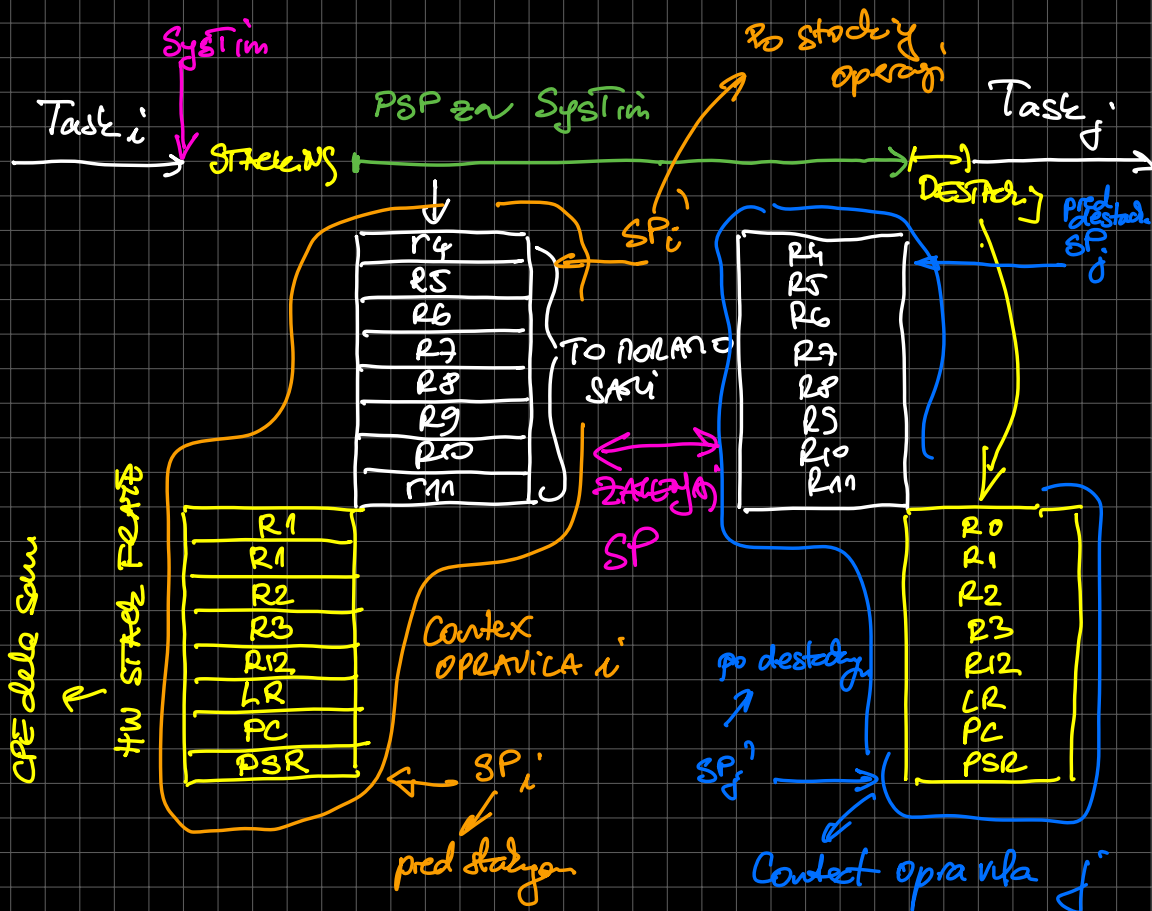
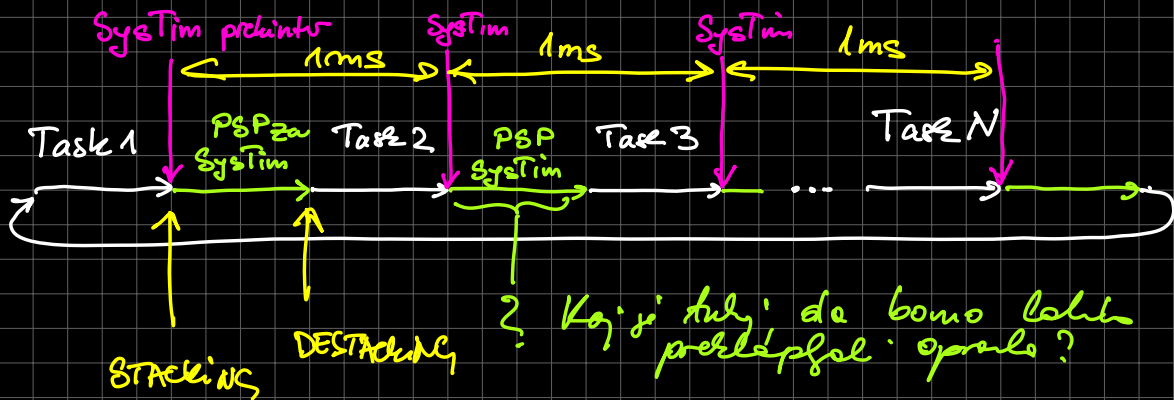
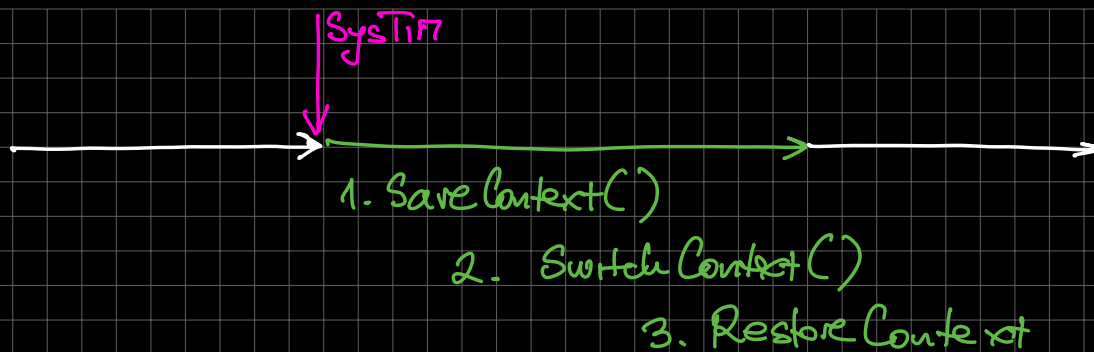


