

TDA, Dec. 15 2020

Stability of persistence diagrams

$S = \{N_1, N_2, \dots, N_n\}$ a sample

$S' = \{N'_1, N'_2, \dots, N'_n\}$ perturbed sample

$\{Rips(S, r)\}_r$

Rips filtrations

$\{Rips(S', r)\}_r$

persistence modules (coefficients in a field)

$U = \{H_p(Rips(S, r))\}$ $U' = \{H_p(Rips(S', r))\}$

interleaving distance: $d_i(U, U') = \text{smallest } \varepsilon$
for which an ε -interleaving exists

$D(S), D(S')$ persistence diagrams corresponding
to the two Rips filtrations

multiset of

point $\{a = (r, s)\}$

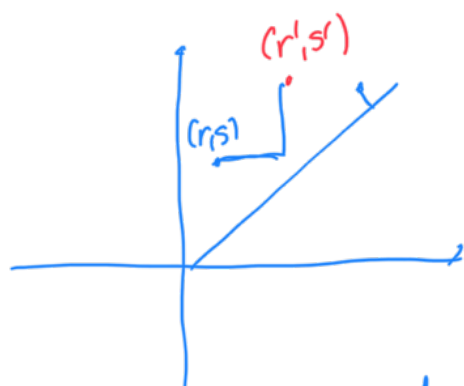
$\leftarrow \# \text{ of generators with this data}$

$\uparrow$ birth-death data

of a homology generator in the filtration

bottleneck distance

$d_b(D(S), D(S')) = \text{minimal cost of matching}$
between $D(S)$ and $D(S')$

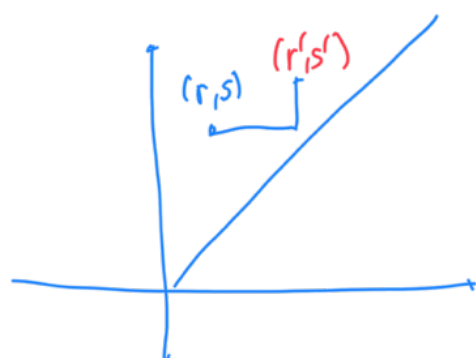
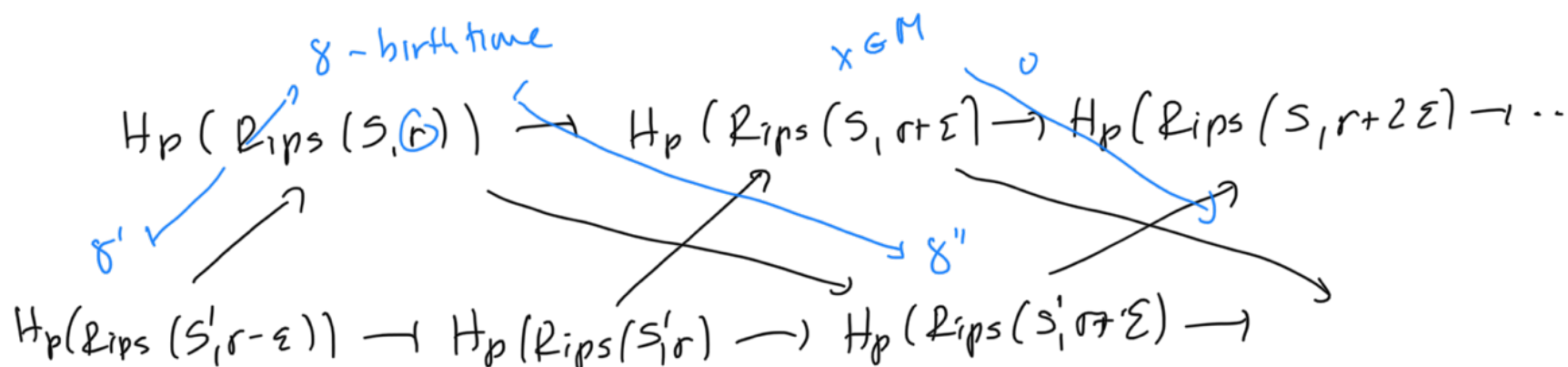


