodic and nonperiodic request queues. The periodic transmit FIFO and queue status register (HPTXSTS) and nonperiodic transmit FIFO and queue status register (HNPTXSTS) are read-only registers which can be used by the application to read the status of each request queue. They contain:

- The number of free entries currently available in the periodic (nonperiodic) request queue (8 max)
- Free space currently available in the periodic (nonperiodic) Tx-FIFO (out-transactions)
- IN/OUT token, host channel number and other status information.

As request queues can hold a maximum of 8 entries each, the application can push to schedule host transactions in advance with respect to the moment they physically reach the

USB for a maximum of 8 pending periodic transactions plus 8 pending nonperiodic transactions.

To post a transaction request to the host scheduler (queue) the application must check that there is at least 1 entry available in the periodic (nonperiodic) request queue by reading the PTXQSAV bits in the OTG\_HS\_HNPTXSTS register or NPTQXSAV bits in the OTG\_HS\_HNPTXSTS register.

# 30.7 SOF trigger

The OTG FS core allows to monitor, track and configure SOF framing in the host and peripheral. It also features an SOF pulse output connectivity.

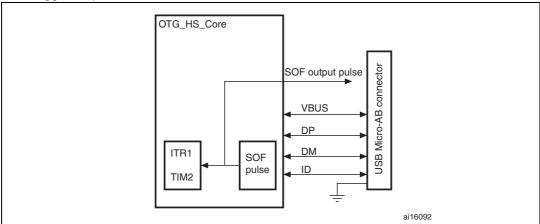
These capabilities are particularly useful to implement adaptive audio clock generation techniques, where the audio peripheral needs to synchronize to the isochronous stream provided by the PC, or the host needs trimming its framing rate according to the requirements of the audio peripheral.

## 30.7.1 Host SOFs

In host mode the number of PHY clocks occurring between the generation of two consecutive SOF (FS) or keep-alive (LS) tokens is programmable in the host frame interval register (OTG\_HS\_HFIR), thus providing application control over the SOF framing period. An interrupt is generated at any start of frame (SOF bit in OTG\_HS\_GINTSTS). The current frame number and the time remaining until the next SOF are tracked in the host frame number register (OTG\_HS\_HFNUM).

An SOF pulse signal is generated at any SOF starting token and with a width of 12 system clock cycles. It can be made available externally on the SOF pin using the SOFOUTEN bit in the global control and configuration register. The SOF pulse is also internally connected to the input trigger of timer 2 (TIM2), so that the input capture feature, the output compare feature and the timer can be triggered by the SOF pulse. The TIM2 connection is enabled through ITR1\_RMP bits of TIM2\_OR register.





# 30.7.2 Peripheral SOFs

In peripheral mode, the start of frame interrupt is generated each time an SOF token is received on the USB (SOF bit in OTG\_HS\_GINTSTS). The corresponding frame number

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can be read from the device status register (FNSOF bit in OTG\_HS\_DSTS). An SOF pulse signal with a width of 12 system clock cycles is also generated and can be made available externally on the SOF pin by using the SOF output enable bit in the global control and configuration register (SOFOUTEN bit in OTG\_HS\_GCCFG). The SOF pulse signal is also internally connected to the TIM2 input trigger, so that the input capture feature, the output compare feature and the timer can be triggered by the SOF pulse (see *Figure*). The TIM2 connection is enabled through ITR1\_RMP bits of TIM2\_OR register.

The end of periodic frame interrupt (GINTSTS/EOPF) is used to notify the application when 80%, 85%, 90% or 95% of the time frame interval elapsed depending on the periodic frame interval field in the device configuration register (PFIVL bit in OTG\_HS\_DCFG).

This feature can be used to determine if all of the isochronous traffic for that frame is complete.

# 30.8 USB\_HS power modes

The power consumption of the OTG PHY is controlled by three bits in the general core configuration register:

- PHY power down (GCCFG/PWRDWN)
  - This bit switches on/off the PHY full-speed transceiver module. It must be preliminarily set to allow any USB operation.
- A-VBUS sensing enable (GCCFG/VBUSASEN)
  - This bit switches on/off the  $V_{BUS}$  comparators associated with A-device operations. It must be set when in A-device (USB host) mode and during HNP.
- B-VBUS sensing enable (GCCFG/VBUSASEN)
  - This bit switches on/off the  $V_{BUS}$  comparators associated with B-device operations. It must be set when in B-device (USB peripheral) mode and during HNP.
  - Power reduction techniques are available in the USB suspended state, when the USB session is not yet valid or the device is disconnected.
- Stop PHY clock (STPPCLK bit in OTG\_HS\_PCGCCTL)
  - When setting the stop PHY clock bit in the clock gating control register, most of the clock domain internal to the OTG high-speed core is switched off by clock gating.
     The dynamic power consumption due to the USB clock switching activity is cut even if the clock input is kept running by the application
  - Most of the transceiver is also disabled, and only the part in charge of detecting the asynchronous resume or remote wakeup event is kept alive.
- Gate HCLK (GATEHCLK bit in OTG\_HS\_PCGCCTL)
  - When setting the Gate HCLK bit in the clock gating control register, most of the system clock domain internal to the OTG\_HS core is switched off by clock gating. Only the register read and write interface is kept alive. The dynamic power consumption due to the USB clock switching activity is cut even if the system clock is kept running by the application for other purposes.
- USB system stop
  - When the OTG\_HS is in USB suspended state, the application can decide to drastically reduce the overall power consumption by shutting down all the clock sources in the system. USB System Stop is activated by first setting the Stop PHY

- clock bit and then configuring the system deep sleep mode in the powercontrol system module (PWR).
- The OTG\_HS core automatically reactivates both system and USB clocks by asynchronous detection of remote wakeup (as an host) or resume (as a Device) signaling on the USB.

# 30.9 Dynamic update of the OTG\_HS\_HFIR register

The USB core embeds a dynamic trimming capability of micro-SOF framing period in host mode allowing to synchronize an external device with the micro-SOF frames.

When the OTG\_HS\_HFIR register is changed within a current micro-SOF frame, the SOF period correction is applied in the next frame as described in *Figure 375*.

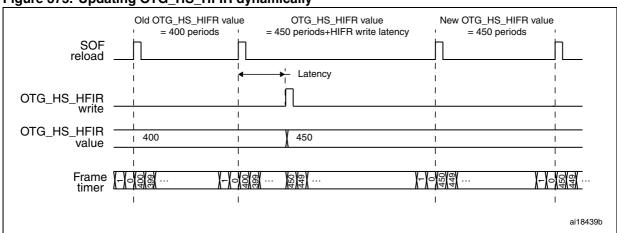


Figure 375. Updating OTG HS HFIR dynamically

# 30.10 FIFO RAM allocation

# 30.10.1 Peripheral mode

### **Receive FIFO RAM**

For Receive FIFO RAM, the application should allocate RAM for SETUP packets: 10 locations must be reserved in the receive FIFO to receive SETUP packets on control endpoints. These locations are reserved for SETUP packets and are not used by the core to write any other data.

One location must be allocated for Global OUT NAK. Status information are also written to the FIFO along with each received packet. Therefore, a minimum space of (Largest Packet Size / 4) + 1 must be allocated to receive packets. If a high-bandwidth endpoint or multiple isochronous endpoints are enabled, at least two spaces of (Largest Packet Size / 4) + 1 must be allotted to receive back-to-back packets. Typically, two (Largest Packet Size / 4) + 1 spaces are recommended so that when the previous packet is being transferred to AHB, the USB can receive the subsequent packet.

Along with each endpoints last packet, transfer complete status information are also pushed to the FIFO. Typically, one location for each OUT endpoint is recommended.

#### **Transmit FIFO RAM**

For Transmit FIFO RAM, the minimum RAM space required for each IN Endpoint Transmit FIFO is the maximum packet size for this IN endpoint.

Note:

More space allocated in the transmit IN Endpoint FIFO results in a better performance on the USB.

#### 30.10.2 Host mode

#### Receive FIFO RAM

For Receive FIFO RAM allocation, Status information are written to the FIFO along with each received packet. Therefore, a minimum space of (Largest Packet Size / 4) + 1 must be allocated to receive packets. If a high-bandwidth channel or multiple isochronous channels are enabled, at least two spaces of (Largest Packet Size / 4) + 1 must be allocated to receive back-to-back packets. Typically, two (Largest Packet Size / 4) + 1 spaces are recommended so that when the previous packet is being transferred to AHB, the USB can receive the subsequent packet.

Along with each host channels last packet, transfer complete status information are also pushed to the FIFO. As a consequence, one location must be allocated to store this data.

### **Transmit FIFO RAM**

For Transmit FIFO RAM allocation, the minimum amount of RAM required for the host nonperiodic Transmit FIFO is the largest maximum packet size for all supported nonperiodic OUT channels. Typically, a space corresponding to two Largest Packet Size is recommended, so that when the current packet is being transferred to the USB, the AHB can transmit the subsequent packet.

The minimum amount of RAM required for Host periodic Transmit FIFO is the largest maximum packet size for all supported periodic OUT channels. If there is at least one High Bandwidth Isochronous OUT endpoint, then the space must be at least two times the maximum packet size for that channel.

Note:

More space allocated in the Transmit nonperiodic FIFO results in better performance on the USB.

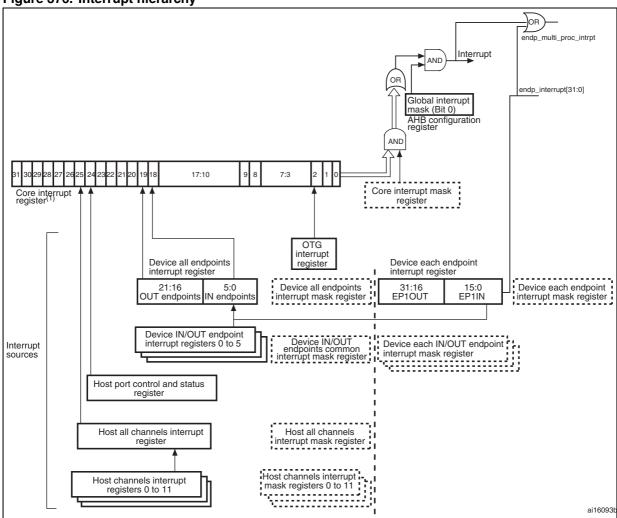
When operating in DMA mode, the DMA address register for each host channel (HCDMAn) is stored in the SPRAM (FIFO). One location for each channel must be reserved for this.

# 30.11 OTG\_HS interrupts

When the OTG\_HS controller is operating in one mode, either peripheral or host, the application must not access registers from the other mode. If an illegal access occurs, a mode mismatch interrupt is generated and reflected in the Core interrupt register (MMIS bit in the OTG\_HS\_GINTSTS register). When the core switches from one mode to the other, the registers in the new mode of operation must be reprogrammed as they would be after a power-on reset.

Figure 376 shows the interrupt hierarchy.

Figure 376. Interrupt hierarchy



1. The core interrupt register bits are shown in OTG\_HS core interrupt register (OTG\_HS\_GINTSTS) on page 1118.

# 30.12 OTG\_HS control and status registers

By reading from and writing to the control and status registers (CSRs) through the AHB slave interface, the application controls the OTG\_HS controller. These registers are 32 bits

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wide, and the addresses are 32-bit block aligned. The OTG\_HS registers must be accessed by words (32 bits). CSRs are classified as follows:

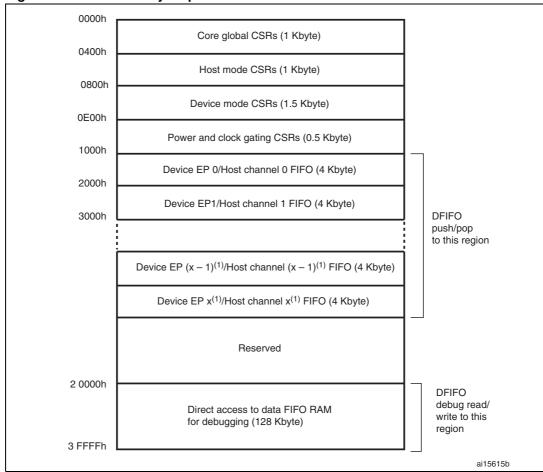
- Core global registers
- Host-mode registers
- Host global registers
- Host port CSRs
- Host channel-specific registers
- Device-mode registers
- Device global registers
- Device endpoint-specific registers
- Power and clock-gating registers
- Data FIFO (DFIFO) access registers

Only the Core global, Power and clock-gating, Data FIFO access, and host port control and status registers can be accessed in both host and peripheral modes. When the OTG\_HS controller is operating in one mode, either peripheral or host, the application must not access registers from the other mode. If an illegal access occurs, a mode mismatch interrupt is generated and reflected in the Core interrupt register (MMIS bit in the OTG\_HS\_GINTSTS register). When the core switches from one mode to the other, the registers in the new mode of operation must be reprogrammed as they would be after a power-on reset.

# 30.12.1 CSR memory map

The host and peripheral mode registers occupy different addresses. All registers are implemented in the AHB clock domain.

Figure 377. CSR memory map



<sup>1.</sup> x = 5 in peripheral mode and x = 11 in host mode.

# Global CSR map

These registers are available in both host and peripheral modes.

Table 160. Core global control and status registers (CSRs)

Acronym	Address offset	Register name
OTG_HS_GOTGCTL	0x000	OTG_HS control and status register (OTG_HS_GOTGCTL) on page 1108
OTG_HS_GOTGINT	0x004	OTG_HS interrupt register (OTG_HS_GOTGINT) on page 1110
OTG_HS_GAHBCFG	0x008	OTG_HS AHB configuration register (OTG_HS_GAHBCFG) on page 1111
OTG_HS_GUSBCFG	0x00C	OTG_HS USB configuration register (OTG_HS_GUSBCFG) on page 1112
OTG_HS_GRSTCTL	0x010	OTG_HS reset register (OTG_HS_GRSTCTL) on page 1115

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Table 160. Core global control and status registers (CSRs) (continued)

Acronym	Address offset	Register name
OTG_HS_GINTSTS	0x014	OTG_HS core interrupt register (OTG_HS_GINTSTS) on page 1118
OTG_HS_GINTMSK	0x018	OTG_HS interrupt mask register (OTG_HS_GINTMSK) on page 1122
OTG_HS_GRXSTSR	0x01C	OTG_HS Receive status debug read/OTG status read and pop registers
OTG_HS_GRXSTSP	0x020	(OTG_HS_GRXSTSR/OTG_HS_GRXSTSP) on page 1125
OTG_HS_GRXFSIZ	0x024	OTG_HS Receive FIFO size register (OTG_HS_GRXFSIZ) on page 1126
OTG_HS_GNPTXFSIZ/ OTG_HS_TX0FSIZ	0x028	OTG_HS nonperiodic transmit FIFO size/Endpoint 0 transmit FIFO size register (OTG_HS_GNPTXFSIZ/OTG_HS_TX0FSIZ) on page 1127
OTG_HS_GNPTXSTS	0x02C	OTG_HS nonperiodic transmit FIFO/queue status register (OTG_HS_GNPTXSTS) on page 1127
OTG_HS_GCCFG	0x038	OTG_HS general core configuration register (OTG_HS_GCCFG) on page 1130
OTG_HS_CID	0x03C	OTG_HS core ID register (OTG_HS_CID) on page 1131
OTG_HS_HPTXFSIZ	0x100	OTG_HS Host periodic transmit FIFO size register (OTG_HS_HPTXFSIZ) on page 1131
OTG_HS_DIEPTXFx	0x104 0x124  0x13C	OTG_HS device IN endpoint transmit FIFO size register (OTG_HS_DIEPTXFx) ( $x = 17$ , where $x$ is the FIFO_number) on page 1131

# **Host-mode CSR map**

These registers must be programmed every time the core changes to host mode.

Table 161. Host-mode control and status registers (CSRs)

Acronym	Offset address	Register name
OTG_HS_HCFG	0x400	OTG_HS host configuration register (OTG_HS_HCFG) on page 1132
OTG_HS_HFIR	0x404	OTG_HS Host frame interval register (OTG_HS_HFIR) on page 1133
OTG_HS_HFNUM	0x408	OTG_HS host frame number/frame time remaining register (OTG_HS_HFNUM) on page 1133
OTG_HS_HPTXSTS	0x410	OTG_HS_Host periodic transmit FIFO/queue status register (OTG_HS_HPTXSTS) on page 1134
OTG_HS_HAINT	0x414	OTG_HS Host all channels interrupt register (OTG_HS_HAINT) on page 1135
OTG_HS_HAINTMSK	0x418	OTG_HS host all channels interrupt mask register (OTG_HS_HAINTMSK) on page 1135

Table 161. Host-mode control and status registers (CSRs) (continued)

Acronym	Offset address	Register name
OTG_HS_HPRT	0x440	OTG_HS host port control and status register (OTG_HS_HPRT) on page 1136
OTG_HS_HCCHARx	0x500 0x520  0x6E0	OTG_HS host channel-x characteristics register (OTG_HS_HCCHARx) $(x = 011, where \ x = Channel_number)$ on page 1138
OTG_HS_HCSPLTx	0x504	OTG_HS host channel-x split control register (OTG_HS_HCSPLTx) $(x = 011$ , where $x = Channel_number$ ) on page 1140
OTG_HS_HCINTx	0x508	OTG_HS host channel-x interrupt register (OTG_HS_HCINTx) ( $x = 011$ , where $x = Channel_number$ ) on page 1141
OTG_HS_HCINTMSKx	0x50C	OTG_HS host channel-x interrupt mask register (OTG_HS_HCINTMSKx) $(x = 011$ , where $x = Channel\_number$ ) on page 1142
OTG_HS_HCTSIZx	0x510	OTG_HS host channel-x transfer size register (OTG_HS_HCTSIZx) $(x = 011$ , where $x = Channel_number$ ) on page 1143
OTG_HS_HCDMAx	0x514	OTG_HS host channel-x DMA address register (OTG_HS_HCDMAx) $(x = 011$ , where $x = Channel_number$ ) on page 1144

# **Device-mode CSR map**

These registers must be programmed every time the core changes to peripheral mode.

Table 162. Device-mode control and status registers

Acronym	Offset address	Register name
OTG_HS_DCFG	0x800	OTG_HS device configuration register (OTG_HS_DCFG) on page 1144
OTG_HS_DCTL	0x804	OTG_HS device control register (OTG_HS_DCTL) on page 1146
OTG_HS_DSTS	0x808	OTG_HS device status register (OTG_HS_DSTS) on page 1148
OTG_HS_DIEPMSK	0x810	OTG_HS device IN endpoint common interrupt mask register (OTG_HS_DIEPMSK) on page 1149
OTG_HS_DOEPMSK	0x814	OTG_HS device OUT endpoint common interrupt mask register (OTG_HS_DOEPMSK) on page 1150
OTG_HS_DAINT	0x818	OTG_HS device all endpoints interrupt register (OTG_HS_DAINT) on page 1151
OTG_HS_DAINTMSK	0x81C	OTG_HS all endpoints interrupt mask register (OTG_HS_DAINTMSK) on page 1151
OTG_HS_DVBUSDIS	0x828	OTG_HS device VBUS discharge time register (OTG_HS_DVBUSDIS) on page 1152
OTG_HS_DVBUSPULSE	0x82C	OTG_HS device VBUS pulsing time register (OTG_HS_DVBUSPULSE) on page 1152

Table 162. Device-mode control and status registers (continued)

Acronym	Offset address	Register name
OTG_HS_DIEPEMPMSK	0x834	OTG_HS device IN endpoint FIFO empty interrupt mask register: (OTG_HS_DIEPEMPMSK) on page 1154
OTG_HS_EACHHINT	0x838	OTG_HS device each endpoint interrupt register (OTG_HS_DEACHINT) on page 1154
OTG_HS_EACHHINTMSK	0x83C	OTG_HS device each endpoint interrupt register mask (OTG_HS_DEACHINTMSK) on page 1155
OTG_HS_DIEPEACHMSK1	0x840	OTG_HS device each in endpoint-1 interrupt register (OTG_HS_DIEPEACHMSK1) on page 1155
OTG_HS_DOEPEACHMSK1	0x880	OTG_HS device each OUT endpoint-1 interrupt register (OTG_HS_DOEPEACHMSK1) on page 1156
OTG_HS_DIEPCTLx	0x920 0x940  0xAE0	OTG device endpoint-x control register (OTG_HS_DIEPCTLx) (x = 07, where x = Endpoint_number) on page 1157
OTG_HS_DIEPINTx	0x908	OTG_HS device endpoint-x interrupt register (OTG_HS_DIEPINTx) (x = 07, where x = Endpoint_number) on page 1164
OTG_HS_DIEPTSIZ0	0x910	OTG_HS device IN endpoint 0 transfer size register (OTG_HS_DIEPTSIZ0) on page 1167
OTG_HS_DIEPDMAx	0x914	OTG_HS device endpoint-x DMA address register (OTG_HS_DIEPDMAx / OTG_HS_DOEPDMAx) (x = 15, where x = Endpoint_number) on page 1171
OTG_HS_DTXFSTSx	0x918	OTG_HS device IN endpoint transmit FIFO status register (OTG_HS_DTXFSTSx) (x = 05, where x = Endpoint_number) on page 1170
OTG_HS_DIEPTSIZx	0x930 0x950  0xAF0	OTG_HS device endpoint-x transfer size register (OTG_HS_DOEPTSIZx) (x = 15, where x = Endpoint_number) on page 1170
OTG_HS_DOEPCTL0	0xB00	OTG_HS device control OUT endpoint 0 control register (OTG_HS_DOEPCTL0) on page 1160
OTG_HS_DOEPCTLx	0xB20 0xB40  0xCC0 0xCE0 0xCFD	OTG device endpoint-x control register (OTG_HS_DIEPCTLx) (x = 07, where x = Endpoint_number) on page 1157

Table 162. Device-mode control and status registers (continued)

Acronym	Offset address	Register name
OTG_HS_DOEPINTx	0xB08	OTG_HS device endpoint-x interrupt register (OTG_HS_DIEPINTx) ( $x = 07$ , where $x = \text{Endpoint}\_\text{number}$ ) on page 1164
OTG_HS_DOEPTSIZx	0xB10	OTG_HS device endpoint-x transfer size register (OTG_HS_DOEPTSIZx) (x = 15, where x = Endpoint_number) on page 1170

# Data FIFO (DFIFO) access register map

These registers, available in both host and peripheral modes, are used to read or write the FIFO space for a specific endpoint or a channel, in a given direction. If a host channel is of type IN, the FIFO can only be read on the channel. Similarly, if a host channel is of type OUT, the FIFO can only be written on the channel.

Table 163. Data FIFO (DFIFO) access register map

FIFO access register section	Address range	Access
Device IN Endpoint 0/Host OUT Channel 0: DFIFO Write Access Device OUT Endpoint 0/Host IN Channel 0: DFIFO Read Access	0x1000-0x1FFC	w r
Device IN Endpoint 1/Host OUT Channel 1: DFIFO Write Access Device OUT Endpoint 1/Host IN Channel 1: DFIFO Read Access	0x2000-0x2FFC	w r
Device IN Endpoint $x^{(1)}$ /Host OUT Channel $x^{(1)}$ : DFIFO Write Access Device OUT Endpoint $x^{(1)}$ /Host IN Channel $x^{(1)}$ : DFIFO Read Access	0xX000h-0xXFFCh	w r

<sup>1.</sup> Where x is 5 in peripheral mode and 11 in host mode.

# Power and clock gating CSR map

There is a single register for power and clock gating. It is available in both host and peripheral modes.

Table 164. Power and clock gating control and status registers

Register name	Acronym	Offset address: 0xE00-0xFFF
Power and clock gating control register	PCGCR	0xE00-0xE04
Reserved		0xE05-0xFFF

# 30.12.2 OTG\_HS global registers

These registers are available in both host and peripheral modes, and do not need to be reprogrammed when switching between these modes.

Bit values in the register descriptions are expressed in binary unless otherwise specified.

# OTG\_HS control and status register (OTG\_HS\_GOTGCTL)

Address offset: 0x000

Reset value: 0x0000 0800

The OTG control and status register controls the behavior and reflects the status of the OTG function of the core.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	BSVLD	ASVLD	DBCT	CIDSTS	Reserved	DHNPEN	HSHNPEN	HNPRQ	HNGSCS	Reserved	SRQ	SRQSCS
	r	r	r	r		rw	rw	rw	r		rw	r

Bits 31:20 Reserved, must be kept at reset value.

#### Bit 19 BSVLD: B-session valid

Indicates the peripheral mode transceiver status.

- 0: B-session is not valid.
- 1: B-session is valid.

In OTG mode, you can use this bit to determine if the device is connected or disconnected.

Note: Only accessible in peripheral mode.

#### Bit 18 ASVLD: A-session valid

Indicates the host mode transceiver status.

- 0: A-session is not valid
- 1: A-session is valid

Note: Only accessible in host mode.

## Bit 17 DBCT: Long/short debounce time

Indicates the debounce time of a detected connection.

- 0: Long debounce time, used for physical connections (100 ms  $+ 2.5 \mu s$ )
- 1: Short debounce time, used for soft connections (2.5 µs)

Note: Only accessible in host mode.

#### Bit 16 CIDSTS: Connector ID status

Indicates the connector ID status on a connect event.

- 0: The OTG\_HS controller is in A-device mode
- 1: The OTG\_HS controller is in B-device mode

Note: Accessible in both peripheral and host modes.

## Bits 15:12 Reserved, must be kept at reset value.

### Bit 11 **DHNPEN:** Device HNP enabled

The application sets this bit when it successfully receives a SetFeature.SetHNPEnable command from the connected USB host.

- 0: HNP is not enabled in the application
- 1: HNP is enabled in the application

Note: Only accessible in peripheral mode.

#### Bit 10 HSHNPEN: Host set HNP enable

The application sets this bit when it has successfully enabled HNP (using the SetFeature.SetHNPEnable command) on the connected device.

0: Host Set HNP is not enabled

1: Host Set HNP is enabled

Note: Only accessible in host mode.

#### Bit 9 HNPRQ: HNP request

The application sets this bit to initiate an HNP request to the connected USB host. The application can clear this bit by writing a 0 when the host negotiation success status change bit in the OTG interrupt register (HNSSCHG bit in OTG\_HS\_GOTGINT) is set. The core clears this bit when the HNSSCHG bit is cleared.

0: No HNP request

1: HNP request

Note: Only accessible in peripheral mode.

#### Bit 8 HNGSCS: Host negotiation success

The core sets this bit when host negotiation is successful. The core clears this bit when the HNP Request (HNPRQ) bit in this register is set.

0: Host negotiation failure

1: Host negotiation success

Note: Only accessible in peripheral mode.

## Bits 7:2 Reserved, must be kept at reset value.

#### Bit 1 SRQ: Session request

The application sets this bit to initiate a session request on the USB. The application can clear this bit by writing a 0 when the host negotiation success status change bit in the OTG Interrupt register (HNSSCHG bit in OTG\_HS\_GOTGINT) is set. The core clears this bit when the HNSSCHG bit is cleared.

If you use the USB 1.1 full-speed serial transceiver interface to initiate the session request, the application must wait until  $V_{BUS}$  discharges to 0.2 V, after the B-Session Valid bit in this register (BSVLD bit in OTG\_HS\_GOTGCTL) is cleared. This discharge time varies between different PHYs and can be obtained from the PHY vendor.

0: No session request

1: Session request

Note: Only accessible in peripheral mode.

### Bit 0 SRQSCS: Session request success

The core sets this bit when a session request initiation is successful.

0: Session request failure

1: Session request success

Note: Only accessible in peripheral mode.

# OTG\_HS interrupt register (OTG\_HS\_GOTGINT)

Address offset: 0x04

Reset value: 0x0000 0000

The application reads this register whenever there is an OTG interrupt and clears the bits in

this register to clear the OTG interrupt.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	DBCDNE	АБТОСНС	HNGDET	Reserved	HNSSCHG	SRSSCHG	Reserved	SEDET	Res.
		rc_ w1	rc_ w1		rc_ w1	rc_ w1		rc_ w1	

Bits 31:20 Reserved, must be kept at reset value.

## Bit 19 DBCDNE: Debounce done

The core sets this bit when the debounce is completed after the device connect. The application can start driving USB reset after seeing this interrupt. This bit is only valid when the HNP Capable or SRP Capable bit is set in the Core USB Configuration register (HNPCAP bit or SRPCAP bit in OTG\_HS\_GUSBCFG, respectively).

Note: Only accessible in host mode.

#### Bit 18 ADTOCHG: A-device timeout change

The core sets this bit to indicate that the A-device has timed out while waiting for the B-device to connect.

Note: Accessible in both peripheral and host modes.

## Bit 17 HNGDET: Host negotiation detected

The core sets this bit when it detects a host negotiation request on the USB.

Note: Accessible in both peripheral and host modes.

Bits 16:10 Reserved, must be kept at reset value.

# Bit 9 HNSSCHG: Host negotiation success status change

The core sets this bit on the success or failure of a USB host negotiation request. The application must read the host negotiation success bit of the OTG Control and Status register (HNGSCS in OTG\_HS\_GOTGCTL) to check for success or failure.

Note: Accessible in both peripheral and host modes.

Bits 7:3 Reserved, must be kept at reset value.

#### Bit 8 SRSSCHG: Session request success status change

The core sets this bit on the success or failure of a session request. The application must read the session request success bit in the OTG Control and status register (SRQSCS bit in OTG\_HS\_GOTGCTL) to check for success or failure.

Note: Accessible in both peripheral and host modes.

## Bit 2 SEDET: Session end detected

The core sets this bit to indicate that the level of the voltage on  $V_{BUS}$  is no longer valid for a B-device session when  $V_{BUS} < 0.8 \text{ V}$ .

Bits 1:0 Reserved, must be kept at reset value.

# OTG\_HS AHB configuration register (OTG\_HS\_GAHBCFG)

Address offset: 0x008 Reset value: 0x0000 0000

This register can be used to configure the core after power-on or a change in mode. This register mainly contains AHB system-related configuration parameters. Do not change this register after the initial programming. The application must program this register before starting any transactions on either the AHB or the USB.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	PTXFELVL	TXFELVL	Reserved	DMAEN	HB	STLE	ΪN	GINT	
r	rw	rw	ъ	rw				rw	Ī

Bits 31:20 Reserved, must be kept at reset value.

#### Bit 8 PTXFELVL: Periodic TxFIFO empty level

Indicates when the periodic TxFIFO empty interrupt bit in the Core interrupt register (PTXFE bit in OTG\_HS\_GINTSTS) is triggered.

0: PTXFE (in OTG\_HS\_GINTSTS) interrupt indicates that the Periodic TxFIFO is half empty 1: PTXFE (in OTG\_HS\_GINTSTS) interrupt indicates that the Periodic TxFIFO is completely empty

Note: Only accessible in host mode.

## Bit 7 TXFELVL: TxFIFO empty level

In peripheral mode, this bit indicates when the IN endpoint Transmit FIFO empty interrupt (TXFE in OTG\_HS\_DIEPINTx.) is triggered.

0: TXFE (in OTG\_HS\_DIEPINTx) interrupt indicates that the IN Endpoint TxFIFO is half empty

1: TXFE (in OTG\_HS\_DIEPINTx) interrupt indicates that the IN Endpoint TxFIFO is completely empty

Note: Only accessible in peripheral mode.

Bit 6 Reserved, must be kept at reset value.

## Bits5 DMAEN: DMA enable

0: The core operates in slave mode

1: The core operates in DMA mode

## Bits 4:1 HBSTLEN: Burst length/type

0000 Single

0001 INCR

0011 INCR4

0101 INCR8

0111 INCR16

Others: Reserved

#### Bit 0 GINT: Global interrupt mask

This bit is used to mask or unmask the interrupt line assertion to the application. Irrespective of this bit setting, the interrupt status registers are updated by the core.

0: Mask the interrupt assertion to the application.

1: Unmask the interrupt assertion to the application

Note: Accessible in both peripheral and host modes.

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# OTG\_HS USB configuration register (OTG\_HS\_GUSBCFG)

Address offset: 0x00C Reset value: 0x0000 0A00

This register can be used to configure the core after power-on or a changing to host mode or peripheral mode. It contains USB and USB-PHY related configuration parameters. The application must program this register before starting any transactions on either the AHB or the USB. Do not make changes to this register after the initial programming.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
CTXPKT	FDMOD	FHMOD	Re	eserv	ed	ULPIIPD	PTCI	PCCI	TSDPS	ULPIEVBUSI	ULPIEVBUSD	NEDICSM	ULPIAR	ULPIFSLS	Reserved	SDATAHA	Reserved		TR	DT		HNPCAP	SRPCAP	Reserved	PHSEL	Re	eserv	ed	Т	OCA	L
rw	rw	rw				rw	rw	rw	rw	rw	rw	rw	rw	rw		rw			r	w		r/rw	r/rw		wo					rw	

#### Bit 31 CTXPKT: Corrupt Tx packet

This bit is for debug purposes only. Never set this bit to 1.

Note: Accessible in both peripheral and host modes.

#### Bit 30 FDMOD: Forced peripheral mode

Writing a 1 to this bit forces the core to peripheral mode irrespective of the OTG\_HS\_ID input pin.

- 0: Normal mode
- 1: Forced peripheral mode

After setting the force bit, the application must wait at least 25 ms before the change takes effect.

Note: Accessible in both peripheral and host modes.

#### Bit 29 FHMOD: Forced host mode

Writing a 1 to this bit forces the core to host mode irrespective of the OTG\_HS\_ID input pin.

- 0: Normal mode
- 1: Forced host mode

After setting the force bit, the application must wait at least 25 ms before the change takes effect.

Note: Accessible in both peripheral and host modes.

### Bits 28:26 Reserved, must be kept at reset value.

## Bit 25 **ULPIIPD:** ULPI interface protect disable

This bit controls the circuitry built in the PHY to protect the ULPI interface when the link tristates stp and data. Any pull-up or pull-down resistors employed by this feature can be disabled. Please refer to the ULPI specification for more details.

- 0: Enables the interface protection circuit
- 1: Disables the interface protection circuit

# Bit 24 PTCI: Indicator pass through

This bit controls whether the complement output is qualified with the internal  $V_{BUS}$  valid comparator before being used in the  $V_{BUS}$  state in the RX CMD. Please refer to the ULPI specification for more details.

- 0: Complement Output signal is qualified with the Internal  $V_{\mbox{\footnotesize BUS}}$  valid comparator
- 1: Complement Output signal is not qualified with the Internal V<sub>BUS</sub> valid comparator

#### Bit 23 PCCI: Indicator complement

This bit controls the PHY to invert the ExternalVbusIndicator input signal, and generate the complement output. Please refer to the ULPI specification for more details.

- 0: PHY does not invert the ExternalVbusIndicator signal
- 1: PHY inverts ExternalVbusIndicator signal

#### Bit 22 TSDPS: TermSel DLine pulsing selection

This bit selects utmi\_termselect to drive the data line pulse during SRP (session request protocol).

- 0: Data line pulsing using utmi\_txvalid (default)
- 1: Data line pulsing using utmi\_termsel

#### Bit 21 **ULPIEVBUSI:** ULPI external V<sub>BUS</sub> indicator

This bit indicates to the ULPI PHY to use an external  $V_{BUS}$  overcurrent indicator.

- 0: PHY uses an internal V<sub>BUS</sub> valid comparator
- 1: PHY uses an external V<sub>BUS</sub> valid comparator

# Bit 20 **ULPIEVBUSD:** ULPI External V<sub>BUS</sub> Drive

This bit selects between internal or external supply to drive 5 V on V<sub>BUS</sub>, in the ULPI PHY.

- 0: PHY drives V<sub>BUS</sub> using internal charge pump (default)
- 1: PHY drives V<sub>BUS</sub> using external supply.

#### Bit 19 ULPICSM: ULPI Clock SuspendM

This bit sets the ClockSuspendM bit in the interface control register on the ULPI PHY. This bit applies only in the serial and carkit modes.

- 0: PHY powers down the internal clock during suspend
- 1: PHY does not power down the internal clock

## Bit 18 ULPIAR: ULPI Auto-resume

This bit sets the AutoResume bit in the interface control register on the ULPI PHY.

- 0: PHY does not use AutoResume feature
- 1: PHY uses AutoResume feature

#### Bit 17 ULPIFSLS: ULPI FS/LS select

The application uses this bit to select the FS/LS serial interface for the ULPI PHY. This bit is valid only when the FS serial transceiver is selected on the ULPI PHY.

- 0: ULPI interface
- 1: ULPI FS/LS serial interface
- Bit 16 Reserved, must be kept at reset value.

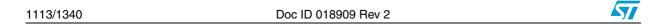
#### Bit 15 PHYLPCS: PHY Low-power clock select

This bit selects either 480 MHz or 48 MHz (low-power) PHY mode. In FS and LS modes, the PHY can usually operate on a 48 MHz clock to save power.

- 0: 480 MHz internal PLL clock
- 1: 48 MHz external clock

In 480 MHz mode, the UTMI interface operates at either 60 or 30 MHz, depending on whether the 8- or 16-bit data width is selected. In 48 MHz mode, the UTMI interface operates at 48 MHz in FS and LS modes.

### Bit 14 Reserved, must be kept at reset value.



#### Bits 13:10 TRDT: USB turnaround time

Sets the turnaround time in PHY clocks.

The formula below gives the value of TRDT:

TRDT = 4 × AHB clock frequency+ 1 PHY clock frequency.

For example:

If AHB clock frequency = 72 MHz (PHY Clock frequency = 48 MHz), the TRDT must be set to q

If AHB clock frequency = 48 Mhz (PHY Clock frequency = 48 MHz), the TRDT must be set to 5.

Note: Only accessible in peripheral mode.

## Bit 9 HNPCAP: HNP-capable

The application uses this bit to control the OTG\_HS controller's HNP capabilities.

0: HNP capability is not enabled

1: HNP capability is enabled

Note: Accessible in both peripheral and host modes.

## Bit 8 SRPCAP: SRP-capable

The application uses this bit to control the OTG\_HS controller's SRP capabilities. If the core operates as a nonSRP-capable B-device, it cannot request the connected A-device (host) to activate  $V_{BUS}$  and start a session.

0: SRP capability is not enabled

1: SRP capability is enabled

Note: Accessible in both peripheral and host modes.

- Bit 7 Reserved, must be kept at reset value.
- Bit 6 PHSEL: USB 2.0 high-speed ULPI PHY or USB 1.1 full-speed serial transceiver select

0: USB 2.0 high-speed ULPI PHY

1: USB 1.1 full-speed serial transceiver

Bits 5:3 Reserved, must be kept at reset value.

#### Bits 2:0 TOCAL: FS timeout calibration

The number of PHY clocks that the application programs in this field is added to the full-speed interpacket timeout duration in the core to account for any additional delays introduced by the PHY. This can be required, because the delay introduced by the PHY in generating the line state condition can vary from one PHY to another.

The USB standard timeout value for full-speed operation is 16 to 18 (inclusive) bit times. The application must program this field based on the speed of enumeration. The number of bit times added per PHY clock is 0.25 bit times.

# OTG\_HS reset register (OTG\_HS\_GRSTCTL)

Address offset: 0x010

Reset value: 0x2000 0000

The application uses this register to reset various hardware features inside the core.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

AHBIDL	DMAREQ	leserved	TXFNUM	TXFFLSH	RXFFLSH	Seerved	FCRST	HSRST	CSRST
r	r	ш.	rw	rs	rs	ш	rs	rs	rs

#### Bit 31 AHBIDL: AHB master idle

Indicates that the AHB master state machine is in the Idle condition.

Note: Accessible in both peripheral and host modes.

#### Bit 30 DMAREQ: DMA request signal

This bit indicates that the DMA request is in progress. Used for debug.

Bits 29:11 Reserved, must be kept at reset value.

#### Bits 10:6 TXFNUM: TxFIFO number

This is the FIFO number that must be flushed using the TxFIFO Flush bit. This field must not be changed until the core clears the TxFIFO Flush bit.

- 00000:
  - Nonperiodic TxFIFO flush in host mode
  - Tx FIFO 0 flush in peripheral mode
- 00001:
  - Periodic TxFIFO flush in host mode
  - TXFIFO 1 flush in peripheral mode
- 00010: TXFIFO 2 flush in peripheral mode

- 00101: TXFIFO 15 flush in peripheral mode
- 10000: Flush all the transmit FIFOs in peripheral or host mode.

Note: Accessible in both peripheral and host modes.

# Bit 5 TXFFLSH: TxFIFO flush

This bit selectively flushes a single or all transmit FIFOs, but cannot do so if the core is in the midst of a transaction.

The application must write this bit only after checking that the core is neither writing to the TxFIFO nor reading from the TxFIFO. Verify using these registers:

- Read: the NAK effective interrupt ensures the core is not reading from the FIFO
- Write: the AHBIDL bit in OTG\_HS\_GRSTCTL ensures that the core is not writing anything to the FIFO

Note: Accessible in both peripheral and host modes.

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#### Bit 4 RXFFLSH: RxFIFO flush

The application can flush the entire RxFIFO using this bit, but must first ensure that the core is not in the middle of a transaction.

The application must only write to this bit after checking that the core is neither reading from the RxFIFO nor writing to the RxFIFO.

The application must wait until the bit is cleared before performing any other operation. This bit requires 8 clocks (slowest of PHY or AHB clock) to be cleared.

Note: Accessible in both peripheral and host modes.

#### Bit 3 Reserved, must be kept at reset value.

#### Bit 2 FCRST: Host frame counter reset

The application writes this bit to reset the frame number counter inside the core. When the frame counter is reset, the subsequent SOF sent out by the core has a frame number of 0.

Note: Only accessible in host mode.

#### Bit 1 HSRST: HCLK soft reset

The application uses this bit to flush the control logic in the AHB Clock domain. Only AHB Clock Domain pipelines are reset.

FIFOs are not flushed with this bit.

All state machines in the AHB clock domain are reset to the Idle state after terminating the transactions on the AHB, following the protocol.

CSR control bits used by the AHB clock domain state machines are cleared.

To clear this interrupt, status mask bits that control the interrupt status and are generated by the AHB clock domain state machine are cleared.

Because interrupt status bits are not cleared, the application can get the status of any core events that occurred after it set this bit.

This is a self-clearing bit that the core clears after all necessary logic is reset in the core. This can take several clocks, depending on the core's current state.

Note: Accessible in both peripheral and host modes.

#### Bit 0 CSRST: Core soft reset

Resets the HCLK and PCLK domains as follows:

Clears the interrupts and all the CSR register bits except for the following bits:

- RSTPDMODL bit in OTG\_HS\_PCGCCTL
- GAYEHCLK bit in OTG\_HS\_PCGCCTL
- PWRCLMP bit in OTG\_HS\_PCGCCTL
- STPPCLK bit in OTG\_HS\_PCGCCTL
- FSLSPCS bit in OTG\_HS\_HCFG
- DSPD bit in OTG\_HS\_DCFG

All module state machines (except for the AHB slave unit) are reset to the Idle state, and all the transmit FIFOs and the receive FIFO are flushed.

Any transactions on the AHB Master are terminated as soon as possible, after completing the last data phase of an AHB transfer. Any transactions on the USB are terminated immediately.

The application can write to this bit any time it wants to reset the core. This is a self-clearing bit and the core clears this bit after all the necessary logic is reset in the core, which can take several clocks, depending on the current state of the core. Once this bit has been cleared, the software must wait at least 3 PHY clocks before accessing the PHY domain (synchronization delay). The software must also check that bit 31 in this register is set to 1 (AHB Master is Idle) before starting any operation.

Typically, the software reset is used during software development and also when you dynamically change the PHY selection bits in the above listed USB configuration registers. When you change the PHY, the corresponding clock for the PHY is selected and used in the PHY domain. Once a new clock is selected, the PHY domain has to be reset for proper operation.

Note: Accessible in both peripheral and host modes.



## OTG\_HS core interrupt register (OTG\_HS\_GINTSTS)

Address offset: 0x014

Reset value: 0x0400 0020

This register interrupts the application for system-level events in the current mode (peripheral mode or host mode).

Some of the bits in this register are valid only in host mode, while others are valid in peripheral mode only. This register also indicates the current mode. To clear the interrupt status bits of the rc w1 type, the application must write 1 into the bit.

The FIFO status interrupts are read-only; once software reads from or writes to the FIFO while servicing these interrupts, FIFO interrupt conditions are cleared automatically.

The application must clear the OTG\_HS\_GINTSTS register at initialization before unmasking the interrupt bit to avoid any interrupts generated prior to initialization.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
WKUINT	SRQINT	DISCINT	CIDSCHG	Reserved	PTXFE	HCINT	HPRTINT	Reserved	DATAFSUSP	IPXFR/INCOMPISOOUT	IISOIXFR	OEPINT	IEPINT	Contract of the contract of th	0000000	EOPF	ISOODRP	ENUMDNE	USBRST	USBSUSP	ESUSP	Reserved		BOUTNAKEFF	GINAKEFF	NPTXFE	RXFLVL	SOF	OTGINT	SIMM	СМОБ
	rc_	w1			r	r	r			rc_	w1	r	r					rc_	w1					r	r	r	r	rc_w1	r	rc_w1	r

## Bit 31 WKUPINT: Resume/remote wakeup detected interrupt

In peripheral mode, this interrupt is asserted when a resume is detected on the USB. In host mode, this interrupt is asserted when a remote wakeup is detected on the USB.

Note: Accessible in both peripheral and host modes.

Bit 30 SRQINT: Session request/new session detected interrupt

In host mode, this interrupt is asserted when a session request is detected from the device. In peripheral mode, this interrupt is asserted when  $V_{BUS}$  is in the valid range for a B-device device. Accessible in both peripheral and host modes.

Bit 29 **DISCINT:** Disconnect detected interrupt

Asserted when a device disconnect is detected.

Note: Only accessible in host mode.

Bit 28 CIDSCHG: Connector ID status change

The core sets this bit when there is a change in connector ID status.

Note: Accessible in both peripheral and host modes.

Bit 27 Reserved, must be kept at reset value.

Bit 26 PTXFE: Periodic TxFIFO empty

Asserted when the periodic transmit FIFO is either half or completely empty and there is space for at least one entry to be written in the periodic request queue. The half or completely empty status is determined by the periodic TxFIFO empty level bit in the Core AHB configuration register (PTXFELVL bit in OTG\_HS\_GAHBCFG).

Note: Only accessible in host mode.

#### Bit 25 **HCINT:** Host channels interrupt

The core sets this bit to indicate that an interrupt is pending on one of the channels of the core (in host mode). The application must read the host all channels interrupt (OTG\_HS\_HAINT) register to determine the exact number of the channel on which the interrupt occurred, and then read the corresponding host channel-x interrupt (OTG\_HS\_HCINTx) register to determine the exact cause of the interrupt. The application must clear the appropriate status bit in the OTG\_HS\_HCINTx register to clear this bit.

Note: Only accessible in host mode.

#### Bit 24 HPRTINT: Host port interrupt

The core sets this bit to indicate a change in port status of one of the OTG\_HS controller ports in host mode. The application must read the host port control and status (OTG\_HS\_HPRT) register to determine the exact event that caused this interrupt. The application must clear the appropriate status bit in the host port control and status register to clear this bit.

Note: Only accessible in host mode.

Bits 23 Reserved, must be kept at reset value.

#### Bit 22 DATAFSUSP: Data fetch suspended

This interrupt is valid only in DMA mode. This interrupt indicates that the core has stopped fetching data for IN endpoints due to the unavailability of TxFIFO space or request queue space. This interrupt is used by the application for an endpoint mismatch algorithm. For example, after detecting an endpoint mismatch, the application:

- Sets a global nonperiodic IN NAK handshake
- Disables IN endpoints
- Flushes the FIFO
- Determines the token sequence from the IN token sequence learning queue
- Re-enables the endpoints
- Clears the global nonperiodic IN NAK handshake If the global nonperiodic IN NAK is cleared, the core has not yet fetched data for the IN endpoint, and the IN token is received: the core generates an "IN token received when FIFO empty" interrupt. The OTG then sends a NAK response to the host. To avoid this scenario, the application can check the FetSusp interrupt in OTG\_FS\_GINTSTS, which ensures that the FIFO is full before clearing a global NAK handshake. Alternatively, the application can mask the "IN token received when FIFO empty" interrupt when clearing a global IN NAK handshake.

## Bit 21 IPXFR: Incomplete periodic transfer

In host mode, the core sets this interrupt bit when there are incomplete periodic transactions still pending, which are scheduled for the current frame.

Note: Only accessible in host mode.

## **INCOMPISOOUT:** Incomplete isochronous OUT transfer

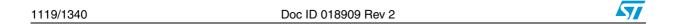
In peripheral mode, the core sets this interrupt to indicate that there is at least one isochronous OUT endpoint on which the transfer is not completed in the current frame. This interrupt is asserted along with the End of periodic frame interrupt (EOPF) bit in this register.

Note: Only accessible in peripheral mode.

## Bit 20 IISOIXFR: Incomplete isochronous IN transfer

The core sets this interrupt to indicate that there is at least one isochronous IN endpoint on which the transfer is not completed in the current frame. This interrupt is asserted along with the End of periodic frame interrupt (EOPF) bit in this register.

Note: Only accessible in peripheral mode.



## Bit 19 **OEPINT:** OUT endpoint interrupt

The core sets this bit to indicate that an interrupt is pending on one of the OUT endpoints of the core (in peripheral mode). The application must read the device all endpoints interrupt (OTG\_HS\_DAINT) register to determine the exact number of the OUT endpoint on which the interrupt occurred, and then read the corresponding device OUT Endpoint-x Interrupt (OTG\_HS\_DOEPINTx) register to determine the exact cause of the interrupt. The application must clear the appropriate status bit in the corresponding OTG\_HS\_DOEPINTx register to clear this bit.

Note: Only accessible in peripheral mode.

#### Bit 18 IEPINT: IN endpoint interrupt

The core sets this bit to indicate that an interrupt is pending on one of the IN endpoints of the core (in peripheral mode). The application must read the device All Endpoints Interrupt (OTG\_HS\_DAINT) register to determine the exact number of the IN endpoint on which the interrupt occurred, and then read the corresponding device IN Endpoint-x interrupt (OTG\_HS\_DIEPINTx) register to determine the exact cause of the interrupt. The application must clear the appropriate status bit in the corresponding OTG\_HS\_DIEPINTx register to clear this bit.

Note: Only accessible in peripheral mode.

## Bits 17:16 Reserved, must be kept at reset value.

#### Bit 15 EOPF: End of periodic frame interrupt

Indicates that the period specified in the periodic frame interval field of the device configuration register (PFIVL bit in OTG\_HS\_DCFG) has been reached in the current frame.

Note: Only accessible in peripheral mode.

## Bit 14 ISOODRP: Isochronous OUT packet dropped interrupt

The core sets this bit when it fails to write an isochronous OUT packet into the RxFIFO because the RxFIFO does not have enough space to accommodate a maximum size packet for the isochronous OUT endpoint.

Note: Only accessible in peripheral mode.

### Bit 13 ENUMDNE: Enumeration done

The core sets this bit to indicate that speed enumeration is complete. The application must read the device Status (OTG\_HS\_DSTS) register to obtain the enumerated speed.

Note: Only accessible in peripheral mode.

## Bit 12 USBRST: USB reset

The core sets this bit to indicate that a reset is detected on the USB.

Note: Only accessible in peripheral mode.

## Bit 11 USBSUSP: USB suspend

The core sets this bit to indicate that a suspend was detected on the USB. The core enters the Suspended state when there is no activity on the data lines for a period of 3 ms.

Note: Only accessible in peripheral mode.

## Bit 10 ESUSP: Early suspend

The core sets this bit to indicate that an Idle state has been detected on the USB for 3 ms.

Note: Only accessible in peripheral mode.

Bits 9:8 Reserved, must be kept at reset value.

#### Bit 7 GONAKEFF: Global OUT NAK effective

Indicates that the Set global OUT NAK bit in the Device control register (SGONAK bit in OTG\_HS\_DCTL), set by the application, has taken effect in the core. This bit can be cleared by writing the Clear global OUT NAK bit in the Device control register (CGONAK bit in OTG\_HS\_DCTL).

Note: Only accessible in peripheral mode.

## Bit 6 GINAKEFF: Global IN nonperiodic NAK effective

Indicates that the Set global nonperiodic IN NAK bit in the Device control register (SGINAK bit in OTG\_HS\_DCTL), set by the application, has taken effect in the core. That is, the core has sampled the Global IN NAK bit set by the application. This bit can be cleared by clearing the Clear global nonperiodic IN NAK bit in the Device control register (CGINAK bit in OTG\_HS\_DCTL).

This interrupt does not necessarily mean that a NAK handshake is sent out on the USB. The STALL bit takes precedence over the NAK bit.

Note: Only accessible in peripheral mode.

## Bit 5 NPTXFE: Nonperiodic TxFIFO empty

This interrupt is asserted when the nonperiodic TxFIFO is either half or completely empty, and there is space in at least one entry to be written to the nonperiodic transmit request queue. The half or completely empty status is determined by the nonperiodic TxFIFO empty level bit in the OTG\_HS\_GAHBCFG register (TXFELVL bit in OTG\_HS\_GAHBCFG).

Note: Only accessible in host mode.

#### Bit 4 RXFLVL: RxFIFO nonempty

Indicates that there is at least one packet pending to be read from the RxFIFO.

Note: Accessible in both host and peripheral modes.

## Bit 3 SOF: Start of frame

In host mode, the core sets this bit to indicate that an SOF (FS), or Keep-Alive (LS) is transmitted on the USB. The application must write a 1 to this bit to clear the interrupt. In peripheral mode, in the core sets this bit to indicate that an SOF token has been received on the USB. The application can read the Device Status register to get the current frame number. This interrupt is seen only when the core is operating in FS.

Note: Accessible in both host and peripheral modes.

#### Bit 2 OTGINT: OTG interrupt

The core sets this bit to indicate an OTG protocol event. The application must read the OTG Interrupt Status (OTG\_HS\_GOTGINT) register to determine the exact event that caused this interrupt. The application must clear the appropriate status bit in the OTG\_HS\_GOTGINT register to clear this bit.

Note: Accessible in both host and peripheral modes.

## Bit 1 MMIS: Mode mismatch interrupt

The core sets this bit when the application is trying to access:

A host mode register, when the core is operating in peripheral mode

A peripheral mode register, when the core is operating in host mode

The register access is completed on the AHB with an OKAY response, but is ignored by the core internally and does not affect the operation of the core.

Note: Accessible in both host and peripheral modes.

## Bit 0 CMOD: Current mode of operation

Indicates the current mode.

0: Peripheral mode

1: Host mode

Note: Accessible in both host and peripheral modes.



## OTG\_HS interrupt mask register (OTG\_HS\_GINTMSK)

Address offset: 0x018 Reset value: 0x0000 0000

This register works with the Core interrupt register to interrupt the application. When an interrupt bit is masked, the interrupt associated with that bit is not generated. However, the Core Interrupt (OTG\_HS\_GINTSTS) register bit corresponding to that interrupt is still set.

30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 IPXFRM/IISOOXFRM ENUMDNEM GONAKEFFM GINAKEFFM FSUSPM IISOIXFRM JSBSUSPM CIDSCHGN SOODRPN NPTXFEM RXFLVLM DISCINT EPMISM EOPFM **ESUSPM** OEPINT USBRST PTXFEM PRTIM OTGINT SRQIM IEPINT SOFM HCIM Reserved Reserved Reserved Reserved rw rw

Bit 31 WUIM: Resume/remote wakeup detected interrupt mask

0: Masked interrupt

1: Unmasked interrupt

Note: Accessible in both host and peripheral modes.

Bit 30 SRQIM: Session request/new session detected interrupt mask

0: Masked interrupt1: Unmasked interrupt

Note: Accessible in both host and peripheral modes.

Bit 29 **DISCINT:** Disconnect detected interrupt mask

0: Masked interrupt1: Unmasked interrupt

Note: Accessible in both host and peripheral modes.

Bit 28 CIDSCHGM: Connector ID status change mask

0: Masked interrupt1: Unmasked interrupt

Note: Accessible in both host and peripheral modes.

Bit 27 Reserved, must be kept at reset value.

Bit 26 PTXFEM: Periodic TxFIFO empty mask

0: Masked interrupt1: Unmasked interrupt

Note: Only accessible in host mode.

Bit 25 HCIM: Host channels interrupt mask

0: Masked interrupt1: Unmasked interrupt

Note: Only accessible in host mode.

Bit 24 PRTIM: Host port interrupt mask

0: Masked interrupt1: Unmasked interrupt

Note: Only accessible in host mode.

Bit 23 Reserved, must be kept at reset value.

Bit 22 FSUSPM: Data fetch suspended mask

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in peripheral mode.

Bit 21 IPXFRM: Incomplete periodic transfer mask

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in host mode.

**IISOOXFRM:** Incomplete isochronous OUT transfer mask

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in peripheral mode.

Bit 20 IISOIXFRM: Incomplete isochronous IN transfer mask

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in peripheral mode.

Bit 19 **OEPINT:** OUT endpoints interrupt mask

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in peripheral mode.

Bit 18 IEPINT: IN endpoints interrupt mask

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in peripheral mode.

Bit 17 EPMISM: Endpoint mismatch interrupt mask

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in peripheral mode.

Bit 16 Reserved, must be kept at reset value.

Bit 15 EOPFM: End of periodic frame interrupt mask

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in peripheral mode.

Bit 14 ISOODRPM: Isochronous OUT packet dropped interrupt mask

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in peripheral mode.

Bit 13 ENUMDNEM: Enumeration done mask

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in peripheral mode.

Bit 12 USBRST: USB reset mask

1123/1340

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in peripheral mode.

## Bit 11 USBSUSPM: USB suspend mask

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in peripheral mode.

## Bit 10 ESUSPM: Early suspend mask

0: Masked interrupt1: Unmasked interrupt

Note: Only accessible in peripheral mode.

## Bits 9:8 Reserved, must be kept at reset value.

## Bit 7 GONAKEFFM: Global OUT NAK effective mask

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in peripheral mode.

#### Bit 6 GINAKEFFM: Global nonperiodic IN NAK effective mask

0: Masked interrupt

1: Unmasked interrupt

Note: Only accessible in peripheral mode.

## Bit 5 NPTXFEM: Nonperiodic TxFIFO empty mask

0: Masked interrupt

1: Unmasked interrupt

Note: Accessible in both peripheral and host modes.