Preklapljanje opravil (task/context switching)



→ vsako opravilo bo imelo svoj sklad v pomnilniku
↳ 1KB



Save Context: na sklad porine registre R4-R11

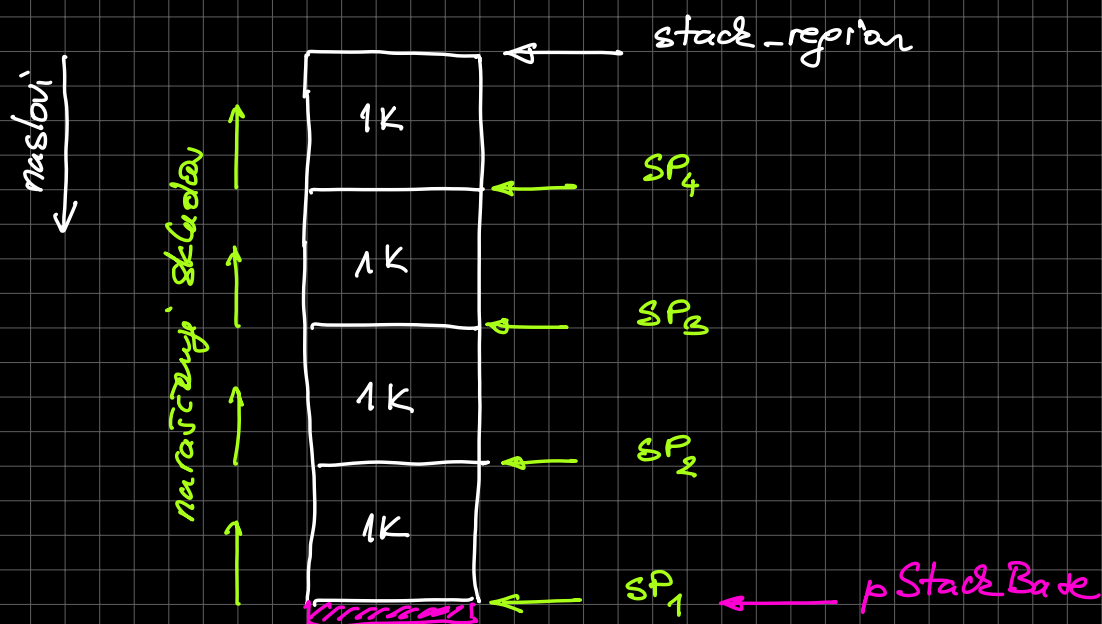
Switch Context: $SP_i \leftrightarrow SP$
 $SP \leftarrow \overline{SP}_j$