$C(M) = \text{max distance between}$
pairs in the matching and
points outside the matching
to the diagonal

$$d_b((r, s), (r', s')) = \max(|r - r'|, |s - r'|)$$

Theorem (stability) $d_b(D(S), D(S')) = d_h(N, N')$

- Sketch of proof: assume that we have an ε -interleaving between N and N' . We want to show: $d_b(D(S), D(S')) \leq \varepsilon$.



$$\max\{|r-r'|, |s-s'|\} < \varepsilon$$

if $\gamma' \neq 0$ then $r-\varepsilon$ is its birth time

(if it appeared at $r-2\varepsilon$, then γ would appear before r)

$\gamma' = 0$, then γ'' is matched to γ
 $\nearrow$ birth time is $r+\varepsilon$

$$|r-r'| \leq \varepsilon$$

similar argument $|s-s'| \leq \varepsilon \Rightarrow$ if there is a

pair $(\gamma, \gamma') \in M$, then the corresponding birth-death points in the persistence diagrams have distance $\leq \varepsilon$.

What about $x \in M \Rightarrow d_p(x, \Delta) < \varepsilon$

$$\Rightarrow d_b(D(S), D(S')) \leq \varepsilon$$

Recall: $S = \{v_1, \dots, v_n\}$ $S' = \{v'_1, \dots, v'_n\}$

if $d(v_i, v'_i) \leq \varepsilon$ for all i (an ε -perturbation)

then this induces a 2ε -interleaving on the
Rips filtrations, which induces a 2ε -interleaving
of the persistence modules

$$U = \{H_p(\text{Rips}(S, r))\}_r \quad U' = \{H_p(\text{Rips}(S', r))\}_r$$

$\Rightarrow$ by the theorem

$$\boxed{d_b(D(S), D(S')) \leq 2\varepsilon}$$

Stability of Rips persistence diagrams,
if the max distance between S and S'
is ε , then the bottleneck distance
between $D(S)$ and $D(S')$ is $\leq 2\varepsilon$.

Rips persistence diagrams are good descriptors
of the data set

Case 2: $f: X \rightarrow \mathbb{R}$

the sublevel set filtration,

$$X_a = \{x \in X; f(x) \leq a\}$$

$$K: K_a = \{\sigma \in K; f(v) \leq a \text{ for all vertices } v \in \sigma\}$$

$$f: |K| \rightarrow \mathbb{R}$$

If f' is a perturbed wep,

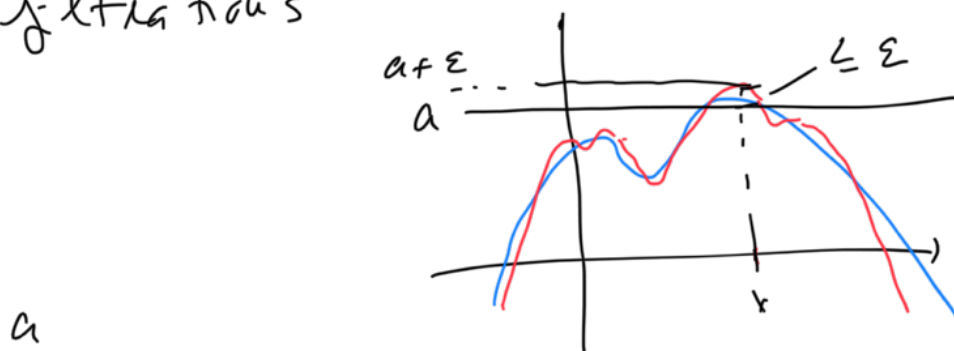
Theorem (stability theorem for maps):

if $\|f - f'\|_\infty = \max_{x \in K} |f(x) - f'(x)| \leq \varepsilon$, then

$$d_b(D(f), D(f')) \leq \varepsilon.$$

↑ ↗
persistence diagrams of the corresponding
sublevel complex filtrations.

Sketch of proof: if $\|f - f'\|_\infty < \varepsilon$, we have
an ε -interleaving on the sublevel complex
filtrations



$$\underbrace{K(f)_a}_{\text{blue}} \subset \underbrace{K(f')_{a+\varepsilon}}_{\text{red}}$$

this gives an ε -interleaving on
the persistence modules which then

implies $d_b(D(f), D(f')) \leq \varepsilon$.