## Bit 4 RXFLVLM: Receive FIFO nonempty mask

0: Masked interrupt

1: Unmasked interrupt

Note: Accessible in both peripheral and host modes.

#### Bit 3 SOFM: Start of frame mask

0: Masked interrupt

1: Unmasked interrupt

Note: Accessible in both peripheral and host modes.

## Bit 2 OTGINT: OTG interrupt mask

0: Masked interrupt

1: Unmasked interrupt

Note: Accessible in both peripheral and host modes.

## Bit 1 MMISM: Mode mismatch interrupt mask

0: Masked interrupt

1: Unmasked interrupt

Note: Accessible in both peripheral and host modes.

Bit 0 Reserved, must be kept at reset value.

## OTG\_HS Receive status debug read/OTG status read and pop registers (OTG\_HS\_GRXSTSR/OTG\_HS\_GRXSTSP)

Address offset for Read: 0x01C Address offset for Pop: 0x020

Reset value: 0x0000 0000

A read to the Receive status debug read register returns the contents of the top of the Receive FIFO. A read to the Receive status read and pop register additionally pops the top data entry out of the RxFIFO.

The receive status contents must be interpreted differently in host and peripheral modes. The core ignores the receive status pop/read when the receive FIFO is empty and returns a value of 0x0000 0000. The application must only pop the Receive Status FIFO when the Receive FIFO nonempty bit of the Core interrupt register (RXFLVL bit in OTG\_HS\_GINTSTS) is asserted.

## Host mode:

 $31 \quad 30 \quad 29 \quad 28 \quad 27 \quad 26 \quad 25 \quad 24 \quad 23 \quad 22 \quad 21 \quad 20 \quad 19 \quad 18 \quad 17 \quad 16 \quad 15 \quad 14 \quad 13 \quad 12 \quad 11 \quad 10 \quad 9 \quad 8 \quad 7 \quad 6 \quad 5 \quad 4 \quad 3 \quad 2 \quad 1 \quad 0$ 

Reserved	PKTSTS	DPID	BCNT	CHNUM
rieserveu	r	r	r	r

Bits 31:21 Reserved, must be kept at reset value.

Bits 20:17 PKTSTS: Packet status

Indicates the status of the received packet

0010: IN data packet received

0011: IN transfer completed (triggers an interrupt)0101: Data toggle error (triggers an interrupt)0111: Channel halted (triggers an interrupt)

Others: Reserved

Bits 16:15 DPID: Data PID

Indicates the Data PID of the received packet

00: DATA0 10: DATA1 01: DATA2 11: MDATA

Bits 14:4 BCNT: Byte count

Indicates the byte count of the received IN data packet.

Bits 3:0 CHNUM: Channel number

Indicates the channel number to which the current received packet belongs.

## Peripheral mode:

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	FRMNUM	PKTSTS	DPID	BCNT	EPNUM
rieserveu	r	ŗ	r	r	r

Bits 31:25 Reserved, must be kept at reset value.

Bits 24:21 FRMNUM: Frame number

This is the least significant 4 bits of the frame number in which the packet is received on the USB. This field is supported only when isochronous OUT endpoints are supported.

Bits 20:17 PKTSTS: Packet status

Indicates the status of the received packet 0001: Global OUT NAK (triggers an interrupt)

0010: OUT data packet received

0011: OUT transfer completed (triggers an interrupt)0100: SETUP transaction completed (triggers an interrupt)

0110: SETUP data packet received

Others: Reserved

Bits 16:15 DPID: Data PID

Indicates the Data PID of the received OUT data packet

00: DATA0 10: DATA1 01: DATA2 11: MDATA

Bits 14:4 BCNT: Byte count

Indicates the byte count of the received data packet.

Bits 3:0 **EPNUM:** Endpoint number

Indicates the endpoint number to which the current received packet belongs.

## OTG\_HS Receive FIFO size register (OTG\_HS\_GRXFSIZ)

Address offset: 0x024

Reset value: 0x0000 0200

The application can program the RAM size that must be allocated to the RxFIFO.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	RXFD
i iesei veu	r/rw

Bits 31:16 Reserved, must be kept at reset value.

Bits 15:0 RXFD: RxFIFO depth

This value is in terms of 32-bit words.

Minimum value is 16 Maximum value is 1024

The power-on reset value of this register is specified as the largest Rx data FIFO depth.

# OTG\_HS nonperiodic transmit FIFO size/Endpoint 0 transmit FIFO size register (OTG\_HS\_GNPTXFSIZ/OTG\_HS\_TX0FSIZ)

Address offset: 0x028

Reset value: 0x0000 0200

## Host mode:

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

NPTXFD	NPTXFSA
r/rw	r/rw

Bits 31:16 NPTXFD: Nonperiodic TxFIFO depth

This value is in terms of 32-bit words.

Minimum value is 16 Maximum value is 1024

Bits 15:0 NPTXFSA: Nonperiodic transmit RAM start address

This field contains the memory start address for nonperiodic transmit FIFO RAM.

## Peripheral mode:

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

TX0FD	TX0FSA
r/rw	r/rw

Bits 31:16 TOXFD: Endpoint 0 TxFIFO depth

This value is in terms of 32-bit words.

Minimum value is 16 Maximum value is 256

Bits 15:0 TX0FSA: Endpoint 0 transmit RAM start address

This field contains the memory start address for Endpoint 0 transmit FIFO RAM.

# OTG\_HS nonperiodic transmit FIFO/queue status register (OTG\_HS\_GNPTXSTS)

Address offset: 0x02C

Reset value: 0x0008 0200

Note: In peripheral mode, this register is not valid.

This read-only register contains the free space information for the nonperiodic TxFIFO and

the nonperiodic transmit request queue.

 $31 \quad 30 \quad 29 \quad 28 \quad 27 \quad 26 \quad 25 \quad 24 \quad 23 \quad 22 \quad 21 \quad 20 \quad 19 \quad 18 \quad 17 \quad 16 \quad 15 \quad 14 \quad 13 \quad 12 \quad 11 \quad 10 \quad 9 \quad 8 \quad 7 \quad 6 \quad 5 \quad 4 \quad 3 \quad 2 \quad 1 \quad 0$ 

ved	NPTXQTOP	NPTQXSAV	NPTXFSAV
Reser	r	r	r

Bit 31 Reserved, must be kept at reset value.

## Bits 30:24 NPTXQTOP: Top of the nonperiodic transmit request queue

Entry in the nonperiodic Tx request queue that is currently being processed by the MAC.

Bits [30:27]: Channel/endpoint number

Bits [26:25]:

- 00: IN/OUT token

01: Zero-length transmit packet (device IN/host OUT)

10: PING/CSPLIT token

11: Channel halt command

Bit [24]: Terminate (last entry for selected channel/endpoint)

## Bits 23:16 NPTQXSAV: Nonperiodic transmit request queue space available

Indicates the amount of free space available in the nonperiodic transmit request queue. This queue holds both IN and OUT requests in host mode. Peripheral mode has only IN requests.

00: Nonperiodic transmit request queue is full

01: dx1 location available10: dx2 locations available

*bxn*: dxn locations available  $(0 \le n \le dx8)$ 

Others: Reserved

#### Bits 15:0 NPTXFSAV: Nonperiodic TxFIFO space available

Indicates the amount of free space available in the nonperiodic TxFIFO.

Values are in terms of 32-bit words. 00: Nonperiodic TxFIFO is full 01: dx1 word available

10: dx2 words available

0xn: dxn words available (where  $0 \le n \le dx1024$ )

Others: Reserved

## OTG\_HS I<sup>2</sup>C access register (OTG\_HS\_GI2CCTL)

Address offset: 0x030 Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	/	ь	5	4	3	2	1	U
BSYDNE	RW	Reserved	12CDATSE0	-	DEV DR	Reserved	ACK	ISCEN			Α	DDF	3					R	EG	ADD	R					ı	RWE	OATA			
rw	rw		rw	rw	rw		rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

## Bit 31 BSYDNE: I2C Busy/Done

The application sets this bit to 1 to start a request on the  $I^2C$  interface. When the transfer is complete, the core deasserts this bit to 0. As long as the bit is set indicating that the  $I^2C$  interface is busy, the application cannot start another request on the interface.

#### Bit 30 RW: Read/Write Indicator

This bit indicates whether a read or write register transfer must be performed on the interface.

0: Write 1: Read

Note: Read/write bursting is not supported for registers.

Bit 29 Reserved, must be kept at reset value.

## Bit 28 I2CDATSE0: I<sup>2</sup>C DatSe0 USB mode

This bit is used to select the full-speed interface USB mode.

0: VP\_VM USB mode 1: DAT\_SE0 USB mode

## Bits 27:26 I2CDEVADR: I2C Device Address

This bit selects the address of the I<sup>2</sup>C slave on the USB 1.1 full-speed serial transceiver corresponding to the one used by the core for OTG signalling.

Bit 25 Reserved, must be kept at reset value.

## Bit 24 ACK: I<sup>2</sup>C ACK

This bit indicates whether an ACK response was received from the I<sup>2</sup>C slave. It is valid when BSYDNE is cleared by the core, after the application has initiated an I<sup>2</sup>C access.

0: NAK 1: ACK

## Bit 23 I2CEN: I<sup>2</sup>C Enable

This bit enables the I<sup>2</sup>C master to initiate transactions on the I<sup>2</sup>C interface.

## Bits 22:16 ADDR: I<sup>2</sup>C Address

This is the 7-bit I<sup>2</sup>C device address used by the application to access any external I<sup>2</sup>C slave, including the I<sup>2</sup>C slave on a USB 1.1 OTG full-speed serial transceiver.

## Bits 15:8 REGADDR: I<sup>2</sup>C Register Address

These bits allow to program the address of the register to be read from or written to.

## Bits 7:0 RWDATA: I<sup>2</sup>C Read/Write Data

After a register read operation, these bits hold the read data for the application.

During a write operation, the application can use this register to program the data to be written to a register.

## OTG\_HS general core configuration register (OTG\_HS\_GCCFG)

Address offset: 0x038

Reset value: 0x0000 0000

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	NOVBUSSENS	SOFOUTEN	VBUSBSEN	VBUSASEN	I2CPADEN	NWGRWY.	Reserved
	rw	rw	rw	rw	rw	rw	

Bits 31:22 Reserved, must be kept at reset value.

## Bit 21 NOVBUSSENS: V<sub>BUS</sub> sensing disable option

When this bit is set,  $V_{BUS}$  is considered internally to be always at  $V_{BUS}$  valid level (5 V). This option removes the need for a dedicated  $V_{BUS}$  pad, and leave this pad free to be used for other purposes such as a shared functionality.  $V_{BUS}$  connection can be remapped on another general purpose input pad and monitored by software.

This option is only suitable for host-only or device-only applications.

- 0: V<sub>BUS</sub> sensing available by hardware
- 1: V<sub>BUS</sub> sensing not available by hardware.
- Bit 20 SOFOUTEN: SOF output enable
  - 0: SOF pulse not available on PAD
  - 1: SOF pulse available on PAD
- Bit 19 VBUSBSEN: Enable the V<sub>BUS</sub> sensing "B" device
  - 0: V<sub>BUS</sub> sensing "B" disabled
  - 1:  $V_{BUS}$  sensing "B" enabled
- Bit 18 VBUSASEN: Enable the V<sub>BUS</sub> sensing "A" device
  - 0: V<sub>BUS</sub> sensing "A" disabled
  - 1: V<sub>BUS</sub> sensing "A" enabled
- Bit 17 **I2CPADEN**: Enable I<sup>2</sup>C bus connection for the external I<sup>2</sup>C PHY interface.
  - 0: I<sup>2</sup>C bus disabled
  - 1: I<sup>2</sup>C bus enabled
- Bit 16 PWRDWN: Power down

Used to activate the transceiver in transmission/reception

- 0: Power down active
- 1: Power down deactivated ("Transceiver active")

Bits 15:0 Reserved, must be kept at reset value.

## OTG\_HS core ID register (OTG\_HS\_CID)

Address offset: 0x03C Reset value:0x0000 1200

This is a read only register containing the Product ID.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

														PF	RODU	JCT_	ID														
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw															

Bits 31:0 **PRODUCT\_ID:** Product ID field

Application-programmable ID field.

## OTG\_HS Host periodic transmit FIFO size register (OTG\_HS\_HPTXFSIZ)

Address offset: 0x100

Reset value: 0x0200 0600

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

							PT	KFD															PT	(SA							
r/r	r/r	r/r	r/r	r/r		r/r	r/r	r/r		r/r		r/r	r/r	r/r		r/r	r/r	r/r	r/r	-	r/r	-	r/r		r/r	r/r	r/r	r/r		r/r	r/r
W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W	W

Bits 31:16 PTXFD: Host periodic TxFIFO depth

This value is in terms of 32-bit words.

Minimum value is 16 Maximum value is 512

Bits 15:0 PTXSA: Host periodic TxFIFO start address

The power-on reset value of this register is the sum of the largest Rx data FIFO depth and largest nonperiodic Tx data FIFO depth.

# OTG\_HS device IN endpoint transmit FIFO size register (OTG\_HS\_DIEPTXFx) $(x = 1..7, where x is the FIFO_number)$

Address offset: 0x104 + (FIFO\_number - 1) × 0x04

Reset value: 0x02000400

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

INEPTXFD	INEPTXSA
r/rw r/rw r/rw r/rw r/rw r/rw r/rw r/rw	r/rw r/rw r/rw r/rw r/rw r/rw r/rw r/rw

Bits 31:16 INEPTXFD: IN endpoint TxFIFO depth

This value is in terms of 32-bit words.

Minimum value is 16 Maximum value is 512

The power-on reset value of this register is specified as the largest IN endpoint FIFO number depth.

## Bits 15:0 INEPTXSA: IN endpoint FIFOx transmit RAM start address

This field contains the memory start address for IN endpoint transmit FIFOx. The address must be aligned with a 32-bit memory location.

## 30.12.3 Host-mode registers

Bit values in the register descriptions are expressed in binary unless otherwise specified.

Host-mode registers affect the operation of the core in the host mode. Host mode registers must not be accessed in peripheral mode, as the results are undefined. Host mode registers can be categorized as follows:

## OTG\_HS host configuration register (OTG\_HS\_HCFG)

Address offset: 0x400

Reset value: 0x0000 0000

This register configures the core after power-on. Do not change to this register after initializing the host.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
														Reserved															FSLSS	ادا دواد	
														ш															r	rw	rw

Bits 31:3 Reserved, must be kept at reset value.

## Bit 2 FSLSS: FS- and LS-only support

The application uses this bit to control the core's enumeration speed. Using this bit, the application can make the core enumerate as an FS host, even if the connected device supports HS traffic. Do not make changes to this field after initial programming.

0: HS/FS/LS, based on the maximum speed supported by the connected device

1: FS/LS-only, even if the connected device can support HS (read-only)

## Bits 1:0 FSLSPCS: FS/LS PHY clock select

When the core is in FS host mode:

01: PHY clock is running at 48 MHz

Others: Reserved

When the core is in LS host mode:

00: Reserved

01: PHY clock is running at 48 MHz.

10: Select 6 MHz PHY clock frequency

11: Reserved

Note: The FSLSPCS bit must be set on a connection event according to the speed of the connected device. A software reset must be performed after changing this bit.

## OTG\_HS Host frame interval register (OTG\_HS\_HFIR)

Address offset: 0x404

Reset value: 0x0000 EA60

This register stores the frame interval information for the current speed to which the

OTG\_HS controller has enumerated.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved								FRI	IVL							
neserveu	rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bits 31:16 Reserved, must be kept at reset value.

#### Bits 15:0 FRIVL: Frame interval

The value that the application programs to this field specifies the interval between two consecutive SOFs (FS), micro-SOFs (HS) or Keep-Alive tokens (LS). This field contains the number of PHY clocks that constitute the required frame interval. The application can write a value to this register only after the Port enable bit of the host port control and status register (PENA bit in OTG\_HS\_HPRT) has been set. If no value is programmed, the core calculates the value based on the PHY clock specified in the FS/LS PHY Clock Select field of the Host configuration register (FSLSPCS in OTG\_HS\_HCFG):

frame duration × PHY clock frequency

Note: The FRIVL bit can be modified whenever the application needs to change the Frame interval time.

## OTG\_HS host frame number/frame time remaining register (OTG\_HS\_HFNUM)

Address offset: 0x408

Reset value: 0x0000 3FFF

This register indicates the current frame number. It also indicates the time remaining (in terms of the number of PHY clocks) in the current frame.

 $31 \quad 30 \quad 29 \quad 28 \quad 27 \quad 26 \quad 25 \quad 24 \quad 23 \quad 22 \quad 21 \quad 20 \quad 19 \quad 18 \quad 17 \quad 16 \quad 15 \quad 14 \quad 13 \quad 12 \quad 11 \quad 10 \quad 9 \quad 8 \quad 7 \quad 6 \quad 5 \quad 4 \quad 3 \quad 2 \quad 1 \quad 0$ 

								FTF	REM															FRN	MUI							
ı	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

#### Bits 31:16 FTREM: Frame time remaining

Indicates the amount of time remaining in the current frame, in terms of PHY clocks. This field decrements on each PHY clock. When it reaches zero, this field is reloaded with the value in the Frame interval register and a new SOF is transmitted on the USB.

## Bits 15:0 FRNUM: Frame number

This field increments when a new SOF is transmitted on the USB, and is cleared to 0 when it reaches 0x3FFF.

## OTG\_HS\_Host periodic transmit FIFO/queue status register (OTG\_HS\_HPTXSTS)

Address offset: 0x410
Reset value: 0x0008 0100

This read-only register contains the free space information for the periodic TxFIFO and the periodic transmit request queue.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

I				PTXC	OTC	>					F	PTXC	QSAV	/									F	TXF	SAV	L						
	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	rw	rw	rw	rw	rw	rw	rw	rw	rw							

## Bits 31:24 PTXQTOP: Top of the periodic transmit request queue

This indicates the entry in the periodic Tx request queue that is currently being processed by the MAC.

This register is used for debugging.

Bit [31]: Odd/Even frame

- 0: send in even (micro) frame
- 1: send in odd (micro) frame

Bits [30:27]: Channel/endpoint number

Bits [26:25]: Type

- 00: IN/OUT
- 01: Zero-length packet
- 11: Disable channel command

Bit [24]: Terminate (last entry for the selected channel/endpoint)

## Bits 23:16 PTXQSAV: Periodic transmit request queue space available

Indicates the number of free locations available to be written in the periodic transmit request queue. This queue holds both IN and OUT requests.

00: Periodic transmit request queue is full

01: dx1 location available 10: dx2 locations available

bxn: dxn locations available ( $0 \le dxn \le PTXFD$ )

Others: Reserved

## Bits 15:0 PTXFSAVL: Periodic transmit data FIFO space available

Indicates the number of free locations available to be written to in the periodic TxFIFO.

Values are in terms of 32-bit words

0000: Periodic TxFIFO is full 0001: dx1 word available 0010: dx2 words available

bxn: dxn words available (where  $0 \le dxn \le dx512$ )

Others: Reserved

## OTG\_HS Host all channels interrupt register (OTG\_HS\_HAINT)

Address offset: 0x414
Reset value: 0x0000 000

When a significant event occurs on a channel, the host all channels interrupt register interrupts the application using the host channels interrupt bit of the Core interrupt register (HCINT bit in OTG\_HS\_GINTSTS). This is shown in *Figure 376*. There is one interrupt bit per channel, up to a maximum of 16 bits. Bits in this register are set and cleared when the application sets and clears bits in the corresponding host channel-x interrupt register.

3	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							Daa																НА	INT							
							nese	erved	l							r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:16 Reserved, must be kept at reset value.

Bits 15:0 HAINT: Channel interrupts

One bit per channel: Bit 0 for Channel 0, bit 15 for Channel 15

## OTG\_HS host all channels interrupt mask register (OTG\_HS\_HAINTMSK)

Address offset: 0x418 Reset value: 0x0000 0000

The host all channel interrupt mask register works with the host all channel interrupt register to interrupt the application when an event occurs on a channel. There is one interrupt mask bit per channel, up to a maximum of 16 bits.

	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Γ								Dage																HAII	NTM							
								nese	ervec	1							rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bits 31:16 Reserved, must be kept at reset value.

Bits 15:0 HAINTM: Channel interrupt mask

0: Masked interrupt1: Unmasked interrupt

One bit per channel: Bit 0 for channel 0, bit 15 for channel 15

## OTG\_HS host port control and status register (OTG\_HS\_HPRT)

Address offset: 0x440 Reset value: 0x0000 0000

This register is available only in host mode. Currently, the OTG host supports only one port.

A single register holds USB port-related information such as USB reset, enable, suspend, resume, connect status, and test mode for each port. It is shown in *Figure 376*. The rc\_w1 bits in this register can trigger an interrupt to the application through the host port interrupt bit of the core interrupt register (HPRTINT bit in OTG\_HS\_GINTSTS). On a Port Interrupt, the application must read this register and clear the bit that caused the interrupt. For the rc\_w1 bits, the application must write a 1 to the bit to clear the interrupt.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	PS	PD		PTO	CTL		PPWR	OFO IO	3	eserved	PRST	PSUSP	PRES	POCCHNG	POCA	PENCHNG	PENA	PCDET	PCSTS
	r	r	rw	rw	rw	rw	rw	r	r	Œ	rw	rs	rw	rc_ w1		rc_ w1	rc_ w0	rc_ w1	r

Bits 31:19 Reserved, must be kept at reset value.

Bits 18:17 PSPD: Port speed

Indicates the speed of the device attached to this port.

00: High speed 01: Full speed 10: Low speed 11: Reserved

### Bits 16:13 PTCTL: Port test control

The application writes a nonzero value to this field to put the port into a Test mode, and the corresponding pattern is signaled on the port.

0000: Test mode disabled 0001: Test\_J mode 0010: Test\_K mode

0011: Test\_SE0\_NAK mode 0100: Test\_Packet mode 0101: Test\_Force\_Enable Others: Reserved

## Bit 12 PPWR: Port power

The application uses this field to control power to this port, and the core clears this bit on an overcurrent condition.

0: Power off1: Power on

## Bits 11:10 PLSTS: Port line status

Indicates the current logic level USB data lines Bit [10]: Logic level of OTG\_HS\_FS\_DP Bit [11]: Logic level of OTG\_HS\_FS\_DM

Bit 9 Reserved, must be kept at reset value.

#### Bit 8 PRST: Port reset

When the application sets this bit, a reset sequence is started on this port. The application must time the reset period and clear this bit after the reset sequence is complete.

0: Port not in reset

1: Port in reset

The application must leave this bit set for a minimum duration of at least 10 ms to start a reset on the port. The application can leave it set for another 10 ms in addition to the required minimum duration, before clearing the bit, even though there is no maximum limit set by the USB standard.

High speed: 50 ms

Full speed/Low speed: 10 ms

## Bit 7 PSUSP: Port suspend

The application sets this bit to put this port in Suspend mode. The core only stops sending SOFs when this is set. To stop the PHY clock, the application must set the Port clock stop bit, which asserts the suspend input pin of the PHY.

The read value of this bit reflects the current suspend status of the port. This bit is cleared by the core after a remote wakeup signal is detected or the application sets the Port reset bit or Port resume bit in this register or the Resume/remote wakeup detected interrupt bit or Disconnect detected interrupt bit in the Core interrupt register (WKUINT or DISCINT in OTG\_HS\_GINTSTS, respectively).

0: Port not in Suspend mode

1: Port in Suspend mode

#### Bit 6 PRES: Port resume

The application sets this bit to drive resume signaling on the port. The core continues to drive the resume signal until the application clears this bit.

If the core detects a USB remote wakeup sequence, as indicated by the Port resume/remote wakeup detected interrupt bit of the Core interrupt register (WKUINT bit in OTG\_HS\_GINTSTS), the core starts driving resume signaling without application intervention and clears this bit when it detects a disconnect condition. The read value of this bit indicates whether the core is currently driving resume signaling.

0: No resume driven

1: Resume driven

## Bit 5 POCCHNG: Port overcurrent change

The core sets this bit when the status of the Port overcurrent active bit (bit 4) in this register changes.

## Bit 4 POCA: Port overcurrent active

Indicates the overcurrent condition of the port.

0: No overcurrent condition

1: Overcurrent condition

## Bit 3 PENCHNG: Port enable/disable change

The core sets this bit when the status of the Port enable bit [2] in this register changes.



#### Bit 2 PENA: Port enable

A port is enabled only by the core after a reset sequence, and is disabled by an overcurrent condition, a disconnect condition, or by the application clearing this bit. The application cannot set this bit by a register write. It can only clear it to disable the port. This bit does not trigger any interrupt to the application.

0: Port disabled

## 1: Port enabled

#### Bit 1 PCDET: Port connect detected

The core sets this bit when a device connection is detected to trigger an interrupt to the application using the host port interrupt bit in the Core interrupt register (HPRTINT bit in OTG\_HS\_GINTSTS). The application must write a 1 to this bit to clear the interrupt.

## Bit 0 PCSTS: Port connect status

0: No device is attached to the port

1: A device is attached to the port

# OTG\_HS host channel-x characteristics register (OTG\_HS\_HCCHARx) (x = 0..11, where x = Channel\_number)

Address offset: 0x500 + (Channel\_number × 0x20)

Reset value: 0x0000 0000

 $31 \quad 30 \quad 29 \quad 28 \quad 27 \quad 26 \quad 25 \quad 24 \quad 23 \quad 22 \quad 21 \quad 20 \quad 19 \quad 18 \quad 17 \quad 16 \quad 15 \quad 14 \quad 13 \quad 12 \quad 11 \quad 10 \quad 9 \quad 8 \quad 7 \quad 6 \quad 5 \quad 4 \quad 3 \quad 2 \quad 1 \quad 0$ 

CHENA	CHDIS	ODDFRM				DAD	ı			М	С	ļ	L L L	LSDEV	leserved	EPDIR		EPN	NUM						N	/IPSI	Z				
rs	rs	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	ш	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

## Bit 31 CHENA: Channel enable

This field is set by the application and cleared by the OTG host.

0: Channel disabled

1: Channel enabled

## Bit 30 CHDIS: Channel disable

The application sets this bit to stop transmitting/receiving data on a channel, even before the transfer for that channel is complete. The application must wait for the Channel disabled interrupt before treating the channel as disabled.

#### Bit 29 ODDFRM: Odd frame

This field is set (reset) by the application to indicate that the OTG host must perform a transfer in an odd frame. This field is applicable for only periodic (isochronous and interrupt) transactions.

0: Even (micro) frame

1: Odd (micro) frame

## Bits 28:22 DAD: Device address

This field selects the specific device serving as the data source or sink.

## Bits 21:20 MC: Multi Count (MC) / Error Count (EC)

- When the split enable bit (SPLITEN) in the host channel-x split control register (OTG\_HS\_HCSPLTx) is reset (0), this field indicates to the host the number of transactions that must be executed per micro-frame for this periodic endpoint. For nonperiodic transfers, this field specifies the number of packets to be fetched for this channel before the internal DMA engine changes arbitration.

00: Reserved This field yields undefined results

01: 1 transaction

b10: 2 transactions to be issued for this endpoint per micro-frame

11: 3 transactions to be issued for this endpoint per micro-frame.

 When the SPLITEN bit is set (1) in OTG\_HS\_HCSPLTx, this field indicates the number of immediate retries to be performed for a periodic split transaction on transaction errors. This field must be set to at least 01.

## Bits 19:18 EPTYP: Endpoint type

Indicates the transfer type selected.

00: Control 01: Isochronous 10: Bulk 11: Interrupt

## Bit 17 LSDEV: Low-speed device

This field is set by the application to indicate that this channel is communicating to a low-speed device.

Bit 16 Reserved, must be kept at reset value.

## Bit 15 EPDIR: Endpoint direction

Indicates whether the transaction is IN or OUT.

0: OUT 1: IN

## Bits 14:11 EPNUM: Endpoint number

Indicates the endpoint number on the device serving as the data source or sink.

#### Bits 10:0 MPSIZ: Maximum packet size

Indicates the maximum packet size of the associated endpoint.



## OTG\_HS host channel-x split control register (OTG\_HS\_HCSPLTx) (x = 0...11, where $x = Channel_number$ )

Address offset: 0x504 + (Channel number × 0x20)

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
SPLITEN							3	Heserved							COMPLSPLT	OCETOW.	VACITOS				HUBADDR							PRTADDR			
rw															rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

#### Bit 31 SPLITEN: Split enable

The application sets this bit to indicate that this channel is enabled to perform split transactions.

Bits 30:17 Reserved, must be kept at reset value.

## Bit 16 COMPLSPLT: Do complete split

The application sets this bit to request the OTG host to perform a complete split transaction.

## Bits 15:14 XACTPOS: Transaction position

This field is used to determine whether to send all, first, middle, or last payloads with each OUT transaction.

11: All. This is the entire data payload of this transaction (which is less than or equal to 188 bytes)

10: Begin. This is the first data payload of this transaction (which is larger than 188 bytes)

00: Mid. This is the middle payload of this transaction (which is larger than 188 bytes)

01: End. This is the last payload of this transaction (which is larger than 188 bytes)

#### Bits 13:7 HUBADDR: Hub address

This field holds the device address of the transaction translator's hub.

## Bits 6:0 PRTADDR: Port address

This field is the port number of the recipient transaction translator.

## OTG\_HS host channel-x interrupt register (OTG\_HS\_HCINTx) (x = 0..11, where $x = Channel\_number$ )

Address offset: 0x508 + (Channel number × 0x20)

Reset value: 0x0000 0000

This register indicates the status of a channel with respect to USB- and AHB-related events. It is shown in *Figure 376*. The application must read this register when the host channels interrupt bit in the Core interrupt register (HCINT bit in OTG\_HS\_GINTSTS) is set. Before the application can read this register, it must first read the host all channels interrupt (OTG\_HS\_HAINT) register to get the exact channel number for the host channel-x interrupt register. The application must clear the appropriate bit in this register to clear the corresponding bits in the OTG\_HS\_HAINT and OTG\_HS\_GINTSTS registers.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

וי	50	23	20	21	20	23	24	20	~~	21	20	13	10	17	10	13	1-	10	12	 10	9	U	,	U	J	-	J	_		U
									Re	eserv	ed									DTERR	FRMOR	нызва	TXERR	NYET	ACK	NAK	STALL	AHBERR	СНН	XFRC
																					rc_ w1		rc_ w1			rc_ w1			rc_ w1	

Bits 31:11 Reserved, must be kept at reset value.

Bit 10 DTERR: Data toggle error

Bit 9 **FRMOR:** Frame overrun
Bit 8 **BBERR:** Babble error

Bit 7 TXERR: Transaction error

Indicates one of the following errors occurred on the USB.

CRC check failure

Timeout Bit stuff error False EOP

Bit 6 NYET: Response received interrupt

Bit 5 ACK: ACK response received/transmitted interrupt

Bit 4 NAK: NAK response received interrupt

Bit 3 STALL: STALL response received interrupt

Bit 2 AHBERR: AHB error

This error is generated only in Internal DMA mode when an AHB error occurs during an AHB read/write operation. The application can read the corresponding DMA channel address register to get the error address.

Bit 1 CHH: Channel halted

Indicates the transfer completed abnormally either because of any USB transaction error or in response to disable request by the application.

Bit 0 XFRC: Transfer completed

Transfer completed normally without any errors.

## OTG\_HS host channel-x interrupt mask register (OTG\_HS\_HCINTMSKx) $(x = 0..11, where x = Channel_number)$

Address offset: 0x50C + (Channel\_number × 0x20)

Reset value: 0x0000 0000

This register reflects the mask for each channel status described in the previous section.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved

Bits 31:11 Reserved, must be kept at reset value.

Bit 10 DTERRM: Data toggle error mask

0: Masked interrupt1: Unmasked interrupt

Bit 9 FRMORM: Frame overrun mask

0: Masked interrupt1: Unmasked interrupt

Bit 8 BBERRM: Babble error mask

0: Masked interrupt1: Unmasked interrupt

Bit 7 TXERRM: Transaction error mask

0: Masked interrupt1: Unmasked interrupt

Bit 6 NYET: response received interrupt mask

0: Masked interrupt1: Unmasked interrupt

Bit 5 ACKM: ACK response received/transmitted interrupt mask

0: Masked interrupt1: Unmasked interrupt

Bit 4 NAKM: NAK response received interrupt mask

0: Masked interrupt1: Unmasked interrupt

Bit 3 STALLM: STALL response received interrupt mask

0: Masked interrupt1: Unmasked interrupt

#### Bit 2 AHBERR: AHB error

This is generated only in Internal DMA mode when there is an AHB error during AHB read/write. The application can read the corresponding channel's DMA address register to get the error address.

#### Bit 1 CHHM: Channel halted mask

0: Masked interrupt1: Unmasked interrupt

### Bit 0 XFRCM: Transfer completed mask

0: Masked interrupt1: Unmasked interrupt

## OTG\_HS host channel-x transfer size register (OTG\_HS\_HCTSIZx) (x = 0..11, where $x = Channel_number$ )

Address offset: 0x510 + (Channel\_number × 0x20)

Reset value: 0x0000 0000

## $31 \quad 30 \quad 29 \quad 28 \quad 27 \quad 26 \quad 25 \quad 24 \quad 23 \quad 22 \quad 21 \quad 20 \quad 19 \quad 18 \quad 17 \quad 16 \quad 15 \quad 14 \quad 13 \quad 12 \quad 11 \quad 10 \quad 9 \quad 8 \quad 7 \quad 6 \quad 5 \quad 4 \quad 3 \quad 2 \quad 1 \quad 0$

Reserved	DF	PID					PKT	CNT													Х	FRS	ΙΖ								
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

## Bit 31 DOPING: Do ping

This bit is used only for OUT transfers. Setting this field to 1 directs the host to do PING protocol.

Note: Do not set this bit for IN transfers. If this bit is set for IN transfers it disables the channel.

#### Bits 30:29 DPID: Data PID

The application programs this field with the type of PID to use for the initial transaction. The host maintains this field for the rest of the transfer.

00: DATA0 01: DATA2 10: DATA1

11: MDATA (noncontrol)/SETUP (control)

### Bits 28:19 PKTCNT: Packet count

This field is programmed by the application with the expected number of packets to be transmitted (OUT) or received (IN).

The host decrements this count on every successful transmission or reception of an OUT/IN packet. Once this count reaches zero, the application is interrupted to indicate normal completion.

## Bits 18:0 XFRSIZ: Transfer size

For an OUT, this field is the number of data bytes the host sends during the transfer. For an IN, this field is the buffer size that the application has reserved for the transfer. The application is expected to program this field as an integer multiple of the maximum packet size for IN transactions (periodic and nonperiodic).

## OTG\_HS host channel-x DMA address register (OTG\_HS\_HCDMAx) (x = 0..11, where x = Channel number)

Address offset: 0x514 + (Channel\_number × 0x20)

Reset value: 0x0000 0000

Bits 31:0 DMAADDR: DMA address

This field holds the start address in the external memory from which the data for the endpoint must be fetched or to which it must be stored. This register is incremented on every AHB transaction.

## 30.12.4 Device-mode registers

## OTG HS device configuration register (OTG HS DCFG)

Address offset: 0x800

Reset value: 0x0220 0000

This register configures the core in peripheral mode after power-on or after certain control commands or enumeration. Do not make changes to this register after initial programming.

31	1 ;	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	/	6	5	4	3	2	1	0
			Rese	erved	I		MINUSCH	5	Reserved					Rese	ervec	I				7,11	PFIVL				DAD				Reserved	NSHOSTZN	מפטט	7 1
							rw	rw												r	w	rw	rw	rw	rw	rw	rw	rw		rw	rw	rw

Bits 31:26 Reserved, must be kept at reset value.

### Bits 25:24 PERSCHIVL: Periodic scheduling interval

This field specifies the amount of time the Internal DMA engine must allocate for fetching periodic IN endpoint data. Based on the number of periodic endpoints, this value must be specified as 25, 50 or 75% of the (micro)frame.

- When any periodic endpoints are active, the internal DMA engine allocates the specified amount of time in fetching periodic IN endpoint data
- When no periodic endpoint is active, then the internal DMA engine services nonperiodic endpoints, ignoring this field
- After the specified time within a (micro)frame, the DMA switches to fetching nonperiodic endpoints

00: 25% of (micro)frame 01: 50% of (micro)frame 10: 75% of (micro)frame

11: Reserved

Bits 23:13 Reserved, must be kept at reset value.

## Bits 12:11 PFIVL: Periodic (micro)frame interval

Indicates the time within a (micro) frame at which the application must be notified using the end of periodic (micro) frame interrupt. This can be used to determine if all the isochronous traffic for that frame is complete.

00: 80% of the frame interval

01: 85% of the frame interval

10: 90% of the frame interval

11: 95% of the frame interval

## Bits 10:4 DAD: Device address

The application must program this field after every SetAddress control command.

Bit 3 Reserved, must be kept at reset value.

## Bit 2 NZLSOHSK: Nonzero-length status OUT handshake

The application can use this field to select the handshake the core sends on receiving a nonzero-length data packet during the OUT transaction of a control transfer's Status stage.

- 1: Send a STALL handshake on a nonzero-length status OUT transaction and do not send the received OUT packet to the application.
- 0: Send the received OUT packet to the application (zero-length or nonzero-length) and send a handshake based on the NAK and STALL bits for the endpoint in the device endpoint control register.

## Bits 1:0 DSPD: Device speed

Indicates the speed at which the application requires the core to enumerate, or the maximum speed the application can support. However, the actual bus speed is determined only after the chirp sequence is completed, and is based on the speed of the USB host to which the core is connected.

00: High speed 01: Reserved 10: Reserved

11: Full speed (USB 1.1 transceiver clock is 48 MHz)



## OTG\_HS device control register (OTG\_HS\_DCTL)

Address offset: 0x804

Reset value: 0x0000 0000

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	POPRGDNE	CGONAK	SGONAK	CGINAK	SGINAK		TCTL		SLSNOÐ	GINSTS	SIQS	RWUSIG
	rw	w	w	w	w	rw	rw	rw	r	r	rw	rw

## Bits 31:12 Reserved, must be kept at reset value.

## Bit 11 POPRGDNE: Power-on programming done

The application uses this bit to indicate that register programming is completed after a wakeup from power down mode.

### Bit 10 CGONAK: Clear global OUT NAK

A write to this field clears the Global OUT NAK.

## Bit 9 SGONAK: Set global OUT NAK

A write to this field sets the Global OUT NAK.

The application uses this bit to send a NAK handshake on all OUT endpoints. The application must set the this bit only after making sure that the Global OUT NAK effective bit in the Core interrupt register (GONAKEFF bit in OTG\_HS\_GINTSTS) is cleared.

#### Bit 8 CGINAK: Clear global IN NAK

A write to this field clears the Global IN NAK.

## Bit 7 SGINAK: Set global IN NAK

A write to this field sets the Global nonperiodic IN NAK. The application uses this bit to send a NAK handshake on all nonperiodic IN endpoints.

The application must set this bit only after making sure that the Global IN NAK effective bit in the Core interrupt register (GINAKEFF bit in OTG\_HS\_GINTSTS) is cleared.

## Bits 6:4 TCTL: Test control

000: Test mode disabled

001: Test\_J mode

010: Test\_K mode

011: Test\_SE0\_NAK mode

100: Test\_Packet mode

101: Test\_Force\_Enable

Others: Reserved

### Bit 3 GONSTS: Global OUT NAK status

0: A handshake is sent based on the FIFO Status and the NAK and STALL bit settings.

1: No data is written to the RxFIFO, irrespective of space availability. Sends a NAK handshake on all packets, except on SETUP transactions. All isochronous OUT packets are dropped.

#### Bit 2 GINSTS: Global IN NAK status

- 0: A handshake is sent out based on the data availability in the transmit FIFO.
- 1: A NAK handshake is sent out on all nonperiodic IN endpoints, irrespective of the data availability in the transmit FIFO.

## Bit 1 SDIS: Soft disconnect

The application uses this bit to signal the USB OTG core to perform a soft disconnect. As long as this bit is set, the host does not see that the device is connected, and the device does not receive signals on the USB. The core stays in the disconnected state until the application clears this bit.

- 0: Normal operation. When this bit is cleared after a soft disconnect, the core generates a device connect event to the USB host. When the device is reconnected, the USB host restarts device enumeration.
- 1: The core generates a device disconnect event to the USB host.

## Bit 0 RWUSIG: Remote wakeup signaling

When the application sets this bit, the core initiates remote signaling to wake up the USB host. The application must set this bit to instruct the core to exit the Suspend state. As specified in the USB 2.0 specification, the application must clear this bit 1 ms to 15 ms after setting it.

Table 165 contains the minimum duration (according to device state) for which the Soft disconnect (SDIS) bit must be set for the USB host to detect a device disconnect. To accommodate clock jitter, it is recommended that the application add some extra delay to the specified minimum duration.

Table 165. Minimum duration for soft disconnect

Operating speed	Device state	Minimum duration
High speed	Not Idle or Suspended (Performing transactions)	125 µs
Full speed	Suspended	1 ms + 2.5 μs
Full speed	Idle	2.5 µs
Full speed	Not Idle or Suspended (Performing transactions)	2.5 µs

## OTG\_HS device status register (OTG\_HS\_DSTS)

Address offset: 0x808 Reset value: 0x0000 0010

This register indicates the status of the core with respect to USB-related events. It must be read on interrupts from the device all interrupts (OTG\_HS\_DAINT) register.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved							FNS	SOF							Reserved	EERR	COMINA	2	SUSPSTS
	r	r	r	r	r	r	r	r	r	r	r	r	r	r		r	r	r	r

Bits 31:22 Reserved, must be kept at reset value.

Bits 21:8 FNSOF: Frame number of the received SOF

Bits 7:4 Reserved, must be kept at reset value.

Bit 3 **EERR:** Erratic error

The core sets this bit to report any erratic errors.

Due to erratic errors, the OTG\_HS controller goes into Suspended state and an interrupt is generated to the application with Early suspend bit of the Core interrupt register (ESUSP bit in OTG\_HS\_GINTSTS). If the early suspend is asserted due to an erratic error, the application can only perform a soft disconnect recover.

## Bits 2:1 ENUMSPD: Enumerated speed

Indicates the speed at which the OTG\_HS controller has come up after speed detection through a chirp sequence.

00: High speed 01: Reserved 10: Reserved

11: Full speed (PHY clock is running at 48 MHz)

Others: reserved

## Bit 0 SUSPSTS: Suspend status

In peripheral mode, this bit is set as long as a Suspend condition is detected on the USB. The core enters the Suspended state when there is no activity on the USB data lines for a period of 3 ms. The core comes out of the suspend:

- When there is an activity on the USB data lines
- When the application writes to the Remote wakeup signaling bit in the Device control register (RWUSIG bit in OTG\_HS\_DCTL).

## OTG\_HS device IN endpoint common interrupt mask register (OTG\_HS\_DIEPMSK)

Address offset: 0x810
Reset value: 0x0000 0000

This register works with each of the Device IN endpoint interrupt (OTG\_HS\_DIEPINTx) registers for all endpoints to generate an interrupt per IN endpoint. The IN endpoint interrupt for a specific status in the OTG\_HS\_DIEPINTx register can be masked by writing to the corresponding bit in this register. Status bits are masked by default.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
										Rese	erved	l										BIM	TXFURM	Reserved	INEPNEM	INEPNMM	ITTXFEMSK	MOT	Reserved	EPDM	XFRCM
																						rw	rw		rw	rw	rw	rw		rw	rw

Bits 31:10 Reserved, must be kept at reset value.

Bit 9 BIM: BNA interrupt mask

0: Masked interrupt

1: Unmasked interrupt

Bit 8 TXFURM: FIFO underrun mask

0: Masked interrupt

1: Unmasked interrupt

Bit 7 Reserved, must be kept at reset value.

Bit 6 INEPNEM: IN endpoint NAK effective mask

0: Masked interrupt

1: Unmasked interrupt

Bit 5 INEPNMM: IN token received with EP mismatch mask

0: Masked interrupt

1: Unmasked interrupt

Bit 4 ITTXFEMSK: IN token received when TxFIFO empty mask

0: Masked interrupt

1: Unmasked interrupt

Bit 3 TOM: Timeout condition mask (nonisochronous endpoints)

0: Masked interrupt

1: Unmasked interrupt

Bit 2 Reserved, must be kept at reset value.

Bit 1 EPDM: Endpoint disabled interrupt mask

0: Masked interrupt

1: Unmasked interrupt

Bit 0 XFRCM: Transfer completed interrupt mask

0: Masked interrupt

1: Unmasked interrupt

## OTG\_HS device OUT endpoint common interrupt mask register (OTG\_HS\_DOEPMSK)

Address offset: 0x814

Reset value: 0x0000 0000

This register works with each of the Device OUT endpoint interrupt (OTG\_HS\_DOEPINTx) registers for all endpoints to generate an interrupt per OUT endpoint. The OUT endpoint interrupt for a specific status in the OTG\_HS\_DOEPINTx register can be masked by writing into the corresponding bit in this register. Status bits are masked by default.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
										Rese	erved	l										≅ BOIM	₹ OPEM	Reserved	≩ B2BSTUP	Reserved	₹ OTEPDM	STUPM ≤	Reserved	₹ EPDM	₹ XFRCM

Bits 31:10 Reserved, must be kept at reset value.

Bit 9 BOIM: BNA interrupt mask

0: Masked interrupt

1: Unmasked interrupt

Bit 8 **OPEM:** OUT packet error mask

0: Masked interrupt1: Unmasked interrupt

Bit 7 Reserved, must be kept at reset value.

Bit 6 B2BSTUP: Back-to-back SETUP packets received mask

Applies to control OUT endpoints only.

0: Masked interrupt

1: Unmasked interrupt

Bit 5 Reserved, must be kept at reset value.

Bit 4 OTEPDM: OUT token received when endpoint disabled mask

Applies to control OUT endpoints only.

0: Masked interrupt

1: Unmasked interrupt

Bit 3 STUPM: SETUP phase done mask

Applies to control endpoints only.

0: Masked interrupt

1: Unmasked interrupt

Bit 2 Reserved, must be kept at reset value.

Bit 1 EPDM: Endpoint disabled interrupt mask

0: Masked interrupt

1: Unmasked interrupt

Bit 0 XFRCM: Transfer completed interrupt mask

0: Masked interrupt

1: Unmasked interrupt

## OTG\_HS device all endpoints interrupt register (OTG\_HS\_DAINT)

Address offset: 0x818 Reset value: 0x0000 0000

When a significant event occurs on an endpoint, a device all endpoints interrupt register interrupts the application using the Device OUT endpoints interrupt bit or Device IN endpoints interrupt bit of the Core interrupt register (OEPINT or IEPINT in OTG\_HS\_GINTSTS, respectively). There is one interrupt bit per endpoint, up to a maximum of 16 bits for OUT endpoints and 16 bits for IN endpoints. For a bidirectional endpoint, the corresponding IN and OUT interrupt bits are used. Bits in this register are set and cleared when the application sets and clears bits in the corresponding Device Endpoint-x interrupt register (OTG\_HS\_DIEPINTx/OTG\_HS\_DOEPINTx).

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

							OEF	PINT															IEP	INT							
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:16 **OEPINT:** OUT endpoint interrupt bits

One bit per OUT endpoint:

Bit 16 for OUT endpoint 0, bit 31 for OUT endpoint 15

Bits 15:0 IEPINT: IN endpoint interrupt bits

One bit per IN endpoint:

Bit 0 for IN endpoint 0, bit 15 for endpoint 15

## OTG\_HS all endpoints interrupt mask register (OTG\_HS\_DAINTMSK)

Address offset: 0x81C Reset value: 0x0000 0000

The device endpoint interrupt mask register works with the device endpoint interrupt register to interrupt the application when an event occurs on a device endpoint. However, the device all endpoints interrupt (OTG\_HS\_DAINT) register bit corresponding to that interrupt is still set.

 $31 \quad 30 \quad 29 \quad 28 \quad 27 \quad 26 \quad 25 \quad 24 \quad 23 \quad 22 \quad 21 \quad 20 \quad 19 \quad 18 \quad 17 \quad 16 \quad 15 \quad 14 \quad 13 \quad 12 \quad 11 \quad 10 \quad 9 \quad 8 \quad 7 \quad 6 \quad 5 \quad 4 \quad 3 \quad 2 \quad 1 \quad 0$ 

							OE	PM															IEF	PM							
rw	rw	rw	rw	rw	rw	rw	rw	rw																							

Bits 31:16 **OEPM:** OUT EP interrupt mask bits

One per OUT endpoint:

Bit 16 for OUT EP 0, bit 18 for OUT EP 3

0: Masked interrupt1: Unmasked interrupt

Bits 15:0 IEPM: IN EP interrupt mask bits

One bit per IN endpoint:

Bit 0 for IN EP 0, bit 3 for IN EP 3

0: Masked interrupt1: Unmasked interrupt

## $OTG_HS$ device $V_{BUS}$ discharge time register ( $OTG_HS_DVBUSDIS$ )

Address offset: 0x0828 Reset value: 0x0000 17D7

This register specifies the V<sub>BUS</sub> discharge time after V<sub>BUS</sub> pulsing during SRP.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved								VBU	SDT							
neserveu	rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bits 31:16 Reserved, must be kept at reset value.

Bits 15:0 VBUSDT: Device V<sub>BUS</sub> discharge time

Specifies the  $V_{BUS}$  discharge time after  $V_{BUS}$  pulsing during SRP. This value equals:

V<sub>BUS</sub> discharge time in PHY clocks / 1 024

Depending on your V<sub>BUS</sub> load, this value may need adjusting.

## OTG\_HS device V<sub>BUS</sub> pulsing time register (OTG\_HS\_DVBUSPULSE)

Address offset: 0x082C Reset value: 0x0000 05B8

This register specifies the V<sub>BUS</sub> pulsing time during SRP.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved			_		_	DVB	USP					
neserveu	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:12 Reserved, must be kept at reset value.

Bits 11:0 **DVBUSP:** Device V<sub>BUS</sub> pulsing time

Specifies the V<sub>BUS</sub> pulsing time during SRP. This value equals:

V<sub>BUS</sub> pulsing time in PHY clocks / 1 024

# OTG\_HS Device threshold control register (OTG\_HS\_DTHRCTL)

Address offset: 0x0830 Reset value: 0x0000 0000

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	ARPEN	Reserved				RX <sup>-</sup>	ΓHRL	-EN				RXTHREN	Reserved				TX1	ΓHRL	.EN				ISOTHREN	NONISOTHREN
	rw		rw	rw	rw	rw	rw	rw	rw	rw	rw	rw		rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:28 Reserved, must be kept at reset value.

# Bit 27 ARPEN: Arbiter parking enable

This bit controls internal DMA arbiter parking for IN endpoints. When thresholding is enabled and this bit is set to one, then the arbiter parks on the IN endpoint for which there is a token received on the USB. This is done to avoid getting into underrun conditions. By default parking is enabled.

Bit 26 Reserved, must be kept at reset value.

# Bits 25: 17 RXTHRLEN: Receive threshold length

This field specifies the receive thresholding size in DWORDS. This field also specifies the amount of data received on the USB before the core can start transmitting on the AHB. The threshold length has to be at least eight DWORDS. The recommended value for RXTHRLEN is to be the same as the programmed AHB burst length (HBSTLEN bit in OTG\_HS\_GAHBCFG).

## Bit 16 RXTHREN: Receive threshold enable

When this bit is set, the core enables thresholding in the receive direction.

Bits 15: 11 Reserved, must be kept at reset value.

# Bits 10:2 TXTHRLEN: Transmit threshold length

This field specifies the transmit thresholding size in DWORDS. This field specifies the amount of data in bytes to be in the corresponding endpoint transmit FIFO, before the core can start transmitting on the USB. The threshold length has to be at least eight DWORDS. This field controls both isochronous and nonisochronous IN endpoint thresholds. The recommended value for TXTHRLEN is to be the same as the programmed AHB burst length (HBSTLEN bit in OTG\_HS\_GAHBCFG).

## Bit 1 ISOTHREN: ISO IN endpoint threshold enable

When this bit is set, the core enables thresholding for isochronous IN endpoints.

# Bit 0 NONISOTHREN: Nonisochronous IN endpoints threshold enable

When this bit is set, the core enables thresholding for nonisochronous IN endpoints.

# OTG\_HS device IN endpoint FIFO empty interrupt mask register: (OTG\_HS\_DIEPEMPMSK)

Address offset: 0x834

Reset value: 0x0000 0000

This register is used to control the IN endpoint FIFO empty interrupt generation

(TXFE\_OTG\_HS\_DIEPINTx).

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved							IN	IEPT	XFE	М						
neserveu	rw	rw	rw	rw	rw	rw	rw	rw	rw							

Bits 31:16 Reserved, must be kept at reset value.

Bits 15:0 INEPTXFEM: IN EP Tx FIFO empty interrupt mask bits

These bits act as mask bits for OTG\_HS\_DIEPINTx.

TXFE interrupt one bit per IN endpoint:

Bit 0 for IN endpoint 0, bit 15 for IN endpoint 15

0: Masked interrupt1: Unmasked interrupt

# OTG\_HS device each endpoint interrupt register (OTG\_HS\_DEACHINT)

Address offset: 0x0838

Reset value: 0x0000 0000

There is one interrupt bit for endpoint 1 IN and one interrupt bit for endpoint 1 OUT.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
						Rese	erved	t						OEP1INT							Re	eserv	⁄ed							IEP1INT	Reserved

Bits 31:18 Reserved, must be kept at reset value.

Bit 17 **OEP1INT:** OUT endpoint 1 interrupt bit

Bits 16:2 Reserved, must be kept at reset value.

Bit 1 IEP1INT: IN endpoint 1 interrupt bit

Bit 0 Reserved, must be kept at reset value.

# OTG\_HS device each endpoint interrupt register mask (OTG\_HS\_DEACHINTMSK)

Address offset: 0x083C Reset value: 0x0000 0000

There is one interrupt bit for endpoint 1 IN and one interrupt bit for endpoint 1 OUT.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved

Reserved

Reserved

Reserved

Reserved

Bits 31:18 Reserved, must be kept at reset value.

Bit 17 OEP1INTM: OUT Endpoint 1 interrupt mask bit

Bits 16:2 Reserved, must be kept at reset value.

Bit 1 IEP1INTM: IN Endpoint 1 interrupt mask bit

Bit 0 Reserved, must be kept at reset value.

# OTG\_HS device each in endpoint-1 interrupt register (OTG\_HS\_DIEPEACHMSK1)

Address offset: 0x844

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
								Rese	erved	I								NAKM		Reserved		BIM	TXFURM	Reserved	INEPNEM	INEPNMM	ITTXFEMSK	MOT	Reserved	EPDM	XFRCM
																		rw				rw	rw		rw	rw	rw	rw		rw	rw

Bits 31:14 Reserved, must be kept at reset value.

Bit 13 NAKM: NAK interrupt mask

0: Masked interrupt1: unmasked interrupt

Bit 12:10 Reserved, must be kept at reset value.

Bit 9 BIM: BNA interrupt mask

0: Masked interrupt1: Unmasked interrupt

Bit 8 TXFURM: FIFO underrun mask

0: Masked interrupt1: Unmasked interrupt

Bit 7 Reserved, must be kept at reset value.

Bit 6 INEPNEM: IN endpoint NAK effective mask

0: Masked interrupt1: Unmasked interrupt

Bit 5 INEPNMM: IN token received with EP mismatch mask

0: Masked interrupt1: Unmasked interrupt

Bit 4 ITTXFEMSK: IN token received when TxFIFO empty mask

0: Masked interrupt1: Unmasked interrupt

Bit 3 TOM: Timeout condition mask (nonisochronous endpoints)

0: Masked interrupt1: Unmasked interrupt

Bit 2 Reserved, must be kept at reset value.

Bit 1 EPDM: Endpoint disabled interrupt mask

0: Masked interrupt1: Unmasked interrupt

Bit 0 XFRCM: Transfer completed interrupt mask

0: Masked interrupt1: Unmasked interrupt

# OTG\_HS device each OUT endpoint-1 interrupt register (OTG\_HS\_DOEPEACHMSK1)

Address offset: 0x884

Reset value: 0x0000 0000

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							Re	eserv	red								NYETM	NAKM	BERRM		Heserved	BIM	TXFURM	Reserved	INEPNEM	INEPNMM	ITTXFEMSK	TOM	Reserved	EPDM	XFRCM
																	rw	rw	rw			rw	rw		rw	rw	rw	rw		rw	rw

Bits 31:15 Reserved, must be kept at reset value.

Bit 14 NYETM: NYET interrupt mask

0: Masked interrupt1: unmasked interrupt

Bit 13 NAKM: NAK interrupt mask

0: Masked interrupt1: Unmasked interrupt

Bit 12 BERRM: Bubble error interrupt mask

0: Masked interrupt1: Unmasked interrupt

Bit 11:10 Reserved, must be kept at reset value.

Bit 9 BIM: BNA interrupt mask

0: Masked interrupt1: Unmasked interrupt

Bit 8 **OPEM:** OUT packet error mask

0: Masked interrupt

1: Unmasked interrupt

Bits 7:3 Reserved, must be kept at reset value.

Bit 2 AHBERRM: AHB error mask

0: Masked interrupt1: Unmasked interrupt

Bit 1 EPDM: Endpoint disabled interrupt mask

0: Masked interrupt1: Unmasked interrupt

Bit 0 XFRCM: Transfer completed interrupt mask

0: Masked interrupt1: Unmasked interrupt

# OTG device endpoint-x control register (OTG\_HS\_DIEPCTLx) (x = 0..7, where $x = Endpoint\_number$ )

Address offset: 0x900 + (Endpoint\_number × 0x20)

Reset value: 0x0000 0000

The application uses this register to control the behavior of each logical endpoint other than endpoint 0.

3	1	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
VIV.	FILINA	EPDIS	SODDFRM	SD0PID/SEVNFRM	SNAK	CNAK		TXF	NUM	l	Stall	Reserved	axtaa		NAKSTS	EONUM/DPID	USBAEP		Rese	erved	d					٨	//PSI	Z				
r	s	rs	w	w	w	w	rw	rw	rw	rw	rw/ rs		rw	rw	r	r	rw					rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

## Bit 31 EPENA: Endpoint enable

The application sets this bit to start transmitting data on an endpoint.