Restore Context: S skladu obnovi R4-R11

① Implementacija skladu za vso opravila

→ za vsa opravila 1K zaporednih pomembnosti besed

→ za vsa opravila (N): $N \times 1K$



uint32_t stack_region [$N \times 256$];
 ↑
 stack offset

inline assembler

asm volatile (asm_template :
 input operands:
 output operands

```
static inline void save_context(void) {
    uint32_t reg;

    asm volatile ( "MRS %0, psp\n\t"
                  "STMFID %0!, {r4-r11}\n\t"
                  "MSR psp, %0\n\t" : "=r" (reg) );
}
```

template : tutaj wpisujemy instrukcje
 assembly

```
static inline void restore_context(void) {
    uint32_t reg;

    asm volatile ( "MRS %0, psp\n\t"
                  "LDMFD %0!, {r4-r11}\n\t"
                  "MSR psp, %0\n\t" : "=r" (reg) );
}
```

Za kod opravila bismo imali:

```
81
82 /*
83  Task state structure.
84  In our simple scheduler, the task is determined only by its stack pointer and
85  the start address (i.e. pointer to a function that implements the task)
86  but in general, the structure should contain a task's state variable (flags)
87  (e.g. RUNNING, STOPPED, ACTIVE, ....)
88
89  Pa3cio Bulić, 9.11.2020
90 */
91 typedef struct {
92     void* sp;           // task's stack pointer
93     void (*pTask)(void); // start address of the task
94     // int flags;
95 } task_table_t;
96
```

skladni: Esalec (frenčna
vredost)

✓
Za kod opravila

Kod opravila

```
void Task (void) {
    while (1) {
        //
    }
}
```

12}

2.11. 2021

Implementacija switch - context ():

↑
"zamenjaj trenutni SP s sledovnim
stolcem nove opravila"

Potreben TABELA OPRAVILA

↪ za vsako opravilo hraniti:
SP in krajec ne
kako opravilo
array struktur task_table_t

1. v tabelo opravil shraniti trenutni PSP
v SP prejšnjega opravila
2. iz URNIKA OPRAVIL dobiti index nove
opravila
3. iz tabele opravil prebrati SP in
PIS: v register PSP

Bravi trenutnega PSP-ja:

```
static inline void* read_process_stack_pointer(void) {  
    void* result;  
    /* read special register PSP into a general-purpose register (picked by compiler)  
       and write it to result: */  
    asm volatile ( "MRS %0, PSP\n\t" : "=r" (result) );  
  
    return (result);  
}
```

Pisane v PSP:

```
static inline void write_process_stack_pointer(void* ptr) {  
    asm volatile ( "MSR PSP, %0\n\t" : : "r" (ptr) );  
}
```

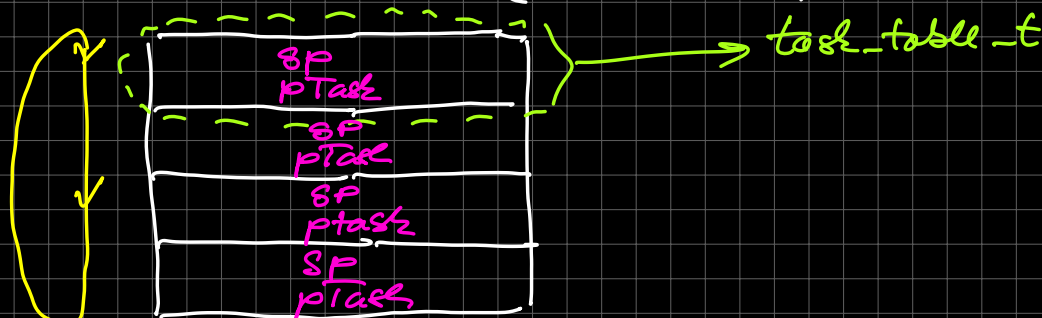
```
void switch_context (void) {  
    // save current stack pointer to current task in task_table:  
    ① task_table[current_task].sp = read_process_stack_pointer();  
  
    // select a new task in a round-robin fashion:  
    ② current_task++;  
    if (current_task == MAX_TASKS) {  
        current_task = 0;  
    }  
  
    // select a new psp:  
    ③ write_process_stack_pointer( task_table[current_task].sp );  
}
```

✓ tabela opravil za trenutno prekinjeno opravilo shraniti PSP

index v tabeli opravil

izberem novo opravilo

TABELA OPRAVIL (task_table)