DISCRETE MORSE THEORY

K simplicial complex (finite)

index the simplices $\sigma \in K$ with a function

$$f: K \rightarrow \mathbb{R}$$

$$\sigma \mapsto f(\sigma) \in \mathbb{R}$$

where f has the following properties:

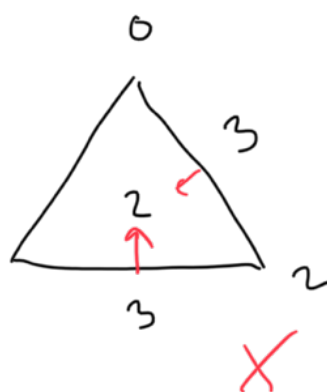
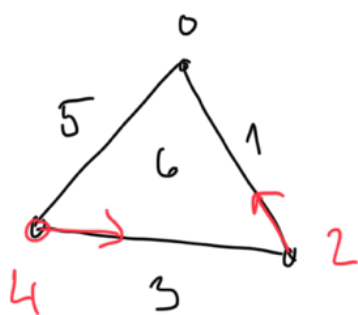
- for any σ^p , the number u_1 of its faces τ^{p-1} such that $f(\tau) \geq f(\sigma)$ is 0 or 1

$$u_1 = 0 \text{ or } 1$$

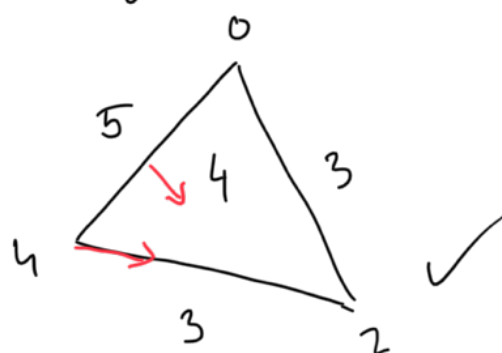
- for any σ^p , the number u_2 of its cofaces ρ^{p+1} such that $f(\rho) \leq f(\sigma)$ is again 0 or 1

The function f is a discrete Morse function

For any σ , values of f strictly decrease toward its faces, with at most one exception and the values strictly increase towards its cofaces, with at most one exception.



triangle has two exceptions



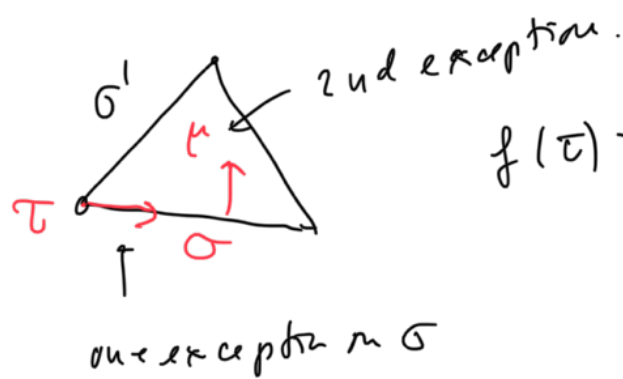
Arrows $\tau^{p-1} \rightarrow \sigma^p$ denote the direction of decreasing function values.

If σ^p is any simplex in K then only one
of the numbers n_1 or n_2 can be 1.

So: for any simplex σ at most one exception
 is allowed.

Proof: assume that a simplex σ is at the same
 time a head and a tail of an arrow, i.e.

$\tau^{p-1} \rightarrow \sigma^p$ and $\sigma^p \rightarrow \mu^{p+1}$ and let $\tau' < \mu^{p+1}$ be a
 different face from τ . Then:



$$f(\tau) > f(\sigma) > f(\mu) > f(\sigma') > f(\tau)$$

$\rightarrow$ a contradiction!

This implies every simplex in K is involved
 in at most one arrow, either as an arrowhead
 or as an arrowtail

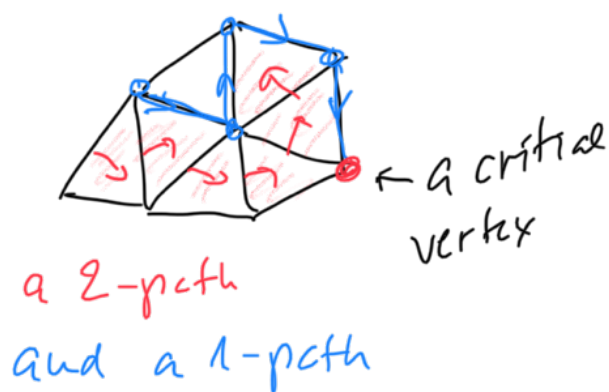
So, we have a partial matching

$\{\tau \rightarrow \sigma\}$ regular pair

$\{\sigma\}$ not in the matching
critical simplex

The collection of arrows $\{\tau \rightarrow \sigma\}$ in a partial matching
 on K is a discrete vector field on K .