The core clears this bit before setting any of the following interrupts on this endpoint:

- SETUP phase done
- Endpoint disabled
- Transfer completed

## Bit 30 EPDIS: Endpoint disable

The application sets this bit to stop transmitting/receiving data on an endpoint, even before the transfer for that endpoint is complete. The application must wait for the Endpoint disabled interrupt before treating the endpoint as disabled. The core clears this bit before setting the Endpoint disabled interrupt. The application must set this bit only if Endpoint enable is already set for this endpoint.

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### Bit 29 SODDFRM: Set odd frame

Applies to isochronous IN and OUT endpoints only.

Writing to this field sets the Even/Odd frame (EONUM) field to odd frame.

### Bit 28 SD0PID: Set DATA0 PID

Applies to interrupt/bulk IN endpoints only.

Writing to this field sets the endpoint data PID (DPID) field in this register to DATA0.

### SEVNFRM: Set even frame

Applies to isochronous IN endpoints only.

Writing to this field sets the Even/Odd frame (EONUM) field to even frame.

### Bit 27 SNAK: Set NAK

A write to this bit sets the NAK bit for the endpoint.

Using this bit, the application can control the transmission of NAK handshakes on an endpoint. The core can also set this bit for OUT endpoints on a Transfer completed interrupt, or after a SETUP is received on the endpoint.

### Bit 26 CNAK: Clear NAK

A write to this bit clears the NAK bit for the endpoint.

### Bits 25:22 TXFNUM: TxFIFO number

These bits specify the FIFO number associated with this endpoint. Each active IN endpoint must be programmed to a separate FIFO number.

This field is valid only for IN endpoints.

## Bit 21 STALL: STALL handshake

Applies to noncontrol, nonisochronous IN endpoints only (access type is rw).

The application sets this bit to stall all tokens from the USB host to this endpoint. If a NAK bit, Global IN NAK, or Global OUT NAK is set along with this bit, the STALL bit takes priority. Only the application can clear this bit, never the core.

Applies to control endpoints only (access type is rs).

The application can only set this bit, and the core clears it, when a SETUP token is received for this endpoint. If a NAK bit, Global IN NAK, or Global OUT NAK is set along with this bit, the STALL bit takes priority. Irrespective of this bit's setting, the core always responds to SETUP data packets with an ACK handshake.

## Bit 20 Reserved, must be kept at reset value.

# Bits 19:18 EPTYP: Endpoint type

This is the transfer type supported by this logical endpoint.

00: Control 01: Isochronous 10: Bulk 11: Interrupt

### Bit 17 NAKSTS: NAK status

It indicates the following:

- 0: The core is transmitting nonNAK handshakes based on the FIFO status.
- 1: The core is transmitting NAK handshakes on this endpoint.

When either the application or the core sets this bit:

For nonisochronous IN endpoints: The core stops transmitting any data on an IN endpoint, even if there are data available in the TxFIFO.

For isochronous IN endpoints: The core sends out a zero-length data packet, even if there are data available in the TxFIFO.

Irrespective of this bit's setting, the core always responds to SETUP data packets with an ACK handshake.

# Bit 16 EONUM: Even/odd frame

Applies to isochronous IN endpoints only.

Indicates the frame number in which the core transmits/receives isochronous data for this endpoint. The application must program the even/odd frame number in which it intends to transmit/receive isochronous data for this endpoint using the SEVNFRM and SODDFRM fields in this register.

0: Even frame

1: Odd frame

**DPID:** Endpoint data PID

Applies to interrupt/bulk IN endpoints only.

Contains the PID of the packet to be received or transmitted on this endpoint. The application must program the PID of the first packet to be received or transmitted on this endpoint, after the endpoint is activated. The application uses the SD0PID register field to program either DATA0 or DATA1 PID.

0: DATA0

1: DATA1

# Bit 15 USBAEP: USB active endpoint

Indicates whether this endpoint is active in the current configuration and interface. The core clears this bit for all endpoints (other than EP 0) after detecting a USB reset. After receiving the SetConfiguration and SetInterface commands, the application must program endpoint registers accordingly and set this bit.

Bits 14:11 Reserved, must be kept at reset value.

## Bits 10:0 MPSIZ: Maximum packet size

The application must program this field with the maximum packet size for the current logical endpoint. This value is in bytes.

# OTG\_HS device control OUT endpoint 0 control register (OTG\_HS\_DOEPCTL0)

Address offset: 0xB00

Reset value: 0x0000 8000

This section describes the device control OUT endpoint 0 control register. Nonzero control endpoints use registers for endpoints 1–15.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

li	EPENA	EPDIS	eserved	SNAK	CNAK	Reserved	Stall	SNPM	EPT	ΥP	NAKSTS	eserved	USBAEP	Reserved	MPS	SIZ	
	w	r	Ж	w	w		rs	rw	r	r	r	ш	r		r	r	

## Bit 31 EPENA: Endpoint enable

The application sets this bit to start transmitting data on endpoint 0.

The core clears this bit before setting any of the following interrupts on this endpoint:

- SETUP phase done
- Endpoint disabled
- Transfer completed

## Bit 30 EPDIS: Endpoint disable

The application cannot disable control OUT endpoint 0.

Bits 29:28 Reserved, must be kept at reset value.

### Bit 27 SNAK: Set NAK

A write to this bit sets the NAK bit for the endpoint.

Using this bit, the application can control the transmission of NAK handshakes on an endpoint. The core can also set this bit on a Transfer completed interrupt, or after a SETUP is received on the endpoint.

## Bit 26 CNAK: Clear NAK

A write to this bit clears the NAK bit for the endpoint.

Bits 25:22 Reserved, must be kept at reset value.

## Bit 21 STALL: STALL handshake

The application can only set this bit, and the core clears it, when a SETUP token is received for this endpoint. If a NAK bit or Global OUT NAK is set along with this bit, the STALL bit takes priority. Irrespective of this bit's setting, the core always responds to SETUP data packets with an ACK handshake.

## Bit 20 SNPM: Snoop mode

This bit configures the endpoint to Snoop mode. In Snoop mode, the core does not check the correctness of OUT packets before transferring them to application memory.

## Bits 19:18 EPTYP: Endpoint type

Hardcoded to 2'b00 for control.

### Bit 17 NAKSTS: NAK status

Indicates the following:

0: The core is transmitting nonNAK handshakes based on the FIFO status.

1: The core is transmitting NAK handshakes on this endpoint.

When either the application or the core sets this bit, the core stops receiving data, even if there is space in the RxFIFO to accommodate the incoming packet. Irrespective of this bit's setting, the core always responds to SETUP data packets with an ACK handshake.

Bit 16 Reserved, must be kept at reset value.

## Bit 15 USBAEP: USB active endpoint

This bit is always set to 1, indicating that a control endpoint 0 is always active in all configurations and interfaces.

Bits 14:2 Reserved, must be kept at reset value.

## Bits 1:0 MPSIZ: Maximum packet size

The maximum packet size for control OUT endpoint 0 is the same as what is programmed in control IN endpoint 0.

00: 64 bytes 01: 32 bytes 10: 16 bytes 11: 8 bytes

# OTG\_HS device endpoint-x control register (OTG\_HS\_DOEPCTLx) (x = 1...3, where $x = Endpoint_number$ )

Address offset for OUT endpoints: 0xB00 + (Endpoint\_number × 0x20)

Reset value: 0x0000 0000

The application uses this register to control the behavior of each logical endpoint other than endpoint 0.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	1	ь	5	4	3	2	ı	U
EPENA	EPDIS	SODDFRM	SD0PID/SEVNFRM	SNAK	CNAK		Rese	erved		Stall	MANS	dALd∃		NAKSTS	EONUM/DPID	USBAEP		Rese	rved						N	/IPSI	Z				
rs	rs	w	w	w	w					rw/ rs	rw	rw	rw	r	r	rw					rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

## Bit 31 EPENA: Endpoint enable

Applies to IN and OUT endpoints.

The application sets this bit to start transmitting data on an endpoint.

The core clears this bit before setting any of the following interrupts on this endpoint:

- SETUP phase done
- Endpoint disabled
- Transfer completed

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## Bit 30 EPDIS: Endpoint disable

The application sets this bit to stop transmitting/receiving data on an endpoint, even before the transfer for that endpoint is complete. The application must wait for the Endpoint disabled interrupt before treating the endpoint as disabled. The core clears this bit before setting the Endpoint disabled interrupt. The application must set this bit only if Endpoint enable is already set for this endpoint.

### Bit 29 SODDFRM: Set odd frame

Applies to isochronous OUT endpoints only.

Writing to this field sets the Even/Odd frame (EONUM) field to odd frame.

### Bit 28 SD0PID: Set DATA0 PID

Applies to interrupt/bulk OUT endpoints only.

Writing to this field sets the endpoint data PID (DPID) field in this register to DATA0.

### SEVNFRM: Set even frame

Applies to isochronous OUT endpoints only.

Writing to this field sets the Even/Odd frame (EONUM) field to even frame.

## Bit 27 SNAK: Set NAK

A write to this bit sets the NAK bit for the endpoint.

Using this bit, the application can control the transmission of NAK handshakes on an endpoint. The core can also set this bit for OUT endpoints on a Transfer Completed interrupt, or after a SETUP is received on the endpoint.

### Bit 26 CNAK: Clear NAK

A write to this bit clears the NAK bit for the endpoint.

## Bits 25:22 Reserved, must be kept at reset value.

## Bit 21 STALL: STALL handshake

Applies to noncontrol, nonisochronous OUT endpoints only (access type is rw).

The application sets this bit to stall all tokens from the USB host to this endpoint. If a NAK bit, Global IN NAK, or Global OUT NAK is set along with this bit, the STALL bit takes priority. Only the application can clear this bit, never the core.

Applies to control endpoints only (access type is rs).

The application can only set this bit, and the core clears it, when a SETUP token is received for this endpoint. If a NAK bit, Global IN NAK, or Global OUT NAK is set along with this bit, the STALL bit takes priority. Irrespective of this bit's setting, the core always responds to SETUP data packets with an ACK handshake.

## Bit 20 SNPM: Snoop mode

This bit configures the endpoint to Snoop mode. In Snoop mode, the core does not check the correctness of OUT packets before transferring them to application memory.

# Bits 19:18 EPTYP: Endpoint type

This is the transfer type supported by this logical endpoint.

00: Control01: Isochronous10: Bulk

11: Interrupt

### Bit 17 NAKSTS: NAK status

Indicates the following:

- 0: The core is transmitting nonNAK handshakes based on the FIFO status.
- 1: The core is transmitting NAK handshakes on this endpoint.

When either the application or the core sets this bit:

The core stops receiving any data on an OUT endpoint, even if there is space in the RxFIFO to accommodate the incoming packet.

Irrespective of this bit's setting, the core always responds to SETUP data packets with an ACK handshake.

## Bit 16 EONUM: Even/odd frame

Applies to isochronous IN and OUT endpoints only.

Indicates the frame number in which the core transmits/receives isochronous data for this endpoint. The application must program the even/odd frame number in which it intends to transmit/receive isochronous data for this endpoint using the SEVNFRM and SODDFRM fields in this register.

0: Even frame

1: Odd frame

**DPID:** Endpoint data PID

Applies to interrupt/bulk OUT endpoints only.

Contains the PID of the packet to be received or transmitted on this endpoint. The application must program the PID of the first packet to be received or transmitted on this endpoint, after the endpoint is activated. The application uses the SD0PID register field to program either DATA0 or DATA1 PID.

0: DATA0 1: DATA1

## Bit 15 USBAEP: USB active endpoint

Indicates whether this endpoint is active in the current configuration and interface. The core clears this bit for all endpoints (other than EP 0) after detecting a USB reset. After receiving the SetConfiguration and SetInterface commands, the application must program endpoint registers accordingly and set this bit.

# Bits 14:11 Reserved, must be kept at reset value.

# Bits 10:0 MPSIZ: Maximum packet size

The application must program this field with the maximum packet size for the current logical endpoint. This value is in bytes.



# OTG\_HS device endpoint-x interrupt register (OTG\_HS\_DIEPINTx) (x = 0...7, where $x = Endpoint_number$ )

Address offset:  $0x908 + (Endpoint number \times 0x20)$ 

Reset value: 0x0000 0080

This register indicates the status of an endpoint with respect to USB- and AHB-related events. It is shown in *Figure 376*. The application must read this register when the IN endpoints interrupt bit of the Core interrupt register (IEPINT in OTG\_HS\_GINTSTS) is set. Before the application can read this register, it must first read the device all endpoints interrupt (OTG\_HS\_DAINT) register to get the exact endpoint number for the device endpoint-x interrupt register. The application must clear the appropriate bit in this register to clear the corresponding bits in the OTG\_HS\_DAINT and OTG\_HS\_GINTSTS registers.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
								Rese	erved	ı								NAK	BERR	PKTDRPSTS	eserved	BNA	TXFIFOUDRN	TXFE	INEPNE	eserved	ITTXFE	TOC	eserved	GSIGAE	XFRC
																					ш			r	rc_ w1 /rw	Ж	rc_ w1	rc_ w1	Н	rc_ w1	rc_ w1

Bits 31:14 Reserved, must be kept at reset value.

## Bit 13 NAK: NAK interrupt

The core generates this interrupt when a NAK is transmitted or received by the device. In case of isochronous IN endpoints the interrupt gets generated when a zero length packet is transmitted due to unavailability of data in the Tx FIFO.

Bit 12 BERR: Babble error interrupt

## Bit 11 **PKTDRPSTS:** Packet dropped status

This bit indicates to the application that an ISOC OUT packet has been dropped. This bit does not have an associated mask bit and does not generate an interrupt.

Bit10 Reserved, must be kept at reset value.

## Bit 9 BNA: Buffer not available interrupt

The core generates this interrupt when the descriptor accessed is not ready for the Core to process, such as host busy or DMA done.

Bit 8 **TXFIFOUDRN**: Transmit Fifo Underrun (TxfifoUndrn) The core generates this interrupt when it detects a transmit FIFO underrun condition for this endpoint.

Dependency: This interrupt is valid only when Thresholding is enabled

# Bit 7 TXFE: Transmit FIFO empty

This interrupt is asserted when the TxFIFO for this endpoint is either half or completely empty. The half or completely empty status is determined by the TxFIFO empty level bit in the Core AHB configuration register (TXFELVL bit in OTG\_HS\_GAHBCFG).

## Bit 6 INEPNE: IN endpoint NAK effective

This bit can be cleared when the application clears the IN endpoint NAK by writing to the CNAK bit in  $OTG_HS_DIEPCTLx$ .

This interrupt indicates that the core has sampled the NAK bit set (either by the application or by the core). The interrupt indicates that the IN endpoint NAK bit set by the application has taken effect in the core.

This interrupt does not guarantee that a NAK handshake is sent on the USB. A STALL bit takes priority over a NAK bit.

Bit 5 Reserved, must be kept at reset value.

## Bit 4 ITTXFE: IN token received when TxFIFO is empty

Applies to nonperiodic IN endpoints only.

Indicates that an IN token was received when the associated TxFIFO (periodic/nonperiodic) was empty. This interrupt is asserted on the endpoint for which the IN token was received.

### Bit 3 TOC: Timeout condition

Applies only to Control IN endpoints.

Indicates that the core has detected a timeout condition on the USB for the last IN token on this endpoint.

Bit 2 Reserved, must be kept at reset value.

### Bit 1 **EPDISD:** Endpoint disabled interrupt

This bit indicates that the endpoint is disabled per the application's request.

## Bit 0 XFRC: Transfer completed interrupt

This field indicates that the programmed transfer is complete on the AHB as well as on the USB, for this endpoint.



# OTG\_HS device endpoint-x interrupt register (OTG\_HS\_DOEPINTx) (x = 0...7, where $x = Endpoint_number$ )

Address offset: 0xB08 + (Endpoint\_number × 0x20)

Reset value: 0x0000 0080

This register indicates the status of an endpoint with respect to USB- and AHB-related events. It is shown in *Figure 376*. The application must read this register when the OUT Endpoints Interrupt bit of the Core interrupt register (OEPINT bit in OTG\_HS\_GINTSTS) is set. Before the application can read this register, it must first read the device all endpoints interrupt (OTG\_HS\_DAINT) register to get the exact endpoint number for the device Endpoint-x interrupt register. The application must clear the appropriate bit in this register to clear the corresponding bits in the OTG\_HS\_DAINT and OTG\_HS\_GINTSTS registers.

;	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
								Re	eserv	red								NYET			Re	eserv	ed			≩ ੬ ਰ B2BSTUP	Reserved		STUP	Reserved		L XFRC

Bits 31:15 Reserved, must be kept at reset value.

## Bit 14 NYET: NYET interrupt

The core generates this interrupt when a NYET response is transmitted for a nonisochronous OUT endpoint.

Bits 13:7 Reserved, must be kept at reset value.

## Bit 6 B2BSTUP: Back-to-back SETUP packets received

Applies to Control OUT endpoint only.

This bit indicates that the core has received more than three back-to-back SETUP packets for this particular endpoint.

Bit 5 Reserved, must be kept at reset value.

## Bit 4 OTEPDIS: OUT token received when endpoint disabled

Applies only to control OUT endpoint.

Indicates that an OUT token was received when the endpoint was not yet enabled. This interrupt is asserted on the endpoint for which the OUT token was received.

## Bit 3 STUP: SETUP phase done

Applies to control OUT endpoints only.

Indicates that the SETUP phase for the control endpoint is complete and no more back-to-back SETUP packets were received for the current control transfer. On this interrupt, the application can decode the received SETUP data packet.

Bit 2 Reserved, must be kept at reset value.

## Bit 1 EPDISD: Endpoint disabled interrupt

This bit indicates that the endpoint is disabled per the application's request.

## Bit 0 XFRC: Transfer completed interrupt

This field indicates that the programmed transfer is complete on the AHB as well as on the USB, for this endpoint.

# OTG\_HS device IN endpoint 0 transfer size register (OTG\_HS\_DIEPTSIZ0)

Address offset: 0x910 Reset value: 0x0000 0000

The application must modify this register before enabling endpoint 0. Once endpoint 0 is enabled using the endpoint enable bit in the device control endpoint 0 control registers (EPENA in OTG\_HS\_DIEPCTL0), the core modifies this register. The application can only read this register once the core has cleared the Endpoint enable bit.

Nonzero endpoints use the registers for endpoints 1–15.

### 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	PKT	CNT	Decembed			Х	FRS	ΙZ		
neserved	rw	rw	Reserved	rw	rw	rw	rw	rw	rw	rw

Bits 31:21 Reserved, must be kept at reset value.

### Bits 20:19 PKTCNT: Packet count

Indicates the total number of USB packets that constitute the Transfer Size amount of data for endpoint 0.

This field is decremented every time a packet (maximum size or short packet) is read from the TxFIFO.

Bits 18:7 Reserved, must be kept at reset value.

## Bits 6:0 XFRSIZ: Transfer size

Indicates the transfer size in bytes for endpoint 0. The core interrupts the application only after it has exhausted the transfer size amount of data. The transfer size can be set to the maximum packet size of the endpoint, to be interrupted at the end of each packet.

The core decrements this field every time a packet from the external memory is written to the TxFIFO.

# OTG\_HS device OUT endpoint 0 transfer size register (OTG\_HS\_DOEPTSIZ0)

Address offset: 0xB10 Reset value: 0x0000 0000

The application must modify this register before enabling endpoint 0. Once endpoint 0 is enabled using the Endpoint enable bit in the device control endpoint 0 control registers (EPENA bit in OTG\_HS\_DOEPCTL0), the core modifies this register. The application can only read this register once the core has cleared the Endpoint enable bit.

Nonzero endpoints use the registers for endpoints 1–15.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

eserved	STUPC NT	Reserved E	PKTCNT	Reserved			X	FRS	ΙΖ		
Œ	rw rw	r	rw		rw	rw	rw	rw	rw	rw	rw

Bit 31 Reserved, must be kept at reset value.

## Bits 30:29 STUPCNT: SETUP packet count

This field specifies the number of back-to-back SETUP data packets the endpoint can receive.

01: 1 packet 10: 2 packets 11: 3 packets

Bits 28:20 Reserved, must be kept at reset value.

### Bit 19 PKTCNT: Packet count

This field is decremented to zero after a packet is written into the RxFIFO.

Bits 18:7 Reserved, must be kept at reset value.

## Bits 6:0 XFRSIZ: Transfer size

Indicates the transfer size in bytes for endpoint 0. The core interrupts the application only after it has exhausted the transfer size amount of data. The transfer size can be set to the maximum packet size of the endpoint, to be interrupted at the end of each packet.

The core decrements this field every time a packet is read from the RxFIFO and written to the external memory.

# OTG\_HS device endpoint-x transfer size register (OTG\_HS\_DIEPTSIZx) $(x = 1..3, where x = Endpoint_number)$

Address offset:  $0x910 + (Endpoint number \times 0x20)$ 

Reset value: 0x0000 0000

The application must modify this register before enabling the endpoint. Once the endpoint is enabled using the Endpoint enable bit in the device endpoint-x control registers (EPENA bit in OTG\_HS\_DIEPCTLx), the core modifies this register. The application can only read this register once the core has cleared the Endpoint enable bit.

### 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

pe	МС	NT			PKTCNT XFRSIZ																										
Reserv	rw/ r/r w	rw/ r/r w	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bit 31 Reserved, must be kept at reset value.

## Bits 30:29 MCNT: Multi count

For periodic IN endpoints, this field indicates the number of packets that must be transmitted per frame on the USB. The core uses this field to calculate the data PID for isochronous IN endpoints.

01: 1 packet 10: 2 packets 11: 3 packets

### Bit 28:19 PKTCNT: Packet count

Indicates the total number of USB packets that constitute the Transfer Size amount of data for this endpoint.

This field is decremented every time a packet (maximum size or short packet) is read from the TxFIFO.

# Bits 18:0 XFRSIZ: Transfer size

This field contains the transfer size in bytes for the current endpoint. The core only interrupts the application after it has exhausted the transfer size amount of data. The transfer size can be set to the maximum packet size of the endpoint, to be interrupted at the end of each packet.

The core decrements this field every time a packet from the external memory is written to the TxFIFO.

r

# OTG\_HS device IN endpoint transmit FIFO status register (OTG\_HS\_DTXFSTSx) (x = 0..5, where x = Endpoint\_number)

Address offset for IN endpoints:  $0x918 + (Endpoint\_number \times 0x20)$  This read-only register contains the free space information for the Device IN endpoint TxFIFO.

# 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0 INEPTFSAV Reserved

31:16 Reserved, must be kept at reset value.

15:0 INEPTFSAV: IN endpoint TxFIFO space avail ()

Indicates the amount of free space available in the Endpoint TxFIFO.

Values are in terms of 32-bit words:

0x0: Endpoint TxFIFO is full 0x1: 1 word available 0x2: 2 words available

0xn: n words available (0 < n < 512)

Others: Reserved

# OTG\_HS device endpoint-x transfer size register (OTG\_HS\_DOEPTSIZx) (x = 1..5, where x = Endpoint\_number)

Address offset: 0xB10 + (Endpoint\_number × 0x20)

Reset value: 0x0000 0000

The application must modify this register before enabling the endpoint. Once the endpoint is enabled using Endpoint Enable bit of the device endpoint-x control registers (EPENA bit in OTG\_HS\_DOEPCTLx), the core modifies this register. The application can only read this register once the core has cleared the Endpoint enable bit.

# 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

rved		PID/S PCNT					PKT	CNT													Х	FRS	ΙZ								
Rese	rw/r/ rw	rw/r/ rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

Bit 31 Reserved, must be kept at reset value.

Bits 30:29 RXDPID: Received data PID

Applies to isochronous OUT endpoints only.

This is the data PID received in the last packet for this endpoint.

00: DATA0 01: DATA2 10: DATA1 11: MDATA

## STUPCNT: SETUP packet count

Applies to control OUT Endpoints only.

This field specifies the number of back-to-back SETUP data packets the endpoint can receive.

01: 1 packet 10: 2 packets 11: 3 packets

## Bit 28:19 PKTCNT: Packet count

Indicates the total number of USB packets that constitute the Transfer Size amount of data for this endpoint.

This field is decremented every time a packet (maximum size or short packet) is written to the RxFIFO.

## Bits 18:0 XFRSIZ: Transfer size

This field contains the transfer size in bytes for the current endpoint. The core only interrupts the application after it has exhausted the transfer size amount of data. The transfer size can be set to the maximum packet size of the endpoint, to be interrupted at the end of each packet.

The core decrements this field every time a packet is read from the RxFIFO and written to the external memory.

# OTG\_HS device endpoint-x DMA address register (OTG\_HS\_DIEPDMAx / OTG\_HS\_DOEPDMAx) (x = 1..5, where x = Endpoint\_number)

Address offset for IN endpoints: 0x914 + (Endpoint\_number × 0x20)

Reset value: 0xXXXX XXXX

Address offset for OUT endpoints: 0xB14 + (Endpoint\_number × 0x20)

Reset value: 0xXXXX XXXX

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
														D١	ΛΑΑΙ	DDR															

## Bits 31:0 DMAADDR: DMA address

This bit holds the start address of the external memory for storing or fetching endpoint data.

Note: For control endpoints, this field stores control OUT data packets as well as SETUP transaction data packets. When more than three SETUP packets are received back-to-back, the SETUP data packet in the memory is overwritten. This register is incremented on every AHB transaction. The application can give only a DWORD-aligned address.

# 30.12.5 OTG\_HS power and clock gating control register (OTG\_HS\_PCGCCTL)

Address offset: 0xE00 Reset value: 0x0000 0000

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This register is available in host and peripheral modes.

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31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved

Reserved

Reserved

Reserved

Reserved

Bit 31:5 Reserved, must be kept at reset value.

## Bit 4 PHYSUSP: PHY suspended

Indicates that the PHY has been suspended. This bit is updated once the PHY is suspended after the application has set the STPPCLK bit (bit 0).

Bits 3:2 Reserved, must be kept at reset value.

## Bit 1 GATEHCLK: Gate HCLK

The application sets this bit to gate HCLK to modules other than the AHB Slave and Master and wakeup logic when the USB is suspended or the session is not valid. The application clears this bit when the USB is resumed or a new session starts.

## Bit 0 STPPCLK: Stop PHY clock

The application sets this bit to stop the PHY clock when the USB is suspended, the session is not valid, or the device is disconnected. The application clears this bit when the USB is resumed or a new session starts.

# 30.12.6 OTG\_HS register map

The table below gives the USB OTG register map and reset values.

Table 166. OTG\_HS register map and reset values

Offset	Register	31	30	29	28	27 26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	1	10	6	8	7	9	2	4	3	2	-	0
0x000	OTG_HS_GOT GCTL					Res	erve	ed					o BSVLD	O ASVLD	o DBCT	- CIDSTS	Re	se	rved	L		OHSHNPEN	o HNPRQ	o HNGSCS		F	Rese	erve	d		o SRQ	SRGSCS
0x004	OTG_HS_GOT GINT					Res	erve	ed					O DBCDNE	<b>(</b> 15	HNGDET 0	•			Reserved		<u> </u>		HNSSCHG	SRSSCHG		Re	ser	ved		o SEDET	Re	
0x008	OTG_HS_GAH BCFG										Re	ser		J	J								U	O PTXFELVL	O TXFELVL		F	Rese	erve			O GINT
0x00C	OTG_HS_GUS BCFG	O CTXPKT	o FDMOD		Res	erved	]	DTCI	PCCI	TSDPS	- ULPIEVBUSI	ULPIEVBUSD	OLPICSM			Reserved		חשא ושפשרו	n I		DT 1	0	1 HNPCAP	o SRPCAP	Reserved	1 PHSEL	Re	serv	/ed	TC 0 I	OCA	\L
0x010		→ AHBIDL o	O DMAREQ													1	0		(FNI	JM		OTXFFLSH	ORXFFLSH	o Reserved	_		O CSRST					

Table 166. OTG\_HS register map and reset values (continued)

Ording   Ordina   O	Table I				_			_			_		_	_			_	_	_	_			1			1	1	1	1	1	,		
Ox018	Offset	Register	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	6	7	9	2	4	ဗ	7	-	0
Ox018	0x014	OTG_HS_GINT STS	WKUINT	SRQINT	DISCINT	CIDSCHG	Reserved	PTXFE	HCINT	HPRTINT	Reserved	DATAFSUSP	XFR/INCOMPISOOUT	IISOIXFR	OEPINT	IEPINT	Beserved		EOPF	ISOODRP	ENOMDNE	USBRST	USBSUSP	ESUSP	Reserved	BOUTNAKEFF	GINAKEFF	NPTXFE	RXFLVL	SOF	OTGINT	MMIS	СМОД
Reservative   0   0   0   0   0   0   0   0   0		Reset value	0	0	0	0	-	1	0	0		0	0	0	0	0			0	0	0	0	0	0		0	0	1	0	0	0	0	0
Reservative   0   0   0   0   0   0   0   0   0	0x018		MUM	SRQIM	DISCINT	CIDSCHGM	Reserved	PTXFEM	HCIM	PRTIM	Reserved	FSUSPM	PXFRM/IISOOXFRM	IISOIXFRM	OEPINT	IEPINT	EPMISM	Reserved	EOPFM	ISOODRPM	ENUMDNEM	USBRST	USBSUSPM	ESUSPM	Reserved	GONAKEFFM	GINAKEFFM	NPTXFEM	RXFLVLM	SOFM	OTGINT	MMISM	Reserved
STSR (Host mode)		Reset value	0	0	0	0	-	0	0	0		0	_	0	0	0	0		0	0	0	0	0	0		0	0	0	0	0	0	0	
Oxform   Content   Oxform   Content   Oxform   Content   Oxform   Content   Oxform   Content   Oxform   Content   Oxform   Oxform   Content   Oxform   Oxf		STSR (Host mode)					Res	serv	/ed					ſ	PKT	STS	6	DF						В	CNT		•	•	•	(	CHN	IUM	
STSR	0x01C													0	0	0	0		0	0	0	0	0	0	0 0	0	0	0	0	0	0	0	
OTG_HS_GRX STS mode)		STSR (peripheral mode)			Re	ser	/ed													0	0	0	Ι ο			Ιn	Ιn	Ιn	Ιn				
STSP (Host mode)										U	U	U		0	U	U	0	U	U	0	U	U	U	U	0 0	U	U	U	U	0	U	U	_
OTG_HS_GNP TXFSIZ (Host mode)		STSP (Host mode)	-				Res	serv	/ed									DF		0	0	0	Ιο			Ιn	10	10	10				1
OX024   FSIZ   Reserved   RXFD   RSYFD   Reserved   RXFD   Reserved   RXFD   Reserved   RXFD   Reserved   RXFD   Reserved   RXFD   O 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	0x020	OTG_HS_GRX STSP (peripheral mode)		Reserved    O   O   O   O   O   O   O   O   O																													
Note										0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0	0	0	0	0	0	0	0
OX028	0x024	FSIZ		Reserved																													
NPTXFSA   NPTXFD   NPTXFSA				Heserved														0	0	0													
Reset value																																	
0x028  OTG_HS_GNP TXFSIZ (peripheral mode)  Reset value  OX02C  OTG_HS_GNP TXSTS  Reset value  OX030  OTG_HS_GIP CTL  OX030  OTG_HS_GIP CTL  OX030  Reset value  OX030  OTG_HS_GIP CTL  OX030  OTG_HS_CIP CTL  OX030  O			_		<u> </u>	I 0		_		^	_	I 0		<u> </u>	_		_	0				_	<u> </u>	_	1 1 0	10	Ι Λ	Ι Λ	Ι Λ	<u> </u>	_	^	0
Mode    Reset value	0x028	OTG_HS_GNP TXFSIZ	U	U	U	U	U	0					0	U	U	<u> </u>	U	U	0	U	U	U	U	U				U	U	U	U	U	
0x02C         OTG_HS_GNP TXSTS         NPTXQTOP         NPTQXSAV         NPTXFSAV           Reset value         0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0		mode)																															
Ox03C   TXSTS   NFTACTOF   NTACTOF   N			0	0		U			_	0	0	0					0	0	0	0	0	0	0	0				0	0	0	0	0	0
0x030         OTG_HS_GI2C CTL         UND SET VALUE         UND SET VALUE<	0x02C	TXSTS	NPTXQTOP NPTQXSAV NPTXFSAV																														
Reset value		Reset value	4	0	0	_			0	0	0	0	0	0	0	1	0	0	0	0	0	0	0	0	1 0	0	0	0	0	0	0	0	0
Reset value	0x030		BSYDNE BS														A																
Reset value		Reset value	0	0		0				0	0	0		0	0	0	0	0															
OVO3C OTG_HS_CID PRODUCT_ID	0x038	FG				F	Rese	rve	d																Rese	erve	ed						
Reset value 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	0,000												U	U	U	U			JCT	_ID													
	0x03C		0	0	0	0	0	0	0	0	0	0	0	0	0	0					0	1	0	0	1 0	0	0	0	0	0	0	0	0

Table 166. OTG\_HS register map and reset values (continued)

			т.	<del>-</del>			, ,					<del>- 1 `</del>	1_			÷			1	, ,				_		, ,	
Offset	Register	31	30	29	28	27	25	24	23	21	50	18	17	16	15	14	10	11	10	6	<b>0</b>	<u> </u>	9	7	t e	7	0
0x100	OTG_HS_HPTX FSIZ							PTX	FD												PTXS	SA					
	Reset value	0	0	0	0	0 1	1	1	0 1	1	0	1 0	0	0	0	0 (	) 1	0	0	0	0 (	0	0 1	C	0	1	0 0
	OTG_HS_DIEP			•			IN	IEPT	XFD			•		•					•	IN	EPT)	xs/	Α				
0x104	TXF1	^				010	141			101	0.1	010			0.1	0 1 /	1 0	Ι.	1 4								0.1.0
	Reset value  OTG HS DIEP	0	0	0	0	0 0			0 0	0	0	0 0	0	0	0	0 (	0 0	0	<u> </u>	0	0 (		0 (	) (	0	0	0 0
0x108	TXF2						IN	IEPT	XFD											IN	EPT	KS/	4				
	Reset value	0	0	0	0	0 0	1	0	0 0	0	0	0 0	0	0	0	0 (	0 0	0	1	0	0 (	0	0 (	) (	0	0	0 0
	OTG_HS_DIEP			•			IN	IEPT	XFD			•		•					•	IN	EPT)	xs/	Α				
0x10C	TXF3 Reset value	_	١.	Ι Λ	0	010				101	Λ.	010	١.	Ι Λ	Λ.	ΛΙ.	1 0	Ι Λ	14						1 0	ΙΛ.	0.1.0
	OTG HS DIEP	0	0	0	0	0 0	1		0 0	0	0	0 0	0	0	0	0 (	0 0	0	<u> </u>	0			0 (	) (	0	0	0 0
0x110	TXF4						IN	IEPT	XFD											IN	EPT	KS/	4				
	Reset value	0	0	0	0	0 0	1	0	0 0	0	0	0 0	0	0	0	0 (	0	0	1	0	0 (	0	0 (	0	0	0	0 0
0x400	OTG_HS_HCF G											Re	ser	ved												FSLSS	FSLSPCS
	Reset value  OTG HS HFIR																				FRIV	<u>'</u>				0	0 0
0x404	Reset value						F	Reser	rved						1	1	1 0	Ι1	0	111	0 0		1   1	ГС	0 0	0	0 0
	OTG_HS_HFN							CTDI							- 1	- 1	~	<u> </u>	ľ	بنا		_					
0x408	UM							FTRI	EIVI												RNU	JIVI					
	Reset value	0	0	0	0	0 0	0	0	0 0	0	0	0 0	0	0	0	0 .	1   1	1	1	1	1 1	1	1 1	1	1	1	1 1
0x410	OTG_HS_HPTX STS			Р	TXQ	ГОР				PΤ	ΓXQ	SAV								Р	XFS	AVI	L				
0.004	Reset value	0	0	0	0	0 0	0	0	YIY	TYT	YΤ	YIY	Υ	Υ	Υ	ΥΙν	/ T Y	ТҮ	ΤΥ	Υ	ΥΙΥ	ΥT	ΥIY	/ T Y	′ T Y	Υ	YIY
	OTG_HS_HAIN			Ŭ		<u> </u>		ŭ		1 - 1	<u> </u>	<u> </u>			- 1	•		<u> </u>	<u> </u>		HAIN			<u> </u>	<u> </u>		· I ·
0x414	T						F	Reser	rved												HAIN	11					
	Reset value														0	0 (	0 0	0	0	0	0 (	0	0 (	0	0	0	0 0
0x418	OTG_HS_HAIN TMSK						_	Reser	ruad											H	TAIN	ГΜ					
0.00410	Reset value	-					Г	iesei	veu						0	0 (	0 0	0	0	0	0 0	0	0 0	) [ (	0 0	0	0 0
0x440	OTG_HS_HPRT					Re	eserv	/ed					PD		PTC	TL	PPWR	i	PLSIS	В	PRST	1000	POCCHNG	POCA	PE	PENA	PCDET PCSTS
	Reset value											0	0	0	0	0 (	0	0	0	0	0 (	0	0 (	) (	0	0	0 0
0x500	OTG_HS_HCC HAR0	CHENA	CHDIS	ODDFRM			DAD			М		EPTYP	LSDEV	Reserved	EPDIR		PNU						MP				
	Reset value	0	0	0	0	0 0	0	0	0 0	0	0	0 0	0	4	0	0 (	0 0	0	0	0	0 (	0	0 (	) (	0	0	0 0
0x520	OTG_HS_HCC HAR1	CHENA	CHDIS	ODDFRM			DAD	)		М	0	EPTYP	LSDEV	Reserved	EPDIR	EI	PNU	М					MP	SIZ			
	Reset value	0	0	0	0	0 0	0	0	0 0	0	0	0 0	0	Н	0	0 (	0 0	0	0	0	0 0	0	0 0	) (	0	0	0 0
0x540	OTG_HS_HCC HAR2	CHENA	CHDIS	ODDFRM			DAD	)		М	С	EPTYP	LSDEV	Reserved	EPDIR	EI	PNU	М					MP	SIZ			
	Reset value	0	0	0	0	0 0	0	0	0 0	0	0	0 0	0	ш	0	0 (	0	0	0	0	0 (	0	0 (	) (	0	0	0 0
0x560	OTG_HS_HCC HAR3	CHENA	CHDIS	ODDFRM			DAD	)		М	С	EPTYP	LSDEV	Reserved	EPDIR	El	PNU	М					MP	SIZ			
	Reset value	0	0	0	0	0 0	0	0	0 0	0	0	0 0	0	æ	0	0 (	0 0	0	0	0	0 (	0	0 (	) (	0	0	0 0
0x580	OTG_HS_HCC HAR4	CHENA	CHDIS	ODDFRM			DAD			М		EPTYP	LSDEV	Reserved	EPDIR		PNU						MP				
	Reset value	0	0	0	0	0 0	0	0	0 0	0	0	0 0	0	_	0	0 (	0 0	0	0	0	0 (	0	0 (	) (	0	0	0 0
0x5A0	OTG_HS_HCC HAR5	CHENA	CHDIS	ODDFRM		<u> </u>	DAD		<u> </u>	МС		EPTYP	, LSDEV	Reserved	EPDIR		PNU				0.1.	<u> </u>	MP				
	Reset value	0	0	0	0	0 0	0	0	0 0	0	0	0 0	0		0	0 (	0 0	0	0	0	0 (	J	0 (	) (	0	0	0 0

Table 166. OTG\_HS register map and reset values (continued)

Offeet		_		_	ter map								, 			1	T			- 1	_	$\neg$
Offset	Register	31		129	28 27 26 26	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	7 2 2	19 12	17	16	_	14	12 2	9	6	8	9	2	4	က	7	- 0
0x5C0	OTG_HS_HCC HAR6	CHENA		ODDFRM	DA	AD	МС	EPTYP	LSDEV	Reserved	EPDIR		EPNUM				N	IPSI	Z			
	Reset value	0	0	0	0 0 0 0	0 0 0	0 0	0 0	0	4	0	0	0 0 0	0	0	0 0	0	0	0	0	0 0	0
0x5E0	OTG_HS_HCC HAR7	CHENA	CHDIS	ODDFRM	DA	AD	МС	EPTYP	LSDEV	Reserved	EPDIR		EPNUM				N	IPSI	Z			
	Reset value	0	0	0	0 0 0 0	0 0 0	0 0	0 0	0	В	0	0	0 0 0	0	0	0 0	0	0	0	0	0 0	0 0
0x600	OTG_HS_HCC HAR8	CHENA	CHDIS	ODDFRM	DA		МС	EPTYP	LSDEV	Reserved	EPDIR		EPNUM					IPSI				
	Reset value	0	0	0 1	0 0 0 0	0 0 0	0 0	+	0		0	0	0 0 0	0	0	0 0	0	0	0	0	0 0	0
0x620	OTG_HS_HCC HAR9	CHENA	CHDIS	ODDFRM	DA	<b>A</b> D	МС	EPTYP	LSDEV	Reserved	EPDIR	ļ	EPNUM				N	IPSI	Z			
	Reset value	0	0	0	0 0 0 0	0 0 0	0 0	+	0	н	0	0	0 0 0	0	0	0 0	0	0	0	0	0 0	0
0x640	OTG_HS_HCC HAR10	CHENA	CHDIS	ODDFRM	DA	AD	МС	EPTYP	LSDEV	Reserved	EPDIR		EPNUM				N	IPSI	Z			
	Reset value	0	0	0	0 0 0 0	0 0 0	0 0	0 0	0	В	0	0	0 0 0	0	0	0 0	0	0	0	0	0 0	0 0
0x660	OTG_HS_HCC HAR11	CHENA	Reserved  Reserved  Reserved  Reserved  Reserved  Reserved															IPSI	Z			
	Reset value	0	0	0	0 0 0 0	0 0 0	0 0	0 0	0		0	0	0 0 0	0	0	0 0	0	0	0	0	0 0	0
0x504	OTG_HS_HCS PLT0	SPLITEN		Wald   DAD   MC   A   A   A   A   A   A   A   A   A															PRT	ADD	ıR	
	Reset value	0									0	0	0 0 0	0	0	0 0	0	0	0		0 0	0
0x508	OTG_HS_HCIN T0					F	Reserve	i											SIALL	CHH		
	Reset value									l⊢.	Ι		Π	U	U	0 0	0	0	0	0	) (	, 0
0x524	OTG_HS_HCS PL1	SPLITEN				Reserved				COMPLSPLT	SOGTOXX		HU	BAC	DDR				PRT	ADD	R	
	Reset value	0								0	0	0	0 0 0	0	0	0 0	+	0	0		0 0	0
0x528	OTG_HS_HCIN T1					F	Reserve	i						O DTERR	o FRMOR	o BBERR o TXERR		o ACK		O SIALL	Ţ	o XFRC
0x544	OTG_HS_HCS PLT2	SPLITEN	Reserved   Line   Complex   Complex																PRT			<u> </u>
	Reset value	0								0			0 0 0	0	0	0 0	0	0	0	0	0 0	0 0
0x548	OTG_HS_HCIN T2					F	Reserve	d						o DTERR	o FRMOR	o BBERR o TXERR		o ACK		O SIALL	CHH	
										COMPLSPLT	XACTBOS	)		<u>, ,                                   </u>	_ ~	-   0	Ť	_ ~	-	-	- 1 -	
0x564	OTG_HS_HCS PLT3	SPLITEN						DDR				PRT										
	Reset value	0								0	0	0	0 0 0	0	0	0 0	0	0	0	0	) (	0

Table 166. OTG\_HS register map and reset values (continued)

Offset	Register	_	1 1 1 2 2 2 2 2 2 2 8 2 2 2 2 8 2 8 2 2 2 2	-	15 4	12 11	10	6	8	7	(O	2	4 (	e c	1 -	0
		က		_								-,				+
0x568	OTG_HS_HCIN T3		Reserved				DTERR	FRMOR	BBERR	TXERR	NYET	ACK	NAK	STALL	SH	XFRC
	Reset value					1	0	0	0	0	0	0	0	0 0		0
0x584	OTG_HS_HCS PLT4	SPLITEN	Reserved	COMPLSPLT	XACTPOS	HUI	BAD	DR				í	PRT	ADDI	3	
	Reset value	0		0	0 0	0 0 0	0	0	0	0	0	0	0	0 0		0
0x588	OTG_HS_HCIN T4		Reserved							TXERR				SIALL		
	Reset value			T <del></del>		<u> </u>	0	0	0	0	0	0	0	0 0	0	0
0x5A4	OTG_HS_HCS PLT5	SPLITEN	Reserved	OCOMPLSPLT	XACTPOS	HUI	BAC	DR				ı	PRT	ADDI	3	
	Reset value	0		0	0 0	0 0 0	0	0	0	0	0	0	0	0 0		0
0x5A8	OTG_HS_HCIN T5		Reserved				DTERR			TXERR	NYET	ACK	NAK	SIALL	CHH	XFRC
	Reset value					ı	0	0	0	0	0	0	0	0 0	0	0
0x5C4	OTG_HS_HCS PLT6	SPLITEN	Reserved	COMPLSPLT	XACTPOS	HUI	BAD	DR				ı	PRT	ADDI	3	
	Reset value	0		0	0 0	0 0 0	0	0	0	0	0	0	0	0 0	_	0
0x5C8	OTG_HS_HCIN T6		Reserved							TXERR				SIALL		XFRC
	Reset value			1		I	0	0	0	0	0	0	0	0 0	0	0
0x5E4	OTG_HS_HCS PLT7	SPLITEN	Reserved	OCOMPLSPLT	XACTPOS	HUI	BAD	DR				ſ	PRT	ADDI	3	
	Reset value	0		0	0 0	0 0 0	0	0	0	0	0	0	0	0 0		0
0x5E8	OTG_HS_HCIN T7		Reserved				o DTERR	o FRMOR	o BBERR	o TXERR	o NYET			O STALL AHBERR		
0x604		SPLITEN	Reserved	COMPLSPLT	XACTPOS	HUI				<u> </u>				ADDI	<u> </u>	1-
	Reset value	0		0	0 0	0 0 0	0	0	0	0	0	0	0	0 0	_	0
0x608	OTG_HS_HCIN T8		Reserved				o DTERR	o FRMOR	o BBERR	o TXERR	o NYET			O STALL AHBERR	HS 0	
				Ļ	(n		-	-				- 1	<u>- 1</u>	- 1		<u> </u>
0x624	OTG_HS_HCS PLT9	SPLITEN	Reserved	COMPLSPLT	XACTPOS	HUI								ADDI		
	Reset value	0		0	0 0	0 0 0	0	0	0	0	0	0	0	0 0	0	0

Table 166. OTG\_HS register map and reset values (continued)

Offset	1	-	0	(a) (a) == 1 == 1 == 1
Onset	Register	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1		0 1 2 3 4 5 6
0x628	OTG_HS_HCIN T9	Reserved	PTERR FRMOR BBERR TXERR	ACK NAK STALL AHBERR CHH XFRC
	Reset value		0 0 0 0	0 0 0 0 0 0 0
0x644	OTG_HS_HCS PLT10	Reserved Reserved HU	JBADDR	PRTADDR
	Reset value	0 0 0 0 0 0	0 0 0 0	
0x648	OTG_HS_HCIN T10	Reserved	PRMOR BBERR TXERR	NYET ACK NAK STALL AHBERR CHH XFRC
	Reset value		0 0 0 0	0 0 0 0 0 0 0
0x664	OTG_HS_HCS PLT11	Heserved NAC X	JBADDR	PRTADDR
	Reset value	0 0 0 0 0 0	0 0 0 0	0 0 0 0 0 0 0
0x668	OTG_HS_HCIN T11	Reserved	FRMOR BBERR TXERR	NYET ACK NAK STALL AHBERR CHH XFRC
	Reset value		0 0 0 0	0 0 0 0 0 0 0
0x50C	OTG_HS_HCIN TMSK0	Reserved	o DTERRIM o FRMORIM o BBERRIM o TXERRIM	NYET ACKM NAKM STALLM CHHM XFRCM
	Reset value			0 0 0 0 0 0 0
0x52C	OTG_HS_HCIN TMSK1	Reserved	o PTERRM o FRMORM o BBERRM o TXERRM	NYET ACKM NAKM STALLM AHBERR CHHM
	Reset value			0 0 0 0 0 0 0
0x54C	OTG_HS_HCIN TMSK2	Reserved	DTERRM FRMORM BBERRM TXERRM	NYET ACKM NAKM STALLM AHBERR CHHM
	Reset value		0 0 0 0	0 0 0 0 0 0 0
0x56C	OTG_HS_HCIN TMSK3	Reserved	o DTERRIM o FRMORIM o BBERRIM o TXERRIM	ACKM NAKM STALLM AHBERR CHHM
	Reset value			0 0 0 0 0 0 0
0x58C	OTG_HS_HCIN TMSK4	Reserved	DTERRM FRMORM BBERRM TXERRM	NYET ACKM NAKM STALLM AHBERR CHHM
	Reset value		0 0 0 0	0 0 0 0 0 0 0
0x5AC	OTG_HS_HCIN TMSK5	Reserved	DTERRM FRMORM BBERRM TXERRM	NYET ACKM NAKM STALLM AHBERR CHHM
	Reset value		0 0 0 0	0 0 0 0 0 0 0
0x5CC	OTG_HS_HCIN TMSK6	Reserved	o DTERRIM o FRMORIM o BBERRIM o TXERRIM	NYET ACKM NAKM STALLM AHBERR CHHM
	Reset value			0 0 0 0 0 0 0
0x5EC	OTG_HS_HCIN TMSK7	Reserved	DTERRM FRMORM BBERRM TXERRM	ACKM NAKM STALLM AHBERR CHHM
	Reset value		0 0 0 0	0 0 0 0 0 0 0

Table 166. OTG\_HS register map and reset values (continued)

Offset	Register	31	30	28	27	26	25	3 4	22	1 2	20	19	18	17	16	15	14	13	12	7	10	6	8	7	9	2	4	က	7	- 0	,
0x60C	OTG_HS_HCIN TMSK8							ı		eser										_	_	FRMORM	BBERRM	TXERRM	NYET	ACKM	NAKM	STALLM	AHBERR	CHHM	
	Reset value																				0	0	0	0	0	0	0	0	0	0 0	1
0x62C	OTG_HS_HCIN TMSK9								R	eser	ved										DTERRM	FRMORM	BBERRM	TXERRM	NYET	ACKM	NAKM	STALLM	AHBERR	CHHM	
	Reset value																				0	0	0	0	0	0	0	0	0	0 0	]
0x64C	OTG_HS_HCIN TMSK10								R	eser	ved										DTERRM	o FRMORM		TXERRM					~	CHHM	
	Reset value																						0	0	0	0	0	0	0	0 0	1
0x66C	OTG_HS_HCIN TMSK11								R	eser	ved													TXERRM					⋖.	CHHM	
	Reset value OTG_HS_HCTS	þ	DPID				DI	CTCN	т				<u> </u>								0	0 RSI	0	0	0	0	0	0	0	0 0	-
0x510	IZ0	Reserved		_				_		Τ.	Ι.	_		_	_	_		_						_		_	I . I				_
	Reset value  OTG_HS_HCTS		0 0	0	0	0			0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	_
0x530	IZ1	Reserved	DPID	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0																XF	RSI	Z									
	Reset value		0 0	DPID PKTCNT 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0														0	0	0	0	0	0	0	0	0	0	0	0 0		
0x550	OTG_HS_HCTS IZ2	rved	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0																XF	RSI	Z										
0,000	Reset value	Rese	0 0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	
0x570	OTG_HS_HCTS IZ3		DPID	DPID PKTCNT  0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0																XF	RSI	Z								1	
0.570	Reset value	Reserved	0 0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	
	OTG_HS_HCTS IZ4	_	DPID				Pł	CTCN	Т	<u> </u>	<u> </u>										XF	RSI	Z				!				1
0x590	Reset value	Reserved	0 0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	_
OvED0	OTG_HS_HCTS IZ5	rved	DPID				Pł	CTCN	Т		<u> </u>										XF	RSI	Z								1
0x5B0	Reset value	Rese	0 0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	1
0x5D0	OTG_HS_HCTS IZ6		DPID				Pł	CTCN	Т												XF	RSI	Z								
UXSDU	Reset value	Reserved	0 0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	
0x5F0	OTG_HS_HCTS IZ7	rved	DPID				Pł	CTCN	Т												XF	RSI	Z								
0,010	Reset value	Reser	0 0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	
0010	OTG_HS_HCTS IZ8	rved	DPID				Pł	CTCN	Т												XF	RSI	Z								
0x610	OTG_HS_HCTS IZ8 Reset value	Resei	0 0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	-
0.000	OTG_HS_HCTS IZ9	rved	DPID				Pł	CTCN	Т	l .											XF	RSI	Z								-
0x630	Reset value	Reserved	0 0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	-
	OTG_HS_HCTS IZ10		DPID				Pł	CTCN	Т												XF	RSI	Z								1
0x650	Reset value	Reserved	0 0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	-
0070	OTG_HS_HCTS IZ11	rved	DPID		1		Pł	CTCN	Т	1	<u> </u>	<u> </u>									XF	RSI	Z		<u> </u>						1
0x670	Reset value	Reser	0 0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	1
0,514	OTG_HS_HCD	4		1	1				<u> </u>	1	1	1	<b>.</b>	DI	MAA	ADD	R								1		<b>.</b>			L	1
0x514	MA0 Reset value	0	0 0	0	0	0	0	0 0	0	0	0	0	0	0	0	0		0	0	0	0	0	0	0	0	0	0	0	0	0 0	-
-				•	•				•		•	•					_									_					_

Table 166. OTG\_HS register map and reset values (continued)

Offset	Dogiotos –		٦	m	_	<u></u>	<u> </u>	1	2 0	σī	_			<u>. T~</u>	<u>,                                    </u>			<u> </u>	<del>≠</del> 1	~	<u> </u>	_			T	T	T	1		Т	Г		
Olisei	Register	30	5	28	2	26	22	ì	ίċ	ĭ	'n	2(	ř	2	1 9	- ;	9	1	14	13	12	÷	10	6	8	7	9	5	4	3	7	1	0
0x524	OTG_HS_HCD MA1															DM	IAA	DDF	}														
	Reset value	0 0	0	0	0	0	0 0	(	) (	)	0	0	0	0	) (	)	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	OTG_HS_HCD															DM	IAA	DDF	₹														
0x544	MA2				_						_	_									_	_											
	Reset value	0 0	0	0	0	0	0 0	1	) (	)	0	0	0	0	) (	)	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x564	OTG_HS_HCD MA3															DM	IΑΑ	DDF	?														
0,504	Reset value	0 0	0	0	0	0	0   0	10	) [ (	) T	0	0	0	0	) [ (	) I	0	0 [	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	OTG_HS_HCD	0 0	Ŭ	Ŭ			<u> </u>	<u> </u>	<u>,                                    </u>			_	Ŭ	Ţ	_			_			Ŭ	_	Ŭ	Ŭ	ľ	ľ	Ů	Ŭ	Ŭ				Ť
0x584	MA4															DIV	IAA	DDF	1														
	Reset value	0 0	0	0	0	0	0 0	(	) (	)	0	0	0	0	) (	)	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	OTG_HS_HCD															DM	IAA	DDF	}														
0x5A4	MA5 Reset value	010	10	I 0 I	Λ.	Λ.	010	1 /	1 I	· T	Λ.	Λ	Ι Λ	10						Λ.	Λ.	Λ	<u> </u>	Ι Λ	Ι Λ	Ι Λ	Ι Λ	Ι Λ	Ι Λ	Ι Λ	Ι Λ	Λ.	٥
	OTG HS HCD	0 0	0	0	0	0	0 0	Ι,	) (	)	0	0	0	0		_			0	0	0	0	0	0	0	0	0	0	0	0	0	0	U
0x5C4	MA6															DM	IAA	DDF	}														
	Reset value	0 0	0	0	0	0	0 0	(	) (	)	0	0	0	0	) (	)	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	OTG_HS_HCD	-							•				•		-	DM	ΙΔΔ	DDF	≀ .					•		•							
0x5E4	MA7										_	_										_											
	Reset value	0 0	0	0	0	0	0 0	(	) (	)	0	0	0	0	) (	)	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x604	OTG_HS_HCD MA8															DM	IAA	DDF	?														
0,000	Reset value	0 0	0	0	0	0	0 0	Т	) I (	) T	0	0	0	Το	) [ (	) [	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	OTG_HS_HCD			11				-		- 1			1					DDF	<u> </u>					1	1	1	<u> </u>		<u> </u>		<u> </u>		
0x624	MA9																		1														
	Reset value	0 0	0	0	0	0	0 0	(	) (	)	0	0	0	0	) (	)	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0044	OTG_HS_HCD MA10															DM	IAA	DDF	?														
0x644	Reset value	0 0	0	0	0	0	0 0	1 (	) (	)	0	0	0	0	1 (	)	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	٥
	OTG HS HCD	0 0	U	U	U	0	0   0	`	<u>,                                    </u>	<u>,                                    </u>	•		0	10		_				٠,	Ü		U	U	10	10	١,٠	U	١٠			U	•
0x664	MA11															DIM	IAA	DDF	1														
	Reset value	0 0	0	0	0	0	0 0	(	) (	)	0	0	0	0	) (	)	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
							PERSCHIVL																							1	NZLSOHSK		
	OTG_HS_						R	Recerved	2												PFIVL					DAD				Reserved	占	DSPD	
0x800	DCFG		Rese	erve	t		RS	g	00					Res	serv	/ed					Н					Δ				ser	LS.	DS	
								ď	-																					Re			
	Reset value						1 0														0	0	0	0	0	0	0	0	0		0	0	0
																						POPRGDNE	ㅊ	¥	$^{\star}$	$\mathbf{x}$				တ	S		IJ
	OTG_HS_DCTL																					9	Ž	Š	Ž	ΑŽ		TCTL		ST	ST	SDIS	S
0x804									Re	sei	rve	d										H	CGONAK	SGONAK	CGINAK	SGINAK		2		GONSTS	GINSTS	JS	RWUSIG
	Reset value																					0	0	0	0	0	0	0	0	0	0		0
																														m	CIMSPI		SUSPSTS
0x808	OTG_HS_DSTS			R	250	rved											F	NSC	ϽF							١,	Res	erve	М	EERR	N N		SPS
0,1000																												0	•	Ш	Z		303
	Reset value									ŀ	0	0	0	0	) (	)	0	0	0	0	0	0	0	0	0	1				0	0	0	0
										_			1												L	_	_	5	쏬				
	OTG_HS_DIEP																							_	TXFURM	Reserved	INEPNEM	INEPNMM	ITTXFEMSK	⋝	pe	M	XFRCM
0x810	MSK									R	ese	erve	ed											BIM	匠	sel	FP	F	Ä	TOM	ē	EPDM	Ĕ
																									F	Re	Ξ	Ξ	Ê		Reserved	ш	×
	Reset value																							0	0	0	0	0	0	0	Ι"	0	0
																									1		Ы	ō	Σ	_			1
	OTG_HS_DOE																							BOIM	OPEM	Reserved	B2BSTUP	Reserved	OTEPDM	STUPM	Reserved	EPDM	XFRCM
0x814	PMSK									R	ese	erve	ed											BC	P	ser	2B(	ese	밀	3T	ser	EP	Ŧ.
	Danaturalura																							Ļ		Be					Re		
	Reset value  OTG_HS_DAIN																							0	0		0	0	0	0		0	0
0x818	T T						OE	ΝI	VΤ																IEF	TNI	Г						
	Reset value	0 0	0	0	0	0	0 0	(	) (	)	0	0	0	0	) (	)	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	OTG_HS_DAIN	ı	•				OE	ΞPI	м.	•				•			1		1					•	IF	PM							
0x81C	TMSK	010			0	0 1				, <sub>T</sub>	0	_	1 ^	1 ^		, ,	$\perp$	<u> </u>	<u>, , </u>	0	0 1	^		I ^							_	0	_
1	Reset value	0 0	0	0	0	0	0 0	Τ(	) (	)	0	0	0	0	) (	י ו	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	U