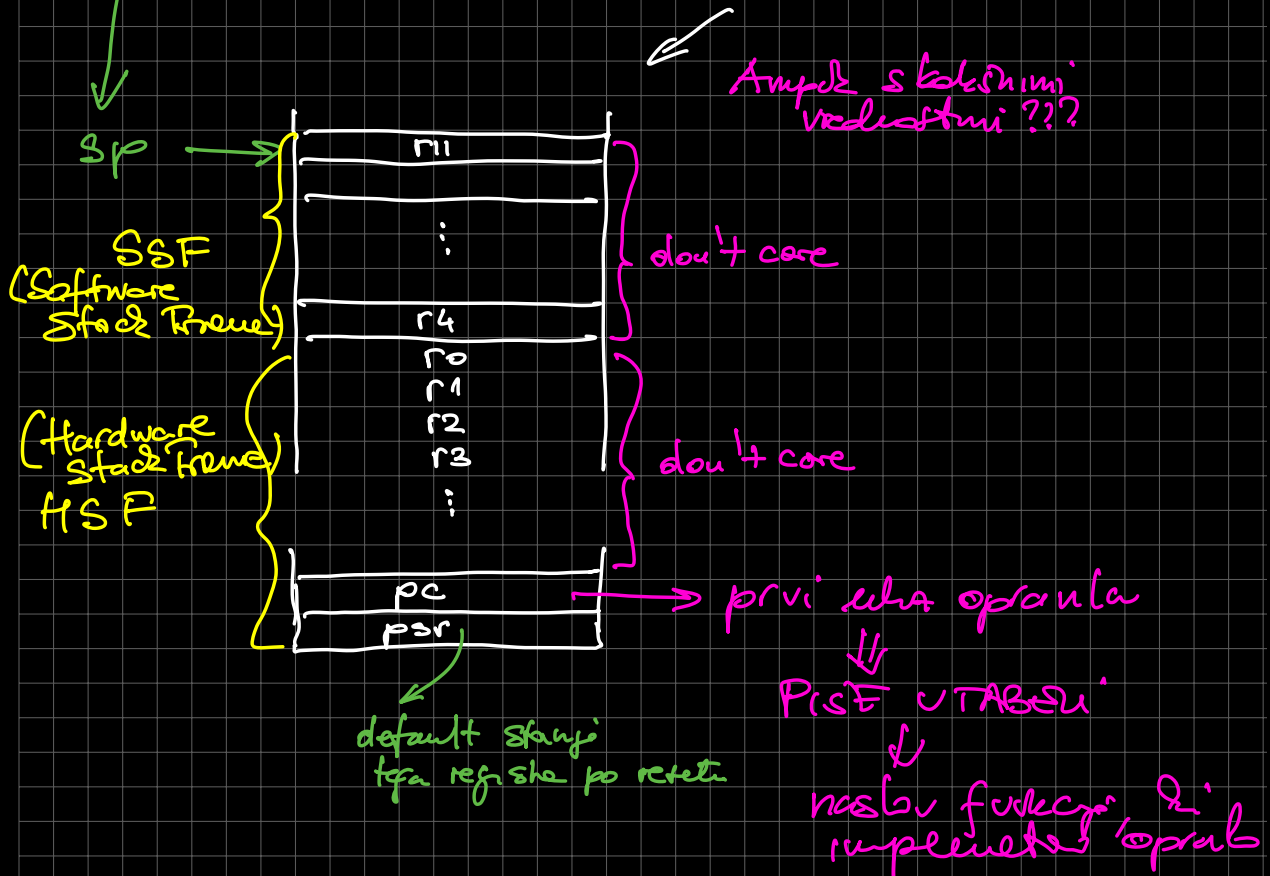


izbirali bomo po pravilu "ROUND ROBIN"

Vsa za opravilo moramo USTVARITI!

1. za vsako opravilo moram imeti nastojen školski katelec SP v upefori strukturi v tabeli.