The simplices of K are subdivided into critical simplices,
arrow heads, arrow tails $K = A \cup B \cup C$
A B C
disjoint union



a sequence of
consecutive arrows is
a path in the discrete vector
field

(forms a path through sx 's
of two consecutive dimensions in K)

$$\tau_1 \mapsto \sigma_1 > \tau_2 \mapsto \sigma_2 > \dots > \tau_r \mapsto \sigma_r > \tau_{r+1}$$

A p-path in K consists of consecutive arrows,
 where $\tau_{i+1}^{p-1} \neq \tau_i^{p-1}$ for every

Function values of f along the path decrease:

$$f(\tau_1) > f(\sigma_1) > f(\tau_2) > f(\sigma_2) > \dots > f(\tau_{r+1})$$

$\uparrow$ arrow $\uparrow$ face-relation

p-paths do not form cycles.

The discrete vector field arising from a discrete
Morse function on K is acyclic.

Critical simplices do not appear in a p-path, except
maybe at the end, as the ending $(p-1)$ -simplex.

Motivation for discrete Morse theory:
critical cells determine the shape of K .

A sublevel complex of K corresponding to $a \in \mathbb{R}$:

$$K_a = \{ \sigma \in K; f(\sigma) \leq a \} \cup \text{all their faces} =$$

$$= \bigcup_{f(\sigma) \leq a} \left(\bigcup_{\tau \leq \sigma} \tau \right)$$

K_i , a value of f on every simplex of K

K_a - sublevel complexes, form a filtration of K .

$$a_1 \leq a_2 \leq a_3 \leq \dots \leq a_n \quad \text{the values of } f$$

assume that f is injective (for now)

$$a_1 < a_2 < a_3 < \dots < a_n$$

sublevel complex filtration

$$K_1 < K_2 < K_3 < \dots < K_n = K$$

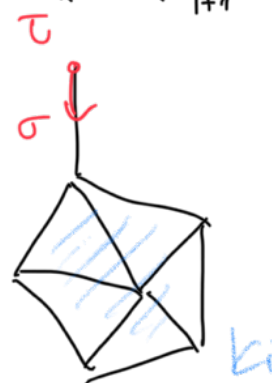
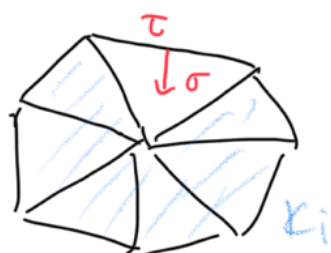
$$K_i = K_{a_i}$$

what happens from K_i to K_{i+1} .

two types of events:

① assume that $a_{i+1} = f(\sigma_{i+1})$ where
where σ_{i+1} belong to a pair $\tau^{p-1} \mapsto \sigma^p$

$$\text{so } K_{i+1} = K_i \cup (\sigma_{i+1}^p \cup \tau^{p-1})$$



$$K_{i+1} = K_i \cup \{\tau \mapsto \sigma\}$$

In this case K_{i+1} collapses onto K_i by
an elementary collapse: $K_{i+1} \searrow K_i$

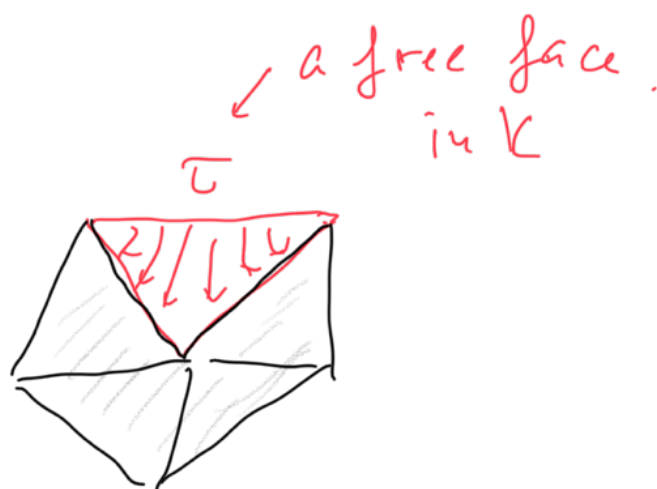
An elementary collapse on a simplicial complex
 K consists of removing a pair of simplices $(\tau^{p-1} \leq \sigma^p)$
where τ is a free face of σ , that is, it is not the
face of any other simplex in K .