Table 166. OTG\_HS register map and reset values (continued)

Offset	Register	31	30	29	27	26	25 23 23	21	20	19	17	16	15	14 13	12	<b>- -</b>	2 6	8	7	9	2	4	3	7	- 0	>
0x828	OTG_HS_DVB USDIS						Reserved										,	VBL	JSD'	Т						
	Reset value												0	0 0	1	0 .	1 1	1	1	1	0	1	0	1	1 1	Π
0x82C	OTG_HS_DVB USPULSE						Rese	erve	h										ı	DVB	USF	•				
OXOZO	Reset value						11000	J. V.	u							0	1 0	1	1	0	1	1	1	0	0 (	
0x830	OTG_HS_DTH RCTL	Re	eser	rved	ARPEN	Reserved	RXT	ΓHR	LEN	I		RXTHREN		Reser	ved				TXT	ΓHR	LEN	l			ISOTHREN	
	Reset value				С		0 0 0 0	0	0	0 0	0	0				(	0	0	0	0	0	0	0	0	0 (	)
0x834	OTG_HS_DIEP EMPMSK						Reserved									_		EP1								
	Reset value OTG_HS_DEA												0	0 0	0	0   0	0 0	0	0	0	0	0	0	0	0 (	
0x838	CHINT					ı	Reserved									F	Reser	ved							200	;
	Reset value										0														0	2
0x83C	OTG_HS_DEA CHINTMSK						Reserved									_	Reser	hav							הפעיי	5
0,000	Reset value					'	iesei veu				0						10301	veu							O O O	)
0x840	OTG_HS_DIEP EACHMSK1						Reserve	d						NAKM	Res	erve	b BIM	TXFURM	Reserved	INEPNEM		Ш		Reserved	EPDM	
	Reset value													0			0	0		0	0	0	0		0 (	)
0x880	OTG_HS_DOE PEACHMSK1						Reserved		NYETM	BERRM	Reserved	BIM	TXFURM	Reserved	INEPNEM	INEPNMM	TTXFEMSK	TOM	Reserved	EPDM						
	Reset value							0 0	0		0	0		0	0	0	0		0 (	)						
0x900	OTG_HS_DIEP CTL0	EPENA	O   SODDFRM   O   SODDFRM   O   O   O   O   O   O   O   O   O																N	IPSI	Z					
	Reset value	0	0			0	0 0 0 0	0		0 0	0	0	0			(	0 0	0	0	0	0	0	0	0	0 (	)
0x918	TG_FS_DTXFS TS0						Reserved										IN	IEP	ΓFS	ΑV						
	Reset value						1			ı			0	0 0	0	0 (	) 1	0	0	0	0	0	0	0	0 (	)
0x920	OTG_HS_DIEP CTL1	EPENA	EPDIS	SODDFRM SDOPID/SEXNERM		CNAK	TXFNUM	Stall	Reserved	ЕРТҮР	NAKSTS	EONUM/DPID	USBAEP	Res	erved					N	IPSI	Z				
	Reset value TG_FS_DTXFS	0	0	0 (	) (	0	0 0 0 0	0		0 0	0	0	0			(	0	0	0	0	0	0	0	0	0 (	)
0x938	TS1						Reserved										IN	IEP								
	Reset value						1			T		1	0	0 0	0	0 (	) 1	0	0	0	0	0	0	0	0 (	)
0x940	OTG_HS_DIEP CTL2			SODDFRM STOPID/SEVNERM			TXFNUM	Stall	Reserved	ЕРТҮР	NAKSTS	EC	USBAEP	Res	erved			10			IPSI			0.1	<u> </u>	
<u> </u>	Reset value TG_FS_DTXFS	0	0	0 (	) (	0	0 0 0 0	0		0 0	0	0	0			(	0 0	0	1	0	0	0	0	U	0 0	4
0x958	TS2 Reset value						Reserved						0	0 0	0	0   0		IEP <sup>-</sup>		AV 0	0	0	0	0	0   0	)
	<u> </u>																- 1	1								_

Table 166. OTG\_HS register map and reset values (continued)

Table I				_			_	and rese								
Offset	Register	31	30	29		27	26	25 23 23 24	21	20	19	17	16	15	13 14 11	10 8 8 7 7 7 7 1 1 1 1 0
0x960		EPENA	EPDIS SODDFRM SDOPID/SEVNFRM SNAK CNAK					TXFNUM	Reserved	ЕРТҮР	NAKSTS	EONUM/DPID	USBAEP	Reserved	MPSIZ	
	Reset value	0	0	0	0	0	0	0 0 0 0	0		0 0	0	0	0		0 0 0 0 0 0 0 0 0 0 0 0
0x978	TG_FS_DTXFS TS3							Reserved								INEPTFSAV
	Reset value													0	0 0 0 0	0 1 0 0 0 0 0 0 0 0 0 0
0x980	OTG_HS_DIEP CTL4	EPENA			SD0PID/SEVNFRM		CNAK	TXFNUM	Stall	Reserved	ЕРТҮР	NAKSTS	EONUM/DPID	USBAEP	Reserved	MPSIZ
	Reset value	0	0	0	0	0	0	0 0 0 0	0		0 0	0	0	0		0 0 0 0 0 0 0 0 0 0 0 0
0x9A0	OTG_HS_DIEP CTL5	EPENA	EPDIS	SODDFRM	SD0PID/SEVNFRM	SNAK	CNAK	TXFNUM	STALL	Reserved	EPTYP	NAKSTS	EONUM/DPID	USBAEP	Reserved	MPSIZ
	Reset value	0	0	0	0	0	0	0 0 0 0	0		0 0	0	0	0		0 0 0 0 0 0 0 0 0 0 0
0x9C0	OTG_HS_DIEP CTL6	EPENA	EPDIS	SODDFRM	SD0PID/SEVNFRM	SNAK	CNAK	TXFNUM	STALL	Reserved	ЕРТҮР	NAKSTS	EONUM/DPID	USBAEP	Reserved	MPSIZ
	Reset value	0	0	0	0	0	0	0 0 0 0	0	1	0 0	0	0	0		0 0 0 0 0 0 0 0 0 0 0 0
0x9E0	OTG_HS_DIEP CTL7	EPENA	EPDIS	SODDFRM	SD0PID/SEVNFRM	SNAK	CNAK	TXFNUM	STALL	Reserved	ЕРТУР	NAKSTS	EONUM/DPID	USBAEP	Reserved	MPSIZ
	Reset value	0	0	0	0	0	0	0 0 0 0	0		0 0	0	0	0		0 0 0 0 0 0 0 0 0 0 0
0xB00	OTG_HS_DOE PCTL0	o EPENA	o EPDIS	Becerved	ט	o SNAK	o CNAK	Reserved	O STALL	O SNPM	EPTY P	ONAKSTS	Reserved	→ USBAEP		Reserved MPSI Z 0 0 0
0xB20	OTG_HS_DOE PCTL1	EPENA	EPDIS	SODDFRM	NFRM		CNAK	Reserved	STALL	SNPM	ЕРТҮР	NAKSTS	EONUM/DPID	USBAEP	Reserved	MPSIZ
	Reset value	0	0	0	0	0	0		0	0	0 0	0	0	0		0 0 0 0 0 0 0 0 0 0 0 0
0xB40	OTG_HS_DOE PCTL2	e EPENA			SD0PID/SEVNFRM		CNAK	Reserved	Stall	SNPM	ЕРТУР	NAKSTS	EONUM/DPID	USBAEP	Reserved	MPSIZ
	Reset value	0	0	0	0	0	0		0	0	0 0	0	0	0		0 0 0 0 0 0 0 0 0 0 0 0

Table 166. OTG\_HS register map and reset values (continued)

Offset	Register	31	30	53	28	27	26	27	22	17	20	6 8	17	16	15	14	3	12	Ξ	10	6	8	7	9	2	4	က	2	_	0
0xB60	OTG_HS_DOE PCTL3	_	SS SD S															erve		MPSIZ										
	Reset value	0	0	0	0	0	0			0	0	0 0	0	0	0	<u> </u>			1	0	0	0	0	0	0	0	0	0	0	0
0x908	OTG_HS_DIEPI NT0		Reserved														ONAK		PKTDRPSTS .	Reserved	BNA	TXFIFOUDRN	TXFE	INEPNE	Reserved	ITTXFE	TOC	Reserved	3	XFRC
	Reset value																	0	0		0	0	1	0		0	0		0	0
0x928	OTG_HS_DIEPI NT1		NAK BERR BERR PKTDRPSTS Reserved														TXFIFOUDRN	TXFE	INEPNE	Reserved	ITTXFE	T0C	Reserved	3	XFRC					
	Reset value																0	-	0		0	0	1	0		0	0		0	0
0x948	OTG_HS_DIEPI NT2		Reserved														NAK	BERR	PKTDRPSTS	Reserved	BNA	TXFIFOUDRN	TXFE	INEPNE	Reserved	ITTXFE	TOC	Reserved	EPDISD	XFRC
	Reset value											0	0	0		0	0	1	0		0	0		0	0					
0x968	OTG_HS_DIEPI NT3		Reserved												NAK	BERR	PKTDRPSTS	Reserved	BNA	TXFIFOUDRN	TXFE	INEPNE	Reserved	ITTXFE	TOC	Reserved	EPDISD	XFRC		
	Reset value	value															0	0	0		0	0	1	0		0	0		0	0
0x988	OTG_HS_DIEPI NT4							R	eserve	ed							NAK	BERR	PKTDRPSTS	Reserved	BNA	TXFIFOUDRN	TXFE	INEPNE	Reserved	ITTXFE	TOC	Reserved	EPDISD	XFRC
	Reset value																0	0	0		0	0	1	0		0	0		0	0
0x9A8	OTG_HS_DIEPI NT5							R	eserve	∍d							NAK		PKTDRPSTS	Reserved	BNA	TXFIFOUDRN	TXFE	INEPNE	Reserved	_	, TOC	Reserved	3	XFRC
	Reset value																0	-	0		0	0	1	0		0	0		0	0
0x9C8	OTG_HS_DIEPI NT6		Reserved													NAK	BERR	PKTDRPSTS	Reserved	BNA	TXFIFOUDRN	TXFE	INEPNE	Reserved	ITTXFE	TOC	Reserved	EPDISD	XFRC	
	Reset value																0	0	0		0	0	1	0		0	0		0	0
0x9E8	OTG_HS_DIEPI NT7		Reserved												NAK		PKTDRPSTS	Reserved	BNA	TXFIFOUDRN	TXFE	INEPNE	Reserved	ITTXFE	TOC	Reserved	3	XFRC		
	Reset value														0	0	0		0	0	1	0		0	0		0	0		
0xB08	OTG_HS_DOE PINT0		Reserved											o NYET	1 m   m   m   m   m   m   m   m   m   m							o EPDISD	O XFRC							
L	i iosot value															J								Ľ		Ľ	<u>'</u>		J	Ŭ

Table 166. OTG\_HS register map and reset values (continued)

Offset	Register	31	0	ဂ္ဂ	œ	7	ဖွ	25	4 9	<u>ლ</u>	2	Σ.	0	6	œ	17	16	15	14	13	12	=	10	6	α	<b>^</b>	9	2	4	ဗ	7	_	0
	3.2.2.	eo	ന	N	7	2	N	N I	2   2	2	7	N	N	듸	드	_		Ξ	1	_		Ξ	_						_	Ë			H
0xB28	OTG_HS_DOE PINT1							ı	Rese	erve	ed								NYET			Re	sei	ved	l		OB2BSTUP	Reserved	OTEPDIS	STUP	Reserved	EPDISD	XFRC
	Reset value																		0									ш.	0	0	ш.	0	0
0xB48	OTG_HS_DOE PINT2							ı	Rese	erve	ed								NYET			Re	ser	ved	J		o B2BSTUP	Reserved	OTEPDIS	STUP	Reserved	EPDISD	XFRC
	Reset value																		0									ш	0	0	ш	0	0
0xB68	OTG_HS_DOE PINT3		Reserved												NYET			Re	ser	ved	l		o B2BSTUP	Reserved	OTEPDIS	STUP	Reserved	EPDISD	XFRC				
	Reset value																		0									_	0	0	_	0	0
0xB88	OTG_HS_DOE PINT4							ı	Res	erve	ed								NYET	Reserved							o B2BSTUP	Reserved	OTEPDIS	STUP	Reserved	EPDISD	XFRC
	Reset value																0								0	$\vdash$	0	0					
0xBA8	OTG_HS_DOE PINT5							ı	Rese	erve	ed								NYET			Re	sei	ved	l		o B2BSTUP	Reserved	OTEPDIS	STUP	Reserved	EPDISD	XFRC
	Reset value																		0									ш	0	0	_	0	0
0xBC8	OTG_HS_DOE PINT6		Reserved													NYET	Reserved						o B2BSTUP	Reserved	OTEPDIS	STUP	Reserved	EPDISD	XFRC				
	Reset value																		0									Ж	0	0	æ	0	0
0xBE8	OTG_HS_DOE PINT7		Reserved												NYET					Reserved				B2BSTUP	Reserved	OTEPDIS	STUP	Reserved	EPDISD	XFRC			
	Reset value														0								0	Œ	0	0	Œ	0	0				
0x910	OTG_HS_DIEP TSIZ0					Res	serv	ed					PK <sup>-</sup> N <sup>-</sup>						Reserved									XI	FRS	ŝΙΖ			
	Reset value												0	0											0	0	0	0	0	0	0		
0x930	OTG_HS_DIEP TSIZ1	rvec	МС	NT				Pl	KTC	TN					1								Х	FR	SIZ								
OXOGO	Reset value	Reserved	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0x934	OTG_HS_DIEP DMA1															DI	MAA	NDD	R					1									
0,004	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	С	0	0	0	0	0	0	0	0
0x93C	OTG_HS_DIEP DMAB1															DN	1AB	ADI	DR														
	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	C	0	0	0	0	0	0	0	0
0x950	OTG_HS_DIEP TSIZ2	erved	МС	NT				P	KTC	TN					l								Х	FR	SIZ								
0.00	Reset value	Rese	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	О	0	0	0	0	0	0	0	0
0x954	OTG_HS_DIEP DMA2																MAA																
	Reset value OTG_HS_DIEP	0	0	0	0	0	0	0	0	0	0	0	0	0	0		0			0	0	0	0	0	C	0	0	0	0	0	0	0	0
0x95C	DMAB2																1AB																
	Reset value  OTG_HS_DIEP	0 p	0	0	0	0	0				0	0	0	0	0	0	0	0	0	0	0	0	0	•			0	0	0	0	0	0	0
0x970	TSIZ3	Reserved	МС	NT	<u></u>			P	KTC	TM					<u>L</u>								Х	FRS	SIZ								
	Reset value	Res	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	C	0	0	0	0	0	0	0	0
0x974	OTG_HS_DIEP DMA3															DI	MAA	DD	R														
	Reset value	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0												0	0	0	C	0	0	0	0	0	0	0	0								
0x97C	OTG_HS_DIEP DMAB3			_	_		_		_	_	_	_	_	_		DN	1AB	ADI	DR	_	-	_		_	_				_	_	_	_	
	Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	С	0	0	0	0	0	0	0	0

Offset				_	<b>~</b>		+   m	<b>A</b> I			<u> </u>	N (0	101=	<u></u>	ΛI –	Jal			Т					_
Offset	Register	31	30	28	27	25	23	22	20	19	18	17 16	15	13	12	: 우 •	8	7	ין פ	5	3	2	1	<u> </u>
0xB10	OTG_HS_DOE PTSIZ0	Reserved	STUP CNT			Rese	erved			PKTCNT			ا	Rese		XFRSIZ								
	Reset value	В	0 0							0								(	)	0 0	0	0	0	0
0xB30	OTG_HS_DOE PTSIZ1	Reserved	RXDPID/ STUPCNT			Pł	(TCN	Т								XFR	SIZ							
	Reset value	ш	0 0	0	0	0 0 0	0 0	0	0 0	0	0	0 0	0 0	0	0 0	0 0	0	0 (	)	0 0	0	0	0	0
0xB34	OTG_HS_DOE PDMA1											DMA	ADDR											
	Reset value	0	0 0	0	0	0 0 0	0 0	0	0 0	0	0	0 0	0 0	0	0 0	0 0	0	0 (	)	0 0	0	0	0	J
0xB3C	OTG_HS_DOE PDMAB1										DMABADDR													
	Reset value	0	0 0	0	0	0 0 0	0 0	0	0 0	0	0	0 0	0 0	0	0 0	0 0	0	0 (	)	0 0	0	0	0	0
0xB50	OTG_HS_DOE PTSIZ2	Reserved	RXDPID/ STUPCNT	STUPCNT PKTCNT												XFR	ISIZ							
	Reset value	ш	0 0	0	0	0 0 0	0 0	0	0 0	0	0	0 0	0 0	0	0 0	0 0	0	0 (	)	0 0	0	0	0	0
0xB54	OTG_HS_DOE PDMA2												ADDR											
	Reset value	0	0 0	0	0	0 0 0	0 0	0	0 0	0	0	0 0	0 0	0	0 0	0 0	0	0 (	)	0 0	0	0	0	0
0xB5C	OTG_HS_DOE PDMAB2												BADDR											
	Reset value	0	0 0	0	0	0 0	0 0	0	0 0	0	0	0 0	0 0	0	0 0	0 0	0	0 (	)	0 0	0	0	0	0
0xB70	OTG_HS_DOE PTSIZ3	Reserved	RXDPID/ STUPCNT			Pł	(TCN	Т								XFR	SIZ							
	Reset value	ш	0 0	0	0	0 0 0	0 0	0	0 0	0	0	0 0	0 0	0	0 0	0 0	0	0 (	)	0 0	0	0	0	0
0xB74	OTG_HS_DOE PDMA3											DMA	ADDR											
	Reset value	0	0 0	0	0	0 0 0	0 0	0	0 0	0	0	0 0	0 0	0	0 0	0 0	0	0 (	)	0 0	0	0	0	0
0xB7C	OTG_HS_DOE PDMAB3											DMAE	BADDR											
OXD7 O	Reset value	0	0 0	0	0	0 0 0	0 0	0	0 0	0	0	0 0	0 0	0	0 0	0 0	0 0	0 (	) [	0 0	0	0	0	0
0xE00	OTG_HS_PCG CCTL					1 1	1		,	Re	serv	ed								PHYSUSP	Becomed	וספסו אפת	GATEHCLK	טוייייייי
	Reset value																							

Table 166. OTG\_HS register map and reset values (continued)

Refer to *Table 2 on page 50* for the register boundary addresses.

# 30.13 OTG\_HS programming model

## 30.13.1 Core initialization

The application must perform the core initialization sequence. If the cable is connected during power-up, the current mode of operation bit in the Core interrupt register (CMOD bit in OTG\_HS\_GINTSTS) reflects the mode. The OTG\_HS controller enters host mode when an "A" plug is connected or peripheral mode when a "B" plug is connected.

This section explains the initialization of the OTG\_HS controller after power-on. The application must follow the initialization sequence irrespective of host or peripheral mode operation. All core global registers are initialized according to the core's configuration:

- 1. Program the following fields in the Global AHB configuration (OTG\_HS\_GAHBCFG) register:
  - DMA mode bit
  - AHB burst length field
  - Global interrupt mask bit GINT = 1
  - RxFIFO nonempty (RXFLVL bit in OTG\_HS\_GINTSTS)
  - Periodic TxFIFO empty level
- 2. Program the following fields in OTG\_HS\_GUSBCFG register:
  - HNP capable bit
  - SRP capable bit
  - FS timeout calibration field
  - USB turnaround time field
- The software must unmask the following bits in the GINTMSK register:

OTG interrupt mask

Mode mismatch interrupt mask

4. The software can read the CMOD bit in OTG\_HS\_GINTSTS to determine whether the OTG\_HS controller is operating in host or peripheral mode.

## 30.13.2 Host initialization

To initialize the core as host, the application must perform the following steps:

- 1. Program the HPRTINT in GINTMSK to unmask
- 2. Program the OTG HS HCFG register to select full-speed host
- 3. Program the PPWR bit in OTG\_HS\_HPRT to 1. This drives V<sub>BUS</sub> on the USB.
- 4. Wait for the PCDET interrupt in OTG\_HS\_HPRT0. This indicates that a device is connecting to the port.
- 5. Program the PRST bit in OTG\_HS\_HPRT to 1. This starts the reset process.
- 6. Wait at least 10 ms for the reset process to complete.
- 7. Program the PRST bit in OTG HS HPRT to 0.
- 8. Wait for the PENCHNG interrupt in OTG\_HS\_HPRT.
- 9. Read the PSPD bit in OTG HS HPRT to get the enumerated speed.
- 10. Program the HFIR register with a value corresponding to the selected PHY clock 1.
- 11. Program the FSLSPCS field in OTG\_FS\_HCFG register according to the speed of the detected device read in step 9. If FSLSPCS has been changed, reset the port.
- 12. Program the OTG\_HS\_GRXFSIZ register to select the size of the receive FIFO.
- 13. Program the OTG\_HS\_GNPTXFSIZ register to select the size and the start address of the nonperiodic transmit FIFO for nonperiodic transactions.
- 14. Program the OTG\_HS\_HPTXFSIZ register to select the size and start address of the periodic transmit FIFO for periodic transactions.

To communicate with devices, the system software must initialize and enable at least one channel.

## 30.13.3 Device initialization

The application must perform the following steps to initialize the core as a device on power-up or after a mode change from host to device.

- 1. Program the following fields in the OTG HS DCFG register:
  - Device speed
  - Nonzero-length status OUT handshake
- 2. Program the OTG\_HS\_GINTMSK register to unmask the following interrupts:
  - USB reset
  - Enumeration done
  - Early suspend
  - USB suspend
  - SOF
- 3. Program the VBUSBSEN bit in the OTG\_HS\_GCCFG register to enable V<sub>BUS</sub> sensing in "B" peripheral mode and supply the 5 volts across the pull-up resistor on the DP line.
- 4. Wait for the USBRST interrupt in OTG\_HS\_GINTSTS. It indicates that a reset has been detected on the USB that lasts for about 10 ms on receiving this interrupt.

Wait for the ENUMDNE interrupt in OTG\_HS\_GINTSTS. This interrupt indicates the end of reset on the USB. On receiving this interrupt, the application must read the OTG\_HS\_DSTS register to determine the enumeration speed and perform the steps listed in *Endpoint* initialization on enumeration completion on page 1215.

At this point, the device is ready to accept SOF packets and perform control transfers on control endpoint 0.

## 30.13.4 DMA mode

The OTG host uses the AHB master interface to fetch the transmit packet data (AHB to USB) and receive the data update (USB to AHB). The AHB master uses the programmed DMA address (HCDMAx register in host mode and DIEPDMAx/DOEPDMAx register in peripheral mode) to access the data buffers.

# 30.13.5 Host programming model

# **Channel initialization**

The application must initialize one or more channels before it can communicate with connected devices. To initialize and enable a channel, the application must perform the following steps:

- 1. Program the GINTMSK register to unmask the following:
- 2. Channel interrupt
  - Nonperiodic transmit FIFO empty for OUT transactions (applicable for Slave mode that operates in pipelined transaction-level with the packet count field programmed with more than one).
  - Nonperiodic transmit FIFO half-empty for OUT transactions (applicable for Slave mode that operates in pipelined transaction-level with the packet count field programmed with more than one).
- Program the OTG\_HS\_HAINTMSK register to unmask the selected channels' interrupts.
- 4. Program the OTG\_HS\_HCINTMSK register to unmask the transaction-related interrupts of interest given in the host channel interrupt register.
- 5. Program the selected channel's OTG\_HS\_HCTSIZx register with the total transfer size, in bytes, and the expected number of packets, including short packets. The application must program the PID field with the initial data PID (to be used on the first OUT transaction or to be expected from the first IN transaction).
- 6. Program the selected channels in the OTG\_HS\_HCSPLTx register(s) with the hub and port addresses (split transactions only).
- 7. Program the selected channels in the HCDMAx register(s) with the buffer start address.
- 8. Program the OTG\_HS\_HCCHARx register of the selected channel with the device's endpoint characteristics, such as type, speed, direction, and so forth. (The channel can be enabled by setting the channel enable bit to 1 only when the application is ready to transmit or receive any packet).

# Halting a channel

The application can disable any channel by programming the OTG\_HS\_HCCHARx register with the CHDIS and CHENA bits set to 1. This enables the OTG\_HS host to flush the posted requests (if any) and generates a channel halted interrupt. The application must wait for the CHH interrupt in OTG\_HS\_HCINTx before reallocating the channel for other transactions. The OTG\_HS host does not interrupt the transaction that has already been started on the USB.

To disable a channel in DMA mode operation, the application does not need to check for space in the request queue. The OTG\_HS host checks for space to write the disable request on the disabled channel's turn during arbitration. Meanwhile, all posted requests are dropped from the request queue when the CHDIS bit in HCCHARx is set to 1.

Before disabling a channel, the application must ensure that there is at least one free space available in the nonperiodic request queue (when disabling a nonperiodic channel) or the periodic request queue (when disabling a periodic channel). The application can simply flush the posted requests when the Request queue is full (before disabling the channel), by programming the OTG\_HS\_HCCHARx register with the CHDIS bit set to 1, and the CHENA bit cleared to 0.

The application is expected to disable a channel on any of the following conditions:

- 1. When an XFRC interrupt in OTG\_HS\_HCINTx is received during a nonperiodic IN transfer or high-bandwidth interrupt IN transfer (Slave mode only)
- 2. When an STALL, TXERR, BBERR or DTERR interrupt in OTG\_HS\_HCINTx is received for an IN or OUT channel (Slave mode only). For high-bandwidth interrupt INs in Slave mode, once the application has received a DTERR interrupt it must disable the channel

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- and wait for a channel halted interrupt. The application must be able to receive other interrupts (DTERR, NAK, Data, TXERR) for the same channel before receiving the halt.
- 3. When a DISCINT (Disconnect Device) interrupt in OTG\_HS\_GINTSTS is received. (The application is expected to disable all enabled channels
- 4. When the application aborts a transfer before normal completion.

# Ping protocol

When the OTG\_HS host operates in high speed, the application must initiate the ping protocol when communicating with high-speed bulk or control (data and status stage) OUT endpoints.

The application must initiate the ping protocol when it receives a NAK/NYET/TXERR interrupt. When the HS\_OTG host receives one of the above responses, it does not continue any transaction for a specific endpoint, drops all posted or fetched OUT requests (from the request queue), and flushes the corresponding data (from the transmit FIFO).

This is valid in slave mode only. In Slave mode, the application can send a ping token either by setting the DOPING bit in HCTSIZx before enabling the channel or by just writing the HCTSIZx register with the DOPING bit set when the channel is already enabled. This enables the HS\_OTG host to write a ping request entry to the request queue. The application must wait for the response to the ping token (a NAK, ACK, or TXERR interrupt) before continuing the transaction or sending another ping token. The application can continue the data transaction only after receiving an ACK from the OUT endpoint for the requested ping. In DMA mode operation, the application does not need to set the DOPING bit in HCTSIZx for a NAK/NYET response in case of Bulk/Control OUT. The OTG\_HS host automatically sets the DOPING bit in HCTSIZx, and issues the ping tokens for Bulk/Control OUT. The HS\_OTG host continues sending ping tokens until it receives an ACK, and then switches automatically to the data transaction.

# Operational model

The application must initialize a channel before communicating to the connected device. This section explains the sequence of operation to be performed for different types of USB transactions.

## Writing the transmit FIFO

The OTG\_HS host automatically writes an entry (OUT request) to the periodic/nonperiodic request queue, along with the last DWORD write of a packet. The application must ensure that at least one free space is available in the periodic/nonperiodic request queue before starting to write to the transmit FIFO. The application must always write to the transmit FIFO in DWORDs. If the packet size is nonDWORD aligned, the application must use padding. The OTG\_HS host determines the actual packet size based on the programmed maximum packet size and transfer size.

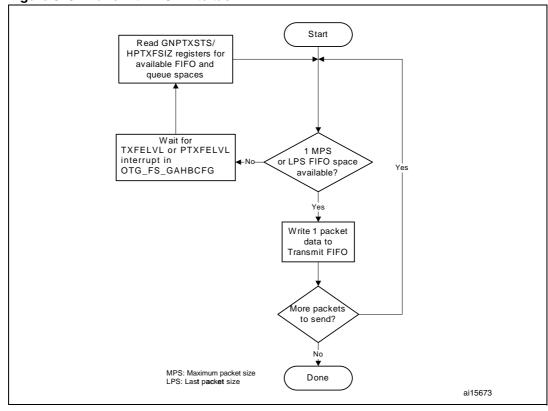
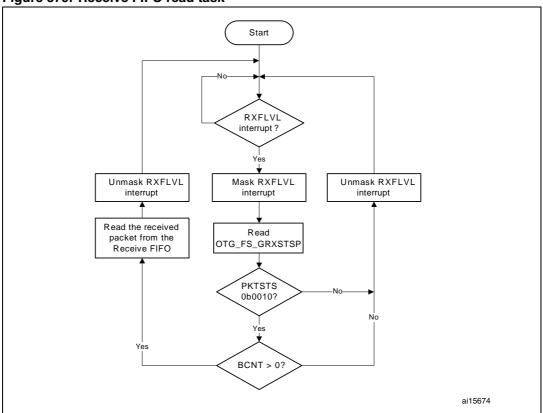


Figure 378. Transmit FIFO write task

## Reading the receive FIFO

The application must ignore all packet statuses other than IN data packet (bx0010).

Figure 379. Receive FIFO read task



### Bulk and control OUT/SETUP transactions

A typical bulk or control OUT/SETUP pipelined transaction-level operation is shown in *Figure 380*. See channel 1 (ch\_1). Two bulk OUT packets are transmitted. A control

SETUP transaction operates in the same way but has only one packet. The assumptions are:

- The application is attempting to send two maximum-packet-size packets (transfer size = 1, 024 bytes).
- The nonperiodic transmit FIFO can hold two packets (128 bytes for FS).
- The nonperiodic request queue depth = 4.

## Normal bulk and control OUT/SETUP operations

The sequence of operations for channel 1 is as follows:

- a) Initialize channel 1
- b) Write the first packet for channel 1
- c) Along with the last DWORD write, the core writes an entry to the nonperiodic request queue
- As soon as the nonperiodic queue becomes nonempty, the core attempts to send an OUT token in the current frame
- e) Write the second (last) packet for channel 1
- f) The core generates the XFRC interrupt as soon as the last transaction is completed successfully
- g) In response to the XFRC interrupt, de-allocate the channel for other transfers
- h) Handling nonACK responses

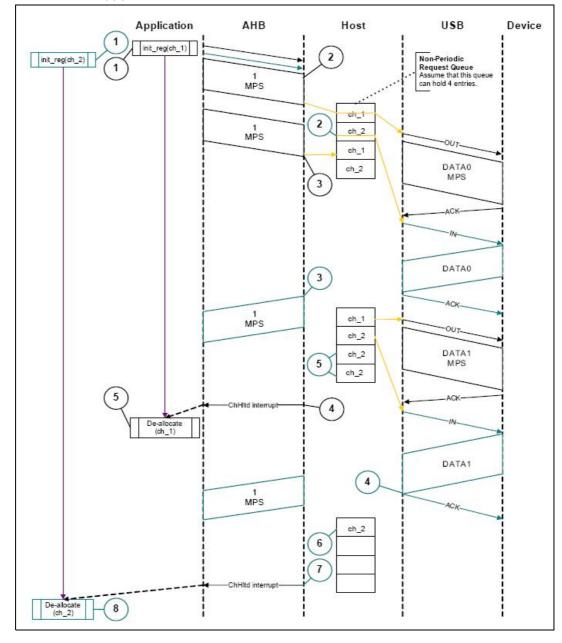


Figure 380. Normal bulk/control OUT/SETUP and bulk/control IN transactions - DMA mode

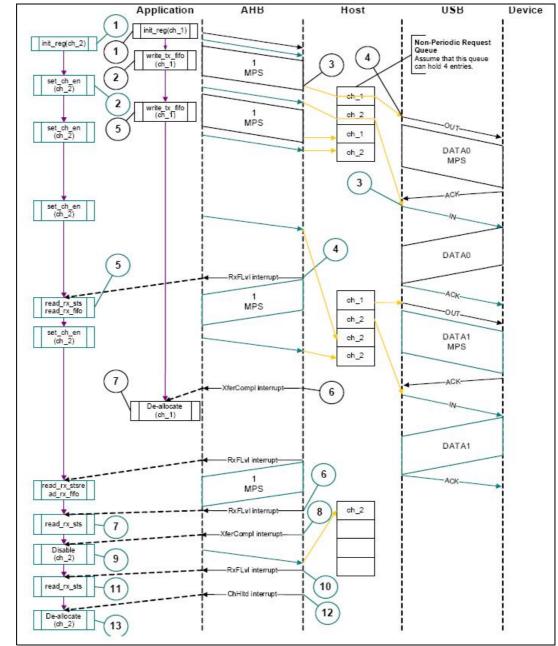


Figure 381. Normal bulk/control OUT/SETUP and bulk/control IN transactions - Slave mode

The channel-specific interrupt service routine for bulk and control OUT/SETUP transactions in Slave mode is shown in the following code samples.

Interrupt service routine for bulk/control OUT/SETUP and bulk/control IN transactions

```
a) Bulk/Control OUT/SETUP
Unmask (NAK/TXERR/STALL/XFRC)
if (XFRC)
{
   Reset Error Count
```

```
Mask ACK
  De-allocate Channel
  }
else if (STALL)
  Transfer Done = 1
  Unmask CHH
  Disable Channel
else if (NAK or TXERR )
  Rewind Buffer Pointers
  Unmask CHH
  Disable Channel
  if (TXERR)
    Increment Error Count
    Unmask ACK
  else
    Reset Error Count
else if (CHH)
  Mask CHH
  if (Transfer Done or (Error count == 3))
    De-allocate Channel
  else
    Re-initialize Channel
  else if (ACK)
    Reset Error Count
    Mask ACK
```

The application is expected to write the data packets into the transmit FIFO as and when the space is available in the transmit FIFO and the Request queue. The application can make use of the NPTXFE interrupt in OTG\_HS\_GINTSTS to find the transmit FIFO space.

### b) Bulk/Control IN

```
Unmask (TXERR/XFRC/BBERR/STALL/DTERR)
if (XFRC)
    {
    Reset Error Count
    Unmask CHH
    Disable Channel
```

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```
Reset Error Count
  Mask ACK
  }
else if (TXERR or BBERR or STALL)
  Unmask CHH
  Disable Channel
  if (TXERR)
    Increment Error Count
    Unmask ACK
else if (CHH)
  Mask CHH
  if (Transfer Done or (Error count == 3))
    De-allocate Channel
  else
    Re-initialize Channel
else if (ACK)
  Reset Error Count
  Mask ACK
else if (DTERR)
  Reset Error Count
```

The application is expected to write the requests as and when the Request queue space is available and until the XFRC interrupt is received.

### Bulk and control IN transactions

A typical bulk or control IN pipelined transaction-level operation is shown in *Figure 382*. See channel 2 (ch\_2). The assumptions are:

- The application is attempting to receive two maximum-packet-size packets (transfer size = 1 024 bytes).
- The receive FIFO can contain at least one maximum-packet-size packet and two status DWORDs per packet (72 bytes for FS).
- The nonperiodic request queue depth = 4.

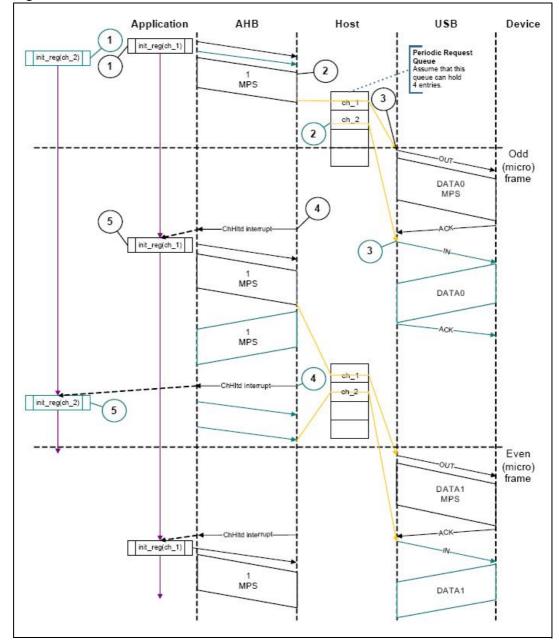


Figure 382. Bulk/control IN transactions - DMA mode

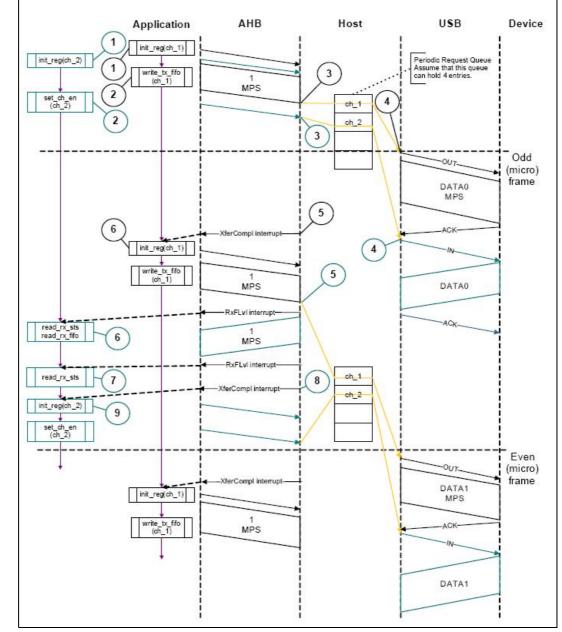


Figure 383. Bulk/control IN transactions - Slave mode

The sequence of operations is as follows:

- a) Initialize channel 2.
- b) Set the CHENA bit in HCCHAR2 to write an IN request to the nonperiodic request queue.
- c) The core attempts to send an IN token after completing the current OUT transaction.
- d) The core generates an RXFLVL interrupt as soon as the received packet is written to the receive FIFO.
- e) In response to the RXFLVL interrupt, mask the RXFLVL interrupt and read the received packet status to determine the number of bytes received, then read the



- receive FIFO accordingly. Following this, unmask the RXFLVL interrupt.
- f) The core generates the RXFLVL interrupt for the transfer completion status entry in the receive FIFO.
- g) The application must read and ignore the receive packet status when the receive packet status is not an IN data packet (PKTSTS in GRXSTSR ≠ 0b0010).
- h) The core generates the XFRC interrupt as soon as the receive packet status is read
- In response to the XFRC interrupt, disable the channel and stop writing the OTG\_HS\_HCCHAR2 register for further requests. The core writes a channel disable request to the nonperiodic request queue as soon as the OTG\_HS\_HCCHAR2 register is written.
- The core generates the RXFLVL interrupt as soon as the halt status is written to the receive FIFO.
- k) Read and ignore the receive packet status.
- The core generates a CHH interrupt as soon as the halt status is popped from the receive FIFO.
- m) In response to the CHH interrupt, de-allocate the channel for other transfers.
- n) Handling nonACK responses

#### Control transactions in slave mode

Setup, Data, and Status stages of a control transfer must be performed as three separate transfers. Setup-, Data- or Status-stage OUT transactions are performed similarly to the bulk OUT transactions explained previously. Data- or Status-stage IN transactions are performed similarly to the bulk IN transactions explained previously. For all three stages, the application is expected to set the EPTYP field in OTG\_HS\_HCCHAR1 to Control. During the Setup stage, the application is expected to set the PID field in OTG\_HS\_HCTSIZ1 to SETUP.

## Interrupt OUT transactions

A typical interrupt OUT operation in Slave mode is shown in *Figure 384*. The assumptions are:

- The application is attempting to send one packet in every frame (up to 1 maximum packet size), starting with the odd frame (transfer size = 1 024 bytes)
- The periodic transmit FIFO can hold one packet (1 KB)
- Periodic request queue depth = 4

The sequence of operations is as follows:

- a) Initialize and enable channel 1. The application must set the ODDFRM bit in OTG HS HCCHAR1.
- b) Write the first packet for channel 1. For a high-bandwidth interrupt transfer, the application must write the subsequent packets up to MCNT (maximum number of packets to be transmitted in the next frame times) before switching to another channel.
- c) Along with the last DWORD write of each packet, the OTG\_HS host writes an entry to the periodic request queue.
- d) The OTG HS host attempts to send an OUT token in the next (odd) frame.
- e) The OTG\_HS host generates an XFRC interrupt as soon as the last packet is transmitted successfully.
- f) In response to the XFRC interrupt, reinitialize the channel for the next transfer.

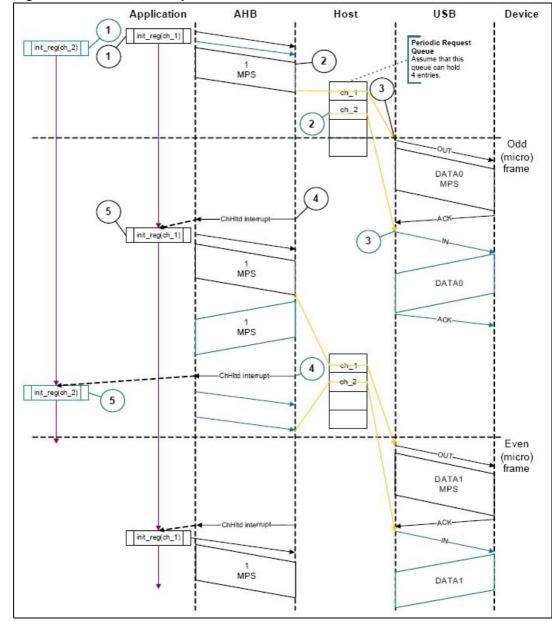


Figure 384. Normal interrupt OUT/IN transactions - DMA mode

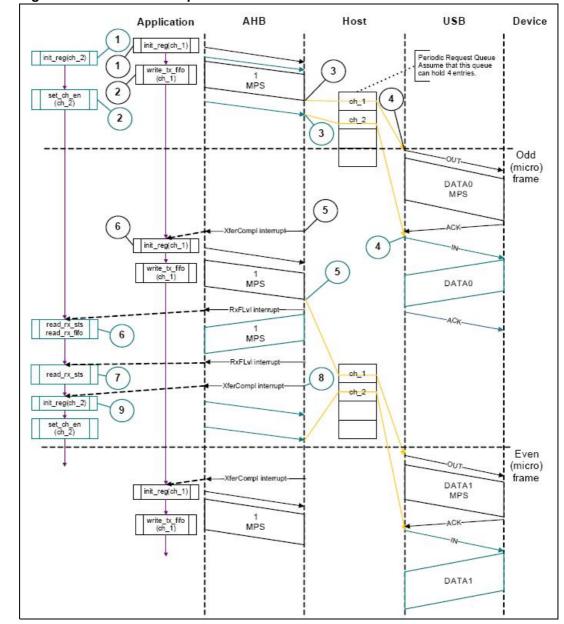


Figure 385. Normal interrupt OUT/IN transactions - Slave mode

## • Interrupt service routine for interrupt OUT/IN transactions

a) Interrupt OUT

```
Unmask (NAK/TXERR/STALL/XFRC/FRMOR)
if (XFRC)
    {
    Reset Error Count
    Mask ACK
    De-allocate Channel
    }
else
    if (STALL or FRMOR)
```

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```
Mask ACK
  Unmask CHH
  Disable Channel
  if (STALL)
    Transfer Done = 1
  }
else
  if (NAK or TXERR)
    Rewind Buffer Pointers
    Reset Error Count
    Mask ACK
    Unmask CHH
    Disable Channel
    }
  else
    if (CHH)
    if (Transfer Done or (Error_count == 3))
       De-allocate Channel
       }
    else
       Re-initialize Channel (in next b interval - 1 Frame)
  else
    if (ACK)
    Reset Error Count
    Mask ACK
```

The application is expected to write the data packets into the transmit FIFO when the space is available in the transmit FIFO and the Request queue up to the count specified in the MCNT field before switching to another channel. The application uses the NPTXFE interrupt in OTG\_HS\_GINTSTS to find the transmit FIFO space.

### b) Interrupt IN

```
Transfer Done = 1
    Unmask CHH
    Disable Channel
  }
else
  if (STALL or FRMOR or NAK or DTERR or BBERR)
    Mask ACK
    Unmask CHH
    Disable Channel
    if (STALL or BBERR)
       Reset Error Count
       Transfer Done = 1
       }
    else
      if (!FRMOR)
      Reset Error Count
else
  if (TXERR)
    {
    Increment Error Count
    Unmask ACK
    Unmask CHH
    Disable Channel
else
  if (CHH)
    Mask CHH
    if (Transfer Done or (Error_count == 3))
      De-allocate Channel
      }
    else
      Re-initialize Channel (in next b interval - 1 /Frame)
  }
else
  if (ACK)
    Reset Error Count
    Mask ACK
```

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}

The application is expected to write the requests for the same channel when the Request queue space is available up to the count specified in the MCNT field before switching to another channel (if any).

# Interrupt IN transactions

The assumptions are:

- The application is attempting to receive one packet (up to 1 maximum packet size) in every frame, starting with odd (transfer size = 1 024 bytes).
- The receive FIFO can hold at least one maximum-packet-size packet and two status DWORDs per packet (1 031 bytes).
- Periodic request queue depth = 4.

### Normal interrupt IN operation

The sequence of operations is as follows:

- a) Initialize channel 2. The application must set the ODDFRM bit in OTG HS HCCHAR2.
- b) Set the CHENA bit in OTG\_HS\_HCCHAR2 to write an IN request to the periodic request queue. For a high-bandwidth interrupt transfer, the application must write the OTG\_HS\_HCCHAR2 register MCNT (maximum number of expected packets in the next frame times) before switching to another channel.
- c) The OTG\_HS host writes an IN request to the periodic request queue for each OTG\_HS\_HCCHAR2 register write with the CHENA bit set.
- d) The OTG\_HS host attempts to send an IN token in the next (odd) frame.
- e) As soon as the IN packet is received and written to the receive FIFO, the OTG\_HS host generates an RXFLVL interrupt.
- f) In response to the RXFLVL interrupt, read the received packet status to determine the number of bytes received, then read the receive FIFO accordingly. The application must mask the RXFLVL interrupt before reading the receive FIFO, and unmask after reading the entire packet.
- g) The core generates the RXFLVL interrupt for the transfer completion status entry in the receive FIFO. The application must read and ignore the receive packet status when the receive packet status is not an IN data packet (PKTSTS in GRXSTSR ≠ 0b0010).
- The core generates an XFRC interrupt as soon as the receive packet status is read.
- i) In response to the XFRC interrupt, read the PKTCNT field in OTG\_HS\_HCTSIZ2. If the PKTCNT bit in OTG\_HS\_HCTSIZ2 is not equal to 0, disable the channel before re-initializing the channel for the next transfer, if any). If PKTCNT bit in

OTG\_HS\_HCTSIZ2 = 0, reinitialize the channel for the next transfer. This time, the application must reset the ODDFRM bit in OTG\_HS\_HCCHAR2.

### Isochronous OUT transactions

A typical isochronous OUT operation in Slave mode is shown in *Figure 386*. The assumptions are:

- The application is attempting to send one packet every frame (up to 1 maximum packet size), starting with an odd frame. (transfer size = 1 024 bytes).
- The periodic transmit FIFO can hold one packet (1 KB).
- Periodic request queue depth = 4.

The sequence of operations is as follows:

- a) Initialize and enable channel 1. The application must set the ODDFRM bit in OTG\_HS\_HCCHAR1.
- b) Write the first packet for channel 1. For a high-bandwidth isochronous transfer, the application must write the subsequent packets up to MCNT (maximum number of packets to be transmitted in the next frame times before switching to another channel.
- c) Along with the last DWORD write of each packet, the OTG\_HS host writes an entry to the periodic request queue.
- d) The OTG\_HS host attempts to send the OUT token in the next frame (odd).
- e) The OTG\_HS host generates the XFRC interrupt as soon as the last packet is transmitted successfully.
- f) In response to the XFRC interrupt, reinitialize the channel for the next transfer.
- g) Handling nonACK responses

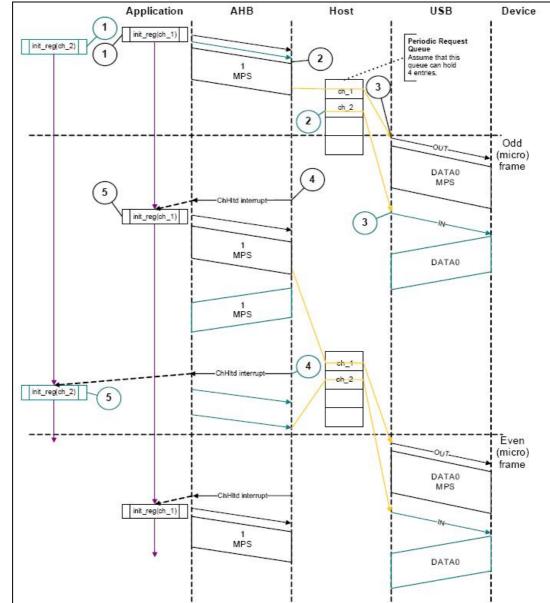


Figure 386. Normal isochronous OUT/IN transactions - DMA mode

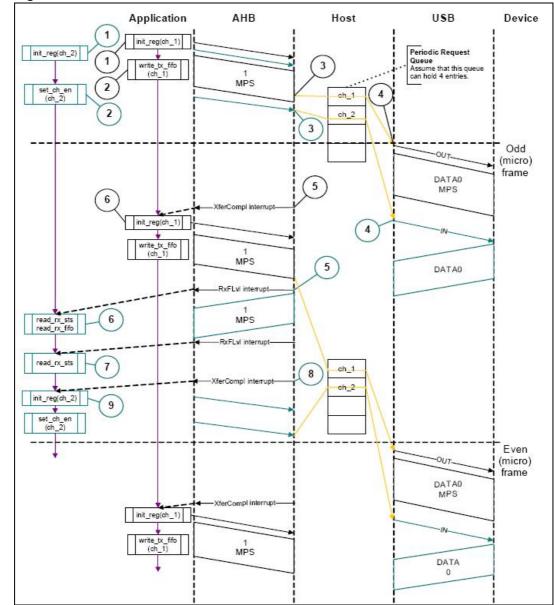


Figure 387. Normal isochronous OUT/IN transactions - Slave mode

# • Interrupt service routine for isochronous OUT/IN transactions

Code sample: Isochronous OUT

```
Unmask (FRMOR/XFRC)
if (XFRC)
   {
   De-allocate Channel
   }
else
   if (FRMOR)
     {
      Unmask CHH
      Disable Channel
   }
```