2. Za vsako opravilo moram imeti
iniciatorno shemo:



✓

Implementacija opravilo

```

void create_task (int task_id, void* task_stack_base_address, void (*pTask)(void)) {
    hw_stack_frame_t* ptask_hw_frame;
    /* set the start address of the HW frame structure */
    ptask_hw_frame = (hw_stack_frame_t *) ((uint8_t *)task_stack_base_address - sizeof(hw_stack_frame_t));
    /* populate the HW stack frame with initial values : */
    /* NOTE: HW stack for Task0 (main()) is ignored as main() uses MAIN STACK */
    ptask_hw_frame -> r0 = 0;
    ptask_hw_frame -> r1 = 0;
    ptask_hw_frame -> r2 = 0;
    ptask_hw_frame -> r3 = 0;
    ptask_hw_frame -> r12 = 0;
    ptask_hw_frame -> lr = 0; // in our simple RTOS the tasks never finish so there is no need to delete the task
    ptask_hw_frame -> pc = (uint32_t)pTask;
    ptask_hw_frame -> psr = 0x01000000;
    // make room for SW stack frame and set the task's stack pointer:
    task_table[task_id].sp = (void *)((uint8_t *)task_stack_base_address
    - sizeof(hw_stack_frame_t)
    - sizeof(sw_stack_frame_t));
}

```

KAZALEC NA HSF

1) **↳ začeti skladu se to opravilo**

2) **↳ začeti poročila v MEM, kjer bo sklad tega opravila**

↳ default vredost na PSR

↳ 3. argument funkcije

↳ bi celotno lahko odstranili

↳ 32B

↳ 32B

Na SSF in HSF lahko gledamo kot na dve podskupini skladu. Če hrenito vsake po 8 32-bitne vredosti

zaradi tega dela naredimo 2. dve strukturi, ki sta preprosta abstrakcija teh delov okoli njer na skladu.

```

/*
Software Stack Frame structure.

Pa3cio Bulic, 9.11.2020
*/
typedef struct {
    uint32_t r4;
    uint32_t r5;
    uint32_t r6;
    uint32_t r7;
    uint32_t r8;
    uint32_t r9;
    uint32_t r10;
    uint32_t r11;
} sw_stack_frame_t;

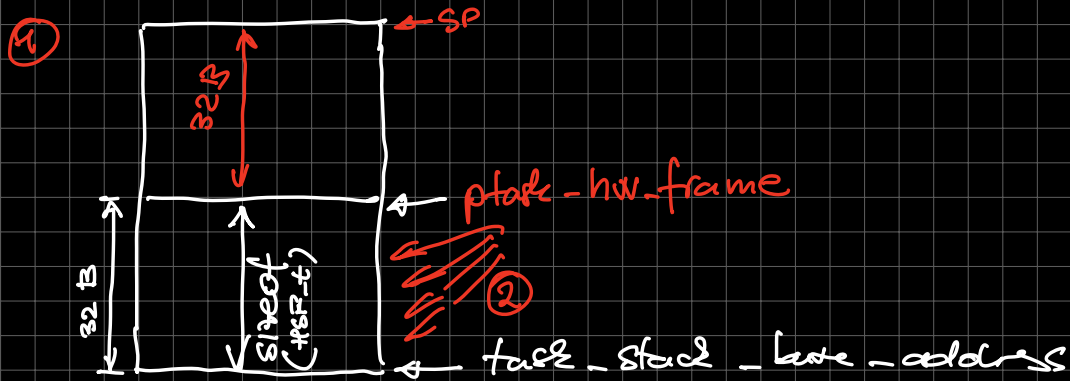
```

```

/*
Hardware Stack Frame structure.

Pa3cio Bulic, 9.11.2020
*/
typedef struct {
    uint32_t r0;
    uint32_t r1;
    uint32_t r2;
    uint32_t r3;
    uint32_t r12;
    uint32_t lr;
    uint32_t pc;
    uint32_t psr;
} hw_stack_frame_t;

```



Zgornjo funkcijo `create_task()` moramo zlasti
 za vsako opravilo posebej:

```
void init_tasks (void) {
    int i;

    /* Task schedule
     sets tasks in the schedule: */
    task_table[0].pTask = &idle; /* this value will be overwritten after first SysTick exception...*/
    task_table[1].pTask = &task1;
    task_table[2].pTask = &idle;
    task_table[3].pTask = &task1;
    task_table[4].pTask = &idle;
    task_table[5].pTask = &task3;
    task_table[6].pTask = &task1;
    task_table[7].pTask = &idle;
    task_table[8].pTask = &task1;
    task_table[9].pTask = &task2;

    /* Create tasks...*/
    for (i = 0; i < MAX_TASKS; i++) {
        create_task (i, (uint8_t *)pTasksStackBase - i*sizeof(uint32_t)*TASKS_STACK_SIZE, task_table[i].pTask);
    }

    /* PSP has not been set until now, so set PSP to point to
     the top of stack of the first interrupted task (Task 0): */
    write_process_stack_pointer(task_table[0].sp);
}
```

transform PC v tabelo
 opravil

VSVAJEN VSA OPRANICA

1 KB