$$K \searrow K - \{\sigma, \tau\}.$$

An elementary collapse
is a homotopy equivalence:

$$|K| \simeq |K - \{\sigma, \tau\}|.$$

A continuous deformation
pushes τ and σ onto the remaining boundary of σ .



Combinatorially, an elementary collapse corresponds to removing the pair $(\tau < \sigma)$ from K .

$$K_1 \subseteq K_2 \subseteq \dots \subseteq K_i \subseteq K_{i+1} \subseteq \dots \subseteq K_n = K$$


 $K_i = K_{a_i}$

if $a_{i+1} = f(\sigma^p)$, where $\tau^p \mapsto \sigma^p$ is a regular pair

then exists an elementary collapse

$$K_{i+1} = K_i \cup \{\sigma, \tau\} \searrow K_i$$

$$K_{i+1} \simeq K_i$$

The next stage K_{i+1} is homotopy equivalent to K_i , the homology groups remain the same.

the step from K_i to K_{i+1} is a regular event in the filtration.

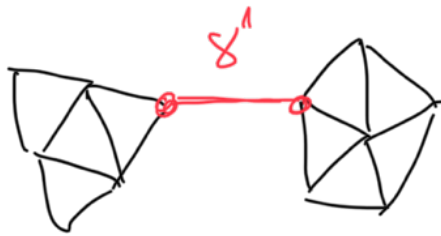
$$\textcircled{2} \quad K_{i+1} = K_i \cup \delta$$

a critical simplex

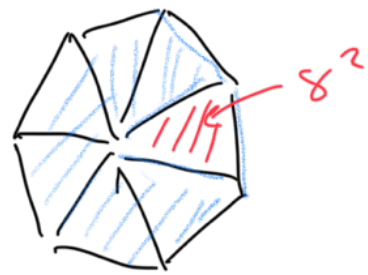
$$a_{i+1} = f(\delta)$$

in this case all the faces of δ are included already in K_i

deaths



δ^1 kills a
homology generator in
dimension $p-1$
a death

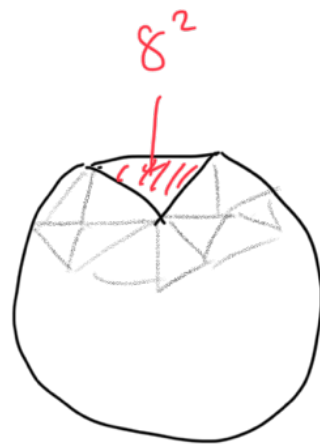


δ^2 kills a
homology generator
in dimension 1
a death

births



δ^1 is the birth of
a new 1-dimensional
homology generator



δ^2 form a new
2-dimensional homology
generator

Adding a critical simplex, $K_{i+1} = K_i \cup \underbrace{\sigma}_{\text{critical}}$

results in a change in the homology from K_i to K_{i+1} :

If σ^p is a p -dimensional critical simplex, then it either:

① completes a new p -cycle, so

$$\beta_p(K_{i+1}) = \beta_p(K_i) + 1 \quad \underline{\text{or}}$$

② forms a new boundary which fills in an existing $(p-1)$ -cycle, so

$$\beta_{p-1}(K_{i+1}) = \beta_{p-1}(K_i) - 1$$

This tells us: critical simplices affect the shape of K .

Special case: K has only one critical simplex

The unique critical simplex σ must correspond to the minimum of f

Since then $f(\sigma) \leq f(\tau)$ for any coface τ of σ ,

σ must be a vertex v of K .