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```
else
  if (CHH)
    {
    Mask CHH
    De-allocate Channel
   Code sample: Isochronous IN
Unmask (TXERR/XFRC/FRMOR/BBERR)
if (XFRC or FRMOR)
  if (XFRC and (OTG_HS_HCTSIZx.PKTCNT == 0))
    Reset Error Count
    De-allocate Channel
    }
  else
    Unmask CHH
    Disable Channel
  }
  else
    if (TXERR or BBERR)
       Increment Error Count
       Unmask CHH
       Disable Channel
       }
    else
       if (CHH)
         Mask CHH
         if (Transfer Done or (Error_count == 3))
           De-allocate Channel
           }
         else
            Re-initialize Channel
```

#### Isochronous IN transactions

The assumptions are:

- The application is attempting to receive one packet (up to 1 maximum packet size) in every frame starting with the next odd frame (transfer size = 1 024 bytes).
- The receive FIFO can hold at least one maximum-packet-size packet and two status DWORDs per packet (1 031 bytes).
- Periodic request queue depth = 4.

The sequence of operations is as follows:

- a) Initialize channel 2. The application must set the ODDFRM bit in OTG HS HCCHAR2.
- b) Set the CHENA bit in OTG\_HS\_HCCHAR2 to write an IN request to the periodic request queue. For a high-bandwidth isochronous transfer, the application must write the OTG\_HS\_HCCHAR2 register MCNT (maximum number of expected packets in the next frame times) before switching to another channel.
- c) The OTG\_HS host writes an IN request to the periodic request queue for each OTG\_HS\_HCCHAR2 register write with the CHENA bit set.
- d) The OTG\_HS host attempts to send an IN token in the next odd frame.
- e) As soon as the IN packet is received and written to the receive FIFO, the OTG\_HS host generates an RXFLVL interrupt.
- f) In response to the RXFLVL interrupt, read the received packet status to determine the number of bytes received, then read the receive FIFO accordingly. The application must mask the RXFLVL interrupt before reading the receive FIFO, and unmask it after reading the entire packet.
- g) The core generates an RXFLVL interrupt for the transfer completion status entry in the receive FIFO. This time, the application must read and ignore the receive packet status when the receive packet status is not an IN data packet (PKTSTS bit in OTG\_HS\_GRXSTSR ≠ 0b0010).
- h) The core generates an XFRC interrupt as soon as the receive packet status is read.
- i) In response to the XFRC interrupt, read the PKTCNT field in OTG\_HS\_HCTSIZ2. If PKTCNT≠ 0 in OTG\_HS\_HCTSIZ2, disable the channel before re-initializing the channel for the next transfer, if any. If PKTCNT = 0 in OTG\_HS\_HCTSIZ2, reinitialize the channel for the next transfer. This time, the application must reset the ODDFRM bit in OTG\_HS\_HCCHAR2.

# Selecting the queue depth

Choose the periodic and nonperiodic request queue depths carefully to match the number of periodic/nonperiodic endpoints accessed.

The nonperiodic request queue depth affects the performance of nonperiodic transfers. The deeper the queue (along with sufficient FIFO size), the more often the core is able to pipeline nonperiodic transfers. If the queue size is small, the core is able to put in new requests only when the queue space is freed up.

The core's periodic request queue depth is critical to perform periodic transfers as scheduled. Select the periodic queue depth, based on the number of periodic transfers scheduled in a micro-frame. In Slave mode, however, the application must also take into account the disable entry that must be put into the queue. So, if there are two nonhigh-bandwidth periodic endpoints, the periodic request queue depth must be at least 4. If at least one high-bandwidth endpoint is supported, the queue depth must be

8. If the periodic request queue depth is smaller than the periodic transfers scheduled in a micro-frame, a frame overrun condition occurs.

### Handling babble conditions

OTG\_HS controller handles two cases of babble: packet babble and port babble. Packet babble occurs if the device sends more data than the maximum packet size for the channel. Port babble occurs if the core continues to receive data from the device at EOF2 (the end of frame 2, which is very close to SOF).

When OTG\_HS controller detects a packet babble, it stops writing data into the Rx buffer and waits for the end of packet (EOP). When it detects an EOP, it flushes already written data in the Rx buffer and generates a Babble interrupt to the application.

When OTG\_HS controller detects a port babble, it flushes the RxFIFO and disables the port. The core then generates a Port disabled interrupt (HPRTINT in OTG\_HS\_GINTSTS, PENCHNG in OTG\_HS\_HPRT). On receiving this interrupt, the application must determine that this is not due to an overcurrent condition (another cause of the Port Disabled interrupt) by checking POCA in OTG\_HS\_HPRT, then perform a soft reset. The core does not send any more tokens after it has detected a port babble condition.

#### Bulk and control OUT/SETUP transactions in DMA mode

The sequence of operations is as follows:

- a) Initialize and enable channel 1 as explained in Section: Channel initialization.
- b) The HS\_OTG host starts fetching the first packet as soon as the channel is enabled. For internal DMA mode, the OTG\_HS host uses the programmed DMA address to fetch the packet.
- c) After fetching the last DWORD of the second (last) packet, the OTG\_HS host masks channel 1 internally for further arbitration.
- d) The HS\_OTG host generates a CHH interrupt as soon as the last packet is sent.
- In response to the CHH interrupt, de-allocate the channel for other transfers.

## NAK and NYET handling with internal DMA

- a) The OTG\_HS host sends a bulk OUT transaction.
- b) The device responds with NAK or NYET.
- c) If the application has unmasked NAK or NYET, the core generates the corresponding interrupt(s) to the application. The application is not required to

- service these interrupts, since the core takes care of rewinding the buffer pointers and re-initializing the Channel without application intervention.
- d) The core automatically issues a ping token.
- e) When the device returns an ACK, the core continues with the transfer. Optionally, the application can utilize these interrupts, in which case the NAK or NYET interrupt is masked by the application.

The core does not generate a separate interrupt when NAK or NYET is received by the host functionality.

### Bulk and control IN transactions in DMA mode

The sequence of operations is as follows:

- a) Initialize and enable the used channel (channel x) as explained in *Section : Channel initialization*.
- b) The OTG\_HS host writes an IN request to the request queue as soon as the channel receives the grant from the arbiter (arbitration is performed in a round-robin fashion).
- c) The OTG\_HS host starts writing the received data to the system memory as soon as the last byte is received with no errors.
- d) When the last packet is received, the OTG\_HS host sets an internal flag to remove any extra IN requests from the request queue.
- e) The OTG\_HS host flushes the extra requests.
- f) The final request to disable channel x is written to the request queue. At this point, channel 2 is internally masked for further arbitration.
- g) The OTG\_HS host generates the CHH interrupt as soon as the disable request comes to the top of the queue.
- h) In response to the CHH interrupt, de-allocate the channel for other transfers.

### Interrupt OUT transactions in DMA mode

- a) Initialize and enable channel x as explained in Section: Channel initialization.
- b) The OTG\_HS host starts fetching the first packet as soon the channel is enabled and writes the OUT request along with the last DWORD fetch. In high-bandwidth

- transfers, the HS\_OTG host continues fetching the next packet (up to the value specified in the MC field) before switching to the next channel.
- c) The OTG\_HS host attempts to send the OUT token at the beginning of the next odd frame/micro-frame.
- After successfully transmitting the packet, the OTG\_HS host generates a CHH interrupt.
- e) In response to the CHH interrupt, reinitialize the channel for the next transfer.

### Interrupt IN transactions in DMA mode

The sequence of operations (channelx) is as follows:

- a) Initialize and enable channel x as explained in Section: Channel initialization.
- b) The OTG\_HS host writes an IN request to the request queue as soon as the channel x gets the grant from the arbiter (round-robin with fairness). In high-bandwidth transfers, the OTG\_HS host writes consecutive writes up to MC times.
- c) The OTG\_HS host attempts to send an IN token at the beginning of the next (odd) frame/micro-frame.
- d) As soon the packet is received and written to the receive FIFO, the OTG\_HS host generates a CHH interrupt.
- e) In response to the CHH interrupt, reinitialize the channel for the next transfer.

### Isochronous OUT transactions in DMA mode

- a) Initialize and enable channel x as explained in Section: Channel initialization.
- b) The OTG\_HS host starts fetching the first packet as soon as the channel is enabled, and writes the OUT request along with the last DWORD fetch. In highbandwidth transfers, the OTG\_HS host continues fetching the next packet (up to the value specified in the MC field) before switching to the next channel.
- c) The OTG\_HS host attempts to send an OUT token at the beginning of the next (odd) frame/micro-frame.
- d) After successfully transmitting the packet, the HS\_OTG host generates a CHH interrupt.
- e) In response to the CHH interrupt, reinitialize the channel for the next transfer.

### Isochronous IN transactions in DMA mode

The sequence of operations ((channel x) is as follows:

- a) Initialize and enable channel x as explained in Section: Channel initialization.
- b) The OTG\_HS host writes an IN request to the request queue as soon as the channel x gets the grant from the arbiter (round-robin with fairness). In high-

- bandwidth transfers, the OTG\_HS host performs consecutive write operations up to MC times.
- c) The OTG\_HS host attempts to send an IN token at the beginning of the next (odd) frame/micro-frame.
- d) As soon the packet is received and written to the receive FIFO, the OTG\_HS host generates a CHH interrupt.
- In response to the CHH interrupt, reinitialize the channel for the next transfer.

### Bulk and control OUT/SETUP split transactions in DMA mode

The sequence of operations in (channel x) is as follows:

- a) Initialize and enable channel x for start split as explained in *Section : Channel initialization*.
- b) The OTG\_HS host starts fetching the first packet as soon the channel is enabled and writes the OUT request along with the last DWORD fetch.
- After successfully transmitting start split, the OTG\_HS host generates the CHH interrupt.
- d) In response to the CHH interrupt, set the COMPLSPLT bit in HCSPLT1 to send the complete split.
- e) After successfully transmitting complete split, the OTG\_HS host generates the CHH interrupt.
- f) In response to the CHH interrupt, de-allocate the channel.

### Bulk/Control IN split transactions in DMA mode

The sequence of operations (channel x) is as follows:

- a) Initialize and enable channel x as explained in Section: Channel initialization.
- b) The OTG\_HS host writes the start split request to the nonperiodic request after getting the grant from the arbiter. The OTG\_HS host masks the channel x internally for the arbitration after writing the request.
- As soon as the IN token is transmitted, the OTG\_HS host generates the CHH interrupt.
- d) In response to the CHH interrupt, set the COMPLSPLT bit in HCSPLT2 and reenable the channel to send the complete split token. This unmasks channel x for arbitration.
- e) The OTG\_HS host writes the complete split request to the nonperiodic request after receiving the grant from the arbiter.
- f) The OTG\_HS host starts writing the packet to the system memory after receiving the packet successfully.
- g) As soon as the received packet is written to the system memory, the OTG\_HS host generates a CHH interrupt.
- h) In response to the CHH interrupt, de-allocate the channel.

## Interrupt OUT split transactions in DMA mode

The sequence of operations in (channel x) is as follows:

- a) Initialize and enable channel 1 for start split as explained in *Section : Channel initialization*. The application must set the ODDFRM bit in HCCHAR1.
- b) The HS\_OTG host starts reading the packet.
- c) The HS\_OTG host attempts to send the start split transaction.
- d) After successfully transmitting the start split, the OTG\_HS host generates the

CHH interrupt.

- e) In response to the CHH interrupt, set the COMPLSPLT bit in HCSPLT1 to send the complete split.
- After successfully completing the complete split transaction, the OTG\_HS host generates the CHH interrupt.
- g) In response to CHH interrupt, de-allocate the channel.

### • Interrupt IN split transactions in DMA mode

The sequence of operations in (channel x) is as follows:

- Initialize and enable channel x for start split as explained in Section: Channel initialization.
- b) The OTG\_HS host writes an IN request to the request queue as soon as channel x receives the grant from the arbiter.
- c) The OTG\_HS host attempts to send the start split IN token at the beginning of the next odd micro-frame.
- d) The OTG\_HS host generates the CHH interrupt after successfully transmitting the start split IN token.
- e) In response to the CHH interrupt, set the COMPLSPLT bit in HCSPLT2 to send the complete split.
- f) As soon as the packet is received successfully, the OTG\_HS host starts writing the data to the system memory.
- g) The OTG\_HS host generates the CHH interrupt after transferring the received data to the system memory.
- h) In response to the CHH interrupt, de-allocate or reinitialize the channel for the next start split.

# Isochronous OUT split transactions in DMA mode

The sequence of operations (channel x) is as follows:

- a) Initialize and enable channel x for start split (begin) as explained in *Section : Channel initialization*. The application must set the ODDFRM bit in HCCHAR1. Program the MPS field.
- b) The HS\_OTG host starts reading the packet.
- c) After successfully transmitting the start split (begin), the HS\_OTG host generates the CHH interrupt.
- d) In response to the CHH interrupt, reinitialize the registers to send the start split (end).
- e) After successfully transmitting the start split (end), the OTG\_HS host generates a CHH interrupt.
- f) In response to the CHH interrupt, de-allocate the channel.

### Isochronous IN split transactions in DMA mode

The sequence of operations (channel x) is as follows:

- a) Initialize and enable channel x for start split as explained in *Section : Channel initialization*.
- b) The OTG\_HS host writes an IN request to the request queue as soon as channel x receives the grant from the arbiter.
- c) The OTG\_HS host attempts to send the start split IN token at the beginning of the next odd micro-frame.

- The OTG\_HS host generates the CHH interrupt after successfully transmitting the start split IN token.
- e) In response to the CHH interrupt, set the COMPLSPLT bit in HCSPLT2 to send the complete split.
- f) As soon as the packet is received successfully, the OTG\_HS host starts writing the data to the system memory.
- g) The OTG\_HS host generates the CHH interrupt after transferring the received data to the system memory. In response to the CHH interrupt, de-allocate the channel or reinitialize the channel for the next start split.

# 30.13.6 Device programming model

# **Endpoint initialization on USB reset**

- 1. Set the NAK bit for all OUT endpoints
  - SNAK = 1 in OTG\_HS\_DOEPCTLx (for all OUT endpoints)
- 2. Unmask the following interrupt bits
  - INEP0 = 1 in OTG\_HS\_DAINTMSK (control 0 IN endpoint)
  - OUTEP0 = 1 in OTG\_HS\_DAINTMSK (control 0 OUT endpoint)
  - STUP = 1 in DOEPMSK
  - XFRC = 1 in DOEPMSK
  - XFRC = 1 in DIEPMSK
  - TOC = 1 in DIEPMSK
- 3. Set up the Data FIFO RAM for each of the FIFOs
  - Program the OTG\_HS\_GRXFSIZ register, to be able to receive control OUT data and setup data. If thresholding is not enabled, at a minimum, this must be equal to 1 max packet size of control endpoint 0 + 2 DWORDs (for the status of the control OUT data packet) + 10 DWORDs (for setup packets).
  - Program the OTG\_HS\_TX0FSIZ register (depending on the FIFO number chosen) to be able to transmit control IN data. At a minimum, this must be equal to 1 max packet size of control endpoint 0.
- 4. Program the following fields in the endpoint-specific registers for control OUT endpoint 0 to receive a SETUP packet
  - STUPCNT = 3 in OTG\_HS\_DOEPTSIZ0 (to receive up to 3 back-to-back SETUP packets)
- 5. In DMA mode, the DOEPDMA0 register should have a valid memory address to store any SETUP packets received.

At this point, all initialization required to receive SETUP packets is done.

# **Endpoint initialization on enumeration completion**

- 1. On the Enumeration Done interrupt (ENUMDNE in OTG\_HS\_GINTSTS), read the OTG\_HS\_DSTS register to determine the enumeration speed.
- 2. Program the MPSIZ field in OTG\_HS\_DIEPCTL0 to set the maximum packet size. This step configures control endpoint 0. The maximum packet size for a control endpoint depends on the enumeration speed.
- 3. In DMA mode, program the DOEPCTL0 register to enable control OUT endpoint 0, to receive a SETUP packet.
  - EPENA bit in DOEPCTL0 = 1

At this point, the device is ready to receive SOF packets and is configured to perform control transfers on control endpoint 0.

## **Endpoint initialization on SetAddress command**

This section describes what the application must do when it receives a SetAddress command in a SETUP packet.

- 1. Program the OTG\_HS\_DCFG register with the device address received in the SetAddress command
- 1. Program the core to send out a status IN packet

### **Endpoint initialization on SetConfiguration/SetInterface command**

This section describes what the application must do when it receives a SetConfiguration or SetInterface command in a SETUP packet.

- When a SetConfiguration command is received, the application must program the endpoint registers to configure them with the characteristics of the valid endpoints in the new configuration.
- 2. When a SetInterface command is received, the application must program the endpoint registers of the endpoints affected by this command.
- 3. Some endpoints that were active in the prior configuration or alternate setting are not valid in the new configuration or alternate setting. These invalid endpoints must be deactivated.
- 4. Unmask the interrupt for each active endpoint and mask the interrupts for all inactive endpoints in the OTG HS DAINTMSK register.
- 5. Set up the Data FIFO RAM for each FIFO.
- 6. After all required endpoints are configured; the application must program the core to send a status IN packet.

At this point, the device core is configured to receive and transmit any type of data packet.

### **Endpoint activation**

This section describes the steps required to activate a device endpoint or to configure an existing device endpoint to a new type.

- Program the characteristics of the required endpoint into the following fields of the OTG\_HS\_DIEPCTLx register (for IN or bidirectional endpoints) or the OTG\_HS\_DOEPCTLx register (for OUT or bidirectional endpoints).
  - Maximum packet size
  - USB active endpoint = 1
  - Endpoint start data toggle (for interrupt and bulk endpoints)
  - Endpoint type
  - TxFIFO number
- Once the endpoint is activated, the core starts decoding the tokens addressed to that endpoint and sends out a valid handshake for each valid token received for the endpoint.

# **Endpoint deactivation**

This section describes the steps required to deactivate an existing endpoint.

- In the endpoint to be deactivated, clear the USB active endpoint bit in the OTG\_HS\_DIEPCTLx register (for IN or bidirectional endpoints) or the OTG\_HS\_DOEPCTLx register (for OUT or bidirectional endpoints).
- Once the endpoint is deactivated, the core ignores tokens addressed to that endpoint, which results in a timeout on the USB.

Note:

The application must meet the following conditions to set up the device core to handle traffic: NPTXFEM and RXFLVLM in GINTMSK must be cleared.

### 30.13.7 Operational model

### **SETUP and OUT data transfers**

This section describes the internal data flow and application-level operations during data OUT transfers and SETUP transactions.

#### Packet read

This section describes how to read packets (OUT data and SETUP packets) from the receive FIFO in Slave mode.

- 1. On catching an RXFLVL interrupt (OTG\_HS\_GINTSTS register), the application must read the Receive status pop register (OTG\_HS\_GRXSTSP).
- 2. The application can mask the RXFLVL interrupt (in OTG\_HS\_GINTSTS) by writing to RXFLVL = 0 (in GINTMSK), until it has read the packet from the receive FIFO.
- 3. If the received packet's byte count is not 0, the byte count amount of data is popped from the receive Data FIFO and stored in memory. If the received packet byte count is 0, no data is popped from the receive data FIFO.
- 4. The receive FIFO's packet status readout indicates one of the following:
  - a) Global OUT NAK pattern:
     PKTSTS = Global OUT NAK, BCNT = 0x000, EPNUM = Don't Care (0x0),
     DPID = Don't Care (0b00).
     These data indicate that the global OUT NAK bit has taken effect.
  - b) SETUP packet pattern:

    PKTSTS = SETUP. BCNT = 0x008. EPNUM = Control EP Num. DPID = D0.

These data indicate that a SETUP packet for the specified endpoint is now available for reading from the receive FIFO.

- c) Setup stage done pattern:
  - PKTSTS = Setup Stage Done, BCNT = 0x0, EPNUM = Control EP Num, DPID = Don't Care (0b00).
  - These data indicate that the Setup stage for the specified endpoint has completed and the Data stage has started. After this entry is popped from the receive FIFO, the core asserts a Setup interrupt on the specified control OUT endpoint.
- d) Data OUT packet pattern:
   PKTSTS = DataOUT, BCNT = size of the received data OUT packet (0 ≤ BCNT ≤ 1 024), EPNUM = EPNUM on which the packet was received, DPID = Actual Data PID.
- e) Data transfer completed pattern:

  PKTSTS = Data OUT Transfer Done, BCNT = 0x0, EPNUM = OUT EP Num on which the data transfer is complete, DPID = Don't Care (0b00).

  These data indicate that an OUT data transfer for the specified OUT endpoint has completed. After this entry is popped from the receive FIFO, the core asserts a Transfer Completed interrupt on the specified OUT endpoint.
- 5. After the data payload is popped from the receive FIFO, the RXFLVL interrupt (OTG\_HS\_GINTSTS) must be unmasked.
- Steps 1–5 are repeated every time the application detects assertion of the interrupt line due to RXFLVL in OTG\_HS\_GINTSTS. Reading an empty receive FIFO can result in undefined core behavior.

Figure 388 provides a flowchart of the above procedure.

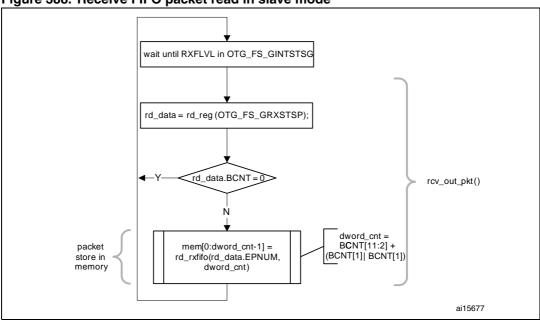
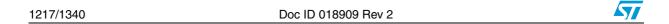


Figure 388. Receive FIFO packet read in slave mode

# SETUP transactions

This section describes how the core handles SETUP packets and the application's sequence for handling SETUP transactions.

Application requirements



- 1. To receive a SETUP packet, the STUPCNT field (OTG\_HS\_DOEPTSIZx) in a control OUT endpoint must be programmed to a nonzero value. When the application programs the STUPCNT field to a nonzero value, the core receives SETUP packets and writes them to the receive FIFO, irrespective of the NAK status and EPENA bit setting in OTG\_HS\_DOEPCTLx. The STUPCNT field is decremented every time the control endpoint receives a SETUP packet. If the STUPCNT field is not programmed to a proper value before receiving a SETUP packet, the core still receives the SETUP packet and decrements the STUPCNT field, but the application may not be able to determine the correct number of SETUP packets received in the Setup stage of a control transfer.
  - STUPCNT = 3 in OTG\_HS\_DOEPTSIZx
- 2. The application must always allocate some extra space in the Receive data FIFO, to be able to receive up to three SETUP packets on a control endpoint.
  - The space to be reserved is 10 DWORDs. Three DWORDs are required for the first SETUP packet, 1 DWORD is required for the Setup stage done DWORD and 6 DWORDs are required to store two extra SETUP packets among all control endpoints.
  - 3 DWORDs per SETUP packet are required to store 8 bytes of SETUP data and 4 bytes of SETUP status (Setup packet pattern). The core reserves this space in the receive data.
  - FIFO to write SETUP data only, and never uses this space for data packets.
- 3. The application must read the 2 DWORDs of the SETUP packet from the receive FIFO.
- 4. The application must read and discard the Setup stage done DWORD from the receive FIFO.

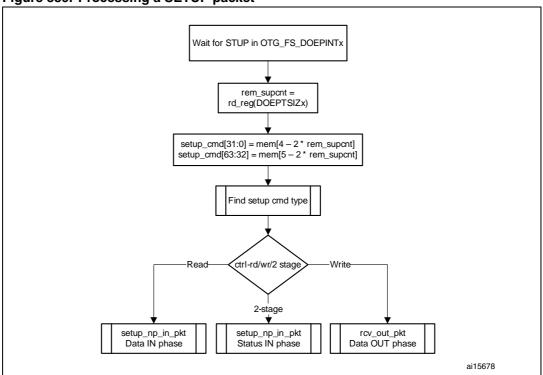
#### Internal data flow

- 5. When a SETUP packet is received, the core writes the received data to the receive FIFO, without checking for available space in the receive FIFO and irrespective of the endpoint's NAK and STALL bit settings.
  - The core internally sets the IN NAK and OUT NAK bits for the control IN/OUT endpoints on which the SETUP packet was received.
- 6. For every SETUP packet received on the USB, 3 DWORDs of data are written to the receive FIFO, and the STUPCNT field is decremented by 1.
  - The first DWORD contains control information used internally by the core
  - The second DWORD contains the first 4 bytes of the SETUP command
  - The third DWORD contains the last 4 bytes of the SETUP command
- When the Setup stage changes to a Data IN/OUT stage, the core writes an entry (Setup stage done DWORD) to the receive FIFO, indicating the completion of the Setup stage.
- 8. On the AHB side, SETUP packets are emptied by the application.
- When the application pops the Setup stage done DWORD from the receive FIFO, the core interrupts the application with an STUP interrupt (OTG\_HS\_DOEPINTx), indicating it can process the received SETUP packet.
  - The core clears the endpoint enable bit for control OUT endpoints.

### Application programming sequence

- 1. Program the OTG\_HS\_DOEPTSIZx register.
  - STUPCNT = 3
- 2. Wait for the RXFLVL interrupt (OTG\_HS\_GINTSTS) and empty the data packets from the receive FIFO.
- 3. Assertion of the STUP interrupt (OTG\_HS\_DOEPINTx) marks a successful completion of the SETUP Data Transfer.
  - On this interrupt, the application must read the OTG\_HS\_DOEPTSIZx register to determine the number of SETUP packets received and process the last received SETUP packet.

Figure 389. Processing a SETUP packet



### Handling more than three back-to-back SETUP packets

Per the USB 2.0 specification, normally, during a SETUP packet error, a host does not send more than three back-to-back SETUP packets to the same endpoint. However, the USB 2.0 specification does not limit the number of back-to-back SETUP packets a host can send to the same endpoint. When this condition occurs, the OTG\_HS controller generates an interrupt (B2BSTUP in OTG\_HS\_DOEPINTx).

## Setting the global OUT NAK

Internal data flow:

- When the application sets the Global OUT NAK (SGONAK bit in OTG\_HS\_DCTL), the
  core stops writing data, except SETUP packets, to the receive FIFO. Irrespective of the
  space availability in the receive FIFO, nonisochronous OUT tokens receive a NAK
  handshake response, and the core ignores isochronous OUT data packets
- 2. The core writes the Global OUT NAK pattern to the receive FIFO. The application must reserve enough receive FIFO space to write this data pattern.

- 3. When the application pops the Global OUT NAK pattern DWORD from the receive FIFO, the core sets the GONAKEFF interrupt (OTG\_HS\_GINTSTS).
- Once the application detects this interrupt, it can assume that the core is in Global OUT NAK mode. The application can clear this interrupt by clearing the SGONAK bit in OTG HS DCTL.

### Application programming sequence

- To stop receiving any kind of data in the receive FIFO, the application must set the Global OUT NAK bit by programming the following field:
  - SGONAK = 1 in OTG\_HS\_DCTL
- Wait for the assertion of the GONAKEFF interrupt in OTG\_HS\_GINTSTS. When asserted, this interrupt indicates that the core has stopped receiving any type of data except SETUP packets.
- The application can receive valid OUT packets after it has set SGONAK in OTG\_HS\_DCTL and before the core asserts the GONAKEFF interrupt (OTG\_HS\_GINTSTS).
- 4. The application can temporarily mask this interrupt by writing to the GINAKEFFM bit in GINTMSK.
  - GINAKEFFM = 0 in GINTMSK
- Whenever the application is ready to exit the Global OUT NAK mode, it must clear the SGONAK bit in OTG\_HS\_DCTL. This also clears the GONAKEFF interrupt (OTG\_HS\_GINTSTS).
  - OTG HS DCTL = 1 in CGONAK
- 6. If the application has masked this interrupt earlier, it must be unmasked as follows:
  - GINAKEFFM = 1 in GINTMSK

### Disabling an OUT endpoint

The application must use this sequence to disable an OUT endpoint that it has enabled.

Application programming sequence:

- Before disabling any OUT endpoint, the application must enable Global OUT NAK mode in the core.
  - SGONAK = 1 in OTG\_HS\_DCTL
- 2. Wait for the GONAKEFF interrupt (OTG\_HS\_GINTSTS)
- 3. Disable the required OUT endpoint by programming the following fields:
  - EPDIS = 1 in OTG\_HS\_DOEPCTLx
  - SNAK = 1 in OTG\_HS\_DOEPCTLx
- 4. Wait for the EPDISD interrupt (OTG\_HS\_DOEPINTx), which indicates that the OUT endpoint is completely disabled. When the EPDISD interrupt is asserted, the core also clears the following bits:
  - EPDIS = 0 in OTG\_HS\_DOEPCTLx
  - EPENA = 0 in OTG\_HS\_DOEPCTLx
- The application must clear the Global OUT NAK bit to start receiving data from other nondisabled OUT endpoints.
  - SGONAK = 0 in OTG HS DCTL
- Generic nonisochronous OUT data transfers

This section describes a regular nonisochronous OUT data transfer (control, bulk, or interrupt).

### Application requirements:

- 1. Before setting up an OUT transfer, the application must allocate a buffer in the memory to accommodate all data to be received as part of the OUT transfer.
- For OUT transfers, the transfer size field in the endpoint's transfer size register must be a multiple of the maximum packet size of the endpoint, adjusted to the DWORD boundary.
  - transfer size[EPNUM] =  $n \times (MPSIZ[EPNUM] + 4 (MPSIZ[EPNUM] \mod 4)$
  - packet count[EPNUM] = n
  - -n>0
- 3. On any OUT endpoint interrupt, the application must read the endpoint's transfer size register to calculate the size of the payload in the memory. The received payload size can be less than the programmed transfer size.
  - Payload size in memory = application programmed initial transfer size core updated final transfer size
  - Number of USB packets in which this payload was received = application programmed initial packet count – core updated final packet count

#### Internal data flow:

- 1. The application must set the transfer size and packet count fields in the endpointspecific registers, clear the NAK bit, and enable the endpoint to receive the data.
- 2. Once the NAK bit is cleared, the core starts receiving data and writes it to the receive FIFO, as long as there is space in the receive FIFO. For every data packet received on the USB, the data packet and its status are written to the receive FIFO. Every packet (maximum packet size or short packet) written to the receive FIFO decrements the packet count field for that endpoint by 1.
  - OUT data packets received with bad data CRC are flushed from the receive FIFO automatically.
  - After sending an ACK for the packet on the USB, the core discards nonisochronous OUT data packets that the host, which cannot detect the ACK, resends. The application does not detect multiple back-to-back data OUT packets on the same endpoint with the same data PID. In this case the packet count is not decremented.
  - If there is no space in the receive FIFO, isochronous or nonisochronous data packets are ignored and not written to the receive FIFO. Additionally, nonisochronous OUT tokens receive a NAK handshake reply.
  - In all the above three cases, the packet count is not decremented because no data are written to the receive FIFO.
- 3. When the packet count becomes 0 or when a short packet is received on the endpoint, the NAK bit for that endpoint is set. Once the NAK bit is set, the isochronous or nonisochronous data packets are ignored and not written to the receive FIFO, and nonisochronous OUT tokens receive a NAK handshake reply.
- 4. After the data are written to the receive FIFO, the application reads the data from the receive FIFO and writes it to external memory, one packet at a time per endpoint.
- 5. At the end of every packet write on the AHB to external memory, the transfer size for the endpoint is decremented by the size of the written packet.

- 6. The OUT data transfer completed pattern for an OUT endpoint is written to the receive FIFO on one of the following conditions:
  - The transfer size is 0 and the packet count is 0
  - The last OUT data packet written to the receive FIFO is a short packet (0 ≤ packet size < maximum packet size)</li>
- 7. When either the application pops this entry (OUT data transfer completed), a transfer completed interrupt is generated for the endpoint and the endpoint enable is cleared.

### Application programming sequence:

- 1. Program the OTG\_HS\_DOEPTSIZx register for the transfer size and the corresponding packet count.
- 2. Program the OTG\_HS\_DOEPCTLx register with the endpoint characteristics, and set the EPENA and CNAK bits.
  - EPENA = 1 in OTG\_HS\_DOEPCTLx
  - CNAK = 1 in OTG\_HS\_DOEPCTLx
- Wait for the RXFLVL interrupt (in OTG\_HS\_GINTSTS) and empty the data packets from the receive FIFO.
  - This step can be repeated many times, depending on the transfer size.
- 4. Asserting the XFRC interrupt (OTG\_HS\_DOEPINTx) marks a successful completion of the nonisochronous OUT data transfer.
- 5. Read the OTG\_HS\_DOEPTSIZx register to determine the size of the received data payload.

#### Generic isochronous OUT data transfer

This section describes a regular isochronous OUT data transfer.

### Application requirements:

- 1. All the application requirements for nonisochronous OUT data transfers also apply to isochronous OUT data transfers.
- 2. For isochronous OUT data transfers, the transfer size and packet count fields must always be set to the number of maximum-packet-size packets that can be received in a single frame and no more. Isochronous OUT data transfers cannot span more than 1 frame.
- The application must read all isochronous OUT data packets from the receive FIFO (data and status) before the end of the periodic frame (EOPF interrupt in OTG\_HS\_GINTSTS).
- 4. To receive data in the following frame, an isochronous OUT endpoint must be enabled after the EOPF (OTG HS GINTSTS) and before the SOF (OTG HS GINTSTS).

### Internal data flow:

- 1. The internal data flow for isochronous OUT endpoints is the same as that for nonisochronous OUT endpoints, but for a few differences.
- 2. When an isochronous OUT endpoint is enabled by setting the Endpoint Enable and clearing the NAK bits, the Even/Odd frame bit must also be set appropriately. The core receives data on an isochronous OUT endpoint in a particular frame only if the following condition is met:
  - EONUM (in OTG\_HS\_DOEPCTLx) = SOFFN[0] (in OTG\_HS\_DSTS)
- 3. When the application completely reads an isochronous OUT data packet (data and status) from the receive FIFO, the core updates the RXDPID field in

OTG\_HS\_DOEPTSIZx with the data PID of the last isochronous OUT data packet read from the receive FIFO.

Application programming sequence:

- 1. Program the OTG\_HS\_DOEPTSIZx register for the transfer size and the corresponding packet count
- 2. Program the OTG\_HS\_DOEPCTLx register with the endpoint characteristics and set the Endpoint Enable, ClearNAK, and Even/Odd frame bits.
  - EPENA = 1
  - CNAK = 1
  - EONUM = (0: Even/1: Odd)
- In Slave mode, wait for the RXFLVL interrupt (in OTG\_HS\_GINTSTS) and empty the data packets from the receive FIFO
  - This step can be repeated many times, depending on the transfer size.
- 4. The assertion of the XFRC interrupt (in OTG\_HS\_DOEPINTx) marks the completion of the isochronous OUT data transfer. This interrupt does not necessarily mean that the data in memory are good.
- 5. This interrupt cannot always be detected for isochronous OUT transfers. Instead, the application can detect the IISOOXFRM interrupt in OTG\_HS\_GINTSTS.
- 6. Read the OTG\_HS\_DOEPTSIZx register to determine the size of the received transfer and to determine the validity of the data received in the frame. The application must treat the data received in memory as valid only if one of the following conditions is met:
  - RXDPID = D0 (in OTG\_HS\_DOEPTSIZx) and the number of USB packets in which this payload was received = 1
  - RXDPID = D1 (in OTG\_HS\_DOEPTSIZx) and the number of USB packets in which this payload was received = 2
  - RXDPID = D2 (in OTG\_HS\_DOEPTSIZx) and the number of USB packets in which this payload was received = 3

The number of USB packets in which this payload was received = Application programmed initial packet count – Core updated final packet count

The application can discard invalid data packets.

### Incomplete isochronous OUT data transfers

This section describes the application programming sequence when isochronous OUT data packets are dropped inside the core.

### Internal data flow:

- For isochronous OUT endpoints, the XFRC interrupt (in OTG\_HS\_DOEPINTx) may not always be asserted. If the core drops isochronous OUT data packets, the application could fail to detect the XFRC interrupt (OTG\_HS\_DOEPINTx) under the following circumstances:
  - When the receive FIFO cannot accommodate the complete ISO OUT data packet, the core drops the received ISO OUT data
  - When the isochronous OUT data packet is received with CRC errors
  - When the isochronous OUT token received by the core is corrupted
  - When the application is very slow in reading the data from the receive FIFO
- 2. When the core detects an end of periodic frame before transfer completion to all isochronous OUT endpoints, it asserts the incomplete Isochronous OUT data interrupt

(IISOOXFRM in OTG\_HS\_GINTSTS), indicating that an XFRC interrupt (in OTG\_HS\_DOEPINTx) is not asserted on at least one of the isochronous OUT endpoints. At this point, the endpoint with the incomplete transfer remains enabled, but no active transfers remain in progress on this endpoint on the USB.

Application programming sequence:

- 1. Asserting the IISOOXFRM interrupt (OTG\_HS\_GINTSTS) indicates that in the current frame, at least one isochronous OUT endpoint has an incomplete transfer.
- 2. If this occurs because isochronous OUT data is not completely emptied from the endpoint, the application must ensure that the application empties all isochronous OUT data (data and status) from the receive FIFO before proceeding.
  - When all data are emptied from the receive FIFO, the application can detect the XFRC interrupt (OTG\_HS\_DOEPINTx). In this case, the application must reenable the endpoint to receive isochronous OUT data in the next frame.
- 3. When it receives an IISOOXFRM interrupt (in OTG\_HS\_GINTSTS), the application must read the control registers of all isochronous OUT endpoints (OTG\_HS\_DOEPCTLx) to determine which endpoints had an incomplete transfer in the current micro-frame. An endpoint transfer is incomplete if both the following conditions are met:
  - EONUM bit (in OTG\_HS\_DOEPCTLx) = SOFFN[0] (in OTG\_HS\_DSTS)
  - EPENA = 1 (in OTG HS DOEPCTLx)
- 4. The previous step must be performed before the SOF interrupt (in OTG\_HS\_GINTSTS) is detected, to ensure that the current frame number is not changed.
- For isochronous OUT endpoints with incomplete transfers, the application must discard
  the data in the memory and disable the endpoint by setting the EPDIS bit in
  OTG HS DOEPCTLx.
- Wait for the EPDIS interrupt (in OTG\_HS\_DOEPINTx) and enable the endpoint to receive new data in the next frame.
  - Because the core can take some time to disable the endpoint, the application may not be able to receive the data in the next frame after receiving bad isochronous data.

### Stalling a nonisochronous OUT endpoint

This section describes how the application can stall a nonisochronous endpoint.

- 1. Put the core in the Global OUT NAK mode.
- 2. Disable the required endpoint
  - When disabling the endpoint, instead of setting the SNAK bit in OTG\_HS\_DOEPCTL, set STALL = 1 (in OTG\_HS\_DOEPCTL).

The STALL bit always takes precedence over the NAK bit.

- 3. When the application is ready to end the STALL handshake for the endpoint, the STALL bit (in OTG HS DOEPCTLx) must be cleared.
- 4. If the application is setting or clearing a STALL for an endpoint due to a SetFeature.Endpoint Halt or ClearFeature.Endpoint Halt command, the STALL bit must be set or cleared before the application sets up the Status stage transfer on the control endpoint.

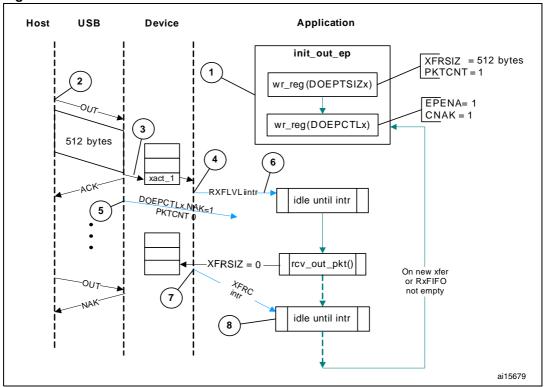
#### **Examples**

This section describes and depicts some fundamental transfer types and scenarios.

Slave mode bulk OUT transaction

*Figure 390* depicts the reception of a single Bulk OUT Data packet from the USB to the AHB and describes the events involved in the process.

Figure 390. Slave mode bulk OUT transaction



After a SetConfiguration/SetInterface command, the application initializes all OUT endpoints by setting CNAK = 1 and EPENA = 1 (in OTG\_HS\_DOEPCTLx), and setting a suitable XFRSIZ and PKTCNT in the OTG\_HS\_DOEPTSIZx register.

- 1. Host attempts to send data (OUT token) to an endpoint.
- 2. When the core receives the OUT token on the USB, it stores the packet in the RxFIFO because space is available there.
- 3. After writing the complete packet in the RxFIFO, the core then asserts the RXFLVL interrupt (in OTG\_HS\_GINTSTS).
- 4. On receiving the PKTCNT number of USB packets, the core internally sets the NAK bit for this endpoint to prevent it from receiving any more packets.
- 5. The application processes the interrupt and reads the data from the RxFIFO.
- 6. When the application has read all the data (equivalent to XFRSIZ), the core generates an XFRC interrupt (in OTG\_HS\_DOEPINTx).
- 7. The application processes the interrupt and uses the setting of the XFRC interrupt bit (in OTG\_HS\_DOEPINTx) to determine that the intended transfer is complete.

#### IN data transfers

#### Packet write

This section describes how the application writes data packets to the endpoint FIFO in Slave mode when dedicated transmit FIFOs are enabled.

- 1. The application can either choose the polling or the interrupt mode.
  - In polling mode, the application monitors the status of the endpoint transmit data FIFO by reading the OTG\_HS\_DTXFSTSx register, to determine if there is enough space in the data FIFO.
  - In interrupt mode, the application waits for the TXFE interrupt (in OTG\_HS\_DIEPINTx) and then reads the OTG\_HS\_DTXFSTSx register, to determine if there is enough space in the data FIFO.
  - To write a single nonzero length data packet, there must be space to write the entire packet in the data FIFO.
  - To write zero length packet, the application must not look at the FIFO space.
- 2. Using one of the above mentioned methods, when the application determines that there is enough space to write a transmit packet, the application must first write into the endpoint control register, before writing the data into the data FIFO. Typically, the application, must do a read modify write on the OTG\_HS\_DIEPCTLx register to avoid modifying the contents of the register, except for setting the Endpoint Enable bit.

The application can write multiple packets for the same endpoint into the transmit FIFO, if space is available. For periodic IN endpoints, the application must write packets only for one micro-frame. It can write packets for the next periodic transaction only after getting transfer complete for the previous transaction.

#### Setting IN endpoint NAK

Internal data flow:

- 1. When the application sets the IN NAK for a particular endpoint, the core stops transmitting data on the endpoint, irrespective of data availability in the endpoint's transmit FIFO.
- 2. Nonisochronous IN tokens receive a NAK handshake reply
  - Isochronous IN tokens receive a zero-data-length packet reply
- The core asserts the INEPNE (IN endpoint NAK effective) interrupt in OTG HS DIEPINTx in response to the SNAK bit in OTG HS DIEPCTLx.
- Once this interrupt is seen by the application, the application can assume that the endpoint is in IN NAK mode. This interrupt can be cleared by the application by setting the CNAK bit in OTG HS DIEPCTLx.

Application programming sequence:

- 1. To stop transmitting any data on a particular IN endpoint, the application must set the IN NAK bit. To set this bit, the following field must be programmed.
  - SNAK = 1 in OTG HS DIEPCTLx
- 2. Wait for assertion of the INEPNE interrupt in OTG\_HS\_DIEPINTx. This interrupt indicates that the core has stopped transmitting data on the endpoint.
- 3. The core can transmit valid IN data on the endpoint after the application has set the NAK bit, but before the assertion of the NAK Effective interrupt.
- 4. The application can mask this interrupt temporarily by writing to the INEPNEM bit in DIEPMSK.
  - INEPNEM = 0 in DIEPMSK
- 5. To exit Endpoint NAK mode, the application must clear the NAK status bit (NAKSTS) in OTG\_HS\_DIEPCTLx. This also clears the INEPNE interrupt (in OTG\_HS\_DIEPINTx).
  - CNAK = 1 in OTG HS DIEPCTLx
- 6. If the application masked this interrupt earlier, it must be unmasked as follows:
  - INEPNEM = 1 in DIEPMSK

#### IN endpoint disable

Use the following sequence to disable a specific IN endpoint that has been previously enabled.

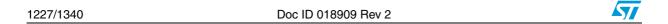
Application programming sequence:

- 1. The application must stop writing data on the AHB for the IN endpoint to be disabled.
- 2. The application must set the endpoint in NAK mode.
  - SNAK = 1 in OTG\_HS\_DIEPCTLx
- 3. Wait for the INEPNE interrupt in OTG\_HS\_DIEPINTx.
- Set the following bits in the OTG\_HS\_DIEPCTLx register for the endpoint that must be disabled.
  - EPDIS = 1 in OTG\_HS\_DIEPCTLx
  - SNAK = 1 in OTG\_HS\_DIEPCTLx
- 5. Assertion of the EPDISD interrupt in OTG\_HS\_DIEPINTx indicates that the core has completely disabled the specified endpoint. Along with the assertion of the interrupt, the core also clears the following bits:
  - EPENA = 0 in OTG HS DIEPCTLx
  - EPDIS = 0 in OTG\_HS\_DIEPCTLx
- 6. The application must read the OTG\_HS\_DIEPTSIZx register for the periodic IN EP, to calculate how much data on the endpoint were transmitted on the USB.
- 7. The application must flush the data in the Endpoint transmit FIFO, by setting the following fields in the OTG\_HS\_GRSTCTL register:
  - TXFNUM (in OTG HS GRSTCTL) = Endpoint transmit FIFO number
  - TXFFLSH in (OTG\_HS\_GRSTCTL) = 1

The application must poll the OTG\_HS\_GRSTCTL register, until the TXFFLSH bit is cleared by the core, which indicates the end of flush operation. To transmit new data on this endpoint, the application can re-enable the endpoint at a later point.

#### Generic nonperiodic IN data transfers

Application requirements:



- 1. Before setting up an IN transfer, the application must ensure that all data to be transmitted as part of the IN transfer are part of a single buffer.
- 2. For IN transfers, the Transfer Size field in the Endpoint Transfer Size register denotes a payload that constitutes multiple maximum-packet-size packets and a single short packet. This short packet is transmitted at the end of the transfer.
  - To transmit a few maximum-packet-size packets and a short packet at the end of the transfer:

```
Transfer size[EPNUM] = x \times MPSIZ[EPNUM] + sp
If (sp > 0), then packet count[EPNUM] = x + 1.
Otherwise, packet count[EPNUM] = x
```

To transmit a single zero-length data packet:

```
Transfer size[EPNUM] = 0
Packet count[EPNUM] = 1
```

 To transmit a few maximum-packet-size packets and a zero-length data packet at the end of the transfer, the application must split the transfer into two parts. The first sends maximum-packet-size data packets and the second sends the zerolength data packet alone.

```
First transfer: transfer size[EPNUM] = x \times MPSIZ[epnum]; packet count = n;
Second transfer: transfer size[EPNUM] = 0; packet count = 1;
```

- 3. Once an endpoint is enabled for data transfers, the core updates the Transfer size register. At the end of the IN transfer, the application must read the Transfer size register to determine how much data posted in the transmit FIFO have already been sent on the USB.
- 4. Data fetched into transmit FIFO = Application-programmed initial transfer size coreupdated final transfer size
  - Data transmitted on USB = (application-programmed initial packet count Core updated final packet count) × MPSIZ[EPNUM]
  - Data yet to be transmitted on USB = (Application-programmed initial transfer size
     data transmitted on USB)

#### Internal data flow:

- 1. The application must set the transfer size and packet count fields in the endpointspecific registers and enable the endpoint to transmit the data.
- 2. The application must also write the required data to the transmit FIFO for the endpoint.
- 3. Every time a packet is written into the transmit FIFO by the application, the transfer size for that endpoint is decremented by the packet size. The data is fetched from the memory by the application, until the transfer size for the endpoint becomes 0. After writing the data into the FIFO, the "number of packets in FIFO" count is incremented (this is a 3-bit count, internally maintained by the core for each IN endpoint transmit FIFO. The maximum number of packets maintained by the core at any time in an IN endpoint FIFO is eight). For zero-length packets, a separate flag is set for each FIFO, without any data in the FIFO.
- 4. Once the data are written to the transmit FIFO, the core reads them out upon receiving an IN token. For every nonisochronous IN data packet transmitted with an ACK

- handshake, the packet count for the endpoint is decremented by one, until the packet count is zero. The packet count is not decremented on a timeout.
- 5. For zero length packets (indicated by an internal zero length flag), the core sends out a zero-length packet for the IN token and decrements the packet count field.
- 6. If there are no data in the FIFO for a received IN token and the packet count field for that endpoint is zero, the core generates an "IN token received when TxFIFO is empty" (ITTXFE) Interrupt for the endpoint, provided that the endpoint NAK bit is not set. The core responds with a NAK handshake for nonisochronous endpoints on the USB.
- 7. The core internally rewinds the FIFO pointers and no timeout interrupt is generated.
- 8. When the transfer size is 0 and the packet count is 0, the transfer complete (XFRC) interrupt for the endpoint is generated and the endpoint enable is cleared.

Application programming sequence:

- 1. Program the OTG\_HS\_DIEPTSIZx register with the transfer size and corresponding packet count.
- 2. Program the OTG\_HS\_DIEPCTLx register with the endpoint characteristics and set the CNAK and EPENA (Endpoint Enable) bits.
- 3. When transmitting nonzero length data packet, the application must poll the OTG\_HS\_DTXFSTSx register (where x is the FIFO number associated with that endpoint) to determine whether there is enough space in the data FIFO. The application can optionally use TXFE (in OTG\_HS\_DIEPINTx) before writing the data.
- Generic periodic IN data transfers

This section describes a typical periodic IN data transfer.

Application requirements:

- 1. Application requirements 1, 2, 3, and 4 of *Generic nonperiodic IN data transfers on page 1227* also apply to periodic IN data transfers, except for a slight modification of requirement 2.
  - The application can only transmit multiples of maximum-packet-size data packets or multiples of maximum-packet-size packets, plus a short packet at the end. To



transmit a few maximum-packet-size packets and a short packet at the end of the transfer, the following conditions must be met:

```
transfer size[EPNUM] = x \times MPSIZ[EPNUM] + sp
(where x is an integer \geq 0, and 0 \leq sp < MPSIZ[EPNUM])
If (sp > 0), packet count[EPNUM] = x + 1
Otherwise, packet count[EPNUM] = x;
MCNT[EPNUM] = packet count[EPNUM]
```

- The application cannot transmit a zero-length data packet at the end of a transfer.
   It can transmit a single zero-length data packet by itself. To transmit a single zero-length data packet:
- transfer size[EPNUM] = 0packet count[EPNUM] = 1MCNT[EPNUM] = packet count[EPNUM]
- 2. The application can only schedule data transfers one frame at a time.
  - $(MCNT 1) \times MPSIZ ≤ XFERSIZ ≤ MCNT \times MPSIZ$
  - PKTCNT = MCNT (in OTG\_HS\_DIEPTSIZx)
  - If XFERSIZ < MCNT × MPSIZ, the last data packet of the transfer is a short packet.
  - Note that: MCNT is in OTG\_HS\_DIEPTSIZx, MPSIZ is in OTG\_HS\_DIEPCTLx, PKTCNT is in OTG\_HS\_DIEPTSIZx and XFERSIZ is in OTG\_HS\_DIEPTSIZx
- 3. The complete data to be transmitted in the frame must be written into the transmit FIFO by the application, before the IN token is received. Even when 1 DWORD of the data to be transmitted per frame is missing in the transmit FIFO when the IN token is received, the core behaves as when the FIFO is empty. When the transmit FIFO is empty:
  - A zero data length packet would be transmitted on the USB for isochronous IN endpoints
  - A NAK handshake would be transmitted on the USB for interrupt IN endpoints
- 4. For a high-bandwidth IN endpoint with three packets in a frame, the application can program the endpoint FIFO size to be 2 × max\_pkt\_size and have the third packet loaded in after the first packet has been transmitted on the USB.

### Internal data flow:

- 1. The application must set the transfer size and packet count fields in the endpointspecific registers and enable the endpoint to transmit the data.
- 2. The application must also write the required data to the associated transmit FIFO for the endpoint.
- 3. Every time the application writes a packet to the transmit FIFO, the transfer size for that endpoint is decremented by the packet size. The data are fetched from application memory until the transfer size for the endpoint becomes 0.
- 4. When an IN token is received for a periodic endpoint, the core transmits the data in the FIFO, if available. If the complete data payload (complete packet, in dedicated FIFO mode) for the frame is not present in the FIFO, then the core generates an IN token received when TxFIFO empty interrupt for the endpoint.
  - A zero-length data packet is transmitted on the USB for isochronous IN endpoints
  - A NAK handshake is transmitted on the USB for interrupt IN endpoints
- 5. The packet count for the endpoint is decremented by 1 under the following conditions:

- For isochronous endpoints, when a zero- or nonzero-length data packet is transmitted
- For interrupt endpoints, when an ACK handshake is transmitted
- When the transfer size and packet count are both 0, the transfer completed interrupt for the endpoint is generated and the endpoint enable is cleared.
- At the "Periodic frame Interval" (controlled by PFIVL in OTG\_HS\_DCFG), when the
  core finds nonempty any of the isochronous IN endpoint FIFOs scheduled for the
  current frame nonempty, the core generates an IISOIXFR interrupt in
  OTG\_HS\_GINTSTS.

#### Application programming sequence:

- Program the OTG\_HS\_DIEPCTLx register with the endpoint characteristics and set the CNAK and EPENA bits.
- 2. Write the data to be transmitted in the next frame to the transmit FIFO.
- 3. Asserting the ITTXFE interrupt (in OTG\_HS\_DIEPINTx) indicates that the application has not yet written all data to be transmitted to the transmit FIFO.
- 4. If the interrupt endpoint is already enabled when this interrupt is detected, ignore the interrupt. If it is not enabled, enable the endpoint so that the data can be transmitted on the next IN token attempt.
- Asserting the XFRC interrupt (in OTG\_HS\_DIEPINTx) with no ITTXFE interrupt in OTG\_HS\_DIEPINTx indicates the successful completion of an isochronous IN transfer. A read to the OTG\_HS\_DIEPTSIZx register must give transfer size = 0 and packet count = 0, indicating all data were transmitted on the USB.
- 6. Asserting the XFRC interrupt (in OTG\_HS\_DIEPINTx), with or without the ITTXFE interrupt (in OTG\_HS\_DIEPINTx), indicates the successful completion of an interrupt IN transfer. A read to the OTG\_HS\_DIEPTSIZx register must give transfer size = 0 and packet count = 0, indicating all data were transmitted on the USB.
- 7. Asserting the incomplete isochronous IN transfer (IISOIXFR) interrupt in OTG\_HS\_GINTSTS with none of the aforementioned interrupts indicates the core did not receive at least 1 periodic IN token in the current frame.

#### Incomplete isochronous IN data transfers

This section describes what the application must do on an incomplete isochronous IN data transfer.

#### Internal data flow:

- 1. An isochronous IN transfer is treated as incomplete in one of the following conditions:
  - a) The core receives a corrupted isochronous IN token on at least one isochronous IN endpoint. In this case, the application detects an incomplete isochronous IN transfer interrupt (IISOIXFR in OTG\_HS\_GINTSTS).
  - b) The application is slow to write the complete data payload to the transmit FIFO and an IN token is received before the complete data payload is written to the FIFO. In this case, the application detects an IN token received when TxFIFO empty interrupt in OTG\_HS\_DIEPINTx. The application can ignore this interrupt, as it eventually results in an incomplete isochronous IN transfer interrupt (IISOIXFR in OTG\_HS\_GINTSTS) at the end of periodic frame.

The core transmits a zero-length data packet on the USB in response to the received IN token.

- 2. The application must stop writing the data payload to the transmit FIFO as soon as possible.
- 3. The application must set the NAK bit and the disable bit for the endpoint.
- 4. The core disables the endpoint, clears the disable bit, and asserts the Endpoint Disable interrupt for the endpoint.

#### Application programming sequence

- 1. The application can ignore the IN token received when TxFIFO empty interrupt in OTG\_HS\_DIEPINTx on any isochronous IN endpoint, as it eventually results in an incomplete isochronous IN transfer interrupt (in OTG\_HS\_GINTSTS).
- Assertion of the incomplete isochronous IN transfer interrupt (in OTG\_HS\_GINTSTS)
  indicates an incomplete isochronous IN transfer on at least one of the isochronous IN
  endpoints.
- 3. The application must read the Endpoint Control register for all isochronous IN endpoints to detect endpoints with incomplete IN data transfers.
- 4. The application must stop writing data to the Periodic Transmit FIFOs associated with these endpoints on the AHB.
- 5. Program the following fields in the OTG\_HS\_DIEPCTLx register to disable the endpoint:
  - SNAK = 1 in OTG\_HS\_DIEPCTLx
  - EPDIS = 1 in OTG HS DIEPCTLx
- 6. The assertion of the Endpoint Disabled interrupt in OTG\_HS\_DIEPINTx indicates that the core has disabled the endpoint.
  - At this point, the application must flush the data in the associated transmit FIFO or overwrite the existing data in the FIFO by enabling the endpoint for a new transfer in the next micro-frame. To flush the data, the application must use the OTG\_HS\_GRSTCTL register.

#### Stalling nonisochronous IN endpoints

This section describes how the application can stall a nonisochronous endpoint.

#### Application programming sequence:

- 1. Disable the IN endpoint to be stalled. Set the STALL bit as well.
- 2. EPDIS = 1 in OTG\_HS\_DIEPCTLx, when the endpoint is already enabled
  - STALL = 1 in OTG\_HS\_DIEPCTLx
  - The STALL bit always takes precedence over the NAK bit
- 3. Assertion of the Endpoint Disabled interrupt (in OTG\_HS\_DIEPINTx) indicates to the application that the core has disabled the specified endpoint.
- 4. The application must flush the nonperiodic or periodic transmit FIFO, depending on the endpoint type. In case of a nonperiodic endpoint, the application must re-enable the other nonperiodic endpoints that do not need to be stalled, to transmit data.
- 5. Whenever the application is ready to end the STALL handshake for the endpoint, the STALL bit must be cleared in OTG\_HS\_DIEPCTLx.
- If the application sets or clears a STALL bit for an endpoint due to a
  SetFeature.Endpoint Halt command or ClearFeature.Endpoint Halt command, the
  STALL bit must be set or cleared before the application sets up the Status stage
  transfer on the control endpoint.

Special case: stalling the control OUT endpoint

The core must stall IN/OUT tokens if, during the data stage of a control transfer, the host sends more IN/OUT tokens than are specified in the SETUP packet. In this case, the application must enable the ITTXFE interrupt in OTG\_HS\_DIEPINTx and the OTEPDIS interrupt in OTG\_HS\_DOEPINTx during the data stage of the control transfer, after the core has transferred the amount of data specified in the SETUP packet. Then, when the application receives this interrupt, it must set the STALL bit in the corresponding endpoint control register, and clear this interrupt.

### 30.13.8 Worst case response time

When the OTG\_HS controller acts as a device, there is a worst case response time for any tokens that follow an isochronous OUT. This worst case response time depends on the AHB clock frequency.

The core registers are in the AHB domain, and the core does not accept another token before updating these register values. The worst case is for any token following an isochronous OUT, because for an isochronous transaction, there is no handshake and the next token could come sooner. This worst case value is 7 PHY clocks when the AHB clock is the same as the PHY clock. When the AHB clock is faster, this value is smaller.

If this worst case condition occurs, the core responds to bulk/interrupt tokens with a NAK and drops isochronous and SETUP tokens. The host interprets this as a timeout condition for SETUP and retries the SETUP packet. For isochronous transfers, the Incomplete isochronous IN transfer interrupt (IISOIXFR) and Incomplete isochronous OUT transfer interrupt (IISOOXFR) inform the application that isochronous IN/OUT packets were dropped.

#### Choosing the value of TRDT in OTG\_HS\_GUSBCFG

The value in TRDT (OTG\_HS\_GUSBCFG) is the time it takes for the MAC, in terms of PHY clocks after it has received an IN token, to get the FIFO status, and thus the first data from the PFC block. This time involves the synchronization delay between the PHY and AHB clocks. The worst case delay for this is when the AHB clock is the same as the PHY clock. In this case, the delay is 5 clocks.

Once the MAC receives an IN token, this information (token received) is synchronized to the AHB clock by the PFC (the PFC runs on the AHB clock). The PFC then reads the data from the SPRAM and writes them into the dual clock source buffer. The MAC then reads the data out of the source buffer (4 deep).

If the AHB is running at a higher frequency than the PHY, the application can use a smaller value for TRDT (in OTG\_HS\_GUSBCFG).

Figure 391 has the following signals:

- tkn rcvd: Token received information from MAC to PFC
- dynced\_tkn\_rcvd: Doubled sync tkn\_rcvd, from PCLK to HCLK domain
- spr\_read: Read to SPRAM
- spr\_addr: Address to SPRAM
- spr\_rdata: Read data from SPRAM
- srcbuf push: Push to the source buffer
- srcbuf\_rdata: Read data from the source buffer. Data seen by MAC

The application can use the following formula to calculate the value of TRDT:

 $4 \times AHB$  clock + 1 PHY clock = (2 clock sync + 1 clock memory address + 1 clock memory data from sync RAM) + (1 PHY clock (next PHY clock MAC can sample the 2 clock FIFO outputs)

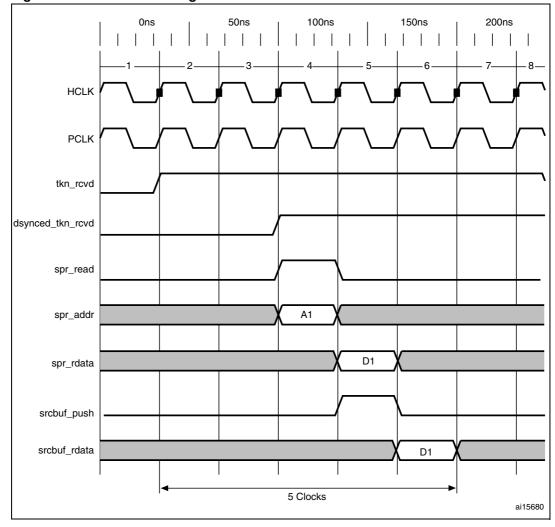


Figure 391. TRDT max timing case

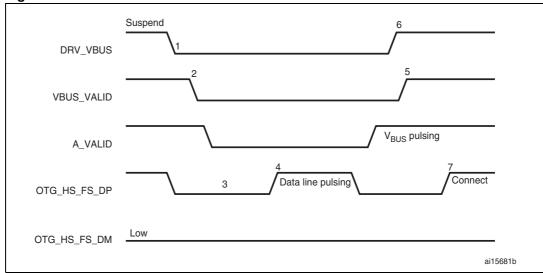
### 30.13.9 OTG programming model

The OTG\_HS controller is an OTG device supporting HNP and SRP. When the core is connected to an "A" plug, it is referred to as an A-device. When the core is connected to a "B" plug it is referred to as a B-device. In host mode, the OTG\_HS controller turns off  $V_{BUS}$  to conserve power. SRP is a method by which the B-device signals the A-device to turn on  $V_{BUS}$  power. A device must perform both data-line pulsing and  $V_{BUS}$  pulsing, but a host can detect either data-line pulsing or  $V_{BUS}$  pulsing for SRP. HNP is a method by which the B-device negotiates and switches to host role. In Negotiated mode after HNP, the B-device suspends the bus and reverts to the device role.

### A-device session request protocol

The application must set the SRP-capable bit in the Core USB configuration register. This enables the OTG\_HS controller to detect SRP as an A-device.

Figure 392. A-device SRP

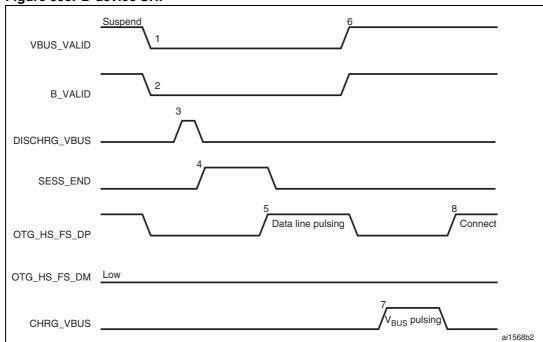


- I. DRV\_VBUS = V<sub>BUS</sub> drive signal to the PHY VBUS\_VALID = V<sub>BUS</sub> valid signal from PHY A\_VALID = A-device V<sub>BUS</sub> level signal to PHY DP = Data plus line DM = Data minus line
- To save power, the application suspends and turns off port power when the bus is idle
  by writing the port suspend and port power bits in the host port control and status
  register.
- 2. PHY indicates port power off by deasserting the VBUS\_VALID signal.
- 3. The device must detect SE0 for at least 2 ms to start SRP when V<sub>BUS</sub> power is off.
- 4. To initiate SRP, the device turns on its data line pull-up resistor for 5 to 10 ms. The OTG\_HS controller detects data-line pulsing.
- 5. The device drives  $V_{BUS}$  above the A-device session valid (2.0 V minimum) for  $V_{BUS}$  pulsing.
  - The OTG\_HS controller interrupts the application on detecting SRP. The Session request detected bit is set in Global interrupt status register (SRQINT set in OTG\_HS\_GINTSTS).
- 6. The application must service the Session request detected interrupt and turn on the port power bit by writing the port power bit in the host port control and status register. The PHY indicates port power-on by asserting the VBUS\_VALID signal.
- When the USB is powered, the device connects, completing the SRP process.

### **B-device session request protocol**

The application must set the SRP-capable bit in the Core USB configuration register. This enables the OTG\_HS controller to initiate SRP as a B-device. SRP is a means by which the OTG\_HS controller can request a new session from the host.





- VBUS\_VALID = V<sub>BUS</sub> valid signal from PHY B\_VALID = B-device valid session to PHY DISCHRG\_VBUS = discharge signal to PHY SESS\_END = session end signal to PHY CHRG\_VBUS = charge V<sub>BUS</sub> signal to PHY DP = Data plus line DM = Data minus line
- 1. To save power, the host suspends and turns off port power when the bus is idle.

The OTG\_HS controller sets the early suspend bit in the Core interrupt register after 3 ms of bus idleness. Following this, the OTG\_HS controller sets the USB suspend bit in the Core interrupt register.

The OTG\_HS controller informs the PHY to discharge V<sub>BUS</sub>.

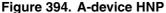
2. The PHY indicates the session's end to the device. This is the initial condition for SRP. The OTG\_HS controller requires 2 ms of SE0 before initiating SRP.

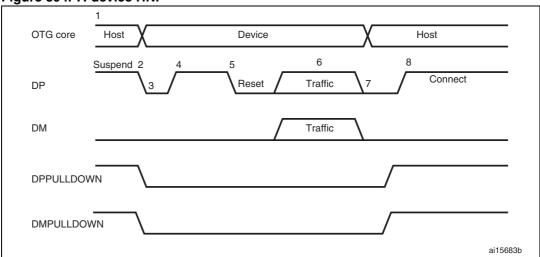
For a USB 1.1 full-speed serial transceiver, the application must wait until  $V_{BUS}$  discharges to 0.2 V after BSVLD (in OTG\_HS\_GOTGCTL) is deasserted. This

- discharge time can be obtained from the transceiver vendor and varies from one transceiver to another.
- 3. The USB OTG core informs the PHY to speed up  $V_{BUS}$  discharge.
- The application initiates SRP by writing the session request bit in the OTG Control and status register. The OTG\_HS controller perform data-line pulsing followed by V<sub>BUS</sub> pulsing.
- The host detects SRP from either the data-line or V<sub>BUS</sub> pulsing, and turns on V<sub>BUS</sub>.
   The PHY indicates V<sub>BUS</sub> power-on to the device.
- 6. The OTG\_HS controller performs V<sub>BUS</sub> pulsing.
  The host starts a new session by turning on V<sub>BUS</sub>, indicating SRP success. The OTG\_HS controller interrupts the application by setting the session request success status change bit in the OTG interrupt status register. The application reads the session request success bit in the OTG control and status register.
- When the USB is powered, the OTG\_HS controller connects, completing the SRP process.

### A-device host negotiation protocol

HNP switches the USB host role from the A-device to the B-device. The application must set the HNP-capable bit in the Core USB configuration register to enable the OTG\_HS controller to perform HNP as an A-device.





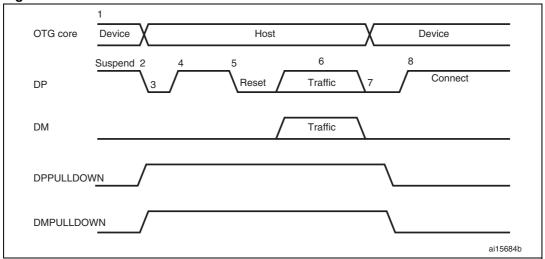
- DPPULLDOWN = signal from core to PHY to enable/disable the pull-down on the DP line inside the PHY. DMPULLDOWN = signal from core to PHY to enable/disable the pull-down on the DM line inside the PHY.
- The OTG\_HS controller sends the B-device a SetFeature b\_hnp\_enable descriptor to enable HNP support. The B-device's ACK response indicates that the B-device supports HNP. The application must set host Set HNP Enable bit in the OTG Control

- and status register to indicate to the OTG\_HS controller that the B-device supports HNP
- 2. When it has finished using the bus, the application suspends by writing the Port suspend bit in the host port control and status register.
- 3. When the B-device observes a USB suspend, it disconnects, indicating the initial condition for HNP. The B-device initiates HNP only when it must switch to the host role; otherwise, the bus continues to be suspended.
  - The OTG\_HS controller sets the host negotiation detected interrupt in the OTG interrupt status register, indicating the start of HNP.
  - The OTG\_HS controller deasserts the DM pull down and DM pull down in the PHY to indicate a device role. The PHY enables the OTG\_HS\_DP pull-up resistor to indicate a connect for B-device.
  - The application must read the current mode bit in the OTG Control and status register to determine peripheral mode operation.
- 4. The B-device detects the connection, issues a USB reset, and enumerates the OTG\_HS controller for data traffic.
- 5. The B-device continues the host role, initiating traffic, and suspends the bus when done.
  - The OTG\_HS controller sets the early suspend bit in the Core interrupt register after 3 ms of bus idleness. Following this, the OTG\_HS controller sets the USB Suspend bit in the Core interrupt register.
- 6. In Negotiated mode, the OTG\_HS controller detects the suspend, disconnects, and switches back to the host role. The OTG\_HS controller asserts the DM pull down and DM pull down in the PHY to indicate its assumption of the host role.
- 7. The OTG\_HS controller sets the Connector ID status change interrupt in the OTG Interrupt Status register. The application must read the connector ID status in the OTG Control and Status register to determine the OTG\_HS controller operation as an Adevice. This indicates the completion of HNP to the application. The application must read the Current mode bit in the OTG control and status register to determine host mode operation.
- 8. The B-device connects, completing the HNP process.

#### **B-device host negotiation protocol**

HNP switches the USB host role from B-device to A-device. The application must set the HNP-capable bit in the Core USB configuration register to enable the OTG\_HS controller to perform HNP as a B-device.

Figure 395. B-device HNP



- DPPULLDOWN = signal from core to PHY to enable/disable the pull-down on the DP line inside the PHY. DMPULLDOWN = signal from core to PHY to enable/disable the pull-down on the DM line inside the PHY.
- The A-device sends the SetFeature b\_hnp\_enable descriptor to enable HNP support.
   The OTG\_HS controller's ACK response indicates that it supports HNP. The application must set the Device HNP enable bit in the OTG Control and status register to indicate HNP support.
  - The application sets the HNP request bit in the OTG Control and status register to indicate to the OTG\_HS controller to initiate HNP.
- 2. When it has finished using the bus, the A-device suspends by writing the Port suspend bit in the host port control and status register.
  - The OTG\_HS controller sets the Early suspend bit in the Core interrupt register after 3 ms of bus idleness. Following this, the OTG\_HS controller sets the USB suspend bit in the Core interrupt register.
  - The OTG\_HS controller disconnects and the A-device detects SE0 on the bus, indicating HNP. The OTG\_HS controller asserts the DP pull down and DM pull down in the PHY to indicate its assumption of the host role.

The A-device responds by activating its OTG\_HS\_DP pull-up resistor within 3 ms of detecting SE0. The OTG\_HS controller detects this as a connect.

The OTG\_HS controller sets the host negotiation success status change interrupt in the OTG Interrupt status register, indicating the HNP status. The application must read the host negotiation success bit in the OTG Control and status register to determine

- host negotiation success. The application must read the current Mode bit in the Core interrupt register (OTG\_HS\_GINTSTS) to determine host mode operation.
- 3. The application sets the reset bit (PRST in OTG\_HS\_HPRT) and the OTG\_HS controller issues a USB reset and enumerates the A-device for data traffic.
- 4. The OTG\_HS controller continues the host role of initiating traffic, and when done, suspends the bus by writing the Port suspend bit in the host port control and status register.
- 5. In Negotiated mode, when the A-device detects a suspend, it disconnects and switches back to the host role. The OTG\_HS controller deasserts the DP pull down and DM pull down in the PHY to indicate the assumption of the device role.
- 6. The application must read the current mode bit in the Core interrupt (OTG\_HS\_GINTSTS) register to determine the host mode operation.
- 7. The OTG\_HS controller connects, completing the HNP process.

# 31 Flexible static memory controller (FSMC)

This section applies to the whole STM32F40x and STM32F41x family, only.

### 31.1 FSMC main features

The FSMC block is able to interface with synchronous and asynchronous memories and 16-bit PC memory cards. Its main purpose is to:

- Translate the AHB transactions into the appropriate external device protocol
- Meet the access timing requirements of the external devices

All external memories share the addresses, data and control signals with the controller. Each external device is accessed by means of a unique chip select. The FSMC performs only one access at a time to an external device.

The FSMC has the following main features:

- Interfaces with static memory-mapped devices including:
  - Static random access memory (SRAM)
  - Read-only memory (ROM)
  - NOR Flash memory/OneNAND Flash memory
  - PSRAM (4 memory banks)
- Two banks of NAND Flash with ECC hardware that checks up to 8 Kbytes of data
- 16-bit PC Card compatible devices
- Supports burst mode access to synchronous devices (NOR Flash and PSRAM)
- 8- or 16-bit wide databus
- Independent chip select control for each memory bank
- Independent configuration for each memory bank
- Programmable timings to support a wide range of devices, in particular:
  - Programmable wait states (up to 15)
  - Programmable bus turnaround cycles (up to 15)
  - Programmable output enable and write enable delays (up to 15)
  - Independent read and write timings and protocol, so as to support the widest variety of memories and timings
- Write enable and byte lane select outputs for use with PSRAM and SRAM devices
- Translation of 32-bit wide AHB transactions into consecutive 16-bit or 8-bit accesses to external 16-bit or 8-bit devices
- A Write FIFO, 2-word long, each word is 32 bits wide, only stores data and not the address. Therefore, this FIFO only buffers AHB write burst transactions. This makes it possible to write to slow memories and free the AHB quickly for other operations. Only one burst at a time is buffered: if a new AHB burst or single transaction occurs while an operation is in progress, the FIFO is drained. The FSMC will insert wait states until the current memory access is complete).
- External asynchronous wait control

The FSMC registers that define the external device type and associated characteristics are usually set at boot time and do not change until the next reset or power-up. However, it is possible to change the settings at any time.

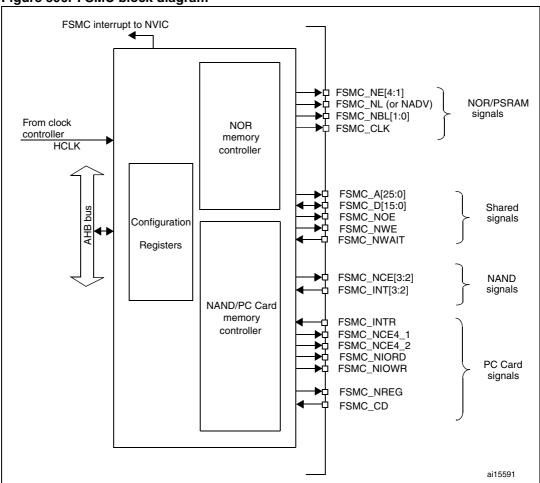
# 31.2 Block diagram

The FSMC consists of four main blocks:

- The AHB interface (including the FSMC configuration registers)
- The NOR Flash/PSRAM controller
- The NAND Flash/PC Card controller
- The external device interface

The block diagram is shown in Figure 396.

Figure 396. FSMC block diagram



### 31.3 AHB interface

The AHB slave interface enables internal CPUs and other bus master peripherals to access the external static memories.

AHB transactions are translated into the external device protocol. In particular, if the selected external memory is 16 or 8 bits wide, 32-bit wide transactions on the AHB are split into consecutive 16- or 8-bit accesses. The Chip Select toggles for each access.

The FSMC generates an AHB error in the following conditions:

- When reading or writing to an FSMC bank which is not enabled
- When reading or writing to the NOR Flash bank while the FACCEN bit is reset in the FSMC\_BCRx register.
- When reading or writing to the PC Card banks while the input pin FSMC\_CD (Card Presence Detection) is low.

The effect of this AHB error depends on the AHB master which has attempted the R/W access:

- If it is the Cortex<sup>™</sup>-M4F CPU, a hard fault interrupt is generated
- If is a DMA, a DMA transfer error is generated and the corresponding DMA channel is automatically disabled.

The AHB clock (HCLK) is the reference clock for the FSMC.

### 31.3.1 Supported memories and transactions

#### General transaction rules

The requested AHB transaction data size can be 8-, 16- or 32-bit wide whereas the accessed external device has a fixed data width. This may lead to inconsistent transfers.

Therefore, some simple transaction rules must be followed:

- AHB transaction size and memory data size are equal There is no issue in this case.
- AHB transaction size is greater than the memory size
   In this case, the FSMC splits the AHB transaction into smaller consecutive memory accesses in order to meet the external data width.
- AHB transaction size is smaller than the memory size
   Asynchronous transfers may or not be consistent depending on the type of external device.
  - Asynchronous accesses to devices that have the byte select feature (SRAM, ROM, PSRAM).
    - a) FSMC allows write transactions accessing the right data through its byte lanes NBL[1:0]
    - b) Read transactions are allowed. All memory bytes are read and the useless ones are discarded. The NBL[1:0] are kept low during read transactions.
  - Asynchronous accesses to devices that do not have the byte select feature (NOR and NAND Flash 16-bit).
    - This situation occurs when a byte access is requested to a 16-bit wide Flash memory. Clearly, the device cannot be accessed in byte mode (only 16-bit words can be read from/written to the Flash memory) therefore:
  - a) Write transactions are not allowed
  - b) Read transactions are allowed. All memory bytes are read and the useless ones are discarded. The NBL[1:0] are set to 0 during read transactions.

### **Configuration registers**

The FSMC can be configured using a register set. See *Section 31.5.6*, for a detailed description of the NOR Flash/PSRAM control registers. See *Section 31.6.8*, for a detailed description of the NAND Flash/PC Card registers.

# 31.4 External device address mapping

From the FSMC point of view, the external memory is divided into 4 fixed-size banks of 256 Mbytes each (Refer to *Figure 397*):

- Bank 1 used to address up to 4 NOR Flash or PSRAM memory devices. This bank is split into 4 NOR/PSRAM regions with 4 dedicated Chip Select.
- Banks 2 and 3 used to address NAND Flash devices (1 device per bank)
- Bank 4 used to address a PC Card device

For each bank the type of memory to be used is user-defined in the Configuration register.

Address Banks Supported memory type 6000 0000h Bank 1 NOR / PSRAM 4 × 64 MB 6FFF FFFFh 7000 0000h Bank 2 4 × 64 MB 7FFF FFFFh NAND Flash 8000 0000h Bank 3 4 × 64 MB 8FFF FFFFh 9000 0000h Bank 4 PC Card 4 × 64 MB 9FFF FFFFh

Figure 397. FSMC memory banks

## 31.4.1 NOR/PSRAM address mapping

HADDR[27:26] bits are used to select one of the four memory banks as shown in *Table 167*.

Table 167. NOR/PSRAM bank selection

HADDR[27:26] <sup>(1)</sup>	Selected bank
00	Bank 1 NOR/PSRAM 1
01	Bank 1 NOR/PSRAM 2

Table 167. NOR/PSRAM bank selection (continued)

HADDR[27:26] <sup>(1)</sup>	Selected bank		
10	Bank 1 NOR/PSRAM 3		
11	Bank 1 NOR/PSRAM 4		

<sup>1.</sup> HADDR are internal AHB address lines that are translated to external memory.

HADDR[25:0] contain the external memory address. Since HADDR is a byte address whereas the memory is addressed in words, the address actually issued to the memory varies according to the memory data width, as shown in the following table.

Table 168. External memory address

Memory width <sup>(1)</sup>	Data address issued to the memory	Maximum memory capacity (bits)
8-bit	HADDR[25:0]	64 Mbytes x 8 = 512 Mbit
16-bit	HADDR[25:1] >> 1	64 Mbytes/2 x 16 = 512 Mbit

In case of a 16-bit external memory width, the FSMC will internally use HADDR[25:1] to generate the
address for external memory FSMC\_A[24:0].
 Whatever the external memory width (16-bit or 8-bit), FSMC\_A[0] should be connected to external memory
address A[0].

#### Wrap support for NOR Flash/PSRAM

Wrap burst mode for synchronous memories is not supported. The memories must be configured in linear burst mode of undefined length.

### 31.4.2 NAND/PC Card address mapping

In this case, three banks are available, each of them divided into memory spaces as indicated in *Table 169*.

Table 169. Memory mapping and timing registers

Start address	End address	FSMC Bank	Memory space	Timing register
0x9C00 0000	0x9FFF FFFF		I/O	FSMC_PIO4 (0xB0)
0x9800 0000	0x9BFF FFFF	Bank 4 - PC card	Attribute	FSMC_PATT4 (0xAC)
0x9000 0000	0x93FF FFFF		Common	FSMC_PMEM4 (0xA8)
0x8800 0000	0x8BFF FFFF	Bank 3 - NAND Flash	Attribute	FSMC_PATT3 (0x8C)
0x8000 0000	0x83FF FFFF	Dalik 3 - NAND Flasii	Common	FSMC_PMEM3 (0x88)
0x7800 0000	0x7BFF FFFF	Bank 2- NAND Flash	Attribute	FSMC_PATT2 (0x6C)
0x7000 0000	0x73FF FFFF	Dalik 2- NAIND Flasii	Common	FSMC_PMEM2 (0x68)

For NAND Flash memory, the common and attribute memory spaces are subdivided into three sections (see in *Table 170* below) located in the lower 256 Kbytes:

- Data section (first 64 Kbytes in the common/attribute memory space)
- Command section (second 64 Kbytes in the common / attribute memory space)
- Address section (next 128 Kbytes in the common / attribute memory space)

Table 170. NAND bank selections

Section name	HADDR[17:16]	Address range
Address section	1X	0x020000-0x03FFFF
Command section	01	0x010000-0x01FFFF
Data section	00	0x000000-0x0FFFF

The application software uses the 3 sections to access the NAND Flash memory:

- To send a command to NAND Flash memory: the software must write the command value to any memory location in the command section.
- To specify the NAND Flash address that must be read or written: the software must write the address value to any memory location in the address section. Since an address can be 4 or 5 bytes long (depending on the actual memory size), several consecutive writes to the address section are needed to specify the full address.
- To read or write data: the software reads or writes the data value from or to any memory location in the data section.

Since the NAND Flash memory automatically increments addresses, there is no need to increment the address of the data section to access consecutive memory locations.

### 31.5 NOR Flash/PSRAM controller

The FSMC generates the appropriate signal timings to drive the following types of memories:

- Asynchronous SRAM and ROM
  - 8-bit
  - 16-bit
  - 32-bit
- PSRAM (Cellular RAM)
  - Asynchronous mode
  - Burst mode
  - Multiplexed or nonmultiplexed
- NOR Flash
  - Asynchronous mode or burst mode
  - Multiplexed or nonmultiplexed

The FSMC outputs a unique chip select signal NE[4:1] per bank. All the other signals (addresses, data and control) are shared.

For synchronous accesses, the FSMC issues the clock (CLK) to the selected external device only during the read/write transactions. This clock is a submultiple of the HCLK clock. The size of each bank is fixed and equal to 64 Mbytes.

Each bank is configured by means of dedicated registers (see Section 31.5.6).

The programmable memory parameters include access timings (see *Table 171*) and support for wait management (for PSRAM and NOR Flash accessed in burst mode).

Table 171. Programmable NOR/PSRAM access parameters

Parameter	Function	Access mode	Unit	Min.	Max.
Address setup	Duration of the address setup phase	Asynchronous	AHB clock cycle (HCLK)	0	15
Address hold	Duration of the address hold phase	Asynchronous, muxed I/Os	AHB clock cycle (HCLK)	1	15
Data setup	Duration of the data setup phase	Asynchronous	AHB clock cycle (HCLK)	1	256
Bust turn	Duration of the bus turnaround phase	Asynchronous and synchronous read	AHB clock cycle (HCLK)	0	15
Clock divide ratio	Number of AHB clock cycles (HCLK) to build one memory clock cycle (CLK)	Synchronous	AHB clock cycle (HCLK)	2	16
Data latency	Number of clock cycles to issue to the memory before the first data of the burst	Synchronous	Memory clock cycle (CLK)	2	17

### 31.5.1 External memory interface signals

*Table 172*, *Table 173* and *Table 174* list the signals that are typically used to interface NOR Flash, SRAM and PSRAM.

Note: Prefix "N". specifies the associated signal as active low.

### NOR Flash, nonmultiplexed I/Os

Table 172. Nonmultipled I/O NOR Flash

FSMC signal name	I/O	Function
CLK	0	Clock (for synchronous burst)
A[25:0]	0	Address bus
D[15:0]	I/O	Bidirectional data bus
NE[x]	0	Chip select, x = 14
NOE	0	Output enable
NWE	0	Write enable
NL(=NADV)	0	Latch enable (this signal is called address valid, NADV, by some NOR Flash devices)
NWAIT	I	NOR Flash wait input signal to the FSMC

NOR Flash memories are addressed in 16-bit words. The maximum capacity is 512 Mbit (26 address lines).

### NOR Flash, multiplexed I/Os

Table 173. Multiplexed I/O NOR Flash

FSMC signal name	I/O	Function
CLK	0	Clock (for synchronous burst)
A[25:16]	0	Address bus
AD[15:0]	I/O	16-bit multiplexed, bidirectional address/data bus
NE[x]	0	Chip select, x = 14
NOE	0	Output enable
NWE	0	Write enable
NL(=NADV)	0	Latch enable (this signal is called address valid, NADV, by some NOR Flash devices)
NWAIT	I	NOR Flash wait input signal to the FSMC

NOR-Flash memories are addressed in 16-bit words. The maximum capacity is 512 Mbit (26 address lines).

### PSRAM/SRAM, nonmultiplexed I/Os

Table 174. Nonmultiplexed I/Os PSRAM/SRAM

FSMC signal name	I/O	Function
CLK	0	Clock (only for PSRAM synchronous burst)
A[25:0]	0	Address bus
D[15:0]	I/O	Data bidirectional bus
NE[x]	0	Chip select, x = 14 (called NCE by PSRAM (Cellular RAM i.e. CRAM))
NOE	0	Output enable
NWE	0	Write enable
NL(= NADV)	0	Address valid only for PSRAM input (memory signal name: NADV)
NWAIT	I	PSRAM wait input signal to the FSMC
NBL[1]	0	Upper byte enable (memory signal name: NUB)
NBL[0]	0	Lowed byte enable (memory signal name: NLB)

PSRAM memories are addressed in 16-bit words. The maximum capacity is 512 Mbit (26 address lines).

### PSRAM, multiplexed I/Os

Table 175. Multiplexed I/O PSRAM

FSMC signal name	I/O	Function			
CLK	0	Clock (for synchronous burst)			
A[25:16]	0	Address bus			
AD[15:0]	I/O	16-bit multiplexed, bidirectional address/data bus			
NE[x]	0	Chip select, x = 14 (called NCE by PSRAM (Cellular RAM i.e. CRAM))			
NOE	0	Output enable			
NWE	0	Write enable			
NL(= NADV)	0	Address valid PSRAM input (memory signal name: NADV)			
NWAIT	I	PSRAM wait input signal to the FSMC			
NBL[1]	0	Upper byte enable (memory signal name: NUB)			
NBL[0]	0	Lowed byte enable (memory signal name: NLB)			

PSRAM memories are addressed in 16-bit words. The maximum capacity is 512 Mbit (26 address lines).

# 31.5.2 Supported memories and transactions

*Table 176* below displays an example of the supported devices, access modes and transactions when the memory data bus is 16-bit for NOR, PSRAM and SRAM. Transactions not allowed (or not supported) by the FSMC in this example appear in gray.

Table 176. NOR Flash/PSRAM supported memories and transactions

Device	Mode	R/W	AHB data size	Memory data size	Allowed/ not allowed	Comments
	Asynchronous	R	8	16	Y	
	Asynchronous	W	8	16	N	
	Asynchronous	R	16	16	Y	
	Asynchronous	W	16	16	Υ	
NOR Flash	Asynchronous	R	32	16	Υ	Split into 2 FSMC accesses
(muxed I/Os and nonmuxed I/Os)	Asynchronous	W	32	16	Y	Split into 2 FSMC accesses
	Asynchronous page	R	-	16	N	Mode is not supported
	Synchronous	R	8	16	N	
	Synchronous	R	16	16	Y	
	Synchronous	R	32	16	Υ	

Table 176. NOR Flash/PSRAM supported memories and transactions (continued)

Device	Mode	R/W	AHB data size	Memory data size	Allowed/ not allowed	Comments
	Asynchronous	R	8	16	Υ	
	Asynchronous	W	8	16	Y	Use of byte lanes NBL[1:0]
	Asynchronous	R	16	16	Υ	
	Asynchronous	W	16	16	Y	
PSRAM	Asynchronous	R	32	16	Υ	Split into 2 FSMC accesses
(multiplexed I/Os and	Asynchronous	W	32	16	Υ	Split into 2 FSMC accesses
nonmultiplexed I/Os)	Asynchronous page	R	-	16	N	Mode is not supported
	Synchronous	R	8	16	N	
	Synchronous	R	16	16	Υ	
	Synchronous	R	32	16	Υ	
	Synchronous	W	8	16	Υ	Use of byte lanes NBL[1:0]
	Synchronous	W	16/32	16	Υ	
	Asynchronous	R	8 / 16	16	Υ	
	Asynchronous	W	8 / 16	16	Υ	Use of byte lanes NBL[1:0]
SRAM and ROM	Asynchronous	R	32	16	Υ	Split into 2 FSMC accesses
	Asynchronous	W	32	16	Y	Split into 2 FSMC accesses. Use of byte lanes NBL[1:0]

# 31.5.3 General timing rules

### Signals synchronization

- All controller output signals change on the rising edge of the internal clock (HCLK)
- In synchronous read and write mode, the output data changes on the falling edge of the memory clock (FSMC\_CLK).

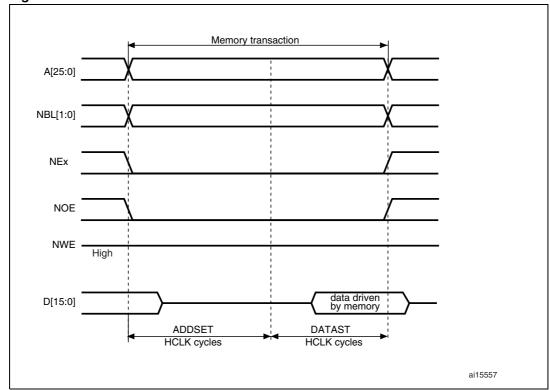
### 31.5.4 NOR Flash/PSRAM controller asynchronous transactions

### Asynchronous static memories (NOR Flash, SRAM)

- Signals are synchronized by the internal clock HCLK. This clock is not issued to the memory
- The FSMC always samples the data before de-asserting the chip select signal NE. This guarantees that the memory data-hold timing constraint is met (chip enable high to data transition, usually 0 ns min.)
- When extended mode is set, it is possible to mix modes A, B, C and D in read and write (it is for instance possible to read in mode A and write in mode B).

### Mode 1 - SRAM/CRAM

Figure 398. Mode1 read accesses



1. NBL[1:0] are driven low during read access.

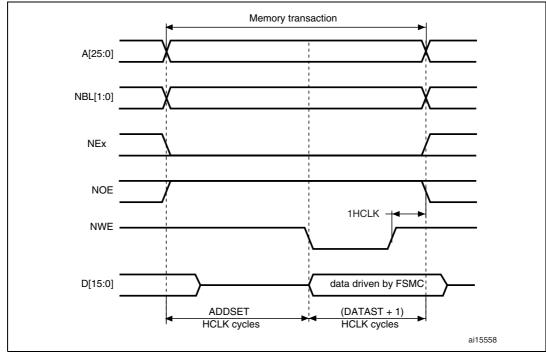


Figure 399. Mode1 write accesses

The one HCLK cycle at the end of the write transaction helps guarantee the address and data hold time after the NWE rising edge. Due to the presence of this one HCLK cycle, the DATAST value must be greater than zero (DATAST > 0).

Table 177. FSMC\_BCRx bit fields

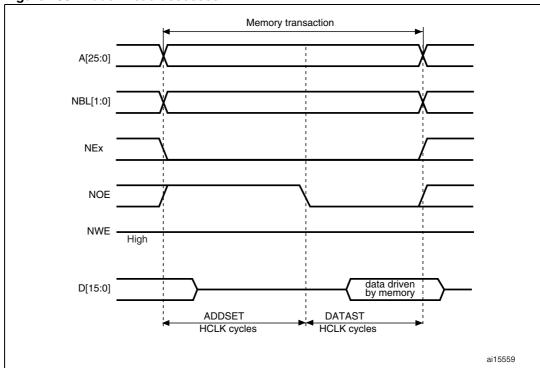
Bit number	Bit name	Value to set
31-16		0x0000
15	ASYNCWAIT	Set to 1 if the memory supports this feature. Otherwise keep at 0.
14-10		0x0
9	WAITPOL	Meaningful only if bit 15 is 1
8	BURSTEN	0x0
7		-
6	FACCEN	-
5-4	MWID	As needed
3-2	MTYP	As needed, exclude 10 (NOR Flash)
1	MUXEN	0x0
0	MBKEN	0x1

Table 178. FSMC\_BTRx bit fields

Bit number	Bit name	Value to set
31-20		0x0000
19-16	BUSTURN	Time between NEx high to NEx low (BUSTURN HCLK)
15-8	DATAST	Duration of the second access phase (DATAST+1 HCLK cycles for write accesses, DATAST HCLK cycles for read accesses).
7-4		0x0
3-0	ADDSET	Duration of the first access phase (ADDSET HCLK cycles). Minimum value for ADDSET is 0.

### Mode A - SRAM/PSRAM (CRAM) OE toggling

Figure 400. ModeA read accesses



1. NBL[1:0] are driven low during read access.

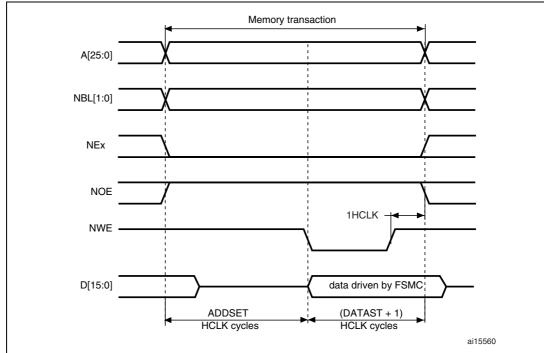


Figure 401. ModeA write accesses

The differences compared with mode1 are the toggling of NOE and the independent read and write timings.

Table 179. FSMC\_BCRx bit fields

Bit number	Bit name	Value to set
31-16		0x0000
15	ASYNCWAIT	Set to 1 if the memory supports this feature. Otherwise keep at 0.
14	EXTMOD	0x1
13-10		0x0
9	WAITPOL	Meaningful only if bit 15 is 1
8	BURSTEN	0x0
7		-
6	FACCEN	-
5-4	MWID	As needed
3-2	MTYP	As needed, exclude 10 (NOR Flash)
1	MUXEN	0x0
0	MBKEN	0x1

Table 180. FSMC\_BTRx bit fields

Bit number	Bit name	Value to set
31-30		0x0
29-28	ACCMOD	0x0
27-20		0x000
19-16	BUSTURN	Time between NEx high to NEx low (BUSTURN HCLK)
15-8	DATAST	Duration of the second access phase (DATAST HCLK cycles) in read.
7-4		0x0
3-0	ADDSET	Duration of the first access phase (ADDSET HCLK cycles) in read. Minimum value for ADDSET is 0.

Table 181. FSMC\_BWTRx bit fields

Bit number	Bit name	Value to set
31-30		0x0
29-28	ACCMOD	0x0
27-20		0x000
19-16	BUSTURN	Time between NEx high to NEx low (BUSTURN HCLK).
15-8	DATAST	Duration of the second access phase (DATAST+1 HCLK cycles) in write.
7-4		0x0
3-0	ADDSET	Duration of the first access phase (ADDSET HCLK cycles) in write Minimum value for ADDSET is 0.

### Mode 2/B - NOR Flash

Figure 402. Mode2/B read accesses

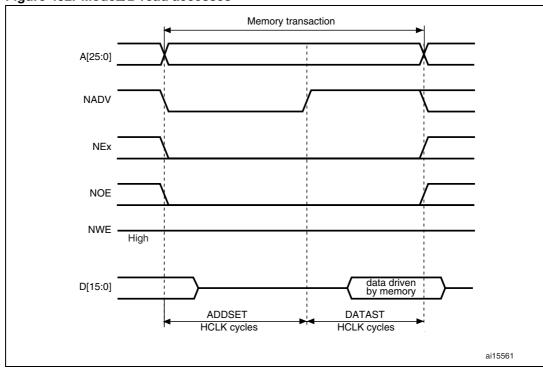
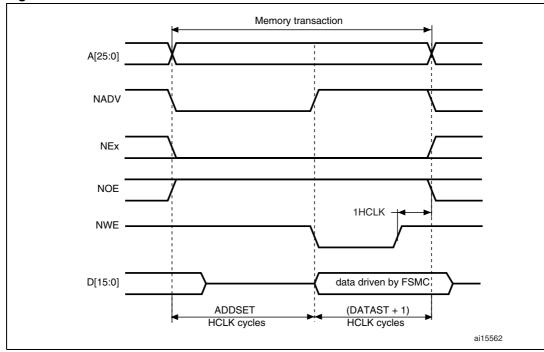


Figure 403. Mode2 write accesses



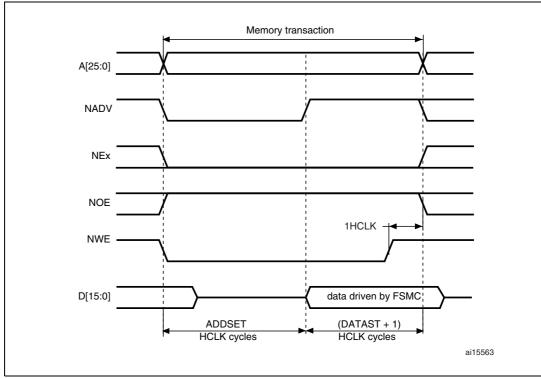


Figure 404. ModeB write accesses

The differences with mode1 are the toggling of NADV and the independent read and write timings when extended mode is set (Mode B).

Table 182. FSMC\_BCRx bit fields

Table 102.	z. FSMC_DChx bit lields	
Bit number	Bit name	Value to set
31-16		0x0000
15	ASYNCWAIT	Set to 1 if the memory supports this feature. Otherwise keep at 0.
14	EXTMOD	0x1 for mode B, 0x0 for mode 2
13-10		0x0
9	WAITPOL	Meaningful only if bit 15 is 1
8	BURSTEN	0x0
7		-
6	FACCEN	0x1
5-4	MWID	As needed
3-2	MTYP	10 (NOR Flash)
1	MUXEN	0x0
0	MBKEN	0x1
		1

Table 183. FSMC\_BTRx bit fields

Bit number	Bit name	Value to set
31-30		0x0
29-28	ACCMOD	0x1 if extended mode is set
27-20		0x000
19-16	BUSTURN	Time between NEx high to NEx low (BUSTURN HCLK)
15-8	DATAST	Duration of the access second phase (DATAST HCLK cycles) in read.
7-4		0x0
3-0	ADDSET	Duration of the access first phase (ADDSET HCLK cycles) in read. Minimum value for ADDSET is 0.

Table 184. FSMC\_BWTRx bit fields

Bit number	Bit name	Value to set
31-30		0x0
29-28	ACCMOD	0x1 if extended mode is set
27-20		0x000
19-16	BUSTURN	Time between NEx high to NEx low (BUSTURN HCLK)
15-8	DATAST	Duration of the access second phase (DATAST+1 HCLK cycles) in write.
7-4		0x0
3-0	ADDSET	Duration of the access first phase (ADDSET HCLK cycles) in write. Minimum value for ADDSET is 0.

Note: The FSMC\_BWTRx register is valid only if extended mode is set (mode B), otherwise all its content is don't care.

### Mode C - NOR Flash - OE toggling

Figure 405. ModeC read accesses

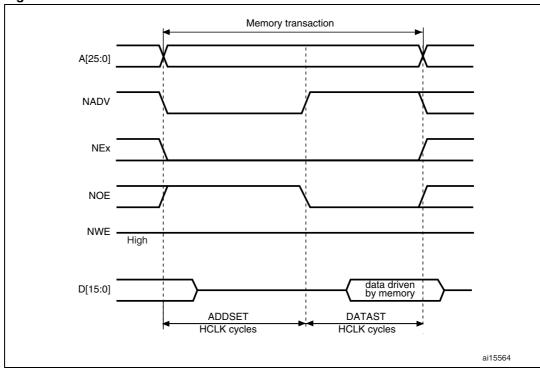
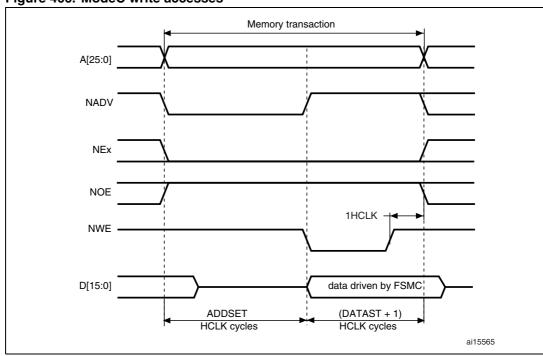


Figure 406. ModeC write accesses



The differences compared with mode1 are the toggling of NOE and NADV and the independent read and write timings.

Table 185. FSMC\_BCRx bit fields

Bit No.	Bit name	Value to set
31-16		0x0000
15	ASYNCWAIT	Set to 1 if the memory supports this feature. Otherwise keep at 0.
14	EXTMOD	0x1
13-10		0x0
9	WAITPOL	Meaningful only if bit 15 is 1
8	BURSTEN	0x0
7		-
6	FACCEN	1
5-4	MWID	As needed
3-2	MTYP	0x02 (NOR Flash)
1	MUXEN	0x0
0	MBKEN	0x1

### Table 186. FSMC\_BTRx bit fields

Bit No.	Bit name	Value to set
31-30		0x0
29-28	ACCMOD	0x2
27-20		0x000
19-16	BUSTURN	Time between NEx high to NEx low (BUSTURN HCLK)
15-8	DATAST	Duration of the second access phase (DATAST HCLK cycles) in read.
7-4		0x0
3-0	ADDSET	Duration of the first access phase (ADDSETHCLK cycles) in read. Minimum value for ADDSET is 0.

### Table 187. FSMC\_BWTRx bit fields

Bit No.	Bit name	Value to set
31-30		0x0
29-28	ACCMOD	0x2
27-20		0x000
19-16	BUSTURN	Time between NEx high to NEx low (BUSTURN HCLK)
15-8	DATAST	Duration of the second access phase (DATAST+1 HCLK cycles) in write.
7-4		0x0
3-0	ADDSET	Duration of the first access phase (ADDSET HCLK cycles) in write. Minimum value for ADDSET is 0.

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## Mode D - asynchronous access with extended address

Figure 407. ModeD read accesses

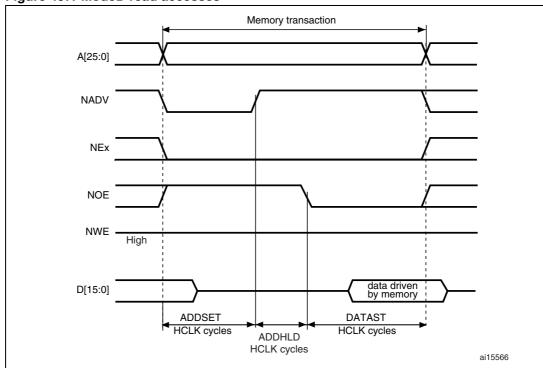
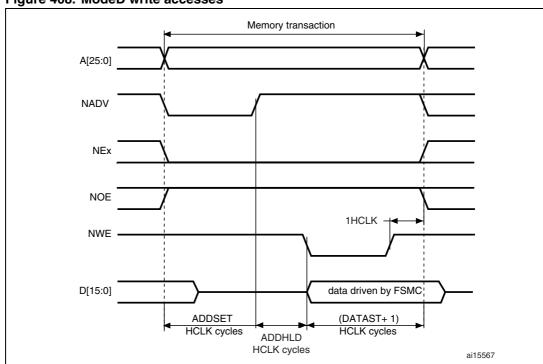


Figure 408. ModeD write accesses



The differences with mode1 are the toggling of NADV, NOE that goes on toggling after NADV changes and the independent read and write timings.

Table 188. FSMC\_BCRx bit fields

Bit No.	Bit name	Value to set					
31-16		0x0000					
15	ASYNCWAIT	et to 1 if the memory supports this feature. Otherwise keep at 0.					
14	EXTMOD	0x1					
13-10		0x0					
9	WAITPOL	Meaningful only if bit 15 is 1					
8	BURSTEN	0x0					
7		-					
6	FACCEN	Set according to memory support					
5-4	MWID	As needed					
3-2	MTYP	As needed					
1	MUXEN	0x0					
0	MBKEN	0x1					

## Table 189. FSMC\_BTRx bit fields

Bit No.	Bit name	Value to set
31-30		0x0
29-28	ACCMOD	0x2
27-20		0x000
19-16	BUSTURN	Time between NEx high to NEx low (BUSTURN HCLK)
15-8	DATAST	Duration of the second access phase (DATAST HCLK cycles) in read.
7-4	ADDHLD	Duration of the middle phase of the read access (ADDHLD HCLK cycles)
3-0	ADDSET	Duration of the first access phase (ADDSETHCLK cycles) in read. Minimum value for ADDSET is 1.

## Table 190. FSMC\_BWTRx bit fields

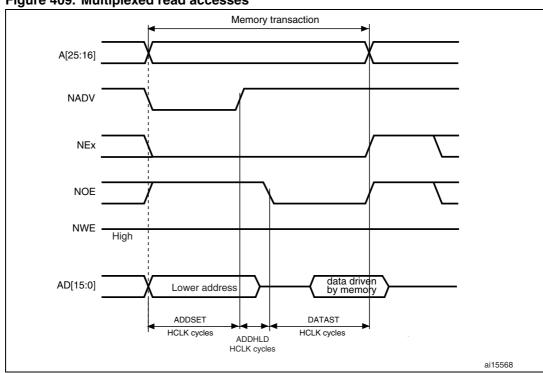
Bit No.	Bit name	Value to set			
31-30		0x0			
29-28	ACCMOD	0x2			
27-20		000			
19-16	BUSTURN	Time between NEx high to NEx low (BUSTURN HCLK)			
15-8	DATAST	Duration of the second access phase (DATAST+1 HCLK cycles) in write.			

Table 190. FSMC\_BWTRx bit fields (continued)

Bit No.	Bit name	Value to set
7-4	ADDHLD	Duration of the middle phase of the write access (ADDHLD HCLK cycles)
3-0	$ \Delta \cap \cup S =  $	Duration of the first access phase (ADDSET HCLK cycles) in write. Minimum value for ADDSET is 1.

## Muxed mode - multiplexed asynchronous access to NOR Flash memory

Figure 409. Multiplexed read accesses



1263/1340 Doc ID 018909 Rev 2

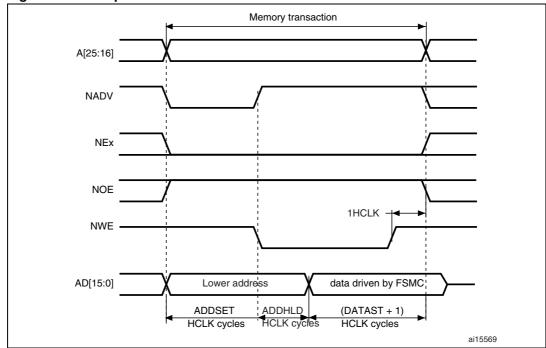


Figure 410. Multiplexed write accesses

The difference with mode D is the drive of the lower address byte(s) on the databus.

Table 191. FSMC\_BCRx bit fields

Table 101. 10Me_Bottx Sit licido								
Bit No.	Bit name	Value to set						
31-16		0x0000						
15	ASYNCWAIT	Set to 1 if the memory supports this feature. Otherwise keep at 0.						
14	EXTMOD	0x0						
13-10		0x0						
9	WAITPOL	Meaningful only if bit 15 is 1						
8	BURSTEN	0x0						
7		-						
6	FACCEN	0x1						
5-4	MWID	As needed						
3-2	MTYP	0x2 (NOR)						
1	MUXEN	0x1						
0	MBKEN	0x1						

Table 192. FSMC\_BTRx bit fields

Bit No.	Bit name	Value to set					
31-20		0000					
19-16	BUSTURN	ime between NEx high to NEx low (BUSTURN HCLK)					

Bit No. Bit name

Value to set

Duration of the second access phase (DATAST HCLK cycles for read accesses and DATAST+1 HCLK cycles for write accesses).

ADDHLD

Duration of the middle phase of the access (ADDHLD HCLK cycles).

Duration of the first access phase (ADDSET HCLK cycles).

Minimum value for ADDSET is 1.

Table 192. FSMC BTRx bit fields (continued)

## WAIT management in asynchronous accesses

If the asynchronous memory asserts a WAIT signal to indicate that it is not yet ready to accept or to provide data, the ASYNCWAIT bit has to be set in FSMC\_BCRx register.

If the WAIT signal is active (high or low depending on the WAITPOL bit), the second access phase (Data setup phase) programmed by the DATAST bits, is extended until WAIT becomes inactive. Unlike the data setup phase, the first access phases (Address setup and Address hold phases), programmed by the ADDSET and ADDHLD bits, are not WAIT sensitive and so they are not prolonged.

The data setup phase ( DATAST in the FSMC\_BTRx register) must be programmed so that WAIT can be detected 4 HCLK cycles before the end of memory transaction. The following cases must be considered:

Memory asserts the WAIT signal aligned to NOE/NWE which toggles:

DATAST 
$$\geq$$
 (4 × HCLK) + max\_wait\_assertion\_time

Memory asserts the WAIT signal aligned to NEx (or NOE/NWE not toggling):

max\_wait\_assertion\_time > address\_phase + hold\_phase

then

 $\label{eq:defDATAST} \mbox{$D$ATAST$} \geq (4 \times \mbox{$HCLK$}) + \left(\mbox{$max$\_wait$\_assertion\_time} - \mbox{$address$\_phase} - \mbox{$hold$\_phase}\right)$  otherwise

$$DATAST \ge 4 \times HCLK$$

where max\_wait\_assertion\_time is the maximum time taken by the memory to assert the WAIT signal once NEx/NOE/NWE is low.

Figure 411 and Figure 412 show the number of HCLK clock cycles that are added to the memory access after WAIT is released by the asynchronous memory (independently of the above cases).

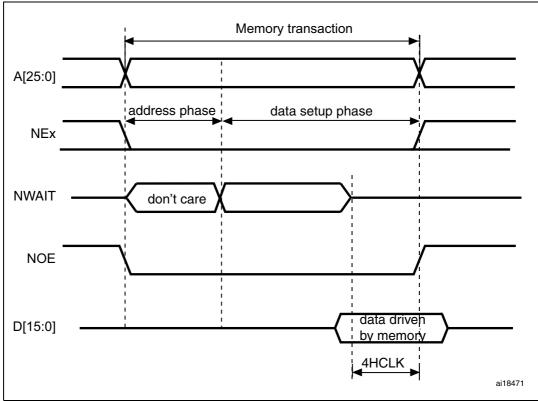
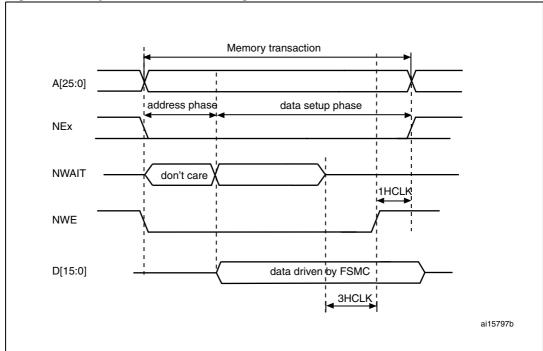


Figure 411. Asynchronous wait during a read access





## 31.5.5 Synchronous burst transactions

The memory clock, CLK, is a submultiple of HCLK according to the value of parameter CLKDIV.

NOR Flash memories specify a minimum time from NADV assertion to CLK high. To meet this constraint, the FSMC does not issue the clock to the memory during the first internal clock cycle of the synchronous access (before NADV assertion). This guarantees that the rising edge of the memory clock occurs *in the middle* of the NADV low pulse.

## Data latency versus NOR Flash latency

The data latency is the number of cycles to wait before sampling the data. The DATLAT value must be consistent with the latency value specified in the NOR Flash configuration register. The FSMC does not include the clock cycle when NADV is low in the data latency count.

## Caution:

Some NOR Flash memories include the NADV Low cycle in the data latency count, so the exact relation between the NOR Flash latency and the FMSC DATLAT parameter can be either of:

- NOR Flash latency = DATLAT + 2
- NOR Flash latency = DATLAT + 3

Some recent memories assert NWAIT during the latency phase. In such cases DATLAT can be set to its minimum value. As a result, the FSMC samples the data and waits long enough to evaluate if the data are valid. Thus the FSMC detects when the memory exits latency and real data are taken.

Other memories do not assert NWAIT during latency. In this case the latency must be set correctly for both the FSMC and the memory, otherwise invalid data are mistaken for good data, or valid data are lost in the initial phase of the memory access.

### Single-burst transfer

When the selected bank is configured in synchronous burst mode, if an AHB single-burst transaction is requested, the FSMC performs a burst transaction of length 1 (if the AHB transfer is 16-bit), or length 2 (if the AHB transfer is 32-bit) and de-assert the chip select signal when the last data is strobed.

Clearly, such a transfer is not the most efficient in terms of cycles (compared to an asynchronous read). Nevertheless, a random asynchronous access would first require to reprogram the memory access mode, which would altogether last longer.

## Wait management

For synchronous burst NOR Flash, NWAIT is evaluated after the programmed latency period, (DATALAT+2) CLK clock cycles.

If NWAIT is sensed active (low level when WAITPOL = 0, high level when WAITPOL = 1), wait states are inserted until NWAIT is sensed inactive (high level when WAITPOL = 0, low level when WAITPOL = 1).

When NWAIT is inactive, the data is considered valid either immediately (bit WAITCFG = 1) or on the next clock edge (bit WAITCFG = 0).

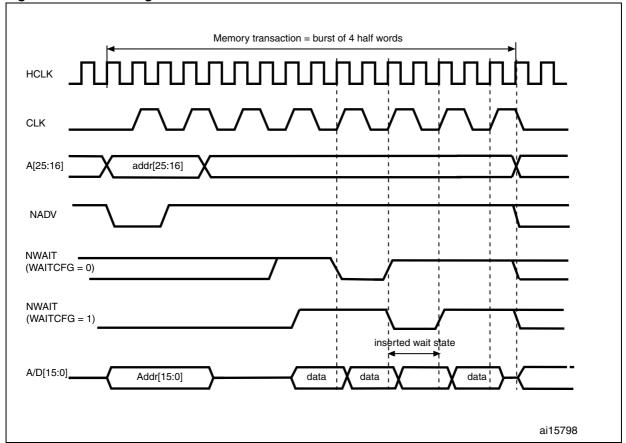
During wait-state insertion via the NWAIT signal, the controller continues to send clock pulses to the memory, keeping the chip select and output enable signals valid, and does not consider the data valid.

There are two timing configurations for the NOR Flash NWAIT signal in burst mode:

- Flash memory asserts the NWAIT signal one data cycle before the wait state (default after reset)
- Flash memory asserts the NWAIT signal during the wait state

These two NOR Flash wait state configurations are supported by the FSMC, individually for each chip select, thanks to the WAITCFG bit in the FSMC\_BCRx registers (x = 0..3).

Figure 413. Wait configurations



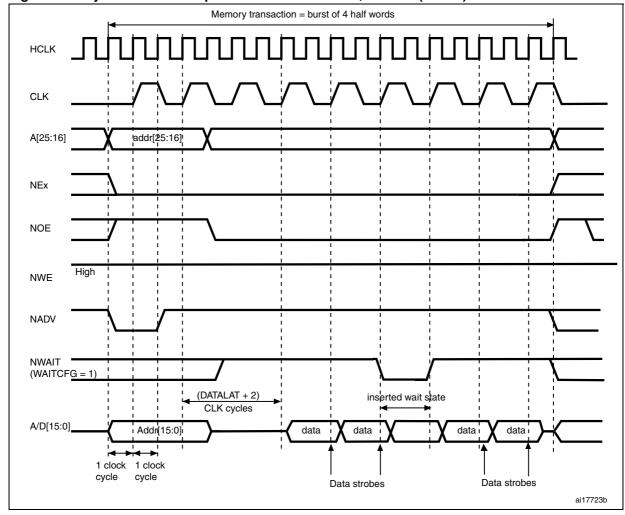


Figure 414. Synchronous multiplexed read mode - NOR, PSRAM (CRAM)

1. Byte lane outputs BL are not shown; for NOR access, they are held high, and, for PSRAM (CRAM) access, they are held low.

Table 193. FSMC\_BCRx bit fields

Bit No.	Bit name	Value to set						
31-20		0x0000						
19	CBURSTRW	No effect on synchronous read						
18-15		0x0						
14	EXTMOD	0x0						
13	WAITEN	When high, the first data after latency period is taken as always valid, regardless of the wait from memory value						
12	WREN	no effect on synchronous read						
11	WAITCFG	to be set according to memory						
10	WRAPMOD	no effect						
9	WAITPOL	to be set according to memory						

Table 193. FSMC\_BCRx bit fields (continued)

Bit No.	Bit name	Value to set
8	BURSTEN	0x1
7	FWPRLVL	Set to protect memory from accidental write access
6	FACCEN	Set according to memory support
5-4	MWID	As needed
3-2	MTYP	0x1 or 0x2
1	MUXEN	As needed
0	MBKEN	0x1

Table 194. FSMC\_BTRx bit fields

Bit No.	Bit name	Value to set
27-24	DATLAT	Data latency
23-20	CLKDIV	0x0 to get CLK = HCLK (not supported) 0x1 to get CLK = 2 × HCLK
19-16	BUSTURN	Time between NEx high to NEx low (BUSTURN HCLK)
15-8	DATAST	no effect
7-4	ADDHLD	no effect
3-0	ADDSET	no effect

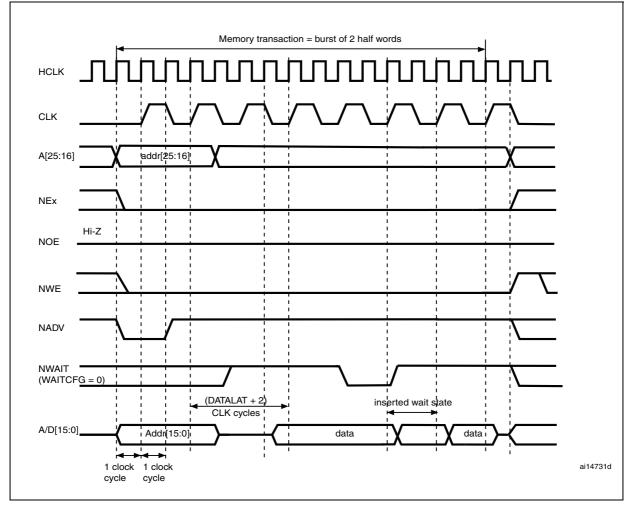


Figure 415. Synchronous multiplexed write mode - PSRAM (CRAM)

<sup>1.</sup> Memory must issue NWAIT signal one cycle in advance, accordingly WAITCFG must be programmed to 0.

<sup>2.</sup> Byte Lane (NBL) outputs are not shown, they are held low while NEx is active.

Table 195. FSMC\_BCRx bit fields

Bit No.	Bit name	Value to set
31-20		0x0000
19	CBURSTRW	0x1
18-15		0x0
14	EXTMOD	0x0
13	WAITEN	When high, the first data after latency period is taken as always valid, regardless of the wait from memory value
12	WREN	no effect on synchronous read
11	WAITCFG	0x0
10	WRAPMOD	no effect
9	WAITPOL	to be set according to memory
8	BURSTEN	no effect on synchronous write
7	FWPRLVL	Set to protect memory from accidental writes
6	FACCEN	Set according to memory support
5-4	MWID	As needed
3-2	MTYP	0x1
1	MUXEN	As needed
0	MBKEN	0x1

Table 196. FSMC\_BTRx bit fields

Bit No.	Bit name	Value to set				
31-30	-	0x0				
27-24	DATLAT	ata latency				
23-20	CLKDIV	0 to get CLK = HCLK (not supported) 1 to get CLK = 2 × HCLK				
19-16	BUSTURN	No effect				
15-8	DATAST	No effect				
7-4	ADDHLD	No effect				
3-0	ADDSET	No effect				

## 31.5.6 NOR/PSRAM control registers

The NOR/PSRAM control registers have to be accessed by words (32 bits).

## SRAM/NOR-Flash chip-select control registers 1..4 (FSMC\_BCR1..4)

Address offset:  $0xA000\ 0000 + 8*(x - 1), x = 1...4$ 

Reset value: 0x0000 30DX