In the sublevel complex filtration

$$\{N_0\} = K_0 \subseteq K_1 \subseteq K_2 \dots$$

all events are regular so there is an elementary collapses $K_{i+1} \searrow K_i$ for all i .

$$\text{So, } K = K_n \simeq K_{n-1} \simeq \dots \simeq \{N_0\}$$

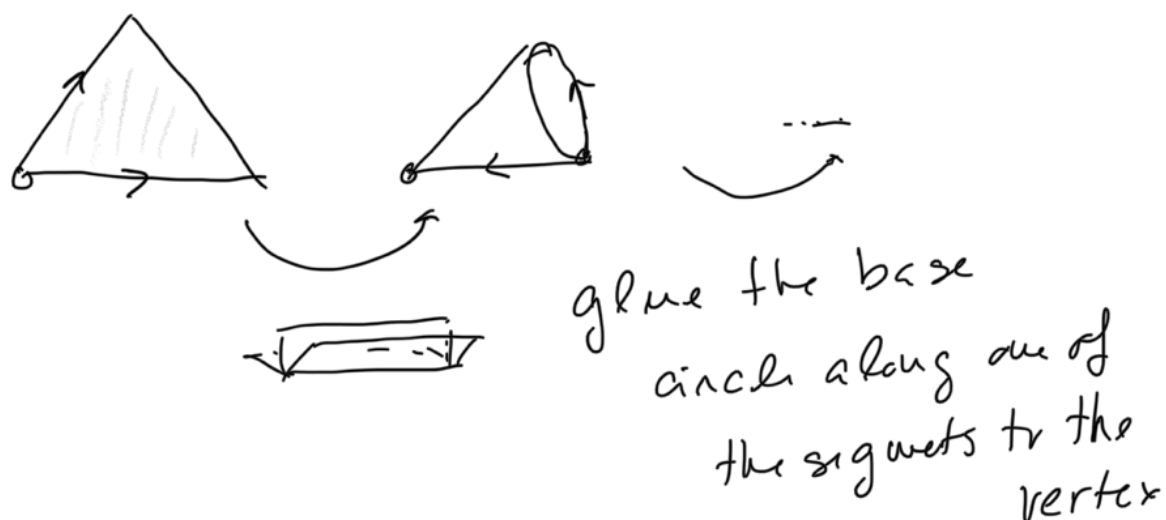
K is contractible.

Definition: A simplicial complex K is collapsible if there exists a sequence of elementary collapses to a point,

$$K \searrow \dots \searrow \{N\}$$

Collapsibility is a stronger property than contractibility: there exist contractible simplicial complexes that are not collapsible.

Example: dunce hat



this is contractible but not collapsible: for any triangulation, no free face exists.

Algorithmically collapsibility is much easier to implement than contractibility: as long as a free face exists, a collapse can be done.

Morse homology of K .

K , discrete Morse function f on K
 $C = \{\text{critical simplices of } f\}$

Recall: to define homology groups we first
 define the simplicial chain complex

$$0 \rightarrow C_k(K) \xrightarrow{\partial_k} C_{k-1}(K) \xrightarrow{\partial_{k-1}} \dots \rightarrow C_1(K) \xrightarrow{\partial_1} C_0(K) \rightarrow 0$$

$C_i(K)$ is an abelian group
 (a vector space if coefficients are $\mathbb{F}$)
 generated by all i -dimensional σ 's

$$\partial_i : C_i(K) \rightarrow C_{i-1}(K) \quad \text{boundary homomorphism}$$

$$\sigma = \langle N_0, N_1, \dots, N_i \rangle$$

$$\partial_i(\sigma) = \sum_{j=0}^i (-1)^j \langle N_0, \dots, N_{j-1}, \hat{N}_j, N_{j+1}, \dots, N_i \rangle$$

Homology groups: $H_i(K) = \frac{Z_i(K)}{B_i(K)} = \frac{\ker \partial_i}{\text{im } \partial_{i+1}}$

$(\partial \circ \partial = 0)$. ↖ simplicial
 homology groups

The Morse chain complex:

$(K, f: K \rightarrow \mathbb{R})$
 C : critical pts of f

$K, f: K \rightarrow \mathbb{R}$ Morse function

C critical simplices of f .

$$0 \rightarrow M_k(K) \xrightarrow{\partial_k} M_{k-1}(K) \xrightarrow{\partial_{k-1}} \dots \rightarrow M_1(K) \xrightarrow{\partial_1} M_0(K) \rightarrow 0$$

where $M_i(K)$, the Morse i -th chain group,
is generated by the critical pts of dimension i

a Morse chain a formal sum

$$\sum_j a_j \gamma_j^i$$

j runs over
all critical pts
of dim i .

What are the boundary homomorphisms?