This register contains the control information of each memory bank, used for SRAMs, ROMs and asynchronous or burst NOR Flash memories.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	CBURSTRW	Re	serv	ed	ASCYCWAIT	EXTMOD	WAITEN	WREN	WAITCFG	WRAPMOD	WAITPOL	BURSTEN	Reserved	FACCEN	CIVALVA	2	GVTM	_	MUXEN	MBKEN
	rw				rw	rw	rw	rw	rw	rw	rw	rw		rw	rw	rw	rw	rw	rw	rw

#### Bit 19 CBURSTRW: Write burst enable.

For Cellular RAM, the bit enables synchronous burst protocol during write operations. For Flash memory access in burst mode, this bit enables/disables the wait state insertion via the NWAIT signal. The enable bit for the synchronous burst protocol during read access is the BURSTEN bit in the FSMC\_BCRx register.

- 0: Write operations are always performed in asynchronous mode
- 1: Write operations are performed in synchronous mode.

### Bit 15 ASYNCWAIT: Wait signal during asynchronous transfers

This bit enables the FSMC to use the wait signal even during an asynchronous protocol.

- 0: NWAIT signal is not taken in to account when running an asynchronous protocol (default after reset)
- 1: NWAIT signal is taken in to account when running an asynchronous protocol

#### Bit 14 **EXTMOD:** Extended mode enable.

This bit enables the FSMC to program inside the FSMC\_BWTR register, so it allows different timings for read and write.

- 0: values inside FSMC\_BWTR register are not taken into account (default after reset)
- 1: values inside FSMC\_BWTR register are taken into account

#### Bit 13 WAITEN: Wait enable bit.

For Flash memory access in burst mode, this bit enables/disables wait-state insertion via the NWAIT signal:

- 0: NWAIT signal is disabled (its level not taken into account, no wait state inserted after the programmed Flash latency period)
- 1: NWAIT signal is enabled (its level is taken into account after the programmed Flash latency period to insert wait states if asserted) (default after reset)

### Bit 12 WREN: Write enable bit.

This bit indicates whether write operations are enabled/disabled in the bank by the FSMC:

- 0: Write operations are disabled in the bank by the FSMC, an AHB error is reported,
- 1: Write operations are enabled for the bank by the FSMC (default after reset).

### Bit 11 WAITCFG: Wait timing configuration.

For memory access in burst mode, the NWAIT signal indicates whether the data from the memory are valid or if a wait state must be inserted. This configuration bit determines if NWAIT is asserted by the memory one clock cycle before the wait state or during the wait state:

- 0: NWAIT signal is active one data cycle before wait state (default after reset),
- 1: NWAIT signal is active during wait state (not for Cellular RAM).

#### Bit 10 **WRAPMOD:** Wrapped burst mode support.

Defines whether the controller will or not split an AHB burst wrap access into two linear accesses. Valid only when accessing memories in burst mode

- 0: Direct wrapped burst is not enabled (default after reset),
- 1: Direct wrapped burst is enabled.

Note: This bit has no effect as the CPU and DMA cannot generate wrapping burst transfers.

#### Bit 9 WAITPOL: Wait signal polarity bit.

Defines the polarity of the wait signal from memory. Valid only when accessing the memory in burst mode:

- 0: NWAIT active low (default after reset),
- 1: NWAIT active high.

#### Bit 8 BURSTEN: Burst enable bit.

Enables the burst access mode for the memory. Valid only with synchronous burst memories:

- 0: Burst access mode disabled (default after reset)
- 1: Burst access mode enable
- Bit 7 Reserved, must be kept at reset value.

#### Bit 6 FACCEN: Flash access enable

Enables NOR Flash memory access operations.

- 0: Corresponding NOR Flash memory access is disabled
- 1: Corresponding NOR Flash memory access is enabled (default after reset)

### Bits 5:4 MWID: Memory databus width.

Defines the external memory device width, valid for all type of memories.

00: 8 bits,

01: 16 bits (default after reset),

10: reserved, do not use,

11: reserved, do not use.

### Bits 3:2 MTYP: Memory type.

Defines the type of external memory attached to the corresponding memory bank:

- 00: SRAM, ROM (default after reset for Bank 2...4)
- 01: PSRAM (Cellular RAM: CRAM)
- 10: NOR Flash/OneNAND Flash (default after reset for Bank 1)
- 11: reserved

### Bit 1 MUXEN: Address/data multiplexing enable bit.

When this bit is set, the address and data values are multiplexed on the databus, valid only with NOR and PSRAM memories:

- 0: Address/Data nonmultiplexed
- 1: Address/Data multiplexed on databus (default after reset)

## Bit 0 MBKEN: Memory bank enable bit.

Enables the memory bank. After reset Bank1 is enabled, all others are disabled. Accessing a disabled bank causes an ERROR on AHB bus.

- 0: Corresponding memory bank is disabled
- 1: Corresponding memory bank is enabled



## SRAM/NOR-Flash chip-select timing registers 1..4 (FSMC BTR1..4)

Address offset:  $0xA000\ 0000 + 0x04 + 8 * (x - 1), x = 1..4$ 

Reset value: 0x0FFF FFFF

This register contains the control information of each memory bank, used for SRAMs, ROMs and NOR Flash memories. If the EXTMOD bit is set in the FSMC\_BCRx register, then this register is partitioned for write and read access, that is, 2 registers are available: one to configure read accesses (this register) and one to configure write accesses (FSMC\_BWTRx registers).

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved		ACCIMOD		TA ITAC	3			2	CLNC CLNC			IGTIEN	NEOLOOG					- V	DAIASI						ADDUA			5	ADDSE	
"	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

#### Bits 29:28 ACCMOD: Access mode

Specifies the asynchronous access modes as shown in the timing diagrams. These bits are taken into account only when the EXTMOD bit in the FSMC\_BCRx register is 1.

00: access mode A 01: access mode B 10: access mode C 11: access mode D

### Bits 27:24 DATLAT: Data latency for synchronous burst NOR Flash memory

For NOR Flash with synchronous burst mode enabled, defines the number of memory clock cycles (+2) to issue to the memory before getting the first data:

0000: Data latency of 2 CLK clock cycles for first burst access

1111: Data latency of 17 CLK clock cycles for first burst access (default value after reset)

Note: This timing parameter is not expressed in HCLK periods, but in Flash clock (CLK) periods. In asynchronous NOR Flash, SRAM or ROM accesses, this value is don't care. In the case of CRAM, this field must be set to '0'.

## Bits 23:20 CLKDIV: Clock divide ratio (for CLK signal)

Defines the period of CLK clock output signal, expressed in number of HCLK cycles:

0000: Reserved

0001: CLK period =  $2 \times HCLK$  periods 0010: CLK period =  $3 \times HCLK$  periods

1111: CLK period = 16 × HCLK periods (default value after reset)

In asynchronous NOR Flash, SRAM or ROM accesses, this value is don't care.

## Bits 19:16 BUSTURN: Bus turnaround phase duration

These bits are written by software to add a delay at the end of a write/read transaction. This delay allows to match the minimum time between consecutive transactions ( $t_{EHEL}$  from NEx high to NEx low) and the maximum time needed by the memory to free the data bus after a read access ( $t_{EHQZ}$ ):

(BUSTRUN + 1)HCLK period  $\geq$  t<sub>EHELmin</sub> and (BUSTRUN + 2)HCLK period  $\geq$  t<sub>EHQZmax</sub> if EXTMOD = '0'

(BUSTRUN + 2)HCLK period  $\geq$  max ( $t_{EHELmin}$ ,  $t_{EHQZmax}$ ) if EXTMOD = '1'.

0000: BUSTURN phase duration = 0 HCLK clock cycle added

...

1111: BUSTURN phase duration = 15 × HCLK clock cycles (default value after reset)

#### Bits 15:8 DATAST: Data-phase duration

These bits are written by software to define the duration of the data phase (refer to *Figure 398* to *Figure 410*), used in SRAMs, ROMs and asynchronous NOR Flash accesses:

0000 0000: Reserved

0000 0001: DATAST phase duration = 1 × HCLK clock cycles

0000 0010: DATAST phase duration = 2 × HCLK clock cycles

...

1111 1111: DATAST phase duration =  $255 \times HCLK$  clock cycles (default value after reset) For each memory type and access mode data-phase duration, please refer to the respective figure (*Figure 398* to *Figure 410*).

Example: Mode1, write access, DATAST=1: Data-phase duration= DATAST+1 = 2 HCLK clock cycles.

Note: In synchronous accesses, this value is don't care.

### Bits 7:4 ADDHLD: Address-hold phase duration

These bits are written by software to define the duration of the *address hold* phase (refer to *Figure 407* to *Figure 410*), used in mode D and multiplexed accesses:

0000: Reserved

0001: ADDHLD phase duration =1  $\times$  HCLK clock cycle

0010: ADDHLD phase duration =  $2 \times HCLK$  clock cycle

1111: ADDHLD phase duration =  $15 \times HCLK$  clock cycles (default value after reset) For each access mode address-hold phase duration, please refer to the respective figure (*Figure 407* to *Figure 410*).

Note: In synchronous accesses, this value is not used, the address hold phase is always 1 memory clock period duration.

### Bits 3:0 ADDSET: Address setup phase duration

These bits are written by software to define the duration of the *address setup* phase (refer to *Figure 398* to *Figure 410*), used in SRAMs, ROMs and asynchronous NOR Flash accesses: 0000: ADDSET phase duration =  $0 \times HCLK$  clock cycle

• • •

1111: ADDSET phase duration =  $1615 \times HCLK$  clock cycles (default value after reset) For each access mode address setup phase duration, please refer to the respective figure (refer to *Figure 398* to *Figure 410*).

Note: In synchronous accesses, this value is don't care.

#### Note:

PSRAMs (CRAMs) have a variable latency due to internal refresh. Therefore these memories issue the NWAIT signal during the whole latency phase to prolong the latency as needed.

With PSRAMs (CRAMs) the filed DATLAT must be set to 0, so that the FSMC exits its latency phase soon and starts sampling NWAIT from memory, then starts to read or write when the memory is ready.

This method can be used also with the latest generation of synchronous Flash memories that issue the NWAIT signal, unlike older Flash memories (check the datasheet of the specific Flash memory being used).

## SRAM/NOR-Flash write timing registers 1..4 (FSMC\_BWTR1..4)

Address offset:  $0xA000\ 0000 + 0x104 + 8*(x - 1), x = 1...4$ 

Reset value: 0x0FFF FFFF



This register contains the control information of each memory bank, used for SRAMs, ROMs and NOR Flash memories. When the EXTMOD bit is set in the FSMC\_BCRx register, then this register is active for write access.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Ī	Res.	AC O			DAT	LAT			CLK	DIV		E	BUST	ruri	١				DAT	AST					ADD	HLD			ADD	SET	
		rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

#### Bits 29:28 ACCMOD: Access mode.

Specifies the asynchronous access modes as shown in the next timing diagrams. These bits are taken into account only when the EXTMOD bit in the FSMC\_BCRx register is 1.

00: access mode A 01: access mode B 10: access mode C 11: access mode D

### Bits 27:24 DATLAT: Data latency (for synchronous burst NOR Flash).

For NOR Flash with Synchronous burst mode enabled, defines the number of memory clock cycles (+2) to issue to the memory before getting the first data:

0000: (0x0) Data latency of 2 CLK clock cycles for first burst access

...

1111: (0xF) Data latency of 17 CLK clock cycles for first burst access (default value after reset)

Note: This timing parameter is not expressed in HCLK periods, but in Flash clock (**CLK**) periods. In asynchronous NOR Flash, SRAM or ROM accesses, this value is don't care. In case of CRAM, this field must be set to 0

### Bits 23:20 CLKDIV: Clock divide ratio (for CLK signal).

Defines the period of CLK clock output signal, expressed in number of HCLK cycles:

0000: Reserved

0001 CLK period =  $2 \times HCLK$  periods

0010 CLK period = 3 × HCLK periods

1111: CLK period = 16 × HCLK periods (default value after reset)

In asynchronous NOR Flash, SRAM or ROM accesses, this value is don't care.

### Bits 19:16 BUSTURN: Bus turnaround phase duration

These bits are written by software to add a delay at the end of a write transaction to match the minimum time between consecutive transactions ( $t_{\text{FHFI}}$  from ENx high to ENx low):

(BUSTRUN + 1) HCLK period ≥ t<sub>EHELmin</sub>.

0000: BUSTURN phase duration = 0 HCLK clock cycle added

...

1111: BUSTURN phase duration = 15 HCLK clock cycles added (default value after reset)

### Bits 15:8 DATAST: Data-phase duration.

These bits are written by software to define the duration of the data phase (refer to *Figure 398* to *Figure 410*), used in SRAMs, ROMs and asynchronous NOR Flash accesses:

0000 0000: Reserved

0000 0001: DATAST phase duration =  $1 \times HCLK$  clock cycles

0000 0010: DATAST phase duration = 2 × HCLK clock cycles

...

1111 1111: DATAST phase duration = 255 × HCLK clock cycles (default value after reset)

Note: In synchronous accesses, this value is don't care.

### Bits 7:4 ADDHLD: Address-hold phase duration.

These bits are written by software to define the duration of the *address hold* phase (refer to *Figure 407* to *Figure 410*), used in SRAMs, ROMs and asynchronous multiplexed NOR Flash accesses:

0000: Reserved

0001: ADDHLD phase duration =  $1 \times HCLK$  clock cycle 0010: ADDHLD phase duration =  $2 \times HCLK$  clock cycle

...

1111: ADDHLD phase duration = 15 × HCLK clock cycles (default value after reset)

Note: In synchronous NOR Flash accesses, this value is not used, the address hold phase is always 1 Flash clock period duration.

## Bits 3:0 ADDSET: Address setup phase duration.

These bits are written by software to define the duration of the *address setup* phase in HCLK cycles (refer to *Figure 407* to *Figure 410*), used in SRAMs, ROMs and asynchronous NOR Flash accessed:

0000: ADDSET phase duration =  $0 \times HCLK$  clock cycle

...

1111: ADDSET phase duration = 15 × HCLK clock cycles (default value after reset)

Note: In synchronous NOR Flash accesses, this value is don't care.

## 31.6 NAND Flash/PC Card controller

The FSMC generates the appropriate signal timings to drive the following types of device:

- NAND Flash
  - 8-bit
  - 16-bit
- 16-bit PC Card compatible devices

The NAND/PC Card controller can control three external banks. Bank 2 and bank 3 support NAND Flash devices. Bank 4 supports PC Card devices.

Each bank is configured by means of dedicated registers (*Section 31.6.8*). The programmable memory parameters include access timings (shown in *Table 197*) and ECC configuration.

Table 197. Programmable NAND/PC Card access parameters

Parameter	Function	Access mode	Unit	Min.	Max.
Memory setup time	Number of clock cycles (HCLK) to set up the address before the command assertion	Read/Write	AHB clock cycle (HCLK)	1	256
Memory wait	Minimum duration (HCLK clock cycles) of the command assertion	Read/Write	AHB clock cycle (HCLK)	2	256
Memory hold	Number of clock cycles (HCLK) to hold the address (and the data in case of a write access) after the command de-assertion	Read/Write	AHB clock cycle (HCLK)	1	255
Memory databus high-Z	Number of clock cycles (HCLK) during which the databus is kept in high-Z state after the start of a write access	Write	AHB clock cycle (HCLK)	0	255

## 31.6.1 External memory interface signals

The following tables list the signals that are typically used to interface NAND Flash and PC Card.

Note: Prefix "N". specifies the associated signal as active low.

## 8-bit NAND Flash

Table 198. 8-bit NAND Flash

FSMC signal name	I/O	Function
A[17]	0	NAND Flash address latch enable (ALE) signal
A[16]	0	NAND Flash command latch enable (CLE) signal
D[7:0]	I/O	8-bit multiplexed, bidirectional address/data bus
NCE[x]	0	Chip select, x = 2, 3
NOE(= NRE)	0	Output enable (memory signal name: read enable, NRE)
NWE	0	Write enable
NWAIT/INT[3:2]	I	NAND Flash ready/busy input signal to the FSMC

There is no theoretical capacity limitation as the FSMC can manage as many address cycles as needed.

## 16-bit NAND Flash

Table 199. 16-bit NAND Flash

FSMC signal name	I/O	Function
A[17]	0	NAND Flash address latch enable (ALE) signal
A[16]	0	NAND Flash command latch enable (CLE) signal
D[15:0]	I/O	16-bit multiplexed, bidirectional address/data bus
NCE[x]	0	Chip select, x = 2, 3
NOE(= NRE)	0	Output enable (memory signal name: read enable, NRE)
NWE	0	Write enable
NWAIT/INT[3:2]	I	NAND Flash ready/busy input signal to the FSMC

There is no theoretical capacity limitation as the FSMC can manage as many address cycles as needed.

## 16-bit PC Card

Table 200. 16-bit PC Card

FSMC signal name	I/O	Function
A[10:0]	0	Address bus
NIORD	0	Output enable for I/O space
NIOWR	0	Write enable for I/O space
NREG	0	Register signal indicating if access is in Common or Attribute space
D[15:0]	I/O	Bidirectional databus
NCE4_1	0	Chip select 1
NCE4_2	0	Chip select 2 (indicates if access is 16-bit or 8-bit)
NOE	0	Output enable in Common and in Attribute space
NWE	0	Write enable in Common and in Attribute space
NWAIT	I	PC Card wait input signal to the FSMC (memory signal name IORDY)
INTR	I	PC Card interrupt to the FSMC (only for PC Cards that can generate an interrupt)
CD	I	PC Card presence detection. Active high. If an access is performed to the PC Card banks while CD is low, an AHB error is generated. Refer to Section 31.3: AHB interface

## 31.6.2 NAND Flash / PC Card supported memories and transactions

*Table 201* below shows the supported devices, access modes and transactions. Transactions not allowed (or not supported) by the NAND Flash / PC Card controller appear in gray.

Table 201. Supported memories and transactions

Table 201.	Cupported in					
Device	Mode	R/W	AHB data size	Memory data size	Allowed/ not allowed	Comments
	Asynchronous	R	8	8	Y	
	Asynchronous	W	8	8	Y	
NAND 0 bit	Asynchronous	R	16	8	Y	Split into 2 FSMC accesses
NAND 8-bit	Asynchronous	W	16	8	Y	Split into 2 FSMC accesses
	Asynchronous	R	32	8	Υ	Split into 4 FSMC accesses
	Asynchronous	W	32	8	Υ	Split into 4 FSMC accesses
	Asynchronous	R	8	16	Υ	
	Asynchronous	W	8	16	N	
NAND 16-bit	Asynchronous	R	16	16	Υ	
INAIND 10-DIL	Asynchronous	W	16	16	Υ	
	Asynchronous	R	32	16	Y	Split into 2 FSMC accesses
	Asynchronous	W	32	16	Υ	Split into 2 FSMC accesses

## 31.6.3 Timing diagrams for NAND and PC Card

Each PC Card/CompactFlash and NAND Flash memory bank is managed through a set of registers:

Control register: FSMC\_PCRx

Interrupt status register: FSMC\_SRx

ECC register: FSMC\_ECCRx

Timing register for Common memory space: FSMC\_PMEMx

Timing register for Attribute memory space: FSMC\_PATTx

Timing register for I/O space: FSMC\_PIOx

Each timing configuration register contains three parameters used to define number of HCLK cycles for the three phases of any PC Card/CompactFlash or NAND Flash access, plus one parameter that defines the timing for starting driving the databus in the case of a write. *Figure 416* shows the timing parameter definitions for common memory accesses, knowing that Attribute and I/O (only for PC Card) memory space access timings are similar.

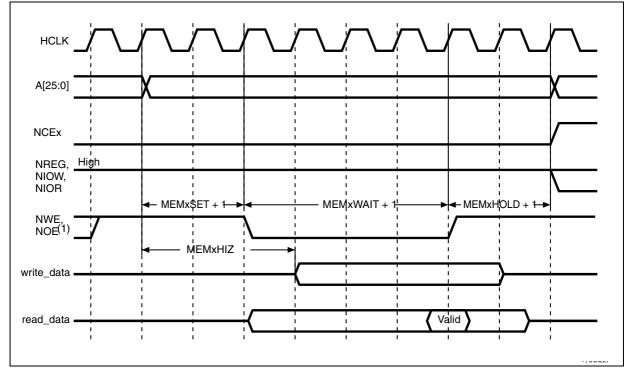


Figure 416. NAND/PC Card controller timing for common memory access

1. NOE remains high (inactive) during write access. NWE remains high (inactive) during read access.

## 31.6.4 NAND Flash operations

The command latch enable (CLE) and address latch enable (ALE) signals of the NAND Flash device are driven by some address signals of the FSMC controller. This means that to send a command or an address to the NAND Flash memory, the CPU has to perform a write to a certain address in its memory space.

A typical page read operation from the NAND Flash device is as follows:

- Program and enable the corresponding memory bank by configuring the FSMC\_PCRx and FSMC\_PMEMx (and for some devices, FSMC\_PATTx, see Section 31.6.5: NAND Flash pre-wait functionality on page 1283) registers according to the characteristics of the NAND Flash (PWID bits for the databus width of the NAND Flash, PTYP = 1, PWAITEN = 1, PBKEN = 1, see section Common memory space timing register 2..4 (FSMC\_PMEM2..4) on page 1289 for timing configuration).
- 2. The CPU performs a byte write in the common memory space, with data byte equal to one Flash command byte (for example 0x00 for Samsung NAND Flash devices). The CLE input of the NAND Flash is active during the write strobe (low pulse on NWE), thus the written byte is interpreted as a command by the NAND Flash. Once the command is latched by the NAND Flash device, it does not need to be written for the following page read operations.
- 3. The CPU can send the start address (STARTAD) for a read operation by writing the required bytes (for example four bytes or three for smaller capacity devices), STARTAD[7:0], STARTAD[15:8], STARTAD[23:16] and finally STARTAD[25:24] for 64 Mb x 8 bit NAND Flash) in the common memory or attribute space. The ALE input of the NAND Flash device is active during the write strobe (low pulse on NWE), thus the written bytes are interpreted as the start address for read operations. Using the

attribute memory space makes it possible to use a different timing configuration of the FSMC, which can be used to implement the prewait functionality needed by some NAND Flash memories (see details in *Section 31.6.5: NAND Flash pre-wait functionality on page 1283*).

- 4. The controller waits for the NAND Flash to be ready (R/NB signal high) to become active, before starting a new access (to same or another memory bank). While waiting, the controller maintains the NCE signal active (low).
- 5. The CPU can then perform byte read operations in the common memory space to read the NAND Flash page (data field + Spare field) byte by byte.
- 6. The next NAND Flash page can be read without any CPU command or address write operation, in three different ways:
  - by simply performing the operation described in step 5
  - a new random address can be accessed by restarting the operation at step 3
  - a new command can be sent to the NAND Flash device by restarting at step 2

## 31.6.5 NAND Flash pre-wait functionality

Some NAND Flash devices require that, after writing the last part of the address, the controller wait for the R/NB signal to go low as shown in *Figure 417*.

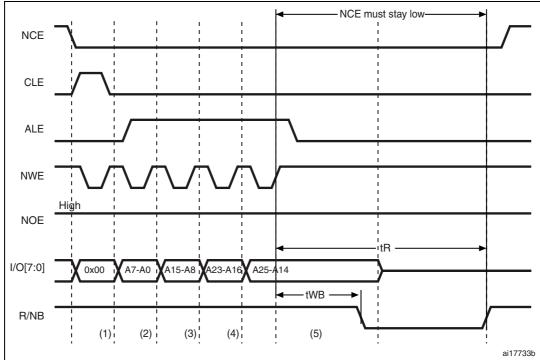


Figure 417. Access to non 'CE don't care' NAND-Flash

- 1. CPU wrote byte 0x00 at address 0x7001 0000.
- 2. CPU wrote byte A7-A0 at address 0x7002 0000.
- 3. CPU wrote byte A15-A8 at address 0x7002 0000.
- 4. CPU wrote byte A23-A16 at address 0x7002 0000.
- 5. CPU wrote byte A25-A24 at address 0x7802 0000: FSMC performs a write access using FSMC\_PATT2 timing definition, where ATTHOLD ≥ 7 (providing that (7+1) × HCLK = 112 ns > t<sub>WB</sub> max). This guarantees that NCE remains low until R/NB goes low and high again (only requested for NAND Flash memories where NCE is not don't care).

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When this functionality is needed, it can be guaranteed by programming the MEMHOLD value to meet the  $t_{WB}$  timing, however any CPU read or write access to the NAND Flash then has the hold delay of (MEMHOLD + 1) HCLK cycles inserted from the rising edge of the NWE signal to the next access.

To overcome this timing constraint, the attribute memory space can be used by programming its timing register with an ATTHOLD value that meets the  $t_{WB}$  timing, and leaving the MEMHOLD value at its minimum. Then, the CPU must use the common memory space for all NAND Flash read and write accesses, except when writing the last address byte to the NAND Flash device, where the CPU must write to the attribute memory space.

## 31.6.6 Error correction code computation ECC (NAND Flash)

The FSMC PC-Card controller includes two error correction code computation hardware blocks, one per memory bank. They are used to reduce the host CPU workload when processing the error correction code by software in the system.

These two registers are identical and associated with bank 2 and bank 3, respectively. As a consequence, no hardware ECC computation is available for memories connected to bank 4

The error correction code (ECC) algorithm implemented in the FSMC can perform 1-bit error correction and 2-bit error detection per 256, 512, 1 024, 2 048, 4 096 or 8 192 bytes read from or written to NAND Flash.

The ECC modules monitor the NAND Flash databus and read/write signals (NCE and NWE) each time the NAND Flash memory bank is active.

The functional operations are:

- When access to NAND Flash is made to bank 2 or bank 3, the data present on the D[15:0] bus is latched and used for ECC computation.
- When access to NAND Flash occurs at any other address, the ECC logic is idle, and does not perform any operation. Thus, write operations for defining commands or addresses to NAND Flash are not taken into account for ECC computation.

Once the desired number of bytes has been read from/written to the NAND Flash by the host CPU, the FSMC\_ECCR2/3 registers must be read in order to retrieve the computed value. Once read, they should be cleared by resetting the ECCEN bit to zero. To compute a new data block, the ECCEN bit must be set to one in the FSMC\_PCR2/3 registers.

## 31.6.7 PC Card/CompactFlash operations

#### Address spaces & memory accesses

The FSMC supports Compact Flash storage or PC Cards in Memory Mode and I/O Mode (True IDE mode is not supported).

The Compact Flash storage and PC Cards are made of 3 memory spaces:

- Common Memory Space
- Attribute Space
- I/O Memory Space

The nCE2 and nCE1 pins (FSMC\_NCE4\_2 and FSMC\_NCE4\_1 respectively) select the card and indicate whether a byte or a word operation is being performed: nCE2 accesses

the odd byte on D15-8 and nCE1 accesses the even byte on D7-0 if A0=0 or the odd byte on D7-0 if A0=1. The full word is accessed on D15-0 if both nCE2 and nCE1 are low.

The memory space is selected by asserting low nOE for read accesses or nWE for write accesses, combined with the low assertion of nCE2/nCE1 and nREG.

- If pin nREG=1 during the memory access, the common memory space is selected
- If pin nREG=0 during the memory access, the attribute memory space is selected

The I/O Space is selected by asserting low nIORD for read accesses or nIOWR for write accesses [instead of nOE/nWE for memory Space], combined with nCE2/nCE1. Note that nREG must also be asserted low during accesses to I/O Space.

Three type of accesses are allowed for a 16-bit PC Card:

- Accesses to Common Memory Space for data storage can be either 8-bit accesses at even addresses or 16 bit AHB accesses.
  - Note that 8-bit accesses at odd addresses are not supported and will not lead to the low assertion of nCE2. A 32-bit AHB request is translated into two 16-bit memory accesses.
- Accesses to Attribute Memory Space where the PC Card stores configuration information are limited to 8-bit AHB accesses at even addresses.
  - Note that a 16-bit AHB access will be converted into a single 8-bit memory transfer: nCE1 will be asserted low, NCE2 will be asserted high and only the even Byte on D7-D0 will be valid. Instead a 32-bit AHB access will be converted into two 8-bit memory transfers at even addresses: nCE1 will be asserted low, NCE2 will be asserted high and only the even bytes will be valid.
- Accesses to I/O Space can be performed either through AHB 8-bit or 16-bit accesses.

Table 202. 16-bit PC-Card signals and access type

nCE2	nCE1	nREG	nOE/nWE	nIORD /nIOWR	A10	6 <b>V</b>	A7-1	0 <b>V</b>	Space	Access Type	Allowed/not Allowed
1	0	1	0	1	Х	Х	X-X	Χ	Common	Read/Write byte on D7-D0	YES
0	1	1	0	1	Χ	Χ	X-X	Χ	Memory	Read/Write byte on D15-D8	Not supported
0	0	1	0	1	Χ	Χ	X-X	0	Space	Read/Write word on D15-D0	YES
Х	0	0	0	1	0	1	X-X	0	Attribute	Read or Write Configuration Registers	YES
Х	0	0	0	1	0	0	X-X	0	Space	Read or Write CIS (Card Information Structure)	YES
1	0	0	0	1	Х	Х	X-X	1	Invalid	Read or Write (odd address)	YES
0	1	0	0	1	Х	Х	X-X	х	Attribute Space	Read or Write (odd address)	YES

nCE2	nCE1	nREG	nOE/nWE	nIORD /nIOWR	A10	6 <b>V</b>	A7-1	0 <b>V</b>	Space	Access Type	Allowed/not Allowed
1	0	0	1	0	Х	Х	X-X	0		Read Even Byte on D7-0	YES
1	0	0	1	0	Х	Х	X-X	1		Read Odd Byte on D7-0	YES
1	0	0	1	0	Х	Х	X-X	0		Write Even Byte on D7-0	YES
1	0	0	1	0	Х	Х	X-X	1	I/O anaga	Write Odd Byte on D7-0	YES
0	0	0	1	0	Х	Х	X-X	0	I/O space	Read Word on D15-0	YES
0	0	0	1	0	Х	Х	X-X	0		Write word on D15-0	YES
0	1	0	1	0	Х	Х	X-X	Χ		Read Odd Byte on D15-8	Not supported
0	1	0	1	0	Х	Х	X-X	Χ		Write Odd Byte on D15-8	Not supported

Table 202. 16-bit PC-Card signals and access type (continued)

The FSMC Bank 4 gives access to those 3 memory spaces as described in *Section 31.4.2:* NAND/PC Card address mapping - Table 169: Memory mapping and timing registers

### **Wait Feature**

The CompactFlash Storage or PC Card may request the FSMC to extend the length of the access phase programmed by MEMWAITx/ATTWAITx/IOWAITx bits, asserting the nWAIT signal after nOE/nWE or nIORD/nIOWR activation if the wait feature is enabled through the PWAITEN bit in the FSMC\_PCRx register. In order to detect the nWAIT assertion correctly, the MEMWAITx/ATTWAITx/IOWAITx bits must be programmed as follows:

Where  $max\_wait\_assertion\_time$  is the maximum time taken by NWAIT to go low once nOE/nWE or nIORD/nIOWR is low.

After the de-assertion of nWAIT, the FSMC extends the WAIT phase for 4 HCLK clock cycles.

## 31.6.8 NAND Flash/PC Card control registers

The NAND Flash/PC Card control registers have to be accessed by words (32 bits).

## PC Card/NAND Flash control registers 2..4 (FSMC\_PCR2..4)

Address offset: 0xA0000000 + 0x40 + 0x20 \* (x - 1), x = 2..4

Reset value: 0x0000 0018

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Reserved	EC	CCP	S		TA	ιR			TC	LR		Res.	ECCEN	PV	/ID	PTYP	PBKEN	PWAITEN	Reserved
	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw		rw	rw	rw	rw	rw	rw	Ä

Bits 19:17 ECCPS: ECC page size.

Defines the page size for the extended ECC:

000: 256 bytes 001: 512 bytes 010: 1024 bytes 011: 2048 bytes 100: 4096 bytes 101: 8192 bytes

Bits 16:13 TAR: ALE to RE delay.

Sets time from ALE low to RE low in number of AHB clock cycles (HCLK).

Time is:  $t_ar = (TAR + SET + 2) \times THCLK$  where THCLK is the HCLK clock period

0000: 1 HCLK cycle (default) 1111: 16 HCLK cycles

Note: SET is MEMSET or ATTSET according to the addressed space.

Bits 12:9 TCLR: CLE to RE delay.

Sets time from CLE low to RE low in number of AHB clock cycles (HCLK).

Time is  $t_{clr} = (TCLR + SET + 2) \times THCLK$  where THCLK is the HCLK clock period

0000: 1 HCLK cycle (default) 1111: 16 HCLK cycles

Note: SET is MEMSET or ATTSET according to the addressed space.

Bits 8:7 Reserved, must be kept at reset value.

Bits 6 ECCEN: ECC computation logic enable bit

0: ECC logic is disabled and reset (default after reset),

1: ECC logic is enabled.

Bits 5:4 PWID: Databus width.

Defines the external memory device width.

00: 8 bits (default after reset)

01: 16 bits (mandatory for PC Card)

10: reserved, do not use11: reserved, do not use

Bit 3 PTYP: Memory type.

Defines the type of device attached to the corresponding memory bank:

0: PC Card, CompactFlash, CF+ or PCMCIA

1: NAND Flash (default after reset)

Bit 2 PBKEN: PC Card/NAND Flash memory bank enable bit.

Enables the memory bank. Accessing a disabled memory bank causes an ERROR on AHB bus

- 0: Corresponding memory bank is disabled (default after reset)
- 1: Corresponding memory bank is enabled
- Bit 1 **PWAITEN:** Wait feature enable bit.

Enables the Wait feature for the PC Card/NAND Flash memory bank:

0: disabled

1: enabled

Note: For a PC Card, when the wait feature is enabled, the MEMWAITx/ATTWAITx/IOWAITx bits must be programmed to a value as follows:

xxWAITx ≥ 4 + max\_wait\_assertion\_time/HCLK

Where max\_wait\_assertion\_time is the maximum time taken by NWAIT to go low once nOE/nWE or nIORD/nIOWR is low.

Bit 0 Reserved, must be kept at reset value.

## FIFO status and interrupt register 2..4 (FSMC\_SR2..4)

Address offset:  $0xA000\ 0000 + 0x44 + 0x20 * (x-1), x = 2..4$ 

Reset value: 0x0000 0040

This register contains information about FIFO status and interrupt. The FSMC has a FIFO that is used when writing to memories to store up to 16 words of data from the AHB.

This is used to quickly write to the AHB and free it for transactions to peripherals other than the FSMC, while the FSMC is draining its FIFO into the memory. This register has one of its bits that indicates the status of the FIFO, for ECC purposes.

The ECC is calculated while the data are written to the memory, so in order to read the correct ECC the software must wait until the FIFO is empty.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
												served													FEMPT	IFEN	NETI	IREN	IFS	STI	IRS
												æ													r	rw	rw	rw	rw	rw	rw

Bit 6 FEMPT: FIFO empty.

Read-only bit that provides the status of the FIFO

0: FIFO not empty

1: FIFO empty

Bit 5 IFEN: Interrupt falling edge detection enable bit

- 0: Interrupt falling edge detection request disabled
- 1: Interrupt falling edge detection request enabled
- Bit 4 ILEN: Interrupt high-level detection enable bit
  - 0: Interrupt high-level detection request disabled
  - 1: Interrupt high-level detection request enabled
- Bit 3 IREN: Interrupt rising edge detection enable bit
  - 0: Interrupt rising edge detection request disabled
  - 1: Interrupt rising edge detection request enabled

Bit 2 IFS: Interrupt falling edge status

The flag is set by hardware and reset by software.

0: No interrupt falling edge occurred

1: Interrupt falling edge occurred

Bit 1 ILS: Interrupt high-level status

The flag is set by hardware and reset by software.

0: No Interrupt high-level occurred

1: Interrupt high-level occurred

Bit 0 IRS: Interrupt rising edge status

The flag is set by hardware and reset by software.

0: No interrupt rising edge occurred

1: Interrupt rising edge occurred

## Common memory space timing register 2..4 (FSMC PMEM2..4)

Address offset: Address:  $0xA000\ 0000 + 0x48 + 0x20 * (x - 1), x = 2..4$ 

Reset value: 0xFCFC FCFC

Each FSMC\_PMEMx (x = 2..4) read/write register contains the timing information for PC Card or NAND Flash memory bank x, used for access to the common memory space of the 16-bit PC Card/CompactFlash, or to access the NAND Flash for command, address write access and data read/write access.

#### 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

	MEMHIZx									М	EMH	IOLE	Эx					N	1EM\	WAIT	x					ľ	ИЕМ	SET:	x		
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw

### Bits 31:24 MEMHIZx: Common memory x databus HiZ time

Defines the number of HCLK clock cycles during which the databus is kept in HiZ after the start of a PC Card/NAND Flash write access to common memory space on socket x. Only valid for write transaction:

0000 0000: (0x00) 0 HCLK cycle (for PC Card)

1111 1111: (0xFF) 255 HCLK cycles (for PC Card) - (default value after reset)

#### Bits 23:16 **MEMHOLDx:** Common memory x hold time

Defines the number of HCLK clock cycles to hold address (and data for write access) after the command deassertion (NWE, NOE), for PC Card/NAND Flash read or write access to common memory space on socket x:

0000 0000: reserved 0000 0001: 1 HCLK cycle

1111 1111: 255 HCLK cycles (default value after reset)

#### Bits 15:8 **MEMWAITx:** Common memory x wait time

Defines the minimum number of HCLK (+1) clock cycles to assert the command (NWE, NOE), for PC Card/NAND Flash read or write access to common memory space on socket x. The duration for command assertion is extended if the wait signal (NWAIT) is active (low) at the end of the programmed value of HCLK:

0000 0000: reserved

0000 0001: 2HCLK cycles (+ wait cycle introduced by deasserting NWAIT)

1111 1111: 256 HCLK cycles (+ wait cycle introduced by the Card deasserting NWAIT)

(default value after reset)

#### Bits 7:0 MEMSETx: Common memory x setup time

Defines the number of HCLK () clock cycles to set up the address before the command assertion (NWE, NOE), for PC Card/NAND Flash read or write access to common memory space on socket x:

0000 0000: 1 HCLK cycle (for PC Card) / HCLK cycles (for NAND Flash)

1111 1111: 256 HCLK cycles (for PC Card) / 257 HCLK cycles (for NAND Flash) - (default

value after reset)

## Attribute memory space timing registers 2..4 (FSMC\_PATT2..4)

Address offset:  $0xA000\ 0000 + 0x4C + 0x20 * (x - 1), x = 2..4$ 

Reset value: 0xFCFC FCFC

Each FSMC\_PATTx (x = 2..4) read/write register contains the timing information for PC Card/CompactFlash or NAND Flash memory bank x. It is used for 8-bit accesses to the attribute memory space of the PC Card/CompactFlash or to access the NAND Flash for the last address write access if the timing must differ from that of previous accesses (for Ready/Busy management, refer to Section 31.6.5: NAND Flash pre-wait functionality).

#### 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

			ATT	HIZx						Α	TTH	OLD	х			ATTWAITx								ATTSETx									
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw		

#### Bits 31:24 ATTHIZx: Attribute memory x databus HiZ time

Defines the number of HCLK clock cycles during which the databus is kept in HiZ after the start of a PC CARD/NAND Flash write access to attribute memory space on socket x. Only valid for write transaction:

0000 0000: 0 HCLK cycle

1111 1111: 255 HCLK cycles (default value after reset)

### Bits 23:16 **ATTHOLDx:** Attribute memory x hold time

Defines the number of HCLK clock cycles to hold address (and data for write access) after the command deassertion (NWE, NOE), for PC Card/NAND Flash read or write access to attribute memory space on socket x

0000 0000: reserved 0000 0001: 1 HCLK cycle

1111 1111: 255 HCLK cycles (default value after reset)

#### Bits 15:8 ATTWAITx: Attribute memory x wait time

Defines the minimum number of HCLK (+1) clock cycles to assert the command (NWE, NOE), for PC Card/NAND Flash read or write access to attribute memory space on socket x. The duration for command assertion is extended if the wait signal (NWAIT) is active (low) at the end of the programmed value of HCLK:

0000 0000: reserved

0000 0001: 2 HCLK cycles (+ wait cycle introduced by deassertion of NWAIT)

1111 1111: 256 HCLK cycles (+ wait cycle introduced by the card deasserting NWAIT)

(default value after reset)

#### Bits 7:0 ATTSETx: Attribute memory x setup time

Defines the number of HCLK (+1) clock cycles to set up address before the command assertion (NWE, NOE), for PC CARD/NAND Flash read or write access to attribute memory space on socket x:

0000 0000: 1 HCLK cycle

1111 1111: 256 HCLK cycles (default value after reset)

## I/O space timing register 4 (FSMC\_PIO4)

Address offset: 0xA000 0000 + 0xB0

Reset value: 0xFCFCFCFC

The FSMC\_PIO4 read/write registers contain the timing information used to gain access to the I/O space of the 16-bit PC Card/CompactFlash.

#### 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

IOHIZx									l	OHO	DLDx	(			IOWAITx								IOSETx									
rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	rw	

### Bits 31:24 IOHIZx: I/O x databus HiZ time

Defines the number of HCLK clock cycles during which the databus is kept in HiZ after the start of a PC Card write access to I/O space on socket x. Only valid for write transaction:

0000 0000: 0 HCLK cycle

1111 1111: 255 HCLK cycles (default value after reset)

### Bits 23:16 IOHOLDx: I/O x hold time

Defines the number of HCLK clock cycles to hold address (and data for write access) after the command deassertion (NWE, NOE), for PC Card read or write access to I/O space on socket x:

0000 0000: reserved 0000 0001: 1 HCLK cycle

1111 1111: 255 HCLK cycles (default value after reset)

### Bits 15:8 IOWAITx: I/O x wait time

Defines the minimum number of HCLK (+1) clock cycles to assert the command (SMNWE, SMNOE), for PC Card read or write access to I/O space on socket x. The duration for command assertion is extended if the wait signal (NWAIT) is active (low) at the end of the programmed value of HCLK:

0000 0000: reserved, do not use this value

0000 0001: 2 HCLK cycles (+ wait cycle introduced by deassertion of NWAIT)

1111 1111: 256 HCLK cycles (+ wait cycle introduced by the Card deasserting NWAIT)

(default value after reset)

## Bits 7:0 IOSETx: I/O x setup time

Defines the number of HCLK (+1) clock cycles to set up the address before the command assertion (NWE, NOE), for PC Card read or write access to I/O space on socket x:

0000 0000: 1 HCLK cycle

1111 1111: 256 HCLK cycles (default value after reset)

## ECC result registers 2/3 (FSMC\_ECCR2/3)

Address offset:  $0xA000\ 0000 + 0x54 + 0x20 * (x - 1), x = 2 \text{ or } 3$ 

Reset value: 0x0000 0000

These registers contain the current error correction code value computed by the ECC computation modules of the FSMC controller (one module per NAND Flash memory bank). When the CPU reads the data from a NAND Flash memory page at the correct address (refer to Section 31.6.6: Error correction code computation ECC (NAND Flash)), the data read from or written to the NAND Flash are processed automatically by ECC computation module. At the end of X bytes read (according to the ECCPS field in the FSMC\_PCRx registers), the CPU must read the computed ECC value from the FSMC\_ECCx registers, and then verify whether these computed parity data are the same as the parity value recorded in the spare area, to determine whether a page is valid, and, to correct it if applicable. The FSMC\_ECCRx registers should be cleared after being read by setting the ECCEN bit to zero. For computing a new data block, the ECCEN bit must be set to one.

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

		ECCx			
		r			

Bits 31:0 ECCx: ECC result

This field provides the value computed by the ECC computation logic. *Table 203* hereafter describes the contents of these bit fields.

Table 203. ECC result relevant bits

ECCPS[2:0]	Page size in bytes	ECC bits
000	256	ECC[21:0]
001	512	ECC[23:0]
010	1024	ECC[25:0]
011	2048	ECC[27:0]
100	4096	ECC[29:0]
101	8192	ECC[31:0]

# 31.6.9 FSMC register map

The following table summarizes the FSMC registers.

Table 204. FSMC register map

Table 2				<del></del>				1		1		1					_					- 1	- 1		_	—
Offset	Register	31 30	29	27	25	23	20	19	18	16	15	14	13	12	7	10	6	8	7	9	2	4	က	7	-	0
0xA000 0000	FSMC_BCR1			Reser	ved			CBURSTRW	Reser	ved	ASYNCWAIT	EXTMOD	WAITEN	WREN	WAITCFG	WRAPMOD	WAITPOL	BURSTEN	Reserved	FACCEN	MMID		MTYP		MUXEN	MBKEN
0xA000 0008	FSMC_BCR2			Reser	ved			CBURSTRW	Reser	ved	ASYNCWAIT ASYNCWAIT	EXTMOD	WAITEN	WREN	WAITCFG	WRAPMOD	WAITPOL	BURSTEN	Reserved	FACCEN	QIMM		MTYP		MUXEN	MBKEN
0xA000 0010	FSMC_BCR3			Reser	ved			CBURSTRW	Reser		ASYNCWAIT	EXTMOD	WAITEN	WREN	WAITCFG	WRAPMOD	WAITPOL	BURSTEN	Reserved	FACCEN	MWID		MTYP	TALLOS IV.	MUXEN	MBKEN
0xA000 0018	FSMC_BCR4			Reser	ved			CBURSTRW	Reser	ved	ASYNCWAIT	EXTMOD	WAITEN	WREN	WAITCFG	WRAPMOD	WAITPOL	BURSTEN	Reserved	FACCEN	MWID		MTYP	14122114	MUXEN	MBKEN
0xA000 0004	FSMC_BTR1	Res.	ACCM OD	DAIL	AT	CLI	KDIV	В	USTUF	RN			ļ	DAT	AST				A	ADD	HLC	)	Α	DDS	3ET	
0xA000 000C	FSMC_BTR2	Res.	ACCM OD	DATL	AT	CLI	KDIV	В	USTUF	RN	DATAST								Α	ADDHLD				ADDSI		
0xA000 0014	FSMC_BTR3	Res.	ACCM OD	DATL	AT	CLI	KDIV	В	BUSTURN					DAT	AST				Α	ADD	HLD	)	Α	DDS	3ET	
0xA000 001C	FSMC_BTR4	Res.	ACCM OD	DATL	AT	CL	KDIV	В	USTUF	RN			ļ	DAT	AST				Α	ADD	HLC	)	Α	DDS	3ET	
0xA000 0104	FSMC_BWTR1	Res.	ACCM OD	DATL	AT	CLI	KDIV	В	USTUF	RN				DAT	AST				Α	ADD	HLD	)	Α	DDS	3ET	
0xA000 010C	FSMC_BWTR2	Res.	ACCM OD	DATL	AT	CLI	KDIV	В	USTUF	RN				DAT	AST				Α	ADD	HLC	)	Α	DDS	3ET	
0xA000 0114	FSMC_BWTR3	Res.	ACCM OD	DATL	AT	CLI	KDIV	BUSTURN					ı	DAT	AST				A	ADD	HLD	)	Α	DDS	3ET	
0xA000 011C	FSMC_BWTR4	Res.	ACCM OD	DATL	AT	CLI	KDIV	BUSTURN					I	DAT	AST	•			A	ADD	HLD	)	Α	DDS	3ET	
0xA000 0060	FSMC_PCR2			Reser	ved			E	CCPS		TA	λR			тс	LR		Re	es.	ECCEN	PW	/ID	РТҮР			Reserved
0xA000 0080	FSMC_PCR3			Reser	ved			E	CCPS		TA	AR			тс	LR		Re	es.	ECCEN	PW	/ID	РТУР		_	Reserved
0xA000 00A0	FSMC_PCR4			Reser	ved		E	CCPS		TAR T					TCLR F				ECCEN	PW	/ID	РТҮР	PBKEN	PWAITEN	Reserved	
0xA000 0064	FSMC_SR2				R	eser	ved		·									ч	IFEN	ILEN	IREN	IFS	ILS	IRS		
0xA000 0084	FSMC_SR3						R	eser	ved											ш	IFEN	ILEN	IREN	IFS	ILS	IRS
0xA000 00A4	FSMC_SR4						R	eser	ved											FEMPT	IFEN	ILEN	IREN	ES C	ILS	IRS
0xA000 0068	FSMC_PMEM2		MEN	ИНІZх			MEM	HOL	Dx	_	MEMWAITx									_	М	EMS	SET:	х		
0xA000 0088	FSMC_PMEM3		MEN	ИНІZх	MEM	HOL	Dx	MEMWAITx								MEMSETx										

Table 204. FSMC register map (continued)

Offset	Register	31 30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	6	8	7	9	5	4	3	2	1	0			
0xA000 00A8	FSMC_PMEM4		М	EMH	ΗZx	(					N	ИЕМI	HOL	Dx					МІ	ΞM\	NAI <sup>-</sup>	Тх					N	/ΕΜ	SE	Tx					
0xA000 006C	FSMC_PATT2		Α	ТТН	IZx						,	ATTH	IOLE	Эх					A	ГТV	VAIT	x					,	ATTS	SET	Гх					
0xA000 008C	FSMC_PATT3	ATTHIZx								ATTHOLDx									ATTWAITx								ATTSETx								
0xA000 00AC	FSMC_PATT4	ATTHIZx							ATTHOLDx									ATTWAITx								ATTSETx									
0xA000 00B0	FSMC_PIO4		IOHIZx						IOHOLDx									IOWAITx									IOSETx								
0xA000 0074	FSMC_ECCR2								ECCx																										
0xA000 0094	FSMC_ECCR3									ECCx																									

Refer to *Table 2 on page 50* for the register boundary addresses.

# 32 Debug support (DBG)

This section applies to the whole STM32F4xx family, unless otherwise specified.

## 32.1 Overview

The STM32F4xx are built around a Cortex<sup>™</sup>-M4F core which contains hardware extensions for advanced debugging features. The debug extensions allow the core to be stopped either on a given instruction fetch (breakpoint) or data access (watchpoint). When stopped, the core's internal state and the system's external state may be examined. Once examination is complete, the core and the system may be restored and program execution resumed.

The debug features are used by the debugger host when connecting to and debugging the STM32F4xx MCUs.

Two interfaces for debug are available:

- Serial wire
- JTAG debug port

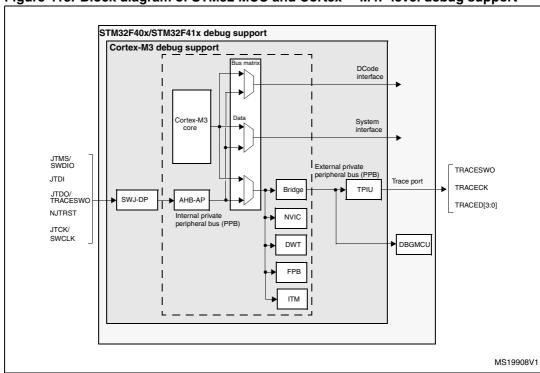


Figure 418. Block diagram of STM32 MCU and Cortex™-M4F-level debug support

Note:

The debug features embedded in the Cortex™-M4F core are a subset of the ARM CoreSight Design Kit.

The ARM Cortex<sup>™</sup>-M4F core provides integrated on-chip debug support. It is comprised of:

- SWJ-DP: Serial wire / JTAG debug port
- AHP-AP: AHB access port
- ITM: Instrumentation trace macrocell
- FPB: Flash patch breakpoint
- DWT: Data watchpoint trigger
- TPUI: Trace port unit interface (available on larger packages, where the corresponding pins are mapped)
- ETM: Embedded Trace Macrocell (available on larger packages, where the corresponding pins are mapped)

It also includes debug features dedicated to the STM32F4xx:

- Flexible debug pinout assignment
- MCU debug box (support for low-power modes, control over peripheral clocks, etc.)

Note:

For further information on debug functionality supported by the ARM Cortex<sup>™</sup>-M4F core, refer to the Cortex<sup>™</sup>-M4F-r0p1 Technical Reference Manual and to the CoreSight Design Kit-r0p1 TRM (see Section 32.2: Reference ARM documentation).

## 32.2 Reference ARM documentation

- Cortex<sup>™</sup>-M4F r0p1 Technical Reference Manual (TRM) (see Related documents on page 1)
- ARM Debug Interface V5
- ARM CoreSight Design Kit revision r0p1 Technical Reference Manual

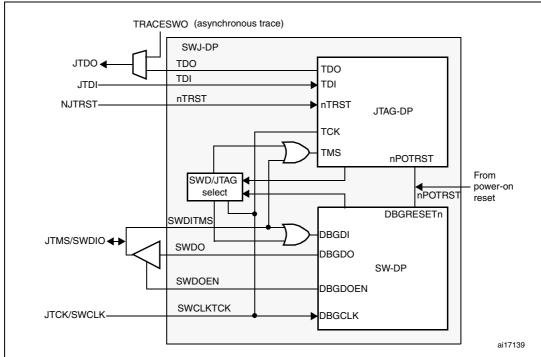
# 32.3 SWJ debug port (serial wire and JTAG)

The STM32F4xx core integrates the Serial Wire / JTAG Debug Port (SWJ-DP). It is an ARM standard CoreSight debug port that combines a JTAG-DP (5-pin) interface and a SW-DP (2-pin) interface.

- The JTAG Debug Port (JTAG-DP) provides a 5-pin standard JTAG interface to the AHP-AP port.
- The Serial Wire Debug Port (SW-DP) provides a 2-pin (clock + data) interface to the AHP-AP port.

In the SWJ-DP, the two JTAG pins of the SW-DP are multiplexed with some of the five JTAG pins of the JTAG-DP.

Debug support (DBG) RM0090



#### Figure 419. SWJ debug port

*Figure 419* shows that the asynchronous TRACE output (TRACESWO) is multiplexed with TDO. This means that the asynchronous trace can only be used with SW-DP, not JTAG-DP.

#### 32.3.1 Mechanism to select the JTAG-DP or the SW-DP

By default, the JTAG-Debug Port is active.

If the debugger host wants to switch to the SW-DP, it must provide a dedicated JTAG sequence on TMS/TCK (respectively mapped to SWDIO and SWCLK) which disables the JTAG-DP and enables the SW-DP. This way it is possible to activate the SWDP using only the SWCLK and SWDIO pins.

This sequence is:

- 1. Send more than 50 TCK cycles with TMS (SWDIO) =1
- 2. Send the 16-bit sequence on TMS (SWDIO) = 0111100111100111 (MSB transmitted first)
- 3. Send more than 50 TCK cycles with TMS (SWDIO) =1

# 32.4 Pinout and debug port pins

The STM32F4xx MCUs are available in various packages with different numbers of available pins. As a result, some functionality (ETM) related to pin availability may differ between packages.

## 32.4.1 SWJ debug port pins

Five pins are used as outputs from the STM32F4xx for the SWJ-DP as *alternate functions* of general-purpose I/Os. These pins are available on all packages.

Table 205. SWJ debug port pins

SWJ-DP pin name	,	JTAG debug port		Pin	
Swo-DP pili flame	Type Description		Туре	Debug assignment	assign ment
JTMS/SWDIO	I	JTAG Test Mode Selection	Ю	Serial Wire Data Input/Output	PA13
JTCK/SWCLK	I	JTAG Test Clock		Serial Wire Clock	PA14
JTDI	I	JTAG Test Data Input	-	-	PA15
JTDO/TRACESWO	0	JTAG Test Data Output	-	TRACESWO if async trace is enabled	РВ3
NJTRST	I	JTAG Test nReset	-	-	PB4

## 32.4.2 Flexible SWJ-DP pin assignment

After RESET (SYSRESETn or PORESETn), all five pins used for the SWJ-DP are assigned as dedicated pins immediately usable by the debugger host (note that the trace outputs are not assigned except if explicitly programmed by the debugger host).

However, the STM32F4xx MCUs offer the possibility of disabling some or all of the SWJ-DP ports and so, of releasing the associated pins for general-purpose IO (GPIO) usage. For more details on how to disable SWJ-DP port pins, please refer to *Section 6.3.2: I/O pin multiplexer and mapping*.

Table 206. Flexible SWJ-DP pin assignment

	SWJ IO pin assigned						
Available debug ports	PA13 / JTMS/ SWDIO	PA14 / JTCK/ SWCLK	PA15 / JTDI	PB3 / JTDO	PB4/ NJTRST		
Full SWJ (JTAG-DP + SW-DP) - Reset State	Χ	Х	Χ	Χ	Х		
Full SWJ (JTAG-DP + SW-DP) but without NJTRST	Х	Х	Х	Х			
JTAG-DP Disabled and SW-DP Enabled	Х	Х			_		
JTAG-DP Disabled and SW-DP Disabled			Rele	ased			

Note:

When the APB bridge write buffer is full, it takes one extra APB cycle when writing the GPIO\_AFR register. This is because the deactivation of the JTAGSW pins is done in two cycles to guarantee a clean level on the nTRST and TCK input signals of the core.

- Cycle 1: the JTAGSW input signals to the core are tied to 1 or 0 (to 1 for nTRST, TDI and TMS, to 0 for TCK)
- Cycle 2: the GPIO controller takes the control signals of the SWJTAG IO pins (like controls of direction, pull-up/down, Schmitt trigger activation, etc.).

## 32.4.3 Internal pull-up and pull-down on JTAG pins

It is necessary to ensure that the JTAG input pins are not floating since they are directly connected to flip-flops to control the debug mode features. Special care must be taken with the SWCLK/TCK pin which is directly connected to the clock of some of these flip-flops.

To avoid any uncontrolled IO levels, the device embeds internal pull-ups and pull-downs on the JTAG input pins:

NJTRST: Internal pull-upJTDI: Internal pull-up

JTMS/SWDIO: Internal pull-upTCK/SWCLK: Internal pull-down

Once a JTAG IO is released by the user software, the GPIO controller takes control again. The reset states of the GPIO control registers put the I/Os in the equivalent state:

NJTRST: AF input pull-up

• JTDI: AF input pull-up

JTMS/SWDIO: AF input pull-upJTCK/SWCLK: AF input pull-down

JTDO: AF output floating

The software can then use these I/Os as standard GPIOs.

Note:

The JTAG IEEE standard recommends to add pull-ups on TDI, TMS and nTRST but there is no special recommendation for TCK. However, for JTCK, the device needs an integrated pull-down.

Having embedded pull-ups and pull-downs removes the need to add external resistors.

## 32.4.4 Using serial wire and releasing the unused debug pins as GPIOs

To use the serial wire DP to release some GPIOs, the user software must change the GPIO (PA15, PB3 and PB4) configuration mode in the GPIO\_MODER register. This releases PA15. PB3 and PB4 which now become available as GPIOs.

When debugging, the host performs the following actions:

- Under system reset, all SWJ pins are assigned (JTAG-DP + SW-DP).
- Under system reset, the debugger host sends the JTAG sequence to switch from the JTAG-DP to the SW-DP.
- Still under system reset, the debugger sets a breakpoint on vector reset.
- The system reset is released and the Core halts.
- All the debug communications from this point are done using the SW-DP. The other JTAG pins can then be reassigned as GPIOs by the user software.

Note: For user software designs, note that:

To release the debug pins, remember that they will be first configured either in input-pull-up (nTRST, TMS, TDI) or pull-down (TCK) or output tristate (TDO) for a certain duration after reset until the instant when the user software releases the pins.

When debug pins (JTAG or SW or TRACE) are mapped, changing the corresponding IO pin configuration in the IOPORT controller has no effect.

#### 32.5 STM32F4xx JTAG TAP connection

The STM32F4xx MCUs integrate two serially connected JTAG TAPs, the boundary scan TAP (IR is 5-bit wide) and the Cortex<sup>™</sup>-M4F TAP (IR is 4-bit wide).

To access the TAP of the Cortex™-M4F for debug purposes:

- 1. First, it is necessary to shift the BYPASS instruction of the boundary scan TAP.
- 2. Then, for each IR shift, the scan chain contains 9 bits (=5+4) and the unused TAP instruction must be shifted in using the BYPASS instruction.
- 3. For each data shift, the unused TAP, which is in BYPASS mode, adds 1 extra data bit in the data scan chain.

Note:

**Important**: Once Serial-Wire is selected using the dedicated ARM JTAG sequence, the boundary scan TAP is automatically disabled (JTMS forced high).

Debug support (DBG) RM0090

STM32F4xxx NJTRST JTMS SW-DP -Selected TMS nTRST TMS nTRST JTDI TDI TDO TDI TDO **Boundary scan** Cortex\_M4F TAP **TAP** IR is 5-bit wide IR is 4-bit wide JTDO

Figure 420. JTAG TAP connections

MS19929V1

## 32.6 ID codes and locking mechanism

There are several ID codes inside the STM32F4xx MCUs. ST strongly recommends tools designers to lock their debuggers using the MCU DEVICE ID code located in the external PPB memory map at address 0xE0042000.

#### 32.6.1 MCU device ID code

The STM32F4xx MCUs integrate an MCU ID code. This ID identifies the ST MCU partnumber and the die revision. It is part of the DBG\_MCU component and is mapped on the external PPB bus (see *Section 32.16 on page 1315*). This code is accessible using the JTAG debug port (4 to 5 pins) or the SW debug port (two pins) or by the user software. It is even accessible while the MCU is under system reset.

#### DBGMCU\_IDCODE

Address: 0xE004 2000

Only 32-bits access supported. Read-only.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	REV_ID														
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Poor	erved							DE	V_ID					
	nese	ei veu		r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:16 REV\_ID(15:0) Revision identifier

This field indicates the revision of the device:

0x1000 = Revision A0x1001 = Revision Z

Bits 15:12 Reserved, must be kept at reset value.

Bits 11:0 **DEV\_ID(11:0)**: Device identifier The device ID is 0x413.

## 32.6.2 Boundary scan TAP

#### JTAG ID code

The TAP of the STM32F4xx BSC (boundary scan) integrates a JTAG ID code equal to 0x06413041.

#### 32.6.3 Cortex<sup>™</sup>-M4F TAP

The TAP of the ARM Cortex<sup>™</sup>-M4F integrates a JTAG ID code. This ID code is the ARM default one and has not been modified. This code is only accessible by the JTAG Debug Port.

This code is **0x4BA00477** (corresponds to Cortex<sup>™</sup>-M4F r0p1, see *Section 32.2: Reference ARM documentation*).

Only the DEV\_ID(11:0) should be used for identification by the debugger/programmer tools.

## 32.6.4 Cortex<sup>™</sup>-M4F JEDEC-106 ID code

The ARM Cortex<sup>™</sup>-M4F integrates a JEDEC-106 ID code. It is located in the 4KB ROM table mapped on the internal PPB bus at address 0xE00FF000\_0xE00FFFFF.

This code is accessible by the JTAG Debug Port (4 to 5 pins) or by the SW Debug Port (two pins) or by the user software.

# 32.7 JTAG debug port

A standard JTAG state machine is implemented with a 4-bit instruction register (IR) and five data registers (for full details, refer to the Cortex<sup>TM</sup>-M4Fr0p1 Techni*cal Reference Manual (TRM), for references, please see Section 32.2: Reference ARM documentation).* 

Table 207. JTAG debug port data registers

IR(3:0)	Data register	Details
1111	BYPASS [1 bit]	
1110	IDCODE [32 bits]	ID CODE 0x4BA00477 (ARM Cortex™-M4F r0p1 ID Code)
1010	DPACC [35 bits]	Debug port access register  This initiates a debug port and allows access to a debug port register.  - When transferring data IN:  Bits 34:3 = DATA[31:0] = 32-bit data to transfer for a write request  Bits 2:1 = A[3:2] = 2-bit address of a debug port register.  Bit 0 = RnW = Read request (1) or write request (0).  - When transferring data OUT:  Bits 34:3 = DATA[31:0] = 32-bit data which is read following a read request  Bits 2:0 = ACK[2:0] = 3-bit Acknowledge:  010 = OK/FAULT  001 = WAIT  OTHER = reserved  Refer to Table 208 for a description of the A(3:2) bits

Table 207. JTAG debug port data registers (continued)

IR(3:0)	Data register	Details
1011	APACC [35 bits]	Access port access register  Initiates an access port and allows access to an access port register.  - When transferring data IN: Bits 34:3 = DATA[31:0] = 32-bit data to shift in for a write request Bits 2:1 = A[3:2] = 2-bit address (sub-address AP registers). Bit 0 = RnW= Read request (1) or write request (0).  - When transferring data OUT: Bits 34:3 = DATA[31:0] = 32-bit data which is read following a read request Bits 2:0 = ACK[2:0] = 3-bit Acknowledge: 010 = OK/FAULT 001 = WAIT OTHER = reserved  There are many AP Registers (see AHB-AP) addressed as the combination of:  - The shifted value A[3:2]  - The current value of the DP SELECT register
1000	ABORT [35 bits]	Abort register  - Bits 31:1 = Reserved  - Bit 0 = DAPABORT: write 1 to generate a DAP abort.

Debug support (DBG) RM0090

Table 208. 32-bit debug port registers addressed through the shifted value A[3:2]

Address	A(3:2) value	Description
0x0	00	Reserved, must be kept at reset value.
0x4	01	DP CTRL/STAT register. Used to:  — Request a system or debug power-up  — Configure the transfer operation for AP accesses  — Control the pushed compare and pushed verify operations.  — Read some status flags (overrun, power-up acknowledges)
0x8	10	DP SELECT register: Used to select the current access port and the active 4-words register window.  - Bits 31:24: APSEL: select the current AP  - Bits 23:8: reserved  - Bits 7:4: APBANKSEL: select the active 4-words register window on the current AP  - Bits 3:0: reserved
0xC	11	DP RDBUFF register: Used to allow the debugger to get the final result after a sequence of operations (without requesting new JTAG-DP operation)

## 32.8 SW debug port

## 32.8.1 SW protocol introduction

This synchronous serial protocol uses two pins:

- SWCLK: clock from host to target
- SWDIO: bidirectional

The protocol allows two banks of registers (DPACC registers and APACC registers) to be read and written to.

Bits are transferred LSB-first on the wire.

For SWDIO bidirectional management, the line must be pulled-up on the board (100  $\mbox{K}\Omega$  recommended by ARM).

Each time the direction of SWDIO changes in the protocol, a turnaround time is inserted where the line is not driven by the host nor the target. By default, this turnaround time is one bit time, however this can be adjusted by configuring the SWCLK frequency.

## 32.8.2 SW protocol sequence

Each sequence consist of three phases:

- 1. Packet request (8 bits) transmitted by the host
- 2. Acknowledge response (3 bits) transmitted by the target
- 3. Data transfer phase (33 bits) transmitted by the host or the target

Table 209. Packet request (8-bits)

Bit	Name	Description
0	Start	Must be "1"
1	APnDP	0: DP Access 1: AP Access
2	RnW	0: Write Request 1: Read Request
4:3	A(3:2)	Address field of the DP or AP registers (refer to Table 208)
5	Parity	Single bit parity of preceding bits
6	Stop	0
7	Park	Not driven by the host. Must be read as "1" by the target because of the pull-up

Refer to the Cortex<sup>™</sup>-M4Fr0p1 *TRM* for a detailed description of DPACC and APACC registers.

The packet request is always followed by the turnaround time (default 1 bit) where neither the host nor target drive the line.

Table 210. ACK response (3 bits)

Bit	Name	Description
02	ACK	001: FAULT 010: WAIT 100: OK

The ACK Response must be followed by a turnaround time only if it is a READ transaction or if a WAIT or FAULT acknowledge has been received.

Table 211. DATA transfer (33 bits)

Bit	Name	Description
031	WDATA or RDATA	Write or Read data
32	Parity	Single parity of the 32 data bits

The DATA transfer must be followed by a turnaround time only if it is a READ transaction.

## 32.8.3 SW-DP state machine (reset, idle states, ID code)

The State Machine of the SW-DP has an internal ID code which identifies the SW-DP. It follows the JEP-106 standard. This ID code is the default ARM one and is set to **0x2BA01477** (corresponding to Cortex<sup>™</sup>-M4F r0p1).

Note:

Note that the SW-DP state machine is inactive until the target reads this ID code.

- The SW-DP state machine is in RESET STATE either after power-on reset, or after the DP has switched from JTAG to SWD or after the line is high for more than 50 cycles
- The SW-DP state machine is in IDLE STATE if the line is low for at least two cycles after RESET state.
- After RESET state, it is mandatory to first enter into an IDLE state AND to perform a READ access of the DP-SW ID CODE register. Otherwise, the target will issue a FAULT acknowledge response on another transactions.

Further details of the SW-DP state machine can be found in the  $Cortex^{TM}$ -M4F r0p1 TRM and the CoreSight Design Kit r0p1 <math>TRM.

#### 32.8.4 DP and AP read/write accesses

- Read accesses to the DP are not posted: the target response can be immediate (if ACK=OK) or can be delayed (if ACK=WAIT).
- Read accesses to the AP are posted. This means that the result of the access is returned on the next transfer. If the next access to be done is NOT an AP access, then the DP-RDBUFF register must be read to obtain the result.
   The READOK flag of the DP-CTRL/STAT register is updated on every AP read access or RDBUFF read request to know if the AP read access was successful.
- The SW-DP implements a write buffer (for both DP or AP writes), that enables it to accept a write operation even when other transactions are still outstanding. If the write buffer is full, the target acknowledge response is "WAIT". With the exception of IDCODE read or CTRL/STAT read or ABORT write which are accepted even if the write buffer is full.
- Because of the asynchronous clock domains SWCLK and HCLK, two extra SWCLK cycles are needed after a write transaction (after the parity bit) to make the write effective internally. These cycles should be applied while driving the line low (IDLE state)

This is particularly important when writing the CTRL/STAT for a power-up request. If the next transaction (requiring a power-up) occurs immediately, it will fail.

## 32.8.5 SW-DP registers

Access to these registers are initiated when APnDP=0

Table 212. SW-DP registers

A(3:2)	R/W	CTRLSEL bit of SELECT register	Register	Notes
00	Read		IDCODE	The manufacturer code is not set to ST code. <b>0x2BA01477</b> (identifies the SW-DP)
00	Write		ABORT	

Table 212. SW-DP registers (continued)

A(3:2)	R/W	CTRLSEL bit of SELECT register	Register	Notes
01	Read/Write	0	DP-CTRL/STAT	Purpose is to:  - request a system or debug power-up  - configure the transfer operation for AP accesses  - control the pushed compare and pushed verify operations.  - read some status flags (overrun, power-up acknowledges)
01	Read/Write	1	WIRE CONTROL	Purpose is to configure the physical serial port protocol (like the duration of the turnaround time)
10	Read		READ RESEND	Enables recovery of the read data from a corrupted debugger transfer, without repeating the original AP transfer.
10	Write		SELECT	The purpose is to select the current access port and the active 4-words register window
11	Read/Write		READ BUFFER	This read buffer is useful because AP accesses are posted (the result of a read AP request is available on the next AP transaction).  This read buffer captures data from the AP, presented as the result of a previous read, without initiating a new transaction

# 32.8.6 SW-AP registers

Access to these registers are initiated when APnDP=1

There are many AP Registers (see AHB-AP) addressed as the combination of:

- The shifted value A[3:2]
- The current value of the DP SELECT register

# 32.9 AHB-AP (AHB access port) - valid for both JTAG-DP and SW-DP

#### Features:

- System access is independent of the processor status.
- Either SW-DP or JTAG-DP accesses AHB-AP.
- The AHB-AP is an AHB master into the Bus Matrix. Consequently, it can access all the data buses (Dcode Bus, System Bus, internal and external PPB bus) but the ICode bus.
- Bitband transactions are supported.
- AHB-AP transactions bypass the FPB.

The address of the 32-bits AHP-AP resisters are 6-bits wide (up to 64 words or 256 bytes) and consists of:

- c) Bits [7:4] = the bits [7:4] APBANKSEL of the DP SELECT register
- d) Bits [3:2] = the 2 address bits of A(3:2) of the 35-bit packet request for SW-DP.

The AHB-AP of the Cortex<sup>™</sup>-M4F includes 9 x 32-bits registers:

Table 213. Cortex<sup>™</sup>-M4F AHB-AP registers

Address offset	Register name	Notes
0x00	AHB-AP Control and Status Word	Configures and controls transfers through the AHB interface (size, hprot, status on current transfer, address increment type
0x04	AHB-AP Transfer Address	
0x0C	AHB-AP Data Read/Write	
0x10	AHB-AP Banked Data 0	
0x14	AHB-AP Banked Data 1	Directly maps the 4 aligned data words without rewriting
0x18	AHB-AP Banked Data 2	the Transfer Address Register.
0x1C	AHB-AP Banked Data 3	
0xF8	AHB-AP Debug ROM Address	Base Address of the debug interface
0xFC	AHB-AP ID Register	

Refer to the *Cortex™-M4F r0p1 TRM* for further details.

# 32.10 Core debug

Core debug is accessed through the core debug registers. Debug access to these registers is by means of the *Advanced High-performance Bus* (AHB-AP) port. The processor can access these registers directly over the internal *Private Peripheral Bus* (PPB).

It consists of 4 registers:

Table 214. Core debug registers

Register	Description
DHCSR	The 32-bit Debug Halting Control and Status Register  This provides status information about the state of the processor enable core debug halt and step the processor
DCRSR	The 17-bit Debug Core Register Selector Register: This selects the processor register to transfer data to or from.
DCRDR	The 32-bit Debug Core Register Data Register: This holds data for reading and writing registers to and from the processor selected by the DCRSR (Selector) register.
DEMCR	The 32-bit Debug Exception and Monitor Control Register: This provides Vector Catching and Debug Monitor Control. This register contains a bit named <b>TRCENA</b> which enable the use of a TRACE.

Note:

**Important**: these registers are not reset by a system reset. They are only reset by a poweron reset.

Refer to the *Cortex™-M4F r0p1 TRM* for further details.

To Halt on reset, it is necessary to:

- enable the bit0 (VC\_CORRESET) of the Debug and Exception Monitor Control Register
- enable the bit0 (C\_DEBUGEN) of the Debug Halting Control and Status Register.

# 32.11 Capability of the debugger host to connect under system reset

The STM32F4xx MCUs' reset system comprises the following reset sources:

- POR (power-on reset) which asserts a RESET at each power-up.
- Internal watchdog reset
- Software reset
- External reset

The Cortex<sup>™</sup>-M4F differentiates the reset of the debug part (generally PORRESETn) and the other one (SYSRESETn)

This way, it is possible for the debugger to connect under System Reset, programming the Core Debug Registers to halt the core when fetching the reset vector. Then the host can release the system reset and the core will immediately halt without having executed any instructions. In addition, it is possible to program any debug features under System Reset.

Note:

It is highly recommended for the debugger host to connect (set a breakpoint in the reset vector) under system reset.

# 32.12 FPB (Flash patch breakpoint)

The FPB unit:

- implements hardware breakpoints
- patches code and data from code space to system space. This feature gives the possibility to correct software bugs located in the Code Memory Space.

The use of a Software Patch or a Hardware Breakpoint is exclusive.

The FPB consists of:

- 2 literal comparators for matching against literal loads from Code Space and remapping to a corresponding area in the System Space.
- 6 instruction comparators for matching against instruction fetches from Code Space.
   They can be used either to remap to a corresponding area in the System Space or to generate a Breakpoint Instruction to the core.

## 32.13 DWT (data watchpoint trigger)

The DWT unit consists of four comparators. They are configurable as:

- a hardware watchpoint or
- a trigger to an ETM or
- a PC sampler or
- a data address sampler

The DWT also provides some means to give some profiling informations. For this, some counters are accessible to give the number of:

- Clock cycle
- Folded instructions
- Load store unit (LSU) operations
- Sleep cycles
- CPI (clock per instructions)
- Interrupt overhead

# 32.14 ITM (instrumentation trace macrocell)

## 32.14.1 General description

The ITM is an application-driven trace source that supports *printf* style debugging to trace *Operating System* (OS) and application events, and emits diagnostic system information. The ITM emits trace information as packets which can be generated as:

- Software trace. Software can write directly to the ITM stimulus registers to emit packets.
- Hardware trace. The DWT generates these packets, and the ITM emits them.
- Time stamping. Timestamps are emitted relative to packets. The ITM contains a 21-bit counter to generate the timestamp. The Cortex™-M4F clock or the bit clock rate of the Serial Wire Viewer (SWV) output clocks the counter.

The packets emitted by the ITM are output to the TPIU (Trace Port Interface Unit). The formatter of the TPIU adds some extra packets (refer to TPIU) and then output the complete packets sequence to the debugger host.

The bit TRCEN of the Debug Exception and Monitor Control Register must be enabled before you program or use the ITM.

#### 32.14.2 Time stamp packets, synchronization and overflow packets

Time stamp packets encode time stamp information, generic control and synchronization. It uses a 21-bit timestamp counter (with possible prescalers) which is reset at each time stamp packet emission. This counter can be either clocked by the CPU clock or the SWV clock.

A synchronization packet consists of 6 bytes equal to 0x80\_00\_00\_00\_00 which is emitted to the TPIU as 00 00 00 00 00 80 (LSB emitted first).

A synchronization packet is a timestamp packet control. It is emitted at each DWT trigger.

For this, the DWT must be configured to trigger the ITM: the bit CYCCNTENA (bit0) of the DWT Control Register must be set. In addition, the bit2 (SYNCENA) of the ITM Trace Control Register must be set.

Note:

If the SYNENA bit is not set, the DWT generates Synchronization triggers to the TPIU which will send only TPIU synchronization packets and not ITM synchronization packets.

An overflow packet consists is a special timestamp packets which indicates that data has been written but the FIFO was full.

Table 215. Main ITM registers

Address	Register	Details
@E0000FB0	ITM lock access	Write 0xC5ACCE55 to unlock Write Access to the other ITM registers
		Bits 31-24 = Always 0
		Bits 23 = Busy
		Bits 22-16 = 7-bits ATB ID which identifies the source of the trace data.
		Bits 15-10 = Always 0
		Bits 9:8 = TSPrescale = Time Stamp Prescaler
		Bits 7-5 = Reserved
@E0000E80	ITM trace control	Bit 4 = SWOENA = Enable SWV behavior (to clock the timestamp counter by the SWV clock).
		Bit 3 = DWTENA: Enable the DWT Stimulus
		Bit 2 = SYNCENA: this bit must be to 1 to enable the DWT to generate synchronization triggers so that the TPIU can then emit the synchronization packets.
		Bit 1 = TSENA (Timestamp Enable)
		Bit 0 = ITMENA: Global Enable Bit of the ITM
		Bit 3: mask to enable tracing ports31:24
@E0000E40	ITM trace privilege	Bit 2: mask to enable tracing ports23:16
₩ E0000E40	ITM trace privilege	Bit 1: mask to enable tracing ports15:8
		Bit 0: mask to enable tracing ports7:0
@E0000E00 ITM trace enable		Each bit enables the corresponding Stimulus port to generate trace.
@E0000000- E000007C	Stimulus port registers 0-31	Write the 32-bits data on the selected Stimulus Port (32 available) to be traced out.

### **Example of configuration**

To output a simple value to the TPIU:

- Configure the TPIU and assign TRACE I/Os by configuring the DBGMCU\_CR (refer to Section 32.17.2: TRACE pin assignment and Section 32.16.3: Debug MCU configuration register)
- Write 0xC5ACCE55 to the ITM Lock Access Register to unlock the write access to the ITM registers
- Write 0x00010005 to the ITM Trace Control Register to enable the ITM with Sync enabled and an ATB ID different from 0x00
- Write 0x1 to the ITM Trace Enable Register to enable the Stimulus Port 0
- Write 0x1 to the ITM Trace Privilege Register to unmask stimulus ports 7:0
- Write the value to output in the Stimulus Port Register 0: this can be done by software (using a printf function)

## 32.15 ETM (Embedded trace macrocell)

## 32.15.1 General description

The ETM enables the reconstruction of program execution. Data are traced using the Data Watchpoint and Trace (DWT) component or the Instruction Trace Macrocell (ITM) whereas instructions are traced using the Embedded Trace Macrocell (ETM).

The ETM transmits information as packets and is triggered by embedded resources. These resources must be programmed independently and the trigger source is selected using the Trigger Event Register (0xE0041008). An event could be a simple event (address match from an address comparator) or a logic equation between 2 events. The trigger source is one of the fourth comparators of the DWT module, The following events can be monitored:

- Clock cycle matching
- Data address matching

For more informations on the trigger resources refer to *Section 32.13: DWT (data watchpoint trigger)*.

The packets transmitted by the ETM are output to the TPIU (Trace Port Interface Unit). The formatter of the TPIU adds some extra packets (refer to *Section 32.17: TPIU (trace port interface unit)*) and then outputs the complete packet sequence to the debugger host.

## 32.15.2 Signal protocol, packet types

This part is described in the chapter 7 ETMv3 Signal Protocol of the ARM IHI 0014N document.

## 32.15.3 Main ETM registers

For more information on registers refer to the chapter 3 of the ARM IHI 0014N specification.

Table 216. Main ETM registers

Address	Register	Details
0xE0041FB0	ETM Lock Access	Write 0xC5ACCE55 to unlock the write access to the other ETM registers.
0xE0041000	ETM Control	This register controls the general operation of the ETM, for instance how tracing is enabled.
0xE0041010	ETM Status	This register provides information about the current status of the trace and trigger logic.
0xE0041008	ETM Trigger Event	This register defines the event that will control trigger.
0xE004101C	ETM Trace Enable Control	This register defines which comparator is selected.
0xE0041020	ETM Trace Enable Event	This register defines the trace enabling event.
0xE0041024	ETM Trace Start/Stop	This register defines the traces used by the trigger source to start and stop the trace, respectively.

## 32.15.4 Configuration example

To output a simple value to the TPIU:

- Configure the TPIU and enable the I/IO\_TRACEN to assign TRACE I/Os in the STM32F4xxdebug configuration register.
- Write 0xC5ACCE55 to the ETM Lock Access Register to unlock the write access to the ITM registers
- Write 0x00001D1E to the control register (configure the trace)
- Write 0000406F to the Trigger Event register (define the trigger event)
- Write 0000006F to the Trace Enable Event register (define an event to start/stop)
- Write 00000001 to the Trace Start/stop register (enable the trace)
- Write 0000191E to the ETM Control Register (end of configuration)

# 32.16 MCU debug component (DBGMCU)

The MCU debug component helps the debugger provide support for:

- Low-power modes
- Clock control for timers, watchdog, I2C and bxCAN during a breakpoint
- Control of the trace pins assignment

## 32.16.1 Debug support for low-power modes

To enter low-power mode, the instruction WFI or WFE must be executed.

The MCU implements several low-power modes which can either deactivate the CPU clock or reduce the power of the CPU.

The core does not allow FCLK or HCLK to be turned off during a debug session. As these are required for the debugger connection, during a debug, they must remain active. The MCU integrates special means to allow the user to debug software in low-power modes.

For this, the debugger host must first set some debug configuration registers to change the low-power mode behavior:

- In Sleep mode, DBG\_SLEEP bit of DBGMCU\_CR register must be previously set by the debugger. This will feed HCLK with the same clock that is provided to FCLK (system clock previously configured by the software).
- In Stop mode, the bit DBG\_STOP must be previously set by the debugger. This will enable the internal RC oscillator clock to feed FCLK and HCLK in STOP mode.

# 32.16.2 Debug support for timers, watchdog, bxCAN and I<sup>2</sup>C

During a breakpoint, it is necessary to choose how the counter of timers and watchdog should behave:

- They can continue to count inside a breakpoint. This is usually required when a PWM is controlling a motor, for example.
- They can stop to count inside a breakpoint. This is required for watchdog purposes.

For the bxCAN, the user can choose to block the update of the receive register during a breakpoint.

For the I<sup>2</sup>C, the user can choose to block the SMBUS timeout during a breakpoint.

## 32.16.3 Debug MCU configuration register

This register allows the configuration of the MCU under DEBUG. This concerns:

- Low-power mode support
- Timer and watchdog counter support
- bxCAN communication support
- Trace pin assignment

This DBGMCU\_CR is mapped on the External PPB bus at address 0xE0042004

It is asynchronously reset by the PORESET (and not the system reset). It can be written by the debugger under system reset.

If the debugger host does not support these features, it is still possible for the user software to write to these registers.

#### DBGMCU\_CR

Address: 0xE004 2004

Only 32-bit access supported

POR Reset: 0x0000 0000 (not reset by system reset)

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
							Rese	erved							
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			Rese	rved				_	CE_ DE :0]	TRACE IOEN	Rese	erved	DBG_ STAND BY	DBG_ STOP	DBG_ SLEEP
								rw	rw	rw			rw	rw	rw

Bits 31:8 Reserved, must be kept at reset value.

#### Bits 7:5 TRACE\_MODE[1:0] and TRACE\_IOEN: Trace pin assignment control

- With TRACE\_IOEN=0:

TRACE\_MODE=xx: TRACE pins not assigned (default state)

- With TRACE\_IOEN=1:
  - TRACE\_MODE=00: TRACE pin assignment for Asynchronous Mode
  - TRACE\_MODE=01: TRACE pin assignment for Synchronous Mode with a TRACEDATA size of 1
  - TRACE\_MODE=10: TRACE pin assignment for Synchronous Mode with a TRACEDATA size of 2
  - TRACE\_MODE=11: TRACE pin assignment for Synchronous Mode with a TRACEDATA size of 4

#### Bits 4:3 Reserved, must be kept at reset value.

#### Bit 2 DBG\_STANDBY: Debug Standby mode

0: (FCLK=Off, HCLK=Off) The whole digital part is unpowered.

From software point of view, exiting from Standby is identical than fetching reset vector (except a few status bit indicated that the MCU is resuming from Standby)

1: (FCLK=On, HCLK=On) In this case, the digital part is not unpowered and FCLK and HCLK are provided by the internal RC oscillator which remains active. In addition, the MCU generate a system reset during Standby mode so that exiting from Standby is identical than fetching from reset

#### Bit 1 DBG\_STOP: Debug Stop mode

0: (FCLK=Off, HCLK=Off) In STOP mode, the clock controller disables all clocks (including HCLK and FCLK). When exiting from STOP mode, the clock configuration is identical to the one after RESET (CPU clocked by the 8 MHz internal RC oscillator (HSI)). Consequently, the software must reprogram the clock controller to enable the PLL, the Xtal, etc.

1: (FCLK=On, HCLK=On) In this case, when entering STOP mode, FCLK and HCLK are provided by the internal RC oscillator which remains active in STOP mode. When exiting STOP mode, the software must reprogram the clock controller to enable the PLL, the Xtal, etc. (in the same way it would do in case of DBG\_STOP=0)

#### Bit 0 DBG\_SLEEP: Debug Sleep mode

0: (FCLK=On, HCLK=Off) In Sleep mode, FCLK is clocked by the system clock as previously configured by the software while HCLK is disabled.

In Sleep mode, the clock controller configuration is not reset and remains in the previously programmed state. Consequently, when exiting from Sleep mode, the software does not need to reconfigure the clock controller.

1: (FCLK=On, HCLK=On) In this case, when entering Sleep mode, HCLK is fed by the same clock that is provided to FCLK (system clock as previously configured by the software).

## 32.16.4 Debug MCU APB1 freeze register (DBGMCU\_APB1\_FZ)

The DBGMCU\_APB1\_FZ register is used to configure the MCU under Debug. It concerns APB1 peripherals. It is mapped on the external PPB bus at address 0xE004 2008.

The register is asynchronously reset by the POR (and not the system reset). It can be written by the debugger under system reset.

Address: 0xE004 2008

Only 32-bits access are supported.

Power on reset (POR): 0x0000 0000 (not reset by system reset)

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
		Reserve	d		DBG_CAN2_STOP	DBG_CAN1_STOP	Reserved	DBG_I2C3_SMBUS_TIMEOUT	DBG_I2C2_SMBUS_TIMEOUT	DBG_I2C1_SMBUS_TIMEOUT			Reserved		
					rw	rw		rw	rw	rw					
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Reserv	ved	DBG_IWDG_STOP	DBG_WWDG_STOP	DBG_RTC_STOP	Reserved	DBG_TIM14_STOP	DBG_TIM13_STOP	DBG_TIM12_STOP	DBG_TIM7_STOP	DBG_TIM6_STOP	TIM6_STOP TIM4_STOP TIM3_STOP			
				rw	rw		rw	rw	rw	rw	rw	rw	rw	rw	rw

Bits 31:27 Reserved, must be kept at reset value.

- Bit 26 DBG\_CAN2\_STOP: Debug CAN2 stopped when Core is halted
  - 0: Same behavior as in normal mode
  - 1: The CAN2 receive registers are frozen
- Bit 25 DBG\_CAN1\_STOP: Debug CAN2 stopped when Core is halted
  - 0: Same behavior as in normal mode
  - 1: The CAN2 receive registers are frozen
- Bit 24 Reserved, must be kept at reset value.
- Bit 23 **DBG\_I2C3\_SMBUS\_TIMEOUT:** SMBUS timeout mode stopped when Core is halted
  - 0: Same behavior as in normal mode
  - 1: The SMBUS timeout is frozen
- Bit 22 DBG\_I2C2\_SMBUS\_TIMEOUT: SMBUS timeout mode stopped when Core is halted
  - 0: Same behavior as in normal mode
  - 1: The SMBUS timeout is frozen
- Bit 21 **DBG\_I2C1\_SMBUS\_TIMEOUT:** SMBUS timeout mode stopped when Core is halted
  - 0: Same behavior as in normal mode
  - 1: The SMBUS timeout is frozen

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- Bit 20:13 Reserved, must be kept at reset value.
  - Bit 12 **DBG\_IWDG\_STOP:** Debug independent watchdog stopped when core is halted
    - 0: The independent watchdog counter clock continues even if the core is halted
    - 1: The independent watchdog counter clock is stopped when the core is halted
  - Bit 11 DBG\_WWDG\_STOP: Debug Window Watchdog stopped when Core is halted
    - 0: The window watchdog counter clock continues even if the core is halted
    - 1: The window watchdog counter clock is stopped when the core is halted
  - Bit 10 DBG\_RTC\_STOP: RTC stopped when Core is halted
    - 0: The RTC counter clock continues even if the core is halted
    - 1: The RTC counter clock is stopped when the core is halted
  - Bit 9 Reserved, must be kept at reset value.
- Bits 8:0 **DBG\_TIMx\_STOP:** TIMx counter stopped when core is halted (x=2..7, 12..14)
  - 0: The clock of the involved Timer Counter is fed even if the core is halted
  - 1: The clock of the involved Timer counter is stopped when the core is halted

## 32.16.5 Debug MCU APB2 Freeze register (DBGMCU\_APB2\_FZ)

The DBGMCU\_APB2\_FZ register is used to configure the MCU under Debug. It concerns APB2 peripherals.

This register is mapped on the external PPB bus at address 0xE004 200C

It is asynchronously reset by the POR (and not the system reset). It can be written by the debugger under system reset.

Address: 0xE004 200C

Only 32-bit access is supported.

POR: 0x0000 0000 (not reset by system reset)

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
					F	Reserved	i						DBG_TIM11 _STOP	DBG_TIM10 _STOP	DBG_TIM9_ STOP
													rw	rw	rw
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	Reserved														DBG_TIM1_ STOP
														rw	rw

Bits 31:19 Reserved, must be kept at reset value.

Bits 18:16 **DBG\_TIMx\_STOP:** TIMx counter stopped when core is halted (x=9..11)

0: The clock of the involved Timer Counter is fed even if the core is halted

1: The clock of the involved Timer counter is stopped when the core is halted

- Bits 15:2 Reserved, must be kept at reset value.
  - Bit 1 DBG\_TIM8\_STOP: TIM8 counter stopped when core is halted
    - 0: The clock of the involved Timer Counter is fed even if the core is halted
    - 1: The clock of the involved Timer counter is stopped when the core is halted
  - Bit 0 DBG\_TIM1\_STOP: TIM1 counter stopped when core is halted
    - 0: The clock of the involved Timer Counter is fed even if the core is halted
    - 1: The clock of the involved Timer counter is stopped when the core is halted

## 32.17 TPIU (trace port interface unit)

#### 32.17.1 Introduction

The TPIU acts as a bridge between the on-chip trace data from the ITM and the ETM.

The output data stream encapsulates the trace source ID, that is then captured by a *trace port analyzer* (TPA).

The core embeds a simple TPIU, especially designed for low-cost debug (consisting of a special version of the CoreSight TPIU).

TRACECLKIN domain CLK domain TPIU TRACECLKIN Asynchronous FTM **FIFO** ► TRACECK **TPIU** Trace out TRACEDATA formatter (serializer) [3:0] Asynchronous ITM **FIFO TRACESWO** External PPB bus

Figure 421. TPIU block diagram

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## 32.17.2 TRACE pin assignment

Asynchronous mode

The asynchronous mode requires 1 extra pin and is available on all packages. It is only available if using Serial Wire mode (not in JTAG mode).

Table 217. Asynchronous TRACE pin assignment

TPUI pin name		Trace synchronous mode	STM32F4xx pin
TPOI pili liaille	Туре	Description	assignment
TRACESWO	0	TRACE Async Data Output	PB3

#### Synchronous mode

The synchronous mode requires from 2 to 6 extra pins depending on the data trace size and is only available in the larger packages. In addition it is available in JTAG mode and in Serial Wire mode and provides better bandwidth output capabilities than asynchronous trace.

Table 218. Synchronous TRACE pin assignment

TPUI pin name		Trace synchronous mode	STM32F4xx pin
Tr of pill flame	Туре	Description	assignment
TRACECK	0	TRACE Clock	PE2
TRACED[3:0]	0	TRACE Sync Data Outputs Can be 1, 2 or 4.	PE[6:3]

#### **TPUI TRACE** pin assignment

By default, these pins are NOT assigned. They can be assigned by setting the TRACE\_IOEN and TRACE\_MODE bits in the *MCU Debug component configuration register*. This configuration has to be done by the debugger host.

In addition, the number of pins to assign depends on the trace configuration (asynchronous or synchronous).

- Asynchronous mode: 1 extra pin is needed
- **Synchronous mode**: from 2 to 5 extra pins are needed depending on the size of the data trace port register (1, 2 or 4):
  - TRACECK
  - TRACED(0) if port size is configured to 1, 2 or 4
  - TRACED(1) if port size is configured to 2 or 4
  - TRACED(2) if port size is configured to 4
  - TRACED(3) if port size is configured to 4

To assign the TRACE pin, the debugger host must program the bits TRACE\_IOEN and TRACE\_MODE[1:0] of the Debug MCU configuration Register (DBGMCU\_CR). By default the TRACE pins are not assigned.

This register is mapped on the external PPB and is reset by the PORESET (and not by the SYSTEM reset). It can be written by the debugger under SYSTEM reset.

DBGMCU\_CR TRACE IO pin assigned register **Pins** assigned for: TRACE TRACE PB3 / JTDO/ PE2/ PE3/ PE5 / PE6 / **PF4** / MODE[1:0] **IOEN TRACESWO** TRACECK|TRACED[0]|TRACED[1]|TRACED[2]|TRACED[3] No Trace Released (1) 0 XX (default state) Asvnchronous Released **TRACESWO** 1 00 Trace (usable as GPIO) Synchronous 1 01 TRACECK TRACED[0] Trace 1 bit Synchronous Released (1) 1 10 TRACECK|TRACED[0]|TRACED[1] Trace 2 bit Synchronous TRACECK|TRACED[0]|TRACED[1]|TRACED[2]|TRACED[3] 1 11 Trace 4 bit

Table 219. Flexible TRACE pin assignment

Note:

By default, the TRACECLKIN input clock of the TPIU is tied to GND. It is assigned to HCLK two clock cycles after the bit TRACE\_IOEN has been set.

The debugger must then program the Trace Mode by writing the PROTOCOL[1:0] bits in the SPP\_R (Selected Pin Protocol) register of the TPIU.

- PROTOCOL=00: Trace Port Mode (synchronous)
- PROTOCOL=01 or 10: Serial Wire (Manchester or NRZ) Mode (asynchronous mode).
   Default state is 01

It then also configures the TRACE port size by writing the bits [3:0] in the CPSPS\_R (Current Sync Port Size Register) of the TPIU:

- 0x1 for 1 pin (default state)
- 0x2 for 2 pins
- 0x8 for 4 pins

## 32.17.3 TPUI formatter

The formatter protocol outputs data in 16-byte frames:

- seven bytes of data
- eight bytes of mixed-use bytes consisting of:
  - 1 bit (LSB) to indicate it is a DATA byte ('0) or an ID byte ('1).
  - 7 bits (MSB) which can be data or change of source ID trace.
- one byte of auxiliary bits where each bit corresponds to one of the eight mixed-use bytes:
  - if the corresponding byte was a data, this bit gives bit0 of the data.
  - if the corresponding byte was an ID change, this bit indicates when that ID change takes effect.

Note: Refer to the ARM CoreSight Architecture Specification v1.0 (ARM IHI 0029B) for further information

<sup>1.</sup> When Serial Wire mode is used, it is released. But when JTAG is used, it is assigned to JTDO.

## 32.17.4 TPUI frame synchronization packets

The TPUI can generate two types of synchronization packets:

The Frame Synchronization packet (or Full Word Synchronization packet)
 It consists of the word: 0x7F\_FF\_FF (LSB emitted first). This sequence can not occur at any other time provided that the ID source code 0x7F has not been used.
 It is output periodically *between* frames.

In continuous mode, the TPA must discard all these frames once a synchronization frame has been found.

The Half-Word Synchronization packet

It consists of the half word: 0x7F\_FF (LSB emitted first).

It is output periodically between or within frames.

These packets are only generated in continuous mode and enable the TPA to detect that the TRACE port is in IDLE mode (no TRACE to be captured). When detected by the TPA, it must be discarded.

### 32.17.5 Transmission of the synchronization frame packet

There is no Synchronization Counter register implemented in the TPIU of the core. Consequently, the synchronization trigger can only be generated by the **DWT**. Refer to the registers DWT Control Register (bits SYNCTAP[11:10]) and the DWT Current PC Sampler Cycle Count Register.

The TPUI Frame synchronization packet (0x7F\_FF\_FF\_FF) is emitted:

- after each TPIU reset release. This reset is synchronously released with the rising edge of the TRACECLKIN clock. This means that this packet is transmitted when the TRACE\_IOEN bit in the DBGMCU\_CFG register is set. In this case, the word 0x7F\_FF\_FF is not followed by any formatted packet.
- at each DWT trigger (assuming DWT has been previously configured). Two cases occur:
  - If the bit SYNENA of the ITM is reset, only the word 0x7F\_FF\_FF is emitted without any formatted stream which follows.
  - If the bit SYNENA of the ITM is set, then the ITM synchronization packets will follow (0x80\_00\_00\_00\_00\_00), formatted by the TPUI (trace source ID added).

## 32.17.6 Synchronous mode

The trace data output size can be configured to 4, 2 or 1 pin: TRACED(3:0)

The output clock is output to the debugger (TRACECK)

Here, TRACECLKIN is driven internally and is connected to HCLK only when TRACE is used.

Note: In this synchronous mode, it is not required to provide a stable clock frequency.

The TRACE I/Os (including TRACECK) are driven by the rising edge of TRACLKIN (equal to HCLK). Consequently, the output frequency of TRACECK is equal to HCLK/2.

## 32.17.7 Asynchronous mode

This is a low cost alternative to output the trace using only 1 pin: this is the asynchronous output pin TRACESWO. Obviously there is a limited bandwidth.

TRACESWO is multiplexed with JTDO when using the SW-DP pin. This way, this functionality is available in all STM32F4xx packages.

This asynchronous mode requires a constant frequency for TRACECLKIN. For the standard UART (NRZ) capture mechanism, 5% accuracy is needed. The Manchester encoded version is tolerant up to 10%.

#### 32.17.8 TRACECLKIN connection inside the STM32F4xx

In the STM32F4xx, this TRACECLKIN input is internally connected to HCLK. This means that when in asynchronous trace mode, the application is restricted to use to time frames where the CPU frequency is stable.

Note: Important: when using asynchronous trace: it is important to be aware that:

The default clock of the STM32F4xx MCUs is the internal RC oscillator. Its frequency under reset is different from the one after reset release. This is because the RC calibration is the default one under system reset and is updated at each system reset release.

Consequently, the trace port analyzer (TPA) should not enable the trace (with the TRACE\_IOEN bit) under system reset, because a Synchronization Frame Packet will be issued with a different bit time than trace packets which will be transmitted after reset release.

# 32.17.9 TPIU registers

The TPIU APB registers can be read and written only if the bit TRCENA of the Debug Exception and Monitor Control Register (DEMCR) is set. Otherwise, the registers are read as zero (the output of this bit enables the PCLK of the TPIU).

Table 220. Important TPIU registers

Address	Register	Description
0xE0040004	Current port size	Allows the trace port size to be selected:  Bit 0: Port size = 1  Bit 1: Port size = 2  Bit 2: Port size = 3, not supported  Bit 3: Port Size = 4  Only 1 bit must be set. By default, the port size is one bit.  (0x00000001)
0xE00400F0	Selected pin protocol	Allows the Trace Port Protocol to be selected:  Bit1:0=  00: Sync Trace Port Mode  01: Serial Wire Output - manchester (default value)  10: Serial Wire Output - NRZ  11: reserved
0xE0040304	Formatter and flush control	Bit 31-9 = always '0  Bit 8 = TrigIn = always '1 to indicate that triggers are indicated  Bit 7-4 = always 0  Bit 3-2 = always 0  Bit 1 = EnFCont. In Sync Trace mode (Select_Pin_Protocol register bit1:0=00), this bit is forced to '1: the formatter is automatically enabled in continuous mode. In asynchronous mode (Select_Pin_Protocol register bit1:0 <> 00), this bit can be written to activate or not the formatter.  Bit 0 = always 0  The resulting default value is 0x102  Note: In synchronous mode, because the TRACECTL pin is not mapped outside the chip, the formatter is always enabled in continuous mode -this way the formatter inserts some control packets to identify the source of the trace packets).
0xE0040300	Formatter and flush status	Not used in Cortex™-M4F, always read as 0x00000008

## 32.17.10 Example of configuration

- Set the bit TRCENA in the Debug Exception and Monitor Control Register (DEMCR)
- Write the TPIU Current Port Size Register to the desired value (default is 0x1 for a 1-bit port size)
- Write TPIU Formatter and Flush Control Register to 0x102 (default value)
- Write the TPIU Select Pin Protocol to select the sync or async mode. Example: 0x2 for async NRZ mode (UART like)
- Write the DBGMCU control register to 0x20 (bit IO\_TRACEN) to assign TRACE I/Os for async mode. A TPIU Sync packet is emitted at this time (FF\_FF\_FF\_7F)
- Configure the ITM and write the ITM Stimulus register to output a value

# 32.18 DBG register map

The following table summarizes the Debug registers

Table 221. DBG register map and reset values

Addr.	Register	31	29	22	26	25	24	23	22	20	19	18	17	16	15	13	12	11	10	6	8	7	9	2	4	3	2	-	0
0xE004 2000	DBGMCU_IDC ODE					R	EV_	.ID							Res	erved							DE\	/_I[	)				
	Reset value <sup>(1)</sup>	хх	XX	Х	Х	ХХ		Х	X	Х	Х	Х	Х	Χ			-	Χ	Χ	Х	Χ	Х	Х	Х	Χ	Х	Х	Х	Χ
0xE004 2004	DBGMCU_CR			Re	eser	ved				DBG_TIM7_STOP	DBG_TIM6_STOP	DBG_TIM5_STOP	DBG_TIM8_STOP	DBG_I2C2_SMBUS_TIMEOUT		Re	ese	rve	d			TR E MC [1	AC E_ DDE :0]	TR AC E_ IO E N	70,000		DBG_STANDBY	DBG_STOP	DBG_SLEEP
	Reset value									0	0	0	0	0								0	0	0			0	0	0
0xE004 2008	DBGMCU_ APB1_FZ	R	eserve	d	DBG_CAN2_STOP	DBG_CAN1_STOP	DBG 1903 SMBHS TIMEOLIT	DBG_I2CS_SIMBUS_TIMEOUT				F	Res	erve	d	OTO COMMISSION	DBG_IWDG_SIOP	DBG_WWDG_STOP	Reserved	DBG_RTC_STOP	DBG_TIM14_STOP	DBG_TIM13_STOP	DBG_TIM12_STOP	DBG_TIM7_STOP	DBG_TIM6_STOP	DBG_TIM5_STOP	DBG_TIM4_STOP	DBG_TIM3_STOP	DBG_TIM2_STOP
	Reset value				0	0	0 (	0 (	0 0								0	0		0	0	0	0	0	0	0	0	0	0
0xE004 200C	DBGMCU_ APB2_FZ				Re	serve	d					OBG_TIM11_STOP	O DBG_TIM10_STOP	O DBG_TIM9_STOP					F	Rese	erve	d						o DBG_TIM8_STOP	O DBG_TIM1_STOP

<sup>1.</sup> The reset value is product dependent. For more information, refer to Section 32.6.1: MCU device ID code.



# 33 Device electronic signature

The electronic signature is stored in the Flash memory area. It can be read using the JTAG/SWD or the CPU. It contains factory-programmed identification data that allow the user firmware or other external devices to automatically match its interface to the characteristics of the STM32F4xx microcontrollers.

# 33.1 Unique device ID register (96 bits)

The unique device identifier is ideally suited:

- for use as serial numbers (for example USB string serial numbers or other end applications)
- for use as security keys in order to increase the security of code in Flash memory while using and combining this unique ID with software cryptographic primitives and protocols before programming the internal Flash memory
- to activate secure boot processes, etc.

The 96-bit unique device identifier provides a reference number which is unique for any device and in any context. These bits can never be altered by the user.

The 96-bit unique device identifier can also be read in single bytes/half-words/words in different ways and then be concatenated using a custom algorithm.

Base address: 0x1FFF 7A10

Address offset: 0x00

Read only = 0xXXXX XXXX where X is factory-programmed

31	30	29	28	27	26	25	24	23	22	21	20	19	18	1/	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	U_ID(31:0)																														
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:0 U\_ID(31:0): 31:0 unique ID bits

Address offset: 0x04

Read only = 0xXXXX XXXX where X is factory-programmed

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	U_ID(63:48)														
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							U_ID(	47:32)							
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:0 U\_ID(63:32): 63:32 unique ID bits

Address offset: 0x08

Read only = 0xXXXX XXXX where X is factory-programmed

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16
	U_ID(95:80)														
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	U_ID(79:64)														
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 31:0 **U\_ID(95:64):** 95:64 Unique ID bits.

## 33.2 Flash size

Base address: 0x1FFF 7A22

Address offset: 0x00

Read only = 0xXXXX where X is factory-programmed

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							F_S	SIZE							
r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r

Bits 15:0 F\_ID(15:0): Flash memory size

This bitfield indicates the size of the device Flash memory expressed in Kbytes. As an example, 0x0400 corresponds to 1024 Kbytes.

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# **Revision history**

Table 222. Document revision history

Date	Version	Changes
15-Sep-2011	1	Initial release.
		Updated reference documents and added <i>Table 1: Applicable products</i> on cover page.
		MEMORY: Updated Section 2.3.1: Embedded SRAM.
		PWR:
		Updated VDDA and VREF+ decoupling capacitor in <i>Figure 3: Powe supply overview</i> .
		Updated case of no external battery in Section 4.1.2: Battery backudomain.
		VOSRDY bit changed to read-only in Section 4.4.2: PWR power control/status register (PWR_CSR).
		Removed VDDA in Section 4.2.3: Programmable voltage detector (PVD) and remove VDDA in PVDO bit description (Section 4.4.2: PWR power control/status register (PWR_CSR)).
19-Oct-2012	2	RCC: Updated Figure 8: Simplified diagram of the reset circuit and minimum reset pulse duration guaranteed by pulse generator restricted to internal reset sources.
		GPIOs:
		Updated Section 6.3.1: General-purpose I/O (GPIO).
		DMA:
		Updated direct mode description in Section 8.2: DMA main feature Updated direct mode description in Section: Memory-to-periphera
		mode, and Section 8.3.12: FIFO/Direct mode.  Modified Stream2 /Channel 2 in Table 22: DMA1 request mapping.
		Added note related to EN bit in Section 8.5.5: DMA stream $x$ configuration register (DMA_SxCR) ( $x = 07$ ). Updated definition on NDT[15:0] bits in Section 8.5.6: DMA stream $x$ number of data register (DMA_SxNDTR) ( $x = 07$ ).
		Updated register access in Section 8.5: DMA registers.
		Interrupts: Updated number of maskable interrupts to 82 in <i>Section 9.1.1: NVI featuress</i> .
		EXTI:
		Updated Section 9.2: External interrupt/event controller (EXTI)

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Table 222. Document revision history

Date	Version	Changes
19-Oct-2012	2 (continued)	ADC: Changed ADCCLK frequency to 30 MHz in Section 10.5: Channel-wise programmable sampling timee. Added recovery from ADC sequence in Section 10.8.1: Using the DMA and Section 10.8.2: Managing a sequence of conversions without using the DMA. Updated AWDIE in Section 10.13.2: ADC control register 1 (ADC_CR1). Added read and write access in Section 10.13: ADC registers.  Advanced control timers (TIM1 and TIM8) Updated 16-bit prescaler range in Section 13.2: TIM1&TIM8 main features. Updated OC1 block diagram in Figure 93: Output stage of capture/compare channel (channel 1 to 3). Updated update event generation in Upcounting mode and Downcounting mode in Section 13.3.2: Counter modes and Section 13.3.3: Repetition counter. Updated bits that control the dead-time generation in Section 13.3.11: Complementary outputs and dead-time insertion. Updated ways to generate a break in Section 13.3.12: Using the break function. Changed OCxREF to ETR in the example given in Section 13.3.13: Clearing the OCxREF signal on an external event and changed OCREF_CLR to ETRF in Figure 103: Clearing TIMx OCxREF. Updated configuration for example of counter operation in encoder interface mode in Section 13.3.16: Encoder interface mode. Added register access in Section 13.4: TIM1&TIM8 registers. Changed definition of ARR[15:0] bits in Section 13.4.12: TIM1&TIM8 auto-reload register (TIMx_ARR). Updated BKE definition in Section 13.4.18: TIM1&TIM8 break and dead-time register (TIMx_BDTR).

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Table 222. Document revision history

Date	Version	Changes
		General purpose timers (TIM2 to TIM5)
19-Oct-2012	2 (continued)	Removed all references to "repetition counter".  Added Figure 113: General-purpose timer block diagram.  Updated 16-bit prescaler range in Section 14.2: TIM2 to TIM5 main features.  External clock mode 2 ETR restricted to TIM2 to TIM4 in Section 14.3.3: Clock selection and Section 14.3.6: PWM input mode.  Updated Section 14.3.9: PWM mode and Section 14.3.11: Clearing the OCxREF signal on an external event.  Updated Figure 153: Master/Slave timer example to change ITR1 to ITR0.  Updated read and write access to registers in Section 14.4: TIM2 to TIM5 registerss.  Restored bits 15 to 8 of TIMx_SMCR as well as Table 62: TIMx internal trigger connection in Section 14.4.3.  Removed note 1 related to OC1M bits in Section 14.4.13: TIMx capture/compare register 1 (TIMx_CCR1).  Updated TIMx_CCER bit description for TIM2 to TIM5 in Section 14.4.9: TIMx capture/compare enable register (TIMx_CCER).  General purpose timers (TIM9 to TIM14)  Updated 16-bit prescaler range in Section 15.2.1: TIM9/TIM12 main features and Section 15.3: TIM10/TIM11 and TIM13/TIM14 main features.  Updated Figure 160: General-purpose timer block diagram (TIM10/11/13/14)) to remove TRGO trigger controller output.  Added register access in Section 15.5: TIM9 and TIM12 registers and Section 15.6: TIM10/11/13/14 registers.  Basic timers (TIM6 and TIM7)  Removed all references to "repetition counter".  Updated 16-bit prescaler range in Section 16.2: TIM6&TIM7 main features.
		Updated Section 21.3.1: Duration of the processing.
		RNG:
		Updated Section 20.1: RNG introduction.

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Table 222. Document revision history

Date	Version	Changes
19-Oct-2012	2 (continued)	RTC:  Updated Figure 216: RTC block diagram.  Added formula to compute fck_apre in Section 22.3.1: Clock and prescalers.  Updated Section 22.3.9: RTC reference clock detection.  Updated Section : RTC register write protection.  Added RTC_SSR shadow register in Section 22.3.6: Reading the calendar.  Updated description of DC[4:0] bits in Section 22.6.7: RTC calibration register (RTC_CALIBR).  Renamed RTC_BKxR into RTC_BKPxR in Table 82: RTC register map and reset values.  Added power-on reset value and changed reset value to system reset value in Section 22.6.11: RTC sub second register (RTC_SSR).  Updated definition of ALARMOUTTYPE in Section 22.6.17: RTC tamper and alternate function configuration register (RTC_TAFCR).  12C:  Modified Section 23.3.7: DMA requests.  Updated bit 14 description in Section 23.6.3: I2C Own address register 1 (I2C_OAR1)).  Updated definition of PE bit and note related to SWRST bit; moved note related to STOP bit to the whole register in Section 23.6.1: I2C Control register 1 (I2C_CR1).  USART:  Section 24.6.6: Control register 3 (USART_CR3)): removed notes related to UART5 in DMAT and DMAR description.  Updated TTable 95: Error calculation for programmed baud rates at fPCLK = 42 MHz or fPCLK = 84 Hz, oversampling by 16 and Table 96: Error calculation for programmed baud rates at fPCLK = 42 MHz or fPCLK = 84 MHz, oversampling by 8.

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Table 222. Document revision history

Date	Version	Changes
19-Oct-2012	2 (continued)	SPI/12S:  Updated Section 25.1: SPI introduction.  Changed I2S simplex communication/mode to half-duplex communication/mode. Updated flags in reception/transmission modes in Section 25.2: I2S features. Added Frame error flag in Table 105: I2S interrupt requests.  Added register access in Section 25.5: SPI and I2S registers.  Updated ERRIE definition in Section 25.5: SPI control register 2 (SPI_CR2).  Renamed TIFRFE to FRE and definition updated in Section 25.5: SPI status register (SPI_SR).  SDIO:  Updated value and description for bits [45:40] and [7:1] in Table 133: R4 response. Updated value at bits [45:40] in Table 135: R5 response.  CAN:  Updated Figure 300: Dual CAN block diagram.  Modified definition of CAN2SB bits in Section: CAN filter master register (CAN_FMR).  Added register access in Section 27.9: CAN registers  ETHERNET:  Updated standard for precision networked clock synchronization in Section 28.1: Ethernet introduction and Section 28.2: MAC core features.  Updated CR bit definition in Section: Ethernet MAC MII address register (ETH_MACMIIAR).  Replace RTPR by PM bit in Table 149: Source address filtering.  USB OTG FS  Updated remote wakeup signaling bit and the resume interrupt in Section: Suspended state.  Added peripheral register access in Section 29.16: OTG_FS control and status registers.  Updated INEPTXSA description in OTG_FS_DIEPTXFX.  Changed PHYSEL from bit 7 to bit 6 of the OTG_FS_GUSBCFG register.

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Table 222. Document revision history

Date	Version	Changes
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19-Oct-2012	_	NAND/PC Card access parameters.  Updated case of synchronous accesses in Section 31.5: NOR Flash/PSRAM controller.  Changed minimum value for ADDSET to 0 in Table 183, Table 185, Table 186, Table 187, Table 189 and Table 191.  Move note from Figure 399: Mode1 write accesses and Figure 398: Mode1 read accesses. Move note from Figure 401: ModeA write accesses to Figure 400: ModeA read accesses.  Updated Section: WAIT management in asynchronous accesses. Added register access in Section 31.5.6: NOR/PSRAM control registers and Section 31.6.2: NAND Flash / PC Card supported memories and transactions.  Removed caution note in Section 31.6.1: External memory interface signalss.  Updated Table 200: 16-bit PC Card.  Updated step 3 in Section 31.6.4: NAND Flash operations.  Updated Figure 417: Access to non 'CE don't care' NAND-Flash and note below in Section 31.6.5: NAND Flash pre-wait functionality.
		DEBUG: Updated Section 32.4.3: Internal pull-up and pull-down on JTAG pins.

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Date	Version	Changes
19-Oct-2012	2 (continued)	Electronic signature  Updated Section 33: Device electronic signature introduction.  Updated REV_ID[15:0] to add revision Z in Section 32.6.1: MCU device ID code.  Updated address and example in Section 33.2: Flash size.

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